

# Repeated Games

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## Part A

# Background Material



## CHAPTER I

### Basic Results on Normal Form Games

Non-cooperative games (or strategic games) are mainly studied through two models: normal form or extensive form. The latter will be presented in ch. II. The former describes the choice spaces of each player and the result of their common choices. This is evaluated in terms of the players' von Neumann-Morgenstern utilities (i.e. the utility of a random variable is its expected utility (von Neumann and Morgenstern, 1944, ch. I, 3.5)), hence the following definition:

A **normal form** game is defined by a set of **players  $\mathbf{I}$** , **strategy spaces**  $S^i$ ,  $i \in \mathbf{I}$ , and **real pay-off functions**  $F^i$ ,  $i \in \mathbf{I}$ , on  $S = \prod_i S^i$ .

It is finite (or a **bi-matrix game** if  $\#\mathbf{I} = 2$ ) if  $\mathbf{I}$  and all  $S^i$  are finite.

Under suitable measurability conditions one defines the **mixed extension** of a game  $G = (\mathbf{I}, (S^i, F^i)_{i \in \mathbf{I}})$  as the game  $\Gamma = (\mathbf{I}, (\Sigma^i, \phi^i)_{i \in \mathbf{I}})$  where  $\Sigma^i$  is the set of probabilities on  $S^i$  and  $\phi^i(\sigma) = \int_S F^i(s) \prod_{i \in \mathbf{I}} \sigma^i(ds^i)$ . Unless explicitly specified (or self-evident) otherwise, the following definitions are always used on the mixed extension of the game.

$s^i$  is a **dominant strategy** of player  $i$  if  $F^i(s^i, s^{-i}) \geq F^i(t^i, s^{-i})$  for all  $t^i$  in  $S^i$ , and  $s^{-i} \in S^{-i} = \prod_{h \neq i} S^h$ .

$s^i$  is **dominated** (resp. strictly dominated) if there exists  $t^i$  with  $F^i(t^i, s^{-i}) > F^i(s^i, s^{-i})$  for some  $s^{-i}$  (resp. all  $s^{-i}$ ) and  $F^i(t^i, \cdot) \geq F^i(s^i, \cdot)$ .

$s^i$  is a  $(\varepsilon)$ -**best reply** to  $s^{-i}$  if  $F^i(s^i, s^{-i}) \geq F^i(t^i, s^{-i}) (-\varepsilon)$  for all  $t^i \in S^i$ .

An  $(\varepsilon)$ -**equilibrium** is an  $\mathbf{I}$ -tuple  $s$  such that for every  $i$ ,  $s^i$  is a  $(\varepsilon)$ -best reply to  $s^{-i}$ .

A **two-person zero-sum game** (or: a **matrix game** if strategy sets are finite) is a normal form game with  $\mathbf{I} = \{\text{I}, \text{II}\}$ ,  $S^{\text{I}} = S, S^{\text{II}} = T$ ,  $F^{\text{I}} = g = -F^{\text{II}}$ . One defines then the **minmax**  $\bar{v}(g) = \inf_{t \in T} \sup_{s \in S} g(s, t)$  and the **maxmin**  $\underline{v}(g) = \sup_{s \in S} \inf_{t \in T} g(s, t)$ .

If they are equal the game has a **value**  $v(g)$ .

$s$  is an  $(\varepsilon)$ -**optimal strategy** for player I if  $g(s, t) \geq \underline{v}(g) (-\varepsilon)$  for all  $t \in T$ .

#### 1. The minmax theorem

A minmax theorem gives conditions under which a two-person zero-sum game  $(S, T; g)$  has a value. We allow here the pay-off function  $g$  to map  $S \times T$  to  $\overline{\mathbb{R}}$ . All minmax theorems in this paragraph will be derived from Theorem 1.6. Prop. 1.8 uses additional convexity assumptions to weaken continuity requirements, and Prop. 1.17 applies the previous result to the mixed extension of the game.

**1.a. Definitions and notations.** We start by introducing basic definitions and notations.

DEFINITION 1.1. A convex set is a convex subset of a vector space on the reals.

DEFINITION 1.2. An admissible topology on a convex set  $S$  is a topology such that, for each  $n$ , and for every  $n$ -tuple  $x_1, \dots, x_n$  of points in  $S$ , the mapping  $\phi_{x_1, \dots, x_n}$  from the  $n-1$  dimensional simplex  $\{p = (p_1, \dots, p_n) \mid p_i \geq 0, \sum_i p_i = 1\}$  to  $S$ , that maps  $p$

to  $\sum_{i=1}^n p_i x_i$ , is continuous, when the simplex is endowed with its usual topology (i.e. the topology induced by the Euclidian norm).

**DEFINITION 1.3.** A convex topological space is a convex set endowed with an admissible topology.

**REMARK 1.1.** Any convex subset of a linear topological (or: topological vector) space is a convex topological space.

**REMARK 1.2.** A compact space is not assumed to be Hausdorff ( $T_2$ ), unless explicitly stated.

**DEFINITION 1.4.** For a topological space  $S$ , a function  $f: S \rightarrow \mathbb{R} \cup \{-\infty\}$  is **upper semi-continuous (u.s.c.)** iff  $\{f \geq x\}$  is closed for every  $x$ .  $f$  is **lower semi-continuous (l.s.c.)** iff  $\{-f\}$  is u.s.c.

**REMARK 1.3.** On a completely regular space, the u.s.c. functions that are bounded from above are the infima of a family of bounded continuous functions.

**DEFINITION 1.5.** For a convex set  $S$ ,  $f: S \rightarrow \overline{\mathbb{R}}$  is **quasi-concave** iff  $\{f \geq x\}$  is convex for every  $x$ .  $f$  is **quasi-convex** iff  $\{-f\}$  is quasi-concave.  $f$  is **concave** (resp. **convex, affine**) iff  $f(\alpha x_1 + (1 - \alpha)x_2) \leq \alpha f(x_1) + (1 - \alpha)f(x_2)$  (resp.  $\geq, =$ ) whenever the right hand member is well defined ( $0 < \alpha < 1$ ).

### 1.b. A basic theorem.

**THEOREM 1.6. (Sion, 1958)** Assume  $S$  and  $T$  are convex topological spaces, one of which is compact. Assume that, for every real  $c$ , the sets  $\{t \mid g(s_0, t) \leq c\}$  and  $\{s \mid g(s, t_0) \geq c\}$  are closed and convex for every  $(s_0, t_0) \in S \times T$ . Then

$$\sup_{s \in S} \inf_{t \in T} g(s, t) = \inf_{t \in T} \sup_{s \in S} g(s, t).$$

If  $S$  (resp.  $T$ ) is compact, then  $\sup$  (resp.  $\inf$ ) may be replaced by  $\max$  (resp.  $\min$ ), i.e. the corresponding player has an **optimal** strategy.

**PROOF.** Obviously, one always has

$$\sup_{s \in S} \inf_{t \in T} g(s, t) \leq \inf_{t \in T} \sup_{s \in S} g(s, t).$$

Suppose therefore, in contrary to the theorem that, for some real number  $c$ ,

$$\sup_{s \in S} \inf_{t \in T} g(s, t) < c < \inf_{t \in T} \sup_{s \in S} g(s, t).$$

Assume, for instance, that  $S$  is compact. The sets  $A_t = \{s \in S \mid g(s, t) < c\}$  form an open covering of  $S$ ; thus, there exist  $t_1, \dots, t_n$  such that the sets  $A_{t_i}$  ( $1 \leq i \leq n$ ) cover  $S$ : We can restrict  $T$  to the convex hull of the set  $T_0$  of points  $t_i$  ( $1 \leq i \leq n$ ) and work now on the  $(n - 1)$ -simplex, using the admissibility of the topology. All our assumptions are still valid. We can now do the same operation with the sets  $B_s = \{t \in T \mid g(s, t) > c\}$ , so that both  $S$  and  $T$  become simplices with vertex sets  $S_0$  and  $T_0$ , and with the property that, for any  $s \in S$ , there exists  $t \in T_0$  such that  $s \in A_t$  and that, for any  $t \in T$ , there exists  $s \in S_0$  such that  $t \in B_s$ . We can further assume that  $S_0$  and  $T_0$  are minimal for this property, eliminating eventually some additional points from  $S_0$  and  $T_0$ .

Let  $s_i$  ( $1 \leq i \leq n$ ) be the points in  $S_0$ , and for each  $i$ ,  $T_i = \{t \in T \mid g(s_i, t) \leq c\}$ . The sets  $T_i$  are compact convex, satisfy  $\bigcap_{i=1}^n T_i = \emptyset$ , and, by the minimality property, for every  $j$ ,  $\bigcap_{i \neq j} T_i \neq \emptyset$ . The next lemma will show that this implies that  $\bigcup_{i=1}^n T_i$  is not convex.

Thus there exists  $t_0 \in T$  such that, for all  $i$ ,  $t_0 \notin T_i$ : we have, for each  $i$ ,  $g(s_i, t_0) > c$ , and thus, for each  $s \in S$ ,  $g(s, t_0) > c$ . The same argument would show dually that there exists  $s_0 \in S$  such that, for each  $t \in T$ ,  $g(s_0, t) < c$ ; we thus have both  $g(s_0, t_0) < c$ , and  $g(s_0, t_0) > c$ , the desired contradiction.

The replacement of  $\sup$  by  $\max$  is possible due to the upper semi-continuity of  $g(\cdot, t)$  on the compact set  $S$  (for each  $t \in T$ ). ■

To complete the proof, we shall thus prove the following lemma:

**LEMMA 1.7.** *Let  $T_i$  ( $1 \leq i \leq n$ ) be compact convex subsets of a Hausdorff locally convex topological vector space, such that  $T = \bigcup_i T_i$  is convex and such that, for every  $i$ ,  $\bigcap_{h \neq i} T_h \neq \emptyset$ . Then  $\bigcap_i T_i \neq \emptyset$ .*

**REMARK 1.4.** For a simpler and more general result, cf. ex. I.4Ex.19 p. 48. Here the only “sport” is to obtain theorem 1.6 using just the separation theorem — and in fact only for polyhedra in finite dimensional space.

**PROOF.** The proof goes by induction on  $n$ . We assume that the lemma is proved up to  $n - 1$ , and is false for  $n$ . Then  $\bigcap_{i < n} T_i$  and  $T_n$  are two disjoint compact convex sets, and can therefore, by the Hahn-Banach theorem (cf. 1.e p. 8), be strongly separated by a closed hyperplane, whose (compact convex) intersection with  $T$  we denote by  $\tilde{T}$ . Let  $\tilde{T}_i = T_i \cap \tilde{T}$ ; the  $\tilde{T}_i$  are compact convex,  $\tilde{T}_n = \emptyset$ , and  $\bigcup_{i < n} \tilde{T}_i = \tilde{T}$ .

Further, for any  $j < n$ ,  $\bigcap_{i \notin \{n, j\}} T_i$ , which is convex, has, by assumption, a non-empty intersection both with  $T_n$  and with  $\bigcap_{i < n} T_i$ , which lie on opposite sides of the hyperplane; therefore,  $\bigcap_{i \notin \{n, j\}} T_i$  has a non-empty intersection with the hyperplane. Thus  $\bigcap_{i \notin \{n, j\}} \tilde{T}_i \neq \emptyset$ .

It follows by the validity of the lemma for  $n - 1$  that  $\bigcap_{i < n} \tilde{T}_i \neq \emptyset$ , i.e.  $(\bigcap_{i < n} T_i) \cap \tilde{T} \neq \emptyset$ ; this yields the contradiction. ■

### 1.c. Convexity.

**PROPOSITION 1.8.** *Assume  $S$  is a compact convex topological space,  $T$  is a convex set, and, for every real  $c$ , and for every  $(s_0, t_0) \in S \times T$ , the sets  $\{s \in S \mid g(s, t_0) \geq c\}$  are closed and convex, and  $g(s_0, t)$  is convex in  $t$ , and  $g < +\infty$ . Then*

$$\max_{s \in S} \inf_{t \in T} g(s, t) = \inf_{t \in T} \max_{s \in S} g(s, t).$$

**PROOF.** As in the proof of the theorem, we can reduce the discussion to the case where  $T$  is a simplex with vertices  $(t_1, \dots, t_k)$  and where

$$\sup_{s \in S} \inf_i g(s, t_i) < c < \inf_{t \in T} \sup_{s \in S} g(s, t).$$

Then, for every  $s \in S$ , the function  $g(s, \cdot)$  is continuous on the interior of the simplex  $T$ , being convex. If  $t_i^n$  are interior points,  $\lim_{n \rightarrow \infty} t_i^n = t_i$ , then, for every  $s$ ,  $\lim_{n \rightarrow \infty} g(s, t_i^n) \leq g(s, t_i)$  by the convexity of  $g$ . Therefore, the sets  $A_{i,n} = \{s \mid g(s, t_i^n) < c\}$ ,  $1 \leq i \leq k$ ,  $n = 1, 2, \dots$ , form an open covering of  $S$ ; extracting a finite subcovering, we see that we can replace  $T$  by some compact polyhedron  $\tilde{T}$  contained in the interior of  $T$ : now we have the continuity and convexity in  $t$  of  $f(t, s)$  for every  $s \in S$ , and we can apply theorem 1.6 p. 4 to yield a contradiction. ■

### 1.d. Mixed strategies.

DEFINITION 1.9. A **regular measure** on a topological space  $X$  is a positive bounded measure  $\mu$  on the Borel sets ( $\sigma$ -algebra generated by the open sets) such that  $\mu(A) = \sup\{\mu(B) \mid B \subseteq A, B \text{ closed and compact}\}$ . Its **support** is the smallest closed subset with negligible complement. The set of those measures which have total mass 1 is denoted by  $\Delta(X)$ .

DEFINITION 1.10.  $\Delta(X)$  is endowed with the **weak $^*$  topology** defined as the weakest topology for which the mapping  $\mu \mapsto \mu(f) = \int f d\mu$  is u.s.c. for all bounded u.s.c.  $f$ .

DEFINITION 1.11.  $f: X \rightarrow Y$  is **Lusin-measurable** if  $\forall \varepsilon > 0, \exists K$  closed, compact with  $\mu(K) > 1 - \varepsilon$  and  $f|_K$  continuous.

REMARK 1.5. The regularity of a measure is equivalent to the Lusin-measurability of all Borel maps with values in separable metric spaces.

PROPOSITION 1.12.  $\Delta(X)$  is always  $T_1$ , and is Hausdorff, resp. compact, resp. completely regular, if  $X$  is so. Further, if  $X$  is completely regular, the above definition coincides with the usual one — using the integral of bounded continuous functions.

PROOF.

- For the  $T_1$  property, i.e. points are closed: If  $\mu_1 \neq \mu_2$ , by regularity there exists an open set  $O$  with  $\mu_1(O) > t > \mu_2(O)$  and  $\mu_2$  does not belong to  $W(\mu_1) = \{\mu \mid \mu(O) > t\}$ .
- For the Hausdorff property, assume  $\mu_1 \neq \mu_2$ . Given a Borel set  $B$  with  $\mu_1(B) > \mu_2(B)$ , choose (regularity) two compacts  $C_1$  (included in  $B$ ) and  $C_2$  (in its complement) satisfying  $\mu_1(C_1) + \mu_2(C_2) > 1$ . Then there exist disjoint open sets  $O_i$ ,  $C_i \subseteq O_i$ ,  $i = 1, 2$  and real numbers  $\alpha_i$ ,  $\alpha_1 + \alpha_2 > 1$  with  $\mu_i(O_i) > \alpha_i$ . Thus the following are disjoint neighbourhoods:  $V_i(\mu_i) = \{\mu \mid \mu(O_i) > \alpha_i\}$ .
- For the compactness, cf. ex. I.1Ex.10 p. 13.
- For the completely regular case, the above remark (sub 1.a p. 3) on u.s.c. functions implies that it is enough to show that  $\int f d\mu = \inf\{\int g d\mu \mid g \text{ continuous, bounded and above } f\}$ . This in turn follows from the fact that  $f$  is Lusin-measurable and from Dini's theorem. ■

PROPOSITION 1.13. (Dudley, 1968) If  $X$  is metrisable and separable, so is  $\Delta(X)$  with the metric  $d(\mu, \nu) = \sup\{|\mu(f) - \nu(f)| \mid f \text{ Lipschitz with constant 1 and bounded by 1}\}$ . Moreover if  $X$  is complete, so is  $\Delta(X)$ .

For more properties, cf. App.9.

FUBINI'S THEOREM 1.14. Given positive bounded regular measures  $\mu_1$  and  $\mu_2$  on two topological spaces  $E_1$  and  $E_2$ , there exists a unique regular measure  $\mu$  on  $E = E_1 \times E_2$  such that  $\mu(K_1 \times K_2) = \mu_1(K_1)\mu_2(K_2)$  for any closed compact subsets  $K_i$  of  $E_i$ , and then for any non-negative  $\mu$ -measurable function  $f$  on  $E_1 \times E_2$ ,  $\int_{E_2} f d\mu_2$  is  $\mu_1$ -measurable and  $\int f d\mu = \int d\mu_1 \int f d\mu_2$ .

PROOF. For existence, extract an increasing sequence of closed compact subsets  $C_i^n$  from  $E_i$  that carry most of the mass. Let  $\mu_i^n$  be the corresponding restriction of  $\mu_i$ . By the validity of the theorem in the compact case (Mertens, 1986a, prop. 2), we obtain corresponding regular measures  $\mu^n = \mu_1^n \otimes \mu_2^n$  on  $C^n = C_1^n \times C_2^n \subseteq E$ . Clearly the  $\mu_i^n$  are increasing, so their limit  $\mu$  is a regular measure on  $E$  for which our condition is easily verified. For any such  $\mu$ , observe that the regularity of  $\mu_1$  and  $\mu_2$  imply then the validity

of our product formula when the  $K_i$  are  $\mu_i$  measurable. Hence the product  $\mathcal{E}_1 \otimes \mathcal{E}_2$  of the  $\sigma$ -fields  $\mathcal{E}_i$  of  $\mu_i$ -measurable sets is contained in the  $\mu$ -measurable sets and  $\mu$  coincides with the product measure there. It follows immediately that  $\mu$  is uniquely determined on closed compact subsets of  $E$ , since those have a basis of neighbourhoods (regularity) that belong to the product of the Borel  $\sigma$ -fields (compactness). Therefore (regularity again),  $\mu$  is unique. ■

**REMARK 1.6.** The above applied inductively yields the existence of a product also for  $n$  factors, and the uniqueness proof remains identical, so the product is “associative”.

**PROPOSITION 1.15.** *The product of regular probabilities is a continuous map from  $\Delta(X) \times \Delta(Y)$  to  $\Delta(X \times Y)$ .*

**PROOF.** For  $O$  open in  $X \times Y$  and  $\mu \in \Delta(X)$ ,  $f(\mu, y) = \mu(O_y)$  is l.s.c. on  $\Delta(X) \times Y$ : choose  $K_0$  closed compact  $\subseteq O_{y_0}$  with  $\mu_0(K_0) \geq \mu_0(O_{y_0}) - \varepsilon$ , then  $O^1$  and  $O^2$  open in  $X$  and  $Y$  respectively with  $K_0 \subseteq O^1, y_0 \in O^2, O^1 \times O^2 \subseteq O$ , then:

$$\liminf_{\mu \rightarrow \mu_0, y \rightarrow y_0} \mu(O_y) \geq \liminf_{\mu \rightarrow \mu_0, y \rightarrow y_0} \mu(O^1) \mathbb{1}_{O^2}(y) \geq \mu_0(O^1) \geq \mu_0(O_{y_0}) - \varepsilon.$$

So  $f(\mu, y) = \sum_{i=1}^n \alpha_i \mathbb{1}_{U^i}$  up to a uniform  $\varepsilon$ , with  $\alpha_i > 0$  and  $U^i$  open in  $\Delta(X) \times Y$ . Hence (Fubini)  $(\mu \otimes \nu)(O) = \int f(\mu, y) \nu(dy) = \sum_{i=1}^n \alpha_i \nu(U_\mu^i)$  up to  $\varepsilon$ , and  $\nu(U_\mu^i)$  is l.s.c. on  $\Delta(X) \times \Delta(Y)$  by our previous argument. ■

**PROPOSITION 1.16.** *Let  $X$  and  $Y$  be topological spaces, with  $Y$  Hausdorff,  $\mu$  be a regular measure on  $X$  and  $f$  be a Lusin-measurable function from  $X$  to  $Y$ .*

- (1) *The image of  $\mu$  by  $f$ ,  $\mu \circ f^{-1}$ , is a regular measure on  $Y$ . Further if  $f$  is continuous, the mapping  $\bar{f}$  from  $\Delta(X)$  to  $\Delta(Y)$ :  $\mu \mapsto \mu \circ f^{-1}$  is continuous (and also denoted  $\Delta(f)$ ).*
- (2) *Let  $(X_k, \rho_{k,\ell})$  be a projective system of Hausdorff spaces ( $\rho_{k,\ell}: X_k \rightarrow X_\ell$  being the continuous projection for  $\ell \leq k$ ), with projective limit  $(X, \rho_k)$ . Given a consistent sequence of regular measures  $\mu_k$  on  $X_k$  (i.e. with  $\bar{\rho}_{k,\ell}(\mu_k) = \mu_\ell$ ) there exists a unique regular measure  $\mu$  on  $X$  with  $\bar{\rho}_k(\mu) = \mu_k$  — the **projective limit** of the sequence  $\mu_k$ .*

**PROOF.** The first point is clear. For the second, cf. Bourbaki (1969, §4, n°3, Théorème 2). ■

In the sequel, we denote by  $\Sigma$  (resp.  $\mathcal{T}$ ) the space  $\Delta(S)$  (resp.  $\Delta(T)$ ), and by  $\mathcal{T}_f$  the space of all probability measures with finite support on  $T$  (points in  $\Sigma$  are **mixed strategies**, points in  $\mathcal{T}_f$  mixed strategies with finite support).

**PROPOSITION 1.17.** *Let  $S$  be a compact topological space,  $T$  any set. Assume that, for each  $t$ ,  $g(\cdot, t)$  is upper semi-continuous in  $s$ . Then*

$$\max_{\sigma \in \Sigma} \inf_{\tau \in \mathcal{T}_f} \int g(s, t) d(\sigma \otimes \tau) = \inf_{\tau \in \mathcal{T}_f} \max_{\sigma \in \Sigma} \int g(s, t) d(\sigma \otimes \tau).$$

**PROOF.**  $\Sigma$  with the weak\* topology is compact convex (1.12) and  $F(\sigma, \tau)$  defined on  $\Sigma \times \mathcal{T}_f$  by

$$F(\sigma, \tau) = \int g(s, t) d(\sigma \otimes \tau) = \int_T d\tau \int_S g(s, t) d\sigma = \int_S d\sigma \int_T g(s, t) d\tau$$

is affine in each variable and upper semi-continuous in  $\sigma$ . We can therefore apply prop. 1.8 p. 5 to yield the equality. ■

**REMARK 1.7.** With an appropriate redefinition of regular measure the above remains true even if  $S$  is countably compact instead of compact (Mertens, 1986a, pp. 243–246, and remark 3 p. 247).

**1.e. Note on the separation theorem.** The Hahn-Banach theorem was used in the first lemma, also in some exercises below, and in many other circumstances. Here follows a short refresher.

*In the following statements,  $E$  is a real topological vector space, all subsets considered are convex,  $U$  denotes an open subset, and letters  $f$  linear functionals.*

The basic result is

**PROPOSITION 1.18.** *If  $0 \notin U$ ,  $\exists f: f(U) > 0$  ( $f$  is then clearly continuous).*

**PROOF.** Apply Zorn's lemma to the open convex subsets disjoint from 0. ■

The basic technique for separating two convex sets is to separate their difference from zero. So

**PROPOSITION 1.19.** *If  $A \cap U = \emptyset$ ,  $\exists f: f(A) \cap f(U) = \emptyset$ .*

[With obvious “refinements”:  $B$  has an interior point, and  $A$  is disjoint from the interior of  $B$ ; also:  $\exists f \neq 0$ , continuous,  $f(A) \geq 0$  iff there exists a non-empty open convex cone disjoint from  $A$  — cf. I.3Ex.12 p.37 for the necessity of the interior point.]

One can always obtain an “algebraic” statement from the above by using the strongest locally convex topology defined as follows:  $A$  is radial at  $x \in A$  iff  $\forall y \in E, \exists \varepsilon > 0: x + \varepsilon y \in A$ ; then  $U$  is open iff it is radial at each of its points. So:

**PROPOSITION 1.20.** *If  $p$  is a sublinear functional ( $p(\lambda x) = \lambda p(x)$  for  $\lambda \geq 0$ ,  $p(x) \in \mathbb{R}$ ,  $p(x+y) \leq p(x) + p(y)$ ),  $V$  a subspace,  $f: V \rightarrow \mathbb{R}$ ,  $f(v) \leq p(v)$  then  $\exists \bar{f}: E \rightarrow \mathbb{R}$ ,  $\bar{f} \leq p$ ,  $\bar{f}(v) = f(v)$  for  $v \in V$ . (Apply 1.19 in  $E \times \mathbb{R}$  to the subsets  $U = \{(x, \varphi) \mid \varphi > p(x), x \in E\}$  and  $A = \{(v, f(v)) \mid v \in V\}$ .)*

*Henceforth  $E$  will be locally convex, subsets closed, and linear functionals continuous and non-zero.*

**PROPOSITION 1.21.** *If  $0 \notin A$ ,  $\exists f: f(A) \geq 1$ .*

**PROOF.** Apply 1.19,  $U$  being a neighbourhood of zero. ■

**COROLLARY 1.22.** *Finite dimensional convex subsets with disjoint relative interiors [i.e., their interiors in the affine subspace they generate] can be separated.*

**PROOF.** It suffices to separate from zero the difference  $D$  of their relative interiors. E.g. by 1.20, we can assume  $E$  is spanned by  $D$ . If 0 is not in the closure of  $D$ , apply 1.21; otherwise  $D$  has non-empty interior and use 1.18]. ■

Otherwise, to apply 1.21 one needs conditions for the difference  $B - C$  of two closed convex sets to be closed. This is true if one is compact; more generally,  $A_B = \bigcap_{\varepsilon > 0} \varepsilon(B - b)$  does not depend on the choice of  $b \in B$ , and is called the **asymptotic cone** of  $B$ ; then:

**PROPOSITION 1.23.** *If  $A_B \cap A_C = \{0\}$ , and  $B$  or  $C$  is locally compact, then  $B - C$  is closed.*

**PROPOSITION 1.24.** *Assume  $B$  and  $C$  are cones, one of them locally compact, and  $B \cap C = \{0\}$ . Then  $\exists f: f(B) \geq 0, f(C) \leq 0$ .*

Taking for  $B$  the polar  $P^0$  ( $= \{ b \mid \langle b, \varphi \rangle \geq 0 \forall \varphi \in P \}$ ) of some closed convex cone  $P$  in the dual, one obtains the alternative: either  $\exists c \in C - \{0\}$ ,  $\langle c, \varphi \rangle \geq 0 \forall \varphi \in P$ , or  $\exists \varphi \in P - \{0\}$ ,  $\langle c, \varphi \rangle \leq 0 \forall c \in C$  — and this alternative holds thus as soon as one of the cones  $C, C^0, P, P^0$  is locally compact.

PROOF. Apply 1.23 to  $B$  and  $C - b$  for  $b \in B - \{0\}$ . ■

COROLLARY 1.25. If in 1.24 both  $B$  and  $C$  are locally compact, and contain no straight lines, then  $\exists f$ :  $f(B \setminus \{0\}) > 0 > f(C \setminus \{0\})$ .

PROOF. Indeed,  $D = B - C$  has then the same properties, which imply that the sets  $\{ d \in D \mid f(d) \leq 1 \}$ , for  $f$  in the dual, form a basis of neighbourhoods of zero in  $D$ . Choose then a  $f$  yielding a compact neighbourhood. ■

### Exercises.

The first series of exercises investigates the general properties of the value operator. For short, we let, for any sets  $S$  and  $T$ , and any function  $g(s, t)$  (with values in  $\mathbb{R} \cup \{+\infty\}$  or in  $\mathbb{R} \cup \{-\infty\}$ )

$$\bar{v}(g) = \inf_{\tau \in \mathcal{T}_f} \sup_{\sigma \in \Sigma_f} g(\sigma, \tau), \quad \underline{v}(g) = \sup_{\sigma \in \Sigma_f} \inf_{\tau \in \mathcal{T}_f} g(\sigma, \tau),$$

where  $g(\sigma, \tau) = \int g(s, t) d(\sigma \otimes \tau)$  and we write  $v(g)$  when they are equal. We start for the record with the obvious properties.

**1.**  $\bar{v} \geq \underline{v}$  and both are **positively homogeneous** of degree one (i.e.,  $\bar{v}(tg) = t\bar{v}(g)$  for  $t \geq 0$ ), monotone and invariant under translation by constant functions (i.e.,  $\bar{v}(g + \lambda) = \bar{v}(g) + \lambda$ ). The latter two properties imply they are **non-expansive**, i.e.  $|\bar{v}(f) - \bar{v}(g)| \leq \|f - g\|$  with  $\|f - g\| = \sup_{s, t} |f(s, t) - g(s, t)|$  (and with the convention  $|\infty - \infty| = 0$ ).

#### 2.

a. Under the assumptions of prop. 1.17 p. 7, the common valued asserted there equals  $\bar{v}(g)$ . We will then use the notation  $\Sigma(g)$  for  $\{ \sigma \in \Sigma \mid g(\sigma, t) \geq \bar{v}(g) \ \forall t \in T \}$  (and similarly for  $\mathcal{T}(g)$  under dual assumptions).  $\Sigma(g)$  is closed (hence compact), convex and non-empty.

b. If a decreasing net  $g_\alpha$  satisfies the assumptions of prop. 1.17, then

- (1)  $g = \lim g_\alpha$  does also
- (2)  $\bar{v}(g) = \lim \bar{v}(g_\alpha)$
- (3) If  $\sigma_\alpha \in \Sigma(g_\alpha)$  — or if only  $g_\alpha(\sigma_\alpha, t) \geq \bar{v}(g) - \varepsilon_\alpha \forall t \in T$ , with  $\varepsilon_\alpha \rightarrow 0$  — then for any limit point  $\sigma$  of  $\sigma_\alpha$  we have  $\sigma \in \Sigma(g)$ .
- (4) If  $\tau \in \mathcal{T}_f$  is  $\varepsilon$ -optimal for  $g$ , it is also so for all  $g_\alpha$  with  $\alpha$  sufficiently large.

HINT. 1 is obvious, and 2 and 3 follow from the fact that if  $f_\alpha$  is a decreasing net of real valued u.s.c. functions on the compact space  $S$  converging to  $f$ , and  $\sigma_\alpha$  s.t.  $f_\alpha(\sigma_\alpha) \geq \lambda - \varepsilon_\alpha$  with  $\varepsilon_\alpha \rightarrow 0$ , then  $f(\sigma) \geq \lambda$  for any limit point  $\sigma$ . This follows in turn from the upper semi-continuity of the  $f_\alpha$  on the compact space  $\Sigma$ , and from  $f_\alpha \rightarrow f$  on  $\Sigma$ .

**3. Continuity.** Let  $S$  and  $T$  be compact. Denote by  $C$  the convex cone of real valued functions on  $S \times T$ , which are u.s.c. in  $s$  and l.s.c. in  $t$ . If  $f_\alpha$  is a net in  $C$ , let

$$\begin{aligned} \varphi_{\alpha_0}(s_0, t_0) &= \limsup_{s \rightarrow s_0} \left[ \sup_{\alpha \geq \alpha_0} f_\alpha(s, t_0) \right], \\ \psi_{\alpha_0}(s_0, t_0) &= \liminf_{t \rightarrow t_0} \left[ \inf_{\alpha \geq \alpha_0} f_\alpha(s_0, t) \right]. \end{aligned}$$

Define the following concept of convergence on  $C$ :  $f_\alpha \rightarrow f$  iff  $\varphi_\alpha$  and  $\psi_\alpha$  converge point-wise to  $f$ . Assume  $f_\alpha \rightarrow f$ .

- a. Then  $\forall t, \exists \alpha_0 : \varphi_\alpha(s, t) < +\infty, \forall s \in S$  and  $\forall \alpha \geq \alpha_0$ , and similarly for  $\psi$ .

b. The convergence is compatible with the lattice structure and with the convex cone structure:

- (1)  $f_\alpha \rightarrow f$  and  $g_\alpha \rightarrow g$  imply  $f_\alpha \wedge g_\alpha \rightarrow f \wedge g$ ,  $f_\alpha \vee g_\alpha \rightarrow f \vee g$ , and  $f_\alpha + g_\alpha \rightarrow f + g$ .
- (2)  $\lambda_\alpha \geq 0$ ,  $\lambda_\alpha \rightarrow \lambda$ , and  $f_\alpha \rightarrow f$  imply  $\lambda_\alpha f_\alpha \rightarrow \lambda f$ .

HINT. For 2, it suffices to prove e.g. convergence from above, and hence to fix  $t_0 \in T$ . Using I.1Ex.3a, subtract then an appropriate constant to reduce (by (1) to the case where  $f_\alpha(s, t_0) \leq 0 \forall \alpha, \forall s$ ). Let  $\mu_{\alpha_0} = \inf_{\alpha \geq \alpha_0} \lambda_\alpha$ : then  $\mu_\alpha \varphi_\alpha(s, t_0)$  are u.s.c.,  $\geq \lambda_\alpha f_\alpha(s, t_0)$ , and decrease to  $\lambda f(s, t_0)$ .

c. Assume  $v(f)$  exists and  $f_\alpha \rightarrow f$ . Then

- (1)  $\bar{v}(f_\alpha)$  and  $\underline{v}(f_\alpha)$  converge to  $v(f)$ .
- (2) Any limit point of  $\varepsilon_\alpha$ -optimal strategies for  $f_\alpha$  ( $\varepsilon_\alpha \rightarrow 0$ ) belongs to  $\Sigma(f)$ .
- (3) Any  $\varepsilon$ -optimal strategy with finite support for  $f$  is so for all  $f_\alpha$  ( $\alpha$  sufficiently large).

HINT. Point 1 allows to use I.1Ex.2b p. 9.

d. If  $S$  and  $T$  are Hausdorff (or just “locally compact”, i.e., such that every point has a basis of compact neighbourhoods), the convergence concept is topological — i.e., derives from a “locally convex” Hausdorff topology on  $C$ .

HINT. Take as subbase of open sets the sets  $V_{t,\lambda,K} = \{g \in C \mid g(s, t) < \lambda \text{ for } s \in K\}$  and  $W_{s,\lambda,K'} = \{g \in C \mid g(s, t) > \lambda \text{ for } t \in K'\}$  with  $\lambda \in \mathbb{R}$ ,  $K$  and  $K'$  compact.

**4.** I.1Ex.3b shows that  $f_\alpha^i \rightarrow f^i$ ,  $\lambda_\alpha^i \rightarrow \lambda^i$ ,  $\lambda_\alpha^i \geq 0$  implies  $\sum_1^n \lambda_\alpha^i f_\alpha^i \rightarrow \sum_1^n \lambda^i f^i$ . Even under the best circumstances, the restriction to  $\lambda_\alpha^i \geq 0$  cannot be dispensed with:

a. Let  $S = \{s_0\}$ ,  $T = [0, 1]$ ,  $f_1(t) = \mathbb{1}_{t>0}$ ,  $f_2(t) = t^{-1}$  for  $t > 0$ ,  $f_2(0) = 0$ :  $f_1$  and  $f_2$  belong to  $C$  and are  $\geq 0$ , also  $-f_1 + \varepsilon f_2$  belongs to  $C$  for  $\varepsilon > 0$  but not for  $\varepsilon = 0$ : we have a monotone straight line whose intersection with  $C$  is not closed.

b. Consider  $g_\varepsilon = (1 - \varepsilon f_2)^2 = (1 - 2\varepsilon f_2 + \varepsilon^2 f_3)$ : this is a nice curve in a fixed plane, with this time  $g_\varepsilon \in C$  for all  $\varepsilon \geq 0$ . Yet  $v(g_\varepsilon) = 0$  for  $\varepsilon > 0$ ,  $v(g_0) = 1$ .

c. Nevertheless, prove that the intersection of  $C$  with a finite dimensional affine function space  $A$  is closed, and that on this intersection the convergence in  $C$  coincides with the usual in  $A$ , if any difference of two functions in  $A$  is separately bounded (i.e., bounded in each variable, the other being fixed).

**5.** To show that monotonicity and compactness in ex. I.1Ex.2 cannot be dispensed with:

a. Let  $S = \overline{\mathbb{N}}$ ,  $T = \{0\}$ ,  $g_k = \mathbb{1}_k$ : the  $g_k$  are a sequence of positive continuous functions converging weakly to zero ( $\sum g_k \leq 1$ ), yet  $v(g_k) = 1$  does not converge to zero.

b. On any infinite compact Hausdorff space  $S$ , the above example can be reduplicated (thus:  $g_k$  continuous,  $\|g_k\| = 1$ ,  $g_k \geq 0$ ,  $g_k(s) > 0 \Rightarrow g_n(s) = 0 \forall n \neq k$ ).

HINT. Show that one can find in  $S$  two disjoint compact sets with non-empty interior, one of which is infinite, and use this inductively.

c. Even with monotonicity, one cannot dispense with compactness: e.g.,  $S = \mathbb{N}$ ,  $g_k(n) = \mathbb{1}_{n \geq k}$ .

d. Cf. also ex. II.2Ex.1 p. 84.

**6. Differentiability.** We keep the notation of ex. I.1Ex.3. Denote by  $G$  the convex cone of functions  $g$  having a non-ambiguous, u.s.c.-l.s.c. extension  $g(\sigma, \tau)$  to  $\Sigma \times \mathcal{T}$  — i.e.,  $g(s, \tau) < +\infty$  is u.s.c. on  $S$  for each  $\tau$ ,  $g(\sigma, t) > -\infty$  is l.s.c. on  $T$  for each  $\sigma$ , and

$$\int g(s, \tau) \sigma(ds) = \int g(\sigma, t) \tau(dt) \text{ for each } (\sigma, \tau).$$

Let also  $V = \{f \in C \mid v(f) \text{ exists}\}.$

For  $f \in V$ ,  $g \in G$ ,  $g|_{\Sigma(f) \times \mathcal{T}(f)} \in V$  (e.g. by theorem 1.6); denote the value of this game by  $v_f(g)$ . We first deal with differentiability along straight lines, then along differentiable curves.

COMMENT 1.8. Theorem 1.6 p. 4 and prop. 2.6 p. 17 give sufficient conditions for  $f \in V$ . Ex. I.1Ex.7 p. 12 and ex. I.2Ex.1 p. 20 give sufficient conditions for  $g \in G$ .

a. For  $f + \varepsilon g \in V$ ,  $g \in G$ ,  $\lim_{\varepsilon \geq 0} [v(f + \varepsilon g) - v(f)]/\varepsilon = v_f(g)$ .

HINT. (cf. Mills, 1956). For  $\tau \in \mathcal{T}(f)$ ,  $\sigma_\varepsilon$  an  $\varepsilon^2$ -optimal strategy with finite support in  $f + \varepsilon g$ , one has  $v(f + \varepsilon g) \leq f(\sigma_\varepsilon, \tau) + \varepsilon g(\sigma_\varepsilon, \tau) + \varepsilon^2$  so  $(v(f + \varepsilon g) - v(f))/\varepsilon \leq g(\sigma_\varepsilon, \tau) + \varepsilon$ .  $g \in G$  implies then  $\limsup_{\varepsilon \geq 0} (v(f + \varepsilon g) - v(f))/\varepsilon \leq \max_{\tilde{\sigma}} g(\tilde{\sigma}, \tau)$  where  $\tilde{\sigma}$  ranges over the limit points of  $\sigma_\varepsilon$ . Ex. I.1Ex.3b and I.1Ex.3c p. 10 implies  $\tilde{\sigma} \in \Sigma(f)$ . Thus  $\limsup_{\varepsilon \geq 0} (v(f + \varepsilon g) - v(f))/\varepsilon \leq \min_{\tau \in \mathcal{T}(f)} \max_{\tilde{\sigma} \in \Sigma(f)} g(\tilde{\sigma}, \tau) = v_f(g)$ .

b. If  $f + \varepsilon g$  is real valued and satisfies the assumptions of theorem 1.6 for  $0 \leq \varepsilon \leq \varepsilon_0$  (e.g. is concave in  $s$  and convex in  $t$ ), one does not need  $g \in G$ : the above argument goes through in pure strategies, interpreting in the result  $\Sigma(f)$  and  $\mathcal{T}(f)$  as pure strategy sets.

HINT. Since  $\Sigma(f)$  and  $\mathcal{T}(f)$  are compact convex,  $(f + \varepsilon g)|_{\Sigma(f) \times \mathcal{T}(f)}$  also satisfies the assumptions of theorem 1.6. But  $f$  is constant on  $\Sigma(f) \times \mathcal{T}(f)$ , so  $g|_{\Sigma(f) \times \mathcal{T}(f)}$  satisfies them. Let now  $h = f + \varepsilon_0 g$ , and use the arguments of I.1Ex.6a for  $f + \varepsilon h$ , plus the above remark and homogeneity.

c. Assume that  $h_\varepsilon \in V$  and  $(h_\varepsilon - h_0)/\varepsilon \rightarrow g \in G$ , in the sense that, like in ex. I.1Ex.3 p. 9, there exists, for each  $\sigma$  and  $\tau$ ,  $\varphi_\varepsilon(s, \tau)$  u.s.c. in  $s$  and decreasing to  $g(s, \tau)$  and  $\psi_\varepsilon(\sigma, t)$  l.s.c. in  $t$  and increasing to  $g(\sigma, t)$  such that  $h_\varepsilon(s, \tau) \leq h_0(s, \tau) + \varepsilon \varphi_\varepsilon(s, \tau)$  and  $h_\varepsilon(\sigma, t) \geq h_0(\sigma, t) + \varepsilon \psi_\varepsilon(\sigma, t)$ . Then  $[v(h_\varepsilon) - v(h_0)]/\varepsilon \rightarrow v_{h_0}(g)$ .

HINT. Argue like in I.1Ex.6a.

d.

- One can use in I.1Ex.6c the homogeneity as was done sub I.1Ex.6b, using the conditions  $h_\varepsilon(s, \tau) \leq (1 - \varepsilon A)h_0(s, \tau) + \varepsilon \varphi_\varepsilon(s, \tau)$  and  $h_\varepsilon(\sigma, t) \geq (1 - \varepsilon A)h_0(\sigma, t) + \varepsilon \psi_\varepsilon(\sigma, t)$  ( $A$  arbitrary) instead of the above with  $A = 0$ , and obtaining then  $(v(h_\varepsilon) - v(h_0))/\varepsilon \rightarrow v_{h_0}(g) - Av(h_0)$  [since  $h_0(\sigma, \tau)$  is not necessarily defined, even on  $\Sigma_{h_0} \times \mathcal{T}_{h_0}$ , one does not obtain as sub I.1Ex.6b that this limit equals  $v_{h_0}(g - Ah_0)$ ].
- Similarly, I.1Ex.6b can be extended: if  $f_\varepsilon$  satisfies the assumptions of theorem 1.6 p. 4 for  $0 \leq \varepsilon < \varepsilon_0$ , and if for  $A$  sufficiently large,  $(f_\varepsilon - f_0)/\varepsilon + Af_0 \rightarrow g + Af_0$  (in the sense of ex. I.1Ex.3), then  $[v(f_\varepsilon) - v(f_0)]/\varepsilon \rightarrow v_{f_0}(g)$ .

e. In fact, closer inspection of the proof shows that much less is needed: assume  $f_\varepsilon(s, t)$  and  $g(s, t)$  are real valued, such that  $f_\varepsilon(0 \leq \varepsilon < \varepsilon_0)$  satisfies the assumptions of theorem 1.6, and such that (letting  $O$  denote an open set):

$$\begin{aligned} \forall t, \forall s_0, \forall \delta > 0, \exists \varepsilon_0 > 0, \exists O : s_0 \in O \subseteq S, \exists A > 0 : \forall s \in O, \forall \varepsilon < \varepsilon_0 \\ [f_\varepsilon(s, t) - f_0(s, t)]/\varepsilon < g(s_0, t) + A[\max_x f_0(x, t) - f_0(s, t)] + \delta \end{aligned}$$

and the dual condition. Then  $g|_{S(f_0) \times T(f_0)}$  satisfies the assumptions of theorem 1.6, and  $[v(f_\varepsilon) - v(f_0)]/\varepsilon \rightarrow v_{f_0}(g)$ .

HINT. Establish first the first statement, next that  $f_\varepsilon \rightarrow f_0$  (ex. I.1Ex.3), next that there exist  $\varphi_{\varepsilon,A}(s,t)$  which are u.s.c. in  $s$ , decreasing in  $\varepsilon$  ( $\varepsilon \rightarrow 0$ ) and  $A(A \rightarrow +\infty)$ , such that  $(f_\varepsilon(s,t) - f_0(s,t))/\varepsilon + Af_0(s,t) - Av(f_0) \leq \varphi_{\varepsilon,A}(s,t)$ , and such that for  $t \in T(f_0)$   $\lim_{\substack{\varepsilon \rightarrow 0 \\ A \rightarrow +\infty}} \varphi_{\varepsilon,A}(s,t) \leq g(s,t)$ . Finally show  $[v(f_\varepsilon) - v(f_0)]/\varepsilon \leq [f_\varepsilon(s_\varepsilon, t) - f_0(s_\varepsilon, t)]/\varepsilon + Af_0(s_\varepsilon, t) - Av(f_0)$ , for  $\varepsilon A \leq 1$ ,  $t \in T(f_0)$  and  $s_\varepsilon \in S(f_\varepsilon)$ , and argue as above.

COMMENT 1.9. Thus I.1Ex.6e is by far the “best” result, and I.1Ex.6c and I.1Ex.6d should be applied only in cases where I.1Ex.6e is not applicable, a.o. because the mixed extension of the game is not defined or lacks the proper u.s.c.-l.s.c. properties. The next exercise shows such cases are bound to be very rare.

**7.** Let  $S$  and  $T$  be compact metric spaces,  $f$  a bounded real valued function on  $S \times T$  such that  $f(\cdot, t)$  is u.s.c. for each  $t$  and  $f(s, \cdot)$  l.s.c. for each  $s$ .

a. Show that  $f$  is Borel measurable. [Hence  $f(\sigma, \tau)$  is unambiguous, also the assumptions of prop. 2.6 p. 17 are satisfied.]

HINT. Find an increasing sequence of functions converging point-wise to  $f$ , where for each function in the sequence there is some partition of  $T$  with the function being constant in  $t$  on each partition element.

b. Show that  $f(\sigma, \tau)$  is u.s.c. on  $\Sigma$  for each  $\tau \in \mathcal{T}$  and l.s.c. on  $\mathcal{T}$  for each  $\sigma \in \Sigma$ .

HINT. Use Fubini’s theorem and Fatou’s lemma.

c. If measurability is known, I.1Ex.7b also follows without the metrisability assumption, but assuming just strict semi-compactness (the closure of a subset is the set of limit points of convergent sequences in the subset). Most compact subsets of topological vector spaces have this property (Grothendieck, 1953).

**8.** Often, the result of ex. I.1Ex.6 strengthens itself in the following way. Assume e.g. we are working on a finite dimensional subspace of games, and we know (e.g. by ex. I.1Ex.1) that the function  $v$  is Lipschitz on this subspace. Then we have: if a Lipschitz function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is such that  $f'_x(y) = \lim_{\varepsilon \geq 0} (f(x + \varepsilon y) - f(x))/\varepsilon$  exists for all  $x$  and  $y$ , then  $f'_x(\cdot)$  is Lipschitz and is a “true” differential, i.e.  $F(y) = f(x + y) - f(x) - f'_x(y)$  is differentiable at zero (with zero differential:  $\lim_{\varepsilon \rightarrow 0} \sup_{0 < \|y\| \leq \varepsilon} |F(y)| / \|y\| = 0$ ).

### 9.

a. The Lipschitz condition in ex. I.1Ex.8 is necessary: on  $[0, 1]^2$ ,  $F(x, y) = x^3y/(x^4 + y^2)$  is analytic except at zero, is Lipschitz at zero:  $0 \leq F(x, y) \leq \frac{1}{2}\|(x, y)\|$ , and has all its directional derivatives zero (thus linear) at zero, yet  $F(t, t^2) = \frac{1}{2}t$ :  $F$  is not differentiable at zero.

b. u.s.c.-l.s.c. functions occur quite easily, e.g. on  $[0, 1]^2 : F(0, 0) = 0$ ,  $F(x, y) = (y - x)/(y + x)$ .

c. The boundedness condition in I.1Ex.7b is really needed: even with separate continuity (where the metrisability condition in ex. I.1Ex.7 become superfluous, cf. ex. I.2Ex.1 p. 20), define  $F$  on  $[0, 1]^2$  by  $F(0, 0) = 0$ ,  $F(x, y) = xy/(x^3 + y^3)$ . Show that the mixed extension (which always exists by I.1Ex.7a using just positivity), although being jointly lower semi-continuous ( $F$  being so), is not separately u.s.c.

Show also that, for real valued measurable functions, boundedness of the function is equivalent to finiteness of the mixed extension (i.e., (absolute) integrability for every product probability).

### 10. Compactness of $\Delta(X)$ for compact $X$ .

(Mertens, 1986a)  
Denote by  $C$  the convex cone of bounded l.s.c. functions on a compact space  $X$  and let  $E = C - C$ . Denote by  $P$  the set of monotone functions on  $C$  satisfying

- $p(tf) = tp(f)$  for  $t \geq 0$  and
- $p(f+g) \leq p(f) + p(g)$ .

$P$  is ordered in the usual way. Let  $M(X)$  be the set of minimal elements of  $P$ .

- a.  $\forall p \in P, \exists \mu \in M(X) : \mu \leq p$ .

HINT. Given  $\alpha \leq 0, f \geq \alpha$  and  $q \leq p$ , one has  $p(f) \geq q(f) \geq q(\alpha) \geq -q(-\alpha) \geq -p(-\alpha)$ , hence the set of possible values for  $q(f)$  is a bounded interval — use Zorn.

- b. Any  $\mu$  in  $M(X)$  can be identified with a positive linear functional on  $E$  satisfying:

$$\mu(f) = \inf\{\mu(g) \mid g \in C, g \geq f\}.$$

HINT.  $\tilde{\mu}$  defined by the above right hand member satisfies i) and ii) and coincides with  $\mu$  on  $C$ . Use Hahn-Banach to get a positive linear functional  $\zeta \leq \tilde{\mu}$ . This  $\zeta$  is unique — otherwise  $\mu$  would not be minimal on  $C$  — hence, using again Hahn-Banach, coincides with  $\tilde{\mu}$ .

- c.  $M(X)$  is the set of regular Borel measures on  $X$ .

HINT. Follows from a Daniell-type extension:

If  $f_n$  is an increasing sequence in  $C$ ,  $\mu(\lim f_n) = \lim \mu(f_n)$ , (use Dini, cf. (Meyer, 1966, X.6) and I.1Ex.10b). Define  $\mu^*$  on the set  $F$  of real bounded functions on  $X$  by  $\mu^*(f) = \inf\{\mu(g) \mid g \in C, g \geq f\}$ . Then:

$$\mu^*(f+g) \leq \mu^*(f \vee g) + \mu^*(f \wedge g) \leq \mu^*(f) + \mu^*(g)$$

and if  $f_n$  is an increasing sequence in  $F$ ,  $\mu^*(\lim f_n) = \lim \mu^*(f_n)$ .

Define  $L = \{f \in F \mid \mu^*(f) + \mu^*(-f) \leq 0\}$ .  $L$  is a vector space and  $\mu^*$  a linear functional on it. Given  $O$  and  $U$  open in  $X$ ,  $\mathbb{1}_O$  and  $\mathbb{1}_{O \cap U}$  are in  $L$ , hence also  $\mathbb{1}_{O \setminus U}$  with:  $\mu^*(O) = \mu^*(O \cap U) + \mu^*(O \setminus U)$ , hence for any subset  $A$ :  $\mu^*(A) \geq \mu^*(A \cap U) + \mu^*(A \setminus U)$ , so that any open  $U$  is  $\mu^*$ -measurable, hence all Borel sets also.

Finally since  $\mu^*(A) = \sup\{\mu(F) \mid F \subseteq A, F \text{ closed}\}$ , and  $X$  is compact,  $\mu^*$  defines a regular Borel measure (unique since equal to  $\mu$  on open sets).

- d. For all  $t > 0$ , the sets  $\{\mu \mid \mu \in \Delta(X), \mu(1) = t\}$  and  $\{\mu \mid \mu \in \Delta(X), \mu(1) \leq t\}$  are closed and compact (Recall 1.12).

HINT. It suffices to prove the compactness of the second set. Given an ultrafilter on it, let  $\varphi$  denote its point-wise limit in the set of positive linear functionals on  $E$ . By I.1Ex.10a there exists  $\nu \in \Delta(X)$  with  $\nu \leq \varphi$  on  $C$ :  $\nu$  is a limit point in  $\Delta(X)$ .

## 2. Complements to the minmax theorem

This section gives a number of more specialised “how to use” tricks and other complements for the minmax theorem of sect. 1 — and its “usual” form (prop. 2.6 p. 17).

**2.a. The topology on  $S$ .** Since there is no Hausdorff requirement, prop. 1.17 just asks that  $S$  be compact when endowed with the coarsest topology for which the functions  $f(\cdot, t)$  are u.s.c. This is equivalent (Mertens, 1986a, remark 1, p. 247) to ask that any point-wise limit of pure strategies — i.e. of functions  $f(s, \cdot)$  — be dominated (i.e., smaller on  $T$ ) by some pure strategy. Using the “countably compact” version, this can even be further weakened to: for any countable subset  $T_0$  of  $T$ , and any sequence  $s_i \in S$ , there exists  $s_0 \in S$  such that, for all  $t \in T_0$ ,  $\liminf_{i \rightarrow \infty} f(s_i, t) \leq f(s_0, t)$  (Mertens, 1986a, remark 3, p. 247).

**2.b. Lack of continuity: regularisation.** Here we consider the case where the u.s.c. or compactness condition is not met.

DEFINITION 2.1. Let  $(\Omega, \mathcal{A})$  denote a measurable space, and  $\Sigma$  a class of probability measures on  $(\Omega, \mathcal{A})$ . The “support function”  $\phi_{\mathcal{F}}(\sigma)$  of a class  $\mathcal{F}$  of extended real valued functions on  $\Omega$  is defined on  $\Sigma$  by

$$\phi_{\mathcal{F}}(\sigma) = \inf_{f \in \mathcal{F}} \int_{\star} f d\sigma$$

where  $\int_{\star}$  denotes the lower integral ( $\int_{\star} f d\sigma = \sup\{\int h d\sigma \mid h \leq f, h \text{ measurable and bounded from above}\}$ ).

DEFINITION 2.2. Let, for each measurable set  $B$ ,  $\mathcal{F}_B = \{f|_B \mid f \in \mathcal{F}\}$ . Denote by  $co(\mathcal{F})$  the convex hull of the set of bounded functions minorated by some element of  $\mathcal{F}$ , and by  $m(\mathcal{F})$  the “monotone (decreasing) class” generated by  $\mathcal{F}$ : the smallest class of functions containing  $\mathcal{F}$  and containing the limit of every decreasing sequence in  $m(\mathcal{F})$ . Let finally  $D(\mathcal{F})$  (or  $D_{\Sigma}(\mathcal{F})$ ) be such that:

$$D(\mathcal{F}) = \{f \text{ bounded from above} \mid \forall \sigma \in \Sigma, \exists B \in \mathcal{A}: \sigma(B) = 1, f|_B \in m(co(\mathcal{F}_B))\}$$

$D(\mathcal{F})$  stands for the decreasing class generated by  $\mathcal{F}$ .

LEMMA 2.3.  $D(\mathcal{F}) \subseteq \{g \text{ bounded above} \mid \int_{\star} g d\sigma \geq \phi_{\mathcal{F}}(\sigma) \quad \forall \sigma \in \Sigma\}$ , with equality if:

- (1) all functions in  $\mathcal{F}$  are  $\sigma$ -measurable and  $\sigma$ -a.e. bounded from above, for all  $\sigma \in \Sigma$ ,
- (2)  $\Sigma$  contains every probability measure which is absolutely continuous w.r.t. some  $\sigma \in \Sigma$ .

PROOF. We first show inclusion. It suffices to show that

$$g \in m(co(\mathcal{F})) \Rightarrow \int_{\star} g d\sigma \geq \phi_{\mathcal{F}}(\sigma)$$

(applying then this result on some set  $B$  with  $\sigma(B) = 1$ ). This follows from the standard properties of the lower integral (note that any decreasing sequence in  $m(co(\mathcal{F}))$  is uniformly bounded above).

We turn now to the other inclusion. It will be sufficient to prove that, for bounded  $g$  satisfying  $\int_{\star} g d\sigma \geq \phi_{\mathcal{F}}(\sigma)$  for all  $\sigma \in \Sigma$ , one has  $g \in D(\mathcal{F})$ , since, for any other  $g$ , one will then have  $g \vee (-n) \in D(\mathcal{F})$  for all  $n$ , hence,  $g \in D(\mathcal{F})$ ,  $D(\mathcal{F})$  being a monotone class. (In fact, assume  $g_n \in D(\mathcal{F})$  decreases to  $g$ ; let  $\sigma(B_n) = 1$ ,  $h_n \in m(co(\mathcal{F}_{B_n})) : h_n \leq g_n|_{B_n}$ . For  $B = \bigcap_n B_n$ , one has  $\sigma(B) = 1$  and  $m(co(\mathcal{F}_B)) \supseteq [m(co(\mathcal{F}_{B_n}))]|_B$ , so  $h_n|_B \in m(co(\mathcal{F}_B))$  and, for some constant  $K$ ,  $h_n|_B \leq g_n|_B \leq K$ , so  $h = \lim_{k \rightarrow \infty} \sup_{n \geq k} h_n|_B \in m(co(\mathcal{F}_B))$  and  $h \leq g|_B$ .)

Fix now  $\sigma$ , and let  $\tilde{g} \in L_{\infty}(\sigma)$  stand for the ( $\sigma$ -)essential supremum of the measurable functions smaller than  $g$ . Note that, for any non-negative measure  $\mu$  in  $L_1(\sigma)$ , i.e. bounded and absolutely continuous w.r.t.  $\sigma$ :  $\int \tilde{g} d\mu \geq \inf_{f \in \mathcal{F}} \int f d\mu$  since, by assumption,  $(\mu/\|\mu\|) \in \Sigma$ .

Now, if  $\mu \in L_1(\sigma)$  has a non-zero negative part  $\mu^-$ , consider the Hahn decomposition  $\mu = \mu^+ - \mu^-$ ,  $\mu^+ \geq 0$ ,  $\mu^- \geq 0$ ,  $\mu^+(B_-) = 0$ ,  $\mu^-(B_+) = 0$ ,  $B_+$  and  $B_-$  measurable,  $B_+ \cap B_- = \emptyset$ . Fix  $f_0 \in \mathcal{F}$  ( $f_0 \leq K$   $\mu$ -a.e.), and let  $f_n = f_0^+ + n\mathbb{1}_{B_-}$ ; we have  $f_n \in \mathcal{L}_{\infty}(\sigma)$ ,  $f_n \geq f_0 \in \mathcal{F}$ , and  $\mu(f_n) = \mu(f_0^+) + n\mu^-(B_-) \rightarrow -\infty$ . Therefore, if  $G = \{g \in \mathcal{L}_{\infty}(\sigma) \text{ and bounded below} \mid \exists f \in \mathcal{F}: f \leq g\}$ , we have, for all  $\mu \in L_1(\sigma)$ ,  $\int \tilde{g} d\mu \geq \inf_G \int g d\mu$ . Thus,  $\tilde{g}$  belongs, by the Hahn-Banach theorem (1.e), to

the  $\sigma(L_\infty, L_1)$ -closed convex hull of  $G$ . Denote by  $G^c$  the convex hull of  $G$ , and by  $G^{c,m}$  the “monotone class” spanned by  $G^c$ , i.e. the smallest class of functions containing  $G^c$  and such that  $f_n \in G^{c,m}$ ,  $f_n$  decreasing to  $f$ ,  $f$  bounded below imply  $f \in G^{c,m}$  (thus  $G^{c,m}$  consists of those functions of  $m(G^c)$  which are bounded below). Note that  $G^{c,m}$  is convex, and  $G^{c,m} + \mathcal{L}_\infty^+(\sigma) \subseteq G^{c,m}$ .

Denote by  $H$  the image of  $G^{c,m}$  in  $L_\infty(\sigma)$ , i.e. the set of equivalence classes of  $G^{c,m}$ .  $H$  is still convex,  $H + L_\infty^+(\sigma) \subseteq H$ , and  $f_n \in H$ ,  $\sup_n \|f_n\|_{L_\infty} < +\infty$  implies  $\limsup_{n \rightarrow \infty} f_n \in H$ . Indeed, if  $g_n \in G^{c,m}$  belongs to the equivalence class of  $f_n$ , if  $M = \sup_n \|f_n\|_{L_\infty}$ , if  $B_n = \{s \mid g_n(s) > M\}$ , and if  $h_n = (\sup_{k \geq n} g_k) \vee (-M)$ , then  $h_n$  decreases everywhere to some element  $h$  of the equivalence class  $\limsup_{n \rightarrow \infty} f_n$ ,  $h$  is bounded below and  $\{h_1 > M\} \subseteq B = \bigcup_n B_n$ . Since  $\sigma(B) = 0$  and since  $h_n \geq g_n \in G^{c,m}$ , we have indeed  $h_n \in G^{c,m}$  and thus  $h \in G^{c,m}$ , so  $\limsup f_n \in H$ .

We have seen that  $\tilde{g}$  belongs to the weak $^\star$ , i.e.  $\sigma(L_\infty, L_1)$ -closure of  $H$ . We will now show that  $H$  is weak $^\star$ -closed, so it will follow that  $\tilde{g}$  is the equivalence class of some element  $\bar{g}$  of  $G^{c,m}$ : there exists  $\bar{g} \in G^{c,m}$ ,  $\bar{g} \leq g$  except on a  $\sigma$ -negligible set.

We want thus to show that a convex subset  $H$  of  $L_\infty$  is weak $^\star$ -closed if ( $f_n \in H$ ,  $\sup_n \|f_n\|_{L_\infty} < +\infty$  implies  $\limsup f_n \in H$ ). By the Krein-Smulian theorem on weak $^\star$ -closed convex sets (Kelley et al., 1963, p. 212), since  $H$  is convex in the dual  $L_\infty$  of the Banach space  $L_1$ , it is sufficient to show that the intersection of  $H$  with any ball is weak $^\star$ -closed: we can assume that  $\sup_{f \in H} \|f\|_\infty = R < +\infty$ . Further, it is sufficient to prove that  $H$  is  $\tau(L_\infty, L_1)$ -closed (Kelley et al., 1963, p. 154, Th. 17.1). But the Mackey-topology  $\tau(L_\infty, L_1)$  is finer than any  $L_p$  topology, hence a fortiori than the topology of convergence in measure. (In fact, they coincide on bounded subsets of  $L_\infty$ , cf. ex. I.2Ex.12 p. 24). Since this topology is metrisable, and since, from any sequence that converges in probability, one can extract an a.e. convergent subsequence, it is sufficient to show that if a sequence  $f_n$  in  $H$  converges a.e. to  $f$  then  $f$  ( $= \limsup f_n$  a.e.) belongs to  $H$ , which is our basic property of  $H$ . Thus, for some  $\bar{g} \in G^{c,m}$ , and some measurable set  $B_0$  with  $\sigma(B_0) = 1$ , we have  $\bar{g} \leq g$  everywhere on  $B_0$ .

Note now that, given a set  $G$  of functions, the union of the monotone classes generated by all countable subsets of  $G$  is a monotone class (because a countable union of countable subsets is still countable), and hence is the monotone class spanned by  $G$ . Thus  $\bar{g}$  belongs to the monotone class spanned by a sequence  $g_n \in G^c$ , and each  $g_n$  is a convex combination of finitely many  $g_{n,i} \in G$ . Since  $g_{n,i} \in \mathcal{L}_\infty(\sigma)$ , there exists  $B_{n,i} \in \mathcal{A}$  with  $\sigma(B_{n,i}) = 1$  such that  $g_{n,i|B_{n,i}}$  is bounded and  $\mathcal{A}$ -measurable. Choose also  $f_{n,i} \in \mathcal{F}$ ,  $f_{n,i} \leq g_{n,i}$ : for  $B = B_0 \cap (\bigcap_{n,i} B_{n,i})$ , we have  $B \in \mathcal{A}$ ,  $\sigma(B) = 1$ ,  $g_{n,i|B}$  is bounded, measurable and  $\geq f_{n,i|B}$ : thus  $g_{n|B} \in co(\mathcal{F}_B)$  and  $\bar{g}|_B \in m(co(\mathcal{F}_B))$  with  $\bar{g}|_B \leq g|_B$ . Since such a construction is possible for each  $\sigma \in \Sigma$ , we have indeed  $g \in D_\Sigma(\mathcal{F})$ . ■

Given a function  $g$  on  $S \times T$ , let, for all  $\tau \in \mathcal{T}_f$ ,  $\phi_\tau(s) = \limsup_{s' \rightarrow s} g(s', \tau)$ , and let  $D(g) = D\{\phi_\tau \mid \tau \in \mathcal{T}_f\}$ . We define similarly  $I(g)$  ( $I$  for increasing), reversing the rôles of  $S$  and  $T$  and the order on the reals.

**PROPOSITION 2.4.** *Let  $S$  be a compact topological space,  $T$  any set. Assume, for all  $t \in T$ ,  $g(\cdot, t) \in D(g)$ . Then,*

$$\max_{\sigma \in \Sigma} \inf_{\tau \in \mathcal{T}_f} \int g(s, t) d(\sigma \otimes \tau) = \inf_{\tau \in \mathcal{T}_f} \sup_{\sigma \in \Sigma} \int g(s, t) d(\sigma \otimes \tau).$$

PROOF. Prop. 1.17 can be applied to the game  $\phi_\tau(s)$ ; one obtains pure strategies  $\tau$  because  $\phi_\tau(s)$  is convex in  $\tau$ . The lemma yields then that the same optimal strategy  $\sigma$  guarantees the same amount against  $g$ , hence the result. ■

COMMENT 2.1. The convexity in  $\tau$  of  $\phi_\tau(s)$  implies that the convexification in the definition of  $D$  will be superfluous: one would equivalently obtain  $D$  if, instead of using  $m(\text{co}(\mathcal{F}))$ , one just stabilised the set of functions  $\{\phi_\tau \mid \tau \in \mathcal{T}_f\}$  under the  $\limsup$  of sequences which are uniformly bounded from above.

COMMENT 2.2. When the pay-off function  $g$  is uniformly bounded from above, as in many applications, one does not even have to stabilise: one could equivalently just define  $G = \{\limsup_{n \rightarrow \infty} \phi_{\tau_n} \mid \tau_n \in \mathcal{T}_f\}$ , and  $D = \{h \mid \forall \sigma \in \Sigma, \exists f \in G: f \leq h \text{ } \sigma\text{-a.e.}\}$ . Indeed, take  $f \in m(\text{co}(\mathcal{F}_B))$ ,  $f \leq h|_B$ ,  $\sigma(B) = 1$ . Then, clearly  $f$  belongs to the closure of  $\text{co}(\mathcal{F}_B)$  for the topology of convergence in measure, i.e., since  $\sigma(B) = 1$ , the equivalence class of  $f$  belongs to the closure of  $\text{co}(\mathcal{F})$ . Thus,  $f$  is the limit in measure of a sequence in  $\text{co}(\mathcal{F})$ , hence the limit of a  $\sigma$ -a.e. convergent sequence  $f_n \in \text{co}(\mathcal{F})$ . Each  $f_n$  is minorated by some  $\phi_{\tau_n}$ , hence  $\limsup_{n \rightarrow \infty} \phi_{\tau_n} \leq h$   $\sigma$ -a.e.. Since the pay-off function is uniformly bounded from above, the sequence  $\phi_{\tau_n}$  is also.

COMMENT 2.3. Further, in most actual applications (cf. exercises), the  $\sigma$ -a.e. aspect in the definition of  $D$  is not needed. Thus this is the form in which the criterion is most often used: show that, for each  $t \in T$ , there exists a sequence  $\tau_n \in \mathcal{T}_f$  such that  $\limsup_{n \rightarrow \infty} \phi_{\tau_n}(\cdot) \leq g(\cdot, t)$  (and such that  $\phi_{\tau_n}(\cdot)$  is uniformly bounded from above if  $g$  is not).

COMMENT 2.4. Only the “obvious part” (monotone convergence theorem) of the lemma was needed. The hard part shows that the above simple use of the monotone convergence theorem (or of Fatou’s lemma) is as powerful as the more sophisticated closure methods, as used for instance in (Karlin, 1950). Indeed, any such closure method will only yield functions satisfying, for all  $\sigma$ ,  $\int_* f d\sigma \geq \inf_\tau \int \phi_\tau d\sigma$ .

**2.c. Lack of compactness: approximation.** When also  $S$  is not necessarily compact, the previous ideas can be combined with an old idea going back to Wald (1950): that compactness is not really necessary, only an appropriate form of precompactness (but in the uniform topology) — and in fact a one-sided form, as was later observed.

Since, however, our typical assumptions are much weaker than a joint continuity of the pay-off function, the typical compactness that we have is not in the uniform topology, not even a one-sided form, so we retain from this precompactness only the one-sided uniform approximation by another game with compact (and not necessarily finite) strategy space: we will let the function  $\phi$  vary with  $\varepsilon$ , and use compact subsets  $\Sigma_\varepsilon$  of  $\Sigma$ .

PROPOSITION 2.5. Assume that, for all  $\varepsilon > 0$ , there exists a compact convex subset  $\Sigma_\varepsilon$  of  $\Sigma$ , and a function  $\phi_\varepsilon: S \times \mathcal{T}_f \rightarrow \mathbb{R} \cup \{-\infty\}$  such that

- (1)  $\phi_\varepsilon(s, \tau)$  is u.s.c. in  $s$  and convex in  $\tau$
- (2) for all  $t \in T$ ,  $g(\cdot, t) \in D_{\Sigma_\varepsilon}\{\phi_\varepsilon(\cdot, \tau) \mid \tau \in \mathcal{T}_f\}$
- (3) for all  $s \in S$ , there exists  $\sigma_s \in \Sigma_\varepsilon$  such that  $g(s, \tau) \leq \int \phi_\varepsilon(\cdot, \tau) d\sigma_s + \varepsilon$ , for all  $\tau$ .

Then

$$\sup_{\Sigma} \inf_{\mathcal{T}_f} \int g(s, t) d(\sigma \otimes \tau) = \inf_{\mathcal{T}_f} \sup_{\Sigma} \int g(s, t) d(\sigma \otimes \tau)$$

PROOF. Apply prop. 1.8 p. 5 to  $\phi_\varepsilon$  on  $\Sigma_\varepsilon \times \mathcal{T}_f$ ; let  $v_\varepsilon$ ,  $\sigma_\varepsilon$  and  $\tau_\varepsilon$  be the corresponding value and optimal strategies. By 2 and lemma 2.3,  $\sigma_\varepsilon$  still guarantees  $v_\varepsilon$  against  $\mathcal{T}_f$  in the

game  $g$ . By 3,  $\tau_\varepsilon$  guarantees  $v_\varepsilon + \varepsilon$  against  $S$  in  $g$ . This being true for all  $\varepsilon$ , the value  $v$  exists and  $\sigma_\varepsilon$  and  $\tau_\varepsilon$  are  $\varepsilon$ -optimal strategies. ■

COMMENT 2.5. Typically, one thinks of  $\Sigma_\varepsilon$  as the set of probabilities on a compact subset  $S_\varepsilon$  of  $S$ .

COMMENT 2.6. The previous proposition was the particular case where  $\sigma_s$  was the unit mass at  $s$  and  $\phi_\varepsilon$  was independent of  $\varepsilon$ .

COMMENT 2.7. Point 3 and the compactness of  $\Sigma_\varepsilon$ , together with the upper semi-continuity of  $\phi_\varepsilon$ , imply that any limit of functions  $g(s, \cdot)$  on  $\mathcal{T}_f$  is  $\varepsilon$ -dominated by some function  $\phi_\varepsilon(\sigma, \cdot)$ .

COMMENT 2.8. This last condition (together with 2) is in principle sufficient, barring some measurability problems: use first prop. 1.17 p. 7 to solve the game where player I's strategy set is the set of all limits of functions  $g(s, \cdot)$ . Let  $\sigma_0$  be his optimal strategy in this game and  $v$  the value. If one could select in a measurable way for each limit function a  $\sigma$  such that  $\phi_\varepsilon(\sigma, \cdot)$   $\varepsilon$ -dominates the function, use this selection to map  $\sigma_0$  to some  $\tilde{\sigma}_0 \in \Sigma$ , which guarantees in the game  $\phi_\varepsilon$  at least  $v - \varepsilon$ . By 2,  $\tilde{\sigma}_0$  will also guarantee  $v - \varepsilon$  in  $g$ , hence be an  $\varepsilon$ -optimal strategy in  $g$ .

Even without such a measurable selection, one might e.g. attempt to define a mixed strategy as some auxiliary probability space (here  $\sigma_0$ ), together with a map from there to the strategy space (here an arbitrary, non-measurable selection), such that player I guarantees himself  $v - \varepsilon$  in the sense of lower integrals.

**2.d. Measurability: symmetric case.** The right-hand member (and thus the left-hand member) in the equality of the above propositions is not increased if  $\Sigma$  is replaced by the space  $\tilde{\Sigma}$  of all order preserving linear functionals of norm 1 on the cone of functions on  $S$  generated by the functions  $g(\cdot, t)$  and the constants. That is to say that quantity is an unambiguous upper bound for any evaluation of the game (because symmetrically player II's strategies are of most restricted type (finite support)). Denote by  $\mathcal{T}_B$  the set of all probability measures on some  $\sigma$ -field  $B$  on  $T$ , and let  $F(\sigma, \tau) = \int_T d\tau \int_S g(s, t) d\sigma$ , where the integral on  $T$  is in the sense of a lower integral. Then

$$\max_{\sigma \in \Sigma} \inf_{\tau \in \mathcal{T}_B} F(\sigma, \tau) = \inf_{\tau \in \mathcal{T}_f} \sup_{\sigma \in \tilde{\Sigma}} F(\sigma, \tau).$$

Thus the possible discrepancy between upper and lower bound apparently depends more on the order of integration than on the allowed strategy spaces.

Although those propositions are a basic tool in proving that a game has a value, the above shows well why they do not assert per se that the game in question has a value: the value might in general depend on the order of integration, or, in other terms, on the greater or lesser generality of mixed strategies allowed for each player. The next propositions show some cases where this ambiguity can be relieved completely. (A less “complete” (cf. ex. I.2Ex.11 p. 23) way of relieving it would be to add to the previous assumptions some measurability requirement and use Fubini's theorem to obtain the minmax theorem directly on the mixed extension). Obviously one would by far prefer to be able to dispense with the hypotheses in the next theorem — cf. (Mertens, 1986a) for the importance of this question, and why this would yield a completely “intrinsic” theorem.

**THEOREM 2.6.** *Assume  $S$  and  $T$  are compact, and  $g$  is real valued and bounded from below or from above. Assume further that  $g(s, \cdot)$  is lower semi-continuous on  $T$  for each*

$s \in S$ , and  $g(\cdot, t)$  is upper semi-continuous on  $S$  for each  $t \in T$ . Then under any one of the following three hypotheses:

- (1)  $g$  is  $\mu \otimes \nu$  measurable for any regular product probability on the Borel sets of  $S \times T$ ,
- (2) one of the two spaces has a countable basis
- (3) one of the two spaces is Hausdorff

one has:

$$\sup_{\sigma \in \Sigma_f} \inf_{t \in T} \int g(s, t) d\sigma = \inf_{\tau \in \mathcal{T}_f} \sup_{s \in S} \int g(s, t) d\tau.$$

Further each player obviously has an optimal strategy in the form of a regular probability.

PROOF. In case 1, prop. 1.17 p. 7 applied both ways, yields the existence of an upper value  $\bar{v}$ , that player I can guarantee with a regular probability  $\mu$ , and player II with probabilities with finite support, and of a lower value  $\underline{v}$ , that player II can guarantee with a regular probability  $\nu$  and player I with probabilities with finite support.  $\int_S g(s, t) \mu(ds) \geq \bar{v}$  implies  $\int_T \nu(dt) \int_S g(s, t) \mu(ds) \geq \bar{v}$  and, similarly,  $\int_T g(s, t) \nu(dt) \leq \underline{v}$  implies  $\int_S \mu(ds) \int_T g(s, t) \nu(dt) \leq \underline{v}$ . By the measurability and boundedness assumptions of  $g$ , we can apply Fubini's theorem:  $\bar{v} \leq \int_T d\nu \int_S g d\mu = \int_S d\mu \int_T g d\nu \leq \underline{v}$ . But, by their very definition  $\underline{v} \leq \bar{v}$  (they are  $\sup \inf$  and  $\inf \sup$  of the game played in strategies with finite support): the proof is complete in this case.

Suppose now that 2 holds; we shall prove that in this case  $g$  is Borel so 1 applies. Assume that  $T$  has a countable basis  $O_n$ , and define:

$$f_n(s, t) = \begin{cases} -\infty, & \text{for } t \notin O_n \\ \inf_{t' \in O_n} f(s, t'), & \text{otherwise.} \end{cases}$$

Then  $f_n$  is Borel (since u.s.c. in  $s$ ) and  $g_n = \max_{k \leq n} f_n$  is an increasing sequence converging to  $f$  (since l.s.c. in  $t$ ).

It remains to consider case 3. Hence assume  $T$  Hausdorff. We will first construct a countable set of best replies. Let  $\mathcal{T}_n = \{\tau \in \mathcal{T}_f \mid \#\text{Supp}(\tau) \leq n\}$ . Denote by  $\Phi_0$  the set of continuous functions  $\varphi$  on  $T$  such that there exists  $s \in S$  with  $f(s, \cdot) \geq \varphi(\cdot)$  on  $T$ . Since  $\bar{v} \leq \inf_{\tau \in \mathcal{T}_n} \sup_{\varphi \in \Phi_0} \int \varphi d\tau$ , the following sets:  $O_{\varphi, k, n} = \{\tau \in \mathcal{T}_n \mid \int \varphi d\tau > \bar{v} - 1/k\}$ , form for  $\varphi \in \Phi_0$  and each fixed  $n$  and  $k$ , an open covering of the compact space  $\mathcal{T}_n$ . Denote by  $\Phi_{k, n}$  the indices of a finite subcovering. Then  $\Phi = \bigcup_{k, n} \Phi_{k, n}$  is a countable subset of  $\Phi_0$ , such that  $\bar{v} = \inf_{\tau \in \mathcal{T}_f} \sup_{\varphi \in \Phi} \int \varphi d\tau$ . We now reduce the situation to one where 2 applies. If  $\varphi_i$  enumerates  $\Phi$ , let us consider on  $T \times T$ ,  $d(t_1, t_2) = \sum_i 2^{-i} |\varphi_i(t_1) - \varphi_i(t_2)| / \|\varphi_i\|$ .  $d$  defines a metrisable quotient space  $\tilde{T}$  of  $T$ , such that, if  $\psi$  denotes the quotient mapping, any  $\varphi \in \Phi$  can be written as  $g \circ \psi$ , for some  $g \in \Psi$ , where  $\Psi$  denote the set of all  $g$  continuous on  $\tilde{T}$ , such that for some  $s \in S$ ,  $g \circ \psi(\cdot) \leq f(s, \cdot)$  on  $T$ . Define  $\tilde{f}$  on  $S \times \tilde{T}$  by  $\tilde{f}(s, \tilde{t}) = \sup\{g(\tilde{t}) \mid g \in C(\tilde{T}), g \circ \psi(\cdot) \leq f(s, \cdot)\}$ . Then we have:

$$\bar{v} \leq \inf_{\tau \in \mathcal{T}_f} \sup_{\varphi \in \Phi} \int \varphi d\tau \leq \inf_{\tau \in \mathcal{T}_f} \sup_{g \in \Psi} \int g \circ \psi(t) d\tau(t) \leq \inf_{\tilde{\tau} \in \widetilde{\mathcal{T}}_f} \sup_{s \in S} \int \tilde{f}(s, \tilde{t}) d\tilde{\tau}(\tilde{t}).$$

Obviously  $\tilde{f}$  is l.s.c. on  $\tilde{T}$  for each  $s \in S$ , and is the largest such function satisfying  $\tilde{f}(s, \psi(t)) \leq f(s, t)$ . Let  $h(s, \tilde{t}) = \inf\{f(s, t) \mid t \in \psi^{-1}(\tilde{t})\}$ : to prove that  $\tilde{f}$  is u.s.c. on  $S$ , we will show that  $\tilde{f} = h$ . This in turn follows from  $h(s, \cdot)$  being l.s.c. on  $\tilde{T}$  for each  $s \in S$ . In fact  $\tilde{T}$  being metrisable, let  $\tilde{t}_i$  be a sequence converging to  $\tilde{t}$ . Choose  $t_i$  such that

$\psi(t_i) = \tilde{t}_i$  and  $f(s, t_i) \leq h(s, \tilde{t}_i) + 1/i$ , and let  $t$  be a limit point of the sequence  $t_i$ : we have  $\psi(t) = \tilde{t}$  and  $h(s, \tilde{t}) \leq f(s, t) \leq \liminf f(s, t_i) \leq \liminf h(s, \tilde{t}_i)$  hence the required property. We then use the result under case 2 for  $f$  on  $S \times \tilde{T}$ . It follows that  $\bar{v} \leq \sup_{\sigma \in \Sigma_f} \inf_{\tilde{t} \in \tilde{T}} \int \tilde{f}(s, \tilde{t}) d\sigma(s) \leq \sup_{\sigma \in \Sigma_f} \inf_{t \in T} \int f(s, t) d\sigma(s)$ , since  $\tilde{f}(s, \psi(t)) \leq f(s, t)$ . This completes the proof of the proposition. ■

REMARK 2.9. When  $g$  is continuous in each variable no further assumptions are needed: either by using ex.I.2Ex.1 p. 20 to show that condition 1 is satisfied or reducing to 3, by using e.g. on  $S$  the coarsest topology for which all functions  $g(\cdot, t)$  are continuous (and going to the quotient space) (cf. ex.I.2Ex.11 p. 23 for an example showing that, even with such assumptions, compactness on both sides is really needed).

We obtain now the analogues to propositions 2.4 p. 15 and 2.5 p. 16.

PROPOSITION 2.7. Let  $S$  and  $T$  be compact. Assume  $g(s, \cdot)$  and  $g(\cdot, t)$  are bounded from below and from above resp. for all  $(s, t) \in S \times T$ . Assume there also exists  $f: S \times T \rightarrow \overline{\mathbb{R}}$ , measurable for any regular product probability and bounded either from below or from above, and such that,

$$\begin{aligned} f(s, \cdot) &\in I(g) & \forall s \in S \\ f(\cdot, t) &\in D(g) & \forall t \in T \end{aligned}$$

Then  $g$  has a value, and both players have  $\varepsilon$ -optimal strategies with finite support.

PROOF. Is the same as the proof of theorem 2.6, but using prop. 2.4 instead of prop. 1.17 ■

COMMENT 2.10. One usually takes  $f$  to be (some regularisation of)  $g$ , cf. exercises.

With the same proof as above one obtains:

PROPOSITION 2.8. Assume, for all  $\varepsilon > 0$ , there exist compact, convex subsets  $\Sigma_\varepsilon$  and  $\mathcal{T}_\varepsilon$  of  $\Sigma$  and  $\mathcal{T}$  and functions  $\phi_\varepsilon: S \times \mathcal{T}_f \rightarrow \mathbb{R} \cup \{-\infty\}$ ,  $\psi_\varepsilon: \Sigma_f \times T \rightarrow \mathbb{R} \cup \{+\infty\}$ ,  $f_\varepsilon: S \times T \rightarrow \overline{\mathbb{R}}$  such that

- (1)  $\phi_\varepsilon$  is u.s.c. in  $s$  and convex in  $\tau$ ;  $\psi_\varepsilon$  is l.s.c. in  $t$  and concave in  $\sigma$ ;
- (2)  $f_\varepsilon$  is measurable for any regular product measure, and bounded from below or from above;
- (3) for all  $s \in S$ , there exists  $\sigma_s \in \Sigma_\varepsilon$  such that, for all  $\tau \in \mathcal{T}_f$ ,

$$g(s, \tau) \leq \phi_\varepsilon(\sigma_s, \tau) + \varepsilon$$

and, for all  $t \in T$ , there exists  $\tau_t \in \mathcal{T}_\varepsilon$  such that, for all  $\sigma \in \Sigma_f$ ,

$$g(\sigma, t) \geq \psi_\varepsilon(\sigma, \tau_t) - \varepsilon;$$

- (4)  $f_\varepsilon(\cdot, t) \in D_{\Sigma_\varepsilon}(\phi_\varepsilon)$  for all  $t \in T$  and  $f_\varepsilon(s, \cdot) \in I_{\mathcal{T}_\varepsilon}(\psi_\varepsilon)$  for all  $s \in S$ .

Then  $g$  has a value, and both players have  $\varepsilon$ -optimal strategies with finite support.

COMMENT 2.11. Theorem 2.6 and propositions 2.7 and 2.8 imply that each player has  $\varepsilon$ -optimal strategies that are safe against any type of mixed strategy of the other player (even finitely additive ones ...) no matter in what order the integrations are performed. Further, those mixed strategies are really “playable” in the sense that one can obviously realise the mixing with a finite number of coin tosses. That is why we say the game has a value.

### 2.e. Pure optimal strategies.

**DEFINITION 2.9.** Call a function  $f$  defined on  $S \times T$  **concave-like** (resp. **convex-like**) if, for any  $\alpha$  ( $0 < \alpha < 1$ ) and for any  $s_1$  and  $s_2$  (resp.  $t_1$  and  $t_2$ ), there exists  $s_0$  (resp.  $t_0$ ) such that, for all  $t$ ,  $f(s_0, t) \geq \alpha f(s_1, t) + (1 - \alpha)f(s_2, t)$  (resp., for all  $s$ ,  $f(s, t_0) \leq \alpha f(s, t_1) + (1 - \alpha)f(s, t_2)$ ).

**PROPOSITION 2.10.** Assume, in addition to the hypotheses of prop. 1.17, that  $g$  is concave-like (resp. convex-like). Then, any strategy in  $\Sigma$  (resp.  $\mathcal{T}_f$ ) is dominated by a pure strategy. In particular the  $(\varepsilon)$ -optimal strategy  $\sigma$  (resp.  $\tau$ ) may be taken as a point mass. In particular, if  $g$  is concave-like, there is an unambiguous value.

**PROOF.** It is sufficient to prove the first statement. Induction on the number of pure strategies used in a mixed strategy with finite support shows immediately that any mixed strategy with finite support is dominated by a pure strategy. This proves the proposition in case  $g$  is convex-like. If  $g$  is concave-like, consider a regular probability  $\mu$  on  $S$ , and a finite subset  $T_0$  of  $T$ : for any  $\varepsilon > 0$ , there exists a probability  $\mu_{\varepsilon, T_0}$  with finite support on  $S$ , such that, for all  $t \in T_0$ ,  $\int_S g(s, t) d\mu_{\varepsilon, T_0} \geq \int_S g(s, t) d\mu - \varepsilon$  — this follows for instance from the strong law of large numbers. But we know that  $\mu_{\varepsilon, T_0}$  is dominated by a point mass, say at  $s_{\varepsilon, T_0}$ : for all  $t \in T_0$ ,  $g(s_{\varepsilon, T_0}, t) \geq \int_S g(s, t) d\mu - \varepsilon$ . Let, for all  $t \in T$ :  $S_{\varepsilon, t} = \{s \in S \mid g(s, t) \geq \int_S g(s, t) d\mu - \varepsilon\}$ ; thus the  $S_{\varepsilon, t}$  are compact subsets of  $S$ , and we have just shown that any finite intersection of them is non-empty: therefore they have a non-empty intersection. Let  $s_\mu \in \bigcap_{\varepsilon, t} S_{\varepsilon, t}$ ; we have, for all  $t \in T$ :  $g(s_\mu, t) \geq \int_S g(s, t) d\mu$ ; the strategy  $\mu$  is thus dominated by the pure strategy  $s_\mu$ . ■

**COMMENT 2.12.** When  $g$  is both concave-like and convex-like, the above result is often referred to as Fan's theorem (1953). One could have deduced prop. 1.17 from it.

**COMMENT 2.13.** Similarly, if  $g$  satisfies only the hypotheses of prop. 2.4, in addition to being concave-like (resp. convex-like), then  $\sigma$  (resp.  $\tau$ ) may still be taken as a point mass. Indeed, the above argument, taking  $\mu$  as the optimal strategy  $\sigma$ , still yields that the sets  $S_{\varepsilon, t} = \{s \in S \mid f(s, t) \geq \bar{v} - \varepsilon\}$  have non-empty finite intersections ( $\bar{v}$  being the (upper) value given in prop. 2.4). Therefore, the compact sets  $\bar{S}_{\varepsilon, t} = \{s \in S \mid \phi_t(s) \geq \bar{v} - \varepsilon\}$  have a non-empty intersection: let  $s_0 \in \bigcap_{\varepsilon, t} \bar{S}_{\varepsilon, t}$ ; then, for all  $t \in T$ ,  $\phi_t(s_0) \geq \bar{v}$ . Therefore, for all  $g \in D(\phi)$ ,  $g(s_0) \geq \bar{v}$  and  $s_0$  is an optimal pure strategy.

### Exercises.

#### 1.

- Let  $S$  and  $T$  be two topological spaces,  $g$  a real valued function defined on  $S \times T$ , continuous in each variable separately. Then  $g$  is  $\mu$ -measurable for any regular probability measure  $\mu$  on the Borel sets of  $S \times T$ .

**HINT.** It is sufficient to consider the case of bounded  $g$ , and (regularity) of compact  $S$  and  $T$ , and further that the topology on  $S$  (resp.  $T$ ) is the coarsest for which every function  $g(\cdot, t)$  (resp.  $g(s, \cdot)$ ) is continuous.

Let  $F$  be any subset of  $T$ , and  $t_0$  in the closure of  $F$ : then there exists (Kelley et al., 1963, th. 8.21) a sequence  $t_i \in F$  converging to  $t_0$ . Therefore,  $\int_S g(\cdot, t_i) d\mu \rightarrow \int_S g(\cdot, t_0) d\mu$  uniformly over every weakly (i.e.  $\sigma(M, M^*)$ ) compact subset of the space  $M$  of all bounded regular measures on  $S$  (using Dunford-Pettis' (e.g. Dunford and Schwartz, 1958, IV.8.11 and V.6.1) equi-integrability criterion for those subsets). Thus the mapping  $\phi$  from  $T$  to the space  $C(S)$  of continuous functions on  $S$  mapping  $t \rightarrow g(\cdot, t)$  is injective and such that  $\phi(\bar{F}) = \overline{\phi(F)}$  when  $C(S)$  is endowed with the topology  $\kappa(C(S), M)$  of uniform convergence on  $\sigma(M, M^*)$  compact subsets of  $M$ :  $\phi$  is continuous, and  $\phi(T)$  compact.

For any fixed probability  $\mu \in M$ , and every  $f \in L_1(\mu)$ , let  $\psi(f) \in C(T)$  be defined by  $[\psi(f)](t) = \int f(s)g(s, t) d\mu(s)$ . Since the measures  $f d\mu$  form a  $\sigma(M, M^*)$  compact subset of  $M$  when  $f$  ranges in the unit ball of  $L_\infty(\mu)$ ,  $\psi$  will, by Ascoli's theorem, map the balls of  $L_\infty(\mu)$  into norm-compact subsets of  $C(T)$ . Those being separable, and  $L_\infty(\mu)$  being dense in  $L_1(\mu)$ ,  $\psi$  will map  $L_1(\mu)$  into a separable subspace of  $C(T)$ , i.e. in a space  $C^\mu(T)$  of all continuous functions for some weaker pseudo-metrizable topology on  $T$ .

Let  $S_\mu = \{s \mid g(s, \cdot) \in C^\mu(T)\}$ , so that  $S_\mu$  is a compact (pseudo) metric subset of  $S$ . Let  $s_0$  be any point in the support of  $\mu$ ; let  $O_\alpha$  be the decreasing net of open neighbourhoods of  $s_0$ , and  $\mu_\alpha(A) = \mu(A \cap O_\alpha)/\mu(O_\alpha)$ ; we thus have  $\lim_\alpha \int g(s, t) d\mu_\alpha = g(s_0, t)$  point-wise, and therefore weakly since each integral is in the closed convex extension of  $\phi(T)$  which is weakly compact. Each approximand being in  $\psi(L_1(\mu))$  and thus in the weakly closed space  $C^\mu(T)$ , we have  $g(s_0, t) \in C^\mu(T)$ ; the support of  $\mu$  is thus contained in  $S_\mu$ . Now  $C^\mu(T)$  is a polish space (cf. App.5) in the norm topology, and thus its Borel subsets for the strong topology and for the weak topology coincide (cf. 5.f). Thus  $\mu$  can be considered as a measure on the Borel subsets of  $C^\mu(T)$  with the strong topology. This being polish, there exists a norm-compact subset  $K_\varepsilon$  of  $C^\mu(T)$  with  $\mu(K_\varepsilon) \geq 1 - \varepsilon$ . The set  $K_\varepsilon^\mu = \{s \mid g(s, \cdot) \in K_\varepsilon\}$  is a compact metric subset of  $S$ , with  $\mu(K_\varepsilon^\mu) \geq 1 - \varepsilon$ , and the functions  $g(s, \cdot)$ ,  $s \in K_\varepsilon^\mu$ , are equicontinuous on  $T$  (Ascoli's theorem): the restriction of  $g$  to  $K_\varepsilon^\mu \times T$  is jointly continuous. For an arbitrary measure  $\mu$  on  $S \times T$ , conclude by considering its marginal  $\tilde{\mu}$  on  $S$ .

b. The above proof showed in particular that, if furthermore  $S$  is compact,  $g$  bounded, and  $\mu$  a regular measure on  $T$ , then  $\int g(s, t) d\mu$  is continuous on  $S$ .

HINT. Apply the proof in case  $T$  is furthermore compact.

c. Deduce from I.2Ex.1a and I.2Ex.1b that, if  $S$  and  $T$  are compact, and  $g$  bounded and separately continuous, the mixed extension  $g(\sigma, \tau)$  is well defined, separately continuous, and bi-linear.

HINT. Use Fubini's theorem.

d. If  $S$  and  $T$  are compact, and  $f_\alpha$  is a net of bounded separately continuous functions decreasing point-wise to zero, the mixed extension  $f_\alpha(\sigma, \tau)$  does also.

HINT. Use I.2Ex.1b.

e. Conclude from I.2Ex.1d that products of regular measures extend naturally to the  $\sigma$ -field generated by the limits of increasing nets of positive, separately continuous functions.

HINT. Consider first the case of  $S$  and  $T$  compact, and use the proof of Riesz's theorem (Kelley et al., 1963, p.127).

**2.** Can one replace “convex” in prop. 1.8 p. 5 by convex-like? (Possibly with additional topological assumptions?)

**3.** Let  $S$  and  $T$  be compact metric spaces,  $g$  a bounded measurable function defined on  $S \times T$ , such that, if  $E = \{(s_0, t_0) \mid g(\cdot, t_0)$  is not continuous in  $s$  at  $s_0$  or  $g(s_0, \cdot)$  is not continuous in  $t$  at  $t_0\}$ , then, for each  $s_0$  and  $t_0$ , the sets  $\{t \mid (s_0, t) \in E\}$  and  $\{s \mid (s, t_0) \in E\}$  consist of at most one point (cf. ex. I.2Ex.1 and I.2Ex.4 for the measurability requirement). For any  $(s_0, t_0) \in S \times T$ , let  $\phi_1(s_0, t_0) = \limsup_{s \neq s_0} g(s, t_0)$ ,  $\phi_2(s_0, t_0) = \liminf_{t \neq t_0} g(s_0, t)$ . If  $\min(\phi_1, \phi_2) \leq g \leq \max(\phi_1, \phi_2)$ , then prop. 2.7 p. 19 applies, with  $f = g$  (using remark 2.3 p. 16). For examples of this type, cf. (Karlin, 1950, examples 1, 2, 3 and remark 3). Even if this last condition does not hold, prop. 2.7 still applies, with  $f = \max(\min(\phi_1, \phi_2), \min(g, \max(\phi_1, \phi_2)))$  (or any  $f$  such that  $\min(\phi_2, \max(g, \phi_1)) \leq f \leq \max(\phi_1, \min(g, \phi_2))$ ). The general game of timing of class II (Karlin, 1959, Vol. II, ch. V, ex. 20) falls in this category.

**4.** Let  $S$  and  $T$  be two topological spaces,  $g$  a real valued function on  $S \times T$ , and denote by  $E$  the set of points of discontinuity of  $g$ . If, for every  $(s_0, t_0) \in S \times T$ , the sets  $\{s \mid (s, t_0) \in E\}$  and  $\{t \mid (s_0, t) \in E\}$  are at most countable, then  $g$  is measurable for any regular product probability  $\sigma \otimes \tau$  on the Borel sets of  $S \times T$ .

HINT. Remark that the set of points of discontinuity of a function is always a  $F_\sigma$ , a countable union of closed sets.

**5.** Let  $S$  and  $T$  be both the unit interval with its usual topology. Let  $g$  be a bounded real valued function on  $S \times T$ , satisfying the condition of ex. I.2Ex.4, and such that for each  $t_0 \in T$ ,  $g(\cdot, t_0)$  is lower semi-continuous from the left in  $s$ :  $\liminf_{s \leq s_0} g(s, t_0) \geq g(s_0, t_0)$ ; for each  $s_0 \in S$ ,  $g(s_0, \cdot)$  is upper semi-continuous from the left in  $t$ :  $\limsup_{t \leq t_0} g(s_0, t) \leq g(s_0, t_0)$ ;  $g(0, t)$  is lower semi-continuous in  $t$ ; and  $g(s, 0)$  is upper semi-continuous in  $s$ . Then prop. 2.7 p. 19 applies (with  $f = g$ ). Example 4 of (Karlin, 1950) is in this category.

**6.** Ex. I.2Ex.5 remains true when  $S$  and  $T$  are compact convex sets in Euclidian space, if  $s < s_0$  (resp.  $t < t_0$ ) is understood coordinatewise, and if we require the lower semi-continuity in  $t$  of  $g(s_0, t)$  for all minimal  $s_0$  (those for which  $\{s \in S \mid s < s_0\}$  is empty), and similarly for the upper semi-continuity in  $s$  of  $g(s, t_0)$ . Those requirements of semi-continuity can be dropped for those minimal  $s_0$  such that  $f(s_0, t)$  is dominated. Similarly, prop. 2.7 applies to the general silent duel (Restrepo, 1957) (even when the accuracy functions  $P(t)$  and  $Q(t)$  are only assumed to be upper semi-continuous and left continuous (and with values in  $[0, 1]$ ,  $P(0) = Q(0) = 0$ )).

**7.** (Sion and Wolfe, 1957) Let  $S = T = [0, 1]$ , and let

$$f(s, t) = \begin{cases} -1 & \text{if } s < t < s + \frac{1}{2}, \\ 0 & \text{if } t = s \text{ or } t = s + \frac{1}{2}, \\ 1 & \text{otherwise.} \end{cases}$$

Show that  $\sup_\sigma \inf_t \int f(s, t) d\sigma = \frac{1}{3}$  and  $\inf_\tau \sup_s \int f(s, t) d\tau = \frac{3}{7}$ .

**8.** Let  $S = T = [0, 1]$ , and let

$$f(s, t) = \begin{cases} 0 & \text{for } 0 \leq s < \frac{1}{2} \text{ and } t = 0, \text{ or for } \frac{1}{2} \leq s \leq 1 \text{ and } t = 1, \\ 1 & \text{otherwise.} \end{cases}$$

Then  $f$  satisfies all conditions of theorem 1.6 except the upper semi-continuity in  $s$  at  $t = 1$  and  $\sup_s \inf_t f = 0$ ,  $\inf_t \sup_s f = 1$  (Sion, 1958). Let  $g(s, t) = tf(s, 1) + (1-t)f(s, 0)$ ;  $g(s, t)$  is linear in  $t$ , satisfies all conditions of the theorem, except the upper semi-continuity in  $s$  for  $t > \frac{1}{2}$  and  $\sup_s \inf_t g = 0$ ,  $\inf_t \sup_s g = \frac{1}{2}$ .

**9. Two Machine Gun Duel.** (Karlin, 1959, Vol. II, pp. 225ff.) Players I and II, possessing quantities  $\alpha$  and  $\beta$  of ammunition approach each other without retreat, using quantities  $\mu(ds)$  and  $\nu(ds)$  of ammunition at a distance between  $s$  and  $s + ds$  ( $\int_0^\infty d\mu = \alpha$ ,  $\int_0^\infty d\nu = \beta$ ,  $\mu \geq 0$ ,  $\nu \geq 0$ ). The probability of scoring a hit at a distance between  $s$  and  $s + ds$ , given they are still alive in this interval, is given by  $\xi(s)\mu(ds)$  and  $\eta(s)\nu(ds)$  respectively.

Strategies of the players are the measures  $\mu$  and  $\nu$ , and the pay-off to player I is 1 if II is destroyed without I being destroyed,  $r$  in case of double survival,  $r'$  in case both players are destroyed, and  $-1$  if I is destroyed without II being ( $-1 \leq r, r' \leq 1$ ). Assume that  $\limsup_{h \geq 0} \xi(s+h) \geq \xi(s)$ ,  $\limsup_{h \geq 0} \eta(s+h) \geq \eta(s)$ , and that  $\xi$  and  $\eta$  are upper semi-continuous and bounded.

a. Show that, whatever a player can guarantee against non-atomic strategies of the other player, he can guarantee against any strategy of the other player (i.e. the “monotone class”,  $D(\cdot)$ , generated by the non-atomic strategies contains all strategies).

Show that, if  $\mu$  and  $\nu$  are non-atomic, the probability that both players are still alive when at a distance  $s$  apart, is given by  $Q_{\tilde{\mu}}(s)Q_{\tilde{\nu}}(s)$ , where  $Q_{\tilde{\mu}}(s) = \exp(-\tilde{\mu}(I_s))$ ,  $Q_{\tilde{\nu}}(s) = \exp(-\tilde{\nu}(I_s))$ ,  $I_s = ]s, \infty[$ ,  $\tilde{\mu}(ds) = \xi(s)\mu(ds)$ ,  $\tilde{\nu}(ds) = \eta(s)\nu(ds)$ . In particular,  $\tilde{\mu}_i \rightarrow \tilde{\mu}_0$  implies  $Q_{\tilde{\mu}_i}(s) \rightarrow Q_{\tilde{\mu}_0}(s)$  at every point of continuity of  $Q_{\tilde{\mu}_0}$ . Further,  $Q_{\tilde{\mu}}$  is a convex function of  $\tilde{\mu}$ .

Show that, if  $\mu$  and  $\nu$  have no common atom, the probability that both players are destroyed is zero. If both are non-atomic, the probability that II is destroyed is  $P(\tilde{\mu}, \tilde{\nu}) = \int_0^\infty Q_{\tilde{\nu}}(s)Q_{\tilde{\mu}}(ds)$ . Remark that  $P(\tilde{\mu}, \tilde{\nu})$  is continuous in one variable as soon as the other is non-atomic.

Show that, for any  $\tilde{\mu}$ ,  $P(\tilde{\mu}, \tilde{\nu})$  is a convex function of  $\tilde{\nu}$ .

Show also that  $P(\tilde{\mu}, \tilde{\nu}) = 1 - Q_{\tilde{\mu}}(0-)Q_{\tilde{\nu}}(0-) - P(\tilde{\nu}, \tilde{\mu})$ , when one of both  $\tilde{\mu}$  and  $\tilde{\nu}$  is non-atomic, and conclude that, for  $\tilde{\nu}$  non-atomic,  $P(\tilde{\mu}, \tilde{\nu})$  is a concave and continuous function of  $\tilde{\mu}$ . For any bounded positive measure  $\sigma$  on the real line, let, for every Borel  $A$ ,  $\sigma_+^\varepsilon(A) = E(\sigma(A - \eta))$  and  $\sigma_-^\varepsilon = E(\sigma(A + \eta))$  where  $\eta$  is a random variable uniform on  $[0, \varepsilon]$  ( $E$  is expectation). For any  $\mu$ ,  $\nu$ , we have  $P(\mu, \nu) = \lim_{\varepsilon \rightarrow 0} P(\mu, \nu_-^\varepsilon)$ ,  $P(\mu, \nu_-^\varepsilon) = \lim_{\eta \rightarrow 0} P(\mu_+^\eta, \nu_-^\varepsilon)$ , both limits being decreasing. The mappings  $\sigma \rightarrow \sigma_+^\varepsilon$  and  $\sigma \rightarrow \sigma_-^\varepsilon$  are linear and continuous for the weak topology on  $\sigma$  and the norm topology on  $\sigma_+^\varepsilon$  (resp.  $\sigma_-^\varepsilon$ ).

Show first that this implies that  $R(\mu, \nu) = P(\tilde{\mu}, \tilde{\nu})$  is upper semi-continuous on the product space, and concave in  $\mu$  for any  $\nu$  (and convex in  $\mu$  for any  $\nu$ ). Show also that it implies that whatever be  $\mu$  and  $\nu$ ,  $P(\tilde{\mu}, \tilde{\nu})$  is the probability that II be destroyed. Since I needs only to consider non-atomic strategies of II, he is faced with the pay-off function  $f(\mu, \nu) = (1 - r)P(\tilde{\mu}, \tilde{\nu}) - (1 + r)P(\tilde{\nu}, \tilde{\mu}) + r$ , which is concave and continuous in  $\tilde{\mu}$ ,  $\tilde{\nu}$  being non-atomic. Thus I has a pure optimal strategy, say  $\tilde{\mu}_0$  (prop. 1.8 p. 5). Similarly, II is faced with the same pay-off function, and thus has a pure optimal strategy, say  $\tilde{\nu}_0$ .

Conclude that the game has a value, and both players have optimal strategies, and none of those depend on  $r'$ .

b. The above solution is for the case of a “silent” duel: none of the players is informed in the course of the game about the quantity of ammunition spent by the other player.

Conclude that the solution remains the same in the noisy duel.

c. Show also that both players have  $\varepsilon$ -optimal strategies that have a bounded density with respect to Lebesgue measure.

d. The above results remain a fortiori true if an additional restriction, say an upper bound on the speed of firing (as used by Karlin) is imposed on the strategy spaces (this restriction may look natural in our interpretation of the model, but may be less natural in an interpretation e.g. in terms of an advertising campaign). What do our results for the unbounded case imply about the behaviour of the value in the bounded case when the bounds get large?

e. What happens to the above results when the accuracy functions  $\xi$  and  $\eta$  are not necessarily bounded? (for instance,  $\lim_{s \geq 0} \xi(s) = \lim_{s \geq 0} \eta(s) = +\infty$ ?)

**10.** Use the results of the present section (notably part 2.b)) to improve those of ex. I.1Ex.2–I.1Ex.6 p. 9. In particular, the convergence concept should allow for convergence of the games  $\exp[-n(s - t)^2]$  to the zero game, not only to the indicator of the diagonal.

**11.** (Kuhn, 1952, p. 118) Player I picks a number  $x$  in  $[0, 1]$ , player II a continuous function  $f$  from  $[0, 1]$  to itself with  $\int_0^1 f(t) dt = 1/2$ . The pay-off is  $f(x)$ . Thus, I’s strategy space is compact metric, II’s strategy space is complete, separable and metric, and the pay-off function is jointly uniformly continuous. Nevertheless, if player I is restricted to mixed strategies with finite support, he can guarantee himself only zero, while otherwise  $1/2$  (and player II guarantees himself  $1/2$  with a pure strategy).

**12.** (Grothendieck, 1953) The Mackey topology  $\tau(L_\infty, L_1)$  coincides on bounded subsets of  $L_\infty$  with the topology of convergence in measure.

HINT. One direction is given in the text in the proof of Lemma 2.3 p. 14. For the other, show first that it is sufficient to prove that a uniformly bounded sequence that converges in measure converges uniformly on weakly compact subsets of  $L_1$ . To obtain this, just use Dunford-Pettis' equi-integrability criterion.

*The next series of exercises concerns applications of the minmax theorem — i.e. the separation theorem — to the problem of how to assign a limit to non-converging sequences, i.e. how to define the pay-off function in an infinitely repeated game.*

**13. Banach Limits.** A Banach limit  $\mathcal{L}$  is a linear functional on  $\ell_\infty$ , such that

$$\mathcal{L}((x_n)_{n \in \mathbb{N}}) \leq \limsup_{n \rightarrow \infty} x_n .$$

We will also write  $\mathcal{L}(x_n)$ .

a. Show (cf. 1.20 p. 8) that Banach limits exist.

b. Banach limits are positive linear functionals of norm 1.

c. Banach limits can equivalently be defined as regular probability measures on the Stone-Čech compactification (Kelley, 1955, p. 152)  $\beta(\mathbb{N})$  of the integers, assigning probability zero to  $\mathbb{N}$ .

d. If  $C$  is a compact, convex subset of a locally convex Hausdorff space, and  $x_n \in C$  for all  $n$ , there exists a unique  $\mathcal{L}(x_n) \in C$  such that, for each continuous linear functional  $\phi$ ,  $\langle \phi, \mathcal{L}(x_n) \rangle = \mathcal{L}(\langle \phi, x_n \rangle)$ .

e. In particular, if  $X_n$  is a uniformly integrable sequence of random variables, there exists a unique random variable  $X_\infty = \mathcal{L}(X_n)$  such that, for each measurable set  $A$ ,  $\int_A X_\infty = \mathcal{L}(\int_A X_n)$ .

f. Similarly, if  $X_n$  is a sequence of r.v. with values in a compact convex subset  $C$  of  $\mathbb{R}^k$ , there exists a unique r.v.  $X_\infty$  with values in  $C$  such that  $\mathcal{L} \mathbf{E}(X_n, Y) = \mathbf{E}(X_\infty, Y)$  for all  $Y \in (L_1)^k$ .

**14. A separation theorem.** (Meyer, 1973)

a. Let  $X$  be a compact convex subset of a locally convex vector space  $V$ .  $\mu$  will stand for an arbitrary regular probability measure on  $X$ . Any  $\mu$  has a barycentre  $b_\mu \in X$  (defined to be such that, for any continuous affine functional  $u$  on  $X$ ,  $u(b_\mu) = \int u d\mu$ ). If  $u_1$  and  $u_2$  are two u.s.c. concave functions on  $X$ , one of them bounded, show that their least upper bound in the set of all concave functions is bounded and u.s.c. Show also (1.21 p. 8) that  $u_1(b_\mu) \geq \int u_1 d\mu$ . Denote by  $\Gamma_-$  the set of bounded concave functions  $u$  on  $X$  that are the least upper bound of a sequence of u.s.c. concave functions, and let  $\Gamma_+ = -\Gamma_-$ . Show that the sequence can, without loss, be assumed monotone, and that:

$$(\star) \quad u(b_\mu) \geq \int u d\mu$$

b. Let  $u \in \Gamma_-, v \in \Gamma_+, u \leq v$ . Show that, for all  $\mu$ , there exist  $u' \in \Gamma_-$  and  $v' \in \Gamma_+$  such that  $u \leq u' \leq v' \leq v$  and  $u' = v'$  a.e.

HINT. Let  $u_n$  u.s.c. concave and bounded increase strictly to  $u$ , and dually for  $v_n$  and  $v$ . Let  $A_n = \{\phi \mid \phi \text{ affine continuous, } u_n \leq \phi \leq v_n\}$ .  $A_n$  decreases and is non-empty (1.21) and convex, so that its closure  $\bar{A}_n$  in  $L_1(\mu)$  is weakly closed and thus (boundedness) weakly compact. Let  $\phi \in \bigcap_n \bar{A}_n$  s.t.  $\|\phi - \phi_n\|_1 \leq 2^{-n}$  with  $\phi_n \in A_n$ ; thus  $\phi_n \rightarrow \phi$   $\mu$ -a.e. Let  $u' = \liminf_{n \rightarrow \infty} \phi_n$ ,  $v' = \limsup_{n \rightarrow \infty} \phi_n$ . Note further that  $u'$  and  $v'$  can be assumed to be in the cones  $\Gamma_-$  and  $\Gamma_+$  for some metrisable quotient space  $\tilde{X}$  (consider the weak topology on  $V$  generated by the functions  $\phi_n$ ).

c. Call a bounded function  $w$  on  $X$  strongly affine if, for any  $\mu$ ,  $w$  is  $\mu$ -measurable and  $w(b_\mu) = \int w d\mu$ . Let  $u \in \Gamma_-$ ,  $v \in \Gamma_+$ ,  $u \leq v$ . Assume the continuum hypothesis, and show that there exists a strongly affine  $w$  such that  $u \leq w \leq v$ .

HINT. By I.2Ex.14b),  $X$  can, without loss of generality, be assumed metrisable. The set of all probabilities on  $X$ , being compact metric, has the power of the continuum. Let thus  $\mu_\alpha$  be an indexing of it by the set of all countable ordinals (continuum hypothesis). Construct by transfinite induction  $u_\alpha \in \Gamma_-$  and  $v_\alpha \in \Gamma_+$  such that  $\alpha < \beta$  implies  $u \leq u_\alpha \leq u_\beta \leq v_\beta \leq v_\alpha \leq v$  and  $u_\alpha = v_\alpha$   $\mu_\alpha$ -a.e. (point I.2Ex.14b)). Since all point masses  $\varepsilon_x$  are among the  $\mu_\alpha$ , the  $u_\alpha$  and the  $v_\alpha$  have a common limit  $w$ , which is strongly affine by  $(\star)$ .

### 15. Medial limits (Mokobodzki). (cf. Meyer, 1973)

a. Show that, under the continuum hypothesis, there exist positive linear functionals  $\ell$  ("medial limits") of norm 1 on  $\ell_\infty$  such that, if  $x_n$  is a convergent sequence,  $\ell(x_n) = \lim_{n \rightarrow \infty} x_n$ , and such that, for any uniformly bounded sequence of random variables  $Z_n(\omega)$ ,  $\ell(Z_n(\omega))$  is measurable and  $E[\ell(Z_n(\omega))] = \ell[E(Z_n(\omega))]$ .

Show that  $\ell$  can even be chosen so as to satisfy  $\ell(x_n) = x_\infty$  for any sequence  $x_n$  converging to  $x_\infty$  in Cesàro's sense (or equivalently Abel's sense (cf. ex. I.2Ex.16 and I.2Ex.17)). This would, in particular, imply that  $\ell$  is translation invariant:  $\ell(x_n) = \ell(x_{n+1})$ .

Even stronger: one can choose  $\ell$  such that, for any  $x \in \ell_\infty$ ,  $\ell(x) = \ell(\bar{x})$ , where  $\bar{x}_n = \frac{1}{n} \sum_{i=1}^n x_i$ .

HINT. Take  $X = [-1, 1]^\mathbb{N}$ . For any  $x \in \ell_\infty$ , let  $\phi(x) = (x - \bar{x})/2$ .  $\phi(X)$  is a compact convex subset of  $X$ . Let  $V_n = \{(t, x) \mid 1 \geq t \geq \sup_{i \geq n} x_i, x \in X\}$ , and  $U_n$  the convex hull of  $-V_n$  and  $[-1, 0] \times \phi(X)$ ,  $u_n(x) = \max\{t \mid (t, x) \in U_n\}$ ,  $v_n(x) = \min\{t \mid (t, x) \in V_n\}$ .

Prove that  $u_n \leq v_n$  (show that  $[V_n + (\varepsilon, 0)] \cap [U_n - (\varepsilon, 0)] = \phi$ ). Let  $u = \lim_{n \rightarrow \infty} u_n$ ,  $v = \lim_{n \rightarrow \infty} v_n$ :  $\liminf_{i \rightarrow \infty} x_i \leq u(x) \leq v(x) = \limsup_{i \rightarrow \infty} x_i$ , and  $x \in \phi(X)$  implies  $u(x) \geq 0$ . Apply ex. I.2Ex.14c to get  $\ell$  strongly affine  $u \leq \ell \leq v$ . Show that  $\ell(0) = 0$ , and extend  $\ell$  by homogeneity to a linear functional on  $\ell_\infty$ . Show that  $\ell$  is positive, of norm 1, and satisfies, for any  $x \in \ell_\infty$ ,  $\liminf_{n \rightarrow \infty} x_n \leq \ell(x) = \ell(\bar{x}) \leq \limsup_{n \rightarrow \infty} x_n$ . By homogeneity, one can assume that  $Z_n$  is an  $X$ -valued random variable  $Z$ . The barycentre of the distribution  $\mu_z$  of  $Z$  on  $X$  is the sequence  $(E(Z_n))$ . The formula  $E[\ell(Z_n(\omega))] = \ell[E(Z_n(\omega))]$  is thus equivalent to the strong affinity of  $\ell$ .

b. Using ex.I.2Ex.13c,  $\ell$  can be extended to all sequences  $x_n$  such either  $(\bar{x})^+$  or  $(\bar{x})^-$  is  $\ell$ -integrable. Show that:

- (1) if a sequence of random variables  $X_n(\omega)$  is bounded in  $L_1$ , then  $\ell(X_n(\omega))$  exists a.e. and is in  $L_1$ .
- (2) if the sequence is uniformly integrable, then  $E[\ell(X_n(\omega))] = \ell[E(X_n(\omega))]$ .
- (3) in particular, if  $X_n(\omega)$  converges weakly in  $L_1$  to  $X(\omega)$ , then  $\ell(X_n(\omega)) = X(\omega)$  a.e.
- (4) if a sequence of random variables converges in probability to a real-valued random variable  $X(\omega)$ , then  $\ell(X_n(\omega)) = X(\omega)$  a.e.

### 16. Abelian theorems.

a. If  $x_n \in \ell_\infty$ ,  $y_n = \sum_i p_{n,i} x_i$ ,  $p_{n,i} \geq 0$ ,  $\sum_i p_{n,i} = 1$ , and if  $\lim_{n \rightarrow \infty} p_{n,i} = 0$ , then  $\limsup y_n \leq \limsup x_n$ .

b. In particular, if also  $p_{n,i} \geq p_{n,i+1}$ , then  $\limsup y_n \leq \limsup \bar{x}_n$ .

HINT. Rewrite the  $y_n$  as convex combinations of the  $\bar{x}_n$ , and use I.2Ex.16a.

c. In particular,  $\limsup_{\lambda \rightarrow 0} \lambda \sum_n (1 - \lambda)^n x_n \leq \limsup_{n \rightarrow \infty} \bar{x}_n$ , (the  $\limsup$  in Abel's sense is smaller than the  $\limsup$  in Cesàro's sense.)

d. Denote by  $P_t^\alpha$  ( $0 \leq \alpha \leq 1$ ) the one-sided stable distribution of index  $\alpha$  (Feller, 1966, Vol. II, XIII.6) (i.e., with Laplace transform  $\exp(-t\lambda^\alpha)$ ). Observe that  $\int_0^\infty P_s^\alpha(\cdot)P_t^\beta(ds) = P_t^{\alpha\beta}(\cdot)$  (subordination, e.g. (Feller, 1966, Vol. II, XIII, 7.e)). Given a bounded measurable function  $x(t)$  on  $\mathbb{R}_+$ , let  $p_\alpha(x) = \limsup_{t \rightarrow \infty} \int_0^\infty x(s)P_t^\alpha(ds)$ . Conclude from the subordination property and from I.2Ex.16a that  $\alpha \leq \beta \Rightarrow p_\alpha(x) \leq p_\beta(x)$ . In particular,  $p_0(x) = \lim_{\alpha \rightarrow 0} p_\alpha(x)$  is a well-defined sublinear functional: more precisely  $p_0(x+y) \leq p_0(x) + p_0(y)$ ,  $p_0(\lambda x) = \lambda p_0(x)$  for  $\lambda \geq 0$ ,  $p_0$  is monotone,  $p_0(1) = 1$ ,  $p_0(-1) = -1$ . And  $p_0 \in \Gamma_+$  (cf. ex. I.2Ex.14), say on the unit ball  $X$  of  $L_\infty$ .

COMMENT 2.14. Observe that  $p_0$  is fully canonical, given the additive semigroup structure and the multiplication by positive scalars on  $\mathbb{R}_+$  — or equivalently, given the additive semigroup structure and the topology. There is also a fully canonical way of transforming a problem of limit of sequences to a problem on  $\mathbb{R}_+$  (as a topological semigroup) (provided one uses on  $\mathbb{R}_+$  only limits  $\ell$  satisfying  $\ell(f(\lambda t)) = \ell(f(t))$  for all  $\lambda > 0$  and  $f$  bounded). This uses the Poisson process: if  $P_t$  denotes the Poisson distribution on  $\mathbb{N}$  at time  $t$ , map first the bounded sequences  $x_n$  to the function  $P_t(x) = \sum_n P_t(n)x_n$ .

COMMENT 2.15. One might wish to add further requirements to a “uniform distribution” on the integers  $\mathcal{L}$ , like  $\mathcal{L}(x_n) \leq \limsup_{\lambda \rightarrow 1} (1-\lambda) \sum_n \lambda^n x_n$ , or like the fact that quotient and rest of a uniform random number upon division by  $k$  are also uniform and independent (given the shift invariance which is guaranteed by any other requirement, the latter requirement amounts to  $\mathcal{L}(x_n) = k \mathcal{L}(y_n)$  if  $(y_n)$  is obtained from  $(x_n)$  by inserting  $(k-1)$  zeros between two successive values).

COMMENT 2.16. What type of limit operations do those considerations lead to? In particular, is there any relationship with the sublinear functional  $q(x) = \limsup_{n \rightarrow \infty} (1/\ln n)(\sum_{i=1}^n x_i/i)$ ? What is the relation between such Banach limits (i.e. satisfying  $\mathcal{L}(x) \leq q(x)$ ) and those for which  $\mathcal{L}(x_n) = \mathcal{L}(\bar{x}_n)$ ?

**17. Hardy and Littlewood's Tauberian theorem.** This provides a partial converse to I.2Ex.16c. Assume the sequence  $x_n$  is bounded from below. If  $\lim_{\lambda \rightarrow 1} (1-\lambda) \sum_n x_n \lambda^n = \ell$  exists, then  $\lim \bar{x}_n = \ell$ .

HINT. By adding a constant, we can assume  $x_n \geq 0$ ; and also  $\ell = 1$  by normalisation.

Show that  $\lim_{\lambda \rightarrow 1} (1-\lambda) \sum_{n=0}^\infty x_n \lambda^n P(\lambda^n) = \int_0^1 P(t) dt$  first for  $P(t) = t^k$ , then for a polynomial, then a continuous function, and finally, for a bounded function whose set of discontinuities has Lebesgue measure zero.

Apply this for  $P(t) = 0$  ( $0 \leq t \leq e^{-1}$ ),  $P(t) = t^{-1}$  ( $t > e^{-1}$ ),  $\lambda = \exp(-1/n)$ .

### 3. The minmax theorem for ordered fields

DEFINITION 3.1. An **ordered field**  $K$  is a (commutative) field  $K$  together with a subset of positive elements  $P$  such that:  $P$  is closed under addition and multiplication; for any element  $x \in K$ , one and only one of the following holds:  $x \in P$ ,  $x = 0$ ,  $-x \in P$ . For two elements  $x$  and  $y$  of  $K$ , we write  $x > y$  iff  $x - y \in P$  and similarly for  $\geq, <, \leq$ .

An ordered field is called **real closed** if it has no ordered algebraic extension. An ordered field is contained in a minimal real closed field called its **real closure** (Jacobson, 1964, Th. 8, p. 285).

Ordered fields arise naturally in studying the asymptotic behaviour of repeated games, e.g. the field of Puiseux series or the real closure of the field of rational fractions in the discount factor  $\lambda$  when studying the asymptotic behaviour of  $v_\lambda$  for stochastic games. Since because of those applications we will have to work with such real closed fields we will use related tools rather than remaining systematically with the elementary methods of ordered fields.

**DEFINITION 3.2.** A  **$K$ -polynomial system** (resp. linear)  $\mathcal{S}$  is a finite number of polynomial (resp. linear) equations and inequalities ( $P = 0, P > 0, P \neq 0$ ) with coefficients in  $K$ .

A **semi-algebraic set** in  $K^n$  ( $K$  a real closed field) is a finite (disjoint) union of sets defined by a  $K$ -polynomial system. Observe that semi-algebraic sets form a Boolean algebra (spanned by the sets  $\{x \mid P(x) > 0\}$  where  $P$  is a polynomial).

A **semi-algebraic function** (correspondence) is one whose graph is semi-algebraic.

**THEOREM 3.3. (Tarski, cf. Jacobson, 1964, Th. 16, p. 312).** Given a  $\mathbb{Q}$ -polynomial system  $\mathcal{S}$  in  $n+m$  variables, there exists a finite family of  $\mathbb{Q}$ -polynomial systems in  $n$  variables,  $\mathcal{T}_i$ , such that for any real closed field  $K$  the projection on  $K^n$  of the semi-algebraic set in  $K^{n+m}$  defined by  $\mathcal{S}$  is the semi-algebraic set in  $K^n$  defined by the  $\mathcal{T}_i$ 's.

**COROLLARY 3.4.** The projection on  $K^n$  of a semi-algebraic set in  $K^{n+m}$  is semi-algebraic. Equivalently one can also allow quantifiers (over variables in  $K$ ) in the sentences defining semi-algebraic sets (besides any logical connectives).

**COROLLARY 3.5.** Images and inverse images of semi-algebraic sets by semi-algebraic functions (correspondences) are semi-algebraic. The closure and the interior of semi-algebraic sets in  $\mathbb{R}^n$  are semi-algebraic.

**COROLLARY 3.6.** If a semi-algebraic set in  $K^{n+m}$  defined by  $\mathbb{Q}$ -polynomial systems  $\mathcal{S}_i$  has a projection on  $K^n$  equals to  $K^n$ , then on every real closed field the semi-algebraic set defined through the  $\mathcal{S}_i$ 's has the same property.

**THEOREM 3.7.** Let  $\tilde{K}$  be an ordered field and  $K$  a subfield. If a  $K$ -linear system has a solution in  $\tilde{K}$ , it has a solution in  $K$ .

**PROOF.** Denote by  $F_i$ ,  $i = 1, \dots, n$  the linear equalities and inequalities in  $m$  unknowns  $x_1, \dots, x_m$  and coefficients in  $K$ . Note that for any  $x \in \tilde{K}$ ,  $x > 0$ , there exists  $\varepsilon \in K$  such that  $0 < \varepsilon < x$ . Consider now a solution  $x^0 = (x_1^0, \dots, x_m^0)$  in  $\tilde{K}^m$ : replace any  $\geq$  (resp.  $\leq$ ) sign in the  $F_i$  by " $>$ " (resp. " $<$ ") or " $=$ " in such a way that  $x^0$  still be a solution. Replace now any " $>$ " (resp. " $<$ ") in the conditions by " $\geq \varepsilon$ " (resp. " $\leq -\varepsilon$ ") with  $\varepsilon > 0$  in  $K$  in such a way that  $x^0$  still be a solution. Take any solution  $y^1$  in  $K^m$  of the equalities of the system. If  $y^1$  does not satisfy the inequalities, then there exists  $\gamma \in \tilde{K}$  such that  $x^1 = \gamma y^1 + (1 - \gamma)x^0$  is still a solution, but one of the  $\geq$  (or  $\leq$ ) becomes an equality at  $x^1$ . Repeat the procedure with  $x^1$ . Since the number of inequalities decreases at every repetition, at some stage  $k$ ,  $y^k$  will satisfy the system, if only because no inequalities are left. ■

**REMARK 3.1.** For a more elementary approach to the above and the next results, cf. comments 3.7 and 3.8 after ex. I.3Ex.4 p. 29. For more powerful versions of the next result, cf. sect. 4 ex. I.4Ex.4 p. 39 and I.4Ex.8 p. 41.

**THEOREM 3.8.** Let  $A = (a_{\ell m})$  be an  $L \times M$ -matrix with elements  $a_{\ell m}$  in an ordered field  $K$ . Then, there exists a unique element  $v$  of  $K$  and there exists  $x_1, \dots, x_L$  and  $y_1, \dots, y_M$  in  $K$  such that  $x_\ell \geq 0 \quad \forall \ell \ 1 \leq \ell \leq L$ ,  $\sum_{\ell=1}^L x_\ell = 1$ ,  $y_m \geq 0 \quad \forall m \ 1 \leq m \leq M$ ,  $\sum_{m=1}^M y_m = 1$ , and

$$\begin{aligned} \sum_{\ell=1}^L x_\ell a_{\ell m} &\geq v, \forall m, 1 \leq m \leq M, \\ \sum_{m=1}^M y_m a_{\ell m} &\leq v, \forall \ell, 1 \leq \ell \leq L. \end{aligned}$$

PROOF. Theorem 1.6 p. 4 implies the result to be true when  $K$  is the real number field. Corollary 3.6 implies that it remains true for any real closed field. Hence that the system of  $2(L + M + 1)$  linear equalities and inequalities in  $(x_\ell, y_m, v)$  with coefficients in  $K$  has a solution in its real closure  $\tilde{K}$  (with unique  $v$ ). Theorem 3.7 implies it has a solution in  $K$ . ■

### Exercises.

**1. Farkas' lemma.** A finite system of linear equations and inequalities (on a vector space on an ordered field) is inconsistent (i.e. has no solution) iff there is an inconsistency proof by a linear combination.

COMMENT 3.2. This is understood modulo the usual rules for inequalities: equalities are preserved after multiplication by any scalar, weak or strict inequalities are preserved (or change sign, or turn into an equality) according to the sign of the scalar; and sums of equations and weak and strict inequalities also follow the usual rules. An inconsistency proof mean that one obtains in this way either  $0 \geq 1$  or  $0 > 0$ .

COMMENT 3.3. One reason for allowing both weak and strict inequalities is to be able to express frequent statements of the type “the system  $f_i(x) \geq a_i$  implies  $f(x) \geq a$ ” in this language as the inconsistency of the system “ $f_i(x) \geq a_i$  and  $f(x) < a$ ”.

COMMENT 3.4. The other reason is that then the “dual” of the system  $f_i(x) \geq a_i, g_j(x) > b_j, h_k(x) = c_k$  — i.e. the system expressing its inconsistency — is the system of the same form

$$\begin{aligned} \sum \lambda_i f_i + \sum \mu_j g_j + \sum \nu_k h_k &= 0, & \lambda_i \geq 0, \mu_j \geq 0, \\ \sum \lambda_i a_i + \sum \mu_j b_j &+ \sum \nu_k c_k \geq 0, \\ \sum \lambda_i a_i + \sum \mu_j b_j &+ \sum \nu_k c_k + \sum \mu_j > 0. \end{aligned}$$

COMMENT 3.5. The duality terminology is justified by the fact that the second dual of a system is the system itself: let  $(e_\alpha)_{\alpha \in A}$  be the inverse images in our vector space  $E$  of a basis of the finite dimensional quotient of  $E$  by the map  $(f, g, h)$ . Then the equalities in our dual system translate to  $\sum \lambda_i f_i(e_\alpha) + \sum \mu_j g_j(e_\alpha) + \sum \nu_k h_k(e_\alpha) = 0 \forall \alpha \in A$ . Thus the second dual is of the form  $\sum \lambda_i f_i(x) + \sum \mu_j g_j(x) + \sum \nu_k h_k(x) + \sum u_i \lambda_i + \sum v_j \mu_j + w[\sum \lambda_i a_i + \sum \mu_j b_j + \sum \nu_k c_k] + r(\sum \lambda_i a_i + \sum \mu_j b_j + \sum \nu_k c_k + \sum \mu_j) = 0$ , with  $x = \sum x_\alpha e_\alpha$ ,  $u_i \geq 0$ ,  $v_j \geq 0$ ,  $w \geq 0$ ,  $r > 0$ . The equation yields  $f_i(x) + u_i + wa_i + ra_i = 0$ ,  $g_j(x) + v_j + wb_j + rb_j + r = 0$ ,  $h_k(x) + wc_k + rc_k = 0$ , thus  $f_i(x) + (w+r)a_i \leq 0$ ,  $g_j(x) + (w+r)b_j \leq -r$ ,  $h_k(x) + (w+r)c_k = 0$ . Let  $y = -x/(w+r)$  (note  $w+r > 0$ ): we have  $f_i(y) \geq a_i$ ,  $g_j(y) \geq b_j + r/(w+r)$ ,  $h_k(y) = c_k$ : it expresses that there exists  $x \in E$  and  $\varepsilon (= r/w+r) > 0$  such that  $f_i(x) \geq a_i$ ,  $g_j(x) \geq b_j + \varepsilon$ ,  $h_k(x) = c_k$ .

The lemma states then that a system is inconsistent iff its dual is consistent.

HINT. It suffices to prove that, if the system is inconsistent, the dual has a solution. As seen sub 3.5 above, it suffices to consider a finite dimensional vector space over the field  $K$ . One gets easily rid of the equations, replacing them by the pair of opposite weak inequalities — this leads in effect to the same dual system. Make the system homogeneous by multiplying the right hand member by an additional variable  $x_0$ , and adding the inequality  $x_0 > 0$ : the system and its dual remain equivalent to the original.

Take now advantage of the homogeneity to replace all inequalities  $g_j(x) > 0$  by  $g_j(x) \geq z$ ,  $z \geq 1$ . Both systems are still equivalent, and are now of the form  $\tilde{f}_i(x) \geq 0$ ,  $x_0 \geq 1$ , and  $\sum \tilde{\lambda}_i \tilde{f}_i(x) + x_0 = 0$ ,  $\tilde{\lambda}_i \geq 0$ . Fixing a basis in the finite dimensional vector space, the  $\tilde{f}_i$  are given by a matrix  $F$  with elements in the ordered field  $K$  (and also  $\tilde{a}_i \in K$ ). We can assume  $K$  real closed (theorem 3.7 p. 27) and then, for any fixed dimensions of  $F$ , uses Tarski's principle (cor. 3.6 p. 27): it suffices to deal with real vector space.

The system being inconsistent means that the subspace  $V = \{(x_0, \tilde{f}_i(x))_{i=1}^{k-1} \mid x \in E\} \subseteq \mathbb{R}^k$  satisfies  $V \cap Q = \emptyset$ , with  $Q = (1, 0, 0, \dots) + \mathbb{R}_+^k$ . This implies  $\delta = d(V, Q) > 0$ : otherwise choose a sequence

$x^n \in V$ ,  $y^n \in Q$  with  $\|x^n - y^n\| \leq n^{-1}$  and  $\|y^n\|$  minimal under those constraints.  $V$  and  $Q$  being closed and disjoint implies  $\|x^n\| \sim \|y^n\| \rightarrow +\infty$ . The minimality implies  $y_0^n = 1$ . Extracting a subsequence, we get that  $x^n/\|x^n\|$  and  $y^n/\|y^n\|$  converge, say to  $z$ , with  $z_0 = 0$ ,  $\|z\| = 1$ ,  $z \in \mathbb{R}_+^k$ , and  $z \in V$ . Hence, for  $n$  sufficiently large,  $x^n - z \in V$  and  $y^n - z \in Q$  contradict the minimality of  $\|y^n\|$ . Thus  $V$  is still disjoint from the open convex set  $Q^\delta = \{y \mid d(y, Q) < \delta\}$ . Separation yields then the result (1.19).

## 2. Strong complementarity. (Bohnenblust et al., 1950)

a. If strategy sets are finite, there exists an optimal strategy pair such that every best reply is used with positive probability.

HINT. Let  $(\sigma, \tau)$  be an optimal pair with maximal support, and assume w.l.o.g. that  $v(g) = 0$  and that  $g(\sigma, t_0) = 0 = \tau(t_0)$ . Thus, the system of linear inequalities in  $\sigma$ :  $[g(\sigma, t) \geq 0]_{t \in T}$  and  $\sigma \geq 0$  implies  $g(\sigma, t_0) \leq 0$ ; hence, by ex. I.3Ex.1 there exist  $\lambda_t \geq 0$  such that  $g(s, t_0) + \sum_t \lambda_t g(s, t) \leq 0$  for all  $s \in S$ : i.e.  $(1 + \lambda_0, \lambda_1, \lambda_2, \dots)/1 + \sum \lambda_t$  is an optimal strategy  $\tilde{\tau}$  of player II with  $\tilde{\tau}(t_0) > 0$ .

Alternative: reduce to a symmetric game (antisymmetric matrix  $A$ ), and express by ex. I.3Ex.1 the inconsistency of  $xA \geq 0$ ,  $x \geq 0$ ,  $xA + x > 0$ : this yields the existence of  $y \geq 0$ ,  $z \geq 0$ ,  $z \neq 0$  s.t.  $A(y + z) + z \leq 0$ . Make a scalar product with  $y + z$ .

b. Hence every pair in the relative interior of the optimal strategy sets has this property. [In particular the barycentre, which has all the symmetries of the game.] And conversely every pair with this property is a relative interior pair.

## 3. Examples.

a. In the game  $\begin{pmatrix} x & -x^2 \\ -1 & 0 \end{pmatrix}$ , where player I chooses  $x \in [0, 1]$  or Bottom, both players have unique optimal strategies, and both have a best reply that is used with probability zero: finiteness of both strategy sets is essential, even for strong complementarity to hold for a single player.

[Ex. I.4Ex.9 p. 42 provides a striking example, where both pure strategy sets are  $[0, 1]$ , and both players have unique optimal strategies, whose support is nowhere dense, while every pure strategy is a best reply. Even better is the polynomial game  $st(t - s)$ .]

b. The “Cobb-Douglas” cone  $C = \{(x, y, z) \mid z^2 \leq xy, x + y \geq 0\} \subseteq \mathbb{R}^3$  is closed, convex and disjoint from the straight line  $D = \{(x, 0, 1)\}$ , yet for any linear functional  $\varphi$  one has  $\varphi(C) \cap \varphi(D) \neq \emptyset$ : there is no hope to obtain anything like ex. I.3Ex.1 in the non-polyhedral case, and the explicit use of the structure of  $\mathbb{R}_+^k$  is necessary at the end of that proof.

**4. Structure of polyhedra.** A (closed, convex) polyhedron is a finite intersection of closed half spaces (in a finite dimensional vector space  $E$  over an ordered field  $K$ ).

a. The projection (hence (finite dimensionality) any affine image) of a polyhedron is a polyhedron (and clearly so are the inverse images).

HINT. The projection being from  $(x, y)$ -space to  $x$ -space, express by ex. I.3Ex.1 that  $x$  does not belong to the image, obtaining thus furthermore a description of the inequalities of the projection as those convex combinations of inequalities of the polyhedron which are independent of  $y$ . Alternatively, eliminating one coordinate at a time is completely elementary.

b. The product of two polyhedra is a polyhedron.

c. The sum (hence the difference) of two polyhedra is a polyhedron. [Use I.3Ex.4a and I.3Ex.4b]. So is the intersection.

d. Define the dual  $P^0$  of a polyhedron  $P$  as  $\{(a, b) \in E^* \times K \mid \langle a, x \rangle \geq b, \forall x \in P\}$ . The dual is a polyhedral cone (in  $E^* \times K$ ).

HINT. It suffices to consider a non-empty polyhedron. Expressing then the inconsistency of  $Mx \geq m$  with  $ax < b$  yields (ex. I.3Ex.1),  $a = (y_1, \dots, y_k)M$ ,  $b = (y_1, \dots, y_k)m - y_0$  for  $y \geq 0$ ; use now I.3Ex.4a.

e. The convex hull of two polyhedra cannot be expressed as the set of solutions to a system of linear inequalities — even allowing strict inequalities —: consider a line and a point. But the closed convex hull of two polyhedra is a polyhedron: consider the intersection I.3Ex.4c of their duals I.3Ex.4d. In particular, the convex hull of finitely many compact polyhedra is a polyhedron (cf. ex. I.3Ex.10d p. 35).

f. A polyhedral cone is defined by finitely many inequalities  $Ax \geq 0$ . (If  $Ax \geq b$ , necessarily  $b \leq 0$ , and  $Ax \geq 0$ ). Conclude that a dual cone is of the form  $\{ (a, b) \mid aX_1 \geq 0, aX_2 \geq b \}$  where  $X_1$  and  $X_2$  are columns of vectors in  $E$ . [Note that  $(0, -1)$  is in the dual cone.]

g. The polyhedra are the sums of a compact polyhedron and a polyhedral cone.

HINT. I.3Ex.4c yields one direction. For the other, use I.3Ex.4f and ex. I.3Ex.1 to express the polyhedron as the set of vectors  $\sum \lambda_i x_i^1 + \sum \mu_j x_j^2$  with  $\lambda_i \geq 0, \mu_j \geq 0, \sum \mu_j = 1$ .

h. The compact polyhedra are the convex hull of their finitely many extreme points, and the polyhedral cones are spanned by finitely many vectors [which can also be chosen extreme (rays) if the cone contains no straight line]. [cf. I.3Ex.4g. Remove from the  $x_j^2$  the non-extremal ones. Use I.3Ex.4a for the opposite direction.] In general, polyhedral cones can be decomposed as the sum of a subspace ( $P \cap (-P)$ ) and of a pointed polyhedral cone (the quotient) — which in turn has a unique decomposition in terms of extreme rays.

i. In  $\mathbb{R}^n$ , show that for any convex set  $B (\neq \emptyset)$  which is the set of solutions of a (finite or infinite) system of (weak or strict) linear inequalities, its asymptotic cone  $A_B$  (cf. 1.22 p. 8) is a closed convex cone, independent of  $b$ .

j.

i. The polyhedral cone in any decomposition as sub I.3Ex.4g of a polyhedron  $P$  is  $A_P$ .

ii. To make the decomposition unique, one can specify the compact polyhedron, when  $P$  contains no straight line, as the convex hull of the extreme points of  $P$ .

For  $x \in P$ , define the dual face  $F_x^0$  of  $x$  as  $\{ (a, b) \in P^0 \mid \langle a, x \rangle = b \}$ , and the face  $F_x$  of  $x$  as  $\{ y \in P \mid \langle a, y \rangle = b, \forall (a, b) \in F_x^0 \}$ .

k. Show that  $F_x^0$  is the convex cone spanned by the vectors  $(A_i, b_i)$  with  $A_i x = b_i$  if  $P = \{ x \mid Ax \geq b \}$ .

l. Show that  $F_x = \{ y \in P \mid x + \varepsilon(x - y) \in P \text{ for some } \varepsilon > 0 \}$ .

m. Show that  $\dim F_x + \dim F_x^0 = \dim E$ .

HINT. E.g., choose  $x$  as origin, and as basis vectors first  $e_1, \dots, e_n \in F_x$ , with  $n = \dim F_x$ , then  $f_1, \dots, f_k \in P$  with  $k+n = \dim(P)$ , finally  $g_1, \dots, g_\ell$ , with  $\ell = \dim E - \dim P$ . The inequalities of  $P$  in this basis must have the form  $P = \{ \sum_i u_i e_i + \sum_j v_j f_j + \sum w_h g_h \mid Au + Bv \leq 1, Cv \geq 0, w = 0 \}$  where  $C$  has full column rank — otherwise one would have  $(0, v, 0) \in F_x$  for some  $v \neq 0$ . Thus  $\dim \{ \mu C \mid \mu \geq 0 \} = k$ , so, by I.3Ex.4k,  $\dim F_x^0 = k + \ell$ . Alternatively, use induction on  $\dim F_x$ , starting with  $F_x = P$ .

n.  $P$  has finitely many dual faces (I.3Ex.4k), hence by definition finitely many faces. If one adds  $\emptyset$  as a face with  $P^0$  as dual face, there is a lattice structure (stability under intersections). (To preserve I.3Ex.4m count as dimension of the empty face 1 less than the dimension of the minimal non-empty faces).

o. The faces of  $P$  are the sets of minimisers of linear functionals. Alternatively, they are the intersections of  $P$  with affine subspaces, whose complement in  $P$  is convex.

p. The dual faces of  $P$  are the faces of  $P^0$ .

q. Think here for simplicity (in order not to deal with  $K$ -valued distances and Lipschitz constants) of  $K$  as being the reals. The distances used on vector spaces are assumed to derive from some norm.

If  $f$  is an affine map from a polyhedron  $P$  to a vector space, then  $f^{-1}$  is Lipschitz as a map from  $f(P)$  to closed (convex) subsets of  $P$ , endowed with the Hausdorff metric  $d(S_1, S_2) =$

$\max[\max_{x \in S_1} d(x, S_2), \max_{x \in S_2} d(x, S_1)]$ . (All norms on finite dimensional vector spaces are equivalent).

HINT. The statement is clearly equivalent to the Lipschitz character of the map  $f^{-1}$  from closed subsets of  $P$  to closed subsets of  $f^{-1}(P)$ , and this property is stable under composition. Assume thus the kernel of  $f$  is 1-dimensional:  $P$  is a polyhedron in  $\mathbb{R}^n \times \mathbb{R}$ , and  $f$  the projection to  $\mathbb{R}^n$ . Consider the  $\mathbb{R}$ -valued functions on  $f(P)$ :  $\bar{u}(x) = \sup\{y \mid (x, y) \in P\}$ ,  $\underline{u}(x) = \inf\{y \mid (x, y) \in P\}$ .  $\bar{u}$  is the minimum of a finite number (possibly zero) of linear functions (look at the inequalities determining  $P$ ), hence Lipschitz, and similarly  $\underline{u}$ .

COMMENT 3.6. I.3Ex.4c immediately implies that two disjoint polyhedra can be strictly separated. Similarly, I.3Ex.4a, I.3Ex.4b and I.3Ex.4c are still true if some of the half spaces in the definition of a polyhedron are allowed to be open — so two such generalised polyhedra  $P_1$  and  $P_2$  that are disjoint can be separated by a linear function  $f$  with  $f(P_1) \cap f(P_2) = \emptyset$  (applying ex. I.3Ex.1 to  $P_1 - P_2$  and  $\{0\}$ ). However ex. I.3Ex.1 yields both those separations directly, with the additional information that  $f$  can be selected to belong to the convex hull of both  $A_1$  and  $A_2$  ( $P_1 = \{x \mid A_1x \geq b_1\}$ ,  $P_2 = \{x \mid A_2x \leq b_2\}$ ).

COMMENT 3.7. The alternative argument of I.3Ex.4a, used for generalised polyhedra, yields a quantifier elimination algorithm for the logic having the linear inequalities as elementary sentences, plus quantifiers and propositional calculus (cf. remark after theorem 3.7 p. 27).

COMMENT 3.8. It also provides an elementary route to Farkas' lemma and the following exercises, without relying on Tarski's theorem. (At the end of ex. I.3Ex.1 take the image of  $Q$  under the quotient mapping by  $V$ , and apply I.3Ex.4a).

COMMENT 3.9. Walkup and Wets (1969) have shown that I.3Ex.4q characterises in fact the polyhedra among all closed convex subsets.

**5. Linear programming.** Consider the general linear program  $\underline{v} = \sup\{xc \mid xA \leq b, x \geq 0\}$  where  $b$  and  $x$  are row vectors,  $c$  a column vector, and  $A$  a matrix [equality constraints can be changed to a pair of inequality constraints, the sign of “ $\geq$ ” inequalities changed, and unrestricted variables  $x$  replaced by  $x^+ - x^-$  to obtain this form]. [If  $E = \emptyset : \sup E = -\infty, \inf E = +\infty$ ].

Consider also the “dual” program  $\bar{v} = \inf\{by \mid Ay \geq c, y \geq 0\}$  [note that from first principles  $\underline{v} \leq \bar{v}$ ] and the symmetric matrix game

$$M = \begin{pmatrix} 0 & -A & c \\ A^t & 0 & -b^t \\ -c^t & b & 0 \end{pmatrix}$$

Apply the strong complementarity property in  $M$  (ex. I.3Ex.2):

a. First to the last strategy, to deduce that either there is an optimal strategy  $(x, y, t)$  with  $t > 0$  in which case  $\bar{v} = \underline{v} \in \mathbb{R}$  and  $x/t$  and  $y/t$  solve  $\underline{v} = xc, xA \leq b, x \geq 0$  and  $\bar{v} = by, Ay \geq c, y \geq 0$  or  $\bar{v} = \underline{v} = +\infty$ , or  $\bar{v} = \underline{v} = -\infty$ , or  $\bar{v} = +\infty, \underline{v} = -\infty$ .

HINT. Note that  $\bar{v} = +\infty$  is equivalent to  $\sup\{xc \mid xA \leq 0, x \geq 0\} = +\infty$ .

b. Then to the other strategies, to obtain the strong complementarity relations for linear programs (Tucker, 1956).

c. Consider a polyhedral game (Wolfe, 1956), i.e. player I and II's strategy sets are polyhedra  $X = \{x \mid xC \leq c\}$  and  $Y = \{y \mid By \geq b\}$ , and the pay-off function is  $xAy$ . Let  $\bar{v}$  and  $\underline{v}$  denote the  $\inf \sup$  and the  $\sup \inf$  respectively. Show that either  $\bar{v} = \underline{v}$ , and then there exist optimal strategies if this value is finite, or  $\bar{v} = +\infty, \underline{v} = -\infty$ .

HINT. The case is clear if  $X$  or  $Y$  or both are empty. Assume thus not. For given strategy  $y$  of II, player I wants to solve the linear program  $\max x(Ay)$  sub  $xC \leq c$ , whose dual is (there is no sign restriction on  $x$ )  $\min cv$  sub  $Cv = Ay, v \geq 0$ , with the same value (cf. I.3Ex.5a) since  $X \neq \emptyset$ . Thus  $y$  guarantees this quantity to player II, and he wants to solve the program  $\min cv$  sub  $Cv = Ay, v \geq 0, By \geq b$  and play the corresponding  $y$ . Similarly player I wants to solve  $\max ub$  sub  $uB = xA, u \geq 0$ ,

$xC \leq c$ . Observe that those two programs are themselves dual [cf. I.3Ex.5d infra] and apply I.3Ex.5a. [Note also that  $\bar{v}$  and  $\underline{v}$  in the sense of the dual programs are always the same as those for the game.]

d. In general, the “dual” of a linear program [ $\underline{v} = \max(xc)$ , subject to a system of equality and inequality constraints] is the expression that  $\underline{v} = \min\{r \mid xc > r\}$  and the system is inconsistent }, where the condition is expressed by means of ex. I.3Ex.1 p. 28, assuming the system itself to be consistent. Verify that this yields indeed the duals used sub I.3Ex.5a and I.3Ex.5c above.

[It is this unverified assumption that causes the possibility of  $\bar{v} \neq \underline{v}$  — as in ( $\max x$  sub  $y \geq 1, -y \geq 1$ ) — while the dual in the sense of ex. I.3Ex.1 always gives a full diagnostic about the original system.]

e. In a polyhedral game, cf. I.3Ex.5c above, with finite value  $v$ ,

- (1) the sets of optimal strategies  $\tilde{X}$  and  $\tilde{Y}$  of players I and II are polyhedra
- (2) (Strong Complementarity) the face (ex. I.3Ex.4)  $\tilde{F}$  of  $X$  spanned by  $\tilde{X}$  (or by a point in the relative interior of  $\tilde{X}$ ) equals  $\{x \in X \mid xAy \geq v \forall y \in \tilde{Y}\}$  — and similarly for the face  $\tilde{G}$  of  $Y$  spanned by  $\tilde{Y}$ .

HINT. Use the dual programs of I.3Ex.5c above, together with ex. I.3Ex.4a for 1 and I.3Ex.5b above for 2.

f.  $\text{codim}_{\tilde{F}}(\tilde{X}) = \text{codim}_{\tilde{G}}(\tilde{Y})$  [same notations as in I.3Ex.5e].

HINT. 2 implies that, if strategy sets are restricted to  $\tilde{F}$  and  $\tilde{G}$  resp.,  $\tilde{X}$  and  $\tilde{Y}$  increase keeping the same dimensions. Assume thus  $\tilde{F} = X$ ,  $\tilde{G} = Y$ : in the affine spaces spanned by those sets, put the origin at some interior optimal strategies: we get a new polyhedral game of the form  $By \geq -1$ ,  $xC \leq 1$ , with pay-off =  $xAy + \alpha y + x\beta + \gamma$ ;  $(0, 0)$  being optimal implies  $\alpha = \beta = 0$ , so the pay-off is  $xAy$ . Thus  $\tilde{X} = \{x \in X \mid xA = 0\}$ ,  $\tilde{Y} = \{y \in Y \mid Ay = 0\}$ ,  $\text{codim}(\tilde{X}) = \text{codim}(\tilde{Y}) = \text{rank}(A)$ .

**6. von Neumann’s model of an expanding economy.** (Thompson, 1956)  $I$  is the set of production processes (activities),  $J$  the set of commodities. The  $I \times J$  matrices  $A$  and  $B$  describe by  $a_{ij}$  (resp.  $b_{ij}$ ) the input (resp. output) of commodity  $j$  corresponding to a unit intensity of process  $i$ . One wants a stationary growth path, i.e., a growth factor  $\lambda (\geq 0)$ , an interest factor  $r > 0$ , an intensity (row-)vector  $x \geq 0$  and a price (column-) vector  $p \geq 0$  such that

$$\begin{aligned} xB &\geq \lambda xA && \text{(the outputs of the current period suffice as inputs for the next period),} \\ (xB - \lambda xA)p &= 0 && \text{(goods in excess supply carry a zero price),} \\ x[Bp - rAp] &\geq \sup_{\tilde{x} \geq 0} \tilde{x}[Bp - rAp] && \text{(profit maximisation: } Bp - rAp \text{ gives the net profit from each activity at } p \text{ and } r) \end{aligned}$$

This last condition yields thus immediately  $Bp \leq rAp$  and  $xBp = rxAp$ . We assume that  $A \geq 0$ ,  $\sum_j a_{ij} > 0$  (no free lunch),  $v(B) > 0$  (there is some combination of activities by which every good can be produced). We require finally that the value of total input be non-zero — to avoid the completely degenerate solutions —:  $xAp > 0$ .

a. The above requirements imply  $r = \lambda$ , and are equivalent to:  $x(B - \lambda A) \geq 0$ ,  $(B - \lambda A)p \leq 0$ ,  $xAp > 0$ , i.e.:  $f(\lambda) = v(B - \lambda A) = 0$ , with as optimal strategies (after normalisation)  $x$  and  $p$ .

[Hence one could equivalently impose in addition (by ex. I.3Ex.2)  $x(B - \lambda A) + p^t > 0$ ,  $(\lambda A - B)p + x^t > 0$ .]

b. The assumptions imply  $f(\lambda)$  is decreasing,  $f(0) > 0 > f(+\infty)$  thus  $\{\lambda \mid f(\lambda) = 0\} = [\underline{\lambda}, \bar{\lambda}]$ ,  $0 < \underline{\lambda} \leq \bar{\lambda} < +\infty$ .

c. To show existence of a solution, show that a strongly complementary pair of optimal strategies (ex. I.3Ex.2) at  $\bar{\lambda}$  (or at  $\Delta$ ) yields  $xAp > 0$  (hence also  $xBp > 0$ ).

HINT. Order the strategies with those used with positive probability first. The matrix  $B - \lambda A$  takes then the form  $(\begin{smallmatrix} Q & T \\ R & S \end{smallmatrix})$ , where  $Q$  is identically zero in  $A$  if  $xAp = 0$ . The strong complementarity property yields  $x(\begin{smallmatrix} T \\ S \end{smallmatrix}) > 0$ : thus  $x$  still guarantees zero in  $B - (\lambda + \varepsilon)A$ .

d. There exists at most  $\min(\#I, \#J)$  solutions  $\lambda$ .

HINT. Consider 2 solutions  $\lambda^1 < \lambda^2$ . For some  $j = j_2$ , we have  $p_j^2 > 0$ ,  $xB_j = \lambda^2 xA_j > 0$ . Then  $xB_j > \lambda^1 xA_j$ , and since  $x$  is still optimal at  $\lambda^1$ , we have  $p_j^1 = 0$ : any  $j_1$  must be different.

e. Assume w.l.o.g. that the solutions  $(x^i, p^i)$  selected at  $\lambda^i$  satisfy the strong complementarity relations, and set, as in d),

$$J_0^2 = \{j \mid p_j^2 > 0, x^2 B_j > 0\}, J_1^2 = \{j \mid p_j^2 > 0, x^2 B_j = 0\}, J_2^2 = \{j \mid p_j^2 = 0\},$$

and similarly for  $J_0^1, J_1^1, J_2^1$  at  $\lambda^1$ , and sets  $I_0^1$ , etc. for player I (e.g.,  $I_0^1 = \{i \mid x_i^1 > 0, B_i p^1 > 0\}$ ).

i. Show that, on  $I_0^k \times J_0^k$ , every row and every column contain some  $a_{ij} > 0$ , and  $a_{ij} = 0$  elsewhere on  $(I_0^k \cup I_1^k) \times (J_0^k \cup J_1^k)$ ,  $k = 1, 2$ . Same conclusions for  $B$  if  $B \geq 0$ .

ii. Show that  $J_0^1 \cup J_1^1 \subseteq J_0^2 \cup J_1^2$  — and similarly  $I_0^2 \cup I_1^2 \subseteq I_0^1 \cup I_1^1$ .

HINT.  $p^1$  is optimal in  $B - \lambda^2 A$

iii. Show that  $J_0^2 \cup J_2^2 \subseteq J_2^1$  — and  $I_2^1 \cup I_0^1 \subseteq I_2^2$ .

HINT. cf. I.3Ex.6d for  $J_0^2$ , use strong complementarity for  $J_2^2$

iv. Deduce from the two last points that  $J_0^1 \cup J_1^1 \subseteq J_1^2, I_0^2 \cup I_1^2 \subseteq I_1^1$ .

v. Conclude from the above, putting the solutions  $\lambda^i$  in increasing order, and using an appropriate ordering on activities and commodities, that  $A$  has the following structure: if there are  $n$  roots  $\lambda^i$ , there exist  $n$  special blocs, any element to the left or below (or both) any special bloc being zero, while in the special blocs every row and every column contains a positive element. If  $B \geq 0$ , it has the same structure, with the same special blocs. The support of the  $i^{\text{th}}$  solution is the  $i^{\text{th}}$  special bloc, together with everything to the left and below.

## 7. Examples on Exercise I.3Ex.6.

a.  $A = (\begin{smallmatrix} 0 & 1 \\ 1 & 0 \end{smallmatrix}), B = (\begin{smallmatrix} 2 & 0 \\ 0 & 1 \end{smallmatrix})$  yields  $\lambda = \sqrt{2}$ : the problem needs a real closed field.

b.  $\#I = \#J = n, a_{ij} = b_{ij} = 0$  for  $i \neq j, a_{ii} = 1, b_{ii} = i$ : there are  $n$  solutions,  $\lambda = 1, 2, \dots, n$ : the upper bound of ex. I.3Ex.6d is attained.

c. The following examples (with 2 goods) show that there is no extension to a continuum of activities. In that case one represents simply the set of all feasible input-output vectors  $(x, y)$  by a closed convex cone in  $\mathbb{R}^2 \times \mathbb{R}^2$ .

(1) First Example.  $x_i \geq 0, y_i \geq 0, x_1(x_2 - y_2) \geq y_2^2, x_1 + x_2 \geq y_1$ .

(2) Second Example. Replace the last inequality in example 1 by the stricter inequality  $x_1(x_2 + y_1) \geq y_1^2$ .

Check that both examples describe convex (cf. below), closed cones  $C$  in  $(\mathbb{R}_+^2)^2$ , which are comprehensive  $[(x, y) \in C, x' \geq x, y' \leq y \Rightarrow (x', y') \in C]$ , contain an interior point of  $\mathbb{R}_+^4$ , say  $P_\varepsilon = (1, \varepsilon(1+\varepsilon); 1+\varepsilon, \varepsilon)$  (check for 2, after having verified that  $C_2 \subseteq C_1$ ), and offer no free lunch ( $x = 0 \Rightarrow y = 0$  — check for 1).

Their efficient frontiers are described, for example 1 by replacing the last two inequalities by equalities, and for example 2 by  $(0, 0, 0, 0)$  and  $\{x_1 = y_1 - y_2 > 0, x_1 x_2 = y_1 y_2, y_2 \geq 0\}$ , which is thus not closed (cf. also ex. I.4Ex.12 p. 42 and I.4Ex.13 p. 43 for more classical, but infinite dimensional examples of a non-closed efficient frontier) — for the closure, put a weak inequality, and add  $x_2 \geq 0$ . They describe a (one-to-one for example 2) map from efficient inputs  $(x_1 > 0, x_2 \geq 0$  or  $(0, 0)$  for example 2 to efficient outputs  $(y_1 > y_2 \geq 0$  or  $(0, 0)$  for example 2 —

in example 2,  $x_1(y) = y_1 - y_2 > 0$ ,  $x_2(y) = y_2 + y_2^2/x_1(y)$ :  $x_1(y)$  being linear, and  $y^2/x$  convex, we obtain indeed that for an average output, one needs less than the average input — hence the convexity of the comprehensive hull. Use a similar argument for the convexity in the first example.

Yet, for both cones  $C$ , there is no solution  $[\lambda, \mu, p, (x^*, y^*)] \in \mathbb{R} \times \mathbb{R} \times \mathbb{R}^2 \times C$  of  $\lambda y^* \geq \mu x^*$ ,  $\langle p, \lambda y - \mu x \rangle \leq 0$ ,  $\forall (x, y) \in C$ ,  $\langle p, y^* \rangle > 0$  (resp.  $\langle p, x^* \rangle > 0$ ) and  $(\lambda, \mu) \neq (0, 0)$  (resp.  $\mu = 0 \Rightarrow \lambda > 0$ ).

HINT.  $\lambda \leq 0$  or  $\mu \leq 0$  is trivial:  $\mu < 0 \Rightarrow p \leq 0$  (free disposal of outputs),  $\mu = 0 < \lambda \Rightarrow p \leq 0$  (free disposal and feasibility of strictly positive output),  $\mu > 0 \geq \lambda \Rightarrow y^* = x^* = 0$  (no free lunch), and  $\lambda < 0 = \mu \Rightarrow y^* = 0$ . Assume thus  $\lambda = 1$ ,  $\mu > 0$ . In example 1 (a fortiori example 2)  $y \geq x$  implies  $y_1 = x_1, y_2 = x_2 = 0$ . Thus  $\mu > 1$  implies  $y^* = x^* = 0$ . Also (free disposal)  $\langle p, y - \mu x \rangle \leq 0$  implies  $p \geq 0$ . And in example 2 (a fortiori example 1)  $\langle p, y - x \rangle \leq 0 \forall (x, y) \in C$  implies then  $p_1 = 0$  (use  $P_\varepsilon$ ). Since ( $P_\varepsilon$  again)  $y_2 - \mu x_2$  can be  $> 0$  for  $\mu < 1$ , it follows that  $p = 0$  for  $\mu < 1$ : in every case  $\langle p, y^* \rangle = 0$ .

**8. Bloc Decomposition.** (Kemeny, cf. Thompson, 1956) Consider the matrix game  $M = \begin{pmatrix} Q & P \\ Q & R \end{pmatrix}$ , in bloc decomposition, where the first bloc consists of zeros only.

a. If  $v(M) = 0$ , then player I has an optimal strategy carried by the first bloc iff  $v(P) \geq 0$ , and similarly for player II iff  $v(Q) \leq 0$ . Also  $v(P) \geq 0 \geq v(Q) \Rightarrow v(M) = 0$ .

b. Denote by  $\tilde{v}(R)$  the value of  $R$  when players I and II are both restricted to optimal strategies in  $Q$  and  $P$  respectively (compare ex. I.1Ex.6 p. 11).

Assume  $v(P) \geq 0 \geq v(Q)$ . Then player I has an optimal strategy in  $M$  that is not carried by the first bloc only iff  $v(Q) = 0$  and either  $v(P) > 0$  or  $\tilde{v}(R) \geq 0$ . (Use ex. I.3Ex.2 p. 29 for the sufficiency of the condition). In particular, both players can use strategies in their second bloc iff  $v(P) = v(Q) = \tilde{v}(R) = 0$ .

c. Conclude from I.3Ex.8a and I.3Ex.8b that every optimal strategy in  $M$  is carried by the first bloc iff  $v(P) > 0 > v(Q)$ .

## 9. Perron Frobenius theorem.

a. In ex. I.3Ex.6 p. 32, take for  $B$  a matrix  $P$  of transition probabilities ( $P_{ij} \geq 0; \sum_j P_{ij} = 1$ ) and  $A = I$ : one obtains the existence of an invariant probability distribution. [One could rescale such as to drop  $\sum_j P_{ij} = 1$ , cf. I.3Ex.9b.]

However, this does not require the strong complementarity, so a direct application of the minmax theorem yields more:

b. Assume  $S$  is compact, and,  $\forall s \in S$ ,  $P_s$  is a non-negative regular measure on  $S$  with  $P_s(f) = \int f(x)P_s(dx)$  continuous whenever  $f$  is continuous. Then there exists a probability measure  $\mu$  on  $S$  with  $\int P_s(A)\mu(ds) = \int_A P_s(1)\mu(ds)$ , for all Borel  $A$ .

HINT. Apply prop. 1.17 p. 7 to the game where  $T$  is the space of continuous functions on  $S$ , and the pay-off is  $P_s(t) - t(s)P_s(1)$ ; consider the minimiser  $s_0$  of  $t$ .

## 10.

a. *Carathéodory.* Any point in the convex hull of a subset  $S$  of  $\mathbb{R}^n$  is a convex combination of  $n+1$  points of  $S$ .

HINT. For  $k = 1, \dots, n$ , let  $\sum_{i=1}^{n+m} \lambda_i x_k^i = a_k$ ,  $\sum_{i=1}^{n+m} \lambda_i = 1$  be the equations: if  $m > 1$ , there are more equations than unknowns, so the set of solutions  $\lambda$  meets the boundary  $\min_i \lambda_i = 0$  of the simplex. Proceed by induction.

b. *Fenchel.* Assume  $S$  has at most  $n$  components, then  $n$  points suffice.

HINT. Assume  $p$  is interior to the simplex  $\{x_1, \dots, x_{n+1}\}$ ,  $x_i$  in  $S$ . Consider the reflexion of this simplex in  $p$  (i.e.,  $y_i = 2p - x_i$ ) and the cones  $C_i$  based at  $p$  with vertices  $(y_1, \dots, y_{i-1}, y_{i+1}, y_{n+1})$ .  $\bigcup_1^{n+1} C_i$  contains  $S$  and one point of  $S$  belongs to the boundary of some  $C_i$ .

c. (cf. ex. I.3Ex.14 p. 37). Carathéodory extends to continuous averages: For a convex subset  $C$  of  $\mathbb{R}^n$ , the expectation of any integrable  $C$ -valued random variable belongs to  $C$ .

REMARK 3.10. Since  $C$  is not necessarily measurable for the distribution  $\mu$  of  $X$ , the above statement is preferable to one that involves the barycentre of measures on  $C$ .

HINT. Assume the expectation  $b \notin C$ , and separate: for a linear functional  $\varphi$  we have  $\varphi(b) \geq \alpha = \sup_{x \in C} \varphi(x)$ . But  $\varphi(b) = \mathbb{E}(\varphi(X))$ , so a.s.  $\varphi(X) = \alpha = \varphi(b)$ :  $X$  is carried by the convex set  $C_1 = \{x \in C \mid \varphi(x) = \varphi(b)\}$ . Proceed by induction on the dimension. More precisely the proof shows that

- the expectation belongs to the relative interior of the convex hull of the support of the distribution,
- this relative interior is included in  $C$ .

d. The convex hull of a compact subset of  $\mathbb{R}^n$  is compact (use I.3Ex.10), just as is in general the convex hull of finitely many compact convex subsets.

**11. Linear programming and polyhedral games—continued.** Consider the program: maximise  $xb$  sub  $xA \leq c$ ,  $x \geq 0$ , and its dual: minimise  $cy$  sub  $Ay \geq b$ ,  $y \geq 0$ . Assume  $\bar{v} \in \mathbb{R}$ .

a. The solution sets  $X$  and  $Y$  are polyhedra, and have an extreme point.

HINT. Ex. I.3Ex.5a p. 31 yields  $X \neq \emptyset$ ,  $Y \neq \emptyset$ . Ex. I.3Ex.4c p. 29 that they are polyhedra. The constraints  $x \geq 0, y \geq 0$  ensure the existence of an extreme point (ex. I.3Ex.4j p. 30).

b. The pairs of extreme points  $(x, y)$  of  $X$  and  $Y$  are the feasible points such that, for some subsets  $I$  of the rows and  $J$  of the columns, one has  $A^{IJ}$  non-singular and  $x^I = c^J(A^{IJ})^{-1}$ ,  $y^J = (A^{IJ})^{-1}b^I$ ,  $x_i = 0$  for  $i \notin I$ ,  $y_j = 0$  for  $j \notin J$  (denoting by  $x^I, y^J, A^{IJ}$ , etc … the restrictions to the corresponding sets of indices). And then  $v = c^J(A^{IJ})^{-1}b^I$ . (The empty  $\times$  empty matrix is non-singular by convention, and then  $v = 0$  as an empty sum).

HINT. Optimality for such a pair  $(x, y)$  follows from  $xb = cy$ . If  $x = \frac{1}{2}(x^1 + x^2)$ ,  $x^k$  feasible ( $k = 1, 2$ ), then  $x_i^k = 0$  ( $i \notin I$ ) and  $x^{k,I}A^{IJ} = c^J$ , hence  $x^k = x$  by the independence of the rows of  $A^{IJ}$ . Similarly  $y$  is extreme. Conversely, for an extreme pair  $(x, y)$ , let  $I_1 = \{i \mid x_i > 0\}$ ,  $J_1 = \{j \mid y_j > 0\}$ ,  $I_2 = \{i \mid A_{ij}y_j = b_i\}$ ,  $J_2 = \{j \mid x_Aj = c_j\}$ :  $I_1 \subseteq I_2$  and  $J_1 \subseteq J_2$ . Further the rows of  $A^{I_1 J_2}$  and the columns of  $A^{I_2 J_1}$  are linearly independent by the extremality of  $x$  and  $y$  resp., cf. supra. Extend thus  $I_1$  (resp.  $J_1$ ) to a basis  $I$  (resp.  $J$ ) of the rows (resp. columns) of  $A^{I_2 J_2}$ . Conclude that  $A^{IJ}$  is itself non-singular.

c. Particularise I.3Ex.11b to a characterisation of the extreme points of a polyhedron  $P = \{x \mid xA \leq b\}$ .

HINT. Rather than first trying to find a system of coordinates — using the extreme points — where  $P \subseteq \mathbb{R}_+^n$ , a more direct approach uses ex. I.3Ex.4m p. 30 and ex. I.3Ex.10 p. 34.

d. (Shapley and Snow, 1950) If  $A$  is a matrix game with value  $v$ , a pair of optimal strategies is an extreme pair iff, for some subsets  $I$  and  $J$  of rows and columns, one has, writing  $\tilde{B}$  for the adjoint of  $B = A^{IJ}$ , and  $1^I, 1^J$ , for appropriate vectors of ones, that  $1^J \tilde{B} 1^I \neq 0$ , and  $x^I = \frac{1^J \tilde{B}}{1^J \tilde{B} 1^I}$ ,  $y^J = \frac{\tilde{B} 1^I}{1^J \tilde{B} 1^I}$ . And then  $v = \det(B)/(1^J \tilde{B} 1^I)$ .

HINT. Assume first  $v > 0$ . Then  $\xi = x/v$  and  $\eta = y/v$  are the solutions of the dual programs minimise  $\xi \cdot 1$  sub  $\xi A \geq 1$ ,  $\xi \geq 0$ , and maximise  $1 \cdot \eta$  sub  $A\eta \leq 1$ ,  $\eta \geq 0$ . Apply I.3Ex.11b to those, and that  $M^{-1} = [\tilde{M}/\det(M)]$ . For other  $v$ , reduce to  $v > 0$  using that adding a constant  $c$  to the entries of  $A$  does not change the solutions  $x$  and  $y$ , and adds  $c$  to  $v$ ; and that, if  $M'$  is obtained by adding  $c$  to all entries of  $M$ , then  $1 \cdot \tilde{M}' = 1 \cdot \tilde{M}$ ,  $\det(M') = \det(M) + c[1 \cdot \tilde{M} \cdot 1]$ . Those equations follow from the previous case for completely mixed matrices  $M$  with  $v > 0$ : indeed, the determinant is clearly an affine function of  $c$ , i) either geometrically, as a volume of a parallelepiped, this volume being clearly bounded by  $K(1 + |c|)$ , and the determinant being polynomial; ii) or analytically, subtracting the first column from the others and expanding. Since  $v = v_0 + c$ , the equation for  $v$  yields then that  $1 \tilde{M} 1$  must be a polynomial of degree 0 in  $c$ , hence constant. Therefore the constancy of  $x$  and  $y$  yields that of  $1 \tilde{M}$  and

of  $\tilde{M}1$ . Those equations for the adjoints and determinants being proved on an open subset of matrices, and being polynomial, hold everywhere.

COMMENT 3.11. The procedure used in the first sentence of the hint is the efficient way to solve a matrix game by linear programming.

e.

i.  $\{c \mid \underline{v}(b, c) > -\infty\} = \{c \mid \exists x \geq 0, xA \leq c\}$  is a polyhedral cone  $\underline{P}$ , independent of  $b$  — cf. ex. I.3Ex.9a p. 34. Similarly  $\overline{P} = \{b \mid \bar{v}(b, c) < +\infty\}$ .

ii.

A.  $\underline{v}(b, c)$  and  $\bar{v}(b, c)$  are convex in  $b$ , concave in  $c$ , and non-decreasing.

B.  $v(b, c)$  is positively homogeneous of degree one (with  $0 \cdot \infty = 0$ ) in  $c$  for  $b \in \overline{P}$  and in  $b$  for  $c \in \underline{P}$ .

C. on  $\bar{P} \times \underline{P}, v \in \mathbb{R}$ ; on  $\overline{P} \times \underline{P}^c, v = -\infty$ ; on  $\overline{P}^c \times \underline{P}, v = +\infty$ ; on  $\overline{P}^c \times \underline{P}^c, \bar{v} = +\infty$  and  $\bar{v} = -\infty$ .

iii. on  $\overline{P} \times \underline{P}$ ,  $v$  is piecewise bi-linear in the sense that there exist subdivisions of  $\overline{P}$  and  $\underline{P}$  into finitely many simplicial cones, such that the restriction of  $v$  is bi-linear on any product of two such cones.

HINT. By I.3Ex.11b, for fixed pair  $(I, J)$ ,  $v$  is bi-linear. I.3Ex.11b is applicable because of I.3Ex.11a. This formula is applicable as long as the corresponding vectors  $x$  and  $y$  are feasible, which translates into a finite system of homogeneous linear inequalities in  $b$  and another one in  $c$ . The union over all possible  $(I, J)$  of those lists of linear functionals cuts up  $\overline{P}$  into finitely many polyhedral cones, add the coordinate mappings to those lists to be sure each of those cones is pointed. A further subdivision (e.g. barycentric) yields then simplicial cones. Similarly for  $\underline{P}$ .

iv. For some  $K > 0$ ,  $v$  is, on  $\overline{P} \times \underline{P}$ , Lipschitz in  $c$  with constant  $K \|b\|$  and in  $b$  with  $K \|c\|$ . (Use I.3Ex.11eiii).

f. The results of I.3Ex.11e remain verbatim true for the  $\inf \sup \bar{v}$  and the  $\sup \inf \underline{v}$  of the polyhedral games of ex. I.3Ex.5c p. 31.

HINT. Rewrite the primal-dual programs of the hint of ex. I.3Ex.5c in the form used in the present exercise, replacing  $y$  and  $x$  by differences of non-negative vectors  $y_1 - y_2$  and  $x_1 - x_2$ , and equations by pairs of opposite inequalities.

g. Those results still remain true, except the monotonicity, if the strategy polyhedra of the polyhedral games have the most general presentation  $X = \{x \mid xC_1 \leq c_1, xC_2 \geq c_2, xC_3 = c_3\}$  and similarly for  $Y$ .

HINT. Rewriting the constraints in the standard way yields  $x\bar{C} \geq \bar{c}$ , where  $\bar{c} = (c_1, -c_2, c_3, -c_3)$  is a linear function of the vector  $c = (c_1, c_2, c_3)$ .

h. Allow now further the pay-off function of the polyhedral game to be bi-affine rather than bi-linear, i.e. of the form  $(1, x) A (1, y)$ . Then  $\bar{v}$  is piecewise bi-linear and concave convex in  $(A_{0,.}, c)$  and  $(A_{.,0}, b)$  respectively.

HINT. Introduce  $u_j$  and  $v_i$  as additional parameters of the strategies of I and II, fixed by constraints to the values  $A_{0,j}$  and  $A_{i,0}$  respectively: the pay-off becomes  $u_0 v_0 + \sum_{j \geq 1} u_j y_j + \sum_{i \geq 1} v_i x_i + \sum_{i,j \geq 1} x_i A_{ij} y_j$  and apply g).

COMMENT 3.12. In the case of polyhedral games, the solution set may have no extreme points. Their analogue would be minimal faces, i.e. faces which are an affine subspace.

COMMENT 3.13. This case includes the general linear programming model (without restrictions), by taking a one-point polyhedron for one of the players.

**12.** To illustrate the need for the internal point conditions in the separation theorems (cf. 1.19 above), even without continuity requirements and with the best other conditions, consider the following example:

a. Denote by  $E = \mathbb{R}[x]$  the vector space of polynomials in  $x$  with real coefficients, ordered by the positive polynomials being those whose coefficient of highest degree is positive. The positive cone  $P$  satisfies  $P \cup (-P) = E$ ,  $P \cap (-P) = \{0\}$  — i.e. the order is total, yet, any non-negative linear function is zero.

b. Deduce that, on any ordered field extension of the reals, there is no positive (i.e; non-negative non-null) linear functional.

c. Show more generally, in part by the above argument, that this still holds for any ordered field extension  $\overline{K}$  of an ordered field  $K$  (viewed as an ordered  $K$ -vector space).

**13.** (Weyl, 1950) Give a direct proof of theorem 3.8 p. 27.

#### 14. Jensen's inequality.

a. Let  $f$  be a convex function from  $\mathbb{R}^n$  to  $\mathbb{R} \cup \{+\infty\}$ , and  $X$  an integrable random variable on  $(\Omega, \mathcal{A}, P)$  with values in  $\mathbb{R}^n$ . Then  $\int_* f(X(\omega))P(d\omega) \geq f(E(X))$ . (Recall that  $\int_*$  denotes lower integral, cf. 2.b p. 14.)

HINT. Reduce first to the case where  $E(X) = 0$  — and where  $X$  is not carried by any strict subspace. In  $\mathbb{R}^n \times \mathbb{R}$ , consider then the convex set  $\overline{C} = \{(x, t) \mid t \geq f(x)\}$  and any point  $(0, \alpha)$  with  $\alpha < f(0)$ : both are disjoint, there exists thus a linear functional  $\varphi$  and  $\lambda \in \mathbb{R}$  such that  $\varphi(x) + \lambda t \geq \lambda\alpha$  for  $(x, t) \in \overline{C}$  (if  $\overline{C} = \emptyset$  there is nothing to prove), and  $\|\varphi\| + |\lambda| > 0$ . Note that  $\lambda \geq 0$ , and if  $\lambda > 0$  then  $f(x) \geq \alpha - \frac{1}{\lambda}\varphi(x)$  yields the desired inequality. So one can assume  $\lambda = 0$ , i.e.  $f(x) < \infty \Rightarrow \varphi(x) \geq 0$ . But since  $X$  is not carried by any strict subspace,  $E(X) = 0$  yields that  $P\{\omega \mid \varphi(X(\omega)) < 0\} > 0$ , hence  $f(x) = +\infty$  on a set of positive (inner) measure: the inequality will also hold, as soon as one obtains an integrable bound for  $f^-(x)$ : for this, reduce again first to the case where zero is in the interior of  $\{x \mid f(x) < +\infty\}$ , and separate as above.

COMMENT 3.14. For  $C$  convex, let  $f(x) = 0$  for  $x \in C$ ,  $= +\infty$  otherwise, and obtain thus ex. I.3Ex.10c p. 35. The same shows there would be no gain in allowing  $f$  to be defined only on a convex subset of  $\mathbb{R}^n$ .

b. *Conditional versions.*

COMMENT 3.15. For conditional versions, we need somewhat stronger measurability requirements — e.g., even in the framework of ex. I.3Ex.10c, we do not know whether, for an integrable random variable  $X$  with values in a convex set  $C$  and  $\mathcal{B} \subseteq \mathcal{A}$ ,  $E(X \mid \mathcal{B}) \in C$  with inner (or even just outer) probability one. Thus:

i. Assume  $\mathcal{B} \subseteq \mathcal{A}$ ,  $X$  is an integrable random variable on  $(\Omega, \mathcal{A}, P)$  with values in  $\mathbb{R}^n$ , and  $g(\omega, x)$  maps  $\Omega \times \mathbb{R}^n$  in  $\mathbb{R} \cup \{+\infty\}$ , is convex in  $x$  for each  $\omega$ , and is  $\mu$ -measurable,  $\mu$  being the distribution of the random variable  $(\omega, X(\omega))$  on  $(\Omega, \mathcal{B}) \otimes (\mathbb{R}^n)$ , Borel sets). Assume also  $g^-(\omega, X(\omega))$  is integrable. Then  $E[g(\cdot, X(\cdot)) \mid \mathcal{B}](\omega) \geq g[\omega, E(X \mid \mathcal{B})(\omega)]$  a.s.

HINT. Show first that  $g(\omega, X(\omega))$  is measurable, so both conditional expectations are well defined. Let then (as in ex. II.1Ex.9 p. 60)  $\mu(dx \mid \omega)$  be a regular version of the conditional distribution of  $X$  given  $\mathcal{B}$ , chosen such that  $\int \|x\| \mu(dx \mid \omega) < +\infty$  everywhere. Show that the measurability conditions on  $g$  assure that  $E(g(\cdot, X(\cdot)) \mid \mathcal{B})(\omega) = \int g(\omega, x) \mu(dx \mid \omega)$  a.e. — and that the right-hand side is everywhere well defined. By (A),  $\int g(\omega, x) \mu(dx \mid \omega) \geq g(\omega, y(\omega))$ , with  $y(\omega) = \int x \mu(dx \mid \omega) (= E(X \mid \mathcal{B})(\omega)$  a.e.).

ii. The measurability requirement on  $g$  can always be relaxed without changing the conclusion by adding to  $\mathcal{B}$  all negligible subsets of  $(\Omega, \mathcal{A}, P)$ . Show however that one cannot add, on  $\Omega \times \mathbb{R}^n$ , all negligible subsets of the distribution of  $(\omega, X(\omega))$  in  $(\Omega, \mathcal{A}) \otimes (\mathbb{R}^n)$ , Borel sets). (The “argument” “just replace  $\Omega$  by  $\Omega \times \mathbb{R}^n$  with that distribution” does not work.)

HINT. Consider some  $g$  taking only the values 0 and  $+\infty$ .

iii. Show that the convexity assumption on  $g$  sub I.3Ex.14bi can be weakened to:  $g(\omega, x)$  has a convex restriction to some convex set  $C_\omega$ , where, for some (and then for any) regular version  $\mu(dx \mid \omega)$  of the conditional distribution of  $X$  given  $\mathcal{B}$ , a.e.  $C_\omega$  has  $\mu(dx \mid \omega)$ -outer measure 1.

HINT. In the proof of I.3Ex.14bi, you must have shown that, a.e.  $g(\omega, x)$  is  $\mu(dx \mid \omega)$ -measurable, so  $\int g(\omega, x)\mu(dx \mid \omega) = \int_{C_\omega} g(\omega, x)\mu^*(dx)$ , with  $\mu^*(dx)$  being the trace of  $\mu(dx \mid \omega)$  on  $C_\omega$ ,  $(\mu^*(B \cap C_\omega) = \mu(B \mid \omega))$ . Apply thus (A) with  $(C_\omega, \mu^*)$  as probability space, modifying  $g$  to  $+\infty$  outside  $C_\omega$ .

COMMENT 3.16. Versions like I.3Ex.14bi above will be needed in(ch.IX, sect. 2. A version with varying  $C_\omega$  (like in I.3Ex.14biii) will be used in lemma 2.16 p. 283. For versions with values in infinite dimensional spaces (e.g. lemma 2.4 p. 275), one has to impose lower-semi-continuity of the convex function, to use the Hahn-Banach theorem. Corresponding conditional versions (like in the use of lemma 2.4 in lemma 3.4 p. 284) can often be obtained by the same technique as above (regular conditionals) — e.g. in the present case, because a (Bochner) integrable random variable with values in a Banach space has (a.e.) its values in a separable subspace, which is polish. (Other techniques involve using the same proof as in the unconditional case — e.g. when the convex function is already the upper bound of a countable family of continuous linear functionals —, or conceivably relying on versions of the martingale convergence theorem, approximating  $\mathcal{B}$  by an increasing sequence of countable partitions  $\mathcal{B}_n$ , e.g. such that  $E(g(X) \mid \mathcal{B})$  is bounded on each partition element, and using the lower semi-continuity in the right-hand side).

### 15. Real valued convex functions.

a. A real valued convex function on a polyhedron is u.s.c.

b. On the space of continuous, real valued convex functions on a compact polyhedron, the topologies of point-wise convergence and of uniform convergence coincide.

HINT. Prove first, using the idea of I.3Ex.15a — plus compactness — that if  $f_\alpha \rightarrow f$  point-wise, then  $(f_\alpha - f)^+$  converges uniformly to zero. So one can assume  $f_\alpha \leq f$ . Next, by Dini's theorem,  $g_\beta \rightarrow f$  uniformly, where  $g_\beta$  denotes the increasing filtering family of all maxima of finitely many affine functions which are strictly smaller than  $f$ . Thus we can assume that  $f > 0$ , and it suffices to show that  $f_\alpha \geq 0$  for  $\alpha \geq \alpha_0$ . Let  $f_\alpha(x_\alpha) \leq 0$ . Then with  $b$  an interior point, and  $y_\alpha = (1 - \varepsilon)x_\alpha + \varepsilon b$ , we have  $f_\alpha(y_\alpha) \leq \varepsilon f(b)$ . Choose  $\varepsilon > 0$  such that  $f(x) > \varepsilon f(b) \forall x$ , and subtract  $\varepsilon f(b)$  from  $f$ : we have  $f > 0$ ,  $f_\alpha(y_\alpha) \leq 0$ , and  $y_\alpha$  is bounded away from the boundary. Thus, for  $M = \max_x f(x)$ , we can set  $\varphi_\alpha$  the convexification of the function which is zero at  $y_\alpha$  and  $M$  elsewhere:  $\varphi_\alpha \geq f_\alpha$  so  $\liminf \varphi_\alpha \geq f > 0$ , and the  $\varphi_\alpha$  are Lipschitz, since  $y_\alpha$  is bounded away from the boundary. Conclude.

c. Define the “lower topology” on the space of convex  $\overline{\mathbb{R}}$ -valued functions on a convex set  $C$  by the neighbourhood system  $V_\varepsilon(f) = \{g \mid \varepsilon + g \geq f\}$ . Show this defines a topology. Prove that, if  $C$  is a compact, convex polyhedron and if a net of convex functions  $f_\alpha$  converges point-wise to an u.s.c. function  $f$ , then  $f_\alpha \rightarrow f$  in the lower topology.

HINT. We claim that  $\varepsilon_\alpha = \max_{p \in C}[f(p) - f_\alpha(p)]$  converges to zero. We can prove this by induction on the dimension, and hence assume that the maximum on the boundary converges to zero. Hence, if the claim were not true we would have at least for some subnet,  $p_\alpha$  in the interior, and  $f_\alpha(p_\alpha) \leq f(p_\alpha) - 3\varepsilon$  for some  $\varepsilon > 0$  (and hence  $f$  is finite everywhere — hence bounded). Extracting a further subnet, we can assume  $p_\alpha \rightarrow p_\infty$ . Denote by  $\bar{f}$  the (convex) extension by continuity of  $f$  from the interior to the boundary [i.e.,  $\bar{f}(p_\infty) = \lim_{\alpha \rightarrow \infty} f(p_\alpha)$ ]. Let  $g$  be the (continuous) convexification of the function with value  $\bar{f}(p_\infty) - 2\varepsilon$  at  $p_\infty$ , and  $\bar{f}$  elsewhere. Denote by  $q_\infty$  some interior point with  $g(q_\infty) < f(q_\infty) - \varepsilon$ : we have  $q_\infty = tq + (1 - t)p_\infty$ , with  $0 < t < 1$ ,  $q$  interior, and  $tf(q) + (1 - t)[\bar{f}(p_\infty) - 2\varepsilon] < f(q_\infty) - \varepsilon$ . Let  $q_\alpha = tq + (1 - t)p_\alpha : q_\alpha \rightarrow q_\infty$ , by convexity  $f_\alpha(q_\alpha) \leq tf_\alpha(q) + (1 - t)f_\alpha(p_\alpha) \leq tf_\alpha(q) + (1 - t)f(p_\alpha) - 3(1 - t)\varepsilon$ , and the right-hand member converges to  $tf(q) + (1 - t)\bar{f}(p_\infty) - 3(1 - t)\varepsilon$ , hence for  $\alpha$  sufficiently large  $f_\alpha(q_\alpha) \leq tf(q) + (1 - t)[\bar{f}(p_\infty) - 2\varepsilon] < f(q_\infty) - \varepsilon$ . But the sequence  $q_\alpha \rightarrow q_\infty$  being compact in the interior, is included in some compact, convex polyhedron  $P$  in the interior of  $C$ . On  $P$ , the functions  $f_\alpha$  and  $f$  are continuous, hence  $f_\alpha \rightarrow f$  implies that  $f_\alpha \rightarrow f$  uniformly on  $P$  (by I.3Ex.15b), contradicting our inequality which implies that, for  $\alpha$  sufficiently large,  $f_\alpha(q_\alpha) < f(q_\alpha) - \varepsilon$ .

#### 4. Equilibrium points

**THEOREM 4.1. (Nash, 1950)(Glicksberg, 1952)** Let  $S^1, \dots, S^n$  be compact topological spaces,  $F^1, \dots, F^n$  continuous real valued functions on  $\prod_{i=1}^n S^i$ . Then, using regular probability measures for mixed strategies, there exists an equilibrium point.

The same result holds true in pure strategies if the sets  $S^i$  are compact convex subsets of locally convex linear spaces and if, for every  $i$  and every real  $\alpha$ , the sets  $A^\alpha(s^j)_{j \neq i} = \{s^i \mid F^i(s^1, \dots, s^n) \geq \alpha\}$  are convex.

**PROOF.** There is no loss in assuming the  $S^i$  to be Hausdorff: first map  $S^i$  to its compact Hausdorff image  $\overline{S}^i$  in the space of continuous  $\mathbb{R}^n$ -valued functions on  $\prod_{j \neq i} S^j$  do the proof on the  $\overline{S}^i$ , and finally extend the resulting  $\overline{\sigma}^i$  on  $\overline{S}^i$  to regular probability measures  $\sigma^i$  on  $S^i$  — such extensions always exist. Let  $\Sigma^i$  be the (convex) set of all regular probabilities on  $S^i$ , endowed with the (compact) topology of weak convergence. Let  $\Phi^i$  be the real valued function defined on  $\Sigma = \prod_i \Sigma^i$  by  $\Phi^i(\sigma^1, \dots, \sigma^n) = \int_{\prod_j S^j} F^i d\prod_j \sigma^j$ . Obviously,  $\Phi^i$  is multilinear. It is continuous by 1.15 p. 7. So  $\Sigma^1, \dots, \Sigma^n$  and  $\Phi^1, \dots, \Phi^n$  satisfy the conditions of the second part of the theorem. Thus there remains only to prove it.

Let  $T$  be the **best reply correspondence**:

$$\begin{aligned} T(s^1, \dots, s^n) &= \prod_{i=1}^n \{ \tilde{s} \in S^i \mid F^i(s^1, \dots, s^{i-1}, \tilde{s}, s^{i+1}, \dots, s^n) \\ &= \max_{s \in S^i} F^i(s^1, \dots, s^{i-1}, s, s^{i+1}, \dots, s^n) \}. \end{aligned}$$

We know that  $T(s^1, \dots, s^n)$  is a non-empty, compact and convex subset of  $S = \prod_i S^i$ . Further the continuity of the  $F^i$  ensures that  $T$  has a closed graph (i.e.,  $G = \{(s, t) \in S^2 \mid t \in T(s)\}$  is closed in  $S^2$ ). By Fan's fixed point theorem (Fan, 1952) (cf. also Glicksberg (1952), and ex.I.4Ex.17 p. 47), the mapping  $T$  has a fixed point. This finishes the proof. ■

#### Exercises.

##### 1. Symmetric equilibria.

a. (Nash, 1951) In the finite case, prove the existence of an equilibrium invariant under all symmetries of the game, i.e., permutations of the player set accompanied by one-to-one mappings between the pure strategy sets of corresponding players, leaving the pay-off functions (or merely the best reply correspondence) invariant.

HINT. Let  $X$  be the set of symmetric  $n$  – tuples of mixed strategies (i.e., invariant under all above symmetries of the game). Prove that  $X$  is convex, compact and non-empty. Show that the set of best replies to a point in  $X$  intersects  $X$  (consider its barycentre).

b. Under the conditions of theorem 4.1 p. 39, assume that  $S^i = S$  and  $F^i(s^1, \dots, s^n) = F^{\theta(i)}(s^{\theta(1)}, \dots, s^{\theta(n)})$  for all  $i$  and all  $\theta$  permutation on  $\{1, \dots, n\}$ . Prove that there exists a symmetric equilibrium. (Use Fan's theorem on  $\Phi(\tau) = \{\sigma \mid F^1(\sigma, \tau, \dots, \tau) \geq F^1(\sigma', \tau, \dots, \tau), \forall \sigma'\}$ ).

2. Under the assumptions of theorem 4.1, show that, if for some player  $i$ ,  $F^i$  is “strictly concave-like”, i.e. any convex combination of two distinct pure strategies of him is strictly dominated by one of his pure strategies, then he uses a pure strategy in any equilibrium point.

3. Let  $S$  and  $T$  be two compact topological spaces,  $f(s, t)$  a continuous real valued function on  $S \times T$ . Deduce the minmax theorem for  $f(s, t)$  from theorem 4.1.

4. The manifold of Nash equilibria. (Kohlberg and Mertens, 1986)

a. We assume finite strategy sets  $S^i$ , and keep the notations of theorem 4.1. Let  $\Gamma^i = \mathbb{R}^{S^i}$  (space of pay-off functions of  $i$ ), then  $\Gamma = \prod_i \Gamma^i$  is the space of games. Let also  $E = \{(G, \sigma) \in \Gamma \times \Sigma \mid \sigma \text{ is an equilibrium of } G\}$  (the equilibrium graph). Denote by  $\bar{\Gamma}$  the one-point compactification of a locally compact space  $\Gamma$ , and by  $\bar{p}: \bar{E} \rightarrow \bar{\Gamma}$  the continuous extension of the projection  $p: E \rightarrow \Gamma$  with  $\bar{p}(\infty) = \infty$ . Prove that  $\bar{p}$  is homotopic to a homeomorphism (under a homotopy mapping  $\infty$  to  $\infty$  and  $E$  to  $\Gamma$ ).

HINT. Let  $T^i = \prod_{h \neq i} S^h$ ; reparameterise  $\Gamma^i$ , the set of all  $S^i \times T^i$ -pay-off matrices  $G_{s,t}^i$  by  $G_{st}^i = \tilde{G}_{s,t}^i + g_s^i$  with  $\sum_t \tilde{G}_{st}^i = 0$ ; letting  $z_s^i = \sigma_s^i + \sum_{t \in T^i} G_{s,t}^i \prod_{j \neq i} \sigma_{t,j}^j$ ,  $(\tilde{G}, z)$  can be viewed as belonging to  $\Gamma$ . Show that  $(G, \sigma) \rightarrow (\tilde{G}, z)$  is a homeomorphism  $\phi$  from  $E$  to  $\Gamma$  (let  $v^i = \min\{\alpha \mid \sum_{s \in S^i} (z_s^i - \alpha)^+ \leq 1\}$ , then  $\sigma_s^i = (z_s^i - v^i)^+$  and  $g_s^i = z_s^i - \sigma_s^i - \sum_{t \in T^i} \tilde{G}_{s,t}^i \prod_{j \neq i} \sigma_{t,j}^j$ ). Then  $p_t(G, \sigma) = (\tilde{G}, tz + (1-t)g)$ , and  $p_t(\infty) = \infty$ , is the desired homotopy; for the continuity at  $\infty$ , show that  $\|z - g\| \leq \|\tilde{G}\| + 1$  in the maximum norm.

b. Deduce that the number of equilibria is generically finite and odd.

## 5.

a. Assume the sets  $S^i$  are finite, and the functions  $F^i$  have values in an ordered field  $K$ . If  $N > 2$ , assume  $K$  is real closed. Show that theorem 4.1 p. 39 remains valid, the  $\mu^i$  being a vector  $(p_1^i, \dots, p_{\#S^i}^i) \in K^{\#S^i}$  such that  $p_s^i \geq 0$  and  $\sum_{s \in S^i} p_s^i = 1$ .

HINT. cf. sect. 3 p. 26.

COMMENT 4.1. This implies the result of sect. 3 (cf. ex. I.4Ex.3).

b. In the following three-person game with integral pay-offs (I chooses the rows, II the columns, III the matrix):

$$\begin{pmatrix} 0,0,0 & 2,1,2 \\ 1,2,2 & 0,0,0 \end{pmatrix} \quad \begin{pmatrix} 2,0,1 & 1,1,0 \\ 0,1,1 & 0,0,0 \end{pmatrix}$$

there is a unique equilibrium with irrational pay-off.

HINT. Consider the golden number as probability of the first strategy, and as pay-off.

c. The symmetric three-person game

$$\begin{pmatrix} 0,0,0 & 2,1,2 \\ 1,2,2 & 0,0,0 \end{pmatrix} \quad \begin{pmatrix} 2,2,1 & 0,0,0 \\ 0,0,0 & 1,1,1 \end{pmatrix}$$

has 9 different equilibria, of which 3 are symmetric, 4 are pure and 2 are completely mixed with the same pay-off.

**6.** Under the assumptions of theorem 4.1 or of ex. I.4Ex.5, show that the support of  $\mu^i$  is contained in the closed set of pure best replies:

$$\begin{aligned} \{s \in S^i \mid \int_{\prod_{h \neq i} S^h} F^i(s^1, \dots, s^{i-1}, s, s^{i+1}, \dots, s^n) d \prod_{h \neq i} \mu^h = \\ \max_{\tilde{s} \in S^i} \int_{\prod_{h \neq i} S^h} F^i(s^1, \dots, s^{i-1}, \tilde{s}, s^{i+1}, \dots, s^n) d \prod_{h \neq i} \mu^h \}. \end{aligned}$$

In particular, the  $\mu^h$  ( $h \neq i$ ) are “equalising” for all pure strategies which are in the support of  $\mu^i$ .

**7. Linear complementarity problem.** This concerns the computation of equilibria of bi-matrix games and related problems in operations research (cf. Lemke and Howson (1964) for corresponding algorithms).

a. *Reduction to symmetric games.* Let  $(A, B)$  be a bi-matrix game. Add constants to the pay-offs such as to have  $\max_x \min_y x A y > 0$ ,  $\max_y \min_x x B y > 0$ , (i.e., each player can guarantee himself a positive amount). Denote by  $(S, S^t)$  the symmetric game  $\begin{pmatrix} 0, 0 & A, B \\ B^t, A^t & 0, 0 \end{pmatrix}$ . Show that there is a one-to-one correspondence between equilibria of  $(A, B)$  and symmetric equilibria of  $(S, S^t)$ .

b. Let  $T = K - S$ , for some constant matrix  $K$  larger than the pay-off  $x^T S x$  of any symmetric equilibrium  $x$  of  $S$ . In terms of  $T$ , the symmetric equilibria of  $S$  are characterised by the linear complementarity problem, denoting by  $z$  a symmetric equilibrium strategy  $x$  divided by  $x^T x$ :

$$Tz = \tilde{1} + u, \quad z \geq 0, u \geq 0, \langle z, u \rangle = 0$$

“ $\tilde{1}$ ” being a column of ones. Show that our assumption on  $K$  can be rewritten in terms of  $T$  as:

$$Tz^+ - z^- = \lambda \cdot \tilde{1}, \quad \lambda \leq 0 \Rightarrow z = \lambda \cdot \tilde{1}.$$

c. Denote by  $\mathcal{S}$  the class of  $n \times n$ -matrices  $S$  such that, for some non-negative, non-singular diagonal matrix  $D$ , one has

$$z \geq 0, z \neq 0 \Rightarrow z^t (SD) z > 0.$$

i. Show that, if  $D_1$  and  $D_2$  are non-negative, non-singular diagonal matrices, then  $S \in \mathcal{S} \Rightarrow D_1 S D_2 \in \mathcal{S}$ , and  $S^t \in \mathcal{S}$ .

ii.  $\mathcal{S}$  contains all sums of a positive definite matrix, a non-negative matrix and an anti-symmetric matrix.

iii. For  $T \in \mathcal{S}$ , the condition sub I.4Ex.7b is satisfied.

iv. For  $T \in \mathcal{S}$ , let  $F: x \rightarrow T x^+ - x^-$ . Show that  $F(\mathbb{R}^n)$  contains the interior of the positive orthant.

HINT. Use I.4Ex.7ci to perform a rescaling from the result of I.4Ex.7b.

d. *A direct approach, compare with ex. I.4Ex.4.* For  $T \in \mathcal{S}$ , the map  $F: x \rightarrow T x^+ - x^-$  (with  $F(\infty) = \infty$ ) is of degree one (homotopic to the identity); in particular  $F(\mathbb{R}^n) = \mathbb{R}^n$ .

HINT. For  $F_t: x \rightarrow (1-t)x + tF(x)$  to be a homotopy, one needs that  $\forall M, \exists K: \|x\| \geq K \Rightarrow \|F_t(x)\| \geq M, \forall t \in [0, 1]$ . By homogeneity (and extracting subsequences), this is equivalent to  $F_t(x) = 0 \Rightarrow x = 0$ , i.e. to  $[x \geq 0, x^t x = 1, T x \geq (x^t T x) \cdot x] \Rightarrow x^t T x > 0$  or equivalently  $x^- = (T + \lambda I)x^+, \lambda \geq 0 \Rightarrow x = 0$ , which clearly holds for  $T \in \mathcal{S}$ . More generally, this holds for all matrices in the connected component of  $I$  in the open set of matrices  $M$  satisfying  $M x^+ = x^- \Rightarrow x = 0$ . [Same argument: Prove openness directly, and use semi-algebraicity to deduce the equivalence of connectedness with pathwise connectedness, and the existence of finitely many connected components]. Show that there is more than one connected component (consider  $-I$ , at least for  $n$  odd.).

e. By rescaling, I.4Ex.7b implies that if  $F_i(z) < 0, \forall i$  implies  $z_i < 0, \forall i$ , and if  $F(z) = 0$  implies  $z = 0$ , then  $F(z) = a$  has a solution whenever  $a_i > 0, \forall i$ . Verify that the condition  $x^- = (T + \lambda I)x^+, \lambda \geq 0 \Rightarrow x = 0$  sub I.4Ex.7d above is weaker.

f. The above existence result remains true over any ordered field.

8. Let  $K$  be a real closed field. Call a subset  $A$  of  $K^\ell$  bounded if there exists  $m \in K$  such that, for all  $x \in A$ ,  $x = (x_1, \dots, x_\ell)$ , for all  $i$  ( $1 \leq i \leq \ell$ ),  $|x_i| \leq m$ . Call a semi-algebraic subset of  $K^\ell$  closed if it can be described using only weak inequalities (i.e.  $\geq$  or  $\leq$ ).

Assume each set  $S^i$  is a closed bounded semi-algebraic subset of  $K^{\ell_i}$ , which is a union of closed algebraic subsets  $A_j^i$  ( $j = 1, \dots, n_i$ ). Assume the restriction of each  $F^i$  to  $\prod_{j=1}^{n_i} A_j^i$  (for all  $j_i$ :  $1 \leq j_i \leq n_i$ ) is a polynomial in the coordinates with coefficients in  $K$ . Show that the game has an equilibrium point in algebraic strategies, i.e. mixed strategies with finite support, the weights of the distinct points in the support being in  $K$  (i.e. as in ex. I.4Ex.5).

HINT. Use first cor. 3.6 p. 27 to reduce the problem to the case  $K = \mathbb{R}$ . Show then that the assumptions of theorem 4.1 p. 39 are satisfied, so that there exists an equilibrium point in mixed strategies  $\mu^1, \dots, \mu^n$ . Note then that the restriction of  $\mu^i$  to  $A_j^i$  intervenes only through a finite number of its moments, so (ex. I.3Ex.10 p. 34) each  $\mu^i$  may be taken with finite support (with at most  $N$  points,  $N$  depending only on the number of  $A_j^i$ 's and on the degrees of the polynomials appearing in the  $F^j$ 's — this is needed to justify the use of Tarski's theorem).

**9.** (Gross, 1957) Let  $S = T = [0, 1]$ , and let  $f(s, t) = \sum_{n=0}^{\infty} 2^{-n}[2s^n - (1 - \frac{s}{3})^n - (\frac{s}{3})^n][2t^n - (1 - \frac{t}{3})^n - (\frac{t}{3})^n]$ . Show that  $f(s, t)$  is a rational pay-off function without singularities on the unit square, i.e. can be written as the ratio of two polynomials with integer coefficients in  $s$  and  $t$ , the denominator of which vanishes nowhere in the unit square.

Show that the value of the corresponding two-person zero-sum game is zero, and that both players have the Cantor distribution  $C(x)$  as unique optimal strategy.

HINT. Note that  $C(x)$  is the unique monotone solution of  $C(x) + C(1-x) = 1$  and  $C(x) = 2C(\frac{x}{3})$ . Deduce that, for any continuous real valued function  $h(x)$ ,  $\int_0^1 [2h(x) - h(1 - \frac{x}{3}) - h(\frac{x}{3})] dC(x) = 0$ . There only remains to show the uniqueness of the solution. Let  $\mu$  be any optimal strategy of player I. Then  $\int_0^1 f(s, t) \mu(ds)$  is an analytical function of  $t$  (integrate the series term by term) that has to vanish on the (infinite) support of the Cantor distribution: it vanishes everywhere. Thus  $\int_0^1 f(s, t) \mu(ds) \mu(dt) = 0 = \sum_{n=0}^{\infty} 2^{-n}(\mu_n)^2$ , so that, for all  $n$ ,  $\int_0^1 [2s^n - (1 - \frac{s}{3})^n - (\frac{s}{3})^n] \mu(ds) = \mu_n = 0$ . This determines inductively in a unique way all moments  $\int_0^1 s^n \mu(ds)$  ( $n \geq 1$ ), and thus the distribution  $\mu$  itself.

COMMENT 4.2. This shows that the solution may be a singular continuous distribution. Glicksberg and Gross (1953) show how to construct rational games over the square whose solutions are purely atomic and have dense support. In view of the importance of the possibility of reducing to mixtures with finite support in the proof of ex.I.4Ex.8, it is not surprising to find (Glicksberg and Gross, 1953) the following example of a rational game over the square with transcendental value, which therefore definitely excludes the possibility of an algebraic solution of rational games.

**10.** Let  $S = T = [0, 1]$ , and  $f(s, t) = (1+s)(1+t)(1-st)/(1+st)^2$ .

Show that the value of the corresponding two-person zero-sum game is  $4/\pi$ , and that each player's optimal strategy is given by the cumulative distribution function on  $[0, 1]$ :  $F(x) = \frac{4}{\pi} \arctan \sqrt{x}$ .

**11.** Can you get similar results to ex. I.4Ex.8 that would include Karlin's 1959 "Generalised Convex Games" and/or "Bell-shaped Games".

**12. Bertrand Competition.** Two firms compete in price over the sale of a product. Assume that production costs are zero (if only by subtracting mentally a constant from prices). Assume the demand  $D(p)$  at any price  $p \geq 0$  satisfies  $0 \leq D(p) < +\infty$ , and  $\limsup_{q \leq p} D(q) \geq D(p) \geq \limsup_{q \geq p} D(q)$  for  $p > 0$ . Typically the demand is decreasing, which implies those relations. We also assume that demand is not identically zero in a neighbourhood of zero. The firm announcing the lowest price serves the whole demand, at that price. If the announced prices are equal, consumers decide indifferently what share of their demand to allocate to each firm. Profit is  $pd$ , if  $p$  is the announced price and  $d$  the demand served, and firms want to maximise expected profits.

- a. Show that  $(0, 0)$  is a Nash equilibrium.
- b. Let  $G(p) = \max_{q \leq p} qD(q)$ . Show that  $G(p)$  is continuous, except possibly at  $p = 0$ .
- c. Show that  $\{p \mid pD(p) < G(p)\}$  is a countable union of disjoint intervals, open on the right.
- d. For any Nash equilibrium, denote by  $\pi_i$  the expected profit of player  $i$ , and by  $F_i(p)$  the probability that his announced price is  $\geq p$ , i.e.  $F_i(p) = \sigma_i([p, +\infty[)$ .

i. Show that, if  $\pi_1 = 0$  or  $\pi_2 = 0$ , we have only the equilibrium sub I.4Ex.12b.

Henceforth we assume  $\pi_1 > 0$ ,  $\pi_2 > 0$ .

ii. Show that  $\pi_1 < +\infty$  and  $\pi_2 < +\infty$ .

iii. Show that  $F_i(p) \leq \pi_j H(p)$ ,  $\forall p$ , with  $H(p) = 1/G(p)$  ( $i \neq j$ ).

iv. Let  $T_j = \{p \mid F_i(p) = \pi_j H(p)\}$ . Show that  $T_j$  is closed, except possibly at zero, and that  $\sigma_j(T_j) = 1$ .

e. Show that  $T_1 = T_2 (= T)$ .

HINT. The complement of  $T_1$  is a countable, disjoint union of open intervals  $\alpha_k, \beta_k$ , (with possibly  $\beta_1 = +\infty$ ), plus either  $\{0\}$  or  $[0, \beta_0[$ . On  $\alpha_k, \beta_k[$  — and on  $]0, \beta_0[$  — we have  $F_1(p) = F_1(\beta_k) \leq \pi_2 H(\beta_k)$  (with  $F(\infty) = H(\infty) = 0$  by convention), and  $\pi_1 H(p) > F_2(p) \geq F_2(\beta_k) = \pi_1 H(\beta_k)$ . Thus  $\pi_2 H(p) > F_1(p)$  hence  $\alpha_k, \beta_k \cap T_2 = \emptyset$ . And certainly  $0 \notin T_1$ ,  $0 \notin T_2$ .

f. Show that any Nash equilibrium is symmetric (thus  $\pi_1 = \pi_2 = \pi$ ,  $\sigma_1 = \sigma_2 = \sigma$ ,  $F_1 = F_2 = F$ ).

HINT. Let  $\tilde{H}(p) = \sup\{H(q) \mid q \in T, q \geq p\}$ , with  $\sup(\emptyset) = 0$ . By I.4Ex.12d and I.4Ex.12e,  $F_i(p) = \pi_j \tilde{H}(p)$ . And  $F_i(0) = 1$ .

g. Show that any Nash equilibrium (other than  $(0, 0)$ ) is nonatomic.

HINT. By I.4Ex.12f every atom is common to both players, and is  $> 0$ . Use then Bertrand's "undercutting" argument.

h. Conclude that either  $T = [a, +\infty[$  with  $a > 0$  or  $T = ]0, +\infty[$ . In particular, let  $\bar{G}(p) = G(p)$  for  $p > 0$ ,  $\bar{G}(0) = \lim_{\varepsilon \geq 0} G(\varepsilon)$ , then  $F(p) = 1 \wedge (\bar{G}(a)/\bar{G}(p))$  for some  $a \geq 0$  with  $0 < \pi = \bar{G}(a) < +\infty$ .

i. Conclude that, for the existence of equilibria  $\neq (0, 0)$ , one needs both  $\bar{G}(\infty) = \infty$  — i.e.  $\limsup_{p \rightarrow \infty} pD(p) = +\infty$  — and  $\bar{G}(0) < +\infty$  — i.e.  $\limsup_{p \rightarrow 0} pD(p) < +\infty$ .

j. Show that, for any  $F(p)$  as sub I.4Ex.12h, one has

- (1)  $(pD(p))F(p) \leq \pi$ ,  $\forall p \geq 0$  and
- (2)  $\sigma\{p \mid (pD(p))F(p) < \pi\} = 0$ .

HINT. For the second point, use I.4Ex.12c.

k. Deduce from I.4Ex.12j that the condition sub I.4Ex.12i is also sufficient: any such  $F$  is an equilibrium, hence the equilibria different from  $(0, 0)$  are under that condition in one-to-one correspondence with the profit levels  $\pi > 0$ ,  $\pi \geq \limsup_{p \rightarrow 0} pD(p)$ .

In particular, if possible profits are bounded in this market, there is a unique equilibrium,  $(0, 0)$ , which involves strategies that are dominated by every other strategy. The next exercise considers better behaved examples of the same situation.

### 13. Variations on Bertrand Competition.

a. Players I and II choose respectively  $x$  and  $y$  in  $[0, 1]$ , and get respectively as pay-offs  $f(x, y) = y^2 x / (2x^2 + y^2)$  and  $f(y, x)$ .

i. The pay-off is continuous, even Lipschitz.

ii. Show that  $(0, 0)$  is the unique Nash equilibrium.

HINT. Observe that, for  $y > 0$  fixed,  $f(x, y)$  decreases after its maximum at  $y/\sqrt{2} < y$ . Do not forget any type of mixed or mixed versus pure equilibrium.

iii. Yet those strategies are dominated by every other strategy: for any other mixed strategy of the opponent any other mixed strategy of the player is strictly better than the equilibrium strategy.

b. Show that, in any game with continuous pay-off function on compact metric strategy spaces, there is some equilibrium whose strategies are limits of undominated strategies. [In particular, the above example of Bertrand's "undercutting" argument is probably the best known case of a limit of undominated points being dominated, cf. also ex. I.3Ex.7 p. 33.]

HINT. Show that each player has a strategy with full support. Show that the perturbed game where each player's choice is played with probability  $1 - \varepsilon$ , (independently of the other players), while a fixed strategy with full support is played otherwise —, satisfies the usual conditions for existence of equilibria. Then let  $\varepsilon \rightarrow 0$ .

c. Show that no example like sub I.4Ex.13a is possible with polynomials — or more generally with separable pay-offs  $f(x, y) = \sum_i^n g_i(x)h_i(y)$  (when the  $g_i$  and  $h_i$  are continuous functions on compact strategy spaces): there cannot exist a unique equilibrium which is dominated by every other strategy.

HINT. The requirement implies the equilibrium is pure — say  $(x_0, y_0)$ . Subtract  $f(x_0, y)$  from player I's pay-off  $f(x, y)$ , and similarly for II's pay-offs: this preserves the separability, and the equilibria: we have now, with  $\varphi_y(x) = f(x, y)$ ,  $\varphi_y(x_0) = 0$ ,  $\varphi_{y_0}(x) = 0$ ,  $\varphi_y(x) \geq 0$ , and  $\|\varphi_y\| > 0$  for  $y \neq y_0$  (if  $\|\varphi_y\| = 0$ ,  $(x_0, y)$  would be another equilibrium, since  $\psi_{x_0}(y) = 0$ ). The functions  $\varphi_y$  vary in a finite ( $n$ )-dimensional space, where all norms are equivalent, and where the unit sphere is compact: construct thus a new game with I's pay-off function  $G(x, y) = [\varphi_y(x)]/\|\varphi_y\|$ , for  $y \neq y_0$ , and take for strategy space of player II the closure of those functions (for computing player II's pay-off, points in the closure are sent to  $y_0$ ). Do now the same with the pay-off function of II. The new game is still continuous (on compact strategy spaces), but now since  $\max_x G(x, y) = 1$ ,  $\forall y$ , player I's minmax pay-off is  $> 0$  (strategy with full support). Hence an equilibrium of this game assigns zero probability to  $x_0$  and to  $y_0$  ( $G(x_0, y) = 0$ ). By the separability, the equilibrium can be chosen with finite support (ex. I.3Ex.10 p. 34). Reconstruct from this an equilibrium of the original game, with same support.

d. Let  $F_1(x, y) = y[x^3 - (1 + 2y)x^2 + (2y + \frac{1}{2}y^2)x]$ ,  $F_2(x, y) = F_1(y, x)$  ( $(x, y) \in [0, 1]^2$ ) be the pay-off functions of players I and II. As polynomials, the pay-offs are a bi-linear function of the first 3 moments  $(\sigma_1, \sigma_2, \sigma_3)$  ( $\sigma^i = \int_0^1 x^i \sigma(dx)$ ) of the strategy  $\sigma$  of I, and of  $(\tau_1, \tau_2, \tau_3)$ , given by the pay-off matrix:

$$\begin{matrix} & \tau_1 & \tau_2 & \tau_3 \\ \sigma_1 & \begin{pmatrix} 0, 0 & 2, -1 & 1/2, 1 \end{pmatrix} \\ \sigma_2 & \begin{pmatrix} -1, 2 & -2, -2 & 0, 0 \end{pmatrix} \\ \sigma_3 & \begin{pmatrix} 1, 1/2 & 0, 0 & 0, 0 \end{pmatrix} \end{matrix}$$

where  $(\sigma_1, \sigma_2, \sigma_3)$  and  $(\tau_1, \tau_2, \tau_3)$  are to be selected in the compact convex "moment space"  $C$  in  $\mathbb{R}^3$  given by

$$C = \{(\mu_1, \mu_2, \mu_3) \mid \mu_2^2 \leq \mu_1\mu_3, (\mu_1 - \mu_2)^2 \leq (1 - \mu_1)(\mu_2 - \mu_3) \text{ together with } 0 \leq \mu_3 \leq 1\}$$

Show that indeed  $C = \{\int_0^1 (x, x^2, x^3) \mu(dx) \mid \mu \in \Delta([0, 1])\}$ .

HINT. A direct proof may go by characterising by hand the supporting affine functionals, hence the extreme rays of the cone of third degree polynomials which are non-negative on  $[0, 1]$ .

Alternatively, write a polynomial as  $A^t \cdot X$  — all vectors being column vectors, with  $A$  the vector of coefficients and  $X$  the successive powers of the variable  $x$ . Then  $A^t X X^t A$  is the square of a polynomial, hence non-negative; since  $\mu$  is carried by  $[0, 1]$ , this remains true if  $XX^t$  is multiplied by powers of  $x$  and/or of  $1 - x$ , yielding a matrix  $Y$ . Thus the expectation of  $Y$  is positive semi-definite. Hence the validity of the inequalities.

Conversely, observe first that on  $C$ , we have  $0 \leq \mu_3 \leq \mu_2 \leq \mu_1 \leq 1$ ; and that  $\{\mu \in C \mid \mu_2^2 = \mu_1\mu_3\}$  is the set of convex combinations of  $(0, 0, 0)$  and a point on the curve  $(x, x^2, x^3)$  — and similarly for the other inequality and the point  $(1, 1, 1)$ . Conclude that the whole of  $C$  belongs to the convex hull of the curve.

More generally, the converse may go by showing that the polynomials we considered span indeed the cone of polynomials non-negative on  $[0, 1]$ : polynomials who are non-negative on  $[0, 1]$  can be factored in a product of quadratic factors which are everywhere non-negative, and of linear factors which are of the

form  $(x + a)$  or  $(a + 1 - x)$  for  $a \geq 0$ . Decompose the quadratic factors as sums of two squares, and the linear factors into the sum of the constant non-negative polynomial  $a$  (a square) and either  $x$  or  $1 - x$ : our whole polynomial is now rewritten as a sum of products of a square with powers of  $x$  and of  $1 - x$ .

e. Show that the game defined sub I.4Ex.13d has a unique Nash equilibrium,  $(0, 0)$ , while strategy 0 is dominated by 1, with strict inequality for all strategies of the opponent different from the equilibrium strategy.

HINT. The dominance is clear. The statement does not depend on the factor  $y$  in  $F(x, y)$ . Let  $\varphi_y(x) = F_1(x, y)/y$ . For  $y > 0$ ,  $\varphi'_y(0) > 0$ ,  $\varphi'_y(y) < 0$ , there is a local maximum  $x(y)$  in  $]0, y[$ , with  $\varphi_y(x(y)) > \varphi_y(y)$ . Also  $\varphi_y(y) - \varphi_y(1) \geq 0$  — thus  $x(y) < y$  is the global maximum. By symmetry, conclude that there is no equilibrium where one player's strategy is pure and  $\neq 0$ . And if the pure strategy is zero, since it is as good a reply as 1, it means the opponent also uses zero with probability one. Thus there only remains to consider equilibria where no player uses a pure strategy. For any such strategy of II, player I's pay-off function is of the form  $Ax^3 + Bx^2 + Cx$ , with  $A > 0$ ,  $C > 0$ : thus if the maximum is not unique, it is a set of the form  $\{x_0, 1\}$  with  $0 < x_0 < 1$ . Thus also II uses a mixture on  $\{y_0, 1\}$  with probabilities  $q$  and  $1 - q$ , with  $0 < y_0 < 1$ ,  $0 < q < 1$ : in player I's pay-off we have  $B = -1 - 2(qy_0 + 1 - q)$ . On the other hand, since I is indifferent between  $x_0$  and 1, we have  $Ax^3 + Bx^2 + Cx = (x - x_0)^2(x - 1) + \text{constant}$ , hence  $B = -1 - 2x_0$ . Thus  $x_0 > y_0$  — and dually  $y_0 > x_0$ .

f. In  $\begin{pmatrix} x, 0 & -x^2, x & 0, 0 \\ 0, 1 & 0, 0 & 1, 0 \end{pmatrix}$  player I also has to choose  $x \in [0, 1]$  if he selects the top row.

This game also has a unique Nash equilibrium, where player I uses a dominated strategy. Further the game is “almost strictly competitive”: the unique equilibrium strategies are also dominant strategies in the game where each player would instead try to minimise his opponent's pay-off — thus we are as close to the zero-sum situation and to finite games as one could possibly get (cf. I.4Ex.13g and I.4Ex.13h)). But the game no longer has the flavour of Bertrand's “undercutting”, player II's strategy is not dominated (cf. I.4Ex.13b and I.4Ex.13g), and the dominating strategy of player I is no longer strictly better for every strategy different from the equilibrium strategy (cf. I.4Ex.13i).

REMARK 4.3. As defined by Aumann (1961a), “almost strictly competitive” means:

- equilibrium pay-offs in  $(g_1, g_2)$  and  $(-g_2, -g_1)$  are the same
- the sets of equilibria in  $(g_1, g_2)$  and  $(-g_2, -g_1)$  intersect.

Then there exists  $(s, t)$  such that  $s$  realises  $\max_s \min_t g_1$  and  $\min_s \max_t g_2$  (and similarly for  $t$ ).  $s$  is not a dominant strategy.

g. If player I has a finite strategy set, his set of undominated mixed strategies is closed and is a union of faces. Indeed, if  $\sigma$  is dominated, say by  $\tilde{\sigma}$ , then any convex combination  $\alpha\sigma + (1 - \alpha)\tau$  ( $0 < \alpha < 1$ ) is also dominated by  $\alpha\tilde{\sigma} + (1 - \alpha)\tau$ . Thus, in a finite dimensional vector space ordered by an arbitrary positive cone, the set of admissible points of a compact polyhedron is a union of faces (Cf. Arrow et al., 1953).

h. In the zero-sum situation, say under the assumptions of Theorem 2.6 p. 17, there always exist undominated optimal strategies.

HINT. Under those assumptions, show that every strategy is dominated by an undominated strategy. (Use Zorn's lemma or an equivalent).

i. If player II's strategy set is finite, player I's is compact, and the pay-off continuous, and if there is an equilibrium where I's strategy is dominated by another one which is strictly better for every strategy of II different from the equilibrium strategy, then there is another equilibrium.

HINT. Since player I maximises, the dominating strategy will then dominate a neighbourhood of his equilibrium strategy. Apply then I.4Ex.13b, only II's strategy has to be perturbed.

**14.** (Dasgupta and Maskin, 1986) Let  $S^i = [0, 1]$  for all  $i$  and let  $\mathcal{F}$  be a finite family of one-to-one continuous mappings from  $[0, 1]$  to  $[0, 1]$ . Let  $A(i) = \{s \mid \exists j \neq i, \exists f \in \mathcal{F}: s^j = f(s^i)\}$ . Assume that  $F^i$  is continuous on  $\mathbb{C}B(i)$  for all  $i$ , with  $B(i) \subseteq A(i)$ .

- a. Define  $G_n$  as the initial game restricted to  $S_n^i = \{k/n \mid k = 0, 1, \dots, n\}$  for all  $i$  and let  $\sigma_n$  be a corresponding equilibrium such that  $\sigma_n$  converges (weak $^\star$ ) to  $\bar{\sigma}$  and  $F^i(\sigma_n)$  converges to some  $\overline{F^i}$ ,  $i = 1, \dots, n$ . Prove that there are at most countably many points  $s^i$  such that  $F^i(s^i, \bar{\sigma}^{-i}) > \overline{F^i}$ . (Prove first that for such an  $s^i$  one has  $\bar{\sigma}^{-i}(\{s^{-i} \mid (s^i, s^{-i}) \in A^i\}) > 0$  and that this implies:  $\exists j \neq i, \exists f \in \mathcal{F}, \sigma^j(f(s^i)) > 0$ ).
- b. Assume moreover  $\sum_{i=1}^n F^i$  u.s.c. and that each  $F^i$  satisfies, for all  $s^i \in B_i(i)$ :  $\exists \lambda \in [0, 1]$  such that for all  $s^{-i}$  with  $s = (s^i, s^{-i}) \in B(i)$ ,  $\lambda \liminf_{t^i \nearrow s^i} F^i(t^i, s^{-i}) + (1 - \lambda) \liminf_{t^i \searrow s^i} F^i(t^i, s^{-i}) \geq F^i(s)$ . Prove that  $\bar{\sigma}$  is an equilibrium with pay-off  $\overline{F}$ .
- c. Extend the result to  $S^i$  compact convex in some  $\mathbb{R}^m$ .
- d. Consider the two-person symmetric game with

$$F^1(s^1, s^2) = \begin{cases} 0, & \text{if } s^1 = s^2 = 1, \\ s^1, & \text{otherwise.} \end{cases}$$

(Note that  $F^1 + F^2$  is not u.s.c. at  $(1, 1)$ ).

**15. Fictitious Play.** (Robinson, 1951) “Fictitious play” is the procedure where each player chooses at every stage a best reply against the distribution of his opponent’s past choices. We prove here that, for zero-sum games, the frequency of moves converges to the optimal strategies.

Given an  $m \times n$  real pay-off matrix  $A$ , write  $A^i$  for the row  $i$  and  $A_j$  for the column. Call admissible a sequence  $(\alpha(t), \beta(t))$  in  $\mathbb{R}^n \times \mathbb{R}^m$ ,  $t \in \mathbb{N}$  satisfying:

- (1)  $\inf_j \alpha_j(0) = \sup_i \beta_i(0)$ .
- (2)  $\forall t \in \mathbb{N}, \exists i: \alpha(t+1) = \alpha(t) + A^i$ , with  $\beta_i(t) = \sup_k \beta_k(t)$ .
- (3)  $\forall t \in \mathbb{N}, \exists j: \beta(t+1) = \beta(t) + A_j$ , with  $\alpha_j(t) = \inf_k \alpha_k(t)$ .

Say that  $i$  (resp.  $j$ ) is useful in  $(t_1, t_2)$  if there exists  $t$  in  $\mathbb{N}$  with  $t_1 \leq t \leq t_2$  and  $\beta_i(t) = \sup_k \beta_k(t)$  (resp.  $\alpha_j(t) = \inf_k \alpha_k(t)$ ).

- a. Prove that if all  $j$  are useful in  $(s, s+t)$  then:

$$\sup_j \alpha_j(s+t) - \inf_j \alpha_j(s+t) \leq 2t \|A\|.$$

Defining  $\mu(t)$  as  $\sup_i \beta_i(t) - \inf_j \alpha_j(t)$ , prove that if all  $i$  and  $j$  are useful in  $(s, s+t)$ , then:

$$\mu(s+t) \leq 4t \|A\|.$$

b. Prove by induction that for every matrix  $A$  and  $\varepsilon > 0$ , there exists  $s$  in  $\mathbb{N}$  such that  $\mu(t) \leq \varepsilon t$  for all  $t \geq s$  and all admissible sequences.

HINT. Let  $r$  be such that  $\mu'(t) \leq \varepsilon t$  for all  $t \geq r$  and all admissible sequences associated to a (strict) submatrix  $A'$  of  $A$  — with corresponding  $\mu'$ . Prove that if  $i$  is not useful in  $(s, s+r)$  one has  $\mu(s+r) \leq \mu(s) + \varepsilon r$ . Finally given  $t = qr+s$  with  $s < r$ , let  $p$  be the largest integer  $\leq q$  such that all  $i, j$  are useful in  $((p-1)r+s, pr+s)$ , and 0 if this never occurs. Show that  $\mu(t) \leq \mu(pr+s) + \varepsilon(q-p)r$  and conclude.

- c. Prove that for all admissible sequences:

$$\lim_{t \rightarrow \infty} (\sup_i \beta_i(t)/t) = \lim_{t \rightarrow \infty} (\inf_j \alpha_j(t)/t) = v(A)$$

(where  $v(A)$  is the value of  $A$ ).

**16.** (Shapley, 1964) Consider fictitious play in the following two-person game:  $S_1 = S_2 = \{1, 2, 3\}$

$$F_1 = \begin{pmatrix} a_1 & c_2 & b_3 \\ b_1 & a_2 & c_3 \\ c_1 & b_2 & a_3 \end{pmatrix} \quad F_2 = \begin{pmatrix} \beta_1 & \gamma_1 & \alpha_1 \\ \alpha_2 & \beta_2 & \gamma_2 \\ \gamma_3 & \alpha_3 & \beta_3 \end{pmatrix}$$

with  $a_i > b_i > c_i$  and  $\alpha_i > \beta_i > \gamma_i$  for all  $i$ .

a. Assume the first choice is  $(1, 1)$ . Prove that after a sequence of  $(1, 1)$  a sequence of  $(1, 3)$  will occur followed by a sequence of  $(3, 3)$  then of  $(3, 2), (2, 2), (2, 1)$  and  $(1, 1)$  again.

b. Denote by  $r_{11}$  the length of a sequence of  $(1, 1)$  and by  $r_{13}$  the length of the following sequence of  $(1, 3)$ . Show that  $r_{13} \geq r_{11}(a_1 - c_1)/(a_3 - b_3)$  and deduce inductively, if  $r'_{11}$  is the length of the next sequence of  $\{1, 1\}$  that:

$$r'_{11} \geq \prod_i \left( \frac{a_i - c_i}{a_i - b_i} \right) \left( \frac{\alpha_i - \gamma_i}{\alpha_i - \beta_i} \right) r_{11}.$$

c. Deduce that the empirical strategies do not converge. Furthermore, no subsequence converges to the (unique) equilibrium pair.

**17. On fixed points.** Note that the quasi-concavity assumption in theorem 4.1 p. 39 was not used in its full force — only the convexity of the best reply sets was used. But even this is too much: given an upper semi-continuous correspondence  $T$  from a compact convex set  $S$  to itself, the condition for  $T$  to have a fixed point should be purely topological. This is the aim of the present exercise. In fact, it will be more convenient to argue in terms of a compact Hausdorff space  $X$  (e.g. the graph of  $T$ ), together with two (continuous) maps  $f$  and  $g$  from  $X$  to  $S$  (e.g., the two projections). We look then for a solution of  $f(x) = g(x)$ , and the best reply sets are the sets  $f^{-1}(s)$ .

We assume also that  $S$  is Hausdorff, that  $f$  is onto and that  $\tilde{H}^q(f^{-1}(s)) = 0$  for all  $q \geq 0$  and  $s \in S$ , where  $\tilde{H}$  denotes reduced Čech-cohomology with coefficients in a module  $G$  (Spanier, 1966). We express this condition shortly by saying that  $f$  is  $(G)$ -acyclic, this is a purely topological condition on the best reply sets, satisfied as soon as each of them is contractible (homotopy invariance). Similarly we call an upper semi-continuous correspondence  $\Gamma$  from  $X$  to  $S$  acyclic if the projection from its graph to  $X$  is so, i.e., if  $\forall x \in X, \tilde{H}^q(\Gamma(x)) = 0$  and  $\Gamma(x) \neq \emptyset$ .

a. Prove that, when  $S$  is a finite dimensional convex set, there exists a solution (of  $f(x) = g(x)$ ).

HINT. Denote by  $\partial S$  the boundary of  $S$ , and let  $\partial X = f^{-1}(\partial S)$ . Apply the Vietoris-Begle mapping theorem (Spanier, 1966, VI.9.15) to  $f: X \rightarrow S$  and to  $f: \partial X \rightarrow \partial S$  to conclude by the “five lemma” (Spanier, 1966, IV.5.11) and exactness (Spanier, 1966, V.4.13) that  $f: (X, \partial X) \rightarrow (S, \partial S)$  is an isomorphism in Čech-cohomology. If  $f(x) \neq g(x)$  for all  $x$ , construct (as in Spanier, 1966, IV.7.5) a map  $\tilde{f}: X \rightarrow \partial S$  which is homotopic to  $f$ . Conclude to a contradiction.

b. To show that acyclic maps are the “right” class of maps for this problem, prove that if  $f: X \rightarrow Y$  and  $g: Y \rightarrow Z$  (all spaces compact) are  $(G)$ -acyclic, so is  $g \circ f$ .

HINT. Clearly  $g \circ f$  is onto. Let  $S = g^{-1}(z) : \tilde{H}^q(S) = 0$ . Since  $f$  as a map from  $\tilde{X} = f^{-1}(S)$  to  $S$  is acyclic,  $f^*: \tilde{H}^q(S) \rightarrow \tilde{H}^q(\tilde{X})$  is an isomorphism (cf. a)):  $\tilde{H}^q(\tilde{X}) = 0$  for all sets  $\tilde{X} = (g \circ f)^{-1}(z)$ .

c. A continuous affine map  $f$  on a compact convex subset  $X$  of a Hausdorff topological vector space is acyclic to  $f(X)$ .

HINT. Use the above mentioned contractibility criterion.

d. Assume  $S$  is a compact convex subset of a topological vector space, such that points of  $S$  are separated by continuous affine functionals on  $S$  — shortly henceforth: a compact convex set. Prove that there exists a solution.

HINT. By the separation property on  $S$  and compactness of  $X$  it suffices to show that  $(p \circ f)(x) = (p \circ g)(x)$  has a solution for every continuous affine map  $p: S \rightarrow \mathbb{R}^k$ . Use I.4Ex.17b and I.4Ex.17c to show that  $p \circ f$  and  $p \circ g$  satisfy the conditions sub I.4Ex.17a.

e. If  $S$  is a retract of a compact convex set  $Y$ , there exists a solution.

HINT. The assumption means that there are continuous maps  $i: S \rightarrow Y$  and  $r: Y \rightarrow S$  such that  $(r \circ i)(s) = s$ ,  $\forall s \in S$ . Let then  $\tilde{X} = \{(x, y) \in X \times Y \mid f(x) = r(y)\}$ ; denote by  $p_1$  and  $p_2$  the two projections, and apply I.4Ex.17d to  $\tilde{f} = p_2$  and  $\tilde{g} = i \circ g \circ p_1$ .

At this stage we have a purely topological statement. To see that those are the right assumptions, note that the statement remains true, but without additional generality, if  $g$  is replaced by an acyclic correspondence  $\Gamma$  from  $X$  to  $S$ . Slightly more generally, since  $f^{-1}$  is an acyclic correspondence from  $S$  to  $X$ , we have:

f. Let  $X_i$  ( $i = 0, \dots, n$ ) be compact Hausdorff spaces, with  $X_0 = X_n$  being a retract of a compact convex set. Let  $\Gamma_i$  ( $i = 1, \dots, n$ ) be  $(G)$ -acyclic correspondences from  $X_{i-1}$  to  $X_i$ . Then their composition  $\Gamma$  has a fixed point  $x \in \Gamma(x)$ .  $[(\Gamma_1 \circ \Gamma_2)(x_0) = \bigcup_{x_1 \in \Gamma_1(x_0)} \Gamma_2(x_1)]$ .

HINT. Let  $Y_j = \{(x_0, \dots, x_j) \in \prod_{i \leq j} X_i \mid x_i \in \Gamma_i(x_{i-1}) \text{ for } 0 < i \leq j\}$ , with projection  $f_j$  to  $Y_{j-1}$ . Then the  $f_j$  are acyclic, hence I.4Ex.17b so is their composition  $f: Y_n \rightarrow X_0$ . Let also  $g: Y_n \rightarrow X_n (= X_0)$  denote the projection to the last coordinate, and apply I.4Ex.17e.

This proof shows, in other words, that the class of correspondences which are the composition of a  $(G)$ -acyclic correspondence with a continuous map is stable under composition.

g. Reverting to our context of games, use Künneth's formula (Spanier, 1966, VI.Ex.E) to show that, if  $G$  is a field, the best reply correspondence will be acyclic if and only if any best reply set  $B$  of any single player satisfies  $\tilde{H}^q(B) = 0$ .

COMMENT 4.4. The above appears as the “correct” formulation for this type of fixed point theorem. Nevertheless, one sees that the crux of the problem lies in the case where  $S$  is an  $n$ -dimensional compact convex set; and there the assumption on  $f$  is only used to ensure that the map  $f^*$  from  $\tilde{H}^n(S, \partial S)$  to  $\tilde{H}^n(X, \partial X)$  is non-zero (and hence that  $f$  is not homotopic to a map to  $\partial S$ ). This assumption is much too stringent for that, and other much more flexible tools are available to prove such things (e.g., homotopy invariance of  $f^*$ ). But for proving existence of (pure strategy) Nash equilibria, this is already a reasonably good tool.

**18. Knaster-Kuratowski-Mazurkiewicz Theorem.** Consider the  $n$ -simplex  $\Delta$  with its faces  $F_i = \{x \in \Delta \mid x_i = 0\}$  ( $i = 0, \dots, n$ ). Let  $C_i$  be closed subsets, with  $F_i \subseteq C_i$ ,  $\bigcup_i C_i = \Delta$ . Then  $\bigcap_i C_i \neq \emptyset$ .

HINT. Let  $f_i(x) = d(x, C_i)$ ,  $f(x) = \sum_0^n f_i(x)$ ,  $\varphi(x) = (f_i(x)/f(x))_{i=0}^n$ :  $\varphi$  is a continuous map from  $\Delta$  to  $\partial\Delta$ , with  $\varphi(F_i) \subseteq F_i$ . View now  $\Delta$  as the unit ball: the map  $x \rightarrow -\varphi(x)$  is a continuous map from  $\Delta$  to itself without fixed point. Apply ex. I.4Ex.17

COMMENT 4.5. Observe that the last part of the proof also shows that the sphere is not a retract of a compact convex set (otherwise  $x \rightarrow -x$  contradicts ex. I.4Ex.17).

**19.** Deduce lemma 1.7 p. 5 from I.4Ex.18.

HINT. Select the extreme point  $e_i$  of  $\Delta$  in  $\bigcap_{j \neq i} T_j$ .

**20.** Also the continuity assumption in theorem 4.1 p. 39 was not the right one — even in the more general framework of ex. I.4Ex.17: it is sufficient to assume for each player  $i$  that his pay-off function  $F^i$  is upper semi-continuous on  $\prod_j S^j$ , and is, for each fixed  $s^i \in S^i$ , continuous in the other variables jointly — and, if one wants a mixed strategy version, like theorem 4.1, and not a pure strategy version like ex. I.4Ex.17, to assume furthermore that  $F^i$  is bounded and the  $S^i$  strictly semi-compact (ex. I.1Ex.7c p. 12), or  $F^i$  continuous in  $s^i$  using ex. I.2Ex.1c p. 21.

HINT. Write shortly  $S$  for  $S^i$ ,  $T$  for  $\prod_{j \neq i} S^j$ ,  $F: S \times T \rightarrow \mathbb{R}$  for  $F^i$ . Let  $R = \{(s, t) \mid F(s, t) = \max_{\tilde{s} \in S} F(\tilde{s}, t)\}$ . Show that  $R$  is closed under the above assumptions. This yields the pure strategy version of ex. I.4Ex.17. For the mixed strategy version, upper semi-continuity and boundedness imply that the mixed extension is well defined. Use then the stated exercises.

**21. Teams.** A team problem (game with finite strategy spaces where all players have the same pay-off function) has finitely many equilibrium pay-offs.

HINT. It suffices to show that there are finitely many pay-offs to completely mixed equilibria. Show that, at such an equilibrium, the gradient of the pay-off function is zero, and use Sard's theorem (e.g. Milnor, 1969, p. 16) to deduce that the set of completely mixed equilibrium pay-offs has Lebesgue measure zero. The semi-algebraicity of this set (cor. 3.4 p. 27) implies it is a finite union of intervals.

**22. Games with pay-off correspondence.** (Simon and Zame, 1990) Consider an  $N$  person game with strategy spaces  $S^i$  and a correspondence  $Q$  from  $S$  to  $\mathbb{R}^N$ . Assume each  $S^i$  compact Hausdorff,  $Q$  non-empty compact convex valued and u.s.c. (3). Prove the existence of a Borel selection  $q$  of  $Q$  such that the game  $(S^i, i = 1, \dots, N; q)$  has an equilibrium.

HINT. For any finite game ( $S_\alpha^i$  finite and included in  $S^i$  and  $q_\alpha$  an arbitrary selection of the graph of  $Q$  above  $S_\alpha$ ) choose equilibrium strategies  $\mu_\alpha = (\mu_\alpha^i)$ . Let  $\mu_\alpha$  and  $q_\alpha d\mu_\alpha$  converge weakly (along some ultrafilter refining the increasing net of all finite subsets  $S_\alpha^i$ ) to  $\mu$  and  $\nu$  and note that  $\nu$  is absolutely continuous w.r.t.  $\mu$ . Let  $q$  be  $\mu$ -measurable with  $qd\mu = d\nu$ . Then  $\mu\{\{s \mid q(s) \notin Q(s)\}\} = 0$ ; otherwise choose  $a, b \in \mathbb{Q}$  and  $x \in \mathbb{Q}^N$  such that  $V = \{s \mid \langle q(s), x \rangle > a > b > \langle t, x \rangle, \forall t \in Q(s)\}$  has positive measure. Use then the regularity of  $\mu$  and Urysohn's lemma to get a contradiction. Change  $q$  on a set of  $\mu$  measure 0 to get a Borel selection and let us still write  $q$  for the mixed extension. Show (as above) that  $T^i = \{s^i \in S^i \mid q^i(s^i, \mu^{-i}) > q^i(\mu)\}$  has  $\mu^i$  measure 0. To get an equilibrium we finally modify  $q$  as follows: let  $p[i]$  be a selection of  $Q$  minimising  $i$ 's pay-off and define

$$\bar{q}(s) = \begin{cases} p[i](s), & \text{if } s^i \in T^i, s^j \notin T^j, j \neq i; \\ q(s), & \text{otherwise.} \end{cases}$$

Note that  $\bar{q}^i(\mu) = q^i(\mu)$  and  $\bar{q}^i(\cdot, \mu^{-i}) = q^i(\cdot, \mu^{-i})$  on  $S^i \setminus T^i$ . Finally on  $T^i$ , since  $Q$  u.s.c. implies  $p[i]^i$  l.s.c. one has:  $\bar{q}^i(t^i, \mu^{-i}) = p[i]^i(t^i, \mu^{-i}) \leq \liminf_\alpha p[i]^i(t_\alpha^i, \mu_\alpha^{-i}) \leq \liminf_\alpha q_\alpha^i(t_\alpha^i, \mu_\alpha^{-i}) \leq \liminf_\alpha q_\alpha^i(\mu_\alpha) \leq q^i(\mu)$ .



## CHAPTER II

### Basic Results on Extensive Form Games

#### 1. The Extensive Form

The extensive form is the most commonly used way to describe or to define a game. We give here a rather general description with some typical variants. The aim is among others to have finite extensive forms, e.g. for repeated games (cf. comment 1.3 below), so as to be able to define them effectively in this way. The flexibility is also needed to describe the solutions of some repeated games with incomplete information as the solutions of auxiliary *finite* games (ch. VIII, sect. 2 p. 361; also ex. VIIIEx.2 p. 393).

##### 1.a. Definitions.

**DEFINITION 1.1.** An **extensive form** game consists of

- (1) a player set  $\mathbf{I} = \{1, \dots, I\}$ ;
- (2) a space of **positions**  $\Omega$ ;
- (3) a **signal** space  $A$  and a function  $\alpha$  from  $\Omega$  to (probabilities on)  $A$ ;
- (4) a partition  $(A^i)_{i \in \mathbf{I}}$  of  $A$ ;
- (5) an **action** space  $S$  and a function  $q$  from  $\Omega \times A \times S$  to (probabilities on)  $\Omega$ ;
- (6) for each  $i$ , a pay-off function  $g^i$  from  $\Omega^\infty$  to  $\mathbb{R}$ ;
- (7) an initial position  $\omega_1$ .

Of course, when dealing with probabilities sub 3 and 5, appropriate precautions will have to be taken, cf. 1.c below. Also when there are no probabilities, the factor  $A$  in 5 is redundant.

Each player  $i \in \mathbf{I}$  has to decide what action  $s \in S$  to choose when he is told  $a$  in  $A^i$ , i.e. to define a function  $\sigma^i$  from  $A^i$  to  $S$ . The game is then played as follows: at stage  $n$ ,  $a_n$  will be chosen according to  $\alpha(\omega_n)$ , and if it belongs to  $A^{i_n}$ , the next position  $\omega_{n+1}$  will be chosen according to  $q(\omega_n, a_n, \sigma^{i_n}(a_n))$ . In other words, if  $a_n$  belongs to  $A^i$ , player  $i$  has the move: in that case, he is told the signal or message  $a_n$  and nothing else: he does not remember his past information or his past choices and not even the stage  $n$  itself, if they are not explicitly included in the message. He selects then an action  $s_n$  as a function of this signal and the new position  $\omega_{n+1}$  is selected using  $q(\omega_n, a_n, s_n)$ . Each player  $i$  receives at the end of the game the pay-off  $g^i(\omega_1, \omega_2, \dots)$ . In some sense, each player  $i$  decides at the beginning of the game on  $\sigma^i$  — “the program of his machine” — and the game then proceeds purely automatically.  $\sigma^i$  is called player  $i$ ’s **pure strategy**. Signals (in the deterministic case) are often identified with their inverse images in the set of positions called **information sets**.

*Some variants.*

- (1) When  $\omega_1$  is chosen at random, add an additional initial state  $\omega_0$ , and extend  $q$  and  $\alpha$  appropriately.
- (2) Often a set of **terminal positions** is needed, where the game stops. This can be formalised by leaving those positions fixed under  $q$ , whatever be the action  $s$

in  $S$ . Conversely, if after a position only one sequence of positions is possible, we can replace it by a terminal position.

- (3) Usually, for any signal  $a$  in  $A$ , only a (non-empty) subset of actions in  $S$  is allowed. This can be taken care of either by allowing in the definition for an  $S$ -valued correspondence on  $A$  and then defining  $q$  only on the relevant set — or by keeping the above definition and extending  $q$  and  $g$  outside this set in a proper way — say by duplication: any additional move is equivalent to some fixed strategy. But the second trick cannot be used when discussing perfect recall and behavioural strategies (cf. below): indeed, if one does not inform the player of which duplicate was used, one loses the perfect recall property, and if one does, one introduces in the new game additional behavioural strategies, which are not behavioural strategies of the original game, but mixed (or “general”) strategies (cf. below for the terminology). Hence in those cases we will write  $\bar{S}$  for the  $S$ -valued correspondence on  $A$ , i.e., its graph in  $A \times S$  (thus:  $\bar{S}_a$  for its value at  $a$ ).  $q$  is then defined on  $\Omega \times \bar{S}$ .
- (4) Another formulation that leads to deterministic signals: take  $\Omega^* = \Omega \times A$  and define  $q^*$  from  $\Omega^* \times S$  to probabilities on  $\Omega^*$  using first  $q$  and then  $\alpha$ .  $\alpha^*$  is then the deterministic projection from  $\Omega^*$  to  $A$ , or simply a sub  $\sigma$ -field  $\mathcal{C}$  on  $\Omega^*$ . The partition  $(A^i)_{i \in \mathbf{I}}$  becomes then a  $\mathcal{C}$ -measurable partition  $\Omega_i^*$  of  $\Omega^*$ , and  $\sigma_i$   $\mathcal{C}$ -measurable on  $\Omega_i^*$ .

Let  $\Omega' = \Omega \times A \times S$  (or  $\Omega \times \bar{S}$  in case of variable action sets), and denote by  $\Omega'_n$  the  $n^{\text{th}}$  copy of  $\Omega'$ . Let  $H_\infty = \prod_1^\infty \Omega'_n$  be the set of **plays**, while the set  $H$  of **histories** is the set of (finite, possibly empty) initial segments of plays.

**REMARK 1.1.** There is no need to allow, like sub 3, for position-dependent sets of possible signals, or similarly for variable sets of possible new positions. Indeed, the transition probabilities  $q$  and  $\alpha$  describe already in some sense what is possible, so one would add a new primitive concept which duplicates to a large extent the information already available in the model. Further, here there is no problem in using the full containing space  $A$  (or  $\Omega$ ), extending the transition probability  $q$  in an arbitrary way when its arguments are not feasible.

The only drawback is that we will in general not have a meaningful definition of a set of **feasible plays**, since it will depend on this “containing space”  $A$ . The set can be defined meaningfully only when all probabilities in the model are purely atomic: then one should define the set as consisting of those plays such that each of their initial segments has positive probability under some strategy vector. Hence, in the general definition below of linear games and of (effectively) perfect recall, we will have to avoid using this concept of feasible play.

**1.b. Finite case.** Here we consider games where there are finitely many different (feasible) plays. As defined above, a **pure strategy** for  $i$  is a mapping from  $A^i$  to  $S$ . A **mixed strategy** will be a probability distribution on those. A **behavioural strategy** is a mapping from  $A^i$  to probability distributions on  $S$ . Finally, a **general strategy** will be a probability distribution on the (compact) set of behavioural strategies. (For the solution of such games, cf. ex. II.1Ex.3 p. 57).

An  $\mathbf{I}$ -tuple of (general) strategies defines a probability distribution on plays. Two strategies of a player are **equivalent** if they induce the same distribution on plays whatever be the strategies of the other players. The **normal form** of the game associates to

every strategy vector the corresponding vector of expected pay-offs (normal form in pure strategies, in mixed —, in behavioural —, in general —).

The above definition of the extensive form does not exclude that a play may produce twice the same message. There could even be two identical positions at two different stages along the same play. We call the game **linear** (Isbell, 1957) (for player  $i$ ) if, a.e. for every strategy vector, every message of player  $i$  occurs at most once along every play. (All the games that we will study in this book will be linear; however, some finite, non-linear games will occur as auxiliary games to solve a class of repeated games with incomplete information (cf. ch. VIII)). The term “linear” is related to the following property:

**THEOREM 1.2. (Isbell, 1957)** *In a linear game (for player  $i$ ), every general strategy of  $i$  is equivalent to a mixed strategy.*

**PROOF.** In fact, given a behavioural strategy  $\tau$  of player  $i$ , we will construct a mixed strategy inducing the same probability on plays, whatever be the other players’ strategies. Let  $\mu$  be defined on each pure strategy  $\sigma$  by:

$$\mu(\sigma) = \Pr_{\mu}(\sigma) = \prod_{a^i \in A^i} \tau(a^i) \{\sigma(a^i)\}.$$

It is then clear that, given any strategies of the other players, the probability of the corresponding sequence of moves in  $S$  will be the same under  $\tau$  and  $\mu$ . Obviously, the same result holds for general strategies, taking the expectation w.r.t. to  $\tau$ . ■

To see the need for the assumption “linear”, cf. ex. II.1Ex.2 p. 56. For a more general case, cf. ex. II.1Ex.10d and II.1Ex.10e p. 63.

The game is said to have **(effectively) perfect recall** for player  $i$  if (knowing the pure strategy he is using), he can deduce from any signal he may get along some feasible play, the sequence of previous messages he got along that play. (All the games we will study will have effectively perfect recall). Observe that a game with effectively perfect recall for  $i$  is necessarily linear for  $i$ . In other words, for having perfect recall, every message must in particular recall both the last message and the last action, while, for having effectively perfect recall, it is sufficient that the player be able to reconstruct the last message, in that case, he can also, knowing his pure strategy, reconstruct his last action, and so on: this shows inductively that he can do as well as if he remembered also his past actions. Formally we have:

**THEOREM 1.3. (Dalkey, 1953)** *In a game with effectively perfect recall for player  $i$ , his pure strategy set is essentially the same (i.e. except for duplication) whether or not he recalls, in addition to his current information, his own past signals and moves.*

**PROOF.** A pure strategy of player  $i$  when he recalls his own past signals and choices is of the form:  $\{s_n = \sigma(a_1, \dots, a_n; s_1, \dots, s_{n-1})\}$  for  $n \geq 1$ , where  $(a_1, \dots, a_n)$  denotes the sequence of signals from  $A^i$  he has heard,  $(s_1, \dots, s_{n-1})$  is the sequence of his past choices in  $S^i$  and  $s_n$  in  $S^i$  is the move to be chosen. Define, by induction on the number  $n$  of signals previously heard, a strategy  $\phi$  that depends only on the current signal: roughly, for each initial signal  $a$  in  $A^i$ ,  $s_1 = \phi(a) = \sigma(a)$ . Inductively, given a signal  $a$ , deduce from  $a$  and  $\phi$  the previous signals, say  $a_1, \dots, a_n$ , and let  $\phi(a) = \sigma(a_1, \dots, a_n, a; \phi(a_1), \dots, \phi(a_n))$ . Whatever be the strategy of the other players, the pure strategy  $\phi$  results in the same probability distribution on plays as the pure strategy  $\sigma$ , hence the proof is complete. ■

For more details and/or more generality, cf. ex. II.1Ex.14 p. 72.

COMMENT 1.2. Observe that, since the set of strategies  $\phi$  is clearly smaller than the set of strategies  $\sigma$  by our construction, several strategies  $\sigma$  will be mapped to the same strategy  $\phi$ . This is the duplication of strategies mentioned above.

COMMENT 1.3. From our previous definitions, it appears that, for a game with effectively perfect recall, we need a very large space of positions, essentially  $H$ . But it follows from the above that, to specify a game with effectively perfect recall, it is sufficient to describe each player's incremental information every time he has to play. This can be arbitrary, in particular does not have to remind him his last move. This is the way we will describe all our models in this book from ch. IV on: the signal  $a \in A^i$  will be player  $i$ 's incremental information. (I.e., his true signals, in the previous general model, are finite sequences in  $A^i$ .) Therefore we will still be able to use small (finite) sets of positions.

The main advantage of games with perfect recall is the following.

**THEOREM 1.4. (Kuhn, 1953)** *In a game with perfect recall for player  $i$ , general, mixed, and behavioural strategies are equivalent for  $i$ .*

**PROOF.** By theorem 1.2 p. 53, it is enough to represent any mixed strategy  $\mu$  by a behavioural strategy  $\tau$ . In fact, given  $\mu$ , probability distribution on pure strategies, compute first, for every initial signal  $a_1$  (cf. ex. II.1Ex.8 p. 59), the marginal distribution  $\mu_{a_1}$  on "strategies after  $a_1$ ", compute then from  $\mu_{a_1}$  the marginal distribution  $\tau_{a_1}$  of  $s_1$ , and the conditional distribution  $\mu_{a_1, s_1}$  on "strategies for the future" given  $s_1$ . Continue then with  $\mu_{a_1, s_1}$  as before with  $\mu$ . ■

As a corollary, by Dalkey's theorem, one can also study games with effectively perfect recall by using behavioural strategies in the corresponding game with perfect recall.

A game is said to have **perfect information** if the position at each stage in a feasible play determines all the previous positions along the play and moreover the signal determines the position. In other words, when he has to play, each player knows the whole past history.

**THEOREM 1.5. (Zermelo, 1913)** *In a game with perfect information, there exists an equilibrium in pure strategies.*

**PROOF.** The proof is very simple using **backwards induction**: Replace each position that is followed only by terminal positions, by a terminal position with, as pay-off, the pay-off corresponding to an optimal choice of action of the player having the move at that position. ■

**1.c. A measurable set up.** We consider now a game where the spaces  $\Omega$ ,  $A$  and  $S$  are measurable as well as the partition  $(A^i)_{i \in \mathbf{I}}$ , the graph  $\bar{S}$  and the map  $g$ ;  $\alpha$  and  $q$  are transition probabilities. A pure strategy, say for player  $i$ , is then a measurable selection from  $\bar{S}$  defined on  $(A^i, \mathcal{A}^i)$ . Similarly, a behavioural strategy will be a transition probability from  $(A^i, \mathcal{A}^i)$  to  $(S, \mathcal{S})$ , assigning probability one to  $\bar{S}$ . Note that there is no adequate measurable structure on those sets, hence it is more appropriate to define a mixed strategy as a measurable selection from  $S$  defined on  $(X^i \times A^i, \mathcal{X}^i \otimes \mathcal{A}^i)$  where  $(X^i, \mathcal{X}^i, Q^i)$  is an auxiliary probability space. One similarly defines general strategies as transition probabilities from  $(X^i \times A^i, \mathcal{X}^i \otimes \mathcal{A}^i)$  to  $(S, \mathcal{S})$  assigning probability one to  $\bar{S}$ .

**PROPOSITION 1.6.** *An  $\mathbf{I}$ -tuple of general strategies induces a unique probability distribution on plays.*

PROOF. Use Ionescu Tulcea's theorem (Neveu, 1970, prop. V.1.1) with  $(\Omega \times A \times S)^\infty$ , then show that  $(\Omega \times \bar{S})^\infty$  has probability one. ■

Hence, the normal form will be well defined (so one can speak of equilibria) e.g. as soon as  $g$  is bounded from below or from above. And, since for each  $x \in \prod_i X^i$  we also have a vector of strategies, the corresponding pay-off  $g(x^1, x^2, \dots)$  is also well defined, and, by the cited theorem, it will be a measurable function on  $\prod_i X^i$ , with  $g(Q^1, Q^2, \dots) = \int g(x^1, x^2, \dots) \prod_{i \in \mathbb{I}} dQ^i$ .

It also follows that the definition of the equivalence of strategies extends to this case. In particular, if general strategies are shown to be equivalent to mixed strategies (cf. ex. II.1Ex.10d p. 63) a strategy vector is an equilibrium iff no player has a profitable pure strategy deviation.

**DEFINITION 1.7.** A game has (**effectively**) **perfect recall** for player  $i$  if, (for every pure strategy of  $i$ ), there exists a measurable map  $\varphi$  from  $A^i$  to  $(A^i \times S^i) \cup \{\iota\}$  such that, for every strategy vector of  $i$ 's opponents a.s.:  $\varphi(a) = \iota$  means that  $a$  is the first signal to this player, otherwise  $\varphi(a)$  is the previous signal to him and the action he took upon it. (For a justification of this definition, cf. ex. II.1Ex.12 p. 64).

**THEOREM 1.8.** Assume player  $i$  has perfect recall and  $(S, \mathcal{S})$  is standard Borel (cf. App.6). Every general strategy of player  $i$  is equivalent to a mixed strategy and to a behavioural strategy.

PROOF. To prove the inclusion of general strategies in the set of mixed strategies, represent  $(S, \mathcal{S})$  as  $[0, 1]$  with the Borel sets; then a general strategy yields a family of cumulative distribution functions  $F_{x,a}(s)$  jointly measurable in  $x, a$  and  $s$  (ex. 1 p. 61). The perfect recall assumption implies there is a measurable map  $n(a)$  (ex. II.1Ex.10 p. 61) from  $A_i$  to  $\mathbb{N}$  that increases strictly along every play. Let then  $u_n$  be an independent sequence of uniform random variables on  $[0, 1]$ , and define  $\mu(u, x, a) = \min\{s \mid F_{x,a}(s) \geq u_{n(a)}\}$ :  $\mu$  is a mixed strategy with as auxiliary space the product  $(X^i, \mathcal{X}^i, Q^i) \times ([0, 1], \mathcal{B}, \lambda)^{\mathbb{N}}$ , where  $\lambda$  denotes Lebesgue measure on the Borel sets  $\mathcal{B}$  of  $[0, 1]$ . And  $\Pr(\mu \leq s \mid x, a, \text{past}) = \Pr(u_{n(a)} \leq F_{x,a}(s) \mid x, a, \text{past}) = F_{x,a}(s)$ , so  $\mu$  induces clearly the same probability distribution on plays as the general strategy.

The proof of the second part is given in ex. II.1Ex.10 p. 61. ■

Actually, the above is the proof of a different statement, cf. ex. II.1Ex.10d. Ex. II.1Ex.10a and II.1Ex.10b give a full proof of theorem 1.8; part II.1Ex.10a being there only to justify the “clearly” in the last sentence of the proof above.

**COMMENT 1.4.** When  $A^i$  is countable (cf. also ex. II.1Ex.14, comment 1.28 p. 73), a behavioural strategy can be identified with a point in  $[S^*]^{A^i}$ , denoting by  $S^*$  the set of probability distributions on  $S$ . The  $\sigma$ -field  $\mathcal{S}$  induces a natural  $\sigma$ -field on  $S^*$  (requiring that the probability of every measurable set be a measurable function on  $S^*$ ), and hence on  $[S^*]^{A^i}$ . Thus general (resp. mixed) strategies are naturally defined as probability distributions on  $[S^*]^{A^i}$ , resp.  $S^{A^i}$ . Those definitions are clearly equivalent to the former (without any standard Borel restrictions): given a probability distribution  $P$  on  $[S^*]^{A^i}$ , use  $([S^*]^{A^i}, [\mathcal{S}^*]^{A^i}, P)$  as  $(X^i, \mathcal{X}^i, Q^i)$  — this also reduces to the former case how to define the probability distributions induced on plays by those new strategies —; conversely, given  $(X^i, \mathcal{X}^i, Q^i)$  there is an obvious measurable map to  $([S^*]^{A^i}, [\mathcal{S}^*]^{A^i})$  defining a  $P$  (and similarly for mixed strategies). It is clear that this transformation preserves the probability

distribution induced on plays. This probability distribution is clearly the average under  $P$  of the distributions induced by the underlying behavioural (resp. pure) strategies. Thus we are led back — at least for linear games (cf. ex. II.1Ex.10c p. 63) — to the general concept (ch. I) of a mixed strategy as a probability distribution over pure strategies, and the corresponding normal forms.

COMMENT 1.5. When  $A$  is countable, and topologies on  $\Omega$  and  $S$  (often: discrete topologies) are given for which  $q$  and  $\alpha$  are continuous, the induced mapping from pure strategies ( $S^A$  with the product topology) to plays (product topology) will be continuous too, so that continuity properties of  $g$  translate immediately into the same (joint) continuity properties of the normal form pay-off function (and compactness of  $S$  yields compactness of  $S^A$ ).

### Exercises.

#### 1.

- a. Show that a game as defined in 1.a p. 51 can also be represented by a deterministic one — i.e.  $\alpha$  and  $q$  functions — by adding nature as player zero (which uses a fixed behavioural strategy).

COMMENT 1.6. This modification is frequently useful in specific situations, because it allows to have a smaller and more manageable set of (feasible) plays, by giving variable actions sets to nature as well as to the other players.

b. Observe that, in such a game with moves of nature, one can always redraw the tree such that the game starts with a move of nature, and is deterministic afterwards, i.e. the only randomness is in the selection of the initial position. (Use theorem 1.4 p. 54 — or its generalisation sub ex. II.1Ex.10 p. 61 below — for player zero, getting a mixed strategy for him.)

c. Similarly, one can also redraw the tree such as to let nature postpone as much as possible its choices (thus histories of the form  $(a_1, s_1, a_2, s_2, a_3, s_3, \dots)$ ), by obtaining this time a behavioural strategy for player zero.

COMMENT 1.7. This transformation is often important: e.g. in ch. VIII, it will reduce a class of games with symmetric information to stochastic games.

2. (Isbell, 1957) In the two-person zero-sum game of Figure 1, show that:

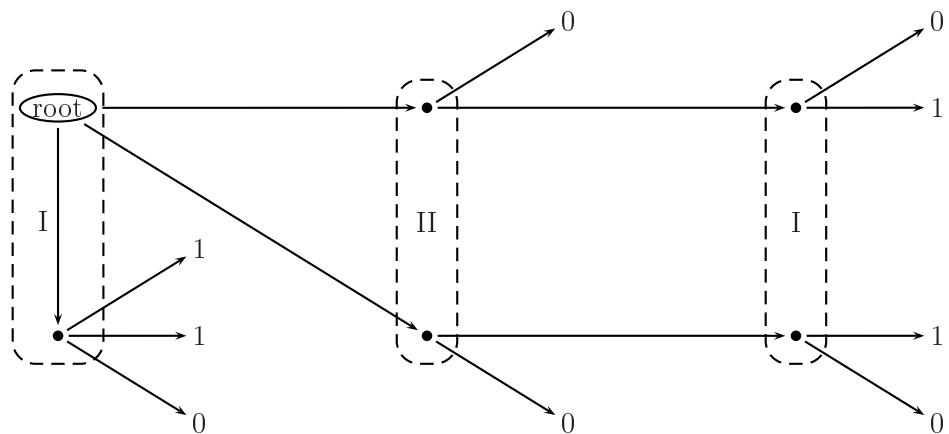


FIGURE 1. A non-linear Game

- by randomising  $(1/2, 1/2)$  between  $(3/4, 0, 1/4; 0, 1)$  and  $(0, 3/4, 1/4; 1, 0)$  player I can guarantee  $9/16$ .

- he cannot guarantee more than  $1/2$  with mixed strategies.
- he cannot guarantee more than  $25/64$  with behavioural strategies.

### 3. Polynomial pay-off functions.

a. Prove that, in a game with finitely many different feasible plays, the pay-off is polynomial in behavioural strategies.

b. Show that any polynomial pay-off function on a product of compact, convex strategy polyhedra can be obtained in this way.

HINT. Reduce first to the case of strategy simplices. Show then how to construct a sum, and reduce thus the problem to the case of the same pay-off function to all players, which is a single monomial.

See ex. I.4Ex.8 p. 41 to show that such games have “finite” solutions.

### 4. An $n$ -dimensional moment problem.

a. Let  $P(x)$  be a polynomial, with  $P(x) > 0$  for  $x > 0$ . Show that, for  $n$  sufficiently large, all coefficients of  $(1+x)^n P(x)$  are positive.

HINT. The binomial formula yields the  $n$ -step distribution of a random walk. Use the corresponding asymptotic approximations. In a more elementary fashion, factor first  $P$  into linear and quadratic factors of the form  $x+q$  and  $(x-p)^2+q$ , with  $q > 0$ , so as to reduce the problem to  $P(x) = (x-p)^2+q$ ,  $p > 0$ ,  $q > 0$ , which can then be handled in a fully elementary way, remaining within the realm of second degree equations.

b. Let  $P(x_1, y_1; x_2, y_2; \dots)$  be a polynomial, homogeneous in each pair  $(x_i, y_i)$  ( $1 \leq i \leq k$ ). Assume  $P > 0$  whenever  $x_i \geq 0, y_i \geq 0, x_i + y_i > 0$  for all  $i$ . Show that, for  $n$  sufficiently large, all coefficients of  $[\prod_i (x_i + y_i)^n] P(x, y)$  are positive.

HINT. Induction on  $k$ . For fixed values  $x_i \geq 0, y_i \geq 0$  for  $i \geq 2$  ( $x_i + y_i > 0$ ), apply II.1Ex.4a to the resulting polynomial in  $x_1/y_1$ . By continuity of the coefficients, they will still be positive in the neighbourhood of  $z = \{ (x_i, y_i) \mid i \geq 2 \}$ ; use compactness.

c. Let the polynomial  $P$  be positive on  $\Delta_k = \{ x \in \mathbb{R}_+^{k+1} \mid \sum x_i = 1 \}$ . Then there is a homogeneous polynomial  $T$  of degree  $d$ , where every possible monomial of degree  $d$  has a positive coefficient, and with  $P(x) = T(x)$  for  $x \in \Delta_k$ . If  $P$  is homogeneous, then  $T = [\sum_i x_i]^n P$ .

HINT. Make first  $P$  homogeneous by multiplying each monomial by the appropriate power of  $\sum x_i$ ; add then  $\varepsilon[1 - \sum_i x_i]$  for  $\varepsilon > 0$  sufficiently small (compactness), such as to have  $P(x) > 0$  on the unit cube. Obtain now  $Q(x, y)$  by replacing in  $P(x)$  every  $x_i^h$  by  $x_i^h (x_i + y_i)^{d_i-h}$ , where  $d_i$  is the maximal power of  $x_i$  in  $P$ : for  $x \in \Delta_k$ ,  $y_i = \sum_{j \neq i} x_j$ ,  $Q(x, y)$  still equals the original  $P(x)$ . Apply II.1Ex.4b to  $Q(x, y)$ , and let  $R(x, y) = [\prod_i (x_i + y_i)^n] Q(x, y)$ : we still have  $R = P$  on  $\Delta_k$ , and all terms of  $R$  are non-negative, hence this remains so when replacing all  $y_i$  by  $\sum_{j \neq i} x_j$  and expanding to obtain  $T(x)$ , which is clearly homogeneous, say of degree  $d$ . And the coefficients of  $x_i^d$  are necessarily strictly positive. So if some other coefficients are still zero, multiplying  $T$  by an appropriate power  $[(k-1)d]$  of  $\sum x_i$  will make them all positive. For the last sentence, set  $n$  equal to the difference in degrees: the two polynomials are then equal on  $\mathbb{R}_+^{k+1}$ , hence are identical.

d. Let  $K$  be a compact convex polyhedron in  $\mathbb{R}^n$  (cf. ex. I.3Ex.4 p. 29) defined by the inequalities  $f_i(x) \geq 0$  ( $i = 1, \dots, k$ ), where the  $f_i$  are affine functionals. Any positive polynomial function  $P$  on  $K$  can be written as a sum with positive coefficients of all monomials of some degree  $d$  in  $f_1, \dots, f_k$ .

HINT. If  $K$  is a singleton, take  $d = 0$ . Since  $K$  is compact, the  $f_i$  separate points. Let  $f = (f_i)_{i=1}^k: \mathbb{R}^n \rightarrow \mathbb{R}^k$ :  $f(K)$  is the intersection of  $\mathbb{R}_+^k$  with an affine subspace, say with equations  $\varphi_i(y) = 0$  ( $i = 0, \dots, h$ ). The  $\varphi_i$  are not all linear, otherwise  $f(K)$  would be unbounded. Assume thus  $\varphi_0(y) = 0$  is of the form  $\sum \mu_i y_i = 1$ . Using this, one can make all other  $\varphi_i$  ( $i = 1, \dots, \ell$ ) linear. Compactness of  $f(K)$  implies it can be separated from  $\sum y_i \geq M$ , which implies (e.g. ex. I.3Ex.1 p. 28) that adding some linear combination of the  $\varphi_j$  ( $j \geq 1$ ) to  $\varphi_0$ , one can assume  $\mu_i > 0$ ,  $i = 1, \dots, k$ . Since the  $f_i$  separate points,  $P$  can be rewritten as a polynomial  $Q$  in the  $f_i(x)$ :  $Q$  is a polynomial on  $\mathbb{R}^k$ ,  $> 0$  on

$f(K)$ . Adding a polynomial  $L \sum_{i \geq 1} [\varphi_i(y)]^2$  for  $L$  sufficiently large will not affect the values on  $f(K)$ , but will make  $Q > 0$  on  $\{y \in \mathbb{R}_+^k \mid \sum \mu_i y_i = 1\}$ . Apply now II.1Ex.4c.

e. Same statement as sub II.1Ex.4d, but if  $K$  is a product of compact polyhedra  $K_j$ , one can use a sum with positive coefficients of all monomials which are separately for each  $j$  of degree  $d_j$  in the  $f_i^j$ .

HINT. Remember from the proof of II.1Ex.4d that for each  $K_j$  ( $\#K_j > 1$ ) we have a relation  $\sum_i \mu_i^j f_i^j = 1$  with  $\mu_i^j > 0$ . Make sure that indeed all possible monomials have strictly positive coefficient.

COMMENT 1.8. Observe those results remain true in some sense for arbitrary finite dimensional compact convex sets: the polynomial function is still strictly positive in some neighbourhood, squeeze a compact convex polyhedron (or a product) in this neighbourhood, containing the given compact set in its interior. One obtains thus a similar representation, with a finite number of affine functionals which are strictly positive on the given compact convex set. Hence it also extends to compact convex sets in topological vector spaces, defining the polynomial functions as the algebra generated by the continuous affine functionals. In particular, if  $C$  is a compact convex set in a Hausdorff locally convex space  $E$ , then any linear functional on the polynomial functions on  $E$  which is non-negative on all finite products of continuous affine functionals which are positive on  $C$ , is a non-negative Radon measure on  $C$ . [! A little additional exercise has to be made, for  $E = \mathbb{R}^n$ , to be sure to prove the result for polynomial functions on  $E$  and not just on  $C$ .]

COMMENT 1.9. It would be interesting to know, even for applications like sub II.1Ex.6 below, what polynomial functions have such a representation when asking only for non-negative coefficients.

**5. Rational pay-off functions.** Consider a finite game — i.e.  $\Omega$ ,  $A$  and  $S$  are finite — with a privileged position  $\omega_0$ , from where play remains in  $\omega_0$ . Assume that any pure strategy vector, and starting from any position, has a positive probability of eventually reaching  $\omega_0$ .

a. Prove that the expected time for reaching  $\omega_0$  is uniformly bounded, over all behavioural (hence also general) strategies (and all possible starting points).

HINT. By finiteness, there exists  $\varepsilon > 0$ , and  $n_0 < \infty$  such that for any pure strategy vector, and any starting point,  $\omega_0$  is reached before  $n_0$  with probability  $> \varepsilon$ . Consider a behavioural strategy vector: by finiteness of the action space  $S$ , there is for each  $a \in A$  an action with probability  $\geq (\#S)^{-1}$ . The probability that the corresponding pure strategy vector be played for the first  $n_0$  stages is therefore  $\geq (\#S)^{-n_0}$ . Thus the probability of reaching  $\omega_0$  in that time is  $\geq \varepsilon (\#S)^{-n_0} = \delta > 0$ , whatever be the starting point and the behavioural strategy vector.

b. Assume a pay-off function is given over  $\Omega$ , zero at  $\omega_0$ , and that the pay-off of the game on  $\Omega^\infty$  is obtained as the sum of the stage pay-offs. Prove that the pay-off is a rational function of the behavioural strategies, with a denominator that vanishes nowhere.

HINT. Denote by  $P$  the transition probability of the resulting Markov chain on  $\Omega$ , by  $u(\omega)$  the pay-off function on  $\Omega$ , and let  $V_\omega$  denote the expected pay-off starting from  $\omega$ . Then  $V = u + PV$  ( $u, V$  column vectors). The equation and the column of  $P$  corresponding to  $\omega_0$  can be deleted, since  $u(\omega_0) = V(\omega_0) = P_{\omega_0}(\mathbb{C}\{\omega_0\}) = 0$ . II.1Ex.5a implies then that the series  $\sum P^n$  is summable, so  $I - P$  has non-zero determinant.

**6. Rational pay-off functions. The converse.** Conversely, show that any normal form game, where each player  $i$  has a product of compact convex polyhedra  $(K_a^i)_{a \in A^i}$  as pure strategy space, and where the pay-off function  $F$  is rational with non-vanishing denominator, arises from some game as sub II.1Ex.5, where  $A^i$  is the set of  $i$ 's signals, where his available actions at  $a \in A^i$  are the extreme points  $v$  of  $K_a^i$ , and where the pay-off to any behavioural strategy vector  $\lambda_v^{i,a}$  ( $\lambda_v^{i,a} \geq 0$ ,  $\sum_v \lambda_v^{i,a} = 1$ ) is given by  $F[(\sum_v \lambda_v^{i,a} \cdot v)_{a \in A^i, i \in I}]$ .

HINT. Assume w.l.o.g. that  $\#K_a^i > 1$  for all  $(i, a)$ , and that the corresponding set of affine constraints  $f_c^{i,a}(x) \geq 0$  satisfies — multiplying each of them by the corresponding  $\mu_c$  (cf. II.1Ex.4d) —, that  $\sum_c f_c^{i,a}(x) = 1$  identically. Denote by  $C^{i,a}$  the set of constraints  $f_c^{i,a}$  of  $K^{i,a}$ , and by  $V^{i,a}$  its set of extreme points.

Use a common denominator  $P$  for all players' pay-off functions. Since it does not vanish, it can be assumed strictly positive, hence in the form sub II.1Ex.4e above. Further, multiplying it — and the numerators —, by a positive constant, we can assume the coefficient of each term is strictly smaller than the coefficient of the corresponding term in the expansion of  $[\prod_{a \in A^i, i \in I} (\sum_{c \in C^{i,a}} f_c^{i,a})]^d = Q(f)$ . Thus this denominator can be written as  $1 - (Q(f) - P(f)) = 1 - R(f)$ , where  $R(f) = \sum_m p_m m(f)$ ,  $m$  runs through all monomials of  $Q(f)$ , and  $0 < p_m < 1$  for all  $m$ . Let  $f_c^{i,a}(v) = q_v^{i,a}(c) : q_v^{i,a}(c) \geq 0$ , and  $\sum_{c \in C^{i,a}} q_v^{i,a}(c) = 1$ . Then, for a mixed action  $\lambda \in \Delta(V^{i,a})$  we have  $f^{i,a}(\lambda) = \sum_{v \in V^{i,a}} \lambda(v) q_v^{i,a} (= q_\lambda^{i,a})$  is a probability distribution on  $C^{i,a}$ , and any monomial  $m$  has the form  $B_n \prod_{i,a,c} [q_\lambda(c)]^{n_c^{i,a}} = m_n(\lambda)$  where  $n = (n_c^{i,a})$  satisfies  $n_c^{i,a} \geq 0$ ,  $\sum_{c \in C^{i,a}} n_c^{i,a} = d$ , and when  $B_n = \prod_{i,a} B_n^{i,a}$ ,  $B_n^{i,a}$  denoting the number of distinct orderings of a set containing  $n_c^{i,a}$  objects of type  $c$  for all  $c \in C^{i,a}$ .

Represent also the numerators as  $\sum_n u_m m(f)$ , with  $u_m \in \mathbb{R}_n^I$  — if necessary by increasing  $d$ . Fix an ordering on  $I$ , and one on each  $A^i$ , obtaining thus an ordering on  $A = \bigcup_{i \in I} A^i$ . Take the successive signals  $a$  in this order, repeating the whole sequence  $d$  times. Every choice of  $v \in V^{i,a}$  is followed by a move of nature selecting  $c \in C^{i,a}$  with probability  $q_v^{i,a}(c)$ . At the end of those  $d \cdot (\#A)$  stages, count as outcome the number of times  $n_c$  each  $c$  has occurred — defining a monomial  $m_n$ . It is clear that, with behavioural strategies  $\lambda^{i,a}$ ,  $\Pr(m_n) = m_n(\lambda)$ . At outcome  $m_n$ , give  $u_{m_n}$  as pay-off (zero pay-off at all previous positions) and use  $p_{m_n}$  as probability of returning to the origin  $\omega_1$ ,  $1 - p_{m_n}$  as probability of going to the cemetery  $\omega_0$ .

COMMENT 1.10. Observe that, since each period has a fixed number  $d \cdot (\#A)$  of stages, and since  $p_m < 1$  for all  $m$ , one could as well — correcting the  $p_m$ 's — view this as a discounted game with small discount factor.

COMMENT 1.11. Observe also that behavioural strategies  $\lambda^{i,a}$  have no influence on history except through their image in  $K^{i,a}$ .

COMMENT 1.12. This provides an extensive form for the game in ex. I.4Ex.9 p. 42.

## 7. Linear games.

a. Consider a game as sub ex. II.1Ex.6, where every return is only to the origin  $\omega_1$ , and where in addition every information set is met only once in between two returns. Then the pay-off is a ratio of multilinear functions in the behavioural strategies. Conversely, every ratio of multilinear functions on a product of simplices, where the denominator never vanishes, is obtained in this way. (It is not clear whether this extends — as in ex. II.1Ex.6 — to compact convex polyhedra instead of simplices.)

b. Any game with compact convex strategy spaces and continuous pay-off functions which are ratios of multilinear functions has a pure strategy equilibrium.

HINT. Prove quasi-concavity and use ex. I.4Ex.20 p. 48.

c. The assumption sub II.1Ex.7b is much stronger than needed for the argument: it would suffice e.g. that each player's pay-off function be the ratio of a numerator which is concave in his own strategy and of a positive denominator which is linear in his own strategy, and even this linearity can be weakened to concavity (resp. convexity) if the numerator is  $\leq 0$  (resp. if this player's minmax value is  $\geq 0$ ).

d. The above yields “equilibria in behavioural strategies” for many games, even without the assumption of perfect recall. Observe however that this assumption is needed to conclude that those are indeed equilibria, via a short dynamic programming argument in each player's tree (cf. ex. II.1Ex.8 below).

## 8. The natural tree structure.

a. One can define the sets of (“feasible”) plays  $H_\infty$  (and hence histories  $H$ ) such that, if the game has perfect recall for player  $i$ , his information sets have a natural tree structure (and such that the “a.s.” clause in the definition of perfect recall becomes superfluous).

b. For an arbitrary game, define the tree of the game by adding an outside observer — a dummy player — who is told (and remembers) everything that happens in the game. Consider his tree. This is  $H$ , with its natural partial order.

**9. Conditional probabilities.** Assume  $(E, \mathcal{E})$ ,  $(X, \mathcal{X})$  and  $(Y, \mathcal{Y})$  are measurable spaces, with  $\mathcal{Y}$  separable and separating, and with  $(X, \mathcal{X})$  standard Borel (cf. App.6). Consider a transition probability  $P$  from  $E$  to  $X$  (denoted  $P_e(dx)$ ), and a measurable function  $g$  from  $E \times X$  to  $Y$ .

a. There exists a transition probability  $Q$  from  $E \times Y$  to  $X$  (denoted  $Q(dx | e, y)$ ) such that, for each  $e \in E$ , and for every positive measurable function  $f$  on  $X$ ,  $\int f(x)Q(dx | e, y)$  composed with  $g$ , is a version of the conditional expectation  $\mathbb{E}_{P_e}(f | \mathcal{F}_e)$  of  $f$  under  $P_e$  given the  $\sigma$ -field  $\mathcal{F}_e$  spanned by  $g(e, \cdot)$ .

HINT. Show first that, for each bounded measurable function  $f$  on  $X$ , there exists a measurable function  $\tilde{f}$  on  $E \times Y$  such that  $\tilde{f} \circ g(e, \cdot) = \mathbb{E}_{P_e}(f | \mathcal{F}_e)(\cdot)$   $P_e$ -a.e.,  $\forall e \in E$ . To this effect, let  $\mathcal{Y}_n$  be an increasing sequence of finite measurable partitions of  $Y$  that spans  $\mathcal{Y}$ , and let  $\mathcal{F}_e^n$ ,  $\tilde{f}^n(e, y)$  be associated with  $\mathcal{Y}_n$  instead of  $\mathcal{Y}$ . Since  $\mathcal{F}_e^n$  increases to  $\mathcal{F}_e$ , the martingale convergence theorem implies  $\tilde{f}(e, y) = \liminf_{n \rightarrow \infty} \tilde{f}^n(e, y)$  will do, provided we know existence of the  $\tilde{f}^n$ : this reduces the problem to the case where  $Y = \{y_1, \dots, y_k\}$  is finite, corresponding to a measurable partition  $B_1, \dots, B_k$  of  $E \times X$ . Set then  $\tilde{f}(e, y) = \sum_{i=1}^k \mathbb{1}_{y=y_i} \int_{B_i} f P_e(dx) / P_e(B_i)$  (with  $0/0 = 0$ ). Measurability of  $\tilde{f}$  will then follow (by composition) if we know that  $\int h(e, x) P_e(dx)$  is measurable for any bounded measurable function  $h$  on  $E \times X$ , cf. (hint of) II.1Ex.9b for this.

Identify now  $(X, \mathcal{X})$  with  $[0, 1]$  and the Borel sets, and select for each rational  $r \in [0, 1]$ ,  $\tilde{f}_r \geq 0$  as above, using  $\mathbb{1}_{[0,r]}$  for  $f$  (and  $\tilde{f}_1 = 1$ ). Let  $F(x, e, y) = \inf_{r > x} \tilde{f}_r(e, y)$ . Clearly  $F$  is measurable,  $0 \leq F \leq 1$ , and for each  $(e, y)$ ,  $F$  is monotone and right continuous in  $x$ : this defines then the transition probability  $Q(dx | e, y)$ . Show that  $\int f(x)Q(dx | e, y)$  composed with  $g$  is indeed a version of  $\mathbb{E}_{P_e}(f | \mathcal{F}_e)(y)$  first for  $f = \mathbb{1}_{[0,\alpha]}$ , then for indicators of finite unions of intervals, then by a monotone class argument for indicators of Borel sets, then for any positive Borel function.

b. For such a  $Q$ , and for a positive measurable function  $f$  on  $X \times E \times Y$ ,  $\int f(x, e, y) Q(dx | e, y)$  is well defined, and measurable on  $E \times Y$ , and one has

$$\int \left[ \int f(x, e, y) Q(dx | e, y) \right] [(P_e \circ g_e^{-1})(dy)] = \int f(e, x, g(e, x)) P_e(dx)$$

where  $P_e \circ g_e^{-1}$  is the distribution of  $g(e, x)$  on  $Y$ ,  $x$  being distributed according to  $P_e$ , and is itself a transition probability.

HINT. Establish the result first for  $f(e, x, y) = \mathbb{1}_{B(x)} \mathbb{1}_{C(e)} \mathbb{1}_{D(y)}$ , then for finite unions of such rectangles, then (monotone class) for measurable indicator functions, then for all positive measurable functions.

c.  $\{(e, y) \mid Q(\{x \mid g(e, x) = y\} \mid e, y) = 1\}$  is measurable, and has for each  $e$  probability one under  $P_e \circ g_e^{-1}$ .

HINT. Show first that the graph  $G$  of  $g$  is measurable: with the same partitions  $\mathcal{Y}_n$  as sub II.1Ex.9a, the  $G_n = \bigcup_{A \in \mathcal{Y}_n} A \times g^{-1}(A)$  are measurable and decrease to  $G$  (recall  $\mathcal{Y}$  is separating). Apply then II.1Ex.9b to the indicator function  $f$  of  $G$ .

COMMENT 1.13. We also proved along the way the following two lemmas, the first at the end of II.1Ex.9a and the second as a restatement of II.1Ex.9b:

- (1) If  $F(e, x)$  is, for each  $e \in E$ , monotone and right continuous in  $x$ , with  $0 \leq F(e, x) \leq F(e, 1) = 1$ , then measurability of  $F$  on  $(E, \mathcal{E})$  for each fixed  $x$  is equivalent to joint measurability of  $F$ , and is still equivalent to  $F$  defining a transition probability from  $(E, \mathcal{E})$  to  $[0, 1]$ .
- (2) A transition probability  $P_e$  from  $(E, \mathcal{E})$  to  $(F, \mathcal{F})$  can equivalently be viewed as a transition probability from  $(E, \mathcal{E})$  to  $(E \times F, \mathcal{E} \otimes \mathcal{F})$ . A measurable function induces a transition probability. The composition of transition probabilities is a transition probability. The composition of a transition probability with a measurable function equals its composition with the induced transition probability.

**10. Sufficiency of mixed and behavioural strategies.** We want to prove theorem 1.8 p. 55.

a.

i. The perfect recall assumption allows in particular to compute a measurable function  $n(a)$ , with  $n(\iota) = 0$ ,  $n(a) = 1 + n(\varphi(a))$ :  $n$  is the “stage for player  $i$ ”.

Let  $(S_n, \mathcal{S}_n)$  be the  $n^{\text{th}}$  copy of  $i$ 's action space  $S^i$ . Denote by  $(A_k, \mathcal{A}_k)$  the subset of  $A^i$  where  $n(a) = k$ , and by  $\varphi_n: A_n \rightarrow A_{n-1} \times S_{n-1}$  the restriction of  $\varphi$ . Allow  $i$ 's behavioural strategies  $P_n$ , as probabilities on  $(S_n, \mathcal{S}_n)$ , to depend not only on  $(A_n, \mathcal{A}_n)$ , but also on the whole past  $\prod_{t < n}(A_t \times S_t, \mathcal{A}_t \times \mathcal{S}_t)$ , as well as on the auxiliary probability space  $(X, \mathcal{X}, P)$  in case of general strategies. At the end, one can always use the maps  $\varphi_n$  to rewrite them if so desired only as functions of  $A_n$  (with an arbitrary probability on  $S$  outside  $\bigcup_n A_n$ ); but in this way, we can forget the maps  $\varphi_n$ .

This class of strategies of player  $i$  is not the most general one, there is no reason not to allow him to use an auxiliary probability space at every stage. Define thus a generalised strategy of player  $i$  as a sequence  $(X_n, \mathcal{X}_n; P_{n-1})_{n=1}^\infty$ , where  $(X_n, \mathcal{X}_n)$  is a measurable space, and  $P_n$  a transition probability from  $\prod_{1 \leq t \leq n}[(S_{t-1}, \mathcal{S}_{t-1}) \times (X_t, \mathcal{X}_t) \times (A_t, \mathcal{A}_t)]$  to  $(S_n, \mathcal{S}_n) \times (X_{n+1}, \mathcal{X}_{n+1})$  ( $S_0 = \{0\}$ ).

COMMENT 1.14. The general strategies correspond then to the case where  $X_n = \{0\}$  for  $n > 1$ , behavioural strategies to  $X_n = \{0\}$ ,  $\forall n$ ; and mixed strategies by definition to general strategies where  $(X_1, \mathcal{X}_1, P_0)$  is the unit interval with the Borel sets and Lebesgue measure, and where the transition probabilities become measurable functions  $\sigma_n$  from  $(X_1, \mathcal{X}_1) \times \prod_{t \leq n}(A_t, \mathcal{A}_t)$  to  $(S_n, \mathcal{S}_n)$ .

COMMENT 1.15. The correspondence we will establish between general(ised) strategies, mixed strategies and behavioural strategies will be completely independent of the surrounding game and depends only on the sequence of spaces  $A_t$  and  $S_t$ . Hence our freedom to modify the game below.

COMMENT 1.16. Generalised strategies make sense only for games with perfect recall: they would turn any other game into a game with perfect recall (cf. also ex. II.1Ex.14 p. 72).

ii. There is no loss in pooling all opponents of  $i$  together, including nature, as a single player, who is always informed of the whole past history, and uses a behavioural strategy.

HINT. Denote by  $H_1$  (resp.  $H_2, H_3$ ) the sets of histories ending with an  $\omega \in \Omega$  (resp.  $a \in A, s \in S$ ). For the new game  $\tilde{\Gamma}$ , let  $\tilde{\Omega} = H$ ,  $\tilde{A}^i = A^i$ . Note  $a$  (= last  $a$ ) is a measurable function on  $H_2$ , let  $\tilde{A}^{\text{opp}} = H \setminus \{\tilde{\omega} \in H_2 \mid a(\tilde{\omega}) \in A^i\}$ ,  $\tilde{A} = \tilde{A}^i \cup \tilde{A}^{\text{opp}}$ , and define  $\tilde{a}$  by  $\tilde{a}(\tilde{\omega}) = \tilde{\omega}$  for  $\tilde{\omega} = \tilde{A}^{\text{opp}}$ ,  $\tilde{a}(\tilde{\omega}) = a(\tilde{\omega})$  otherwise. The set of actions  $\tilde{S}^i$  of player  $i$  equals  $S$ , and let  $\tilde{S}_1 = A$ ,  $\tilde{S}_2 = S$  and  $\tilde{S}_3 = \Omega$  be the sets of actions of the opponent corresponding to  $H_1, H_2$  and  $H_3$ . The map  $\tilde{q}$  is defined in the obvious way. Set  $\tilde{S}^{\text{opp}} = \tilde{S}_1 \times \tilde{S}_2 \times \tilde{S}_3 \times \{1, 2, 3\}$  to have a single action space for the opponent — only the relevant component will be used. (One could similarly use now  $\tilde{S} = \tilde{S}^i \times \tilde{S}^{\text{opp}}$ , but we have to preserve the property that  $\tilde{S}^i$  is standard Borel.) To make sure the opponent uses a behavioural strategy, include in the above construction the selection by any player  $j \neq i$  of a point in one of his auxiliary spaces as a move in the game, if he uses a mixed (or a general) strategy, so all those choices appear in the space  $H$  of histories,

hence in  $\tilde{A}^{\text{opp}}$ . Use the product of all auxiliary spaces, in order to stay with a single action space for the opponent. Observe finally there is a Borel map from the space of all such plays to plays of the original game: for legal plays,  $\tilde{s}$  with index  $i \in \{1, 2, 3\}$  can occur only at time equal to  $i \bmod 3$ , and  $s$  only at time  $2 \pmod 3$ . On this subset, keep  $s_t$  and replace  $\tilde{s}_t$  by its coordinate  $i$ . Map the (measurable) remainder to some given play. This map is such that, if a strategy vector of the original game is transformed in the obvious way to a strategy pair in the new game, the map will transform correctly the induces probability distributions on histories.

COMMENT 1.17. Note we did not change the sets  $A^i$  and  $S^i$  in the construction, hence also not the sets of strategies of player  $i$ . We can now view the set of histories as the set of sequences  $(\tilde{s}_1, \tilde{s}_2, \dots, \tilde{s}_t, s_{t+1}, \tilde{s}_{t+2}, \dots)$ , with a map  $\alpha$  from the set  $H$  of all such finite sequences  $h$  to  $A^i \cup \{\text{opp}\}$ . If  $\alpha(h) = \text{opp}$ , the opponent picks the next point  $\tilde{s}$  using a transition probability that depends on the whole  $h$ , otherwise player  $i$  picks the next  $s$  as a function of his own past history.  $\alpha$  maps the empty sequence to “opp”, and depends only on the last coordinate of the sequence since, if it is  $s$ , or if  $\tilde{s}$  with index  $\in \{2, 3\}$ , it is mapped to “opp”, while otherwise  $\tilde{s}$  is already an element of  $A$ .

iii. Let  $(T_n, \mathcal{T}_n)$  be the set of possible “histories between player  $i$ ’s stage  $n - 1$  and stage  $n$ ”. The opponent’s strategy (and the rules of the game) yields a sequence of transition probabilities  $Q_n$  from  $\prod_{t < n}(T_t \times A_t \times S_t, \mathcal{T}_t \otimes \mathcal{A}_t \otimes \mathcal{S}_t)$  to  $(T_n \times A_n, \mathcal{T}_n \times \mathcal{A}_n)$ . (Observe  $Q_1$  is just a probability on  $T_1 \times A_1$ .)

HINT. We have now a natural player partition  $(H^0, H^i)$  of  $H$ . We drop tildes. The opponents’ strategy is now a transition probability  $R$  from  $H^0$  to  $H$ , where the  $S^0$  coordinate is just appended to the history. Complete  $R_1$  in a transition probability  $\bar{R}$  from  $H$  to  $H$  by sending all points in  $H^i$  to themselves. By Ionescu-Tulcea’s theorem, iterative use of  $\bar{R}$  induces a transition probability  $\tilde{R}$  from  $H$  to  $(H)^\infty$ , hence to the space  $(T, \mathcal{T}) = H^i \cup H_\infty$ . Take  $(T_n, \mathcal{T}_n)$  to be a copy of  $(T, \mathcal{T})$ . Define  $Q_n$  — using the given map  $q$  from  $H^i$  to  $A^i$  — in the following way: if  $t_{n-1} \in H_\infty$ ,  $t_n = t_{n-1}$  and  $a_n$  is any given point in  $A^i$ ; otherwise use  $\tilde{R}$  from  $(t_{n-1}, s_{n-1})$  (which belongs to  $H^0$ ) to  $T_n$ .  $Q_1$  is defined by applying  $\tilde{R}$  to the empty sequence. Again, like at the end of II.1Ex.10aii, note that the set of plays of the previous game is the set of all sequences in  $[S^i \cup S^0]^\infty$  where every element of  $S^i$  has a predecessor in the subset  $S^{00}$  of  $S^0$  where  $\alpha \neq \text{opp}$ . And obtain an appropriate measurable map from the set of plays in our new game to that set: the set being measurable in  $[S^0 \cup S^i]^\infty$ , it suffices to get an appropriate measurable map to the latter. (Be careful again not to use conditions like “if  $t_n$  extends  $t_{n-1}$ ”, since such sets are not necessarily measurable given absence of separability and separateness.)

b. We keep the notations and definitions from II.1Ex.10a. Assume the spaces  $(S_n, \mathcal{S}_n)$  are standard Borel. For every general(ised) strategy, there exists both a mixed strategy and a behavioural strategy such that, for every  $(T_n, \mathcal{T}_n, Q_n)_{n=1}^\infty$ , the induced probability on  $\prod_n(T_n \times A_n \times S_n)$  is the same.

- HINT.
- (1) Set  $(Y_n, \mathcal{Y}_n) = \prod_{t \leq n}[(X_{t+1}, \mathcal{X}_{t+1}) \times (S_t, \mathcal{S}_t) \times (A_t, \mathcal{A}_t)]$ ,  $Y_0 = \{0\}$ , and  $f_n$  the projection from  $Y_n$  to  $S_n$ . Define the transition probability  $R_n$  from  $(Y_{n-1}, \mathcal{Y}_{n-1}) \times (A_n, \mathcal{A}_n)$  to  $(Y_n, \mathcal{Y}_n)$  from  $P_n$  (2 p. 61),  $R_1$  is induced by  $P_0, P_1$  and the identity on  $A_1$ .
  - (2) Denote by  $\mathcal{Y}_n^0$  the trivial  $\sigma$ -field on  $Y_n$ , and by  $\mathcal{Y}_n^k \supseteq \mathcal{Y}_n^{k-1}$  a separable sub  $\sigma$ -field of  $\mathcal{Y}_n$  making  $f_n: (Y_n, \mathcal{Y}_n^k) \rightarrow (S_n, \mathcal{S}_n)$  and  $R_{n+1}: (Y_n, \mathcal{Y}_n^k) \times (A_{n+1}, \mathcal{A}_{n+1}) \rightarrow (Y_{n+1}, \mathcal{Y}_{n+1}^{k-1})$  measurable. Indeed, if  $R$  is a transition probability from  $(E \times F, \mathcal{E} \times \mathcal{F})$  to  $(G, \mathcal{G})$  and  $\mathcal{G}$  is separable, there exists separable sub  $\sigma$ -fields  $\mathcal{E}_0$  and  $\mathcal{F}_0$  of  $\mathcal{E}$  and  $\mathcal{F}$  for which  $R$  is still a transition probability: denote by  $G_i$  a sequence generating  $\mathcal{G}$ , by  $B_{ij} \in (\mathcal{E} \otimes \mathcal{F})$  the inverse images by  $R(G_i \mid e, f)$  of the rational interval  $I_j$  and by  $E_{i,j,k} \times F_{i,j,k}$  a sequence of  $(\mathcal{E} \times \mathcal{F})$ -measurable rectangles generating a  $\sigma$ -field containing  $B_{ij}$ : span  $\mathcal{E}_0$  (resp.  $\mathcal{F}_0$ ) by the  $E_{ijk}$  (resp.  $F_{ijk}$ ). It follows that the  $\sigma$ -fields  $\mathcal{Y}_n^\infty$  spanned by the  $\mathcal{Y}_n^k$  are separable, and for those the  $R_n$  are still transition probabilities and the  $f_n$  measurable.
  - (3) Let now  $\tilde{Y}_n$  be the quotient of  $Y_n$  by the equivalence relation  $y \sim y'$  if they belong to the same elements of  $\mathcal{Y}_n^\infty$ . The  $\mathcal{Y}_n^\infty$  can still be viewed as a separable  $\sigma$ -field on  $\tilde{Y}_n$ , for which the atoms are the singletons. And  $R_n$  and  $f_n$  are already defined on the quotients  $\tilde{Y}_n$ . In short, we can now assume the  $\sigma$ -fields  $\mathcal{Y}_n$  are separable, and separate points of  $Y_n$ . And similarly for

the  $(A_n, \mathcal{A}_n)$ , we can construct a measurable map  $p_n$  to a space  $(B_n, \mathcal{B}_n)$  which is separable and separating, and such that  $R_n$  is a transition probability from  $(Y_{n-1}, \mathcal{Y}_{n-1}) \times (B_n, \mathcal{B}_n)$  to  $(Y_n, \mathcal{Y}_n)$ .

- (4) If  $\mathcal{Z}$  is a separable and separating  $\sigma$ -field on  $Z$ , then  $(Z, \mathcal{Z})$  can be identified with a subset of the Cantor set  $C = \{0, 1\}^\infty$ , endowed with its Borel sets  $\mathcal{C}$  — a standard Borel space — by the map  $z \mapsto (\mathbb{1}_{Z_i}(z))_{i=1}^\infty$ , where  $Z_i$  denotes a sequence that generates  $\mathcal{Z}$  (the subset being endowed with the trace  $\sigma$ -field). And any measurable map  $f$  from  $(Z, \mathcal{Z})$  to  $[0, 1]$  can be extended to  $C$ : indeed, this holds, by definition of the trace  $\sigma$ -field, for indicator functions, therefore for their convex combinations the step functions, hence if  $f_n$  is a sequence of step functions converging to  $f$ , with extensions  $\bar{f}_n$ , let  $\bar{f} = \liminf_{n \rightarrow \infty} \bar{f}_n$ .
- (5) Using 4, we can view the  $Y_n$  and the  $B_n$  as subsets of  $\bar{Y}_n$  and  $\bar{B}_n$  which are copies of  $([0, 1], \mathcal{B})$ ; the maps  $f_n$  have an extension to  $\bar{Y}_n$ , and the transition probabilities  $R_n$  can be viewed as transition probabilities from  $Y_{n-1} \times B_n$  to  $\bar{Y}_n$ , assigning outer probability one to  $Y_n$ . Thus to show that the  $R_n$  too have also an extension  $\bar{R}_n$  as transition probabilities from  $\bar{Y}_{n-1} \times \bar{B}_{n-1}$  to  $\bar{Y}_n$ , it suffices (by 4) to show that the space  $(M, \mathcal{M})$  of probability measures on  $[0, 1]$  is a standard Borel space (in the weak<sup>\*</sup>-topology) — which is obvious, as a compact metric space —, and to use 9.e).
- (6) We can thus think of all  $(Y_n, \mathcal{Y}_n)$  and  $(B_n, \mathcal{B}_n)$  as being copies of  $([0, 1], \mathcal{B})$ . Introduce for each  $n$  a copy  $(U_n, \mathcal{U}_n, \lambda_n)$  of  $([0, 1], \mathcal{B}, \lambda)$ ,  $\lambda$  being Lebesgue measure, and replace  $R_n$  by a measurable map  $h_n$  from  $(Y_{n-1}, \mathcal{Y}_{n-1}) \times (B_n, \mathcal{B}_n) \times (U_n, \mathcal{U}_n)$  to  $(Y_n, \mathcal{Y}_n)$  (cf. proof of theorem 1.8 p. 55). Composition of the  $h_t$  ( $t \leq n$ ) and of  $f_n$  yields a description of the strategy by Borel maps  $g_n$  from  $\prod_{t \leq n} [(B_t, \mathcal{B}_t) \times (U_t, \mathcal{U}_t)]$  to  $(S_n, \mathcal{S}_n)$ . Since  $\prod_{n=1}^\infty (U_n, \mathcal{U}_n, \lambda_n)$  is itself Borel isomorphic to  $([0, 1], \mathcal{B}, \lambda)$ , we can as well think of  $g_n$  as Borel maps from  $([0, 1], \mathcal{B}) \times \prod_{t \leq n} (B_t, \mathcal{B}_t)$  to  $(S_n, \mathcal{S}_n)$ : this (together with the  $p_n: A_n \rightarrow B_n$ ) is the mixed strategy.
- (7) Check that, in each of the previous steps, the probability distribution induced on  $\prod_n (T_n \times A_n \times S_n)$  did not change.

There only remains therefore to replace the mixed strategy by a behavioural strategy:

- (8) Use now ex. II.1Ex.9 p. 60 inductively for  $n = 1, 2, 3, \dots$ ; using  $(A_n, \mathcal{A}_n) \times [\prod_{t < n} (A_t \times S_t, \mathcal{A}_t \times \mathcal{S}_t)] \times$  for  $(E, \mathcal{E}) (= (E_n, \mathcal{E}_n))$ ,  $(S_n, \mathcal{S}_n)$  for  $(Y, \mathcal{Y})$ , the mixed strategy  $\sigma_n$  for  $g$ ,  $([0, 1], \mathcal{B})$  for  $(X, \mathcal{X})$ ,  $P^n$  for  $P$  with  $P^1 = \lambda$  and yielding  $Q(B | e, y)$  for  $P^{n+1}$  and  $P_e \circ g_e^{-1}$  as transition probability  $\bar{P}^n$  from  $(E_n, \mathcal{E}_n)$  to  $(S_n, \mathcal{S}_n)$ .
- (9) The sequence  $\bar{P}_n$  forms the required behavioural strategy. Check that here too the probability distributions induced on plays are unaffected.

COMMENT 1.18. In particular, when games with perfect recall are presented in terms of incremental information (remark 1.3 after Theorem 1.3), this proof shows that the correspondence between behavioural, mixed and general(ised) strategies is completely independent of the game — it depends only on  $(A_i, \mathcal{A}_i)$  and  $(S_i, \mathcal{S}_i)$ . It also shows that, in terms of mixed strategies, there is no need to remember past actions.

c. Show, without drastic modifications in the proof, that the above remains valid if the “set of dates for player  $i$ ”, instead of  $\mathbb{N}$  is allowed to be any countable well-ordered set.

d. Consider a “countably linear” game for player  $i$ , i.e. there exists a countable measurable partition of  $(A_i, \mathcal{A}_i)$  such that each partition element is met at most once along any play. Prove like in 2 p. 62 that general strategies for  $i$  can be replaced by mixed strategies.

e. If  $A^i$  is countable, and the game is linear for  $i$ , then every general strategy for  $i$  is equivalent to a probability distribution over pure strategies (without any standard Borel restriction on the action sets  $S_a^i$ ).

HINT. Construct first a probability distribution over  $X \times \sum^i = X \times (\prod_{a \in A^i} S_a^i)$ , treating the different factors  $S_a^i$  as conditionally independent given  $X$ .

f. Show the above results remain valid with variable action sets, provided player  $i$  has a behavioural strategy (resp. a pure strategy). (Apply the result with the embedding space  $S$ , then modify the obtained strategy where it is not carried by  $\bar{S}$ .)

g. Show that, without any assumption on the game, a generalised strategy can always be represented by  $(X, \mathcal{X}, x_1, P, \sigma)$ , where  $x_1$  is the initial state in the auxiliary space  $(X, \mathcal{X})$ ,  $P$  is a transition probability from  $(X, \mathcal{X}) \times (A^i, \mathcal{A}^i)$  to  $(X, \mathcal{X})$ , and  $\sigma$  is a measurable map from  $(X, \mathcal{X})$  to  $(S^i, \mathcal{S}^i)$  (such that  $\forall x, \forall a, P[\sigma^{-1}(\bar{S}_a^i) | x, a] = 1$ ).

### 11. Best Replies.

a. Still in the context of ex. II.1Ex.10a p. 61, even the concept of a generalised strategy is not necessarily satisfactory as a concept of reply (say in the definition of an equilibrium). Indeed, for a reply, the game and the strategies of the others — i.e.  $(T_n, \mathcal{T}_n, Q_n)_{n=1}^\infty$  — are given, so every time player  $i$  has to randomise according to  $P_n$ , the full probability measure on the conditioning space is known. It becomes then more natural to require the measurability of  $P_n$  only with respect to this measure and maybe also to require equalities to hold only a.e.

b. Assume  $Q$  is a “ $P$ -transition probability” from a probability space  $(E, \mathcal{E}, P)$  to a standard Borel space  $(F, \mathcal{F})$ , i.e.  $\forall A \in \mathcal{F}, Q_e(A)$  is  $P$ -measurable and  $P$ -a.e.  $\in [0, 1]$ ,  $Q_e(F) = 1$   $P$ -a.e., and if  $A_i \in \mathcal{F}$  is a disjoint sequence, then  $P$ -a.e.  $\sum_i Q_e(A_i) = Q_e(\bigcup_i A_i)$ . Show that there exists a unique probability  $P \cdot Q$  on the product, satisfying  $(P \cdot Q)(A \times B) = \int_A Q_e(B)P(de), \forall A \in \mathcal{E}, \forall B \in \mathcal{F}$ , and that there exists a transition probability  $\bar{Q}$ , with  $P \cdot \bar{Q} = P \cdot Q$ .

HINT. E.g., fix a lifting (cf. ex. II.1Ex.15 p. 73)  $\varrho$  on  $L_\infty(E, \mathcal{E}, P)$ , view  $F$  as a compact metric space with its Borel sets, define  $\bar{Q}_e(\varphi) = \varrho[Q_e(\varphi)]$  for every continuous function  $\varphi$  on  $F$ , and use Riesz's theorem (cf. also ex. II.1Ex.16 p. 75).

c. The above would e.g. allow to use, in “general replies”,  $P$ -transition probabilities from the past to  $(S_n, \mathcal{S}_n)$  — but shows that it would not matter — the same probability distributions on histories will be generated as with true strategies.

d. Show that such probability distributions have the property that the conditional distribution on  $T_n \times A_n$  given the past is given by  $Q_n$  and that the conditional distribution on  $S_n$  given {the past and  $T_n \times A_n$ } depends only on  $A_n \times \prod_{t < n} (A_t \times S_t)$ , i.e.,  $S_n$  and  $\prod_{t \leq n} T_t$  are conditionally independent given  $A_n \times \prod_{t < n} (A_t \times S_t)$ , and clearly any reasonable concept of reply of player  $i$  has to lead to distributions on histories having those two properties.

e. Conversely, show that any distribution on histories having those two properties can be generated by a behavioural strategy of player  $i$  (use II.1Ex.11b). There is thus no need to look for a concept of reply wider than the concept of strategy.

### 12. The definition of perfect recall.

a. *Justifying the measurability assumption.*

i. The definition of a game with perfect recall for  $i$  given in 1.c p. 54 seems completely unsatisfactory; the conceptually correct definition is: for every feasible  $a \in A^i$  (i.e., occurring along some feasible play), there exists a unique pair  $\varphi(a) = (\beta(a), \sigma(a)) \in A^i \times S^i \cup \{\iota\}$  such that, for any feasible play where  $a$  occurs, either this is the first move of  $i$  along the play and then  $(\beta(a), \sigma(a)) = \iota$  or the previous move of  $i$  was  $\sigma(a)$  at  $\beta(a)$ . The maps  $\beta$  and  $\sigma$  are then completely defined by the model, and are in no sense primitive data, hence their measurability should follow from assumptions on the primitive data of the model. To give a precise meaning to the above notion of “feasible play”, reduce first, as in ex. II.1Ex.1a p. 56, the game to a deterministic one, by adding a player 0. Feasible plays are then all those compatible with the restrictions on action sets, and with the given initial state.

ii. Denote by  $A_0^i$  the feasible part of  $A^i$ , and assume the game has perfect recall in the above sense, i.e., the map  $(\alpha, \sigma)$  is well defined on  $A_0^i$ . If  $(A_0^i, \mathcal{A}_0^i)$  is separable and separated, and  $(S^i, \mathcal{S}^i)$  is standard Borel, then  $(\alpha, \sigma)$  is measurable, and  $(A_0^i, \mathcal{A}_0^i)$  is a Blackwell space, i.e. isomorphic to an analytic subset of  $[0, 1]$  (cf. App.6).

- HINT.
- (1) Show that one can always replace  $\Omega$  by  $X = \bigcup_{n \geq 0} S^n$  — the measurable maps of the model induce the required measurable maps from  $X \times S$  to  $X$ , from  $X$  to  $A$  and from  $X$  to  $\Omega$  — hence from  $(X)^\infty$  to pay-offs (which in fact factorises through  $S^\infty$ ). Note  $X$  is standard Borel by our assumptions. Similarly, since  $(A_0, \mathcal{A}_0)$  is separable and separated, it can be identified with a subset of  $([0, 1], \mathcal{B})$ , cf. ex. 4 p. 63, hence with  $([0, 1], \mathcal{B})$  itself since the statement involves only  $A_0$  which is unaffected. Extend the measurable partition  $(A^i)_{i \in I}$  to a measurable partition of  $[0, 1]$ . Assume thus  $(A, \mathcal{A})$  standard Borel.
  - (2) Let  $\theta$  be the above map from  $X$  to  $A$ . Let  $\psi: X \times S^\infty \rightarrow A^\infty$  be defined by  $\psi(x, s_1, s_2, s_3, \dots) = (\theta(x), \theta(x, s_1), \theta(x, s_1, s_2), \dots)$ :  $\psi$  is a measurable map between standard Borel spaces, so its graph  $G \subseteq X \times (A \times S)^\infty$  is measurable, hence standard Borel.
  - (3) Given a Borel set  $B$  in  $A^i \times S (\subseteq A \times S)$ ,  $\tilde{B} = G \cap [X \times B \times (A \times S)_{n=1}^\infty]$  is Borel in the standard Borel space  $G$ . Let  $T: G \rightarrow [\mathbb{N} \setminus \{0\}] \cup \{+\infty\}$ :  $T(g) = \min\{n \geq 1 \mid a_n \in A^i\}$ , with  $\min \emptyset = +\infty$ , for  $g = (x; a_0, s_0; a_1, s_1; \dots)$ .  $T$  is measurable. Thus  $f(g) = a_{T(g)}$  for  $T(g) < \infty$ ,  $f(g) = \text{"out"}$  for  $T(g) = +\infty$  is also measurable, as composition: thus,  $\overline{B} = f(\tilde{B}) \setminus \{\text{"out"}\}$  is analytic. And  $\overline{B} = (\beta, \sigma)^{-1}(B)$ .
  - (4) Thus the inverse image of any Borel set of  $A^i \times S$  is analytic. The same holds for the inverse image of  $\iota$  — just drop in the above argument the first factor  $X \times A \times S$ . Our map  $(\beta, \sigma)$  from  $A_0^i$  to  $(A^i \times S) \cup \{\iota\}$  is such that the inverse image of any measurable set is analytic:  $A_0^i$  is analytic and, by the separation theorem for analytic sets (3.h),  $(\beta, \sigma)$  is measurable.

iii. The assumption sub II.1Ex.12aii that  $\mathcal{A}$  is separating is crucial: consider the one person game with  $A = X$  consisting of 4 points:

iv. The separability assumption is also crucial: otherwise consider the same structure, but where initially nature picks a point in  $[0, 1]$ , in period 1 player I observes the Borel  $\sigma$ -field, and in period 2 he remembers only the  $\sigma$ -field generated by the singletons.

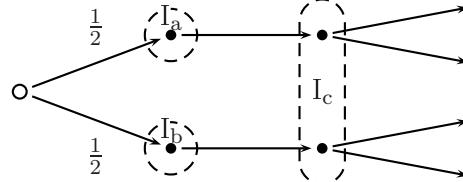


FIGURE 2. The need for separating  $\sigma$ -fields.

COMMENT 1.19. It follows that the measurability assumption on  $(\beta, \sigma)$  is indeed a correct generalisation of the definition to cases where  $\mathcal{A}$  is possibly not separated or separating — while still remaining in a framework with unambiguous set of feasible plays. (Possibly one has to reduce first  $A$  to some subset  $A^0$ , and similarly  $\Omega$  to its feasible subset  $\Omega^0$  (with an appropriate separable sub  $\sigma$ -field), in such a way that, with  $A^0$  and  $\Omega^0$ , all assumptions of the model are still valid, that all strategies remain strategies, and that now  $A$  and  $\Omega$  are Blackwell spaces, and each of their points if feasible.) Observe that, with this definition, we can not only, as in 1 p. 65 above, identify  $\Omega$  with  $X$ , but also  $A$  (or rather  $A^0$ ) with  $X - \mathcal{A}$  being a sub  $\sigma$ -field. The canonical map  $\varphi$  from  $(X, \mathcal{A}_i)$  to  $(X, \mathcal{A}_i \otimes \mathcal{S})$  is then always well defined, whatever be the game, and the whole perfect recall assumption is about its measurability. Still equivalently, it could be phrased in terms of an increasing sequence of  $\sigma$ -fields of player  $i$  in  $H_\infty$ .

We consider now how to extend the definition to the general model, with moves of nature, where, as mentioned before, there is no well defined set of feasible plays. We will therefore require the map  $\varphi$  to be defined a.e., and also only to be true a.e. — in order not to have to reduce first  $A$  to some subset  $A^0$ , as above, and to have a definition independent of the set of feasible plays used.

Observe that, if there exists a strategy  $\mathbf{I}^{\text{tuple}}$  or profile, there exists a behavioural strategy profile, and that a set of plays which is negligible for every behavioural strategy profile is negligible for every strategy profile. So in the following, “a.s.” will mean “a.s. for every (behavioural) strategy profile”.

#### b. The General Case.

- i. For every separable sub  $\sigma$ -field  $\mathcal{B}$  of  $\mathcal{H}_\infty$ , any given set of measurable maps  $\varphi^i: A^i \rightarrow \overline{S}^i = [(A^i \times S) \cap \overline{S}] \cup \{\iota\}$ , and for every sequence of (general) strategies, there exist separable

sub  $\sigma$ -fields  $\tilde{\mathcal{A}}$  and  $\tilde{\mathcal{O}}$  on  $A$  and  $\Omega$  such that  $\mathcal{B}$  is included in the corresponding product  $\sigma$ -field, such that all measurability assumptions of the model are still satisfied with the  $\sigma$ -fields  $\tilde{\mathcal{A}}$  and  $\tilde{\mathcal{O}}$ , such that all given strategies are still strategies, and such that the  $\varphi^i$  are still measurable.

HINT. Observe that any element of a product  $\sigma$ -field belongs to a product of separable sub  $\sigma$ -fields. Start then with separable sub  $\sigma$ -fields  $\mathcal{A}_0$  and  $\mathcal{O}_0$  on  $A$  and  $\Omega$  such that  $\mathcal{B}$  is included in the corresponding product  $\sigma$ -field, such that the  $A^i$  are measurable, as well as the maps  $g^i$  and the set  $\bar{S} \subseteq A \times S$ . Define inductively  $\mathcal{A}_{n+1} \supseteq \mathcal{A}_n$  and  $\mathcal{O}_{n+1} \supseteq \mathcal{O}_n$  as separable sub  $\sigma$ -fields making all maps  $\varphi^i$  and all transition probabilities (those of nature and of the players) measurable when using  $\mathcal{A}_n$  and  $\mathcal{O}_n$  on the image.  $\bigcup_n \mathcal{A}_n$  and  $\bigcup_n \mathcal{O}_n$  answer the question.

ii. Prove that the following conditions on a measurable map  $\varphi^i: A^i \rightarrow \bar{S}^i$ , defined a.e. on  $A^i$  (i.e.,  $\varphi^i$  is well-defined along a.e. play) are equivalent, letting, for  $\omega \in H_\infty$ ,  $T_n^i(\omega) = \inf\{t > T_{n-1}^i(\omega) \mid a_t(\omega) \in A^i\}$  (and  $T_0^i = 0$ ,  $\inf(\emptyset) = +\infty$ ):

- $\forall B \in \bar{\mathcal{S}}^i$ ,  $\mathbb{1}_B \circ \varphi^i(a_{T_{n+1}^i}) = \mathbb{1}_B(s_{T_n^i})$  a.s.  $\forall n \geq 1$  and  $\varphi^i(a_{T_1^i}) = \{\iota\}$  a.s.  
 $s_t$  is the pair in  $\bar{S}_t \subseteq (A \times S)_t$
- $\forall C \in \bar{\mathcal{S}}^i \otimes \bar{\mathcal{S}}^i$  such that  $(s_1, s_2) \in C \Rightarrow s_1 \neq s_2$ ,  $\mathbb{1}_C(s_{T_n^i}, \varphi^i(a_{T_{n+1}^i})) = 0$  a.s.  
 $\forall n \geq 1$ , and  $\varphi^i(a_{T_1^i}) = \{\iota\}$ .

Call such a function  $\varphi^i$  a recall function, and say then that player  $i$  has perfect recall. Prove also that, given a recall function for each player in  $\mathbf{I}_0 \subseteq \mathbf{I}$ , the set of plays where all those recall functions are well defined and exact has outer probability one for every strategy vector.

HINT. Observe that II.1Ex.12bi remains valid if the  $\varphi^i$  are only defined a.e. — except that then, for the separable sub  $\sigma$ -fields,  $\varphi^i$  will be defined a.e. only for the strategies in the prescribed sequence. Use this, and that when  $\mathcal{A}^i$  is separable and separating,  $\{\omega \mid \varphi^i(a_{T_{n+1}^i}) = s_{T_n^i}\}$  is measurable, and has probability one by the first condition.

iii. Given an extensive form, for which there exists a strategy  $\mathbf{I}^{\text{tuple}}$ , assume  $\omega_1 \in \tilde{\Omega} \subseteq \Omega$  and  $\tilde{A} \subseteq A$  are such that the restriction to  $\tilde{\Omega}$  and  $\tilde{A}$  is also an extensive form, i.e. the outer probability of  $\tilde{A}$  is one for every  $\omega \in \tilde{\Omega}$  and the outer probability of  $\tilde{\Omega}$  is one for every  $(\omega, s) \in \tilde{\Omega} \times \tilde{S}$ . Then if  $i$  has perfect recall in  $\Gamma$ , he also has in  $\tilde{\Gamma}$ .

HINT. Show first that the behavioural strategy profile of  $\tilde{\Gamma}$  are the restrictions of those in  $\Gamma$  — using that a behavioural strategy is a measurable map to the standard Borel space of probabilities on  $S$ , that measurable maps to a standard Borel space can always be extended, and that the set where such an extension is not a strategy is measurable, so that the extension can be replaced there by the given strategy profile of  $\Gamma$ . Show then that, for any strategy profile, the set of plays of  $\tilde{\Gamma}$  is a subset with outer probability one of the set of plays of  $\Gamma$ , and define  $\tilde{\varphi}^i$  as an appropriate restriction of  $\varphi^i$ .

iv. The converse of II.1Ex.12biii holds if  $\tilde{A}$  is universally measurable (cf. 4.d) in  $A$ , or if  $(A, \mathcal{A})$  is standard Borel.

HINT. In the first case, there is no need to extend  $\tilde{\varphi}$ ; in the second case, as reminded sub 3 p. 65,  $\tilde{\varphi}$  has an extension  $\varphi$  (with values in  $\bar{S} \subseteq A \times S$ ) to some Borel subset of  $A$  containing  $\tilde{A}$ .

COMMENT 1.20. II.1Ex.12biii and II.1Ex.12biv show thus that the definition of perfect recall is, as required, essentially independent of the surrounding space (“essentially” because of the small restrictions in the converse — which seem unavoidable). The definition is also flexible enough to include all cases studied sub II.1Ex.12a p. 64, without having to restrict first artificially the sets  $A$  and  $\Omega$  to appropriate subsets  $A_0$  and  $\Omega_0$ .

### 13. A converse to Exercise II.1Ex.12a.

a.

i. Show that, under the assumptions of ex. II.1Ex.12a p. 64, one obtains a stronger form of perfect recall: the “a.s.” qualifications do not only hold for every strategy vector of the players, but even for every (pseudo-) strategy vector of players that would be fully informed about the whole past (i.e., in some sense, “whatever players do” instead of “whatever be their strategies”).

COMMENT 1.21. The purpose of the present exercise is to show that, in this form, the statement is an equivalence.

ii. If  $(A, \mathcal{A})$  is separable (and separated), and if a subset  $R$  of players has perfect recall in the strong a.s. sense as above, there exist measurable restrictions on the actions of nature — i.e. measurable subsets  $\Omega^0$  of  $\Omega$ ,  $C$  of  $\Omega^0 \times A$  and  $K$  of  $\Omega^0 \times A \times S \times \Omega^0$ , such that  $\omega_1 \in \Omega^0$ ,  $\alpha(C_\omega) = 1$ ,  $\forall \omega \in \Omega^0$  and such that  $(\omega, a, s, \omega') \in K \Rightarrow (\omega, a) \in C$ ,  $(a, s) \in \overline{S}$ , and  $\forall (\omega, a, s) : (\omega, a) \in C, (a, s) \in \overline{S} \Rightarrow q_{\omega, a, s}(K_{\omega, a, s}) = 1$  — for which every feasible play (cf. ex. II.1Ex.12ai p. 64) exhibits perfect recall for all players in  $R$  [in the sense that, for some measurable functions  $\varphi^i$  ( $i \in R$ ) from a subset  $\tilde{A}^i$  of  $A^i$  to  $\{(a, s) \in \overline{S} \mid a \in \tilde{A}^i\} \cup \{\iota\}$  (which coincide with the given functions on their common domain),  $\varphi^i(a_t)$  is the previous pair  $(a, s)$  encountered by  $i$  along that play (or  $\iota$  if none), for any  $a_t \in A^i$  occurring along that play]. (Outside  $\Omega^0$ , one can use the original restrictions, if any).

Actually, to avoid hiding any “perfect recall” type of assumptions, we will assume the strong “a.s.” property only for “pseudo-strategy vectors” which are transition probabilities depending only on current state and signal.

To make the definition of perfect recall non-vacuous, we also assume that there exists at least one (pseudo-)strategy vector  $\sigma_0$  for the game — which can then be assumed to be behavioural.

The proof itself of the above statement is given sub II.1Ex.13b below.

iii. The following example shows the above strong form of perfect recall to be definitely stronger than the definition: Nature picks  $x_1$  in  $[0, 1]$  with Lebesgue measure. Player II chooses  $x_2$  in  $[0, 1]$  after having observed  $x_1$ ; then player I chooses  $x_3$  after having observed  $a = x_2 + \xi$  ( $\text{mod } 1$ ), where the noise  $\xi$  is uniformly distributed on  $[-\varepsilon, \varepsilon]$ . Later player I has to move again, and will be reminded of  $(a, x_3)$ , except if  $x_3 = x_1$ , in which case he gets a blank signal.

Clearly player I has perfect recall in this game, but not according to the strong variant of the definition. (The only purpose of player II’s presence, instead of letting player I just move after nature, without information, is to prevent an objection that by interchanging the ordering of the moves (cf. ex. II.1Ex.1c p. 56) one would obtain perfect recall in the strong sense).

iv. It follows that the only relaxation involved in our definition of perfect recall vis-à-vis the strong definition of ex. II.1Ex.12a lies in restricting (for the definition of “a.s.”) the players to use only true strategies — which seems natural enough not to require further justification, and is indeed sufficient for the proof of Kuhn’s theorem.

b. *General revision exercise* — « *La méthode des épeluchures* ». We turn now to the proof of the statement sub II.1Ex.13aii. Actually, it is often more convenient to formulate a game already with restrictions on the moves of nature so the transition probabilities only have to be defined on the relevant sets. We therefore allow those, and require the sets  $\Omega^0, C$  and  $K$  to be co-analytic — and will construct smaller measurable sets  $\Omega^0, C$  and  $K$  for which every feasible play exhibits perfect recall.

Sub II.1Ex.13c below some required lemmas are given, and it is shown how to reduce the problem to the case where the spaces  $A$  and  $\Omega$  are standard Borel, where  $\varphi$  is defined on the whole of  $A$ , and where the sets  $\Omega^0, C$  and  $K$  are Borel.

The central iteration of the proof is given in part II.1Ex.13d, and the final argument sub II.1Ex.13e.

c. *Preliminaries*. We first reduce (sub II.1Ex.13ci, II.1Ex.13cii, II.1Ex.13ciii) to the standard Borel case.

i. If  $\mathcal{A}$  is not separating, pass to the quotient — all assumptions remain valid, and measurable restrictions with the quotient induce such with the original.

ii. Choose a separable sub- $\sigma$ -field  $\mathcal{O}_0$  on  $\Omega$  such that

- the pay-off function is still measurable on the infinite product
- the sets  $\Omega^0, C$  and  $K$  are still co-analytic.
- the transition probability  $\alpha$  to  $(A, \mathcal{A})$  is still measurable
- the given pseudo-strategy is still a transition probability from  $C$  to  $S$ .

HINT. Use that the transition probabilities are to separable spaces, that a measurable subset of a product is generated by countably many products of measurable sets, and that a Souslin scheme involves only countably many measurable sets.

Similarly, let now  $\mathcal{O}_{n+1}$  be a separable sub- $\sigma$ -field containing  $\mathcal{O}_n$  and such that  $q_{\omega,a,s}(A)$  is  $\mathcal{O}_{n+1} \otimes \mathcal{A} \otimes \mathcal{S}$ -measurable  $\forall A \in \mathcal{O}_n$ . The  $\sigma$ -field  $\mathcal{O}_\infty$  spanned by  $\bigcup_n \mathcal{O}_n$  is then a separable sub- $\sigma$ -field on  $\Omega$  for which all assumptions remain true. Conclude that we can also assume  $(\Omega, \mathcal{O})$  to be separable and separated.

iii.  $A$  and  $\Omega$  can now be viewed as subsets of standard Borel spaces  $\tilde{A}$  and  $\tilde{\Omega}$ . Using 9.a and 9.e, select a measurable extension of the pay-off function to  $\tilde{\Omega}^\infty; \tilde{S}, \tilde{\Omega}^0, \tilde{C}$  and  $\tilde{K}$  as measurable or co-analytic extensions of  $\overline{S}, \Omega^0, C$  and  $K$  in  $\tilde{A} \times S, \tilde{\Omega}, \tilde{\Omega} \times \tilde{A}, \tilde{\Omega} \times \tilde{A} \times S \times \tilde{\Omega}; \tilde{\alpha}$  as a transition probability from  $\tilde{\Omega}$  to  $\tilde{A}$ ,  $\tilde{q}$  as a transition probability from  $\tilde{\Omega} \times \tilde{A} \times S$  to  $\tilde{\Omega}$ ,  $\tilde{\sigma}$  as a transition probability from  $\tilde{\Omega} \times \tilde{A}$  to  $S$ ,  $(\tilde{A}^i)_{i \in \mathbb{I}}$  as a measurable extension of the player partition, and finally  $(\tilde{\varphi}^i)_{i \in R}$  as Borel maps from  $\tilde{A}^i$  to  $\tilde{S} \cap [\tilde{A}^i \times S] \cup \{i\}$ .

iv. Before continuing we need 4 lemmas:

A. Let  $D = \{(\omega, a) \in \tilde{\Omega} \times \tilde{A} \mid \tilde{\sigma}_{\omega,a}(\tilde{S}_a) = 1\}$ .  $D$  is a Borel set containing  $C$ . We will define strategies in the enlarged game as transition probabilities from  $D$  to  $S$ , such that  $\tilde{S}$  has probability 1.

- Note first that the restriction of any strategy in the enlarged game is a pseudo-strategy in the original; conversely, any strategy  $\tau$  of the original game is such a restriction (consider any Borel extension  $\tilde{\tau}$ , and let  $\tilde{\tau}_{\omega,a} = \tilde{\tau}_{\omega,a}$  if  $\tilde{\tau}_{\omega,a}(\tilde{S}_a) = 1$ ,  $\tilde{\tau}_{\omega,a} = \tilde{\sigma}_{\omega,a}$  otherwise).
- To construct (by Ionescu Tulcea's theorem) the probability distribution  $P_{\tilde{\tau}}$  induced on plays  $(\tilde{H}_\infty = (\tilde{\Omega} \times \tilde{A} \times S)^\infty)$  in the enlarged game by a strategy  $\tilde{\tau}$  in that game, extend first  $\tilde{\tau}$  in an arbitrary Borel way on  $(\tilde{\Omega} \times \tilde{A}) \setminus D$ . Denote also  $P_\tau$  the probability distribution induced on  $H_\infty = \{(\omega_t, a_t, s_t)_{t=1}^\infty \mid (\omega_t, a_t) \in C, (a_t, s_t) \in \overline{S}, \omega_1 = \omega_1\}$  by the restriction  $\tau$  of  $\tilde{\tau}$  to the original game — i.e. to  $C$ .

Show that, for any measurable set  $B$  in  $\tilde{H}_\infty$ ,  $P_{\tilde{\tau}}(B) = P_\tau(B \cap H_\infty)$ .

HINT. Since the right hand member is also a measure on  $\tilde{H}_\infty$ , it suffices to show equality on generators. Hence we can by induction assume that equality holds for all measurable  $B_1$  in the product of the first  $n$  factors and prove that  $P_{\tilde{\tau}}(B_1 \times B_2) = P_\tau(B_1 \times B_2 \cap H_\infty)$  for measurable  $B_2$  in  $X_{n+1}$  — this is elementary.

- In particular  $P_{\tilde{\tau}}$  depends only on  $\tau$ , not on its extension.
- It follows also that the negligible subsets of  $H_\infty$  are the traces of the negligible subsets of  $\tilde{H}_\infty$  (negligible: of probability zero for every  $P_\tau$  — resp.  $P_{\tilde{\tau}}$ ). And universally measurable — in particular analytic (4.d.1) — subsets of  $\tilde{H}_\infty$  are negligible iff their trace on  $H_\infty$  is so.
- Call a subset  $X$  of  $\tilde{\Omega} \times \tilde{A} \times S \times \tilde{\Omega}$  negligible (and similarly for the original space), if for all  $n$   $\prod_{t < n} (\tilde{\Omega} \times \tilde{A} \times S)_t \times (X \times \tilde{A}_{n+1} \times S_{n+1}) \times \prod_{t > n+1} (\tilde{\Omega} \times \tilde{A} \times S)$  is negligible. It follows that such an analytic  $X$  is negligible iff its trace on  $\Omega \times A \times S \times \Omega$  is so. The same holds for subsets of  $\tilde{\Omega}$ , of  $\tilde{\Omega} \times \tilde{A}$ , and of  $\tilde{\Omega} \times \tilde{A} \times S$ .

B. The set  $\Pi$  of probability distributions over plays  $(\tilde{\Omega} \times \tilde{A} \times S)^\infty$  induced by (pseudo)-strategies is Borel. So the corresponding sets of feasible distributions on  $\tilde{\Omega}, \tilde{\Omega} \times \tilde{A}, \tilde{\Omega} \times \tilde{A} \times S, \tilde{\Omega} \times \tilde{A} \times S \times \tilde{\Omega}$  — which are the union over  $t \geq 1$  of the marginals of  $\Pi$  on the relevant factor of the set of plays, are analytic. (Thus, the negligible subsets of one of those sets are those which are negligible for every feasible distribution).

HINT. Like in ex. II.1Ex.11d and II.1Ex.11e p. 64 and by II.1Ex.13civA p. 68  $\Pi$  consists of those distributions on the standard Borel set of plays such that (α) the marginal on  $\Omega_1$  is the unit mass at  $\omega_1$  (a closed subset); (β)  $(a_t, s_t) \in \tilde{S} \forall t$  — a Borel condition; (γ) the conditional probabilities of nature are correct — i.e., given the factors  $X_1, X_2, X_3, \dots$  require that for every  $n$ , and a countable generating family of Borel sets  $A$  in  $\prod_1^n X_t$  and  $B$  in  $X_{n+1}$ , one has  $\mu(A \times B) = \int_A r(B)d\mu$  (denoting by  $r$  the specified transition probability of nature) — this is again a countable set of Borel conditions; (δ) the required conditional independence conditions hold for the (“pseudo-”) players — i.e.  $s_t$  is independent of the past given  $(\omega_t, a_t)$ . Use ex. II.1Ex.9 p. 60 to show that, if  $X, Y$  and  $Z$  are standard Borel, the set of probability distributions on  $X \times Y \times Z$  such that  $X$  and  $Y$  are conditionally independent given  $Z$  is a Borel set (take as parameter space  $E$  the space of all probabilities  $X \times Y \times Z$ ). So this condition too determines a Borel subset. Finally (ε), the conditional distribution  $\sigma$  of  $s_t$  given  $(\omega_t, a_t)$  is the same for all  $t$ . For this it is sufficient to show that the set of pairs of distributions  $P, Q$  on  $X \times Y$  ( $X, Y$  standard Borel) that have the same conditional on  $Y$  given  $X$  is a Borel set. By ex. II.1Ex.9 p. 60, there exists a conditional probability of  $Y$  given  $X$  for  $\frac{1}{2}(P + Q)$ , denoted by  $R(dy \mid x, P + Q)$ , which is jointly measurable in  $X$  and  $(P + Q)$ . The condition is then that, for measurable subsets  $X_1$  of  $X$  and  $Y_1$  of  $Y$ ,  $\int_{X_1} R(Y_1 \mid x, P + Q)P(dx) = P(X_1 \times Y_1)$ ,  $\int_{X_1} R(Y_1 \mid x, P + Q)Q(dx) = Q(X_1 \times Y_1)$ . For  $X_1$  and  $Y_1$  in countable generating sub-algebras, one gets in this way a countable family of equations, both members of which are Borel functions of  $(P, Q)$ . Hence lemma II.1Ex.13civB.

C. An analytic subset  $X$  of  $\tilde{\Omega} \times \tilde{A} \times S \times \tilde{\Omega}$  is negligible iff the analytic set  $\{(\omega, a, s) \mid \tilde{q}_{\omega, a, s}(X) > 0\}$  is so.

HINT. 9.f.

D. An analytic subset of  $\tilde{\Omega} \times \tilde{S}$  is negligible iff its analytic projection on  $\tilde{\Omega} \times \tilde{A}$  is.

HINT. Clearly the projection is analytic, and the condition sufficient. Consider thus an analytic subset  $X$  of  $\tilde{\Omega} \times \tilde{S}$ , with analytic projection  $Y$  on  $\tilde{\Omega} \times \tilde{A}$ , and assume  $Y$  is not negligible: for some strategy  $\tilde{\tau}$ , and for some  $t$ ,  $P_{\tilde{\tau}}((\omega_t, a_t) \in Y) > 0$ . By 7.j, there exists a universally measurable map  $s$  from  $Y$  to  $S$ , with  $(y, s(y)) \in X \forall y \in Y$ . So there exists a Borel subset  $B$  of  $Y$ , such that the restriction of  $s$  to  $B$  is Borel measurable, and such that  $P_{\tilde{\tau}}((\omega_t, a_t) \in B) > 0$ . Let  $\bar{s}(\omega, a)$  be the unit mass at  $s(\omega, a)$  for  $(\omega, a) \in B$ ,  $\bar{s}(\omega, a) = \tilde{\sigma}(\omega, a)$  on  $D \setminus B$  and consider the behavioural strategy  $\tilde{s} = \frac{1}{2}\bar{s} + \frac{1}{2}\tilde{\tau}$ : it has probability  $2^{-t}$  of playing like  $\tilde{\tau}$  before  $t$  and like  $\bar{s}$  at  $t$ , so  $P_{\tilde{s}}((\omega_t, a_t, s_t) \in X) \geq 2^{-t}P_{\tilde{\tau}}((\omega_t, a_t) \in B) > 0$ :  $X$  is not negligible.

E. An analytic subset  $X$  of  $\tilde{\Omega} \times \tilde{A}$  is negligible iff the analytic set  $\{\omega \mid \tilde{\alpha}_\omega(X) > 0\}$  is so.

HINT. 9.f.

F. Negligible analytic subsets of  $\tilde{\Omega} \times \tilde{A} \times S \times \tilde{\Omega}$ , etc. are contained in negligible Borel sets.

HINT. Use II.1Ex.13civB p. 69 and 4.d.3.

G. For any negligible analytic subset  $N$  of  $\tilde{\Omega} \times \tilde{A} \times S \times \tilde{\Omega}$ , there exist Borel restrictions  $\Omega^0 \subseteq \tilde{\Omega}^0$ ,  $C^0 \subseteq \tilde{C}$ ,  $K^0 \subseteq \tilde{K}$  such that  $(\omega, a) \in C^0 \Rightarrow \tilde{\sigma}_{\omega, a}(\tilde{S}_a) = 1$ , and  $K^0 \cap N = \emptyset$ .

REMARK 1.22. It follows that, with those restrictions, the extended model (i.e. with  $\sim$ ) is a true extensive form, with  $S_a \neq \emptyset$  for every feasible  $a$ , and with  $\tilde{\sigma}$  an everywhere defined pseudo-strategy.

HINT. Let  $\tilde{C}_1 = \tilde{C} \cap D$ , and add to  $N$  the complements of  $\tilde{K}$ , of  $\tilde{C}_1 \times S \times \tilde{\Omega}$  and of  $\tilde{\Omega}^0 \times \tilde{S} \times \tilde{\Omega}^0$ , obtaining thus  $\tilde{N}$ .  $\tilde{C}_1$  is still co-analytic, as a Borel subset of the co-analytic  $\tilde{C}$ , and all sets added are negligible by II.1Ex.13civA p. 68. So  $\tilde{N}$  is a negligible analytic subset. Apply then II.1Ex.13civF p. 69 to

get  $K_1$  Borel disjoint from  $\tilde{N}$  and with negligible complement, next II.1Ex.13civC p. 69 to get that the Borel set  $B_1 = \{(\omega, a, s) \in \tilde{\Omega} \times \tilde{S} \mid \tilde{q}_{\omega, a, s}(K_1) < 1\}$  is negligible, so its analytic projection  $P_1$  on  $\tilde{\Omega} \times \tilde{A}$  is (II.1Ex.13civD p. 69) also negligible, hence (II.1Ex.13civF) disjoint from a Borel set  $C_1 \subseteq \tilde{C}_1$  with negligible complement, so finally, by II.1Ex.13civE p. 69,  $\Omega_1 = \{\omega \mid \tilde{\alpha}_\omega(C_1) = 1\}$  is a Borel set with negligible complement — and clearly included in  $\tilde{\Omega}_0$ . Let now  $K_2 = \{(\omega, a, s, \omega') \in K_1 \mid \omega, \omega' \in \Omega_1, (\omega, a) \in C_1\}$  and continue inductively, obtaining decreasing sequences of Borel sets  $K_n, C_n, \Omega_n$  with  $\tilde{\alpha}_\omega(C_n) = 1 \forall \omega \in \Omega_n$ ,  $(\omega, a) \in C_n \Rightarrow [\tilde{\sigma}_{\omega, a}(\tilde{S}_a) = 1 \text{ and } s \in \tilde{S}_a \Rightarrow \tilde{q}_{\omega, a, s}(K_n) = 1], (\omega, a, s, \omega') \in K_{n+1} \Rightarrow \omega' \in \Omega_n$ . Denote their intersection  $K^0, C^0, \Omega^0$ . Observe  $\omega_1 \in \Omega^0$  because  $\Omega^0$  has negligible complement.

v.

A. If we obtain a system of restrictions as in II.1Ex.13civG p. 69 where the perfect recall relations implied by the  $\tilde{\varphi}^i$  are exact for every feasible path, the same will be true for the traces of  $\Omega^0, C^0$  and  $K^0$  on the original spaces  $\Omega$  and  $A$ , so the result will be proved [restricting also  $\tilde{\varphi}^i$  to  $(\tilde{\varphi}^i)_i^{-1}(S \cup \{\iota\})$ ].

B. Because such restrictions are measurable, it remains true (as in ex. II.1Ex.13civA p. 68) that analytic subsets of  $K^0, C^0$ , etc., or the space of feasible paths are negligible iff their trace on the original model is negligible. (And recall from ex. II.1Ex.12 p. 64 that, when  $\mathcal{A}$  is separable,  $\varphi^i(a_{T_{n+1}}) = (a_{T_n}, s_{T_n})$  with probability one.)

C. Therefore, we can henceforth assume II.1Ex.13civG p. 69 has been applied a first time, say with  $N = \emptyset$ , so we are in a true extensive form model  $(\Omega^0, C^0, K^0, \tilde{A}, (\tilde{A}^i)_{i \in \mathbf{I}}, S, \tilde{S}, \tilde{q}, \tilde{\alpha}, (\tilde{\varphi}^i)_{i \in R})$ , which is fully standard Borel, and where (pseudo-) strategies are defined everywhere on  $C^0$  — in particular  $\overline{S}_a \neq \emptyset \forall (\omega, a) \in C^0$ . Similarly  $\tilde{\varphi}^i$  is now defined on the whole of  $\tilde{A}^i$ ,  $\tilde{\alpha}$  on  $\Omega^0$ ,  $\tilde{q}$  on  $(C^0 \times S) \cap (\Omega^0 \times \tilde{S})$ .

We can therefore henceforth drop the  $\sim$  and the superscripts  $^0$ , and will iteratively impose smaller and smaller restrictions, using each time without further reference II.1Ex.13cvB above and II.1Ex.13civG p. 69 — without writing new superscripts.

D. Let  $\Omega^i = \{\omega \mid \alpha_\omega(A^i) = 0\}$  and remove  $\bigcup_i \Omega^i \times A^i$  from  $C$ . If  $\Omega \times A^i$  is negligible, remove it also. For  $i \in R$ , let also  $A_0^i = (\varphi^i)^{-1}(\iota)$ ,  $A_{n+1}^i = (\varphi^i)^{-1}(A_n^i \times S)$ ,  $A_\infty^i = A^i \setminus \bigcup_n A_n^i$ . The  $A_n^i$  are Borel, and  $A_\infty^i$  negligible — remove  $\Omega \times \bigcup_{i \in R} A_\infty^i$  from  $C$ : now  $A^i = \bigcup_n A_n^i$ . Define thus inductively the Borel map  $h^i$  on  $S^i = (A^i \times S) \cap \overline{S}$  by  $h^i(a, s) = (h^i(\varphi^i(a)), a, s)$  — with  $h^i(\iota) = \iota$  to initialise. And let  $f^i(a) = h^i(\varphi^i(a))$ .

d. *The main iteration.* The rest of the proof is to be done for all  $i \in R$  in succession. So  $i$  is fixed henceforth.

i. Let  $M = \bigcup_{n \geq 0} (S^i)^n$  — where every sequence (including the empty) is preceded by  $\iota$ :  $h^i$  and  $f^i$  map  $S^i$  and  $A^i$  into the standard Borel space  $M$ .

Define first inductively the increasing sequences of analytic (9.f) subsets  $D_1^n \subseteq \tilde{D}_1^n \subseteq \Omega$ ,  $D_2^n \subseteq C$ ,  $D_3^n \subseteq (\Omega \times \overline{S}) \cap (C \times S)$  with  $D_j^0 = \tilde{D}_j^0 = \emptyset$ ,  $D_2^n = \text{Proj}(D_3^n)$ ,  $D_1^n = \{\omega \mid \alpha_\omega(D_2^n) > 0\}$ ,  $D_3^{n+1} = [(\Omega \times S^i) \cap (C \times S)] \cup \{(\omega, a, s) \mid q_{\omega, a, s}(D_1^n) > 0\}$ ,  $\tilde{D}_j^n = \text{Proj}(\{(\omega, a, s, \omega') \in K \mid \omega' \in D_1^n\})$ ; ( $j = 1, 2, 3$ ). Like in the proof of II.1Ex.13civD p. 69, observe that, by induction, for every initial probability distribution  $\mu$  on  $D_j^n$ , there exists a pseudo-strategy vector  $\tau$  such that, starting with  $\mu$  and following  $\tau$ , player  $i$  will have to play with positive probability before stage  $n+1$  — and conversely, if  $\mu$  assigns probability zero to  $D_j^n$ , then for any  $\tau$ , player  $i$  will not have to play before stage  $n+1$ .

(\*) In particular, by the same trick, for any non-negligible analytic subset of  $D_j^n$ , there exists a pseudo strategy vector  $\tau$  under which first this subset will be hit with positive probability, and next player  $i$  will with positive probability have to play at most  $n$  stages later.

It follows also that when stronger restrictions  $\Omega, C$  and  $K$  are imposed, the new  $D_j^n$  are just the restrictions of the old ones (while the  $\tilde{D}_j^n$  may shrink).

ii. Define now inductively Borel functions  $\psi_j$  from  $\tilde{D}_j^n$  to  $M$ , with  $\psi_3(\omega, a, s) = (f^i(a), a, s)$  on  $D_3^1$ ,  $\psi_2(\omega, a) = f^i(a)$  on  $D_2^1$ ,  $\psi_2(\omega, a) = \psi_3(\omega, a, s)$  if  $a \notin A^i$  and  $(\omega, a, s) \in \tilde{D}_3^n$ ,  $\psi_1(\omega) = \psi_2(\omega, a)$  for  $(\omega, a) \in \tilde{D}_2^n$ ,  $\psi_3(\omega, a, s) = \psi_1(\omega')$  for  $\omega' \in K_{\omega, a, s} \cap D_1^{n-1}$  — all the time imposing further restrictions  $(\Omega, C, K)$ .

HINT. Assume by induction the  $\psi_j$  are already defined with the above properties on the  $\tilde{D}_j^n$ . Consider  $Q_{(\omega, a, s)}(B) = q_{\omega, a, s}[D_1^n \cap \psi_1^{-1}(B)]$  for any Borel set  $B$  in  $M$ , and prove that  $N = \{(\omega, a, s) \mid Q_{\omega, a, s} \neq 0 \text{ and } Q_{\omega, a, s} \text{ is not concentrated on a single point}\}$  is analytic. [ $D_1^n$  is the projection of a Borel set  $H$  in  $\Omega \times [0, 1]$ . The set of probabilities  $\mu$  on  $\Omega \times [0, 1]$  such that  $\mu(H) > 0$ , and such that the image of  $\mu|_H$  by  $\psi_1 \circ \pi$  ( $\pi$  being the projection) is not concentrated on a single point, is Borel, as the inverse image by a Borel map of the Borel set of positive measures on  $M$  which are not concentrated on a single point. Thus it is standard Borel, so has an analytic projection in the space of probabilities on  $\Omega$ : (using 7.j)  $N$  is the inverse image of this analytic set by the measurable 9.e map  $q$ .] By (\*) and by the properties of  $\psi$  on the  $D_j^n$ , it follows that  $N$  is negligible, so can be neglected. Now  $q_{\omega, a, s}(D_1^n) > 0 \Rightarrow Q_{(\omega, a, s)}(B)$  is concentrated at a single point, say  $g(\omega, a, s)$ . For any Borel set  $B$  in  $M$ ,  $g^{-1}(B) = \{(\omega, a, s) \mid q_{\omega, a, s}[D_1^n \cap \psi_1^{-1}(B)] > 0\}$  is analytic. So its complement  $g^{-1}(B^c)$  is analytic too, so both are Borel in the analytic set  $g^{-1}(M)$  (3.f). Thus  $g$  is a Borel function on its domain. Therefore  $N = \{(\omega, a, s, \omega') \in (g^{-1}(M) \times D_1^n) \cap K \mid g(\omega, a, s) \neq \psi_1(\omega')\}$  is analytic, as a Borel subset of an analytic set, and clearly negligible: we neglect it too. To obtain that  $g$  coincides with  $\psi_3$  on  $D_3^n$ , it suffices now to obtain that  $q_{\omega, a, s}(D_3^n) > 0$ ,  $a \in A^i \Rightarrow g(\omega, a, s) = (f^i(a), a, s)$ . Let thus  $N = \{(\omega, a, s) \in g^{-1}(M) \cap D_3^n \mid g(\omega, a, s) \neq \psi_3(\omega, a, s)\}$ : it is (Borel in) analytic, and, by (\*) and the properties of  $\psi$ , negligible. Neglecting it, we can use  $g$  to extend  $\psi_3$  to  $D_3^{n+1}$ . Consider now the analytic set  $G = \{(\omega, a, m) \mid a \notin A^i, \exists s: (\omega, a, s) \in D_3^{n+1}, \psi_3(\omega, a, s) = m\}$ ; then  $N = \{(\omega, a) \mid \#G_{\omega, a} > 1\}$  is analytic: viewing  $M$  as  $[0, 1]$ , one has  $N = \bigcup_{r \in \mathbb{Q}} \{(\omega, a) \mid \exists m > r, (\omega, a, m) \in G\} \cap \{(\omega, a) \mid \exists m < r, (\omega, a, m) \in G\}$ . For the same reason, there exist, using 7.j, two universally measurable functions  $s_1$  and  $s_2$  defined on  $N$ , with  $(\omega, a, s_i(\omega, a)) \in D_3^{n+1}$  and  $\psi_3(\omega, a, s_1(\omega, a)) \neq \psi_3(\omega, a, s_2(\omega, a))$ . As in (\*), construct now a strategy vector that reaches  $N$  with positive probability, in  $N$  plays Borel modifications of  $s_1$  and  $s_2$  with positive probability each, and afterwards reaches player  $i$  with positive probability in  $n+1$  steps: this is impossible, so  $N$  is negligible. Neglecting it, we obtain that  $G$  is the (analytic) graph of a function  $g$  from  $\{(\omega, a) \in D_2^{n+1} \mid a \notin A^i\}$  to  $M$ , so  $g$  is Borel (3.f). This defines thus  $\psi_2$  on the whole of  $D_2^{n+1}$  — there is clearly no problem of compatibility. To define now  $\psi_1$  on  $D_1^{n+1}$ , use the same argument as used above for  $\psi_3$  on  $D_3^{n+1}$  — only view, for the proof of analyticity,  $\alpha$  as a transition probability from  $\Omega$  to  $\Omega \times A$  instead of to  $A$ ; and notice also that there is no problem of compatibility. Finally, to extend the domains of definition to the  $\tilde{D}_j^{n+1}$ , notice that  $\psi_1$  has a Borel extension to the whole of  $\Omega$  (9.a). Then  $\{(\omega, a) \in D_2^{n+1} \mid \psi_1(\omega) \neq \psi_2(\omega, a)\}$  is (Borel in) analytic, and clearly negligible since for those  $\omega$ ,  $\alpha_\omega(D_2^{n+1}) = 0$ . Neglecting it, we have  $\psi_1(\omega) = \psi_2(\omega, a)$  on  $\text{Proj } D_3^{n+1}$ . Similarly,  $\{(\omega, a, s, \omega') \in K \mid \omega' \in D_1^n, \psi_1(\omega') \neq \psi_1(\omega), a \notin A^i \text{ or } \psi_1(\omega') \neq (\psi_1(\omega), a, s), a \in A^i\}$  is analytic, and is now also negligible, since for those,  $q_{\omega, a, s}(D_1^n) = 0$  (because if  $q_{\omega, a, s}(D_1^n) > 0$ , then both  $\psi_1(\omega) = \psi_2(\omega, a)$  and  $\psi_3(\omega, a, s) = \psi_1(\omega')$  for  $\omega' \in K_{\omega, a, s} \cap D_1^n$ ). Neglecting it too finishes the induction.

iii.

A. Once this induction is over, consider (with the new restrictions  $\Omega, C, K$ ), the increasing sequences of analytic sets  $D_j^n$  and  $\tilde{D}_j^n$ , and denote by  $D_j$  and  $\tilde{D}_j$  their union. The functions  $\psi_j$  are well defined on  $\tilde{D}_j$  ( $\supseteq D_j$ ) — show they are Borel.

HINT. Functions with analytic graph on an analytic space are Borel.

B. We have now that  $D_3 = [(\Omega \times S^i) \cap (C \times S)] \cup \{(\omega, a, s) \mid (\omega, a) \in C, (a, s) \in \overline{S}, q_{\omega, a, s}(D_1) > 0\}$ ,  $D_2 = \text{Proj}_C(D_3)$ ,  $D_1 = \{\omega \mid \alpha_\omega(D_2) > 0\}$ , and  $\tilde{D}_j = \text{Proj}\{(\omega, a, s, \omega') \in K \mid \omega' \in D_1\}$ . And  $\psi_1: \tilde{D}_1 \rightarrow M$  is Borel such that, for  $(\omega, a, s, \omega') \in K$ ,  $\omega' \in D_1$ ,  $\psi_1(\omega') = \psi_1(\omega)$  if  $a \notin A^i$  and  $\psi_1(\omega') = (\psi_1(\omega), a, s)$ ,  $\psi_1(\omega) = f^i(a)$  if  $a \in A^i$ .

e. End of the proof: last iteration.

i. Using II.1Ex.13cvD p. 70, show that  $\text{Proj}[(\Omega \times A^i) \cap C] \subseteq D_1$  and that if  $D_1 \neq \emptyset$ ,  $\omega_1 \in D_1$ .

ii. Show that, if  $\omega_1 \in D_1$ ,  $\psi_1(\omega_1) = \iota$ , and that, along every feasible play, as long as  $\omega_t$  remains in  $\tilde{D}_1$ ,  $\psi_1(\omega_t)$  is the correct sequence of past  $(a_n, s_n)$  ( $n < t, a_n \in A^i$ ). Thus player  $i$  can receive incorrect messages (from  $\varphi^i$ , or from  $\psi_1$ ) only on feasible paths that have first left  $\tilde{D}_1$ , next reentered  $D_1$  (using II.1Ex.13ei) — which they can do only by first passing through  $\tilde{D}_1 \setminus D_1$ .

iii. Denote by  $H_\infty$  the (standard Borel) space of feasible plays. Let  $\eta: H_\infty \times \mathbb{N} \rightarrow M$ , where  $\eta(h, t)$  denotes the correct sequence of past  $(a_n, s_n)$  ( $n < t, a_n \in A^i$ ):  $\eta$  is Borel. Let  $E = \{(h, t) \mid \eta(h, t) \neq \psi_1(\omega_t), \omega_t (= \omega_t(h)) \in \tilde{D}_1\}$ :  $E$  is (Borel in) analytic. Do now inductively for  $t = 3, 4, 5, \dots$  the following, with as inductive assumption  $\{(h, n) \in E \mid n < t\} = \emptyset$  [for  $t = 3$ , this follows from II.1Ex.13eii].

iv. Let  $N_t = \{\omega_t(h) \mid (h, t) \in E\}$ :  $N_t$  is analytic (App.6) and, using II.1Ex.13eii and the inductive assumption,  $N_t \subseteq \tilde{D}_1 \setminus D_1$ . Thus  $L = \{(\omega, a, s, \omega') \in K \mid \omega \in N_t, \omega' \in D_1\}$  is analytic, and negligible (because  $L = \{(\omega, a, s, \omega') \in L \mid q_{\omega, a, s}(D_1) = 0\} \cup \{(\omega, a, s, \omega') \in L \mid (\omega, a) \in D_2, \alpha_\omega(D_2) = 0\}$  since  $N_t \cap D_1 = \emptyset$ ): neglect it. The inductive assumption is now satisfied for  $t + 1$ .

v. At the end of the iterations,  $E = \emptyset$ ; since  $\text{Proj}[(\Omega \times A^i) \cap C] \subseteq D_1$ , this implies that  $\varphi^i$  recalls the correct last information to player  $i$  along every feasible play.

**14. Effectively perfect recall.** Assume player  $i$  has effectively perfect recall, and  $A^i$  is countable. Then every generalised strategy (cf. ex. II.1Ex.10ai p. 61) of player  $i$  is equivalent to a probability distribution over pure strategies.

COMMENT 1.23. Since generalised strategies allow the player to recall all his past signals and moves, the statement is a strong way of expressing both that the terminology “effectively perfect recall” is justified, and that there is nothing more general than the usual mixed strategies for such games. Also, applied to pure strategies, it yields Theorem 1.3 p. 53.

COMMENT 1.24. Player  $i$  being fixed, we will systematically drop the superscript  $i$ , and write for instance  $A$  for  $A^i$ .

a. *Equivalence of generalised strategies and distributions over policies.* Call policy a pure strategy of fictitious player  $i$  who would recall his past signals. Thus the space of policies is  $\Theta = \prod_{n=1}^{\infty} (\prod_{a \in A_n} \bar{S}_a)^{\prod_{t < n} A_t}$  (with the corresponding product  $\sigma$ -field), letting  $A_n$  denote the  $n^{\text{th}}$  copy of  $A$ .

Prove the equivalence, for  $A^i$  countable.

HINT. Given a generalised strategy  $(X_n, \mathcal{X}_n, P_{n-1})_{n=1}^{\infty}$ , where  $P_n$  is a transition probability from  $\prod_{1 \leq t \leq n} (S_{t-1} \times X_t \times A_t)$  to  $S_n \times X_{n+1}$ , with  $S_0 = \{0\}$ , construct by induction probabilities  $Q_n$  on  $Y_n = X_1 \times \prod_{t=1}^n [\prod_{a \in k_t} (\bar{S}_{t,a} \times X_{t+1})]^{\prod_{s \leq t} A_s}$  for  $n = 0, \dots, \infty$ . [Using  $Q_{n-1}, P_n$  and the map  $h_n: Y_{n-1} \times \prod_{t \leq n} A_t \rightarrow \prod_{t \leq n} (S_{t-1} \times X_t \times A_t)$ . Treat the factors corresponding to different elements of  $\prod_{t \leq n} A_t$  as conditionally independent.] Use Ionescu Tulcea’s theorem again to obtain  $Q_\infty$  on  $Y_\infty$ . Take then the marginal distribution of  $Q_\infty$  on  $\Theta$  forgetting the factors  $X_t$ .

b. *A “ $\mu$ -completely mixed” behavioural strategy of the opponents.* Given a probability distribution  $\mu$  over  $\Theta$ , there exists a behavioural strategy vector  $\tau_\mu$  of the opponents such that, for each general strategy  $\tau$  of the opponents there is a  $\mu$ -negligible subset  $N_\tau$  of  $\Theta$  such that  $\forall \theta \notin N_\tau, \forall n, \forall (a_1, \dots, a_n) \in A^n, P_{\theta, \tau}(a_1, \dots, a_n) > 0 \Rightarrow P_{\theta, \tau_\mu}(a_1, \dots, a_n) > 0$ . ( $P(a_1, \dots, a_n)$  is the probability that the  $n$  first signals received by  $i$  are  $a_1, \dots, a_n$  in this order.)

HINT. Let  $\Theta_{a_1, \dots, a_n}^\tau = \{\theta \mid P_{\theta, \tau}(a_1, \dots, a_n) > 0\}$ . If  $\tau_k$  are behavioural strategies, and  $\tau = \sum 2^{-k} \tau_k$ , prove that  $\Theta^\tau \supseteq \bigcup_k \Theta^{\tau_k}$ . Deduce that there exists a behavioural strategy  $\tau_{\mu, a_1, \dots, a_n}$  such that  $\Theta_{a_1, \dots, a_n}^{\tau_{\mu, a_1, \dots, a_n}} = \mu\text{-ess sup}\{\Theta_{a_1, \dots, a_n}^\tau \mid \tau \text{ behavioural strategy}\}$  and again that there exists  $\tau_\mu$  such that  $\forall n, \forall (a_1, \dots, a_n), \forall \tau \text{ behavioural strategy}, \Theta_{a_1, \dots, a_n}^\tau \subseteq \Theta_{a_1, \dots, a_n}^{\tau_\mu}$   $\mu$ -a.e. Extend this conclusion to general strategy vectors  $\tau$  using Fubini’s theorem on the product of  $(\Theta, \mu)$  and the auxiliary spaces of the opponents. Let finally  $N_\tau = \bigcup_{\{\alpha \in \bigcup_n A^n\}} [\Theta_\alpha^\tau \setminus \Theta_\alpha^{\tau_\mu}]$ .

c. *A measurable recall function.* There exists a universally measurable function  $\alpha: \Sigma \times A \rightarrow A \cup \{\iota\}$  such that,  $\forall \sigma \in \Sigma$ ,  $P_{\sigma, \tau_\mu}$ -a.e. one has  $\alpha(\sigma, a_1(\omega)) = \iota$  and  $\alpha(\sigma, a_{n+1}(\omega)) = a_n(\omega)$  for  $n \geq 1$ . ( $\Sigma$  denotes the pure strategy space  $\Sigma^i$  of player  $i$  — thus  $\prod_{a \in A} \overline{S}_a$ . And  $a_n(\omega)$  is the  $n^{\text{th}}$  signal received by player  $i$ .)

HINT. Use von Neumann's selection theorem (7.j) on the measurable set

$$G = \{ (\sigma, f) \in \Sigma \times (A \cup \{\iota\})^A \mid P_{\sigma, \tau_\mu}\text{-a.s.}: f(a_1(\omega)) = \iota, f(a_{n+1}(\omega)) = a_n(\omega) \forall n \geq 1 \}$$

|Show  $G = \bigcap_{\substack{(a, a') \in A^2 \\ n \geq 1}} \left( [\{\sigma \mid P_{\sigma, \tau_\mu}(a_{n+1} = a \text{ and } a_n = a') = 0\}] \times (A \cup \{\iota\})^A \right) \cup [\Sigma \times \{f \mid f(a) = a'\}] \right)$ , still intersected with a similar term for  $a_1$ .|

d. *End of the proof.*

i. Define inductively universally measurable maps  $\alpha^n: \Sigma \times A \rightarrow (A \cup \{\iota\})^{n+1}$  — using  $\alpha(\sigma, \iota) = \iota$  — by  $\alpha^0(\sigma, a) = a$ ,  $\alpha^{n+1}(\sigma, a) = \langle \alpha(\sigma, \alpha^n(\sigma, a)), \alpha^n(\sigma, a) \rangle$  (where  $\langle \cdot, \cdot \rangle$  stands for concatenation). Let  $B_n = (\alpha^n)^{-1}(A^{n+1})$  universally measurable, and define inductively universally measurable maps  $\sigma_n: \Theta \rightarrow \Sigma$  with  $[\sigma_{n+1}(\theta)](a) = \theta_{n+1}[\alpha^n(\sigma_n(\theta), a)]$  if  $(\sigma_n(\theta), a) \in B_n$ ,  $= [\sigma_n(\theta)](a)$  otherwise. Let finally  $[\sigma_\infty(\theta)](a) = \lim_{n \rightarrow \infty} [\sigma_n(\theta)](a)$  if, for all sufficiently large  $n$ ,  $[\sigma_n(\theta)](a) \notin B_n$ ,  $[\sigma_\infty(\theta)](a) = [\sigma_1(\theta)](a)$  otherwise. Show  $\sigma_\infty$  is universally measurable too.

ii. Show by induction over  $n$  that  $\forall \tau, \forall \theta \notin N_\tau, \forall k \geq n-1$ ,  $P_{\sigma_k(\theta), \tau}$  and  $P_{\theta, \tau}$  coincide on histories up to  $a_n$  (the  $n^{\text{th}}$  signal to  $i$ ) — including for  $k = \infty$ .

HINT. Assume the statement for  $n$ . Then  $\forall a \in A^n, \forall \theta \notin N_\tau, \forall k \geq n-1$ ,  $(P_{\theta, \tau}(a) =) P_{\sigma_k(\theta), \tau}(a) > 0 \Rightarrow \theta \in \Theta_a^{\tau_\mu} \Rightarrow \alpha^{n-1}(\sigma_k(\theta), a_n) = a$ , so  $(\sigma_k(\theta), a_n) \in B_{n-1} \setminus \bigcup_{m \geq n} B_m$ , and  $[\sigma_k(\theta)](a_n) = \theta_n(a) \quad \forall k \geq n$ . Conclude.

iii. The image distribution of  $\mu$  by  $\sigma_\infty$  is the required distribution on  $\Sigma$ .

COMMENT 1.25. Observe that the map we obtained from  $\mu$  to  $\sigma_\infty(\mu)$  is neither canonical, nor linear:  $\sigma_\infty$  itself depends on  $\mu$ , through  $\alpha$  and hence  $\tau_\mu$ .

COMMENT 1.26. The proof shows that the apparently weaker definition of “effectively perfect recall” that for each pure strategy  $\sigma$  of  $i$  and every behavioural strategy vector  $\tau$  of the opponents, there exists a corresponding recall function  $\alpha_{\sigma, \tau}$  of player  $i$ , is already sufficient: use  $\tau_\sigma$  for  $\tau$  (cf. II.1Ex.14b p. 72) to obtain that  $\alpha_\sigma$  can be selected independently of  $\tau$ .

COMMENT 1.27. What is the “right” definition of “effectively perfect recall” when  $A^i$  is standard Borel? I.e., that would be equivalent to the present definition for  $A^i$  countable, and for which the present result would generalise for  $(S^i, \mathcal{S}^i)$  standard Borel. (For part II.1Ex.14a p. 72, one can use ex. 7 p. 63).

COMMENT 1.28. Because of this theorem, one can treat as (linear) games with countable  $A^i$  — and hence use remarks 1.4 and 1.5 p. 55 — a number of situations where the true signals do not fall in a countable set, e.g. because they are in fact the full past history of signals and moves — and  $S^i$  is not necessarily countable. This applies in particular to all games as modelled according to remark 1.3 p. 54.

*In the next 5 exercises, we fix a complete probability space  $(\Omega, \mathcal{A}, P)$ .*

**15. Liftings.** A lifting  $\varrho$  is a homomorphism of rings with unit from  $L_\infty(\Omega, \mathcal{A}, P)$  to  $\mathcal{L}_\infty(\Omega, \mathcal{A}, P)$  such that  $\varrho(f) \in f$  (recall the points of  $L_\infty$  are equivalence classes in  $\mathcal{L}_\infty$ ).

a. There exists a lifting.

HINT. Call “linear lifting” a positive linear map from  $L_\infty$  to  $\mathcal{L}_\infty$  such that  $\varrho(f) \in f$  and  $\varrho(1) = 1$ . Clearly every lifting is linear. We first show the existence of a linear lifting. Consider the family of pairs formed by a sub  $\sigma$ -field  $\mathcal{B}$  (containing all null sets) and a linear lifting  $\varrho$  on  $L_\infty(\Omega, \mathcal{B}, P)$ , ordered by  $(\mathcal{B}_1, \varrho_1) \leq (\mathcal{B}_2, \varrho_2)$  iff  $\mathcal{B}_1 \subseteq \mathcal{B}_2$  and  $\varrho_2$  extends  $\varrho_1$ . The family is non-empty (trivial  $\sigma$ -field), and the order is inductive, i.e. any totally ordered subfamily  $(\mathcal{B}_\alpha, \varrho_\alpha)$  is majorated: indeed, either there is no countable co-final set of indices, and then  $\mathcal{B}_\infty = \bigcup_\alpha \mathcal{B}_\alpha$  is a  $\sigma$ -field, with an obvious linear lifting  $\varrho_\infty$ , or there is one, say  $\alpha_i$  — let then  $\mathcal{B}_\infty$  be the  $\sigma$ -field spanned by  $\bigcup_\alpha \mathcal{B}_\alpha$ , and  $\varrho_\infty(f) = \lim_{\mathcal{U}} \varrho_{\alpha_i}[\mathbb{E}(f | \mathcal{B}_{\alpha_i})]$ , where  $\mathcal{U}$  is an ultrafilter on the integers:  $\varrho_\infty$  is again a linear lifting on  $\mathcal{B}_\infty$ , by the martingale convergence theorem. So, by Zorn’s lemma, there exists a maximal pair  $(\mathcal{B}, \varrho)$ . If  $A \in \mathcal{A} \setminus \mathcal{B}$ , construct a linear lifting extending  $\varrho$  on the  $\sigma$ -field spanned by  $\mathcal{B}$  and  $A$ . Hence the existence of a linear lifting  $\varrho$ . Let then  $R$  denote the set of all linear liftings  $\varrho'$  such that  $[\varrho(f)](\omega) = 1$  implies  $[\varrho'(f)](\omega) = 1$  for all  $\omega$  and every indicator function  $f$ .  $R$  is clearly convex; to prove its compactness in  $\mathbb{R}^{\Omega \times L_\infty}$  it suffices to show that for any limit point  $\tilde{\varrho}$ ,  $\tilde{\varrho}(f) \in f$  for all  $f$ . By linearity and positivity (which implies uniform continuity), it suffices to do this for indicator functions, say  $f = \mathbb{1}_A$ . Then  $\{\omega \mid [\varrho(f)](\omega) = 1\} = A$  a.e., so  $\tilde{\varrho}(f) \geq f$  a.e., and similarly  $\tilde{\varrho}(1-f) \geq 1-f$  a.e., hence the result. So by Krein-Milman,  $R$  has an extreme point  $\tilde{\varrho}$ .  $\tilde{\varrho}$  is then still extreme in the convex set of all linear liftings: if a convex combination of two linear liftings belongs to  $R$ , each one does. Hence  $\tilde{\varrho}$  is a lifting: for  $f \in L_\infty$ ,  $0 \leq f \leq 1$ , let  $T_f: L_\infty \rightarrow \mathcal{L}_\infty: g \mapsto \tilde{\varrho}(fg) - \tilde{\varrho}(f)\tilde{\varrho}(g)$ ,  $T^1(g) = \tilde{\varrho}(g) + T_f(g)$ ,  $T^2(g) = \tilde{\varrho}(g) - T_f(g)$ . An immediate computation shows the  $T^i$  are linear liftings, with  $\tilde{\varrho}$  as average. By extremality,  $T_f = 0$ , i.e.  $\tilde{\varrho}(fg) = \tilde{\varrho}(f)\tilde{\varrho}(g)$ . By linearity, this extends then to all  $f$ .

b.  $L_\infty$  is a complete lattice: i.e., if  $(f_\alpha)_{\alpha \in A}$  is a uniformly bounded decreasing net in  $L_\infty$ , it has a greatest lower bound:  $\text{ess inf}_{\alpha \in A} f_\alpha$  in  $L_\infty$ .

HINT. Consider  $\inf_i f_{\alpha_i}$ , where the sequence  $\alpha_i$  is chosen such that  $\inf_i \int f_{\alpha_i} dP = \inf_\alpha \int f_\alpha dP$ .

c. For a lifting  $\varrho$ , and for  $f = \text{ess inf}_{\alpha \in A} f_\alpha$ , one has  $[\inf_{\alpha \in A} \varrho(f_\alpha)] \in f$ .

HINT. The left-hand member is  $\geq \varrho(f)$  and  $\leq \inf_i \varrho(f_{\alpha_i})$ , and both those bounds are measurable with the same integral. Use again completeness.

REMARK 1.29. A lifting maps indicator functions to indicator functions, hence induces a map from subsets to subsets.

d. If  $\Omega$  is a Hausdorff space, and  $P$  a regular probability, call a lifting strong if  $\varrho(f) \leq \hat{f}$  on the support of  $P$  for all  $f \in L_\infty$ , where  $\hat{f}(\omega) = \limsup_{V \in \mathcal{V}_\omega} \frac{1}{P(V)} \int_V f dP$ , and where  $\mathcal{V}_\omega$  denotes the decreasing net of all open neighbourhoods of  $\omega$ .

Observe (regularity) that  $\hat{f}$  is, outside the atoms of  $P$ , the smallest u.s.c. function majorating  $f$  a.e. on  $\text{Supp } P$ , so it is equivalent to require that  $\varrho_\omega(f) \leq f(\omega)$  for  $f$  u.s.c. and  $\omega \in \text{Supp } P$ , or still that  $U \subseteq \varrho(U)$  for every open subspace  $U$  of  $\text{Supp } P$ .

Assume there is a sequence of compact metric subsets  $K_n$  such that  $P(\bigcup_n K_n) = 1$ . Then any (linear) lifting  $\varrho$  is a.e. equal to a strong (linear) lifting  $\bar{\varrho}$ .

HINT. Let  $N = \bigcup_n N_n \cup \complement \bigcup_n K_n$ , where  $N_n = \{\omega \mid \exists f \in C(K_n): \varrho_\omega(\bar{f}) \neq \bar{f}(\omega)$ , where  $\bar{f}(\omega) = f(\omega)$  for  $\omega \in K_n$ ,  $= 0$  else $\}$ .  $N_n$ , hence  $N$ , is negligible, by the separability of  $C(K_n)$ . For  $\omega$  in  $N \cap \text{Supp}(P)$ , let  $\bar{\varrho}_\omega$  be an extreme point (Krein-Milman) of the non-empty (Hahn-Banach) set of all positive linear functionals  $p$  on  $L_\infty$  satisfying  $p(\mathbb{1}_O) = 1$  for every open set  $O$  containing  $\omega$ . By the same argument as sub(1),  $\bar{\varrho}_\omega$  is multiplicative. Let  $\bar{\varrho}_\omega = \varrho_\omega$  elsewhere:  $\bar{\varrho}$  is a lifting, a.e. equal to  $\varrho$ . Remains to show that for  $O$  open,  $\omega \in O \cap (K_n \setminus N_n)$ ,  $\varrho_\omega(O) = 1$ : this follows from the definition of  $N_n$ , selecting  $f \in C(K_n)$  with  $0 \leq f \leq \mathbb{1}_O$  and  $f(\omega) = 1$ .

REMARK 1.30. Regular probabilities on most classical function spaces satisfy the above assumption (cf. e.g. Edwards, 1965), even when compact subsets are not metrisable.

e. A lifting  $\varrho$  on  $(\Omega, A, P)$  operates on a Baire-measurable map  $f$  with relatively compact values in a completely regular space  $E$  by  $h(\varrho(f)) = \varrho(h \circ f)$  for every continuous function  $h$  on  $E$ . Prove the following:

- i. Existence and uniqueness of  $\varrho(f)$ .
- ii.  $\varrho(f)$  depends only on the equivalence class of  $f$  (i.e., all  $g$  such that,  $\forall h, h \circ g = h \circ f$  a.e.), and belongs to it.
- iii.  $\varrho(f)$  is Borel-measurable, and the image measure on the Borel sets is regular.
- iv. If  $f_\alpha$  are such maps to spaces  $E_\alpha$ , and  $f: \prod_\alpha E_\alpha \rightarrow E$  is continuous,  $\varrho(f \circ \prod_\alpha f_\alpha) = f \circ \prod_\alpha (\varrho(f_\alpha))$ .

One way to prove II.1Ex.15eiii and II.1Ex.15eiv is by the following two points:

v.  $\varrho$  defines a map from  $\Omega$  into the (compact) space of characters  $S$  of  $L_\infty$ . The map is by definition Baire-measurable. For the induced probability  $\bar{P}$  on  $S$ , use II.1Ex.15b p. 74 to show that any closed set with empty interior is contained in a closed Baire set with measure zero and also that the closure of an open set is open (i.e., the space is hyperstonian). Deduce that  $\bar{P}$  is “completion regular”: the completion of  $\bar{P}$  on the Baire  $\sigma$ -field is a regular probability on the Borel sets. Hence the map  $\varrho$  is Borel measurable, and induces a regular probability.

vi. A map  $f$  as in II.1Ex.15e induces an algebra-homomorphism from the space of continuous functions  $C(E)$  into  $C(S)$ , hence by transposition a continuous map from  $S$  to  $E$ .  $\varrho(f)$  can equivalently be defined as the composition of this map with the map  $\varrho$  from  $\Omega$  to  $S$ .

## 16. Regular conditional probabilities.

a.

i. Let  $u$  be a positive linear map from  $C(K)$  to  $L_\infty(\Omega, \mathcal{A}, P)$ , where  $K$  is compact. Assume  $u(1) = 1$ . Then there exists a transition probability  $Q$  from  $(\Omega, \mathcal{A}, P)$  to  $(K, \mathcal{B})$ , where  $\mathcal{B}$  is the Borel  $\sigma$ -field on  $K$ , such that, for each  $\omega$ ,  $Q_\omega$  is a regular Borel measure on  $K$ , and such that  $Q(f) \in u(f) \ \forall f \in C(K)$ .

HINT. Let  $Q_\omega(f) = [\varrho(u(f))](\omega)$  for  $f \in C(K)$ , and use Riesz’s theorem. For  $C$  compact in  $K$ , denote by  $f_\alpha$  the decreasing net of all continuous functions  $\geq 1_C$ , and use II.1Ex.15c p. 74 to prove the measurability of  $Q_\omega(C)$  — so  $Q_\omega$  has indeed the required properties.

ii. Those  $Q$ ’s have the additional property that, if  $f_\alpha$  is a decreasing net of u.s.c. functions converging to  $f$ , then  $Q(f) = \text{ess inf}_\alpha Q(f_\alpha)$ .

HINT. Observe that  $Q(f_\alpha) \geq \varrho(Q(f_\alpha))$ .

iii. Using II.1Ex.15e p. 74, one can even require  $Q$  to be a Borel-measurable map to  $\Delta(K)$ , inducing a regular image measure.

b. Conversely, it is clear that any such transition probability  $Q$  defines in this way a unique  $u$  — but it is more instructive to follow the longer route:

i.  $P$  and  $Q$  determine uniquely a probability distribution  $P \otimes Q$  on  $(\Omega, \mathcal{A}) \otimes (K, \mathcal{B})$  — which has  $P$  as marginal.

ii. For any probability distribution  $R$  on  $(\Omega, \mathcal{A}) \otimes (K, \mathcal{B})$  having  $P$  as marginal, let  $u(f) = \mathbf{E}_R(f | \mathcal{A})$  for any  $f \in C(K)$ : this is the required  $u$ .

iii. Although, given  $u$ , many  $Q$ ’s will satisfy the requirements of II.1Ex.16a, they will all induce the same  $R$ .

HINT. By II.1Ex.16aii, for two such  $Q$ ’s,  $Q^1$  and  $Q^2$ , one will have  $Q^1(f) = Q^2(f)$  a.e. for any positive Borel function  $f$  on  $K$ .

II.1Ex.16aii yields also that:

iv. Those  $R$ 's — being in one-to-one correspondence with their restriction to product sets — are exactly the positive measures on  $(\Omega, \mathcal{A})$  with values in the space of regular Borel measures on  $(K, \mathcal{B})$ , such that  $[R(A)](K) = P(A)$ .

c. Conclude that, for any Hausdorff topological space  $K$ , with its Borel  $\sigma$ -field  $\mathcal{B}$ , and any measurable space  $(\Omega, \mathcal{A})$ , any probability  $R$  on the product whose marginal distribution on  $(K, \mathcal{B})$  is regular can be decomposed into its marginal  $P$  on  $(\Omega, \mathcal{A})$  and a regular conditional (in the sense of II.1Ex.16aiii)  $Q$  on  $(K, \mathcal{B})$  given  $(\Omega, \mathcal{A})$ .

HINT. If  $R(\Omega \times B)$  is regular, so is  $R(A \times B)$  for each  $A \in \mathcal{A}$ . Consider first the case  $K$  compact; for arbitrary  $K$ , use the regularity of  $R(\Omega \times B)$  to find a  $K_\sigma$  carrying this measure.

COMMENT 1.32. In typical applications,  $Q$  becomes a behavioural strategy. The equivalent  $R$  is sometimes called a **distributional strategy**.

d. In the same setup, with  $K$  compact, conclude also that any probability measure on  $\mathcal{A} \otimes \mathcal{B}_a$  — where  $\mathcal{B}_a$  denotes the Baire  $\sigma$ -field, spanned by the continuous functions — has a (clearly unique) regular extension to  $\mathcal{A} \otimes \mathcal{B}$ .

### 17. Convergence of transition probabilities.

a. Tychonoff's theorem yields that the point-wise convergence topology (with the weak<sup>\*</sup> topology on  $L_\infty$ ) is compact on the set of  $u$ 's defined sub II.1Ex.16b.

b. Define, for regular transition probabilities  $Q$  from  $(\Omega, \mathcal{A}, P)$  to  $(K, \mathcal{B})$  (i.e.,  $\forall \omega, Q_\omega$  is a regular Borel probability, and  $\forall B \in \mathcal{B}, Q_\omega(B)$  is measurable), the convergence concept  $Q^\alpha \rightarrow Q$  iff  $Q_\omega^\alpha(f) \rightarrow Q_\omega(f)$   $\sigma(L_\infty, L_1)$ ,  $\forall f \in C(K)$  — i.e., iff  $\forall f \in C(K), \forall A \in \mathcal{A}, \int_A Q_\omega^\alpha(f)P(d\omega) \rightarrow \int_A Q_\omega(f)P(d\omega)$ . Conclude that this convergence concept derives from a compact Hausdorff topology on (equivalence classes of)  $Q$ 's.

c. If  $K$  is metrisable, this topology is also strictly semi-compact (ex. I.1Ex.7c p. 12). If in addition  $\mathcal{A}$  is separable, the topology is metrisable.

### 18. Operator interpretation of Liftings.

a.

i. The assumptions sub II.1Ex.16a imply that  $u \in \mathcal{L}(C(K), L_\infty)$ , i.e.  $u$  is a (norm)-continuous linear map from  $C(K)$  to  $L_\infty$ .

ii. For any such map  $u$ , there exists a kernel  $Q$  — i.e.,  $\forall \omega, Q_\omega$  is a regular Borel measure on  $K$ , and  $\forall B \in \mathcal{B}, Q_\omega(B)$  is measurable —, such that  $\|Q_\omega\| \leq \|u\| \forall \omega, Q_\omega(f) \in u(f) \forall f \in C(K)$ , and the marginal on  $K$  of  $|P \otimes Q|$  is a regular Borel measure.

HINT. Proceed as sub II.1Ex.16ai. For the last point, which corresponds to II.1Ex.16aii, define  $R^+ = [P \otimes Q]^+$ ,  $R^- = [P \otimes Q]^-$ ,  $u^+(f)$  as the Radon-Nikodym derivative of  $\mu^+(B) = \int_{B \times K} f dR^+$  with respect to  $P$  and similarly for  $u^-(f)$ .  $u^+$  and  $u^-$  are positive linear maps from  $C(K)$  to  $L_\infty$ , with  $u = u^+ - u^-$ . Hence we get  $Q^+$  and  $Q^-$ , each one verifying II.1Ex.16aiii, and, by the linearity of the lifting,  $Q = Q^+ - Q^-$  — so  $P \otimes Q^+ + P \otimes Q^- \geq |P \otimes Q|$ .

iii. Since the converse to II.1Ex.18aii is obvious (cf. hint above), we get an isometry between  $\mathcal{L}(C(K), L_\infty)$  and a space  $L_\infty^{[C(K)]'}$  of equivalence classes of kernels satisfying the properties sub II.1Ex.16aiii. We use this notation because the space of regular Borel measures can be seen as the dual  $[C(K)]'$  of  $C(K)$ , and a kernel as a bounded (“scalarly”-) measurable map with values in this space (in fact, even Pettis-integrable when taking the space of bounded universally measurable functions on  $K$  as dual of the space of regular measures) (or again, using II.1Ex.15e, as a Borel map to  $[C(K)]'$ , with regular image). The equivalence classes are derived from the norm, where  $\|Q\| = \sup\{\|Q_\omega(f)\|_\infty \mid f \in C(K), \|f\| \leq 1\}$ .

b.

- i. Any kernel  $Q_\omega$  with  $\sup_\omega \|Q_\omega\| < \infty$  defines a continuous linear functional  $\varphi_u$  on the space  $L_1^{C(K)}$  of (equivalence classes of) Bochner-integrable functions with values in  $C(K)$  — by  $\varphi_u(f) = \int f d(P \otimes Q)$ . And  $\|\varphi_u\| \leq \sup_\omega \|Q_\omega\|$ .

HINT. Bochner-integrable functions  $f$  are such that  $\|f - f_n\| \rightarrow 0$ , for measurable step functions  $f_n$  with values in  $C(K)$ , and where  $\|g\| = \int^* \|g(\omega)\| P(d\omega)$  (where  $\int^*$  denotes the upper integral). Deduce (Egorov), by selecting  $\|f - f_n\|$  summable, that one can have in addition  $\|f_n(\omega) - f(\omega)\| \rightarrow 0$   $P$ -a.e., so in particular  $f$  is measurable on  $\Omega \times K$ , and similarly  $Q_\omega(f(\omega))$  is measurable and  $\|f(\omega)\|$  is even integrable since  $\|f(\omega)\| \leq \|f_n(\omega)\| + \|f(\omega) - f_n(\omega)\|$ . Since  $\|Q_\omega\| \leq K < \infty$ , we have then also  $|Q_\omega(f(\omega))| \leq K \|f(\omega)\|$ , so  $Q_\omega(f(\omega))$  is integrable. Let then  $\varphi_u(f) = \int Q_\omega(f(\omega)) P(d\omega)$ : linearity is obvious, and  $|\varphi_u(f)| \leq \int K \|f(\omega)\| P(d\omega) = K \|f\|$ , so  $\|\varphi_u\| \leq K$ . We have then also  $\varphi_u(f_n) = \int f_n d(P \otimes Q)$  for all  $n$ , hence the  $f_n$  form a Cauchy sequence in  $L_1(P \otimes Q)$ , that converges point-wise to  $f$ :  $f \in L_1(P \otimes Q)$ , so  $\varphi_u(f) = \int f d(P \otimes Q)$ .

- ii. Any continuous linear functional  $\varphi$  on  $L_1^{C(K)}$  defines a continuous linear map  $u: C(K) \rightarrow L_\infty$ , with  $\|u\| \leq \|\varphi\|$ .

HINT.  $u(x) \in L_\infty: L_1 \rightarrow \mathbb{R}$ ,  $f \mapsto \varphi(xf)$ .

- iii. The maps sub II.1Ex.18bi and II.1Ex.18bii (under the identification sub II.1Ex.18a above of  $u$  and  $Q$ ) define an isometry between  $[L_1^{C(K)}]'$  and  $\mathcal{L}(C(K), L_\infty)$ .

HINT.  $u \xrightarrow{(A)} \varphi \xrightarrow{(B)} u$  is obviously the identity. That  $\varphi \xrightarrow{(B)} u \xrightarrow{(A)} \varphi$  is also, is sufficient to check on step functions, hence on functions  $\mathbb{1}_A \cdot x$  ( $A \in \mathcal{A}, x \in C(K)$ ). The isometric aspect follows then because we have shown (using also II.1Ex.18aiii) that all maps decrease the norm.

- iv. The isometries sub II.1Ex.18aiii and II.1Ex.18biii, given II.1Ex.18ai, allow to view (equivalence classes of) transition probabilities as a subset of the dual of  $L_1^{C(K)}$ . Show that the topology introduced sub II.1Ex.17 p. 76 is the weak\* topology.

c.

- i. Any Banach space  $E$  can be viewed isometrically as a subspace of a space  $C(K)$ , taking for  $K$  the unit ball of the dual (Banach Alaoglu).

ii. Deduce from the above (and the Hahn-Banach theorem) the existence of (canonical) isometries between  $(L_1^E)'$ ,  $L_\infty^{E'}$  and  $\mathcal{L}(E, L_\infty)$ .  $[L_\infty^{E'}]$  is the set of equivalence classes of bounded maps  $f$  from  $\Omega$  to  $E'$ , which are scalarly measurable in the sense that,  $\forall x \in E$ ,  $\langle f(\omega), x \rangle$  is measurable, (or, which are Borel-measurable with regular image), and with  $\|f\| = \sup_{\|x\| \leq 1} \|\langle f(\omega), x \rangle\|_\infty$ .

HINT. Consider the barycentre of  $Q_\omega$ .

A direct proof, from II.1Ex.15 p. 73, is just as easy, considering the embeddings  $L_1 \rightarrow (x \cdot L_1)$  of  $L_1$  into  $L_1^E$ , for all  $x \in E$ .

Show that those spaces are also isometric to the space  $B(L_1, E)$  of continuous bi-linear functionals on  $L_1 \times E$ .

- iii. Show that  $L_1^E$  is isometric to the completed tensor product  $L_1 \hat{\otimes} E$  [for  $z \in E \otimes F$ ,  $\|z\| = \inf\{ \sum_{i=1}^n \|x_i\| \cdot \|y_i\| \mid x_i \in E, y_i \in F, \sum_{i=1}^n x_i \otimes y_i = z \}$ ].

HINT. The “step functions” are dense in  $L_1 \otimes E$ . Also  $L_1^E$  is complete by the usual argument: if  $f_n$  is a Cauchy sequence, extract a subsequence with  $\sum_n \|f_{n+1} - f_n\| < \infty$  (cf. hint sub II.1Ex.18bi p. 77). So both spaces can be viewed as a completion of the step functions: it suffices to prove there is an isometry on step functions.

iv. So the only additional information we got in the case of  $E = C(K)$  is that  $\forall \tilde{f} \in L_\infty^{E'}$  there exists  $f \in \tilde{f}$  [with  $\|f(\omega)\| \leq \|\tilde{f}\| \forall \omega$  and] which is “scalarly well integrable” even when taking as dual of  $E'$  all bounded universally measurable functions.

Show also (cf. II.1Ex.18aii p. 76) that when  $E = C(K)$ , our isometries are also isomorphisms of complete Banach lattices. And an additional isometry is obtained with the space of measures on  $\Omega \times K$  whose absolute value has a regular projection on  $K$  and a projection on  $\Omega$  majorated by  $\lambda \cdot P$  for some  $\lambda > 0$ . Observe finally that, if  $K$  is metrisable,  $\|Q\| = \|(\|Q_\omega\|)\|_\infty$ .

v. Extend II.1Ex.18cii and II.1Ex.18ciii above to the case where  $E$  is a locally convex space. (Then  $L_1^E$  is endowed with the semi-norms  $\bar{p}(f) = \int^* p(f(\omega))P(d\omega)$ , for every continuous semi-norm  $p$  on  $E$ . And the polars of the neighbourhoods of zero in  $E$  are the equicontinuous subsets of  $E'$ , which play therefore the rôle of the balls above).

**19. Strassen's Theorem.** Assume  $(\Omega, A, P)$  complete, and for each  $\omega$  let  $p_\omega$  be a sublinear functional (i.e.,  $p_\omega(\lambda x) = \lambda p_\omega(x)$  and  $p_\omega(x+y) \leq p_\omega(x) + p_\omega(y)$  for  $\lambda \geq 0, x, y \in E$ ) on a Banach space  $E$ , which is bounded —  $\sup_\omega \sup_{\|x\| \leq 1} p_\omega(x) < \infty$  — and weakly measurable —  $p_\omega(x)$  is measurable  $\forall x \in E$ . Let also  $\bar{p}(x) = \int p_\omega(x)P(d\omega)$ , and assume further that either  $E$  is separable, or  $P$  is a regular Borel probability on the Hausdorff space  $\Omega$ , such that  $P(\bigcup_n K_n) = 1$  for appropriate compact metric subsets  $K_n$ , and  $p_\omega(x)$  is upper semi-continuous  $\forall x \in E$ .

Then any linear functional  $\bar{\varphi}$  on  $E$  which is majorated by  $\bar{p}$  can be written as  $\bar{\varphi}(x) = \int \varphi_\omega(x)P(d\omega)$ , where  $\varphi$  is a bounded, Borel-measurable map to  $E'$  with the weak\* topology, having a regular image measure, and satisfying  $\varphi_\omega(x) \leq p_\omega(x)$  everywhere.

HINT. Define  $\tilde{p}$  on  $L_1^E$  (ex. II.1Ex.18c p. 77) by  $\tilde{p}(f) = \int p_\omega(f(\omega))P(d\omega)$ . Identifying  $E$  with the constant functions in  $L_1^E$ ,  $\tilde{p}$  extends  $\bar{p}$ . By the separation theorem 1.20 p. 8, take an extension  $\tilde{\varphi}$  of  $\bar{\varphi}$  with  $\tilde{\varphi} \leq \tilde{p}$ : then (ex. II.1Ex.18cii p. 77)  $\tilde{\varphi}(f) = \int \varphi_\omega(f(\omega))P(d\omega)$  in particular, for  $x \in E$  and  $A \in \mathcal{A}$  we get, with  $f = x1_A$ ,  $\int_A \varphi_\omega(x)P(d\omega) \leq \int_A p_\omega(x)P(d\omega)$ , hence  $\varphi_\omega(x) \leq p_\omega(x)$  a.e. If  $E$  is separable, use a dense sequence to find that, outside a negligible set,  $\varphi_\omega(x) \leq p_\omega(x) \forall x \in E$ , and use the separation theorem to redefine  $\varphi_\omega$  on this negligible set such as to have the inequality everywhere. In the other case, assume by (ex. II.1Ex.15d p. 74) that  $\varphi$  was obtained by a strong lifting  $\varrho$ : then we obtain  $\forall x, \forall \omega: \varphi_\omega(x) = \varrho(\varphi_\omega(x)) \leq \varrho(p_\omega(x)) \leq p_\omega(x)$ , the last inequality because  $p_\omega(x)$  is u.s.c. and the lifting is strong.

**20. The Blackwell-Stein-Sherman-Cartier-Fell-Meyer-Ionescu-Tulcea Theorem.** (Ionescu Tulcea and Ionescu Tulcea, 1969)

a. Let  $X$  and  $Y$  be two compact Hausdorff spaces,  $\varphi: X \rightarrow Y$  a continuous map,  $\mathcal{S} \subseteq C(Y)$  a convex cone such that  $1 \in \mathcal{S}$  and  $f, g \in \mathcal{S} \Rightarrow \min(f, g) \in \mathcal{S}$ . Let  $\nu$  and  $\mu \neq 0$  be regular non negative measures on  $X$  and  $\varphi(X)$  respectively, such that  $\int_X (f \circ \varphi)d\nu \leq \int_Y f d\mu \forall f \in \mathcal{S}$ . Assume  $\mu(Y \setminus \bigcup_n K_n) = 0$  for appropriate compact metric subsets  $K_n$ . Then there is a Borel-measurable map  $\lambda: y \mapsto \lambda_y$  from  $Y$  with the  $\mu$ -measurable sets to the space  $\mathcal{M}^+(X)$  of non-negative regular Borel measures on  $X$  with the weak\* topology, such that  $\mu$  has a regular image measure by  $\lambda$  with  $\nu(B) = \int \lambda_y(B)\mu(dy)$  for every Borel set  $B$ , and  $\lambda_y(f \circ \varphi) \leq f(y), \forall y \in \varphi(X), \forall f \in \mathcal{S}$ .

HINT. Letting  $\lambda_y = 0$  for  $y \notin \varphi(X)$  we can assume  $\varphi$  onto. Let then  $p_y(f) = \inf\{h(y) \mid h \in \mathcal{S}, h \circ \varphi \geq f\}$  for  $f \in C(X)$ , and  $\bar{p}(f) = \int_Y p_y(f)\mu(dy)$ . Observe (regularity of  $\mu$ ) that  $\bar{p}(f) \geq \nu(f)$ , and apply Strassen's theorem ( $\lambda_y \geq 0$  follows from  $\lambda_y(f) \leq p_y(f)$  for  $f \leq 0$ ).

b. For any Borel-measurable map  $\lambda: \omega \mapsto \lambda_\omega$  from a complete probability space  $(\Omega, \mathcal{A}, P)$  to  $\Delta(X)$ , where  $X$  is a Hausdorff  $K_\sigma$ , such that  $P$  has a regular image measure by  $\lambda$ , let  $\bar{P}(B) = \int \lambda_\omega(B)P(d\omega)$  for every Borel set  $B$ . Then  $\bar{P}$  is well defined, and  $\bar{P} \in \Delta(X)$ .

REMARK 1.33. The theorem is II.1Ex.20a. II.1Ex.20b is a comment on it — c.f. e.g. the proof of the claim in theorem 1.2 p. 111 below for an example of its use.

REMARK 1.34. A classical case of this theorem is where  $\varphi$  is the identity and more specifically  $X = Y$  is compact, convex in some locally convex space, and  $\mathcal{S}$  is the cone of concave functions.

This case is closely related to the concavification operators to be encountered later in ch. V and VI.

Another is the following, which gives the “right form” (i.e. as in ex. II.1Ex.9 p. 60, except for the parameterisation by an auxiliary measurable space) of ex. II.1Ex.16 p. 75.

**21. Disintegration of measures.** Let  $\varphi: X \rightarrow Y$  be a continuous map between Hausdorff spaces,  $\nu \in \Delta(X)$ ,  $\mu = \varphi(\nu)$  (1 p. 7). Assume that  $\mu(\bigcup_n K_n) = 1$  for appropriate compact metric subsets  $K_n$ . Then there is a Borel measurable map  $\lambda: y \mapsto \lambda_y$  from  $Y$  with the  $\mu$ -measurable sets to  $\mathcal{M}^+(X)$  with the weak\* topology such that  $\lambda_y \in \Delta(X)$  for  $y \in \varphi(X)$ ,  $\text{Supp } \lambda_y \subseteq \varphi^{-1}(y)$   $\forall y \in Y$ , and  $\mu$  has a regular image measure on  $\Delta(X)$  by  $\lambda$ , with  $\nu(B) = \int_Y \lambda_y(B) \mu(dy)$  for every Borel set  $B$ .

HINT. For  $X$  compact, this follows straight from ex. II.1Ex.20 p. 78, with  $\mathcal{S} = C(Y)$ . The only problem here is to preserve the measurability properties of  $\lambda$  in the generalisation. By regularity, there is no loss to assume the sequence  $K_n$  to be disjoint, and that  $\forall n K_n \subseteq \varphi(X)$ , with  $\mu(K_n) > 0$ ;  $\lambda_y$  can then be defined separately on each  $K_n$ , selecting an arbitrary point mass in  $\varphi^{-1}(y)$  for  $y \in \varphi(X) \setminus \bigcup_n K_n$ , and setting  $\lambda_y = 0$  for  $y \notin \varphi(X)$ . After replacing  $X$  by  $\varphi^{-1}(K_n)$  and renormalising  $\mu$  (9.b.2 p. 428), this reduces, the problem to the case where  $\varphi$  is onto, and  $Y$  is compact metric.

Similarly select a disjoint sequence of compact subsets  $X_n$  of  $X$ , with  $\alpha_n \nu_n = \nu|_{X_n}$ ,  $\nu_n \in \Delta(X_n)$ ,  $\alpha_n > 0$ ,  $\sum_n \alpha_n = 1$ ,  $X_n = \text{Supp } \nu_n$ . Repeating the above argument, we can assume that  $\varphi(\bigcup_n X_n) = Y$ , and hence that  $X$  is the disjoint union of the  $X_n$ . Finally let  $\mu_n = \varphi(\nu_n)$ ,  $f_n = \alpha_n d\mu_n/d\mu$ :  $\sum_n f_n = 1$  a.e. Use Lusin’s theorem to find disjoint compact subsets  $Y_k$  of  $Y$ , with  $\mu(\bigcup_k Y_k) = 1$ , and such that  $f: y \mapsto (f_n(y))_{n=1}^\infty$  is continuous from  $Y_k$  to  $\Delta(\mathbb{N})$ . Repeating our first argument, it suffices to work on each  $Y_k$  separately — hence we can further assume  $f: Y \rightarrow \Delta(\mathbb{N})$  is continuous. Fix now a strong lifting  $\varrho$  on  $(Y, \mu)$ . Observe  $O_n = \{y \mid f_n(y) > 0\}$  is open, and that  $\mu_n$  and  $\mu$  are mutually absolutely continuous on  $O_n$ , while  $\mu_n$  vanishes outside. So  $\varrho$  is still a strong lifting on  $(O_n, \mu_n)$ , and extends thus to a strong lifting  $\varrho_n$  on  $(\bar{O}_n, \mu_n)$ . Since  $\varphi(X_n) = \bar{O}_n$  [because  $X_n = \text{Supp } \nu_n$ ], use the compact case mentioned above to construct, with  $\varrho_n$ , a regular conditional  $\lambda_n$  from  $\bar{O}_n$  to  $\Delta(X_n)$ . Let  $\bar{\lambda}_n(y) = \lambda_n(y)$  for  $y \in O_n$ ,  $\bar{\lambda}_n(y) = \nu_n$  for  $y$  outside  $O_n$ : now  $\bar{\lambda}_n$  is  $\mu$ -measurable from  $Y$  to  $\Delta(X_n)$ . Use ex. II.1Ex.15e p. 74 to let  $\tilde{\lambda}_n = \varrho(\bar{\lambda}_n)$  — in particular,  $\tilde{\lambda}_n(y) = \lambda_n(y)$  for  $y \in O_n$ . Since also  $f = \varrho(f)$ , ex. II.1Ex.15e yields that the map  $y \mapsto [f(y), (\tilde{\lambda}_n(y))_{n \in \mathbb{N}}]$  from  $(Y, \mu)$  to  $\Delta(\mathbb{N}) \times \prod_n \Delta(X_n)$  with the Borel sets is measurable, with a regular image measure. Since the map  $\beta, (\mu_n)_{n \in \mathbb{N}} \mapsto \sum_n \beta_n \mu_n$  from  $\Delta(\mathbb{N}) \times \prod_n \Delta(X_n)$  to  $\Delta(X)$  is continuous, it follows that the composition  $y \mapsto \lambda_y = \sum_n f_n(y) \tilde{\lambda}_n(y) = \sum_n f_n(y) \lambda_n(y)$  is measurable from  $(Y, \mu)$  to  $\Delta(X)$  with the Borel sets, with a regular image measure [and even satisfies  $\lambda = \varrho(\lambda)$ ]. It clearly satisfies our requirements.

**22. Games with almost perfect information.** (Birch, 1955) Consider a game where the position at each stage in a feasible play determines all the previous positions along the play, i.e. the history. An history  $h$  induces a **subgame** if for any feasible play consistent with  $h$ , any signal  $a_i$  to some player  $i$  on this play after (and including)  $h$ , will identify  $h$ . A game has **almost perfect information** for player  $i$  if every history where he is playing induces a subgame, as well as the histories ending by one of his moves.

Show that a game with finitely many plays has an equilibrium where the players who have almost perfect information use pure strategies.

## 2. Infinite Games

**2.a. Infinite games with perfect information.** These games can be described as follows. We are given a set  $\Omega$  (with discrete topology). Player I chooses a point  $\omega_1$  in  $\Omega$  and this choice is told to player II who now selects some  $\omega_2$  in  $\Omega$ .  $\omega_2$  is then announced to player I who chooses  $\omega_3$  and so on. Hence both players play sequentially knowing all previous moves. The game  $\Gamma$  corresponds thus to an infinite tree (with perfect information), that we identify with  $H = \bigcup_{n \geq 0} \Omega^n$ , set of histories or positions, with  $\Omega^0 = \{\emptyset\}$ .  $H_\infty = \Omega^\infty$  is the set of plays. We shall write  $\prec$  for the natural partial order on  $H \cup H_\infty$

where  $h \prec h'$  iff  $h'$  extends  $h$ , and a base of open sets in  $H_\infty$  (for the product topology) is given by the sets  $\{h' \in H_\infty \mid h \prec h'\}$ ,  $h \in H$ . As for the pay-off, we are given a real valued function  $g$  on  $H_\infty$ . A pair of pure strategies  $(\sigma, \tau)$  in  $\Sigma \times \mathcal{T}$  (i.e. mappings from  $H$  to  $\Omega$ ) induces in a natural way  $(h_1 = \sigma(\phi), h_{2n+1} = (h_{2n}, \sigma(h_{2n})), h_{2n+2} = (h_{2n+1}, \tau(h_{2n+1})))$ , a point  $h_\infty = (\omega_1, \omega_2, \dots, \omega_n, \dots)$  in  $H_\infty$  and we define  $\gamma(\sigma, \tau) = g(h_\infty)$ . Given  $h$  in  $H$ , the **subgame** starting from  $h$  is played like  $\Gamma$ , I (resp. II) moving first if  $h$  has even (resp. odd) length and the play  $h_\infty$  induces the pay-off  $g(h, h_\infty)$ , (where  $(h, h_\infty)$  is  $h$  followed by  $h_\infty$ ).  $H(\sigma)$  is the set of positions that can be reached if player I plays  $\sigma$  and  $H_\infty(\sigma)$  is the corresponding set of plays.  $H' \subseteq H$  is a I-sub-tree iff there exists  $\sigma$  of player I with  $H(\sigma) \subseteq H'$ . We then say that  $\sigma$  is compatible with  $H'$ .

The game  $\Gamma$  is open, closed,  $G_\delta, \dots$ , Borel, when  $g$  is the indicator function of an open, closed,  $G_\delta, \dots$ , Borel subset  $W$  of  $H_\infty$ .

In this framework, the game is **determined** if one of the players has a **winning** pure strategy; i.e. either player I can force  $h_\infty$  to be in  $W$  (there exists  $\sigma$  such that  $H_\infty(\sigma) \subseteq W$ ) or player II can force  $h_\infty$  to belong to its complement  $W^c$ .

The first result is

**PROPOSITION 2.1. (Gale and Stewart, 1953)** *Open and closed games are determined.*

**PROOF.** Assume first  $W$  is open. Suppose that player I has no winning strategy in  $\Gamma$ . This implies that, for every  $\omega_1$ , there exists  $\omega_2$  such that player I has still no winning strategy in the subgame starting from  $(\omega_1, \omega_2)$  (i.e. in which a play  $h_\infty$  induces a pay-off  $g(\omega_1, \omega_2, h_\infty)$ ). This defines inductively, for every  $n$  and every position  $h_{2n+1} = (\omega_1, \dots, \omega_{2n+1})$  a move  $\omega_{2n+2}$  such that player I has still no winning in the subgame following  $\omega_1, \dots, \omega_{2n+2}$ . Define  $\tau$  to be a strategy of player II that makes the above choice of  $\{\omega_{2n+2}\}_{n \geq 0}$ . Then  $\tau$  is winning. Otherwise, for some  $\sigma$ ,  $(\sigma, \tau)$  would generate a play  $h_\infty$  in  $W$ .  $W$  being open, there exists  $n$  in  $\mathbb{N}$  such that  $(h_{2n+1}, h'_\infty)$  belongs to  $W$  for every  $h'_\infty$  in  $H_\infty$ , contradicting the choice of  $\omega_{2n+2}$ .

Similarly, if  $W$  is closed and player II has no winning strategy, there exists a first choice  $\omega'_1$  of player I such that player II has still no winning strategy in the subgame following  $\omega'_1$ . Reversing the rôle of the players implies that this subgame is open, hence determined so that player I has a winning strategy  $\sigma'$  for it. It follows then that  $\sigma$  defined by: play  $\omega'_1$  then use  $\sigma'$ , is a winning strategy for player I in the original game. ■

**REMARK 2.1.** The above result can be written in the following equivalent way: Given  $W$  closed in  $H_\infty$ , let  $W^\prec = \{h \in H \mid \exists h' \in W, h \prec h'\}$ . Let  $L^i(W)$ ,  $i = \text{I, II}$ , be the set of positions from where player  $i$  can force the set  $W$ . If  $\emptyset \in L^i(W)$ , then  $L^i(W)$  is included in  $W^\prec$  and is a  $i$ -winning sub-tree, i.e., every strategy compatible with  $L^i(W)$  is winning.

**PROPOSITION 2.2. (Gale and Stewart, 1953)** *Under the Axiom of Choice, there exists an undetermined game.*

**PROOF.** Take  $\Omega = \{0, 1\}$ . Note first that, for each  $\sigma$  of player I, the cardinality of the set  $H_\infty(\sigma)$  is  $\wp(\aleph_0)$ . In fact, player II can follow on his stages any sequence of 0 and 1; on the other hand, one obviously has  $\#\mathcal{T} \leq \wp(\aleph_0)$ . Let us now construct a winning set  $W$ . Let  $\alpha$  be the least ordinal such that there are  $\wp(\aleph_0)$  ordinals less than  $\alpha$ . We can thus index the players' strategies as  $\sigma_\beta, \tau_\beta$  with  $\beta \in \{\gamma \mid \gamma < \alpha\}$ . Choose  $y$  in  $H_\infty(\tau_0)$  and then  $x$  in  $H_\infty(\sigma_0)$  with  $x \neq y$ . Inductively, if  $y_\gamma$  and  $x_\gamma$  have been chosen for  $\gamma < \beta < \alpha$ , the set  $\{x_\gamma \mid \gamma < \beta\}$  has less than  $\wp(\aleph_0)$  elements hence  $H_\infty(\tau_\beta) \setminus \{x_\gamma \mid \gamma < \beta\}$  is not empty.

Choose  $y_\beta$  in it. Similarly  $H_\infty(\sigma_\beta) \setminus \{y_\gamma \mid \gamma \leq \beta\}$  is not empty and we take  $x_\beta$  in it. We now claim that the sets  $X = \{x_\beta \mid \beta < \alpha\}$  and  $Y = \{y_\beta \mid \beta < \alpha\}$  are disjoint. Assume not and let  $x_\beta = y_\gamma$ . If  $\gamma \leq \beta$  (resp.  $\gamma > \beta$ ), we have a contradiction by the choice of  $x$  (resp.  $y$ ). Choose finally  $W$  with  $Y \subseteq W$  and  $W \cap X = \emptyset$ . Consider now  $\sigma$  in  $\Sigma$ . Then  $\sigma$  corresponds to some index, say  $\beta$ . By construction, there exists a play  $x_\beta$  in  $H_\infty(\sigma_\beta) \cap X$  hence a strategy  $\tau$  of player II inducing against  $\sigma$  a play in  $X$ . Player I has no winning strategy, moreover the maxmin (supinf) is 0. Similarly, player II has no winning strategy and the minmax is 1.  $\blacksquare$

Prop. 2.1 has been gradually extended to larger classes of games. The more general result is (Martin, 1975, 1985):

**THEOREM 2.3.** *Borel games are determined.*

We first introduce some notations and definitions.

Given a tree  $H$ , we denote also by  $[H]$  the corresponding set of plays (that was also previously denoted  $H_\infty$ ).

$(H^*, \pi, \phi)$  is a **covering** of the tree  $H$  if

- (1)  $H^*$  is a tree,  $\pi$  is a mapping from  $[H^*]$  to  $[H]$  and  $\phi$  a mapping from  $\Sigma^*$  (resp.  $\mathcal{T}^*$ ) to  $\Sigma$  (resp.  $\mathcal{T}$ ) such that  $\pi$  and  $\phi$  commute in the following sense: if  $h \in [H(\phi(\sigma^*))]$ , there exists  $h^* \in [H^*(\sigma^*)]$  with  $\pi(h^*) = h$ , and similarly for  $\tau^*$ .

**LEMMA 2.4.** *Let  $(H^*, \pi, \phi)$  be a covering of  $H$  and  $W \subseteq [H]$ . If the game  $(H^*, \pi^{-1}(W))$  is determined, so is the game  $(H, W)$ .*

**PROOF.** In fact, let  $\sigma^*$  be a winning strategy in  $(H^*, \pi^{-1}(W))$  and let  $\sigma = \phi(\sigma^*)$ . If  $h$  is a play in  $[H]$  compatible with  $\sigma$ , there exists, by 1, a play  $h^*$  in  $[H^*]$  compatible with  $\sigma^*$  and with image  $h$ . Since  $h^*$  belongs to  $\pi^{-1}(W)$ ,  $\sigma$  is winning.  $\blacksquare$

The idea of the proof of the theorem is roughly to prove inductively that one can construct, for any Borel set  $W$ , a covering with  $\pi^{-1}(W)$  closed. In fact, we need to work with more specific coverings. An  **$n$ -covering** is a covering that satisfies moreover:

- (2) For all  $k$ , if  $h|_k$  denotes the restriction of  $h$  to  $\tilde{H}_k$  (histories of length at most  $k$ ) and similarly for strategies, then  $\pi(h^*)|_k$  depends only on  $h|_k$  and  $\phi(\sigma^*)|_k$  depends only on  $\sigma^*|_k$ .
- (3)  $\tilde{H}_n = \tilde{H}_n^*$  and  $\phi(\sigma^*)|_n = \sigma^*|_n$ , and similarly for  $\tau^*$ .

We can now define a projective limit of  $n$ -coverings by the following

**LEMMA 2.5.** *Given  $H^0$ , assume that, for every integer  $k$ ,  $(H^{k+1}, \pi^{k+1}, \phi^{k+1})$  is an  $(n+k)$ -covering of  $H^k$ . Then, there exists  $H^*$  and, for each  $k$ , a  $(n+k)$ -covering  $(H^*, \pi^{*k}, \phi^{*k})$  of  $H^k$  such that, for all  $k$  and all  $h^*$  in  $[H^*]$ :*

$$\pi^{*k}(h^*) = \pi^{k+1} \circ \pi^{k+1}(h^*) .$$

**PROOF.** By hypothesis, one has, for  $\ell \geq k$ ,  $\tilde{H}_{n+k}^\ell = \tilde{H}_{n+k}^k$ . Define then  $H^* = \bigcup_k \tilde{H}_{n+k}^k$ , hence  $H^*$  is also a tree. For  $h^*$  in  $[H^*]$ , we define  $\pi^{*k}(h^*)$  as follows: by composing the mappings  $\pi$  and  $\phi$ , one has, for  $\ell \geq k$ , an  $n$ -covering  $(H^\ell, \pi^{k\ell}, \phi^{k\ell})$  of  $H^k$ . Let us now choose  $h^\ell$  in  $[H^\ell]$  such that  $h_{n+\ell}^* = h_{n+\ell}^\ell$  and put  $\pi^{*k}(h^*) = \lim_{\ell \rightarrow \infty} \pi^{k\ell}(h^\ell)$ . Similarly, for  $\sigma^*$  in  $\Sigma^*$ , choose, for each  $\ell$ ,  $\sigma^\ell$  in  $\Sigma^\ell$  that coincides with  $\sigma^*$  on  $\tilde{H}_{n+\ell}^\ell$  and let

$\phi^{\star k}(\sigma^\star) = \lim_{\ell \rightarrow \infty} \phi^{k\ell}(\sigma^\ell)$ . It is easy to see that  $(\pi^{\star k}, \phi^{\star k})$  are well-defined and  $(H^\star, \pi^{\star k}, \phi^{\star k})$  satisfies the requirements.  $\blacksquare$

Given a tree  $H, W \subseteq [H]$  is **standard** if, for every integer  $n$ , there exists a  $n$ -covering  $(H^\star, \pi, \phi)$  of  $H$  such that  $H^{-1}(W)$  is clopen (open and closed) in  $[H^\star]$ . Theorem 2.3 will then follow from prop. 2.1 and the following lemma.

LEMMA 2.6. *A closed set is standard.*

PROOF OF THE THEOREM. In fact, assuming this result, we first prove by induction that Borel sets are standard. Let  $\Sigma_1$  be the class of open sets and, for every countable ordinal  $\alpha$  let  $\Pi_\alpha = \{W \mid W^c \text{ is in } \Sigma_\alpha\}$ ,  $\Sigma_\alpha = \{W \mid W \text{ is a countable union of sets in } \bigcup_{\beta < \alpha} \Pi_\beta\}$ . Assume that for all trees, all  $W$  in  $\Pi_\beta$ ,  $\beta < \alpha$  are standard, and let  $W = \bigcup W_k$  with  $W_k$  in  $\Pi_{\beta_k}$  with  $\beta_k < \alpha$ . Given  $n$ , let  $(T_1, \pi_1, \phi_1)$  be an  $n$ -covering of  $H$  such that  $H_1^{-1}(W_1)$  is clopen. Let now  $(T_2, \pi_2, \phi_2)$  be an  $(n+1)$ -covering of  $T_1$  such that  $\pi_2^{-1} \circ \pi_1^{-1}(W_2)$  is clopen and so on inductively. Using then lemma 2.5, one gets  $(H^\star, \pi^\star, \phi^\star)$  being an  $n$ -covering of  $H$  for which every  $H^{\star-1}(W_k)$  is clopen and hence  $\pi^{\star-1}(W)$  is open in  $[H^\star]$ . Just note that the complement of a standard set is standard, hence Borel sets are standard. Lemma 2.4 proves then the theorem.  $\blacksquare$

PROOF OF THE LEMMA. Let  $W$  closed in  $[H]$  and  $k$  an integer. A position  $h^* \in H^\star$  corresponds to a position  $h$  in  $H$  except at stages  $2k+1$  and  $2k+2$  where:  $\omega_{2k+1}^* = (\omega_{2k+1}, H_I)$ ,  $H_I$  being a I-sub-tree starting from  $(\omega_1, \dots, \omega_{2k+1})$ .  $\omega_{2k+2}^*$  is either of the form:  $(\omega_{2k+2}, \langle 0, u \rangle)$ ,  $u$  being an history in  $H$  with even length extending  $(\omega_1, \dots, \omega_{2k+2})$  and belonging to  $H_I \setminus W^\prec$  (in this case, the moves after stage  $2k+2$  have to extend  $u$ ), or  $(\omega_{2k+2}, \langle 1, H_{II} \rangle)$  where  $H_{II}$  is a II-sub-tree of  $H_I$  contained in  $W^\prec$  (the moves from  $2k+2$  on have then to respect  $H_{II}$ ). Note that since  $W$  is closed this set of choices is non-empty (prop. 2.1).  $\pi$  is obviously the natural projection on  $H$  and, by definition,  $h^*$  belongs to  $\pi^{-1}(W)$  if  $(\omega_1, \dots, \omega_n) \in W^\prec$  for all  $n$ , hence if  $\omega_{2k+2}^*$  can be written as  $(\cdot, \langle 1, \cdot \rangle)$ , so that  $H^{-1}(W)$  is clopen. It is easily seen that  $H^\star$  is a tree; there thus remains to define  $\phi$ .

- (1) Consider first  $\Sigma^\star$ . On  $\tilde{H}_{2k}$ ,  $\sigma$  coincides with  $\sigma^\star$ . Assume now  $\sigma^\star(\omega_1, \dots, \omega_{2k}) = (\omega_{2k+1}, H_I)$  and consider the  $H_I$  game starting from  $(\omega_1, \dots, \omega_{2k+1})$  where  $W^c$  is winning. This game being open, is determined.
  - (a) If player I has there a winning strategy, say  $\sigma'$ , let him use it until the (finite) stage where an even history  $h$  in  $H_I \setminus W^\prec$  is reached. Consider now the corresponding position  $h^*$  in  $H^\star$  with  $\pi(h^*) = h$  and  $\omega_{2k+2}^* = (\omega_{2k+2}, \langle 0, h \rangle)$ . This history is compatible with  $\sigma^\star$  and I plays now in  $H$  (after  $h$ ) by following  $\sigma'$  in  $H^\star$  after  $h^*$ .
    - (b) Otherwise,  $(\omega_1, \dots, \omega_{2k+1}) \in L^{\text{II}}(W)$ . At every history  $h$  in  $H$ , either there exists  $h^*$  in  $H^\star$  with  $\omega_{2k+2}^* = (\omega_{2k+2}, \langle 1, L^{\text{II}}(W) \rangle)$  and  $\pi(h^*) = h$ , then  $\sigma$  follows  $\sigma^\star$  at  $h^*$ ; otherwise this means that, at some stage, the partial history  $h'$  is no longer in  $L^{\text{II}}(W)$ . This implies that, after  $h'$ , player I can force  $W^c$  and we use the construction above (1a).
- (2) For  $\tau^\star$  in  $\mathcal{T}^\star$ ,  $\phi(\tau^\star) = \tau$  is defined as  $\tau^\star$  on  $\tilde{H}_{2k}$ . Consider now, given  $(\omega_1, \dots, \omega_{2k+1})$ , the game starting from this history where II is winning iff he reaches an even history  $h$  for which there exists a I-sub-tree, say  $H_I(h)$ , satisfying  $\tau^\star(\omega_1, \dots, \omega_{2k}, (\omega_{2k+1}, H_I(h))) = (\omega_{2k+2}, \langle 0, h \rangle)$ . Note that in this game, the set  $V$  of winning plays of I is closed. Consider again two cases:

- (a) If player II has a winning strategy in this game, he uses it to reach  $h = (\omega_1, \dots, \omega_{2\ell})$ , then follows  $\tau^*$  from  $(\omega_1, \dots, (\omega_{2k+1}, H_I(h)), (\omega_{2k+2}, \langle 0, h \rangle))$  on.
- (b) Else  $(\omega_1, \dots, \omega_{2k+1}) \in L^I(V)$ . Obviously  $\tau^*(\omega_1, \dots, (\omega_{2k+1}, L^I(V)))$  is of the form  $(\omega_{2k+2}, \langle 1, H_{II} \rangle)$  by the definition of the previous game.  $\tau$  follows then  $\tau^*$  as long as  $h$  is the image of some possible  $h^*$ . If not, this means that  $h$  is no longer in  $L^I(V)$  so that, after  $h$ , II has a winning strategy and we are back to (2a).

It is now easy to check that  $\phi$  is well-defined and  $(H^*, \pi, \phi)$  is a  $k$ -covering of  $H$ . ■

**REMARK 2.2.** The theorem states that, if  $W$  is a “Borel property”, the negation of the infinite sentence “ $\exists \omega_1 \forall \omega_2 \exists \omega_3 \forall \omega_4, \dots$  such that  $W$ ” is the sentence “ $\forall \omega_1 \exists \omega_2 \forall \omega_3 \exists \omega_4, \dots$  such that [not  $W$ ]”. It yields thus an extension of the usual rule of negation to infinite sentences. Hence the interest of logicians in this question (cf. e.g. Moschovakis, 1980).

**2.b. Comments: Infinite Games without Perfect Information.** Without the previous perfect information assumption the analysis is much harder and only very partial results are available.

$\Gamma$  is the game with perfect recall described as follows:  $S$  and  $T$  are finite sets and both players choose at stage  $n$  independently moves  $s_n$  and  $t_n$  in  $S$  or  $T$ .  $\omega_n = (s_n, t_n)$  in  $\Omega = S \times T$  is then announced to both and they proceed to the next stage.  $g$  is a real bounded measurable function on  $(H_\infty, \mathcal{H}_\infty)$  (with the product  $\sigma$ -algebra induced by the discrete topology on each factor). Strategies are defined as mappings from  $H$  to  $\Delta(S)$  or  $\Delta(T)$  and the pay-off if given by  $\gamma(\sigma, \tau) = E_{\sigma, \tau}(g)$ . By the minmax theorem (e.g. prop. 1.17 p. 7) and remark 1.5 p. 56 we have immediately:

**LEMMA 2.7.** *If  $g$  is l.s.c., the game has a value, and player II has an optimal strategy.*

**REMARK 2.3.** This implies in particular that, if  $g = \mathbb{1}_W$  with  $W$  open or closed, the game has a value (note nevertheless that one player may not have a winning strategy).

One can go one step further to obtain (Blackwell, 1969) (cf. also Orkin (1972b) for an extension to the Boolean algebra generated by the  $G_\delta$ 's):

**PROPOSITION 2.8.** *If  $W$  is a  $G_\delta$ , the game has a value.*

**PROOF.** Note first that  $W$  is a  $G_\delta$  in  $H_\infty$  iff there exists a subset  $Y$  of  $H$  such that  $h_\infty$  belongs to  $W$  iff  $h_n$  belongs to  $Y$  for infinitely many  $n$  in  $\mathbb{N}$  (i.e.  $h_\infty$  hits  $Y$  infinitely often). For any position  $x$  in  $H$ , denote by  $\bar{v}(x)$  the minmax of the subgame  $\Gamma(x)$  starting from  $x$  (and  $\bar{v} = \bar{v}(\emptyset)$ ).

We now introduce  $\Gamma'(x)$  as the subgame starting from  $x$  and with pay-off  $\bar{v}(y)$ , where  $y$  is the entrance position in  $Y$  after  $x$ . Formally, let  $\theta_x(h_\infty) = \min(\{n \mid 1 \leq n < \infty, (x, h_n) \in Y\} \cup \{\infty\})$ . Then the pay-off of  $\Gamma'(x)$  is  $f$  with  $f(h_\infty) = \mathbb{1}_{n < \infty} \bar{v}(x, h_n)$  where  $n = \theta_x(h_\infty)$ . It is clear that  $f$  is l.s.c. hence the previous lemma implies that  $\Gamma'(x)$  has a value  $v'(x)$  and player II has an optimal strategy  $\tau'(x)$ . Let us first remark that:

$$(1) \quad v'(x) \geq \bar{v}(x) \quad \forall x \in H$$

(consider the following strategy for player II in  $\Gamma(x)$ : play according to  $\tau'(x)$  until  $Y$  is reached; if  $Y$  is reached at position  $y$ , play then an  $\varepsilon$ -optimal strategy in  $\Gamma(y)$ ).

Let us prove now that player I can guarantee  $\bar{v}$ , so that  $\Gamma$  has a value, namely  $\bar{v}$ . Let  $\varepsilon_m = \varepsilon/2^{m+1}$ , and define inductively  $\sigma$  as follows: Play  $\varepsilon_0$ -optimally in  $\Gamma'(\phi)$  until  $Y$  is reached. If this happens at position  $y_1$ , play then  $\varepsilon_1$ -optimally in  $\Gamma'(y_1)$  until  $Y$  is reached

again. If this occurs at  $y_2$ , play  $\varepsilon_2$ -optimally in  $\Gamma'(y_2)$  and so on. Given  $\sigma$  and any  $\tau$  in  $\mathcal{T}$ , define a sequence of pay-offs  $\{\rho_m\}$  by:  $\rho_0 = \bar{v}$  and  $\rho_m = \bar{v}(y_m)$  if  $y_m$  is defined (i.e. if  $h_\infty$  hits  $Y$  at least  $m$  times) and 0 otherwise. Then we have:

$$(2) \quad E_{\sigma,\tau}(\rho_{m+1} \mid \mathcal{H}_m) \geq \rho_m - \varepsilon_m$$

where  $\mathcal{H}_m$  is the  $\sigma$ -algebra generated by  $(y_1, \dots, y_m)$  on  $H_\infty$ . This is clear if  $\rho_m = 0$ . Otherwise, after the position  $y_m$ , player I was playing  $\varepsilon_m$  optimally in  $\Gamma'(y_m)$ , meaning that his pay-off, which is precisely  $\rho_{m+1}$ , has a conditional expectation greater than  $v'(y_m) - \varepsilon_m$ , hence the inequality (2), using (1). Taking expectation and summing in (2), we obtain:

$$E_{\sigma,\tau}(\rho_m) \geq \rho_0 - (\varepsilon_1 + \dots + \varepsilon_m) \geq \bar{v} - \varepsilon.$$

Since  $0 \leq \rho_m \leq 1$  and  $\rho_m$  is 0 unless  $Y$  is hit  $m$  times, this implies that:

$$P_{\sigma,\tau}(Y \text{ is hit } m \text{ times}) \geq \bar{v} - \varepsilon$$

and letting  $m \rightarrow \infty$  yields

$$(3) \quad P_{\sigma,\tau}(Y \text{ is hit infinitely often}) = P_{\sigma,\tau}(h_\infty \in W) \geq \bar{v} - \varepsilon. \quad \blacksquare$$

### Exercises.

The games in ex. II.2Ex.1–II.2Ex.3 are as in sect. 2.b.

**1. Counterexample in approximation.** (Orkin, 1972a)(cf. also ex. I.1Ex.4 p. 10 and I.1Ex.5 p. 10). Let  $S, T = \{0, 1\}$ ,  $\Omega = S \times T$ . Define  $Y'_n = \{\omega \in \Omega \mid \exists i \leq n \text{ with } \omega_i = (1, 1)\}$  and  $Z = \{\omega \in \Omega \mid t_i = 0 \forall i\}$ . Show that the game with (closed) winning set  $Y_n = Y'_n \cup Z$  is determined with value 0 and player II has a winning strategy. Prove that the game with winning set  $Y = \bigcup_{n \geq 1} Y_n$  has value 1, and player I has a winning strategy.  $Y$  is the union of an open and a closed set, hence both a  $G_\delta$  and an  $F_\sigma$ . (This shows that one can not use an approximation argument like in ex. I.1Ex.2 p. 9 to prove the previous result by using open sets — or dually.)

**2.** (Orkin, 1972b) Take  $S, T = \{0, 1\}$ ,  $\Omega = S \times T$ . The winning set is  $Z = X \cup Y$  with  $X = \{\omega \mid \omega_n = (0, 0) \text{ for infinitely many } n \text{ and } \omega_n = (1, 1) \text{ for infinitely many } n\}$  and  $Y = \{\omega \mid \omega_n = (0, 0) \text{ for at most finitely many } n \text{ and } \omega_n = (1, 1) \text{ for at most finitely many } n\}$ .  $X$  is a  $G_\delta$  and  $Y$  an  $F_\sigma$ . Show that the following is a winning strategy for I: he chooses 1 as long as II chooses 0, and reverses his behaviour just after II is doing so.

**3. A  $G_\delta$  game.** (Orkin, 1972b)  $\Omega = S \times T = \{0, 1\}^2$ . As soon as I chooses 1, the game ends and player I wins if player II chooses 1 at that move and loses otherwise. If player I always plays 0, he wins iff player II chooses infinitely often 0 (cf. the related “Big Match”, ch. VII, ex. VIIEx.4 p. 346). Show that the value of the game is 1 but player I has no optimal strategy.

HINT. An  $\varepsilon$ -optimal strategy (due to Blackwell) of player I is as follows: take  $N > 1/\varepsilon$ ; let  $N_j = 2^j N$ ,  $j = 1, \dots$  and divide the play into successive blocs of length  $N_j$ . On each bloc  $j$ , player I selects a stage  $n_j$  uniformly distributed on this bloc and plays 1 at this stage iff the previous  $n_j - 1$  moves of player II in this bloc were 1. Otherwise, he plays always 0 on this bloc.

**4. Non-zero-sum Borel Games (Mertens & Neyman).** (cf. Mertens, 1986b) We consider games with perfect information like in sect. 2.a p. 79.

a. Prove that, if the pay-off function  $g$  is Borel, the game has a value (consider the level sets).

b. Note that, if  $g$  takes finitely many values, the players have optimal strategies and they can even be improved such as to remain still optimal at any position.

c. Assume now a set  $\mathbf{I}$  of players (of whatever cardinality) and that each pay-off function  $g^i$  satisfies II.2Ex.4a and II.2Ex.4b. Define, for each  $i$ ,  $\{\sigma^j(i), j \in \mathbf{I}\}$  to be a strategy vector as improved sub II.2Ex.4b in the two-person zero-sum game  $(i, \mathbb{C}\{i\})$  with pay-off  $g^i$ . Let finally  $\tau^i$  be defined as: play  $\sigma^i(i)$  as long as the other players do so and switch to  $\sigma^i(j)$  if player  $j$  deviates. This induces a pure equilibrium.

d. Deduce the existence of pure  $\varepsilon$ -equilibria for bounded Borel pay-off functions.

**5. The second separation theorem.** (after Blackwell, 1967a) [For notations and definitions below, cf. App.1 and 2.a.]

Let  $\mathcal{P}$  be a paving on a set  $X$ , with  $\mathcal{P}_c \subseteq \mathcal{P}_s$ . Given  $C_n \in \mathcal{P}_{sc}$ , there exists a sequence  $B_n \in \mathcal{P}_{sc}$ ,  $B_n \subseteq C_n$ , which forms a partition of  $\bigcup_n C_n$ .

a. There is a loss in assuming  $\mathcal{P}$  to be a ( $\sigma$ )-field.

HINT. Use the stability of the Souslin operation ( $\mathcal{P}_s = \mathcal{P}_{ss}$ ) — which implies immediately  $\mathcal{P}_s = \mathcal{P}_{s\sigma} = \mathcal{P}_{s\delta}$  since clearly  $\mathcal{P}_\sigma \subseteq \mathcal{P}_s$ ,  $\mathcal{P}_\delta \subseteq \mathcal{P}_s$ :  $\mathcal{B} = \mathcal{P} \cap \mathcal{P}_{sc}$  is thus a  $\sigma$ -field, with  $\mathcal{P} \subseteq \mathcal{B}$  (assumption  $\mathcal{P}_c \subseteq \mathcal{P}_s$ ), hence  $\mathcal{P}_s \subseteq \mathcal{B}_s \subseteq \mathcal{P}_{ss} = \mathcal{P}_s$ .

b. Let  $A_n = C_n^c = \complement C_n$  have the Souslin scheme  $A_n = \bigcup_{\sigma \in \mathbb{N}^\mathbb{N}} \bigcap_{k \in \mathbb{N}} P_{\sigma_k}^n$  with  $P_{\sigma_k}^n$  in the boolean algebra  $\mathcal{P}$ , and where  $\sigma_k$  denotes the initial segment of length  $k$  of  $\sigma \in \mathbb{N}^\mathbb{N}$ . Set  $P_{\sigma_0}^n = X$  ( $\sigma_0 = \emptyset$ ). For  $x \in X$ , define the game  $\Gamma_x$  with players  $n \in \mathbb{N}$  each picking an integer in the following order: 1,1,2,1,2,3,1,2,3,4,1,2,.... At any stage  $t$ , denote by  $h_t(n)$  the sequence of past choices of player  $n$ . The game continues until at some stage  $t$ ,  $x \notin P_{h_t(n)}^n$  for some player  $n$ , who is then declared the loser. The sets  $L_n$  of plays where player  $n$  loses are disjoint open sets in  $\mathbb{N}^\mathbb{N}$ . Let  $L_0 = (\bigcup_{n \geq 1} L_n)^c$ , and observe  $L_0 = \emptyset$  for  $x \notin A = \bigcap_n A_n$ .

Let  $A_n' = \{x \mid \text{player } n \text{ can avoid losing in } \Gamma_x\}$

$A_n'' = \{x \mid \text{player } n \text{ cannot force a loss upon his opponents in } \Gamma_x\}$ .

Show that:

- (1)  $A_n \subseteq A_n'$
- (2)  $A_n' \cap A_n'' = A$
- (3)  $A_n' \cup A_n'' = X$
- (4)  $\bigcup_n A_n'' = X$  (if everyone forces a loss upon his opponents, one will lose ...)  
(by 2 and 3, this is equivalent to  $\bigcap_n A_n' = A$ ).
- (5)  $A_n'$  and  $A_n''$  belong to  $\mathcal{P}_s$ .

HINT. E.g., for  $A_n''$ , the set of  $x$  such that the opponents have a joint strategy  $\tau$  that forces the closed set  $L_0 \cup L_n$ , i.e. such that, for any finite history  $h$ , either  $h$  is not compatible with  $\tau$ , or no opponent has lost along  $h$ . Observe the set of histories is countable, hence the set of strategies can be viewed as  $\mathbb{N}^\mathbb{N}$ : our condition on the pair  $(\tau, x)$  is thus the intersection over countably many histories  $h$  of conditions “either  $\tau$  belongs to some clopen set or  $x$  belongs to some set in the field  $\mathcal{B}$ ”:  $\{(\tau, x)\} = \bigcap [(K_i \times X) \cup (\mathbb{N}^\mathbb{N} \times B_i)]$  with  $K_i$  clopen,  $B_i \in \mathcal{B}$ . Use now the projection property — or, to keep the proof fully self contained, write the corresponding Souslin scheme:  $A_n'' = \bigcup_{\tau \in \mathbb{N}^\mathbb{N}} \bigcap_{k \in \mathbb{N}} Q_{\tau_k}^n$ , with  $Q_{\tau_k}^n = \bigcap \{B_i \mid i \leq k, \sigma \in K_i \Rightarrow \sigma_k \neq \tau_k\} \in \mathcal{B}$ .

c. Set now  $B_n = (A_n')^c \cap \bigcap_{i < n} (A_i'')^c$  to finish the proof.

**6. Borel sets via games.** (Blackwell, 1981) Let  $H = \bigcup_{n \geq 0} \mathbb{N}^n$ .  $X \subseteq H$  is a stop rule if for all infinite sequences  $h_\infty$ , there exists a unique  $x$  in  $X$  with  $x \prec h_\infty$ . Given a stop rule  $X$ , a function on  $X$  whose values are intervals of  $\mathbb{R}$  and a real number  $u$ , consider the game  $G(X, f, u)$  with perfect information on  $\Omega = \mathbb{N}$  where I plays first and wins if  $u \in f(\omega_1, \dots, \omega_k)$  where  $(\omega_1, \dots, \omega_k)$  is in  $X$ . Let  $B(X, f)$  the set of  $u$ 's for which I wins. Prove that the family  $\mathcal{B}$  of all  $B(X, f)$  is the  $\sigma$ -algebra of Borel sets on  $\mathbb{R}$ .

HINT. Note that  $\mathcal{B}$  includes intervals. Prove that  $\mathcal{B}$  is stable by countable unions and intersections. (Consider the extended game where I or II chooses first which game  $G(X_n, f_n, u)$  to play.) For the converse note that  $B(X, f)$  is analytic (as in ex. II.2Ex.5) but  $G(X, f, u)$  is clopen hence by prop. 2.1 p. 80,  $B(X, f)$  is coanalytic and use 3.h.

**7. Analytic sets via games.** (Dellacherie, 1980) Let  $X$  be a compact metric space.  $\mathcal{C} \subseteq \mathcal{P}(X)$  is a capacitance if

- (1)  $A \in \mathcal{C}, A \subseteq B \Rightarrow B \in \mathcal{C}$
- (2)  $A_n \nearrow A$  and  $A \in \mathcal{C} \Rightarrow \exists n, A_n \in \mathcal{C}$ .

Define a game  $G_A$  with perfect information as follows: I chooses  $\mathcal{C}_1$  with  $A \in \mathcal{C}_1$ , II chooses  $A_1 \subseteq A$ ,  $A_1 \in \mathcal{C}_1$ , I chooses  $\mathcal{C}_2$  with  $A_1 \in \mathcal{C}_2, \dots$  and so on. II wins if  $\bigcap_n \bar{A}_n \subseteq A$ , where  $\bar{A}_n$  is the closure of  $A_n$ . Let  $\mathcal{D}$  be the family of sets  $A$  where II wins in  $G_A$ . Prove that  $\mathcal{D}$  is the family  $\mathcal{A}$  of analytic subsets of  $X$ .

HINT. (1)  $\mathcal{A} \subseteq \mathcal{D}$ .

- (a) Let first  $A = \bigcap_n \bigcup_m K_m^n$ ,  $K_m^n$  compact increasing in  $m$  for fixed  $n$ . Given  $\mathcal{C}_1$  with  $A \in \mathcal{C}_1$  there exists  $m_1$  with  $A \cap K_{m_1}^1 = \tau(\mathcal{C}_1)$  in  $\mathcal{C}_1$ . Given  $\mathcal{C}_2$  with  $\tau(\mathcal{C}_1) \in \mathcal{C}_2$ , there exists  $m_2$  with  $\tau(\mathcal{C}_1) \cap K_{m_2}^2 = \tau(\mathcal{C}_1, \mathcal{C}_2) \in \mathcal{C}_2$ . Since  $\cap K_{m_i}^i \subseteq A$ ,  $\tau$  is winning.
- (b)  $\mathcal{D}$  is stable by continuous functions: let  $\mathcal{C}' = \{B \mid f(B) \in \mathcal{C}\}$  and define, if  $\tau'$  is winning for  $A'$ ,  $f(A') = A$ ,  $\tau(\mathcal{C}_1) = f(\tau'(\mathcal{C}_1'))$ ,  $\tau(\mathcal{C}_1, \mathcal{C}_2) = f(\tau'(\mathcal{C}_1', \mathcal{C}_2'))$  and so on. Since  $f(\cap K_n) = \cap f(K_n)$ ,  $K_n$  compact,  $\tau$  is winning for  $A$ .
- (c) Any analytic set can be obtained as  $f(A)$ ,  $f$  continuous,  $A$  like sub 1a — 3.e.

- (2)  $\mathcal{D} \subseteq A$ .

Let  $A$  be in  $\mathcal{D}$  and  $\tau$  a winning strategy for II. Note that the class of sets that cannot be written as  $\bigcup_n B_n, B_n \subseteq \tau(\mathcal{C}_n)$ ,  $\mathcal{C}_n$  capacitance containing  $A$ , is a capacitance, hence does not contain  $A$ . Thus  $A = \bigcup_n \tau(\mathcal{C}_n)$  for some sequence  $\mathcal{C}_n$ . Similarly, for each  $\omega_1$  in  $\mathbb{N}$ , the family of sets, that cannot be written as  $\bigcup_n B_n$  with  $B_n \subseteq \tau(\mathcal{C}_{\omega_1}, \mathcal{C}_n)$  for capacitances  $\mathcal{C}_n$  containing  $A_{\omega_1} = \tau(\mathcal{C}_{\omega_1})$ , does not contain  $A_{\omega_1}$ . Define then two mappings from finite sequences in  $\mathbb{N}$  to capacitances and subsets of  $A$  such that, writing  $\omega_k$  for the first  $k$  terms of some  $\omega$  in  $\mathbb{N}^\mathbb{N}$ :

- (a)  $A_{\omega_k} = \tau(\mathcal{C}_{\omega_1}, \dots, \mathcal{C}_{\omega_k})$
- (b)  $A_{\omega_k} = \bigcup_n A_{\omega_k, n}$  with  $A_\emptyset = A$ .

Deduce that  $\bigcap_k \bar{A}_{\omega_k} \subseteq A$  ( $\tau$  is winning); hence, by 2b,  $A = \bigcup_\omega \bigcap_k A_{\omega_k} = \bigcup_\omega \bigcap_k \bar{A}_{\omega_k}$ . Hence  $A$  is analytic (Souslin scheme) (3.b).

**8. Topological games and the Baire property (Choquet, Christensen, Saint Raymond).** (Choquet, 1969; Saint-Raymond, 1983) Given a topological space  $X$  consider the game with perfect information  $G$  where I chooses a non-empty open set  $U_1$ , then II chooses a non-empty open set  $V_1 \subseteq U_1$ , then I picks  $U_2 \subseteq V_1$  and so on. II wins in  $G$  if  $\bigcap_n V_n \neq \emptyset$ . Say that  $X$  is of type  $B$  (resp.  $A$ ) if I cannot win in  $G$  (resp. II wins).

Define similarly  $G'$  with moves  $U_n$  and  $(V_n, x_n)$ ,  $U_n, V_n$  open,  $x_n$  in  $X$ ,  $U_n \subseteq V_n \subseteq U_{n-1}$ , where II wins if the sequence  $x_n$  has an accumulation point in  $\bigcap_n V_n$ , and  $G''$  with moves  $(U_n, x_n)$  and  $V_n, U_n, V_n$  open,  $x_n$  in  $U_n$  and  $V_{n-1} \supseteq U_n \supseteq V_n \ni x_n$ , where II wins if  $\bigcap_n V_n \neq \emptyset$ . We introduce similarly types  $B', \dots, A''$ .

Observe that property  $A$  (resp.  $A', A''$ ) always implies property  $B$  (resp.  $B', B''$ ), and that  $A'' \Rightarrow A' \Rightarrow A, B'' \Rightarrow B' \Rightarrow B$ .

a. The following spaces are of type  $A''$ : complete metric spaces, locally compact Hausdorff spaces, products of spaces of type  $A''$ .

b. Prove that  $X$  is of type  $B$  if and only if  $X$  is a Baire space.

HINT. If  $X$  is not Baire, let  $U$  be a non-empty open set  $U \subseteq \bigcup_n F_n$ ,  $F_n$  closed sets with empty interior. Let  $\sigma(\phi) = U$ ,  $\sigma(V_1, \dots, V_n) = V_n \setminus F_n$ , then  $\sigma$  is winning. If  $\sigma$  is winning,  $\sigma(\phi) = U_1$  is of first category. In fact define  $I_1 = \{1\}$ ,  $U_1^1 = U_1$  and recursively a maximal family of non-empty open sets  $(U_i^n, V_i^{n-1})_{i \in I_n}$  with:

- (1)  $(U_i^n)_{i \in I_n}$  are pair-wise disjoint.
- (2)  $\forall i \in I_{n+1} \exists j \in I_n, V_i^{n+1} \subseteq U_j^n$ .

- (3)  $U_{i_1}^1 \supseteq U_{i_2}^2 \supseteq \dots \supseteq U_{i_n}^n$  implies  $U_{i_n}^n = \sigma(V_{i_2}^1, \dots, V_{i_n}^{n-1})$ .  
(4)  $U_i^n \subseteq V_i^{n-1}$

Let then  $W_n = \bigcup_{i \in I_n} U_i^n$  and show inductively that  $W_n$  is dense in  $U_1$ : Assume it is true for  $n$  and let  $W$  be open and disjoint from  $W_{n+1}$ ; let  $i_n \in I_n$  with  $W \cap U_{i_n}^n \neq \emptyset$ . This determines a unique  $(i_1, \dots, i_n)$  by 3 and  $V^n = W \cap U_{i_n}^n$ ,  $U^{n+1} = \sigma(V_{i_2}^1, \dots, V_{i_n}^{n-1})$  contradicts the maximality of the family. Finally  $\bigcap W_n = \emptyset$  since if it contains  $a$ ,  $a$  is compatible with a play of  $\sigma$  and  $\sigma$  is winning.

c. A metrisable space is of type  $A''$  iff it is topologically complete.

HINT. It suffices to show that if  $X$  is of type  $A''$ ,  $X$  is a  $G_\delta$  in its completion  $(E, d)$ .

Show there exists a family of indices  $I_{\dots}$  and open sets  $O_{\dots}$  in  $E$  constructed inductively by:

- (1)  $\omega_1 \in I, \dots, \omega_{n+1} \in I_{\omega_1, \dots, \omega_n}$ .  
(We shall write  $\omega$  for such a compatible sequence and  $\omega(n)$  for the first  $n$  terms).  
(2)  $O_{\omega(n), \omega_{n+1}} \subseteq O_{\omega(n)}$ ,  $X \cap (\bigcup_{\omega_{n+1}} O_{\omega(n), \omega_{n+1}}) = X \cap O_{\omega(n)}$ ,  $d(O_{\omega(n)}) < 2^{-n}$ .  
(3)  $\bigcap_n O_{\omega(n)}$  is a point of  $X$ .

Fix a well ordering on each  $I$  and define inductively  $A_{\omega_1} = \bigcup_{\theta_1 < \omega_1} O_{\theta_1}$ ,  $A_{\omega(n), \omega_{n+1}} = A_{\omega(n)} \cup (\bigcup_{\theta_{n+1} < \omega_{n+1}} O_{\omega(n), \theta_{n+1}})$ , then  $B_{\omega(n)} = O_{\omega(n)} \setminus A_{\omega(n)}$  and  $X_n = \bigcup_{\omega} B_{\omega(n)}$ ,  $Y = \bigcap_n X_n$ . Note that  $X \subseteq X_n$  hence  $X \subseteq Y$ . On the other hand for every  $\omega(n)$ , the  $B_{\omega(n), \omega_{n+1}}$  are included in  $B_{\omega(n)}$  and disjoint, so that  $Y = \bigcup_{\omega} \bigcap_n B_{\omega(n)} \subseteq \bigcup_{\omega} \bigcap_n O_{\omega(n)} \subseteq X$ . It suffices to show that  $X_n$  is a  $G_\delta$  in  $E$ . Since  $E$  is metric proceed by localisation:  $C \subseteq E$  is a  $G_\delta$  iff for every open covering  $(F_i)$  of  $C$ ,  $C \cap F_i$  is a  $G_\delta$ . (Use an open covering  $G_j$  of  $\bigcup F_i$ , finer than  $F_i$  and locally finite. Each  $C \cap G_j$  is the intersection of a decreasing sequence  $G_{j,n}$  of open subsets of  $G_j$ . Let  $\Omega_n = \bigcup_j G_{j,n}$  since  $G_j$  is locally finite one has  $\bigcap_n \Omega_n = \bigcup_j \bigcap_n G_{j,n} = \bigcup_j C \cap G_j = C$ ).

The proof then follows by induction.

d. (Hausdorff) Deduce that: if  $X$  is topologically complete and  $f$  an open continuous mapping in a metrisable space,  $f(X)$  is topologically complete.

e. If  $X$  separable is of type  $B$  it is of type  $B'$ .

HINT. Let  $(a_n)$  be dense in  $X$ , and given  $\sigma'$  in  $G'$ , let  $\sigma'(V_1, \dots, V_n) = \sigma'((V_1, a_1), \dots, (V_n, a_n))$ . If  $\tau$  wins against  $\sigma$ ,  $\cap V_n$  contains an accumulation point of  $(a_n)$ , hence  $\sigma'$  is losing against  $(V_n, a_n)$ .

f. A Hausdorff space  $X$  is a Namioka space if for all compact  $Y$ , all metrisable  $Z$  and all  $f$  from  $X \times Y$  to  $Z$ , separately continuous on  $X$  and  $Y$ , there exists a dense  $G_\delta$  of  $X$ ,  $A$ , such that  $f$  is continuous at each point of  $A \times Y$ .

Prove that if  $X$  is of type  $B'$ ,  $X$  is a Namioka space.

HINT. Else one can assume  $Z = [-1, 1]$  (consider the function  $d(f(x, y_1), f(x, y_2))$  on  $X \times (Y \times Y)$ ), and, using II.2Ex.8b, that on an open set  $W$  the oscillation of  $F: x \mapsto f(x, \cdot)$  is  $> \delta$ . Let for each  $k \geq 1$ ,  $(P_j^k)_{j \geq 1}$  be a dense sequence in  $\mathcal{P}^k$ , the set of real continuous functions on  $[-1, 1]^k$ . Define  $\sigma$  by  $\sigma(\phi) = W$  and  $\sigma((V_1, x_1), \dots, (V_n, x_n)) = U_{n+1} = V_n \setminus \bigcup_{j+k \leq n} C_{k,j}$  with  $C_{k,j} = \{x \mid \|F(x) - \varphi_{k,j}\| \leq \delta/3\}$  and  $\varphi_{k,j}(y) = P_j^k(f(x_1, y), \dots, f(x_k, y))$ . Since the diameter of  $F(C_{k,j})$  is less than  $2\delta/3$ ,  $C_{k,j}$  has an empty interior in  $W$ , hence  $U_{n+1} \neq \emptyset$ . Let  $(V_n, x_n)_{n \geq 1}$  be winning against  $\sigma$ . Write  $x_\infty$  for an accumulation point of  $(x_n)$  in  $\bigcap_n V_n$ . Let  $\phi: Y \rightarrow [-1, 1]^\mathbb{N}$ ,  $\phi(y) = (f(x_n, y))_n$ . Since  $\phi(y) = \phi(y')$  implies  $f(x_\infty, y) = f(x_\infty, y')$  there exists  $\varphi$  continuous on the compact set  $\phi(Y)$  with  $F(x_\infty) = \varphi \circ \phi$  hence by Urysohn's theorem a continuous function  $\psi$  on  $[-1, 1]^\mathbb{N}$  that coincides with  $\varphi$  on  $\phi(Y)$ . Let  $\psi_k$  on  $[-1, 1]^k$  defined by  $\psi_k(u_1, \dots, u_k) = \psi(u_1, \dots, u_k, 0, 0, \dots)$  and  $\pi_k$  the projection of  $[-1, 1]^\mathbb{N}$  on  $[-1, 1]^k$ . By uniform continuity of  $\psi$  there exists  $k$  with  $\|\psi - \psi_k \circ \pi_k\| \leq \delta/12$ . Choose  $j$  with  $\|\psi_k - P_j^k\| \leq \delta/4$ , then  $\|F(x_\infty) - \varphi_{k,j}\| \leq \delta/3$  hence  $x_\infty \in C_{k,j}$ , hence  $x_\infty \notin \bigcap_n V_n$ .

**9. Games without value.** (Davis, 1964) Consider a game with perfect information where player I chooses a finite sequence in  $\Omega = \{0, 1\}$ , then II chooses a point in  $\Omega$ , then I a finite sequence and so on. Given  $W$  in  $\Omega^\infty$ , player I wins if the play belongs to  $W$ .

a. Prove that I wins iff  $W$  contains a perfect set.

HINT. If  $\sigma$  is winning, the set of plays consistent with  $\sigma$  is a perfect set in  $W$ .

Conversely, let  $P$  be perfect in  $W$ . Given  $h$  in  $H = \bigcup_n \Omega^n$ ,  $f(h) = \{x \in \Omega^\infty \mid h \prec x\}$  has an empty or perfect intersection with  $P$ . Deduce that  $Q = \{h \in H \mid f(h, 0) \cap P \text{ and } f(h, 1) \cap P \text{ are perfect}\}$  is non-empty and induces a winning I-sub-tree.

b. Show that player II wins iff  $W$  is countable.

HINT. If  $W$  is countable  $= \{w^n\}_{n \in \mathbb{N}}$ , player II can force a play with  $w^n \notin f(h_{2n})$ ,  $\forall n$ .

Show that if II has a winning strategy he has also a winning strategy that depends only on the position at stage  $n$  (and not on the history: sequence of previous positions).

Thus, if  $\tau$  is winning, for each  $w$  in  $W$ , there exists  $N$  such that  $n \geq N$  implies  $w \notin f(w_n, \tau(w_n))$ . Denote by  $N(w)$  the smallest such  $N$ , and let  $W_k = \{w \in W \mid N(w) \leq k\}$ . Note that for any  $h$  in  $\Omega^k$ ,  $f(h) \cap W_k$  contains at most one point hence  $W_k$  is finite.

**10. An operator solution proof of prop. 2.8.** (Blackwell, 1989) For each function  $u$  in  $\mathcal{U}$ , i.e. defined on  $H$  with  $0 \leq u \leq 1$  and each position  $x$  let  $\Gamma(u; x)$  be the game starting from  $x$  and with pay-off  $u(y)$  where  $y$  is the entrance position in  $Y$  after  $x$ . Note that  $\Gamma(u; x)$  has a value, say  $Tu(x)$ ; this defines an operator  $T$  from  $\mathcal{U}$  to itself.

Show  $T$  has a fixed point  $u^*$  and  $u^*(\emptyset)$  is the value of the original  $G_\delta$  game, say  $\Gamma(\emptyset)$ .

HINT. For each countable ordinal  $\alpha$  define  $u^\alpha$  in  $\mathcal{U}$  by:

$$u^0 \equiv 1, u^{\alpha+1} = Tu^\alpha, \text{ and if } \alpha \text{ is a limit ordinal } u^\alpha = \inf\{u^\beta \mid \beta < \alpha\}.$$

To prove that I can guarantee  $u^*(\emptyset)$  let him play optimally in a sequence of games  $\Gamma(u^*; y_m)$  (cf. the proof of prop. 2.8 p. 83). For II, prove inductively that for any countable ordinal  $\alpha$  any position  $x$  and any  $\varepsilon > 0$ , he can obtain  $u^\alpha(x) + \varepsilon$  in  $\Gamma(x)$  (play first optimally in some  $\Gamma(u^\beta; x)$  then obtain  $u^\beta(y)$  in  $\Gamma(y)$ ).

### 3. Correlated equilibria and extensions

We follow here roughly the approach of Forges (1986a).

**3.a. Correlated equilibria.** A correlation device  $c$  (for the player set  $\mathbf{I}$ ) is a probability space  $(E, \mathcal{E}, P)$  together with sub  $\sigma$ -fields  $(\mathcal{E}^i)_{i \in \mathbf{I}}$  of  $\mathcal{E}$ . The extension  $\Gamma_c$  of a game  $\Gamma$  by  $c$  is the game where first nature selects  $e \in E$  according to  $P$ , next each player  $i \in \mathbf{I}$  is informed of the events in  $\mathcal{E}^i$  which contain  $e$ , then  $\Gamma$  is played (every player remembering all along his information about  $e$ ).

A correlated equilibrium of  $\Gamma$  (Aumann, 1974) is a pair  $(c, \text{equilibrium of } \Gamma_c)$ .

In such a correlated equilibrium, each player  $i \in \mathbf{I}$  has a private probability space  $(X^i, \mathcal{X}^i, Q^i)$  to do his own randomisation (cf. sect. 1); one can replace  $E$  by its product with all  $X^i$ 's, such as to reduce mixed (resp. general) strategy correlated equilibria to pure (resp. behavioural) strategy correlated equilibria. Further, if the action sets are standard Borel and  $\Gamma$  is countably linear, ex. II.1Ex.10d p. 63 allows to reduce correlated equilibria in general strategies to the pure strategy case. Similarly, if the game is linear and  $A$  is countable, the method of ex. II.1Ex.10e p. 63 [with  $(E, \mathcal{E}^i)$  instead of  $(X, \mathcal{X})$ ], reduces behavioural strategy correlated equilibria (hence, by the above, also general strategy correlated equilibria) to the pure strategy case. (One cannot use directly the statement of ex. II.1Ex.10e, since the set of signals of  $\Gamma_c$  is not countable.)

Henceforth we assume  $\Gamma$  is a linear game with countable signal space (recall ex. II.1Ex.14, comment 1.28 p. 73).

**THEOREM 3.1.** The correlated equilibria of  $\Gamma$  have distributions on the space  $(\Sigma, \mathcal{S}) = \prod_i (\Sigma^i, \mathcal{S}^i)$  of pure strategy  $\mathbf{I}$ -tuples of  $\Gamma$ . This set  $C$  of distributions is convex.

**PROOF.** According to remark 1.4 p. 55, pure strategy correlated equilibria can now also be seen as an  $\mathbf{I}^{\text{tuple}}$  of measurable maps  $\hat{\sigma}^i$  from  $(E, \mathcal{E}^i, P)$  to  $i$ 's pure strategy space  $(\Sigma^i, \mathcal{S}^i)$  in  $\Gamma$ . Hence the first sentence.

Convexity follows from the fact that a lottery between two correlated equilibria yields again a correlated equilibrium, with device: ■

**DEFINITION 3.2.**  $C$  is called the set of **correlated equilibrium distributions**.  $C_0$  is the set of corresponding pay-offs. A **canonical correlated equilibrium** is one where  $c = ((\Sigma, \mathcal{S}), (\tilde{\mathcal{S}}^i)_{i \in \mathbf{I}}, P)$  — where  $(\Sigma, \mathcal{S}) = \prod_{i \in \mathbf{I}} (\Sigma^i, \mathcal{S}^i)$  and  $\tilde{\mathcal{S}}^i$  is the  $\sigma$ -field on  $\Sigma$  spanned by  $\mathcal{S}^i$  —, and where the equilibrium strategies are the projections from  $\Sigma$  to  $\Sigma^i$  ( $i \in \mathbf{I}$ ). (The corresponding devices  $c$  are called canonical correlation devices). Thus canonical correlated equilibria can be identified with their distributions.

**THEOREM 3.3. (Aumann, 1974)**  $C$  is the set of canonical correlated equilibria.

**PROOF.** The maps  $\hat{\sigma}_i$  of the previous proof allow to take equivalently as correlation device  $(E \times \Sigma, \mathcal{E} \times \mathcal{S}, (\mathcal{E}^i \times \tilde{\mathcal{S}}^i)_{i \in \mathbf{I}})$  with the probability induced by  $P$  and the maps  $\hat{\sigma}_i$ , and the projections to  $\Sigma^i$  as equilibrium. Thus each player  $i$  is told by the device some  $\sigma^i \in \Sigma^i$ , plus some additional information about  $\omega$ , and he uses the recommended  $\sigma^i$ . Hence he would a fortiori still be in equilibrium if he were only told  $\sigma^i$  — with less information, he has less strategies to deviate to. So the marginal on  $(\Sigma, \mathcal{S}, (\tilde{\mathcal{S}}_i)_{i \in \mathbf{I}})$  is a canonical correlated equilibrium. ■

**THEOREM 3.4.**  $C$  is weak $^*$ -closed if the action sets of  $\Gamma$  are separable metric spaces and each pay-off function  $g^i(\sigma^i, \sigma^{-i})$  is bounded and continuous on  $\Sigma$ .

**PROOF.**  $\Sigma$  is now also a separable metric space; the result follows then (cf. ex.II.3Ex.1d p. 96) from the fact that  $P \in C$  can be equivalently characterised by  $\int \Sigma \varphi(\sigma^i) [g^i(\sigma^i, \sigma^{-i}) - g^i(\tau^i, \sigma^{-i})] P(d\sigma) \geq 0$ ,  $\forall \tau^i \in \Sigma^i$ ,  $\forall i \in \mathbf{I}$ , and for every bounded, positive continuous function  $\varphi$  on  $\Sigma^i$ . ■

**COROLLARY 3.5.** If  $\Sigma$  is finite,  $C$  is a compact, convex polyhedron.

**PROOF.** In that case, the above system of linear inequalities in  $P$  becomes finite, taking into account that any positive  $\varphi$  is a positive linear combination of the (finitely many) indicator functions of singletons. ■

**REMARK 3.1.** The correlated equilibrium concept is purely non-cooperative — it does not require any binding commitments, and players could build for themselves, during the pre-play communication stage, a device that would make the required randomisation and signal to each player, once he has left the room and this stage is over, his recommended strategy.

**REMARK 3.2.** The assumption that the player remembers throughout the game his signal from the device — even in a game without perfect recall — fits completely with the standard interpretation for such games: it is the player (the “strategist”) who is sitting in the room, and gets the signal from the device. Afterwards he gives his instructions (including the signal if he wishes) to all his agents manning the different information sets. In the same spirit, and with the same motivation, is the classical assumption that each player knows what pure (or behavioural) strategy he uses — cf. e.g. the definition of effectively perfect recall.

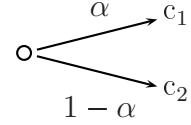


FIGURE  
3. Convexity  
of the Correla-  
ted Equilibria

**REMARK 3.3.** In another interpretation,  $E$  stands for the set of all states of the world — including therefore the strategies  $\sigma^i(e)$  players are going to use in the game —, and  $\mathcal{E}^i$  stands for player  $i$ 's private information (which includes  $\sigma^i(e)$ ). The “Harsanyi doctrine” requires that each player's subjective probability on  $E$  be the conditional distribution derived from a common prior  $P$  given his private information  $\mathcal{E}^i$ . From individual decision theory, each player maximises the expectation (with his subjective probability) of his utility function. It follows that  $(E, (\mathcal{E}^i)_{i \in \mathbf{I}}, P, (s^i)_{i \in \mathbf{I}})$  is a correlated equilibrium (Aumann, 1987).

**3.b. Multistage games, extensive form correlated equilibria.** Here we assume the game  $\Gamma$  under consideration is a **multistage game**. Intuitively, at each period, players play simultaneously, and they have effectively perfect recall. The easiest way to formalise this is to use a framework like that of sect. 1, except that now  $A = \prod_{i \in \mathbf{I}} A^i$ ,  $S = \prod_{i \in \mathbf{I}} S^i$ . Thus, at position  $\omega_n$ , a vector of signals  $(a_n^i)_{i \in \mathbf{I}}$  is selected according to  $\alpha_{\omega_n}$ , each player  $i$  is told his signal  $a_n^i$  and selects then his action  $s_n^i$  — then  $\omega_{n+1}$  is selected according to  $q_{\omega_n, a_n, s_n}$ .

Since the game has effectively perfect recall, we can use remark 1.3 p. 54, and assume that  $a^i$  contains just incremental information (it does not even have to contain the date, which the player knows by looking at the length of the sequence of past signals).

It is easy to see how to rewrite such a game in the general formalism of Section 1, subdividing each period into  $\mathbf{I}$  sub-periods where players play in turn (so all results and definitions of sect. 1 remain applicable). For the present problem, physical time is crucial however (cf. ex. II.3Ex.3 p. 96), and is more convenient to handle in the present formalism. For the need for effectively perfect recall, cf. ex. II.3Ex.6 p. 98.

[In the context we are going to deal with — correlation and communication — the differences between the general definitions of perfect recall sub II.1Ex.12b p. 65 and II.1Ex.13 p. 66 would conceivably matter, since the former apparently relies on players using independent strategies. And it is not clear what (intermediate?) form would be the “right” one. However, introducing as we do here effectively perfect recall as incremental information avoids such difficulties, and the sufficiency of mixed strategies in all extended games follows directly from ex. II.1Ex.10a and II.1Ex.10b p. 61 — one can stop at 7 p. 63, cf. also comment at the end of II.1Ex.10b p. 62].

An **autonomous device**  $a$  (for the player set  $\mathbf{I}$ ) is a probability space  $(E, \mathcal{E}, P)$  together with random variables  $m_n^i$  with values in a message space  $(M^i, \mathcal{M}^i)$ . The extension  $\Gamma_a$  of  $\Gamma$  by  $a$  is the game where first nature selects  $e \in E$  according to  $P$ , then before each stage  $n$  each player  $i$  is informed of  $m_n^i(e)$ . Formally  $\Gamma_a$  is the multistage game with  $(\tilde{\Omega}, \tilde{\mathcal{A}}) = (\Omega, \mathcal{A}) \otimes (E, \mathcal{E}) \otimes \mathbb{N}$ , with an initial position  $\tilde{\omega}_1$  added, from which nature selects (by  $\tilde{\alpha}$ ) the point  $(\omega_1, e, 1)$  ( $\omega_1$  deterministically,  $e$  according to  $P$ ). The  $e$  coordinate is preserved throughout the game, while every selection by  $q$  of a new state  $\omega$  is accompanied by a unit increment in the last coordinate. We have  $(\tilde{A}^i, \tilde{\mathcal{A}}^i) = (A^i, \mathcal{A}^i) \otimes (M^i, \mathcal{M}^i) \otimes \mathbb{N}$ , where at  $(\omega, e, n)$ , the first coordinate is selected according to  $\alpha_\omega$ , and the second is  $(m_n^i(e), n)$ .

An **extensive form correlated equilibrium** of  $\Gamma$  is a pair  $(a, \text{equilibrium of } \Gamma_a)$ . In such an extensive form correlated equilibrium, each player  $i \in \mathbf{I}$  has a private probability space  $(X^i, \mathcal{X}^i, Q^i)$  to do his own randomisation (cf. sect. 1); one can replace  $E$  by its product  $\overline{E}$  with all  $X^i$ 's, adding at every stage  $x^i$  to the current message  $m_n^i$ ; so as to obtain an equivalent behavioural strategy equilibrium. Further one can use ex. II.1Ex.10d

p.63 if the action set  $(S^i, \mathcal{S}^i)$  is standard Borel, or the method of ex.II.1Ex.10e p.63 if the set  $A^i$  is countable, to obtain an equivalent pure strategy equilibrium.

As in sect.3.a p.88, we assume henceforth  $\Gamma$  has countably many signals, and obtain:

**THEOREM 3.6.** *The extensive form correlated equilibria have distributions on the space  $(\Sigma, \mathcal{S}) = \prod_{i \in I} (\Sigma^i, \mathcal{S}^i)$  of pure strategy  $I$ -tuples of  $\Gamma$ . This set  $E$  of distributions is convex. ( $E_0$  will denote the set of pay-offs corresponding to  $E$ ).*

**PROOF.** Now there are also countably many finite sequences of signals, and recall from sect. 1 that pure strategies depend only on past signals — so one can still view the pure strategy space as a countable product of the action spaces. In particular, the pure strategy correlated equilibrium yields measurable maps  $\hat{\sigma}_n^i$  from  $\overline{E}$  to  $\Sigma_n^i = (S_n^i)^{\prod_{m \leq n} A_m^i}$ . The rest of the proof is the same as before. ■

**A canonical extensive form correlated equilibrium** is one where the device  $a = [\Sigma, \mathcal{S}, P; (\hat{\sigma}_n^i)_{n \in \mathbb{N}, i \in I}, (\Sigma_n^i)_{n \in \mathbb{N}, i \in I}]$  where  $\hat{\sigma}_n^i$  are the projections, and where the equilibrium strategies are the identity maps from messages to  $\Sigma_n^i$ . They are identified with their distributions  $P$ .

**THEOREM 3.7.**  *$E$  is the set of canonical extensive form correlated equilibria.*

**PROOF.** As for theorem 3.3 p.89. ■

**COMMENT 3.4.** In fact, as well as for correlated equilibria, to get a truly canonical representation, without redundancies, one should still eliminate duplicate strategies from  $\Sigma$  — i.e. identify any two equivalent pure strategies of  $\Sigma^i$  — it is useless for the device to tell the player what to do in case he has deviated from previous recommendation; this irrelevant (and potentially dangerous) information should be deleted for a canonical device.

**THEOREM 3.8.**  *$E$  is weak\*-closed if the action sets of  $\Gamma$  are separable metric spaces and each pay-off function  $g^i(\sigma^i, \sigma^{-i})$  is continuous on  $\Sigma$  and bounded.*

**PROOF.** The proof of theorem 3.4 p.89 has to be slightly modified:  $P \in E$  can be equivalently characterised by the system of inequalities:

$$\int [g^i(\hat{\sigma}) - g^i(\bar{\sigma}^i(\hat{\sigma}^i), \hat{\sigma}^{-i})] P(d\hat{\sigma}) \geq 0$$

where  $\bar{\sigma}_n^i$  is a continuous map from  $\hat{\sigma}_1^i \cdots \hat{\sigma}_n^i$  to probabilities over a fixed (i.e., depending only on  $\bar{\sigma}^i$ ) finite subset of  $\Sigma_n^i$ . Indeed, by Lebesgue's dominated convergence theorem, the same inequality will still hold if the  $\bar{\sigma}_n^i$  are just Borel instead of continuous — because when stabilising under point-wise limits the continuous maps from a separable metric space to a simplex one obtains all Borel maps. Hence the inequality is true whenever the  $\bar{\sigma}_n^i$  are Borel maps from  $\hat{\sigma}_1^i \cdots \hat{\sigma}_n^i$  to  $\Sigma_n^i$  taking only finitely many values. Since any measurable map to a separable metric space is a point-wise limit of measurable maps taking only finitely many values, another use of Lebesgue's theorem yields the inequality whenever the  $\bar{\sigma}_n^i$  are Borel maps from  $\prod_{k \leq n} \Sigma_k^i$  to  $\Sigma_n^i$ . Then it expresses that, in the extended game  $\Gamma_a$ , player  $i$  has no profitable pure strategy deviation, i.e.  $P \in E$ . The result follows, since all our linear inequalities are obviously weak\*-continuous. ■

**COROLLARY 3.9.** *If  $\Sigma$  is finite,  $E$  is a compact, convex polyhedron.*

**PROOF.** In this case, characterise  $E$  by the same set of inequalities as above, but where  $\bar{\sigma}^i$  is an arbitrary map from  $\prod_{k \leq n} \Sigma_k^i$  to  $\Sigma_n^i$ . ■

**3.c. Communication equilibria.** We still assume  $\Gamma$  a multistage game. But here we want the device not only to send messages to the players before every stage (“sun-spots”), but also to receive information from them, such as to embody the most general communication possibilities (e.g., letters given to a notary to be delivered at a later date, conditional to a specific event, to a specific subset of players). We still want the device to be completely outside the original game: it receives no direct information whatsoever about what is going on in the game, only through the players.

Thus, we think of the following scenario: At each stage  $n$ , first players receive their signal in  $\Gamma$ , then send some input to the device, then the device selects a vector of messages to the players, finally players choose an action in  $\Gamma$  (an illegal input by some player or an absence of input will be treated as a fixed input). Formally, we introduce the following:

**DEFINITION 3.10.** A tuple  $d = \{E, \mathcal{E}, e_1, [(H_n^i, \mathcal{H}_n^i)_{i \in \mathbf{I}}, (M_n^i, \mathcal{M}_n^i)_{i \in \mathbf{I}}, P_n]_{n \in \mathbb{N}}\}$  is called a **communication device**. Here  $e_1 \in E$  is the initial state of the device,  $H_n^i$  the space of inputs of player  $i$  at stage  $n$  and  $M_n^i$  the space of messages to him, while  $P_n$  is a transition probability from  $(E, \mathcal{E}) \otimes \bigotimes_{i \in \mathbf{I}} (H_n^i, \mathcal{H}_n^i)$  to  $(E, \mathcal{E}) \otimes \bigotimes_i (M_n^i, \mathcal{M}_n^i)$  that selects, given the current state  $e_n$  and the current inputs  $h_n^i$ , the current messages  $m_n^i$  and new state  $e_{n+1}$ .

**DEFINITION 3.11.** The extension  $\Gamma_d$  of  $\Gamma$  by  $d$  is the induced multistage game between players in  $\mathbf{I}$ . Observe however one could also consider it as a multistage game with  $\mathbf{I} + 1$  players, the last being the device, which has  $\prod_{i \in \mathbf{I}} (H_n^i, \mathcal{H}_n^i)$  as space of signals  $A_n^{\text{dev}}$  and  $\prod_{i \in \mathbf{I}} (M_n^i, \mathcal{M}_n^i)$  as action space  $S_n^{\text{dev}}$  — and uses the generalised strategy (cf. ex. II.1Ex.10ai p. 61)  $(E, \mathcal{E}, e_1, (P_n)_{n \in \mathbb{N}})$  (give him zero pay-off).

[It is easily seen that any generalised strategy can be written in the above form, taking for  $E$  the space of all finite histories of internal choices, inputs and messages.]

**DEFINITION 3.12.** A **communication equilibrium** of  $\Gamma$  is a pair  $(d, \text{equilibrium of } \Gamma_d)$ .

**DEFINITION 3.13.** A communication device is **standard** if  $(H_n^i, \mathcal{H}_n^i) = (A_n^i, \mathcal{A}_n^i)$ ,  $(M_n^i, \mathcal{M}_n^i) = (S_n^i, \mathcal{S}_n^i)$ .

**DEFINITION 3.14.** A **standard communication equilibrium** is one where the device is standard, and the equilibrium strategies are the identity maps. It is identified with the corresponding (generalised) strategy of the standard device.

**THEOREM 3.15.** Every communication equilibrium is equivalent to (induces the same probability distribution on plays as) a standard communication equilibrium.

**PROOF.** Given the communication device, and the equilibrium strategies of the players, we describe the construction of the new (and “larger”) standard device informally as follows (to avoid heavy notation, that would only obscure things). Think of the players’ strategies as personal devices (which remember all their past choices, inputs, and outputs), that would do all randomisations and computations for the player: the player would just have to instruct it at each stage of the signals he gets from the underlying game, to receive a recommended input to be sent to the central device, and of the message from the central device, to receive a recommended action to be taken in the underlying game. Certainly the player cannot deviate profitably, given such a setup, from being truthful to his own device and following its recommendations, since otherwise the composition of his deviation strategy and of his personal device would describe a profitable deviation from the given communication equilibrium. Assume now that he connects directly his personal

device to the central device, such that the personal device sends inputs directly to the central device, without even informing him, and receives directly the messages from the central device, without informing him either. Then he has even less moves to deviate to (he can no longer act as a middleman between his personal device and the central device), and less information — so certainly he has no profitable deviation. The central device together with all connected personal devices can now be seen as one single big standard device, to which players report truthfully their signals, and of which they follow the recommended actions. To construct formally the big device, just take care that the simultaneous randomisations in the different personal devices are done independently of each other; and take as state space the space of all finite histories of inputs, outputs, and internal choices of both the central device and all personal devices. ■

**COMMENT 3.5.** Now we have really an  $(\mathbf{I}+1)$ -person multistage game, since the sets of signals and actions of all  $(\mathbf{I}+1)$ -players are fixed, and no longer variable as in the general definition. And the communication equilibria are the equilibria of this game, while the standard communication equilibria are those where the strategies of the original  $\mathbf{I}$  players are the identity maps.

**COMMENT 3.6.** To write the multistage game, we have however to subdivide the stages of the original game — each stage has to be subdivided at least in 3. Actually, to define properly the relevant probabilities by Ionescu-Tulcea's theorem, it is more convenient to subdivide stage  $t$  into  $t_1 < t_2 < t_3 < t_4 < t_5$ :

- at  $t_1$ , nature chooses the  $a_i$  and informs the players
- at  $t_2$ , the players report to the device
- at  $t_3$ , the device selects its new internal state and sends its messages to the players
- at  $t_4$ , the players take an action in the game
- at  $t_5$ , a new position is chosen by nature.

**COROLLARY 3.16.** If the spaces  $(S^i, \mathcal{S}^i)$  are standard Borel, or if the sets  $A^i$  are countable, then

- (1) it suffices to consider pure strategy deviations by the players to determine the standard communication equilibria
- (2) standard communication equilibria have a representation in the form of a mixed strategy [with auxiliary space  $([0, 1], \lambda)$ ] for the device — in the countable case (“first canonical representation”), in the form of a probability distribution over  $\hat{\Sigma} = \prod_{n \in \mathbb{N}} (\prod_{a \in A_n} \overline{S}_{n,a}^{\prod_{t < n} A_t})$ , the space of pure standard devices (“joint pure strategies” of the players).

**PROOF.** Use ex. II.1Ex.10e p. 63 in the countable case, and ex. II.1Ex.10d p. 63 in the standard Borel case — for the players to obtain 1, and for the device (the  $(\mathbf{I}+1)^{\text{th}}$  player) for 2. ■

**COROLLARY 3.17.** If the spaces  $(S^i, \mathcal{S}^i)$  are standard Borel, standard communication equilibria have a second canonical representation in the form of a behavioural strategy for the device, i.e. a sequence of transition probabilities  $P_n$  from  $(\prod_{t < n} \overline{S}_t) \times A_n$  to  $S_n$  (carried by  $\overline{S}_{n,a_n}$ ).

**PROOF.** Use ex. II.1Ex.10a and II.1Ex.10b p. 61 for the device. ■

Standard communication equilibria in (first or second) canonical representation are called **canonical communication equilibria** (possibly with qualifier “first” or “second”).

Define  $D$  as the set of generalised strategies of the device in standard communication equilibria in the general case — in the standard Borel case, use the second canonical representation, and in the countable case, the first canonical representation.

Denote by  $D_0$  the set of corresponding pay-off vectors.

**THEOREM 3.18.**  *$D$  is convex (not in the second canonical representation).*

**PROOF.** A random selection between two generalised strategies is again a generalised strategy. ■

**THEOREM 3.19.** *In the countable case, if the spaces  $(S^i, \mathcal{S}^i)$  are separable metric spaces,  $D$  is weak\*-closed, if each player's pay-off function  $g^i$  is continuous and bounded over  $\hat{\Sigma}$ .*

**PROOF.** Use cor. 3.16.1 p. 93, and argue as in theorem 3.8 p. 91. (Note the additional action sets  $A^i$  of the player — for his inputs to the device — are also separable metric.) ■

**COROLLARY 3.20.** *If  $\Sigma$  is finite,  $D$  is a compact, convex polyhedron.*

**PROOF.** Note  $\hat{\Sigma}$  is then also finite, as well as the strategy sets of the players in the canonical  $\Gamma_d$ . Then as in cor. 3.9 p. 91. ■

**3.d. Finite games.** Here we assume  $\Sigma$  finite, and use Corollaries 3.5, 3.9 and 3.20. We now want to construct even nicer such finite devices, that are in particular independent of the particular equilibrium selected.

**THEOREM 3.21.** *For each finite game there exists:*

- (1) *A universal autonomous device, whose output is, in period zero a public integer and  $I$  private random variables uniform on a finite set  $A$ , all  $(I+1)$  of those independent; and in each of the following periods, a sequence of independent, public random variables, uniform on  $A$ . With this device, all extensive form correlated equilibria of the original game become feasible as pure strategy (Nash) equilibria of the extended game.*
- (2) *Similarly, a universal correlation device.*
- (3) *A universal communication device, which is a finite automaton all whose signals are public, and with which all communication equilibria of the original game become feasible as Nash equilibria of the extended game — alternatively one could use a randomising automaton, and pure Nash equilibria. The first alternative may require the game to have rational pay-offs.*

**REMARK 3.7.** In particular, there is no specific problem with those concepts of “bargaining about the device”, it is just the general problem of bargaining about (Nash) equilibria that one reencounters.

**PROOF.** The set of possible signals from the device to the players being finite, one could select an independent random permutation of this finite set for each day and each player; the device could then announce publicly every day the whole vector of encoded signals, provided it informs each player privately before the start of play of his sequence of decoding keys. In this way, all devices give an independent and uniformly distributed sequence of public signals from the start of the game on — the only private information is, as in a correlated equilibrium, before the start of play.

The device could even select an independent set of such keys for each extreme point of the polyhedra of communication equilibria or of extensive form correlated equilibria —

and simulate each of the extreme points in parallel independently, giving as single output every period the whole vector indexed by the extreme points. In the case of extensive form correlated equilibria, the players can very well be told initially their keys for all extreme points, since the irrelevant keys will just give them information independent of the true game and of the extreme point they want to play. But in the case of communication equilibria, they should receive only the relevant decoding keys. In that case however, they can be given each an additional finite set of inputs to the device, to be used at stage 1, by which they inform the device of which extreme point they want to play. Each player receives the decoding key relative to the extreme point he announced; his inputs are transmitted as is to the sub-device simulating that extreme point, and for the other sub-devices the first input is always selected as a fictitious input for him. So — even if, in a game with incomplete information, all this communication happens after players know their true types — they have no incentive to deviate from announcing the right extreme point to the device, assuming all others do so. Once the device knows the extreme point, it gives to each one his decoding key (together with his stage 1 output). Those decoding keys can now themselves be encoded, as before with the key for this encoding being given before the start of the game, to preserve the property that from stage 1 on, all the device's announcements are public, independent and uniform.

For correlated equilibria, it suffices to simulate independently and in parallel the devices corresponding to each of the extreme points. All this can be realised by a fixed, finite automaton containing a fixed, finite lottery mechanism — and the required finite sets of input buttons and output signals for each player. With this universal (for the game in question) device, all extreme point equilibria can be realised.

Use now ex. II.3Ex.2 p. 96 to add to it a coin tossing mechanism, by which it will output publicly, at the start of the game, a sequence of one's, deciding each time independently with probability  $2^{-n}$  to stop ( $2^n$  larger than the number of extreme points) — which it signals by appending a zero: we still have a finite automaton, and now players can use the number of 1's in the sequence to select an extreme point with whatever probability distribution they want. Thus, for each game, there is a fixed, finite automaton which is a universal communication device or autonomous device.

Finally, for the communication device, one can now use the possibility of inputs by the players to the device to make all outputs of the device public, and to dispense with the need for a lottery mechanism in the device, having the players themselves input their decoding keys to the device, and generate the required randomness by a jointly controlled lottery mechanism (ex. II.3Ex.5) — at least when the parameters of the game are rational. There is no problem for the sequence of 0's and 1's needed to select the extreme point: players successively choose simultaneously at random a 0 or a 1, and the device just re-transmits to each the sum mod 2: after a finite number of trials, players will know the corresponding integer, hence the corresponding extreme point, and will transmit this to the automaton. After that, players have to generate the finite lottery required by the device: if this is rational, it can be replaced by a finite lottery over a bigger number of events, say  $n$ , with equal probability: the players pick to this effect a uniform number in  $\mathbb{Z}_n$ , the machine uses the sum. ■

### Exercises.

- 1. A direct proof of existence of correlated equilibrium (finite case).** (Hart and Schmeidler, 1989) Consider a two-person zero-sum game  $G$  where player I chooses a point  $s$  in

$S = \prod S^i$ , player II chooses a triple  $(i, r^i, t^i)$ ,  $i \in \mathbf{I}$ ,  $r^i, t^i \in S^i$  and the corresponding pay-off is  $h^i(s^{-i}, r^i) - h^i(s^{-i}, t^i)$  if  $s^i = r^i$  and 0 otherwise.

- a. Prove that, if  $v(G) \geq 0$ , an optimal strategy of I induces a correlated equilibrium.
- b. Prove then that, given any mixed strategy  $y = (y^i(r^i, t^i))$  of player II, player I can get 0.

HINT. By ex.I.3Ex.9 p. 34, given non-negative numbers  $a_{\ell m}$ ,  $\ell, m \in M$ , there exists  $\alpha \in \Delta(M)$  such that, for any  $\beta$  in  $\mathbb{R}^M$ ,  $\phi(\alpha, \beta) = \sum_{\ell} \alpha_{\ell} \sum_m a_{\ell m} (\beta_{\ell} - \beta_m) = 0$ . Take then, for  $i$  fixed,  $a_{\ell m} = y^i(r^i, t^i)$ ,  $\alpha_{\ell} = x^i(r^i)$ ,  $\beta_{\ell} = h^i(s^{-i}, r^i)$ ,  $\beta_m = h^i(s^{-i}, t^i)$  and deduce that  $x$  defined by  $x(r) = \prod_i x^i(r^i)$  gives 0 in  $G$ .

c. (Peleg, 1969) Consider now the following game  $\Gamma$ : the set of players is  $\mathbb{N}$  and each player's strategy set is  $\{0, 1\}$ . The pay-off function is given by

$$g^i(s) = \begin{cases} s^i, & \text{if } \sum_j s^j < \infty \text{ (finitely many ones);} \\ -s^i, & \text{otherwise.} \end{cases}$$

Prove that there exists no Nash equilibrium.

HINT. Use the zero-one law.

Let  $s(j)$  in  $S$  defined by  $s^i(j) = 1$  iff  $i \leq j$ . Define  $P_1$  with support on the  $s(j)$ 's by  $P_1(s(j)) = (1/j) - 1/(j+1)$  and  $P_0$  product of its marginals on the  $S^i$ 's with  $P_0(s^i = 1) = 1/i$ . Prove then that  $(P_0 + P_1)/2$  is a correlated equilibrium.

d. The above method also yields the existence of correlated equilibria for continuous pay-off functions on compact strategy spaces.

HINT. Either use a direct approximation, or let player II choose a triple of a player  $i$ , a strategy  $s_0 \in S^i$  and a continuous non-negative function  $\varphi$  on  $S^i$ , with pay-off  $[h^i(s^{-i}, s^i) - h^i(s^{-i}, s_0)]\varphi(s^i)$ . Note that in step II.3Ex.1a, prop. 1.17 p. 7, which uses strategies with finite support for player II, reduces the problem to one with finite strategy spaces. Continuity of the pay-off function is used to pass from the non-profitability of step-function deviations to that of arbitrary deviations.

COMMENT 3.8. In whatever way one uses such a method, one will need the upper semi-continuity of  $h^i(s^{-i}, s^i) - h^i(s^{-i}, s_0)$  on  $S$  (player I's strategy space in the fictitious game), for each  $i$  and  $s_0 \in S^i$ . (Observe also that this is exactly what is needed for the upper semi-continuity of the best reply correspondence). This is equivalent to  $h^i(s^{-i}, s^i) = f^i(s^{-i}, s^i) + g^i(s^{-i})$ , where  $f^i$  is u.s.c. on  $S$  and continuous in  $s^{-i}$  for each  $s^i$ .  $g^i$  does not affect Nash equilibria or correlated equilibria, and we know existence of Nash equilibria for  $f$  — ex.I.4Ex.20 p. 48 —, at least in the compact metric case. (And note that the continuity assumption in II.3Ex.1d implies one can reduce the problem to the compact metric case, e.g. by Stone-Weierstrass.) Thus refinements of this method are unlikely to yield existence of correlated equilibria under much wider conditions than those known for the existence of equilibria.

## 2. (Blackwell, 1953)

a. Prove that, for  $\lambda \leq \frac{1}{n}$ ,  $\lambda_m = \lambda(1 - \lambda)^{m-1}$  and  $z$  in the simplex of  $\mathbb{R}^n$ , there exists a partition of  $\{1, 2, \dots, m, \dots\}$ , say  $N_i$ ,  $i = 1, \dots, n$  with  $\sum_{m \in N_i} \lambda_m = z_i$ .

b. Prove that if  $\eta_m$  is any other probability definition with the same property as  $\lambda_m$  above sub a), and if  $\lambda = \frac{1}{n}$ , then  $\sum_{n \leq N} \eta_m \leq \sum_{m \leq N} \lambda_m$  for all  $N$ .

c. Prove that, if, along every play, player  $i$  moves at most once, with at most  $n$  actions, then any mixed (or behavioural) strategy of  $i$  is equivalent to a countable mixture of pure strategies with weights  $\lambda_m$  and  $\lambda = \frac{1}{n}$ .

## 3.

a. Prove in the following game that  $(2, 2)$  is an extensive form correlated equilibrium pay-off but not a correlated equilibrium pay-off.

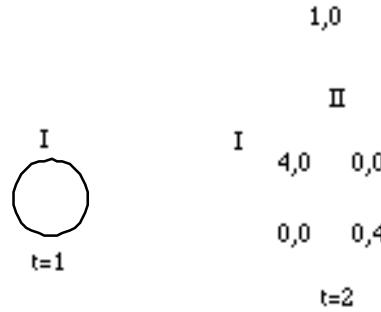


FIGURE 4. Extensive form Correlated Equilibria are not Correlated Equilibria

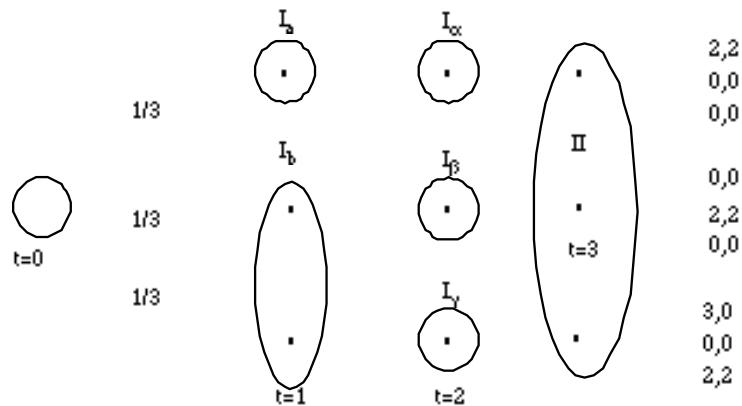


FIGURE 5. Necessity of the timing structure

b. Consider the multistage game of figure 5.

Show that  $(2, 2)$  is a communication equilibrium. It is no longer so if the information is given at once at time  $t = 1$ . This shows the additional structure of multistage games — i.e., the time function — is essential for this concept.

#### 4. (Forges, 1990a)

a. In a correlated equilibrium of two-person zero-sum game with finite strategy sets, the conditional probability over the opponent's actions given a pure strategy having positive probability is an optimal strategy of the opponent.

b.  $\begin{pmatrix} \frac{1}{3} & \frac{1}{3} & 0 \\ \frac{1}{3} & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$  is a correlated equilibrium of the game  $\begin{pmatrix} 0 & 0 & 1 \\ 0 & 0 & -1 \\ -1 & 1 & 0 \end{pmatrix}$ , while every convex combination of pairs of optimal strategies with  $p(2, 2) = 0$  satisfies  $p(1, 1) \geq \frac{1}{2}$ .

c. The pay-off to any communication equilibrium in a two-person zero-sum game is the value.

5. Define a general jointly controlled lottery over a finite set of alternatives  $A$  as a finite game — with elements of  $A$  instead of pay-off vectors — together with an  $n$ -tuple of strategies such that no player can affect the resulting probability distribution on  $A$  by unilateral deviation.

a. Show that this implies that knowledge of his own action gives a player no information about the outcomes in  $A$ .

b. Show that general joint lotteries can only generate algebraic, hence countably many, distributions on  $A$  (and with two players only rational ones).

HINT. Use ex. I.4Ex.21 p. 49.

c. Show that any rational distribution on  $A$  can be realised by an  $n$ -player lottery for any  $n \geq 2$ , which is fully symmetric in players and strategies, and where no proper coalition can affect the outcome.

HINT. Let  $k$  be the smallest common denominator of the probabilities: it suffices to select a uniform point in  $\mathbb{Z}_k$ . Each player does and the results are added.

d. General joint lotteries which are stable against deviations by proper sub-coalitions (like in II.3Ex.5c) yield rational probability distributions.

HINT. For  $a_0 \in A$ , use the minmax theorem for the games between one player and the opposing coalition with as pay-off 1 if  $a_0$ , 0 otherwise.

**6.** Consider a multistage game:

At stage 1 players I and II play the following game with incomplete information, player I being informed:

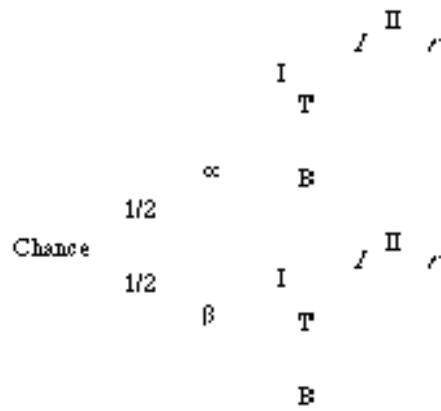


FIGURE 6. Stage 1 of the multistage game

At stage 2, player I plays in the game of Figure 7, where the pay-off of player II is identically

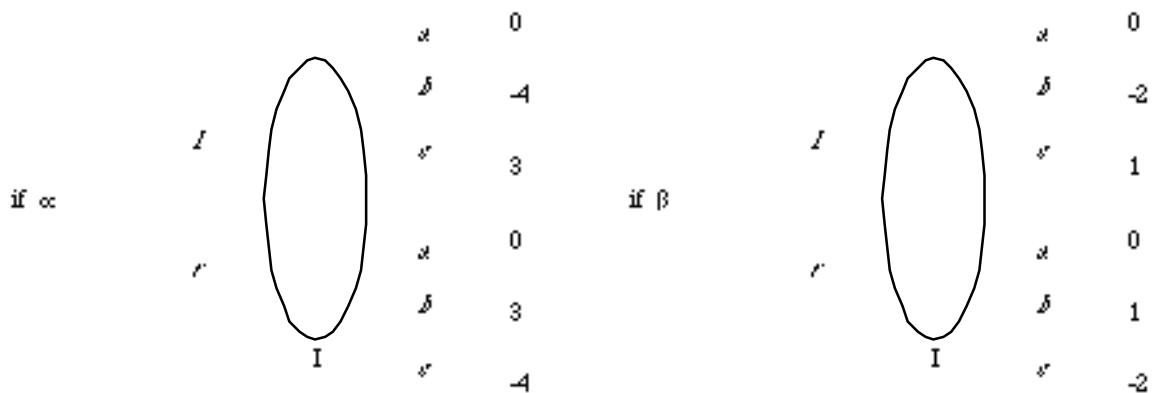


FIGURE 7. Stage 2 of the multistage game

- a. Show that the following distribution  $P$  on moves at time 1:

$$\begin{array}{cc} \ell & r \\ \begin{matrix} TT \\ TB \\ BT \\ BB \end{matrix} & \begin{pmatrix} 1/9 & 2/9 \\ 5/9 & 5/18 \\ 2/9 & 4/9 \\ 1/9 & 1/18 \end{pmatrix} \end{array} \quad (TB \text{ is } T \text{ if } \alpha, B \text{ if } \beta)$$

with player I choosing  $a$  at  $t = 2$ , defines an extensive form correlated equilibrium if the players do not recall their signal at stage 2 (but perfect recall of actions is still assumed).

b. Add a third player playing at stage 1, informed of the chance move. His pay-offs are independent of  $a, b, c$ . Define 9 strategies as follows: strategy 0 correspond to the above game and gives him pay-off 0. Strategy  $\alpha T\ell$  defines a pay-off  $x_{\alpha T\ell}$  for III if  $(\alpha T\ell)$ ,  $-1$  otherwise, with expectation zero under  $P$  and pay-off 0 for I and II, and similarly for all seven other outcomes. Add finally a dummy player with pay-off 1 if I chooses  $a$  and III uses 0, and 0 otherwise. Show that no extensive form correlated equilibria (with perfect recall) achieves the pay-off obtained by  $P$ .

- c. Prove the same result for communication equilibria.

HINT. Compute first the conditional distributions on  $(T\ell, Tr, B\ell, Br)$  given  $\alpha$  and  $\beta$ , and show that for any other conditional distributions, player III would deviate. Show that if the posteriors on  $\ell$  computed by player I given his signals are requested to be in  $[1/3, 2/3]$ ,  $P$  is the only correlated distribution compatible with the above conditionals. Check finally that with these conditional distributions player I will have an incentive to lie about his type.

COMMENT 3.9. It is clear that any extensive form correlated equilibrium would still be one (in terms of the induced distribution of plays) if the players did not recall previous messages — let just the device repeat them. So if this modification of the definition would not increase the set of equilibrium pay-offs, for games with perfect recall, it might have given a plausible extension of the definition to games without perfect recall. The above example (which has probably generic variants) shows this is definitely not so. Even in the most standard cases, it yields pay-offs that are not even communication equilibrium pay-offs: the most basic incentive constraints are violated. Hence the reason for sticking in the definition of multistage games to the perfect recall assumption: although technically one could do with much less, it is the only framework where we feel for the moment completely comfortable with the assumption of perfect recall of past signals — indeed, the justification in remark 3.2 p. 89 no longer applies; on the other hand, parallelism with the correlated equilibrium situation is an additional reason for maintaining perfect recall of past signals.

The example shows also that a variant applicable to any extensive form game, where the device would send a separate signal to every agent — i.e., correlated equilibria of the agent normal form —, would a fortiori be too large. This is an additional reason (besides ex. II.3Ex.3 p. 96) for sticking for the time being to the framework of multistage games.

## 7. Protocols and correlated equilibria. (Barany, 1992)

- a. Let  $\#\mathbf{I} = 4$ ,  $S^i$  finite,  $S = \prod S^i$  and  $E$  a finite set endowed with a partition  $\{E_s\}_{s \in S}$  and projections  $\text{Proj}_i: E \rightarrow S^i$  with  $\text{Proj}_i(e) = s_i$  if  $e \in E_s$ .

We describe here a procedure (protocol) of communication between the players according to which a point  $e$  in  $E$  will be chosen uniformly at random in  $E$ , each player being only informed of  $\text{Proj}_i(e)$ . Moreover the probabilities are the same under non-detectable deviations.

- (1) Recall that any subset of at least 2 players can generate in a non manipulable way a discrete random variable  $X$  (say, the choice at random of a permutation of a finite set) that is common knowledge (ex. II.3Ex.5 p. 97). (Note that with 4 players there exists a sequential procedure of binary communications where at the end only I and II know  $X$ :

let III and IV choose  $X_{\text{III}}$  resp.  $X_{\text{IV}}$  and inform I and II — who check that the signal is the same and then use  $X_{\text{III}} \circ X_{\text{IV}}$ ).

- (2) For all  $i$  and all  $j$ , let  $\alpha_i$  (resp.  $\beta_{i,j}$ , resp.  $\gamma$ ) be a permutation of  $E$  chosen at random and known by  $-i$  (i.e.  $\mathbf{I} - \{i\}$ ) (resp.  $\{i, j\}$ , resp.  $\{\mathbf{I}, \mathbf{II}\}$ ). Let also  $f$  be chosen at random in  $E$  and known by  $\{\text{III}, \text{IV}\}$ .

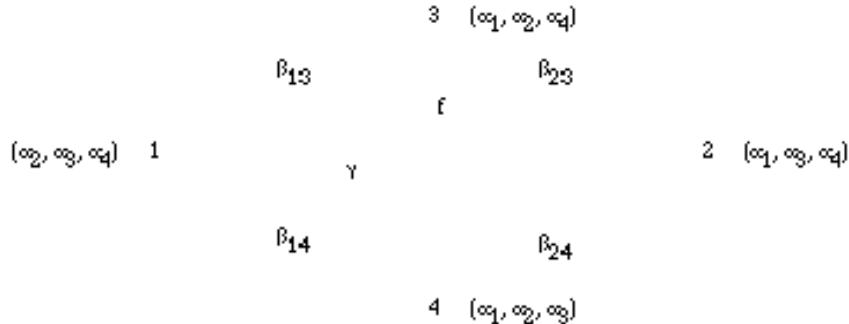


FIGURE 8. The Protocol

At stage 1, I is informed of the following messages:  $\beta_{\text{II},\text{III}} \circ \alpha_I$  and  $\beta_{\text{II},\text{IV}} \circ \alpha_I$  from II,  $\beta_{\text{II},\text{III}} \circ \alpha_I$  from III,  $\beta_{\text{II},\text{IV}} \circ \alpha_I$  from IV. A similar procedure is applied to each player  $i = \text{II}, \text{III}, \text{IV}$ .

At stage 2, III receives  $\beta_{\text{I},\text{IV}} \circ \gamma$  from I and  $\beta_{\text{II},\text{IV}} \circ \gamma$  from II (and symmetrically for IV).

At stage 3, I receives  $\beta_{\text{II},\text{IV}} \circ \gamma(f)$  from III and  $\beta_{\text{II},\text{III}} \circ \gamma(f)$  from IV (and dually for II).

At stage 4, I receives  $\text{Proj}_I \circ \beta_{\text{II},\text{III}}^{-1}$  from II and III, II gets  $\text{Proj}_{\text{II}} \circ \beta_{\text{I},\text{IV}}^{-1}$  from I and IV, III obtains  $\text{Proj}_{\text{III}} \circ \beta_{\text{II},\text{IV}}^{-1}$  from II and IV, and IV is told  $\text{Proj}_{\text{IV}} \circ \beta_{\text{I},\text{II}}^{-1}$  by I and II. The chosen point is  $e = \gamma(f)$ . Note that each player  $i$  knows  $\text{Proj}_i(e)$ , that at each stage the players can check detectable deviations (i.e. false messages). Show that the above procedure answers the question.

b. Deduce that in any game with *at least 4 players* and rational pay-offs any correlated equilibrium distribution can be realised as a Nash equilibrium in the game extended by finite pre-play communication with recording (i.e. each player has a move (STOP) that reveals all past events).

**8. Communication equilibria and correlated equilibria in games of information transmission.** (Forges, 1985) Let  $G$  be a two-person game with incomplete information (cf. sect. 4 p. 131), with (finite) types set  $K$  for player I and (finite) actions set  $S$  for player II. The types sets of player II and the actions sets of player I are assumed to be singletons. The game  $G_M$  obtained by allowing player I to send a message in  $M$  to player II after the choice of nature is thus a “game of information transmission”. Let  $q$  be a canonical communication equilibrium of  $G$  (described by a transition probability  $q$  from  $K$  to  $S$ ).

There exists a correlated equilibrium of  $G_M$  (associated with a finite set  $M$ ) that induces the same conditional distributions on  $S$  (given  $k \in K$ ) than  $q$ .

HINT. (1) Let  $P$  be the product probability distribution induced by  $q$  on  $S^K$ :

$$P(f) = \prod_k q(f(k) | k).$$

Let  $M_0$  be a copy of  $K \times A$ . We first define a correlation device for  $G_{M_0}$ , which transmits  $\sigma_0 \in M_0^K$  to player I and  $\tau_0 \in S^{M_0}$  to player II. The device selects a permutation  $\pi_0$  of  $M_0$ ,

uniformly and independently of  $f$ . Let then

$$\begin{cases} \tilde{f}(k) &= (k, f(k)), k \in K \\ \sigma_0 &= \pi_0 \circ \tilde{f} \\ \tau_0 &= \text{Proj}_S \circ \pi_0^{-1} \end{cases}$$

Let  $Q_0$  be the induced probability distribution over  $K \times M_0^K \times S^{M_0}$ .  $Q$  has the following properties:

- (a)  $\tau_0 \circ \sigma_0$  is independent of  $\sigma_0$  and of  $\tau_0$ ;  $\tau_0 \circ \sigma_0)(k)$  is independent of  $(\tau_0 \circ \sigma_0)(\ell)$  if  $k \neq \ell$ .
  - (b)  $Q_0((\tau_0 \circ \sigma_0)(k) = s) = q(s | k) \forall k \in K, \forall s \in S$
  - (c)  $Q_0(k | \sigma_0(k) = m, \tau_0) = p^k q(\tau_0(m) | k) / \sum_\ell p^\ell q(\tau_0(m) | \ell)$ , i.e., the posterior probability distributions over  $K$  derived from the communication equilibrium and from the present scheme are the same.
  - (d)  $Q_0(\tau_0(m) = s | \sigma_0) = q(s | k)$  if  $m = \sigma(k)$ ,  $k \in K$ .
- (2) The previous scheme can be completed so as to deter player I from sending  $m \notin \sigma(K)$ . Let  $M$  be a finite set (of at least  $\#(K \times A)$  elements). Let the correlation device choose  $M_0$  uniformly in  $M$  and  $s \in S$  according to some  $\rho \in \Delta(S)$ , independently of all previous choices. Define  $\sigma \in M^K$  by  $\sigma(k) = \sigma_0(k) \forall k \in K$  and  $\tau \in S^M$  by

$$\begin{aligned} \tau(m) &= \tau_0(m) && \text{if } m \in M_0 \\ &= s && \text{if } m \notin M_0. \end{aligned}$$

The correlation device only transmits  $\sigma$  (resp.  $\tau$ ) to player I (resp. II). In particular, player I does not know  $M_0$ . Let  $Q$  denote the induced probability over  $K \times M^K \times S^M$ . Then:

- properties 1a–1d still hold with  $Q, \sigma, \tau$  instead of  $Q_0, \sigma_0, \tau_0$
- $\pi$  and the size of  $M$  can be chosen so that  $Q(\tau(m) = s | \sigma) = \frac{1}{\#K} \sum_k q(s | k)$  if  $m \notin \sigma(K)$ : let  $\rho(s) = \frac{1}{\#M - \#M_0} \left[ \frac{\#M}{\#K} \sum_k q(s | k) - \#K \right]$  with  $\#M \geq \max_s \left[ \frac{(\#K)^2}{\sum_k q(s | k)} \right]$

## 9. Universal mechanisms and communication equilibria. (Forges, 1992)

The previous result can be generalised to an  $\#\mathbf{I}$ -person game  $G$  with incomplete information ( $\#\mathbf{I} \geq 4$ ). Let  $K^i$  (resp.  $S^i$ ) be player  $i$ 's finite set of types (resp. actions).

a. Let  $q$  be a canonical communication equilibrium in  $G$  (described by a transition probability from  $K = \prod_i K^i$  to  $S = \prod_i S^i$ ). Let us extend  $G$  by two stages of pre-play communication after the choice of nature.

**Stage 1:** Every player  $i$  sends a public message in a copy  $L^i$  of  $K^i$  (let  $L = \prod_i L^i$ ).

**Stage 2:** Every player  $i$  sends a message in  $M^j = L \times S^j$  to every player  $j$  in  $\mathbf{I}^i = \{i+1, i+2, i+3\} \pmod{\mathbf{I}}$ .

The correlation device first selects uniformly, independently of each other, the bijections  $\gamma^i: L^i \rightarrow K^i$  ( $i \in \mathbf{I}$ ) and independently of this choice,  $f \in S^K$  according to the probability distribution induced by  $q$  (cf. ex. II.3Ex.8). Let  $\gamma = (\gamma^i)_{i \in \mathbf{I}}$ , let  $g$  be defined by  $g \circ \gamma = f$  and let  $g^i = \text{Proj}_i \circ g$  (where  $\text{Proj}_i$  is the projection on  $S^i$ ).

Given  $g^i \in (S^i)^L$ , the device chooses  $\sigma^{ij} \in (M^j)^L$  and  $\tau^i \in (A^i)^{M^i}$  ( $i \in \mathbf{I}, j \in \mathbf{I}^i$ ) exactly as  $(\sigma_0, \tau_0)$  given  $f$  in ex. II.3Ex.8. The strategy for  $i$  is, after receiving  $(\gamma^i, \sigma^{ij}, \tau^i)$  (and  $k^i$ ), to announce  $\gamma^i(k^i)$  then to send  $\sigma^{ij}(\ell)$  to  $j \in \mathbf{I}^i$  if  $\ell$  is the message at stage one. Finally if at least two messages received by  $i$  coincide, say to  $m^i$ , he plays  $\tau^i(m^i)$ .

Check that the above mechanism is in fact an equilibrium.

HINT. Note that the information during the pre-play communication phase is non-revealing, and that given  $\tau^i(m^i) = s^i$  it is the same as in the initial communication scheme. Moreover to announce  $\ell^i \neq \gamma^i(k^i)$  and to play  $\bar{s}^i \neq s^i$  is the same as using  $\bar{k}^i = (\gamma^i)^{-1}(\ell^i)$  and  $\bar{s}^i$  in the communication scheme.

COMMENT 3.10. The above mechanism is universal (i.e. independent of the game  $G$ ) and does not require recording (compare with the previous ex. II.3Ex.7).

COMMENT 3.11. 3 players suffice if an alarm is allowed or messages in  $[0, 1]$  can be used. (Inform player  $j$  of a code (in  $[0, 1]$ ) on  $L \times M^j$  while  $i$  knows only its restriction to the graph of  $\sigma^{ij}$ ,  $j \in \mathbf{I}_i$ ).

b. Deducer that if the pay-offs are rational in  $G$ , any communication equilibrium distribution can be obtained as a Nash equilibrium of the game extended by pre-play communication (after chance's move).

#### 4. Vector pay-offs

In this paragraph, we consider an analogue due to Blackwell (1956a) of the minmax theorem for matrix games with vector pay-offs. Let  $(S, \Sigma)$  and  $(T, \mathcal{T})$  be two-measurable spaces,  $\phi$  a measurable mapping from  $(S, \Sigma) \times (T, \mathcal{T})$  into the set of all probability distributions on Euclidian space  $\mathbb{R}^k$  having finite first order moments. At any stage  $n$  ( $n = 1, \dots$ ), player I chooses a point  $s_n$  in  $S$  and II chooses simultaneously a point  $t_n$  in  $T$ ; a point  $x_n$  in  $\mathbb{R}^k$  is then chosen at random according to the distribution  $\phi(s_n, t_n)$ , independently of all other choices.

Both players may get, at every stage, some information, which — for I — includes at least  $x_n$  (and all  $x_i$ ,  $i \leq n$ ). Let  $\bar{x}_n = \frac{1}{n} \sum_1^n x_i$ .

DEFINITION 4.1. A set  $C$  in  $\mathbb{R}^k$  is **approachable** for a player if he has a strategy  $\sigma$  such that, for each  $\varepsilon > 0$ , there exists an integer  $N$  such that, for every strategy  $\tau$  of his opponent,

$$\Pr_{\sigma, \tau}(\sup_{n \geq N} \delta_n \leq \varepsilon) \geq 1 - \varepsilon \quad \text{and} \quad \sup_{n \geq N} E_{\sigma, \tau}(\delta_n) \leq \varepsilon$$

where  $\delta_n = d(\bar{x}_n, C)$ .

DEFINITION 4.2. A set  $C$  in  $\mathbb{R}^k$  is **excludable** for a player if, for some  $\varepsilon > 0$ , the set  $\bar{C}^\varepsilon$  is approachable for that player, where  $\bar{C}^\varepsilon = \{x \mid d(x, C) \geq \varepsilon\}$ .

REMARK 4.1. A set  $C$  is approachable iff its closure  $\bar{C}$  is.

Denote by  $f(s, t)$  the barycentre (mean value) of  $\phi(s, t)$ , and assume the  $\phi(s, t)$  have second-order moments uniformly bounded by  $K$ . Denote also by  $P$  (resp.  $Q$ ) the set of one-stage mixed strategies of I (resp. II) (i.e. probabilities on  $(S, \Sigma)$ , resp.  $(T, \mathcal{T})$ ). For any  $p \in P$ , denote by  $R(p)$  the convex hull of  $\{\int_S f(s, t) dp \mid t \in T\}$ .

THEOREM 4.3. Let  $C$  be any closed set in  $\mathbb{R}^k$ . Assume that, for every  $x \notin C$ , there is a  $p (= p(x)) \in P$  such that the hyperplane through  $y(x)$  (a closest point in  $C$  to  $x$ ), and perpendicular to the line segment  $(x, y)$  separates  $x$  from  $R(p)$ . Then  $C$  is approachable by I with the strategy  $f = (f_n)_{n \in \mathbb{N}}$ , where

$$f_{n+1} = \begin{cases} p(\bar{x}_n) & \text{if } n > 0 \text{ and } \bar{x}_n \notin C, \\ \text{arbitrary} & \text{otherwise.} \end{cases}$$

With that strategy, one has

$$E(\delta_n^2) \leq 4K/n \quad \text{and} \quad \Pr(\sup_{n \geq N} \delta_n \geq \varepsilon) \leq 8K/\varepsilon^2 N.$$

PROOF. Assume I uses the above described strategy. Let  $y_n = y(\bar{x}_n)$  and  $u_n = \bar{x}_n - y_n$ . The definition of  $f$  implies that (note that  $u_n = 0$  if  $\bar{x}_n \in C$ ):

$$(1) \quad E(\langle u_n, x_{n+1} \rangle \mid x_1, \dots, x_n) \leq \langle u_n, y_n \rangle ; \quad \text{and:}$$

$$(2) \quad \delta_{n+1}^2 \leq \|\bar{x}_{n+1} - y_n\|^2 = \|\bar{x}_n - y_n\|^2 + 2\langle \bar{x}_n - y_n, \bar{x}_{n+1} - \bar{x}_n \rangle + \|\bar{x}_{n+1} - \bar{x}_n\|^2$$

Since

$$\bar{x}_{n+1} - \bar{x}_n = (x_{n+1} - \bar{x}_n)/(n+1) = ((x_{n+1} - y_n) - (\bar{x}_n - y_n))/(n+1) ,$$

one gets from (2), taking conditional expectations and using (1):

$$(3) \quad \mathbb{E}(\delta_{n+1}^2 | \mathcal{F}_n) \leq (1 - 2/(n+1)) \delta_n^2 + w_n/(n+1)^2$$

with  $\mathcal{F}_n$  the  $\sigma$ -field spanned by  $x_1, \dots, x_n$ , and  $w_n = \mathbb{E}(\|x_{n+1} - \bar{x}_n\|^2 | \mathcal{F}_n)$ . Note that

$$\mathbb{E}(w_n) \leq 2 \mathbb{E} \|x_{n+1}\|^2 + 2 \mathbb{E} \|\bar{x}_n\|^2 \leq 4K .$$

We claim that, for any sequence of random variables satisfying (3), where  $w_n \geq 0$  and  $\delta_n$  are  $\mathcal{F}_n$ -measurable and where  $\mathbb{E}(w_n) \leq 4K$ , one has

$$\mathbb{E}(\delta_n^2) \leq \frac{4K}{n} \quad \text{and} \quad \Pr(\sup_{n \geq N} \delta_n \geq \varepsilon) \leq \frac{8K}{\varepsilon^2 N} .$$

Let first  $e_n = \mathbb{E}(\delta_n^2/4K)$ ; from (3),  $e_n \leq (1 - 2/n)e_{n-1} + 1/n^2$ . It is readily checked that this implies, by induction for  $n \geq 2$ , that  $e_n \leq 1/n$  (which, by the way, obviously holds for  $n = 1$ ). So  $\mathbb{E}(\delta_n^2) \leq 4K/n$ . This implies in particular that  $\delta_n$  converges in probability to zero. To show that it converges to zero with probability one, let  $Z_n = \delta_n^2 + \mathbb{E}(\sum_{i=1}^{\infty} w_i/(i+1)^2 | \mathcal{F}_n)$ . Replacing, in (3),  $(1 - 2/(n+1))$  by 1, one sees that (3) implies  $\mathbb{E}(Z_{n+1} | \mathcal{F}_n) \leq Z_n$ ; thus  $Z_n$  is a positive supermartingale, and

$$\mathbb{E}(Z_n) \leq \mathbb{E}(\delta_n^2) + 4K \sum_{i=1}^{\infty} 1/i^2 \leq 4K/n + 4K/n = 8K/n .$$

The supermartingale convergence theorem implies then that  $Z_n$  goes to zero with probability one — and a fortiori  $\delta_n$ , since  $0 \leq \delta_n^2 \leq Z_n$ . More precisely, we get from Doob's maximal inequality for supermartingales (Neveu, 1970, prop. IV.5.2) that

$$\Pr(\exists n \geq N : Z_n \geq \varepsilon) \leq \mathbb{E}(Z_N)/\varepsilon \leq \frac{8K}{\varepsilon N} ,$$

and thus the result follows. ■

**COMMENT 4.2.** The approaching player has to know only  $\bar{x}_n$  and not  $x_1, \dots, x_n$  at stage  $n$ .

**COMMENT 4.3.** Theorem 4.3 obviously remains true under the condition that, for any  $x \notin C$ , there is a sequence  $p_i \in P$  such that

$$\lim_{i \rightarrow \infty} \sup_{z \in R(p_i)} \langle z - y(x), x - y(x) \rangle \leq 0 .$$

(Choosing, at stage  $n$ , a  $p_i$  such that this supremum is  $\leq \delta/n^2$ , where  $\delta$  is some small positive number, and increasing  $K$  by  $\delta$ ).

**COROLLARY 4.4.** For any  $p \in P$ ,  $R(p)$  is approachable by I using the constant strategy  $f_n = p$ .

For any  $q \in Q$ , denote by  $T(q)$  the closed convex hull of  $\{ \int_T f(s, t) dq \mid s \in S \}$ .

**THEOREM 4.5.** Let  $C$  be a closed convex set in  $\mathbb{R}^k$ , and assume that the game with pay-off function  $\langle u, f(s, t) \rangle$  has a value  $v(u)$  for all  $u$  in  $\mathbb{R}^k$  such that  $\inf_{x \in C} \langle u, x \rangle > -\infty$ . Then:

- (1)  $C$  is approachable by I iff, for all  $q \in Q$ ,  $T(q) \cap C \neq \emptyset$ ;

- (2) if  $T(q_0) \cap C = \emptyset$ , then  $C$  is excludable by II (using the constant strategy  $q_0$ );
- (3) the condition of remark 4.3 above is necessary and sufficient for approachability of  $C$ .

PROOF. The second statement follows right away from the previous corollary. There remains only to be shown that, if the condition of remark 4.3 is not satisfied, there exists a  $q_0$  such that  $T(q_0) \cap C = \emptyset$ . Let  $x_0 \notin C$  be such that the condition fails, and let  $u = y(x_0) - x_0$ . Then there exists  $\varepsilon > 0$  such that, for each  $p \in P$ ,  $\inf_{x \in R(p)} \langle u, x \rangle < \min_{x \in C} \langle u, x \rangle - \varepsilon = M$ . Thus  $v(u) \leq M$  and therefore there exists  $q_0 \in Q$  such that  $\max_{x \in T(q_0)} \langle u, x \rangle \leq M + \varepsilon/2 = \min_{x \in C} \langle u, x \rangle - \varepsilon/2$ . The result follows. ■

COMMENT 4.4. The condition on having a value will be automatically satisfied for instance if  $S$  and  $T$  are compact and  $f(s, t)$  is continuous in each variable (cf. remark after theorem 2.6 p. 17). If the  $u$ 's such that  $\inf_{x \in C} \langle u, x \rangle > -\infty$  have non-negative coordinates for instance, one may even use propositions 2.4 p. 15 or 2.7 p. 19 to establish the existence of a value.

COMMENT 4.5. Theorem 4.5 implies that any convex set is either approachable or excludable. This is analogous to the usual minmax theorem which in effect states the same thing for sets of the form  $[a, +\infty[$  or  $-\infty, a]$  in  $\mathbb{R}^1$ .

COROLLARY 4.6. Under the hypotheses of theorem 4.5, a convex set  $C$  is approachable iff, for all  $u$ ,

$$v(u) \geq \inf_{x \in C} \langle u, x \rangle$$

Even without those hypotheses, this condition is, for convex  $C$ , equivalent to the sufficient condition of approachability (theorem 4.3 p. 102) if  $v(u)$  stands for the  $\sup \inf$  value of the game, and its negation equivalent to a sufficient condition for excludability if  $v(u)$  stands for the  $\inf \sup$  value of the game.

### Exercises.

1. (Blackwell, 1956a) Assume  $f(s, t)$  is continuous in each variable,  $S$  and  $T$  compact. Any closed set in  $\mathbb{R}^1$  is either approachable or excludable.

HINT. Let  $v$  and  $v'$  be the values of the games with pay-off functions  $f(s, t)$  and  $f(t, s)$  respectively. Prove that a closed set  $C$  is approachable iff  $[v', v] \cap C \neq \emptyset$  in case  $v' \leq v$  and  $[v, v'] \subseteq C$  in case  $v \leq v'$ , and that  $C$  is excludable otherwise: in  $\mathbb{R}$ , any closed set is either approachable or excludable.

2. (Blackwell, 1956a) In  $\mathbb{R}^2$ , there are sets neither approachable nor excludable.

HINT. Consider  $A = \begin{pmatrix} (0, 0) & (0, 0) \\ (1, 0) & (1, 1) \end{pmatrix}$  and the set  $C = \{(1/2, x_2) \mid 0 \leq x_2 \leq 1/4\} \cup \{(1, x_2) \mid 1/4 \leq x_2 \leq 1\}$ .

- (1) To show that  $C$  is not excludable, consider the strategy of I of playing row 2 for  $n$  stages and continuing for the next  $n$  stages with either row 1 or row 2 according to whether  $(1/n) \sum_1^n x_2(i)$  is smaller or larger than  $1/2$  (show that  $\bar{x}(2n) \in C$ ).
- (2) To show that  $C$  is not approachable, observe that II may play, at each stage  $i$ , column 1 or 2 according to whether  $\bar{x}_1(i)$  is larger or smaller than  $3/4$ .
- (3) Define a set to be weakly approachable (by I) if for each  $\varepsilon > 0$ , there exists an integer  $N$  and for all  $n \geq N$  a strategy  $\sigma_n$  of I such that for all  $\tau$  of II:  $P_{\sigma_n, \tau}(\delta_n \leq \varepsilon) \geq 1 - \varepsilon$  and  $E_{\sigma_n, \tau}(\delta_n) \leq \varepsilon$ . Prove that if a closed set  $D$  intersects the graph of any increasing function  $f$  from  $[0, 1]$  to  $[0, 1]$  with  $f(x) - f(y) \leq x - y$ , it is weakly approachable by I, and that otherwise the closure of its complement is weakly approachable by II.

**3.** Assume  $S$  and  $T$  are finite. Show theorem 4.3 p. 102 remains valid if the  $\phi(t)$  only have finite moments of order 1.

HINT. Use a truncation method.

**4. A strong law of large numbers for martingales.** Assume that  $f(s, t) = 0$ . Then any strategy pair approaches  $\{0\}$ . (Thus, if  $X_n$  is a martingale with  $E[(X_{n+1} - X_n)^2 | \mathcal{F}_n] \leq K$ , then  $X_n/n \rightarrow 0$  a.s.)

HINT. Let  $\Pi$  choose both  $s$  and  $t$  and apply theorem 4.3 p. 102, or ex. II.4Ex.3 if  $S$  and  $T$  are finite.

**5.** The same theorems hold if I, instead of being informed at each stage of  $x_n$ , is informed of  $f(s_n, t_n)$  — for instance he is informed of the moves of both players in the extensive form.

HINT. Let I use the same strategy, with  $f(s_n, t_n)$  replacing  $x_n$ , and use ex. II.4Ex.4.

**6.** Under the assumptions of theorem 4.3 p. 102,  $E(\sup_{n \geq N} \delta_n) \leq 4\sqrt{2K/N}$ .

HINT. Only the last inequality of the theorem is needed.

**7. Bayes strategies.** (Hannan, 1957; Blackwell, 1956b) — (cf. also Luce and Raiffa, 1957, App. A.8.6) Let  $S$  and  $T$  be finite,  $k = 1$ , and let I be informed, after every stage, of  $\Pi$ 's pure strategy choice. If  $f_n$  denotes the frequency used by  $\Pi$  in the first  $n$  stages, and if  $\phi(q)$  denotes, for any probability  $q$  on  $T$ , the quantity  $\max_s \int_t f(s, t) dq$ , define then I's regret  $r_n$  at stage  $n$  to be the difference between  $\phi(f_n)$  and his average actual pay-off up to stage  $n$ . Then, I has a strategy in the infinitely repeated game such that, for any strategy of  $\Pi$ ,  $r_n$  goes to zero, more precisely  $r_n$  is of the order of  $1/\sqrt{n}$ .

HINT. Take the actual pay-off as first coordinate of a vector pay-off, having  $\#T$  additional coordinates that count the number of times every column is used.

**8.** (Mertens, 1980) Consider a game in normal form. Let for any  $T \subseteq \mathbf{I}$ :

$$v_\alpha(T) = \{x \in \mathbb{R}^T \mid \exists \sigma \in \Delta(\prod_{i \in T} S_i): \forall \tau \in \Delta(\prod_{i \notin T} S_i), \forall i \in T \quad \int_{\prod_i S_i} F_i(s) d\sigma d\tau \geq x_i\},$$

$$v_\beta(T) = \{x \in \mathbb{R}^T \mid \forall \tau \in \Delta(\prod_{i \notin T} S_i), \exists \sigma \in \Delta(\prod_{i \in T} S_i): \forall i \in T \quad \int_{\prod_i S_i} F_i(s) d\sigma d\tau \geq x_i\}.$$

a. Show that always  $v_\alpha \subseteq v_\beta$ .

b. For any  $\lambda \in \mathbb{R}^\mathbf{I}$ ,  $\lambda \geq 0$ , let  $v_\lambda(T)$  be the value of the two-person zero-sum game having as players  $T$  and  $\complement T$  and, as pay-off function,  $\sum_{i \in T} \lambda_i F_i$ . Show that

$$v_\beta(T) = \{x \in \mathbb{R}^T \mid \forall \lambda \in \mathbb{R}^\mathbf{I}, \lambda \geq 0 \quad \sum_{i \in T} \lambda_i x_i \leq v_\lambda(T)\}.$$

c. Give an example where  $v_\alpha \neq v_\beta$ .

d. Consider the repeated game, where, after every stage, any coalition  $T$  is informed of the pay-offs accruing to its members in that stage. Denote by  $v_\alpha^n$ ,  $v_\beta^n$  the corresponding “characteristic functions”  $v_\alpha$  and  $v_\beta$  when the game is repeated  $n$  times, with the average pay-off as pay-off function — and similarly for  $v_\alpha^\infty$  and  $v_\beta^\infty$ . Show that  $v_\beta^n = v_\beta^\infty = v_\alpha^\infty \supseteq v_\alpha^n$ .

e. *Superadditivity.* (Aumann, 1961b) Deduce from II.4Ex.8d that, for any game,  $v_\alpha(T)$  and  $v_\beta(T)$  are convex and, if  $T_1 \cap T_2 = \emptyset$ , then

$$v_\alpha(T_1) \times v_\alpha(T_2) \subseteq v_\alpha(T_1 \cup T_2), \text{ and } v_\beta(T_1) \times v_\beta(T_2) \subseteq v_\beta(T_1 \cup T_2).$$



## CHAPTER III

### The Beliefs Space

In this chapter we give a formal treatment to the basic concepts of models with incomplete information, namely beliefs, types, consistency of beliefs, common knowledge, etc. Given a parameter space  $K$ , we first construct the Universal Beliefs Space  $\Omega$ . A point  $\omega$  in  $\Omega$ , which we call a state of the world consists (in addition to the value of the parameters) of the specification of the “state of mind” of each of the players (i.e. the probabilities generated by beliefs on  $K$ , beliefs on these beliefs and so on.) We study the mathematical structure and properties of  $\Omega$ , introducing the notion of beliefs subspace which is closely related to the concept of common knowledge. We next define the concept of consistent beliefs as beliefs which can be derived as conditional probabilities from some common prior distribution, given each player’s private information. We characterise the consistent states of the world (i.e. states of the world in which the players’ beliefs are consistent). An approximation theorem is then presented stating roughly that any incomplete information situation can be approximated by one in which there are finitely many potential types of each player. In the last section we discuss models and strategic equilibria of games with incomplete information, based on the structure of the beliefs space developed in this chapter.

#### 1. The universal beliefs space

When considering a situation involving a finite player set  $\mathbf{I}$  facing some uncertainty about a space  $K$  (which we refer to as the space of **states of nature**), one is naturally led to handle **infinite hierarchies of beliefs**: Adopting a Bayesian approach, each player will base his decision on some subjective beliefs (i.e. probability measure) on  $K$ . Since the outcome is determined not only by the player’s own actions but the other players’ actions as well, and those are influenced by their beliefs on  $K$ , each player must also have beliefs on other players’ beliefs on  $K$ . By the same argument he must have beliefs on other players’ beliefs on his own beliefs on  $K$ , beliefs on other players’ beliefs on his beliefs on their beliefs on  $K$ , etc. Thus it seems unavoidable to have an infinite hierarchy of beliefs for each player. These hierarchies are linked together by the fact that each belief of a player is also the subject of beliefs for the other players. The object of this section is to construct the space of these hierarchies.

##### 1.a. States of the world and types.

**THEOREM 1.1.** (1) Given a Hausdorff space  $K$  and a player set  $\mathbf{I}$  there exists a Hausdorff space  $\Theta(K)$  and a homeomorphism  $\delta_K$  from  $\Theta(K)$  to  $\Delta(K \times [\Theta(K)]^{\#\mathbf{I}-1})$  such that, letting  $\Theta^i(i \in \mathbf{I})$  denote a copy [cf. 4] of  $\Theta(K)$  and  $\delta^i: \Theta^i \rightarrow \Delta(K \times \prod_{j \neq i} \Theta^j)$  the corresponding copy [cf. 4] of  $\delta_K$ , the following property holds:

(P): Given topological spaces  $(\Sigma^i)_{i \in \mathbf{I}}$  and  $\tilde{K}$ , and continuous maps  $\sigma^i$  from  $\Sigma^i$  to  $\Delta(\tilde{K} \times \prod_{j \neq i} \Sigma^j)$  and  $f: \tilde{K} \rightarrow K$ , there exists a unique  $\mathbf{I}$ -tuple of (say

universally measurable) maps  $\Theta^i(f)$  from  $\Sigma^i$  to  $\Theta^i$  such that the following diagrams ( $i \in \mathbf{I}$ ) commute — and those  $\Theta^i(f)$  are continuous:

$$\begin{array}{ccc} \Sigma^i & \xrightarrow{\sigma^i} & \Delta(\tilde{K} \times \prod_{j \neq i} \Sigma^j) \\ \Theta^i(f) \downarrow & & f \downarrow \\ \Theta^i & \xrightarrow{\delta^i} & \Delta(K \times \prod_{j \neq i} \Theta^j) \end{array}$$

- (2) Property (P) characterises the spaces  $\Theta(K)$  and maps  $\delta_K$ : If satisfied by topological spaces  $\tilde{\Theta}^i$  and continuous maps  $\tilde{\delta}^i$  — including that the right hand arrow maps regular measures to regular measures — (even just when the  $\Sigma^i$  are assumed Hausdorff, the  $\sigma^i$  and  $f$  homeomorphisms, and uniqueness is only used within the category of continuous maps), then the  $\Theta^i(id_K): \tilde{\Theta}^i \rightarrow \Theta^i$  are canonical homeomorphisms.
- (3) Let  $\#\Theta_{-1} = 1$ , and define inductively maps  $p_n^i, q_n^i$  and spaces  $\Theta_n^i$  by the commutativity of (all  $i \in \mathbf{I}$ , and  $n \geq -1$ , but using the lower part of the diagram only for  $n \geq 0$ ):

$$\begin{array}{ccc} \Theta^i & \xrightarrow{\delta^i} & \Delta(K \times \prod_{j \neq i} \Theta^j) \\ p_{n+1}^i \downarrow & = & id_K \downarrow \\ \Theta_{n+1}^i & = & \Delta(K \times \prod_{j \neq i} \Theta_n^j) \\ q_{n+1}^i \downarrow & = & id_K \downarrow \\ \Theta_n^i & = & \Delta(K \times \prod_{j \neq i} \Theta_{n-1}^j) \end{array}$$

Then the  $(\Theta_n^i, q_n^i)$  form a projective system with limit  $\Theta^i$  and  $q_n^i \circ p_n^i = p_{n-1}^i$ . Further the maps  $p_n^i$  are onto and have continuous selections  $r_n^i$  (i.e.  $p_n^i \circ r_n^i = id_{\Theta_n^i}$ ): Defining inductively  $s_n^i$  by the commutative diagrams ( $n \geq 0$ ,  $s_0^i$  arbitrary):

$$\begin{array}{ccc} \Theta_{n+1}^i & = & \Delta(K \times \prod_{j \neq i} \Theta_n^j) \\ s_{n+1}^i \uparrow & & id_K \uparrow \\ \Theta_n^i & = & \Delta(K \times \prod_{j \neq i} \Theta_{n-1}^j) \end{array}$$

— so one gets inductively  $q_n^i \circ s_n^i = id_{\Theta_{n-1}^i}$  — one can then define  $r_n^i$  as the projective limit [i.e.  $r_{n-1}^i = r_n^i \circ s_n^i$ ].

- (4) The “copies” sub 1 are in the following sense: Given bijections  $\varphi_i: \mathbb{C}\{i\} \rightarrow \#\mathbf{I}-1$ , there exists a unique system of homeomorphisms  $h_i: \Theta(K) \rightarrow \Theta^i$  such that the

following diagrams commute:

$$\begin{array}{ccc} \Theta(K) & \xrightarrow{\delta_K} & \Delta(K \times [\Theta(K)]^{\#I-1}) \\ \downarrow h_i & & \downarrow id_K \\ \Theta^i & \xrightarrow{\delta_i} & \Delta(K \times \prod_{j \neq i} \Theta^j) \end{array}$$

with  $\psi_i(\theta_1, \dots, \theta_{\#I-1}) = (h_j(\theta_{\varphi_i(j)}))_{j \neq i}$ .

And the top line of the diagram is constructed by a projective limit as sub 3.

PROOF. Recall that the space  $\Delta(X)$  of regular probability measures with the weak\* topology is Hausdorff if  $X$  is so (1.12 p.6), and that a product of Hausdorff spaces is so. One defines thus inductively a projective system of Hausdorff spaces  $\Theta_n^i$  and continuous projections  $q_n^i$  by the lower part of the first diagram sub 3 — using 1 p.7 to show that the maps are well defined and continuous. This defines a Hausdorff limit  $\Theta^i$ , together with projections  $p_n^i$  to  $\Theta_n^i$ . Hence the diagonal maps in the upper part of that diagram associate to every  $\theta \in \Theta^i$  a projective system of regular probabilities on the spaces  $X_n^i = K \times \prod_{j \neq i} \Theta_n^j$ , hence (2 p.7) a regular probability  $\delta^i(\theta)$  on the projective limit  $X^i = K \times \prod_{j \neq i} \Theta^j$  — such that the upper part of the diagram also commutes. To show the continuity of  $\delta^i$ , observe that the topology on  $\Delta(X^i)$  is the weakest topology for which the integrals of bounded l.s.c. functions  $f$  are l.s.c.. Adding a constant, we can assume  $f \geq 0$ , and hence the limit of an increasing sequence of positive linear combinations of indicators of open sets — which can further be approximated from inside by a finite union of basic open sets. Those are of the form  $\{x \mid x_\ell \in U_\ell \ \forall \ell \leq \ell_0\}$  with  $U_\ell$  open in  $X_\ell^i$ . By continuity of the projection maps from  $X_{\ell_0}^i$  to  $X_\ell^i$ , it suffices to take  $\{x \mid x_{\ell_0} \in U_{\ell_0}\}$  as basic open sets. By the same argument, finite unions of such sets are again of this form. The result follows then from the definition of the topology on  $\Delta(X_{\ell_0}^i)$  and from the continuity of  $p_{\ell_0+1}^i$ .

To establish that  $\delta^i$  is a homeomorphism, consider the map  $g: \Delta(X^i) \rightarrow \Theta^i$ :  $g: \mu \mapsto [\mu \circ (\psi_n^i)^{-1}]_{n \geq 0}$ , where  $\psi_n^i: X^i \rightarrow X_n^i$  is the natural projection from  $X^i$  to  $X_n^i$ . The map  $g$  is well defined and continuous to  $\prod_n \Delta(X_n^i) = \prod_n \Theta_n^i$  (by 1 p.7). And commutativity of the diagram sub 3 yields that  $q_{n+1}^i([g(\mu)]_{n+1}) = [g(\mu)]_n$  — so  $g(\mu) \in \Theta^i$ . There remains thus only to show that  $\delta^i$  and  $g$  are inverse to each other:  $g \circ \delta^i$  being the identity is obvious from the definitions, while for  $\delta^i \circ g$ : let  $\theta = g(\mu)$  (i.e.  $\theta_n = \mu \circ (\psi_n^i)^{-1}$ ),  $\bar{\mu} = \delta^i(\theta)$ . Then, by definition of  $\delta^i$ ,  $\bar{\mu}$  is the only point in  $\Delta(X^i)$  with  $\bar{\mu} \circ (\psi_n^i)^{-1} = \theta_n \ \forall n$ . But  $\mu$  is another such point, so  $\mu = \bar{\mu}$ :  $\delta^i \circ g = id_{\Delta(X^i)}$ .

We have constructed our spaces  $\Theta^i$  and homeomorphisms  $\delta^i$ , such that the first part of 3 holds — and the remaining part (from “Further ...” on) is now obvious, as well as 4. So there only remains to establish property (P), and then to prove point 2.

Define thus continuous maps  $\varphi_n^j$  inductively by the commutativity of the diagrams ( $n \geq 0$ ) (using 1 p.7)

$$\begin{array}{ccc} \Sigma^i & \xrightarrow{\sigma^i} & \Delta(\tilde{K} \times \prod_{j \neq i} \Sigma^j) \\ \varphi_n^i \downarrow & & \downarrow f \\ \Theta_n^i & = & \Delta(K \times \prod_{j \neq i} \Theta_{n-1}^j) \\ & & \downarrow \varphi_{n-1}^j \end{array}$$

$q_n^i \circ \varphi_n^i = \varphi_{n-1}^i$  also follows inductively, so  $(\varphi_n^i)_{n \geq -1} = \varphi^i = \Theta^i(f): \Sigma^i \rightarrow \Theta^i$  is well defined and continuous. The required commutative diagram follows now by going to the limit in the above. Finally uniqueness of the  $\Theta^i(f)$  follows by taking the composition of the diagram with that sub 3: by induction, all  $p_n^i \circ (\Theta^i(f))$  are uniquely defined.

2 is immediate: if  $\varphi^i$  and  $\psi^i$  denote respectively  $\Theta^i(id_K)$  from  $\tilde{\Theta}^i$  to  $\Theta^i$  (using property (P) for the space  $\Theta$ ) and from  $\Theta^i$  to  $\tilde{\Theta}^i$  (using the property for  $\tilde{\Theta}$ ), then  $\varphi^i \circ \psi^i$  yields a commutative diagram

$$\begin{array}{ccc} \Theta^i & \xrightarrow{\delta^i} & \Delta(K \times \prod_{j \neq i} \Theta^j) \\ \varphi^i \circ \psi^i \downarrow & & \downarrow \text{id}_k \\ \Theta^i & \xrightarrow{\delta^i} & \Delta(K \times \prod_{j \neq i} \Theta^j) \end{array}$$

so  $\varphi^i \circ \psi^i$  is the identity, by the uniqueness part of property (P), since the identity also makes the diagram commute. Similarly  $\psi^i \circ \varphi^i$  is the identity — using now uniqueness for  $\tilde{\Theta}^i$ . So  $\Theta^i(id_K)$  is a (canonical — uniqueness) homeomorphism. ■

COMMENT 1.1. A reinterpretation of the result is the following: consider the category whose objects are systems  $S = [(\Sigma^i, \delta^i)_{i \in \mathbf{I}}, K]$  — where the  $\Sigma^i$  and  $K$  are topological spaces and  $\delta^i: \Sigma^i \rightarrow \Delta(K \times \prod_{j \neq i} \Sigma^j)$  continuous maps — and where a morphism  $\sigma: \tilde{S} \rightarrow S$  consists of continuous maps  $\sigma^i: \tilde{\Sigma}^i \rightarrow \Sigma^i$  and  $\bar{\sigma}: \tilde{K} \rightarrow K$  such that the following diagrams commute for all  $i \in \mathbf{I}$

$$\begin{array}{ccc} \tilde{\Sigma}^i & \xrightarrow{\tilde{\delta}^i} & \Delta(\tilde{K} \times \prod_{j \neq i} \tilde{\Sigma}^j) \\ \sigma^i \downarrow & & \downarrow \bar{\sigma} \\ \Sigma^i & \xrightarrow{\delta^i} & \Delta(K \times \prod_{j \neq i} \Sigma^j) \end{array}$$

which includes the requirement, if  $S$  is not Hausdorff, that the Borel measure  $\mu \circ [\bar{\sigma} \times \prod_{j \neq i} \sigma^j]^{-1}$  is regular for every  $\mu \in \Delta(\tilde{K} \times \prod_{j \neq i} \tilde{\Sigma}^j)$ . The fact that this is a category is immediate. Then  $\Theta$  is a covariant functor from the category of Hausdorff spaces and continuous maps to the category of systems. And the system  $\Theta_K$  is characterised by the fact that, for any system  $\tilde{S}$ , any continuous map  $\bar{\sigma}: \tilde{K} \rightarrow K$  extends to a unique morphism  $\sigma (= [(\Theta^i(\bar{\sigma}))_{i \in \mathbf{I}}, \bar{\sigma}])$  from  $\tilde{S}$  to  $\Theta_K$ .

COMMENT 1.2. In a similar vein, point 4 of the theorem incites to study the functor  $\Theta$  as a contravariant functor of player sets: a map  $\varphi: \mathbf{I}_1 \rightarrow \mathbf{I}_2$  defines continuous maps  $h_i: \Theta_2^{\varphi(i)} \rightarrow \Theta_1^i$  such that

$$\begin{array}{ccc} \Theta_1^i & \xrightarrow{\delta_2^i} & \Delta(K \times \prod_{j \neq i} \Theta_1^j) \\ h_i \uparrow & & \uparrow id_K \\ \Theta_2^{\varphi(i)} & \xrightarrow{(\delta_2^i, id_{\Theta_2^{\varphi(i)}})} & \Delta(K \times \prod_{j \in \mathbf{I}_2} \Theta_2^j) \end{array}$$

commutes, with  $\psi_i: (\theta_j)_{j \in \mathbf{I}_2} \mapsto [h_j(\theta_{\varphi(j)})]_{j \in \mathbf{I}_1 \setminus \{i\}}$

**THEOREM 1.2.** (1) If  $K$  is separable metric, compact, completely regular,  $K$ -analytic,  $K$ -Lusin,  $\mathcal{Z}_{c\delta}$  in a compact, analytic, Lusin, Polish, quasi-Radon,  $\tau$ -Radon, or countably Radon space, then so is  $\Theta(K)$ .

- (2) In property (P), if the spaces  $\Sigma^i$  and  $\tilde{K}$  are compact and non-empty, and the maps  $\sigma^i$  and  $f$  onto, then so is  $\Theta^i(f)$ .
- (3) (a) If  $f: K_1 \rightarrow K_2$  is one-to-one, or an inclusion (of a closed subset, of a  $\mathcal{Z}$ -subset, of a  $\mathcal{Z}_{c\delta}$ -subset), so is  $\Theta(f): \Theta(K_1) \rightarrow \Theta(K_2)$ .
- (b) If  $K_1 \subseteq K_2$ , then  $\Theta(K_1) = \bigcap_n A_n$ , with  $A_{-1} = \Theta(K_2)$ , and  $A_{n+1} = \{ \theta \in \Theta(K_2) \mid (\delta(\theta))[K_1 \times (A_n)^{\# \mathbf{I}^{-1}}] = 1 \}$  (in the sense of inner measure if necessary).
- (c) If  $K_1$  is  $K$ -analytic, and  $f: K_1 \rightarrow K_2$  is onto, so is  $\Theta(f)$ .

**COMMENT 1.3.** For the definitions, cf. App.3, App.5 and App.9.

**PROOF.** Point 1 follows from the fact that each of those categories of spaces is stable under countable products, closed subspaces, and the operation  $X \rightarrow \Delta(X)$  (cf. a.o. 3.a and 9.c, 10.a, 10.b) and from theorem 1.1.3 p. 108.

2: Assume first just that the spaces are  $K$ -analytic, and let  $\bar{\Theta}^j = [\Theta^j(f)](\Sigma^j) \subseteq \Theta^j$ : the  $\bar{\Theta}^j$  are non-empty and  $K$ -analytic,  $\delta^i$  maps  $\bar{\Theta}^i$  into  $\Delta(K \times \prod_{j \neq i} \bar{\Theta}^j) \subseteq \Delta(K \times \prod_{j \neq i} \Theta^j)$  (9.b.2), and the map is onto using 9.b.3 — hence it is a homeomorphism since  $\delta^i$  itself is one (theorem 1.1, part 1): we have reduced the problem to the case where  $f$  is the identity, the  $\sigma^i$  homeomorphisms, and the  $\Theta^i(f)$  inclusions. It follows then, inductively using 9.b.3 in the first diagram of theorem 1.1 part 3, that  $p_n^i(\bar{\Theta}^i) = \Theta_n^i$  — which implies in turn, by definition of the projective limit topology (theorem 1.1 part 3), that  $\bar{\Theta}^i$  is dense in  $\Theta^i$ . Hence, in the compact case,  $\bar{\Theta}_i$  is also compact, and being dense, equals  $\Theta^i$ .

3. Given  $f: K_1 \rightarrow K_2$ , observe from part 3 of theorem 1.1 that for inductively defined continuous maps  $\varphi_n^i$ , the following diagram commutes:

$$\begin{array}{ccc}
 \Theta^i(K_1) & \xrightarrow{\delta_1^i} & \Delta[K_1 \times \prod_{j \neq i} \Theta^j(K_1)] \\
 \downarrow p_n^i & \Theta_n^i(K_1) = \Delta[K_1 \times \prod_{j \neq i} \Theta_{n-1}^j(K_1)] & \downarrow f \\
 \Theta^i(K_2) & = \Delta[K_2 \times \prod_{j \neq i} \Theta_{n-1}^j(K_2)] & \downarrow f \\
 \downarrow p_n^i & & \downarrow \Theta^j(f) \\
 \Theta^i(K_2) & \xrightarrow{\delta_2^i} & \Delta[K_2 \times \prod_{j \neq i} \Theta^j(K_2)]
 \end{array}$$

Indeed, induction in the central part of the diagram defines uniquely the (continuous) maps  $\varphi_n^i$ . By induction again, they yield commutativity of:

$$\begin{array}{ccc}
 \Theta_n^i(K_1) & \xrightarrow{q_n^i} & \Theta_{n-1}^i(K_1) \\
 \downarrow \varphi_n^i & & \downarrow \varphi_{n-1}^i \\
 \Theta_n^i(K_2) & \xrightarrow{q_n^i} & \Theta_{n-1}^i(K_2)
 \end{array}$$

Hence they define a unique — and continuous — map  $\varphi^i: \Theta^i(K_1) \rightarrow \Theta^i(K_2)$  between the projective limits, such that  $p_n^i \circ \varphi^i = \varphi_n^i \circ p_n^i$  for all  $n$ . So, using also part 3 of theorem

1.1, we get commutativity of:

$$\begin{array}{ccc}
 \Theta^i(K_1) & \xrightarrow{\delta_1^i} & \Delta[K_1 \times \prod_{j \neq i} \Theta^j(K_1)] \\
 \downarrow p_n^i & & \downarrow \text{id}_{K_1} & \downarrow p_{n-1}^j \\
 \Theta_n^i(K_1) & = & \Delta[K_1 \times \prod_{j \neq i} \Theta_{n-1}^j(K_1)] \\
 \varphi^i \downarrow & \varphi_n^i \downarrow & f \downarrow & \varphi_{n-1}^j \downarrow \\
 \Theta_n^i(K_2) & = & \Delta[K_2 \times \prod_{j \neq i} \Theta_{n-1}^j(K_2)] \\
 \downarrow \nearrow p_n^i & & \uparrow \text{id}_{K_2} & \uparrow p_{n-1}^j \\
 \Theta^i(K_2) & \xrightarrow{\delta_2^i} & \Delta[K_2 \times \prod_{j \neq i} \Theta^j(K_2)]
 \end{array}$$

This yields the commutativity of:

$$\begin{array}{ccc}
 \Theta^i(K_1) & \xrightarrow{\delta_2^i} & \Delta(K_1 \times \prod_{j \neq i} \Theta^j(K_1)) \\
 \varphi^i \downarrow & & \downarrow f & \downarrow \varphi^j \\
 \Theta^i(K_2) & \xrightarrow{\delta_2^i} & \Delta(K_2 \times \prod_{j \neq i} \Theta^j(K_2))
 \end{array}$$

Indeed, if for some point in  $\Theta^i(K_1)$ , the two paths yielded two different measures on  $K_2 \times \prod_{j \neq i} \Theta^j(K_2)$ , the projections (by  $\text{id}_{K_2} \times \prod_{j \neq i} p_n^j$ ) on  $K_2 \times \prod_{j \neq i} \Theta_n^j(K_2)$  would already be different for some  $n$  (uniqueness of projective limits for regular probability measures), contradicting the commutativity of the previous diagram. Hence, by the uniqueness part in property (P), we have  $\varphi^i = \Theta^i(f)$ , thus establishing our full diagram.

If now  $f$  is one-to-one, an inclusion (of a closed subset, or of a  $\mathcal{L}_{c\delta}$ -subset), or onto, the same induction that defines the  $\varphi_n^i$  shows, using 9.b that the  $\varphi_n^i$  have the same property — in the onto case, using also the  $K$ -analyticity of the  $\Theta_n^i(K_1)$ , which follows e.g. from that of  $\Theta^i(K_1)$  (point 1) and the onto character of  $p_n^i$  (part 3 of theorem 1.1) by 3.a. Hence the conclusion for (3a), going to the projective limit (by our above diagrams)  $\Theta^i(f): \Theta^i(K_1) \rightarrow \Theta^i(K_2)$ .

For (3b), observe that, by (3a) and 2.b,  $\Theta_1^i \subseteq \Theta^i$  is such that  $\delta^i$  induces a homeomorphism between  $\Theta_1^i$  and  $\{\mu \in \Delta(K_2 \times \prod_{j \neq i} \Theta_2^j) \mid \mu(K_1 \times \prod_{j \neq i} \Theta_1^i) = 1\}$ . It follows then inductively that  $\Theta_1^i \subseteq A_n$ , hence  $\Theta_1^i \subseteq A_\infty = \bigcap_n A_n$ . On the other hand, it is by definition clear that  $\delta^i$  maps  $A_\infty$  onto  $\Delta(K_1 \times (A_\infty)^{\#I^{-1}})$  (hence homeomorphically). So by our diagram above, it follows inductively that  $p_n^i(A_\infty) = \Theta_n^i(K_1)$ : the image of some point in  $A_\infty$  by  $[\Delta(\text{id}_{K_2} \times \prod_{j \neq i} p_{n-1}^j)] \circ \delta_n^i$  is by the induction hypothesis some element of  $\Delta(K_1 \times \prod_{j \neq i} \Theta_{n-1}^j(K_1))$  — remember  $f$  and  $\varphi_{n-1}^j$  are inclusions — i.e. of  $\Theta_n^i(K_1)$ , (or rather its image by  $\varphi_n^i$ ). Since the  $\varphi_n^i$  are inclusions, it follows that, for  $\theta \in A_\infty$ , the sequence  $(p_n^i(\theta))_{n=-1}^\infty$  is a consistent sequence in the  $\Theta_n^i(K_1)$ , hence stems from some point in  $\Theta^i(K_1)$  — so  $A_\infty = \Theta^i(K_1)$ .

Remains only to prove (3c). We start with some preliminaries. Given continuous maps  $f_i: X_i \rightarrow Y$  ( $i = 1, 2$ ; all spaces Hausdorff), the **fibered product** of  $f_1$  and  $f_2$  is the space

$Z = \{(x_1, x_2) \in X_1 \times X_2 \mid f_1(x_1) = f_2(x_2)\}$  (together with the projections  $p_1$  and  $p_2$  to  $X_1$  and  $X_2$ ). Let also  $\Delta(f_i)$  denote the induced map between regular probabilities, etc. Then we claim:

The map  $(\Delta(p_1), \Delta(p_2))$  from  $\Delta(Z)$  to the fibered product of  $\Delta(f_1)$  and  $\Delta(f_2)$  is onto.

(The map is clearly into since  $\Delta(f \circ g) = [\Delta(f)] \circ [\Delta(g)]$ ). We first reduce the problem to the case of compact spaces. Assume thus this case solved, and fix  $\mu_i \in \Delta(X_i)$  with  $[\Delta(f_i)](\mu_i) = \nu$ . Take  $K_i$  compact in  $X_i$ , with  $\mu_i(K_i) \geq 1 - \varepsilon$ ; let  $\nu_i$  be the image by  $f_i$  of  $\mu_{i|K_i}$ , with  $h_i = d\nu_i/d\nu$ : we have  $\int h_i d\nu \geq 1 - \varepsilon$  with  $0 \leq h_i \leq 1$ , so with  $h = \min(h_1, h_2)$  we have  $\int h d\nu \geq 1 - 2\varepsilon$ . Let then  $g_i = (h/h_i) \circ f_i: X_i \rightarrow [0, 1]$ , and  $\sigma_i(B) = \int_{K_i \cap B} g_i(x) \mu_i(dx)$ : then  $\sigma_i \leq \mu_i$  has compact support, and  $[\Delta(f_i)](\sigma_i) = h d\nu$ . Continuing in the same way with the measures  $\bar{\mu}_i = \mu_i - \sigma_i$ , we see that we can write  $\mu_i = \sum_{k=1}^{\infty} \alpha^k \mu_i^k$ , where  $\mu_i^k \in \Delta(X_i)$  has compact support,  $[\Delta(f_i)](\mu_i^k) = \nu^k$ , and  $\alpha^k > 0$ . Hence if we get  $\lambda^k \in \Delta(Z)$  with  $[\Delta(p_i)](\lambda^k) = \mu_i^k$ ,  $\lambda = \sum_k \alpha^k \lambda^k$  will solve the problem. Assume thus  $X_i = \text{Supp}(\mu_i)$  and  $Y = \text{Supp}(\nu)$  are compact. We now show how to reduce the problem to the case where  $Y$  is in addition metrisable: denote by  $A$  the increasing filtering family of all separable closed sub-algebras of  $C(Y)$  that contain the constants. For  $\alpha \in A$ , denote also by  $Y_\alpha$  the corresponding (compact metric) quotient space, with  $\varphi_\alpha: Y \rightarrow Y_\alpha$  as quotient mapping, and  $Z_\alpha \subseteq X_1 \times X_2$  the fibered product of  $\varphi_\alpha \circ f_1$  and  $\varphi_\alpha \circ f_2$ . Since  $Y_\alpha$  is metrisable, there exists by assumption  $\lambda_\alpha \in \Delta(Z_\alpha)$  with  $[\Delta(p_i)](\lambda_\alpha) = \mu_i$ . Since  $(Z_\alpha)_{\alpha \in A}$  are a decreasing filtering family of compact sets, with  $Z$  as intersection, any limit point  $\lambda$  of the  $\lambda_\alpha$  will belong to  $\Delta(Z)$ , and is mapped to  $\mu_i$  by  $\Delta(p_i)$ .

So assume  $Y$  compact metric. By ex.II.1Ex.21 p.79, let  $\lambda_y^i$  be a measurable map from  $(Y, \nu)$  to  $\Delta(X_i)$  (observe  $f_i$  is onto) with the Borel sets, such that the image measure of  $\nu$  on  $\Delta(X_i)$  is regular, with  $\text{Supp}(\lambda_y^i) \subseteq f_i^{-1}(y) \quad \forall y$  and  $\mu^i(B) = \int_Y \lambda_y^i(B) \nu(dy)$  for every Borel set  $B$ .  $\lambda^1$  and  $\lambda^2$  can be constructed with the same strong lifting  $\varrho$  — hence  $\lambda^i = \varrho(\lambda^i)$  yields (ex.II.1Ex.15e p.74) that  $(\lambda^1, \lambda^2)$  is a measurable map from  $(Y, \nu)$  to  $\Delta(X_1) \times \Delta(X_2)$  with the Borel sets, such that the image measure of  $\nu$  is regular. By the continuity of the product of two measures — from  $\Delta(X_1) \times \Delta(X_2)$  to  $\Delta(X_1 \times X_2)$  (1.15 p.7), we obtain that, with  $\lambda_y = \lambda_y^1 \otimes \lambda_y^2$ ,  $\lambda$  is a measurable map from  $(Y, \nu)$  to  $\Delta(X_1 \times X_2)$  with the Borel sets, inducing a regular image measure — in particular, by ex.II.1Ex.20b p.78,  $\lambda_y(B)$  is measurable for every Borel set  $B$ , and, with  $\int \lambda_y(B) \nu(dy) = \sigma(B)$ ,  $\sigma \in \Delta(X_1 \times X_2)$ . Using  $Z$  for the Borel set  $B$ , we have  $\lambda_y(Z) = 1 \quad \forall y$ , hence  $\sigma(Z) = 1$ , so  $\sigma \in \Delta(Z)$ . Finally, with  $B = p_i^{-1}(\tilde{B})$ ,  $\tilde{B}$  Borel in  $X_i$ , we obtain  $\lambda_y(B) = \lambda_y^i(\tilde{B})$ , hence  $\sigma(B) = \int \lambda_y^i(\tilde{B}) \nu(dy) = \mu^i(\tilde{B}) : [\Delta(p_i)](\sigma) = \mu^i$ . This proves the claim.

For (3c), consider the commutative diagram (cf. above):

$$\begin{array}{ccc} \Theta_{n+1}^i(K_1) & \xrightarrow{q_n^1} & \Theta_n^i(K_1) \\ \varphi_{n+1}^i \downarrow & & \downarrow \varphi_n^i \\ \Theta_{n+1}^i(K_2) & \xrightarrow{q_n^2} & \Theta_n^i(K_2) \end{array}$$

Such a diagram corresponds by definition to a continuous map  $\psi_n = \varphi_{n+1}^i \times q_n^1$  from  $\Theta_{n+1}^i(K_1)$  to the fibered product  $Z_n$  of  $q_n^2$  and  $\varphi_n^i$ . It suffices to prove that those maps  $\psi_n$  are onto: indeed, given then  $\theta^2$  in  $\Theta^i(K_2)$ , let  $\theta_\ell^2 = p_\ell^i(\theta^2)$  so  $\theta_\ell^2 = q_\ell^2(\theta_{\ell+1}^2)$ , and assume by induction we have already  $\theta_\ell^1 \in \Theta_\ell^i(K_1)$  for  $\ell \leq n$ , with  $\varphi_\ell^i(\theta_\ell^1) = \theta_\ell^2$  and

$q_{\ell-1}^1(\theta_\ell^1) = \theta_{\ell-1}^1$ . Then  $(\theta_{n+1}^2, \theta_n^1) \in Z_n$ , hence the surjectivity of  $\psi_n$  will yield the existence of  $\theta_{n+1}^1 \in \Theta_{n+1}^i(K_1)$  with  $\varphi_{n+1}^i(\theta_{n+1}^1) = \theta_{n+1}^2$  and  $q_n^1(\theta_{n+1}^1) = \theta_n^1$ : we obtain a full sequence  $(\theta_n^1)_{n=-1}^\infty$ , i.e. a point  $\theta^1 \in \Theta^i(K_1)$  with,  $\forall n$ ,  $p_n^i(\theta^1) = \theta_n^1$ , so  $(\varphi_n^i \circ p_n^i)(\theta^1) = p_n^i(\theta^2) \quad \forall n$  — hence by our diagram in the beginning of the proof of (3),  $[\Theta^i(f)](\theta^1) = \theta^2$ .

We will prove the surjectivity of  $\psi_n$  by induction. So assume  $n \geq 0$ , and  $\psi_{n-1}$  is onto. Let  $X_n^i = K_i \times \prod_{j \neq i} \Theta_n^j(K_j)$ , also  $g_n = f \times \prod_{j \neq i} \varphi_n^j: X_n^1 \rightarrow X_n^2$ , and  $h_n^i = id_{K_i} \times \prod_{j \neq i} q_n^j: X_{n+1}^i \rightarrow X_n^i$ . In the commutative diagram:

$$\begin{array}{ccc} X_n^1 & \xrightarrow{h_{n-1}^1} & X_{n-1}^1 \\ g_n \downarrow & & \downarrow g_{n-1} \\ X_n^2 & \xrightarrow{h_{n-1}^2} & X_{n-1}^2 \end{array}$$

the surjectivity of  $\psi_{n-1}$  on each of the factors  $\Theta^j(j \neq i)$  separately yields that of the induced map  $\chi_n = g_n \times h_{n-1}^1$  from  $X_n^1$  to the fibered product  $P_n$  of  $h_{n-1}^2$  and  $g_{n-1}$ . And, by our diagram in the beginning of the proof of (3), proving the surjectivity of  $\psi_n$  amounts to prove that in the diagram:

$$\begin{array}{ccc} \Delta(X_n^1) & \xrightarrow{\Delta(h_{n-1}^1)} & \Delta(X_{n-1}^1) \\ \Delta(g_n) \downarrow & & \downarrow \Delta(g_{n-1}) \\ \Delta(X_n^2) & \xrightarrow{\Delta(h_{n-1}^2)} & \Delta(X_{n-1}^2) \end{array}$$

the induced map from  $\Delta(X_n^1)$  to the fibered product of  $\Delta(h_{n-1}^2)$  and  $\Delta(g_{n-1})$  is onto. But this is the composition of  $\Delta(\chi_n)$  with the map from  $\Delta(P_n)$  to the fibered product of  $\Delta(h_{n-1}^2)$  and  $\Delta(g_{n-1})$ . The first is onto by 9.b.3 [K-analyticity of  $X_n^1$  follows from point 1 of the present theorem, the onto character of  $p_n^i$ , theorem 1.1 part 3, and 3.a], and the second by our above claim. ■

COMMENT 1.4. Compactness is clearly necessary in point 2, e.g.,  $K$ -analyticity, to generalise at the same time point (3c), would be far insufficient. Consider for example any  $K$ -analytic set  $K$ , on which there exists a non-constant continuous function. Then there is a proper  $\mathcal{Z}$ -subset, say  $Z$ , so by (3a)  $\Theta^i(Z)$  is a  $\mathcal{Z}$ -subset of  $\Theta^i(K)$ : let  $\bar{\Theta}_0^i = \Theta^i(K) \setminus \Theta^i(Z)$ ,  $\bar{\Theta}_{n+1}^i = (\delta^i)^{-1}(\Delta(K \times \prod_{j \neq i} \bar{\Theta}_n^j))$ . Observe that, using inductively 9.b.2, the  $\bar{\Theta}_n^i$  are  $\mathcal{Z}_{c\delta}$ -subsets of  $\Theta^i(K)$ . Further  $\bar{\Theta}_1^i \subseteq \bar{\Theta}_0^i$  — must clearly be disjoint from  $\Theta^i(Z)$  —, so by induction the  $\bar{\Theta}_n^i$  form a decreasing sequence. Let  $\bar{\Theta}^i = \bigcap_n \bar{\Theta}_n^i$ : clearly  $\delta^i(\bar{\Theta}^i) = \Delta(K \times \prod_{j \neq i} \bar{\Theta}^j)$ , further the  $\bar{\Theta}^i$  are  $\mathcal{Z}_{c\delta}$ -subsets of  $\Theta^i(K)$ , strictly included in it, and non-empty e.g. because they contain (induction again)  $\Theta^i(K \setminus Z)$ .

COMMENT 1.5. The previous comment implies in particular that the homeomorphism between  $\bar{\Theta}^i$  and  $\Delta(K \times \prod_{j \neq i} \bar{\Theta}^j)$  is far from sufficient to characterise the spaces  $\Theta^i(K)$  — even when the  $\bar{\Theta}^i$  are restricted to be (nice —  $\mathcal{Z}_{c\delta}$ ) subsets of  $\Theta^i(K)$ . There is a different sense in which too it is insufficient and this even with all spaces compact: one can have, with  $K$  compact, compact spaces  $\bar{\Theta}^i$  and homeomorphisms  $\sigma_i$  from  $\bar{\Theta}^i$  to  $\Delta(K \times \prod_{j \neq i} \bar{\Theta}^j)$ , without  $\bar{\Theta}^i$  being in any sense homeomorphic to  $\Theta^i(K)$ . For instance, let  $X_1 = [0, 1]$ ,  $X_{n+1} = \Delta(X_n)$ ,  $q_1: X_2 \rightarrow X_1$  maps to the barycentre,

$q_n = \Delta(q_{n-1}) : X_{n+1} \rightarrow X_n$  for  $n > 1$ . Obviously the  $X_n$  are compact metric, and the  $q_n$  continuous and onto (induction). Let  $X$  be the projective limit (compact metric), with projections  $p_n$  onto  $X_n$ . For  $x \in X$ , the  $p_{n+1}(x)$  define a projective system of probabilities on the  $X_n$  ( $n \geq 1$ ), hence a probability  $\delta(x) \in \Delta(X)$ . The map  $\delta$  is clearly one-to-one, continuous, and onto, so a homeomorphism from  $X$  to  $\Delta(X)$  by compactness. Let then, for  $\mathbf{I} = \{1, 2\}$ ,  $\#K = 1$ ,  $\bar{\Theta}^i = X$ , with  $\delta^i$  a copy of  $\delta$ ,  $\delta^i : \bar{\Theta}^i \rightarrow \Delta(K \times \bar{\Theta}^j)$  ( $i \neq j$ ). So our  $\bar{\Theta}^i, \delta^i$  have all required properties for the one-point space  $K$ , still are in no sense homeomorphic to the one-point space  $\Theta^i(K)$ .

We prove now that property  $(P)$  still holds when the maps are only assumed measurable. Actually, our solution does not seem exactly right, since we have to restrict slightly the class of topological spaces considered: it would be more natural to investigate, for given spaces, for which class of maps the property still holds, rather than, for a given measurability concept, what class of spaces can be used.

**THEOREM 1.3.** Assume that in property  $(P)$ , the spaces  $\Sigma^i$  and  $\tilde{K}$  are quasi-Radon (10.a), and the maps  $\sigma^i$  and  $f$  universally measurable. Then there exist unique maps  $\Theta^i(f)$  making the diagram commutative in the sense that for every Borel set  $B$  in  $K \times \prod_{j \neq i} \Theta^j$  such that  $\tilde{B} = [f \times \prod_{j \neq i} \Theta^j(f)]^{-1}(B)$  is  $\sigma^i(s_i)$  measurable,  $[\sigma^i(s_i)](\tilde{B}) = [\delta^i[(\Theta^i(f))(s_i)]](B)$ . Those  $\Theta^i(f)$  are universally measurable — so the diagram is unrestrictedly commutative (even for all universally measurable  $B$ ).

**PROOF.** Is the same as in theorem 1.1, proving inductively that the  $\varphi_n^j$  are universally measurable (and uniquely defined) — using 10.a.5, 9.d.1, 10.a.4 and again 9.d.1 at every step of the induction. Then by induction, we must have  $p_n^i \circ \Theta^i(f) = \varphi_n^i$  for all  $n$ :  $\Theta^i(f)$  is thereby uniquely defined, and its universal measurability follows by writing it as the composition (9.d.1) of the inclusion of  $\Sigma^i$  as the diagonal of  $(\Sigma^i)^\infty$  and the universally measurable  $\prod_n \varphi_n^i : (\Sigma^i)^\infty \rightarrow \prod_n \Theta_n^i$  (9.d.1) — together with the remark that a universally measurable map is still universally measurable to every subspace of the image space that contains the range. Commutativity of the diagram follows now for every Borel set  $B$  of the form  $(id_K \times \prod_{j \neq i} p_n^j)^{-1}(\tilde{B})$  with  $\tilde{B}$  Borel, hence by a monotone class argument for all Borel sets (clearly any open set belongs to this monotone class, being a union of basic open sets). This extends then immediately to the universally measurable  $B$ 's. ■

**COMMENT 1.6.** In particular, if two quasi-Radon spaces are “universally measurably isomorphic”, one obtains a similar isomorphism between their corresponding systems: thus, the  $(\Theta^i(K), \delta^i)_{i \in \mathbf{I}}$  and their “universally measurable structure” (all universally measurable maps to and from all topological spaces) are uniquely determined by the universally measurable structure of  $K$ : in some sense, the solution is “purely” measure theoretic.

Also, one encounters often the variant of property  $(P)$  where the map  $f$  is replaced by a transition probability. Then

**THEOREM 1.4.** Assume in property  $(P)$  that  $K$  is quasi-Radon,  $f$  is a continuous map to  $\Delta(K)$ , and the maps  $\sigma^i$  are continuous. Then there are unique (in the same sense as in theorem 1.3) continuous maps  $\Theta^i(f)$  making the diagram commutative, where the right hand arrow is interpreted as  $\beta \circ [\Delta(f \times \prod_{j \neq i} \mu^j \circ \Theta^j(f))]$ , with the barycentric map  $\beta$  as in 10.a and  $\mu^j : \Theta^j \rightarrow \Delta(\Theta^j)$  mapping every point to the unit mass at this point.

**PROOF.** By the continuity of the product of probabilities (1.15 p. 7), and 10.a.6 (the spaces are quasi-Radon as in theorem 1.2, point 1), the proof is now as in theorem 1.1

p. 107 with the modified interpretation of the right hand arrow. Only some more care is required for the proof that  $q_n^i \circ \varphi_n^i = \varphi_{n-1}^i$  also follows inductively. ■

**COMMENT 1.7.** With  $\tilde{K} = \Delta(K)$ ,  $\Sigma^i = \Theta^i(\Delta(K))$ , and  $f$  the identity one obtains the canonical continuous map from  $\Theta^i(\Delta(K))$  to  $\Theta^i(K)$ . So in the general case, when  $f$  has values in  $\Delta(K)$  and all maps are universally measurable, one still obtains the result, assuming all spaces are quasi-Radon:  $\Theta^i(f)$  is universally measurable from  $\Sigma^i$  to  $\Theta^i(K)$ , and obtained by composing the map from  $\Sigma^i$  to  $\Theta^i(\Delta(K))$  obtained in Theorem 1.3 with the canonical map above from  $\Theta^i(\Delta(K))$  to  $\Theta^i(K)$ . Of course, a direct definition, like in theorem 1.4, is also possible.

**COMMENT 1.8.** In the line of the remark after theorem 1.1, there is a category of quasi-Radon spaces, with universally measurable transition probabilities as morphisms; indeed,  $f: X \rightarrow \Delta(Y)$  and  $g: Y \rightarrow \Delta(Z)$  compose as  $g * f = \beta \circ (\Delta(g)) \circ f$  (5.d and 5.f). Remark 1.7 shows then that  $\Theta^i$  is a (covariant) functor from this category to the category of quasi-Radon systems with universally measurable morphisms.

**COROLLARY 1.5.** (1) If  $K$  is countably Radon, or analytic, then for any other countably Radon or analytic topology on  $K$  with the same Borel  $\sigma$ -field or the same  $\mathcal{B}_u$ , one obtains another countably Radon or analytic topology on  $\Theta(K)$ , with the same map  $\delta$ , and with the same Borel  $\sigma$ -field or same  $\mathcal{B}_u$ .

(2) Assume  $K$  is countably Radon, or analytic.

Let  $(\Sigma_i, \mathcal{S}_i)$  and  $(\Omega, \mathcal{A})$  be measurable spaces, together with a transition probability  $Q$  from  $(\Omega, \mathcal{A})$  to  $K$  with the Borel sets, and transition probabilities  $\sigma_i$  from  $(\Sigma_i, \mathcal{S}_i)$  to  $(\Omega, \mathcal{A}) \times \prod_{j \neq i} (\Sigma_j, \mathcal{S}_j)$ . Then there exists a unique system of measurable maps  $\Theta^i(Q)$  from  $(\Sigma_i, \mathcal{S}_i)$  to  $\Theta^i(K)$  with the Borel sets such that the diagrams of property (P) commute.

(3) The measurable space  $\Theta^i(K)$  with the Borel sets is uniquely characterised by the above property.

**PROOF.** Apply the above, and 10.b — including remark 1.7 p. 116. ■

In view of theorem 1.1 p. 107 we refer to  $\Theta$  as the **universal type space**. It is determined by  $K$  and  $\mathbf{I}$  only. The space  $\Theta^i$  is called the **type set** of player  $i$ .

We define now the space  $\Omega$  by:  $\Omega = K \times \prod_{i \in \mathbf{I}} \Theta^i$ . This space is called the **universal beliefs space** and its elements are called **states of the world**.

The definition of  $\Omega$  and the basic property of the type space:  $\Theta^i$  homeomorphic to  $\Delta(K \times \prod_{j \neq i} \Theta^j)$ , are the formal expression of Harsanyi's approach in modelling games with incomplete information namely: A state of the world consists of a state of nature together with a list of types; one for each player. A type of a player defines a joint probability distribution on the states of nature and the types of the other players.

Note that our construction is based on some implicit assumptions on the players' beliefs: The subject of each player's beliefs are the states of nature and the beliefs of the others. He 'knows' his own beliefs. Second, the definition of the projections  $\varphi_n^i$  implies that each level of player  $i$ 's beliefs is compatible with his lower levels beliefs and he believes that this is so for all players and so on. An alternative approach is to make the player's beliefs be on the states of nature and all lower levels beliefs, including his own. If we do so then the above mentioned coherency conditions become explicit restrictions on the "admissible" beliefs. This is done in ex. IIIEx.2 p. 141.

Given  $\theta \in \Theta^i$ , let  $\bar{\theta} = \delta^i(\theta) \times \varepsilon_\theta \in \Delta(\Omega)$  — where  $\varepsilon_\theta$  is the unit mass at  $\theta$ . Recall that the product of regular probabilities is a continuous map (1.15 p. 7): the map  $\theta \mapsto \bar{\theta}$  is continuous. So is the projection map  $\theta^i: \Omega \rightarrow \Theta^i$ , and hence also the map  $\omega \mapsto \bar{\theta}_\omega^i$  from  $\Omega$  to  $\Delta(\Omega)$ .  $\theta^i(\omega)$  or  $(\bar{\theta}_\omega^i)$  is referred to as **the type of player  $i$  at  $\omega$** . Note that if  $\tilde{\omega} \in \text{Supp}(\bar{\theta}_\omega^i)$  then  $\bar{\theta}_{\tilde{\omega}}^i = \bar{\theta}_\omega^i$ . In words: at any state of the world, each player knows his own type.

**1.b. Beliefs subspaces.** The universal beliefs space  $\Omega$  consists of all possible informational configurations regarding  $K$ . Any private or “public” information about  $K$  may restrict the conceivable states of the world to a subset of  $\Omega$  containing all the beliefs of all players. This motivates the following definition.

**DEFINITION 1.6.** A **beliefs subspace** (*BL*-subspace) is a non-empty subset  $Y$  of  $\Omega$  which satisfies:

$$\forall \omega \in Y, \bar{\theta}_\omega^i(Y) = 1 \quad \forall i$$

A point of a *BL*-subspace  $Y$  can be written as  $\omega = (k; \theta^1(\omega), \dots, \theta^i(\omega), \dots)$ . The set  $\theta^i(Y)$  is the **type set** of player  $i$  in the *BL*-subspace  $Y$ . Clearly  $\Omega$  itself is a *BL*-subspace.

The following are examples of *BL*-subspaces corresponding to well known classes of games.

- (1) Games with complete information.

For fixed  $k \in K$  let  $Y = \{\omega\}$  where  $\omega = (k; \varepsilon_\omega, \dots, \varepsilon_\omega)$ , and  $\varepsilon_\omega$  is the measure on  $Y$  assigning unit mass to  $\omega$ . This is a usual situation of complete information.

- (2) Games with a random move to choose the state of nature.

For  $k_1$  and  $k_2$ , two distinct elements in  $K$ , let

$$\begin{aligned} Y &= \{\omega_1, \omega_2\} \\ \omega_1 &= (k_1; (p, 1-p), \dots, (p, 1-p)) \\ \omega_2 &= (k_2; (p, 1-p), \dots, (p, 1-p)) \end{aligned}$$

Here there is only one type of each player:  $(p, 1-p)$  which assigns probability  $p$  to  $\omega_1$  and  $(1-p)$  to  $\omega_2$ . This is a situation in which a random move chooses the state of nature  $k_1$  or  $k_2$  with probabilities  $(p, 1-p)$ .

In the following examples there are two players, I and II and the elements of  $K$  are “games” denoted by  $G^1, G^2, \dots$  or  $G^{12}, G^{22}, \dots$  etc.

- (3) Games with incomplete information on one side.

$$\begin{aligned} Y &= \{\omega_1, \dots, \omega_K\} \text{ where for } k = 1, \dots, K, \\ \omega_k &= (G^k; \varepsilon_{\omega_k}, (p^1, \dots, p^K)) \end{aligned}$$

That is, the game  $G^k \in \{G^1, \dots, G^K\}$  is chosen according to the probability vector  $p = (p^k)_{k \in K}$ . Player I is informed which game was chosen and player II is not.

- (4) Games with incomplete information “on  $1\frac{1}{2}$  sides” (cf. ex. VIEx.8 p.318).

$$\begin{aligned} Y &= \{\omega_1, \omega_2, \omega_3\} \text{ where} \\ \omega_1 &= (G^1; (1/3, 2/3, 0), (1, 0, 0)) \\ \omega_2 &= (G^1; (1/3, 2/3, 0), (0, 1/2, 1/2)) \\ \omega_3 &= (G^2; (0, 0, 1), (0, 1/2, 1/2)) \end{aligned}$$

To see the main feature of this *BL*-subspace assume that the actual state of the world is  $\omega_2$ . While player I knows that the game is  $G^1$ , he is uncertain about the “state of mind” of player II; whether he knows that it is  $G^1$  (state  $\omega_1$ ) or his beliefs are  $(1/2, 1/2)$  (state  $\omega_2$ ). Player I assigns to these two possibilities probabilities  $1/3$  and  $2/3$  respectively.

- (5) Incomplete information on two sides: the independent case.

$$\begin{aligned} Y &= \{\omega_{11}, \omega_{12}, \omega_{21}, \omega_{22}\} \text{ where} \\ \omega_{11} &= (G^{11}; (1/2, 1/2, 0, 0), (1/3, 0, 2/3, 0)) \\ \omega_{12} &= (G^{12}; (1/2, 1/2, 0, 0), (0, 1/3, 0, 2/3)) \\ \omega_{21} &= (G^{21}; (0, 0, 1/2, 1/2), (1/3, 0, 2/3, 0)) \\ \omega_{22} &= (G^{22}; (0, 0, 1/2, 1/2), (0, 1/3, 0, 2/3)) \end{aligned}$$

Player I can be of two types:  $I_1 = (1/2, 1/2, 0, 0)$  or  $I_2 = (0, 0, 1/2, 1/2)$  while player II can be either of type  $\Pi_1 = (1/3, 0, 2/3, 0)$  or  $\Pi_2 = (0, 1/3, 0, 2/3)$ . Note that each player knows his own type. An equivalent description of this informational structure is the following: Chance chooses a pair of types from  $\{I_1, I_2\} \times \{\Pi_1, \Pi_2\}$  according to the probability distribution:

$$\begin{array}{cc} \Pi_1 & \Pi_2 \\ I_1 & \left( \begin{array}{cc} 1/6 & 1/6 \\ 1/3 & 1/3 \end{array} \right) \\ I_2 & \end{array}$$

If the selected pair of types is  $(I_r, \Pi_s)$ , each player is informed of his type (only) and then they proceed to play the game  $G^{rs}$ . Note that in this representation the types of the players are chosen independently.

- (6) Incomplete information on two sides: the dependent case.

$$\begin{aligned} Y &= \{\omega_{11}, \omega_{12}, \omega_{21}, \omega_{22}\} \text{ where} \\ \omega_{11} &= (G^{11}; (1/2, 1/2, 0, 0), (1/3, 0, 2/3, 0)) \\ \omega_{12} &= (G^{12}; (1/2, 1/2, 0, 0), (0, 1/7, 0, 6/7)) \\ \omega_{21} &= (G^{21}; (0, 0, 1/4, 3/4), (1/3, 0, 2/3, 0)) \\ \omega_{22} &= (G^{22}; (0, 0, 1/4, 3/4), (0, 1/7, 0, 6/7)) \end{aligned}$$

This situation can be interpreted and represented in the same way as the previous example with the joint probability on types being:

$$\begin{array}{cc} \Pi_1 & \Pi_2 \\ I_1 & \left( \begin{array}{cc} 1/10 & 1/10 \\ 1/5 & 3/5 \end{array} \right) \\ I_2 & \end{array}$$

Thus, unlike the previous example, the types of the two players are not independent.

- (7) Incomplete information on two sides: the inconsistent case.

Consider a variant of the previous example in which the second type of player II is  $\Pi_2 = (0, 1/5, 0, 4/5)$  instead of  $(0, 1/7, 0, 6/7)$ . With this modification the players’ beliefs can no longer be viewed as conditional distributions, given each player’s type, derived from some prior joint distribution on  $\{I_1, I_2\} \times \{\Pi_1, \Pi_2\}$ .

Such situations of incomplete information are called inconsistent cases (cf. example in the next section).

The following are some general constructions of beliefs subspaces.

- (1) Any union of beliefs subspaces is one.
- (2) In particular, the union  $\Omega_f$  of all finite beliefs subspaces.
- (3) Any countable intersection of beliefs subspaces is one.
- (4) Any intersection of closed beliefs subspaces is one.
- (5) In particular,  $\forall \omega \in \Omega$  there is a smallest closed beliefs subspace containing  $\omega$ . It is obvious how to “construct” it directly (by possibly transfinite induction ...).
- (6) (a) As a cautionary remark, observe that one could also use in 4 above “separately closed”, in the sense that a subset is separately closed if the inverse image of every point in any factor  $\Theta^i(i \in \mathbf{I})$  is closed. One obtains thus also an analogous statement to 5, but associating with every  $\omega \in \Omega$  a typically smaller *BL*-subspace.  
(b) In the same vein, one could also define a concept of “topological *BL*-subspace” by requiring that  $\forall \omega \in Y, \forall i \in \mathbf{I}, \text{Supp}(\bar{\theta}_\omega^i) \subseteq Y$ . Clearly, topological *BL*-subspaces are closed under arbitrary unions and intersections. In particular, for every  $\omega \in \Omega$  there is a smallest topological *BL*-subspace containing it — which should be (typically) even smaller than the one obtained sub 6a.
- (7) Given  $\Omega_0 \subseteq \Omega$ , let  $\Omega_{n+1} = \{\omega \in \Omega_n \mid \bar{\theta}_\omega^i(\Omega_n) = 1 \ \forall i \in \mathbf{I}\}$ . Then  $\Omega_\infty = \bigcap_n \Omega_n$  is a beliefs subspace. It is the subspace where  $\Omega_0$  is “common knowledge”.  
(a)  $\Omega_\infty$  is the maximal (cf. 1) beliefs subspace contained in  $\Omega_0$ .  
(b) When  $\Omega_0$  is closed, or a  $G_\delta$ , (or Borel, or universally measurable, or has any of the properties mentioned in theorem 1.2 part 1 p. 111), then so is  $\Omega_\infty$ .  
(c) Observe that the above generalises theorem 1.2 part 3b p. 111. Cf. also ex. IIIEx.3 p. 142 for another application.
- (8) Let  $\Omega_0 = \{\omega \in \Omega \mid \omega \in \text{Supp}(\bar{\theta}_\omega^i) \ \forall i \in \mathbf{I}\}$ .  
(a) If  $\Omega$  is metrisable, then  $\Omega_0$  is a  $G_\delta$  in  $\Omega$ . Indeed, for a metric space  $E$ , the set  $\{(e, \mu) \in E \times \Delta(E) \mid d(e, \text{Supp}(\mu)) \geq \varepsilon\}$  is closed.  
(b) In particular by 6b and Theorem 1.2 part 1 p. 111, if  $K$  is separable metric, then  $\Omega_\infty$  is a  $G_\delta$ .

COMMENT 1.9. It is the beliefs subspace where it is common knowledge that no player believes a neighbourhood of the true state to be impossible. Consider an interpretation of  $\Omega$  as a model for modal logic, where (Borel) subsets of  $\Omega$  would correspond to propositions, and where  $B_i(p)$  — player  $i$  believes  $p$  — would be interpreted as  $\theta_\omega^i$  assigns probability one to the (Borel) set  $p$ . Observe that, when  $p$  is a closed set,  $B_i(p)$  is one too. Restricting the model to  $\Omega_\infty$  corresponds then to require the beliefs-operators to be knowledge operators on closed sets. (One could also restrict propositions to the Boolean algebra generated by the closed sets.) (Thus, independently of the class of subsets to which one restricts propositions, there is a variant concept of  $B_i(p)$  “behind the scenes”, which is that player  $i$  assigns probability one to some closed subset of  $p$ .)

Other examples will be given in the next section.

## 2. Consistency and common knowledge

The types of the players at a given state  $\omega$  are probability distributions (beliefs) on  $\Omega$ . A natural question is: Can these be conditional probability distributions derived from some prior probability on  $\Omega$  given the players' types?

**DEFINITION 2.1.** A probability distribution  $Q \in \Delta(\Omega)$  is **consistent** if for all  $i \in \mathbf{I}$  and every Borel set  $B$  of  $\Omega$ :

$$(1) \quad Q(B) = \int_{\Omega} \bar{\theta}_{\omega}^i(B) Q(d\omega)$$

In words:  $Q$  is consistent if it equals the average (according to  $Q$ ) of each of the beliefs  $\bar{\theta}_{\omega}^i$ . The following lemma proves that definition 2.1 indeed captures the intuitive meaning of consistency.

**LEMMA 2.2.** If  $Q \in \Delta(\Omega)$  is consistent then for all  $i \in \mathbf{I}$ , and for any Borel — or universally measurable subset  $A$  of  $\Omega$ ,

$$(2) \quad \bar{\theta}_{\omega}^i(A) = Q(A \mid \mathcal{T}^i)(\omega) \quad Q\text{-a.e.}$$

where  $\mathcal{T}^i$  is the sub  $\sigma$ -field on  $\Omega$  generated by  $\theta^i$  and the Borel — or universally measurable —  $\sigma$ -field on  $\Theta^i$ .

**PROOF.** We have to prove that for any measurable set  $A$  in  $\Omega$  and for any  $B \in \mathcal{T}^i$ :

$$\int_B \bar{\theta}_{\omega}^i(A) Q(d\omega) = \int_B \mathbb{1}_A Q(d\omega).$$

The right hand side is equal to  $Q(A \cap B)$  so that the equation will follow from (1) (applied to the measurable set  $A \cap B$ ) if we show that  $\mathbb{1}_B(\omega) \bar{\theta}_{\omega}^i(A) = \bar{\theta}_{\omega}^i(A \cap B)$  for any  $B \in \mathcal{T}^i$ . This follows from the fact that  $\bar{\theta}_{\omega}^i$  is constant on the support of  $\bar{\theta}_{\omega}^i$ , so that the full support is either in  $B$  or disjoint from  $B$ . ■

**COMMENT 2.1.** If  $K$  is completely regular, it suffices already that equation (1) be satisfied for bounded continuous functions instead of Borel sets. Indeed, since the completely regular spaces are the subspaces of compact spaces,  $\Theta^i$  and hence  $\Omega$  are then completely regular too, so the equation passes from continuous functions to indicator functions of compact sets (regularity of  $Q$  and  $\bar{\theta}_{\omega}^i$ ), and then to Borel sets (regularity of  $Q$ ).

**COMMENT 2.2.** Equation (1) remains valid with non-negative universally measurable functions instead of the Borel set  $B$ .

Denote by  $\mathcal{P}$  the set of all consistent probabilities on  $\Omega$ .

**THEOREM 2.3.** The set  $\mathcal{P}$  is closed and convex in  $\Delta(\Omega)$ .

**PROOF.** Convexity of  $\mathcal{P}$  is obvious — equation (1) is linear. Let thus  $P_{\alpha} \in \mathcal{P}$  converge to  $P \in \Delta(\Omega)$ . Let  $\mu(B) = \int \bar{\theta}_{\omega}^i(B) P(d\omega)$ . Observe that, for any compact set  $C$ ,  $\mu(C) = \inf\{\mu(\bar{O}) \mid \bar{O} \text{ closure of an open set } O \text{ containing } C\}$ . Indeed, since  $\Omega$  is  $T_2$  and  $C$  compact, the intersection of those sets  $\bar{O}$  equals  $C$  — so the  $\bar{\theta}_{\omega}^i(\bar{O})$  form a decreasing net of u.s.c. functions converging to  $\bar{\theta}_{\omega}^i(C)$  (regularity of  $\bar{\theta}_{\omega}^i$ ), hence, by the regularity of  $P$ ,  $\mu(\bar{O})$  decreases to  $\mu(C)$ .

By the upper semi-continuity of  $\bar{\theta}_{\omega}^i(\bar{O})$ ,  $\limsup_{\alpha} \int \bar{\theta}_{\omega}^i(\bar{O}) P_{\alpha}(d\omega) \leq \int \bar{\theta}_{\omega}^i(\bar{O}) P(d\omega)$ , i.e.,  $\limsup_{\alpha} P_{\alpha}(\bar{O}) \leq \mu(\bar{O})$ . But also  $\liminf P_{\alpha}(O) \geq P(O) \geq P(C)$ . Hence  $P(C) \leq \mu(\bar{O})$  for every  $\bar{O}$ , so  $P(C) \leq \mu(C)$  for every compact set  $C$ . The regularity of  $P$  implies then  $P = \mu$ . ■

COROLLARY 2.4. All topological properties of  $K$  mentioned in theorem 1.2.1 p. 111 are inherited by  $\Omega$  and by  $\mathcal{P}$ .

PROOF. By theorem 1.2.1, they are inherited by  $\Theta^i$  — and by  $\Delta(K)$ , by countable products and by closed subspaces. Therefore they are also inherited by  $\Omega = K \times \prod_i \Theta^i$ , hence by  $\Delta(\Omega)$ , hence by the closed subspace  $\mathcal{P}$  of  $\Delta(\Omega)$ . ■

In general, a  $BL$ -subspace will fail to have a consistent distribution.

EXAMPLE 2.3. Consider a situation of two players each of which has two types. The  $BL$ -subspace  $Y$  has four points corresponding to the four possible couples of types:

$$Y = \begin{pmatrix} 11 & 12 \\ 21 & 22 \end{pmatrix}$$

At each point of  $Y$  (an entry of the matrix), the first digit denotes the type of player I and the second is that of player II. Similarly we denote the subjective probabilities of the players by:

$$\text{For I: } \begin{matrix} \text{II}_1 & \text{II}_2 \\ \begin{pmatrix} p_1 & 1-p_1 \\ p_2 & 1-p_2 \end{pmatrix} \end{matrix} \quad \text{For II: } \begin{matrix} \text{II}_1 & \text{II}_2 \\ \begin{pmatrix} q_1 & q_2 \\ 1-q_1 & 1-q_2 \end{pmatrix} \end{matrix}$$

It is easily verified that there is a consistent distribution iff  $q_1(1-p_1)/[p_1(1-q_1)] = q_2(1-p_2)/[p_2(1-q_2)]$ .

### Examples on $BL$ -subspaces (continued).

- (1) If  $K$  is a separable metric space, or is analytic, so is  $\mathcal{P}$  (cor. 2.4). Hence  $\mathcal{P}$  has a countable dense sequence  $(P_n)_{n=1}^\infty$ . By theorem 2.3,  $P_\infty = \sum_1^\infty 2^{-n} P_n$  belongs to  $\mathcal{P}$ . Further  $\text{Supp}(P_n) \subseteq \text{Supp}(P_\infty)$ . Since the sequence  $P_n$  is dense, we obtain  $\text{Supp}(P) \subseteq \text{Supp}(P_\infty) \forall P \in \mathcal{P}$ : The set of supports of consistent probabilities has a maximal element ( $\text{Supp}(P_\infty)$ ), say  $S$ . Even for general  $K$ , one can always define  $S$  as  $\cup\{\text{Supp}(P) \mid P \in \mathcal{P}\}$ .
- (2) For any  $P \in \mathcal{P}$ ,  $\text{Supp}(P)$  is a (closed)  $BL$ -subspace. This can be combined with the constructions of sect. 1.b, like 2, 3, 4, or 8a — to yield other  $BL$ -subspaces.
- (3) A variant of 5 p. 119 becomes now available, assuming  $K$  compact: For every decreasing net of supports of consistent probabilities, the intersection contains a minimal such set (using compactness of  $\mathcal{P}$ ).
- (4) Given  $\Omega_0$  as in example 7 p. 119, and  $P \in \mathcal{P}$ , if  $P(\Omega_0) = 1$ , then  $P(\Omega_\infty) = 1$ .
- (5) With  $\Omega_0$  as in example 8a p. 119, and  $P \in \mathcal{P}$ , one has  $P(\Omega_0) = 1$ .
- (6) Combining all the above, let  $S_0 = S \cap \Omega_\infty$ , with  $\Omega_\infty$  as in example 8a p. 119, and  $S$  as in example 1 above. Then  $S_0$  is a canonical  $BL$ -subspace, with  $P(S_0) = 1$ ,  $\forall P \in \mathcal{P}$ , and expresses common knowledge of both consistency of beliefs and that nobody believes a neighbourhood of the true state to be impossible.
- (7) Assume  $K$  is standard Borel. So by Corollary 1.5 p. 116 the Borel spaces  $\Theta_K^i$  are uniquely determined, and are standard Borel too.
  - (a) By ex. II.1Ex.9 p. 60, there exists for every player  $i$  a Borel measurable transition probability  $q_i(dk \mid \theta)$  from  $\prod_j \Theta^j$  to  $K$ , such that for every  $\theta_i \in \Theta^i$ , and every non-negative Borel function  $f$  on  $\Omega$ ,  $\int_\Omega f(\omega) \bar{\theta}(d\omega) = \int [\int f(k, \theta_{-i}, \theta_i) q_i(dk \mid \theta_{-i}, \theta_i)] \bar{\theta}_i(d\theta_{-i})$  — i.e.  $q_i$  is the conditional probability on  $K$  given  $\theta$ , under the distribution  $\bar{\theta}_i$ . Indeed, use  $E = \Theta^i$ ,  $Y = \Theta^{-i} (\equiv \prod_{j \neq i} \Theta^j)$ ,  $X = Y \times K$ , and  $g$  the projection. And since  $g(e, x)$  is onto

$\forall e \in E$ , we can obtain in II.1Ex.9c p. 60 that  $Q(\{x \mid g(e, x) = y\} \mid e, y) = 1$   $\forall e, \forall y$ , by modifying  $Q$  on the negligible set. Therefore  $Q$  can be viewed as a transition probability from  $E \times Y$  to  $K$ .

- (b) It follows that  $q_i$  is also,  $\forall P \in \mathcal{P}$ , a version of the conditional probability on  $K$  given  $\theta \in \prod_j \Theta^j$ . In particular,  $\Omega_0 = \{\theta \mid q_i = q_j \forall i, j\}$  is a Borel set with  $P(\Omega_0) = 1 \forall P \in \mathcal{P}$ .
- (c) Thus, the corresponding beliefs space  $\Omega_\infty$  is a Borel subset of  $\prod_i \Theta^i$ , with  $P(\Omega_\infty) = 1 \forall P \in \mathcal{P}$ , and there is a Borel measurable transition probability  $q$  from  $\Omega_\infty$  to  $K$  which is a version of the conditional probability on  $K$  given  $\prod_i \Theta^i$  under every  $P \in \mathcal{P}$ , and under every  $\bar{\theta}_\omega^i, \forall i \in \mathbf{I}$  and  $\forall \omega \in \Omega_\infty$ .
- (d) In the two player case, one could do a bit better: Modify  $q^1$  by setting it equal to  $q^2$  on  $\Theta_0^1 \times \Theta^2$ , where  $\Theta_0^1 = \{\theta^1 \mid \bar{\theta}^1(\Omega_0) = 1\}$ . Do then the same with  $q^2$  and  $\Theta_0^2$ . (Instead of iterating the procedure, into the transfinite?), one can now choose for  $\Omega_0$  the set  $\Theta_0^1 \times \Theta_0^2$ , and obtain then a similar product set  $\Theta_\infty^1 \times \Theta_\infty^2$  for  $\Omega_\infty$ . Using remark 1.2 p. 110, one can now symmetrise the situation under permutations of the players, and obtain a single Borel set  $\Theta_\infty$ , which has probability one under any  $\theta \in \Theta_\infty$ , such that  $\Theta_\infty \times \Theta_\infty$  has probability one under every  $P \in \mathcal{P}$ , and a single Borel transition  $q(dk \mid \theta^1, \theta^2)$  on  $(\Theta_\infty \times \Theta) \cup (\Theta \times \Theta_\infty)$ , which is symmetric in  $(\theta^1, \theta^2)$ , and is the conditional probability under any  $P \in \mathcal{P}$  and under  $\theta \in \Theta_\infty$ . And when  $K$  is finite, one could further symmetrise under all permutations of  $K$ .

COMMENT 2.4. Are there any canonical results in this vein? E.g., does the union of such beliefs spaces cover the set  $S$ ? Is there a canonical such beliefs space — even replacing Borel-measurability by universal-measurability, or by  $\mu$ -measurability for every consistent prior and for each belief in the beliefs space?

COMMENT 2.5. The above is frequently used, in the context of sect. 4 below. In that case, to every  $k \in K$  there corresponds a game (pay-off functions)  $g^k$ . Averaging them under such a common conditional  $q$  yields a model where one has a single game for every  $\mathbf{I}$ -tuple of types. If further the (consistent) joint distribution is absolutely continuous on  $\prod_{i \in \mathbf{I}} \Theta^i$  w.r.t. some product measure — then it is a.c. w.r.t. the product of its marginals —, one can, using  $f(\theta)$  for its density with respect to the product of the marginals, replace every every game  $g^\theta$  by  $f(\theta) \cdot g^\theta$ . In this way, one obtains a model (information scheme as below) where furthermore all players' private signals are independent. It would seem worthwhile to check what such transformations yield for the corresponding canonical consistent distributions (theorem 2.5 part 2 below).

**Information Schemes.** In games with incomplete information, the information structure is defined by an **information scheme**,  $\mathcal{J}$ , that consists of a probability space  $\mathbf{E} = (E, \mathcal{E}, Q)$  and sub  $\sigma$ -fields  $\mathcal{E}^i$  (for each  $i$  in  $\mathbf{I}$ ), with a measurable mapping  $\mathbf{k}_E$  from  $E$  to the space  $K$  of states of nature. The sub  $\sigma$ -field  $\mathcal{E}^i$  describes the information of player  $i$ . The following theorem states that any information scheme has a canonical representation as a consistent probability measure on  $\Omega$ . Denoting by  $\mathbf{k}$  the canonical projection from  $\Omega$  to  $K$  we have:

THEOREM 2.5. Assume an information scheme  $\mathcal{J}$  such that  $\mathbf{k}_E$  is Borel-measurable and  $Q$  has a regular image measure  $\mu$ .

- (1) Then there exists a  $Q$ -measurable map  $\varphi$  from  $(E, \mathcal{E})$  to  $\Omega$  with the Borel sets such that:

- (a)  $(\mathbf{k} \circ \varphi)^{-1}(B) = (\mathbf{k}_E)^{-1}(B)$   $Q$ -a.e., for every Borel subset  $B$  of  $K$ .
  - (b)  $Q \circ \varphi^{-1} \in \mathcal{P}$  [or just:  $\in \Delta(\Omega)$ ].
  - (c)  $\bar{\theta}^i(\varphi(e))(B) = (Q[\varphi^{-1}(B) \mid \mathcal{E}_i])(e)$   $Q$ -a.e., for every Borel subset  $B$  of  $\Omega$  and all  $i \in \mathbf{I}$ .
- (2) If  $\tilde{\varphi}$  is another such map, then  $\varphi^{-1}(B) = \tilde{\varphi}^{-1}(B)$   $Q$ -a.e., for every Borel set  $B$  of  $\Omega$ . (In particular,  $Q \circ \varphi^{-1}$  is uniquely defined).
- (3) Assume there exists  $K_0 \subseteq K$  with  $\mu(K_0) = 1$  such that points of  $K_0$  are separated by some sequence of  $\mu$ -measurable sets. Then:
- (a) Point 2 implies that  $\tilde{\varphi} = \varphi$   $Q$ -a.e.
  - (b) In (1a) one can require  $\mathbf{k} \circ \varphi = \mathbf{k}_E$ .
  - (c) One can further require that  $\bar{\theta}^i(\varphi(e))$  is  $\mathcal{E}^i$ -measurable.

PROOF. During the proof of (1) and (2) we will assume  $\mathcal{E}^i$  contains all negligible subsets of  $(E, \mathcal{E}, Q)$ . We start with the following:

LEMMA 2.6. Let  $\varphi$  and  $\tilde{\varphi}$  two random variables on a probability space  $(\Omega, \mathcal{B}, P)$  with values in a measurable space  $(S, \mathcal{S})$ .

- (1) Assume  $S$  is completely regular and  $\mathcal{S}$  the Borel  $\sigma$ -field, the distributions of  $\varphi$  and  $\tilde{\varphi}$  are regular, or just  $\tau$ -smooth (cf. 10.a) and that for every real bounded continuous function  $f$  on  $S$ ,  $f \circ \varphi = f \circ \tilde{\varphi}$  a.e. Then the same holds true for any real Borel function  $f$ .
- (2) Assume  $\mathcal{S}$  separable and separating and that for every  $T \in \mathcal{S}$ ,  $\varphi^{-1}(T) = \tilde{\varphi}^{-1}(T)$  a.e., then  $\varphi = \tilde{\varphi}$  a.e. If  $\mathcal{A} \subseteq \mathcal{B}$  is such that every element of  $\mathcal{B}$  is a.e. equal to some element of  $\mathcal{A}$  then there exists  $\bar{\varphi}$  which is  $\mathcal{A}$ -measurable and is a.e. equal to  $\varphi$ .

PROOF. 1. Going to the Stone-Čech compactification, one can assume the image space is compact. One obtains then first, by a monotone class argument, that  $f \circ \varphi = f \circ \tilde{\varphi}$  a.e. for every Baire function  $f$ . So the distributions coincide on the Baire  $\sigma$ -field, hence (regularity) on all compact sets, — since those have a basis of open Baire neighbourhoods, hence (regularity) on all Borel sets. Let  $P$  be this distribution. For  $K$  compact, let  $O_n$  be, as just said, an open Baire set with  $K \subseteq O_n$ ,  $P(O_n \setminus K) \leq n^{-1}$ ; let  $\hat{K} = \bigcap_n O_n$ :  $\hat{K}$  is a Baire set with  $K \subseteq \hat{K}$ ,  $P(\hat{K} \setminus K) = 0$ . Hence  $\varphi^{-1}(K) =_{a.e.} \varphi^{-1}(\hat{K}) =_{a.e.} \tilde{\varphi}^{-1}(\hat{K}) =_{a.e.} \tilde{\varphi}^{-1}(K)$ . so we still have  $\varphi^{-1}(C) = \tilde{\varphi}^{-1}(C)$  a.e., for every  $K_\sigma$ -set  $C$ . The same conclusion follows now in the same way for Borel sets, approximating them by regularity from inside by  $K_\sigma$ -sets. Hence the result.

2. Let  $T_k$  be a sequence of generators of  $\mathcal{S}$ , and neglect the points where, for some  $k$ ,  $\varphi^{-1}(T_k) \neq \tilde{\varphi}^{-1}(T_k)$ . Now  $(\varphi, \tilde{\varphi})^{-1}(R_k) = \phi$  with  $R_k = [T_k \times \mathbb{C} T_k] \cup [\mathbb{C} T_k \times T_k]$ . Since  $\bigcup_k R_k$  is the complement of the diagonal in  $S \times S$ , the first assertion follows. For the second, reduce to the case where  $(S, \mathcal{S})$  is a subset of  $[0, 1]$  with the Borel subsets and approximate  $\varphi$  uniformly by a sequence of step functions  $\varphi_n$  and let  $\bar{\varphi} = \liminf \varphi_n$ . ■

For the theorem, we first reduce the proof of (1) and (2) to the case where  $K$  is compact. Find a sequence  $C_n$  of disjoint compact subsets of  $K$  with  $\mu(\bigcup_n C_n) = 1$ . Denote by  $\bar{K}$  the (locally compact) disjoint union of the  $C_n$ , and by  $\hat{K}$  its one-point compactification. For point (2), observe first that  $\varphi$  (and  $\tilde{\varphi}$ ) can be modified on a null set such as to have values in  $\Omega_{\bigcup_n C_n} \subseteq \Omega_K$ , by theorem 1.2 part 3b p. 111. Indeed, assume  $\varphi$  has already been modified such as to have values in  $[(\bigcup_k C_k) \times (A_n)]^\mathbf{I}$ , then, taking for  $B$  the latter subset in 1c, we get that  $\bar{\theta}_{\varphi(e)}^i(B) = 1$  a.e. — hence the induction step. Next, since  $\bar{K}$  is a  $\mathcal{Z}_c$ -set

in  $\hat{K}$ ,  $\Omega_{\bar{K}}$  is a  $\mathcal{X}_{c\delta}$  set (theorem 1.2 part 3a p. 111 in the compact space (theorem 1.2 part 1)  $\Omega_{\hat{K}}$ , hence  $K$ -Lusin (3.d, 3.e). Hence the map from  $\Omega_{\bar{K}}$  to  $\Omega_{\bigcup_n C_n}$  is continuous, one-to-one and onto (theorem 1.2 parts 3a and 3c p. 111). Therefore  $\varphi$  and  $\tilde{\varphi}$  still have the same properties when viewed as maps to  $\Omega_{\bar{K}}$  by (9.b.4 and 9.b.3). And clearly  $\mathbf{k}_E$  can be viewed as  $\bar{K}$ -valued. Finally, theorem 1.2 part 3a allows to assume that the values are in  $\Omega_{\hat{K}}$  and in  $\hat{K}$ .

For point 1 of theorem 2.5, our assumptions are clearly still satisfied with  $\hat{K}$  instead of  $K$ . Obtaining thus an appropriate  $\Omega_{\hat{K}}$ -valued  $\varphi$ , and letting  $B_0 = \{\omega \in \Omega_{\hat{K}} \mid \mathbf{k}_\omega \in \bar{K}\}$ ,  $B_{n+1} = \{\omega \in B_n \mid \bar{\theta}_\omega^i(B_n) = 1 \forall i\}$ , then, by theorem 2.5 parts 1a and 1c and inductively, we get  $Q(\varphi^{-1}(B_n)) = 1 \forall n$ , hence, after modifying  $\varphi$  on a null set,  $\varphi: E \rightarrow \Omega_{\bar{K}} = \bigcap_n B_n$  (theorem 1.2 part 3b). So, by theorem 1.2 part 3a, composing  $\varphi$  with the continuous one-to-one map from  $\Omega_{\bar{K}}$  to  $\Omega_K$  will yield the desired result. Hence, for parts 1 and 2 of the theorem we can assume  $K$  compact.

We now start the proof of 1. Let  $\varrho$  be a lifting on  $(E, \mathcal{E}, Q)$ ; and let  $\bar{\mathbf{k}} = \varrho(\mathbf{k}_E)$  (ex. II.1Ex.15e p. 74). Then, by the above lemma, we have  $\mathbf{k}_E^{-1}(B) = \bar{\mathbf{k}}^{-1}(B)$   $Q$ -a.e., for every Borel set  $B$ . So it suffices to prove 1 with  $\bar{\mathbf{k}}$  instead of  $\mathbf{k}_E$  — i.e., we can assume that  $\mathbf{k}_E = \varrho(\mathbf{k}_E)$ . Assume, for  $n \geq -1$ , that  $\varphi_n = (\mathbf{k}_E, (t_n^i)_{i \in \mathbf{I}}): E \rightarrow \Omega_n = K \times \prod_i \Theta_n^i$  is well-defined, and  $Q$ -measurable, such that  $\varphi_n = \varrho(\varphi_n)$  — hence  $\varphi_n$  is measurable to the Borel  $\sigma$ -field, with  $Q \circ \varphi_n^{-1} \in \Delta(\Omega_n)$ , by (ex. II.1Ex.15eiii p. 75) —, and such that  $t_n^i(e)(B) = \mathbb{E}[(\varphi_{n-1}^i)^{-1}(B) \mid \mathcal{E}_i](e)$   $Q$ -a.e., for every Borel subset  $B$  of  $X_{n-1}^i = K \times \prod_{j \neq i} \Theta_{n-1}^j$ , — where  $\varphi_n^i = \text{Proj}_{X_n^i} \circ \varphi_n$ . Assume finally that  $\varphi_{n-1} = \text{Proj}_{\Omega_{n-1}} \circ \varphi_n$ , where  $\text{Proj}_{\Omega_{n-1}} = \text{id}_K \times \prod_i q_n^i$ . Map then  $Q$  to its image on  $(E, \mathcal{E}_i) \times X_n^i$  and apply II.1Ex.16c to this image measure, yielding an  $\mathcal{E}_i$ -measurable map  $t_{n+1}^i$  to  $\Theta_{n+1}^i = \Delta(X_n^i)$ , such that  $t_{n+1}^i = \varrho(t_{n+1}^i)$ . Then  $\varphi_{n+1} = (\mathbf{k}_E, (t_{n+1}^i)_{i \in \mathbf{I}}): E \rightarrow \Omega_{n+1}$  satisfies  $\varphi_{n+1} = \varrho(\varphi_{n+1})$  by ex. II.1Ex.15eiv p. 75. By definition  $t_{n+1}^i(e)(B) = \mathbb{E}[(\varphi_n^i)^{-1}(B) \mid \mathcal{E}_i](e)$   $Q$ -a.e. for every Borel set  $B$  of  $X_n^i$ . In particular, for  $B$  Borel in  $X_{n-1}^i$ , and  $h_n^i: X_n^i \rightarrow X_{n-1}^i$  the projection, we have

$$\begin{aligned} [q_{n+1}^i(t_{n+1}^i(e))](B) &= t_{n+1}^i(e)[(h_n^i)^{-1}(B)] \\ &= \mathbb{E}[(h_n^i \circ \varphi_n^i)^{-1}(B) \mid \mathcal{E}_i](e) \\ &= \mathbb{E}[(\varphi_{n-1}^i)^{-1}(B) \mid \mathcal{E}_i](e) = t_n^i(e)[B] \quad Q\text{-a.e.} \end{aligned}$$

So, as maps into  $\Delta(X_{n-1}^i)$ , for every continuous linear functional  $\psi$ , we have  $\psi \circ q_{n+1}^i \circ t_{n+1}^i = \psi \circ t_n^i$   $Q$ -a.e. Hence, by the Stone-Weierstrass theorem, the same holds true for any continuous function  $\psi$  on  $\Delta(X_{n-1}^i)$ . Therefore, by ex. II.1Ex.15ei and II.1Ex.15eii p. 75  $\varrho[q_{n+1}^i \circ t_{n+1}^i] = \varrho(t_n^i) = t_n^i$ . But by ex. II.1Ex.15eiv,  $\varrho(q_{n+1}^i \circ t_{n+1}^i) = q_{n+1}^i \circ \varrho(t_{n+1}^i) = q_{n+1}^i \circ t_{n+1}^i$ . So  $t_n^i = q_{n+1}^i \circ t_{n+1}^i$ , and hence  $\varphi_n = \text{Proj}_{\Omega_n} \circ \varphi_{n+1}$ . This finishes the induction step. So we have  $\forall i$ ,  $t^i(e) = (t_n^i(e))_{n=1}^\infty$  an  $\mathcal{E}^i$ -measurable map to  $(\Theta^i, \text{Borel sets})$  — by the relation  $t_{n-1}^i = q_n^i \circ t_n^i$  and because the Borel  $\sigma$ -field of  $\Theta^i$  is generated by the projections  $p_n^i$  (since every closed set  $F$  equals  $\bigcap_n p_n^i(F)$  by definition of the product topology). Also  $t^i = \varrho(t^i)$  by ex. II.1Ex.15eiv p. 75, and so similarly  $\varphi = (\mathbf{k}_E, (t^i)_{i \in \mathbf{I}}) = \varrho(\varphi): E \rightarrow \Omega$ . In particular, those maps are Borel measurable with  $Q \circ \varphi^{-1} \in \Delta(\Omega)$ . Further, for every Borel set  $B$  in  $X^i$ , with  $\psi_n^i$  the projection from  $X^i$  to  $X_n^i$ , we have  $t^i(e)[(\psi_n^i)^{-1}(B)] = \mathbb{E}[(\varphi^i)^{-1}[(\psi_n^i)^{-1}(B)] \mid \mathcal{E}_i](e)$   $Q$ -a.e. By our above remark concerning the Borel  $\sigma$ -field on a projective limit, we obtain thus by a monotone class argument that [since  $t^i(e) = \delta^i(\theta_{\varphi(e)}^i)$ ], for every Borel set  $B$  in  $X^i$ ,

$$\delta^i(\theta_{\varphi(e)}^i)[B] = \mathbb{E}[(\varphi^i)^{-1}(B) \mid \mathcal{E}_i](e) \quad Q\text{-a.e.}$$

i.e.,  $\bar{\theta}_{\varphi(e)}^i(B) = \mathbb{E}[\varphi^{-1}(B) \mid \mathcal{E}_i](e)$   $Q$ -a.e., for every set  $B = \Theta^i \times \tilde{B}$ , with  $\tilde{B}$  Borel in  $X^i$ . Or, since  $\theta_{\varphi(e)}^i$  is  $\mathcal{E}_i$ -measurable, we get, for any Borel sets  $B_1$  in  $\Theta^i$  and  $B_2$  in  $X^i$  that

$$\begin{aligned}\bar{\theta}_{\varphi(e)}^i(B_1 \times B_2) &= \mathbb{1}_{\theta_{\varphi(e)}^i \in B_1} \cdot \bar{\theta}_{\varphi(e)}^i(\Theta^i \times B_2) = \mathbb{1}_{\theta_{\varphi(e)}^i \in B_1} \mathbb{E}[\varphi^{-1}(\Theta^i \times B_2) \mid \mathcal{E}_i](e) \\ &= \mathbb{E}[\mathbb{1}_{\theta_{\varphi(e)}^i \in B_1} \cdot \mathbb{1}_{\varphi(e) \in \Theta^i \times B_2} \mid \mathcal{E}_i](e) = \mathbb{E}[\varphi^{-1}(B_1 \times B_2) \mid \mathcal{E}_i](e) \quad Q\text{-a.e.}\end{aligned}$$

By additivity, the same formula remains then true for finite disjoint unions of sets  $B_1 \times B_2$ , and then by a monotone class argument we still get

$$\bar{\theta}_{\varphi(e)}^i(B) = \mathbb{E}[\varphi^{-1}(B) \mid \mathcal{E}_i](e)$$

for every set  $B$  in the product of the Borel  $\sigma$ -fields on  $\Theta^i$  and on  $X^i$ . Now,  $\bar{\theta}_{\varphi(e)}^i$ , as the composition of the continuous map  $\omega \mapsto \bar{\theta}_\omega^i$  (cf. supra) and of the map  $\varphi$  satisfying  $\varphi = \varrho(\varphi)$ , is a map from  $(E, \mathcal{E}_i)$  to  $\Delta(\Omega)$  with the Borel sets that satisfies  $\bar{\theta}_{\varphi(e)}^i = \varrho[\bar{\theta}_{\varphi(\cdot)}^i](e)$  and has a regular image measure (ex. II.1Ex.15eiii and II.1Ex.15eiv p. 75). This one has a barycentre  $\bar{P} \in \Delta(\Omega)$  since, by 10.a.2, the compact space  $\Omega$  is quasi-Radon. I.e., for every Borel set  $B$  in  $\Omega$ ,  $\bar{\theta}_{\varphi(e)}^i(B)$  is measurable and  $\int \bar{\theta}_{\varphi(e)}^i(B)Q(de) = \bar{P}(B)$ . Also  $Q \circ \varphi^{-1} = P \in \Delta(\Omega)$ , i.e., for every Borel set  $B$ ,  $P(B) = \int \mathbb{E}(\varphi^{-1}(B) \mid \mathcal{E}_i)dQ$ . Similarly, for  $A \in \mathcal{E}$ , let  $\bar{P}_A(B) = \int_A \bar{\theta}_{\varphi(e)}^i(B)Q(de)$  and  $P_A(B) = \int_A \mathbb{E}(\varphi^{-1}(B) \mid \mathcal{E}_i)dQ$ : since  $\bar{P}_A \leq \bar{P}$  and  $P_A \leq P$ , both are regular measures, which coincide by our previous formula on the product of the Borel  $\sigma$ -fields on  $\Theta^i$  and on  $X^i$ . By regularity, this implies that they coincide on all compact sets — since those have a basis of open neighbourhoods belonging to the product  $\sigma$ -field —, and hence on all Borel sets. Thus, for every Borel set  $B$  in  $\Omega$ , we have that

$$\int_A \bar{\theta}_{\varphi(e)}^i(B)Q(de) = \int_A [\mathbb{E}(\varphi^{-1}(B) \mid \mathcal{E}_i)]dQ \quad \forall A \in \mathcal{E}.$$

This means the two measurable functions are equal a.e.: hence 1c p. 123. By 1c we have, for  $B$  Borel,  $P(B) = Q(\varphi^{-1}(B)) = \int \mathbb{E}(\varphi^{-1}(B) \mid \mathcal{E}_i)dQ = \int \bar{\theta}_{\varphi(e)}^i(B)Q(de) = \int \bar{\theta}_\omega^i(B)P(dw)$ . Since also  $P \in \Delta(\Omega)$ , this yields 1b: point 1 is fully proved.

For point 2, denote by  $p_n$  the projection from  $\Omega$  to  $\Omega_n$ . Assume that  $\varphi^{-1}[p_n^{-1}(B)] = \tilde{\varphi}^{-1}[p_n^{-1}(B)]$   $Q$ -a.e., for every Borel set  $B$  in  $\Omega_n$ . [By 1a, this inductive assumption holds for  $n = -1$ ]. Then by 1c we obtain that  $\bar{\theta}_{\varphi(e)}^i(p_n^{-1}(B)) = \bar{\theta}_{\tilde{\varphi}(e)}^i(p_n^{-1}(B))$   $Q$ -a.e., for every such  $B$ . Let thus  $P_e = \bar{\theta}_{\varphi(e)}^i \circ p_n^{-1}$ ,  $\tilde{P}_e = \bar{\theta}_{\tilde{\varphi}(e)}^i \circ p_n^{-1}$ , both  $\in \Delta(\Omega_n)$ : then  $P_e(B) = \tilde{P}_e(B)$   $Q$ -a.e.,  $\forall B$  Borel. Hence for every continuous linear functional  $\psi$  on  $\Delta(\Omega_n)$  we have  $\psi \circ P_e = \psi \circ \tilde{P}_e$   $Q$ -a.e., so, by Stone-Weierstrass, this remains true for every continuous function  $\psi$  on  $\Delta(\Omega_n)$ . Take now a continuous function  $F$  on  $\Omega_{n+1} = K \times \prod_i \Theta_{n+1}^i$  — it is, by Stone-Weierstrass again, approximated by linear combinations of functions  $F = f_0 \times \prod_i \psi_i$ , where  $f_0 \in C(K)$  and  $\psi_i$  is continuous on  $\Theta_{n+1}^i = \Delta(X_n^i)$ . Hence  $F \circ p_{n+1} \circ \varphi = f_0(\mathbf{k}_E \circ \varphi) \times \prod_i \psi_i(\mathbf{Proj}_{X_n^i}(P^i))$  — here the argument of  $\psi_i$  equals  $\mathbf{Proj}_{X_n^i}(P^i) = \mathbf{Proj}_{X_n^i}(\bar{\theta}_{\varphi(e)}^i \circ p_n^{-1}) = \theta_{n+1}^i(e)$ . Since  $\psi_i \circ \mathbf{Proj}_{X_n^i}$  is continuous on  $\Delta(\Omega_n)$ , we obtain equality a.e. when replacing  $\varphi$  by  $\tilde{\varphi}$ . So  $F \circ p_{n+1} \circ \varphi = F \circ p_{n+1} \circ \tilde{\varphi}$   $Q$ -a.e., for every  $F \in C(\Omega_{n+1})$ . Therefore, by the lemma,  $\varphi^{-1}[p_{n+1}^{-1}(B)] = \tilde{\varphi}^{-1}(p_{n+1}^{-1}(B))$   $Q$ -a.e., for every Borel set  $B$  in  $\Omega_{n+1}$ . This completes the induction. It follows that,  $\forall n, \forall F \in C(\Omega_n)$ ,  $(F \circ p_n) \circ \varphi = (F \circ p_n) \circ \tilde{\varphi}$   $Q$ -a.e. Since the functions  $F \circ p_n$  are dense in  $C(\Omega)$  (Stone-Weierstrass), the lemma yields us finally that  $\varphi^{-1}(B) = \tilde{\varphi}^{-1}(B)$   $Q$ -a.e., for every Borel set  $B$  in  $\Omega$ . This proves point 2.

As for 3, we first prove the following claim:

*Under the assumptions of 3, there exists a sequence of disjoint compact metric subsets  $K_n$  of  $K$  such that  $\mu(\bigcup_n K_n) = 1$ .*

Let  $C_k$  be a sequence of disjoint compact subsets of  $K_0$ , with  $\mu(\bigcup_k C_k) = 1$  — hence we can assume  $K_0 = \bigcup_k C_k$ . Let also  $M_n$  be the sequence of  $\mu$ -measurable sets. Then  $M_n \cap C_k$  differs from some Baire subset  $B_{n,k}$  of  $C_k$  by a negligible set  $N_{n,k}$ . Let  $f_{k,i}$  be a sequence of continuous functions on  $C_k$  that generates all  $B_{n,k}$  ( $n = 1, 2, \dots$ ). Let also  $C_{k,j}$  be a sequence of disjoint compact subsets of  $C_k \setminus (\bigcup_n N_{n,k})$  with  $\mu(\bigcup_j C_{k,j}) = \mu(C_k)$ . Then the points of  $C_{k,j}$  are separated by the sequence of continuous functions  $f_{k,i}$  ( $i = 1, 2, \dots$ ), so  $C_{k,j}$  is metrisable, and  $\mu(\bigcup_{k,j} C_{k,j}) = 1$ .

So let,  $L = \bigcup_n K_n$ . Then  $L$  is Lusin (e.g., 5.f), so  $\Omega_L \subseteq \Omega_K$  is Lusin too (points 1 and 3 of theorem 1.2 p. 111). Further, as observed in the beginning of the proof of the theorem,  $\varphi$  (and  $\tilde{\varphi}$ ) can be modified on a null set such as to have values in  $\Omega_L$  — and similarly  $\mathbf{k}_E$  can be assumed to have values in  $L$ . The result follows then from point 2 of lemma 2.6 p. 123 (separability is by definition — App.5). This finishes the proof of theorem 2.5 ■

**COMMENT 2.6.** To facilitate the interpretation of 1a and 2 of theorem 2.5 p. 122, if  $\varphi$  and  $\tilde{\varphi}$  are two Borel measurable functions from a probability space to a Hausdorff space  $X$ , which have regular image measures and are such that, for every Borel set  $B$ ,  $\varphi^{-1}(B) = \tilde{\varphi}^{-1}(B)$  a.e., then the distribution of  $(\varphi, \tilde{\varphi})$  on  $X \times X$  endowed with the product of the Borel  $\sigma$ -fields is the restriction to this  $\sigma$ -field of some regular measure carried by the diagonal.

**COMMENT 2.7.** If there exists a strong lifting for  $(K, \mu)$  (ex. II.1Ex.15d p. 74), one can replace 1a p. 123 by  $\mathbf{k} \circ \varphi = \mathbf{k}_E$ . Indeed, composing this lifting with the map  $\mathbf{k}_E$  yields a lifting on a sub  $\sigma$ -field of  $(E, \mathcal{E})$ , which can then be extended to a lifting  $\varrho$  for  $(E, \mathcal{E}, Q)$  [cf. (proof of) II.1Ex.15a p. 74]. For this lifting, one will have  $\varrho(\mathbf{k}_E) = \mathbf{k}_E$  a.e. — they differ only on the negligible set where  $\mathbf{k}_E$  takes values outside the support of  $\mu$ . Hence one obtains, at the end of the construction, that  $\mathbf{k} \circ \varphi = \mathbf{k}_E$   $Q$ -a.e.. Changing the  $K$ -coordinate of the map  $\varphi$  on the exceptional set yields then the result. But the condition in (II.1Ex.15d p. 74) for the existence of a strong lifting is (cf. above claim) equivalent to the condition of theorem 2.5.3 p. 123 — where our direct argument yields more.

**COMMENT 2.8.** One can always, at the end of the construction, change the  $K$ -coordinate of  $\varphi$  and set it equal to  $\mathbf{k}_E$  — obtaining thus a map  $\bar{\varphi}$ . However in the conclusions  $\bar{\varphi}^{-1}$  is then restricted to the product of the Borel  $\sigma$ -field on  $K$  and that on  $\prod_{i \in I} \Theta^i$ .

The notion of *BL*-subspace is closely related to that of consistency and common knowledge. In the remainder of this section we study these relations. In doing this we shall consider from now on only **finite** *BL*-subspaces, and assume that all beliefs  $\bar{\theta}_\omega^i$  we use have finite support, i.e., we restrict ourselves to the space  $\Omega_f$  of Example 2 p. 119.

**DEFINITION 2.7.** A *BL*-subspace  $Y$  is **consistent** if there is a consistent distribution  $P$  with  $\text{Supp}(P) = Y$ .

**DEFINITION 2.8.** A state of the world  $\omega \in \Omega$  is **consistent** if it belongs to some consistent *BL*-subspace — i.e., to  $\Omega_f \cap S$ .

**LEMMA 2.9.** If  $\omega$  is consistent then  $\omega \in \text{Supp}(\bar{\theta}_\omega^i) \quad \forall i \in I$ .

**PROOF.** Let  $\omega \in Y$  and  $P \in \mathcal{P}$  such that  $\text{Supp}(P) = Y$ . Since  $P(\omega) > 0$  and  $\omega \in \mathcal{T}^i(\omega) = \{\tilde{\omega} \mid \theta^i(\tilde{\omega}) = \theta^i(\omega)\}$  it follows that  $\bar{\theta}_\omega^i(\omega) = P(\omega \mid \mathcal{T}^i(\omega)) > 0$ . ■

The term “common knowledge” is often used in game theory (including in this book) in an informal sense: Something is common knowledge among the players if every player

knows it, every player knows that every player knows it, every player knows that every player knows that every player knows it, etc. The most intuitive way to think of a common knowledge event is as an event which is announced (or shown) publicly. A formal definition of common knowledge was provided in (Aumann, 1976): A measurable space  $(Y, \bar{\mathcal{Y}})$  of states of the world is given together with sub  $\sigma$ -fields  $\mathcal{Y}^i$  of  $\bar{\mathcal{Y}}$ , one for each player. An event  $A$  is common knowledge at  $\omega$  if there exists  $B \in \bigcap_{i \in \mathbf{I}} \mathcal{Y}^i$  with  $\omega \in B \subseteq A$ .

In our model which is a model of beliefs, the natural analogue for “player  $i$  believes  $A$ ” at  $\omega$  is  $\bar{\theta}_\omega^i(A) = 1$ . And he would “know”  $A$  if further  $\omega \in A$ . In particular “player believes  $A$ ” does not necessarily imply that  $A$  is true. And of the “knowledge” operator as defined is somehow pathological, since when a player believes  $A$ , he has no way to know whether he actually knows it. This is because we did not impose any consistency conditions on the players’ beliefs. To illustrate this consider the following  $BL$ -subspace.

EXAMPLE 2.9. In a two-player situation let  $Y = \{\omega_1, \omega_2, \omega_3\}$  where

$$\begin{aligned}\omega_1 &= (k_1; (1, 0, 0), (0, 1, 0)) \\ \omega_2 &= (k_2; (1, 0, 0), (0, 1, 0)) \\ \omega_3 &= (k_3; (1, 0, 0), (0, 1, 0))\end{aligned}$$

If the true state of the world is  $\omega_3$  then player I “believes” that the state of nature is  $k_1$ , player II “believes” that it is  $k_2$  while the true state of nature is  $k_3$ . This extremely inconsistent beliefs system which can be found in the universal  $BL$ -space does not make sense as a usual knowledge system. Nevertheless we can use this notion of “beliefs” to derive what each player  $i$  believes to be “common knowledge”.

For  $\omega \in Y$  and  $i \in \mathbf{I}$  define:  $C_{\omega,1}^i = \text{Supp}(\bar{\theta}_\omega^i)$ , and inductively for  $r = 1, 2, \dots$

$$C_{\omega,r+1}^i = C_{\omega,r}^i \cup \left( \bigcup_{\tilde{\omega} \in C_{\omega,r}^i} \bigcup_j \text{Supp}(\bar{\theta}_{\tilde{\omega}}^j) \right).$$

The interpretation of this sequence is straightforward: According to player  $i$ ’s beliefs at  $\omega$ , any  $BL$ -subspace containing the state of the world, (which he generally does not know), must contain all states which he believes are possible, that is  $C_{\omega,1}^i$ . Furthermore, it must contain all states considered possible by some player in any of the states in  $C_{\omega,1}^i$ . This is  $C_{\omega,2}^i$ , and so on. Denote the union of this sequence by  $C_\omega^i$ . This is, according to player  $i$ ’s beliefs at  $\omega$ , the minimal  $BL$ -subspace containing the real state of the world but in fact it may not contain it. However it is a  $BL$ -subspace and if  $Y$  is a  $BL$ -subspace then  $C_\omega^i \subseteq Y$  for all  $i$  and for all  $\omega$  in  $Y$ . For  $\omega \in \Omega$  denote by  $Y(\omega)$  the smallest  $BL$ -subspace containing  $\omega$ .

LEMMA 2.10. If  $\omega$  is consistent then:

- (1)  $\omega \in C_\omega^i = Y(\omega)$  for all  $i$ .
- (2) There exists a unique consistent probability  $P_\omega$  with  $\text{Supp}(P_\omega) = Y(\omega)$ .

PROOF. 1. is obvious: Since  $C_\omega^i$  is a  $BL$ -subspace, and contains  $\omega$  by lemma 2.9, it contains  $Y(\omega)$ . And as observed above,  $C_\omega^i \subseteq Y(\omega)$  always.

2. For any  $\omega$  let  $\mathcal{T}^i(\omega) = \{\tilde{\omega} \mid \theta^i(\tilde{\omega}) = \theta^i(\omega)\}$ . For any  $Q \in \mathcal{P}$  we have by the consistency that for any  $j$  and any  $\tilde{\omega} \in \text{Supp}(\bar{\theta}_z^j)$ ,  $Q(\tilde{\omega}) = \bar{\theta}_z^j(\tilde{\omega})Q(\mathcal{T}^j(z))$ , hence:

$$Q(z) > 0 \text{ and } \tilde{\omega} \in \text{Supp}(\bar{\theta}_z^j) \implies \frac{Q(\tilde{\omega})}{Q(z)} = \frac{\bar{\theta}_z^j(\tilde{\omega})}{\bar{\theta}_z^j(z)} > 0$$

So by induction, if  $Q(\omega) > 0$  then  $Q(\tilde{\omega})$  is uniquely determined and  $> 0 \forall \tilde{\omega} \in C_{\omega,r}^i$  — hence  $\forall \tilde{\omega} \in C_{\omega}^i = Y(\omega)$ . Take now any  $Q \in \mathcal{P}$  which contains  $\omega$  in its support, then  $\text{Supp}(Q) \cap Y(\omega)$  is non-empty since it contains  $\omega$  (by (1)) hence  $Q(Y(\omega)) > 0$ . The required consistent  $P_{\omega}$  is

$$(3) \quad P_{\omega}(\tilde{\omega}) = \begin{cases} 0 & \text{for } \tilde{\omega} \notin Y(\omega) \\ Q(\tilde{\omega})/Q(Y(\omega)) & \text{for } \tilde{\omega} \in Y(\omega) \end{cases} \quad \blacksquare$$

COMMENT 2.10. In view of the last lemma it makes sense to think of the consistent distribution  $P_{\omega}$  on  $Y(\omega)$  as a **prior distribution**, not only because it is so mathematically speaking, but also because it is considered ‘common knowledge’ by all players: each player can first compute  $Y(\omega)$  — as  $C_{\omega}^i$  —, and then compute  $P_{\omega}$  and test for consistency by checking whether  $P_{\omega}$  is consistent and whether  $Y_{\omega} = \text{Supp}(P_{\omega})$ , using just his own knowledge  $\theta_{\omega}^i$ . The key point in the proof of lemma 2.10 p. 127 is that a consistent state is in the support of each player’s beliefs (lemma 2.9 p. 126). When this is not satisfied for a certain player, he may reach a wrong conclusion (in his test) that the current state is consistent, but his subjective probability of committing such a “type II” error is always zero. And he can never reach the wrong conclusion that the state is inconsistent when it is in fact consistent (“type I” error). The following examples show various types of (‘objective’) errors which may be committed by the players when  $\omega \notin \text{Supp}(\bar{\theta}_{\omega}^i)$ .

EXAMPLE 2.11. Consider a *BL*-subspace of two players each of which has two types thus

$$Y = \begin{pmatrix} 11 & 12 \\ 21 & 22 \end{pmatrix}$$

The subjective probability of each player on the types of the other player are given by:

$$\begin{matrix} & \text{II}_1 & \text{II}_2 \\ \text{I}_1 & \left( \begin{matrix} (1, 3/5) & (0, 0) \\ (2/3, 2/5) & (1/3, 1) \end{matrix} \right) \\ \text{I}_2 & & \end{matrix}$$

This is to be read as follows: Player  $\text{II}_1$  assigns probability  $3/5$  to the state  $11$  and  $2/5$  to the state  $21$ . Player  $\text{I}_2$  assigns probability  $2/3$  to the state  $21$  and  $1/3$  to the state  $22$ , etc. If the actual state of the world is  $\omega = 12$ , then  $\text{Supp}(\bar{\theta}_{\omega}^{\text{I}}) = \{11\}$  and  $\text{Supp}(\bar{\theta}_{\omega}^{\text{II}}) = \{22\}$ . Both players will find the *BL*-subspace  $\{11, 21, 22\}$  with the (only) consistent probability  $(1/2, 1/3, 1/6)$ . So both players will conclude that the state is consistent, committing type II error. However each player assigns to this event zero subjective probability. Note that in spite of its being inconsistent, the state  $\omega = 12$  led both players to the same ‘consistent’ space  $Y(\omega) = \{11, 21, 22\}$ . The next example shows that this is not guaranteed in an inconsistent state.

EXAMPLE 2.12. Consider the previous example with different subjective probabilities:

$$\begin{matrix} & \text{II}_1 & \text{II}_2 \\ \text{I}_1 & \left( \begin{matrix} (1, 1) & (0, 0) \\ (0, 0) & (1, 1) \end{matrix} \right) \\ \text{I}_2 & & \end{matrix}$$

If  $\omega = 12$ , player I will find the “consistent”  $C_{\omega}^{\text{I}} = \{11\}$  with probability  $Q(11) = 1$  while player II will find  $C_{\omega}^{\text{II}} = \{22\}$  with  $Q(22) = 1$ .

EXAMPLE 2.13. A *BL*-subspace  $Y$  consists of 20 states with four types of player I and five types of player II. With the same notation as in the previous examples  $Y$  is given by:

$$\begin{array}{ccccc} & \text{II}_1 & \text{II}_2 & \text{II}_3 & \text{II}_4 & \text{II}_5 \\ \text{I}_1 & (0, 0) & (1, 3/5) & (0, 0) & (0, 0) & (0, 0) \\ \text{I}_2 & (1/3, 1) & (2/3, 2/5) & (0, 0) & (0, 0) & (0, 0) \\ \text{I}_3 & (0, 0) & (0, 0) & (0, 1/2) & (1/4, 1/2) & (3/4, 1/3) \\ \text{I}_4 & (0, 0) & (0, 0) & (0, 1/2) & (1/4, 1/2) & (3/4, 2/3) \end{array}$$

If the actual state of the world is  $\omega = 13$ , player I finds  $C_\omega^I = \{21, 22, 12\}$  with the consistent probability distribution  $Q = (1/6, 1/3, 1/2)$  hence he will mistakenly conclude that the state is consistent. Player II on the other hand will find  $C_\omega^{II} = \{33, 34, 35, 43, 44, 45\}$  with no consistent  $Q$  on it. He will therefore correctly conclude that the state is inconsistent. Note that  $\omega \notin \text{Supp}(\bar{\theta}_\omega^I)$ . Unlike in previous examples player II reaches a correct negative conclusion despite of  $\omega \notin \text{Supp}(\bar{\theta}_\omega^{II})$  but this is just a matter of accident.

EXAMPLE 2.14. Consider the following *BL*-subspace with 16 states and four types for each of the players I and II.

$$\begin{array}{ccccc} & \text{II}_1 & \text{II}_2 & \text{II}_3 & \text{II}_4 \\ \text{I}_1 & (1, 3/5) & (0, 0) & (0, 0) & (0, 0) \\ \text{I}_2 & (1/3, 2/5) & (2/3, 0) & (0, 0) & (0, 0) \\ \text{I}_3 & (0, 0) & (0, 1/2) & (3/5, 1/2) & (2/5, 1/3) \\ \text{I}_4 & (0, 0) & (0, 1/2) & (3/7, 1/2) & (4/7, 2/3) \end{array}$$

If the state of the world is  $\omega = 13$ , it is inconsistent and we expect player I to come to this conclusion. In fact he will compute  $C_\omega^I = \{11, 21, 22, 32, 33, 34, 42, 43, 44\}$  but no consistent distribution on it. (To see that: by lemma 2.2 p. 120, any consistent  $Q$  must have  $Q(11) = 0$  since  $\bar{\theta}_\omega^{II}(11) = 0$  but also  $Q(11) > 0$  since  $\bar{\theta}_\omega^I(11) > 0$ ). So player I will in fact conclude that the state of the world is inconsistent. On the other hand player II will compute  $C_\omega^{II} = \{33, 34, 43, 44\}$  with consistent distribution  $Q = (1/4, 1/6, 1/4, 1/3)$  on it and mistakenly conclude that the state is consistent.

### 3. An approximation theorem

Most incomplete information models in the game theory literature, including most examples in this book, consist of **finite** *BL*-subspaces. In fact any consistent *BL*-subspace can be approximated by a finite consistent *BL*-subspace. This is the content of the following theorem:

**THEOREM 3.1.** *The probabilities in  $\mathcal{P}$  with finite support are dense in  $\mathcal{P}$ .*

**PROOF.** We first consider the case in which  $K$  is a compact metric space. 1) Let  $Q$  be a consistent probability on  $\Omega = K \times \prod_i \Theta^i$ . Let  $\mathcal{K}_m$  and  $\mathcal{C}_m^i$  be increasing sequences of measurable finite partitions of  $K$  and  $\Theta^i$  respectively, with diameter of each partition element less than  $1/m$ . Let  $\mathcal{E}_m$  be the following information scheme:  $\omega$  is chosen according to  $Q$ , each player  $i$  is informed of the atom of  $\mathcal{C}_m^i$  that contains  $\theta^i(\omega)$  and the state of nature is some given element of the atom of  $\mathcal{K}_m$  containing  $k(\omega)$ . Denote by  $Q_m$  the canonical consistent probability associated to this information scheme (theorem 2.5 p. 122). We shall prove that  $Q_m$  converges weakly to  $Q$ .

2) We first need a preliminary lemma. Recall that given a random variable  $X$  from a probability space  $(B, \mathcal{B}, P)$  to some compact metric space  $C$  with the Borel  $\sigma$ -field, we

can consider the conditional distribution of  $X$  given a sub  $\sigma$ -algebra  $\mathcal{A} \subseteq \mathcal{B}$  as transition probability from  $(B, \mathcal{A})$  to  $C$  (cf. ex. II.1Ex.16c p. 76).

LEMMA 3.2. Let  $X_m$  be a family of random variables from a probability space  $(B, \mathcal{B}, P)$  to some compact metric space and let  $\mathcal{B}_m$  be an increasing sequence of sub  $\sigma$ -algebras converging to  $\mathcal{B}_\infty$ . Assume that  $X_m$  converges a.e. to some  $X_\infty$  and let  $P_m$  denote the conditional distribution of  $X_m$  given  $\mathcal{B}_m$ . Then a.e.  $P_m$  converges weakly to  $P_\infty$ .

PROOF. Let  $\mathcal{F}$  be a countable dense subset of continuous functions on  $C$ . For  $f$  in  $\mathcal{F}$  let  $Y_m = f(X_m)$ ,  $\bar{Y}_m = \sup_{n \geq m} Y_n$ ,  $\underline{Y}_m = \inf_{n \geq m} Y_n$ .  $\bar{Y}_m$  is a decreasing sequence converging a.e. to  $Y_\infty = f(X_\infty)$ . Let  $\bar{Z}_m = E(\bar{Y}_m | \mathcal{B}_m)$ .  $\bar{Z}_m$  is a bounded supermartingale which converges a.e. to some  $\bar{Z}$ . Note that  $\bar{Z} \geq Y = E(Y_\infty | \mathcal{B}_\infty)$  a.s. since  $E(\bar{Y}_m | \mathcal{B}_m) \geq E(Y_\infty | \mathcal{B}_m)$  a.s. and this last term converges a.e. to  $Y$ . Similarly  $\underline{Z} \leq Y$  a.e. But by the supermartingale property  $E(\bar{Z} - \underline{Z}) \leq E(\bar{Z}_m - \underline{Z}_m) = E(\bar{Y}_m - \underline{Y}_m)$ , which tends to zero by the dominated convergence theorem. Hence  $\bar{Z} = \underline{Z} = Y$  a.e. Hence  $E(f(X_m) | \mathcal{B}_m)$  converges a.e. to  $E(f(X_\infty) | \mathcal{B}_\infty)$  and the result follows. ■

3) Let  $\varphi^m$  be the map of theorem 3.3 p. 123 corresponding to  $\mathcal{E}_m$ , from  $\Omega$  to itself, and let  $\varphi_n^m$  be the composition of  $\varphi^m$  with the projection  $p_n$  from  $\Omega$  to  $\Omega_n = K \times \prod_i \Theta_n^i$ . Note that  $\varphi_n^m$  converges (as  $m \rightarrow \infty$ ) a.e. to  $p_n$ , inductively, since it is true on  $K$ , and by the previous lemma the conditional distribution of  $\text{Proj}_{X_n^i} \circ \varphi_n^m$  given  $\mathcal{C}_m^i$  (i.e. the projection on  $\Theta_{n+1}^i$ ) converges a.e. to the conditional distribution of  $\text{Proj}_{X_n^i} \circ p_n$  given  $\mathcal{C}_\infty^i$  (which is precisely  $p_{n+1}^i \circ \theta^i(\omega) = \theta_{n+1}^i(\omega)$  — indeed the  $\sigma$ -field  $\mathcal{C}^i$  generated by all  $\mathcal{C}_m^i$  is a separable sub  $\sigma$ -field of the Borel  $\sigma$ -field on a compact metric space and has the points as atoms, hence it is the Borel  $\sigma$ -field on  $\Theta^i$  (App.6)). It follows that  $\varphi^m$  converges  $Q$ -a.e. to the identity on  $\Omega$  and hence the claim.

4) For a general compact  $K$  note first that for any continuous function  $f$  on  $\Omega$  there exists a metrisable quotient space  $\tilde{K}$  of  $K$  such that if  $p$  denotes the induced map from  $\Omega$  to  $\Omega_{\tilde{K}}$  (i.e.,  $p = [g_n \times \prod_i \Theta^i(g)]$  where  $g$  is the quotient map), then  $f$  factorises through  $p$  (use that continuous functions on a product of compact sets can be approximated by polynomials in continuous functions on the factors and that continuous functions on the space of probabilities on a compact set  $S$  can be approximated by polynomials in the integrals of finitely many continuous functions on  $S$ ). Next, given continuous functions  $f_1, \dots, f_n$  on  $\Omega$ , construct such a metrisable quotient  $\tilde{K}$  with the above property valid for all  $f_j$ , and use the previous case on  $\tilde{K}$  to find a consistent approximation with finite support  $\tilde{Q}_\varepsilon$  on  $\Omega_{\tilde{K}}$  approximating the image  $p(Q)$  up to  $\varepsilon$  w.r.t. the integral of each  $f_j$ . Finally, lift the  $\tilde{K}$  coordinates in the support of  $\tilde{Q}_\varepsilon$  back to  $K$  to get  $Q_\varepsilon$  ( $= [\ell \times \prod_i \Theta^i(\ell)](\tilde{Q}_\varepsilon)$  where  $\ell$  is the lifting map — any map having the right values and defined on the finitely many points in  $\tilde{K}$  that have positive probability under  $\tilde{Q}_\varepsilon$ ) on  $\Omega$ .

5) In the general case, let first  $K_n$  ( $n = 1, \dots, \infty$ ) be a sequence of disjoint compact subsets of  $K$  with  $Q(L) = 1$  for  $L = \bigcup_n K_n$ . Observe that  $\Omega$  together with  $Q$  can be viewed as an information scheme for  $L$  — by modifying the projection map to  $K$  on a null set —, so theorem 2.5 p. 122 yields that it is represented by some consistent probability  $\bar{Q}$  on  $\Omega_L \subseteq \Omega_K$ .  $\bar{Q}$  is then also a good representation for this information scheme when viewed as with values in  $K$  — i.e., the map  $\varphi$  still satisfies conditions 1a, 1b and 1a of theorem 2.5 if the range space is  $K$ . But  $Q$  itself with the identity map also satisfies them, hence by part 2 of the theorem we have  $Q = \bar{Q}$  — i.e.,  $Q(\Omega_L) = 1$ . So we can

assume  $K = \bigcup_n K_n$ . Denote by  $\overline{K}$  the disjoint union of the  $K_n$ : then as observed in the proof of theorem 2.5 p. 122,  $\Omega_{\overline{K}}$  and  $\Omega_K$  are  $K$ -Lusin, and the induced map  $f$  from  $\Omega_{\overline{K}}$  to  $\Omega_K$  is continuous, one-to-one, onto, and a Borel isomorphism, so  $\Delta(f)$  is also continuous and onto:  $Q$  can be viewed as a consistent probability on  $\Omega_{\overline{K}}$  (formally, the one obtained from theorem 2.5 p. 122 by viewing the information scheme as having values in  $\overline{K}$ ). And by the continuity of  $\Delta(f)$ , it suffices now to prove the result on  $\Omega_{\overline{K}}$ : we can assume  $K$  is a disjoint union of compact sets  $K_n$  — in particular locally compact, with one-point compactification  $\hat{K}$ . Since  $\Omega_K \subseteq \Omega_{\hat{K}}$  (point 3a of theorem 1.2 p. 111),  $Q$  can be viewed as a consistent probability for  $\hat{K}$ , assigning probability zero to the point at infinity “ $\infty$ ”. Then our previous construction yields indeed the desired result: taking care in part (4) to add to our sequence of continuous functions the indicators of the  $K_n$  yields that no point of  $K$  is identified with  $\infty$  under the quotient mapping; if we further choose in (2) the singleton  $\{\infty\}$  as an element of every partition  $\mathcal{K}_m$ , we obtain indeed that  $\infty$  is not in the support of any of the approximands. ■

COMMENT 3.1. Cf. (Mertens and Zamir, 1985) for a proof that the set of finite beliefs subspaces is dense in the space of all closed beliefs subspaces, with the Hausdorff topology.

#### 4. Games with incomplete information

**4.a. The model.** The objective of this chapter is to suggest a mathematical structure for the beliefs space in a situation of incomplete information involving several players. It is clear however that the main purpose of this structure is its application to games. Let us therefore conclude by showing briefly how this beliefs space is used in modelling **games with incomplete information**. To define a game we obviously have to add a few more ingredients to our model: Each player  $i$  in  $\mathbf{I}$  has an action set  $S^i$  (without loss of generality this may be assumed to be independent of player  $i$ 's type. One can achieve this by taking as  $S^i$  the product of the type dependent action set over all types). For  $i$  in  $\mathbf{I}$  and  $\omega$  in  $\Omega$  there is a utility function  $u_\omega^i$  which is a real valued function defined in the product of the action sets;  $S = \prod_{i \in \mathbf{I}} S^i$ . Given any finite  $BL$ -subspace  $Y$ , we first define a **vector pay-off game**  $\Gamma_Y$  in which:

- The player set is  $\mathbf{I}$ .
- The strategy set  $\Sigma^i$  of player  $i$  is the set of mappings  $\sigma^i: Y \rightarrow S^i$  which are  $\theta^i$ -measurable.
- The pay-off to player  $i$  resulting from the strategy profile  $\sigma = (\sigma^i)_{i \in \mathbf{I}}$  is the vector pay-off:  $U_i(\sigma) = (U_{\theta^i}(\sigma))_{\theta^i \in \Theta^i}$  (i.e. a pay-off for each type  $\theta^i$ ) where

$$U_{\theta^i}(\sigma) = \int u_\omega^i(\sigma(\omega)) d\bar{\theta}^i(\omega).$$

Note that  $U_{\theta^i}$  is  $\theta^i$ -measurable as it should be. Although this is not a game in the usual sense, the concept of equilibrium can be defined in the usual way, namely:

DEFINITION 4.1. The strategy profile  $\sigma = (\sigma^i)_{i \in \mathbf{I}}$  is an **equilibrium** in  $\Gamma_Y$  if for all  $\omega$  in  $Y$  and for all  $i$  in  $\mathbf{I}$ ,

$$U_{\theta^i(\omega)}(\sigma) \geq U_{\theta^i(\omega)}(\tilde{\sigma}^i, \sigma^{-i}) \text{ holds for all } \tilde{\sigma}^i \in \Sigma^i,$$

COMMENT 4.1. Note that the above game is an  $\mathbf{I}$ -person game in which the pay-off for player  $i$  is a vector with a number of coordinates equal to  $\#\Theta^i$ . It is easily seen that as far as equilibria are concerned, this game is equivalent to an ordinary  $\prod_{i \in \mathbf{I}} \#\Theta^i$ -person game in which each “player”  $\theta^i \in \Theta^i$  selects an action in  $S^i$  and then nature selects his  $\mathbb{C}\{i\}$

partners, one from each  $\Theta^j$ ,  $j \neq i$ , according to the distribution  $\theta^i$  on  $Y$ . (This is a “Selten game”, cf. (Harsanyi, 1968b, sect. 15, p. 496)). We can also define an **I**-person game with the above strategy sets where the pay-off function for player  $i$  is  $\bar{U}^i = \sum_{\theta^i \in \Theta^i} \gamma_{\theta^i} U_{\theta^i}$  where for each  $\theta^i \in \Theta^i$ ,  $\gamma_{\theta^i}$  is a strictly positive constant. Clearly, independently of the constants  $\gamma_{\theta^i}$  used, this game has the same equilibria as the vector pay-off game (and hence the corresponding Selten game).

For a consistent *BL*-subspace one has the following theorem, which permits, in looking for equilibria, to replace the normal form game by an equivalent extensive form game, called its **standard form** by Harsanyi.

**THEOREM 4.2. (Harsanyi, 1968b)** *Let  $Y$  be a consistent *BL*-subspace of  $\Omega$ . Let  $Q$  be a consistent probability distribution whose support is  $Y$ . Then the game  $\Gamma_Y$  is equivalent to the following extensive form game:*

- A chance move chooses  $\omega \in Y$  according to  $Q$ , then each player  $i \in \mathbf{I}$  is informed of his type  $\theta^i(\omega)$ .
- Each player  $i \in \mathbf{I}$  chooses an action  $s^i \in S^i$  and receives a pay-off  $U_{\theta^i(\omega)}^i(\sigma)$ .

PROOF. Follows from the definitions and the fact that  $\text{Supp}(Q) = Y$ . ■

The above theorem is especially appealing if one recalls that when  $Y$  is minimal consistent then there is a unique consistent probability distribution  $P$  on it.

**4.b. Two-person zero-sum case.** Now we start moving beyond considering the purely topological or even measure-theoretic aspects of beliefs, to richer structures, more intimately tied to game-theoretic applications. Indeed, the canonical homeomorphism between  $\Theta^i$  and  $\Delta(K \times \prod_{j \neq i} \Theta^j)$  endows  $\Theta$  with a canonical convex structure — which comes in addition to the canonical convex structure on  $\mathcal{P}$ . Further, given such a convex structure, one can consider various convex cones of continuous functions on spaces like  $K \times \prod_j \Theta^j$  — e.g. those functions which are jointly concave in the  $\theta$ -vector, those which are separately convex, or one of the above with the further requirement of being independent of  $K$ . And, according to ex. II.1Ex.20 p. 78, each such convex cone yields an ordering on the space of probability measures on this product — hence orderings on the spaces  $\Theta$  and  $\mathcal{P}$ . In addition, given such orderings, one can define other cones — and hence other orderings, etc. — by putting monotonicity requirements (possibly in addition to the concavity requirements) on the continuous functions.

The game-theoretic significance of some of those constructs appears already in (one-shot) games with incomplete information — which is our subject in this section — and much more seems to appear in the repeated case, which is the subject of ch. V and VI. But all this is still clearly at a very exploratory stage — one might argue that we will know (a first) part of what is to be understood only when we will be able to develop fully ch. VI for general “entrance laws” (and in turn such general entrance laws in ch. VI are needed for the simplest case in ch. IX) —, and a lot remains to be clarified, even for one-shot zero-sum games with finite strategy sets. Some insight however can already be gained from what follows.

We assume here that the state space  $K$  and the strategy spaces  $S$  and  $T$  of players I and II respectively are compact, and that the pay-off function  $g^k(s, t)$  on  $K \times S \times T$  is jointly continuous. (The reason for going beyond finite  $S$  and  $T$  is that in the next chapter we will want to apply those results to the discounted games.)

We first consider an auxiliary family of games, indexed by  $\mathbf{P}$  in  $\mathcal{P}$  and corresponding to the “canonical information scheme”. For each  $\mathbf{P}$  in  $\mathcal{P}$  a (one-stage) zero-sum game

$[g, \mathbf{P}]$  is now defined as follows: first a triple  $(k, \theta^I, \theta^{II})$  is chosen in  $\Omega$  according to  $\mathbf{P}$ . Next  $\theta^i$  is told to player  $i$  and finally each player selects simultaneously a move  $s$  in  $S$  (resp.  $t$  in  $T$ ). The resulting pay-off is then  $g_{st}^k$ . (The above description is known by both players).

We shall prove here some properties of  $[g, \mathbf{P}]$ .

**PROPOSITION 4.3.**  $[g, \mathbf{P}]$  has a value  $v_g(\mathbf{P})$  and both players have, for all  $\varepsilon > 0$ ,  $\varepsilon$ -optimal strategies with a finite support consisting of pure strategies taking finitely many values. They also have  $\varepsilon$ -optimal strategies which are continuous behavioural strategies with support in a fixed, finite subset of their strategy space. Moreover both players also have optimal strategies that are regular transition probabilities from  $\Theta$  to  $S$  (resp.  $T$ ).

Note that the above statement indicates that  $[g, \mathbf{P}]$  has a value in an unambiguous sense.

**PROOF.** We still denote by  $\mathbf{P}^i$  the marginal of  $\mathbf{P}$  on  $\Theta^i$ , space of types of player  $i$ . As mixed strategy space of the players we consider the compact (cf. ex. II.1Ex.17 p. 76) space  $\Sigma$  (resp.  $\mathcal{T}$ ) of regular transition probabilities from  $\Theta^I$  (resp.  $\Theta^{II}$ ) to  $S$  (resp.  $T$ ). The pay-off is then in  $L_\infty(\Omega \times S \times T, P \times \sigma \times \tau)$  and its expectation  $E_{\sigma, \tau, \mathbf{P}}(g) = \gamma(\sigma, \tau)$  is a bi-linear separately continuous (by ex. II.1Ex.18b p. 76 — recall that a jointly continuous function on a product of two compact spaces can also be viewed as a continuous map from one of the factors into the Banach space of continuous functions on the other factor) function on  $\Sigma \times \mathcal{T}$ . Theorem 1.6 p. 4 implies the existence of a value and of optimal strategies.

By the same property of jointly continuous maps, one can now find a finite subset  $S_0$  of  $S$  and a Borel map  $\varphi: S \rightarrow S_0$  such that  $|g_k(s, t) - g_k(\varphi(s), t)| \leq \varepsilon$  on  $K \times S \times T$ . The image by  $\varphi$  of player I's optimal strategy is then clearly  $\varepsilon$ -optimal, and is carried by  $S_0$ .

For the second part consider  $\sigma \mapsto \gamma(\sigma, \tau)$  as a mapping from  $S_0$ -valued transition probabilities, endowed with convergence in probability, to the space  $\mathbb{R}^{\mathcal{T}}$ , with uniform convergence (where  $\mathcal{T}$  is the set of behavioural strategies of player II), and note that this mapping is continuous. It is then possible to obtain from our previous  $\sigma$ , an  $\varepsilon$ -approximation in probability by a step function (resp., by a continuous function), say  $\sigma_\varepsilon$ , which will be  $\varepsilon$ -optimal. Finally remark that  $\sigma_\varepsilon$  can be written as a finite convex combination of pure strategies, hence the claim. ■

We now study  $v_g(\mathbf{P})$  as a function of  $\mathbf{P}$  on  $\mathcal{P}$ . Recall (theorem 2.3 p. 120 and cor. 2.4 p. 121) that this set of consistent probabilities is closed and convex.

**PROPOSITION 4.4.**  $v_g(\mathbf{P})$  is continuous and affine on  $\mathcal{P}$ .

**PROOF.** Let us first prove continuity. Denote by  $\sigma = \{\sigma(\theta) \mid \sigma(\theta) \in \Delta(S), \theta \in \Theta\}$ , an  $\varepsilon$ -optimal strategy of player I in  $[g, \mathbf{P}]$ , that we can assume, by prop. 4.3, continuous in  $\theta$ . If player II is of type  $\theta^{II}$  and uses the move  $t$ , his expected pay-off against  $\sigma$  will be

$$\int_{K \times \Theta^I} \sigma(\theta^I) g_t^k \theta^{II}(dk, d\theta^I)$$

(with  $\sigma(\theta^I) g_t^k = \sum_s \sigma(\theta^I)(s) g_{s,t}^k$ ) which is a continuous linear function of  $\theta^{II}$  by the above remarks. A best reply of player II is then obviously to choose, given  $\theta^{II}$ , some  $t$  in  $T$  minimising the above expression.

Let now  $\mathbf{P}_\alpha$  converge to  $\mathbf{P}$  in  $\mathcal{P}$ . It follows that  $\sigma$  guarantees to player I in  $[g, \mathbf{P}_\alpha]$  an expected pay-off of:

$$\int_{\Theta^{\text{II}}} \left\{ \min_{t \in T} \int_{K \times \Theta^{\text{I}}} \sigma(\theta^{\text{I}}) g_t^k \theta^{\text{II}}(dk, d\theta^{\text{I}}) \right\} \mathbf{P}_\alpha^{\text{II}}(d\theta^{\text{II}})$$

where  $\mathbf{P}_\alpha^{\text{II}}$  denotes the marginal of  $\mathbf{P}_\alpha$  on  $\Theta^{\text{II}}$  (recall that  $\mathbf{P}_\alpha$  is consistent). The integrand being continuous, so will be the integral. Note that this pay-off is a lower bound for  $v_g(\mathbf{P}_\alpha)$ , hence by the choice of  $\sigma$  we obtain that  $v_g$  is lower semi-continuous at  $\mathbf{P}$ . Permuting the rôle of the players gives upper semi-continuity hence the claim.

As for linearity, take  $\mathbf{P}'$  and  $\mathbf{P}''$  in  $\mathcal{P}$  and let  $z$  in  $(0, 1)$ . Consider the “compound” game  $[g; \mathcal{I}]$  corresponding to the following information scheme  $\mathcal{I}$ : first  $\mathbf{P}'$  and  $\mathbf{P}''$  are chosen with probabilities  $z$  and  $1 - z$ , next the players are informed of the result and then the corresponding canonical information scheme is used. A map  $\varphi$  as in theorem 2.5 p. 122, is obviously constructed using the identity on both canonical information schemes, i.e. forgetting the first lottery. Since the conditions of theorem 2.5 are obviously satisfied, it follows that the canonical consistent probability associated with  $\mathcal{I}$  is precisely  $z\mathbf{P}' + (1 - z)\mathbf{P}''$ . Now it is clear that the value of  $[g; \mathcal{I}]$  is  $\alpha v_g(\mathbf{P}') + (1 - \alpha)v_g(\mathbf{P}'')$ , thus the following prop. 4.5 implies the linearity. ■

Recalling the description of the information structure before prop. 4.3 p. 133, note that we could also define the game by starting with an information scheme  $\mathcal{I}$  consisting of some abstract probability space  $\mathbf{E} = (E, \mathcal{E}, Q)$  equipped with a random variable  $\tilde{k}$  with values in  $K$  and sub  $\sigma$ -fields  $\mathcal{E}^{\text{I}}$  and  $\mathcal{E}^{\text{II}}$ . Denote by  $[g; \mathcal{I}]$  this game and by  $v(g; \mathcal{I})$  its value. Then we have the following equivalence lemma:

**PROPOSITION 4.5.** *Let  $\tilde{\mathbf{P}}$  denote the canonical consistent probability corresponding to the scheme  $\mathcal{I}$  (theorem 2.5.1b p. 123). Then:*

$$v(g; \mathcal{I}) = v_g(\tilde{\mathbf{P}})$$

Moreover,  $\varepsilon$ -optimal strategies in  $[g, \tilde{\mathbf{P}}]$  induce, by the projection mapping  $\varphi$ ,  $\varepsilon$ -optimal strategies in  $[g; \mathcal{I}]$ .

**PROOF.** Consider a strategy  $\tilde{\sigma}$  of player I in  $[g, \tilde{\mathbf{P}}]$ . Denote by  $\sigma$  the strategy induced in  $[g; \mathcal{I}]$  by the projection mapping  $\varphi$ . It will be sufficient to show that  $\sigma$  guarantees to player I as much in  $[g; \mathcal{I}]$  as  $\tilde{\sigma}$  in  $[g, \tilde{\mathbf{P}}]$ .

Let us write the expected pay-off for player II in  $[g; \mathcal{I}]$  against  $\sigma$ , given his information  $\mathcal{E}^{\text{II}}$  and his move  $t$ . Using the structure of  $\sigma$  we obtain:

$$\int_C \tilde{\sigma}(\theta^{\text{I}} \circ \varphi(\cdot)) g_t^{\tilde{k}(\cdot)} Q(\cdot \mid \mathcal{E}^{\text{II}})$$

Now if  $\theta^{\text{II}}$  denotes the image of  $e$  in player II’s space of types, parts 1a and 1c of Theorem 2.5 p. 122 imply that the above equals:

$$\int_{K \times \Theta} \tilde{\sigma}(\theta^{\text{I}}) g_t^k \bar{\theta}^{\text{II}}(dk, d\theta^{\text{I}})$$

It follows that player’s II best replies in  $[g; \mathcal{I}]$  against  $\sigma$  at  $e$  are the same as his best replies to  $\tilde{\sigma}$  in  $[g, \tilde{\mathbf{P}}]$  at  $\theta^{\text{II}}$  ( $= \theta^{\text{II}}(\varphi(e))$ ) and both lead to the same pay-off. Finally this pay-off is only a function of  $\theta^{\text{II}}$  hence the average pay-offs under  $Q$  and  $\tilde{\mathbf{P}}$  will also coincide. ■

It is now easy to compare two consistent probabilities when they come from comparable information schemes.

**DEFINITION 4.6.** Given  $\mathbf{P}_1$  and  $\mathbf{P}_2$  in  $\mathcal{P}$ ,  $\mathbf{P}_1$  is **more informative** to player I than  $\mathbf{P}_2$ , whenever there are two information schemes  $\mathcal{I}_1$  and  $\mathcal{I}_2$  of which they are canonical representations and such that  $\mathcal{I}_1 = ((C), k, m_1^I, m^{\text{II}})$ ,  $\mathcal{I}_2 = ((C), k, m_2^I, m^{\text{II}})$ , with  $m_2^I = f(m_1^I)$ . (i.e. player I has more information in  $\mathcal{I}_1$ ).

**PROPOSITION 4.7.** If  $\mathbf{P}_1$  is more informative to player I than  $\mathbf{P}_2$  then  $v_g(\mathbf{P}_1) \geq v_g(\mathbf{P}_2)$ .

**PROOF.** The result follows from prop. 4.5 p. 134 and the relation  $v(g; \mathcal{I}_1) \geq v(g; \mathcal{I}_2)$  which is obvious since player I has more strategies in the first case and the pay-offs are the same in both. ■

**COMMENT 4.2.** It is clear that much sharper results are needed along this line, basically analytic characterisations, in the vein of Blackwell's "comparison of experiments". Observe there are two different sets of orderings one may be interested to characterise: What is the transitive closure of the above? Call it  $P_1 \succ_1 P_2$ . Then define  $P_1 R P_2$  as either  $P_1 \succ_1 P_2$  or  $P_2 \prec_2 P_1$  ( $P_2$  is less informative to player II than  $P_1$ ). What is the transitive closure of  $R$ ? Call it  $P_1 \succ P_2$ . In addition to the analytic characterisations, does the latter order allow for a converse of prop. 4.7? (E.g., is every continuous affine function on  $\mathcal{P}$  which is monotone for  $\succ$  a uniform limit of functions  $v_g(P)$ , where  $g$  varies over finite games?)

**4.c. “Approachability” in one-shot games.** We finally turn to a theorem that sharpens all the above. Henceforth  $g$  is fixed and will accordingly be dropped from the notation.

**DEFINITION 4.8.**  $\mathbf{T} = \{ \mathbf{t}: \Theta^I \rightarrow \mathbb{R} \mid \exists \text{ a continuous behavioural strategy of player II as in prop. 4.3 such that } \mathbf{t} = \mathbf{t}_\tau: \theta^I \mapsto \max_{s \in S} \int_{K \otimes \Theta^{\text{II}}} g_{s,t}^k \tau(dt \mid \theta^{\text{II}}) \theta^I(dk, d\theta^{\text{II}}) \}$ .

**DEFINITION 4.9.** For  $\mu \in \Delta(\Theta^I)$ , denote by  $P_\mu$  the canonical consistent probability corresponding to the scheme  $\mathcal{I}_\mu$  (theorem 2.5, part 1b p. 123), where  $\mathcal{I}_\mu = (\Theta^I \times \Theta^{\text{II}} \times K, \text{Borel sets}, \theta^I(d\theta^{\text{II}}, dk)\mu(d\theta^I), \Theta^I, \Theta^{\text{II}}, \text{Proj}_K)$ .

**THEOREM 4.10.** (1)  $\int_{\Theta^I} \mathbf{t} d\mu \geq v(P_\mu) \forall \mathbf{t} \in \mathbf{T}, \forall \mu \in \Delta(\Theta^I)$ .

(2) For every l.s.c., convex function  $F$  on  $\Delta(\Theta^I)$ , with  $F(\mu) > v(P_\mu) \forall \mu, \exists \mathbf{t} \in \mathbf{T}$  with  $\int \mathbf{t} d\mu < F(\mu) \forall \mu$ .

**PROOF.** (1) A strategy  $\tau$  of player II as in prop. 4.3 p. 133 guarantees that,  $\forall \theta^I \in \Delta(\Theta^{\text{II}} \times K)$ , player I's maximal expected pay-off under  $\theta^I$  equals  $\mathbf{t}_\tau(\theta^I)$ . Thus, under the scheme  $\mathcal{I}$ , the maximal expected pay-off equals  $\int \mathbf{t}_\tau(\theta^I) \mu(d\theta^I)$  (using a measurable selection theorem (7.j) to select  $\mu$ -measurably for each  $\theta^I$  an  $(\varepsilon)$ -optimal  $s$ ) — hence the integral exceeds  $v(g, \mathcal{I})$ , and so the result, by prop. 4.5 p. 134.

(2) Consider the game where player II chooses  $\mathbf{t}_\tau \in \mathbf{T}$ , player I chooses  $\mu \in \Delta(\Theta^I)$ , and the pay-off equals  $\int_{\Theta^I} \mathbf{t}_\tau d\mu - F(\mu)$ . Player I's strategy space is compact [by theorem 1.2 part 1 p. 111, remember our above assumption of compactness of  $K$ ] and convex, and player II's set of  $\tau$ 's is convex. Further, the pay-off is  $< +\infty$ , is concave and u.s.c. in  $\mu$ , and is convex in  $\tau$  since, for  $\tau = \alpha\tau_1 + (1 - \alpha)\tau_2$ ,  $\mathbf{t}_\tau \leq \alpha\mathbf{t}_{\tau_1} + (1 - \alpha)\mathbf{t}_{\tau_2}$ , because it is always beneficial for player I to know the result of the coin toss with probabilities  $\alpha$  and  $1 - \alpha$  rather than not. So all the assumptions of the minmax theorem (prop. 1.8 p. 5) are satisfied. Finally, by prop. 4.5 p. 134,  $F(\mu) > v(P_\mu) = v(g, \mathcal{I})$ . Hence approximating an

$\varepsilon$ -optimal  $\tau$  in the game  $[g, \mathcal{I}]$  by a continuous  $\tau$  as in prop. 4.3 p. 133, which, as shown sub 1 above, guarantees exactly  $\int \mathbf{t}_\tau d\mu$ , we obtain that the  $\supinf$  of our game is negative. Hence the result by the minmax theorem. ■

We reformulate now the above result without using functions as general as  $F$  above.

- COROLLARY 4.11. (1) The function  $v(P_\mu)$  is concave and u.s.c. on  $\Delta(\Theta^I)$ .  
(2) For every continuous (or l.s.c.) function  $f$  on  $\Theta^I$ , with  $\int f d\mu > v(P_\mu) \forall \mu, \exists \mathbf{t} \in \mathbf{T}$  with  $\mathbf{t} < f$ .  
(3)  $v(P_\mu) = \inf_{\mathbf{t} \in \mathbf{T}} \int \mathbf{t} d\mu$ .

PROOF. Let in theorem 4.10 part 2  $F(\nu) = +\infty$  for  $\nu \neq \mu$ ,  $F(\mu) = v(P_\mu) + \varepsilon$ : we get  $v(P_\mu) \geq \inf_{\mathbf{t} \in \mathbf{T}} \int \mathbf{t} d\mu$ . By part 1 of theorem 4.10, we get then equality, i.e.(3), and hence (1). And (2) is obviously a particular case of part 2 of theorem 4.10. ■

REMARK 4.3. The corollary has the full force of the theorem, by the separation theorem (1.21 and 1.23 p. 8). Therefore it also implies the following:

COROLLARY 4.12. Each player has continuous behavioural strategies (or strategies with finite support) as in prop. 4.3 p. 133, which are  $\varepsilon$ -optimal for all  $P \in \mathcal{P}$ .

PROOF. The maps from  $\mathcal{P}$  to the corresponding marginals on  $\Theta^I$  and on  $\Theta^{II}$  are affine, continuous, and one to one, so they are affine homeomorphisms of compact (cor. 2.4 p. 121), convex sets. So when  $\mu$  varies over this set of marginals, part 1 of Corollary 4.11 yields that  $v(P)$  is concave and u.s.c. on  $\mathcal{P}$ . Hence by duality it is affine and continuous. Define then  $F(\mu) = v(P) + \varepsilon$  when  $\mu$  is the projection of some  $P \in \mathcal{P}$ ,  $F(\mu) = +\infty$  otherwise:  $F$  is convex and l.s.c. Apply now part 2 of the theorem. ■

COMMENT 4.4. The corollary “confirms” the comment after lemma 2.10 p. 127: That lemma is restricted to finite *BL*-spaces, and states only that, under consistency, the minimal *BL*-subspace containing  $\omega$  is common knowledge, so if players’ beliefs satisfy the finiteness assumption and consistency, but especially if they believe that the true state of the world is always generated by the minimal consistent probability under which it has positive probability<sup>1</sup>, then they would be able to know the true game, and hence to play correctly, knowing just their own types. The corollary implies that knowing one’s own type is sufficient to be able to play correctly, without any assumption beyond consistency.

COMMENT 4.5. Further, the continuity of the behavioural strategies guarantees that it suffices for the player to know his own type approximately.

COROLLARY 4.13.  $v(P)$  is continuous and affine on  $\mathcal{P}$ .

PROOF. Was derived in the proof of cor. 4.12. Alternatively,  $v(P)$  is the uniform limit (cor. 4.12) of the continuous affine functions  $\int \mathbf{t}(\theta_\omega^I) P(d\omega)$ . ■

COMMENT 4.6. The convexity of the functions  $\mathbf{t} \in \mathbf{T}$  seems closely related to prop. 4.7 p. 135, cf. also the comment following this proposition. But this relation remains to be elucidated.

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<sup>1</sup>Such an assumption would be quite inconsistent *in se* since there are quite obvious information schemes (like the one used in the proof of the second half of prop. 4.4 p. 133) that can generate  $\omega$  without being minimal with this property. And when leaving the finiteness restriction, it becomes even formally nonsense, since for many states of the world, there may be so many consistent *BL*-subspaces containing it in their support (in whatever sense — cf. e.g. examples 6 or 8a p. 119) that there may very well even be no minimal one despite 3 p. 121, or that there may be several minimal ones.

COMMENT 4.7. A (related) question to be elucidated is to state separate properties of the map  $\mu \mapsto P_\mu$  and of the map  $P \mapsto v(P)$  which ensure that the composition is concave and u.s.c. (part 1 of cor. 4.11).

COMMENT 4.8. And possibly the composition is continuous after all? Observe that modulo a proof that the set of functions  $v_g(P)$ , where  $g$  varies, separates points of  $\mathcal{P}$ , this would immediately imply the continuity of  $\mu \mapsto P_\mu$ , thus giving a first partial answer to remark 4.7. The continuity of this map would also be of great help whenever handling concavification and convexification operators (cf. below). Finally, the continuity of  $v(P_\mu)$  would be a first step in showing that, with  $\mathbf{T}_m$  being the set of minimal elements of  $\{f: \Theta^I \rightarrow \mathbb{R} \mid f \text{ is convex and continuous, and } \int f d\mu \geq v(P_\mu) \forall \mu\}$ , every  $\mathbf{t} \in \mathbf{T}$  is minorated by some  $f \in \mathbf{T}_m$ , and  $\mathbf{T}_m$  is equicontinuous. [It is a situation like this that we are going to encounter in ch. V and VI, but in a more restrictive framework, where the  $\mu$ 's are restricted to be carried by those  $\theta^I$ 's that assign probability one to a fixed, finite subset of  $\Theta^I$ , and where the game is repeated, so those functions in  $\mathbf{T}_m$  can be realised by some strategy of II (and the function guaranteed by any strategy of II is minorated by some  $f \in \mathbf{T}_m$ ). Further there one tries to define those sets without using *any* topological restriction.]

COMMENT 4.9. Observe that already now we could define some  $\mathbf{T}_m$  by asking only upper semi-continuity instead of continuity of the functions  $f$ , using e.g. Zorn's lemma, and the regularity of the measures. We would indeed have that any convex u.s.c. function  $f$  satisfying the inequalities, and thus a fortiori any l.s.c.  $f$  (convex or not) (part 2 of cor. 4.11), and a fortiori any  $f \in \mathbf{T}$  is minorated by one in  $\mathbf{T}_m$ . But nothing more; even if one were to define  $\tilde{\mathbf{T}}_m$  analogously but deleting the convexity restriction on the functions  $f$ , we do not know we would get the same set (cf. also ex. I.3Ex.15c p. 38 for how such a problem is handled in a finite dimensional case).

COMMENT 4.10. The problem of equicontinuity in remark 4.9 is the occasion to mention two other such problems (conceivably somewhat related). Lipschitz properties are of crucial importance in ch. V and VI.

- (1) One expects a conditional expectation operator to be a smoothing operator, e.g. not to increase the Lipschitz constant for appropriate distances. Assume thus a distance on  $K$ , or even, take  $K$  finite, with distance 1 between any two distinct points. Consider the set  $E_0$  of Lipschitz functions with Lipschitz constant 1 on  $K$ , viewed as functions on  $\Omega$ , and then the smallest lattice  $E$  of functions on  $\Omega$  that contains  $E_0$  and contains the function  $\omega \mapsto \bar{\theta}_\omega^i(f) \forall i \in \mathbf{I}, \forall f \in E$ . Is  $E$  equicontinuous? (This is equivalent to the existence of a distance with the above mentioned property).
- (2) For  $K$  finite, is the set of functions  $v_g(P)$  on  $\mathcal{P}$  equicontinuous? Here  $S$  and  $T$  vary over all finite sets, and  $g$  over all games with  $|g_{s,t}^k| \leq 1 \forall k, \forall s, \forall t$ . The answer is affirmative when  $P$  is restricted to those consistent probabilities that project to a point mass on one of the factors (cf. Mertens, 1986b).

**4.d. Concavification and Convexification.** We introduce here two types of definitions, motivated by part 1 of cor. 4.11 p. 136, and which will be fundamental in the sequel.

DEFINITION 4.14. A real function  $f$  on  $\mathcal{P}$  is *concave w.r.t. I* ( $\mathbf{Cav}_I$ ) [resp. *convex w.r.t. II* ( $\mathbf{Vex}_{II}$ )] if  $\mu \mapsto f(P_\mu)$  is concave on  $\Delta(\Theta^I)$  (resp. convex on  $\Delta(\Theta^{II})$ , with the obvious definition of  $P_\mu$ ).

COMMENT 4.11. This is just a “template” definition: A number of variants are possible, e.g. concave and u.s.c., or a minimum of integrals of convex continuous functions, etc. When one will be able to develop ch. VI entirely with arbitrary entrance laws (and after clearing some of the open problems in this section) one will also see which exactly of the variants is the most useful.

With the same precautions, one can also define the following:

- DEFINITION 4.15. (1) For any  $P \in \mathcal{P}$ , denote by  $\mathcal{I}_P$  the corresponding canonical information scheme (i.e., like  $\mathcal{I}_\mu$  in definition 4.9 p. 135, but with  $P$  as probability measure). For any information scheme  $\mathcal{I}$ , let  $P_{\mathcal{I}}$  denote the corresponding canonical consistent probability as in definition 4.9 p. 135. Let also  $\mathcal{I}_c$  denote the corresponding canonical information scheme.  
(2) Given a function  $f$  on  $\mathcal{P}$ , and an information scheme  $\mathcal{I}$ , define  $\bar{f}(\mathcal{I})$  as  $f(P_{\mathcal{I}})$ . With  $\mathcal{I} = (E, \mathcal{E}, Q, (\mathcal{E}_i)_{i \in \mathbf{I}}, \tilde{k}_E)$  define also  $\mathcal{I}_g$ ,  $\forall g \in L_1(E, \mathcal{E}_1, Q)$  with  $g \geq 0$ ,  $\|g\|_1 = 1$  as the information scheme  $(E, \mathcal{E}, g dQ, (\mathcal{E}_i)_{i \in \mathbf{I}}, \tilde{k}_E)$ , (so  $\mathcal{I} = \mathcal{I}_1$ ). Define finally  $\hat{f}(\mathcal{I})$  as  $\sup\{\sum \sigma_n \bar{f}(\mathcal{I}_{g_n}) \mid \sigma \in \Delta(N), \sum_n \sigma_n g_n = 1, N \text{ finite}\}$ .  
(3) Define now  $f$  on  $\mathcal{P}$  as strongly concave (w.r.t. I) if  $\bar{f}(\mathcal{I}) = \hat{f}(\mathcal{I})$  for every information scheme  $\mathcal{I}$ .

REMARK 4.12. Equivalently,  $\hat{f}(\mathcal{I}_g)$  is the smallest concave function of  $g$  that majorates  $\bar{f}(\mathcal{I}_g)$ . And  $f$  is strongly concave if  $g \rightarrow \bar{f}(\mathcal{I}_g)$  is concave for every  $\mathcal{I}$ .

DEFINITION 4.16. A function  $f$  on  $\mathcal{P}$  is monotone w.r.t. I, if for any information scheme  $\mathcal{I}$ , and any event  $A \in \mathcal{E}_I$ , the information scheme  $\mathcal{I}_A$  obtained by adding  $\forall i$ ,  $A$  to  $\mathcal{E}_i$  satisfies  $\bar{f}(\mathcal{I}) \geq \bar{f}(\mathcal{I}_A)$ .

COMMENT 4.13. In reference to our comment after prop. 4.7 p. 135, the above ordering corresponds (for two-person games) to a particular case of the monotonicity w.r.t. the ordering “ $P \prec_2 P_A$ ”. The particular case appears less natural than the general concept, since there is no good reason in general to restrict II’s additional information to come from player I. However, it will be sufficient for our purposes in this book. In particular, in ch. V, player I is fully informed, so the distinction becomes immaterial, and in ch. VI, “signalling matrices” are independent of the state of nature, so the only additional information a player can get must come from the other player. But again, the general meaning of those concepts, and their relation, remains to be elucidated.

- PROPOSITION 4.17. (1) A strongly concave function w.r.t. I is concave w.r.t. I.  
(2) The value function  $v(P)$  of a game is strongly concave w.r.t. I.  
(3) More precisely, if a concave real valued function  $f$  on  $\mathcal{P}$  is monotone w.r.t. I, it is strongly concave w.r.t. I.  
(4) For every real valued function  $f$  on  $\mathcal{P}$ , there is a smallest (strongly) concave function w.r.t. I that majorates it, denoted by  $\text{Cav}_I(f)$  [resp.  $S\text{Cav}_I(f)$ ].

PROOF. (1) follows from the Radon-Nikodym theorem: if  $\mu = \sum_n \sigma_n \mu_n$ , let  $g_n = d\mu_n/d\mu$ , and apply remark 4.12 using  $\mathcal{I}_\mu$  (definition 4.9) for the scheme  $\mathcal{I}$ .

(2) follows from (3) by propositions 4.4 p. 133 and 4.7 p. 135 and remark 4.13.

(3) If  $\mathcal{I}$  is an information scheme, with  $g_n \in L_1^+(E, \mathcal{E}_I, Q)$   $\forall n \in N$ ,  $\|g_n\|_1 = 1$ , and  $\sum_n \sigma_n g_n = 1$ , define the information scheme  $\mathcal{I}_{\sigma,g}$  as  $[(E \times N, \mathcal{E} \times N, Q_{\sigma,g}), (\mathcal{E}_i \times N)_{i \in \mathbf{I}}, \tilde{k}_E \circ \text{Proj}_E]$  with  $Q_{\sigma,g}$  defined by  $\sigma_n$  being the marginal probability of  $n$ , and  $g_n(e)Q(de)$  the conditional probability on  $E$  given  $n$ . We first claim that, if  $f$  is monotone w.r.t. I, then

$\bar{f}(\mathcal{J}) \geq \bar{f}(\mathcal{J}_{\sigma,g})$ . Indeed, let  $\bar{\mathcal{J}}_{\sigma,g}$  denote the same information scheme as  $\mathcal{J}_{\sigma,g}$ , except that players  $i$  different from I have only  $\mathcal{E}_i$  [i.e.  $(\text{Proj}_E)^{-1}(\mathcal{E}_i)$ ] as private information. The monotonicity implies that  $\bar{f}$  decreases every time when, starting from  $\bar{\mathcal{J}}_{\sigma,g}$ , one adds to all players' private information the events  $n = 1$ , then  $n = 2$ , etc.; so at the end  $\bar{f}(\bar{\mathcal{J}}_{\sigma,g}) \geq \bar{f}(\mathcal{J}_{\sigma,g})$ , and there remains only to show that  $\bar{f}(\bar{\mathcal{J}}_{\sigma,g}) = \bar{f}(\mathcal{J})$ . To prove that  $\mathcal{J}_{\sigma,g}$  and  $\mathcal{J}$  have the same canonical distribution it suffices to show that:

LEMMA 4.18. *Let  $\mathcal{J}$  be an information scheme  $[(E, \mathcal{E}, Q), (\mathcal{E}_i)_{i \in \mathcal{I}}, \tilde{k}_E]$ , and similarly for  $\tilde{\mathcal{J}}$ . Let  $\psi: E \rightarrow \tilde{E}$  be a map, which is  $(\mathcal{E}, \tilde{\mathcal{E}})$ -measurable and also  $(\mathcal{E}_i, \tilde{\mathcal{E}}_i)$ -measurable  $\forall i \in \mathbf{I}$ , at least when all  $Q$ -negligible sets have been added to  $\mathcal{E}$  and to all  $\mathcal{E}_i$ . Assume that  $\tilde{Q} = Q \circ \psi^{-1}$ , that for every Borel set  $B$  in  $K$ ,  $\tilde{k}_E^{-1}(B) = (\tilde{k}_{\tilde{E}} \circ \psi)^{-1}(B)$   $Q$ -a.e., and that,  $\forall i \in \mathbf{I}$ ,  $\mathcal{E}$  and  $\mathcal{E}_i$  are conditionally independent given  $\psi^{-1}(\tilde{\mathcal{E}}_i)$ . Then  $P_{\mathcal{J}} = P_{\tilde{\mathcal{J}}}$  — more precisely, if one of the schemes satisfies the assumptions of Theorem 2.5 p. 122 so does the other, and  $\phi_{\tilde{\mathcal{J}}} \circ \psi$ , where  $\phi_{\tilde{\mathcal{J}}}$  is the map associated to  $\tilde{\mathcal{J}}$  by part 1 of theorem 2.5, is such a map for  $\mathcal{J}$ .*

REMARK 4.14. The map  $\phi$  itself from  $\mathcal{J}$  to  $\mathcal{J}_c$  has the above properties (part 1 of theorem 2.5 p. 122). It is thus a “maximal” map with those properties.

PROOF. Obvious verification. ■

This lemma is indeed sufficient, because the  $\mathcal{E}_I$ -measurability of  $g_n$  implies that, for  $f \geq 0$ ,  $\mathcal{E}$ -measurable,  $E_Q(f \mid \mathcal{E}_I) = E_{g_n dQ}(f \mid \mathcal{E}_I)$ , and hence the required conditional independence of  $\mathcal{E}$  and  $\mathcal{E}_I \times N$  given  $\mathcal{E}_I$ . Thus we have proved our claim that, under I-monotonicity,  $\bar{f}(\mathcal{J}) \geq \bar{f}(\mathcal{J}_{\sigma,g})$ .

There just remains to show that, under concavity,  $\bar{f}(\mathcal{J}_{\sigma,g}) \geq \sum_n \sigma_n \bar{f}(\mathcal{J}_{g_n})$ . This follows from the following lemma, which implies that  $P_{\mathcal{J}_{\sigma,g}} = \sum_n \sigma_n P_{\mathcal{J}_{g_n}}$  (and which was already used implicitly in the proof of prop. 4.4 p. 133).

LEMMA 4.19. *Given an information scheme  $\mathcal{J}$ , satisfying the assumptions of Theorem 2.5 p. 122, denote by  $\bar{\mathcal{E}}_i$  the  $\sigma$ -field generated by  $\mathcal{E}_i$  and all  $Q$ -null sets, and let  $\mathcal{E}_{ck} = \bigcap_{i \in \mathbf{I}} \bar{\mathcal{E}}_i$ . Given a countable  $\mathcal{E}_{ck}$ -measurable partition  $A_n$ , let  $\sigma_n = Q(A_n)$ , and, for  $\sigma_n > 0$ ,  $Q_n(B) = Q(B \mid A_n)$  — otherwise e.g.  $Q_n = Q$ . Denote by  $\mathcal{J}_{A_n}$  the same scheme with  $Q_n$  instead of  $Q$ . Then  $P_{\mathcal{J}} = \sum_n \sigma_n P_{\mathcal{J}_{A_n}}$ .*

PROOF. Clearly  $\mathcal{J}_{A_n}$  also satisfies the assumptions of Theorem 2.5. Let  $\phi_n$  be a map from  $\mathcal{J}_{A_n}$  to  $\Omega$  as in part 1 of theorem 2.5. Define  $\phi$  by letting it equal  $\phi_n$  on  $A_n$ . So  $P_{\mathcal{J}_{A_n}} = \phi_n(Q_n) = \phi(Q_n) \in \mathcal{P}$ , and thus (closedness and convexity of  $\mathcal{P}$ ) also  $\sum_n \sigma_n P_{\mathcal{J}_{A_n}} \in \mathcal{P}$ . But  $\sum_n \sigma_n P_{\mathcal{J}_{A_n}} = \sum_n \sigma_n \phi(Q_n) = \phi(\sum \sigma_n Q_n) = \phi(Q)$ . So  $\phi(Q) \in \mathcal{P}$ . The other two properties of the map  $\phi$  required in part 1 of theorem 2.5 are also an obvious verification. The result follows then from part 2 of the same theorem. ■

This finishes the proof of point 3.

For point 4, just observe that, using remark 4.12 for Strong Concavity, both concavity and strong concavity w.r.t. I are defined in terms of usual concavity of auxiliary functions, and that a lower bound of usually concave functions is still so.

This finishes the proof of prop. 4.17. ■

COMMENT 4.15. Observe thus that, given a real function  $f$  on  $\mathcal{P}$ , we can define the following “concavifications”:

- $C_0(f)$  as the concavification over the convex set  $\mathcal{P}$

- $C_4(f)$  as  $\text{Cav}_I(f)$
- $C_5(f)$  as  $S\text{Cav}_I(f)$
- $C_1(f)(P) = \hat{f}(\mathcal{J}_P)$
- $C_2(f)(P) = \sup\{\hat{f}(\mathcal{J}_\mu) \mid P_\mu = P\}$
- $C_3(f)(P) = \sup\{\hat{f}(\mathcal{J}) \mid P_{\mathcal{J}} = P\}$

COMMENT 4.16. All those functions are concave functions on the convex set  $\mathcal{P}$ . This is obvious for  $C_0, C_1, C_4$  and  $C_5$ . For  $C_2(f)$ , let  $\hat{f}(\mathcal{J}_{\mu_i}) \simeq C_2(f)(P_i)$ , with  $P = \alpha P_1 + (1-\alpha)P_2$ ,  $P_i = P_{\mu_i}$ . Choose Borel isomorphisms  $\varphi_1$  and  $\varphi_2$  of  $\Theta^{\text{II}}$  with disjoint Borel subsets of  $\Theta^{\text{II}}$ , and let, for  $\theta \in \Theta^I = \Delta(K \times \Theta^{\text{II}})$ ,  $\bar{\varphi}_i(\theta)$  be the image measure of  $\theta$  under  $id_K \times \varphi_i$ , and let  $\nu_i$  be the image measure of  $\mu_i$  under  $\bar{\varphi}_i$ . Then  $\mathcal{J}_{\nu_i}$  is Borel-isomorphic to  $\mathcal{J}_{\mu_i}$ , so (lemma 4.18 p. 139)  $P_{\nu_i} = P_{\mu_i} = P_i$  and  $\hat{f}(\mathcal{J}_{\nu_i}) = \hat{f}(\mathcal{J}_{\mu_i})$ . Further, with  $\nu = \alpha\nu_1 + (1-\alpha)\nu_2$ , lemma 4.19 p. 139 yields that  $P_\nu = \alpha P_{\nu_1} + (1-\alpha)P_{\nu_2} = P$ . Thus  $\hat{f}(\mathcal{J}_\nu) \geq \alpha\hat{f}(\mathcal{J}_{\nu_1}) + (1-\alpha)\hat{f}(\mathcal{J}_{\nu_2}) = \alpha\hat{f}(\mathcal{J}_{\mu_1}) + (1-\alpha)\hat{f}(\mathcal{J}_{\mu_2}) \simeq \alpha[C_2(f)](P_1) + (1-\alpha)[(C_2(f))(P_2)]$  — hence the result. Finally, for  $C_3(f)$ , the argument is similar — there is no need for the Borel isomorphisms, just construct the obvious bigger information scheme  $\mathcal{J}$  where  $\mathcal{J}_1$  or  $\mathcal{J}_2$  is selected at random with probability  $\alpha$  or  $(1-\alpha)$ , and the choice is told to both players.

COMMENT 4.17. We have the following obvious relations between the  $C_i$ :  $C_0(f) \leq C_1(f) \leq C_2(f) \leq C_3(f) \leq C_5(f)$ , and  $C_2(f) \leq C_4(f) \leq C_5(f)$ . There is obviously a reasonable hope that  $C_2 = C_3$  (remember the Borel isomorphisms — there might be a similar trick that produces, for any  $\mathcal{J}$ , some  $\mathcal{J}_\mu$  which is “sufficiently isomorphic”). A fortiori one “should” have  $C_3 \leq C_4$ . Also the obvious reason for having defined  $C_2$  and  $C_3$  is that they were the natural candidates for being equal resp. to  $C_4$  and  $C_5$ . By our above inequalities, either of those equalities would imply  $C_4 \leq C_3$ . Finally, the reason to have introduced  $C_1$  is that, in all our applications in ch. V and VI,  $C_1(f) = C_5(f)$ . But this might depend on some monotonicity properties of the functions  $f$  considered; in particular the “ $u$ ”-functions used there are monotone. Since we also use iterated operators like  $\text{Vex}_{\text{II}}\text{Cav}_I(f)$ , it would be important to know what monotonicity properties of  $f$  are preserved by what concavification operators.

COMMENT 4.18. Consider the following cautionary example, that there may be difficulties even when everything is finite. Let  $K = \{A, B\}$ , fix  $\theta^\alpha \neq \theta^\beta$  in  $\Theta^{\text{II}}$ , let  $\theta^1 = \frac{1}{2}\delta_{(A,\theta^\alpha)} + \frac{1}{2}\delta_{(B,\theta^\beta)} \in \Theta^I$  and  $\theta^2 = \frac{1}{2}\delta_{(B,\theta^\alpha)} + \frac{1}{2}\delta_{(A,\theta^\beta)} \in \Theta^I$ , and let  $\mu = \frac{1}{2}\delta_{\theta^1} + \frac{1}{2}\delta_{\theta^2} \in \Delta(\Theta^I)$ . Then  $P = P_\mu$  is (e.g. by Lemma 4.18 p. 139) the canonical distribution of the information scheme where it is common knowledge that  $A$  and  $B$  are selected with equal probability, i.e.  $P = \frac{1}{2}\delta_{\omega_A} + \frac{1}{2}\delta_{\omega_B}$ , with  $\omega_A = (\theta^I, \theta^{\text{II}}, A)$ ,  $\omega_B = (\theta^I, \theta^{\text{II}}, B)$ ,  $\bar{\theta}^I = \bar{\theta}^{\text{II}} = P$ . So, since  $P$  projects as a unit mass on  $\Theta^I$ , we have, for any function  $f$ ,  $C_0(f)(P) = C_1(f)(P) = f(P)$ .

Let also  $Q$  be the canonical distribution of the information scheme where  $A$  and  $B$  are selected with equal probability, while player II is informed of this choice and player I not. I.e.,  $Q = \frac{1}{2}\delta_{\omega_a} + \frac{1}{2}\delta_{\omega_b}$ , with  $\omega_a = (\theta, \theta_a, A)$ ,  $\omega_b = (\theta, \theta_b, B)$ ,  $\theta = Q$ ,  $\theta_a = \delta_{\omega_a}$ ,  $\theta_b = \delta_{\omega_b}$ .

Consider now, for the scheme  $\mathcal{J}_\mu$ , the functions  $g_1 = 2 \cdot \mathbb{1}_{\{\theta^1\}}$ ,  $g_2 = 2 \cdot \mathbb{1}_{\{\theta^2\}}$ , and  $\sigma_1 = \sigma_2 = \frac{1}{2}$ . Then  $P_{(\mathcal{J}_\mu)_{g_1}} = P_{(\mathcal{J}_\mu)_{g_2}} = Q$ , so  $\hat{f}(\mathcal{J}_\mu) \geq f(Q)$ . Hence if  $f(Q) > f(P)$ , we have  $C_2(f) \neq C_1(f)$ .

COMMENT 4.19. Observe that the above example relies on a non-monotonicity of  $f$ :  $f(Q) > f(P)$ . To finish confusing this issue, consider now a variant of the above where  $\mu_\varepsilon = \frac{1}{2}(1+\varepsilon)\delta_{\theta^1} + \frac{1}{2}(1-\varepsilon)\delta_{\theta^2}$ . Then  $P_\varepsilon = P_{\mu_\varepsilon}$  is the unique consistent probability on (dropping superscripts  $\varepsilon$  on all  $\omega$ 's and  $\theta$ 's)  $\{\omega_{11}, \omega_{12}, \omega_{21}, \omega_{22}\}$  with  $\omega_{11} = (A, \theta^1, \theta^\alpha)$ ,

$\omega_{12} = (B, \theta^1, \theta^\beta)$ ,  $\omega_{21} = (B, \theta^1, \theta^\beta)$ ,  $\omega_{22} = (A, \theta^2, \theta^\beta)$ ,  $\theta^1 = \frac{1}{2}\delta_{\omega_{11}} + \frac{1}{2}\delta_{\omega_{12}}$ ,  $\theta^2 = \frac{1}{2}\delta_{\omega_{21}} + \frac{1}{2}\delta_{\omega_{22}}$ ,  $\theta^\alpha = \frac{1}{2}(1 + \varepsilon)\delta_{\omega_{11}} + \frac{1}{2}(1 - \varepsilon)\delta_{\omega_{21}}$ , and  $\theta^\beta = \frac{1}{2}(1 + \varepsilon)\delta_{\omega_{12}} + \frac{1}{2}(1 - \varepsilon)\delta_{\omega_{22}}$ . (And  $\theta^1 \neq \theta^2$ ,  $\theta^\alpha \neq \theta^\beta$ ).

We claim that  $\omega_{ij}^\varepsilon \rightarrow \omega_{ij}^\infty$ , and  $P_\varepsilon \rightarrow P_\infty = P$ , where  $\omega_{11}^\infty = \omega_{22}^\infty = \omega_A$ ,  $\omega_{12}^\infty = \omega_{21}^\infty = \omega_B$ . Otherwise, extract ( $\Omega$  and  $\mathcal{P}$  are compact metric) a subsequence along which  $\omega_{ij}^\varepsilon$  converges and  $P_\varepsilon$  converges, but the limits are not the above. By this convergence, our relations between the  $\omega_{ij}^\varepsilon$  and the  $\theta^{*,\varepsilon}$  pass to the limit, yielding that  $P_\infty$  is a consistent probability on  $\{\omega_{11}^\infty, \omega_{12}^\infty, \omega_{21}^\infty, \omega_{22}^\infty\}$  with  $\omega_{11} = (A, \theta^{1,\infty}, \theta^{\alpha,\infty})$ , etc.,  $\theta^{1,\infty} = \frac{1}{2}\delta_{\omega_{11}^\infty} + \frac{1}{2}\delta_{\omega_{12}^\infty}$ ,  $\theta^{2,\infty} = \dots$ ,  $\theta^{\alpha,\infty} = \frac{1}{2}\delta_{\omega_{11}^\infty} + \frac{1}{2}\delta_{\omega_{21}^\infty}$ ,  $\theta^{\beta,\infty} = \dots$ . If either  $\theta^{1,\infty} \neq \theta^{2,\infty}$  or  $\theta^{\alpha,\infty} \neq \theta^{\beta,\infty}$ , the four points  $\omega_{ij}^\infty$  are different, so the canonical information scheme associated with  $P_\infty$  is isomorphic to our original  $\mathcal{I}_\mu$  (of remark 4.18), so they have the same consistent distribution  $P$ . Since  $P_\infty \in \mathcal{P}$ , it equals its own consistent distribution, so  $P_\infty = P$ . While if both  $\theta^{1,\infty} = \theta^{2,\infty}$  and  $\theta^{\alpha,\infty} = \theta^{\beta,\infty}$ , then  $\omega_{11}^\infty = \omega_{22}^\infty$  and  $\omega_{12}^\infty = \omega_{21}^\infty$ , and the canonical information scheme associated with  $P_\infty$  is isomorphic to that of  $P$ , hence again  $P_\infty = P$ , by the same argument. So, any way  $P_\infty = P$ . It follows then that  $\omega_{11}^\infty = \omega_{22}^\infty$ , being in the support of  $P$  and being mapped to  $A$ , must equal  $\omega_A$ , and similarly  $\omega_{12}^\infty = \omega_{21}^\infty = \omega_B$ . Hence our convergence; so  $\theta^{\alpha,\varepsilon} \rightarrow \theta^{\text{II}}$ ,  $\theta^{\beta,\varepsilon} \rightarrow \theta^{\text{II}}$ ,  $\theta^{1,\varepsilon} \rightarrow \theta^{\text{I}}$ ,  $\theta^{2,\varepsilon} \rightarrow \theta^{\text{I}}$ .

For the canonical scheme  $\mathcal{I}_{P_\varepsilon}$ , let  $g_1^\varepsilon = \frac{2}{1+\varepsilon}\mathbb{1}_{\theta^{1,\varepsilon}}$ ,  $g_2^\varepsilon = \frac{2}{1-\varepsilon}\mathbb{1}_{\theta^{2,\varepsilon}}$ ,  $\sigma_1^\varepsilon = \frac{1+\varepsilon}{2}$ ,  $\sigma_2^\varepsilon = \frac{1-\varepsilon}{2}$ . Then, as above,  $P_{(\mathcal{I}_{P_\varepsilon})_{g_1^\varepsilon}} = P_{(\mathcal{I}_{P_\varepsilon})_{g_2^\varepsilon}} = Q$ , so  $\hat{f}(\mathcal{I}_{P_\varepsilon}) \geq f(Q)$ , hence  $[C_1(f)](P_\varepsilon) \geq f(Q)$ .

Since  $P_\varepsilon \rightarrow P$ , any u.s.c. function  $\varphi$  on  $\mathcal{P}$  with  $\varphi \geq C_1(f)$  — in particular  $C_1(f)$  itself when u.s.c. — satisfies  $\varphi(P) \geq f(Q)$ . So it appears that the relations between the different  $C_i(f)$ 's may be substantially strengthened, not only under monotonicity assumptions of  $f$ , but also possibly under topological assumptions, either on  $f$ , or on  $C_i(f)$ , or as a variant of the concavification operators (cf. remark 4.11, and recall also the u.s.c. in part 1 of cor. 4.11 p. 136, and that all “ $u$ ”-functions are continuous on  $\mathcal{P}$ ).

### Exercises

- If  $f: K_1 \rightarrow K_2$  is Borel and surjective between analytic spaces, so is  $\Theta(f)$ .

HINT. 7.j.

**2. An alternative construction of the beliefs space.** (Mertens and Zamir, 1985) If we take the beliefs of a player to be a joint probability distribution on  $K$  and the beliefs of all players (including his own) we make the beliefs spaces to be the same for all players. The coherency of the beliefs are imposed as restrictions on these larger dimensional probabilities. This can be done as follows:

- Define the sequence  $\{Y_r\}_{r=0}^\infty$  of beliefs spaces by:  $Y_0 = K$  and for  $r = 1, 2, \dots$ ,

$$Y_r = \{ \omega_r \in Y_{r-1} \times [\Delta(Y_{r-1})]^{\mathbf{I}} \mid \text{the two conditions below are satisfied } \}.$$

- For all  $i \in \mathbf{I}$ , the marginal distribution of  $\theta_r^i(\omega_r)$  on the  $i^{\text{th}}$  copy of  $\Delta(Y_{r-1})$  is the Dirac mass at  $\theta_{r-1}^i(\rho_{r-1}(\omega_r))$ .
- For all  $i \in \mathbf{I}$ , the marginal distribution of  $\theta_r^i(\omega_r)$  on  $Y_{r-2}$  is  $\theta_{r-2}^i(\rho_{r-1}(\omega_r))$ . Here  $\rho_{r-1}$  and  $\theta_r^i$  are the projections from  $Y_r$  to  $Y_{r-1}$  and the  $i^{\text{th}}$  copy of  $\Delta(Y_{r-1})$  respectively.

Note that 1 imposes the condition that each player ‘knows’ his own beliefs while 2 is the condition that each level of beliefs is compatible with the lower levels. The fact that  $\theta_r^i$  is a mapping to  $\Delta(Y_{r-1})$  implies that it is common knowledge that 1 and 2 are satisfied.

- Define the universal beliefs space  $\Omega$  to be the projective limit of  $\{Y_r\}_{r=0}^\infty$  with respect to the projections  $\rho_{r-1}: Y_r \rightarrow Y_{r-1}$ .

c. Prove that  $\rho_r(Y_{r+1}) = Y_r$  for all  $r = 0, 1, 2, \dots$ . Hence  $\rho_r(\Omega) = Y_r$  for all  $r = 0, 1, 2, \dots$  and in particular  $\Omega \neq \emptyset$ .

d. Prove that each  $\omega \in \Omega$  and  $i \in \mathbf{I}$  uniquely determine a probability measure  $\bar{\theta}_\omega^i$  in  $\Delta(\Omega)$  and the mapping  $\bar{\theta}^i: \Omega \rightarrow \Delta(\Omega)$  is continuous.

e. Prove that as a consequence of conditions 1 and 2, for any  $\omega \in \Omega$  and any  $i \in \mathbf{I}$ ;

$$\text{if } \tilde{\omega} \in \text{Supp}(\bar{\theta}_\omega^i) \text{ then } \bar{\theta}_{\tilde{\omega}}^i = \bar{\theta}_\omega^i.$$

f. Let  $\Theta^i = \bar{\theta}^i(\Omega) \subseteq \Delta(\Omega)$  and prove that:

- (1) The space  $\Omega$  is homeomorphic to  $K \times \prod_{i \in \mathbf{I}} \Theta^i$ .
- (2) The space  $\Theta^i$  is homeomorphic to  $\Delta(K \times \prod_{j \neq i} \Theta^j)$ .
- (3) Property (P) of theorem 1.1 p. 107 is satisfied.

### 3. Universal *BL*-space in games with private information.

a. Show that the formalism introduced in sect. 1 p. 107 can also be used in situations in which players have some exogenous private information (i.e. a  $\sigma$ -field  $\mathcal{K}^i$  on  $K$ ).

HINT. Construct  $\Omega$  as in sect. 1 and define  $\Omega_0$  as follows:

$$\Omega_0 = \{\omega \mid \bar{\theta}_\omega^i[X] = 1 \forall X: k(\omega) \in X \in \mathcal{K}^i, \forall i \in \mathbf{I}\}$$

Use example 7 p. 119 to obtain  $\Omega_\infty$  as the universal *BL*-space for this game with private information.

COMMENT 4.20. The above is just a definition. To be useful, it should be accompanied by theorems — a.o. representation theorems like theorem 2.5 —. And those may require some additional assumptions (e.g. universal measurability of  $\mathcal{K}^i$ , topological assumptions . . . ).

b. (Böge and Eisele, 1979);(Aumann, 1974), (Aumann, 1985) Consider the situation in which  $K = K_0 \times U \times S$  (where  $K_0$  is the space of states,  $S$  is the action space and  $U$  is the set of utility functions —  $\mathbf{I}$ -tuples of real valued functions on  $K_0 \times S$ ). For each player let  $\mathcal{K}^i$  be the partition of  $K$  generated by the projection on his own utility and action space (and their Borel  $\sigma$ -field). One can then apply the above construction to obtain the appropriate universal *BL*-space.

Next consider only *BL*-subspaces where each player maximises his expected utility and apply again a procedure like that in IIIEx.3a to construct a universal such *BL*-subspace. Every *BL*-subspace of this universal space then describes a game with incomplete information together with one of its equilibria. If  $K_0$  and  $U$  are singletons, one obtains Aumann's Bayesian interpretation of correlated equilibria. (cf. remark 3.3 p. 90).

**4. Lower semi-continuity of pay-off in terms of information.** (Forges, 1988b) The following illustrates the decision-theoretic relevance of the weak\* topology on information that was used in this chapter (cf. also, in this vein, prop. 4.4 p. 133):

Let  $K$  be a finite set,  $U$  a separable metric space, and  $f: U \times \Delta(K) \rightarrow \overline{\mathbb{R}}_+$  be lower semi-continuous, and convex on  $\Delta(K)$  for each  $u \in U$ .

Let  $\phi(P) = \int_U f(u, [P(k \mid u)]_{k \in K}) P(du)$  for  $P \in \Delta(K \times U)$ .

Then  $\phi$  is convex and lower semi-continuous, and  $\{(P, \phi(P)) \mid P \text{ has finite support}\}$  is dense in the graph of  $\phi$ .

COMMENT 4.21. Regularity of  $P$  plays no rôle — the theorem remains true over the set of all Borel probability measures.

COMMENT 4.22. A decision-maker who observes  $u$  before making a decision  $d$ , whose pay-off  $\alpha_d^k(u)$  depends on the unknown state of nature  $(k, u)$ , will obtain  $f(u, P(k \mid u)) = \sup_d \sum_k P(k \mid u) \alpha_d^k(u)$ , a convex, continuous function  $f$  of  $P(k \mid u)$ , depending also on  $u$ . Hence his (ex ante) expected pay-off is  $\phi(P)$ . Lower semi-continuity expresses the fact that there can be a loss of information (and never a gain) in going to the limit on an observed random variable — e.g., let

$u_n(\omega) = [u_1(\omega)]/n$ :  $u_\infty$  is zero, and loses all the information. An application will be found in ch. IX sect. 3 p. 408.

COMMENT 4.23. Assume the function  $f$  is extended by homogeneity of degree 1 from  $\Delta(K)$  to  $\mathbb{R}_+^K$ . Then

$$\phi(P) = \int_U f(u, (P(k, du))_{k \in K})$$

in the sense that for any measure  $\lambda$  on  $U$  such that all measures  $P_k(du) = P(k, du)$  ( $k \in K$ ) are absolutely continuous w.r.t.  $\lambda$ ,

$$\phi(P) = \int_U f[u, (h_k(u))_{k \in K}] \lambda(du)$$

where  $h_k(u)$  is the Radon-Nikodym derivative of  $P(k, du)$  w.r.t.  $\lambda$  — thus justifying the notation  $\phi(P) = \int_U f[u, (h_k(u)\lambda(du))_{k \in K}] = \int_U f[u, (P(k, du))_{k \in K}]$  (cf. e.g. Edwards, 1965, IV.15.11).

HINT. Prove first that the integral does not depend on  $\lambda$  ( $g_u$  also is positively homogeneous of degree 1). Next set  $\lambda$  equal to the marginal distribution of  $u$ .

a. *Reduction to the case where  $U$  is compact metric,  $f$  is independent of  $u$ , and a maximum of finitely many linear functions on  $\mathbb{R}_+^K$ .*

HINT. Let  $(\overline{U}, d)$  be a compact metric space containing  $U$ . For fixed  $\tilde{u} \in U$ , the separation theorem yields that  $f(\tilde{u}, \cdot)$  is the supremum of the linear functions on  $\Delta(K)$  that are strictly smaller — and those can further be chosen with rational coefficients. Let  $\varphi$  be such a linear function, and use the compactness of  $\Delta(K)$  and the lower semi-continuity of  $f$  to show that, for some  $\varepsilon_0 > 0$ , one still has  $f(u, x) > \varphi(x) \forall x \in \Delta(K)$  whenever  $d(u, \tilde{u}) < \varepsilon_0$ . If  $u_n$  is a dense sequence in  $U$ , and  $\varepsilon$  is rational with  $\frac{1}{2}\varepsilon_0 < \varepsilon < \varepsilon_0$ , and one chooses  $n_0$  such that  $d(u_{n_0}, \tilde{u}) < \frac{\varepsilon}{k+1}$ , then  $[1 - \frac{k+1}{k\varepsilon}d(u, u_{n_0})]^+[\varphi(x)]^+ = F_{k, n_0, \varepsilon, \varphi}(u, x)$  is everywhere  $\leq f$ , and at  $\tilde{u}$  it is  $\geq (1 - \frac{1}{k})[\varphi(x)]^+$ . Hence  $f$  is the supremum of the set of all functions  $F_{k, n, \varepsilon, \varphi}$  ( $k$  and  $n$  integers,  $\varepsilon$  rational,  $\varphi$  with rational coefficients) that lie everywhere below. Let  $F_i$  enumerate this set of functions: the sequence  $f_j = \max_{i \leq j} F_i$  increases to  $f$ , and consists of bounded, Lipschitz functions on  $\overline{U} \times \Delta(K)$  which are, for every  $u \in \overline{U}$ , piecewise linear and convex on  $\Delta(K)$ . Hence (monotone convergence theorem) it suffices to prove the property for such functions. In particular we can henceforth assume  $U$  is compact metric.

Fix now  $\varepsilon > 0$ , and for  $u \in U$  let  $O_u$  be an open neighbourhood such that  $u' \in O_u \implies |f(u', x) - f(u, x)| < \varepsilon \forall x \in \Delta(K)$ . Let  $u_i$  ( $i = 1, \dots, n$ ) be a finite subset such that the  $O_{u_i}$  cover  $U$ , and let  $g$  be a corresponding continuous partition of unity — i.e.  $g_i$  is continuous, with values in  $[0, 1]$ , vanishes outside of  $O_i$  and  $\sum_{i=1}^n g_i(u) = 1, \forall u \in U$ .

Then  $h(u, x) = \sum_{i=1}^n g_i(u)f(u_i, x)$  is uniformly  $\varepsilon$ -close to  $f$ , so it suffices to prove the result for  $h$ . Hence it suffices to prove it for one function  $g_i(u)f(u_i, x)$ : we can assume our function has the form  $g(u)h(x)$ , where  $g: U \rightarrow \mathbb{R}_+$  is continuous, and  $h: \Delta(K) \rightarrow \mathbb{R}_+$  is a maximum of finitely many linear functions.

Define now a map  $P \mapsto \tilde{P}$  from  $\Delta(K \times U)$  to the space  $M$  of non-negative, bounded measures on  $K \times U$ , by  $\tilde{P}(k \times B) = \int_B g(u)P(k, du)$  for every Borel set  $B$  in  $U$ . The map is clearly linear and continuous. Further  $\tilde{P}(k \mid u) = P(k \mid u)$  a.e., so we get  $\phi(P) = \int h([\tilde{P}(k \mid u)]_{k \in K})\tilde{P}(du)$ . Since  $P \mapsto \tilde{P}$  is linear and continuous, it suffices therefore to prove lower semi-continuity and convexity on  $M$  of the map  $P \mapsto \int h([P(k \mid u)]_{k \in K})P(du)$ . The map being positively homogeneous of degree one, it suffices for this to prove lower semi-continuity and convexity on  $\Delta(K \times U)$ .

b. *Lower semi-continuity: Reduction to the case where furthermore the marginal of  $P$  on  $K$  is fixed.*

HINT. Since  $\Delta(K \times U)$  is metrisable, consider a sequence  $P_n \rightarrow P$  we want to show that  $\liminf \phi(P_n) \geq \phi(P)$ . Let  $p(k)$  (resp  $p_n(k)$ ) be the marginals on  $K$ : then  $p_n \rightarrow p$ . So, for  $\varepsilon > 0$ , choose  $n_0$  such that  $(1 + \varepsilon)p_n(k) \geq p(k) \forall k, \forall n \geq n_0$ . Extracting a subsequence, we can assume this holds for all  $n$ , and also that  $\{k \mid p_n(k) > 0\}$  is independent of  $n$  — hence we can assume it is the whole of  $K$ :  $p_n(k) > 0 \forall n, \forall k$ . Then  $P_n(k, du)$  can be written as  $p_n(k)g_{n,k}(u)P_n(du)$ . If  $p(k) > 0$ , then  $g_{n,k}(u)P_n(du) \rightarrow P(du \mid k)$  weakly. Assume that everywhere  $g_{n,k} \geq 0$ ,  $\sum_k p_n(k)g_{n,k}(u) = 1$  — so  $P_n(k \mid u) = p_n(k)g_{n,k}(u)$  and  $\phi(P_n) = \int f[(p_n(k)g_{n,k}(u))]_{k \in K}P_n(du)$ . Let  $\overline{P}_n(k, du) = p(k)g_{n,k}(u)P_n(du)$ :

$\bar{P}_n \rightarrow P$  weakly, and  $\bar{P}_n(du) = h_n(u)P_n(du)$  with  $h_n(u) = \sum_k p_k g_{n,k}(u)$  — so  $\bar{P}_n(k | u) = p_k \bar{g}_{n,k}(u)$  with  $\bar{g}_{n,k}(u) = g_{n,k}(u)/h_n(u)$ . Thus  $\phi(\bar{P}_n) = \int f[(p(k)\bar{g}_{n,k}(u))_{k \in K}]h_n(u)P_n(du)$ .

Now  $(1 + \varepsilon)p_n \geq p$  yields  $h_n(u) \leq 1 + \varepsilon$ . Since  $\int h_n(u)P_n(du) = 1$ , this yields that  $h_n(u) \geq 1 - \sqrt{\varepsilon}$  with  $P_n(du)$ -probability at least  $1 - \sqrt{\varepsilon}$ . In that case, we get  $(1 + \varepsilon)p_n(k)g_{n,k}(u) \geq p(k)g_{n,k}(u) \geq (1 - \sqrt{\varepsilon})p(k)\bar{g}_{n,k}(u)$ . Since both points belong to  $\Delta(K)$ , the arguments of  $f$  in the expressions of  $\phi(P_n)$  and  $\phi(\bar{P}_n)$  differ by less than  $2\sqrt{\varepsilon}$  in norm. Using the Lipschitz character (with constant  $L$ ) of  $f$ , and again that  $h_n \leq 1 + \varepsilon$ , we obtain that  $(1 + \varepsilon)[\phi(P_n) + 2L\sqrt{\varepsilon}] \geq \phi(\bar{P}_n)$ . So it suffices to prove that  $\liminf \phi(\bar{P}_n) \geq \phi(P)$ .

c. *Use of the Dudley-Skhorod Theorem.*

HINT. Since the marginal  $p(k)$  on  $K$  is fixed — and can be assumed strictly positive — we know that  $P_n(du | k) \rightarrow P(du | k)$  weakly, for each  $k$ . Apply thus the Dudley Skhorod theorem for each of those conditional distributions to construct, a probability space  $(\Omega, A, Q)$  together with a random variable  $k(\omega)$  to  $K$  and a sequence  $u_n(\omega)$  converging a.e. to  $u(\omega)$  such that  $P_n$  is the distribution of  $(k(\omega), u_n(\omega))$  and  $P$  that of  $(k(\omega), u(\omega))$ . Let  $J_k = \{\omega | k(\omega) = k\}$ .

Thus we want to show that, if  $u_n(\omega) \rightarrow u_\infty(\omega)$  a.s., and  $J_k$  is a finite measurable partition, then

$$\liminf_{n \rightarrow \infty} \mathbb{E} f([Q(J_k | u_n(\omega))]_{k \in K}) \geq \mathbb{E} f([Q(J_k | u_\infty(\omega))]_{k \in K}).$$

d. *Lemma.* Let  $(\Omega, \mathcal{A}, Q)$  be a probability space, with a finite measurable partition  $(J_k)_{k \in K}$ , and a sequence of random variables  $u_n(\omega) \rightarrow u_\infty(\omega)$  with values in a separable metric space  $U$ . Let  $q_n^k = Q(\mathbf{1}_k | u_n)$ . Then:

- (1) weak $^*$  limits  $q$  of the  $q_n$  exist and satisfy  $\mathbb{E}(q | u_\infty) = q_\infty$
- (2) for any weak $^*$  limit  $q$  of the  $q_n$  there exists a sequence of convex combinations  $r_i = \sum_n \alpha_n^i q_n$ , with  $\alpha_n^i = 0$  for  $i$  sufficiently large, and  $r_i \rightarrow q$  a.e. — hence  $\mathbb{E}(r_i | u_\infty) \rightarrow \mathbb{E}(q | u_\infty) = q_\infty$  a.s.

HINT. The existence follows from Banach-Alaoglu. For 2: use ex. I.2Ex.12 p. 24, Egorov's theorem, and (Kelley et al., 1963, p. 212). For 1, consider  $f \in C(U)$  — then  $f(u_n)$  is uniformly bounded and converges a.e. to  $f(u_\infty)$  — so the convergence is uniform on weakly compact subsets of  $L_1$  — in particular, extracting first a weakly (or weak $^*$ )-convergent subsequence from the  $q_n$ ,  $\mathbb{E}(q^k f(u_\infty)) = \lim_{n \rightarrow \infty} \mathbb{E}(q_n^k f(u_n)) = \lim_{n \rightarrow \infty} \mathbb{E}[\mathbb{1}_{J_k} f(u_n)] = \mathbb{E}(\mathbb{1}_{J_k} f(u_\infty))$ . Extend this equation to positive Borel  $f$ , and conclude.

e. *Lower semi-continuity.*

HINT. By IIIEx.4a, IIIEx.4b, IIIEx.4c and the notations of IIIEx.4d, it suffices (taking a subsequence) to show that

$$\begin{aligned} \mathbb{E} f(q_\infty) &= \lim_{i \rightarrow \infty} \mathbb{E} f[\mathbb{E}(r_i | u_\infty)] \\ &\leq \limsup_{n \rightarrow \infty} \mathbb{E} f(E(q_n | u_\infty)) \\ &\leq \limsup_{n \rightarrow \infty} \mathbb{E} f(q_n), \end{aligned}$$

which follows from IIIEx.4d and the convexity of  $f$  (ex. I.3Ex.14 p. 37).

f. *Density.*

HINT. Make a Borel partition of  $U \times \Delta(K)$  into finitely many Borel sets of diameter  $< \varepsilon$  each, map this partition back into  $U$ , and map the whole mass of each partition element to some point of that element. This yields measures with finite support  $P_\varepsilon$  that converge weak $^*$  to  $P$  and such that  $\phi(P_\varepsilon) \rightarrow \phi(P)$ .

g. *Convexity.*

HINT. Prove first (using Remark 4.23 p. 143) convexity over the set of probabilities with finite support: this boils down to the convexity of  $f$ . Consider now  $P$  (and similarly  $Q$ ) arbitrary. Apply then IIIEx.4f and use the lower semi-continuity.

REMARK 4.24. We shall use later the following result in the specific case where  $U$  is the set of continuous convex functions from  $\Delta(K)$  to  $\mathbb{R}_+$ , with  $f(u, \cdot) = u(\cdot)$ .

**5. Lower semi-continuity (continued).** In the same situation as above, decompose measures  $P$  on  $K \times U$  into their marginal  $p$  on  $K$  and a conditional  $\tau_k(du)$  on  $U$  given  $K$ . Let  $\psi(\tau): p \mapsto \phi(p \otimes \tau)$ .

Then  $\psi$  is a convex, l.s.c. map with values in the set  $\tilde{C}$  of l.s.c. convex maps from  $\Delta(K)$  to  $\overline{\mathbb{R}}_+$  (l.s.c. of  $\psi$  means:  $\forall g$  continuous on  $\Delta(K)$ ,  $\{\tau \mid \psi(\tau) > g\}$  is weak $^*$ -open).

HINT. Convexity follows from IIIEx.4. For the lower semi-continuity, use compactness of  $\Delta(K)$ .



## CHAPTER IV

### General Model of Repeated Games

#### 1. The model

In this chapter we introduce formally the general model of repeated games.

We start with a non-cooperative game  $G$  and define a new game  $\Gamma_\infty$  a play of which is an infinite sequence of plays of  $G$ .

In fact it appears in many applications that current moves not only influence the current pay-off but also the future play hence some state variable of the model. This is the reason why stochastic games appear in a natural way.

Moreover we have to describe the information available to the players. There may be some differences between their initial knowledge of the characteristics: initial state, preferences, even transition law. This is taken into account in the framework of games with incomplete information.

Finally it is necessary for a full description of the game to specify what additional information is transmitted to the players after each stage of the play. It is easy to see that assuming the knowledge of the other players' strategies is unrealistic and at most the actual moves may be observed. But even, it may be useful to look at the case where only the individual player's pay-off is known to him, or even less, hence no full monitoring of the previous moves, or even outcomes or own pay-offs is possible. This leads to the notion of signals that may depend in a random way on the actual moves and state.

To integrate all such effects it is sufficient to define a state and move dependent lottery that selects at every stage the signals for the players, their pay-offs and the next state of nature. If one wants in addition to incorporate the effect of information lags this transition may also depend on the past events. In fact we will see that this quite huge construction can be reduced to a simple and convenient form (cf. prop. 2.3 p. 156).

It follows from the above presentation that this model is an adequate description of a stationary multistage game in the sense that its formulation is time shift invariant (adding new states and pay-offs if necessary) and needs only some counter, to let the stage of the game be known to the players.

We will give more formally a first model of the game and introduce explicitly the main definitions.

**DEFINITION 1.1.** A **repeated game** is a finite multistage game, where for  $\omega \in \Omega^\infty$ ,  $g^i(\omega) = \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{k \leq n} h^i(\omega_k)$ , for some function  $h^i$  from  $\Omega$  to  $\mathbb{R}$  (cf. 1.a p. 51 and 3.b p. 90).

Here "finite" means all sets  $(\Omega, A, S)$  are finite. The assumption on the pay-off function gives also a finite description for the latter, the set  $\Omega$  being finite. There only remains the need to be a bit more explicit about the solution concepts, since the pay-off function  $g^i$  is not always everywhere defined.

In fact we will consider here an equivalent and more convenient model where  $q$  is independent of  $A$  (1.a.4 p. 52) and  $h$  is defined on  $\Omega \times S$ . Other equivalent representations will be studied in sect. 2 p. 153.

**1.a. States, signals and transitions.** We are given a finite set of **states**  $K$  and a finite set of players  $\mathbf{I}$  (there will be no confusion to write also sometimes  $\mathbf{I}$  for its cardinality). For each player  $i$  in  $\mathbf{I}$ , let  $S^i$  be its finite set of **moves** (or **actions**) and  $A^i$  its finite set of **signals**. ( $K$  being finite, both sets can be assumed to be independent of the state  $k$  in  $K$ , duplicating eventually moves or signals). Denote as usual by  $S$  and  $A$  respectively, the products of the  $S^i$ , resp.  $A^i$ , over  $\mathbf{I}$ .

$P$  is an (initial) probability on  $K \times A$  and  $Q$  is a transition probability from  $K \times S$  to  $W \times K \times A$  where  $W$  is some compact set in  $\mathbb{R}^{\mathbf{I}}$  (pay-offs are uniformly bounded). Let  $G$  be the marginal distribution induced by  $Q$  on  $W$  so that  $G^k(s^1, \dots, s^{\mathbf{I}})$  denotes the distribution of the vector pay-off in state  $k$  given the vector of moves  $s$ . Similarly the marginal distribution of  $Q$  on  $K$  corresponds to the transition on the state space and the marginal distribution of  $Q$  on  $A$  determines the distribution of the signals.

The **repeated game**  $\Gamma_\infty$  is played as follows:

- At stage 0 a point  $(k_1, a_1)$  is chosen in  $K \times A$  according to  $P$  and  $a_1^i$  is announced to player  $i$ , for each  $i$ .
- At each stage  $n$ ,  $n \geq 1$ , each player  $i$  chooses independently a move  $s_n^i$  in  $S^i$ . The distribution  $Q(k_n, s_n)$  is used to choose a point  $(g_n, k_{n+1}, a_{n+1})$  in  $W \times K \times A$ . The new state is  $k_{n+1}$ . The signal  $a_{n+1}^i$  is announced to player  $i$  and  $g_n^i$  is his pay-off, for all  $i$  in  $\mathbf{I}$ . (Remark that the pay-off is not explicitly observed, it may be deducible from the signal).

The above description is known to the players. To complete explicitly the information structure of the game we still have to assume that each player remembers all information he receives in the past (effectively perfect recall).  $\Gamma_\infty$  is thus a game in extensive form as defined in ch. II p. 51. (To obtain explicitly the same description, divide each stage into  $\mathbf{I}$  substages and let the players play in order, extending the transition on states, signals and pay-offs in the obvious way). The specificity of this game is its stationary aspect, both of the transitions and of the pay-offs (cf. 1.b p. 52).

A **play** of  $\Gamma_\infty$  is then identified with an infinite sequence  $(k_1, a_1, s_1, g_1, k_2, a_2, \dots, s_n, g_n, k_{n+1}, a_{n+1}, \dots)$ . As before we denote by  $H_\infty$  the set of plays. An initial sequence of length  $n$  of a play, i.e. ending by  $(\dots, k_n, a_n)$ , will be called an  $n$ -history and  $H_n$  will denote the corresponding set. The set of all **histories** (or positions) is  $H = \bigcup_{n \geq 1} H_n$ .

If on some play and for some  $n$  the sequence  $\ell^i = (a_1^i, \dots, a_n^i)$  determines the sequence  $m^{-i} = (a_1^{-i}, s_1^{-i}, \dots, s_{n-1}^{-i}, a_n^{-i})$ , for all  $i$ , we define the **subgame** from position  $\ell = (a_1, \dots, a_n)$  like the game  $\Gamma_\infty$  (i.e. same moves, signals, transitions and pay-offs) but starting with an initial distribution on  $K \times A^n$  defined by  $\ell$  and the law of  $k_n$  given  $\ell$  (which is known by all players). Finally given  $h_\infty$  in  $H_\infty$ , we write  $h_n$  for its projection on  $H_n$ ,  $\mathcal{H}_n$  is the  $\sigma$ -algebra on  $H_\infty$  generated by  $H_n$  and  $\mathcal{H}_\infty = \bigvee_{n \geq 1} \mathcal{H}_n$  (product  $\sigma$ -algebra).  $\mathcal{H}$  will denote the induced  $\sigma$ -algebra on  $H$ .

We will also write  $g$  for the stream of pay-offs  $(g_1, \dots, g_n, \dots)$ .

**1.b. Strategies and pay-offs.** By the above description the information of player  $i$  before stage  $n$  is a vector  $(a_1^i, \dots, a_n^i)$  in  $\prod_{m=1}^n (A^i)_m$ .

Denote by  $\mathcal{H}_n^i$  the  $\sigma$ -field induced by this set on  $H_\infty$ . The restriction to each  $H_n$  defines a measurable structure on  $H$ , written  $\mathcal{H}^i$ , that describes  $i$ 's **information partition** on  $H$ .

A **pure strategy**  $s^i$  is thus a measurable mapping from  $(H, \mathcal{H}^i)$  into  $S^i$ .

(Note that the pure strategy set of player  $i$  is a product of finite sets.)

A **mixed strategy**  $\sigma^i$  is then as usual a probability on this compact set (with the product topology induced by the discrete topology on each factor), and the corresponding set will be denoted by  $\Sigma^i$ .

One can always add player  $i$ 's last move to his signal (theorem 1.3 p. 53) and then perfect recall implies (theorem 1.8 p. 55) that we can equivalently view  $\sigma^i$  as a mapping from  $(H, \mathcal{H}^i)$  into  $\Delta(S^i) = X^i$ . According to this interpretation it is also useful to think of  $\sigma^i$  as a sequence  $(\sigma_n^i)_{n \geq 1}$  where  $\sigma_n^i$  is the restriction of  $\sigma^i$  to  $H_n$  and corresponds to the "strategy at stage  $n$ ". One can as well consider each  $\sigma_n^i$  as being defined on  $H_\infty$  with  $\sigma_n^i(h_\infty) = \sigma^i(h_n)$ .

As pay-off function for the game we will consider the Cesàro limit of the sequence of stage pay-offs  $g_n$ . This may not be defined for every feasible play (i.e. the game is not a well specified game in normal form) but we will take care of this fact when defining the solutions and show that this specification of the pay-offs is unambiguous and sufficient.

Note that this model can also be used to study two other important classes of games namely:

**Discounted games:**  $\Gamma_\lambda$  has pay-off function  $\bar{g}_\lambda = \lambda \sum_{n=1}^{\infty} (1 - \lambda)^{n-1} g_n$  with  $\lambda \in (0, 1]$ .

In order to reduce this game to the previous model, add for each point in  $K \times S$  a new absorbing state with pay-off  $E(G^k(s))$  forever. Now the new transition will operate in two stages, the first one being like the old transition and the second one choosing with probability  $\lambda$  the absorbing state corresponding to  $(k, s)$  and with probability  $(1 - \lambda)$  keeping the same state and giving a zero pay-off. Remark that one could also work with different discount rates for each player.

**Finite games:**  $\Gamma_N$  has pay-off function  $\bar{g}_N = 1/N \sum_{n=1}^N g_n$ ,  $N \in \mathbb{N}$ .

In this case a reduction similar to the previous one can be done, after replacing  $K$  by its product by  $\{1, \dots, N\}$  and using "time dependent parameters"  $\lambda_n = 1/(N - n + 1)$ .

Basically the reason for using these averages is that they preserve the stationary character of the game (cf. ex. IV.1Ex.1 p. 152).

**1.c. Zero-sum case.** For solution concepts, in the two-person, zero-sum case, an unambiguous solution is provided by the value, together with a corresponding pair of  $\varepsilon$ -optimal strategies.

Before defining these concepts in our framework, let us introduce some notational principles. In general, in the two-person case we will write:  $s, \sigma, a$ , (respectively  $t, \tau, b$ ) for the moves, strategies and signals of player I (the maximiser), (respectively player II), and  $S, \Sigma, A$ , (respectively  $T, \mathcal{T}, B$ ) for the corresponding sets.

Given  $\sigma$  and  $\tau$ ,  $P_{\sigma, \tau}$  is the probability induced on  $(H_\infty, \mathcal{H}_\infty)$  by  $P, Q, \sigma, \tau$ , (cf. 1.6 p. 54) and  $E_{\sigma, \tau}$  is the corresponding expectation.

**DEFINITION 1.2.** Player I can guarantee  $d$  if:

$$(i) \quad \forall \varepsilon > 0, \exists \sigma_\varepsilon, \exists N, \text{ such that : } E_{\sigma_\varepsilon, \tau}(\bar{g}_n) \geq d - \varepsilon, \quad \forall \tau, \forall n \geq N$$

Player II can defend  $d$  if:

$$(ii) \quad \forall \varepsilon > 0, \forall \sigma, \exists \tau, \exists N, \text{ such that : } E_{\sigma, \tau}(\bar{g}_n) \leq d + \varepsilon, \quad \forall n \geq N$$

$\underline{v}$  is the **maxmin** of  $\Gamma_\infty$  if player I can guarantee  $\underline{v}$  and player II can defend  $\underline{v}$ .

In this case a strategy  $\sigma_\varepsilon$  associated to  $\underline{v}$  is  **$\varepsilon$ -optimal**. A strategy is **optimal** if it is  $\varepsilon$ -optimal for all  $\varepsilon$ .

The **minmax**  $\bar{v}$  and ( $\varepsilon$ -) optimal strategies for player II are defined in a dual way.

The game has a **value**, denoted by  $v_\infty$ , iff  $\underline{v} = \bar{v}$  with then  $v_\infty = \underline{v} = \bar{v}$ .

REMARK 1.1. Whenever possible, for example in stochastic games (cf. ch. VII p. 327), one may also require (i) and its dual for  $N = \infty$ , namely:

$$(iii) \quad E_{\sigma_\varepsilon, \tau}(\liminf_{n \rightarrow \infty} \bar{g}_n) \geq \underline{v} - \varepsilon, \quad \forall \tau$$

$$(iv) \quad E_{\sigma, \tau_\varepsilon}(\limsup_{n \rightarrow \infty} \bar{g}_n) \leq \underline{v} + \varepsilon, \quad \forall \sigma$$

REMARK 1.2. Note that with the above definitions the existence of  $\underline{v}$  and  $\bar{v}$  has to be proved. Remark also that the definitions provide insights in the study of the long but finite games  $\Gamma_n$  or of the games with small discount factor  $\Gamma_\lambda$ . For the first class it is clear that  $\sigma_\varepsilon$  guarantees a pay-off  $\underline{v} - \varepsilon$  in any  $\Gamma_n$  with  $n \geq N$ . As for the second it follows clearly from (i) that for all discount factor  $\lambda$  smaller than some  $\bar{\lambda}$ , function of  $N$  and the pay-off range  $W$  only, one has:  $\forall \tau, E_{\sigma_\varepsilon, \tau}(\bar{g}_\lambda) \geq \underline{v} - \varepsilon$  as well. So that  $\sigma_\varepsilon$  guarantees also  $\underline{v} - \varepsilon$  in any  $\Gamma_\lambda$  with  $\lambda \leq \bar{\lambda}$ . Moreover in both cases, by (ii),  $\underline{v}$  is the best that can be achieved by strategies that do not depend on the exact specification of the duration or of the discount factor of the game.

REMARK 1.3. Condition (i) is a uniform (in  $\tau$ ) property on the liminf of the average expected pay-off. On the other hand (iii) corresponds to a pay-off function defined on plays and would be a desirable property.

The following game shows that it may not hold.

EXAMPLE 1.4. Consider a zero-sum game with 2 states and pay-offs matrices  $G^1 = (1, 0)$ ,  $G^2 = (0, 1)$  (player I is a dummy player). Assume  $\Pr\{k_n = 1, \forall n\} = \Pr\{k_n = 2, \forall n\} = 1/2$ . Obviously  $E_{\sigma, \tau}(\bar{g}_n) = 1/2$ , for all  $\sigma, \tau$ , but if player II plays left and right in alternating blocs of size  $L_m$  with  $L_m/L_{m-1} \rightarrow \infty$ , then  $E_{\sigma, \tau}(\liminf \bar{g}_n) = 0$ .

Another approach would then be to define a pay-off for pure strategy pairs, (i.e. taking expectation with respect to all random parameters of the game or all corresponding plays), by associating to the sequence of stage pay-offs some limit, say liminf (hence obtaining 1/2 in the previous example). In this case too, the normal form game may have no value satisfying (i) to (iv):

EXAMPLE 1.5. Consider the same model as above with now:

$$G^1 = \begin{pmatrix} 2 & 0 \end{pmatrix} \quad G^2 = \begin{pmatrix} 0 & 1 \\ -1 & 2 \end{pmatrix}$$

and player I knows the state. The moves of player II are not announced to I and the signals to II are  $a$  or  $b$  according to the following probability distributions:

$$Q^1 = ((1/3, 2/3) \quad (1/3, 2/3)) \quad Q^2 = \begin{pmatrix} (0, 1) & (0, 1) \\ (2/3, 1/3) & (2/3, 1/3) \end{pmatrix}$$

(e.g., if  $k = 2$ , I plays B and II L,  $a$  will be announced to II with probability  $2/3$ ). Note that by playing  $\sigma_0 = (1/2, 1/2)$  i.i.d. if  $k = 2$ , player I obtains a constant expected pay-off  $3/4$ , and this is the only way “not to reveal information” about the state.

Considering the measure  $\mu_\sigma$  induced by any strategy  $\sigma$  of I on  $\mathcal{H}^{\text{II}}$  ( $\sigma$ -field of information of II), one then obtains that: if  $\sigma$  guarantee  $3/4 - \delta$  then  $\|\mu_\sigma - \mu_{\sigma_0}\| \leq (16/3)\delta$  (cf. ex. VEx.13 p. 262). Assume now  $\delta$  small.

Let  $N_n = 2^{2^n}$  and define a strategy  $\tau$  of II as follows:

Play arbitrarily at stage 1. If the first signal is  $a$ , play  $N_1$  times R, then  $N_2$  times L and so on, and the reverse if  $b$  is announced at stage 1.

Since the frequency of the signal  $a$  under  $\mu_\sigma$  is near  $1/3$  (being near the one under  $\mu_{\sigma_0}$ ), the frequency of B under  $\sigma$ , if  $k = 2$ , is near  $1/2$ ; hence denoting by  $f_m^1$  and  $f_m^2$  the frequencies of L until the end of bloc  $m$ , and by  $\varphi_m$  the corresponding average expected pay-off, one obtains that  $\varphi_m$  is near  $3/4 + E(f_m^1 - f_m^2)$  for  $m$  large enough and almost all pure strategies in the support of  $\sigma$ .

Note that  $E(f_m^1)$  is near  $2/3$  if  $m$  is odd and near  $1/3$  if  $m$  is even. As for the expectation of  $f_m^2$  we obtain if the first move is T, 1 if  $m$  is odd and 0 if even, and similarly after B  $1/3$  (odd) and  $2/3$  (even). It follows that in each case  $\liminf \varphi_m$  is  $3/4 - 1/3$ , hence  $E_{\sigma,\tau}(\liminf \bar{g}_n)$  is near  $5/12$  for such strategies. (See ex. VEx.13 p. 262 for precise computations).

Remaining with the finite games  $\Gamma_n$  or the discounted games  $\Gamma_\lambda$ , it is clear by prop. 2.6 p. 17 that they possess a value, denoted by  $v_n$  or  $v_\lambda$ , since our finiteness assumptions imply that the pay-offs  $\bar{g}_n(\sigma, \tau) = E_{\sigma,\tau}(\bar{g}_n)$  and  $\bar{g}_\lambda(\sigma, \tau) = E_{\sigma,\tau}(\bar{g}_\lambda)$  are continuous on the product of the compact pure strategies spaces.

It follows from the previous definitions that the existence of  $v_\infty$  also implies that both  $\lim_{n \rightarrow \infty} v_n$  and  $\lim_{\lambda \rightarrow 0} v_\lambda$  exist and satisfy:  $v_\infty = \lim_{n \rightarrow \infty} v_n = \lim_{\lambda \rightarrow 0} v_\lambda$ . In this situation,  $\varepsilon$ -optimal strategies are also  $\varepsilon$ -optimal in any sufficiently long game as well as in any game with sufficiently small discount factor.

Now when  $\underline{v} < \bar{v}$  (recall that the definitions clearly imply  $\underline{v} \leq \bar{v}$ ) another main question of research will be the study of the nets  $v_n$  and  $v_\lambda$  (as  $n \rightarrow \infty$  and  $\lambda \rightarrow 0$ ) and of the asymptotic behaviour of the corresponding optimal strategies.

Before going to the non-zero-sum case, let us recall that the study of  $\bar{v}$  is also of first interest in this case because it defines the individually rational level, i.e. what player I can obtain in the worst situation where the other players are considered as a single player trying to minimise his pay-off (cf. sect. 4 p. 162).

**1.d. Non-zero-sum case.** In the non-zero-sum case we will be mainly interested in equilibria and correlated or communication equilibria (cf. sect. I.4 p. 39 and II.3 p. 88).

Recall that the latter form a larger set but present some conceptual advantages and have a much nicer mathematical structure. Furthermore their natural extensions to repeated games seem ideally suited to take into account the relation between the initial information of the players and the one they obtain through the correlation device. (The equilibria of a one-stage game  $\Gamma_1$  with state independent pay-offs are correlated equilibria in the underlying game with no signals). Both those concepts have been criticised and several refinements have been proposed but many further complications arise because we consider infinite games.

To define equilibria and equilibrium pay-offs we proceed as follows: first we ask for the analogue of the uniformity condition for the value.

DEFINITION 1.3.  $\sigma$  is a **uniform equilibrium** if  $\bar{\gamma}_n^i(\sigma)$  converges to some  $\bar{\gamma}^i(\sigma)$  and for all  $\varepsilon > 0$ , there exists  $N$  such that  $n \geq N$  implies  $\bar{\gamma}_n^i(\sigma_{-i}, \tau^i) \leq \bar{\gamma}^i(\sigma) + \varepsilon$ , for all  $i$  and all  $\tau^i$ .

In order to avoid the difficulties due the lack of a well specified pay-off one also introduces a family of games: for each Banach limit  $\mathcal{L}$  (cf. ex. I.2Ex.10 p. 23) we define a normal form game  $\Gamma_{\mathcal{L}}$  by the strategy sets  $\Sigma^i$  and the pay-off  $\mathcal{L}(E_{\sigma}(\bar{g}_n^i))$ ,  $i \in \mathbf{I}$ .

DEFINITION 1.4. If the set of equilibrium pay-offs of  $\Gamma_{\mathcal{L}}$  does not depend on  $\mathcal{L}$  and moreover each of these pay-offs can be realised by a uniform equilibrium, we define it as the set  $E_{\infty}$  of **equilibrium pay-offs** in  $\Gamma_{\infty}$ . Similarly  $C_{\infty}$  and  $D_{\infty}$  correspond to correlated and communication equilibrium pay-offs (cf. sect. 3 p. 88).

Note that the concept of uniform equilibrium is rather strong: it corresponds to a strategy independent of  $\varepsilon$  in (i) p. 149. It implies in particular that any equilibrium pay-off can be sustained by a strategy vector which is in equilibrium in every  $\Gamma_{\mathcal{L}}$ .

For weaker conditions, needed for stochastic games and related to  $E_0$ : pay-offs  $\varepsilon$ -close to  $\varepsilon$ -equilibrium pay-offs, for all  $\varepsilon > 0$ , (cf. sect. 4 p. 341).

Finally, since strategies in  $\Gamma_{\infty}$  induce strategies in any subgame of it, we can define a **subgame perfect equilibrium** as strategies inducing an equilibrium in any subgame.

The above arguments (interest of the zero-sum case, use of it for the non-zero-sum setup, conceptual and mathematical complications in the latter) indicates why we will largely concentrate in this chapter as well as in this book on the zero-sum case (cf. nevertheless sect. 4 p. 102, 2 p. 153 and 4 p. 162, ch. VII p. 327 and IX p. 397).

**1.e. Stochastic games and games with incomplete information.** These two classes of repeated games will be extensively studied and a first presentation follows, according to our general formulation.

In a **stochastic game**, the current state and the stage pay-off are known to all players: for all  $i$  and all  $n$ , the signal  $a_{n+1}^i$  differs on two histories where  $k_{n+1}$  or  $g_n$  are not the same. The main case is obtained when the signal is the same for all players and consists of  $(s_n, k_{n+1})$  at stage  $n + 1$ , cf. ch. VII p. 327.

In a **game with incomplete information** the state is constant along the play:  $k_n = k_1$ , for all  $n$ , but unknown to at least one player, cf. ch. V p. 183 and VI p. 273.

It follows that in the first case the main goal is to control the transitions on the state space, while in the second the decisive aspects concern the transmission of information. Nevertheless we will see that the two fields are deeply related and are better understood when studied in parallel.

### Exercises.

**1. Recursive utilities.** Cf. e.g. references in (Kreps and Porteus, 1979), and (Becker et al., 1989)

a. If  $x_n$  denotes the outcome at stage  $n$ , “recursive utilities” are defined recursively by  $U_n(x_n, x_{n+1}, \dots) = \Phi_n(x_n, U_{n+1}(x_{n+1}, \dots))$  — where  $\Phi_n$  is non-decreasing in the second variable. Show that, for the  $U_n$  to be von Neumann-Morgenstern utilities one needs:

$$\Phi_n(x, U) = \alpha_n(x) + \beta_n(x)U, \quad \beta_n \geq 0$$

b. Consider a repeated game with pay-offs defined recursively by  $U_n = \alpha_n + \beta_n \cdot U_{n+1}$ , where  $\alpha_n$  and  $\beta_n$  are stationary functions of  $(a_{n-m}, k_{n-m}, s_{n-m}, \dots, a_n, k_n, s_n, a_{n+1}, k_{n+1})$  and  $\beta_n \leq 1$ . Show how to obtain an equivalent game with enlarged space  $\tilde{K}$  where the pay-off is  $\sum_n g(\tilde{k}_n)$  (or  $\lim(1/n) \sum_n g(\tilde{k}_n)$ ).

## 2. Equivalent representations

We will describe here different variants in the formalisation that may make the model more tractable in analysing some of its properties.

### 2.a. Simple transformations.

- (1) Let us first remark that information lags are easily incorporated in our model; more precisely consider a stationary bounded recall transition, namely  $Q$  defined on  $H$  where for some fixed  $m$ , for any  $n$  and any  $h_n$  in  $H_n$  (resp.  $h_m$  in  $H_m$ ), if  $h_{n+m}$  denotes  $(h_n, h_m)$ ,  $Q(h_{n+m})$  is only a function of  $h_m$ . It is then enough to redefine a new set of states — say as  $K \times H_m$  — and to extend in an obvious way  $P$  and  $Q$  to remain in the same class. One can similarly deal with non-stationary transitions with finite memory.
- (2) To get rid of the initial lottery and of the initial information of the players, it is enough to add a new initial state from where, whatever being the actions of the players, the pay-off is 0 and the new state as well as the signals to the players are selected according to the initial distribution. (This will shift one stage further all the pay-offs but does not influence long term average).
- (3) As for the pay-offs one can first assume that  $W$  is included in  $\{g \in \mathbb{R}^I \mid g_i \geq 0, \forall i \in I, \sum_i g_i \leq 1\}$  by adding some constant and then rescaling.

One can now replace the distribution on the pay-offs by any other distribution having the same expectation: for example take a deterministic pay-off, or a probability with support on the extreme points of the simplex  $\Delta(I)$  and the zero vector in  $\mathbb{R}^I$ . In fact this change has no information effect on the game hence does not influence the strategies. Moreover the expected pay-off at each stage remains the same. It now follows from ex. II.4Ex.4 p. 105 that for any choice of  $\sigma$  the difference of the average pay-offs  $\bar{g}_n$  in both formulation will converge to zero a.s.. Hence the asymptotic properties of the game are not affected by this change.

- (4) Taking the second variant above we have now a finite set of pay-offs. Redefine then the set of states as to include the old states, the vector pay-offs and the vector of signals, namely  $K \times W \times A$ , and extend the transition on the new  $K$  in the natural way. We obtain then a model where the game starts at some initial state and after each stage, a new state is chosen at random as a function of the old stage and of the actions of the players at that stage. Signals to the players and pay-offs are now functions of the new state.

One can also, without affecting the asymptotic behaviour, shift the pay-offs one stage further and assume that the pay-off is only a function of the current state. Remark now that the players' signals can be viewed as a partition of  $K$  such that before each stage each player is told the element of his partition that contains the true state.

The game is thus described by the following elements:

- a finite set of states  $K$  with an initial state  $k_0$

- a finite set of moves  $S^i$  for each player  $i$  in  $\mathbf{I}$
- a transition probability  $P$  from  $K \times S$  to  $K$
- a partition  $\mathbf{K}^i$  of  $K$ , for each  $i$  in  $\mathbf{I}$
- a partition  $\mathbf{W} = \{W^0, W^1, \dots, W^{\mathbf{I}}\}$  of  $K$

(where  $W^0$  corresponds to the set of states with zero pay-off and  $W^i$  to those with pay-off 1 to player  $i$ ).

Before each stage, every player  $i$  is told in which element of  $\mathbf{K}^i$  the current state  $k$  is. He then chooses an action in  $S^i$  and receives a pay-off 1 iff  $k$  belongs to  $W^i$ . Then  $P$  selects a new state and the game proceeds to the next stage. Note that  $W^i$  is the set of winning states for player  $i$ ,  $W^0$  corresponds to a draw and each player maximises his expected winning frequency. In the zero-sum case one can scale player I's pay-off to lie between 0 and 1, and obtain then a partition of  $K$  in  $W^{\mathbf{I}} \cup W^{\mathbf{II}}$ : there is no draw.

We have thus proved:

**PROPOSITION 2.1.** *The games described in sect. 1.a p. 148 and 2.a p. 153 have the same asymptotic properties.*

**2.b. A deterministic framework.** We will assume here that all the coefficients defining the previous transition probability  $P$  are rational and we will reduce the model to deterministic  $P$ . The construction will be done here for the two-person zero-sum case (for the general case, cf. (Mertens, 1986b)). The purpose of this transformation is to have a better feeling of the essential structure of the problem while adding a mild assumption.

So let  $m$  be a common denominator to all rational coefficients that appear in  $P$ . Let  $\mathbb{Z}_m = \mathbb{Z}/m\mathbb{Z}$ . If  $(X_n)_{n \geq 0}$  is a sequence of i.i.d. random variables uniformly distributed on  $\mathbb{Z}_m$ , then  $P$  can be represented as a function from  $K \times S \times T \times \mathbb{Z}_m$  to  $K$  such that:  $k_{n+1} = P(k_n, s_n, t_n, x_n)$ . In fact we will show that the players can generate themselves such a sequence.(cf. also ex. II.3Ex.5 p. 97).

Assume that each player  $i$  chooses at each stage  $n$ , besides his action, an element  $z_n^i$  in  $\mathbb{Z}_m$  and take as random variables the sum in  $\mathbb{Z}_m$  of these two choices. Formally we have a new game  $\tilde{\Gamma}$  where  $\tilde{S} = S \times \mathbb{Z}_m$ ,  $\tilde{T} = T \times \mathbb{Z}_m$  and  $k_{n+1} = \tilde{P}(k_n, \tilde{s}_n, \tilde{t}_n) = P(k_n, s_n, t_n, z_n^{\mathbf{I}} + z_n^{\mathbf{II}})$ . (We keep the same signalling structure: no player is ever informed of any of the past auxiliary choices of his opponent.)

**PROPOSITION 2.2.**  *$\Gamma$  and  $\tilde{\Gamma}$  have the same asymptotic properties.*

**PROOF.** Let us prove this first for  $\underline{v}$ : existence, value and  $\varepsilon$ -optimal strategies.

We denote by  $\mathbb{Z}_m^\infty$  the compact group  $\prod_1^\infty (\mathbb{Z}_m)_n$ , by  $z$  a generic element in it and by  $\mu$  the Haar measure on it.

Given  $\tilde{\sigma}$  in  $\tilde{\Sigma}$  (strategy set of player I in  $\tilde{\Gamma}$ ) and  $z$  in  $\mathbb{Z}_m^\infty$  define  $\tilde{\sigma}_z$  by  $(s_n, z_n^{\mathbf{I}} + z_n)$  at each stage  $n$  ( $\tilde{\tau}_z$  is defined similarly). Considering  $\tilde{\sigma}_z$  as a transition probability from  $\mathbb{Z}_m^\infty$  to pure strategies in  $\tilde{\Gamma}$  one can introduce  $\tilde{\sigma} = \int \tilde{\sigma}_z \mu(dz)$  which is a mixed strategy in  $\tilde{\Gamma}$  and denote by  $\underline{\sigma}$  its projection (marginal distribution) in  $\Gamma$ .

Now for any  $\underline{\sigma}$  in  $\Sigma$  we can define  $\mu \otimes \underline{\sigma}$  in  $\tilde{\Sigma}$  and it follows from the previous construction that  $\tilde{\sigma}$  and  $\mu \otimes \underline{\sigma}$  induce the same behavioural strategies, hence by Kuhn's theorem (theorem 1.4 p. 54), the same distribution on histories in  $\tilde{\Gamma}$  for every  $\tilde{\tau}$ .

Note finally that the map  $\tilde{\sigma} \mapsto \tilde{\sigma}_z$  is a permutation of  $\tilde{\Sigma}$ , the map  $\tilde{\sigma} \mapsto \underline{\sigma}$  from  $\tilde{\Sigma}$  to  $\Sigma$  is onto and that  $(\tilde{\sigma}_z, \tilde{\tau}_{-z})$  induces the same distribution on histories that  $(\tilde{\sigma}, \tilde{\tau})$ .

Given any bounded measurable function  $f$  on  $(H_\infty, \mathcal{H}_\infty)$ , (as  $\bar{g}_n$  or  $\liminf \bar{g}_n$ ) we therefore have, for any  $\tilde{\sigma}$ :

$$f_{\tilde{\sigma}} \equiv \inf_{\tilde{\tau}} E_{\tilde{\sigma}, \tilde{\tau}}(f) = \inf_{\tilde{\tau}} E_{\tilde{\sigma}_z, \tilde{\tau}_{-z}}(f) = \inf_{\tilde{\tau}} E_{\tilde{\sigma}_z, \tilde{\tau}}(f) = f_{\tilde{\sigma}_z}$$

hence

$$E_{\underline{\sigma}, \tilde{\tau}}(f) = \int \mu(dz) E_{\tilde{\sigma}_z, \tilde{\tau}}(f) \geq f_{\tilde{\sigma}}$$

thus also

$$E_{\mu \otimes \underline{\sigma}, \tilde{\tau}}(f) \geq f_{\tilde{\sigma}}$$

and

$$f_\mu \otimes \underline{\sigma} \geq f_{\tilde{\sigma}}$$

Defining similarly in  $\Gamma$ ,  $f_\sigma \equiv \inf_{\tau} E_{\sigma, \tau}(f)$ , we claim that:  $f_{\underline{\sigma}} = f_\mu \otimes \underline{\sigma}$ .

In fact, using the properties of the Haar measure:

$$\begin{aligned} E_{\mu \otimes \underline{\sigma}, \tilde{\tau}}(f) &= E_{(\mu \otimes \underline{\sigma})_{-z}, \tilde{\tau}_z}(f) = E_{\mu \otimes \underline{\sigma}, \tilde{\tau}_z}(f) = \int \mu(dz) E_{\mu \otimes \underline{\sigma}, \tilde{\tau}_z}(f) \\ &= E_{\mu \otimes \underline{\sigma}, \tilde{\tau}}(f) = E_{\mu \otimes \underline{\sigma}, \mu \otimes \underline{\tau}}(f) \end{aligned}$$

Now under  $\mu \otimes \underline{\sigma}$  and  $\mu \otimes \underline{\tau}$  the random variables  $z^I$  and  $z^{II}$  are independent and their marginals uniform on  $\mathbb{Z}_m$  hence:  $E_{\mu \otimes \underline{\sigma}, \mu \otimes \underline{\tau}} = E_{\underline{\sigma}, \underline{\tau}}$ .

We obtain thus:

$$(1) \quad E_{\mu \otimes \underline{\sigma}, \tilde{\tau}}(f) = E_{\underline{\sigma}, \underline{\tau}}(f)$$

and also:

$$(2) \quad E_{\underline{\sigma}, \underline{\tau}}(f) = E_{\tilde{\sigma}, \mu \otimes \underline{\tau}}(f)$$

hence the claim.

Coming back to definition 1.2 p. 149 it follows from  $f_{\underline{\sigma}} \geq f_{\tilde{\sigma}}$  and  $f_{\underline{\sigma}} = f_\mu \otimes \underline{\sigma}$  that  $(\tilde{\sigma}, w)$  satisfies (i) p. 149 in  $\tilde{\Gamma}$  only if  $(\underline{\sigma}, w)$  satisfies it in  $\Gamma$ , and the latter holds iff  $(\mu \otimes \underline{\sigma}, w)$  satisfies (i) in  $\tilde{\Gamma}$ .

Similarly (2) implies that if  $w$  satisfies (ii) p. 150 in  $\Gamma$ , then it also does in  $\tilde{\Gamma}$ . Finally (1) shows the reverse inequality since  $\mu \otimes \Sigma$  is included in  $\tilde{\Sigma}$ . This implies that the maxmin will exists in  $\Gamma$  iff it exists in  $\tilde{\Gamma}$ . In this case its value will be the same and  $\varepsilon$ -optimal strategies can be mapped through  $\underline{\sigma} \mapsto \mu \otimes \underline{\sigma}$ .

It is easy to see that analogous results for  $v_\lambda$  or  $v_n$  follow from the same arguments. ■

**2.c. A combinatorial form.** The aim of this last transformation is to push one step further towards a simple sequential combinatorial form.

We first rank the players in cyclical order (identifying  $\mathbf{I}$  and  $\mathbb{Z}_{\mathbf{I}}$ ) and subdivide each stage into  $\mathbf{I}$  substages, player  $i$  being the only one to choose an action at each stage  $n \equiv i \pmod{\mathbf{I}}$ . The new set of states  $F$  is partitioned into subsets  $F^1, \dots, F^{\mathbf{I}}$  and an element in  $F^i$  is a point in  $E$  together with the last actions of players  $j, 1 \leq j \leq i-1$ .

After each substage  $\mathbf{I}$ ,  $P$  is used to compute an element of  $F$  (in fact, of  $F^1 = E$ ), as a function of the old element in  $E = F^1$  and the vector of actions, or equivalently as a function of the element in  $F^{\mathbf{I}}$  and the last action of player  $\mathbf{I}$ .

Since all players' pay-offs are functions only of the point in  $F^1$ , included in  $F^i$  for all  $i$ , one can delay them so that player  $i$ 's winning set  $W^i$  will be a subset of  $F^i$ .

Each player's action set  $S^i$  can now be considered as a subset of the set of functions from  $F^i$  to  $F^{i+1}$ . In fact, to modelise his information one can replace  $S^i$  by the set of  $\mathbf{K}^i$  — measurable functions from  $F^i$  to  $S^i$  — hence to  $F^{i+1}$ ; and give him his private information in  $\mathbf{K}^i$  about the true state only after his move.

We thus obtain the following combinatorial form:

- a finite set of states  $F$  with an initial state  $f_0$  and a partition  $F^1, \dots, F^I$  of  $F$ .
- for each  $F^i$  a partition  $\mathcal{F}^i$  and a set  $Y^i$  of functions from  $F^i$  to  $F^{i+1}$
- a subset  $W$  of  $F$ .

The play of the game is as follows, starting from  $f_0$ . At the current state  $f$  in  $F^i$ , player  $i$  selects a function  $y^i$  in  $Y^i$ . He then gets a pay-off 1 iff  $f$  is in  $W$ , zero otherwise, and is told the element of  $\mathcal{F}^i$  that contains  $f$ . The new state is now  $y^i(f)$  in  $F^{i+1}$ . The players want to maximise their expected winning frequencies.

Replacing now  $Y^i$  by  $Y = \prod_{i \in I} Y^i$ , and letting  $\mathcal{F}$  be the partition  $\bigcup_{i \in I} \mathcal{F}^i$  of  $F$ , we reach the following description, coming back to our standard notations:  $K$  is the state space,  $S$  is a set of functions from  $K$  to  $K$ ,  $A$  is a partition on  $K$  and  $a(k)$  is the point in  $A$  containing  $k$ . Players play in cyclical order, choosing at each time  $n$  an element  $s_n$  of  $S$ , being then informed of  $a(k_n)$ . The new state is  $k_{n+1} = s_n(k_n)$ , and the player who moved receives 1 if  $k_n \in W$ , zero otherwise. In the two-person zero-sum case, he receives 1 from his opponent if  $k_n \in W$ , and pays him 1 otherwise.

We finally show how to reduce oneself further (by enlarging  $K$ ) to the case where  $\#S = \#A = 2$ . Fix a two-element set,  $\{f, g\}$ , and a map from  $\{f, g\}^l$  onto  $S$ . Fix also a map from  $A$  into  $\{a_1, a_2\}^h$ . Replace every stage of the above game by a block of  $l \cdot I$  substages followed by a bloc of  $h \cdot I$  substages. If this was a stage for player  $i$ , at every substage corresponding to a player  $j \neq i$  this player's choice in  $\{f, g\}$  has no effect, and this player is told  $a_1$ . While on the first  $l$  substages corresponding to player  $i$ , this player is told  $a_1$ , and his sequence of choices in  $\{f, g\}$  gives a point in  $\{f, g\}^l$ , and hence determines the map  $s \in S$  to be used. The signal in  $A$  he has to hear is then encoded in  $\{a_1, a_2\}^h$ , and told to him over the last  $h$  substages, where his choices in  $\{f, g\}$  have no effect. As for pay-offs, all substages except the last substage of player  $i$  are outside  $W$ , on the last substage the correct set  $W$  is used. Thus we now have a model where  $S = \{f, g\} \subseteq K^K$ , and where  $A = \{a_1, a_2\}$ , i.e. corresponds also to a subset of  $K$ .

We thus have proved:

**PROPOSITION 2.3.** *If  $P$  (as obtained in 2.a p. 153) is rational the games described in sect. 1.a p. 148 and 2.c p. 155 have the same asymptotic properties.*

### 3. Recursive structure

In this section we will study what is the natural space in which one can model the information obtained during the game while keeping the stationary aspects.

Basically we are looking for entrance laws that could allow us to start the study at any stage of the game: we will see that a basic structure like common knowledge of a prior on the state space and private partitions is not enough.

In this part again, we will concentrate on two-person zero-sum games but the analysis can be extended to more general cases (cf. ch. III and (Mertens, 1986b)).

**3.a. A canonical representation.** Consider a pair of strategies in  $\Gamma$ . Together with the description of the game, including the initial conditions, they determine a probability distribution on  $(H_\infty, \mathcal{H}_\infty)$ , space of plays, that both players can compute. By the previous reductions, at stage  $n$ , the future play of the game depends only on the current state  $k_n$ . Each player  $i$ ,  $i = I, II$ , has in addition accumulated some signals that can be modeled by means of some set  $M^i$ . More precisely  $k_n$  is a random variable from  $(H_\infty, \mathcal{H}_\infty)$  to  $K$  and the information of player  $i$  is a random variable  $m^i$  from  $(H_\infty, \mathcal{H}_\infty)$  to  $M^i$ . In our framework, there exists moreover an underlying probability  $Q$  on  $(H_\infty, \mathcal{H}_\infty)$  defined by  $P$  and the strategies. If we let  $(E, \mathcal{E}) = (H_\infty, \mathcal{H}_\infty)$  and  $\mathcal{E}^i$  be the sub  $\sigma$ -field generated by  $m^i$  we obtain an information scheme, hence by Theorem 2.5 p. 122 there exists a canonical representation on  $\Omega$  with  $\mathbf{P}_n$  in  $\mathcal{P}$ : note that this probability contains all relevant information about the past, more precisely the future aspects of the game should be the same if it was starting at stage  $n$  using  $\mathbf{P}_n$  to choose the state and the information to the players (we use here the zero-sum assumption).

It is clear that typically these  $\mathbf{P}_n$  will have specific properties, basically finite support in our finite framework, but one sees easily that generically their size cannot be bounded, hence the advantage of working directly with the closure  $\mathcal{P}$  in order to have a stationary set of “state variables” which are the **entrance laws**  $\mathbf{P}_n$ . Before seeing in the next paragraph the merits of such an approach, let us remark at this point that  $\mathbf{P}_{n+1}$  is easily constructed using  $\mathbf{P}_n$  and the behavioural strategies at stage  $n$  (i.e. the restriction of  $\sigma$  and  $\tau$  to  $H_n$ ) and that similarly, according to 2.b p. 154, the pay-off at stage  $n$  for player  $i$  is simply  $E_{\sigma, \tau}(g_n^i) = \mathbf{P}_n(W^i)$ .

**3.b. The recursive formula.** We first want to be able to apply the results of sect. 4.b p. 132 to the finite game  $\Gamma_n$  and the discounted game  $\Gamma_\lambda$ .

**PROPOSITION 3.1.** *The results of sect. 4.b p. 132 apply to  $v_n$  and  $v_\lambda$ , strategies being behavioural strategies.*

**PROOF.** Indeed both  $\Gamma_n$  and  $\Gamma_\lambda$  have  $K$  finite, and  $S$  and  $T$  compact metric, with  $g_k(s, t)$  continuous. Remains to show the results are still true with behavioural strategies. This follows from II.1Ex.10, except possibly for the continuity. For this aspect, let the finite subset  $S_0$  of  $S$  consist, for  $\Gamma_n$  of the whole of  $S$ , and for  $\Gamma_\lambda$ , of all pure strategies which from a certain stage  $n_0$  on play always the first pure strategy ( $n_0$  is fixed). There is no problem in requiring further in sect. 4.b that every pure strategy in  $S_0$  has strictly positive probability. Then the map from mixed strategies to behavioural strategies is continuous (with the above set  $S_0$ )). ■

We are now going to use fully the structure of “entrance laws” that we introduced in 3.a p. 157. To obtain a nice recursive formula it will be convenient to keep with the framework of 2 p. 153 where the current pay-off is only a function of the state at this stage and where the information is given at the end of each stage.

Assume then some entrance law  $\mathbf{P}$  and behavioural strategies  $x$  and  $y$  of the players for the first stage, namely measurable mappings from  $\Theta$  to  $\Delta(S)$  or  $\Delta(T)$ .

We define  $\mathcal{I}(\mathbf{P}, x, y)$  as the following information scheme: first a triple  $(k, \theta^I, \theta^{II})$  in  $\Omega$  is selected according to  $\mathbf{P}$  and player  $i$  is informed of  $\theta^i$ . Both players select then independently moves according to  $x(\theta^I)$  or  $y(\theta^{II})$ . Finally a new state and random signals  $(\tilde{k}, a^I, a^{II})$  are selected in the game as usual as a function of the old state and of the pair

of moves. Formally:

$$\mathcal{I}(\mathbf{P}, x, y) = (\Omega; \mathbf{P}_{x,y}; \tilde{k}; (\theta^I, s, a^I); (\theta^{II}, t, a^{II}))$$

Finally  $\mathbf{P}[x, y]$  in  $\mathcal{P}$  will denote the correspondent canonical probability.

We can now justify “le bien fondé” of the previous point of view by stating our main result:

**THEOREM 3.2** (Recursive Formula). (1) Let  $v(\mathbf{P})$  stands either for  $v_\lambda(\mathbf{P})$  or  $v_n(\mathbf{P})$ . Then both  $\max_x \min_y v(\mathbf{P}[x, y])$  and  $\min_y \max_x v(\mathbf{P}[x, y])$  exist and are equal.  
(2) Denoting by  $V(v[\mathbf{P}])$  this saddle point value we have:

$$\begin{aligned} v_\lambda(\mathbf{P}) &= \lambda \mathbf{E}_\mathbf{P}(g_1) + (1 - \lambda)V(v_\lambda[\mathbf{P}]) \\ nv_n(\mathbf{P}) &= \mathbf{E}_\mathbf{P}(g_1) + (n - 1)V(v_{n-1}[\mathbf{P}]) \end{aligned}$$

(3) In the space of all bounded functions on  $\mathcal{P}$ ,  $v_\lambda$  and  $v_n$  are uniquely determined by the above formulae.

**PROOF.** Let  $\sigma = (\sigma_n)_{n \geq 1}$  and  $\tau = (\tau_n)_{n \geq 1}$  be behavioural strategies of players I and II in  $\Gamma_\lambda$  and assume  $\sigma$  optimal. Let us write  $\sigma_+$  or  $\tau_+$  for the sequences  $(\sigma_n)_{n \geq 2}$  or  $(\tau_n)_{n \geq 2}$  and note that both  $\sigma_+$  and  $\tau_+$  are behavioural strategies in  $[\Gamma_\lambda; \mathcal{I}(\mathbf{P}, \sigma_1, \tau_1)]$ . Denoting the corresponding expected pay-off by  $\bar{g}_\lambda(\mathcal{I}(\mathbf{P}, \sigma_1, \tau_1); \sigma_+, \tau_+)$  we obtain thus, writing the total pay-off in  $\Gamma_\lambda$  as the sum of the first stage pay-off and of the remaining one:

$$\begin{aligned} v_\lambda(\mathbf{P}) &\leq \bar{g}_\lambda(\mathbf{P}; \sigma, \tau) \\ &= \lambda \mathbf{E}_\mathbf{P}(g_1) + (1 - \lambda)\bar{g}_\lambda(\mathcal{I}(\mathbf{P}, \sigma_1, \tau_1); \sigma_+, \tau_+). \end{aligned}$$

taking the infimum in  $\tau_+$  on the right hand side then the maximum in  $\sigma_+$  we also obtain:

$$v_\lambda(\mathbf{P}) \leq \lambda \mathbf{E}_\mathbf{P}(g_1) + (1 - \lambda) \sup_{\sigma_+} \inf_{\tau_+} \bar{g}_\lambda(\mathcal{I}(\mathbf{P}, \sigma_1, \tau_1); \sigma_+, \tau_+)$$

Using propositions 4.5 p. 134 and the above 3.1, the  $\sup \inf$  is precisely  $v_\lambda[\mathcal{I}(\mathbf{P}, \sigma_1, \tau_1)]$ , i.e.  $v_\lambda(\mathbf{P}[\sigma_1, \tau_1])$ . Thus:

$$v_\lambda(\mathbf{P}) \leq \lambda \mathbf{E}_\mathbf{P}(g_1) + (1 - \lambda)v_\lambda(\mathbf{P}[\sigma_1, \tau_1])$$

So,  $\tau_1$  being arbitrary:

$$v_\lambda(\mathbf{P}) \leq \lambda \mathbf{E}_\mathbf{P}(g_1) + (1 - \lambda) \inf_{\tau_1} v_\lambda(\mathbf{P}[\sigma_1, \tau_1])$$

It remains to prove that we can actually replace the inf by a min.

For this purpose consider now the  $\lambda$ -discounted game where player I is restricted to use  $\sigma_1$  in the first stage. Redefining the state space it is easily seen that this game belongs to the same family, hence by prop. 3.1 p. 157 has a value and optimal strategies. It follows that even if player I was informed about the first stage strategy  $\tau_1$  of player II the value would be the same and the optimal strategy still optimal. But the value of this variant is precisely the previous right hand member where now player II has an optimal strategy, hence:

$$v_\lambda(\mathbf{P}) \leq \lambda \mathbf{E}_\mathbf{P}(g_1) + (1 - \lambda) \min_{\tau_1} v_\lambda(\mathbf{P}[\sigma_1, \tau_1])$$

Denote by  $\sigma_1^*$  and  $\tau_1^*$  the first stage components of optimal strategies of both players in  $\Gamma_\lambda$ . We have shown that, for all  $\sigma_1^*$ ,  $\min_{\tau_1} v_\lambda(\mathbf{P}[\sigma_1^*, \tau_1])$  exists and:

$$v_\lambda(\mathbf{P}) \leq \lambda \mathbf{E}_\mathbf{P}(g_1) + (1 - \lambda) \min_{\tau_1} v_\lambda(\mathbf{P}[\sigma_1^*, \tau_1])$$

Since the dual result holds, we obtain finally:

$$\max_{\sigma_1} v_\lambda(\mathbf{P}[\sigma_1, \tau_1^*]) \leq \min_{\tau_1} v_\lambda(\mathbf{P}[\sigma_1^*, \tau_1])$$

so that  $\sigma_1^*, \tau_1^*$  is a saddle point of  $v_\lambda(\mathbf{P}[\cdot, \cdot])$ , hence  $V(v_\lambda[\mathbf{P}])$  exists and satisfies:

$$v_\lambda(\mathbf{P}) = \lambda \mathbf{E}_\mathbf{P}(g_1) + (1 - \lambda)V(v_\lambda[\mathbf{P}])$$

The results concerning  $v_n(\mathbf{P})$  are obtained in the same way.

To prove point 3 p. 158 note first that  $v_n$  is uniquely determined given  $v_{n-1}$  and that  $v_1$  is well defined.

For  $v_\lambda$  replace  $V$  by the previous  $\sup \inf$  in the equation. Note that the operator involved is contracting, hence use Picard's contraction principle in the space of all bounded functions on  $\mathcal{P}$  (lemma 1.1 p. 327) to get a unique solution. This completes the proof of the theorem. ■

Define thus an operator  $\Psi$  on bounded functions on  $\mathcal{P}$  by:

$$[\Psi(f)](\mathbf{P}) = \mathbf{E}_\mathbf{P}(g_1) + V[f(\mathbf{P})]$$

(where to avoid ambiguity  $V$  is chosen to be the  $\inf \sup$ ), and note that  $\Psi$  is non-expansive.

Denoting  $V_\lambda = (1/\lambda)v_\lambda$  and  $V_n = nv_n$  one has:

$$V_\lambda = \Psi[(1 - \lambda)V_\lambda] \text{ and } V_n = \Psi[V_{n-1}].$$

This recursive formula (or rather its particular incarnations) will play a fundamental rôle in the next chapters, in proving asymptotic properties of  $\Gamma_n$  or  $\Gamma_\lambda$ . It will also allow us to get precise bounds on the speed of convergence of  $v_n$  and  $v_\lambda$  to their limits.

Another use of it is to show that we have reduced the problem of studying  $\Gamma_n$  or  $\Gamma_\lambda$  to the same problem for a class of stochastic games. Indeed denote by  $Y^I$  and  $Y^{II}$  the set of Borel functions from  $\Theta^I$  or  $\Theta^{II}$  to  $\Delta(S)$  or  $\Delta(T)$ . Let us define  $\Gamma^*$  as the stochastic game with continuous state and actions spaces  $\mathcal{P}, Y^I, Y^{II}$ , played as follows: if  $\mathbf{P}$  is the current state and  $y^I, y^{II}$  the actions selected by the players, the current pay-off is  $\mathbf{E}_\mathbf{P}(g_1)$  and the new state is  $\mathbf{P}[y^I, y^{II}]$ .

$\Gamma^*$  is thus a "deterministic stochastic" game where the current pay-off is solely a function of the current state. Write  $\Gamma_n^*(\mathbf{P})$  and  $\Gamma_\lambda^*(\mathbf{P})$  for the finite and discounted versions of  $\Gamma^*$  starting at state  $\mathbf{P}$ . Then we have:

**PROPOSITION 3.3.**  $\Gamma_n^*(\mathbf{P})$  and  $\Gamma_\lambda^*(\mathbf{P})$  have values  $v_n^*(\mathbf{P})$  and  $v_\lambda^*(\mathbf{P})$  and optimal strategies. Moreover:

$$v_\lambda^* = v_\lambda, \quad v_n^* = v_n.$$

**PROOF.** We will prove that there exist optimal Borel pure strategies which are at each stage versions of the corresponding component of optimal strategies in  $\Gamma_\lambda$ , resp.  $\Gamma_n$ .

For  $\Gamma_n^*$  the proof follows immediately by induction on  $n$ .

For  $\Gamma_\lambda^*$  assume that both players use the above described strategies. The recursive formula implies obviously that the pay-off will be  $v_\lambda$ . Assume now that player II is allowed to choose arbitrarily his strategy during the first  $n$  stages, but has to play the above strategy thereafter. At stage  $n$ , using the recursive formula his optimal choice is again to use an optimal strategy in  $\Gamma_\lambda(\mathbf{P}_n)$ . By induction his optimal play will always have this property: he cannot get below  $v_\lambda(\mathbf{P})$  within this class of strategies. Finally, since the pay-off depends only up to  $\varepsilon$  on the play of the game after stage  $n$ , for  $n$  large enough, it follows that the prescribed strategy of player I guarantees him  $v_\lambda(\mathbf{P})$ . A dual statement for player II now implies the result. ■

The purpose of the two next sections will be to present some classes of games belonging to this general model while having a very specific structure (other examples will be found in exercises). The hope is that this will help the reader to get a better feeling of the basic aspects of repeated games and to see how different presentations can be handled in the general framework.

### Exercises.

**1. Gleason's game.** Three positions, say  $A, B, C$ , with corresponding pay-offs  $1, 2, -3$  are arranged in cyclic order and alternatively each player tells the referee whether he wants to move clockwise or counterclockwise, after which the referee tells him his position.

a. Prove that there exists a sequence  $w_n$  with  $|v_n(p) - w_n| \leq K/n$ , for some  $K > 0$ , where  $v_n(p)$  is the value of a  $n$  stage game with initial probability  $p$  on the states and any kind of initial signals. Moreover each player can guarantee  $w_n$  up to  $O(1/n)$  by a strategy independent of his initial information.

HINT. Consider the least favourable situation for I: II knows the initial position and  $p$  is public knowledge. Write  $v_n(p^A, p^B, p^C)$  for the corresponding value and let  $u_A(q) = v_n(0, q, 1-q)$ ,  $u_B(q) = v_n(1-q, 0, q)$ ,  $u_C(q) = v_n(q, 1-q, 0)$ ,  $u(q)$  is the median of those three numbers and  $w_n = \min_q u(q) = u(q_0)$ . Given the strategy of player II and his own last position, say  $D$ , player I can compute the last mixed action  $q$  of player II. Let him play  $(1/2, 1/2)$  i.i.d. until the first stage where  $u_D(q) \geq u(q)$  and from then on optimally in the remaining game.

Similarly in the most favourable situation for player I, player II can play alternatively and independently  $(1/2, 1/2)$  and  $(q_0, 1-q_0)$  until being informed, after a random choice  $(q_0, 1-q_0)$ , of a last position  $B$  or  $C$  (assuming  $u_A(q_0) > u(q_0)$ ) and from then on play optimally.

Finally remark that the pay-off guaranteed by a strategy of player I in the first situation is linear in  $p$  and consider an optimal strategy for some interior point  $p$ .

b. Deduce from IV.3Ex.1a that player I can guarantee  $\limsup w_n$ .

c. Conclude that  $\Gamma$  has a value  $v$  and both players have optimal strategies independent of the initial information.

REMARK 3.1. The value of  $v$  and the existence of optimal stationary strategies are still open problems.

**2. A game with two-move information lag.** (Dubins, 1957), (Karlin, 1957), (Ferguson, 1967) We consider the following game  $\Gamma$ . Player II moves on the integers by choosing at each time  $n$  an element  $t_n \in \{-1, +1\}$ . His position after time  $n$  is  $x_n = \sum_{i=1}^n t_i$  and  $y_n$  is the corresponding history:  $\{x_1, \dots, x_n\}$ . The information of player I at time  $n$  is  $y_{n-2}$  and his aim is to guess the actual position of player II, i.e.  $x_n$ .

Formally let  $\theta$  be the time where player I tries to guess the position of player II by choosing some  $x$ . The corresponding pay-off is then  $\mathbb{1}_{\theta < \infty} \times \mathbb{1}_{x_\theta = x}$ .

a. Let  $\Gamma_n$  be the  $n$  stage game (where the pay-off is identically 0 after stage  $n$ ) and denote by  $v_n$  its value.

Prove that  $\Gamma$  has a value  $v$  with  $v = \lim \uparrow v_n$  and that player II has an optimal strategy  $\tau^*$  in  $\Gamma$  (use ex. I.1Ex.2b p. 9).

We want now to compute explicitly  $v$ , and to describe  $\varepsilon$ -optimal strategies. In IV.3Ex.2b and IV.3Ex.2c two alternative approaches are used to obtain  $v$  and  $\tau^*$ . In IV.3Ex.2d it is proved that player I has no optimal strategy. In IV.3Ex.2e a more general model is analysed and an  $\varepsilon$ -optimal strategy of player I is obtained.

b.

i. Given an history  $y$  of length  $n$  we introduce the positions that can occur at time  $n+2$  namely:  $x^0(y) = x_n, x^1(y) = x_n + 2, x^{-1}(y) = x_n - 2$ .

Let  $\tau$  be a strategy of player II and  $Y(\tau)$  the set of paths having positive probability under  $\tau$ . For  $y$  in  $Y(\tau)$ ,  $Q^i(\tau, y)$  is the conditional probability, given  $\tau$  and  $y$  that  $\{x_{n+2} = x^i(y)\}, i = -1, 0, 1$ . Show the following:

There is no  $\tau$  such that for all  $y$  in  $Y(\tau)$   $Q^i(\tau, y) \leq 1/3$ .

$\tau$  is optimal iff for all  $y$  in  $Y(\tau)$ ,  $Q^i(\tau, y) \leq v, i = -1, 0, 1$ .

ii. Given  $\tau$  and  $y$  of length  $n$ , let  $\tau(y)$  denote the probability that  $t_{n+1} = 1$ . Define  $\tau'$  by  $\tau'(\emptyset) = 1/2$  and for  $y$  of length  $n$ ,  $\tau'(y) = \rho$  if  $t_n = 1, \tau'(y) = 1 - \rho$  if  $t_n = -1$ . Prove that there exists  $\rho$  such that the corresponding  $\tau' (= \tau'(\rho))$  guarantees  $u = (3 - \sqrt{5})/2$  to player II. Deduce that player II has no Markov  $\varepsilon$ -optimal strategy (cf. sect. 5 p. 175).

iii. Take now some optimal  $\tau$  and assume  $\tau(y) = \alpha_1 > \rho$  for some  $y \in Y(\tau)$ . Define  $\alpha_2 = 1 - \tau(y, -1)$  where  $(y, x, x', \dots)$  is the path  $y$  followed by  $x$  then  $x'$  and so on. Let  $\beta_1 = \alpha_1, \beta_2 = (1 - 2v)/(1 - \beta_1)$ . Prove that  $\alpha_2 \geq \beta_2$  and construct inductively  $\alpha_n, \beta_n$  with  $1 \geq \alpha_n \geq \beta_n = (1 - 2v)/(1 - \beta_{n-1})$  and  $\beta_n$  increasing. Deduce a contradiction by going to the limit.

iv. Prove that  $\tau$  is optimal iff  $\tau(\emptyset) \in [1 - \rho, \rho]$  and  $\tau$  coincides with  $\tau'$  thereafter. Deduce that  $v = u$ .

c. Let  $w_n(\alpha)$  be the value of the game  $\Gamma_n$  where moreover player II is restricted to use strategies  $\tau$  satisfying  $\tau(\emptyset) = \alpha$ . (Put  $w_1 \equiv 0$ ).

i. Prove that (recursive formula):

$$w_n(\alpha) = Tw_{n-1}(\alpha) \equiv \min_{0 \leq \beta, \gamma \leq 1} \max \begin{cases} \beta\alpha \\ (1 - \beta)\alpha + \gamma(1 - \alpha) \\ (1 - \alpha)(1 - \gamma) \\ \alpha w_{n-1}(\beta) + (1 - \alpha)w_{n-1}(\gamma) \end{cases}$$

and that:

$$v_n = \min_{0 \leq \alpha \leq 1} w_n(\alpha) = w_n(1/2)$$

ii. Prove that  $w_n$  converges uniformly to some continuous function  $w$  satisfying  $w = Tw$  and that  $v = \min_{0 \leq \alpha \leq 1} w(\alpha)$ .

iii. Let  $\delta = \max\{\alpha \mid w(\alpha) = v\}$  and let  $(\beta, \gamma)$  be the corresponding values achieving the minimum in  $Tw(\delta)$ . Show that  $(\beta, \gamma)$  belongs to  $[1 - \delta, \delta]$  and that  $v \geq 1 - \delta$ . Prove then that  $v > 1 - \delta$  implies  $v > u$  and use IV.3Ex.2bii to get finally  $\beta = 1 - \gamma = \delta$  and  $v = u$ . Deduce that  $\tau'$  is optimal.

d.

i. Assume that  $\sigma$  is optimal and let  $y_{n-2}$  be a minimal history such that  $P_\sigma(\theta = n \mid y_{n-2}) = a > 0$ . Prove then that player I has to choose  $x = x_{n-2}$ .

ii. Show now that if there exists  $\alpha$  such that:  $a(1 - \alpha) + (1 - a)w(\alpha) < v$ ,  $\sigma$  is not optimal. Prove then inductively on  $n$  that there exists  $L > 0$  and for every  $n, \varepsilon_n > 0$  such that for  $\varepsilon < \varepsilon_n$ :

$$w(\rho + \varepsilon) - w(\rho) \leq L\rho^n \varepsilon$$

Deduce finally the above inequality.

iii. Alternative approach: Let  $b$  be the probability given  $y_{n-1} = (y_{n-2}, 1)$  that  $\theta = n+1$  and  $x = x_{n-1}^1$  and similarly  $c$  for  $y_{n-1} = (y_{n-2}, -1)$  and  $x = x_{n-1}^{-1}$ . Prove by letting player II play some  $\tau'(\rho + \varepsilon)$ , that  $b$  and  $c$  are greater than  $a/\rho$  and deduce inductively a contradiction.

e. We consider now an extension of  $\Gamma$  where player II is moving on a graph satisfying the following conditions:

- no edge joins a vertex to itself
- each vertex is joined to  $k + 1$  vertices.
- there are no four sided figures

i. Define a strategy  $\tau^*$  of player II as follows:  $\tau^*(\emptyset)$  is uniform on the  $(k + 1)$  adjacent vertices. Given  $y$  of length  $n$ ,  $\tau$  goes to  $x_{n-1}$  with probability  $1 - kp$  and to each other vertices with probability  $p$ . (Note that condition c) above implies that the positions after 2 stages differ if the first move is not the same, unless they are both the initial one). Show that there exists  $p$  such that the corresponding  $\tau^*$  guarantees  $u_k = (k^2 + 2 - k\sqrt{k^2 + 4})/2$ .

Let  $\theta < u_k$ . We want to construct  $\sigma$  that guarantees  $\theta$  to player I.

ii. Define first inductively a sequence  $\{q_n\}$  by:

- $q_0 = 1, q_1/(q_0 + (k^2 + 1)q_1) = \theta$ , then  $q_n$  satisfying:
- $(q_0 + \dots + q_{n-2} + q_n)/(q_0 + (k^2 + 1)(q_0 + \dots + q_{n-1} + q_n)) = \theta$ .

Let  $R = (1 - (k^2 + 1)\theta)^{-1}$  and prove that:

- $q_n = D[(R + \sqrt{R^2 - 4R})/2]^n + (1 - D)[(R - \sqrt{R^2 - 4R})/2]^n$  with
- $2D = 1 - ((k^2 - 1)R + 2)/((k^2 + 1)\sqrt{R^2 - 4R})$ , and that  $1 < R < 4$  for  $\theta$  sufficiently close to  $u_k$ .

Show then that there exists a first  $n$ , say  $N$ , with  $q_n \leq q_{n-1}$ . Define now  $q = \sum_{n=0}^N q_n$  and  $p_n = q_n/q$ .

iii. Define finally  $\sigma$  by:

- initial procedure:  $P_\sigma(\theta = 2, x = 0) = p_0$ , and given  $x_1$ ,  $P_\sigma(\theta = 3, x) = p_1$  for all  $x$  such that there exists a feasible path  $y_3 = (x_1, x_2, x_3)$  with  $x_2 \neq 0$  and  $x_3 = x$ .
- if  $x_2 \neq 0$  start the initial procedure at stage 2 from  $x_2$
- else let  $P_\sigma(\theta = 4, x) = p_2$ , for all  $(k^2 + 1)$  positions  $x$  such that there exists  $y_4 = (x_1, 0, x_3, x_4)$  with  $x_3 \neq x_1$  and  $x_4 = x$  (i.e. an history leading to  $x$  without passing by  $x_1$  again).
- if  $x_3 \neq x_1$ , start the initial procedure at stage 3 from  $x_3$ .
- else let  $P_\sigma(\theta = 5, x) = p_3$  for all  $x$  such that there exists  $y_5 = (x_1, 0, x_1, x_4, x_5)$  with  $x_4 \neq 0$ , and  $x_5 = x$ , and so on...

Prove that, for all  $\tau$ :  $P_\sigma(\theta < \infty) = 1$  and that  $\sigma$  guarantees  $\theta$ .

#### 4. Supergames

We will consider in this section a particular version of the general model where there is only one state. In this case the repeated game is called the **supergame** associated to the one-shot game.

The simplest framework corresponds to standard signalling (all the players are told the previous moves): we obtain a game with **complete information** and **full monitoring**.

**4.a. Standard signalling.** We first introduce some notations:

$D$  is the set of **feasible pay-offs** (with correlated strategies in  $\Gamma_1$ ), i.e. the convex hull of the set of pay-offs attainable with pure strategies in the one-shot game. (Recall that the pay-offs are uniformly bounded by some constant  $C$ ).

The **minmax level** for player  $i$  is defined by  $v^i = \min_{X^{-i}} \max_{X^i} \gamma_1^i(x^i, x^{-i})$ .

$x^{-i}(i)$  denotes a point in  $X^{-i}$  realising the above minimum.

$v$  with component  $v^i$  is the **threat point**.

The set of **feasible and individually rational** (i.r.) pay-offs is defined by:  $E = \{d \in D \mid d^i \geq v^i \forall i \in \mathbf{I}\}$

We denote by  $E_\infty, E_\lambda, E_n$  the set of equilibrium pay-offs in  $\Gamma_\infty, \Gamma_\lambda, \Gamma_n$ , respectively.

The following basic result, known as the Folk theorem, is the starting point of the theory of supergames. It states that the set of equilibrium pay-offs in the infinitely repeated game coincides with the set of feasible and i.r. pay-offs of the one-shot version.

**THEOREM 4.1.**  $E_\infty = E$ .

**PROOF.** The inclusion  $E_\infty \subseteq E$  is easy. First notice that each stage pay-off  $g_n$  is in  $D$  (closed and convex) hence also average, expectation and limits; thus any equilibrium pay-off (in fact any feasible pay-off in  $\Gamma_\infty, \Gamma_\lambda$  or  $\Gamma_n$ ) is in  $D$ . To prove that the pay-off is necessarily i.r., recall that full monitoring is assumed, hence given any history  $h$  and the vector of mixed strategies of his opponents  $\sigma^{-i}$ , player  $i$  has a reply to the corresponding vector of mixed moves  $\sigma^{-i}(h)$  that gives him, at that stage a pay-off greater than  $v^i$ .

The proof that any point in  $E$  corresponds to an equilibrium pay-off relies on two basic tools: **plan** and **punishment**.

A plan is a play,  $h$ , that leads to a specified pay-off.

A punishment is a strategy that dictates to play  $x^{-i}(i)$  i.i.d. as soon as player  $i$  deviates: a **deviation** means that the actual history  $h'$  is not an initial part of the play  $h$ , and denoting by  $n$  the first stage where this holds, player  $i$  is the first (in some order) among the players whose moves at that stage do not coincide with the one defined by  $h$ .

It is now clear how to define through a plan and punishments an  $\mathbf{I}$ -tuple of strategies: the players are requested to follow  $h$  and to punish the first deviator (if any). It follows that every play leading to an i.r. pay-off will correspond to an equilibrium since by the above description any potential deviation of  $i$  (leading to a one-shot bounded gain) would induce a future expected stage pay-off at most  $v^i$ , hence a limiting average pay-off less or equal than  $v^i$ .

It remains thus to remark that the repetition of the game allows to convexify the set of feasible pay-offs: in fact given  $d$  in  $D$ , there exist actions  $\{s_t^i\}$ ,  $i = 1, \dots, \mathbf{I}$ ,  $t = 1, \dots, \mathbf{I}+1$ , and barycentric coefficients  $\mu_t$  such that:  $d = \sum_t \mu_t \gamma_1(s_t)$ . Let  $p_t^n/q^n$  be rational approximations of  $\mu_t$  converging, as  $n \rightarrow \infty$ , to  $\mu_t$  (in the simplex of dimension  $\mathbf{I}+1$ ). The plan is now defined by a sequence of blocs indexed by  $n$ . On the  $n^{\text{th}}$  bloc (of length  $q^n$ ) the play consists of  $p_1^n$  times  $s_1$ , then  $p_2^n$  times  $s_2, \dots$  and so on. The pay-off associated to  $h$  is clearly  $d$ . This proves the theorem. ■

Note that one could as well define the plan by choosing at each stage  $n$ , the (first in some order) vector of pure moves that minimises the distance from the new average pay-off  $\bar{g}_n$  to  $d$ . This yields a Borel map from  $E$  to equilibria.

**REMARK 4.1.** We prove the above result by using expected stage pay-off. It is worthwhile to notice that it still holds if one considers the pay-off on the play namely the random variables  $\bar{g}_n$  and ask for non-profitable deviation and  $\bar{g}_\infty = d$  a.e. (Use ex. II.4Ex.4 p. 105).

We now turn to similar properties for  $\Gamma_\lambda$  and  $\Gamma_n$ .

Concerning the discounted game, the asymptotic set of equilibria may differ from  $E$  (Forges et al., 1986), as shown by the following 3 person game where player III is a dummy:

$$\begin{pmatrix} (1, 0, 0) & (0, 1, 0) \\ (0, 1, 0) & (1, 0, 1) \end{pmatrix}$$

This being basically a zero-sum game between players I and II the only equilibrium (optimal) strategies are  $(1/2, 1/2)$  i.i.d. in  $\Gamma_\lambda$  as well as in  $\Gamma_n$ , hence the only equilibrium pay-off is  $(1/2, 1/2, 1/4)$ . On the other hand  $E$  contains the point  $(1/2, 1/2, 1/2)$ .

Nevertheless the following generic result holds:

**THEOREM 4.2.** *Assume that there exists some  $d$  in  $E$  with  $d^i > v^i$  for all  $i$ . Then  $E_\lambda$  converges (in the Hausdorff topology) to  $E$  as  $\lambda$  goes to 0.*

**PROOF.** By the hypothesis and the convexity of  $E$ , it is enough to prove that any point  $d$  in  $E$  with  $d^i > v^i$  belongs to  $E_\lambda$  for  $\lambda$  small enough. The idea of the proof is then very similar to the previous one.

We first construct a play leading to  $d$ . Decompose  $d$  according to the extreme points of  $D$  (attainable through pure moves), to get  $d = \sum_{t=1}^{I+1} \mu_t d_t$ . (Note that an approximation of  $d$  for  $\lambda$  small is easy to obtain as in the previous proof but we will obtain here an exact representation — for the use of this result cf. ex. IV.4Ex.7 p. 172, cf. also ex. II.3Ex.2 p. 96).

Assume  $\lambda \leq 1/(I+1)$  then one of the  $\mu_t$ , say  $\mu_1$ , is larger than  $\lambda$  and we can write  $d = \lambda d_1 + (1-\lambda)d(2)$  with  $d(2)$  in  $D$ , or more precisely in  $\{d_t\}$ . Doing the same decomposition with  $d(2)$  we obtain inductively a sequence  $d^n$  in  $\{d_t\}$  with  $d = \sum_{n=1}^{\infty} \lambda(1-\lambda)^{n-1} d^n$ .  $h$  is then defined at stage  $n$  by the moves in  $S$  leading to  $d^n$ , and  $\bar{\gamma}_\lambda(h) = \sum_n \lambda(1-\lambda)^{n-1} d^n = d$ .

Moreover if in the previous construction we use a greedy algorithm choosing at each stage a vector of moves such that  $d(n+1)$  is the closest to  $d$ , we obtain that the sequence  $\{d^n\}_{n \geq m}$  is still a good approximation of  $d$ . Formally:  $\forall d \in D, \forall \varepsilon > 0, \exists \bar{\lambda}, \forall \lambda \leq \bar{\lambda}, \exists h$ , such that  $\gamma_\lambda(h) = d$  and for any decomposition  $h = (h', h'')$  with  $h'$  in  $H$ ,  $\|d - \bar{\gamma}_\lambda(h'')\| \leq \varepsilon$ . ( $h$  is  $\varepsilon$ -adapted to  $d$ ). To finish the proof let  $2\varepsilon = \min_i \|d^i - v^i\|$ , and  $\lambda$  small enough so that the one-shot gain by deviation, at most  $2\lambda C$ , is less than the loss due to the punishment, at least  $(1-\lambda)\varepsilon$ . ■

**REMARK 4.2.** One can show that  $\bar{\gamma}_\lambda(\Sigma) = D$  as soon as  $\lambda \leq 1/I$  (cf. ex. IV.4Ex.1 p. 172).

**REMARK 4.3.** It is also easy to see that the result holds without restriction for  $I = 2$ , or more generally if:  $\exists d \in D, d^i > v^i$  for all but one player  $i$ . (ex. IV.4Ex.6 p. 172).

We finally consider the finitely repeated games, where no similar generic result holds, the classical counter example being the Prisoner's Dilemma described by the following pay-off matrix:

$$\begin{pmatrix} (3, 3) & (0, 4) \\ (0, 4) & (1, 1) \end{pmatrix}$$

and where  $E_n = (\{1, 1\})$  for all  $n$ .

In fact we have (recall that  $v$  is the threat point):

**PROPOSITION 4.3.** *Assume  $E_1 = \{v\}$  then  $E_n = \{v\}$  for all  $n$ .*

**PROOF.** Given an equilibrium  $\sigma$ , let  $m$  be the smallest integer such that after each history of length  $m$  compatible with  $\sigma$ ,  $\sigma$  induces the pay-off  $v$  at each of the remaining stages. By the hypothesis  $m \leq n - 1$ . If  $m > 0$ , consider an history of length  $m - 1$  compatible with  $\sigma$  and where at stage  $(m - 1)$ ,  $\sigma$  does not lead to  $v$ . Clearly one of the players has then a profitable deviation at that stage and cannot be punished later. ■

A sufficient condition for convergence to  $E$  is given by the following:

**THEOREM 4.4. (Benoît and Krishna, 1987)** Assume that for all  $i$  there exists  $e(i)$  in  $E_1$  with  $e^i(i) > v^i$ . Then  $E_n$  converges in the Hausdorff topology to  $E_\infty$  as  $n$  goes to infinity.

**PROOF.** The idea is to avoid backwards induction effects by ending the play by a phase of rewarding or punishment consisting of a fixed number of stages: the influence on the total pay-off will be negligible as  $n$  goes to infinity.

Given a play  $h$  this phase will be a sequence of  $R$  cycles of mixed moves leading to the pay-offs  $(e(1), \dots, e(\mathbf{I}))$  if the previous history follows  $h$  and a sequence of i.i.d. punishments  $x^{-i}(i)$  in case of previous deviation of  $i$ .

Now given  $d$  in  $E$ , let  $\delta = \min_i(e^i(i) - v^i)$  and  $\varepsilon \leq \delta/2$ . Choose an  $\varepsilon$ -rational approximation  $d$  in  $E$  as in the proof of the Folk theorem, corresponding to an history  $h'$  of length  $N$ . Let  $R > 2C/\delta$ , then for  $n \geq 2C(R\mathbf{I} + N)/\varepsilon$ , the strategies defined by a play  $h$  consisting of a cycle of histories  $h'$  (until stage  $n - R\mathbf{I}$ ) followed by the last phase defined above, clearly induce an equilibrium with pay-off within  $2\varepsilon$  of  $d$ . In fact any deviation when some  $h'$  is played will be observed, hence punished. On the other hand if  $h$  is followed during the first phase the second one consists of a sequence of one-shot equilibria where no deviation is profitable. ■

**REMARK 4.4.** It is clear that a sufficient condition for the previous result to hold is:  $\forall i, \exists n_i, \exists e(i) \in E_{n_i}$  with  $e^i(i) > v^i$ .

More precise results can be obtained for  $\mathbf{I} = 2$  (cf. Benoît and Krishna, 1987).

For related results with subgame perfect equilibria, cf. ex. IV.4Ex.3–IV.4Ex.8.

**4.b. Partial monitoring.** Most of the results of this section are due to Lehrer (1990, 1991, 1992a, 1992c).

**4.b.1. Notations and definitions.** Coming back to the general model we assume still here that there is only one state but after each stage  $n$  each player  $i$  is told  $Q^i(s_n)$  where  $Q^i$  is a mapping from  $S$  to  $A^i$ . Note that equivalently  $Q^i(s^i, \cdot)$  can be viewed as a partition of  $S^{-i}$ ,  $Q^i(s^i, s^{-i})$  being the partition element containing  $s^{-i}$ .

Let us first provide a general characterisation of uniform equilibria in supergames:

**PROPOSITION 4.5.**  $d$  is a uniform equilibrium pay-off iff there exists a sequence  $\varepsilon_m$  decreasing to 0,  $N_m$  and  $\sigma_m$ , such that  $\sigma_m$  is an  $\varepsilon_m$  equilibrium in  $\Gamma_{N_m}$  leading to a pay-off within  $\varepsilon_m$  of  $d$ .

**PROOF.** The condition is obviously necessary.

For the sufficiency define super-blocs  $M_m$  as a sequence of  $l_m$  blocs of size  $N_m$  and let  $\sigma$  be: play  $\sigma_m$  on  $M_m$  (i.e. starting with an empty history after each cycle of  $N_m$  moves). Choose  $l_m$  such that  $2CN_{m+1}/(\ell_m \times N_m) \leq \varepsilon_m$ . It follows easily that if  $n \in M_{m+1}$ ,  $\sigma$  is a  $2(\sum_{k \leq m} \varepsilon_k \ell_k N_k / \sum_{k \leq m} \ell_k N_k)$  equilibrium in  $\Gamma_n$ , hence the result.

For correlated equilibria define the auxiliary space as  $\Omega = \prod_m (\Omega_m)^{l_m}$  with the product probability induced by the correlation device on each factor. Due to the independence it is easily verified that announcing the signals at the beginning of the game or at the beginning of each bloc gives the same result. ■

**REMARK 4.5.** The same requirement in a general repeated game would lead to  $E_0$  (cf. VII).

We first consider the two-person case and assume **non-trivial signalling structure** hence for each player  $i = I, II$  and  $j \neq i$  there exists  $s^i$  in  $S^i$  and  $s^j, t^j$  in  $S^j$  satisfying:

$$(\star) \quad Q^i(s^i, s^j) \neq Q^i(s^i, t^j)$$

so that the players can communicate through their moves (the other case is much simpler to analyse (cf. ex. IV.4Ex.9 p. 173)).

Since in general the signals are not common knowledge, equilibrium strategies  $\sigma$  do not induce, conditionally to an atom of the  $\sigma$ -algebra of common knowledge events (finest  $\sigma$ -algebra containing all  $\mathcal{H}^i$ ), having positive probability under  $\sigma$ , an equilibrium, but rather a correlated equilibrium. In fact those are much easier to deal with.

**DEFINITION 4.6.** We define two relations between moves by:

$$s^i \sim t^i \Leftrightarrow Q^j(s^i, s^j) = Q^j(t^i, s^j) \text{ for all } s^j \text{ in } S^j$$

(this means that after one stage, player  $j$  has no possibility to distinguish whether  $i$  was playing  $s^i$  or  $t^i$ ) and:

$$\begin{aligned} s^i \succ t^i &\Leftrightarrow [(s^i \sim t^i) \text{ and } (Q^i(t^i, s^j) \neq Q^i(t^i, t^j) \text{ implies} \\ &Q^i(s^i, s^j) \neq Q^i(s^i, t^j) \text{ for all } s^j, t^j \text{ in } S^j)] \end{aligned}$$

(player  $i$  is always getting more information on  $j$ 's move by playing  $s^i$  rather than  $t^i$ ).

Then we have:

**LEMMA 4.7.** Given a pure strategy  $\sigma^i$ , at each history  $h$  player  $i$  can use any action  $t^i$  rather than  $\sigma^i(h) = s^i$  with  $t^i \succ s^i$ , while still inducing the same probability distribution on  $\mathcal{H}^j$ .

**PROOF.** By playing  $t^i$  the signal to player  $j$  will be the same. Now for the next stage since the partition on player  $j$ 's moves is finer with  $t^i$  than with  $s^i$ , player  $i$  can deduce what would have been his signal if he had played  $s^i$  and play accordingly in the future. ■

Let  $\Delta(S)$  be the set of probabilities on  $S$  (correlated moves) and extend the pay-off to  $\Delta(S)$  by integration. The sets of equilibrium pay-offs will be characterised through the following sets:

$$\begin{aligned} \mathbf{A}^i &= \{ P \in \Delta(S) \mid t^i \succ s^i \Rightarrow \sum_{s^j} P(s^i, s^j) G^i(s^i, s^j) \geq \sum_{s^j} P(s^i, s^j) G^i(t^i, s^j) \} \\ \mathbf{B}^i &= \mathbf{A}^i \cap X = \{ x \in X \mid t^i \succ s^i \Rightarrow x^i(s^i) G^i(s^i, x^j) \geq x^i(s^i) G^i(t^i, x^j) \} \end{aligned}$$

**REMARK 4.6.**  $\mathbf{A}^i = \text{Co}(\mathbf{B}^i)$ . In fact  $\mathbf{A}^i$  is convex and contains  $\mathbf{B}^i$ . Now given  $P$  in  $A^i$ , for any  $s^i$  with  $\rho(s^i) = \sum_{s^j} P(s^i, s^j) > 0$ ,  $(s^i, y^j) \in \mathbf{B}^i$  with  $y^j(s^j) = P(s^i, s^j)/\rho(s^i)$ .

**REMARK 4.7.** Note that like in the Folk theorem these sets are defined through the one-shot game only.

**REMARK 4.8.** It will turn out that the sets of equilibrium pay-offs for our rather strong definition in sect. 1.d p. 151 will exist:  $E_\infty$  for equilibria and  $C_\infty, D_\infty$  for correlated and communication equilibria. (For liminf pay-offs, cf. ex. IV.4Ex.8 p. 173).

4.b.2. *Correlated and communication equilibria.* Let us characterise correlated and communication equilibrium pay-offs. We write  $IR$  for the set of individually rational pay-offs; it will obviously contain any equilibrium pay-off.

**THEOREM 4.8.**  $D_\infty = C_\infty = G(\bigcap_i \mathbf{A}^i) \cap IR$

**PROOF.** We first prove inclusion of  $D_\infty$  in the right hand set. Assume that  $d = (d^I, d^{II})$  is an  $\mathcal{L}$ -equilibrium pay-off not in  $G(\bigcap_i \mathbf{A}^i)$  (the inclusion in  $IR$  is standard, cf. theorem 4.1 p. 163 above).

Denoting by  $P_n$  the correlated distribution on  $S$  induced by the equilibrium strategies at stage  $n$  and  $\bar{P}_n = (1/n) \sum_1^n P_m$ , one has  $d = G(\bar{P})$ , with  $\bar{P} = \mathcal{L}(\bar{P}_n)$ .

Given any  $P$  in  $\Delta(S)$  let us define  $P^i$  in  $\mathbf{A}^i$  as follows: First introduce a map  $\phi^i$  from  $S^i$  to  $S^i$  such that  $\sum_{s^j} P(s^i, s^j) G^i(\phi(s^i), s^j)$  maximises  $\sum_{s^j} P(s^i, s^j) G^i(t^i, s^j)$  on the set  $\{t^i \succ s^i\}$  and let then  $P^i(t^i, s^j) = \sum_{s^i, \phi(s^i)=t^i} P(s^i, s^j)$ .

In words, we replace any action of  $i$  in the support of  $P$  by a best reply against  $P$ , among the more informative moves.

Assuming  $\bar{P} \notin \mathbf{A}^I$ , define an alternative strategy of player I as follows: At each stage use  $\phi^I(s^I)$  (computed for  $\bar{P}$ ) rather than  $s^I$  and keep for the following stages the signal that would have been obtained by using  $s^I$ .

We obtain thus (using lemma 4.7 p. 166):

$$\mathcal{L}(\bar{\gamma}_n^I(\tau^I, \sigma^{II})) = G^I(\bar{P}^I) > G^I(\bar{P})$$

hence the contradiction, and the required inclusion.

(Note that the same inclusion holds for more than 2 players with the appropriate extension of  $\succ$ ).

Consider now  $P$  in  $\bigcap_i \mathbf{A}^i$  with  $G(P)$  in  $IR$ . By the previous prop. 4.5 p. 165, it is enough to construct, for any  $\varepsilon_0 > 0$ , an  $\varepsilon_0$ -equilibrium in a finite game with pay-off within  $\varepsilon_0$  of  $G(P)$ . Let  $\varepsilon_0 = 8\varepsilon$ .

Using the hypothesis of non-trivial signalling structure we can define an injective mapping from  $A^i$  to finite sequences of signals for  $j$  (for example a finite sequence of elements in  $\{Q^j(s^i, t^j), Q^j(t^i, t^j)\}$  satisfying  $(\star)$ ), so that both players have a code to report their signal at some stage in a bounded (say by  $B$ ) number of stages.

The strategies will be defined on blocs of stages as follows: Let  $\eta < \varepsilon/4C$ ,  $n$  such that  $(2n + 2B)/2^n \leq \eta$  and define  $N_1 = 2^n + 2n + 2B$ . We now describe the behaviour on a bloc of size  $N_1$ .

During the first  $2^n$  stages the players are requested to play according to some correlation device  $\bar{R}$ . Define first a probability  $R$  on  $\Omega^I \times \Omega^{II} = S^I \cup (S^I \times S^{II}) \times (S^{II} \cup (S^I \times S^{II}))$  by the following procedure: Take the convex combination of the uniform distribution on  $S$  (with coefficient  $\eta$ ) and  $P$  (coefficient  $(1 - \eta)$ ) and independently announce with probability  $\eta$  to one of the players the move of his opponent. Namely:

$$\begin{aligned} R(s) &= (\eta / (\#S) + (1 - \eta)P(s)) / (1 + 2\eta) \\ R(s^i, \{s^i, s^j\}) &= \eta R(s) \end{aligned}$$

Then  $\bar{R}$  is the product of  $2^n$  independent copies of  $R$ , a signal is selected in  $(\Omega^I \times \Omega^{II})^{2^n}$  according to it, its component on  $(\Omega^i)^{2^n}$  is transmitted to  $i$  who is supposed to follow the projection on  $(S^i)^{2^n}$ . Note that with positive probability at each stage every move is played and every move that player  $i$  has to play is announced with positive probability to player  $j$ .

During the next  $n$  stages player I plays an i.i.d. mixture  $(1/2, 1/2)$  on the moves  $(s^I, t^I)$  and player II uses  $s^{II}$  with  $Q^{II}(s^I, s^{II}) \neq Q^{II}(t^I, s^{II})$  and symmetrically for the following next  $n$  stages. These random moves are used to generate random times say  $\theta^{II}$  and  $\theta^I$  independent and uniformly distributed on the previous  $2^n$  stages and public knowledge.

Finally during the last  $B$  stages the previously defined code is alternatively used by each player  $i$  to report the signal he got at stage  $\theta^i$ .

This ends the description of the strategies “on the equilibrium path” on these  $N^1$  stages and remark that the corresponding pay-off is within  $4\eta C$  (hence less than  $\varepsilon$ ) of  $G(P)$ .

Consider now a collection of  $M$  blocs of size  $N_1$ , call it a super-bloc  $N_2$  and finally define  $N$  as a family of  $M'$  super-blocs  $N_2$ .

Define strategies in  $\Gamma_N$  as independent sequences of strategies as above on each bloc of size  $N_1$ . Namely the signals are chosen according to the product of independent probabilities and announced at the beginning of the game. On each bloc of size  $N_1$  the players play according to the corresponding component of their sequence of signals during the first  $2^n$  stages, then as described above. We shall say that a deviation  $\tau^i$  for player  $i$  is  $\varepsilon$  profitable in  $\Gamma_N$  if  $\gamma_N^i(\tau^i, \sigma^j) > G^i(P) + \varepsilon$  and prove by specifying the strategies “in case of detected deviation”, that for  $N$  large enough there is no  $6\varepsilon$  profitable strategies, hence  $\sigma$  is a  $\varepsilon_0$  equilibria.

We first show that on blocs of size  $N_1$  any  $3\varepsilon$  profitable deviation has a strictly positive probability  $\pi$  of being detected. Remark in fact that on each such bloc, a deviation near the end (i.e. during the last  $2n + 2B$  stages) modifies the pay-off on this block by less than  $\varepsilon/2$ . Hence we only consider deviation at other stages. Note that on these stages the move of each player is independent of the past, hence it is sufficient to consider history-independent deviation. Now by the choice of  $R$  if  $\tau^i$  gives  $2\varepsilon$  more than  $\sigma^i$  against  $R$ , the gain is at least  $\varepsilon$  against  $P$ . Recall that by playing  $t^i$  rather than  $s^i$  and  $t^i \succ s^i$ , the pay-off is not increased since  $P$  belongs to  $A^i$ . On the other hand if  $t^i \not\succ s^i$  there is a positive probability that player  $j$  is told  $i$ 's recommendation i.e.,  $s^i$  and is using at this stage a revealing move namely  $t^j$  with either a)  $Q^j(s^i, t^j) \neq Q^j(t^i, t^j)$ , or b)  $Q^i(s^i, t^j) \neq Q^i(t^i, t^j)$ . In case a)  $j$  observes  $i$ 's deviation at that stage; in case b), there is a positive (and independent) probability that player  $i$  will be asked to report his signal at that stage and will give a wrong answer observable by  $j$ .

Obviously the strategy of each player in case of wrong signal is to punish for ever, hence to reduce the pay-off to at most  $\gamma^i(P)$  (recall that  $\gamma(P)$  is  $IR$ ).

Define now  $M$  (the number of blocs  $N_1$  in  $N_2$ ) such that  $(1 - \pi)^{\varepsilon M/C} \leq \varepsilon/2C$ . It follows then from the above computations that if  $\tau^i$  is  $4\varepsilon$  profitable on a  $N_2$  bloc, a deviation will be detected with probability at least  $1 - \varepsilon/2C$ .

Define finally  $M'$  so that the relative size of a bloc  $N_2$  in games of length  $N = M'N_2$  is at most  $\varepsilon/2C$ . We obtain:

$$\gamma_N^i(\tau^i, \sigma^j) \leq \varepsilon + (1 - \varepsilon/2C)(\gamma^i(P) + 4\varepsilon + 2C/M') \leq \gamma^i(P) + 6\varepsilon$$

hence the result. ■

The main difficulties in trying to extend the previous result to equilibria are:

- The lack of common knowledge events on which to condition the analysis, while remaining in an equilibrium framework, i.e. without being led to correlated equilibria.

- The fact that one cannot restrict the players to use pure actions and the right equivalence classes of mixed moves are hard to define.

A simple and useful result uses the following set:

$$\mathbf{C}^i = \{x \in X \mid y^i \in X^i, Q^j(y^i, \cdot) = Q^j(x^i, \cdot) \Rightarrow G^i(y^i, x^{-i}) \leq G^i(x^i, x^{-i})\}$$

( $Q^j(x^i, \cdot)$  is a mapping from  $S^j$  to probabilities on  $j$ 's signals.)

Note that  $\mathbf{C}^i \subseteq \mathbf{B}^i$ .

**PROPOSITION 4.9.** *All points of  $\text{Co } G(\bigcap_i \mathbf{C}^i) \cap IR$  are uniform equilibrium pay-offs.*

**PROOF.** We first prove that  $d$  in  $G(\bigcap_i \mathbf{C}^i) \cap IR$  is a uniform equilibrium pay-off. So fix  $x$  in  $\bigcap_i \mathbf{C}^i$  with  $G(x) = d$ . Each player is required to play  $x_\varepsilon^i$  i.i.d. where  $x_\varepsilon^i$  is  $(1-\varepsilon)x^i + \varepsilon\bar{y}^i$  ( $\bar{y}^i$  uniform on  $S^i$ ). The checking is done at the end of blocs of increasing sizes, say  $N_\ell = 2^\ell$ . Player  $i$  is punished for  $N_\ell C/\varepsilon$  stages if the empirical distribution of player  $j$ 's signals on bloc  $\ell$  differs by more than  $\varepsilon$  from what it should be according to  $Q^j(x_\varepsilon^i, \cdot)$ , namely if  $\max_{t^j} \max_{a^j} |Q^j(x_\varepsilon^i, t^j)(a^j) - \#\{1 \leq m \leq 2^\ell; a_m^j = a^j; t_m^j = t^j\}/\#\{1 \leq m \leq 2^\ell; t_m^j = t^j\}| \geq \varepsilon$ . After the punishment phase one plays on bloc  $\ell + 1$ .

Let  $\varepsilon_0 = 6C\varepsilon$ . It is easy to check that the strategies described above will induce an  $\varepsilon_0$ -equilibrium with pay-off within  $\varepsilon_0$  of  $d$  in a sufficiently long game, hence the result by prop. 4.5 p. 165. (See sect. 3.c p. 195. for similar computations made in more detail).

If  $d \in \text{Co } G(\bigcap_i \mathbf{C}_i) \cap IR$ , let just alternate between plays defined as above, corresponding to different points in  $G(\bigcap_i \mathbf{C}_i)$ . ■

We will study now specific cases.

4.b.3. *Observable pay-offs.* We still consider 2 players but we assume here that the signal reveals the pay-off, namely:

$$G^i(s^i, s^j) \neq G^i(s^i, t^j) \Rightarrow Q^i(s^i, s^j) \neq Q^i(s^i, t^j) \text{ for all } i, s^i, s^j, t^j.$$

This signalling structure implies specific properties on the pay-offs like the following.

**LEMMA 4.10.**  $\forall x \in X$ , (resp.  $\mathbf{C}^i, \mathbf{B}^i$ ),  $\exists y^j$  such that  $(x^i, y^j) \in X \cap \mathbf{C}^j$ , (resp.  $\mathbf{C}^i \cap \mathbf{C}^j, \mathbf{B}^i \cap \mathbf{C}^j$ ) and  $G^i(., x^j) = G^i(., y^j)$ .

**PROOF.** Given  $x^j$ , let  $y^j$  satisfying  $Q^i(x^j, \cdot) = Q^i(y^j, \cdot)$  and  $(x^i, y^j) \in \mathbf{C}^j$ . The previous equality implies that  $i$ 's pay-off against  $x^j$  or  $y^j$  is the same whatever be his strategy. Hence if  $(x^i, x^j) \in \mathbf{C}^i$  (resp.  $\mathbf{B}^i$ ) we still have  $(x^i, y^j) \in \mathbf{C}^i$  (resp.  $\mathbf{B}^i$ ). ■

Recall that  $v^i$  is the minmax for  $i$  and that  $x^{-i}(i)$  realises it.

**LEMMA 4.11.** *There exists a point  $d_0$  in  $G(\bigcap_i \mathbf{C}^i)$  with  $d_0^i \leq v^i$ ,  $i = I, II$ .*

**PROOF.** Start with  $x$  defined by the punishing strategy  $x^i(j)$ ,  $i = I, II$  and use lemma 4.10 for both. ■

**LEMMA 4.12.**  $\text{Co } G(\bigcap_i \mathbf{B}^i) \cap IR = \bigcap_i \text{Co } G(\mathbf{B}^i) \cap IR$ .

**PROOF.** Consider  $d$  in  $\bigcap_i \text{Co } G(\mathbf{B}^i)$  with  $d^j$  maximal among the points in this set with the same  $d^i$ . Decompose then  $d$  as a barycentre of pay-offs from points in  $\mathbf{B}^i$ , say  $d = \sum_k \lambda_k G(x_k)$ . We now prove that  $x_k$  is in  $\mathbf{B}^j$  for all  $k$ . In fact like in lemma 4.10 above, one could otherwise define  $y_k$  in  $\mathbf{B}^i$  inducing the same pay-off to  $i$  and a strictly better one to  $j$ . Averaging over the  $\lambda$ 's would give a point in  $\text{Co } G(\bigcap_i \mathbf{B}^i)$  with a higher  $d^j$ .

Now any point  $d$  in  $\bigcap_i \text{Co } G(\mathbf{B}^i) \cap IR$  is in the convex hull of  $\bar{d}$  (same  $d^i$  and  $d^j$  maximal on  $\bigcap_i \text{Co } G(\mathbf{B}^i)$ ),  $\underline{d}$  (defined dually) and  $d_0$  (introduced in lemma 4.11). These three points being in  $G(\bigcap_i \mathbf{B}^i)$  the result follows. ■

Given a set  $D \subseteq \mathbb{R}^n$ , its admissible points are those  $x \in D$ , such that  $y \geq x$ ,  $y \in D$  implies  $y = x$ .

LEMMA 4.13. *The set  $Ad$  of admissible points of  $G(\bigcap_i \mathbf{B}^i)$  is included in  $G(\bigcap_i \mathbf{C}^i)$ .*

PROOF. Use again, like in lemma 4.10, the fact that if  $x$  is not in  $\mathbf{C}^i$  one can modify  $x^i$  to  $y^i$  such that  $(y^i, x^j) \in \mathbf{C}^i \cap \mathbf{B}^j$ , player  $i$ 's pay-off is increased, player  $j$ 's pay-off is the same. ■

THEOREM 4.14.

$$\begin{aligned} E_\infty &= C_\infty = D_\infty = \text{Co } G(\bigcap_i \mathbf{B}^i) \cap IR \\ &= \bigcap_i \text{Co } G(\mathbf{B}^i) \cap IR \end{aligned}$$

PROOF. The inclusion is clear from theorem 4.8 since  $\text{Co } G(\mathbf{B}^i)$  equals  $G(\mathbf{A}^i)$ .

We will represent all points in  $\text{Co } G(\bigcap_i \mathbf{B}^i) \cap IR$  as barycentres of points where "simple" strategies will be used. We have already two components:  $d_0$  (lemma 4.11) and the admissible part  $Ad$  (lemma 4.13).

Consider now pay-offs on the other part of the boundary:

LEMMA 4.15. *Let  $d$  be an extreme point of  $\text{Co } G(\bigcap_i \mathbf{B}^i) \setminus Ad$  such that  $d^i$  is maximal among the points having in this set the same  $d^j$ . Then  $d = G(x)$  for some  $x$  with:*

- (1)  $x^j = s^j$  is a pure move.
- (2)  $x^i$  is a best reply to  $s^j$  among the set of  $y^i$  satisfying

$$Q^j(y^i, s^j) = Q^j(x^i, s^j).$$

- (3)  $t^j \succ s^j$  implies that  $t^j$  is a duplicate of  $s^j$  (i.e. same signals and pay-offs to both players).

PROOF. Let  $d = G(y)$  with  $y \in \bigcap_i \mathbf{B}^i$ . Note first that  $y \in \mathbf{C}^i$  (otherwise like in lemma 4.10 p. 169 one could find  $y'$  in  $\mathbf{C}^i \cap \mathbf{B}^j$  with same pay-off for  $j$  and higher pay-off for  $i$ . Let us now prove that one can take  $y^j$  pure. Otherwise for each  $t^j$  in the support of  $y^j$ ,  $(y^i, t^j) \in \mathbf{B}^j$ . By lemma 4.10 again there exists  $x^i(t^j)$  such that  $(x^i(t^j), t^j) \in \mathbf{C}^i \cap \mathbf{B}^j$ , keeping the same pay-off for  $j$  and increasing  $i$ 's pay-off. Hence if  $y^j = \sum \alpha(t^j)t^j$ , one obtains  $d^j = \sum \alpha(t^j)G^j(x^i(t^j), t^j)$  and  $d^i \leq \sum \alpha(t^j)G^i(x^i(t^j), t^j)$ .  $d$  being a point of  $\text{Co } G(\bigcap_i \mathbf{B}^i)$  with maximal  $d^i$  on the  $d^j$  line implies that the second relation is an equality, hence  $(y^i, t^j) \in \mathbf{B}^i \cap \mathbf{B}^j$  for all  $t^j$ ;  $d$  being extreme finally implies  $d = G(y^i, t^j)$ .

We can moreover assume that no  $s^j$  satisfies  $s^j \succ t^j$  and  $t^j \not\succ s^j$ . Otherwise denote again here by  $s^j$  a maximal point for this preorder, maximising  $G^j(y^i, \cdot)$ , among those. If  $G^j(y^i, s^j) = G^j(y^i, t^j)$  take  $(y^i, s^j)$  as new initial point (obviously in  $\mathbf{C}^i \cap \mathbf{B}^j$ ). If not, then  $G^j(y^i, s^j) < G^j(y^i, t^j)$ , (recall that  $(y^i, t^j)$  belongs to  $\mathbf{B}^j$ ), and  $(y^i, s^j)$  in  $\mathbf{C}^i \cap \mathbf{B}^j$  induces a pay-off  $d'$  with  $d_i = d'_i$  and  $d_j > d'_j$  contradicting the choice of  $d$  ( $d \notin Ad$  and is extreme).

Coming back to our pair  $(y^i, t^j)$ , assume that there exists  $x^i$  with  $Q^j(x^i, t^j) = Q^j(y^i, t^j)$  and  $G^i(x^i, t^j) > G^i(y^i, t^j)$ . One can then even assume  $(x^i, t^j) \in \mathbf{C}^i$  (like in lemma 4.10). The choice of  $d$  implies then  $(x^i, t^j) \notin \mathbf{B}^j$ .

By the above remark on  $t^j$  letting  $(x^i, s^j) \in \mathbf{B}^j$  with  $s^j \succ t^j$ , one has also  $t^j \succ s^j$ , so that:  $Q^j(x^i, t^j) = Q^j(y^i, t^j)$  implies  $Q^j(x^i, s^j) = Q^j(y^i, s^j)$ . Now  $Q^j(x^i, u^j) = Q^j(y^i, u^j)$  implies

$G^j(x^i, u^j) = G^j(y^i, u^j)$  for all  $u^j$ , hence  $G^j(x^i, t^j) < G^j(x^i, s^j)$  implies  $G^j(y^i, t^j) < G^j(y^i, s^j)$  and this again contradicts the fact that  $(y^i, t^j)$  is in  $\mathbf{B}^j$ . ■

We can now describe the equilibrium strategies: decompose an *IR* pay-off  $d$  as a (finite) convex combination of pay-offs induced by elements in  $\mathbf{C}^i \cap \mathbf{C}^j$  or extreme points of  $\text{Co } G(\bigcap_i \mathbf{B}^i) \setminus Ad$ , say  $d = \sum_k \alpha(k)d(k)$ . A play corresponding to the pay-off  $d(k)$  will be used on a fraction  $\alpha(k)$  of the stages (using as usual rational approximation). It is thus sufficient to describe these plays: for points in  $G(\bigcap_i \mathbf{C}^i)$  use prop. 4.9 p. 169, for the other points we use the previous lemma 4.15: let  $d = G(x^i, s^j)$  satisfying the corresponding properties. Player  $j$  will be asked to play always  $s^j$  and player I to play i.i.d. some perturbation  $x_\varepsilon^i$  of  $x^i$  with strictly positive probability on each move.

It follows then easily from the properties of  $(x^i, s^j)$  that player  $i$  has no profitable non-detectable deviations (like in the previous theorem,  $j$  uses  $x^j(i)$  for finitely many stages if the empirical distribution of signals differs too much from  $Q^j(x^i, s^j)$ ).

On the other hand to check that  $j$  does not deviate, since  $i$  is playing completely mixed, the arguments in the proof of the previous theorem 4.8 p. 167 still apply: player  $j$  will be repeatedly asked to report his signal at some random move and in case of profitable deviation with positive probability his answer will be wrong. ■

4.b.4. “semi-standard” information. We end now this section by presenting a result concerning the **I** person case where the signal received by each player is public and independent of his own move.

The actions set  $S^i$  are equipped with a partition  $\tilde{S}^i$  and after each stage every player is only informed of the element of the product partition that contains the vector of moves. Denote by  $\tilde{x}^i$  the probability induced by  $x^i$  on  $\tilde{S}^i$  then the previous sets  $\mathbf{B}^i$  are now:

$$\mathbf{D}^i = \{x \in X \mid G^i(x) \geq G^i(y^i, x^{-i}) \text{ for all } y^i \text{ with } \tilde{y}^i = \tilde{x}^i\}$$

and they coincide with the previous  $\mathbf{C}^i$ .

PROPOSITION 4.16.  $E_\infty = \text{Co } G(\bigcap_i \mathbf{D}^i) \cap IR$

PROOF. The proof that any pay-off in the right hand side set can be achieved as an equilibrium pay-off is obtained as in prop. 4.9 p. 169 above, with a simpler proof. (Note that the statistics on the signals send by  $i$  is common knowledge and there is no need for  $x_\varepsilon$ )

To get the other inclusion we explicitly use the fact that there exists a “common knowledge”  $\sigma$ -algebra, conditionally on which the moves of the players are still independent. Indeed, the initial strategy  $\sigma$  can be replaced by  $\tilde{\sigma}$  where at each stage  $n$  and for each player  $i$ ,  $\sigma_n^i(h_n)$  is modified to  $\tilde{\sigma}_n^i(h_n) = E(s_n^i \mid \tilde{h}_n)$  with  $h_n$  in  $\tilde{h}_n$ , without changing the pay-off or the equilibrium condition.

On the corresponding events  $\tilde{h}_n$ , if  $\tilde{\sigma}$  is not in  $\mathbf{D}^i$ , player  $i$  can profitably deviate, without being detected (as in lemma 4.7 p. 166). Denoting by  $\mathbf{D}_\varepsilon^i$  an  $\varepsilon$  neighbourhood of  $\mathbf{D}^i$  the equilibrium condition leads to  $\mathcal{L}(\Pr(\tilde{\sigma}_n \notin \mathbf{D}_\varepsilon^i)) = 0$ , for all  $i$  and all positive  $\varepsilon$ . So that  $\mathcal{L}(\Pr(\tilde{\sigma}_n \notin \bigcap_i \mathbf{D}_\varepsilon^i)) = 0$ , hence also  $\mathcal{L}(\Pr(\tilde{\sigma}_n \notin (\bigcap_i \mathbf{D}^i)_\varepsilon)) = 0$ , for all positive  $\varepsilon$ .

Letting  $\sigma^*$  such that  $\sigma^*(h)$  is in  $\mathbf{D}^i$  a closest point to  $\tilde{\sigma}(h)$ , for all  $h$ , one obtains:  $\mathcal{L}(\gamma_n(\sigma^*)) = \mathcal{L}(\gamma_n(\tilde{\sigma}))$ , hence the result since the first term is in  $\text{Co } G(\bigcap_i \mathbf{D}^i)$  by ex. I.2Ex.13 p. 24 and ex. I.3Ex.10d p. 35). ■

REMARK 4.9. As the partitions become finer the equilibrium set increases (basically the set of non-detectable deviations is smaller): one goes from the convex hull of  $E_1$  (equilibrium pay-offs set of  $\Gamma_1$ ) to the set of feasible i.r. pay-offs (theorem 4.1 p. 163).

REMARK 4.10. Note that  $E_\infty$  may differ from  $CE_\infty$  (take a game where  $\text{Co } E_1 \neq CE_1$  and let  $\tilde{S}^i = \{S^i\}$ ).

### Exercises.

*Full monitoring is assumed in exercises IV.4Ex.1–IV.4Ex.7.*

**1.** Use the fact that  $\gamma_1(X)$  is connected and ex. I.3Ex.10 p. 34. to prove that  $D_\lambda = D$  for  $\lambda \leq 1/\#\mathbf{I}$ . Show that it is the best bound.

**2.** Prove theorem 4.2 p. 164 in the general case for  $\#\mathbf{I} = 2$ .

*In the following exercises on subgame perfect equilibria we will denote by  $E'_\infty$ ,  $E'_n$ ,  $E'_\lambda$  the set of subgame perfect equilibrium pay-offs in  $\Gamma_\infty$ ,  $\Gamma_n$ ,  $\Gamma_\lambda$ .*

**3. Perfect Folk Theorem.** (Aumann and Shapley, 1976), (Rubinstein, 1977) Prove that  $E'_\infty = E_\infty = E$ .

HINT. In the proof of theorem 4.1 p. 163 it suffices to punish the last deviator (say at stage  $n$ ) during  $n$  stages, then revert to the original plan (ignoring deviations during punishment phases).

**4. A property of subgame perfect equilibria in discounted multi-move games.** Say that  $\tau^i$  is a **one-stage deviation** from  $\sigma^i$  if  $\tau^i$  coincides with  $\sigma^i$  except at some history  $h_n$ .

Prove that an  $n$ -tuple  $\sigma$  is a subgame perfect equilibrium in  $\Gamma_\lambda$  iff there is no profitable one-stage deviation in the subgame starting at that stage.

HINT. No stationary structure is required: we only need the total pay-off to be the discounted sum of the uniformly bounded stage pay-off, in fact simply continuous.

Use the continuity to reduce to the case where  $\sigma$  is always played from some stage on and look at the last stage where a deviation is still profitable.

**5. A recursive formula for subgame perfect equilibria in discounted games.** Given a bounded set  $F$  of  $\mathbb{R}^{\mathbf{I}}$ , let  $\phi_\lambda(F)$  be the set of equilibrium pay-offs of the one-shot games with pay-off  $\lambda\gamma_1 + (1 - \lambda)f$ , where  $f$  is a mapping from  $S$  (histories at stage 2) to  $F$ . Prove that  $E'_\lambda$  is the largest (for inclusion) bounded fixed point of  $\phi_\lambda$ .

HINT. Assume  $F \subseteq \phi_\lambda(F)$  and construct inductively a sequence of future expected pay-offs and adapted equilibria. Prove that the strategies defined by this sequence induce the same future expected pay-off so that  $F \subseteq E'_\lambda$  by using the previous exercise.

**6.** Prove theorem 4.2 p. 164 under the assumption:  $\exists d \in D, d^i > v^i$ , for all  $i \neq 1$ .

HINT. Show that either the condition of theorem 4.2 holds or there exists no feasible pay-off with  $d^1 > v^1$ .

**7.** (Fudenberg and Maskin, 1986) Consider the following 3 person game (player I chooses the row, player II the column, player III the matrix):

$$\begin{pmatrix} (1, 1, 1) & (0, 0, 0) \\ (0, 0, 0) & (0, 0, 0) \end{pmatrix} \quad \begin{pmatrix} (0, 0, 0) & (0, 0, 0) \\ (0, 0, 0) & (1, 1, 1) \end{pmatrix}$$

Prove that any point  $(z, z, z)$  in  $E'_\lambda$  (or in  $E'_n$ ) satisfies  $z \geq 1/4$ , for all  $\lambda \in (0, 1)$  (and all  $n \geq 1$ ). Compute  $E$  and compare with theorem 4.2 p. 164.

**8. Perfect equilibria in discounted games.** (Fudenberg and Maskin, 1986) Assume  $E$  with non-empty interior. Prove that  $E'_\lambda$  converges to  $E$ .

HINT. Consider  $d$  in  $D$  such that the ball centred at  $d$  with radius  $3\varepsilon$  is included in  $E$  and let  $h(\varepsilon, \lambda)$  be  $\varepsilon$ -adapted to  $d$  (i.e., such that  $\gamma_\lambda(h) = d$  and  $|\gamma_\lambda(h'') - d| \leq \varepsilon$  for any decomposition  $h = (h', h'')$ ). Define  $R$  such that  $2C < R\varepsilon$  and for each  $i$  and every history  $\ell$  of length  $R$ , let  $d(\ell, i)$  be defined by  $d^j = \sum_{n=1}^R \lambda(1-\lambda)^{n-1} \gamma_n^j(\ell_n) + (1-(1-\lambda)^R)d^j(\ell, i)$  for  $j \neq i$ , and by  $d^i(\ell, i) = d^i - 2\varepsilon$ . Observe that for  $\lambda$  small enough one has  $d(\ell, i) \in E$  and  $|d^j(\ell, i) - d^j| \leq \varepsilon$ , and prove that for  $\lambda$  small enough, there exists  $\tilde{h}(\ell, i) (= \tilde{h})$   $\varepsilon$ -adapted to  $d(\ell, i)$  with moreover  $\bar{\gamma}_\lambda^i(\tilde{h}'') \geq d^i(\ell, i)$  for all decompositions  $\tilde{h} = (\tilde{h}', \tilde{h}'')$ .

Consider now the following strategies  $\sigma$ : Play  $h$ , if  $i$  deviates use  $R$  times  $x^{-i}(i)$  then if  $\ell$  have been achieved follow  $\tilde{h}(\ell, i)$ . Inductively if  $\tilde{h}(\ell', j)$  is played and player  $k$  deviates use  $R$  times  $x^{-k}(k)$  then if  $\ell''$  results follow  $\tilde{h}(\ell'', k)$ . Deviations during the  $R$  stages where some punishment strategy is used are ignored.

**9. Lower equilibrium pay-offs (Nontrivial signalling).** (Lehrer, 1989) Say that  $\sigma$  is a lower equilibrium if:  $\bar{\gamma}_n(\sigma)$  converges to some  $\gamma(\sigma)$  and for all  $i$  and all  $\tau^i$ ,  $\liminf \bar{\gamma}_n^i(\tau^i, \sigma^{-i}) \geq \gamma^i(\sigma)$ . Denote by  $lE_\infty$  and  $lC_\infty$  the corresponding sets of equilibrium and correlated equilibrium pay-offs.

a. Prove that  $lC_\infty \subseteq G(\mathbf{A}^i)$ .

HINT. Prove like in theorem 4.8 that otherwise player  $i$  has profitable deviations on a set of stages with positive lower density.

b. Prove that  $\bigcap_i \text{Co } G(\mathbf{B}^i) \cap IR \subseteq lE_\infty$ .

HINT. Define a sequence of blocs of length  $\ell_n$  with  $\ell_n / \sum_{k \leq n} \ell_k \rightarrow \infty$ . On odd blocs approximate  $d$  in  $\bigcap_i \text{Co } G(\mathbf{B}^i)$  by a rational combination of points in  $g(\mathbf{B}^i)$ , where moreover player I uses a pure move. Let then player II use an i.i.d. sequence of perturbations with full support of his previous mixed move. Player I is thus checked on odd blocs (deviations where  $t^i \not\sim s^i$  will be detected with high probability and punished during a large finite number of stages). For the other deviations he is asked to report, using the usual code, at stage  $n^2$  the signal he got at stage  $n$ . A similar construction holds for II on even blocs and punishments are forever.

c. Deduce that  $lE_\infty = lC_\infty = \bigcap_i G(\mathbf{A}^i) \cap IR$

d. Show that  $\bigcap_i G(\mathbf{A}^i) \cap IR \neq G(\bigcap_i \mathbf{A}^i) \cap IR$  and similarly for  $\mathbf{B}^i$ .

HINT. Take  $\begin{pmatrix} (2,2) & (1,1) \\ (1,1) & (0,0) \end{pmatrix}$  and white signals (duplicating strategies if necessary).

**10. Trivial signalling.** Prove that if player I has trivial information (i.e.,  $Q^I(s^I, \cdot)$  constant on  $S^{\text{II}}$  for all  $s^I$ ) the previous results (theorem 4.8 p. 167, Theorem 4.14 p. 170 and ex. IV.4Ex.9 p. 173) hold with  $\mathbf{A}^i$  replaced by:

$$\tilde{\mathbf{A}}^i = \{ P \in \Delta(S) \mid \sum_{s^j} P(s^i, s^j) G^i(s^i, s^j) \geq \sum_{s^j} P(s^i, s^j) G^i(t^i, s^j) \text{ for all } s^i, t^i \text{ with } t^i \sim s^i \}$$

and  $\mathbf{B}^i$  by  $\mathbf{C}^i$ .

**11. Correlated equilibria with semi-standard information.** (Naudé, 1991) Prove that, in the framework of 4.b.4 p. 171,  $C_\infty = D_\infty = G(\bigcap_i \mathbf{A}^i) \cap IR$ .

**12. A constructive approach to  $E_\infty$ .** (Lehrer, 1992b) Given  $\varepsilon > 0$ , define:

$$C_\varepsilon^i = \{x \in X \mid G^i(x^i, x^j) \geq G^i(y^i, x^j) - \varepsilon \text{ for all } y^i \in X^i \text{ with } Q^j(x^i, \cdot) = Q^j(y^i, \cdot)\}$$

and let  $C_\varepsilon = \bigcap_i C_\varepsilon^i$ .

Define similarly  $C_\varepsilon^{(n)}$  for the  $n$ -stage game  $G_n$  viewed as a one-shot game in normal form with the natural extension of the signalling function  $Q$ .

$$\text{Prove that: } E_\infty = \bigcap_\varepsilon \left( \overline{\bigcup_n G_n(C_\varepsilon^{(n)})} \right).$$

HINT. Use the ideas of the proof of prop. 4.9 p. 169.

**13. Internal correlation.** (Lehrer, 1991)

a. Assume that two players can communicate through the following public signalling matrix  $Q = \begin{pmatrix} a & b \\ b & b \end{pmatrix}$ , where  $a$  and  $b$  are two arbitrary signals.

Consider now any  $S \times T$  correlation matrix  $M$  corresponding to a canonical correlation device and with rational entries, say  $r_{st}/r$ . We will describe a procedure and strategies  $(\sigma, \tau)$  of both players that will generate  $M$  and such that unilateral deviations  $\sigma'$  satisfying  $\sigma' \sim \sigma$  (with respect to the distribution on signals induced by  $Q$ ) will still mimic  $M$  in terms of probabilities and information.

- (1) Let  $R$  be the l.c.m. of the  $r_{st}$ . We define a  $(S \times R) \times (T \times R)$  matrix  $\Phi$  of zeros and ones as follows: the “bloc”  $(s, t)$  is a  $R \times R$  matrix of the form:

$$\begin{pmatrix} "1" & "0" & \dots & "0" \\ "0" & "1" & \dots & "0" \\ \vdots & \vdots & \ddots & \vdots \\ "0" & "0" & \dots & "1" \end{pmatrix}$$

where “1” (resp. “0”) stands for a  $r_{st} \times r_{st}$  matrix of ones (resp. zeros).

Denote by  $(\alpha, \beta)$  an entry of  $\Phi$  and note that if  $\alpha \in s$  (i.e.  $\alpha$  is a line in the bloc  $s$ ), then:  $\sum_{\beta \in t} \Phi(\alpha, \beta) = r_{st}$ . Assume now that  $\alpha$  and  $\beta$  are chosen at random uniformly. Then one has  $P(\alpha \in s, \beta \in t \mid \Phi(\alpha, \beta) = 1) = r_{st}/r$  and  $P(\beta \in t \mid \alpha, \Phi(\alpha, \beta) = 1) = \frac{r_{st}}{\sum_k r_{sk}}$ .

The matrix  $Q$  is now used to check whether  $(\alpha, \beta) \in Z = \{(\alpha, \beta) \mid \Phi(\alpha, \beta) = 0\}$ : Top (resp. Left) meaning “yes, I uses  $\alpha$ ” (resp. II uses  $\beta$ ), Bottom (resp. Right) meaning “no”, and recall that an answer “no” is not informative.

- (2) The strategy  $\sigma$  ( $\tau$  is similar) is now formally defined as follows:

**Step 1:** Choose  $\alpha$  uniformly among the  $S \times R$  lines of  $\Phi$ .

**Step 2:** Given an enumeration of the elements of  $Z$ , answer by yes each time an element  $(\alpha', \beta')$  is checked, with  $\alpha = \alpha'$ .

**Step 3:** Finally report: “I chose  $\alpha$ ”, by using a code as usual (cf. the proof of theorem 4.8 p. 167).

A pure strategy for I, say  $\omega$ , is thus defined by a couple  $(f_\omega, \theta_\omega)$ , where  $f_\omega$  is a mapping from  $Z$  to {yes, no} and  $\theta_\omega$  is a line of  $\Phi$ . An alternative strategy  $\sigma'$  is thus a probability, say  $P'$  on such  $\omega$ 's.

- (3) The procedure works as follows: there is first a checking phase corresponding to answers to an enumeration of  $Z$ . As soon as the entry  $a$  of  $Q$  appears (corresponding under  $\sigma$  and  $\tau$  to a double “yes”), the procedure starts again. If this occurs at a stage where  $(\alpha, \beta)$  is checked, we call this event: a failure at  $(\alpha, \beta)$ . If the previous phase generates a sequence of  $b$ 's, called it a success, one proceeds to the report phase.

Clearly, under  $\sigma$  and  $\tau$ , the procedure will produce an outcome  $(\alpha, \beta)$  after a random time with finite expectation. Prove that if  $\sigma' \sim \sigma$ , it will induce the same distribution on outcomes, given  $\tau$ .

HINT. Show that  $\forall \beta: P'\{\omega \mid f_\omega(\cdot, \beta)$  leads to failure at  $\alpha\} = \frac{1}{S \times R}$ ,  $\forall \alpha: (\alpha, \beta) \in Z$ ; and that  $P'\{\omega \mid \theta_\omega = \alpha, f_\omega(\cdot, \beta)$  yields success  $\} = \frac{\mathbb{1}_{(\alpha, \beta) \notin Z}}{\#\{\alpha \mid \Phi(\alpha, \beta) = 1\}}$ ; which in turn implies  $P'\{\omega \mid \theta_\omega = \alpha\} = \frac{1}{S \times R}$ .

Finally, given  $\theta_\omega = \alpha$ , prove that:  $\{\beta \mid \Phi(\alpha, \beta) = 1\} = \{\beta \mid f_\omega(\cdot, \beta)$  induces a success  $\}$  and conclude.

b. Consider a two-person game with the following signalling structure: given any pair of moves  $(s, t)$ , the signal is public and is either  $(s, t)$  or a constant, say  $\zeta$ .

Say that  $s$  is non-revealing if the corresponding line of signals contains only  $\zeta$ , and revealing otherwise. In case of only non-revealing moves,  $E_\infty$  is obviously the convex hull of  $E_1$ . We now assume the existence of revealing moves.

Then  $E_\infty = C_\infty = D_\infty = G(\bigcap_i \mathbf{A}^i) \cap IR$ .

HINT. Prove the result directly if one player has only revealing strategies.

Otherwise note that there exists a submatrix of signals like the  $Q$  described in part IV.4Ex.13a. By ex. IV.4Ex.12 it is then enough to show that for any pay-off  $d$  in  $G(\bigcap_i \mathbf{A}^i)$  and any  $\varepsilon > 0$  there exists  $n$  and  $(\tilde{\sigma}, \tilde{\tau})$  in  $G_n$  inducing  $d'$  in  $C_\varepsilon^{(n)}$ ,  $\varepsilon$ -close to  $d$ .

Given  $M$  in  $\bigcap_i \mathbf{A}^i$ , use the strategies  $(\sigma, \tau)$  defined in part IV.4Ex.13a, adding between step 2 and step 3 a large (compared to the expected length of the checking and report phases) number of stages, say  $L$ , during which  $(s, t)$  is played if  $(\alpha \in s, \beta \in t)$ .

[Note that if  $\tilde{\sigma}' \sim \tilde{\sigma}$  and  $\tilde{\sigma}'$  produces some  $\alpha \in s$ , there is no gain in playing  $s' \neq s$  during the above  $L$  stages — in fact either player I's signal is  $\zeta$  or his move is revealed.]

## 5. Recursive games

Recursive games were first defined and studied by Everett (1957). They are two-person zero-sum stochastic games where the pay-off is either 0 or absorbing. More precisely we are given a finite set  $J \cup K$  of states, sets of actions  $S, T$ , a transition probability  $Q$  from  $K \times S \times T$  to  $J \cup K$  and some real function  $G$  on  $J$ . The game  $\Gamma$  is now played as follows: given some state  $k_n$  in  $K$  at stage  $n$ , both players choose simultaneously their moves  $(s_n, t_n)$  and a new state  $k_{n+1}$  is selected according to  $Q$ , the current stage pay-off being 0. If  $k_n$  is in  $J$  the pay-off is  $g_n = G(k_n)$  for each following stage. It follows that given any play  $h_\infty$  we can associate to the stream of pay-offs  $g$  its Cesàro limit  $\bar{g}$ , i.e.  $\lim_{n \rightarrow \infty} \bar{g}_n$  exists.  $\Gamma$  is thus a “well defined” game in normal form with vector pay-off function  $\gamma(\sigma, \tau) = E_{\sigma, \tau}(\bar{g})$ ,  $(\gamma^k(\sigma, \tau)$  being the pay-off in  $\Gamma^k$ , i.e.  $\Gamma$  where the initial state is  $k$ ). Obviously we are interested only in  $\Gamma^k$  for  $k$  in  $K$  and we will just consider those.

Note now that if we define the stopping time  $\theta$  on  $H_\infty$  by:  $\theta(h) = \min(\{n \mid k_n \in J\} \cup \{\infty\})$  the pay-off is given by  $\bar{g}(h) = \mathbb{1}_{\theta(h) < \infty} G(k_{\theta(h)})$ .

The analysis is somehow easier if we use another representation. First we will shift the pay-offs one stage backwards, so that at stage  $n$ , given  $(k_n, s_n, t_n)$  in  $K \times S \times T$  the pay-off is  $\sum_{j \in J} Q(j; k_n, s_n, t_n)G(j) \equiv E(f_n)$  where  $f_n = G(k_{n+1})\mathbb{1}_J(k_{n+1})$ .

We can now let  $f_n$  be 0 if  $k_n$  is in  $J$  and define the pay-off up to stage  $n$   $\tilde{f}_n$  as the sum of the previous pay-offs,  $\tilde{f}_n = \sum_{m=1}^n f_m$ . It is then clear that  $\tilde{f}_n$  converges to some  $\tilde{f}$  and that starting from  $k$  in  $K$ ,  $\bar{g}$  and  $\tilde{f}$  coincide.

We will also write  $P$  for the restriction of  $Q$  to  $K$ , hence for all  $s, t$ ,  $P(s, t)$  is a positive kernel on  $K$  with mass less or equal to 1.

Given  $\alpha$  in  $\mathbb{R}^K$ , let  $G^k(\alpha)$  be the one-stage game obtained through  $\Gamma$  starting from  $k$  and with an absorbing pay-off of  $\alpha^\ell$  if  $\ell$  in  $K$  is the state at stage 2. If this game has a value, we denote it by  $U^k(\alpha)$ .

Note that the recursive formula 3.2 p. 158 says that if  $\Gamma$  has a value  $w$  it verifies  $w = U(w)$ .

A strategy is Markov (resp. stationary) if it depends only on the current state and stage (resp. state). We can now state the main result due to Everett (1957) (cf. also Orkin, 1972c):

**THEOREM 5.1.** *If  $U$  exists on  $K \times \mathbb{R}^K$  the recursive game have a value. Moreover both players have  $\varepsilon$ -optimal Markov strategies. (In particular if  $S$  and  $T$  are finite and then the above strategies are stationary).*

**PROOF.** We first prove two lemmas, the first being straightforward.

- LEMMA 5.2.**
- (1)  $\|U(\alpha) - U(\beta)\| \leq \|\alpha - \beta\|$
  - (2) if  $\alpha \geq \beta$  in  $\mathbb{R}^K$  then  $U(\alpha) \geq U(\beta)$  in  $\mathbb{R}^K$

We define now the following sets:

$$\begin{aligned} C_1 &= \{ \alpha \in \mathbb{R}^K \mid U^k(\alpha) \geq \alpha^k \text{ and } U^k(\alpha) > \alpha^k \text{ if } \alpha^k > 0 \} \\ C_2 &= \{ \alpha \in \mathbb{R}^K \mid U^k(\alpha) \leq \alpha^k \text{ and } U^k(\alpha) < \alpha^k \text{ if } \alpha^k < 0 \} \end{aligned}$$

The vectors in  $C_1$  minorate the maxmin of  $\Gamma$  since we have:

**LEMMA 5.3.** *Player I can guarantee any  $\alpha$  in  $C_1$  with Markov strategies.*

**PROOF.** Given  $\alpha$  in  $C_1$ , define  $K(\alpha) = \{ k \in K \mid \alpha^k > 0 \}$ . Let  $\delta = \min\{ U^k(\alpha) - \alpha^k \mid k \in K(\alpha) \}$ , hence  $\delta > 0$ , and finally let  $e = \mathbb{1}_{K(\alpha)}$  in  $\mathbb{R}^K$ . Given  $\varepsilon > 0$ , denote by  $x_n^k$  an  $(\varepsilon/2^n)$  optimal strategy of player I in  $G^k(\alpha)$  and define  $\sigma$  as: play according to  $x_n^k$  in state  $k$ , at stage  $n$ . For any strategy  $\tau$  of player II, one has using the definition of  $\sigma$  and the choice of  $\alpha$ :

$$\mathbb{E}_{\sigma,\tau}(\sum_{m=1}^n f_m + \alpha^{k_{n+1}} \mid \mathcal{H}_n) \geq \sum_{m=1}^{n-1} f_m + \alpha^{k_n} + \delta e^{k_n} - \varepsilon/2^n$$

Hence by recursion:

$$\mathbb{E}_{\sigma,\tau}(\sum_{m=1}^n f_m + \alpha^{k_{n+1}}) \geq \alpha^{k_1} + \delta \mathbb{E}(\sum_{m=1}^n e^{k_m}) - \varepsilon$$

Now, since  $\alpha \leq \delta M e$  for  $M$  large enough we obtain first:

$$\varphi_n(\sigma, \tau) = \mathbb{E}_{\sigma,\tau}(\tilde{f}_n) \geq \alpha^{k_1} + \delta \mathbb{E}(\sum_{m=1}^n e^{k_m} - e^{k_{n+1}}) - \varepsilon$$

Thus  $\mathbb{E}(\sum e^{k_m})$  converges, so  $\mathbb{E}(e^{k_{n+1}})$  goes to 0 hence:

$$(1) \quad \gamma(\sigma, \tau) = \lim \varphi_n(\sigma, \tau) \geq \alpha - \varepsilon. \quad \blacksquare$$

So it will thus suffice to show by induction on  $\#K$ , the number of active states, that  $\bar{C}_1 \cap \bar{C}_2 \neq \emptyset$ , where  $\bar{C}_i$  denotes the closure of  $C_i$ .

Assume first  $\#K = 1$  and consider  $G^1(\alpha)$ . By lemma 5.2.1 p. 176  $U^1$  is a non-expansive mapping from  $[-C, C]$  to itself, where as usual  $C$  is an uniform bound on the pay-offs. It follows that  $U^1$  has a non-empty closed interval of fixed points and we will write  $\alpha^*$  for one of its elements with smallest norm.  $\alpha^*$  belongs to  $\bar{C}_1$ : if  $\alpha^* \leq 0$ , because it is fixed point of  $U^1$ ; while if  $\alpha^* > 0$ , let  $\alpha < \alpha^*$  then  $U^1(\alpha) > \alpha$ , implying  $\alpha$  is in  $C_1$  hence  $\alpha^* \in \bar{C}_1$ . Dually  $\alpha^* \in \bar{C}_2$ .

Assume now that  $\bar{C}_1 \cap \bar{C}_2 \neq \emptyset$  for all games with strictly less than  $\#K$  active states. Obviously this set is then reduced to the value vector.

For each real  $\alpha$ , define the game  $\Gamma_1(\alpha)$  as a recursive game with  $\#K - 1$  active states deduced from  $\Gamma$  by adding to  $J$  state 1 with an absorbing pay-off  $\alpha$ . By induction it has a value for all initial  $k$ . Write  $\Phi_1(\alpha)$  for the vector in  $\mathbb{R}^K$  of its components in  $K$ . Consider

now  $G^1(\Phi_1(\alpha))$  and write  $u(\alpha) = U^1(\Phi_1(\alpha))$  for its value. Obviously  $u$  is a non-expansive mapping from  $[-C, C]$  to itself and we choose again  $\alpha^*$  to be a fixed point with minimum norm.

We claim that  $\Phi_1(\alpha^*)$  belongs to  $\bar{C}_1$ . Consider the 2 cases:

If  $\alpha^* > 0$ , we can choose as above for every positive and small enough  $\varepsilon$ ,  $\alpha = \alpha^* - \varepsilon > 0$  such that  $u(\alpha) > \alpha$ . By induction there exists, for all  $\delta > 0$  a  $K - 1$  dimensional vector  $\beta$  such that  $|\beta^k - \Phi_1^k(\alpha)| \leq \delta$  and  $\beta$  belongs to the set  $C_1$  of the reduced game  $\{\Gamma_1^k(\alpha) \mid k \in K, k \neq 1\}$ . Using lemma 5.2 we obtain by continuity that  $U^1(\alpha, \beta) > \alpha$  for  $\delta$  small enough. Note that for  $k \neq 1$ ,  $U^k(\alpha, \beta)$  is also the value of  $G'(\beta)$ , where  $G'$  is the one-shot game related to  $\{\Gamma_1^k(\alpha) \mid k \in K, k \neq 1\}$ . Thus  $(\alpha, \beta)$  belongs to the set  $C_1$  for the original game, hence the claim.

If  $\alpha^* \leq 0$ , let  $L$  be the set of states  $k$  in  $K$  for which  $\Phi_1^k(\alpha^*) = \alpha^*$  and denote by  $M$  its complement in  $K$ . Consider the recursive game  $\Gamma_L(\alpha)$  with active states set  $M$ , where the states in  $L$  are now absorbing with same pay-off  $\alpha$  and write  $\Phi_L(\alpha)$  for its vector of values on  $K$ . Note that  $U(\Phi_1(\alpha^*)) = \Phi_1(\alpha^*) = \Phi_L(\alpha^*)$ , for all  $k$  in  $K$  (recursive formula 3.2).

Then we have  $\Phi_L^k(\alpha^*) - \Phi_L^k(\alpha) < \alpha^* - \alpha$  for all  $\alpha < \alpha^*$  and all  $k$  in  $M$ . Indeed, given  $\sigma^*$  (resp.  $\tau$ )  $\varepsilon$ -optimal for player I (resp. player II) in  $\Phi^k(\alpha^*)$  (resp.  $\Phi^k(\alpha)$ ), let  $\pi = \text{Pr}_{\sigma^*, \tau}(\exists n; k_n \in L)$ . The pay-off corresponding to  $(\sigma^*, \tau)$  in  $\Phi_L^k(\alpha^*)$  (resp.  $\Phi_L^k(\alpha)$ ) can be written as  $\pi\alpha^* + (1 - \pi)c$  (resp.  $\pi\alpha + (1 - \pi)c$ ) hence we obtain:  $\pi(\alpha^* - \alpha) \geq \Phi_L^k(\alpha^*) - \Phi_L^k(\alpha) - 2\varepsilon$ . Letting  $\varepsilon \rightarrow 0$ , this implies first the weak inequality above and that equality would yield  $\pi \rightarrow 1$ , so  $\Phi_L^k(\alpha^*) = \alpha^*$ , contradicting the definition of  $M$ .

It follows that for all  $\varepsilon > 0$ , there exists  $\delta > 0$  such that  $\Phi_L^k(\alpha) - \delta \geq \Phi_L^k(\alpha^*) - \varepsilon$  for  $\alpha = \alpha^* - \varepsilon$  and all  $k$  in  $M$ . Let now  $\xi$  in  $\mathbb{R}^M$  be such that:

$$(2) \quad |\xi^k - \Phi_L^k(\alpha)| < \delta \quad k \in M$$

$$(3) \quad \xi \text{ belongs to the set } C_1 \text{ of the reduced game } \{\Gamma_L^k(\alpha) \mid k \in M\}$$

One has  $U^k(\Phi_L(\alpha^*) - \varepsilon) \geq \alpha^* - \varepsilon$  for all  $k$  in  $L$ . By monotonicity this implies  $U^k((\Phi_L(\alpha) - \delta)\mathbf{1}_M, \alpha\mathbf{1}_L) \geq \alpha$  for all  $k$  in  $L$ , hence using (2) we obtain:

$$(4) \quad U^k(\xi, \alpha\mathbf{1}_L) \geq \alpha \quad \text{for all } k \text{ in } L$$

Since  $\alpha \leq 0$ , (3) and (4) just imply that  $(\xi, \alpha\mathbf{1}_L)$  is in the set  $C_1$  for the original game. Using the continuity (lemma 5.2 again) we finally get that  $\Phi_1(\alpha^*)$  belongs to the closure of  $C_1$  and this proves the theorem. ■

### Exercises.

#### 1. Applications of recursive games.

a. (Orkin, 1972c) Consider the games defined in 2.b p. 83. Let  $T_1, \dots, T_n$  be disjoint sets of positions of length  $k$  with corresponding pay-offs  $c_1, \dots, c_n$ . Let  $T = \bigcup_{1 \leq i \leq n} T_i$ ,  $\theta$  be the entrance time in  $T$  after any position:  $\theta = \min(\{n \mid (\omega_{n-k+1}, \dots, \omega_n) \in T\} \cup \{\infty\})$ , and define the pay-off as  $g(h_\infty) = \mathbb{1}_{\theta < \infty} \cdot c_i \mathbb{1}_{h_\infty \in T_i}$ .

Prove that the game has a value.

b. Consider a finite (i.e.  $S, T, K$ , finite) two-person zero-sum discounted stochastic game  $\Gamma_\lambda$  (at each stage both players know the previous history). Using theorem 5.1 p. 176 prove that  $v_\lambda$  exists.

**2. Ruin games.** (Milnor and Shapley, 1957) Let  $G$  be an  $S \times T$  real matrix, and  $0 \leq r \leq R$ . The ruin game associated,  $\Gamma(r)$  is a repeated game where both players choose moves  $(s_n, t_n)$  at stage  $n$  inducing a new fortune  $r_n = r_{n-1} + G_{s_n t_n}$  with  $r_0 = r$ . The pay-off is 1 (resp. 0) on every play where  $[R, +\infty]$  (resp.  $(-\infty, 0]$ ) is hit first and some function  $Q$  on  $H_\infty$  otherwise, with  $0 \leq Q \leq 1$ .

a. *Preliminary results.* Given a real function  $f: \mathbb{R} \rightarrow \mathbb{R}$  we introduce  $W_f: \mathbb{R} \rightarrow \mathbb{R}$  where  $W_f(r)$  is the value of the  $S \times T$ matrix  $B_f(r)$  with coefficients  $f(r + G_{st})$ .

i. Let:

$$(Tf)(r) = \begin{cases} 1 & r \geq R \\ W(r) & \text{and } R > r > 0 \\ 0 & 0 \geq r. \end{cases}$$

Prove that if player I can guarantee  $f(r)$  in  $\Gamma(r)$ , for all  $r \in \mathbb{R}$ , he can also guarantee  $(Tf)(r)$ .

ii. Let  $f_0(r) = \mathbb{1}_{[R, +\infty]}(r)$  and  $f_n = Tf_{n-1}$ . Prove that  $f_n \uparrow \underline{f}$ , that  $\underline{f}$  is increasing and that player I can guarantee it. Define similarly  $f^0(r) = \mathbb{1}_{]0, +\infty[}$  and  $\bar{f} = \lim \downarrow T^n f^0$ .

iii. Prove that if  $\Gamma(r)$  has a value  $v(r)$ , then  $v$  satisfies:

$$(\star) \quad v = Tv \quad \text{and} \quad \underline{f} \leq v \leq \bar{f}$$

Deduce from IV.5Ex.2aii that if  $(\star)$  has a unique solution,  $w$ , then  $w$  is the value of  $\Gamma$  and it is independent of  $Q$ .

b. *Special case: coefficients in  $\mathbb{Z}$ .*

i. Let  $G = \begin{pmatrix} rr^T & -b \\ -c & d \end{pmatrix}$  with  $b, c, d \in \mathbb{N}^*$ . Prove that  $\underline{f}$  is strictly monotone on the integers in  $(0, R)$ .

HINT. Let  $w$  be another solution of  $(\star)$  and  $k = \max\{n \mid |w(n) - \underline{f}(n)| = \delta = \max_m |w(m) - \underline{f}(m)|\}$ . Assume  $w(k) > \underline{f}(k)$ . Let  $x$  (resp.  $y$ ) be an optimal strategy for player I (resp. II) in  $B_w(k)$  (resp.  $\bar{B}_f(k)$ ). Compute the pay-off associated to  $(x, y)$  to prove  $\delta = 0$  and conclude.

ii. Assume  $Q$  constant and prove directly that  $\Gamma$  has a value by using theorem 5.1.

c. *Further results.*

i. Given  $w$  bounded solution of  $(\star)$  define a  $w$ -local strategy of player I as follows: given the fortune  $r_n \in (0, 1)$  at stage  $n$ , play at stage  $n+1$  an optimal strategy in  $B_w(r_n)$ . Let  $\sigma$  be a  $w$ -local strategy. Prove that, for every  $\tau$ ,  $W_w(r_n)$  is a bounded submartingale and  $P_{\sigma, \tau}(r_n > 0, \forall n) \geq w(r_0)$ .

Deduce that: if  $Q \equiv 1$  (resp.  $Q \equiv 0$ ),  $\Gamma$  has a value  $\bar{f}$ , (resp.  $\underline{f}$ ); moreover if  $w$  is strictly monotone and  $G_{st} \neq 0, \forall (s, t)$  the game has a value, independent of  $Q$ .

ii. Properties of  $\bar{f}$ . Prove that the following conditions are equivalent:

- (1)  $\bar{f}$  is continuous at  $R$
- (2)  $\bar{f}(r) = 1$  on  $(0, +\infty)$
- (3)  $G$  has a non-negative row.

iii. Properties of  $\underline{f}$ . Prove that the following conditions are equivalent:

- (1)  $\underline{f}$  is continuous at  $R$
- (2)  $\underline{f}(r) = 1$  on  $(0, +\infty)$
- (3) Every subset of columns of  $G$  has a non-zero, non-negative row.

HINT. To prove  $3 \Rightarrow 2$ , let  $q$  be the smallest non-zero  $|G_{st}|$  and  $r$  such that  $\underline{f}(r) < \underline{f}(r+q)$ . Consider then an optimal strategy of player I in  $B_f(r)$ .

d. A special case: if  $G$  is zero-free, then  $\Gamma$  has a value, independent of  $Q$ .

HINT. We can assume that the value of  $G$  is positive and by IV.5Ex.2ciii that  $\underline{f}$  is discontinuous at  $R$ . We are going to construct a strictly monotone solution of some approximation  $T_\varepsilon$  of  $T$ . Let  $C = \max |G_{st}|$  and define:

$$w_0(r) = \begin{cases} \varepsilon(r - R - C) & r < R \\ 1 + \varepsilon(r - R - C) & r \geq R \end{cases} \quad \text{and} \quad w_n(r) = \begin{cases} \varepsilon(r - R - C) & r \leq 0 \\ Bw_{n-1}(r) & r \in (0, R) \\ 1 + \varepsilon(r - R - C) & r \geq R \end{cases}$$

- (1) Prove that  $w_n(r) \uparrow$  to some  $w(r)$ .
- (2) Prove that for  $\varepsilon$  small enough  $w_n(r) - \varepsilon r$  is monotone for all  $n$ . (For the case  $0 < r < R \leq s$ , prove that  $w_n \leq \underline{f}$  and choose  $\varepsilon$  such that  $\varepsilon C \leq 1 - \underline{f}(R^-)$ ).
- (3) Deduce then the result.

For the existence of optimal strategies, cf. (Milnor and Shapley, 1957).

e. General case. We assume here the following restriction on  $Q$ : whenever  $r_n$  converges, say to  $r \in (0, R)$ , then the value of  $Q$  on the corresponding history is only a function, say  $P^*$ , of  $r$ . We extend now  $P^*$  to  $\mathbb{R}$  by putting  $P^* = 0$  on  $(-\infty, 0]$  and  $= 1$  on  $[R, +\infty)$  and we define  $Q^*$  as the restriction of  $Q$  to non-convergent pay-offs.

If  $P^*$  is increasing then  $\Gamma$  has a value, independent of  $Q^*$ .

HINT. (1) Consider the following auxiliary game  $\Gamma_w^*(r)$ , for  $w$  bounded and  $r \in (0, R)$ . If  $G_{st} \neq 0$  the pay-off is  $w(r + G_{st})$  and the game ends. If  $G_{st} = 0$  the game is repeated. The pay-off corresponding to a non-terminating play is  $P^*(r)$ . Use theorem 5.1 p. 176 to prove that  $\Gamma_w^*(r)$  has a value  $V_w^*(r)$ .

- (2) Define a new operator  $T^*$  by:  $T^*w(r) = \begin{cases} V_w^*(r) & r \in (0, R) \\ P^*(r) & r \notin (0, R) \end{cases}$  and introduce  $\underline{f}^*$  as in IV.5Ex.2aii.

Prove that there exists a monotone solution to:

$$(\star\star) \quad T^*w = w$$

- (3) Assume that  $(\star\star)$  has a strictly monotone solution  $w$ . Prove that  $w$  is the value of  $\Gamma$ . [Let  $\varepsilon_m = \varepsilon/2^{m+1}$ ,  $n_m = \min\{l \geq n_{m-1} \mid r_l \neq r_{n_{m-1}}\}$ ; player I plays an  $\varepsilon_m$ -optimal strategy in  $\Gamma_w^*(r_{n_m})$  between stages  $n_m$  and  $n_{m+1}-1$ .]
- (4) Define now, as in IV.5Ex.2d),  $P_\varepsilon^*$  as being equal to  $P^*$  on  $(0, R)$  and

$$P_\varepsilon^*(r) = \begin{cases} \varepsilon(r - R - C) & r \leq 0 \\ 1 + \varepsilon(r - R - C) & r \geq R \end{cases}$$

Prove then that the corresponding equation  $(\star\star)$  possesses, if the value of  $G$  is positive and  $\underline{f}_*$  is discontinuous at  $R$ , a strictly monotone solution.

- (5) Prove finally that if  $\underline{f}$  and  $P^*$  have jumps at  $R$  so does  $\underline{f}^*$  and conclude by approximating  $P^*$ .

### 3. A game with no value. (Zamir, 1971–1972)

a. Consider the game with incomplete information, where  $k \in K = \{1, 2\}$  is chosen according to  $p = (1/2, 1/2)$  and remains fixed, no player being informed of it:

$$G^1 = \begin{pmatrix} 0 & 8 \\ 0 & 8 \end{pmatrix}, \quad A^1 = \begin{pmatrix} a & a \\ b & c \end{pmatrix}, \quad G^2 = \begin{pmatrix} 8 & 0 \\ 8 & 0 \end{pmatrix}, \quad A^2 = \begin{pmatrix} a & a \\ d & c \end{pmatrix}.$$

After each move  $(s, t)$ ,  $A_{st}^k$  is announced to both players. Prove that  $\lim v_n = v_\infty = 4$ .

b. Consider now  $\Gamma$ , played as in IV.5Ex.3a, with  $K = \{1, 2, 3\}$ ,  $p = (1/4, 1/4, 1/2)$ ,  $G^k, A^k$ ,  $k = 1, 2$  as above and

$$G^3 = \begin{pmatrix} 0 & -4 \\ 0 & 0 \end{pmatrix}, \quad A^3 = \begin{pmatrix} a & a \\ e & f \end{pmatrix}.$$

- i. Prove that  $\bar{v} = \lim v_n = 1$ .

HINT. Let II play  $(1/2, 1/2)$  as long as  $a$  is announced and optimally thereafter. For player I take  $s_n = 1$  if  $P_\tau(t_n = 1) \geq 1/2$ , 2 otherwise.

- ii. Prove that  $\underline{v} = 0$ .

HINT. Let  $q_n = P_\sigma(s_n = 1 \mid a_m = a, \forall m = 1, \dots, n-1)$  and  $q = \prod_1^\infty q_n$ . If  $q = 0$ , player II plays  $t_n = 1$  as long as  $a$  appears. If  $q > 0$ , player II plays  $t_n = 1$  up to some large stage  $N$  and  $t_n$  thereafter.

**4. Duels.** We consider here a noisy duel (cf. ex. I.2Ex.9 p. 22) between time 0 and 1, with symmetric accuracy function  $Q(t) = t$ , and where player I (resp. II) has  $m$  bullets (resp.  $n$ ). Let  $v(m, n)$  be the corresponding value.

- a. Prove that  $v(1, 1) = 1/2$  and that both players have optimal strategies.

b. Prove by induction (recursive formula) that:  $v(m, n) = (m - n)/(m + n)$ , that player I has an optimal strategy as long as  $m > n$  (shoot at time  $1/m + n$  if player II did not shoot before) and only  $\varepsilon$ -optimal strategy otherwise (shoot at random on a small interval around  $1/m + n$ , if player II did not shoot before).

## Part B

### The Central Results



## CHAPTER V

### Full Information on one Side

We start now to study repeated games with incomplete information. In the present chapter we consider the simplest class of those games namely two-person zero-sum games in which one player, say player I, is fully informed about the state of nature while the other player, player II, knows only the prior distribution according to which the state is chosen.

#### 1. General properties

In this section we prove some general properties of a one-shot game with incomplete information which will be later on applied to various versions of the game: finitely or infinitely repeated games or discounted games. The game considered here is a two-person zero-sum game of the following form: chance chooses a state  $k$  from a finite set  $K$  of states (games) according to some probability  $p \in \Pi = \Delta(K)$ . Player I (the maximiser) is informed which  $k$  was chosen but player II is not. Players I and II then choose simultaneously  $\sigma^k \in \Sigma$  and  $\tau \in \mathcal{T}$  respectively and finally  $G^k(\sigma^k, \tau)$  is paid to player I by player II. The sets  $\Sigma$  and  $\mathcal{T}$  are some convex sets of strategies and the pay-off functions  $G^k(\sigma^k, \tau)$  are bi-linear and uniformly bounded on  $\Sigma \times \mathcal{T}$ .

In normal form this is a game in which the strategies are  $\sigma \in \Sigma^K$  and  $\tau \in \mathcal{T}$  respectively and the pay-off function is  $G^p(\sigma, \tau) = \sum_k p^k G^k(\sigma^k, \tau)$ . Denote this game by  $\Gamma(p)$ .

**THEOREM 1.1.**  $\overline{w}(p) = \inf_{\tau} \sup_{\sigma} G^p(\sigma, \tau)$  and  $\underline{w}(p) = \sup_{\sigma} \inf_{\tau} G^p(\sigma, \tau)$  are concave.

**PROOF.** The proof is the same for both functions. We write it for  $\overline{w}(p)$ . Let  $(p_e)_{e \in E}$  be finitely many points in  $\Delta(K)$ , and let  $\alpha = (\alpha_e)_{e \in E}$  be a point in  $\Delta(E)$  such that  $\sum_{e \in E} \alpha_e p_e = p$ , we claim that  $\overline{w}(p) \geq \sum_{e \in E} \alpha_e \overline{w}(p_e)$ . To see that consider the following two-stage game: A chance move chooses  $e \in E$  according to the probability distribution  $(\alpha_e)_{e \in E}$ , then  $k \in K$  is chosen according to  $p_e$ , the players choose  $\sigma^k \in \Sigma$  and  $\tau \in \mathcal{T}$  respectively and the pay-off is  $G^k(\sigma^k, \tau)$ . We consider two versions in both of which player I is informed of everything (both  $e$  and  $k$ ) while player II may or may not be informed of the value of  $e$  (but in any case he is not informed of the value of  $k$ ).

Now if player II is informed of the outcome  $e$ , the situation following the first lottery is equivalent to  $\Gamma(p_e)$ . Thus, the  $\inf_{\tau} \sup_{\sigma}$  for the game in which player II is informed of the outcome of the first stage is  $\sum_{e \in E} \alpha_e \overline{w}(p_e)$ . This game is more favourable to II than the game in which he is not informed of the value of  $e$ , which is equivalent to  $\Gamma(\sum_e \alpha_e p_e) = \Gamma(p)$ . Therefore we have:

$$(1) \quad \overline{w}(p) \geq \sum_{e \in E} \alpha_e \overline{w}(p_e). \quad \blacksquare$$

**REMARK 1.1.** Since player II's strategy space and the pay-off function are completely general, nothing prevents the game to be a normalised form game in which player II first observes the result of a  $k$ -dependent lottery, i.e. his own “type”, and then chooses

his action. We obtain thus the concavity of  $\bar{w}(p)$  and  $\underline{w}(p)$  in games with incomplete information on both sides, when  $p$  is restricted to the subset of the simplex where player I's conditional probability on the state  $k$ , given his own type, is fixed.

**REMARK 1.2.** The concavity of  $\underline{w}(p)$  can also be proved constructively. For this let us first prove the following proposition which we shall refer to as the **splitting procedure**:

**PROPOSITION 1.2.** Let  $(p_e)_{e \in E}$  be finitely many points in  $\Delta(K)$ , and let  $\alpha = (\alpha_e)_{e \in E}$  be a point in  $\Delta(E)$  with  $\sum_{e \in E} \alpha_e p_e = p$ , then there are vectors  $(\mu^k)_{k \in K}$  in  $\Delta(E)$  such that the probability distribution  $P$  on  $K \times E$  obtained by the composition of  $p$  and  $(\mu^k)_{k \in K}$  (that is:  $k \in K$  is chosen according to  $p$  and then  $e \in E$  is chosen according to  $\mu^k$ ,) satisfies: For all  $e \in E$ ,

$$P(\cdot | e) = p_e \quad \text{and} \quad P(e) = \alpha_e$$

**PROOF.** If  $p^k = 0$ ,  $\mu^k$  can be chosen arbitrarily in  $\Delta(E)$ . If  $p^k > 0$ ,  $\mu^k$  is given by  $\mu^k(e) = \alpha_e p_e^k / p^k$ . Using Bayes' formula the required properties are directly verified. ■

**COROLLARY 1.3.** The function  $\underline{w}(p)$  is concave.

**PROOF.** Let  $p = \sum_e \alpha_e p_e$ . Let player I use the above described lottery and then use  $\varepsilon$ -optimal strategy in  $\Gamma(p_e)$ . In this way he obtains at least  $\sum_e \alpha_e \underline{w}(p_e) - \varepsilon$ , even if player II were informed of the outcome of the lottery. So  $\underline{w}(p)$  is certainly larger than that. ■

**PROPOSITION 1.4.** For any  $\tau \in \mathcal{T}$ , the function  $\sup_\sigma G^p(\sigma, \tau)$  is linear in  $p$ .

**PROOF.** Indeed, player I just optimises given  $k$  (and  $\tau$ ), yielding:

$$(2) \quad \sup_\sigma G^p(\sigma, \tau) = \sum_k p^k \sup_{\sigma^k} G^k(\sigma^k, \tau). \quad \blacksquare$$

**COMMENT 1.3.** The last proposition provides another proof of the concavity of  $\bar{w}(p)$ .

**COROLLARY 1.5.** If  $\Sigma = \Delta(S)$  and  $\mathcal{T} = \Delta(T)$  where  $S$  and  $T$  are finite, then the function  $\bar{w}(p)$  is piecewise linear.

**PROOF.** Knowing already that  $\bar{w}(p)$  is concave, we have to show that the set

$$A = \{ \alpha = (\alpha_k)_{k \in K} \in \mathbb{R}^K \mid \exists \tau \text{ such that } G^k(s^k, \tau) \leq \alpha_k ; \forall s^k \in S \ \forall k \}$$

is a convex polyhedron. This is true since it is written as the projection of the polyhedron (in  $(\alpha, \tau) \in \mathbb{R}^K \times \mathbb{R}^T$ ):

$$\{ (\alpha, \tau) \mid \sum_{t \in T} \tau_t G^k(s^k, t) - \alpha_k \leq 0 ; \forall s^k \in S \ \forall k \}$$

Indeed since obviously  $A \subseteq \{ \alpha \in \mathbb{R}^K \mid \alpha_k \geq -C \ \forall k \}$ , for sufficiently large constant  $C$ , any point in  $A$  is minorated by a convex combination of extreme points. Hence we have  $\bar{w}(p) = \min_\alpha \sum_k \alpha_k p_k$  where  $\alpha$  varies over the finitely many extreme points. ■

Another general property worth mentioning is the Lipschitz property of all functions of interest in particular  $\bar{w}(p)$ . This follows just from the uniform boundedness of the pay-offs, and hence it is valid for any repeated game as defined in ch. IV p. 147.

**THEOREM 1.6.** The function  $\bar{w}(p)$  is Lipschitz with constant  $C$  (the bound on the absolute value of pay-offs).

**PROOF.** Indeed the pay-off functions of two games  $\Gamma(p_1)$  and  $\Gamma(p_2)$  differ by at most  $C \|p_1 - p_2\|_1$ . ■

**DEFINITION 1.7.** Given any real valued function  $f$  on  $\Pi = \Delta(K)$ , we denote by  $\text{Cav } f$  the (point-wise) minimal function  $g$ , concave and greater than  $f$  on  $\Pi$ .

## 2. Elementary tools and the full monitoring case

In this section we introduce some elementary tools frequently used in repeated games with incomplete information. We do this by studying first the relatively simple special case of **full monitoring**. This is the case in which the moves (and only the moves) of the players at each stage are observed by both of them and hence they serve as the (only) device for transmitting information about the state of nature. This important special case will serve to show the results and the main ideas in a relatively simpler framework. The more complex general case will be treated in later sections.

The game considered here is a special case of the general model introduced in ch.IV p.147: There are two players, I and II with finite action sets  $S$  and  $T$  respectively. The state space is a finite set  $K$  on which the prior probability distribution is  $p \in \Delta(K)$ . The pay-off in state  $k \in K$  is given by the  $(S \times T)$  matrix  $G^k$  with elements  $G_{st}^k$ . Let  $C = \max_{k,s,t} |G_{st}^k|$ .

The repeated game  $\Gamma(p)$  is played as follows:

- At stage 0 a chance move chooses  $k \in K$  with probability distribution  $p \in \Delta(K)$ . The result is told to player I, the row chooser, but not to player II who knows only the initial probability distribution  $p$ .
- At stage  $m$ ;  $m = 1, 2, \dots$  player I chooses  $s_m \in S$  and II chooses  $t_m \in T$  and  $(s_m, t_m)$  is announced. The pay-off  $g_m$ , for player I, is  $G_{s_mt_m}^k$ .

Denote the  $n$ -stage game by  $\Gamma_n(p)$  and its value by  $v_n(p)$  (cf. sect. 1.b p. 148). The infinite  $\lambda$ -discounted game is denoted by  $\Gamma_\lambda(p)$  and its value by  $v_\lambda(p)$ . We also consider the infinitely repeated game  $\Gamma_\infty(p)$  without specifying the pay-off function and with the usual definitions of minmax, maxmin and value  $v_\infty(p)$  (cf. sect. 1.c p. 149).

To fix ideas think of the following example:

**EXAMPLE 2.1.** Consider a game with two states  $K = \{1, 2\}$  in which the pay-offs and the probability are given by:

$$G^1 = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \quad G^2 = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} \quad p = (1/2, 1/2)$$

The main feature of these games is that the informed player's moves will typically depend (among other things) on his information (i.e. on the value of  $k$ ). Since these moves are observed by the uninformed player, they serve as a channel which can transfer information about the actual state  $k$ . This must be taken into account by player I when choosing his strategy. In our example for instance playing at some stage the move  $s = 1$  if  $k = 1$  and  $s = 2$  if  $k = 2$  is a dominant strategy as far as the pay-offs at that stage are concerned. However such behaviour will reveal the value of  $k$  to player II and by that enable him to reduce the pay-offs to 0 in all subsequent stages. This is of course very disadvantageous in the long run and player I would be better off even by simply ignoring his information: playing the mixed move  $(1/2, 1/2)$  at each stage independently of the value of  $k$  guarantees an expected pay-off of at least  $1/4$  per stage. We shall see that this is in fact the best he can do in the long run in this game.

**2.a. Posterior probabilities and non-revealing strategies.** For  $n = 1, 2, \dots$ , let  $H_n^{\text{II}} = [S \times T]^{n-1}$  be the set of histories for player II, or II-histories at stage  $n$ . (An element  $h_n \in H_n^{\text{II}}$  is a sequence  $(s_1, t_1; s_2, t_2; \dots; s_{n-1}, t_{n-1})$  of moves of both players in the

first  $n - 1$  stages of the game; it is known by both players at stage  $n$ .) Compared with the similar notations in (IV.1.a), a II-history does not contain the state  $k$  and the set of plays is  $H_\infty = K \times (S \times T)^\infty$ , with the usual  $\sigma$ -algebra  $\mathcal{H}_\infty$ . Let  $\mathcal{H}_n^{\text{II}}$  be the  $\sigma$ -algebra on  $H_\infty$  generated by the cylinders above  $H_n^{\text{II}}$  and corresponding to II's information partition (sect. 1.b p. 148).

Any strategies  $\sigma$  and  $\tau$ , of players I and II respectively and  $p \in \Delta(K)$  induce a probability distribution  $P_{\sigma,\tau}^p$  on the measurable space of plays. This will be our basic probability space and we will simply write  $P$  or  $E$  for probability or expectation when no confusion can arise. We denote by  $p_n$  the conditional probability on  $K$  given  $\mathcal{H}_n^{\text{II}}$ , that is

$$p_n^k = P(k \mid \mathcal{H}_n^{\text{II}}), \quad \forall k \in K.$$

This random variable on  $\mathcal{H}_n^{\text{II}}$  has a clear interpretation: This is player II's **posterior probability distribution** on  $K$  at stage  $n$  given the history up to that stage. Let  $p_1 \equiv p$  by definition. These posterior probabilities turn out to be the natural **state variable** of the game and therefore play a central rôle in our analysis.

First observe that the sequence  $(p_n)_{n=1}^\infty$  is a  $(\mathcal{H}_n^{\text{II}})_{n=1}^\infty$  martingale, being a sequence of conditional probabilities with respect to an increasing sequence of  $\sigma$ -fields, i.e.

$$E(p_{n+1} \mid \mathcal{H}_n^{\text{II}}) = p_n, \quad \forall n = 1, 2, \dots$$

In particular this implies  $E(p_n) = p, \forall n$ . Furthermore, since this martingale is uniformly bounded, we have the following bound on its variation in the  $L_1$ -norm  $\|\cdot\|$ :

LEMMA 2.1.

$$\frac{1}{n} \sum_{m=1}^n E \|p_{m+1} - p_m\| \leq \sum_k \sqrt{\frac{p^k(1-p^k)}{n}}$$

PROOF. Note that for a martingale  $(p_m)_{m=1}^n$  and for any  $k \in K$ , and  $m = 1, \dots, n$

$$\begin{aligned} E(p_{m+1}^k - p_m^k)^2 &= E(E[(p_{m+1}^k - p_m^k)^2 \mid \mathcal{H}_m^{\text{II}}]) \\ &= E(E[(p_{m+1}^k)^2 \mid \mathcal{H}_m^{\text{II}}] - 2p_m^k E[p_{m+1}^k \mid \mathcal{H}_m^{\text{II}}] + (p_m^k)^2) \\ &= E(E[(p_{m+1}^k)^2 \mid \mathcal{H}_m^{\text{II}}] - (p_m^k)^2) \\ &= E((p_{m+1}^k)^2 - (p_m^k)^2) \end{aligned}$$

Hence

$$E \sum_{m=1}^n (p_{m+1}^k - p_m^k)^2 = E((p_{m+1}^k)^2 - (p_m^k)^2) \leq p_1^k(1-p_1^k)$$

Use now the fact that  $\|x\|_1 \leq \|x\|_2$  on the probability space in which  $m$  is chosen, independently of  $h_n$  in  $H_n^{\text{II}}$ , uniformly between 1 and  $n$  to obtain:

$$E \frac{1}{n} \sum_{m=1}^n |p_{m+1}^k - p_m^k| \leq \sqrt{\frac{1}{n} E \sum_{m=1}^n (p_{m+1}^k - p_m^k)^2} \leq \sqrt{\frac{p^k(1-p^k)}{n}}$$

The proof is concluded by summing on  $k \in K$ . ■

COMMENT 2.2. Note that  $\sum_k \sqrt{p^k(1-p^k)} \leq \sqrt{\#K - 1}$  since the left hand side is maximised by  $p^k = 1/(\#K)$  for all  $k$ . Intuitively, lemma 2.1 above means that at “most of the stages”  $p_{m+1}$  cannot be very different from  $p_m$ .

By induction one can compute explicitly  $p_m$ : Given a strategy  $\sigma$  of player I, a stage  $n$  and a II-history  $h_n \in H_n^{\text{II}}$  let  $\sigma(h_n) = (x_n^k)_{k \in K}$  denote the vector of mixed moves of player I at that stage. Namely he uses the mixed move  $x_n^k = (x_n^k(s))_{s \in S} \in X = \Delta(S)$  in the game  $G^k$ . Given  $p_n(h_n) = p_n$ , let  $\bar{x}_n = \sum_{k \in K} p_n^k x_n^k$  be the (conditional) average mixed move of

player I at stage  $n$ . The (conditional) probability distribution of  $p_{n+1}$  can now be written by Bayes' formula;  $\forall s \in S$  such that  $\bar{x}_n(s) > 0$  and  $\forall k \in K$ :

$$(1) \quad p_{n+1}^k(s) = P(k | h_n, s_n = s) = \frac{p_n^k x_n^k(s)}{\bar{x}_n(s)}$$

It follows that if  $x_n^k = \bar{x}_n$  whenever  $p_n^k > 0$ , then  $p_{n+1} = p_n$ , hence:

**PROPOSITION 2.2.** *Given any II-history  $h_n$ , the posterior probabilities do not change at stage  $n$  if player I's mixed move at that stage is independent of the state  $k$  over all values of  $k$  for which  $p_n^k > 0$ .*

In such a case we shall say that player I plays **non-revealing** at stage  $n$  and we define the corresponding set:

**DEFINITION 2.3.**

$$\text{NR} = \{x \in X^K \mid x^k = x^{k'} \ \forall k, k' \in K\}$$

Due to the full monitoring assumption, not revealing the information is equivalent to not using the information. But then the outcome of the initial chance move (choosing  $k$ ) is not needed during the game: this lottery can as well be made at the end, just to compute the pay-off.

**DEFINITION 2.4.** For  $p$  in  $\Delta(K)$ , the non-revealing game at  $p$ , denoted by  $D(p)$ , is the (one-shot) two-person zero-sum game with pay-off matrix

$$D(p) = \langle p, G \rangle = \sum_{k \in K} p^k G^k.$$

Let  $u(p)$  denote the value of  $D(p)$ .

**REMARK 2.3.** Clearly  $u$  is a continuous function on  $\Delta(K)$  (and Lipschitz with constant  $C$ .)

Coming back to the martingale generated by player I's strategy, we shall now see that at the stages  $m$  in which  $p_{m+1}$  is close to  $p_m$ , player I is not playing "very revealing" and player II can hold his maximal pay-off close to  $u$ .

Given a strategy  $\sigma$  of player I, let  $\sigma_n = (\sigma_n^k)_{k \in K}$  be "the strategy at stage  $n$ " (cf. sect. 1.b p. 148). Its average (over  $K$ ) is the random variable  $\bar{\sigma}_n = \mathbb{E}(\sigma_n | \mathcal{H}_n^{\text{II}}) = \sum_k p_n^k \sigma_n^k$ . Note that  $\bar{\sigma}_n$ , like  $\sigma_n$ , is a function on  $H_n$  and that it has values in NR.

A very crucial element in the theory is the following property: It is intuitively clear that if the  $\sigma_n^k$  are close (i.e. all near  $\bar{\sigma}_n$ ),  $p_{n+1}$  will be close to  $p_n$ . In fact a much more precise relation is valid; namely these two distances are equal:

**LEMMA 2.5.** *For any strategies  $\sigma$  and  $\tau$  of the two players*

$$\mathbb{E}(\|\sigma_n - \bar{\sigma}_n\| \mid \mathcal{H}_n^{\text{II}}) = \mathbb{E}(\|p_{n+1} - p_n\| \mid \mathcal{H}_n^{\text{II}}).$$

**PROOF.** Let  $\zeta$  denote a generic point in  $H_n^{\text{II}}$ . In accordance with our previous notation we write  $x_n$  for  $\sigma_n(\zeta)$  and hence  $\bar{x}_n$  for  $\bar{\sigma}_n(\zeta)$ . For  $s \in S$ , the  $s$  coordinates of these vectors are denoted by  $x_n(s)$  and  $\bar{x}_n(s)$  respectively. Evaluating the above left hand side expectation at  $\zeta$  we have:

$$\mathbb{E}(\|\sigma_n - \bar{\sigma}_n\| \mid \zeta) = \sum_k p_n^k(\zeta) \|x_n^k - \bar{x}_n\| = \sum_k p_n^k(\zeta) \sum_s |x_n^k(s) - \bar{x}_n(s)|.$$

On the other hand by the definition of  $p_n$

$$\mathbb{E}(\|p_{n+1} - p_n\| \mid \zeta) = \sum_s P(s \mid \zeta) \|p_{n+1}(\zeta, s) - p_n(\zeta)\|,$$

where by (1) p. 187:

$$P(s \mid \zeta) = \bar{x}_n(s) \text{ and } p_{n+1}^k(\zeta, s) = \frac{p_n^k(\zeta)x_n^k(s)}{\bar{x}_n(s)} \text{ if } \bar{x}_n(s) > 0.$$

So we obtain

$$\begin{aligned} \mathbb{E}(\|p_{n+1} - p_n\| \mid \zeta) &= \sum_s \bar{x}_n(s) \sum_k \left| \frac{p_n^k(\zeta)x_n^k(s)}{\bar{x}_n(s)} - p_n^k(\zeta) \right| \\ &= \sum_k p_n^k(\zeta) \sum_s |x_n^k(s) - \bar{x}_n(s)| \\ &= \mathbb{E}(\|\sigma_n - \bar{\sigma}_n\| \mid \zeta). \end{aligned}$$
■

We observe now that the distance between pay-offs is bounded by the distance between the corresponding strategies. In fact given  $\sigma$  and  $\tau$  let  $\rho_n(\sigma, \tau) = \mathbb{E}(g_n \mid \mathcal{H}_n^{\text{II}})$ , and define  $\tilde{\sigma}(n)$  to be the same as the strategy  $\sigma$  except for stage  $n$  where  $\tilde{\sigma}_n(n) = \bar{\sigma}_n$ , then we have:

LEMMA 2.6. *For any  $\sigma$  and  $\tau$ ,*

$$|\rho_n(\sigma, \tau) - \rho_n(\tilde{\sigma}(n), \tau)| \leq C \mathbb{E}(\|\sigma_n - \bar{\sigma}_n\| \mid \mathcal{H}_n^{\text{II}}).$$

PROOF. Note that  $p_n$  is the same under  $\sigma$  and under  $\tilde{\sigma}(n)$ . Again let  $\zeta \in H_n^{\text{II}}$  and write  $\tau_n(\zeta) = y_n$ . We have

$$\begin{aligned} \rho_n(\sigma, \tau)(\zeta) &= \sum_k p_n^k(\zeta)x_n^k G^k y_n \text{ and} \\ \rho_n(\tilde{\sigma}(n), \tau_n)(\zeta) &= \sum_k p_n^k(\zeta)\bar{x}_n G^k y_n. \end{aligned}$$

So we get

$$\begin{aligned} |\rho_n(\sigma, \tau) - \rho_n(\tilde{\sigma}(n), \tau)|(\zeta) &\leq C \sum_k p_n^k(\zeta) \|x_n^k - \bar{x}_n\| \\ &= C \mathbb{E}(\|\sigma_n - \bar{\sigma}_n\| \mid \zeta). \end{aligned}$$
■

**2.b.  $\lim v_n(p)$  and  $v_\infty(p)$ .** We state first a property valid for the general case (sect. 3 p. 191):

PROPOSITION 2.7. *In any version of the repeated game  $(\Gamma_n(p), \Gamma_\lambda(p)$  or  $\Gamma_\infty(p)$ ) if player I can guarantee  $f(p)$  then he can also guarantee  $\text{Cav } f(p)$ .*

PROOF. The proof is similar to that of cor. 1.3 p. 184: Given  $\varepsilon > 0$  and  $p$  choose  $(p_e)_{e \in E}$  in  $\Delta(K)$  with  $\#E \leq \#K + 1$  and  $\alpha \in \Delta(E)$ , such that  $p = \sum \alpha_e p_e$ ,  $\text{Cav } f(p) \leq \sum \alpha_e f(p_e) + \varepsilon$  (Carathéodory). Then player I performs the lottery described in prop. 1.2 p. 184 and guarantees  $f(p_e)$  in  $\Gamma(p_e)$ . This gives the proof for  $\Gamma_n(p)$  and  $\Gamma_\lambda(p)$ . As for  $\Gamma_\infty(p)$  note that if the strategy which  $\varepsilon/\#E$ -guarantees  $f(p)$  in  $\Gamma_\infty(p_e)$  corresponds to  $N_e$  then the above described strategy  $2\varepsilon$ -guarantees  $\text{Cav } f(p)$  with the corresponding  $N = \max N_e$ . ■

PROPOSITION 2.8. *Player I can guarantee  $\text{Cav } u(p)$  in  $\Gamma_\infty(p)$ . Moreover,  $v_n(p)$  and  $v_\lambda(p)$  are both at least  $\text{Cav } u(p)$  (for all  $n$  and all  $\lambda > 0$ ).*

PROOF. If player I uses NR moves at all stages, the posterior probabilities remain constant. Hence the (conditional) pay-off at each stage can be computed from the NR game  $D(p)$ . In particular by playing an optimal strategy in  $D(p)$  player I can obtain an expected pay-off of at least  $u(p)$  at each stage, hence  $v_n(p) \geq u(p)$  and  $v_\lambda(p) \geq u(p)$  and player I can guarantee  $u(p)$  also in  $\Gamma_\infty(p)$ . The result follows now from prop. 2.7 p. 188. ■

PROPOSITION 2.9. For all  $p$  in  $\Delta(K)$  and all  $n$ ,

$$v_n(p) \leq \text{Cav } u(p) + \frac{C}{\sqrt{n}} \sum_k \sqrt{p^k(1-p^k)}$$

PROOF. Making use of the minmax theorem it is enough to prove that for any strategy  $\sigma$  of player I in  $\Gamma_n(p)$ , there exists a strategy  $\tau$  of player II such that:

$$\bar{\gamma}_n(\sigma, \tau) \leq \text{Cav } u(p) + \frac{C}{\sqrt{n}} \sum_k \sqrt{p^k(1-p^k)}$$

Given  $\sigma$  let  $\tau$  be the following strategy of player II: At stage  $m$ , given  $h_m = \zeta$ , compute  $p_m(\zeta)$  and play a mixed action  $\tau_m(\zeta)$  which is optimal in  $D(p_m(\zeta))$ .

By lemma 2.6 p. 188 and lemma 2.5 p. 187, for  $m = 1, \dots, n$

$$\rho_m(\sigma, \tau) \leq \rho_m(\tilde{\sigma}(m), \tau) + C \mathbb{E}(\|p_{m+1} - p_m\| \mid \mathcal{H}_m^{\text{II}}).$$

Now

$$\rho_m(\tilde{\sigma}(m), \tau) = \sum_k p_m^k \bar{\sigma}_m G^k \tau_m,$$

with  $\bar{\sigma}_m \in \text{NR}$  and  $\tau_m$  optimal in  $D(p_m)$ , hence

$$\rho_m(\tilde{\sigma}(m), \tau) \leq u(p_m) \leq \text{Cav } u(p_m),$$

which yields

$$\mathbb{E}(g_m \mid \mathcal{H}_m^{\text{II}}) = \rho_m(\sigma, \tau) \leq \text{Cav } u(p_m) + C \mathbb{E}(\|p_{m+1} - p_m\| \mid \mathcal{H}_m^{\text{II}}).$$

Averaging on  $m = 1, \dots, n$  and over all possible histories  $\zeta \in H_n^{\text{II}}$  we obtain (using  $\mathbb{E} \text{Cav } u(p_m) \leq \text{Cav } u(p)$ , by Jensen's inequality):

$$\bar{\gamma}_n(\sigma, \tau) \leq \text{Cav } u(p) + \frac{C}{n} \sum_{m=1}^n \mathbb{E} \|p_{m+1} - p_m\|.$$

The claimed inequality now follows from lemma 2.1 p. 186. ■

The main result of this section can thus be written as:

THEOREM 2.10. For all  $p \in \Delta(K)$ ,  $\lim_{n \rightarrow \infty} v_n(p)$  exists and equals  $\text{Cav } u(p)$ . Furthermore the speed of convergence is bounded by:

$$0 \leq v_n(p) - \text{Cav } u(p) \leq \frac{C}{\sqrt{n}} \sum_k \sqrt{p^k(1-p^k)}.$$

PROOF. Follows from prop. 2.8 p. 188 and prop. 2.9 p. 189. ■

COROLLARY 2.11.  $\lim_{\lambda \rightarrow 0} v_\lambda(p)$  exists and equals  $\text{Cav } u(p)$ ; and the speed of convergence satisfies:

$$0 \leq v_\lambda(p) - \text{Cav } u(p) \leq C \sqrt{\frac{\lambda}{2-\lambda}} \sum_k \sqrt{p^k(1-p^k)}.$$

PROOF. The bound in this case follows also from the strategy in prop. 2.9 p. 189 since

$$\begin{aligned} \sum_1^\infty \lambda(1-\lambda)^{m-1} \mathbb{E} \|p_{m+1}^k - p_m^k\| &\leq \mathbb{E} \left[ \left( \sum_1^\infty \lambda^2(1-\lambda)^{2(m-1)} \cdot \sum_1^\infty (p_{m+1}^k - p_m^k)^2 \right)^{1/2} \right] \\ &\leq \left( \frac{\lambda^2}{1-(1-\lambda)^2} \right)^{1/2} \left( \mathbb{E} \sum_{m=1}^\infty (p_{m+1}^k - p_m^k)^2 \right)^{1/2} \\ &\leq \sqrt{\frac{\lambda}{2-\lambda}} \sqrt{p^k(1-p^k)} \end{aligned}$$
■

Having proved the existence of the asymptotic value we are now in the position to establish the value of the infinite game  $\Gamma_\infty(p)$ :

**THEOREM 2.12.** *For all  $p \in \Delta(K)$  the value  $v_\infty(p)$  of  $\Gamma_\infty(p)$  exists and equals  $\text{Cav } u(p)$ .*

**PROOF.** Use prop. 2.8 p. 188, theorem 2.10 p. 189 and theorem 3.1 p. 191.  $\blacksquare$

A proof for the general case is given in theorem 3.5 p. 195 below. For the full monitoring case we provide now an alternative proof by constructing optimal strategies for both players. In fact prop. 2.8 p. 188 (through prop. 2.7 p. 188) provides an optimal strategy for player I. An optimal strategy for player II is given in the next section. It is based on the notion of approachability (cf. sect. 4 p. 102) which plays a very central rôle in repeated games with incomplete information.

**2.c. Approachability strategy.** Let  $\ell = (\ell^k)_{k \in K}$  be a supporting hyperplane to  $\text{Cav } u(p)$  at  $p$  (recall that  $u$  is continuous) i.e.

$$\begin{aligned} \text{Cav } u(p) &= \langle \ell, p \rangle = \sum_k \ell^k p^k \\ \text{and} \quad u(q) &\leq \langle \ell, q \rangle \quad \forall q \in \Delta(K). \end{aligned}$$

Consider now  $\Gamma_\infty(p)$  as a game with vector pay-offs in  $\mathbb{R}^K$  (cf. sect. 4 p. 102). The  $k^{\text{th}}$  coordinate being the pay-off according to  $G^k$ .

**PROPOSITION 2.13.** *The set*

$$M = \{ m \in \mathbb{R}^K \mid m^k \leq \ell^k, \forall k \in K \}$$

*is approachable by player II.*

**PROOF.** Since  $M$  is convex, using cor. 4.6 p. 104, it suffices to prove that for all  $z \in \mathbb{R}^K$

$$w(z) \geq \inf_{m \in M} \langle m, z \rangle,$$

where  $w(z)$  is the value of the game with pay-off  $\sum_k z^k G^k$  in which player II is the maximiser. The above inequality is obviously satisfied if  $z^k > 0$  for some  $k \in K$  or if  $z = 0$ . Otherwise let  $q \in \Delta(K)$  be the normalisation (to a unit vector) of  $-z$ .

Since  $-w(-q) = u(q)$  the condition becomes

$$u(q) \leq \langle \ell, q \rangle \quad \text{for all } q \in \Delta(K)$$

and follows from the property of  $\ell$  as a supporting hyperplane. This completes the proof.  $\blacksquare$

Note that the corresponding approachability strategy is then an optimal strategy of player II since for every  $n$  the average expected pay-off up to stage  $n$  satisfies (cf. sect. 4 p. 102):

$$\bar{\gamma}_n = \langle p, \bar{g}_n \rangle \leq \langle p, \ell \rangle + \frac{K}{\sqrt{n}} = \text{Cav } u(p) + \frac{K}{\sqrt{n}}$$

hence completing the proof of theorem 2.12 p. 190. This proof leads to an explicit optimal strategy for the uninformed player using the sufficient condition for the approachability of convex sets (cf. ex. VEx.2 p. 253).

**Example 2.1 p. 185 revisited.** In this example  $D(p)$  is the matrix game:

$$p \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} + (1-p) \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} p & 0 \\ 0 & 1-p \end{pmatrix}$$

and its value is  $u(p) = p(1-p)$ . Since this is a concave function of  $p$ ,  $\text{Cav } u(p) = u(p) = p(1-p)$  and we have:

$$\lim_{n \rightarrow \infty} v_n(p) = v_\infty(p) = p(1-p)$$

(Thus  $v_\infty(1/2) = 1/4$ ). So asymptotically the value is that of the game in which none of the players is informed about the value of  $k$ . In other words the informed player has an advantage only in games of finite length. This advantage can be measured by  $v_n(p) - v_\infty(p)$ . By theorem 2.10 p. 189 this is bounded by:

$$v_n(p) - p(1-p) \leq \frac{2\sqrt{p(1-p)}}{\sqrt{n}} \leq \frac{1}{\sqrt{n}}$$

Later (cf. prop. 5.7 p. 251) we shall see that for this specific game this bound can be improved and that in fact

$$v_n(p) - p(1-p) = O\left(\frac{\ln n}{n}\right).$$

### 3. The general case

We proceed now to generalise the model by dropping the assumption of full monitoring: We no longer assume that the moves are announced after each stage but rather some individual message is transmitted to each player. We will prove that our main result so far, namely the existence of  $v_\infty$ , extends to this case. However, several significantly new ideas will be required.

The model we consider here is the same as that defined in the previous section but with no restrictions on the signalling. These are described by two finite sets of signals  $A$  (for player I),  $B$  (for player II) and transition probability  $Q$  from  $K \times S \times T$  to  $A \times B$ . We denote by  $Q_{s,t}^k$  the probability distribution at  $(k, s, t)$ .

The repeated game  $\Gamma(p)$  is played as in the full monitoring case except that at each stage  $n$ , ( $n \geq 1$ ), as player I chooses  $s_n$  and player II chooses  $t_n$ , the distribution  $Q_{s_n t_n}^k$  is used to choose  $(a_n, b_n)$  in  $A \times B$ . The signal  $a_n$  is announced to player I and  $b_n$  is announced to player II. The games  $\Gamma_n(p)$ ,  $\Gamma_\infty(p)$  and  $\Gamma_\lambda(p)$  are defined as usually based on the pay-off sequence  $(g_n)_{n=1}^\infty$ . Occasionally we write  $\Gamma$  or  $\Gamma(p)$  just for the data of the game (excluding or including  $p$ ).

**3.a.  $\lim v_n(p)$  and  $v_\infty(p)$ .** The first result for all repeated games with incomplete information on one side is the convergence of  $v_n$  and  $v_\lambda$ :

**THEOREM 3.1.**  $v_n$  and  $v_\lambda$  converge uniformly (as  $n \rightarrow \infty$  and  $\lambda \rightarrow 0$  respectively) to the same limit which can moreover be guaranteed by player II.

**PROOF.** Let  $\tau_n$  be an  $n^{-1}$ -optimal strategy of player II in  $\Gamma_n(p)$ , let  $v_{n_i}(p)$  converge to  $\liminf_{n \rightarrow \infty} v_n(p)$ . Let  $\tau$  be the following strategy of player II: For  $i = 1, 2, \dots$ , play  $n_{i+1}$  times  $\tau_{n_i}$  (thus during  $n_i n_{i+1}$  stages) before increasing  $i$  by 1.

Let us prove that this strategy  $\tau$  guarantees player II  $\liminf_{n \rightarrow \infty} v_n(p)$ . It is sufficient to show that on each bloc where (according to  $\tau$ ) player II has to play  $\tau_{n_i}$ , the average

pay-off per stage is at most  $v_{n_i}(p)$ . This follows by first computing conditional expectation (of the average pay-off) with respect to player I's  $\sigma$ -field of information  $\mathcal{H}_{n_i}^I$  and then taking expectation.

It follows that  $v_n(p)$  converges (uniformly, by theorem 1.6 p. 184).

As to the convergence of  $v_\lambda$ , the above described strategy of player II proves that

$$\limsup_{\lambda \rightarrow 0} v_\lambda(p) \leq \lim_{n \rightarrow \infty} v_n(p)$$

To complete the proof we shall prove that  $\lim_{n \rightarrow \infty} v_n(p) \leq \liminf_{\lambda \rightarrow 0} v_\lambda(p)$  by showing that  $\lim_{n \rightarrow \infty} v_n(p) \leq v_\lambda(p)$ , for any  $\lambda > 0$ . In fact, given  $\lambda > 0$  let  $\tau_\lambda$  be an optimal strategy of player II in the  $\lambda$ -discounted game and consider the following strategy (for player II): Start playing  $\tau_\lambda$  and at each stage restart  $\tau_\lambda$  with probability  $\lambda$  and with probability  $(1 - \lambda)$  continue playing the previously started  $\tau_\lambda$ . With this strategy, for any  $\varepsilon > 0$ , we have  $E(\bar{g}_n) \leq v_\lambda + \varepsilon$  for all  $n$  sufficiently large (compared to  $1/\lambda$ ). It follows that  $\lim_{n \rightarrow \infty} v_n(p) \leq v_\lambda(p)$ . ■

**REMARK 3.1.** The argument used here for the convergence of  $v_n$  is even more general, and will be used later in different context (cf. Gleason's game, ex. IV.5Ex.1 p. 177). It does not require for instance that state remains fixed throughout the game but just that player II can "almost" control it (in particular this is the case in irreducible stochastic games).

**REMARK 3.2.** If we interpret the discounted game as a repeated game with probability  $\lambda$  of stopping after each stage then the convergence of  $v_\lambda$  can be generalised as follows: Let  $a = \{a_n\}_{n=1}^\infty$  be a probability distribution on the positive integers with  $a_n \geq a_{n+1}$ . Let  $\Gamma_a$  have  $\sum_n a_n g_n$  as pay-off function, and  $v_a(p)$  as value. If  $\{a^\ell\}_{\ell=1}^\infty$  is a sequence of such distributions with  $a_1^\ell$  going to zero then  $\lim_{\ell \rightarrow \infty} v_{a^\ell}(p) = \lim_{n \rightarrow \infty} v_n(p)$ .

**3.b. The non-revealing game.** With the concavity properties proved in sect. 1, the next step is to get lower bounds for the various values i.e. to get the results of prop. 2.8 p. 188 for the general case. To do that we need to extend the notion of the non-revealing game.

The main feature of repeated games with incomplete information is the possibility of a player to collect information about the state of nature along the play of the game. This information he deduces from the sequence of signals he receives. In the games considered in this chapter, the uninformed player, player II, tries to learn about  $k$  from the signals  $(b_1, b_2, \dots)$  which he receives in stages 1, 2, ... respectively. In the full monitoring case  $b_n$  is just  $(s_n, t_n)$ , the (pure) moves of the players at stage  $n$ . In the general case  $b_n$  is a random variable whose distribution is the marginal distribution on  $B$  of  $Q_{s_n, t_n}^k$  where  $k$  is the state chosen at stage 0. Since player I knows  $k$  and his moves typically depend on this knowledge, his moves may be revealing to player II, i.e. they may enable him to learn something about  $k$  via the signals. This motivates the following definition:

As usually  $x = (x^k)_{k \in K} \in X^K$  denotes a strategy of player I in the one-stage game  $\Gamma_1(p)$ , where  $X = \Delta(S)$  is the set of his mixed moves.

**DEFINITION 3.2.**  $x \in X^K$  is called **non-revealing** at  $p \in \Delta(K)$  if

$$(1) \quad \text{For any } t \in T : \quad P_{x,t}^p(b) > 0 \implies P_{x,t}^p(k | b) = p^k \quad \forall k \in K, \forall b \in B.$$

Denote by  $\text{NR}(p)$  the set of non-revealing strategies  $p$ .

To obtain an operational expression for  $\text{NR}(p)$ , let  $Q^{\Pi,k}$  be the  $S \times T$  matrix whose  $st$  element, denoted by  $Q_{st}^{\Pi,k}$  is the marginal distribution on  $B$  of  $Q_{st}^k$ . The elements  $Q_{st}^{\Pi,k}$  are probability vectors in the simplex  $\Delta(B)$ . If in  $\Gamma_1(p)$  player I uses  $x = (x^k)_{k \in K}$  and player II uses  $y \in \Delta(T)$ , then if the state is  $k$ , the probability distribution of the signal  $b$  received by II is  $x^k Q^{\Pi,k} y$  ( $x^k$  is thought of as a row vector and  $y$  as a column vector.)

LEMMA 3.3.

$$\text{NR}(p) = \{ x \in X^K \mid x^k Q^{\Pi,k} = x^{k'} Q^{\Pi,k'} \text{ whenever } p^k > 0 \text{ and } p^{k'} > 0 \}$$

(For any move  $t \in T$  of player II, the distribution of  $b$  (induced by  $t$  and  $x^k$ ) in the  $k^{\text{th}}$  state, is the same for all  $k$  for which  $p^k > 0$ ).

PROOF.  $x \in \text{NR}(p)$  means that  $\forall b \in B$

$$p^k > 0, p^{k'} > 0 \implies P_{p,x,t}(k \mid b) = P_{p,x,t}(k' \mid b)$$

But since

$$P_{p,x,t}(k \mid b) = \frac{p^k (x^k Q_t^{\Pi,k})(b)}{\sum_{k' \in K} p^{k'} (x^{k'} Q_t^{\Pi,k'})(b)}$$

$P_{p,x,t}(k \mid b) = P_{p,x,t}(k' \mid b)$  is equivalent to  $P_{p,x,t}(b \mid k) = P_{p,x,t}(b \mid k')$ . ■

COMMENT 3.3. Note that if  $p_0$  is an extreme point of  $\Delta(K)$  then  $\text{NR}(p_0) = X^K$ . This follows readily from the definition of  $\text{NR}(p)$  and is intuitively obvious: An extreme point of  $\Delta(K)$  corresponds to a situation of complete information, where  $k$  is known to both players, hence every strategy of I is non-revealing since there is nothing to reveal. Note also that by definition 3.2 above, the posterior probabilities remain constant as a result of any sequence of moves in  $\text{NR}(p)$ . This will be intensively used in the sequel. Finally remark that  $\text{NR}(p)$  depends only on the support  $K(p)$  of  $p$ .

DEFINITION 3.4. Given  $\Gamma$  and  $p$ , the **non-revealing game** ( $\text{NR}$ -game), denoted by  $D(p)$ , is the one-stage game in which player I's strategy set is restricted to  $\text{NR}(p)$ .

We denote by  $u(p)$  the value of  $D(p)$  and refer to it as the  $\text{NR}$ -value. If  $\text{NR}(p) = \emptyset$  we define  $u(p) = -\infty$ .

REMARK 3.4.  $\text{Cav } u$  is Lipschitz with constant  $C$ .

In the following examples (for  $x \in [0, 1]$  we abbreviate  $(1 - x)$  by  $x'$ ):

$$K = \{1, 2\} \quad B = \{a, b, c, \dots\} \quad (p, p') \in \Delta(K)$$

It will be convenient to denote elements of  $\Delta(B)$  by  $a$  (meaning: signal  $a$  with probability 1),  $b$ ,  $c$ ,  $\dots$  or by  $\frac{1}{2}b + \frac{1}{2}c$  (meaning signals  $b$  or  $c$  with probability 1/2 each) etc.

EXAMPLE 3.5. The first important example is that of games with full monitoring treated in previous section. This is the special case in which for all  $k$ ,  $s$ , and  $t$ ,  $Q_{st}^k$  is a unit mass at the signal  $b = (s, t)$ . It worth recalling the main properties of this case for future reference:

- (1) The projection of  $\text{NR}(p)$  on the support of  $p$  lies on the diagonal of  $X^K$  i.e. player I plays independently of his information about  $k$ .
- (2) The signals received by player II are independent of his own moves, therefore player II has no “information incentive” to play a certain action.
- (3) The information that player II can get on  $k$  is only through the moves of player I (since the signals are state independent.)

As we will see later, dropping some of the above properties imply adding substantial difficulty and complexity to the analysis.

EXAMPLE 3.6. Consider the game with  $K = \{1, 2\}$ , pay-off matrices

$$G^1 = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \quad G^2 = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$$

and with signalling matrices:

$$Q^1 = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \quad Q^2 = \begin{pmatrix} c & d \\ a & b \end{pmatrix}$$

For  $0 < p < 1$ ,  $x \in \text{NR}(p)$  has to satisfy:

$$\begin{aligned} x_1 a + x'_1 c &\stackrel{*}{=} x_2 c + x'_2 a \\ x_1 b + x'_1 d &\stackrel{*}{=} x_2 d + x'_2 b \end{aligned}$$

which implies  $x_1 = x'_2$  and so for  $0 < p < 1$ ,

$$\text{NR}(p) = \{ ((\alpha, \alpha'), (\alpha', \alpha)) \mid 0 \leq \alpha \leq 1 \}$$

and

$$D(p) = \begin{pmatrix} p & p' \\ 0 & 0 \end{pmatrix} \quad u(p) = \min(p, p')$$

The equation for  $u(p)$  is valid for all  $p \in [0, 1]$  since  $u(1) = u(0) = 0$ .

Remark that unlike in the full monitoring case where non-revealing strategies were strategies that did not use the information about  $k$ , in this example, in order to play non-revealingly, player I has to use his information about the state  $k$ : He has to play differently in the two games. For example the optimal strategy in  $D(p)$  for  $0 < p < 1$  is to play the top row if  $k = 1$  and the bottom row if  $k = 2$ .

EXAMPLE 3.7. The same game as in previous example with

$$Q^1 = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \quad Q^2 = \begin{pmatrix} a & e \\ f & d \end{pmatrix}$$

For  $0 < p < 1$ , a strategy  $x \in \text{NR}(p)$  has to satisfy:

$$\begin{aligned} x_1 a + x'_1 c &\stackrel{*}{=} x_2 a + x'_2 f \\ x_1 b + x'_1 d &\stackrel{*}{=} x_2 e + x'_2 d \end{aligned}$$

which is impossible hence  $\text{NR}(p) = \emptyset$  for  $0 < p < 1$  and therefore

$$u(p) = \begin{cases} 0 & \text{if } pp' = 0 \\ -\infty & \text{if } pp' > 0 \end{cases}$$

EXAMPLE 3.8. Let  $K = \{1, 2, 3\}$  and let the signalling matrices for both players be

$$Q^1 = Q^2 = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \quad Q^3 = \begin{pmatrix} a' & b' \\ c' & d' \end{pmatrix}$$

The signals following the first stage are **partially revealing**: If the signals are in  $\{a', b', c', d'\}$  it will become common knowledge that  $k = 3$  and the game is reduced to a

complete information game in which  $G^3$  is played repeatedly. If on the other hand the signals are in  $\{a, b, c, d\}$  the game will be reduced to  $\hat{\Gamma}(\hat{p})$  in which

$$\hat{K} = \{1, 2\} \quad \hat{p} = (p^1, p^2)/(p^1 + p^2)$$

and pay-off matrices are  $G^1, G^2$ . In either case, the game resulting after the first stage is one with full monitoring.

**3.c. Study of  $v_\infty(p)$ .** We intend to prove here:

**THEOREM 3.5.**  $v_\infty(p) = \text{Cav } u(p)$  for all  $p$  in  $\Pi$ .

**PROOF.** By theorem 3.1 p. 191 we know that  $\lim v_n(p)$  (and similarly  $\lim_{\lambda \rightarrow 0} v_\lambda$ ) exist and player II can guarantee this limit in  $\Gamma_\infty(p)$ . As for player I, observe first that in the general case also player I can guarantee  $u(p)$  in  $\Gamma_n(p)$  (as well as in  $\Gamma_\infty(p)$  and  $\Gamma_\lambda(p)$ ). This is obvious if  $\text{NR}(p) = \emptyset$  (since then  $u(p) = -\infty$ ), otherwise it is achieved by playing i.i.d. at each stage a fixed optimal strategy in  $D(p)$ . Combined with prop. 2.7 p. 188 this yields the analogue of prop. 2.8 p. 188 namely:

**PROPOSITION 3.6.** For all  $p \in \Delta(K)$  Player I can guarantee  $\text{Cav } u(p)$  in  $\Gamma_\infty(p)$ . Moreover,  $v_n(p)$  and  $v_\lambda(p)$  are both at least  $\text{Cav } u(p)$  (for all  $n$  and all  $\lambda > 0$ ).

It follows from the above results that the existence of  $v_\infty(p)$  will be established if we prove that  $\lim_{n \rightarrow \infty} v_n(p) = \text{Cav } u(p)$ . For this we shall prove the counterpart of the above prop. 3.6.

**PROPOSITION 3.7.**  $\limsup v_n(p) \leq \text{Cav } u(p)$ .

**PROOF.** In proving this inequality, we need the posterior probabilities  $p_n^k$  to be known by both players. To achieve this we modify the game so that after each stage, both the signal of player II and the pair of moves are communicated also to player I (in addition to his own signal). With this new signalling structure  $\mathcal{H}^{\text{II}} \subseteq \mathcal{H}^{\text{I}}$ , and hence  $\mathcal{H}^{\text{II}}$  is public knowledge. Since, being in favour of player I, this modification can only increase  $v_n$  it is sufficient to prove the above inequality for the modified game.

Let  $\zeta$  denote a II-history in  $H^{\text{II}}$ , i.e. a sequence of signals to II where the signal includes II's move. As in the previous section we denote by  $p_n$  the conditional probability on  $K$  given II's information up to stage  $n$  (i.e.  $p_n^k = P(k | \mathcal{H}_n^{\text{II}}) \forall k \in K$ ). For the evaluation at  $\zeta$  we write  $p_n(\zeta)$  (if  $n = 1$ , as the history  $\omega$  consists of one signal  $b$ , we simply write  $p(b)$ ).

As we already observed, the key object in measuring the amount of information revealed by player I at stage  $n$  is  $E\{\|p_{n+1} - p_n\| \mid \mathcal{H}_n^{\text{II}}\}$ . This quantity depends not only on the mixed move of player I at that stage but also on  $p_n$ . Intuitively, he may play 'very differently' at state  $k$  with very small  $p^k$  and still reveal 'very little'. This indicates that the appropriate space to work with is the space of probabilities on the product space  $K \times S$ . So let  $Z = \Delta(K \times S)$ . Given  $z \in Z$  denote its marginal on  $K$  by  $p_z = (p_z^k)_{k \in K} \in \Delta(K)$ , and its conditional on  $S$  given  $K$  by  $x_z \in X^K$  where, for  $k \in K$  with  $p_z^k > 0$  and  $s \in S$ ,  $x_z^k(s) = z(k, s)/p_z^k$ . Hence any pair  $(p, x)$  with  $p \in \Delta(K)$  and  $x \in X^K$  determines uniquely a point  $z(p, x) \in Z$ .

Consider the following subset of  $Z$ :

$$Z_0 = \{z \in Z \mid x_z \in \text{NR}(p_z)\}$$

For  $z \in Z$  denote by  $d(z, Z_0)$  its (Euclidian) distance from  $Z_0$ . For  $z \in Z$  and  $y \in Y$  define

$$e(z, y) = E_{x_z, y}(\|p_z(b)\|^2 - \|p_z\|^2)$$

which can also be written as

$$e(z, y) = \mathbb{E}(\|p_z(b) - p_z\|^2) = \mathbb{E} \sum_k (p_z^k(b) - p_z^k)^2.$$

LEMMA 3.8. *Given a completely mixed  $y$ , then  $\forall \xi > 0$ ,  $\exists \eta > 0$  such that*

$$e(z, y) < \eta \Rightarrow d(z, Z_0) < \xi$$

PROOF. If the lemma is false then there exist  $(\eta_j > 0)_{j=1}^\infty$  satisfying  $\lim_{j \rightarrow \infty} \eta_j = 0$  and  $(z_j)_{j=1}^\infty$  such that  $\forall j$ :

$$e(z_j, y) < \eta_j \quad \text{and} \quad d(z_j, Z_0) > \xi$$

We may assume without loss of generality that  $(z_j)$  converges, say to  $z$  which (by continuity of  $e(\cdot, y)$  and  $d(\cdot, Z_0)$ ) will then satisfy

$$e(z, y) = 0 \quad \text{and} \quad d(z, Z_0) \geq \xi,$$

in contradiction with lemma 3.2 p. 192, since if  $y$  is completely mixed (i.e.  $y(t) > 0$  for all  $t \in T$ ) then  $e(z, y) = 0$  if and only if  $z \in Z_0$ .  $\blacksquare$

We need a version of the last lemma in which  $\eta$  does not depend on  $y$  provided it is in  $Y_\varepsilon = \{y \in Y \mid y(t) \geq \varepsilon\}$ .

COROLLARY 3.9. *Given  $\xi > 0$  and  $\varepsilon > 0$ ,  $\exists \eta > 0$  such that  $\forall y \in Y_\varepsilon$*

$$e(z, y) < \eta \Rightarrow d(z, Z_0) < \xi$$

PROOF. Note that  $e(z, y)$  is linear in  $y$ , so if  $y^0$  is the uniform distribution on  $T$  we have:

$$y \in Y_\varepsilon \Rightarrow e(z, y) \geq \varepsilon e(z, y^0)$$

For  $\xi$  and  $y^0$ , let  $\tilde{\eta}$  be determined by lemma 3.8 p. 196, and take  $\eta = \varepsilon \tilde{\eta}$ .  $\blacksquare$

Given  $t \in T$  and  $z$ , we consider  $G_t = (G_{st}^k)_{s \in S}^{k \in K}$  as a point in  $\mathbb{R}^{K \times S}$ . The expected pay-off can then be written as  $\langle z, G_t \rangle = \sum_k p^k x^k G_t^k$ . For  $y \in Y$  and  $\varepsilon < 1/\#T$  define the  $\varepsilon$ -perturbation of  $y$  (denoted by  $y^\varepsilon$ ) as follows:  $y^\varepsilon(t) = (1 - \#T\varepsilon)y(t) + \varepsilon$ ,  $\forall t \in T$ . We say that  $y \in Y$  is a best reply to  $z$  if

$$\langle z, G_t \rangle < \langle z, G_{t'} \rangle \Rightarrow y(t') = 0$$

LEMMA 3.10. *There exists a constant  $\tilde{C}$  such that  $\forall \varepsilon > 0$ ,  $\exists \eta > 0$  such that if  $y$  is a best reply to  $z$  and  $e(z, y^\varepsilon) < \eta$  then*

$$\gamma(z, y^\varepsilon) = \langle z, G y^\varepsilon \rangle \leq \mathsf{Cav} u(p) + \tilde{C}\varepsilon.$$

PROOF. Recall that  $C$  is an upper bound for all pay-offs' absolute values and take  $\tilde{C} = 1 + C(1 + \#S)$ . Let

$$f(z) = \min_{t \in T} \gamma(z, t) = \min_{t \in T} \langle z, G_t \rangle.$$

This is a continuous function on  $Z_0$  as the minimum of finitely many such functions. By definition of  $Z_0$ ,

$$z \in Z_0 \Rightarrow x_z \in \mathbf{NR}(p_z) \Rightarrow f(z) \leq u(p_z)$$

On the other hand since  $y$  is a best reply to  $z$ ,  $\gamma(z, y) = f(z)$ . But  $\gamma(z, y)$  is linear in  $y$  and hence Lipschitz with constant  $C$ , so we have:

$$\gamma(z, y^\varepsilon) \leq f(z) + C\varepsilon.$$

Let  $\psi(\varepsilon)$  be the modulus of uniform continuity of  $f$ . For  $\varepsilon > 0$  apply cor. 3.9 p. 196 with  $\xi = \min(\psi(\varepsilon), \varepsilon)$  and  $\varepsilon$  to obtain  $\eta > 0$ . To see that it has the required property, since  $e(z(p, x), y^\varepsilon) < \eta$  it follows from Corollary 3.9 that  $\exists z_0 \in Z_0$  such that  $\|z(p, x) - z_0\| < \psi(\varepsilon)$  and hence  $|f(z(p, x)) - f(z_0)| < \varepsilon$ . So we obtain:

$$\gamma(x, y^\varepsilon) \leq f(z(p, x)) + C\varepsilon \leq f(z_0) + \varepsilon + C\varepsilon \leq u(p_{z_0}) + (1 + C)\varepsilon.$$

Finally note that, if  $z_0 = (p_0, x_0)$ , then  $\|p - p_0\| \leq \#S \|z(p, x) - z_0\| < \#S\varepsilon$ . Also, since  $\mathbf{Cav} u(p)$  is Lipschitz with constant  $C$ ,

$$\mathbf{Cav} u(p_0) \leq \mathbf{Cav} u(p) + C \|p - p_0\| \leq \mathbf{Cav} u(p) + C \#S\varepsilon$$

We conclude that

$$(2) \quad \gamma(x, y^\varepsilon) \leq \mathbf{Cav} u(p) + C \#S\varepsilon + (1 + C)\varepsilon = \mathbf{Cav} u(p) + \tilde{C}\varepsilon. \quad \blacksquare$$

Given any strategy  $\sigma$  of player I and any  $\varepsilon > 0$  we construct now the following reply  $\tau = \tau(\sigma, \varepsilon)$  for player II: At any stage  $m$ , given a II-history  $\zeta$  play a mixed move  $y_m(\zeta)$  which is an  $\varepsilon$ -perturbation of a best reply to  $(p_m(\zeta), x_m(\zeta))$ . By lemma 3.10 p. 196  $\exists \eta > 0$  such that for any stage  $m$

$$e_m(\zeta) \stackrel{\text{def}}{=} e(p_m(\zeta), x_m(\zeta), y_m(\zeta)) < \eta \implies g_m(\zeta) \leq \mathbf{Cav} u(p_m(\zeta)) + 2\tilde{C}\varepsilon$$

Now for any  $n$ :

$$\#K \geq \mathbb{E}(\|p_{n+1}(\zeta)\|^2 - \|p\|^2) = \sum_{m=1}^n \mathbb{E}(\|p_{m+1}(\zeta)\|^2 - \|p_m(\zeta)\|^2) = \sum_{m=1}^n \mathbb{E}(e_m(\zeta)).$$

It follows that  $\mathbb{E}(e_m(\zeta)) \leq \varepsilon\eta$  for at least  $[n - \#K/(\varepsilon\eta)]$  stages, and on these stages  $P_{\sigma, \tau}^p \{e_m(\zeta) \geq \eta\} < \varepsilon$ . On this last event we have (remark that  $C \leq \tilde{C}$ )

$$g_m(\zeta) \leq \mathbf{Cav} u(p_m(\zeta)) + 2\tilde{C}\varepsilon,$$

while on the complement (by lemma 3.10 p. 196)

$$g_m(\zeta) \leq \mathbf{Cav} u(p_m(\zeta)) + 2\tilde{C}\varepsilon.$$

Thus, for stages  $m$  such that  $\mathbb{E}(e_m(\zeta)) \leq \varepsilon\eta$  we get

$$\gamma_m \leq \mathbb{E}(\mathbf{Cav} u(p_m(\zeta)) + 4\tilde{C}\varepsilon) \leq \mathbf{Cav} u(p) + 4\tilde{C}\varepsilon.$$

On the other stages we majorate the expected pay-off by  $\gamma_m \leq 2\tilde{C}$ . Taking the average on  $n$  stages we have:

$$\bar{\gamma}_n \leq \mathbf{Cav} u(p) + 4\tilde{C}\varepsilon + 2C \frac{\#K}{n\varepsilon\eta}.$$

For  $n \geq N > \#K/(\varepsilon^2\eta)$  the last term is less than  $2\tilde{C}\varepsilon$ , hence we conclude:

For any strategy of player I and for any  $\varepsilon > 0$  there is a reply strategy of player II and  $N$  such that for  $n \geq N$

$$\bar{\gamma}_n \leq \mathbf{Cav} u(p) + 6\tilde{C}\varepsilon$$

Hence for  $n \geq N$

$$v_n(p) \leq \mathbf{Cav} u(p) + 6\tilde{C}\varepsilon$$

implying  $\limsup_{n \rightarrow \infty} v_n(p) \leq \mathbf{Cav} u(p)$ , and hence  $\lim_{n \rightarrow \infty} v_n(p) = \mathbf{Cav} u(p)$ . This completes the proof of prop. 3.7 p. 195, hence of theorem 3.5 p. 195.  $\blacksquare$

**3.d. Optimal strategy for the uninformed player.** We shall now provide an approachability strategy for player II which guarantees  $\text{Cav } u(p)$ . This generalises the (full monitoring) strategy in sect. 2.c p. 190.

The main feature of the optimal strategy for player II in the full monitoring case (cf. the proof of theorem 2.12 p. 190) is that it is based on the statistics  $\bar{g}_n = \frac{1}{n} \sum_{m=1}^n g_m$  viewed as a vector pay-off function on histories. This is the average of the stage vector-pay-offs  $(G_{s_m t_m}^k)_{k \in K}$  which are observable by player II in the full monitoring case since he observes the moves  $(s_m, t_m)$  for  $m = 1, \dots, n$ . In the general case  $\bar{g}_n$  being no longer observable by player II, another optimal strategy is to be provided which is based only on the II-history  $(b_1, \dots, b_n)$  available to him at each stage. Before defining formally this strategy let us discuss first the principal ideas involved.

For any signal  $b \in B$ , any move  $t \in T$  and any stage  $n$ , let  $\beta_n^{tb}$  be the proportion of stages, up to stage  $n$ , in which  $b$  was obtained by player II following a move  $t$ , out of all stages in which move  $t$  was played, i.e.

$$\beta_n^{tb} = \frac{\#\{m \mid m \leq n, b_m = b, t_m = t\}}{\#\{m \mid m \leq n, t_m = t\}}$$

The vector  $\beta_n = (\beta_n^{tb})_{t \in T, b \in B}$ , which is observable by player II after each stage  $n$ , is the basis for his strategy. The vector pay-off  $\xi_n$  which plays the rôle of the non-observable  $g_n$  is, roughly speaking, the worst vector pay-off which is compatible (up to a small deviation  $\delta$ ) with the observed vectors  $\beta_1, \dots, \beta_n$  and with the assumption that player I was playing i.i.d. To this vector pay-off we shall apply Blackwell's approachability theorem. The definition of  $\xi_n$  and the strategy of player II will be such that:

- The  $\xi$ -pay-off will be as close as we wish to  $\text{Cav } u(p)$ .
- The actual unobserved pay-off will not exceed the observed  $\xi$ -pay-off by more than arbitrarily small  $\varepsilon$ .

**3.d.1. The strategy construction.** Let us start by some notations which will be used in the construction of the strategy.

Let  $\mathcal{B} = [\Delta(B)]^T$ . This is the set which contains all possible  $\beta_n$  vectors.

For  $y \in \Delta(T)$  let  $\tilde{y}$  be the strategy (in  $\Gamma_n$  or  $\Gamma_\infty$ ) which plays the mixed move  $y$  repeatedly and independently at each stage.

For  $(s, t) \in S \times T$  we denote by  $f_n^s$  and  $\varphi_n^{st}$  respectively, the frequencies of  $s$  and  $(s, t)$  up to stage  $n$ :

$$f_n^s = \frac{1}{n} \#\{m \mid m \leq n, s_m = s\} \quad f_n = (f_n^s)_{s \in S}$$

$$\varphi_n^{st} = \frac{1}{n} \#\{m \mid m \leq n, s_m = s, t_m = t\}$$

LEMMA 3.11. Given  $\tilde{y}$  and  $\sigma$  in  $\Gamma_\infty(p)$  then  $\forall (s, t) \in S \times T$

$$P_{\sigma \tilde{y}} \left( \lim_{n \rightarrow \infty} (\varphi_n^{st} - f_n^s y_t) = 0 \right) = 1$$

PROOF. For fixed  $(s_0, t_0) \in S \times T$ , consider the game in which at each stage  $m$  player I chooses  $s_m \in S$  and the pay-off is  $R(s_m)$  where  $R(s) = \mathbb{1}_{s=s_0, t=t_0} - \mathbb{1}_{s=s_0} y_{t_0}$ , and  $t$  is a random variable with distribution  $y$ , (thus player II is dummy.) Since

$$\mathbb{E}(\mathbb{1}_{s=s_0, t=t_0}) = P(s = s_0)P(t = t_0 \mid s = s_0) = \mathbb{E}(\mathbb{1}_{s=s_0})y_{t_0},$$

we have  $\mathbb{E}(R(s)) = 0 \ \forall s \in S$ .

It follows from ex. II.4Ex.4 p. 105 that the set  $\{0\}$  is approachable by any strategy  $\sigma$  of player I. The approachability in this case amounts to:

$$P_{\sigma, \tilde{y}} \left( \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{m=1}^n R(s_m) = 0 \right) = 1$$

which is equivalent to:

$$(3) \quad P_{\sigma, \tilde{y}} \left( \lim_{n \rightarrow \infty} (\varphi_n^{s_0 t_0} - f_n^{s_0} y_{t_0}) = 0 \right) = 1 \quad \blacksquare$$

COMMENT 3.9. As noted in ex. II.4Ex.4 p. 105, we applied here a version of a strong law of large numbers. The one-player auxiliary game that we defined is actually a situation in which a gambler (player I) is to choose a betting system (cf. e.g. Feller, 1966, vol. I, p. 199): At each stage  $m$ , based on his past gamble experience at stages  $1, \dots, m-1$  only, he decides either not to participate in the gamble (he chooses  $s \neq s_0$ ) or to participate (he chooses  $s = s_0$ ) in which case he either wins  $1 - y_{t_0}$  (if  $t = t_0$ ) or loses  $y_{t_0}$  (if  $t \neq t_0$ ). The form of the strong law of large numbers states that whatever betting system he chooses, his average net profit tends to zero with probability 1 exactly as it does if he participates in all stages (by the usual strong law of large numbers).

LEMMA 3.12. Let  $\tilde{y}$  be a stationary strategy of player II such that  $y(t) > 0 \ \forall t \in T$  then for any strategy  $\sigma$  of player I and for all  $k \in K$ :

$$(4) \quad P_{\sigma^k, \tilde{y}}^k \left( \lim_{n \rightarrow \infty} (\bar{g}_n - f_n G y) = 0 \right) = 1$$

$$(5) \quad P_{\sigma^k, \tilde{y}}^k \left( \lim_{n \rightarrow \infty} (\beta_n - f_n Q) = 0 \right) = 1$$

PROOF. By definition  $\bar{g}_n^k = \sum_{s \in S} \sum_{t \in T} \varphi_n^{st} G_{st}^k$ , so (4) follows from lemma 3.11 p. 198. Now when the state is  $k$  and moves  $(s, t)$  are played, the random signal  $b$  to player II has the distribution  $Q_{st}^k$ . So

$$(f_n Q^k)^{tb} = \sum_{s \in S} f_n^s Q_{st}^k(b) \ \forall n, \forall t \in T, \forall b \in B$$

On the other hand  $\beta_n^{tb}$  can be written as:

$$\beta_n^{tb} = \sum_{s \in S} \left( \varphi_n^{st} \frac{\#\{m \mid m \leq n, s_m = s, t_m = t, b_n = b\}}{\#\{m \mid m \leq n, s_m = s, t_m = t\}} \right) / \frac{\#\{m \mid m \leq n, t_m = t\}}{n}$$

Since player II is playing  $\tilde{y}$ , by the strong law of large numbers:

$$\lim_{n \rightarrow \infty} \frac{\#\{m \mid m \leq n, t_m = t\}}{n} = y_t > 0 \text{ a.s.}$$

and given  $k$ :

$$\lim_{n \rightarrow \infty} \varphi_n^{st} \left[ \frac{\#\{m \mid m \leq n, s_m = s, t_m = t, b_n = b\}}{\#\{m \mid m \leq n, s_m = s, t_m = t\}} - Q_{st}^k(b) \right] = 0 \text{ a.s.}$$

It follows that with probability 1 and  $\forall k, \forall b$ ,

$$\lim_{n \rightarrow \infty} (\beta_n^{tb} - (f_n Q^k)^{tb}) = \lim_{n \rightarrow \infty} \sum_{s \in S} (\varphi_n^{st} / y_t - f_n^s) Q_{st}^k(b),$$

and this limit is 0 with probability 1 by lemma 3.11 p. 198.  $\blacksquare$

COMMENT 3.10. The meaning of lemma 3.12 is that when player II uses for a long period a stationary strategy which assigns positive probability to all pure moves then, up to an arbitrarily small error he may assume, both for pay-offs considerations and for information considerations, that player I has been also using a stationary strategy, namely

the stationary strategy  $f_n = (f_n^s)_{s \in S}$ . Note furthermore that if we use the version of the law of large numbers derived from Blackwell's approachability theorem (4.3 p. 102 and ex. II.4Ex.4 p. 105) we obtain an exact bound for the speed of convergence which can be used to compute the length of the period required for a given level of error. This bound does not depend on  $y$  in the first formula and depends on  $1/y_t$  in the second one.

For the rest of the construction of an optimal strategy for player II it will be convenient to have the following modifications of the game. Consider the game  $\Gamma_\infty(p_0)$ . Let  $\ell \in \mathbb{R}^K$  be a supporting hyperplane to  $\text{Cav } u$  at  $p_0$ , i.e.  $\langle \ell, p_0 \rangle = \text{Cav } u(p_0)$  and  $\text{Cav } u(p) \leq \langle \ell, p \rangle \forall p \in \Delta(K)$ . By subtracting  $\ell^k$  from all entries of  $G^k$  we get a new game for which  $\text{Cav } u(p_0) = 0$  and  $u(p) \leq 0, \forall p \in \Delta(K)$ . Clearly any strategy that guarantees 0 in the new game guarantees  $\text{Cav } u(p_0)$  in the original game. Next, by dividing all pay-offs in the (modified) game by an appropriate positive number we may assume without loss of generality that  $|G_{st}^k| \leq 1$ , for all  $k, s$  and  $t$ .

In view of lemma 3.12 p. 199 let us define the following functions: For  $\delta > 0$ ,  $\beta \in \mathcal{B}$  and  $y \in \Delta(T)$  let:

$$\begin{aligned} F(k, \beta, \delta) &= \left\{ f \in \Delta(S) \mid \|fQ^k - \beta\| \leq \delta \right\} \\ \xi^k(\beta, y, \delta) &= \max\{ fG^ky \mid f \in F(k, \beta, \delta) \} \end{aligned}$$

Here  $\|\cdot\|$  is the  $\ell_1$  norm in  $\mathbb{R}^{T \times B}$ . If  $F(k, \beta, \delta) = \emptyset$ , we define  $\xi^k(\beta, y, \delta) = -\infty$ . The meaning of these functions is quite straightforward: If  $\beta$  is the observed vector of frequencies of signals to player II, then  $F(k, \beta, \delta)$  is the set of all stationary strategies  $f$  in  $\Delta(S)$  of player I which could have yielded in state  $k$  a frequency vector which is  $\delta$ -close to the observed  $\beta$ . Consequently, when it is finite,  $\xi^k(\beta, y, \delta)$  is the worst pay-off that player II might have paid in state  $k$  while playing  $\tilde{y}$ . The following proposition formalises this interpretation:

**LEMMA 3.13.** *Given  $y$  with  $y(t) > 0$  for all  $t \in T$ , and given  $\eta > 0$  and  $\delta > 0$ , there exists an  $M = M(\eta, \delta)$  such that for any strategy  $\sigma$  of player I and any  $k \in K$ ,  $m \geq M$  implies:*

$$P_{\sigma^k, \tilde{y}}\{\bar{g}_m^k > \xi^k(\beta_m, y, \delta) + \eta\} < \eta.$$

**PROOF.** Since convergence with probability 1 implies convergence in probability, by lemma 3.12 p. 199 there is an  $M = M(\eta, \delta)$  such that for any strategy  $\sigma$  of player I,  $m \geq M$  implies:

$$\begin{aligned} P_{\sigma^k, \tilde{y}}\{f_m \notin F(k, \beta_m, \delta)\} &< \eta/2 \\ P_{\sigma^k, \tilde{y}}\{\bar{g}_m^k > f_m G^k y + \eta\} &< \eta/2 \end{aligned}$$

Since in the intersection of the complements of the above two events  $\xi^k(\beta_m, y, \delta) \geq f_m G^k y$  and  $\bar{g}_m^k \leq f_m G^k y + \eta$  which imply  $\bar{g}_m^k \leq \xi^k(\beta_m, y, \delta) + \eta$ , it follows that

$$(6) \quad P_{\sigma^k, \tilde{y}}\{\bar{g}_m^k > \xi^k(\beta_m, y, \delta) + \eta\} < \eta. \quad \blacksquare$$

Note that by the remark following lemma 3.12 p. 199, the dependence of the constant  $M$  on  $y$  is through  $1/y_t$  therefore if for  $\varepsilon > 0$  we let

$$\Delta_\varepsilon(T) = \{ y \in \Delta(T) \mid y_t \geq \frac{\varepsilon}{\#T} \ ; \ \forall t \in T \},$$

we have:

COROLLARY 3.14. In lemma 3.13 there is an  $M$  satisfying the statement uniformly for all  $y \in \Delta_\varepsilon(T)$ .

We denote now

$$\begin{aligned} X(k, \beta, \delta) &= \{x \in X^K \mid x^k \in F(k, \beta, \delta)\} \\ \delta\text{NR}(p) &= \bigcup_{\beta \in \mathcal{B}} \bigcap_{k: p^k > 0} X(k, \beta, \delta) \end{aligned}$$

The set  $\delta\text{NR}(p)$  is compact since it is the projection on  $X^K$  of the compact set

$$\left\{(x, \beta) \in X^K \times \mathcal{B} \mid \|x^k Q^k - \beta\| \leq \delta \forall k \in K(p)\right\}$$

and by definition  $\bigcap_{\delta > 0} \delta\text{NR}(p) = \text{NR}(p)$ .

It follows that  $\max_{x \in \delta\text{NR}(p)} \langle c, x \rangle$  converges uniformly to  $\max_{x \in \text{NR}(p)} \langle c, x \rangle$  as  $\delta$  goes to 0, for any  $c$  in the unit ball of  $\mathbb{R}^{K \times S}$ . Since there are finitely many distinct sets  $\text{NR}(p)$  we obtain the following:

LEMMA 3.15. For every  $\varepsilon > 0$  there exists  $\bar{\delta} = \bar{\delta}(\varepsilon)$  such that  $\delta \in (0, \bar{\delta}]$  implies

$$\max_{x \in \delta\text{NR}(p)} \langle c, x \rangle \leq \max_{x \in \text{NR}(p)} \langle c, x \rangle + \varepsilon,$$

for all  $c \in \mathbb{R}^{K \times S}$  with  $\|c\| \leq 1$  and for all  $p$ .

For  $p \in \Delta(K)$  let  $Y(p) = \{y \in \Delta(T) \mid y \text{ is optimal for player II in } D(p)\}$ .

LEMMA 3.16. For each  $\varepsilon > 0$  there exists  $\bar{\delta} = \bar{\delta}(\varepsilon)$  such that  $\forall \beta \in \mathcal{B}$

$$y \in Y(p) \implies \langle p, \xi(\beta, y, \delta) \rangle \leq \varepsilon \quad \forall \delta \in (0, \bar{\delta}]$$

PROOF.

$$\begin{aligned} \langle p, \xi(\beta, y, \delta) \rangle &= \sum_{k \in K} p^k \max_{f \in F(k, \beta, \delta)} f G^k y \\ &= \max_{x \in \bigcap_{k: p^k > 0} X(k, \beta, \delta)} \sum_{k \in K} p^k x^k G^k y \end{aligned}$$

So by definition of  $\delta\text{NR}(p)$ ,

$$\langle p, \xi(\beta, y, \delta) \rangle \leq \max_{x \in \delta\text{NR}(p)} \sum_{k \in K} p^k x^k G^k y$$

The proof then follows from lemma 3.15 p. 201 and from the optimality of  $y$  in  $D(p)$ . ■

We will use now the approachability theorem to prove that player II can decrease  $\bar{g}_n$  down to 0 or, in view of the last lemma, to decrease  $\xi(\beta, y, \delta)$  down to  $\varepsilon$ . Since  $\xi$  is a majoration to the real vector pay-off  $\bar{g}_n$  only if player II uses a strictly mixed and stationary strategy, we divide the stages of the game  $\Gamma_\infty(p)$  into blocs consisting of a large number of stages each. Within the  $m^{\text{th}}$  bloc, player II plays a certain, strictly mixed, stationary strategy  $\hat{y}^m$ . Our approximation of the average vector pay-off in the  $m^{\text{th}}$  bloc will be  $\xi_m = \xi(\beta_m, y^m, \delta)$ , for a certain  $\delta$ . The strategies  $y^m$  will be chosen so as to decrease  $\langle p_0, \xi_m \rangle$  to zero.

For any given  $\delta > 0$  consider the game with vector pay-offs where, at each stage, player I chooses  $\beta \in \mathcal{B}$ , player II chooses  $y \in \Delta_\varepsilon(T)$  and the pay-off is  $\xi(\beta, y, \delta)$ :

PROPOSITION 3.17. For every  $\varepsilon > 0$  there exists  $\bar{\delta} = \bar{\delta}(\varepsilon)$  and  $N = N(\varepsilon)$  such that for any sequence  $\beta_m$  there is a sequence  $y^m = y^m(\xi_1, \dots, \xi_m)$  in  $\Delta_\varepsilon(T)$  such that:

$$n \geq N(\varepsilon) \Rightarrow E(\bar{\xi}_n^k) \leq 3\varepsilon \quad \forall k \in K, \text{ where } \bar{\xi}_n^k = \frac{1}{n} \sum_{m=1}^n \xi_m^k$$

PROOF. We prove that the set  $M_{2\varepsilon} = \{x \in \mathbb{R}^K \mid x^k \leq 2\varepsilon \forall k\}$  is approachable in the auxiliary game for player II using strategies in  $\Delta_\varepsilon(T)$  only. Since  $|G_{st}^k| \leq 1$ , it is enough to prove that the set  $M_\varepsilon$  is approachable using any strategy in  $\Delta(T)$ . Since  $M_\varepsilon$  is convex, using cor. 4.6 p. 104, it is enough to prove that for all  $z \in \mathbb{R}^K$

$$w(z) \geq \inf_{m \in M_\varepsilon} \langle m, z \rangle$$

where  $w(z)$  is the value of the game with pay-off  $\sum_k z^k \xi^k$  in which player II is the maximiser. The above inequality is obviously satisfied if  $z^k > 0$  for some  $k \in K$  or if  $z = 0$ . Otherwise let  $q \in \Delta(K)$  be the normalisation (to a unit vector) of  $-z$ . Since  $-w(-q) = u(q)$  and since for  $z > 0$  the right hand side minimiser in  $M_\varepsilon$  is at  $m = (\varepsilon, \dots, \varepsilon)$ , the condition becomes

$$\max_{\beta} \min_{y \in \Delta_\varepsilon(T)} \langle q, \xi(\beta, y, \delta) \rangle \leq \varepsilon$$

which is true by lemma 3.16 p. 201. ■

We shall now use the above result to provide an optimal strategy for the uninformed player. Given  $\varepsilon > 0$  let  $\bar{\delta}, N$  be determined by prop. 3.17 p. 201. Let  $M$  be determined by cor. 3.14 p. 201 for  $\delta = \bar{\delta}$  and  $\eta = \varepsilon/N$ .

Consider the following strategy  $\tau_\varepsilon$ : Player II plays in blocs of length  $M$  each. In all the stages of the  $m^{\text{th}}$  block he plays i.i.d. the same mixed move  $y^m \in \Delta_\varepsilon(T)$  which is determined as follows:  $y^1$  is arbitrary. At the beginning of the  $m^{\text{th}}$  bloc ( $m \geq 1$ ), player II computes  $\xi_{m-1}$  for the last bloc, and plays  $y^m(\xi_1, \dots, \xi_{m-1})$  — determined also by prop. 3.17 p. 201 — i.i.d. in that bloc.

By our construction, on a segment of  $N$  blocs (with duration  $n_\varepsilon = N \cdot M$ ), the expected average pay-off per stage is  $\leq 4\varepsilon$ , whatever be  $k$ .

Using this for a sequence  $\varepsilon_i$  decreasing to 0 (with the corresponding  $n_{\varepsilon_i}$ ) and repeating the argument in the proof of prop. 4.5 p. 165 we can construct a single strategy  $\tau$  such that for any  $\varepsilon$  there exists  $N(\varepsilon)$  satisfying  $\bar{\gamma}_n^k \leq \varepsilon$  for all  $\sigma, k$  and  $n \geq N(\varepsilon)$ . This is an optimal strategy of player II.

Summing up, the main results so far are the existence of  $\lim v_n(p)$ ,  $\lim v_\lambda(p)$  and  $v_\infty(p)$  which are given explicitly and depend on the signalling matrices of player II only (and not on those of player I). By ex. VEx.3 p. 253 this value, as a function of  $p$ , is semi-algebraic.

**3.e. Approachability.** The section is independent of the rest of the chapter, except for (a particular case of) prop. 3.7 p. 195.

In the next theorem, we use  $B$  as a short notation for  $B \times T$  (when the signals do not inform player II of his move).

**THEOREM 3.18.** *Let  $\varphi$  be a convex function  $\geq u$  on  $\Pi$ . For any sequence  $\varepsilon_n > 0$  converging to zero there is a strategy  $\tau$  of player II in  $\Gamma_\infty$  and a map  $\mathbf{l}: \bigcup_n B^n \rightarrow \mathbb{R}^K$  for which*

- (1)  $\varphi_n(p) = \max_{b \in B^n} \langle \mathbf{l}(b), p \rangle$  is decreasing, with  $\lim_{n \rightarrow \infty} \varphi_n(p) \leq \varphi(p)$ ,  $\varphi_n(p) \leq 2C$  and  $\forall b \in B^n$  there exists  $p$  strictly positive with  $\langle \mathbf{l}(b), p \rangle = \varphi_n(p)$
- (2) Given a play  $\omega$ , define  $\mathbf{l}_n(\omega) = \mathbf{l}^k(b_1, \dots, b_n)$ . Let  $E_n = \bar{g}_n - \mathbf{l}_n$ , and  $N = \sup\{n \mid E_n > 0\}$  ( $\sup \emptyset = 0$ ). Then  $\sup_{\sigma, k} P_{\sigma, \tau}^k(N \geq n) \leq \exp(-n\varepsilon_n) \quad \forall n \geq 0$ .
- (3) Let  $M_n = \sup_{\substack{\sigma, k \\ m \leq n}} \|E_m\|_{L_\infty(\sigma, \tau, k)}$ . Then  $\varepsilon_n(M_n - 3C) \leq 1$

**COMMENT 3.11.** We will be interested in sequences converging slowly to 0 ( $n\varepsilon_n \rightarrow \infty$ , cf. 2) and 3 means that one can select bounds for  $M_n$  that increase arbitrarily slowly to  $+\infty$ , above 3.c p. 195.

COMMENT 3.12. The geometric meaning of 1 is that  $\mathbf{l}(B^n)$  is the efficient frontier of its convex, comprehensive hull, and the sequence of those sets is decreasing. The point  $p = p_b \in \text{int}(\Delta(K))$  is such that  $\langle \mathbf{l}(b), p \rangle \geq \langle \mathbf{l}(b'), p \rangle \forall b' \in \bigcup_{m \geq n} B^m$ . It might possibly be interpreted therefore as a kind of “objective” posterior of player II on the states of nature (i.e. independent of any prior on  $\sigma$ ).

COMMENT 3.13. When the signals do not inform player II of his moves, and  $\mathbf{l}$  is thus a map on  $\bigcup_n (B \times T)^n$ , we have a game of essentially perfect recall, where (1 p. 51) one has to think of  $\tau$  as a mixed strategy of II, i.e. a probability distribution over pure strategies. Every pure strategy  $\tilde{t}$  selected by  $\tau$  allows then to compute player II’s moves in terms of his past signals, and generates thus from  $\mathbf{l}$  the map  $\mathbf{l}_{\tilde{t}}$  from  $\bigcup_n B^n$  to  $\mathbb{R}^K$ . Hence  $\tau$  can be viewed as a probability distribution over pairs  $(\tilde{t}, \mathbf{l}_{\tilde{t}})$ , and  $\mathbf{l}_{\tilde{t}}$  yields indeed an estimated vector pay-off to player II solely on the basis of his signals.

PROOF. The proof is subdivided in three parts.

#### PART A. The building blocs.

PROPOSITION 3.19. Let  $\varphi$  be a continuous convex function,  $\varphi > u$  on  $\Pi$ . Then there exists  $N$ , a strategy  $\tau$  of player II in  $\Gamma_N$ , and a map  $\mathbf{l}: B^N \rightarrow \mathbf{L}_\varphi = \{x \in \mathbb{R}^K \mid \langle p, x \rangle \leq \varphi(p) \forall p \in \Delta(K)\}$  such that for every strategy  $\sigma$  of player I in  $\Gamma_N$  and for all  $k \in K$ ,  $E_{\sigma, \tau}^k(E_N) \leq 0$ .

PROOF.

##### STEP 1. Simplification of the game.

Replace first the game  $\Gamma$  by a game  $\Gamma'$  more favourable to player I, replacing  $S$  by  $K \times S$ : an action in  $\Gamma'$  is the choice of a type and of an action of that type in  $\Gamma$ , with the corresponding pay-offs and signals, except that the pay-off equals  $-M$  whenever player I “lies” about his type. Since  $u'$ , as a function of  $M$ , decreases to  $u$ , we can find  $M$  sufficiently large such that still  $u' < \varphi$  everywhere, by compactness. In  $\Gamma'$ ,  $Q$  is independent of the state of nature  $k$ , and by theorem 1.3 p. 53 we can assume that  $Q$  informs each player of his own pure action in  $\Gamma'$ .

Replace now  $\Gamma'$  by its  $\delta$ -modification  $\Gamma''$ , even more favourable to player I, where II’s choice of  $t$  results in  $t$  with probability  $1 - \delta$ , in which case his only signal is  $(t, \text{blank})$ , and player I’s is selected according to  $Q$ , and results with probability  $\delta$  in a uniform distribution over  $T$  choosing  $t'$ , in which case his signal is  $(t, b)$  and player I’s is  $a$ , where  $(a, b)$  is selected according to  $Q_{s, t'}$  (recall  $b$  includes  $t'$  and  $a$  includes  $s$ ). For  $\delta$  sufficiently small, we will still have  $u'' < \varphi$  everywhere, and now  $Q^{\text{II}}$  is independent of  $t$ .

Consider now  $\Gamma'''$ , with  $u''' = u'' < \varphi$ , even more favourable to I, obtained by informing player I at each stage of player II’s signal. Then every information player I gets in addition to player II’s signal and his own pure action choice is irrelevant — he can simulate it by doing some additional randomisation himself. Discarding it, we obtain a model  $\Gamma^4$  where  $Q_s(b)$  selects a signal  $b$  just as a function of player I’s pure action choice, player II is informed of  $(t, b)$  and player I of  $(s, t, b)$ .

In this game, let  $\mathbf{NR} = \bigcap_p \mathbf{NR}(p) = \mathbf{NR}(p_0)$  for any interior point  $p_0$ . Then for all  $p$ ,  $\mathbf{NR}(p)$  and  $\mathbf{NR}$  have the same projection on  $\{k \mid p_k > 0\}$  — in particular  $u(p)$  can be computed by restricting player I to the compact, convex polyhedron  $\mathbf{NR}$ . Replace now player I’s pure strategy set by the set of extreme points  $e$  of  $\mathbf{NR}$  (this includes all former strategies): by definition of  $\mathbf{NR}$ ,  $Q_e(b)$  is well defined and independent of  $k \in K$ . And

player II's message will be of the form  $(t, b)$  and player I's of the form  $(e, s, t, b,)$  — but repeating our previous argument leading from  $\Gamma'''$  to  $\Gamma^4$  we can assume player I just hears  $(e, t, b)$ : Now, in  $\Gamma^5$ , we have in addition the property that  $\forall x \in \text{NR}, \exists \sigma_x \in \Delta(S)$  such that, for all  $k \in K$ ,  $Q_{\sigma_x}(b) = Q_{x^k}(b) \forall b \in B$  and  $\sigma_x G^k = x^k G^k$ . In particular  $u(p) = \text{Val}\langle p, G \rangle$ .

STEP 2. *Construction of  $(N_0, \tau)$  ( $\tau$  mixed strategy of player II in  $\Gamma_{N_0}$ ) such that  $(\bar{\gamma}_{N_0}^k(\sigma, \tau))_{k \in K} \in \mathbf{L}_{\varphi-\varepsilon}$  for all non-revealing strategies  $\sigma$  of I in  $\Gamma'_{N_0}$ .*

We first argue in the game  $\Gamma^5$ . Since  $\varphi > u$ , one can find  $\varepsilon > 0$  and  $\varphi'$ , with  $u < \varphi' < \varphi - \varepsilon$  and  $\varphi'$  a maximum of finitely many linear functions — i.e. it is the convexification of a function with values  $\varphi_i$  at  $p_i$  ( $i \in I$ ,  $I$  finite), and  $(+\infty)$  elsewhere.

Consider then the game  $\bar{\Gamma}$  where player I is initially informed of  $i \in I$  (chosen with probability  $\lambda_i$ ), next  $G^i = \langle p_i, G \rangle$  is to be played, with the same  $Q$  as above: we have  $\bar{u}(\lambda) = u(\sum \lambda_i p_i) < \varphi'(\sum \lambda_i p_i) \leq \sum \lambda_i \varphi_i$ , for all  $\lambda$  (using  $u(p) = \text{Val}\langle p, G \rangle$ ).

Hence, by prop. 3.7 p. 195 (in the game  $\bar{\Gamma}$ ), there exists  $N_0$  such that  $\text{Val}(\bar{\Gamma}_{N_0}(\lambda)) < \sum \lambda_i \varphi_i$  for all  $\lambda$ . Apply the minmax theorem in the finite game  $\tilde{\Gamma}$  where player I chooses his type  $i$  and a pure strategy in  $\bar{\Gamma}_{N_0}$ , player II chooses a pure strategy in  $\bar{\Gamma}_{N_0}$  and the pay-off equals  $-\varphi_i$  plus the pay-off in  $\bar{\Gamma}_{N_0}^i$ : there exists a strategy  $\tau$  of II in  $\bar{\Gamma}_{N_0}$  — i.e. in  $\Gamma_{N_0}^5$  — that guarantees him zero in this game. This means that, for any  $i \in I$ , and every strategy  $\sigma$  of player I in  $\Gamma_{N_0}^5$  that is independent of  $k$ ,  $E_{\sigma, \tau}^{p_i}(\bar{g}_{N_0}) \leq \varphi_i$ . Since, for fixed  $(\sigma, \tau)$ ,  $E_{\sigma, \tau}^p(\bar{g}_{N_0})$  is linear in  $p$ , this means therefore that  $(\bar{\gamma}_{N_0}^k(\sigma, \tau))_{k \in K} \in \mathbf{L}_{\varphi-\varepsilon}$  for all strategies  $\sigma$  in  $\Gamma_{N_0}^5$  which are independent of  $k$ .

$\tau$  is also a strategy in  $\Gamma''$  (player II's action sets and information are the same), and it is clear what is the corresponding (“generalised” cf. ex. II.1Ex.10ai p. 61) strategy in  $\Gamma'_{N_0}$ , which, by ex. II.1Ex.14 p. 72, can be assumed to be also a mixed strategy in the original game  $\Gamma_{N_0}$ , i.e. depending only on his signals in that game.

Consider a non-revealing strategy  $\sigma$  of player I in  $\Gamma'_{N_0}$ . It is still a strategy, and non-revealing, in  $\Gamma'''_{N_0}$ .

To go to  $\Gamma_{N_0}^4$ , change it by letting player I do his own randomisations, on an auxiliary probability space: we still have a non-revealing (generalised) strategy in  $\Gamma_{N_0}^4$ . Write it as a behavioural strategy  $\sigma_n^k$ . Modify it now in the following way: after stage 1, player I uses the conditional probability on  $s_1$  given  $k$  and  $b_1$  to average his “strategy for the future” w.r.t.  $s_1$ ; rewrite this new strategy for the future as a behavioural strategy, which now depends no longer on  $s_1$ . The joint distribution of  $(k, s_2, t_2, b_2, \dots)$  is the same as before hence the expected pay-off from stage 2 on is not affected — and clearly the expected pay-off in stage 1 is also the same — so the modification does not affect expected pay-offs in the game (here we use crucially the additive separability of pay-offs). For the same reason, the new strategy is still non-revealing. Do now the same with this new “strategy for the future” and  $s_2$ , and so on: we obtain a non-revealing  $\sigma$  of player I such that  $\sigma_n^k$  depends only on the past signals of II, and which yields the same pay-offs as the previous one.

$\sigma$  being non-revealing in  $\Gamma^4$  implies now that every history  $h$  of signals of II that has positive probability under  $\sigma$  for some  $k \in K$  (and some  $\tau$ ) has so for every  $k \in K$  — and that for every such  $h$  of length  $n$ ,  $\sigma_n(h) \in \text{NR}$ . There is no problem in modifying  $\sigma$  for the other histories, thus  $\sigma$  is a map from histories of player II to  $\text{NR}$ . Hence, in  $\Gamma^5$ , we can find an equivalent map from histories of player II to  $\Delta(S)$ : in particular, this is independent of  $k \in K$ .

STEP 3. *End of the proof.*

Write now  $\Gamma'_{N_0}$  as a one-shot game, with a single strategy —  $\tau$  — for player II, and where every pure strategy  $s$  of player I yields a joint distribution of pay-offs and signals — being the joint distribution of the average pay-off up to stage  $N_0$  and the sequence of  $N_0$  pairs of signals it induces against  $\tau$  in  $\Gamma'_{N_0}$ . In this game, by step 2, every non-revealing strategy of player I yields a pay-off in  $\mathbf{L}_{\varphi-\varepsilon}$  — i.e., the new  $u$ -function is still  $\leq \varphi - \varepsilon$ . Further, there is no loss in adding again to player I's signal knowledge of his own pure strategy choice and of player II's signals. Denote this game by  $\tilde{\Gamma}$ , by  $\tilde{B}$  the corresponding alphabet for player II, and  $\tilde{A}$  for I.

Denote by  $D \subseteq \Delta(\tilde{B})$  the (compact) convex hull of the points  $(\tilde{Q}_s^{\text{II}})_{s \in \tilde{S}}$ , the marginal distribution on  $\tilde{B}$  of  $\tilde{Q}_s$ . For every  $\pi \in \Delta(\tilde{B})$ , denote by  $\bar{\pi}$  the closest point in  $D$ , and let  $\hat{\mathbf{l}}^k(\pi) = \max \left\{ \sum_{s \in \tilde{S}} x_s \tilde{g}_s^k \mid x \in \Delta(\tilde{S}), \sum_{s \in \tilde{S}} x_s q_s^{\text{II}}(b) = \bar{\pi}(b) \forall b \in \tilde{B} \right\}$  ( $\tilde{g}_s^k$  is the expected pay-off in  $\tilde{\Gamma}^k$  induced by action  $s \in \tilde{S}$ ).

Observe that, for all  $\pi$ ,  $\hat{\mathbf{l}}(\pi)$  is the pay-off to a non-revealing strategy ( $x^k$  induces for all  $k$  the same distribution of signals  $\bar{\pi}$ ), hence  $\hat{\mathbf{l}}: \Delta(\tilde{B}) \rightarrow \mathbf{L}_{\varphi-\varepsilon}$ . Further  $\hat{\mathbf{l}}$  is clearly Lipschitz, by ex.I.3Ex.4q p. 30.

Fix now  $k \in K$ , and consider  $\tilde{\Gamma}^k$  as a game with vector pay-offs in  $\mathbb{R} \times \Delta(\tilde{A} \times \tilde{B})$  — action  $s$  yielding  $\tilde{g}_s^k$  and the random pair of signals generated by  $\tilde{Q}_s$ . The set  $C = \{(r, \pi) \in \mathbb{R} \times D \mid r \leq \hat{\mathbf{l}}^k(\pi)\}$  (i.e., one neglects the coordinate in  $\tilde{A}$ , which was kept only in order that player I's strategy set remains the same in  $\tilde{\Gamma}$ ) is closed and convex, as the (comprehensive hull of the) linear image of  $\Delta(\tilde{S})$ , and is, by 4.4 p. 103, approachable by player II: i.e., there exist constants  $M_k$  such that for all  $n$ , all  $\sigma$  and all  $k$ ,  $\mathbb{E}_{\sigma^k} d[(\bar{g}_n^k, f_n), C] \leq \frac{M_k}{\sqrt{n}}$  (4.3 p. 102), where  $f_n$  denotes the empirical frequency on  $\tilde{B}$  until stage  $n$ . The Lipschitz character of  $\hat{\mathbf{l}}$  implies therefore there exists  $M > 0$  such that for all  $\sigma, k$  and  $n$ ,  $\mathbb{E}_{\sigma^k} (\bar{g}_n^k - \hat{\mathbf{l}}^k(f_n))^+ \leq M/\sqrt{n}$ . In particular, let  $N_1 \geq (M/\varepsilon)^2$  and define  $\tilde{\mathbf{l}}: \tilde{B}^{N_1} \rightarrow \mathbf{L}_\varphi$  as  $(\hat{\mathbf{l}}^k(f_{N_1}) + \varepsilon)_{k \in K}$ ; then  $\mathbb{E}_{\sigma^k} (\bar{g}_{N_1}^k - \tilde{\mathbf{l}}_{N_1}^k) \leq 0$  for all  $\sigma$  and  $k$ .

Revert now to  $\Gamma'$ , with  $N = N_0 \cdot N_1$ , so  $B^N = \tilde{B}^{N_1}$ : we have a strategy  $\tau$  in  $\Gamma'_N$ , and  $\mathbf{l}: B^N \rightarrow \mathbf{L}_\varphi$  such that, for all  $\sigma$  and  $k$ ,  $\mathbb{E}_{\sigma, \tau}^k (\bar{g}_N - \mathbf{l}_N) \leq 0$ . This remains then true in the original game  $\Gamma$ , which differs from  $\Gamma'$  only by the fact that player I has fewer strategies  $\sigma$  (as mentioned, if in  $\Gamma$  player II's signal does not inform him of his move,  $\mathbf{l}$  becomes a map from  $(B \times T)^N$  to  $\mathbf{L}_\varphi$ ).

This finishes the proof of prop. 3.19 p. 203. ■

## PART B. Construction of the strategy, and points 1 and 3 of the theorem.

Denote, for any  $p_0 \in \Delta(K)$ , by  $\psi_{p_0}$  the convexification of the function having value  $u(p_0)$  at  $p_0$ ,  $C^k = \max_{s,t} G_{s,t}^k$  at the  $k^{\text{th}}$  extreme point, and  $+\infty$  elsewhere. Observe that  $\psi_{p_0} \geq u$  everywhere: player II guarantees it in  $D(p)$  by playing an optimal strategy in  $D(p_0)$  (any strategy if  $\text{NR}(p_0) = \emptyset$ ). It follows that, for any convex function  $\varphi \geq u$ , denoting by  $\tilde{\varphi}$  the function equal to  $C^k$  at the  $k^{\text{th}}$  extreme point and to  $\varphi$  elsewhere, the function  $\bar{\varphi} = \min(\varphi, \text{Vex } \tilde{\varphi})$  is convex, with  $u \leq \bar{\varphi} \leq \varphi$  and has value  $\leq C^k$  at the  $k^{\text{th}}$  extreme point. It suffices therefore to prove the result for functions  $\varphi$  which have value  $\leq C^k$  at the  $k^{\text{th}}$  extreme point. In particular such functions have values in  $[-\infty, C]$  and are thus upper semi-continuous. Therefore, if  $f_\ell$  is a decreasing sequence of continuous functions converging to  $\varphi$ , and  $\psi_\ell = \ell^{-1} + \text{Vex}(f_\ell)$ , the  $\psi_\ell$  are continuous convex functions  $> u$  and decreasing to  $\varphi$ .

Hence we have, by prop. 3.19, for each  $\ell$ , an integer  $N_\ell$ , a strategy  $\tau_\ell$  of player II in  $\Gamma_{N_\ell}$ , and a function  $\hat{\mathbf{l}}_\ell$  from  $B^{N_\ell}$  to  $\mathbf{L}_{\psi_\ell}$  such that, for all strategies  $\sigma$  of player I in  $\Gamma_{N_\ell}$  and all  $k \in K$

$$\mathsf{E}_{\sigma, \tau_\ell}^k(\bar{g}_{N_\ell} - \hat{\mathbf{l}}_\ell) \leq 0$$

(with the slight abuse of notation  $\hat{\mathbf{l}}_\ell(\omega) = \hat{\mathbf{l}}_\ell^{k(\omega)}(b_1(\omega), \dots, b_{N_\ell}(\omega))$ ). There is no loss in assuming  $N_\ell > N_{\ell-1}$ .

We show first that, if necessary by extracting a subsequence, we can further assume that  $\varphi_\ell(p) = \max_{b \in B^{N_\ell}} \langle \hat{\mathbf{l}}_\ell(b), p \rangle$  decreases to  $\varphi$ . For this, observe first (e.g. by Zorn's lemma, but this is not needed: since  $\varphi < +\infty$ , it suffices to minimise on rational points) that  $\varphi \geq \bar{\varphi} \geq u$ , where  $\bar{\varphi}$  is a minimal element of the set of convex functions  $\geq u$ . So we can assume  $\varphi$  itself is minimal. Next, observe that  $\varphi_\ell$  is a convex function  $\geq u$ : our inequality implies  $\mathsf{E}_{\sigma, \tau_\ell}(\bar{g}_{N_\ell}) \leq \mathsf{E}_{\sigma, \tau_\ell}[\mathsf{E}_{\sigma, \tau_\ell}(\hat{\mathbf{l}}_\ell | \mathcal{H}_{N_\ell}^\Pi)] = \mathsf{E}_{\sigma, \tau_\ell}(\sum_k p_{N_\ell}^k \hat{\mathbf{l}}_\ell^k) \leq \mathsf{E}_{\sigma, \tau_\ell} \varphi_\ell(p_{N_\ell})$ , so if  $\sigma$  is an optimal strategy in  $D(p)$  — repeated independently stage after stage —, then  $p_{N_\ell} = p$ , and  $\mathsf{E}_{\sigma, \tau_\ell}(\bar{g}_{N_\ell}) \geq u(p)$ , hence  $u(p) \leq \varphi_\ell(p)$ . Since  $\varphi_\ell \leq \psi_\ell$  and  $\psi_\ell$  converges to  $\varphi$ , we obtain that  $\limsup \varphi_\ell \leq \varphi$ . But  $\limsup \varphi_\ell$  is convex and  $\geq u$ , so by the minimality of  $\varphi$  we have  $\limsup \varphi_\ell = \varphi$ . Since the same argument applies along any subsequence, we obtain  $\lim_{\ell \rightarrow \infty} \varphi_\ell = \varphi$ . This implies that  $\varepsilon_\ell = \max_p [\varphi(p) - \varphi_\ell(p)]$  converges to zero by ex. I.3Ex.15c p.38. Let us thus add  $\eta_\ell = \sup_{i \geq \ell} \varepsilon_i + \ell^{-1}$  to the function  $\hat{\mathbf{l}}_\ell$ ; we obtain now that  $\varphi_\ell > \varphi$ , and the other properties are still valid. Define finally a subsequence  $\ell_i$  inductively by  $\ell_1 = \min\{\ell \mid \psi_\ell + \eta_\ell \leq 2C\}$ ,  $\ell_{i+1} = \min\{\ell \mid \psi_\ell + \eta_\ell < \varphi_{\ell_i}\}$  [this exists, by compactness and by continuity of the  $\psi_\ell$  and  $\varphi_{\ell_i}$ , since  $\varphi_{\ell_i} > \varphi$  and  $\psi_\ell + \eta_\ell$  decreases to  $\varphi$ ]. Now  $\varphi_{\ell_i} > \varphi_{\ell_{i+1}}$ . Assume thus the original sequence satisfies this. So we can henceforth replace the original sequence  $\psi_\ell$  by  $\psi_\ell(p) = \max_{b \in B^{N_\ell}} \langle \hat{\mathbf{l}}_\ell(b), p \rangle$ : this one also decreases to  $\varphi$ . Replace finally  $\hat{\mathbf{l}}_1$  by the constant function  $2C$ . Choose also  $C_\ell = \max(C_{\ell-1}, C + \max_{b \in B^{N_\ell}, k \in K} |\hat{\mathbf{l}}_\ell^k(b)|)$ , ( $C_0 = 0$ ).

Denote then by  $S_\ell^k$  the pure strategy set of player I in  $\Gamma_{N_\ell}^k$ , by  $P_{\ell,s}^k$  the joint distribution (on  $\mathbb{R} \times A^{N_\ell}$ , given  $\tau_\ell$ ,  $k \in K$  and  $s \in S_\ell^k$ ) of  $\bar{g}_{N_\ell}^k - \hat{\mathbf{l}}_\ell^k$  and of the signals of player I (in  $A^{N_\ell}$ ). Denote also by  $\mu_{\ell,s}^k$  the marginal of  $P_{\ell,s}^k$  on  $\mathbb{R}$ . Each  $\mu_{\ell,s}^k$  has support in  $[-C_\ell, C_\ell]$  and barycentre  $e_{\ell,s}^k \leq 0$ .

Define now the strategy  $\tau$  of player II in  $\Gamma_\infty$  as follows: for an appropriate sequence of positive integers  $R_\ell$ , play for  $R_1$  successive blocs of length  $N_1$  the strategy  $\tau_1$ , then for  $R_2$  blocs  $\tau_2$ , etc. Define also, at any stage  $n$ , given a decreasing sequence  $\eta_n$  converging to zero to be specified later,  $\mathbf{l}_n(b_1, \dots, b_n)$  as  $\eta_n$  plus  $n^{-1}$  times the sum over all previous blocs of the corresponding value of  $\hat{\mathbf{l}}$  multiplied by the corresponding bloc-length, plus  $C$  times the number of stages in the last, incomplete bloc.

Let  $n = \sum_{j=1}^{i-1} R_j N_j + m N_i + r$ , with  $0 \leq m < R_i$  and  $0 \leq r < N_i$ . Then  $\langle \mathbf{l}_n, p \rangle - \eta_n \leq \frac{1}{n} \left[ \sum_{j=1}^{i-1} R_j N_j \psi_j(p) + m N_i \psi_i(p) + r C \right] = \varphi_n(p)$ .

Denote by  $f_n(p)$  the same function, but where  $C$  is replaced by  $\psi_i(p)$ . Since  $\psi_\ell(p)$  decreases to  $\varphi(p)$ ,  $f_n(p)$  are a decreasing sequence of convex, Lipschitz functions converging to  $\varphi$ . And  $\varphi_n(p) \leq \frac{r}{n}[C - \psi_i(p)] + f_n(p)$ . Now  $C_i \geq \max_p |C - \psi_i(p)|$ , so let  $\delta_m = \sup_{n \geq m} (\frac{r}{n} C_i)$  (recall  $r$  and the index  $i$  vary with  $n$ ).

For  $\frac{r}{n} C_i$  to converge to zero, it suffices to consider values of  $n$  where  $r = N_i - 1$ , and then where  $m = 0$ : thus one needs  $N_\ell C_\ell / \left( \sum_{i=1}^{\ell-1} N_i R_i + N_\ell \right)$  to converge to zero. For this it suffices that  $R_{\ell-1} \geq N_\ell C_\ell$  (since  $N_{\ell-1} \rightarrow +\infty$ ). Assume the sequence  $R_\ell$  satisfies this

condition. Then  $\delta_n$  converges to zero. Let then  $\bar{\varphi}_n(p) = f_n(p) + \delta_n + \eta_n$ :  $\bar{\varphi}_n$  is a decreasing sequence of convex, Lipschitz functions converging to  $\varphi$ , and for all  $n$ ,  $\mathbf{l}_n: B^n \rightarrow \mathbf{L}_{\bar{\varphi}_n}$ .

We show now how to modify this function  $\mathbf{l}$  such as to satisfy also 1. With the above notation, let  $n_0 = \sum_{j=1}^{i-1} R_j N_j + m N_i$ . Observe  $\varphi_{n_0}(p) = f_{n_0}(p) = \max_{b \in B^{n_0}} \langle \mathbf{l}_{n_0}(b), p \rangle - \eta_{n_0}$ : since over each bloc there exists such a maximising string of signals, it suffices to put those together. Define now  $\bar{\mathbf{l}}_n$  as  $\mathbf{l}_n$ , but replacing the term  $C$  in the last incomplete bloc by a repetition of the last  $\hat{\mathbf{l}}_i$  estimated ( $\hat{\mathbf{l}}_{i-1}$  if  $m = 0$ ): the same argument shows that  $g_n(p) = \max_{b \in B^n} \langle \bar{\mathbf{l}}_n(b), p \rangle - \eta_n = \frac{1}{n} \left[ \sum_{j=1}^{i-1} R_j N_j \psi_j(p) + (m N_i + r) \psi_i(p) \right]$  (here again,  $\psi_{i-1}(p)$  if  $m = 0$ ) — and as before the functions  $g_n$  decrease to  $\varphi$ . So by adding  $\delta_n$  to  $\bar{\mathbf{l}}_n$  and to  $g_n$ , those properties are preserved and now  $\bar{\mathbf{l}} \geq \mathbf{l}$ , so all our estimates in part C for  $\mathbf{l}$  will a fortiori apply to  $\bar{\mathbf{l}}$ .

Finally, to have all properties of 1, increase still, for each  $b \in B^n$ , the coordinates of  $\bar{\mathbf{l}}_n(b)$  such as to obtain  $\langle \bar{\mathbf{l}}_n(b), p \rangle = g_n(p) + \eta_n$  for some interior  $p$ . Finally, for the inequality  $\varphi_n(p) \leq 2C$ , replace  $\bar{\mathbf{l}}$  by the identically zero function if  $C = 0$ , otherwise, since  $\max_{b \in B^n} \langle \bar{\mathbf{l}}_n(b), p \rangle$  decreases to  $\varphi(p)$ , and since by minimality  $\varphi(p) \leq C$  for all  $p$  as seen above, there exists  $n_0$  with  $\bar{\mathbf{l}}_{n_0}(b) < 2C$  for all  $b \in B^{n_0}$ . Just set  $\bar{\mathbf{l}}_n = 2C$  for  $n < n_0$  — this preserves all other properties. This yields the function of the theorem — for part C it will however be sufficient to deal with the original function  $\mathbf{l}$ .

For point 3, observe that now for all  $n$ ,  $E_n \geq -3C$ , and  $M_1 \leq 3C$ . Choose each  $R_\ell$  sufficiently large such that  $\varepsilon_{R_\ell N_\ell}(C_{\ell+1} - 3C) \leq 1$ . Then 3 follows immediately for the original function  $\mathbf{l}$ . Since the true function  $\bar{\mathbf{l}}$  is larger or equal to  $\mathbf{l}$ , 3 follows a fortiori for  $\bar{\mathbf{l}}$  (the true  $E_n$  being anyway  $\geq -3C$ ).

We also proved along the way:

**COROLLARY 3.20.** *Every convex function  $\varphi \geq u$  is minorated by a minimal such function, which has value  $\leq \max_{s,t} G_{s,t}^k$  at the  $k^{\text{th}}$  extreme point.*

### PART C. Point 2 of the Theorem.

**STEP 1.** Given  $F(x)$  convex increasing with  $\frac{F(x)}{x} \rightarrow \infty$ ,  $\frac{F(x)}{x^2} \rightarrow 0$ , one can select  $R_\ell$  such as to have for  $X$  standard normal,  $\mathbb{E}_{\sigma,\tau}^k f(E_n) \leq \mathbb{E} f\left(\frac{\sqrt{F(n)}}{n} X - \eta_n\right)$  for all  $\sigma, k, n$  and for all increasing, convex  $f$ .

Let  $W_i = \sum_{j=1}^i N_{\ell_j} X_{s_j}$ ,  $T_i = \sum_{j=1}^i N_{\ell_j}$ , where  $s_j \in S_{\ell_j}^k$  denotes the pure strategy choice of player I in bloc  $j$  of length  $N_{\ell_j}$ , and  $X_{s_j}$  is selected according to  $\mu_{\ell_j, s_j}^k$ . Then, during bloc  $i+1$ , i.e. for  $T_i \leq n < T_{i+1}$ , we have  $\bar{g}_n^k - \mathbf{l}^k(b_1, \dots, b_n) \leq \frac{W_i}{n} - \eta_{T_i}$ . We first try to replace  $W_i$  by random variables  $\tilde{W}_i$  such that  $\mathbb{E} f(W_i) \leq \mathbb{E} f(\tilde{W}_i)$  for any convex increasing  $f$ .

Thus, by monotonicity of  $f$ , we can replace each  $X_{s_j}$  by  $X_{s_j} - e_{\ell_j, s_j}^k$ , since  $e_{\ell_j, s_j}^k \leq 0$ . I.e., we are reduced to the case where each  $\mu_{\ell,s}^k$  has barycentre zero, and support included in  $[-2C_\ell, 2C_\ell]$ . Consider now random variables  $Y'_j$ , whose conditional distribution given all other variables in the problem — including the other  $Y$ 's — is carried by  $\{-2C_\ell, 2C_\ell\}$  and has expectation  $X_{s_j}$ : by the convexity of  $f$ , and Jensen's inequality, we can replace the  $X_{s_j}$  by  $Y'_j$ . I.e., we can assume that each  $\mu_{\ell,s}^k$  not only has expectation zero, but also is carried by  $\{-2C_\ell, 2C_\ell\}$ . Thus it assigns probability  $\frac{1}{2}$  to each of those points. In particular, all strategies of player I, and all  $k \in K$  reduce now to the same problem; if we let  $Y'_j = 2C_\ell Y_j$ , then the  $Y_j$  are an i.i.d. sequence, uniform on  $\{-1, 1\}$ . Let now  $V_j$  be an i.i.d.

sequence of standard normal random variables (i.e., with mean zero and variance one), independent of all others. Since  $E(|V_j|) = \sqrt{\frac{2}{\pi}}$ , the  $V'_j = Y_j | V_j | \sqrt{\frac{\pi}{2}}$  have  $Y_j$  as conditional expectation given all other variables — so we can replace the  $Y_j$  by those, and are i.i.d. normal with standard deviation  $\sqrt{\frac{\pi}{2}}$ . To sum up, we can define  $\tilde{W}_i$  by replacing  $X_{s_j}$  by  $\sqrt{2\pi}C_{\ell_j}X_j$ , where the  $X_j$  are an i.i.d. sequence of standard normal variables.

Let  $D_\ell = \sqrt{2\pi}C_\ell$ . The variance of  $\tilde{W}_i$  equals  $V(T_i) = \sum_{j=1}^i (D_{\ell_j}N_{\ell_j})^2$ . During super-bloc  $\ell$ , the function  $V$  has therefore slope  $N_\ell D_\ell^2$ , which increases to  $+\infty$ .

Fix now a convex, increasing function  $F$  from  $\mathbb{R}_+$  to itself such that  $\lim_{x \rightarrow \infty} F(x)/x = +\infty$  and  $\lim_{x \rightarrow \infty} F(x)/x^2 = 0$ , and define the successive  $R_\ell$  as follows: continue super-bloc  $\ell$  as long as needed such that from the endpoint of the graph of  $(V(T), T)$  reached at the end of super-bloc  $\ell$ , the straight line with slope  $N_{\ell+1}D_{\ell+1}^2$  will lie everywhere below the graph of  $F(T)$ . (And increase  $R_\ell$  if necessary still some more to satisfy also our previous condition  $R_\ell \geq N_{\ell+1}C_{\ell+1}$ ). This defines now fully the strategy  $\tau$ , given the function  $F$ . And by construction, we have now  $V(T) \leq F(T)$  for all  $T \geq R_1N_1$ .

Hence, for all  $i \geq R_1$ , we can still add (Jensen again) to  $\tilde{W}_i$  an independent normal variable with mean zero and variance  $F(T_i) - V(T_i)$ , so the  $\tilde{W}_i$  become normal with mean zero and variance  $F(T_i)$ .

Finally, since for  $T_i \leq n < T_{i+1}$  we have  $E_n \leq \frac{W_i}{n} - \eta_{T_i}$ , since  $\eta_n \leq \eta_{T_i}$ , and since  $F(n) \geq F(T_i)$  we can conclude, by a last use of Jensen's inequality, that for all  $n \geq R_1N_1$ , and all convex increasing functions  $f$ ,  $E[f(E_n)] \leq E[f(\frac{W_n}{n} - \eta_n)]$ , where  $W_n$  is normal  $(0, F(n))$ , i.e.  $E f(E_n) \leq E f\left(\frac{\sqrt{F(n)}}{n}X - \eta_n\right)$  where  $X$  is standard normal: this finishes step 1.

STEP 2. Two lemmas.

LEMMA 3.21. Consider a couple of random variables  $(X, Y)$ , where  $Y$  is standard normal. Assume  $E f(X) \leq E f(Y)$  for all convex increasing functions  $f$  ( $f$  Lipschitz and bounded below). Then  $\Pr(X \geq \lambda) < \text{Erf}(\lambda - \frac{1}{\lambda})$  for all  $\lambda > 0$  (with  $\text{Erf}(\mu) = \Pr(Y \geq \mu)$ ).

PROOF. Let  $f(x) = (x - r)^+$ :  $f(\lambda) \Pr(f(X) \geq f(\lambda)) \leq f(X)$ . So for  $r < \lambda$ :

$$\Pr(X \geq \lambda) \leq \frac{1}{\lambda - r} \int_r^\infty (y - r) \frac{1}{\sqrt{2\pi}} e^{-y^2/2} dy = \frac{1}{\lambda - r} \left[ \frac{1}{\sqrt{2\pi}} e^{-r^2/2} - r \text{Erf}(r) \right]$$

The derivative of the right-hand member w.r.t.  $r$  equals

$$\frac{1}{(\lambda - r)^2} \left[ \frac{1}{\sqrt{2\pi}} e^{-r^2/2} - \lambda \text{Erf}(r) \right].$$

So choose  $r$  such that  $\lambda = e^{-r^2/2} / (\sqrt{2\pi} \text{Erf}(r)) = F(r)$ .

[There always exists a unique such  $r$ , because the right-hand member is strictly increasing from 0 ( $r = -\infty$ ) to  $+\infty$  ( $r = +\infty$ ). This follows from the inequality  $\text{Erf}(r) < e^{-r^2/2}/(r\sqrt{2\pi})$  for  $r > 0$ , which insures both that the derivative of the right-hand member is positive for  $r > 0$  — and it is obviously so for  $r \leq 0$  — and that the right-hand member tends to  $+\infty$  with  $r$ . The zero limit at  $r = -\infty$  is obvious. The inequality  $r \text{Erf}(r) < \frac{1}{\sqrt{2\pi}} e^{-r^2/2}$  follows in turn by expressing  $E(Y - r)^+ > 0$ , and implies also that the chosen  $r$  is  $< \lambda$ .]

Substituting thus  $\lambda \text{Erf}(r)$  for  $\frac{1}{\sqrt{2\pi}} e^{-r^2/2}$  in our majoration yields

$$P(X \geq \lambda) \leq \text{Erf}(r).$$

There only remains to show that  $r > \lambda - 1/\lambda$ , i.e., by the previously checked strict monotonicity of  $F$ , that  $\lambda > F(\lambda - 1/\lambda)$  — or, letting  $x = \lambda - 1/\lambda$  and hence ( $\lambda > 0$ )  $\lambda = \frac{x+\sqrt{x^2+4}}{2}$ , that  $\frac{x+\sqrt{x^2+4}}{2} > F(x)$ , or  $\sqrt{2\pi}\text{Erf}(x) - \frac{2e^{-x^2/2}}{x+\sqrt{x^2+4}}$  is positive. Since the limit at  $+\infty$  is clearly zero, it suffices to check the function is decreasing. Taking the derivative, this amounts to  $x\sqrt{x^2+4} < x^2 + 2$ , which is obvious. ■

LEMMA 3.22. For every sequence  $\varepsilon_n$  converging to zero there exists a sequence  $\delta_n$  converging to zero such that  $\forall n \geq 1$

$$\sum_{m=n}^{\infty} \text{Erf} \sqrt{m\delta_m} \leq \exp(-n\varepsilon_n).$$

PROOF. By the bound  $\text{Erf}(x) \leq \frac{1}{\sqrt{2\pi}x} e^{-x^2/2}$  (e.g. lemma 3.21), it suffices to get  $n\delta_n \geq 1/(2\pi)$  and  $\sum_{m=n}^{\infty} e^{-m\delta_m/2} \leq \exp(-f(n))$ , with  $f(n) \geq n\varepsilon_n$ . For the latter inequality, it suffices to have

$$e^{-n\delta_n/2} \leq e^{-f(n)} - e^{-f(n+1)} \quad (\text{so } f(n) < f(n+1)), \text{ i.e.}$$

$$\frac{\delta_n}{2} \geq \frac{-1}{n} \ln \left( e^{-f(n)} - e^{-f(n+1)} \right) = \frac{f(n+1)}{n} - \frac{1}{n} \ln \left( e^{f(n+1)-f(n)} - 1 \right)$$

Hence for such  $\delta_n \rightarrow 0$  to exist it suffices that  $\frac{f(n)}{n} \rightarrow 0$  (so also  $\frac{f(n+1)}{n} \rightarrow 0$ ) and  $\liminf_{n \rightarrow \infty} \frac{1}{n} \ln(f(n+1) - f(n)) \geq 0$  (since  $e^x - 1 \geq x$ ). So it suffices that  $f(n+1) - f(n) \geq 1/[n(n+1)]$ . Hence  $f(n) = 1 - n^{-1} + \max_{i \leq n} i\varepsilon_i$  will do. ■

STEP 3. End of the proof.

From step 1, let  $f\left(\frac{\sqrt{F(n)}}{n}X - \eta_n\right) = g(X)$ , so  $f(E_n) = g\left[\frac{n}{\sqrt{F(n)}}(E_n + \eta_n)\right]$ , and  $\mathbb{E}g\left(\frac{n}{\sqrt{F(n)}}(E_n + \eta_n)\right) \leq \mathbb{E}g(X)$  for all  $n$  and all convex increasing  $g$ . Hence, by lemma 3.21,

$$\Pr\left[\frac{n}{\sqrt{F(n)}}(E_n + \eta_n) \geq \lambda\right] \leq \text{Erf}\left(\lambda - \frac{1}{\lambda}\right) \quad \text{for all } \lambda > 0,$$

i.e., for  $\lambda' = \frac{\sqrt{F(n)}}{n}\lambda - \eta_n$ ,

$$\Pr(E_n \geq \lambda') \leq \text{Erf}\left[\frac{n}{\sqrt{F(n)}}(\lambda' + \eta_n) - \frac{\sqrt{F(n)}}{n(\lambda' + \eta_n)}\right] \quad \text{for all } \lambda' \geq 0.$$

For a sequence  $\delta_n$  decreasing to zero, choose now  $\eta_n$  and  $F$  such that

$$\frac{n}{\sqrt{F(n)}}(\lambda + \eta_n) - \frac{\sqrt{F(n)}}{n(\lambda + \eta_n)} \geq \sqrt{n\delta_n}(\lambda + 1) \quad \text{for all } \lambda \geq 0, \text{ i.e. for } \lambda = 0.$$

E.g., take  $\eta_n = (\frac{2}{n} + \delta_n)^{\frac{1}{3}}$ ,  $F$  the convexification of  $n/\eta_n$ . Those satisfy all our requirements, so, for all  $n \geq R_1 N_1$ , we have  $\Pr(E_n \geq \lambda) \leq \text{Erf}[\sqrt{n\delta_n}(\lambda + 1)]$  for all  $\lambda \geq 0$ . To obtain this also for the other values of  $n$ , just increase  $\eta_n$  for those such as to have  $I_n^k(b) > C \forall b \in B^n, \forall n \leq R_1 N_1, \forall k \in K$  — then  $\Pr(E_n \geq 0) = 0$ .

Now,  $\Pr(N \geq n) \leq \sum_{m=n}^{\infty} \Pr(E_n > 0) \leq \sum_{m=n}^{\infty} \text{Erf} \sqrt{m\delta_m}$  — so point 2 of the theorem follows from lemma 3.22. ■

This finishes the proof of theorem 3.18. ■

### 3.f. The errors $E_n^+$ in the approachability theorem.

LEMMA 3.23. Let  $\Psi = \{ \psi: \mathbb{N} \times \mathbb{R}_+ \rightarrow \mathbb{R}_+ \mid \forall n, \psi(n, x) \text{ is non-decreasing in } x \text{ and } \forall x, \lim_{n \rightarrow +\infty} \frac{\psi(n, x)}{n} = 0 \}$ . And let  $\Psi_0 = \{ \psi \in \Psi \mid \psi(0, x) = 0 \}$ .

- (1) Every sequence  $\psi_i \in \Psi$  is majorated by an additively separable one:

$$\psi_i(n, x) \leq K_i + g(x) + f(n)$$

(i.e.,  $g: \mathbb{R}_+ \rightarrow \mathbb{R}_+$  is non-decreasing and  $\lim_{n \rightarrow \infty} \frac{f(n)}{n} = 0$ ).  
[Hence by adding e.g.  $\sqrt{x}$  to  $g$  and  $\sqrt{n}$  to  $f$  we will have  $\psi_i(n, x) \leq f(n) + g(x)$  except in a bounded region.]

- (2)  $\forall \psi \in \Psi_0, \exists f, h: f(nh(x)) \in \Psi_0 \text{ and } \psi(n, x) \leq f(nh(x))$ .

PROOF. 1) Let  $j$  enumerate the set of pairs  $(i, m)$ , with  $f_j(n) = \psi_i(n, m)$ , and let  $f(n) = \max\{f_j(n) \mid \forall k \leq j, \forall \ell \geq n, f_k(\ell) \leq \ell/j\}$ : then  $f_j(n) \leq f(n)$  for  $n$  large enough, so there exists  $\varphi(i, m)$  with  $\psi_i(n, m) \leq \varphi(i, m) + f(n)$ . Let  $g_0(m) = \max_{i \leq m} \varphi(i, m)$ :  $\varphi(i, m) \leq g_0(m), \forall m \geq i$ , so  $\varphi(i, m) \leq g_0(m) + K_i$ , hence with  $g(x) = \max_{m < x+1} g_0(m)$  we obtain 1).

2) By 1), we assume  $\psi(n, x) = f(n) + g(x)$  for  $n \geq 1$ . Replace  $f(n)$  by  $\max_{i \leq n} f(i) + \ln n$ , so  $f(kn) \geq f(n) + \ln k$ , and extend  $f$  by linear interpolation to  $\mathbb{R}_+$ , with  $f(0) = 0$ . Now  $h(x) = \sup_{n \geq 1} \frac{1}{n} f^{-1}(f(n) + g(x))$  is finite ( $\leq 1 + \exp g(x)$ ) and monotone, and  $f(n) + g(x) \leq f(nh(x))$  for all  $x$  and all  $n \geq 1$ . This finishes the proof since  $\psi(0, x) = 0$ . ■

REMARK 3.14. We will be basically interested in  $\Psi$  or  $\Psi_0$ , which is a convex cone and lattice, etc. The lemma gives us convenient co-final sets to work with.

REMARK 3.15. non-decreasing functions  $g$  from  $\mathbb{R}_+$  to itself can always be majorated by very “nice” ones — e.g. that have an everywhere convergent power-expansion with all coefficients positive. Similarly  $f$  with  $\lim_{x \rightarrow \infty} \frac{f(x)}{x} = 0$  can be replaced by its concavification, and a number of further conditions can be imposed.

### COROLLARY 3.24. Main Corollary

- (1) For every  $\psi \in \Psi_0$ ,  $(\tau, \mathbf{l})$  can be chosen in theorem 3.18 p. 202 such that  $\forall k, \sigma, \lambda \geq 0$ ,

$$P_{\sigma, \tau}^k [\psi(N, M_N) \geq \lambda] \leq e^{-\lambda}.$$

- (2) Equivalently (lemma 3.23 p. 210), for every non-decreasing functions  $h$  and  $f$  from  $\mathbb{R}_+$  to itself such that  $\lim_{x \rightarrow \infty} \frac{f(x)}{x} = 0$  there exists  $(\tau, \mathbf{l})$  such that  $\forall k, \sigma, \lambda \geq 0$

$$P_{\sigma, \tau}^k (Nh(M_N) > \lambda) \leq \exp -f(\lambda).$$

PROOF. Assume by lemma 3.23  $\psi(n, x) \leq f(n) + g(x)$ . By point 3 of the theorem, we can select  $M_n$  to increase as slowly to  $+\infty$  as desired — in particular such that  $g(M_n) \leq Af(n)$ ; then  $K(n) = \psi(n, M_n)$  still satisfies  $\frac{K(n)}{n} \rightarrow 0$ . We have then to show that we can find  $(\tau, \mathbf{l})$  such that  $P(N \geq n)$  ( $= P(K(N) \geq K(n)) \leq \exp -K(n) \forall n \geq 1$ ). This is point 2 of the theorem. ■

COROLLARY 3.25. (1)  $\forall \psi \in \Psi_0$  there exists  $(\tau, \mathbf{l})$  such that in addition the Laplace transform  $\mathcal{L}_{\tau, \mathbf{l}, \psi}^k(\alpha) = \sup_{\sigma} E_{\sigma, \tau}^k \exp[\alpha \psi(N, M_N)]$  is finite.

- (2) Given a sequence  $\psi_i \in \Psi_0$ , choose  $\psi_0$  such that  $\psi_i \leq \psi_0 + K_i$  (cf. lemma 3.23 p. 210). Let  $\psi$  in  $\Psi_0$  such that  $\psi(n, x) \geq \psi_0(n, x) + \ell(n)$  with  $\ell(n) \rightarrow \infty$ . For any

$(\tau, \mathbf{l})$  as in theorem 3.18 p. 202, which further satisfies, for each  $k$  where possible, that  $\forall \sigma, P_{\sigma,\tau}^k(N=0) = 1, \exists L_i$  such that

$$\mathcal{L}_{\tau,\mathbf{l},\psi_i}^k(\alpha) \leq L_i \mathcal{L}_{\tau,\mathbf{l},\psi}^k(\alpha) \text{ for all } \alpha, i \text{ and } k$$

(and the same inequality holds with  $L_i = 1$  for all  $\alpha$  sufficiently large).

PROOF. 1) Apply lemma 3.23.1 p. 210 to the sequence  $i\psi \in \Psi_0$ , so  $i\psi \leq \psi' + K_i$  with  $\psi' \in \Psi$ ; apply then the main corollary to  $\psi'$  to conclude that  $\sup_{\sigma,k} \mathbf{E} e^{\frac{1}{2}\psi'(N,M_N)} < +\infty$ , hence  $\mathcal{L}^k(\alpha) < \infty$  since for  $i \geq 2\alpha$ ,  $\alpha\psi \leq \frac{1}{2}\psi' + \frac{1}{2}K_i$ .

2) We first show that the full statement follows from the parenthesis. It suffices to show that the inequality holds for  $\alpha \leq 0$  — since it also holds for  $\alpha$  sufficiently large, increasing  $L_i$  will make it hold everywhere. For  $\alpha \leq 0$ , we have  $\mathcal{L}_{\tau,\mathbf{l},\psi_i}^k(\alpha) \leq 1$ , and  $\mathcal{L}_{\tau,\mathbf{l},\psi}^k(\alpha) \geq \inf_\sigma P_{\sigma,\tau}^k(N=0) > 0$  (since  $\varepsilon_1 > 0$  in the theorem). Hence the claim.

For the parenthesis, there is something to prove only if  $\mathcal{L}_{\tau,\mathbf{l},\psi}^k(\alpha) < \infty$  everywhere (in particular,  $P_{\sigma,\tau}^k(N < \infty) = 1$  for all  $\sigma$ ). Fix thus such  $(k, \tau, \mathbf{l})$ . Observe that  $\mathcal{L}_{\tau,\mathbf{l},\psi_i}^k(\alpha) \leq e^{K_i\alpha} \mathcal{L}_{\tau,\mathbf{l},\psi_0}^k(\alpha)$  (for  $\alpha \geq 0$ ), so it suffices to show that  $\frac{1}{\alpha} [\ln \mathcal{L}_{\tau,\mathbf{l},\psi}^k(\alpha) - \ln \mathcal{L}_{\tau,\mathbf{l},\psi_0}^k(\alpha)] \rightarrow +\infty$ .  $(\tau, \mathbf{l})$  determines the sequence  $M_n$ ; let thus  $G(n) = \exp \psi_0(n, M_n)$ ,  $H(n) = \exp \psi(n, M_n)$ : we have  $G(n) \geq 1$ ,  $\frac{H(n)}{G(n)} \geq e_n = \exp(\ell_n) \rightarrow +\infty$ , and we want to show that  $\frac{\sup_\sigma \|H(N)\|_\alpha}{\sup_\sigma \|G(N)\|_\alpha}$  goes to  $+\infty$  with  $\alpha$ .

Observe first we can assume there exists some  $\sigma_1$  under which the distribution  $p_1$  of  $N$  has infinite support. Otherwise  $\exists n_0 : P_{\sigma,\tau}^k(N \leq n_0) = 1 \forall \sigma$  (if  $P_{\sigma,n,\tau}^k(N \geq n) > 0 \forall n$ , then  $\sigma = \sum 2^{-n} \sigma_n$  would give infinite support to  $N$ ), so that setting  $\mathbf{l}_n = 2C$  for  $n \leq n_0$  would show that one can have  $P_{\sigma,\tau}^k(N=0) = 1 \forall \sigma$  — which implies  $\mathcal{L}_{\tau,\mathbf{l},\psi}^k$  is identically one for any  $\psi \in \Psi_0$ , so the conclusion is obvious.

It suffices to prove the convergence along any subsequence  $\alpha_i$  with  $\alpha_i \geq 2^i$  ( $i \geq 2$ ). Let thus  $p_i$  be the distribution of  $N$  induced (under  $k, \tau, \mathbf{l}$ ) by some  $\sigma$  that (approximately) maximises  $\|G(N)\|_{\alpha_i}$ , and let  $p = \sum_{i \geq 1} 2^{-i} p_i$ .

$p$  is induced by some strategy (convexity). Also  $p \geq 2^{-i} p_i$  yields  $\|G\|_{L_{\alpha_i}(p)} \geq (2^{-i})^{1/\alpha_i} \|G\|_{L_{\alpha_i}(p_i)}$ ; since  $(2^{-i})^{1/\alpha_i} \rightarrow 1$  it follows that it suffices to show that  $\frac{\|H\|_\alpha}{\|G\|_\alpha} \rightarrow +\infty$  for the probability distribution  $p$  with infinite support on  $\mathbb{N}$ .

If  $\|G\|_\alpha$  is bounded, the proof is finished, because  $H(n) \geq G(n) \cdot e_n \geq e_n$  converges to  $+\infty$ , so  $\|H\|_\alpha$  converges to  $+\infty$  because  $p$  has infinite support (recall  $\|X\|_\alpha$  is monotone and converges to  $\|X\|_\infty$ ), and  $\|G\|_\alpha \geq 1$ . Otherwise (same monotonicity)  $\|G\|_\alpha$  converges to  $+\infty$ , which means that, letting for  $\lambda > 0$ ,  $N_\lambda = \{n \mid G(n) \geq \lambda\}$  and  $e_\lambda = \min\{e_n \mid n \in N_\lambda\}$ , we have  $p(N_\lambda) > 0$  and  $e_\lambda \rightarrow +\infty$ . Then

$$\frac{\|H\|_\alpha}{\|G\|_\alpha} \geq \frac{e_\lambda \|G \mathbf{1}_{N_\lambda}\|_\alpha}{\|G \mathbf{1}_{N_\lambda}\|_\alpha [1 + \frac{1}{p(N_\lambda)}]^{1/\alpha}}.$$

Since the bracket converges to 1 as  $\alpha \rightarrow +\infty$ , we obtain  $\liminf_{\alpha \rightarrow \infty} \frac{\|H\|_\alpha}{\|G\|_\alpha} \geq e_\lambda$ , and since  $e_\lambda \rightarrow \infty$  this finishes the proof. ■

COMMENT 3.16.  $f(\alpha) = \exp(\alpha x)$  is log-convex (has convex logarithm). Check that an average of two log-convex functions still has the same property (reduce first to the case where the two are exponential functions, next to the verification that  $\ln(1 + e^x)$  is convex). Going to the limit conclude that a Laplace transform  $\mathbf{E} \exp(\alpha X)$  is log-convex; finally a supremum of convex functions still being so,  $\mathcal{L}_{\tau,\alpha,\psi}^k(\alpha)$  is log-convex.

COMMENT 3.17. Point 1 of the corollary has the same force as the main corollary: fix such a  $(\tau, \mathbf{l})$ , and fix also  $k$ . Let  $\varphi(n) = \psi(n, M_n)$ ,  $F_\sigma(\alpha) = E_\sigma \exp(\alpha \varphi(N))$ ,  $F(\alpha) = \sup_\sigma F_\sigma(\alpha) < \infty$ . Then, for  $\alpha > 0$ ,  $e^{\alpha\lambda} P_\sigma(\varphi(N) \geq \lambda) \leq F_\sigma(\alpha)$ , so  $\sup_\sigma P_\sigma(\psi(N, M_N) \geq \lambda) \leq \exp[\ln F(\alpha) - \alpha\lambda]$  for all  $\alpha \geq 0$ . By the convexity of the bracket (comm. 3.16), the minimum over  $\alpha$  of the right-hand member is easily obtained as being  $\exp[-\int_{\lambda_0}^\lambda G(s)ds]$  for  $\lambda \geq \lambda_0$ , where the function  $G(\lambda)$  gives the root of the derivative of the bracket w.r.t.  $\alpha$  (make a picture!).  $G$  could have horizontal and vertical segments, and could take the value  $+\infty$ , but is monotone and converges to  $+\infty$ .  $\lambda_0$  is (any) root of  $G(\lambda_0) = 0$ , i.e.  $\lambda_0 = E_\sigma \varphi(N)$  for some strategy  $\sigma$ . So  $H(\lambda) = \int_{\lambda_0}^\lambda G(s)ds$  is convex (with minimum  $H(\lambda_0) = 0$ ) and satisfies  $H(\lambda)/\lambda \rightarrow +\infty$ , in particular  $H(\lambda) \geq \lambda$  for  $\lambda \geq \lambda_1$  — hence  $\sup_\sigma P_\sigma(\psi(N, M_N) \geq \lambda) \leq e^{-\lambda}$  for  $\lambda \geq \lambda_1$ . This bounds correctly the probability for all  $N \geq n_1$ , and it is trivial to remember that one can always increase  $\mathbf{l}$  to  $2C$  for  $1 \leq n \leq n_1$  such as to give zero probability to those values of  $N$ .

COMMENT 3.18. The main corollary — or (cf. remark 3.17) the above — imply our Gaussian bound (on  $E_n$ ) proved during the proof of the theorem. Indeed, given  $\varepsilon_n$ , let  $h(x) = (x+1)^4 \mathbb{1}_{x>0}$ , and  $f(x) = \frac{1}{2} + \frac{1}{2}(\sqrt{x}\delta(x)+1)^2$ , where  $\delta(x) \rightarrow 0$  is such that  $x\delta(x^4)$  is non-decreasing and  $\delta(n) \geq \sqrt{\varepsilon_n}$  (thus  $\delta(x) = x^{-1/4} \max_{n \leq x} n^{1/4} \sqrt{\varepsilon_n}$ ). Then,  $\forall M_n \geq 0$ ,  $n^{1/4}(M_n + 1)\delta(n(M_n + 1)^4) \geq n^{1/4}\sqrt{\varepsilon_n}$ , so  $f(nh(M_n)) \geq \frac{1}{2} + \frac{1}{2}[\sqrt{n\varepsilon_n}(M_n + 1) + 1]^2$ , hence  $\exp(-f(nh(M_n))) \leq \text{Erf}(\sqrt{n\varepsilon_n}(M_n + 1))$  [using — cf. proof of lemma 3.21 —  $\text{Erf}(x) \geq \frac{2}{\sqrt{2\pi}} \frac{e^{-x^2/2}}{x+\sqrt{x^2+4}} \geq \frac{1}{\sqrt{2\pi}} \frac{e^{-x^2/2}}{x+1} \geq \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2}-x} \geq \exp[-(\frac{1}{2}x^2+x+1)]$ .  $\delta(x) \rightarrow 0$  yields  $\frac{f(x)}{x} \rightarrow 0$ , so let, by the main corollary,  $(\tau, \mathbf{l})$  be such that  $\Pr(Nh(M_N) \geq x) \leq \exp - f(x)$  for all  $x > 0$  and all  $\sigma, k$ . Then, for  $\lambda \geq M_n$ , we have  $\Pr(E_n > \lambda) = 0 \leq \text{Erf}(\sqrt{n\varepsilon_n}(\lambda+1))$ , and for  $0 \leq \lambda \leq M_n$  we have  $\Pr(E_n > \lambda) \leq \Pr(N \geq n) = \Pr(Nh(M_N) \geq nh(M_n)) \leq \exp - f(nh(M_n)) \leq \text{Erf}(\sqrt{n\varepsilon_n}(M_n + 1)) \leq \text{Erf}(\sqrt{n\varepsilon_n}(\lambda+1))$ .

[It is therefore clear such a Gaussian bound applies not only to  $E_n$ , but also to any function  $h(E_n)$ .]

COMMENT 3.19. In typical applications,  $\psi(n, M_n)$  will take the form  $f(nh(M_n))$ , where the random variable  $Nh(M_N)$  bounds random variables like  $\sum_n h(E_n)$  ( $h(x) = 0$  for  $x \leq 0$ , recall  $E_n$  is the error at stage  $n$ ), or like  $\sup_n nE_n$  — the maximum non-normalised error (for  $h(x) = x^+$ ), or  $N$  itself...

COMMENT 3.20. Even for repeated coin tosses — i.e., since we do not have random pay-offs in the model, a single state of nature, “matching pennies”, where both players use their i.i.d. optimal strategy —, the above results cannot be improved, and would not be true for  $f(nh(x))$  with  $f$  the identity — thus for  $\psi(N, M_N) = N$  itself.

Indeed, denote by  $\bar{g}_n$  the frequency of successes, and  $T_\eta = \#\{n \mid \bar{g}_n > \eta\}$  for  $1/2 < \eta < 1$  (i.e. we do not even ask that  $\eta_n \rightarrow \frac{1}{2}$ , and  $T_\eta$  is much smaller than  $N$ ). Even in this most favourable case it is not true that  $T_\eta$  has a finite Laplace transform; the region where it is finite shrinks to zero as  $\eta \rightarrow \frac{1}{2}$  — so one cannot relax the requirement that  $\frac{f(x)}{x} \rightarrow 0$ . To see this, apply first th. 2 p. 399 in (Feller, 1966), next lemma 1 p. 398 (Feller, 1966) to the random walk  $n(\bar{g}_n - \eta)$ , to conclude letting  $n \rightarrow \infty$  that  $\Pr(T_\eta = k) = q_\infty p_k$ , hence  $E s^{T_\eta} = q_\infty p(s)$  will be finite if  $p(s)$  is so, i.e. by formulae 7.13 and 7.2 (ibidem), if  $\tau(s) < 1$  with  $\tau(s) = \sum_n s^n P(\bar{g}_1 \leq \eta, \bar{g}_2 \leq \eta, \dots, \bar{g}_{n-1} \leq \eta, \bar{g}_n > \eta)$ . The probability in the right-hand member is  $\geq P(\bar{g}_k \leq \eta \ \forall k)$ .  $P(\bar{g}_n > \eta \mid \bar{g}_{n-1} \leq \eta) \geq \frac{1}{2} P(T_\eta = 0) P(\bar{g}_{n-1}$  takes its largest possible value  $\leq \eta$ ).

Evaluating this last probability by Stirling's formula we get  $1 \geq [2\sqrt{2\pi\eta(1-\eta)}]^{-1} \times P(T_\eta = 0) \sum_n \frac{1}{\sqrt{n}} \left(\frac{s}{2\eta^\eta(1-\eta)^{1-\eta}}\right)^n$ , so this convergence requires  $s \leq 2\eta^\eta(1-\eta)^{1-\eta}$  — which converges indeed to 1 as  $\eta \rightarrow \frac{1}{2}$ . The result becomes even stronger: let  $S_\eta = \min\{n \geq 1 \mid \bar{g}_n \leq \eta\}$ ; the same conclusions apply even to  $S_\eta$ . Indeed, using formula 7.15, p. 397 (Feller, 1966), we obtain  $\Pr(T_\eta = k) = q_\infty \Pr(S_\eta > k)$ , hence  $q_\infty E \exp(\alpha S_\eta) = \Pr(T_\eta = 0) + (e^\alpha - 1) E \exp(\alpha T_\eta) = +\infty$  for  $\alpha \geq \varphi(\eta) = \ln(2\eta^\eta(1-\eta)^{1-\eta})$ .

Let us show that, if  $\eta_n \rightarrow \frac{1}{2}$ , the above implies that, even in our "typical applications" of the previous remark, the Laplace transform will be infinite for all  $\alpha > 0$  if we let  $f(x) = x$ . There is no loss in assuming  $\eta_n$  to be decreasing. Consider first the random variable  $N$  — or just  $T$ , the total number of errors. Given  $\alpha > 0$ , let  $n_0 = \min\{n \mid \varphi(\eta_n) \leq \alpha\}$ .

To prove that  $E \exp(\alpha T) = \infty$ , it suffices to show that  $E[\exp \alpha T \mid \bar{g}_{n_0} = 1] = \infty$ , since the condition has positive probability. But the conditional distribution of  $T$ , given  $\bar{g}_{n_0} = 1$ , is clearly (stochastically) larger than  $S_{\eta_{n_0}}$  — hence the result. Consider now the random variable  $X = \sum_n h(E_n)$ , with  $E_n = (\bar{g}_n - \eta_n)^+$  and  $h: [0, \frac{1}{2}] \rightarrow \mathbb{R}_+$  non-decreasing and not identically zero. Let  $x_0 < \frac{1}{2}$  be such that  $h(x_0) > 0$ , and choose  $n_0$  large such that  $\eta_{n_0} < \frac{1}{2} - x_0$ . Then  $\sum_n h(E_n) \geq h(x_0) \#\{n \geq n_0 \mid \bar{g}_n \geq x_0 + \eta_{n_0} + \frac{1}{2}\}$ . Hence, by the same conditional bound as above,  $E \exp(\alpha X) = +\infty$  for  $\alpha \geq \varphi(x_0 + \eta_{n_0} + \frac{1}{2})$  — i.e., letting  $n_0 \rightarrow \infty$ , for  $\alpha > \varphi(x_0 + \frac{1}{2})$ . Thus here also the Laplace transform of  $X$  is never finite for  $\alpha \geq \ln(2)$ , and is infinite for all  $\alpha > 0$  if  $x > 0$  implies  $h(x) > 0$ .

Consider finally  $Y = \sup_n n(\bar{g}_n - \eta_n)^+$ . For  $1 > \eta > \eta_{n_0}$ , we have, conditionally to  $\bar{g}_{n_0} = 1$ ,  $Y \geq (\eta - \eta_{n_0})S_\eta$ , hence  $E \exp(\alpha Y) = \infty$  if  $E \exp[\alpha(\eta - \eta_{n_0})S_\eta] = \infty$ , hence if  $\alpha(\eta - \eta_{n_0}) \geq f(\eta)$ . I.e., for  $\eta_{n_0} = \frac{1}{2} + \frac{1}{4}\varepsilon < \frac{3}{4}$  and  $\eta = 2\eta_{n_0} - \frac{1}{2}$ , if  $\alpha \geq \frac{2}{\varepsilon}[(1+\varepsilon)\ln(1+\varepsilon) + (1-\varepsilon)\ln(1-\varepsilon)]$  — hence for all  $\alpha > 0$  by letting  $n_0 \rightarrow \infty$  and thus  $\varepsilon \rightarrow 0$ .

**COMMENT 3.21.** The main corollary is essentially equivalent to the theorem: it obviously implies point 2 (with  $h(x) = \mathbb{1}_{x>0}$ ), but also (with any  $h$  converging to  $+\infty$ ) point 3 — as soon as in at least one state of nature some randomness is present like in "matching pennies".

Indeed, it can be written  $P(N \geq n) \leq \exp -f(nh(M_n))$  — but if the "matching pennies" aspect is present, we obtain a lower bound as sub comm. 3.20: for  $n \geq n_0$ ,  $P(N \geq n) \geq \Pr(\bar{g}_n \text{ is the first value exceeding } \eta_{n_0}) \sim \frac{K(\eta_{n_0})}{\sqrt{n}} \exp(-n\delta(\eta_{n_0}))$ , where  $\eta_n$  decreases to  $\frac{1}{2}$  — so  $\delta(\eta_n) = \ln[2\eta_n^{\eta_n}(1-\eta_n)^{1-\eta_n}]$  decreases to zero. Hence for  $n \geq n'_0$  we have  $f(nh(M_n)) \leq 2n\delta(\eta_{n_0})$ , so  $\limsup \frac{1}{n} f(nh(M_n)) \leq 2\delta(\eta_{n_0})$ , hence  $\frac{1}{n} f(nh(M_n)) \rightarrow 0$ . In particular, for  $n \geq n_0$ ,  $h(M_n) \leq \frac{1}{n} f^{-1}(n)$ . Since, given an arbitrary sequence  $\rho_n$  converging to  $+\infty$ ,  $h(\rho_n)$  converges to  $+\infty$ , and having minorated  $h$  by a strictly (for  $x \geq x_0$ ) monotone one converging to  $+\infty$ , we can then choose  $f$  such that  $\frac{1}{n} f^{-1}(n) \leq h(\rho_n)$ , so conclude that  $M_n \leq \rho_n$  for  $n$  sufficiently large. To obtain the inequality also for  $n \leq n_0$  it suffices to set  $\mathbf{l}_n = 2C$  for  $n \leq n_0$ .

**COMMENT 3.22.** The same proof as in cor. 3.25 p. 210 yields even an improved result: let  $\bar{\psi}(n, x) = 2[\psi(n, x) + \ell_n]$ . The proof shows that there exists a single strategy  $\sigma$  such that  $E_{\sigma, \tau} \exp[2^i \bar{\psi}(N, M_N)] \geq 2^{-i} \mathcal{L}_{\tau, \mathbf{l}, \bar{\psi}}^k(2^i) \geq 2^{-i} \mathcal{L}_{\tau, \mathbf{l}, \psi+\ell}^k(\alpha)$  for  $2^i \leq \alpha \leq 2^{i+1}$ , so  $E_{\sigma, \tau} \exp[\alpha \bar{\psi}(N, M_N)] \geq \frac{1}{\alpha} \mathcal{L}_{\tau, \mathbf{l}, \psi+\ell}^k(\alpha)$  ( $\alpha \geq 4$ ). But it also shows that  $\frac{1}{\alpha} [\ln \mathcal{L}_{\tau, \mathbf{l}, \psi+\ell}^k(\alpha) - \ln \mathcal{L}_{\tau, \mathbf{l}, \psi}^k(\alpha)] \rightarrow +\infty$ , and since  $\frac{1}{\alpha} \ln \alpha \rightarrow 0$  it follows that, for  $\alpha \geq \alpha_0$ ,  $\frac{1}{\alpha} \mathcal{L}_{\tau, \mathbf{l}, \psi+\ell}^k(\alpha) \geq \mathcal{L}_{\tau, \mathbf{l}, \psi}^k(\alpha)$ . Using again the fact that  $E_{\sigma, \tau} \exp[\alpha \bar{\psi}(N, M_N)]$  is bounded away from zero, we obtain that  $\mathcal{L}_{\tau, \mathbf{l}, \psi}^k(\alpha) \leq L E_{\sigma, \tau} \exp[\alpha \bar{\psi}(N, M_N)]$ . And therefore, for all  $i$ ,

$\mathcal{L}_{\tau,1,\psi_i}^k(\alpha) \leq L_i E_{\sigma,\tau}^k \exp[\alpha \bar{\psi}(N, M_N)]$ : now the right-hand member is a true Laplace transform, for a fixed distribution induced by a single strategy  $\sigma$ .

COMMENT 3.23. The fact that the expectation of a random variable  $X$  under  $E_{\sigma,\tau}^k$  is uniformly bounded over  $\sigma$  does not a priori imply that the dominated convergence theorem for variables  $\leq X$  holds uniformly in  $\sigma$ . Still it is this type of result that we need, e.g. to conclude that  $\sup_\sigma \|\sup_{m \geq n} m E_m^+\|_{L_p(\sigma)}$  (cf. remark 3.18 p. 212) converges to zero when  $n \rightarrow \infty$ , since all our concepts in this book are based on errors converging to zero uniformly over  $\sigma$ .

However, in the present case this implication holds: If  $X_i(e_1, e_2, \dots)$  converges pointwise to zero, then for all  $n$ , since there are only finitely many histories of length  $n$ , there are only finitely many possible sequences  $(e_1, e_2, \dots, e_n, 0, 0, 0, \dots)$ , hence  $X_i$  converges to zero uniformly on them, so that, for  $i$  sufficiently large,  $\{X_i \geq \varepsilon\}$  is included in  $\{N \geq n\}$ . And our results (main corollary, or remark 3.17 p. 212) imply that  $\sup_\sigma \int_{\{N \geq n\}} X$  converges to zero (say for  $X \in \Psi$ , and  $(\tau, \mathbf{l})$  as in the main corollary).

Further, for a given point-wise convergent sequence  $X_i$ , our results yield an easy criterion for the existence of such an  $X$  and  $(\tau, \mathbf{l})$  (we assume  $X_i$  are just functions of the errors  $e_1, \dots, e_n, \dots$ ): let  $Y = \sup_i X_i$  ( $Y$  is finite since  $X_i$  is convergent), and  $\varphi(n, m) = \max\{Y(e_1, \dots, e_n, 0, 0, 0, \dots) \mid e_i \leq m \forall i\}$ : we need  $\frac{1}{n} \ln \varphi(n, m) \rightarrow 0$  for all  $m$ .

Now we turn to another estimate, relative to the norm-summability of the errors in  $L_p$  (cf. remark 3.26 below).

LEMMA 3.26. (1) Let  $f_i: \mathbb{R}_+ \rightarrow \mathbb{R}$  be such that  $f_i(x)/x$  is bounded and converges to zero at  $\infty$ . Then there exists  $f: \mathbb{R}_+ \rightarrow \mathbb{R}$  concave, Lipschitz, with  $f(0) = 0 = \lim_{x \rightarrow \infty} f(x)/x$  such that, for some function  $r: \mathbb{R}_+ \rightarrow \mathbb{R}$  satisfying  $r(x) \geq 1$  and  $\lim_{x \rightarrow \infty} r(x) = +\infty$  one has  $f(\prod_{i=1}^n x_i) \geq r(\sum_i x_i) \prod_{i=1}^n f_i(x_i)$ .  
(2) Let also  $\ell: \mathbb{R}_+ \rightarrow \mathbb{R}$  be locally bounded, with  $\lim_{x \rightarrow \infty} \ell(x)/x = 0$  and  $\lim_{x \rightarrow 0} \ell(x) = -\infty$ . Fix  $\underline{r} \geq 1$  and assume  $f_1(x) \geq 0$ .

Then the function  $f$  can moreover satisfy  $f(x_1 x_2) \geq f_1(x_1)[\tilde{r}(x_1 + x_2) + \ell(x_2)]$  with  $\lim_{x \rightarrow +\infty} \tilde{r}(x) = +\infty$  and  $\tilde{r}(x) \geq \underline{r}$ .

PROOF. 1) Let  $g = \text{Cav}(\max_i |f_i|)$ :  $g$  is Lipschitz, with  $g(0) = 0 = \lim_{x \rightarrow \infty} g(x)/x$ , and we can assume  $f_i = g$ . Let  $\bar{r}(x) = \sqrt{x/g(x)}$ ,  $\bar{g}(x) = \bar{r}(x)g(x)$ :  $\bar{g}$  is like  $g$ , and  $\bar{r}$  is monotone, strictly positive and converges to  $+\infty$ . Let  $f_0(u) = \sup_{\prod x_i = u} \prod_{i=1}^n \bar{g}(x_i)$ : if  $L$  is the Lipschitz constant of  $\bar{g}$ , we have  $\frac{1}{u} f_0(u) = \sup_{\prod x_i = u} \prod_{i=1}^n \frac{\bar{g}(x_i)}{x_i} \leq L^n$ , and if  $u_k = \prod_{i=1}^n x_k^i \rightarrow \infty$  then at least one coordinate, say  $x_k^1$ , converges to  $\infty$  (along a subsequence), so  $\frac{1}{u_k} f_0(u_k) \leq \frac{\bar{g}(x_k^1)}{x_k^1} L^{n-1}$  converges to zero. So  $\frac{1}{u} f_0(u)$  is bounded and converges to zero; and  $f_0(\prod x_i) \geq \prod_{i=1}^n \bar{r}(x_i) \prod_{i=1}^n g(x_i)$ . Thus  $f = \left(\frac{1}{\bar{r}(0)}\right)^n \text{Cav } f_0$  satisfies our requirements, with  $r(s) = \left(\frac{1}{\bar{r}(0)}\right)^n \min_{\sum x_i = s} \prod \bar{r}(x_i)$ .

2) Replace  $\ell$  by  $\text{Cav}(\ell)$ , and add  $\sqrt{x}$  to be sure that  $\ell(x) \rightarrow +\infty$ . Let  $f_2(\cdot) = \text{Cav}([\underline{r} + \ell(\cdot)]^+)$ ;  $f(x_1 x_2) \geq \bar{r}(x_1 + x_2) f_1(x_1) f_2(x_2)$  with  $\bar{r}(x) \geq 1$  converging to  $\infty$ . Let  $\tilde{r}(u) = \min_{x \leq u} (\bar{r}(u)f_2(x) - \ell(x))$ : it suffices to show that  $\tilde{r}(x) \geq \underline{r}$  and converges to  $+\infty$  as  $x$  goes to  $+\infty$ . The first point is immediate, so there remains to show that  $\bar{r}(u_i)f_2(x_i) - \ell(x_i) \leq K$  for  $0 \leq x_i \leq u_i$  implies  $u_i$  is bounded. Observe first that  $x_i \geq \underline{x} > 0$  (with  $\ell(\underline{x}) = -K$ ). If  $x_i \geq \ell^{-1}(0)$  then  $K \geq \bar{r}(u_i)f_2(x_i) - \ell(x_i) \geq [\bar{r}(u_i) - 1]\ell(x_i) + \bar{r}(u_i)$  implies  $\bar{r}(u_i) \leq K$ . And if  $-K \leq \ell(x_i) \leq 0$  then  $\bar{r}(u_i) \leq K/f_2(\underline{x})$ . ■

COROLLARY 3.27. Let  $h: \mathbb{R}_+ \rightarrow \mathbb{R}_+$  be monotone, with  $h(0) = 0$ . Assume  $p_n/n \rightarrow 0$ ,  $p_n \geq 0$ . Let also, by the above lemma (2),  $f: \mathbb{R}_+ \rightarrow \mathbb{R}_+$  be concave, Lipschitz, with  $f(0) = 0 = \lim_{x \rightarrow \infty} f(x)/x$ , and such that  $f(nx) > p_n[\ln(n) + r_n]$  for some sequence  $r_n \geq \underline{r}$  converging to  $+\infty$ . For any such  $f$ , and for  $(\tau, \mathbf{l})$  satisfying the main corollary with  $(f, h)$ , one has  $\sup_{\sigma, k} \|h(E_n)\|_{p_n} \leq e^{-r_n}$  — hence the sequence converges to zero and its maximum can be made arbitrarily small. In particular,  $\sum_n \sup_{\sigma, k} \|h(E_n)\|_{p_n}^{\ln(n)} < +\infty$ .

PROOF. For such  $(\tau, \mathbf{l})$  and all  $\sigma, k$ , we have

$$\begin{aligned} \|h(E_n)\|_{p_n} &\leq h(M_n)[\Pr(E_n > 0)]^{1/p_n} \leq h(M_n)[\Pr(N \geq n)]^{1/p_n} \\ &\leq h(M_n) \exp[-f(nh(M_n))/p_n] \end{aligned}$$

[the last inequality because  $f$  is increasing (if  $p_n$  is identically zero, the conclusion is obvious), and noting that, if  $h(M_n) = 0$ , there is nothing to prove]. So  $\|h(E_n)\|_{p_n} \leq e^{-r_n}$ . The “in particular” clause is obvious. ■

COMMENT 3.24. A direct proof that some  $(\tau, \mathbf{l})$  guarantees  $\sup_{\sigma, k} \|h(E_n)\|_{p_n} \rightarrow 0$  is also immediate (i.e., without the lemma) from the theorem. The only additional point here is that it follows already from the main corollary alone, and just by the choice of an appropriate  $f$ .

COMMENT 3.25. Since one can always assume  $h(x) \geq 1$  for  $x > 0$ , the apparently more general statement with  $[\mathbb{E}(h(E_n))^{q_n}]^{1/p_n}$  where  $q_n/n \rightarrow 0$  and  $p_n/n \rightarrow 0$  is equivalent to the present one: indeed, if  $h(x) \neq 0 \Rightarrow h(x) \geq 1$ , the terms only increase if  $q_n$  is increased, so we can assume  $q_n \geq p_n$ ; next if a sequence converges to zero (or is summable), all but finitely many terms are increased — so  $q_n = p_n$ .

COMMENT 3.26. In particular, for all  $p > 0$ ,  $q > 0$ ,  $(\sup_{\sigma, k} \|h(E_n)\|_p)_{n=1}^\infty \in \ell_q$ . (Take  $q_n = p$ ,  $p_n = \frac{p}{q} \ln(n)$  in comm. 3.25 — thus it follows from the summability in cor. 3.27 with  $\frac{p_n}{\ln(n)} \rightarrow +\infty$ .)

COMMENT 3.27. Even for matching pennies (cf. remark 3.20), with  $h: [0, \frac{1}{2}] \rightarrow \mathbb{R}_+$  not identically zero, one needs  $p_n/n \rightarrow 0$ : indeed, for fixed  $\delta$  ( $\frac{1}{2} < \delta < 1$ ), the probability in  $n$  fair coin tosses of obtaining the smallest possible frequency  $\geq \delta$  is, by Stirling's formula (cf. remark 3.20 above)  $\frac{K(\delta)}{\sqrt{n}} e^{-n\varphi(\delta)}$ , with  $K(\delta) = [2\pi\delta(1-\delta)]^{-1/2} > 0$  and  $\varphi(\delta) = \ln[2\delta^\delta(1-\delta)^{1-\delta}] > 0$ . Hence if  $h(x_0) > 0$ ,  $x_0 < \frac{1}{2}$  we have, for  $\eta_{n_0} < \frac{1}{2} - x_0$  and letting  $\delta \frac{1}{2} + x_0 + \eta_{n_0}$  that, for all  $n \geq n_0$ ,  $\mathbb{E}[h(E_n)]^{p_n} \geq [h(x_0)]^{p_n} \frac{K(\delta)}{\sqrt{n}} e^{-n\varphi(\delta)}$ , so for  $p_{n_i} \geq \varepsilon n_i$ ,  $\|h(E_n)\|_{p_{n_i}} \geq h(x_0) \left( \frac{K(\delta)}{\sqrt{n_i}} \right)^{1/\varepsilon n_i} \exp \frac{-\varphi(\delta)}{\varepsilon}$  for all  $i$  sufficiently large — and the limit  $h(x_0) \exp \frac{-\varphi(\delta)}{\varepsilon}$  of the lower bound is  $> 0$ .

COMMENT 3.28. For  $h$  such that  $h(x) \neq 0 \Rightarrow h(x) \geq 1$  one reobtains point 2 of the theorem: say  $\underline{r} = 1$ , then  $\|h(E_n)\|_{p_n} \leq e^{-1}$  implies  $\Pr(E_n > 0) \leq e^{-p_n}$ , so lemma 3.22 p. 209 yields the conclusion.

COROLLARY 3.28. One can further require in the theorem that, for all  $n$ ,  $\sigma$  and  $k$

$$\mathbb{E}_{\sigma, \tau}^k(E_n) \leq 0.$$

PROOF. Since (cor. 3.27),  $\mathbb{E}_{\sigma, \tau}^k(E_n^+) \leq \delta_n$  where  $\delta_n$  decreases to zero, it suffices to add  $\delta_n$  to the function  $\mathbf{l}_n$  (and again replace  $\mathbf{l}_n$  by  $2C$  if  $\mathbf{l}_n^k(b) > 2C$  for some  $b \in B^n$  and  $k \in K$ ). ■

COMMENT 3.29. In the present setup, bounding the error  $E_n^+$  is the only game in town: the other side of the coin, the speed of convergence of  $\varphi_n$  to  $\varphi$ , seems much more difficult to access with the present methods of proof — it would probably require a direct proof, like in 4 p. 102 —, and is even hard to formulate in the present framework, since when  $\varphi$  is not continuous the convergence is not uniform — and  $\varphi$  may even take at some points the value  $-\infty$ .

Anyway, one can study this other aspect only under some constraints on the errors  $E_n^+$ , so the present subsection investigating what can be achieved in this respect seems a necessary first step. Another clear prerequisite, even for the case of linear functions  $\varphi$ , is to have some improvements on the results of sect. 5.b p. 248, and to know exactly what speed of convergence of  $v_n(p)$  to  $v_\infty(p)$  can be achieved.

COMMENT 3.30. But when the function  $u$  is continuous — as happens for instance when the distribution of player II's signals is independent of the state of nature, one sees at least how to formulate precisely this other aspect.

Observe first that each minimal  $\varphi$  is then continuous, so (e.g. by Dini's theorem) the convergence of  $\varphi_n(p)$  to  $\varphi(p)$  is uniform:  $\varphi_n(p) \leq \varphi(p) + \delta_n$ , where  $\delta_n$  is a positive sequence decreasing to zero — the corresponding speed of convergence.

In this framework, it may be more natural to lower the function  $\mathbf{I}$  to  $\hat{\mathbf{I}}$ , a map from  $\bigcup_n B^n$  to (the efficient frontier of)  $\mathbf{L}_\varphi = \{x \in \mathbb{R}^K \mid \langle p, x \rangle \leq \varphi(p) \forall p \in \Delta(K)\}$ , and let  $\mathbf{l}_n = \hat{\mathbf{l}}_n + \delta_n$ . The problem becomes then the optimal tradeoff between the speed of convergence  $\delta_n$  and the size of the estimation errors  $E_n^+$ . [From prop. 5.1 p. 247 below it is clear that in general we will not be able to get  $\delta_n$  better than  $n^{-1/3}$ ; in a “matching pennies”-like case this yields  $\Pr(E_n > 0) \sim \exp(-2n^{1/3})$ , hence sharp bounds (of the same order) can still be obtained (by summing) for  $\Pr(N \geq n)$ ].

Further, observe this error term  $\delta_n$  can be selected independent of  $\varphi$ : the minimal  $\varphi$  are all Lipschitz with constant  $C$ , and  $\leq C$  in absolute value. Hence select, by compactness, for each  $\varepsilon = 3^{-\ell}$ , a finite subset  $\varphi_1, \dots, \varphi_{k_\varepsilon}$  such that every minimal  $\varphi$  has thus one of those  $\varphi_i$ 's with  $\varphi + 3^{-\ell} \leq \varphi_i \leq \varphi + 3.3^{-\ell}$ , hence if such one is selected for each  $\ell$ , they will indeed form a sequence decreasing to  $\varphi$ . And since for each  $\ell$ , we have a finite number of them, we can take, in the proof of the theorem, for  $N_\ell$  the maximum of the corresponding quantities (and  $C_\ell = 2C$ ), so all quantities in the proof of the theorem, and hence the corresponding final estimates, will be independent of the particular  $\varphi$ : all  $\varphi$  will be approximated with the same speed of convergence  $\delta_n$  and the same bounds on errors  $E_n^+$ .

[Even if one wants to make sure to obtain the same monotonicity (point 1 in the theorem), just replace “3” by “4” above, and use the additional playroom to make sure, by increasing  $N_\ell$  if necessary (apply the present theorem to approximating  $\varphi_i$ ) that  $\max_{b \in B^{N_\ell}} \langle \mathbf{l}(b), p \rangle$  is between  $\varphi_i$  and  $\varphi_i - 4^{-\ell}$ ].

COMMENT 3.31. Another case, that could be investigated separately, and that would lead to weakening the above bounds, is that where the pay-offs are random (and selected by the same lottery as the one selecting the signals) and have finite first (or second) moments. In such a case, one could presumably use a similar method as in the present theorem: use the present theorem for the game where the random pay-offs have been replaced by their conditional expectation given the state of nature, the pair of moves, and the pair of signals (the present theorem applies even with random pay-offs as long as they have bounded support). Then imagine, as in the present proof, that one is only interested

in convex functions of the errors, and make a dilatation of each of the conditional distributions of pay-offs (given state, actions and signals) to some common distribution (still with zero mean): then the differences between the actual pay-offs and the fictitious pay-offs used in the application of the theorem are an i.i.d. sequence with zero expectation. Hence, this reduces the problem to that of such a sequence  $X_n$ : finding bounds on the errors  $\bar{X}_n - \eta_n$ , for an appropriate  $\eta_n$ , using then a similar technique as in Lemma 3.21 p. 208 to get rid of the restriction to convex functions by increasing  $\eta_n$ , so that one gets bounds for the differences between the original actual pay-offs and the fictitious pay-off, and finally adding  $\eta_n$  to  $\mathbf{l}_n$ .

### 3.g. Implications of the approachability theorem.

COROLLARY 3.29. Given  $(\tau, \mathbf{l})$ , define for every strategy  $\sigma$  of player I,  $N_\sigma (= N_{\sigma, \tau, \mathbf{l}})$  on II-histories by:

$$N_\sigma(\zeta) = \sup\{ n \mid \exists k: E_{\sigma, \tau}^k(\bar{g}_n \mid \mathcal{H}_n^{\text{II}})(\zeta) \geq \mathbf{l}^k(b_1, \dots, b_n) \}.$$

(Recall  $\mathcal{H}_n^{\text{II}}$  is the  $\sigma$ -field spanned by  $\prod_1^n(B \times T)$ ).

- (1) For  $n > N_\sigma$ ,  $E_{\sigma, \tau}(\bar{g}_n \mid \mathcal{H}_n^{\text{II}}) \leq \varphi_n(p_n)$ .
- (2) In theorem 3.18 p. 202 (and in its corollaries), one can further require of  $(\tau, \mathbf{l})$  that

$$\sup_{\sigma, k} P_{\sigma, \tau}^k(N_\sigma \geq n) \leq \exp(-n\varepsilon_n) \quad \forall n \geq 0.$$

- (3) As a consequence, one has then, for every weak $^*(\sigma(L_\infty, L_1))$ -limit point  $\bar{g}_\infty$  of  $\bar{g}_n$ , and every  $\sigma$ :

$$E_{\sigma, \tau}(\bar{g}_\infty \mid \mathcal{H}_\infty^{\text{II}}) \leq \limsup_{n \rightarrow \infty} E_{\sigma, \tau}(\bar{g}_n \mid \mathcal{H}_n^{\text{II}}) \leq \varphi(p_\infty).$$

PROOF. 1) follows by averaging over  $k$  (point 1 of Theorem 3.18 p. 202), recalling that  $p_n(k) = P(k \mid \mathcal{H}_n^{\text{II}})$ , and that  $\{ E^k(\bar{g}_n^k \mid \mathcal{H}_n^{\text{II}})(\zeta) < \mathbf{l}_n^k(b_1, \dots, b_n) \forall k \in K \}$  is  $\mathcal{H}_n^{\text{II}}$ -measurable and contains  $\{N_\sigma < n\}$ .

2) From  $P_{\sigma, \tau}^k(E_n > 0) \leq e^{-f(n)}$  and  $E_n \leq M_n$  we obtain

$$P_{\sigma, \tau}^k[E(E_n^+ \mid \mathcal{H}_n^{\text{II}}) \geq \lambda_n] \leq \frac{1}{\lambda_n} M_n e^{-f(n)} \leq e^{-n\delta_n}$$

where  $\delta_n \geq \frac{f(n)}{n} + \frac{1}{n} \ln \lambda_n - \frac{1}{n} \ln M_n$  is an arbitrary sequence converging to zero provided one chooses (point 3 of the theorem)  $\ln M_n \leq n\delta_n$ , (point 2)  $f(n) = 3n\delta_n$  and  $\ln \lambda_n \geq -n\delta_n$ ,  $\lambda_n$  decreasing to zero. (All this is completely independent of  $\sigma$  and  $k$ ). Hence by adding now  $\lambda_n$  to  $\mathbf{l}_n$  one obtains

$$\Pr_{\sigma, \tau}^k(E(\bar{g}_n \mid \mathcal{H}_n^{\text{II}}) \geq \mathbf{l}_n^k) \leq e^{-n\delta_n},$$

so  $\Pr_{\sigma, \tau}^k(N_\sigma \geq n) \leq \sum_n^\infty e^{-\ell\delta_n}$ : use of the proof of lemma 3.22 p. 209 finishes the proof.

3) 1 yields that  $E(\bar{g}_n \mid \mathcal{H}_n^{\text{II}}) \leq \varphi_{n_0}(p_n)$  for  $n \geq n_0 \vee N_\sigma$  (point 1 of theorem) so (martingale convergence theorem for  $p_n$  and continuity of  $\varphi_{n_0}$ ):  $X = \limsup_{n \rightarrow \infty} E(\bar{g}_n \mid \mathcal{H}_n^{\text{II}}) \leq \varphi_{n_0}(p_\infty)$ , hence the second inequality when  $n_0 \rightarrow +\infty$ . And, by Fatou's lemma

$$E(X \mid \mathcal{H}_{n_0}^{\text{II}}) \geq \limsup_{n \rightarrow \infty} E(\bar{g}_n \mid \mathcal{H}_{n_0}^{\text{II}}) \geq E(\bar{g}_\infty \mid \mathcal{H}_{n_0}^{\text{II}})$$

(finiteness of  $\mathcal{H}_{n_0}^{\text{II}}$ ), so letting  $n_0 \rightarrow \infty$  yields the first inequality. ■

There is obviously also a converse to theorem 3.18 p. 202:

DEFINITION 3.30. For any strategy  $\tau$  of player II, and any Banach limit  $\mathcal{L}$  (I.2Ex.13 p. 24), let  $h_{\tau,\mathcal{L},\sigma}(p) = \mathcal{L}(\bar{\gamma}_n(p, \sigma, \tau))$ , and  $\varphi_{\tau,\mathcal{L}}(p) = \sup_{\sigma \in [\text{NR}(p)]^\infty} h_{\tau,\mathcal{L},\sigma}(p)$  (where  $\sigma \in [\text{NR}(p)]^\infty$  means  $\sigma$  is the i.i.d. repetition of some element of  $\text{NR}(p)$ ).

PROPOSITION 3.31. (1)  $h_{\tau,\mathcal{L},\sigma}$  is linear, and  $\varphi_{\tau,\mathcal{L}}$  is convex and  $\geq u$ .

(2) If  $\tau$  corresponds by theorem 3.18 to some convex  $\varphi \geq u$ ,  $\varphi_{\tau,\mathcal{L}} \leq \varphi$  for every  $\mathcal{L}$ ; more precisely:  $\limsup_{n \rightarrow \infty} \sup_{\sigma \in \text{NR}^n(p)} \bar{\gamma}_n(p, \sigma, \tau) \leq \varphi(p)$  (where  $\text{NR}^n(p)$  is the set of non-revealing strategies of player I in  $\Gamma_n(p)$ ).

PROOF. 1) The linearity of  $h$  is obvious. For each face  $F$  of  $\Delta(K)$ , let  $\sum_F = [\text{NR}(p)]^\infty$  for  $p$  interior in  $F$ , and for  $\sigma \in \sum_F$ , let  $\underline{h}_{\tau,\mathcal{L},\sigma,F}(p) = h_{\tau,\mathcal{L},\sigma}(p)$  for  $p \in F$ ,  $\underline{h}_{\tau,\mathcal{L},\sigma,F}(p) = -\infty$  for  $p \notin F$ . Then  $\underline{h}$  is convex, and  $\varphi_{\tau,\mathcal{L}} = \sup_{\sigma,F} \underline{h}_{\tau,\mathcal{L},\sigma,F}$  is thus also convex. Finally  $\varphi_{\tau,\mathcal{L}}(p) \geq u(p)$  is obvious, taking  $\sigma$  optimal in  $D(p)$ .

2) The better inequality follows directly from part 2 of cor. 3.29 p. 217, since  $p_n = p$  for  $\sigma \in \text{NR}^n(p)$ . ■

In the particular case of linear functions, our results yield the following.

DEFINITION 3.32. Let  $\bar{Z} = \{ z \in \mathbb{R}^K \mid \exists \tau: \forall \sigma, \forall k, \limsup_{n \rightarrow \infty} \bar{\gamma}_n^k(\sigma, \tau) \leq z^k \}$  and  $\underline{Z} = \{ z \in \mathbb{R}^K \mid \exists \tau: \forall \sigma \exists \mathcal{L}: \forall k \mathcal{L}(\bar{\gamma}_n^k(\sigma, \tau)) \leq z^k \}$ .

REMARK 3.32. The existence of  $\mathcal{L}$  just means that the convex hull of the limit points intersects  $z - \mathbb{R}_+^K$ .

COROLLARY 3.33. (1)  $\bar{Z} = \underline{Z} = Z = \{ z \in \mathbb{R}^K \mid \langle p, z \rangle \geq u(p) \forall p \in \Delta(K) \}$ .

In particular the set is closed, convex, compactly generated (by  $Z \cap [-C, C]^K$ ) and has  $v_\infty(p) = \text{Cav } u(p)$  as support function.

$Z$  is therefore called the set of **approachable vectors**.

(2) The strategies  $(\tau, \mathbf{l})$  of theorem 3.18 p. 202 corresponding to  $z \in Z$  can be taken of the form  $\mathbf{l}_n = z + \varepsilon_n$ , where the sequence  $\varepsilon_n \in \mathbb{R}_+$  decreasing to zero is independent of  $z \in Z$ .

COMMENT 3.33. In a game with vector pay-offs, as in 4 p. 102, the signals are the moves, and players use a single strategy — independent of the coordinate. In our context, this is equivalent to saying that player I uses a non-revealing strategy. So our results (including 3.30 p. 216) include those of 4 p. 102 [at least for convex sets  $C$  satisfying  $C - D = C$ , where  $D$  is the positive orthant for some ordering on  $\mathbb{R}^n$ , i.e.  $D$  is a closed convex cone with non-empty interior] — except of course for the explicit  $\sqrt{n}$  rate of convergence.

In general however, the result is easier to formulate, as done here, in terms of approaching convex functions than in terms of approaching convex sets, since when the convex function is discontinuous there is no clear corresponding convex set.

COMMENT 3.34. A caricatural example of later applications (chiefly to correlated equilibria in non-zero-sum games) is the following: Imagine player I is not initially informed about the true state of nature, but receives some private information about it. Imagine also the signals are the moves. Then player I's initial information is characterised by some  $p \in \Delta(K)$ , while player II's is by some probability distribution  $\mu$  over  $\Delta(K)$  — and in the game,  $\mu$  is first used to select  $p$ , I is told of  $p$ , then  $p$  is used to select  $k \in K$ , no player being informed. To model this as in this chapter, we would need to introduce each  $p$  as a different state of nature, with the corresponding average game, and then compute the  $u$ -function and concavify it over this infinite dimensional simplex.

But by theorem 3.18 p. 202 (and prop. 3.31 p. 218), the best player II can do is to select some convex function  $\varphi$  on  $\Delta(K)$  with  $\varphi$  larger or equal to the usual  $u$ -function, and guarantee that player I will not get more than  $\varphi(p)$  whatever be his type  $p$ . Thus the optimal  $\varphi$  (and hence the optimal strategy) will be the one minimising  $\int_{\Delta(K)} \varphi(p) \mu(dp)$  over all convex  $\varphi \geq u$ . (See also comment 3.44 p. 226).

**3.h. A continuum of types.** We extend here (theorem 3.39 p. 224) theorem 3.18 p. 202 to the case of a continuum of types of player I or equivalently — as in next chapter — to the case where for each of his types there are several different pay-off matrices, and where player II wants to guarantee some vector pay-off in those.

Assume a repeated game with incomplete information  $\Gamma$  described as follows. Player II has no initial private information. Nature first chooses  $i \in I$  ( $I$  finite), which determines the pure strategy set  $S_i$  of player I, and sets  $\Omega_i$  of types of I and  $K_i$  of  $S_i \times T$  pay-off matrices.  $K_i$  is finite, and  $\Omega_i$  is a measurable space. The  $S_i \times T$  signalling matrix  $Q_\omega$  of probability distributions on  $A \times B$  depends only on  $\omega \in \Omega_i$ , and the set  $B$  is finite, while  $A$  may be any measurable space. The marginal distribution  $Q_i^{\text{II}}$  of  $Q_\omega$  on  $B$  is independent of  $\omega \in \Omega_i$ , hence will be denoted  $Q_i^{\text{II}}$ . Given  $i \in I$  nature selects  $(\omega, k) \in \Omega_i \times K_i$ , and informs player I of  $\omega$ . The game is then played forever as usually.

More generally, the pay-off functions in  $K_i$  could be functions on  $S_i \times T \times B$  or, if  $Q_\omega$  is independent of  $\omega \in \Omega_i$  (so:  $Q_i$ ) and  $A$  is finite, on  $S_i \times T \times A \times B$ : we only need that every history in  $S_i \times T \times A \times B$  determines a history of vector pay-offs in  $\mathbb{R}^{K_i}$ .

We will also write  $\Omega$  for the (disjoint) union  $\bigcup_{i \in I} \Omega_i$ . The above can be viewed as the extension of our standard model with incomplete information on one side to the case of a continuum of types of player I — but sticking to the assumption of finitely many signalling matrices  $Q_i^{\text{II}}$ . Indeed, finiteness of  $K_i$  is no real restriction (as long as pay-offs are uniformly bounded), since anyway the space of  $S_i \times T$  pay-off matrices is finite dimensional.

The function  $u$  is the same as usually; however since  $\Omega$  may be infinite, it is more convenient to view  $u$  as being defined on  $M = \Delta(P)$  with  $P = \bigcup_{i \in I} \Delta(K_i)$  — the point in  $P$  being player I's probability distribution over the pay-off matrices. For  $\mu \in M$  let  $I_\mu = \{ i \in I \mid \mu(\Delta(K_i)) > 0 \}$ .

**PROPOSITION 3.34.** (1) (a)  $u$  is u.s.c. and has a Lipschitz restriction to each open face  $M_{I_0} = \{ \mu \in M \mid I_\mu = I_0 \}$ .

$u$  is continuous if the  $S_i$  and the  $Q_i^{\text{II}}$  are independent of  $i \in I$ .

(b)  $u$  is monotone: define  $\mu_1 \leq \mu_2$  iff  $\int f(p) \mu_1(dp) \leq \int f(p) \mu_2(dp)$  for every real valued function  $f$  whose restriction to every  $\Delta(K_i)$  is convex. Then  $\mu_1 \leq \mu_2$  implies  $u(\mu_1) \leq u(\mu_2)$ .

(2) By adding finitely many convex combinations (depending just on the  $Q_i^{\text{II}}$  — thus independently of the sets  $K_i$  or the corresponding pay-off matrices) to the pure action sets  $S_i$  one obtains that,  $\forall \mu \in M$ , both best replies and optimal strategies in  $\text{NR}(\mu)$  are given by mixtures of pure strategies in  $\text{NR}(\mu)$ .

(3) Assume the actions mentioned sub 2 have been added. For every  $i \in I$ , let  $V_i = \{ (Q_i^{\text{II}}(s, t))_{t \in T} \mid s \in S_i \}$ . For every vector  $v = (v_t)_{t \in T}$  of probability distributions on  $B$ , every  $\tau \in \Delta(T)$  and  $p \in \Delta(K_i)$  let, with  $\max \emptyset = -\infty$ :

$$F_v(p, \tau) = \max \left\{ \sum_{t \in T, k \in K_i} p^k G_{s,t}^k \tau(t) \mid s \in S_i: Q_i^{\text{II}}(s, t) = v_t \quad \forall t \in T \right\}$$

Let also  $f_v(\tau) = \int_P F_v(p, \tau) \mu(dp)$ , and  $f(\tau) = \max_v f_v(\tau) = \max\{f_v(\tau) \mid v \in \bigcap_{I_\mu} V_i\}$ : then  $u(\mu) = \min_\tau f(\tau)$ .

**REMARK 3.35.** Thus  $u(\mu)$  is the value of a game, depending linearly on  $\mu$ , where player I has finitely many strategies (say  $\bigcup_{i \in I} V_i$ ), and the pay-off function is convex and continuous (even Lipschitz) on the strategy space  $\Delta(T)$  of player II.

**REMARK 3.36.** For the order in prop. 3.34.1b, cf. remark after ex. II.1Ex.20 p. 78.

**PROOF.** We start with 2 and 3. For  $i \in I$ , let  $W_i$  be the convex hull of  $V_i$ . Take  $\mu \in M_{I_0}$ . A behavioural strategy  $[x_i(p)]_{i \in I}$  belongs to  $\text{NR}(\mu)$  iff  $\exists v \in \bigcap_{I_0} W_i$  such that  $\sum_{s \in S_i} x_i(p)(s) Q_i^{\text{II}}(s, t) = v_t$ ,  $\mu$ -a.e.  $\forall t \in T$ . Modifying it on a null set, we can assume  $\sum_{s \in S_i} x_i(p)(s) Q_i^{\text{II}}(s, t) = v_t \forall t \in T, \forall p \in K_i, \forall i \in I_0$ . Observe  $\text{NR}(\mu)$  is a (weak\*) closed, hence compact, convex subset of the strategy space of player I (e.g. ex. II.1Ex.19 p. 78), so the minmax theorem 1.6 p. 4 applies and  $u(\mu)$  is well defined. Therefore

$$u(\mu) = \min_{\tau \in \Delta(T)} \max_{v \in \bigcap_{I_0} W_i} \int_P \mu(dp) G_v(h(p, \tau))$$

where, for  $p \in \Delta(K_i)$ ,  $[h(p, \tau)]_s = \sum_{t \in T} p^k G_{s,t}^k \tau(t)$  for  $s \in S_i$ ,  $= 0$  otherwise, and  $G_v(h) = \max\{ \sum_{\bigcup_{I_0} S_i} x(s) h(s) \mid \forall i \in I_0, (x_s)_{s \in S_i} \in \Delta(S_i) \text{ and } \sum_{s \in S_i} x(s) Q_i^{\text{II}}(s, t) = v_t \forall t \in T \}$ . (The measurable selection is trivial).

By ex. I.3Ex.11h p. 36, the function  $G_v(h)$  is, still for fixed  $i \in I_0$ , concave in  $v$ , convex in  $h$ , and piecewise bi-linear in  $(v, h)$ . I.e., there exists a triangulation of the polyhedron  $\bigcap_{i \in I_0} W_i$  such that, whatever be  $h$  vanishing outside  $\bigcup_{I_0} S_i$ ,  $G_v(h)$  is linear in  $v$  on each simplex of this triangulation. Hence for each  $\tau \in \Delta(T)$ , a best reply  $x(p) \in \text{NR}(\mu)$  — which determines some  $v \in \bigcap_{I_0} W_i$  — can be viewed equivalently as the random selection of some vertices  $v_\alpha$  of the sub-polyhedron containing  $v$  (and with  $v$  as expectation), followed by the use of a maximiser in the definition of  $G_{v_\alpha}[h(p, \tau)]$ , which (for  $p \in \Delta(K_i)$ ) can be selected as one of the finitely many extreme points of the polyhedron  $X_{v_\alpha}^i = \{x \in \Delta(S_i) \mid \sum_{s \in S_i} x(s) Q_i^{\text{II}}(s, t) = v_{\alpha,t} \forall t \in T\}$ . Add therefore, for all  $i \in I_0$ , to  $S_i$  all extreme points of all  $X_{v_\alpha}^i$  — for every vertex  $v_\alpha$  of our subdivision of  $\bigcap_{I_0} W_i$  — and repeat the same thing for each of the finitely many subsets  $I_0$  of  $I$ . Now,  $\forall \mu \in M$  and  $\forall \tau \in \Delta(T)$ , player I has a pure strategy best reply in  $\text{NR}(\mu)$ . Hence  $G_v(h(p, \tau)) = F_v(p, \tau)$ , and 3 follows.

For 2, let us still show that,  $\forall \mu \in M$ , player I has optimal strategies which are mixtures of pure strategies in  $\text{NR}(\mu)$ . Since he has best replies in this set, it suffices to show that the minmax theorem applies with this strategy set. For each  $v$ , denote by  $\Sigma_v$  the set of behavioural strategies  $x(p)$  of player I whose support is compatible with  $v$ , i.e. such that,  $\forall i \in I_0, \forall p \in K_i, \forall s \in S_i : [x(p)](s) > 0$  one has  $(Q_i^{\text{II}}(s, t))_{t \in T} = v$ . Since  $\Sigma_v$  is the set of behavioural strategies of player I in a game with finite action sets, it is compact and convex (ex. II.1Ex.19 p. 78). Denote by  $\Sigma_{I_0}$  the convex hull of  $\bigcup \Sigma_v$ : since only finitely many of them are non-empty,  $\Sigma_{I_0}$  is still compact and convex, and every strategy in  $\Sigma_{I_0}$  is a (finite) mixture of behavioural strategies. Hence the minmax theorem (theorem 1.6 p. 4) applies with  $\Sigma_{I_0}$  as strategy space, showing that player I can guarantee  $u(\mu)$  with a strategy in  $\Sigma_{I_0}$ . To prove our claim, it suffices therefore to show that every behavioural strategy in  $\Sigma_v$  can be replaced by a mixture of pure strategies in  $\Sigma_v$  — which follows from ex. II.1Ex.10 p. 61.

1 follows now immediately from 3. ■

REMARK 3.37. If  $W_{i_1} = W_{i_2}$  (recall  $W_i$  is the convex hull of  $V_i$ ), one can add convex combinations of pure strategies to the sets  $S_{i_1}$  and  $S_{i_2}$ , such that, after this addition, one obtains  $V_{i_1} = V_{i_2}$ . By duplicating then pure strategies if necessary, one can get that in addition each point  $v$  in  $V_{i_1} = V_{i_2}$  is generated by as many pure strategies in  $S_{i_1}$  as in  $S_{i_2}$ . Hence, renumbering now the strategies will yield  $Q_{i_1}^{\text{II}} = Q_{i_2}^{\text{II}}$ : one can now pool the two indices  $i_1$  and  $i_2$  into one, with the disjoint union of  $K_{i_1}$  and  $K_{i_2}$  as set of games.

REMARK 3.38. Extending the strategy sets as in the above proposition, one sees that the restriction of  $u$  to an open face  $M_{I_0}$  has a Lipschitz extension  $u_{I_0}$  to its closure  $\overline{M}_{I_0} = \{ \mu \in M \mid I_\mu \subseteq I_0 \}$ :  $\min_\tau \max_{\bigcap_{i \in I_0} V_i} f_v(\tau)$ .

This is itself the  $u$ -function of some game in our class, where in addition  $S_i$  and  $Q_i^{\text{II}}$  are independent of  $i$ : the formula shows that it suffices to delete all  $i \notin I_0$ , and all  $s \in S_i$  that do not lead to some  $v \in \bigcap_{i \in I_0} V_i$ . As shown in remark 3.37, we obtain then after duplicating strategies that  $Q_i^{\text{II}}$  is independent of  $i \in I (= I_0)$ .

We now extend prop. 3.19 p. 203.

PROPOSITION 3.35. Let  $\varphi$  be a l.s.c. convex function on  $M$ ,  $\varphi > u$ . Then there exists  $N$ , a strategy  $\tau$  of  $\text{II}$  in  $\Gamma_N$ , and a map  $\mathbf{l}: B^N \rightarrow C(P)$  ( $C(P)$  denoting the set of continuous functions on  $P$  with convex (piecewise linear)restriction to each  $\Delta(K_i)$ ) such that

- (1)  $\forall b \in B^N, \forall \mu \in M, \langle \mu, \mathbf{l}(b) \rangle \leq \varphi(\mu)$
- (2) On a history in  $\Gamma_N$  starting with  $p \in \Delta(K_i)$  — and where the choice of  $k \in K_i$  happens at the end of history, so  $(\bar{g}_N^k)_{k \in K_i}$  is a random vector —, define the random variable

$$(E_N = ) E = \max_{q \in \Delta(K_i)} \left[ \sum_{k \in K_i} q^k \bar{g}_N^k - [\mathbf{l}(b)](q) \right].$$

Then

$$\mathbb{E}_{\sigma, \tau}^p(E) \leq 0 \text{ for all } \sigma \text{ and } p.$$

PROOF. We just mention the differences with the proof of prop. 3.19 p. 203. First increase player II's information by informing him in addition after each stage of his move; next increase player I's information by informing him in addition of player II's signal and of his own move. At this stage, player I's old signal is uncorrelated (given player I's other information) with anything in the game — it thus serves just to describe a generalised strategy of player I, which can equivalently be described as a (behavioural, or mixed) strategy (ex. II.1Ex.10 p. 61) that uses just player II's past signals, player I's past moves, and his own type. The old signals can therefore be discarded (except in the case where the pay-off depends on them).

Thus, at this stage the signalling matrices  $Q_i$  select a message for player II, that in particular informs him of his last move, and player I's message contains player II's message together with his own past move;  $Q_i$  depends only on  $i$ , not on the type of player I.

Then, one can replace  $\Gamma$  by  $\Gamma'$ , replacing each  $S_i$  by  $\bigcup_j S_j$ .  $u'$  still decreases to  $u$ : there is something to prove only if the lower bound  $v$  of the  $u'_M(\mu)$  is  $> -\infty$ ; in that case — cf. the proof in the previous proposition that the minmax theorem holds for those games — the sets  $\Sigma_M$  of behavioural strategies of player I in  $\text{NR}(\mu)$  that guarantee him at least  $v$  are a decreasing sequence of non-empty compact sets; by the monotone convergence theorem, any strategy in the intersection guarantees  $v$  with  $M = \infty$ , hence is a non-revealing strategy in the original game that guarantees  $v$ . Hence the convergence. Since  $\varphi$  is l.s.c.

$> u$ , and  $u'_M$  are, by the previous proposition, u.s.c. and decrease to  $u$ , there will be  $M$  sufficiently large such that  $\varphi > u'_M$  (compactness). It suffices to prove the proposition for this game  $\Gamma'$ .

As in remark 3.37 above, since  $Q_i$  is now independent of  $i$ , one can pool all  $K_i$ 's together into their disjoint union  $K$  — setting  $\varphi(\mu) = +\infty$  for every  $\mu$  on  $\Delta(K)$  which is not carried by  $\bigcup_i \Delta(K_i)$  preserves convexity and lower semi-continuity: so if we prove the result on  $\Delta(K)$ , it will suffice at the end to replace the maps  $1(b)$  on  $\Delta(K)$  by their restriction to  $\bigcup_i \Delta(K_i)$ .

We have thus reduced the problem to that of a single  $i$  — i.e. a single set  $S$ , a single  $K$ , and a single  $Q$ . Now the function  $u$  is, by the previous proposition, continuous, and even Lipschitz.

Consider now an increasing sequence  $L_n$  of finite Borel partitions of  $\Delta(K)$  such that the maximum diameter  $\delta_n$  of the partition elements tends to zero. For each  $n$ , and every  $\mu$  on  $\Delta(K)$ , denote by  $\bar{\mu}_n$  the corresponding point in  $\Delta(L_n)$ ; also for every  $\ell \in L_n$  let  $\bar{G}_{s,t}^\ell = \sup_{p \in \ell} \sum_k p^k G_{s,t}^k$ . This game  $\bar{\Gamma}^n$  can be viewed as a game covered by Theorem 3.18 p. 202, i.e. with finite set of types  $L_n$ . But it can equivalently be viewed as the original game  $\Gamma$ , but where the pay-off function  $\langle p, G \rangle$  has been increased to  $\bar{G}^{\ell(p)}$ . By our assumptions on the sequence  $L_n$ , this implies that  $\bar{u}^n(\bar{\mu}_n)$  decreases uniformly to  $u(\mu)$ . Let also  $\bar{\varphi}^n(\bar{\mu}_n) = \inf\{\varphi(\mu) \mid \mu(\ell) = \bar{\mu}_n(\ell) \forall \ell \in L_n\}$ : since the map  $\mu \rightarrow \bar{\mu}_n$  is linear,  $\bar{\varphi}^n$  is convex on  $\Delta(L_n)$ ; since  $L_n$  is increasing, so is  $\bar{\varphi}^n(\bar{\mu}_n)$ ; finally  $\varphi > u$ , together with l.s.c. of  $\varphi$ , u.s.c. of  $u$ , and compactness of  $M$  imply there exists  $\varepsilon > 0$  such that  $d(\mu_1, \mu_2) \leq \varepsilon$  implies  $\varphi(\mu_1) \geq u(\mu_2) + \varepsilon$ . Choose e.g. as distance  $d(\mu_1, \mu_2) = \sup\{\langle \mu_1 - \mu_2, f \rangle \mid f \text{ has Lipschitz constant 1 on } \Delta(K)\}$ . Then  $\bar{\mu}_n = \bar{\nu}_n$  implies  $d(\mu, \nu) \leq \delta_n$  — so choosing  $\nu_n$  with  $\varphi(\nu_n)$  approximating  $\bar{\varphi}^n(\bar{\mu}_n)$  and  $\nu_n(\ell) = \bar{\mu}_n(\ell) \forall \ell \in L_n$  yields that  $\bar{\varphi}^n(\bar{\mu}_n) \geq u(\mu) + \varepsilon$  for all  $\mu$  and all  $n \geq n_0$  (with  $\delta_{n_0} \leq \varepsilon$ ). Uniform convergence of  $\bar{u}^n(\bar{\mu}_n)$  to  $u(\mu)$  implies thus that, for all  $n \geq n_1$ ,  $\bar{\varphi}^n \geq \bar{u}^n + \frac{1}{2}\varepsilon$  on  $\Delta(L_n)$ . Hence, even if  $\bar{\varphi}^n$  was not l.s.c., the same inequality would, by continuity of  $\bar{u}^n$ , still hold for the l.s.c. regularisation of  $\bar{\varphi}^{n_1}$ . Therefore, to those two we can apply prop. 3.19 p. 203 and obtain the existence of  $N$  and of a strategy  $\tau$  of player II in the  $N$ -stage repetition of  $\bar{\Gamma}^{n_1}$  such that, for all non-revealing strategies  $\sigma$  of I in this game,  $\sum_{\ell \in L_{n_1}} q_\ell \bar{\gamma}_N^\ell(\sigma, \tau) \leq \bar{\varphi}^{n_1}(q) - \varepsilon/4$  for all  $q \in \Delta(L_{n_1})$ .

Since  $\bar{\Gamma}^{n_1}$  differs from  $\Gamma'$  only by having a larger pay-off function, and since  $\bar{\varphi}^{n_1}(\bar{\mu}^{n_1}) \leq \varphi(\mu)$  for all  $\mu \in \Delta(M)$ , we obtain a fortiori that  $\bar{\gamma}_N^\mu(\sigma, \tau) = \int \bar{\gamma}_N^p(\sigma, \tau) \mu(dp) \leq \varphi(\mu) - \varepsilon/4$  for all  $\mu$ , and every non-revealing strategy  $\sigma$  of player I in  $\Gamma'_N$ .

This finishes the analogue of step 2 in the proof of prop. 3.19.

For step 3, consider now  $\Gamma'_N$  as a one-shot game  $\tilde{\Gamma}$ , with a single strategy  $\tau$  for player II. The  $u$ -function  $\tilde{u}(\mu)$  of  $\tilde{\Gamma}$  is thus convex, continuous, and strictly smaller than  $\varphi$ . But  $\tilde{\Gamma}$  can also be viewed as a game with (random) vector pay-offs, in  $\mathbb{R}^K \times \Delta(\tilde{B})$ , where every history in  $\Gamma'_N$  is mapped to the corresponding pay-off in  $\mathbb{R}^K$  and the corresponding string of messages  $[ \in B^N = \tilde{B}$ , identified with the extreme points of  $\Delta(\tilde{B})$ ], and where therefore pure strategy  $\tilde{s}$  yields a random outcome in  $\mathbb{R}^K \times \Delta(\tilde{B})$  whose distribution  $\pi_{\tilde{s}}$  is induced by the distribution of histories under  $\tilde{s}$  and  $\tau$ . Denote the barycentre of  $\pi_{\tilde{s}}$  by  $(f_{\tilde{s}}, \beta_{\tilde{s}}) \in \mathbb{R}^K \times \Delta(\tilde{B})$ , and let  $D$  be the convex hull of those points. By theorem 4.1 p. 236,  $D$  is approachable in  $\tilde{\Gamma}$  — i.e. for some constant  $M$ , and every strategy  $\sigma$  of player I (as before, there is no loss in assuming that, in  $\tilde{\Gamma}$ , player I is informed after each stage of the full vector pay-off and of his pure strategy  $\tilde{s}$ ),  $E_\sigma[d[(f_n, \beta_n), D]] \leq M/\sqrt{n}$ ,

where  $(f_n, \beta_n)$  denotes the random average pay-off after  $n$  repetitions of  $\tilde{\Gamma}$ . Denote by  $D'$  the projection of  $D$  on  $\Delta(\tilde{B})$ ; for  $\beta \in \Delta(\tilde{B})$  denote by  $\bar{\beta}$  its projection on  $P$ ; and for  $\beta \in P$  let  $D_\beta = \{f \in \mathbb{R}^K \mid (f, \beta) \in D\}$ . Then the map  $\beta \rightarrow D_\beta$  is Lipschitz by Ex. I.3Ex.4q p.30, say with constant  $L$ , and  $\beta \rightarrow \bar{\beta}$  is clearly also Lipschitz with constant 1. So  $d[f_n, D_{\bar{\beta}_n}] \leq (L+1)d[(f_n, \beta_n), D]$ , hence with  $M' = M(L+1)$  we have  $\mathbb{E}_\sigma d(f_n, D_{\bar{\beta}_n}) \leq M'/\sqrt{n}$  for all  $n$  and  $\sigma$ .

Denote by  $C'$  the maximum absolute value of pay-offs in  $\Gamma'$ , i.e.  $C' = \max_{k,s,t} |G_{s,t}^k|$  ( $G$  being the expected pay-off matrix of  $\Gamma'$  now). Let also  $[\mathbf{l}(\beta_n)](p) = \max\{\sum_k p^k \varphi_k \mid \varphi \in D_{\bar{\beta}_n}\}$ . Clearly  $\mathbf{l}(\beta_n)$  is a convex function on  $\Delta(K)$ , piecewise linear since  $D_{\bar{\beta}_n}$  is a polyhedron, and with Lipschitz constant and uniform norm  $\leq C'$ . Further, using  $\tilde{G}^p$  for  $\sum_k p^k \tilde{G}^k$ ,  $[\mathbf{l}(\beta_n)](p) = \sum_{\tilde{s} \in \tilde{S}} x(p)(\tilde{s}) \tilde{G}_{\tilde{s}}^p = \max\{\sum_{\tilde{s}} x(\tilde{s}) \tilde{G}_{\tilde{s}}^p \mid x \in \Delta(\tilde{S}), \sum x(\tilde{s}) Q_{\tilde{s}}^\Pi(b) = \bar{\beta}_n(b) \quad \forall b \in \tilde{B}\}$  — the existence of a measurable selection  $x(p)$  is trivial. Hence  $x(p)$  is a non-revealing strategy of player I in  $\tilde{\Gamma}$ , since it yields, for every type  $p$ , the same distribution  $\bar{\beta}$  on the signals of  $\Pi$ . Therefore  $\int [\mathbf{l}(\beta_n)](p) \mu(dp) \leq \tilde{u}(\mu)$  for all  $\mu$  and all  $\beta_n$ . Finally  $\max_{p \in \Delta(K)} [\tilde{g}_n^p - [\mathbf{l}(\beta_n)](p)] \leq C''d(f_n, D_{\bar{\beta}_n})$ , so for  $\tilde{M} = C'' \cdot M'$ ,  $\mathbb{E}_\sigma \max_{p \in \Delta(K)} [\tilde{g}_n^p - [\mathbf{l}(\beta_n)](p)] \leq \tilde{M}/\sqrt{n}$  for all  $n$  and  $\sigma$ . Thus, choosing  $0 < \delta < \min_\mu (\varphi(\mu) - \tilde{u}(\mu))$ , adding  $\delta$  to every function  $\mathbf{l}(\beta_n)$ , and choosing  $N_1 > [\tilde{M}/\delta]^2$  we obtain that in  $\tilde{\Gamma}_{N_1}$ ,  $\mathbb{E}_\sigma \max_{p \in \Delta(K)} [\tilde{g}_{N_1}^p - [\mathbf{l}(\beta_{N_1})](p)] \leq 0 \quad \forall \sigma$  and  $\int [\mathbf{l}(\beta_{N_1})](p) \mu(dp) < \varphi(\mu) \quad \forall \mu, \forall \beta_{N_1}$ .

Reverting now to the original game finishes the proof of the proposition. ■

**REMARK 3.39.** The proof also shows that, if in  $\Gamma$  all  $S_i$  and  $Q_i^\Pi$  are identical, then  $\mathbf{l}(b)$  has, for all  $b \in B^N$ , Lipschitz constant and uniform norm  $\leq C$ .

**LEMMA 3.36.** Denote by  $M$  the space of all probability measures on a separable metric space  $P$ . Endow  $M$  with the metric (for fixed positive constants  $C_u, C_\ell$ )  $d(\mu, \nu) = \sup\{\langle \mu - \nu, f \rangle \mid f \text{ is Lipschitz with constant } C_\ell \text{ and } \sup_x f(x) - \inf_x f(x) \leq C_u\}$ .

- (1) all those metrics — when  $C_u$  and  $C_\ell$  vary, or when the distance on  $P$  is changed to an equivalent distance — are equivalent, and induce the weak $^*$ -topology. In particular they admit the same class of Lipschitz functions.
- (2) For any extended real valued function  $u$  on a convex set  $C$ , any convex function  $\varphi \geq u$  is minorated by a minimal such function.
- (3) If  $u: M \rightarrow \mathbb{R}$  is Lipschitz with constant  $C$ , any minimal convex function  $\geq u$  has the same Lipschitz constant.

**PROOF.** 1) is obvious, except perhaps that the metric induces the weak $^*$ -topology. If a sequence  $\mu_n$  converges to  $\mu$  according to the metric, this implies  $\int f d\mu_n \rightarrow \int f d\mu$  for all bounded Lipschitz functions, hence for every bounded continuous function, since this is the limit both of an increasing and of a decreasing net (or even sequence) of bounded Lipschitz functions. Conversely, such a set of Lipschitz functions is compact in the uniform topology, i.e. the topology of uniform convergence on the unit ball of the dual of the bounded continuous functions. So, by Ascoli's theorem, it is equicontinuous on this unit ball. This yields the converse.

2) follows straight from Zorn's lemma.

3) Let  $\varphi$  be a function  $\geq u$ . Define the Lipschitz regularisation  $\hat{\varphi}$  of  $\varphi$  as the largest function with Lipschitz constant  $C$  which is  $\leq \varphi$  — this is well defined, because  $u$  is such a function, and the supremum of a family of functions with Lipschitz constant  $C$  has the same constant. Assume  $\varphi$  was convex: we claim  $\hat{\varphi}$  is convex. Indeed  $\hat{\varphi}(\mu) = \inf_\nu [\varphi(\nu) + Cd(\mu, \nu)]$ , so it suffices to prove convexity of  $\varphi(\nu)$  — which is assumed

— and of  $d(\mu, \nu)$ , which is obvious from its definition as a supremum of continuous linear functionals  $\int f d(\mu - \nu)$ .

Thus if  $\varphi$  is minimal convex,  $\varphi = \hat{\varphi}$ : it has Lipschitz constant  $C$ . ■

**COROLLARY 3.37.** (1) Every minimal convex function  $\varphi \geq u$  is the limit of a decreasing sequence of convex continuous functions.

(2) Its restriction to every open face has Lipschitz constant  $C$ .

**PROOF.** For each open face  $I_0$  — in the sense of remark 3.38 — we can by 2 find a minimal convex function  $\varphi_{I_0}$  which is  $\leq \varphi$  everywhere and which is, on  $\overline{M}_{I_0}$ ,  $\geq u_{I_0}$  (thus,  $\varphi_{I_0} = -\infty$  outside  $\overline{M}_{I_0}$ ). This is because, by upper semi-continuity of  $u$  (prop.3.34 p. 219),  $u_{I_0} \leq u \leq \varphi$  on  $\overline{M}_{I_0}$ . Since  $u_{I_0}$  is Lipschitz with constant  $C$ ,  $\varphi_{I_0}$  has the same Lipschitz constant (lemma 3.36 p. 223) on  $\overline{M}_{I_0}$ . In particular,  $\varphi_{I_0}$  is convex and u.s.c., so  $\psi = \max_{I_0} \varphi_{I_0}$  is convex and  $u \leq \psi \leq \varphi : \varphi = \max_{I_0} \varphi_{I_0}$  by minimality. Hence 2, and for 1, it suffices to construct such a decreasing sequence for each  $\varphi_{I_0}$ : taking term by term maxima will then yield the result. Let thus e.g.  $\bar{\varphi}_{I_0} = +\infty$  outside  $\overline{M}_{I_0}$ ,  $= \varphi_{I_0}$  on  $\overline{M}_{I_0}$ , and consider its Lipschitz regularisation  $\hat{\varphi}_{I_0}$ , as in the proof of lemma 3.36 p. 223:  $\hat{\varphi}_{I_0}$  is convex, Lipschitz, and coincides with  $\varphi_{I_0}$  on  $\overline{M}_{I_0}$ . Thus  $\varphi_n(\mu) = n^{-1} + \hat{\varphi}_{I_0}(\mu) - n\mu(\bigcup_{i \notin I_0} \Delta(K_i))$  satisfies our requirements. ■

**COROLLARY 3.38.** Assume  $\varphi$  is a minimal convex function  $\geq u$ , and that the convex functions  $\varphi_n \geq u$  are such that  $\limsup \varphi_n \leq \varphi$ . Then  $e_n = \sup_\mu [\varphi(\mu) - \varphi_n(\mu)]$  converges to zero.

**PROOF.** Since  $\limsup \varphi_n$  is convex and  $\geq u$ , minimality of  $\varphi$  implies  $\limsup \varphi_n = \varphi$ . Since this remains true for any subsequence,  $\varphi_n \rightarrow \varphi$  point-wise. In particular (without even taking  $\sup$ ),  $\liminf e_n \geq 0$ . In the other direction, if we minorate each  $\varphi_n$  by a minimal convex function  $\geq u$  using 2 p. 223, our assumption remains valid. So we can assume the  $\varphi_n$  are minimal convex  $\geq u$ . By 2, their restrictions to every open face  $M_{I_0}$  have Lipschitz constant  $C$ . Since they converge point-wise to  $\varphi$ , this convergence is uniform on every  $M_{I_0}$  — hence on  $M$ . ■

We are now ready for the generalisation of theorem 3.18 p. 202::

**THEOREM 3.39.** In the above described game, given a convex function  $\varphi \geq u$ , and a sequence  $\varepsilon_n > 0$  converging to zero, there exists a strategy  $\tau$  of player II and a map  $1: \bigcup_n B^n \rightarrow C(P)$  such that

- (1)  $\varphi_n(\mu) = \max_{b \in B^n} \langle \mu, 1(b) \rangle$  is decreasing, with  $\lim_{n \rightarrow \infty} \varphi_n(\mu) \leq \varphi(\mu)$ ,  $\varphi_n \leq 2C$ , and  $\forall n, \forall b \in B^n$ ,  $1(b)$  is a maximal (measurable) function such that  $\langle \mu, 1(b) \rangle \leq \varphi_n(\mu)$ .
- (2) Define for each  $n$ ,  $E_n$  as in 3.35 — this is on a space of histories that includes the initial choice of  $i \in I$ , but not the choice of the type of player I or of the pay-off matrix. Let  $N = \sup \{ n \mid E_n > 0 \}$ , with  $\sup(\emptyset) = 0$ .

Then  $\sup_{i,p \in \Delta(K_i), \sigma} P_{\sigma, \tau}^{i,p}(N \geq n) \leq \exp(-n\varepsilon_n) \forall n \geq 0$ .

- (3) Let  $M_n = \sup_{\substack{i,p,\sigma \\ m \leq n}} \|E_m\|_{L_\infty(\sigma^p, \tau, i)}$ . Then  $\varepsilon_n(M_n - 3C) \leq 1$ .

**REMARK 3.40.** The interpretation of the theorem is the same as for Theorem 3.18 p. 202. But now,  $N$  becomes the last stage where the vector pay-off does not belong to the convex set having  $1_n(b)$  as support function.

**PROOF.** There remains to do the analogue of part B of the proof of theorem 3.18 (part C will be the same). Reduce first, as in the beginning of the proof of prop.3.35 p. 221, to

the case where  $Q_i$  depends only on  $i$ . Now the suprema, in points 2 and 3 of the theorem, can be taken just over  $\sigma^p$  instead of over  $\omega$  and  $\sigma$ .

Use 2 p. 223 to replace  $\varphi$  by a minimal convex function, then 1 p. 224 to obtain a strictly decreasing sequence  $\psi_\ell$  of convex continuous functions converging to  $\varphi$ , then prop. 3.35, for each  $\ell$ , to obtain corresponding  $\tau_\ell$ ,  $N_\ell$ , and  $\hat{\mathbf{l}}_\ell: B^{N_\ell} \rightarrow C(P)$ . Without loss of generality, we can assume  $N_\ell > N_{\ell-1}$ . Then  $\varphi_\ell(\mu) = \max_{b \in B^{N_\ell}} \langle \mu, \hat{\mathbf{l}}_\ell(b) \rangle$  is convex and  $\leq \psi_\ell$  by prop. 3.35. Also  $\varphi_\ell \geq u$ , as already argued in the beginning of part B of theorem 3.18 p. 202 (by letting player I play independently stage after stage some fixed optimal strategy in  $D(\mu)$  — the non-revealing aspect of this strategy is used in that, in  $\int [E_{\sigma^p}(\hat{\mathbf{l}}_\ell(b))(p)] \mu(dp)$ , the distribution of  $b \in B^{N_\ell}$  is independent of  $p$ ). Thus, let  $\delta_\ell = \ell^{-1} + \sup_{\ell' \geq \ell} \sup_{\mu} [\varphi(\mu) - \varphi_\ell(\mu)]$ : by Corollary 3.38 p. 224  $\delta_\ell$  converges to zero. Hence, adding the constant  $\delta_\ell$  to the function  $\mathbf{l}(b)$  for each  $b \in B^{N_\ell}$ , and to  $\psi_\ell$ , we still have  $\psi_\ell$  decreasing to  $\varphi$ , and all conclusions of prop. 3.35, but now  $\varphi_\ell > \varphi \forall \ell$ . Hence by compactness, we can extract a subsequence  $\ell_i$  such that  $\varphi_{\ell_i}$  is decreasing to  $\varphi$  — and such that further  $\varphi_{\ell_1} \leq 2C$ . [If  $C = 0$ , the theorem is obvious; otherwise, choosing  $\ell_1$  such that  $\delta_{\ell_1} \leq C$  does the job, since before the addition  $\varphi_{\ell_1} \leq C$ . Next, since the  $\psi_\ell$  are continuous and decreasing to  $\varphi$ , and since  $\varphi_{\ell_i}$  is continuous and  $> \varphi$ , define inductively  $\ell_{i+1} = \min\{ \ell > \ell_i \mid \psi_\ell < \varphi_{\ell_i} \}$ .] Thus, we can assume that, for our original sequence,  $2C \geq \psi_\ell(p) = \max_{b \in B^{N_\ell}} \langle \mu, \hat{\mathbf{l}}_\ell(b) \rangle$ . Replace now also  $\hat{\mathbf{l}}_1$  by the constant function  $2C$ , and let  $C_\ell = \max(C_{\ell-1}, C + \max_{b \in B^{N_\ell}} \max_{p \in P} |\hat{\mathbf{l}}_\ell(b)(p)|)$ ,  $C_0 = 0$ .

Denote then by  $S_\ell^i$  the pure strategy set of player I in  $\Gamma_{N_\ell}^i$ , where he is not informed of the choice of  $p \in \Delta(K_i)$ ; let  $P_{\ell,s}^i$  be the joint distribution (on  $\mathbb{R} \times A^{N_\ell}$ , given  $\tau_\ell$ ,  $i \in I$  and  $s \in S_\ell^i$ ) of  $E_{N_\ell}$  and of the signals of player I (in  $A^{N_\ell}$ ).

The rest of part B is now as before — with the obvious changes in notation, like  $p \rightarrow \mu$ . Just for the matter of increasing the “coordinates” of  $\bar{\mathbf{l}}_n(b)$  such as to obtain equality with  $g_n + \eta_n$ , observe that, when increasing them thus in a maximal way, one obtains a convex combination of other functions  $\bar{\mathbf{l}}_n(b')$ , and thus still a function in  $C(P)$ . Indeed,  $B^{N_\ell}$  is finite, so we are as in a polyhedral case.

Part C is identical to the previous case — only in the beginning the inequality  $E_n \leq \frac{W_i}{n} - \eta_{T_i}$  is justified by the fact that the maximum of a sum is less or equal to the sum of the maxima. ■

**COMMENT 3.41.** Some work was required to be able to deal with any convex function  $\varphi \geq u$ , and to get the sequence  $\varphi_n$  to be decreasing — but the prize is worthwhile: not only does it yield the strongest possible form of convergence for the  $\varphi_n$ , and the weakest possible assumptions on  $\varphi$ , but chiefly it makes the statement completely “non-topological” — no weak\* or other topology on  $M$  appears explicitly or implicitly in the statement, which remains therefore just as valid on  $\Omega$  instead of  $P$ , taking for  $M$  an arbitrary convex set of probability measures on  $\Omega$ : Make first the reduction to the case  $Q_i$  depending only on  $i \in I$ , let then  $p_\omega \in P$  be a version of the conditional probability on  $\bigcup_{i \in I} K_i$  given  $\omega$ , and use it to map all  $\mu \in M$  to their images  $\tilde{\mu}$  in  $\Delta(P)$ . Observe that  $\tilde{\mu}_1 = \tilde{\mu}_2 \Rightarrow u(\mu_1) = u(\mu_2)$ , and that since  $\mu \rightarrow \tilde{\mu}$  is linear,  $\tilde{\varphi}(\nu) = \inf\{ \varphi(\mu) \mid \tilde{\mu} = \nu \}$  (with  $\inf \emptyset = +\infty$ ) is a convex function on  $\Delta(P)$ ,  $\geq u$ . Apply the theorem to that function. Transform finally the function  $\mathbf{l}(b) \in C(P)$  to the function  $[\mathbf{l}(b)](p_\omega)$  on  $\Omega$ . Then  $\sup_{i,p \in \Delta(K_i), \sigma}$  is  $\sup_{\omega, \sigma}$ ,  $P_{\sigma, \tau}^{i,p}$  is  $P_{\sigma, \tau}^\omega$  and  $L_\infty(\sigma^p, \tau, i)$  becomes  $L_\infty(\sigma^\omega, \tau)$ .

**COMMENT 3.42.** In this vein, the sets  $K_i$  are often a nuisance, and it is sometimes more convenient to argue directly in a straight model with incomplete information on one

side, with one pay-off matrix  $G_\omega$  ( $S_i \times T$ ) for each  $\omega \in \Omega_i$  (i.e., the matrix  $\sum_{K_i} p_\omega^k G^k$ ). As long as pay-offs are uniformly bounded, this model is perfectly equivalent, since one can then always include all possible  $S_i \times T$  matrices  $G_\omega$  in a large simplex, and take its set of vertices for  $K_i$ .

In such a framework, the functions  $\mathbf{l}(b) \in C(P)$  become piecewise affine, convex functions of  $G_\omega$ . It seems one should be able to do somewhat better, and obtain them as piecewise linear functions minus a constant — with other properties like monotonicity, invariance under addition of constants, etc.

Still in this framework, a more canonical representation is possible, viewing  $\mu$  as a probability measure over pay-off matrices — or more generally as a bounded, positive measure, extending everything by homogeneity of degree 1. In addition to the previously found properties — like the above mentioned invariance under addition of constants, monotonicity w.r.t. the usual order, the monotonicity of prop. 3.34 p. 219, etc. one obtains now that two measures determine the same model if they differ only by rescaling some matrices  $G_\omega$ , and rescaling in a compensatory way the mass attached to them (i.e., if the vector measures  $G_\omega^1 \mu^1(d\omega)$  and  $G_\omega^2 \mu^2(d\omega)$  are the same). This can be used to rescale all matrices  $G \neq 0$  onto the unit sphere (for some norm) — and hence to get rid of the assumption of uniformly bounded pay-offs.

**COMMENT 3.43.** Could the theorem in some sense be “decomposed”, e.g. into:

- an intrinsic characterisation (independent of pay-offs . . . but this may be a fallacy, since pay-off matrices are the linear functionals on  $\Delta(S \times T)$ ) of “approachable” sets, for instance in  $\prod_i \Delta(S^i \times T)$  or  $\Delta(B)$ .
- a proof that, for given pay-off matrices
  - the above “approachable” sets are characterised by convex functions  $\geq u$
  - the corresponding “approach strategies” yield also the function  $\mathbf{l}(b)$

**COMMENT 3.44. Statistically perfect monitoring.**

A particular case where things simplify, and that may give a better feel for the theorem, is that where, for each  $i$  and  $k \in K_i$ , every column of  $G^k$  is a linear combination of the columns  $Q_t^i(b)$  ( $t \in T, b \in B$ ). (This is substantially more general than asking that the  $S^i \times (T \times B)$ -matrix  $Q_{s,t}^i(b)$  be of full row rank, since whenever the game arises from some extensive form where several mixed strategies of player I correspond to the same behavioural strategies, those several mixed strategies will induce the same distribution on  $B$  for each  $t \in T$  — so there is no full row rank —, but will also induce the same pay-offs for each  $t \in T$ ).

Indeed, in such a case, expanding the pure strategy sets  $S_i$  as in prop. 3.34 p. 219, our assumption implies that all pure strategies corresponding to some  $v \in V_i$  are duplicates (induce the same pay-offs) — thus, after identification, one can think of  $V_i$  itself as the pure strategy set  $S_i$ . But then the “max” in  $F_v(p, \tau)$  disappears — so  $F_v$  is bi-linear in  $(p, \tau)$ , hence  $f_v$  is linear in  $\tau$  and depends only on the “barycentre”  $\bar{\mu} \in \Delta(K)$  of  $\mu$  — i.e. the induced probability distribution  $K = \bigcup_i K_i$ . So  $u$  becomes a function on  $\Delta(K)$  — the value of an “average game”. Further, since  $\mu \rightarrow \bar{\mu}$  is linear, if  $\varphi$  is a convex function on  $M$  with  $\varphi(\mu) \geq u(\bar{\mu})$ , then  $\bar{\varphi}: \Delta(K) \rightarrow \mathbb{R}$  defined by  $\bar{\varphi}(\pi) = \inf\{\varphi(\mu) \mid \bar{\mu} = \pi\}$  is convex on  $\Delta(K)$  and  $\geq u$ . Thus we can limit attention to those convex functions  $\varphi$  on  $M$  which arise from some convex function  $\bar{\varphi} \geq u$  on  $\Delta(K)$  by  $\varphi(\mu) = \bar{\varphi}(\bar{\mu})$ . For each  $b \in \bigcup_n B^n$ , let then  $\bar{\mathbf{l}}(b)$  denote the (linear) concavification of  $\mathbf{l}(b)$  (thus,  $\bar{\mathbf{l}}(b) \in \mathbb{R}^K$ ). Denote also, for  $\mu \in P$  — or for  $\bar{\mu} \in \Delta(K)$  —, by  $\tilde{\mu} \in M$  the corresponding measure carried by the extreme points of the  $\Delta(K_i)$ . Then for each  $\mu$ ,  $\langle \mu, \bar{\mathbf{l}}(b) \rangle = \langle \bar{\mu}, \bar{\mathbf{l}}(b) \rangle = \langle \tilde{\mu}, \mathbf{l}(b) \rangle \leq \varphi_n(\tilde{\mu}) =$

$\overline{\varphi}_n(\bar{\mu})$  by definition. So 1, used for  $\tilde{\mu}$ , yields that  $\overline{\varphi}_n(\bar{\mu}) = \max_{b \in B^n} \langle \bar{\mu}, \bar{\mathbf{I}}(b) \rangle$ , hence the  $\overline{\varphi}_n$  are piecewise linear and convex on  $\Delta(K)$ . It yields also that  $\overline{\varphi}_n(\bar{\mu})$  decreases and has limit  $\leq \overline{\varphi}(\bar{\mu}) = \varphi(\mu)$ . Finally,  $\overline{\varphi}_n \leq 2C$  and the maximality of  $\bar{\mathbf{I}}$  follow also immediately from the corresponding properties of  $\varphi_n$  and  $\mathbf{I}$ . For 2 and 3, note that, by increasing thus  $\mathbf{I}$  to  $\bar{\mathbf{I}}$ , we have only decreased  $E_n$ , so our bounds are a fortiori valid. And now  $N$  is simply the last time where  $\bar{\mathbf{I}}_n(b)$  is not an upper bound for the vector pay-off in  $\mathbb{R}^{K_i}$ . Thus, the theorem implies that, under our assumptions, the maps  $\mathbf{I}(b)$  can all be taken linear on each  $\Delta(K_i)$ . The  $\varphi_n$  become then convex functions on  $\Delta(K)$ , and the theorem becomes an approachability theorem of the convex sets in  $\mathbb{R}^K$  having the  $\varphi_n$  as support functions.

COMMENT 3.45. Still under the same assumptions as in the previous comment, it is clear that, faced with such a strategy  $\tau$ , the best player I can obtain, when constrained by such an approachable set, if he is of type  $p \in \Delta(K_i)$ , is the amount  $\varphi_n(p)$  in the  $n$ -stage game (here  $\Delta(K_i)$  is identified with the corresponding face of  $\Delta(K)$ ). Hence, player II will guarantee himself in the  $n$ -stage game the amount  $\int \varphi_n(p) \mu(dp)$ , hence in the infinite game  $\int \varphi_\infty(p) \mu(dp)$  by monotone convergence — i.e.  $\int \overline{\varphi}(p) \mu(dp)$  if he has chosen  $\overline{\varphi}$  minimal convex above  $u$  on  $\Delta(K)$ . So an optimal strategy of player II is to minimise  $\int \overline{\varphi}(p) \mu(dp)$  among all minimal convex functions on  $\Delta(K)$  which are  $\geq u$ , next to approach this  $\overline{\varphi}$  as in the theorem (cf. also comment after prop. 3.44 p. 230).

COMMENT 3.46. Corollaries 2 p. 224 or 3.38 p. 224 provide a compact  $T_1$ -topology (with countable basis) on the space of minimal convex functions  $\varphi$  (or  $\overline{\varphi}$  as in the above remark) with are  $\geq u$ : take as basis of neighbourhoods of  $\varphi$  all  $\varphi'$  which are uniformly  $> \varphi - \varepsilon$  (cf. ex. I.3Ex.15c p. 38). In this topology, expressions like  $\int \varphi(p) \mu(dp)$  above are l.s.c. — hence the relevant minima are achieved. Thus the theorem can be viewed as providing compact spaces of “sufficient” strategies for the infinite game. (Recall from sect. 1 p. 3 that the minmax theorem does not require the  $T_2$ -assumption on strategy spaces.)

**3.i. Implications of the approachability theorem (bis).** We extend here the results of sect. 3.g to the setup of sect. 3.h.

COROLLARY 3.40. Cor. 3.29 p. 217 remains word for word true in the present setup, when replacing  $k \in K$  by  $p \in P$ , and the posterior probability  $p_n$  on  $\Delta(K)$  (including  $n = \infty$ ) by the posterior probability  $\mu_n$  on  $\Delta(P)$ .

PROOF. Use the same proof. Observe that the martingale convergence theorem holds for probability distributions on a compact metric space, using the weak\* topology [reduce to the scalar case by considering a countable dense set of continuous functions]. ■

DEFINITION 3.41. Replace  $p$  by  $\mu$  in definition 3.30 p. 218.

PROPOSITION 3.42. Prop. 3.31 p. 218 remains word for word true replacing  $p \in \Delta(K)$  by  $\mu \in \Delta(P)$ . In 2 we have further that  $\varphi_\tau = \varphi_{\tau, \mathcal{L}}$  is independent of  $\mathcal{L}$ , and is a minimal convex function  $\geq u$  having all properties of cor. 3.37 p. 224. And we have both  $\limsup_{n \rightarrow \infty} \sup_{\sigma \in \text{NR}^n(\mu)} \overline{\gamma}_n(\mu, \sigma, \tau) \leq \varphi_\tau(\mu)$  and the existence of  $\sigma \in \text{NR}^\infty(\mu)$  (in fact,  $\sigma \in \Delta([\text{NR}(\mu)]^\infty)$ ) such that  $\overline{\gamma}_n(\mu, \sigma, \tau) \rightarrow \varphi_\tau(\mu)$ .

PROOF. For the faces  $F$  of  $\Delta(K)$ , use the  $\overline{M}_{I_0}$  of remark 3.36 p. 220. Since  $\tau$  is in fact derived from  $\varphi_0$ , some minimal convex function  $\leq \varphi$  and  $\geq u$ , we obtain  $\varphi_{\tau, \mathcal{L}} \leq \varphi_0$ , hence by minimality of  $\varphi_0$ ,  $\varphi_{\tau, \mathcal{L}} = \varphi_0 \forall \mathcal{L}$ : there only remains to construct  $\sigma \in \text{NR}^\infty(\mu)$ . Consider the game ( $\mu$  and  $\tau$  fixed) when I chooses  $\sigma \in [\text{NR}(\mu)]^\infty$  and II chooses  $\mathcal{L}$ , with

pay-off  $h_{\tau, \mathcal{L}, \sigma}(\mu) = f(\sigma, \mathcal{L})$ . We have just shown that  $\max_{\sigma} f(\sigma, \mathcal{L}) = \varphi_{\tau}(\mu) \forall \mathcal{L}$ . Further, player II has a compact convex strategy set (in  $\ell'_{\infty}, \sigma(\ell'_{\infty}, \ell_{\infty})$ ), such that,  $\forall \sigma$ ,  $f(\sigma, \mathcal{L})$  is affine and continuous in  $\mathcal{L}$ . So, by the minmax theorem 1.8 p. 5, there exist convex combinations  $\sigma_k$  of strategies in  $[\text{NR}(\mu)]^{\infty}$  such that

$$\min_{\mathcal{L}} f(\sigma_k, \mathcal{L}) > \varphi_{\tau}(\mu) - 1/k, \quad \text{i.e. } \liminf_{n \rightarrow \infty} \bar{\gamma}_n(\mu, \sigma_k, \tau) > \varphi_{\tau}(\mu) - 1/k.$$

Choose thus  $n_k > n_{k-1}$  such that for  $n \geq n_k$   $\sup_{\sigma \in \text{NR}^n(\mu)} \bar{\gamma}_n(\mu, \sigma, \tau) \leq \varphi_{\tau}(\mu) + k^{-1}$ , and  $\bar{\gamma}_n(\mu, \sigma_k, \tau) \geq \varphi_{\tau}(\mu) - k^{-1}$ .

Select now  $N_0 = 0$ ,  $N_k \geq n_{k+1}$ ,  $\frac{N_{k+1}}{k+1} \geq 3 \frac{N_k}{k}$ , such that the  $N_k$  are “end of bloc”-dates for the strategy  $\tau$  of player II, i.e. such that his strategy after  $N_k$  is independent of the history up to  $N_k$ . And define  $\sigma \in \text{NR}^{\infty}(\mu)$  as using  $\sigma_k$  at all dates  $n$  with  $N_{k-1} < n \leq N_k$  (using e.g. independent realisations of all  $\sigma_k$ ).

Assume, by induction, that  $\bar{\gamma}_{N_k}(\mu, \sigma, \tau) \geq \varphi_{\tau}(\mu) - 2k^{-1}$ ; so for  $N_k \leq n \leq N_{k+1}$ ,

$$\begin{aligned} \bar{\gamma}_n(\mu, \sigma, \tau) &\geq \bar{\gamma}_n(\mu, \sigma_{k+1}, \tau) - \frac{N_k}{n} (\bar{\gamma}_{N_k}(\mu, \sigma_{k+1}, \tau) - \bar{\gamma}_{N_k}(\mu, \sigma, \tau)) \\ &\geq \bar{\gamma}_n(\mu, \sigma_{k+1}, \tau) - \frac{3N_k}{kn} \end{aligned}$$

using our induction assumption and the first inequality determining  $n_k (\leq N_k)$ . Letting thus  $n = N_{k+1}$  and using  $\frac{N_{k+1}}{k+1} \geq \frac{3N_k}{k}$  and the second inequality determining  $n_{k+1} (\leq N_{k+1})$  we obtain  $\bar{\gamma}_{N_{k+1}}(\mu, \sigma, \tau) \geq [\varphi_{\tau}(\mu) - (k+1)^{-1}] - (k+1)^{-1}$ : this is the induction step. Since for  $k = 1$  the inequality follows from  $N_1 \geq n_1$ , our induction is proved for all  $k$ . Our formula yields then, for  $N_k \leq n \leq N_{k+1}$ ,  $\bar{\gamma}_n(\mu, \sigma, \tau) \geq \varphi_{\tau}(\mu) - \frac{1}{k+1} - \frac{3N_k}{kn}$  since  $N_k \geq n_{k+1}$ , so  $\bar{\gamma}_n(\mu, \sigma, \tau) \geq \varphi_{\tau}(\mu) - \frac{4}{k}$ : thus  $\liminf_{n \rightarrow \infty} \bar{\gamma}_n(\mu, \sigma, \tau) \geq \varphi_{\tau}(\mu)$ . ■

**REMARK 3.47.** Since our existence proof of  $\sigma \in \text{NR}^{\infty}(\mu)$  depends explicitly of the form of the  $\tau$ 's constructed in theorem 3.39 p. 224, it would be better, instead of the last part of the statement, to have that  $\sup_{\sigma \in \text{Co}[\text{NR}(\mu)]^{\infty}} \liminf_{n \rightarrow \infty} \bar{\gamma}_n(\mu, \sigma, \tau) = \varphi(\mu)$  whenever  $\varphi$  is minimal convex  $\geq u$ , and for every strategy  $\tau$  of player II such that  $\varphi_{\tau, \mathcal{L}} \leq \varphi \forall \mathcal{L}$ . (Here  $\sigma \in \text{Co}[\text{NR}(\mu)]^{\infty}$  means  $\sigma$  is a convex combination of strategies in  $[\text{NR}(\mu)]^{\infty}$  — in particular  $\sigma \in \text{NR}^{\infty}(\mu)$ .)

We turn now to affine functions  $\geq u$ , and the value of the game.

**PROPOSITION 3.43.** Let  $\mathbf{H} = \{ h: P \rightarrow [-C, C] \mid h|_{\Delta(K_i)} \text{ is convex } \forall i, h \text{ has Lipschitz constant } \leq C, \int_P h(p)\mu(dp) \geq u(\mu) \forall \mu \in \Delta(P) \}$  and let  $\mathbf{H}_0$  denote the minimal elements of  $\mathbf{H}$ . Then

- (1) Every (extended real valued) affine function  $\geq u$  is minorated by a minimal such function.
- (2) The minimal such functions are the functions  $\int_P h(p)\mu(dp)$ , for  $h \in \mathbf{H}_0$ .
- (3)  $\mathbf{H}$  is compact and convex in the uniform norm, and  $\mathbf{H}_0$  is compact in the  $T_1$  topology with countable basis where the sets  $\{ h \in \mathbf{H}_0 \mid h > h_0 - \varepsilon \}$  form a basis of neighbourhoods of  $h_0 \in \mathbf{H}_0$ .
- (4)  $\mathbf{H}_0$  is a  $G_{\delta}$  in  $\mathbf{H}$ , and the inclusion map is a Borel isomorphism. Further, there exists a Borel map  $r: \mathbf{H} \rightarrow \mathbf{H}_0$  such that  $r(h) \leq h \forall h$ .

**PROOF.** 1 follows from Zorn's lemma.

For 2, assume  $\varphi$  is minimal affine; construct by lemma 3.36 p. 223 a minimal convex  $\psi$  with  $u \leq \psi \leq \varphi$ . By cor. 3.37 p. 224,  $\psi$  is u.s.c., and its restriction to every open face  $M_{I_0}$

has Lipschitz constant  $C$ . And  $\varphi$  is minimal affine  $\geq \psi$ . Let then  $(B_1, \dots, B_n)$  be a Borel partition of  $P$  into subsets of diameter  $\leq \varepsilon$  — in particular every  $B_j$  is contained in a single  $\Delta(K_i)$ . For  $\mu \in \Delta(P)$ , let  $\alpha_j = \mu(B_j)$ ,  $\mu_j(B) = (\mu(B \cap B_j))/\alpha_j$  if  $\alpha_j > 0$ , and  $\mu_j$  is an arbitrary probability on  $B_j$  otherwise. Then  $\varphi(\mu) = \sum_j \alpha_j \varphi(\mu_j) \geq \sum_j \alpha_j \psi(\mu_j)$ . And since  $\mu_j$  and every unit mass  $\delta_p$  for  $p \in B_j$  are contained in the same open face, and are  $\varepsilon$ -distant from each other, we have  $|\psi(\mu_j) - \int \psi(\delta_p) \mu_j(dp)| \leq C\varepsilon$ , thus  $\varphi(\mu) \geq \int \psi(\delta_p) \mu(dp) - C\varepsilon$ , hence  $\varphi(\mu) \geq \int \psi(\delta_p) \mu(dp)$ . Let  $h(p) = \psi(\delta_p)$ . By 2,  $h$  has Lipschitz constant  $C$ . Further, by convexity of  $\psi$ ,  $\psi(\mu) \leq \sum \alpha_j \psi(\mu_j) \leq \sum \alpha_j \int \psi(\delta_p) \mu_j(dp) + C\varepsilon = \int h(p) \mu(dp) + C\varepsilon$  by the same inequality as above, hence  $\psi(\mu) \leq \int h(p) \mu(dp)$ . So by minimality of  $\varphi$  we obtain  $\varphi(\mu) = \int h(p) \mu(dp)$ . Further, by theorem 3.39 p. 224, let  $\psi^n(\mu) = \max_{b \in B^n} \int_P a_b(p) \mu(dp)$ : we have  $\psi^n$  decreases to  $\psi$ . In particular,  $h^n(p) = \psi^n(\delta_p) = \max_{b \in B^n} a_b(p)$  decreases to  $h(p)$ . So convexity of  $a_b(p)$  implies that of  $h^n$  and hence of  $h$ . This proves 2.

Compactness and convexity of  $\mathbf{H}$  are obvious. Let  $V_\varepsilon(h_0) = \{h \in \mathbf{H}_0 \mid h > h_0 - \varepsilon\}$ : obviously  $h \in V_\varepsilon(h_0) \Rightarrow \exists \eta > 0: V_\eta(h) \subseteq V_\varepsilon(h_0)$ , so the  $V_\varepsilon(h)$  form the basis of a topology on  $\mathbf{H}_0$ , which is equivalent to the specified basis of neighbourhoods. This topology is compact because, for any ultrafilter  $\mathcal{U}$  on  $\mathbf{H}_0$ , if  $h$  denotes the limit of  $\mathcal{U}$  in the (compact) space  $\mathbf{H}$ , and  $h_0 \leq h$ ,  $h_0 \in \mathbf{H}_0$ , then obviously  $h_0$  is a limit point of  $\mathcal{U}$  in  $\mathbf{H}_0$ . It is  $T_1$  because for  $h_i \in \mathbf{H}_0$  ( $i = 1, 2$ ), if  $h_1 \neq h_2$  there exists  $\varepsilon > 0$  such that  $h_2 \notin V_\varepsilon(h_1)$  by minimality of  $h_2$ . Finally, it is second countable: Since  $\mathbf{H}_0 \subseteq \mathbf{H}$ , we can find a sequence  $h_i \in \mathbf{H}_0$  which is dense in  $\mathbf{H}_0$  in the uniform topology. Consider the sequence of open sets  $U_{k,i} = V_{k^{-1}}(h_i)$ : we have to show that, given  $h \in \mathbf{H}_0$  and  $\varepsilon > 0$ ,  $\exists(k, i)$  with  $h \in U_{k,i} \subseteq V_\varepsilon(h)$ : choose  $0 < d < k^{-1}$ ,  $d + k^{-1} < \varepsilon$ , and  $h_i$   $d$ -close to  $h$  in the uniform distance. This proves 3.

For  $f, h \in \mathbf{H}$ , let  $S_f(h)$  denote the convexification of  $f \wedge h$ . Thus  $S_f(h)$  is convex, has Lipschitz constant  $C$ , and is  $\leq h$  — so if  $\int [S_f(h)](p) \mu(dp) \geq u(\mu) \forall \mu$ , we will have  $S_f(h) \in \mathbf{H}$ . But, since  $S_f: h \rightarrow S_f(h)$  is continuous (Lipschitz constant 1 in the uniform topology),  $\mathbf{H}_f = \{h \in \mathbf{H} \mid S_f(h) \in \mathbf{H}\}$  is closed (compactness of  $\mathbf{H}$ ). Let thus  $T_f(h) = S_f(h)$  for  $h \in \mathbf{H}_f$ ,  $T_f(h) = h$  otherwise: we have  $T_f: \mathbf{H} \rightarrow \mathbf{H}$  is Borel, with  $T_f(h) \leq h \forall h$ , and  $T_f(h) \leq f$  if  $\exists h_0 \in \mathbf{H}: h_0 \leq h \wedge f$ . Let now  $f_i$  denote a dense sequence in  $\mathbf{H}$ , and add to it all  $S_C(f_n + \varepsilon)$  for  $\varepsilon \geq 0$  rational. Define Borel maps  $R_n: \mathbf{H} \rightarrow \mathbf{H}$  by  $R_0(h) = h$ ,  $R_n = T_{f_n} \circ R_{n-1}$ . Observe that  $R_n(h)$  is decreasing in  $\mathbf{H}$ ; compactness of  $\mathbf{H}$  yields thus  $R_n(h) \rightarrow R_\infty(h)$  point-wise, so  $R_\infty$  is Borel also. Clearly  $R_\infty(h) \leq h$ , further either  $R_n(h) \leq f_n$ , so  $R_\infty(h) \leq f_n$ , or there is no  $h_0 \in \mathbf{H}$  with  $h_0 \leq R_n(h)$ ,  $h_0 \leq f_n$  — so certainly not with  $h_0 \leq R_\infty(h)$ ,  $h_0 \leq f_n$ . Thus, for all  $n$ ,  $T_{f_n}(h_0) = h_0$  for  $h_0 = R_\infty(h)$ . This means  $h_0 \leq f_n$  is there exists  $h_1 \in \mathbf{H}$  with  $h_1 \leq h_0 \wedge f_n$ . So assume  $g \in \mathbf{H}$ ,  $g \leq h_0$ ,  $g \neq h_0$ . The sequence  $f_n$  being dense, we can extract from it a subsequence  $n_i$  with  $\|f_{n_i} - g\| < \varepsilon_{n_i}$ . Choose  $\varepsilon_{n_i}$  rational, since  $f_{n_i} + \varepsilon_{n_i} > g$  we can extract a further subsequence such as to make them decrease to  $g$ . Then also  $S_C(f_{n_i} + \varepsilon_{n_i})$  decreases to  $g$ : there exists a subsequence of the  $f_n$  that decreases to  $g$ . In particular, there is  $f_{n_0} \geq g$  such that  $f_{n_0}$  is not  $\geq h_0$ . But since  $g \in \mathbf{H}$ ,  $g \leq f_{n_0} \wedge h_0$ , this means  $T_{f_{n_0}}(h_0) \neq h_0$ : a contradiction. Thus  $h_0$  is minimal in  $\mathbf{H}$ :  $R_\infty(\mathbf{H}) \subseteq \mathbf{H}_0$ . Since  $R_\infty(h) \leq h$  it follows that  $R_\infty(\mathbf{H}) = \mathbf{H}_0$ , and  $R_\infty$  is the identity on  $\mathbf{H}_0$ . This establishes the existence of our Borel map  $r$  — of which we have further shown that it is even Borel as a map from  $\mathbf{H}$  to itself. On the other hand, let  $B = \{(h_1, h_2) \in \mathbf{H} \times \mathbf{H} \mid h_2 \leq h_1, h_2 \neq h_1\}$ :  $B$  is a  $K_\sigma$  in a compact metric space, as a difference of two closed sets. And  $\mathbf{C}\mathbf{H}_0$  being the projection of  $B$  on the first factor is therefore also a  $K_\sigma$ . Thus  $\mathbf{H}_0$  is a  $G_\delta$  in  $\mathbf{H}$ .

Thus, to finish the proof, there remains to show that the two measurable structures on  $\mathbf{H}_0$  coincide, i.e. that the inclusion map  $\mathbf{H}_0 \subseteq \mathbf{H}$  is Borel (because clearly the subspace topology on  $\mathbf{H}_0$  is stronger than the topology on  $\mathbf{H}_0$ ). Since  $\mathbf{H}$  is compact metric, its Borel sets are generated by the evaluation maps  $h \rightarrow h(p)$  ( $p \in P$ ) (using either the first separation theorem, or Stone-Weierstrass). Thus it suffices to show that  $h \rightarrow h(p)$  is Borel-measurable on  $\mathbf{H}_0$  — which is obvious since  $\{h \mid h(p) > \alpha\}$  is clearly open. ■

COMMENT 3.48. (3) and (4) imply that every probability measure on the Borel sets of  $\mathbf{H}_0$  is regular (cf. 1.d p. 6): the Borel set is also Borel in  $\mathbf{H}$ , and the measure a (regular) measure on  $\mathbf{H}$ , so the Borel set can be approximated from inside by a compact subset of  $\mathbf{H}$ ; being contained in  $\mathbf{H}_0$ , this subset is closed (and clearly compact) in  $\mathbf{H}_0$ .

- PROPOSITION 3.44.
- (1) Define  $(\mathbf{Cav} u)(\mu) = \sup\{\sum_{i=1}^n \alpha_i u(\mu_i) \mid \alpha_i \geq 0, \sum \alpha_i = 1, \sum \alpha_i \mu_i = \mu\}$  (i.e., it is the smallest concave function which is  $\geq u$ ). Then  $(\mathbf{Cav} u)(\mu) = \max\{\int u(\nu)\rho(d\nu) \mid \rho \text{ probability measure}, \int \nu \rho(d\nu) = \mu\}$ , and has Lipschitz constant  $C$ .
  - (2)  $\int h(p)\mu(dp)$  is (“uniformly”) l.s.c. on  $\mathbf{H}_0 \times M$ .
  - (3)  $(\mathbf{Cav} u)(\mu) = \min_{h \in \mathbf{H}_0} \int h(p)\mu(dp)$ .
  - (4) Every convex function  $\geq \mathbf{Cav} u$  is minorated by some function  $\int h(p)\mu(dp)$  with  $h \in \mathbf{H}_0$ .

PROOF. 1 Since  $\{\rho \mid \int \nu \rho(d\nu) = \mu\}$  is compact, and since  $u$  is u.s.c. (1 p. 224), it is clear that the maximum is achieved, and is  $\geq (\mathbf{Cav} u)(\mu)$ . Conversely, given  $\rho$ , consider for every  $I_0 \subseteq I$ ,  $I_0 \neq \emptyset$ , the restriction  $\rho_{I_0}$  of  $\rho$  to the open face  $M_{I_0}$ , and some measure  $\tilde{\rho}_{I_0}$  with finite support on  $M_{I_0}$ , such that  $\tilde{\rho}_{I_0}$  has same mass and same barycentre as  $\rho_{I_0}$ , and is  $\varepsilon$ -close to  $\rho_{I_0}$ . [E.g., find for every  $\mu \in \overline{M}_{I_0}$  a closed neighbourhood of diameter  $\leq \varepsilon$  of the form  $\{\nu \in \overline{M}_{I_0} \mid \nu(f_i) \geq 0\}$  where the  $f_i$  are finitely many continuous functions on  $P$ . Extract a finite open covering by the interiors of those sets, and let  $(g_j)_{j=1}^k$  enumerate the finitely many continuous functions on  $P$  thus obtained. Let  $g: M_{I_0} \rightarrow \{-1, 0, 1\}^k$ ,  $g(\nu) = (\text{sign}(\nu(g_j)))_{j=1}^k$ , and let  $B_\ell = g^{-1}(\ell)$  for  $\ell \in \{-1, 0, 1\}^k$ . The  $B_\ell$  are a Borel partition of  $M_{I_0}$  into convex sets of diameter  $\leq \varepsilon$ ; define  $\tilde{\rho}_{I_0}$  by assigning mass  $\rho_{I_0}(B_\ell)$  to the barycentre of the normalised restriction of  $\rho_{I_0}$  to  $B_\ell$ .] Let also  $\tilde{\rho} = \sum_{I_0} \tilde{\rho}_{I_0}$ . Then, since  $u$  has Lipschitz constant  $C$  on every  $M_{I_0}$  (1 p. 219), we have  $|\int u(\nu)\rho(d\nu) - \int u(\nu)\tilde{\rho}(d\nu)| \leq C\varepsilon$ . Since  $\tilde{\rho}$  has finite support and has barycentre  $\mu$ , this proves our second formula for  $\mathbf{Cav} u$ . Observe that, with the u.s.c. of  $u$ , this implies immediately that  $\mathbf{Cav} u$  is u.s.c. — since the set of probability measures  $\rho$  on  $M$  is compact. The Lipschitz aspect will follow from 3 since a Lipschitz constant is preserved when taking minima.

2 Obviously we have  $\int h d\tilde{\mu} > \int h_0 d\mu - 2\varepsilon$  for  $h \in V_\varepsilon(h_0)$  ( $h, h_0 \in \mathbf{H}_0$ ), if  $d(\tilde{\mu}, \mu) < \frac{\varepsilon}{C}$ , using the Lipschitz property of  $h_0$ .

3 The minimum is achieved by 2 and compactness of  $\mathbf{H}_0$  (3 p. 228). It is obviously  $\geq (\mathbf{Cav} u)(\mu)$ . So to prove equality, consider the locally convex space  $F = E \times \mathbb{R}$ , where  $E$  is the set of all bounded measures on  $P$  with the weak $^*$  topology. Let  $K = \{(\mu, x) \mid \mu \in M, (\mathbf{Cav} u)(\mu) \geq x \geq -C\} \subseteq F$ . By the compactness of  $M$  in  $E$  and  $\mathbf{Cav} u$  being u.s.c. (cf. supra),  $K$  is compact — and clearly convex — in  $F$ . Consider  $x_0 > (\mathbf{Cav} u)(\mu_0)$ :  $(\mu_0, x_0) \notin K$ , so it can be (strictly) separated from  $K$  by a continuous linear functional on  $F$  (cf. 1.21 p. 8), which takes the form  $\langle \mu, f \rangle + \alpha x$ , for some continuous function  $f$  on  $P$ . I.e., we have  $\langle \mu_0, f \rangle + \alpha x_0 > \max\{\langle \mu, f \rangle + \alpha x \mid (\mu, x) \in K\}$ . Let  $x_1 = (\mathbf{Cav} u)(\mu_0)$ : we have  $x_0 > x_1$ , and  $\langle \mu_0, f \rangle + \alpha x_0 > \langle \mu_0, f \rangle + \alpha x_1$ , hence  $\alpha > 0$  — so

we can divide  $f$  by  $\alpha$ , and obtain  $\langle \mu_0, f \rangle + x_0 > \max\{ \langle \mu, f \rangle + (\text{Cav } u)(\mu) \mid \mu \in M \} = \beta$ . Let thus  $g = \beta - f$ : we have  $(\text{Cav } u)(\mu) \leq \int g d\mu$ , and  $\int g d\mu_0 < x_0$ . By prop. 3.43 p. 228, it follows that  $\inf_{h \in \mathbf{H}_0} \int h d\mu_0 < x_0$ . Since  $x_0 > (\text{Cav } u)(\mu_0)$  was arbitrary, the conclusion follows.

4 By lemma 3.36 p. 223, assume  $\varphi$  minimal, hence Lipschitz —  $\text{Cav } u$  being Lipschitz. Separate then, as above,  $K$  from  $\{ \mu, x \mid x \geq \varphi(\mu) + n^{-1} \}$ , yielding an affine function  $\psi_n$  between  $\text{Cav } u$  and  $\varphi + n^{-1}$ . Take a limit  $\psi$  following some ultrafilter: it is affine, and  $\text{Cav } u \leq \psi \leq \varphi$ . Conclude by prop. 3.43 p. 228. ■

COMMENT 3.49. In the situation of comments 3.44 and 3.45 after theorem 3.39 p. 224, we claim that viewing  $u$  as a function on  $\Delta(K)$ ,  $(\text{Cav } u)(\mu) = \max\{ \int_{\Delta(K)} u d\nu \mid \nu \preceq \mu \}$ , where the order  $\nu \preceq \mu$  means  $\int \psi d\nu \leq \int \psi d\mu$  for all convex real valued  $\psi$  on  $\Delta(K)$ .

Indeed, observe first the maximum in the right-hand member is achieved: since every  $\psi$  is the limit of a decreasing sequence of  $\psi$ 's which are furthermore continuous, the definition of the order would remain the same if one required in addition continuity of  $\psi$ . Therefore,  $\{ \nu \mid \nu \preceq \mu \}$  is closed, and the maximum is achieved.

Next, observe that the functions  $h \in \mathbf{H}_0$  are those convex functions such that  $\int_P h(p) \mu(dp) \geq u(\bar{\mu})$  for every  $\mu$  on  $P$  with barycentre  $\bar{\mu}$ . Let  $p_i \in \Delta(K_i)$  be the barycentre of the restriction of  $\mu$  to  $\Delta(K_i)$ : by convexity, it suffices that  $\sum_i \alpha_i h(p_i) \geq u(\sum_i \alpha_i p_i)$  whenever  $p_i \in \Delta(K_i)$ ,  $\alpha_i \geq 0$ ,  $\sum_i \alpha_i = 1$ . This means that, denoting by  $\bar{h}$  the convexification of  $h$  [i.e., the largest convex function on  $\Delta(K)$  that coincides with  $h$  on each  $\Delta(K_i)$ ],  $\bar{h}(q) \geq u(q) \quad \forall q \in \Delta(K)$ :  $\mathbf{H}_0$  can be identified with the minimal convex functions on  $\Delta(K)$  that majorate  $u$ , i.e., it coincides with the minimal functions  $\bar{\varphi}$  of comment 3.44 (but the functions  $h$  apply to measures  $\mu \in M$  through  $\int h(q) \mu(dq)$ , while for  $\varphi$  it is  $\varphi(\bar{\mu})$ ,  $\bar{\mu}$  being the barycentre of  $\mu$ ).

Assume now  $\nu \preceq \mu$ : then for  $h \in \mathbf{H}_0$ , we have  $\int u d\nu \leq \int h d\nu \leq \int h d\mu$ , so by 3 p. 230 we have  $\int u d\nu \leq (\text{Cav } u)(\mu)$ . So the maximum is less or equal to  $(\text{Cav } u)(\mu)$ . Conversely, choose by prop. 1  $\sum_i \alpha_i \mu_i = \mu$ ,  $\sum_i \alpha_i u(\bar{\mu}_i) \geq (\text{Cav } u)(\mu) - \varepsilon$ . Let  $\nu = \sum_i \alpha_i \delta_{\bar{\mu}_i}$ ; we have  $\int u d\nu \geq (\text{Cav } u)(\mu) - \varepsilon$ , and clearly, by Jensen's inequality,  $\int \psi d\nu \leq \int \psi d\mu$  for any convex  $\psi$  — i.e.,  $\nu \preceq \mu$ .

PROPOSITION 3.45. (1) There exists a Borel map  $\mu \mapsto \overline{Q}_\mu$  from  $M$  to  $\Delta(M)$  such that, for every  $\mu$ ,  $\mu = \int \nu \overline{Q}_\mu(d\nu)$ , i.e.  $\mu$  is the barycentre of  $\overline{Q}_\mu$ , and  $\int u(\nu) \overline{Q}_\mu(d\nu) = (\text{Cav } u)(\mu)$ .

- (2) For every such map  $\overline{Q}_\mu$ , there exists  $Q_\mu(d\nu \mid p)$ , s.t.  $(\mu, p) \mapsto Q_\mu(\cdot \mid p)$  is Borel from  $M \times P$  to  $\Delta(M)$  and is, for every  $\mu$ , a version of the conditional distribution of  $\nu$  given  $p$  under  $\nu(dp) \overline{Q}_\mu(d\nu)$  — i.e.  $\int f(p, \nu) \nu(dp) \overline{Q}_\mu(d\nu) = \int f(p, \nu) Q_\mu(d\nu \mid p) \mu(dp)$  for every non-negative Borel function  $f$  on  $P \times M$ . In other words,  $\forall \nu$ ,  $\nu$  in  $\Delta(P)$  is the posterior on  $P$  given that  $\nu \in M$  was observed — when  $\nu \in M$  was selected according to  $Q_\mu(d\nu \mid p)$  and the true state  $p \in P$  according to  $\mu(dp)$ .
- (3) There exists a Borel map  $\sigma(\mu, p)$  on  $M \times P$  such that, for  $p \in \Delta(K_i)$ ,  $\sigma(\mu, p) \in \Delta(S_i)$  and such that,  $\forall \mu$ ,  $\sigma(\mu, \cdot)$  is an optimal strategy in  $\text{NR}(\mu)$ .
- (4)  $\forall \mu$ , the strategy of player I consisting in selecting, given his observation  $p \in P$ ,  $\nu \in M$  according to  $Q_\mu(d\nu \mid p)$ , next using  $\sigma(\nu, p)$  independently at each stage is an optimal strategy in  $\Gamma^\infty(\mu)$  (written as a generalised strategy — cf. ex. II.1Ex.10 p. 61) — i.e., guarantees him  $(\text{Cav } u)(\mu)$  at every stage  $n$ :  $E_{\mu, \sigma, \tau}(\gamma_n) \geq (\text{Cav } u)(\mu) \quad \forall n, \forall \tau$ . [Thus, this strategy (or the corresponding mixed strategy) depends in a Borel way on the parameter  $\mu$ .]

PROOF. 1 The set  $E$  of pairs  $(\mu, \bar{Q})$  satisfying the requirements is a closed subset of  $M \times \Delta(M)$  by 1 p. 230 and 1 p. 224, non-empty above every  $\mu$ . Use thus 7.i.

2 follows from ex. II.1Ex.9 p. 60.

3 It suffices to define  $\sigma$  on each open face  $M_{I_0}$  separately. By remark 3.38 p. 221, this reduces the problem to that of a single  $i$ , thus  $\#I = 1$ . So one can take **NR** independent of  $\mu$ . Further, by 2, we can reduce **NR** to the mixtures of pure strategies in **NR**: the set  $S$  is partitioned into subsets  $(S_v)_{v \in V}$ , and  $Q_{s,t}^{\text{II}}$  depends only on  $v$ : we can write  $Q_{v,t}^{\text{II}}$  — i.e., optimal strategies in **NR** can be written as  $\alpha_v \in \Delta(V)$ , together with, for each  $v \in V$ , a type dependent selection of  $s \in S_v$  — i.e.,  $\forall v \in V$ , a Borel map  $\sigma_v$  from  $\Delta(K)$  to  $\Delta(S_v)$ . Let us first view  $\bar{\sigma}$  as the measure  $\mu(dp)\alpha(v)\sigma_{v,p}(ds)$  on  $\Delta(K) \times S$ . As we said,  $v$  and  $\Delta(K)$  must be independent, i.e. we need  $\bar{\sigma}(A \cap S_v) = \mu(A)\bar{\sigma}(S_v)$  for every Borel set  $A$  in  $\Delta(K)$  and  $v \in V$ . This set of pairs  $(\mu, \bar{\sigma})$  is clearly weak\*-compact — the Borel sets  $A$  can be replaced by continuous functions. And an optimal strategy is one that maximises  $\min_{t \in T} \int \bar{\sigma}(dp, ds) G_{s,t}^p$ , which is weak\*-continuous. An u.s.c. correspondence like that from  $\mu$  to the set of compatible  $\bar{\sigma}$  is a Borel map to compact subsets with the Hausdorff topology (7.b). But again the correspondence from compact subsets with the Hausdorff topology to the corresponding set of maximisers of some continuous function is u.s.c., hence Borel — so the composite map, from  $\mu$  to the set of optimal  $\bar{\sigma}$  at  $\mu$ , is a Borel map to compact sets, hence admits a Borel selection  $\bar{\sigma}_\mu$  (7.i). Apply now e.g. ex. II.1Ex.9 p. 60 to obtain non-negative Borel functions  $(\sigma_s(\mu, p))_{s \in S}$  with  $\sum_s \sigma_s = 1$  such that  $\sigma_s(\mu, \cdot)$  is a Radon-Nikodym derivative of  $\bar{\sigma}(A \times \{s\})$  w.r.t.  $\mu = \bar{\sigma}(A \times S)$ .

4 is obvious. ■

COMMENT 3.50. Non-revealing strategies have also a basically finite dimensional representation: the pay-off  $\gamma(\bar{\sigma}, \tau)$  resulting from  $\bar{\sigma}$  and  $\tau$  equals, as seen above,  $\int \bar{\sigma}(dp, ds) G_{s,\tau}^p$ . Since  $G^p$  is linear in  $p$ , this means only the barycentre  $\bar{p}_s$  of the distribution  $\nu_s(dp)$  of  $p$  given  $s$  matters, together with  $\lambda_v(s)$ , the probability of  $s$  given  $S_v$ , and  $\alpha_v$ . Our above condition on  $\bar{\sigma}$  becomes:  $\forall v, \sum_{S_v} \lambda_v(s) \nu_s = \mu$ . The condition that there exist such  $\nu_s$  with barycentre  $\bar{p}_s$  is that  $\mu$  is obtained by dilatation of the measure  $\sum_{S_v} \lambda_v(s) \delta_{\bar{p}_s}$ , i.e., by Strassen's theorem (ex. II.1Ex.20 p. 78), that  $\sum_{S_v} \lambda_v(s) \varphi(\bar{p}_s) \leq \int \varphi(p) \mu(dp)$  for every convex (continuous) function  $\varphi$  on  $\Delta(K)$  and  $v \in V$ .

So this is the system of inequalities describing the constraints on the  $\lambda_v$  and the  $\bar{p}_s$  — and in terms of those variables, letting  $\lambda(s) = \alpha_v \lambda_v(s)$  for  $s \in S_v$ , we have  $\gamma(\lambda, \bar{p}; \tau) = \sum_s \lambda(s) G_{s,\tau}^{\bar{p}_s}$ .

We turn now to the optimal strategies of player II.

PROPOSITION 3.46. (1) There exists a Borel map  $h: \mu \mapsto h_\mu$  from  $M$  to  $\mathbf{H}_0$  such that  $(\mathbf{Cav} u)(\mu) = \int h_\mu(p) \mu(dp)$ .

(2) For every sequence  $\varepsilon_n$  converging to zero there exists a sequence  $\delta_n$  with  $2 \geq \delta_n \geq \delta_{n+1}$  converging to zero and a continuous map  $\tau: h \mapsto \tau_h$  from  $\mathbf{H}$  to behavioural strategies of player II in  $\Gamma_\infty$  (or to mixed strategies, with the weak\* topology), such that  $\forall h \in \mathbf{H}, \forall n \geq 0, \forall i \in I, \forall p \in \Delta(K_i), \forall \sigma, P_{\sigma, \tau_h}^{i,p}(N \geq n) \leq \exp(-n\varepsilon_n)$ , where  $N = \sup\{n \mid E_n > C\delta_n\}$  (and  $\sup(\emptyset) = 0$ ), and where  $E_n = \max_{q \in \Delta(K_i)} [\sum_{k \in K_i} q^k \bar{g}_n^k - h(q)]$ .

Further  $\delta_n$  can be chosen to depend only on the sequence  $\varepsilon_n$  and on the  $Q_{s,t}^i(b)$  ( $b \in B, s \in S_i, t \in T, i \in I$ ), provided one sets  $E_n = d(\bar{g}_n, \mathbf{L}_h)$ , with  $\mathbf{L}_h = \{x \in \mathbb{R}^{K_i} \mid \langle q, x \rangle \leq h(q) \forall q \in \Delta(K_i)\}$  and with  $d$  being the Euclidian distance and provided one interprets the constant  $C$  as  $\max_{i,s,t} d[(G_{s,t}^k)_{k \in K_i}, 0]$ .

COMMENT 3.51. It seems plausible one might be able to select  $\tau_h$  independently of the sequence  $\varepsilon_n$ . This would be equivalent to requiring  $\limsup_{n \rightarrow \infty} [F_n(y)]^{1/n} < 1 \forall y > 0$ , where  $F_n(y) = \sup_{i,p,\sigma} P_{\sigma,\tau_h}^{i,p}(E_n > y)$ . It would thus lead at the same time to a much simpler and a much sharper statement. It would however probably require a completely different proof, where the strategy is not build up from longer and longer blocs. [Even with a single state of nature, assume we have for each bloc of length  $n$  a strategy  $\tau_n$  such that  $E_n = +1$  or  $-1$  with probability  $\frac{1}{2}$  each (thus no “error term”). If we use successively such longer and longer blocs, this will force  $\limsup_{n \rightarrow \infty} [F_n(y)]^{1/n} = 1$ .]

COMMENT 3.52. The reinterpretation with the Euclidian distances in the “further” clause is in order to obtain a sequence  $\delta_n$  independent of the dimensions  $\#K_i$  of the space of vector pay-offs.

PROOF. 1 Let  $\mathbf{H}(\mu) = \{h \in \mathbf{H} \mid \int h d_\mu = (\text{Cav } u)(\mu)\}$ ,  $\mathbf{H}_0(\mu) = \mathbf{H}(\mu) \cap \mathbf{H}_0$ . By 2 p. 230,  $\mathbf{H}_0(\mu) \neq \emptyset$ . By the (joint) continuity of  $\int h d\mu$  and the continuity of  $\text{Cav } u$  (1 p. 230), the correspondence  $\mathbf{H}(\mu)$  is u.s.c. Hence, by 7.i, there is a measurable selection; composing this with the map  $r$  of 4 p. 228 yields the result.

2 Fix  $h \in \mathbf{H}$ , and apply theorem 3.39 p. 224 to get a corresponding  $(\tau, a)$ . Observe that  $\varphi_n(\mu) = \max_{b \in B^n} \int [\mathbf{l}(b)](p) \mu(dp)$  are a decreasing sequence of continuous functions, with limit less or equal to the continuous function  $\varphi(\mu) = \int h(p) \mu(dp)$ . So, by Dini’s theorem,  $\delta_n = \max_\mu (\varphi_n(\mu) - \varphi(\mu))$  decreases to zero. So  $\forall n, \forall b \in B^n, \forall \mu \in M$  we have  $\int [\mathbf{l}(b)](p) \mu(dp) \leq \int h(p) \mu(dp) + \delta_n$ , thus  $[\mathbf{l}(b)](p) \leq h(p) + \delta_n$ : we can assume without loss of generality that  $[\mathbf{l}(b)](p) = h(p) + \delta_n$ , since  $h \in C(P)$ . By theorem 2 p. 224, we have thus  $P_{\sigma,\tau}^{i,p}(E_n > \delta_n) \leq \exp(-n\varepsilon_n) \forall i \in I, \forall p \in \Delta(K_i), \forall n, \forall \sigma$ . Since  $|E_n| \leq 2C$  we obtain  $E_{\sigma,\tau}^{i,p}(E_n) \leq \delta_n + 2C \exp(-n\varepsilon_n) = \delta'_n(h)$ .

Let  $f(n, \tau, h) = \sup_{i,p,\sigma} E_{\sigma,\tau}^{i,p}(E_n)$  — where  $h \in \mathbf{H}$  appears in  $E_n$ , and  $\tau$  varies over all strategies of player II. Clearly  $f$  has Lipschitz constant 1 with respect to  $h$ . Let  $g(n, h) = \inf_\tau f(n, \tau, h)$ :  $g$  also has Lipschitz constant 1, and we have just seen that,  $\forall h, g(n, h) \leq \delta'_n(h)$ , so  $\limsup_{n \rightarrow \infty} g(n, h) \leq 0$ . By compactness of  $\mathbf{H}$ , we have uniform convergence: there exists a sequence  $\delta''_n$  decreasing to zero such that,  $\forall n, \forall h \in \mathbf{H}, \exists \tau_{n,h}: f(n, \tau_{n,h}, h) < \delta''_n$ . Further, we can assume  $\tau_{n,h}$  completely mixed since player II has a finite pure strategy set in  $\Gamma_n$ . By the Lipschitz property of  $f$ , there is an open neighbourhood  $U_h$  of  $h$  such that  $f(n, \tau_{n,h}, h') < \delta''_n \forall h' \in U_h$ . Extract a finite covering  $(U_{h_j})_{j \in J}$  from this open covering, and consider a corresponding continuous partition of unity  $(\varphi: \mathbf{H} \rightarrow \Delta(J)$  such that  $\varphi_j(h) > 0 \Rightarrow h \in U_{h_j}$ ). Let  $\tau_n(h) = \sum_{j \in J} \varphi_j(h) \tau_{n,h_j}$  (the convex structure here being that of mixed strategies). Since  $f$  is a supremum of linear functions of  $\tau$ ,  $\{\tau \mid f(n, \tau, h) < \delta''_n\}$  is convex, and since  $\forall h, \varphi_j(h) > 0 \Rightarrow f(n, \tau_{n,h_j}, h) < \delta''_n$ , we obtain  $f(n, \tau_n(h), h) < \delta''_n \forall n, \forall h$ . Clearly  $\tau_n(h)$  depends continuously on  $h$ ; since it is completely mixed for all  $h$ , the corresponding behavioural strategy is also a continuous function of  $h$ . And  $\forall h \in \mathbf{H}, \forall i \in I, \forall p \in \Delta(K_i), \forall n > 0, \forall \sigma \quad E_{\sigma,\tau_n(h)}^{i,p}(E_n) \leq \delta''_n$ .

There only remains to repeat, with this sequence of strategies  $\tau_n(h)$  (and with  $[\mathbf{l}(b)](p) = h(p) + \delta''_n \quad \forall b \in B^n$ ), the last part of the proof of theorem 3.39 p. 224 — i.e. of theorem 3.18 p. 202.

To obtain also the “further” clause, consider first the case where the pay-offs  $G_{s,t}^k$  are non-random — i.e. depend only on the action pair  $(s, t)$ . Observe it suffices to obtain a single sequence  $\delta_n$  valid for all games (and all sets  $K_i$ ) with  $C = 1$ . Fix then a subset  $L_i$  of  $\mathbb{R}^{S_i \times T}$  with  $\#L_i = \#(S_i \times T) + 1$  such that the simplex spanned by  $L_i$  contains the

unit ball (for the maximum norm), and is contained in a ball of minimal diameter, say  $\zeta(\#(S_i \times T))$ . Let  $\bar{C} = \max_{i \in I} \zeta(\#(S_i \times T))$ .

Construct as above a sequence  $\delta_n$  and a map  $\tau_h$  for the game with the sets  $L_i$  instead of  $K_i$ : we claim this is the required sequence; there only remains to construct a corresponding map  $\bar{\tau}$  for the game with sets  $K_i$ .

By construction we have, for each  $k \in K_i$ , a probability distribution  $\pi_k$  over  $L_i$  such that  $G^k = \sum_{\ell \in L_i} \pi_{k\ell} G^\ell$ . This induces an affine map  $\varphi$  from  $P = \bigcup_i \Delta(K_i)$  to  $Q = \bigcup_i \Delta(L_i)$ , with  $\varphi(p) = \sum_{k \in K_i} p^k \pi_k$  for  $p \in \Delta(K_i)$ ,  $\ell \in L_i$ . Denote the  $u$ -functions of the games with  $K_i$  and  $L_i$  by  $u$  and  $v$  respectively, and observe that  $u(\mu) = v(\varphi(\mu)) \quad \forall \mu \in \Delta(P)$ .

Indeed, best replies in  $\Gamma_{NR}$  can be chosen to be  $\varphi(p)$ -measurable, so the result follows using the minmax theorem with this  $\sigma$ -field. Given  $h: P \rightarrow \mathbb{R}$  in  $\mathbf{H}$  ( $= \mathbf{H}_K$ ), let  $\bar{h}(q) = \min\{h(p) \mid \varphi(p) = q\}$  (with  $\min \emptyset = +\infty$ ). Then  $\bar{h}: Q \rightarrow \mathbb{R}$  is convex and l.s.c., and for every  $\nu \in \Delta(Q)$  for which  $\int \bar{h}(q)\nu(dq) < \infty$  there exists (measurable selection theorem)  $\mu \in \Delta(P)$  with  $\varphi(\mu) = \nu$  and with  $\int \bar{h}(q)\nu(dq) = \int h(p)\mu(dp) \geq u(\mu) = v(\nu)$ . So  $\int \bar{h}(q)\nu(dq)$  is an affine function  $\geq v$ . Denote by  $\varphi^*(h): Q \rightarrow \mathbb{R}$  the Lipschitz regularisation with constant  $\bar{C}$  (as in the proof of 3 p. 223) of  $\text{Vex } \min(\bar{h}, \bar{C})$ : by 1 and 2 p. 228, we also have that  $\int [\varphi^*(h)](q)\nu(dq)$  is an affine function  $\geq v$ , so  $\varphi^*(h)$  belongs to  $\mathbf{H}_L$  (for the game with  $L_i$ ). The map  $\varphi^*$  is clearly continuous (from  $\mathbf{H}_K$  to  $\mathbf{H}_L$ ), with  $h(p) \geq [\varphi^*(h)](\varphi(p)) \quad \forall p \in P, \forall h \in \mathbf{H}_K$ . Define then  $\bar{\tau}_h$  as  $\tau_{\varphi^*(h)}$ : the result is now clear, since the random variables  $E_n$  (with  $K$ ) are less or equal to the corresponding  $E_n$  with  $L$ .

Consider finally the general case. Consider the auxiliary game  $\tilde{\Gamma}$  where the random pay-offs  $G_{s,t,a,b}^k$  have been replaced by their barycentre  $G_{s,t}^k$  and on every history, consider besides the actual average vector pay-off  $\bar{g}_n$  its analogue  $\tilde{g}_n$  that would arise from  $\tilde{\Gamma}$ . Observe that  $\max_{k,s,t} |G_{s,t}^k| \leq C$ ; so let  $\tau$  and  $\delta_n$  be those constructed above for the game  $\tilde{\Gamma}$ : now (2) holds, with  $E_n = \max_{q \in \Delta(K_i)} \left[ \sum_{k \in K_i} q^k \tilde{g}_n^k - h(q) \right]$ . Then  $d(\tilde{g}_n, \mathbf{L}_h) \leq E_n \sqrt{\dim}$ ,  $\dim$  being the dimension of the space. But, as seen above, this dimension is in fact bounded by  $\#(S_i \times T)$  instead of  $\#K_i$ . Hence, after multiplying the sequence  $\delta_n$  by  $\max_{i \in I} \sqrt{\#(S_i \times T)}$ , we obtain indeed the desired statement, but still with  $\tilde{g}_n$  instead of  $\bar{g}_n$ . Observe now that, by ex. II.4Ex.6 p. 105 (in the game with vector pay-offs  $G_{s,t,a,b}^k - G_{s,t}^k$ ),  $\mathbb{E} \sup_{m \geq n} d(\bar{g}_n, \tilde{g}_n) \leq 4\sqrt{2} \frac{C}{\sqrt{n}}$ .

Let thus  $E_n = d(\bar{g}_n, \mathbf{L}_h)$ ,  $\tilde{E}_n = d(\tilde{g}_n, \mathbf{L}_h)$ : we have  $\mathbb{E}_{\tau, \sigma}(E_n) \leq \mathbb{E}(\tilde{E}_n) + 4\sqrt{2} \frac{C}{\sqrt{n}} \leq 4\sqrt{2} \frac{C}{\sqrt{n}} + C\delta_n + C \exp(-n\varepsilon_n) = C\eta_n$ , where the sequence  $\eta_n$  converges to zero and depends only on the  $Q_{s,t}^i(b)$ . Since also  $0 \leq E_n \leq C$ , it suffices now again to repeat part C of the proof of theorem 3.18 p. 202, with  $N_\ell = \ell$ ,  $\tau_\ell$  the restriction of  $\tau$  to the first  $\ell$  stages, and the variables  $\frac{E_\ell}{C}$ . ■

**COMMENT 3.53.** Given the above results, definition 3.32 and cor. 3.33 p. 218 generalise now directly, interpreting this time  $\overline{Z}$  and  $\underline{Z}$  as sets of affine functions on  $M$  — and  $Z$  is compactly generated by  $\mathbf{H}_0$ .

**COMMENT 3.54.** If one thinks of  $\Delta(K_i)$  as a neighbourhood of zero in the space of all  $S_i \times T$  matrices, then, for  $h \in \mathbf{H}_0$ ,  $h|_{\Delta(K_i)}$  becomes a convex function on this neighbourhood. Its minimality implies immediately it is positively homogeneous of degree one — so extends to the whole space —, similarly it is monotone, and invariant under addition of constants. [The latter two properties imply the Lipschitz property and uniform norm with constant  $C$ .]

COMMENT 3.55. When  $\#I = 1$ , prop.3.46.2 p. 232 becomes an extension of the approachability theorem of sect. 4 p. 102 to the case where the game with vector pay-offs has signals — i.e., the game is described by having, for every pure strategy pair  $s, t$ , a probability distribution over triplets formed of a signal of player I, a signal of player II, and a vector pay-off in the ball of radius  $C$ . The sets  $S, T$ , and  $B$  must be finite.

Indeed, one can always assume that II's signals inform him of his own action; the worst case is then when player I's signal consists of his own move, player II's signal, and the actual vector pay-off — any additional information is irrelevant. Formally the previous treatment requires player I's set of signals to be finite, but one sees that the above argument still applies in this case.

Also, at first sight this concerns only approachability of convex sets  $A$  such that  $A - \mathbb{R}_+^K \subseteq A$  — but this yields immediately the general result by mapping  $\mathbb{R}^K$  into  $\mathbb{R}^{2K}$  by the map  $\varphi_{2i-1}(x) = x_i, \varphi_{2i}(x) = -x_i$ : for any arbitrary convex set  $A$  and point  $x$  in  $\mathbb{R}^K$ , let  $\tilde{x} = \varphi(x)$ ,  $\tilde{A} = \varphi(A) - \mathbb{R}_+^{2K}$ : then  $d(x, A) \leq d(\tilde{x}, \tilde{A}) \leq \sqrt{2}d(x, A)$ , so our approachability theorem for the set  $\tilde{A}$  immediately implies the corresponding result for the general compact convex set  $A$  (and so for closed convex  $A$ , replacing it by its intersection with the ball of radius  $C$ ).

COMMENT 3.56. The corresponding excludability criterion is obvious (e.g. from proposition 4 p. 231): for any other convex set  $C$ , the corresponding function  $h_C$  is such that  $\int h_C(p)\mu(dp) < u(\mu)$  for some  $\mu \in \Delta(P)$  — which (continuity of  $u$ ) can be assumed with finite support. I.e., there exists  $\varepsilon > 0$ ,  $(p_j)_{j \in J}$  in  $\Delta(K)$  — with  $J$  finite — and  $\sigma_j \in \Delta(S)$  such that, for any one-stage strategy of player II, all  $\sigma_j$  induce the same probability distribution on signals for player II, and such that, for any strategy  $\tau$  of player II in  $\Gamma_\infty$ , if  $z_n(j)$  denotes the expected vector pay-off generated by  $\tau$  and by the i.i.d. play of  $\sigma_j$  in  $\Gamma_n$ , then  $\max_J [\langle p_j, z_n(j) \rangle - \sup_{z \in C} \langle p_j, z \rangle] \geq \varepsilon$  for all  $n$ .

#### 4. The rôle of the normal distribution

The appearance of  $1/\sqrt{n}$  in the speed of convergence may have the following probabilistic interpretation: Consider a game with  $K = \{1, 2\}$  and both  $G^1$  and  $G^2$  are  $2 \times 2$  games. Suppose that  $\text{Cav } u(p) = u(p)$ . So  $\lim_{n \rightarrow \infty} v_n(p) = u(p)$  which means that in the limit as  $n \rightarrow \infty$  player I can obtain no more than  $u(p)$  which is what he can guarantee by ignoring his additional information and playing identically in both games. Nevertheless player I can generally obtain more than  $u(p)$  in  $\Gamma_n(p)$  for any finite  $n$ . In order to do this he has to play differently in the two possible games. In other words he has to deviate from this non-revealing strategy. How much can he deviate and how much can he gain by this deviation? Let  $(s, 1-s)$  be player I's optimal mixed move in  $D(p)$  (i.e. play the first pure move with probability  $s$  and the second with probability  $1-s$  ( $0 \leq s \leq 1$ )). By definition, if player I plays this mixed move repeatedly in all  $n$  stages of  $\Gamma_n(p)$  he guarantees  $u(p)$ . If player I in fact does this, the actual choice of pure move made by him can be regarded as a Bernoulli trial with probabilities  $(s, 1-s)$ . By the Central Limit Theorem, the proportion of times that the first pure move is played in  $n$  such trials is approximately normally distributed around the mean  $s$  with a standard deviation of the order of  $1/\sqrt{n}$ . Therefore if player I wishes to “cheat” player II by making use of his additional information (i.e. playing different mixed moves in the two possible games) without enabling him to detect it, he has to do this in such a way that this proportion will fall within the limits of few standard deviations i.e.  $\alpha/\sqrt{n}$  from  $s$ . Using again the minmax theorem, any deviation of higher order will be detected and used by player II to hold the pay-off to a number smaller

than  $u(p)$ . Clearly a deviation of order not higher than  $1/\sqrt{n}$  from  $(s, 1-s)$  will make a deviation in the pay-off which is also of order not higher than  $1/\sqrt{n}$ . The existence of games with **error term**  $\delta_n(p) = v_n(p) - v_\infty(p) \approx (1/\sqrt{n})$  implies that there are games in which player I can exhaust the whole probabilistic deviation mentioned above. In other games he may be able to exhaust only a small part of it such as  $\approx (\frac{\ln n}{n})$  or  $\approx (1/n)$ .

In this section we shall see that the connection of the error term to the Central Limit Theorem is even much closer than what is outlined above. The normal distribution appears explicitly in the asymptotic behaviour of  $v_n(p)$ . In the next section we prove some more results about the speed of convergence of  $v_n(p)$ .

**EXAMPLE 4.1.** Consider the following game with  $K = \{1, 2\}$ , full monitoring and pay-off matrices:

$$G^1 = \begin{pmatrix} 3 & -1 \\ -3 & 1 \end{pmatrix} \quad G^2 = \begin{pmatrix} 2 & -2 \\ -2 & 2 \end{pmatrix}$$

and the prior probability distribution on  $K$  is  $(p, 1-p)$ .

We know that  $v_n(p)$ , which equals in this case to the error term  $\delta_n(p) = v_n(p) - v_\infty(p)$ , is bounded by (by theorem 2.10 p. 189):

$$0 \leq v_n(p) \leq \frac{6\sqrt{pp'}}{\sqrt{n}}$$

For this specific game sharper bounds can be obtained:  $v_n(p) \leq \frac{\sqrt{pp'}}{\sqrt{n}}$  (See ex. VEx.10a p. 258) and  $v_n(p) \geq \frac{pp'}{\sqrt{n}}$ , cf. (Zamir, 1971–1972). It turns out however that a much stronger result can be proved namely:

**THEOREM 4.1.** For all  $p \in [0, 1]$ ,

$$\lim_{n \rightarrow \infty} \sqrt{n} v_n(p) = \phi(p)$$

where

$$\phi(p) = \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2}x_p^2} \text{ and } \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{x_p} e^{-\frac{1}{2}x^2} dx = p$$

In words: the limit of  $\sqrt{n}v_n(p)$  is the standard normal density function evaluated at its  $p$ -quantile.

Another way to state this theorem is: The coefficient of the leading term (i.e. of  $1/\sqrt{n}$ ) in the expansion of  $v_n(p)$  is  $\phi(p)$ . Before proving this puzzling result it may be helpful to outline the heuristic arguments that lead to it. For this we shall need the following recursive formula for  $v_n$  which will be useful also in later sections.

**LEMMA 4.2. (Recursive Formula)** For any two-person zero-sum game with incomplete information on one side (player II is uninformed) and with full monitoring, the following holds for all  $p \in \Delta(K)$  and for all  $n \geq 1$ ,

$$(1) \quad v_{n+1}(p) = \frac{1}{n+1} \max_{x \in X^K} \left( \min_{t \in T} \sum_{k \in K} p^k x^k G_t^k + n \sum_s \bar{x}_s v_n(p_s) \right),$$

Where  $\bar{x} = \sum_{k \in K} p^k x^k$ , and the probabilities  $p_s = (p_s^k)_{k \in K} \in \Delta(K)$  are given by  $p_s^k = p^k x_s^k / \bar{x}_s$  for all  $s$  such that  $\bar{x}_s > 0$ , and  $G_t^k$  is the  $t^{\text{th}}$  column of  $G^k$ .

PROOF. This is clearly a special case of the general recursive formula 3.2 (cf. ex. VEx.6 p. 255), nevertheless we give here an elementary proof which is available in this case: A strategy of player I in  $\Gamma_{n+1}(p)$  includes an element  $x = (x^k)_{k \in K}$  of  $X^K$  where if the state is  $k$ , player I plays in the first stage the mixed move  $x^k = (x_s^k)_{s \in S}$  in  $\Delta(S)$ . If player II plays at first stage his pure move  $t \in T$ , the expected first stage pay-off is  $\sum_{k \in K} p^k x^k G_t^k$ . Given  $x$  and given  $s_1 = s$ , the conditional probability distribution on  $K$  is  $p_s = P(\cdot | s)$ . By playing optimally in  $\Gamma_n(p_s)$  player I can guarantee a conditional expected cumulative pay-off of at least  $n v_n(p_s)$  for the last  $n$  stages. Since the probability of  $s$  is  $\bar{x}_s = \sum_{k \in K} p^k x^k$  we conclude that  $x$  followed by an optimal strategy in the resulting  $\Gamma_n(p_s)$  guarantees player I at least

$$\frac{1}{n+1} \min_{t \in T} \left( \sum_{k \in K} p^k x^k G_t^k + n \sum_s \bar{x}_s v_n(p_s) \right),$$

which, since the  $p_s$  are independent of  $t$ , equals to

$$\frac{1}{n+1} \left( \min_{t \in T} \sum_{k \in K} p^k x^k G_t^k + n \sum_s \bar{x}_s v_n(p_s) \right)$$

Taking the maximum over  $X^K$  it follows that  $v_{n+1}(p)$  is greater or equal to the right hand side of equation (1). The other direction of the inequality is obtained using the minmax theorem: Given  $\sigma$  (and hence  $x$ ), player II can minimise his expected pay-off at the first stage and then, following a move  $s$ , compute  $p_s$  and play optimally in  $\Gamma_n(p_s)$ . Doing that he guarantees the maximand  $(\dots)$  in the right hand side of equation (1). This proves the inequality:  $v_{n+1} \leq \dots$  and completes the proof of the lemma. ■

**4.a. The heuristics of the result.** Our departure point is the recursive formula (1) which for this game reads:

$$(2) \quad v_{n+1}(p) = \frac{1}{n+1} \max_{0 \leq x, y \leq 1} \left\{ \min [3p(x - x') + 2p'(y - y'), p(x' - x) + 2p'(y' - y)] + n(\bar{x}v_n(p_T) + \bar{x}'v_n(p_B)) \right\}$$

where  $(x, x')$ ,  $(y, y')$  are the mixed strategies of player I in  $G^1$  and  $G^2$  respectively,

$$(3) \quad (\bar{x}, \bar{x}') = (px + p'y, px' + p'y')$$

and

$$(4) \quad p_T = \frac{px}{\bar{x}} \quad p_B = \frac{px'}{\bar{x}'}$$

Assume (as often turn out to be the case in our proofs) that the  $\max_{0 \leq x, y \leq 1}$  in that formula is achieved when  $x$  and  $y$  equalise player II's expected pay-off for left and right first move, i.e.

$$3p(x - x') + 2p'(y - y') = p(x' - x) + 2p'(y' - y)$$

which is  $y = 1/2 - p(x - x')/(2p')$ . Changing variable to  $\xi = p(x - x')$  and denoting  $U_n(p) = \sqrt{n}v_n(p)$  we obtain (using the notation  $a \wedge b = \min(a, b)$ ):

$$(5) \quad U_{n+1}(p) = \frac{1}{\sqrt{n+1}} \max_{0 \leq \xi \leq p \wedge p'} \left\{ \xi + \frac{\sqrt{n}}{2} (U_n(p + \xi) + U_n(p - \xi)) \right\}$$

Suppose now that  $\lim_{n \rightarrow \infty} U_n(p)$  exists and equals to  $\varphi(p)$ . Letting  $\xi = \alpha_n/\sqrt{n}$  we then have:

$$\begin{aligned}\sqrt{1 + \frac{1}{n}} \varphi(p) &\simeq \max_{\alpha_n} \left\{ \frac{\alpha_n}{n} + \frac{1}{2} [\varphi(p + \frac{\alpha_n}{\sqrt{n}}) + \varphi(p - \frac{\alpha_n}{\sqrt{n}})] \right\} \\ &\simeq \max_{\alpha_n} \left\{ \frac{\alpha_n}{n} + \varphi(p) + \frac{\alpha_n^2}{2n} \varphi''(p) \right\} \\ &= \varphi(p) + \frac{1}{n} \max_{\alpha_n} \left\{ \alpha_n + \frac{\alpha_n^2}{2} \varphi''(p) \right\} \\ &= \varphi(p) - \frac{1}{2n \varphi''(p)}\end{aligned}$$

On the other hand  $(\sqrt{1+1/n})\varphi \simeq \varphi + \varphi/(2n)$  thus  $\varphi = -1/\varphi''$ . We conclude that if  $\sqrt{n}v_n(p)$  converges, the limit is a solution of the differential equation:

$$(6) \quad \varphi(p)\varphi''(p) + 1 = 0$$

To solve equation (6) we first have

$$\varphi'(p) = - \int_{1/2}^p \frac{1}{\varphi(p)} dp$$

Here we have chosen 1/2 as the lower bound of the integration so as to have  $\varphi'(p) = 0$  which is implied by the symmetry of  $\varphi(p)$  about  $p = 1/2$ . (This is a property of all  $U_n(p)$  and can be easily proved by inducting using (5) and the symmetry of the arbitrary  $U_0(p)$ .)

Now let

$$z(p) = -\varphi'(p) = \int_{1/2}^p \frac{1}{\varphi(p)} dp$$

We have  $z'(p) = 1/\varphi(p)$  and thus  $\varphi = dp/dz$ . Replacing the variable  $p$  by  $z$  we get:

$$\varphi'_z = \varphi'_p \frac{dp}{dz} = \varphi'_p \varphi = -z\varphi$$

and thus  $\ln \varphi = K - z^2/2$  or

$$(7) \quad \varphi = \frac{A}{\sqrt{2\pi}} e^{-\frac{1}{2}z^2}$$

for some constant  $A$ . Since  $\varphi = dp/dz$ ,

$$(8) \quad p = B + \int_{-\infty}^{z(p)} \frac{A}{\sqrt{2\pi}} e^{-\frac{1}{2}x^2} dx$$

for some constant  $B$ . Denoting by  $F(x)$  the cumulative standard normal distribution we have therefore:

$$\begin{aligned}\varphi(z) &= AF'(z) \\ p &= B + AF(z)\end{aligned}$$

Now  $\varphi \geq 0$  and  $\varphi \not\equiv 0$  imply  $A > 0$  from which it follows by (8) that  $z(p)$  is increasing in  $p$ . Since  $\varphi(0) = \varphi(1) = 0$  we have by (7):  $z(0) = -\infty$  and  $z(1) = +\infty$ , and therefore by (8) we have

$$0 = B \quad 1 = B + A$$

So  $\varphi(p) = F'(z)$  and  $p = F(z)$ , i.e.  $\varphi(p)$  is the standard normal density evaluated at its  $p$ -quantile, that is  $\phi(p)$ .

COMMENT 4.2. Note that the above heuristics has nothing to do with the intuitive argument in the beginning of this section. In other words the normal distribution appeared not through the Central Limit Theorem but rather as the solution of a certain differential equation. As we shall see later in this section, this is not an isolated result for this specific example only but rather a general one. At least for  $2 \times 2$  games we know so far the following: Whenever the leading term in the expansion of the error term is  $1/\sqrt{n}$  the coefficient is an appropriately scaled  $\phi$  function.

**4.b. Proof of Theorem 4.1.** For the formal proof of theorem 4.1 p. 236 we need the following general result about martingales in  $[0, 1]$ . Let  $\mathcal{X}_p^n = \{X_m\}_{m=1}^n$  denote an  $n$ -martingale bounded in  $[0, 1]$  with  $E(X_1) = p$ , and let  $V(\mathcal{X}_p^n)$  denote its  $L_1$  variation i.e.

$$V(\mathcal{X}_p^n) = \sum_{m=1}^{n-1} E(|X_{m+1} - X_m|),$$

then we have:

THEOREM 4.3. (*The  $L_1$  variation of a bounded martingale.*)

$$\lim_{n \rightarrow \infty} \sup_{\mathcal{X}_p^n} \left[ \frac{1}{\sqrt{n}} V(\mathcal{X}_p^n) \right] = \phi(p)$$

The proof of this theorem uses some constructions and several lemmas (two of which are proved in the appendix to this chapter) and it is concluded after Lemma 4.1 p. 236.

For  $p \in [0, 1]$  let  $S(p) = \{(\xi, \eta) \mid 0 \leq \xi \leq p'; 0 \leq \eta \leq p\}$  and define two sequences of functions on  $[0, 1]$ ,  $\{\varphi_n\}$  and  $\{\psi_n\}$  by  $\varphi_1 \equiv \psi_1 \equiv 0$ , and for  $n = 1, 2, \dots$ :

$$(9) \quad \sqrt{n+1} \varphi_{n+1}(p) = \max_{(\xi, \eta) \in S(p)} \left[ \sqrt{n} \left( \frac{\eta}{\xi + \eta} \psi_n(p + \xi) + \frac{\xi}{\xi + \eta} \psi_n(p - \eta) \right) + \frac{2\xi\eta}{\xi + \eta} \right]$$

$$(10) \quad \psi_{n+1} = \text{Cav } \varphi_{n+1}$$

In the recursive formula (9) the expression in the square brackets is defined to be  $\sqrt{n}\psi_n(p)$  for  $\xi = \eta = 0$ . We first observe that

$$\varphi_n(0) = \psi_n(0) = \varphi_n(1) = \psi_n(1) = 0 \quad \text{for all } n.$$

It also follows readily from the definitions that

$$(11) \quad \varphi_2(p) = \psi_2(p) = 2p(1-p) \quad \text{for } p \in [0, 1].$$

REMARK 4.3. All functions  $\{\varphi_n\}_{n=1}^\infty$  and  $\{\psi_n\}_{n=1}^\infty$  are symmetric about  $p = 1/2$ . This is easily proved by induction using (9) and (10) and observing that the  $\text{Cav}$  operator preserves symmetry about  $p = 1/2$ .

LEMMA 4.4.

$$\sup_{\mathcal{X}_p^2} V(\mathcal{X}_p^2) = \psi_2(p) = 2p(1-p).$$

PROOF. Follows readily by Jensen's inequality since  $|x - p|$  is convex. ■

LEMMA 4.5. *For all  $n$  and for all  $p \in [0, 1]$*

$$\sup_{\mathcal{X}_p^n} \left[ \frac{1}{\sqrt{n}} V(\mathcal{X}_p^n) \right] \leq \psi_n(p)$$

PROOF. For  $n = 1$  both sides are 0 and for  $n = 2$  the inequality follows from lemma 4.7 p. 241 and equation (11). Proceeding by induction, assume it is true for  $n \leq m - 1$  and let us prove it for  $n = m$ . Since  $\psi_m = \mathbf{Cav} \varphi_m$  it is enough to prove that for  $p \in [0, 1]$

$$\sup_{\mathcal{X}_p^m} \left[ \frac{1}{\sqrt{m}} V(\mathcal{X}_p^m) \right] \leq \varphi_m(p)$$

To prove this let  $\Omega_L = \{X_2 \leq p\}$ ,  $\Omega_R = \{X_2 > p\}$ ,  $\lambda = P(\Omega_R)$ . For any  $\mathcal{X}_p^m$  we have

$$\begin{aligned} V(\mathcal{X}_p^m) &= \sum_2^m \mathbb{E}(|X_i - X_{i-1}|) \\ &= \mathbb{E}(|X_2 - p|) + \lambda \sum_3^m \mathbb{E}(|X_i - X_{i-1}| \mid \Omega_R) + (1 - \lambda) \sum_3^m \mathbb{E}(|X_i - X_{i-1}| \mid \Omega_L) \end{aligned}$$

which is, by induction hypothesis, with  $X_L = \mathbb{E}(X_2 \mid \Omega_L)$ ,  $X_R = \mathbb{E}(X_2 \mid \Omega_R)$ ,

$$\leq \lambda(X_R - p) + (1 - \lambda)(p - X_L) + \sqrt{m-1} [\lambda\psi_{m-1}(X_R) + (1 - \lambda)\psi_{m-1}(X_L)]$$

Let  $X_R = p + \xi$ ,  $X_L = p - \eta$ : then  $\xi \geq 0$ ,  $\eta \geq 0$ ,  $\lambda = \eta/(\xi + \eta)$  and the last inequality becomes:

$$V(\mathcal{X}_p^m) \leq \sqrt{m-1} \left[ \frac{\eta}{\xi + \eta} \psi_{m-1}(p + \xi) + \frac{\xi}{\xi + \eta} \psi_{m-1}(p - \eta) \right] + \frac{2\xi\eta}{\xi + \eta}.$$

Hence by definition  $V(\mathcal{X}_p^m) \leq \sqrt{m}\varphi_m(p)$ , concluding the proof of lemma 4.5.  $\blacksquare$

Proceeding in the proof of theorem 4.3 p. 239 define the sequence  $\{(\tilde{\phi}_n)\}_{n=1}^\infty$  of functions on  $[0, 1]$  by  $\tilde{\phi}_1 \equiv 0$  and for  $n = 1, 2, \dots$ ,

$$(12) \quad \sqrt{n+1}\tilde{\phi}_{n+1}(p) = \max_{(\xi, \eta) \in S(p)} \left[ \sqrt{n} \left( \frac{\eta}{\xi + \eta} \tilde{\phi}_n(p + \xi) + \frac{\xi}{\xi + \eta} \tilde{\phi}_n(p - \eta) \right) + \frac{2\xi\eta}{\xi + \eta} \right]$$

Here again the expression in the square brackets is defined to be  $\sqrt{n}\tilde{\phi}_n(p)$  if  $\xi = \eta = 0$ . It follows readily from the definitions that for all  $p$  and all  $n$ ,  $\psi_n(p) \geq \tilde{\phi}_n(p)$ .

LEMMA 4.6. For all  $n$  and all  $p \in [0, 1]$

$$\sup_{\mathcal{X}_p^n} \frac{1}{\sqrt{n}} V(\mathcal{X}_p^n) \geq \tilde{\phi}_n(p)$$

PROOF. For each  $n = 2, 3, \dots$  and  $p \in [0, 1]$  we construct a martingale  $\mathcal{X}_p^n$  s.t.

$$(13) \quad \frac{1}{\sqrt{n}} V(\mathcal{X}_p^n) = \tilde{\phi}_n(p)$$

We do the construction inductively on  $n$ : For  $n = 2$  let  $P\{X_2 = 0\} = p'$ ,  $P\{X_2 = 1\} = p$ ; then  $V(\mathcal{X}_p^2) = 2pp' = \tilde{\phi}_2(p)$ . Assume that for each  $m \leq n$  we constructed  $\mathcal{X}_p^m = \{X_i(p)\}_{i=1}^m$  having the right variation and let us construct  $\mathcal{X}_p^{n+1}$ . Let  $(\xi_n, \eta_n)$  be a point at which the maximum is attained in (12). Define the martingale  $\mathcal{X}_p^{n+1} = \{Z_i(p)\}_{i=1}^{n+1}$  as follows:  $Z_1(p) \equiv p$

$$(14) \quad P\{Z_2(p) = p + \xi_n\} = \frac{\eta_n}{\xi_n + \eta_n}; \quad P\{Z_2(p) = p - \eta_n\} = \frac{\xi_n}{\xi_n + \eta_n}$$

and for  $i = 2, \dots, n$ :

$$(Z_{i+1}(p) \mid Z_2(p) = p + \xi_n) = X_i(p + \xi_n)$$

$$(Z_{i+1}(p) \mid Z_2(p) = p - \eta_n) = X_i(p - \eta_n)$$

It follows from (12) and (14) that

$$\begin{aligned} \frac{1}{\sqrt{n+1}} V(\mathcal{X}_p^{n+1}) &= \frac{1}{\sqrt{n+1}} [2\xi_n \eta_n + \eta_n V(\mathcal{X}_{p+\xi_n}^n) + \xi_n V(\mathcal{X}_{p-\eta_n}^n)] / (\xi_n + \eta_n) \\ &= \frac{1}{\sqrt{n+1}} [2\xi_n \eta_n + \sqrt{n} (\eta_n \tilde{\phi}_n(p + \xi_n) + \xi_n \tilde{\phi}_n(p - \eta_n))] / (\xi_n + \eta_n) \\ &= \tilde{\phi}_{n+1}(p) \end{aligned}$$

This completes the proof of lemma 4.6.  $\blacksquare$

To proceed we need the following two properties of the normal distribution  $\phi(p)$ :

LEMMA 4.7. There exists a constant  $c > 0$  such that for  $p \in [0, 1]$  and for all  $n \geq 1$ :

$$(15) \quad \frac{1}{\sqrt{n+1}} \max_{0 \leq x \leq p \wedge p'} \left[ \frac{\sqrt{n}}{2} (\phi(p+x) + \phi(p-x)) + x \right] \geq \phi(p) - \frac{c}{\sqrt{n^2}}$$

Here  $p \wedge p' = \min(p, p')$ .

LEMMA 4.8. There exists a constant  $K > 0$  such that for  $p \in [0, 1]$  and for all  $n \geq 1$ :

$$(16) \quad \frac{1}{\sqrt{n+1}} \max_{(\xi, \eta) \in S(p)} \left[ \sqrt{n} \left( \frac{\eta}{\xi + \eta} \phi(p + \xi) + \frac{\xi}{\xi + \eta} \phi(p - \eta) \right) + \frac{2\xi\eta}{\xi + \eta} \right] \leq \phi(p) + \frac{K}{n^2}$$

The rather lengthy and technical proofs to these lemmas are given in the Appendix to this chapter.

LEMMA 4.9. There exists  $\alpha > 0$  such that for all  $n \geq 1$  and for all  $p \in [0, 1]$ ,

$$(17) \quad \tilde{\phi}_n(p) \geq \phi(p) - \frac{\alpha}{\sqrt{n}}.$$

PROOF. We first prove by induction on  $k$  that for  $n \geq 1$ ,

$$(18) \quad \tilde{\phi}_{n+k}(p) \geq \phi(p) - \frac{1}{\sqrt{n+k}} \left[ \frac{\sqrt{n}}{2} + \sum_{i=n}^{n+k} \frac{4c}{i\sqrt{i}} \right]$$

for  $k = 0, 1, \dots$ , where  $c$  is the constant in lemma 4.7. In fact for  $k = 0$ ,  $\tilde{\phi}_n(p) \geq 0 \geq \phi(p) - 1/2$ . For the next step notice that in lemma 4.7 we may replace  $c/n^2$  by  $4c/(n+1)^2$ . Assume that (18) holds for  $k$ , then, by Lemma 4.7,  $\tilde{\phi}_{n+k+1}(p) =$

$$\begin{aligned} &= \frac{1}{\sqrt{n+k+1}} \max_{(\xi, \eta) \in S(p)} \left[ \sqrt{n+k} \left( \eta \tilde{\phi}_{n+k}(p + \xi) + \xi \tilde{\phi}_{n+k}(p - \eta) \right) + 2\xi\eta \right] / (\xi + \eta) \\ &\geq \frac{1}{\sqrt{n+k+1}} \max_{0 \leq x \leq p \wedge p'} \left[ \frac{\sqrt{n+k}}{2} \left( \tilde{\phi}_{n+k}(p+x) + \tilde{\phi}_{n+k}(p-x) \right) + x \right] \\ &\geq \frac{1}{\sqrt{n+k+1}} \max_{0 \leq x \leq p \wedge p'} \left[ \frac{\sqrt{n+k}}{2} (\phi(p+x) + \phi(p-x)) + x - \left( \frac{\sqrt{n}}{2} + \sum_{i=n}^{n+k} \frac{4c}{i\sqrt{i}} \right) \right] \\ &\geq \phi(p) - \frac{4c}{(n+k+1)^2} - \frac{1}{\sqrt{n+k+1}} \left( \frac{\sqrt{n}}{2} + \sum_{i=n}^{n+k} \frac{4c}{i\sqrt{i}} \right) \\ &= \phi(p) - \frac{1}{\sqrt{n+k+1}} \left( \frac{\sqrt{n}}{2} + \sum_{i=n}^{n+k+1} \frac{4c}{i\sqrt{i}} \right) \end{aligned}$$

completing the proof of (18), from which we obtain for  $n = 1$

$$\tilde{\phi}_{k+1}(p) \geq \phi(p) - \frac{1}{\sqrt{k+1}} \left( \frac{1}{2} + \sum_{i=1}^{\infty} \frac{4c}{i\sqrt{i}} \right) = \phi(p) - \frac{\alpha}{\sqrt{k+1}}$$

where  $\alpha = 1/2 + 4c \sum_{i=1}^{\infty} i^{-3/2}$ . Since this holds for all  $k$ , the proof of lemma 4.9 is completed. ■

LEMMA 4.10. *There exists  $\beta > 0$  such that for all  $n \geq 1$  and for all  $p \in [0, 1]$ ,*

$$(19) \quad \psi_n(p) \leq \phi(p) + \frac{\beta}{\sqrt{n}}.$$

PROOF. The proof is almost identical to that of the previous lemma. First  $\psi_2(p) = 2pp' \leq \phi(p) + 1/2$ . Next use Lemma 4.8 and (9) p. 239 to prove that

$$(20) \quad \psi_{k+1}(p) \leq \phi(p) + \frac{1}{\sqrt{k+1}} \left( \frac{1}{2} + \sum_{i=1}^{k+1} \frac{K}{i\sqrt{i}} \right)$$

implies

$$(21) \quad \varphi_{k+2}(p) \leq \phi(p) + \frac{1}{\sqrt{k+2}} \left( \frac{1}{2} + \sum_{i=1}^{k+2} \frac{K}{i\sqrt{i}} \right)$$

where  $K$  is the constant in lemma 4.8. Observing that the right hand side of the last inequality is a concave function of  $p$ , it implies:

$$(22) \quad \psi_{k+1} \leq \text{Cav } \varphi_{k+2}(p) \leq \phi(p) + \frac{1}{\sqrt{k+2}} \left( \frac{1}{2} + \sum_{i=1}^{k+2} \frac{K}{i\sqrt{i}} \right)$$

proving that (20) holds for all  $k$ , and hence the lemma — with  $\beta = 1/2 + K \sum_{i=1}^{\infty} i^{-3/2}$ . ■

We are now ready to conclude the proof of theorem 4.3 p. 239: By Lemmas 4.6 p. 240 and 4.9 p. 241

$$\liminf_{n \rightarrow \infty} \sup_{\mathcal{X}_p^n} \left[ \frac{1}{\sqrt{n}} V(\mathcal{X}_p^n) \right] \geq \liminf_{n \rightarrow \infty} \tilde{\phi}_n(p) \geq \phi(p)$$

and by lemma 4.5 p. 239 and lemma 4.10,

$$(23) \quad \limsup_{n \rightarrow \infty} \sup_{\mathcal{X}_p^n} \left[ \frac{1}{\sqrt{n}} V(\mathcal{X}_p^n) \right] \leq \limsup_{n \rightarrow \infty} \psi_n(p) \leq \phi(p) \quad ■$$

COROLLARY 4.11. *The martingale  $\mathcal{X}_p^n$  constructed in the proof of Lemma 4.6 p. 240 satisfies*

$$\lim_{n \rightarrow \infty} \left[ \frac{1}{\sqrt{n}} V(\mathcal{X}_p^n) \right] = \phi(p) \quad \text{for } p \in [0, 1],$$

*with speed of convergence of the order of  $1/\sqrt{n}$ .*

PROOF. By lemmas 4.9 and 4.10

$$(24) \quad \phi(p) - \frac{\alpha}{\sqrt{n}} \leq \tilde{\phi}_n(p) \leq \psi_n(p) \leq \phi(p) + \frac{\beta}{\sqrt{n}}. \quad ■$$

PROOF OF THEOREM 4.1 P. 236. Given any strategy  $\sigma$  of player I in  $\Gamma_n(p)$  let  $\tau_\sigma$  be the reply of player II consisting of playing optimally in  $D(p_m)$  at stage  $m$  for  $m = 1, 2, \dots$ . Then the conditional expected pay-off at stage  $m$  satisfies

$$(25) \quad \rho_m(\sigma, \tau_\sigma) \leq \mathbb{E}(|p_{m+1} - p_m| \mid \mathcal{H}_m^{\text{II}})$$

This is a somewhat stronger inequality than the general one derived in the proof of prop. 2.9 p. 189, the difference being that in the specific game we are dealing with, the constant  $2C = 2 \max_{k,s,t} |G_{st}^k|$  can be replaced by 1 (just by estimating errors more tightly, cf. ex. VEx.10a p. 258).

From (25), taking average over  $m$  and expectation over histories we have:

$$\bar{\gamma}_n(\sigma, \tau_\sigma) \leq \frac{1}{n} \sum_{m=1}^n \mathbb{E}(|p_{m+1} - p_m|)$$

(where  $p_1 = p$ ) and hence

$$(26) \quad v_n(p) \leq \max_{\sigma} \frac{1}{n} \sum_{m=1}^n \mathbb{E}(|p_{m+1} - p_m|)$$

Since the sequence  $\{p_m\}_{m=1}^\infty$  is such a martingale, we combine theorem 4.3 p. 239 with inequality (26) to obtain:

$$(27) \quad \limsup_{n \rightarrow \infty} \sqrt{n} v_n(p) \leq \phi(p)$$

Next we shall find a lower bound for  $\liminf_{n \rightarrow \infty} \sqrt{n} v_n(p)$ . To do this we denote  $w_n(p) = \sqrt{n} v_n(p)$  and obtain, as we did for the heuristic argument (i.e. use the recursive formula (2) p. 237, restrict the first stage strategy by  $y = 1/2 - p(x-x')/(2p')$  and let  $\xi = p(x-x')$ ), the recursive inequality:

$$(28) \quad w_{n+1}(p) \geq \frac{1}{\sqrt{n+1}} \max_{0 \leq \xi \leq p \wedge p'} \left\{ \xi + \frac{\sqrt{n}}{2} (w_n(p+\xi) + w_n(p-\xi)) \right\}$$

Define now the sequence of functions  $\{U_n(p)\}_{n=0}^\infty$  by  $U_0(p) \equiv 0$  and the recursive relation (5) p. 237, then clearly  $U_n(p) \leq w_n(p)$  for all  $p \in [0, 1]$  and for all  $n$ . Therefore to complete the proof of the theorem it suffices to prove:

$$(29) \quad \liminf_{n \rightarrow \infty} U_n(p) \geq \phi(p)$$

First we claim that for any  $n$  the inequality

$$(30) \quad U_{n+m}(p) \geq \phi(p) - \frac{\sqrt{n}}{2\sqrt{n+m}} - \sum_{i=n}^{n+m} \frac{c}{i^2}$$

holds for  $m = 0, 1, 2, \dots$ , where  $c$  is a constant in lemma 4.7 p. 241. This is proved by induction on  $m$ ; It is clearly true for  $m = 0$  since  $U_n(p) \geq 0 \geq \phi(p) - 1/2$ . Assume it holds for  $m$ , using (5) p. 237 we have

$$\begin{aligned} U_{n+m+1}(p) &= \frac{1}{\sqrt{n+m+1}} \max_{0 \leq \xi \leq p \wedge p'} \left\{ \xi + \frac{\sqrt{n+m}}{2} (U_{n+m}(p+\xi) + U_{n+m}(p-\xi)) \right\} \\ &\geq \frac{1}{\sqrt{n+m+1}} \max_{0 \leq \xi \leq p \wedge p'} \left\{ \xi + \frac{\sqrt{n+m}}{2} (\phi(p+\xi) - \frac{\sqrt{n}}{2\sqrt{n+m}} - \sum_{i=n}^{n+m} \frac{c}{i^2}) \right. \\ &\quad \left. + \phi(p-\xi) - \frac{\sqrt{n}}{2\sqrt{n+m}} - \sum_{i=n}^{n+m} \frac{c}{i^2}) \right\} \\ &\geq \frac{1}{\sqrt{n+m+1}} \max_{0 \leq \xi \leq p \wedge p'} \left\{ \xi + \frac{\sqrt{n+m}}{2} (\phi(p+\xi) + \phi(p-\xi)) \right\} \\ &\quad - \frac{\sqrt{n}}{2\sqrt{n+m+1}} - \sum_{i=n}^{n+m} \frac{c}{i^2} \end{aligned}$$

So by (15) p. 241, we get

$$\begin{aligned} U_{n+m+1}(p) &\geq \phi(p) - \frac{c}{(n+m+1)^2} - \frac{\sqrt{n}}{2\sqrt{n+m+1}} - \sum_{i=n}^{n+m} \frac{c}{i^2} \\ &= \phi(p) - \frac{\sqrt{n}}{2\sqrt{n+m+1}} - \sum_{i=n}^{n+m+1} \frac{c}{i^2} \end{aligned}$$

which establishes (30) from which we get

$$\liminf_{n \rightarrow \infty} U_n(p) = \liminf_{m \rightarrow \infty} U_{n+m}(p) \geq \phi(p) - \sum_{i=n}^{\infty} \frac{c}{i^2}$$

Since the last inequality must hold for every  $n$  and since  $c \sum_{i=n}^{\infty} i^{-2} < \infty$ , this completes the proof of (29) and hence the proof of theorem 4.1. ■

**4.c. More general results.** In this section we present results showing that the appearance of the normal distribution in the asymptotic value of the game in example 4.1 p. 236 is not an isolated incident but rather part of a general phenomenon. Unfortunately the proofs of these results are technically quite involved and lengthy which make them unaffordable in the framework of this book. They can be found in (Mertens and Zamir, 1995). (See also (De Meyer, 1996b), (De Meyer, 1996c).)

We consider here any two  $2 \times 2$  pay-off matrices  $G^1$  and  $G^2$ . To state the results we need some notations. Let  $G^k = (G_{st}^k)$  for  $k = 1, 2$ ,  $s = 1, 2$ ,  $t = 1, 2$ , and denote:

- For  $k = 1, 2$ ,  $\Delta_k = G_{11}^k + G_{22}^k - G_{12}^k - G_{21}^k$ .
- Define  $\delta(G^1, G^2)$  by:

$$\delta(G^1, G^2) = \begin{cases} 1 & \text{If player I has a strategy which is optimal in } D(p) \forall p \in [0, 1] \\ 0 & \text{Otherwise} \end{cases}$$

- For  $p \in [0, 1]$  let  $(t_p, t'_p)$  be an optimal strategy of player II in  $D(p)$  (we use  $t_2$  for  $t_0$  since it is optimal in  $G^2$ ).
- For  $p \in [0, 1]$  (given  $t_1, t_2$  and  $t_p$ ) let  $q = q(p)$  be defined by  $t_p = qt_1 + q't_2$ .
- Finally we define

$$K(G^1, G^2, p) = \delta(G^1, G^2) |t_1 - t_2| \sqrt{ss'} |p\Delta_1 + p'\Delta_2|,$$

where  $s$  is the uniformly optimal move of player I in  $D(p)$  (relevant only if it exists and is unique i.e. when  $\delta(G^1, G^2) \neq 0$ .)

Our result can be now stated as follows:

**THEOREM 4.12.** *In any game with incomplete information on one side with full monitoring,  $K = \{1, 2\}$  and pay-off matrices  $G^1$  and  $G^2$ , there exists  $A > 0$  such that for all  $n$  and for all  $p \in [0, 1]$ :*

$$(31) \quad \left| v_n(p) - \mathbf{Cav} u(p) - K(G^1, G^2, p) \phi(q(p)) \frac{1}{\sqrt{n}} \right| \leq \frac{A}{n^{2/3}}$$

where  $\phi$  is the normal density function defined in theorem 4.1 p. 236.

**REMARK 4.4.** The theorem provides the highest order term (in  $n$ ) in the expansion of the error term  $v_n(p) - \mathbf{Cav} u(p)$  namely that of order  $1/\sqrt{n}$ : The coefficient of this term is a multiple of an appropriately scaled  $\phi$  function.

**REMARK 4.5.** Both the constant  $K(G^1, G^2, p)$  and the argument  $q(p)$  are, as they should be, invariant under addition of a constant to all pay-offs.

REMARK 4.6. Inequality (31) states that the error term is of order  $1/\sqrt{n}$  (i.e.  $K(G^1, G^2, p) \neq 0$ ) if and only if player I has the same strictly mixed optimal strategy in  $G^1$  and in  $G^2$  (and therefore in  $D(p)$  for all  $p \in [0, 1]$ ) while player II's optimal strategies are different in  $G^1$  and in  $G^2$ .

REMARK 4.7. A priori  $t_1$ ,  $t_2$  and  $s$  may not be unique which would mean that  $K(G^1, G^2, p)$  is not well defined. However observe that in a (non-constant)  $2 \times 2$  game player I has a strictly mixed uniformly optimal move  $s$  in  $D(p)$  only when player II has a unique optimal move in each game. If  $s$  is (mixed and) not unique, player II must have the same optimal pure move in both games and hence  $(t_1 - t_2) = 0$ . This means that whenever  $|t_1 - t_2| \neq 0$ ,  $s$  is unique and if  $\delta(G^1, G^2)\sqrt{ss'} \neq 0$ , then  $t_1$  and  $t_2$  are unique. Hence  $K$  is always well defined.

REMARK 4.8. The existence of a (unique)  $q = q(p)$  satisfying  $t_p = qt_1 + q't_2$  is part of the statement of the theorem i.e. whenever  $K(G^1, G^2, p) \neq 0$ ,  $q(p)$  is well defined.

REMARK 4.9. Note finally that in the example in the previous section we have:

$$\delta(G^1, G^2) = 1, s = 1/2, t_1 = 1/4, t_2 = 1/2, \Delta_1 = \Delta_2 = 8$$

and hence  $K(G^1, G^2, p) = 1$ . Also since

$$t_p = \frac{p + 2p'}{4} = pt_1 + p't_2$$

we have  $q(p) = p$  for all  $p \in [0, 1]$  and therefore  $K(G^1, G^2, p)\phi(q(p)) = \phi(p)$ .

Let us indicate now the main steps in the proof in order to see how the special example in theorem 4.1 p. 236 evolves to the general result in theorem 4.12 p. 244.

STEP 1. If the pay-off matrices are

$$G^1 = \begin{pmatrix} \theta'a & -\theta'a' \\ \theta a & \theta a' \end{pmatrix} \text{ and } G^2 = \begin{pmatrix} \theta'b & -\theta'b' \\ \theta b & \theta b' \end{pmatrix}$$

for  $a, b$  and  $\theta$  are in  $[0, 1]$ , then there exists  $A > 0$  such that for all  $n$  and for all  $p \in [0, 1]$ :

$$(32) \quad \left| \sqrt{n}v_n(p) - |a - b|\sqrt{\theta\theta'}\phi(p) \right| \leq C \frac{\ln n}{n}$$

- In these games

$$u(p) \equiv 0, \delta(G^1, G^2) = 1, s = \theta, t_1 = a', t_2 = b', \Delta_1 = \Delta_2 = 1$$

and so  $K(G^1, G^2, p) = |a - b|\sqrt{\theta\theta'}$ . Also the optimal move of player II in  $D(p)$  is  $t_p = pa' + p'b' = pt_1 + p't_2$  implying  $q(p) = p$ . So inequality (32) is in fact a special case of (31).

- This is a first step in generalising Example 4.1 p. 236. In that example  $a = 3/4$ ,  $b = 1/2$ ,  $\theta = 1/2$  and all pay-offs are multiplied by 8.
- The proof of Step 1 is basically the same as the first part of the proof of theorem 4.1 p. 236 namely the proof of inequality (27) p. 243.

STEP 2. Extend the result in Step 1 by applying two standard transformations to the pay-off matrices; adding a constant to all pay-offs (this adds the same constant to all value functions of the games) and multiplying each matrix by a positive constant adds

the same constant to all value functions of the games: For  $\alpha > 0$ ,  $\beta > 0$ , if  $\tilde{G}^1 = \alpha G^1$  and  $\tilde{G}^2 = \beta G^2$  then  $\forall p \in [0, 1]$  and  $\forall n$

$$(33) \quad \tilde{u}(p) = (p\alpha + p'\beta)u\left(\frac{p\alpha}{p\alpha + p'\beta}\right),$$

and similar transformations hold for  $\text{Cav } u(p)$  and  $v_n(p)$ .

The proof of (33) is straightforward. The consequence of this step is that if a game satisfies (31) then so does any game obtainable from it by standard transformations. In particular: all games obtainable from the games considered in Step 1 satisfy (31).

STEP 3. If  $u(p) = pv(G^1) + p'v(G^2)$  then:

- If  $K(G^1, G^2, p) \neq 0$  for some  $p$ ,  $0 < p < 1$ , then the game is obtainable from the class of games in Step 1 by standard transformations. Hence it satisfies (31).
- If  $K(G^1, G^2, p) = 0$  for all  $p \in [0, 1]$ , then there exists  $A > 0$  such that

$$v_n(p) - u(p) \leq \frac{A}{n} \min(p, p')$$

STEP 4. If  $K(G^1, G^2, p) \neq 0$  for some  $p \in [0, 1]$  then  $u$  is linear.

By Step 3 all games with linear  $u$  satisfy (31) and by Step 4, whenever  $u$  is not linear then  $K(G^1, G^2, p) \equiv 0$ , so the last step is:

STEP 5. If  $u$  is not linear in  $[0, 1]$  then

$$\exists A > 0 \text{ such that } \forall n \text{ and } \forall p \in [0, 1] \quad v_n(p) - \text{Cav } u(p) \leq \frac{A}{n^{2/3}}$$

The proof of this step is rather lengthy and involves some non-trivial arguments. However all these concern terms of order lower than  $1/\sqrt{n}$ . As far as the  $1/\sqrt{n}$  term and the normal function  $\phi$  are concerned these are completely covered by Steps 1 to 4.

## 5. The speed of convergence of $v_n$

In this section we study more closely the nature of the uniform convergence of the values  $v_n(p)$  of the finite games  $\Gamma_n(p)$  to  $v_\infty(p)$ , the value of  $\Gamma_\infty(p)$ . For the case of full monitoring, by theorem 2.10 p. 189, the speed of this convergence is bounded from above by  $O(1/\sqrt{n})$ . (As we shall prove in theorem 6.2 p. 304 this also true for the case of incomplete information on two sides.) This is the least upper bound for that case since in theorem 4.1 p. 236 we provided a game with  $v(p) \equiv 0$  and  $\lim_{n \rightarrow \infty} \sqrt{n}v_n(p) = \phi(p)$ .

**5.a. State independent signalling.** Without the full monitoring assumption,  $O(1/\sqrt{n})$  is no longer an upper bound for the speed of convergence of  $v_n(p)$ . The next level of generality is the case of state independent signals i.e. the signalling matrices  $Q^k$  are the same for all  $k \in K$ . For this case it will be proved (6 p. 304) that even for incomplete information on two sides the speed of convergence is bounded above by  $O(1/\sqrt[3]{n})$ . We shall now prove that this is in fact the least upper bound.

EXAMPLE 5.1. Let  $K = \{1, 2\}$ . The pay-off matrices  $G^1$ ,  $G^2$  and the signalling matrices  $Q^1$ ,  $Q^2$  (to player II) are given by:

$$G^1 = \begin{pmatrix} 1 & 3 & -1 \\ 1 & -3 & 1 \end{pmatrix} \quad G^2 = \begin{pmatrix} 1 & 2 & -2 \\ 1 & -2 & 2 \end{pmatrix} \quad Q^1 = Q^2 = \begin{pmatrix} a & c & d \\ b & c & d \end{pmatrix}$$

Observe that deleting the left strategy of player II and changing the signalling matrices so as to provide full monitoring we obtain the game in example 4.1 p. 236 for which we have just proved  $\delta_n(p) \approx (1/\sqrt{n})$ . However, in our present example player II is not informed of the last move of his opponent unless he chooses his left strategy which is strictly dominated in terms of pay-offs. In other words player II has to pay 1 unit whenever he wants to observe his opponent's move. Since observing the moves of the informed player is his only way to collect information about the state  $k$  it is not surprising that his learning process will be slower and/or more costly than in example 4.1 p. 236. This will mean a slower rate of convergence of  $v_n$  to  $v_\infty$ .

Keeping the same notations as in example 4.1 p. 236, the set of non-revealing moves is readily seen to be

$$\text{NR}(p) = \{(x, y) \mid x = y ; 0 \leq x \leq 1\} \text{ if } 0 < p < 1$$

and the non-revealing game is therefore

$$D(p) = \begin{pmatrix} 1 & 3p + 2p' & -p - 2p' \\ 1 & -3p - 2p' & p + 2p' \end{pmatrix}$$

with value  $u(p) = 0$  for  $0 < p < 1$ . Since also  $v(G^1) = v(G^2) = 0$  we conclude that  $u$ , and therefore  $v_\infty$ , are the same as in Example 4.1 p. 236 namely:

$$v(p) = \text{Cav } u(p) = u(p) = 0 \quad \forall p \in [0, 1]$$

**PROPOSITION 5.1.** *For the game in example 5.1 the following holds*

$$\forall n \geq 1 \text{ and } \forall p \in [0, 1] : \quad v_n(p) \geq \frac{p(1-p)}{\sqrt[3]{n}}$$

**PROOF.** The recursive formula from lemma 4.2 p. 236 yields here:

$$(1) \quad \begin{aligned} v_{n+1} &= \frac{1}{n+1} \max_{0 \leq x, y \leq 1} \min_{0 \leq \varepsilon \leq 1} \{ \varepsilon \\ &\quad + (1-\varepsilon) \min[3p(x-x') + 2p'(y-y'), p(x'-x) + 2p'(y'-y)] \\ &\quad + n(1-\varepsilon)v_n(p) + n\varepsilon[\bar{x}v_n(p_T) + \bar{x}'v_n(p_B)] \} \end{aligned}$$

where  $\bar{x}$  is given by (3) p. 237 and  $p_T, p_B$  are given by (4) p. 237.

For  $n = 1$ ,  $v_1(p) = \min(p, p') \geq pp'$  in accordance with the proposition. We proceed by induction on  $n$ : Assume the inequality holds for  $n$  and for all  $p \in [0, 1]$ . In the recursive formula, restrict  $x$  and  $y$  by  $p(x-x') = p'(y-y') \stackrel{\text{def}}{=} \delta$  we obtain, by the induction hypothesis:

$$(2) \quad v_{n+1}(p) \geq \max_{|\delta| \leq \min(p, p')} \min_{0 \leq \varepsilon \leq 1} \frac{1}{n+1} [\varepsilon + pp'n^{2/3} + (1-\varepsilon)\delta - \varepsilon n^{2/3}\delta^2]$$

Since the function  $[\dots]$  is concave in  $\delta$  (and linear in  $\varepsilon$ ) we apply the minmax theorem to first maximise on  $\delta$ . The maximum is at  $\delta = \min[pp', (1-\varepsilon)/(2\varepsilon n^{2/3})]$ . Letting  $A = pp'$  and noticing that  $(n+1)^{2/3} - n^{2/3} \leq 2/(3n^{1/3})$  it then suffices to prove that:

$$n \geq 1, 0 \leq \varepsilon \leq 1, 0 \leq A \leq \frac{1}{4}, \text{ and } \delta = \min[A, \frac{1-\varepsilon}{2\varepsilon n^{2/3}}] \text{ imply :}$$

$$\varepsilon + (1-\varepsilon)\delta - \varepsilon n^{2/3}\delta^2 \geq \frac{2A}{3n^{1/3}}$$

If  $\delta = A$ , the left hand side is linear (and increasing) in  $\varepsilon$ , and the inequality obviously holds for  $\varepsilon = 0$ . So it suffices to check it at the maximum value of  $\varepsilon$ , where

$\delta = (1 - \varepsilon)/(2\varepsilon n^{2/3})$ . Thus we may anyway assume that  $\delta = (1 - \varepsilon)/(2\varepsilon n^{2/3}) \leq A$  and we have then to show that  $0 < \varepsilon$  and  $A \leq 1/4$  imply

$$\varepsilon + \frac{(1 - \varepsilon)^2}{4\varepsilon n^{2/3}} \geq \frac{2A}{3n^{1/3}},$$

and it is clearly enough to prove this for  $A = 1/4$ . Thus letting  $\xi = 1/n^{1/3}$ , we have to prove that  $\xi \leq 1, 0 < \varepsilon \leq 1$  imply  $\frac{(1-\varepsilon)^2}{\varepsilon}\xi^2 - \frac{2}{3}\xi + 4\varepsilon \geq 0$ . The unconstrained minimum in  $\xi$  is at  $\xi_0 = \frac{\varepsilon}{3(1-\varepsilon)^2}$ , where the value of the function is  $\varepsilon(4 - \frac{1}{9(1-\varepsilon)^2})$ . Thus the inequality is proved for  $\varepsilon \leq 5/6$ . For  $5/6 < \varepsilon \leq 1$  the minimum is at  $\xi = 1$  where the inequality is obviously satisfied for these values of  $\varepsilon$ . This completes the proof of the proposition. ■

Combined with the result of theorem 6.1 p. 304 we conclude (again since  $v_n(p) = \delta_n(p)$ ):

COROLLARY 5.2. *For the game in example 5.1 p. 246 the speed of convergence of  $v_n$  is  $(1/\sqrt[3]{n})$ .*

**5.b. State dependent signalling.** In the previous sections we were able to determine the least upper bound for the speed of convergence of  $v_n(p)$  for the full monitoring case and for the case of state independent signalling. The results are much less complete for games with state dependent signalling. Here we have only partial results for the case of two states and common signals to both players. So from here on in this section we assume  $K = \{1, 2\}$ .

Recall that  $\text{NR}(p)$  depends only on the support of  $p$ . So it is the same for  $0 < p < 1$ . Denote this set simply by  $\text{NR}$ .

THEOREM 5.3. *If  $\text{NR} = \emptyset$  then for some  $C_1 > 0$*

$$v_n(p) \leq \text{Cav } u(p) + C_1 \sqrt[4]{pp'/n}$$

for all  $p \in [0, 1]$  and for all  $n$ .

PROOF. In establishing upper bounds for  $v_n$  we first modify the game in favour of player I so as to have  $\mathcal{H}_n^{\text{II}} \subseteq \mathcal{H}_n^{\text{I}}$ , (cf. sect. 3.c p. 195), hence the posterior probabilities  $p_n$  (given  $\mathcal{H}_n^{\text{II}}$ ) are common knowledge. The proof is carried in four steps some of which will be used also in the next theorem.

STEP 1.  $v_1(p) \leq \text{Cav } u(p) + 2C_1 pp'$ .

This is a consequence of the Lipschitz property of  $v_1(p)$ :

$$(3) \quad v_1(p) \leq \min[v_1(0) + Cp, v_1(1) + Cp'] \leq pv_1(1) + p'v_1(0) + 2C_1 pp' \\ = pu_1(1) + p'u_1(0) + 2C_1 pp' \leq \text{Cav } u(p) + 2C_1 pp'$$

STEP 2. *For any strategy  $\sigma$  of player I, if at stage  $m$  player II plays his uniformly mixed move  $y^0$  (i.e.  $y^0(t) = 1/\#T$ ,  $\forall t \in T$ ) then*

$$(4) \quad \mathbb{E}(|p_{m+1} - p_m| \mid \mathcal{H}_m^{\text{II}}) = p_m p'_m d(x),$$

where

$$d(x) = \frac{1}{\#T} \sum_t \sum_b |x^1 Q_t^1(b) - x^2 Q_t^2(b)|.$$

Here  $x = (x^1, x^2)$  is the strategy of player I at stage  $m$  and  $Q_t^k$  is the  $t^{\text{th}}$  column of  $Q^k$ . Note that  $d(x) = 0$  if and only if  $x \in \text{NR}$ .

To see this observe that when  $x$  and  $t \in T$  are played, the probability of a signal  $b \in B$  is

$$P(b) = p_m x^1 Q_t^1(b) + p'_m x^2 Q_t^2(b)$$

In this event, the new conditional probability of  $\{k = 1\}$  will be:

$$p_{m+1}(b) = \frac{p_m x^1 Q_t^1(b)}{p_m x^1 Q_t^1(b) + p'_m x^2 Q_t^2(b)}.$$

STEP 3. If  $\text{NR} = \emptyset$  then there exists a constant  $\eta > 0$  such that  $d(x) \geq \eta; \forall x \in X^2$ .

In fact  $d(x)$  is a non-negative continuous function of  $x$  on the compact set  $X^2$  and therefore attains its minimum there. This minimum cannot be 0 since this would imply  $\text{NR} \neq \emptyset$ .

STEP 4. In  $\Gamma_n(p)$ , given any strategy  $\sigma$  of player I, at each stage  $m$  let player II play  $y^0$  with probability  $\varepsilon_n$  (to be determined later) and with probability  $1 - \varepsilon_n$  play the strategy that guarantees (Step 1)  $\mathbb{E}(g_m | \mathcal{H}_m^{\text{II}}) \leq \text{Cav } u(p_m) + 2C p_m p'_m$ . By Steps 2 and 3,

$$\mathbb{E}(|p_{m+1} - p_m| | \mathcal{H}_m^{\text{II}}) \geq \eta \varepsilon_n p_m p'_m.$$

Hence we have for all  $m$ :

$$\begin{aligned} \mathbb{E}(g_m | \mathcal{H}_m^{\text{II}}) &\leq \text{Cav } u(p_m) + \varepsilon_n C + 2(1 - \varepsilon_n) C p_m p'_m \\ &\leq \text{Cav } u(p_m) + \varepsilon_n C + \frac{2C}{\eta \varepsilon_n} \mathbb{E}(|p_{m+1} - p_m| | \mathcal{H}_m^{\text{II}}) \end{aligned}$$

Taking average over the  $n$  stages and expectation on  $\mathcal{H}_m^{\text{II}}$  we obtain (using lemma 2.1 p. 186):

$$\bar{\gamma}_n \leq \text{Cav } u(p) + \varepsilon_n C + \frac{2C}{\eta \varepsilon_n} \sqrt{\frac{pp'}{n}}.$$

Choosing  $\varepsilon_n = \sqrt{2/\eta} \sqrt[4]{pp'/n}$  we conclude, using  $C_1 = C \sqrt{8/\eta}$ , that

$$(5) \quad v_n(p) \leq \text{Cav } u(p) + C_1 \sqrt[4]{pp'/n}, \quad \blacksquare$$

To bound the error term for the case in which  $\text{NR} \neq \emptyset$  we need the following:

LEMMA 5.4. Let  $\Delta$  be a non-empty polyhedron in  $\mathbb{R}^m$  and let  $L_i$  with  $i = 1, \dots, d$  be linear functionals. Define

$$\text{NR} = \{x \in \Delta \mid L_i x = 0; i = 1, \dots, d\}$$

If  $\text{NR} \neq \emptyset$ , then there exists a constant  $\beta > 0$  such that for all  $x \in \Delta$

$$\max_i |L_i x| \geq \beta \|x - \text{NR}\|$$

(Here  $\|x - \text{NR}\| = \min_{x' \in \text{NR}} \|x - x'\|$ ,  $\|\cdot\|$  being a norm in  $\mathbb{R}^m$ .)

PROOF. This is a consequence of ex.I.3Ex.4q p.30 which states that if  $L$  is an affine map from  $\Delta$  to  $\mathbb{R}^d$ , then  $L^{-1}(x)$  is a Lipschitz function (in the Hausdorff distance) of  $x \in L(\Delta)$  (where distance is derived from the given norms).  $\blacksquare$

THEOREM 5.5. For any  $\Gamma_n(p)$  there exist a constant  $C_2$  such that for all  $n$  and for all  $p \in [0, 1]$

$$v_n(p) \leq \text{Cav } u(p) + C_2 \sqrt[6]{pp'/n}.$$

PROOF. If  $\text{NR} = \emptyset$  the claimed inequality follows from theorem 5.3 p. 248. Assume therefore that  $\text{NR} \neq \emptyset$  then by lemma 5.4 p. 249 there exists  $\beta > 0$  such that for all  $x \in X^2$

$$(6) \quad d(x) \geq \beta \|x - \text{NR}\|$$

Notice that, given any strategy of player I, player II guarantees, by playing at stage  $m$  an optimal move in  $D(p_m)$ :

$$\mathbb{E}(g_m | \mathcal{H}_m^{\text{II}}) \leq \text{Cav } u(p) + C \|x - \text{NR}\| \leq \text{Cav } u(p) + Cd(x)/\beta.$$

Combined with Step 1 of theorem 5.3 this yields: Given  $\sigma$ , player II can guarantee at each stage  $m$

$$\begin{aligned} \mathbb{E}(g_m | \mathcal{H}_m^{\text{II}}) &\leq \text{Cav } u(p_m) + C \min(2p_m p'_m, d(x)/\beta) \\ &\leq \text{Cav } u(p_m) + C \sqrt{2p_m p'_m d(x)/\beta}. \end{aligned}$$

The proof now proceeds as that of the previous theorem: At each stage  $m$  player II plays with probability  $\varepsilon_n$  the uniform move  $y^0$  and with probability  $(1 - \varepsilon_n)$  he plays the above move. This guarantees (using Step 2)

$$\begin{aligned} \mathbb{E}(g_m | \mathcal{H}_m^{\text{II}}) &\leq \text{Cav } u(p_m) + C\varepsilon_n + (1 - \varepsilon_n)C \sqrt{2p_m p'_m d(x)/\beta} \\ &\leq \text{Cav } u(p_m) + C\varepsilon_n + C \sqrt{\frac{2}{\beta\varepsilon_n}} \sqrt{\mathbb{E}(|p_{m+1} - p_m| | \mathcal{H}_m^{\text{II}})} \end{aligned}$$

Taking averages over all  $n$  stages and expectation on  $\mathcal{H}_m^{\text{II}}$  we get (using the Cauchy-Schwartz inequality and theorem 2.10 p. 189):

$$\begin{aligned} \bar{\gamma}_n &\leq \text{Cav } u(p) + C\varepsilon_n + \frac{C}{\sqrt{n}} \sqrt{\frac{2}{\beta\varepsilon_n}} \sqrt{\sum_1^n \mathbb{E}(|p_{m+1} - p_m| | \mathcal{H}^{\text{II}})} \\ &\leq \text{Cav } u(p) + C\varepsilon_n + \sqrt{\frac{2C}{\beta\varepsilon_n}} \sqrt[4]{\frac{pp'}{n}}. \end{aligned}$$

Choosing  $\varepsilon_n = \sqrt[3]{4/\beta} \sqrt[6]{pp'/n}$  we conclude

$$v_n(p) \leq \text{Cav } u(p) + C_2 \sqrt[6]{pp'/n},$$

where  $C_2 = 4C/\sqrt[3]{2\beta}$ , completing the proof. ■

**5.c. Games with error term  $\approx (\ln n)/n$ .** We start by proving a lemma which provides a sufficient condition for the error term to be bounded by  $(\ln n)/n$ . For the sake of simplicity only, the lemma is stated and proved for two states. It can be done for any finite set of states  $K$  (cf. ex. VEx.10c p. 258).

**LEMMA 5.6.** *Let  $\Gamma(p)$  be a game with full monitoring and two states. If  $u(p)$  is twice differentiable and  $\exists \eta > 0$  such that  $\forall p \in [0, 1]$ ,  $u''(p) < -\eta$  then for some constant  $A > 0$ ,  $\delta_n(p) \leq A(\ln n)/n$  for all  $p \in [0, 1]$ .*

PROOF. Note first that if player II plays an optimal strategy in  $D(p)$  he guarantees a first stage pay-off of at most  $u(p) + 2C \mathbb{E} |p_2 - p|$  (where  $p_2$  is the conditional probability of  $\{k = 1\}$  after the first stage). This follows from lemma 2.6 p. 188 and lemma 2.5 p. 187 (and can be easily verified directly for this simple case.) In particular this implies that  $\forall x \in X^2$ .

$$(7) \quad \min_{t \in T} (px^1 G_t^1 + p'x^2 G_t^2) \leq u(p) + 2C \mathbb{E} |p_2 - p|$$

For any (finite) distribution of  $p_2$  (with  $\mathbb{E}(p_2) = p$ ) we have

$$(8) \quad v_\infty(p) - \mathbb{E} v_\infty(p_2) \geq \frac{\eta}{2} \mathbb{E} |p_2 - p|^2$$

This follows by using Taylor's expansion of  $v_\infty(q)$  for each value  $q$  of  $p_2$ .

$$\begin{aligned} v_\infty(q) &= v_\infty(p) + (q - p)v'_\infty(p) + \frac{1}{2}(q - p)^2 v''_\infty(\xi) \\ &\leq v_\infty(p) + (q - p)v'_\infty(p) - \frac{\eta}{2}(q - p)^2, \end{aligned}$$

where  $\xi = p + \theta(q - p)$ ,  $0 \leq \theta \leq 1$ . Taking expectation on  $p_2$  we obtain inequality (8).

Since  $\sum_{r=1}^n \frac{1}{r} = O(\ln n)$ , the proof of the lemma will follow if we prove the existence of  $A > 0$  such that for all  $p \in [0, 1]$  and all  $n$ , the functions  $V_n(p) = nv_n(p)$  satisfy:

$$(9) \quad V_n(p) \leq nv_\infty(p) + A \sum_{r < n} r^{-1}$$

We prove this by induction on  $n$ : Choose  $A_1$  large enough to make (9) true for  $n = 1$  and let  $A = \max(A_1, C^2/(2\eta))$ . By the recursive formula (1) p. 236, inequality (7) p. 250 and the induction hypothesis we have:

$$\begin{aligned} V_{n+1}(p) &\leq \max_{x \in X^2} \{u(p) + 2C \mathbb{E} |p_2 - p| + \mathbb{E} V_n(p_2)\} \\ &\leq \max_{x \in X^2} \{v_\infty(p) + 2C \mathbb{E} |p_2 - p| + n \mathbb{E}(v_\infty(p_2)) + A \sum_{r < n} r^{-1}\} \\ &= \max_{x \in X^2} \{(n+1)v_\infty(p) + A \sum_{r < n} r^{-1} + 2C \mathbb{E} |p_2 - p| - n(v_\infty(p) - \mathbb{E} v_\infty(p_2))\} \end{aligned}$$

and by (8):

$$\begin{aligned} V_{n+1}(p) &\leq (n+1)v_\infty(p) + A \sum_{r < n} r^{-1} + \max_{x \in X^2} \left\{ 2C \mathbb{E} |p_2 - p| - \frac{n\eta}{2} \mathbb{E} |p_2 - p|^2 \right\} \\ &\leq (n+1)v_\infty(p) + A \sum_{r < n} r^{-1} + \max_{x \in X^2} \left\{ 2C \mathbb{E} |p_2 - p| - \frac{n\eta}{2} (\mathbb{E} |p_2 - p|)^2 \right\} \end{aligned}$$

Since in the right hand side only  $p_2$  depends on  $x$ , maximisation with respect to  $\mathbb{E} |p_2 - p|$  yields

$$\begin{aligned} V_{n+1}(p) &\leq (n+1)v_\infty(p) + A \sum_{r < n} r^{-1} + \frac{2C^2}{\eta n} \\ &\leq (n+1)v_\infty(p) + A \sum_{r < n} r^{-1}, \end{aligned}$$

Completing the proof of lemma 5.6. ■

For instance, in example 2.1 p. 185 observe that  $v_\infty(p) = u(p) = pp'$  implying  $v''(p) = -2 \quad \forall p \in [0, 1]$ . Consequently, the error term of this game is bounded by  $(O(\ln n)/n)$ . We shall now prove that  $(O(\ln n)/n)$  is also a lower bound.

**PROPOSITION 5.7.** *In the game in example 2.1 p. 185*

$$\delta_n(p) = v_n(p) - v_\infty(p) \approx \left( \frac{\ln n}{n} \right)$$

**PROOF.** Letting  $(\alpha, \alpha')$   $(\beta, \beta')$  be the first stage mixed moves in  $G^1$  and  $G^2$  respectively, The recursive formula (1) p. 236 gives for this game:  $v_{n+1}(p) =$

$$\max_{\substack{0 \leq \alpha \leq 1 \\ 0 \leq \beta \leq 1}} \frac{1}{n+1} \left\{ \min(\alpha p, \beta' p') + n \left[ (\alpha p + \beta p') v_n \left( \frac{\alpha p}{\alpha p + \beta p'} \right) + (\alpha' p + \beta' p') v_n \left( \frac{\alpha' p}{\alpha' p + \beta' p'} \right) \right] \right\}$$

Letting  $V_n(p) = nv_n(p)$ ,  $x = \alpha p$ ,  $y = \beta' p'$ , we rewrite this as:

$$V_{n+1}(p) = \max_{\substack{0 \leq x \leq p \\ 0 \leq y \leq p'}} \left\{ \min(x, y) + (p' + x - y)V_n\left(\frac{x}{p' + x - y}\right) + (p - x + y)V_n\left(\frac{p - x}{p - x + y}\right) \right\}$$

We shall prove that for some large  $N$ , the inequality

$$V_n(p) \geq C_n pp' - D_n \stackrel{\text{def}}{=} F_n(p)$$

is satisfied for all  $n \geq N$  with  $C_n = n + \frac{1}{4} \ln n$  and  $B_n \stackrel{\text{def}}{=} D_{n+1} - D_n = (\ln n)/(8n)^2$ . Since  $B_n$  is summable this will imply  $V_n(p) \geq C_n pp' - D$  for some constant  $D$  for  $n \geq N$  and hence for all  $n$ , by choosing  $D$  sufficiently large. This will establish the required lower bound  $v_n(p) - v_\infty(p) \geq (pp'/4)(\ln n)/n - D/n$ .

Using the recursive formula, the induction step is to prove that for  $n \geq N$  and  $p$  with  $0 \leq p \leq 1/2$  (by the symmetry of  $u$  and  $v_n$  about  $p = 1/2$ ):

$$F_{n+1}(p) = \max_{\substack{0 \leq x \leq p \\ 0 \leq y \leq p'}} \left\{ \min(x, y) + (p' + x - y)F_n\left(\frac{x}{p' + x - y}\right) + (p - x + y)F_n\left(\frac{p - x}{p - x + y}\right) \right\}$$

Replacing  $F$  by its formula and setting  $x = y = pp'(1 + \delta)$  with  $\delta = \min\left(\frac{p}{p'}, \frac{1}{2C_n}\right)$ , we want to prove:

$$C_{n+1}pp' - D_{n+1} \leq C_n pp' - D_n + pp'(1 + \delta - C_n \delta^2)$$

Using  $\ln(1 + \frac{1}{n}) \leq \frac{1}{n}$  and  $C_n \geq n$  it is enough to prove (setting  $z = p/p'$ ):

$$\frac{1}{4n} - \frac{B_n}{pp'} \leq \begin{cases} z - C_n z^2 & \text{if } z \leq (2C_n)^{-1} \\ (4C_n)^{-1} & \text{if } z \geq (2C_n)^{-1} \end{cases}$$

- For the first case it will suffice to have  $\frac{1}{4n} + C_n z^2 \leq \frac{B_n}{z}$  and by the monotonicity it is enough if it is satisfied for  $z = 1/(2C_n)$ , that is if  $1 + C_n/n \leq 8B_n C_n^2$ . This is clearly true for  $n$  sufficiently large since the left hand side is bounded and the right hand side tends to infinity.
- For the second case it is enough to check the inequality for  $p = 1/2$ , i.e.  $16B_n \geq 1/n - 1/C_n$ . In fact:

$$\frac{1}{n} - \frac{1}{C_n} = \frac{\ln n}{4nC_n} \leq \frac{\ln n}{4n^2} = 16B_n$$

This completes the proof of the proposition. ■

### Exercises

#### 1. Subadditivity of $V_n = nv_n$ .

- a. Using the argument in the proof of theorem 3.1 p. 191 prove that for any  $n$  and  $m$

$$V_{n+m} \leq V_n + V_m$$

In particular  $v_{2n} \leq v_n$ .

- b. Deduce that  $v_n$  converges.

**2. Optimal strategy for player II: Explicit construction.** In the case of full monitoring, prove that the following is an optimal strategy for player II in  $\Gamma_\infty(p_0)$ :

Let  $c \in \mathbb{R}^K$  be a supporting hyperplane to the graph of  $\text{Cav } u$  at  $p_0$ , and  $D = \{x \in \mathbb{R}^K \mid x \leq c\}$ . At stage  $n+1$ ,  $n = 1, 2, \dots$ , compute the average vector pay-off  $\bar{g}_n$  up to that stage. If  $\bar{g}_n \in D$ , play arbitrarily. If  $\bar{g}_n \notin D$ , let  $\xi \in D$  be the closest point to  $\bar{g}_n$  in  $D$ , compute  $q = (\bar{g}_n - \xi) / \|\bar{g}_n - \xi\| \in \Delta(K)$  and play an optimal mixed move in  $D(q)$ .

HINT. Use sect. 4 p. 102.

**3. Semi-algebraicity of  $v_\infty$ .** Prove that  $v_\infty(p)$  is semi-algebraic.

HINT.  $u(p)$  is piecewise rational,  $v_\infty(p)$  is linear on  $\{p \mid \text{Cav } u(p) > u(p)\}$ , and the boundary of  $\{p \mid \text{Cav } u(p) = u(p)\}$  is determined by polynomial equations.

**4. Non-existence of Markovian equilibrium in  $\Gamma_n(p)$ .** Consider the game with two states, full monitoring and pay-off matrices:

$$G^1 = \begin{pmatrix} 2 & 0 \\ 0 & -1 \end{pmatrix} \quad G^2 = \begin{pmatrix} -1 & 0 \\ 0 & 2 \end{pmatrix}$$

a. Prove that

$$\begin{aligned} V_1(p) &= 2 \min(p, p') = 2(p \wedge p') \\ V_2(p) &= \min\{3(p \wedge p'), 1 + (p \wedge p')/2\} \\ V_3(p) &= \begin{cases} \frac{1}{2} + \frac{9}{4}p & \text{for } \frac{2}{5} \leq p \leq \frac{6}{13} \\ \frac{4}{5}(1 + 2p) & \text{for } \frac{6}{13} \leq p \leq \frac{1}{2} \end{cases} \end{aligned}$$

where  $(p, p')$  is the probability distribution on the states and  $V_n(p) = nv_n(p)$ .

b. Prove that for any optimal strategy of player I in  $\Gamma_2(\frac{1}{2})$  the posterior probability after stage 1 equals  $\frac{1}{2}$  a.s. Conclude that (for any optimal strategy of player I) any Markovian strategy of player II (i.e. a strategy in which each stage behaviour is the same for two histories leading to the same posterior probability), is history-independent.

HINT. In any of his optimal strategies, player I plays  $(\frac{1}{2}, \frac{1}{2})$  at the first stage in both states.

c. Prove that with history-independent strategies, player II can guarantee not more than  $2V_1(1/2) = 2$  which is worse than  $V_2(1/2)$ . Conclude that for any optimal strategy of player I, no Markovian strategy of player II is a best reply. Consequently, there does not exist a pair of Markovian optimal strategies in  $\Gamma_2(\frac{1}{2})$ .

d. Prove that the same negative conclusion is valid also for  $\Gamma_3(p)$  for the whole interval  $\frac{2}{5} < p < \frac{6}{13}$ .

HINT. Any optimal strategy of player I leads with positive probability, after stage 1, to the posterior probability  $\frac{1}{2}$ .

**5. Proving  $\lim v_n(p) = \text{Cav } u(p)$  by the recursive formula.** The recursive formula (lemma 4.2 p. 236) which we used mainly to study the speed of convergence of  $v_n$  can be used also to prove the convergence itself namely to prove  $\lim_{n \rightarrow \infty} v_n(p) = \text{Cav } u(p)$ . In the proof outlined here we make use of the following result (Kohlberg and Neyman, 1981):

Let  $X$  be a Banach space with norm  $\|\cdot\|$  and denote its dual by  $X^*$ . Let  $\Psi: X \rightarrow X$  be a non-expansive mapping i.e.

$$(10) \quad \|\Psi x - \Psi y\| \leq \|x - y\| \quad \text{for all } x \text{ and } y \text{ in } X.$$

Let  $\rho = \inf_{x \in X} \|\Psi x - x\|$  then:

$$(11) \quad \forall x \in X, \lim_{n \rightarrow \infty} \left\| \frac{\Psi^n x}{n} \right\| = \rho$$

$$(12) \quad \forall x \in X \exists f_x \in X^*: \|f_x\| = 1 \text{ and } f_x(\Psi^n x - x) \geq n\rho \quad \forall n$$

To use this result let  $X$  be the set of continuous functions on  $\Delta(K)$  endowed with the maximum norm and let  $\Psi$  be the mapping defined (cf. 3.b p. 157) by the recursive formula (lemma 4.2 p. 236) (without the factor  $1/(n+1)$ ) i.e. for  $w \in X$ :

$$(13) \quad (\Psi w)(p) = \max_{x \in X^K} \left( \min_t \sum_k p^k x^k G_t^k + \mathbb{E}_\sigma w(p_s) \right)$$

where  $X^K$  and  $p_s$  are as in lemma 4.2.

a. Prove that  $\lim v_n(p)$  exists by proving the following:

i. If we let  $v_0 = \mathbf{0}$  (the 0-function in  $X$ ) then  $\forall n \geq 0$ ,  $\Psi^n \mathbf{0} = nv_n$ .

ii. The mapping  $\Psi$  is non-expansive.

iii. By (11),  $\max_p |v_n(p)|$  converges as  $n \rightarrow \infty$ . By adding the same constant to all  $G^k$  we may assume w.l.o.g. that  $v_n(p) \geq 0 \forall p \in \Delta(K)$  hence  $\lim_{n \rightarrow \infty} \max_p v_n(p)$  exists.

iv. Adding a constant  $-\alpha^k$  to all entries of  $G^k$  we have that

$$\lim_{n \rightarrow \infty} \max_p (v_n(p) - \langle \alpha, p \rangle) \text{ exists } \forall \alpha \in \mathbb{R}^K.$$

So if we denote

$$\phi_n(\alpha) = \max_p (v_n(p) - \langle \alpha, p \rangle)$$

then, for all  $\alpha$ , the sequence  $\phi_n(\alpha)$  converges, to say  $\phi(\alpha)$ . However  $\phi_n(\alpha)$  are all Lipschitz with constant 1 hence the convergence  $\phi_n(\alpha) \rightarrow \phi(\alpha)$  is uniform and  $\phi(\alpha)$  is continuous.

v. Show that  $v_n(p) = \min_\alpha (\phi_n(\alpha) - \langle \alpha, p \rangle)$ .

HINT. The inequality  $\leq$  follows from the definition of  $\phi_n(\alpha)$ . As for the other direction, since  $v_n$  is concave, let  $\alpha_0$  be a supporting hyperplane to  $v_n$  at  $p_0$  then  $\phi_n(\alpha_0) = 0$  and

$$v_n(p_0) = \langle \alpha_0, p_0 \rangle = \phi_n(\alpha_0) + \langle \alpha_0, p_0 \rangle \geq \min_\alpha (\phi_n(\alpha) + \langle \alpha, p_0 \rangle)$$

vi. Conclude that:

$$\lim_{n \rightarrow \infty} v_n(p) = \min_\alpha (\phi(\alpha) - \langle \alpha, p \rangle) \stackrel{\text{def}}{=} v_\infty(p)$$

Since  $v_n$  are all Lipschitz with the same constant, the convergence is uniform and  $v_\infty(p)$  is concave.

b. Prove that  $v_\infty(p) = \text{Cav } u(p)$  by showing:

i. Since  $v_\infty$  is concave and  $v_\infty(p) \geq u(p)$  (this holds for all  $v_n$ ) we have  $v_\infty(p) \geq \text{Cav } u(p)$ . Therefore to complete the proof it is enough to show that  $v_\infty(p) = u(p)$  at each point  $p$  of strict concavity of  $v_\infty$  (i.e. at the extreme points of the epigraph of  $v_\infty$ ,  $\{(p, x) \mid x \leq v_\infty(p)\}$ ).

ii. At each point  $p_0$  of strict concavity of  $v_\infty$  there is a supporting hyperplane  $\alpha$  such that the maximum of  $v_\infty(p) + \langle \alpha, p \rangle$  is attained only at  $p_0$ . Adding  $\alpha^k$  to all entries of  $G^k$ , the  $u$  function becomes  $u(p) + \langle \alpha, p \rangle$ , the  $v_\infty$  function becomes  $v_\infty(p) + \langle \alpha, p \rangle$  and  $p_0$  is the only maximum of the new  $v_\infty$ .

iii. Take  $x = \mathbf{0}$  in (12) and let  $\mu$  be a regular Borel measure of norm 1 on  $\Delta(K)$  representing  $f_{\mathbf{0}}$  i.e.

$$f_{\mathbf{0}}(w) = \int w(p)\mu(dp) \quad \forall w \in X.$$

So by (12) and VEx.5ai:

$$f_{\mathbf{0}}(\Psi^n \mathbf{0} - \mathbf{0}) = \int nv_n(p)\mu(dp) \geq n\rho$$

Hence

$$(14) \quad \int v_\infty(p)\mu(dp) \geq \rho$$

But by (11) and VEx.5ai and VEx.5bii:

$$\rho = \lim_{n \rightarrow \infty} \left\| \frac{1}{n} \Psi^n \mathbf{0} \right\| = \|v_\infty\| = v(p_0)$$

Since  $p_0$  is the unique maximum of  $v_\infty$ , inequality (14) implies that  $\mu = \delta_{p_0}$  (unit mass at  $p_0$ ).

iv.  $\lim_{n \rightarrow \infty} \left\| \frac{1}{n} \Psi^n w - v_\infty \right\| = 0$  for all  $w$ .

HINT. Since  $\lim_{n \rightarrow \infty} \|v_n - v_\infty\| = 0$ , and (non-expansiveness)  $\|\Psi^n w - \Psi^n \mathbf{0}\| \leq \|w\|$ .

v. For any  $w$  and its  $f_w$  (according to (12)):

$$f_w\left(\frac{\Psi^n w}{n} - v_\infty\right) + f_w(v_\infty - \frac{w}{n}) = f_w\left(\frac{\Psi^n w}{n} - \frac{w}{n}\right) \geq \rho$$

So by VEx.5biv (since  $f_w(\mathbf{0}) = 0$ ),  $f_w(v_\infty) \geq \rho$ , therefore if  $f_w$  is represented by the measure  $\mu_w$ , then  $\mu_w = \delta_{p_0}$  for every  $w$ .

vi. From the definitions of  $\Psi$  and  $u$  it follows that  $\Psi w \geq u + w$  for all  $w$ . Hence  $\rho \geq \|u\|$  and by VEx.5bv:

$$(15) \quad (\Psi w - w)(p_0) \geq \rho \geq \|u\|$$

vii. Choose  $w \in X$  such that  $w(p) \in [0, 1] \ \forall p$ ,  $w(p_0) = 1$ , and  $w(p) = 0 \ \forall p$  such that  $\|p - p_0\| > \varepsilon$ . Since

$$\mathsf{E}_x \|p - p_s\| = \mathsf{E}_p \|x - \bar{x}\| \quad \text{where } \bar{x}^k = \sum_k p^k x^k, \quad \forall k,$$

we have from (13) p. 254 that

$$\Psi w(p_0) \leq u(p_0) + w(p_0) + c\varepsilon$$

Hence by (15):

$$u(p_0) + c\varepsilon \geq (\Psi w - w)(p_0) \geq \rho \geq \|u\|$$

Since this must hold for all  $\varepsilon > 0$  conclude that  $u(p_0) = \|u\| = \rho = v_\infty(p_0)$ , completing the proof.

## 6. Recursive formula.

a. Prove that in a game in which the signals to player I are **more informative** than those to player II (i.e.  $\mathcal{H}^I \supseteq \mathcal{H}^{II}$ ), the following generalisation of the recursive formula of Lemma 4.2 p. 236 holds:

$$v_{n+1}(p) = \frac{1}{n+1} \max_{x \in X^K} \min_{t \in T} \left\{ \sum_{k \in K} p^k x^k G_t^k + n \sum_{b \in B} q_{bt} v_n(p_{bt}) \right\}$$

where if  $\tilde{Q}_t$  is the probability distribution on  $K \times S \times A \times B$  induced by  $p$  and  $x$ , and given that player II uses move  $t \in T$ , then  $q_{bt} = \tilde{Q}_t(b)$  (that is  $q_{bt} = \sum_k \sum_a p^k x^k Q_t^k(a, b)$ ) and  $p_{bt}$  is the conditional probability distribution on  $K$  given  $b$  (and  $t$ ). In terms of the functions  $V_n = nv_n$ , made positively homogeneous of degree 1, and the variable  $z \in \Delta(K \times S)$  this formula can be rewritten as:

$$V_{n+1}(p) = \max_{\bar{z}^k = p^k} \min_{t \in T} \left\{ z G_t + \sum_{b \in B} V_n[(z^k Q_t^k(b))_{k \in K}] \right\}$$

Here  $\bar{z}^k = \sum_{s \in S} z_s^k$  is the marginal probability of  $k$  according to  $z$ .

b. Prove that this formula is a special case of the general recursive formula (theorem 3.2 p. 158).

HINT. A game with incomplete information on one side with a finite state space  $K$  corresponds to a finite  $BL$ -subspace  $Y_p$  with a consistent probability  $p \in \mathcal{P}$  which is also an element of  $\Delta(K)$  (cf. example 3 p. 117).

Given a pair of strategies  $(x, y)$  of the two players in the first stage, any  $p \in \Delta(K)$  is mapped into a probability distribution  $P[x, y] \in \Delta(\Delta(K))$  which is the distribution on  $\Omega$  of the posterior probability distribution on  $K$  given the signal to the uninformed player. Hence in this special case the state space  $\Omega$  in Theorem 3.2 p. 158 may be restricted to a much smaller space namely  $\Delta(K)$ .

Note that, unlike in the more special lemma 4.2 p. 236, the operation  $\min_{t \in T}$  cannot be applied to the first stage term only, since now  $q_{bt}$  and  $p_{bt}$  also depend on  $t$ .

**7. Monotonicity of  $v_n$ .** Assuming  $\mathcal{H}^I \supseteq \mathcal{H}^{II}$ , use the recursive formula of ex. VEx.6 to prove that the sequence of value functions  $v_n$  is decreasing that is:

$$v_{n+1}(p) \leq v_n(p) \quad \forall n \geq 1, \forall p \in [0, 1]$$

HINT. The proof is by induction on  $n$ . For  $n = 1$  it follows from ex. VEx.1 p. 252. Make use of the concavity of  $v_n$  and the fact that  $p_{bt}$  is a martingale.

**8. Non-monotonicity of  $v_n$ .** (Lehrer, 1987) An intuitive argument supporting the monotonicity of  $v_n$  is the following: Player I has, at the beginning of the game some information not available to player II. This advantage can only diminish as the game progresses since player II (and only him) can gain more information about the state. The following example shows this intuitive argument to be false unless player I always knows whatever player II knows:

Consider the game in which  $K = \{1, 2, 3\}$ ,  $p = (1/3, 1/3, 1/3)$ . The moves of player I are  $S = \{\alpha, \beta\}$  and of player II are:  $T = \{t_1, t_2, t_3, t_4, t_5\}$ . The pay-off matrices are

$$G^1 = \begin{pmatrix} 100 & 4 & 0 & 4 & 4 \\ 100 & 0 & 4 & 4 & 4 \end{pmatrix} \quad G^2 = \begin{pmatrix} 0 & 100 & 4 & 4 & 4 \\ 4 & 100 & 0 & 4 & 4 \end{pmatrix} \quad G^3 = \begin{pmatrix} 4 & 0 & 100 & 4 & 4 \\ 0 & 4 & 100 & 4 & 4 \end{pmatrix}$$

The signalling matrices for player I:

$$A^1 = A^2 = A^3 = \begin{pmatrix} 1 & 2 & 3 & \theta & \theta \\ 1 & 2 & 3 & \theta & \theta \end{pmatrix}$$

The signalling matrices for player II:

$$B^1 = \begin{pmatrix} \theta & \theta & \theta & a & b \\ \theta & \theta & \theta & c & d \end{pmatrix} \quad B^2 = \begin{pmatrix} \theta & \theta & \theta & c & e \\ \theta & \theta & \theta & f & b \end{pmatrix} \quad B^3 = \begin{pmatrix} \theta & \theta & \theta & f & d \\ \theta & \theta & \theta & a & e \end{pmatrix}$$

Recall that in addition to his signal each player knows his own move. In words the signalling structure is as follows: Player I gets to know the move of player II except for moves  $t_4$  and  $t_5$

between which he cannot distinguish. Player II gets no information neither about  $k$  nor about the move of player I whenever he chooses  $t_1$ ,  $t_2$  or  $t_3$ . When he plays  $t_4$  or  $t_5$  he gets a signal which enables him to exclude one of the states.

Prove that  $v_1 \geq v_2 < v_3$ .

HINT. (1) Clearly  $v_1 \leq 4$ .

(2) The following strategy for player II in  $\Gamma_2$  shows that  $v_2 \leq 3$ :

At first stage play  $\frac{1}{2}t_4 + \frac{1}{2}t_5$  and at the second stage

- Play  $t_1$  if received signal  $e$  or  $f$ .
- Play  $t_2$  if received signal  $a$  or  $d$ .
- Play  $t_3$  if received signal  $c$  or  $b$ .

(3) The following strategy of player I in  $\Gamma_3$  shows that  $v_3 \geq 10/3$ :

At the first two stages play  $(1/2, 1/2)$  (i.e.  $\frac{1}{2}\alpha + \frac{1}{2}\beta$ ). If both signal received were  $\theta$  play  $(1/2, 1/2)$  at the third stage as well. Otherwise let  $t$  be the first signal different from  $\theta$  (i.e. 1, 2, or 3). At the third stage

- When  $k = 1$  play  $(1/2, 1/2)$  if  $t = 1$ ,  $\alpha$  if  $t = 2$  and  $\beta$  if  $t = 3$ .
- When  $k = 2$  play  $\alpha$  if  $t = 1$ ,  $(1/2, 1/2)$  if  $t = 2$  and  $\beta$  if  $t = 3$ .
- When  $k = 3$  play  $\alpha$  if  $t = 1$ ,  $\beta$  if  $t = 2$  and  $(1/2, 1/2)$  if  $t = 3$ .

Conclude that the game under consideration does not have a recursive formula as in ex. VEx.6 p. 255. How do the conditions for that formula fail in this example?

**9. The impact of the signals on the value.** In sect. 5.b p. 248 we saw that the signalling structure may affect the speed of convergence of  $v_n$ . Work out the details of the following example (due to Ponssard and Sorin) to show that the effect of the signalling may occur already in  $v_2$ .

Let  $K = \{1, 2\}$ . The pay-off matrices are:

$$G^1 = \begin{pmatrix} 0 & 2 & 2 \\ 0 & 2 & -3 \end{pmatrix} \quad G^2 = \begin{pmatrix} 2 & 0 & -3 \\ 2 & 0 & 2 \end{pmatrix}$$

and consider the two cases where the signalling matrices to player II are:

$$(Case 1) \quad Q^1 = Q^2 = \begin{pmatrix} a & b & c \\ a & b & f \end{pmatrix}$$

$$(Case 2) \quad \tilde{Q}^1 = \tilde{Q}^2 = \begin{pmatrix} a & b & c \\ d & e & f \end{pmatrix}$$

In both cases player I is informed of the moves of both players. Show that:

a. In both cases  $\text{NR}(p)$  (for  $0 < p < 1$ ) is the set of moves independent of  $k$ .

b. In both cases

$$u(p) = \begin{cases} 2p & \text{if } p \in [0, 2/7] \\ 2 - 5p & \text{if } p \in [2/7, 1/2] \\ u(1-p) & \text{if } p \in [1/2, 1] \end{cases}$$

c. In both cases

$$v_\infty(p) = \lim v_n(p) = 2 \min(p, \frac{2}{7}, 1-p).$$

d. In both cases  $v_1(p) = 2 \min(p, 1-p)$  for  $p \in [0, 1]$ .

e. For  $p \notin [2/7, 5/7]$ ,  $v_n(p) = v_1(p) \forall n$  (use the monotonicity of  $v_n$ .)

f. Using the recursive formula show that in Case 1

$$v_2(p) = v_1(p) = 2 \min(p, 1-p) \quad \text{for } p \in [0, 1],$$

while in Case 2

$$v_2(p) = (3p + 2)/5 \quad \text{for } p \in [2/7, 5/7].$$

REMARK 5.2. The last point provides an example to the fact that  $v_1 = v_2$  does not imply that  $v_n$  is constant in  $n$ : In Case 1,  $v_1 = v_2$  but  $\lim v_n(p) \neq v_1(p)$  for  $p \in [2/7, 5/7]$ .

### 10. On the speed of convergence of $v_n$ .

a. For the game in example 4.1 p. 236 do the majorations more carefully to reduce the Lipschitz constant to  $1/2$  the maximum difference between pay-offs. Conclude that for this game  $v_n(p) \leq \sqrt{pp'/n}$

b. For the game in example 5.1 p. 246 prove directly (by induction using the recursive formula) that for all  $n$  and for all  $p \in [0, 1]$ :

$$v_n(p) \leq \frac{\alpha \sqrt{pp'}}{\sqrt[3]{n}} \quad \text{with } \alpha = \sqrt[3]{192}$$

c. Prove Lemma 5.6 p. 250 for any finite state set  $K$  of size  $k$ : If  $u(p)$  is twice differentiable let  $u''(p)$  be the  $(k-1) \times (k-1)$  matrix with elements

$$(u''(p))_{ij} = \frac{\partial^2 u(p)}{\partial p^i \partial p^j} \quad \text{for } 1 \leq i \leq k-1 ; 1 \leq j \leq k-1$$

If  $\exists \eta > 0$  such that  $e u''(p) \tilde{e} \leq -\eta$ , for all  $p \in \Delta(K)$  and for any unit vector  $e \in \mathbb{R}^{k-1}$  ( $\tilde{e}$  is the transposition of  $e$ ), then for some constant  $A > 0$ ,  $\delta_n(p) \leq A \frac{\ln n}{n}$  for all  $p \in \Delta(K)$ .

d.

i. Show that for the game of example 2.1 p. 185:

$$V_n(p) \leq \left( n + 1 + \frac{1}{4} \ln(n+1) \right) pp'$$

HINT. Cf. the proof of prop. 5.7 p. 251, let

$$F(\alpha) = \max_{\substack{0 \leq x \leq p \\ 0 \leq y \leq p'}} \frac{1}{pp'} \{ \min(x, y) + \alpha [\phi(x, p' - y) + \phi(p - x, y) - \phi(p, p')] \} \quad \text{with } \phi(p, q) = \frac{pq}{p+q}. \quad \text{Show that}$$

$F(\alpha) = 2 - \alpha$  for  $\alpha \leq 1/2$  and  $F(\alpha) = 1 + \frac{1}{4\alpha}$  for  $\alpha \geq 1/2$ , and that if  $C_0 \geq 0$  and  $C_{n+1} - C_n \geq F(C_n)$  then  $V_n(p) \leq C_n pp'$ .

ii. In prop. 5.7 p. 251 it was shown that  $V_n(p) \geq (n + 1 + \frac{1}{4} \ln(n+1)) pp' - D$  for some  $D$ .

Show there is no  $\tilde{D}$  such that  $V_n(p) \geq \left[ n + 1 + \frac{1}{4} \ln(n+1) \right] pp' - \tilde{D}$ .

HINT. Use theorem 1.6 p. 184.

REMARK 5.3. One should be able, even in general, to get by analytic means (i.e., without having to care about the combinatorial end effect) the expansion of  $v_n$  up to terms of the order of  $Kpp'$  (i.e., an expression  $\phi_n(p)$  with  $|V_n(p) - \phi_n(p)| \leq Kpp'$ ) and one needs such if the expansion is to be useful also for small values of  $p$ . VEx.10dii shows therefore that even for this specific game we still need some improvement.

iii. Deduce from the above that for this game no separable expression of the form  $\phi_n(p) = \alpha_n \phi(p)$  will satisfy  $|V_n(p) - \phi_n(p)| \leq Kpp'$ .

e. Prove the following sufficient condition for the error term to be of the order  $O(1/n)$ . Such an error term means that the cumulative excess of pay-off to  $\text{Cav } u$  is bounded:

In a game with  $K = \{1, 2\}$ , if  $\exists \alpha > 0$  such that for all  $p \in [0, 1]$  and for all  $n \geq 1$ :

$$v_\infty(p) - u(p) \geq \alpha pp'$$

then

$$v_n(p) \leq v_\infty(p) + A \frac{\sqrt{pp'}}{n}$$

for some constant  $A > 0$ .

This result can also be proved for any finite  $K$  (cf. Zamir, 1971–1972).

HINT. • For  $x = (x^1, x^2) \in X^2$  let  $\|x^1 - x^2\| = \sum_{s \in S} |x_s^1 - x_s^2|$  and prove:

$$\min_{t \in T} (px^1 G_t^1 + p' x^2 G_t^2) \leq u(p) + Cpp' \|x^1 - x^2\|_1$$

- Given  $x \in X^2$ , if  $p_s$  is the conditional probability of  $\{k = 1\}$  resulting from move  $s$  then  $\mathbb{E} \sqrt{p_s p'_s} = \sqrt{pp'} \sum_s \sqrt{x_s^1 x_s^2}$ , and, since  $\sqrt{z^1 z^2} \leq \bar{z} - (z^1 - z^2)^2/(8\bar{z})$  with  $\bar{z} = (z_1 + z_2)/2$ ,

$$\mathbb{E} \sqrt{p_s p'_s} \leq (1 - \frac{1}{8} \|x^1 - x^2\|_2^2) \sqrt{pp'}.$$

- Use the recursive formula to prove the result — with  $A = (\#S)C^2/\alpha$  — by induction on  $n$ . The main inductive step being — with  $E_n(p) = n(v_n(p) - v_\infty(p))$  assumed  $\leq A\sqrt{pp'}$ :

$$E_{n+1}(p) \leq \max_{\delta \geq 0} [-\alpha pp' + (\#S)^{\frac{1}{2}} Cpp' \delta + A\sqrt{pp'}(1 - \frac{\delta^2}{8})]$$

using  $\delta$  for  $\|x^1 - x^2\|_2$ .

REMARK 5.4. For the game

$$G^1 = \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix} \quad G^2 = \begin{pmatrix} -2 & 0 \\ 0 & 2 \end{pmatrix}$$

$u(p) = \max(-p, -q^2/p)$ , with  $q = 1 - p$ , and hence  $v_\infty(p) = 0$ . So, point VEx.10e is not applicable. In fact it was proved in (Mertens and Zamir, 1995) that for this game (and whenever  $\min(\#S, \#T) \leq 2 = \#K$ , and  $u(p) < v_\infty(p)$  for  $0 < p < 1$ )

$$v_n(p) - v_\infty(p) \leq (K \ln n)/n.$$

f. To show that this cannot be improved, show that, for this game,  $v_n(p) \geq \frac{K}{n} W_n(q)$  with  $W_n(q) = g_n(q) - qg_n(1)$ ,  $g_n(x) = \ln(1 + n^{1/3}x)$  — (and e.g.  $K = \frac{1}{3}$ ).

HINT. (1) It suffices by the recursive formula to show that

$$W_{n+1}(q) \leq \max_{\substack{0 \leq u \leq p \\ 0 \leq v \leq q}} \left[ \frac{2v - 2u + p - 2q}{K} + (u + v)W_n\left(\frac{v}{u+v}\right) + (1 - u - v)W_n\left(\frac{q-v}{1-u-v}\right) \right]$$

i.e. that

$$\min_{\substack{0 \leq u \leq v + \frac{1}{2}p - q \\ 0 \leq v \leq q}} \frac{W_{n+1}(q) - (u + v)W_n\left(\frac{v}{u+v}\right) - (1 - u - v)W_n\left(\frac{q-v}{1-u-v}\right)}{2v - 2u + p - 2q} \leq K^{-1}.$$

Now  $W_{n+1}(q) - W_n(q) = f_n(q) - qf_n(1)$ , with  $f_n(x) = \ln(1 + \frac{[(1+n^{-1})^{1/3}-1]q}{n^{-1/3}+q})$ . Since  $f_n$  is monotone, we have

$$\frac{f_n(q) - qf_n(1)}{1 - q} \leq f_n(q) \leq \frac{[(1+n^{-1})^{1/3}-1]q}{n^{-1/3}+q} \leq \frac{q}{3n(n^{-1/3}+q)},$$

so

$$W_{n+1}(q) - W_n(q) \leq \frac{1 - q}{3n(1 + \frac{1}{n^{1/3}q})}.$$

Thus we have to show that

$$\min_{\substack{0 \leq u \leq v + \frac{1}{2}p - q \\ 0 \leq v \leq q}} \frac{\varphi_n(p, q) + g_n(q) - (u + v)g_n\left(\frac{v}{u+v}\right) - (1 - u - v)g_n\left(\frac{q-v}{1-u-v}\right)}{2v - 2u + p - 2q} \leq K^{-1}$$

with  $\varphi_n(p, q) = p/[3n(1 + 1/qn^{1/3})]$ .

(2) Case  $q \leq \frac{1}{2}$ .

Let  $u = \frac{1}{2} - q$ ,  $v = q - \frac{1}{2}z$ . Our problem becomes, with  $\psi_n(q, z) = 2g_n(q) - (1 - z)g_n\left(\frac{2q-z}{1-z}\right) - (1 + z)g_n\left(\frac{z}{1+z}\right)$ , to show that  $F_n(q) = \min_{0 \leq z \leq q} \frac{\varphi_n(p, q) + \frac{1}{2}\psi_n(q, z)}{q - z} \leq K^{-1}$ . Now

$$\begin{aligned} \psi_n(q, z) &= \ln \left[ 1 + \frac{1}{\left( \frac{q+n^{-1/3}}{q-z(1+n^{-1/3})} \right)^2 - 1} \right] + z \ln \left[ 1 + 2 \frac{qn^{1/3} - z(1 + n^{1/3})}{1 + z(1 + n^{1/3})} \right] \\ &\quad + (1 - z) \ln(1 - z) + (1 + z) \ln(1 + z) \\ &\leq \ln \left[ 1 + \frac{1}{\left( \frac{q+n^{-1/3}}{q-z} \right)^2 - 1} \right] + z \ln \left[ 1 + 2 \frac{q - z}{n^{-1/3} + z} \right] + (2 \ln 2)z^2 \\ &\leq \frac{1}{\left( \frac{q+n^{-1/3}}{q-z} \right)^2 - 1} + \frac{2z(q - z)}{z + n^{-1/3}} + (2 \ln 2)z^2 \\ &\leq \frac{1}{\left( \frac{q+n^{-1/3}}{q-z} \right)^2 - 1} + 2(q - z) + (2 \ln 2)q^2. \end{aligned}$$

So, with  $B = q + n^{-1/3}$ ,  $A = q^2 \ln 2 + pq/(3nB)$ ,  $w = (q - z)/B$ ,

$$F_n(q) \leq 1 + \frac{1}{B} \min_{0 \leq w \leq q/B} \left[ \frac{A}{w} + \frac{1}{2} \frac{w}{1 - u^2} \right]$$

Now  $w = \sqrt{2A}/(\sqrt{2A} + 1)$  [whenever  $w \leq q/B = \delta$ , i.e. whenever  $2A \leq (\frac{\delta}{1-\delta})^2$ ] yields a minimum  $\leq \sqrt{2A} + A$ , and  $w = \delta$  yields  $A/\delta + \frac{1}{2} \frac{\delta}{1-\delta^2}$  whenever  $(\frac{\delta}{1-\delta})^2 \leq 2A$  — hence  $\frac{1}{2} \frac{\delta}{1-\delta^2} \leq A/\delta$ . So anyway  $F_n(q) \leq 1 + \frac{1}{B} \max(2AB/q, \sqrt{2A} + A)$ .

Now  $A = q^2 \ln 2 + \frac{1}{3n}pq/(q + n^{-1/3}) \leq q^2 \ln 2 + \frac{1}{3}pq/(q + 1) \leq \frac{1}{4} [\ln 2 + \frac{1}{3}\frac{2}{3}] \leq \frac{1}{4}$ . So  $\sqrt{2A} + A \leq \frac{5}{2}\sqrt{A}$ , and we get  $F_n(q) \leq 1 + \max(2A/q, 5\sqrt{A}/2B)$ .

Further  $2A/q \leq (2 \ln 2)q + \frac{4}{3}/(q + 1) - \frac{2}{3} \leq \ln 2 + \frac{2}{9} < 2$ .

And  $\frac{5}{2} \frac{\sqrt{A}}{B} \leq \frac{5}{4} \sqrt{\frac{\ln 2}{(\varepsilon+1/2)^2} + \frac{\varepsilon^3/3}{(\varepsilon+1/2)^3}} \leq \frac{5}{2} \sqrt{\ln 2} < 2$ , so:  $F_n(q) \leq 3 \ \forall q \leq \frac{1}{2}, \forall n$ .

(3) Case  $q \geq \frac{1}{2}$ .

Let then  $u = 0$ ,  $v = 1 - \frac{5}{4}p$ . Then we get  $F_n(q) = [\varphi_n(p, q) + \chi_n(p, q)]/(p/2)$  with

$$\chi_n(p, q) = \ln \frac{1 + n^{1/3}q}{1 + n^{1/3}} + \frac{5}{4}p \ln \frac{1 + n^{1/3}}{1 + n^{1/3}/5} \leq \frac{p/q}{1 + n^{1/3}} - p + \frac{5p}{4} \ln 5 - \frac{5p}{n^{1/3} + 5}.$$

Hence

$$\begin{aligned} F_n(q) &\leq \frac{2}{3n(1 + 1/(qn^{1/3}))} + \frac{2}{q(1 + n^{1/3})} - \frac{10}{5 + n^{1/3}} - 2 + \frac{5}{2} \ln 5 \\ &\leq \frac{2}{3n + 6n^{2/3}} + \frac{4}{1 + n^{1/3}} - \frac{10}{5 + n^{1/3}} - 2 + \frac{5}{2} \ln 5 \leq \frac{5}{2} \ln 5 - \frac{13}{9} \leq 3 \\ &\forall q \geq \frac{1}{2}, \forall n \geq 1 \end{aligned}$$

(4)  $K = \frac{1}{3}$  also satisfies the inequality for  $n = 1$ .

**11. A game with incomplete information played by “non-Bayesian players”.** (Megiddo, 1980), cf. also VII.5Ex.6 p. 356 Consider a game with incomplete information on one side in which the uninformed player II is told his pay-off at each stage.

- a. Prove that  $u(p)$  is linear.

HINT.  $\text{NR}(p)$  is the set of strategies of player I which yield, against any strategy of player II, the same expected pay-off in all games in the support of  $p$ .

- b. Prove that if  $\text{Cav } u$  is linear then every optimal strategy of player II at an interior point of  $\Delta(K)$  is optimal for all  $p$ .

HINT. Use prop. 1.4 p. 184.

- c. *Extension to countable  $K$ .* Prove that for countable  $K$  player II has still a strategy which guarantees  $v(G^k)$  for all  $k$ , by playing optimally in a sequence of finite games.

- d. Consider now a situation in which the pay-off is according to some matrix  $G$ , player II knows only the set of his moves (columns) and is told his pay-off at each stage. Prove that for this situation VEx.11c still holds.

HINT. Observe that in VEx.11c, a strategy is still optimal if it responds to an unexpected pay-off as to a closest possible pay-off.

## 12. Discounted repeated games with incomplete information. (Mayberry, 1967)

Consider the game  $\Gamma_\lambda(p)$  with value  $v_\lambda(p)$ .

- a. Prove that

$$(1) \quad v_\lambda(p) = \max_x \left\{ \lambda \min_t \sum_k p^k x^k G_t^k + (1 - \lambda) \sum_s \bar{x}_s v_\lambda(p_s) \right\}$$

- b. For game in example 2.1 p. 185 prove that

$$(2) \quad v_\lambda(p) = \max_{s,t} \left\{ \lambda \min(ps, p't') + (1 - \lambda) (\bar{s}v_\lambda(ps/\bar{s}) + \bar{s}'v_\lambda(ps'/\bar{s}')) \right\}$$

Here  $(s, t) \in X^2$  is the pair of mixed moves used by player I in first stage and  $\bar{s} = ps + p't$ .

- c. Using the concavity of  $v_\lambda$  prove that the maximum in (2) is obtained at  $ps = p't'$ , hence denoting  $ps$  by  $x$  and using the symmetry of  $v_\lambda$  obtain

$$(3) \quad v_\lambda(p) = \begin{cases} \max_{0 \leq x \leq p} \{ \lambda x + (1 - \lambda)(pv_\lambda((p - x)/p) + p'v_\lambda(x/p')) \} & \text{for } 0 \leq p \leq \frac{1}{2} \\ v_\lambda(p') & \text{for } \frac{1}{2} \leq p \leq 1 \end{cases}$$

- d. For  $2/3 \leq \lambda \leq 1$ , use the concavity and the Lipschitz property of  $v_\lambda$  to simplify (3) further to:

$$(4) \quad v_\lambda(p) = \begin{cases} \lambda p + (1 - \lambda)p'v_\lambda(p/p') & \text{for } 0 \leq p \leq \frac{1}{2} \\ v_\lambda(p') & \text{for } \frac{1}{2} \leq p \leq 1 \end{cases}$$

- e. Observe that equation (4) reduces the problem of computing  $v_\lambda$  at any rational  $p = n/m \leq 1/2$  to the problem of computing  $v_\lambda$  at some other rational  $q$  with smaller denominator. Use (4) to compute  $v_\lambda(p)$  at some rational values of  $p$ :

$$\begin{aligned} v_\lambda(1/2) &= \lambda/2 \\ v_\lambda(1/3) &= \lambda(2 - \lambda)/3 \\ v_\lambda(1/4) &= \lambda(3 - 3\lambda + \lambda^2)/4 \\ v_\lambda(1/5) &= \lambda(4 - 6\lambda + 4\lambda^2 - \lambda^3)/5 \\ v_\lambda(2/5) &= 2/5 - (2 - 4\lambda + 4\lambda^2 - \lambda^3)/5 \end{aligned}$$

- f. By differentiating (4) obtain (letting  $v'_\lambda = \frac{dv_\lambda}{dp}$ ):

$$(5) \quad v'_\lambda(p) = (1 - \lambda)(1 - p/p')v'_\lambda(p/p') - (1 - \lambda)v_\lambda(p/p')$$

From this conclude (using the symmetry of  $v_\lambda$ ) that for  $2/3 < \lambda < 1$ , and at  $p = 1/2$ , the function has left derivative and right derivative but they are not equal.

g. Prove by induction on the denominator that for any rational  $p$ , the sequence of derivatives obtained by repeated use of equation (5) leads to  $v'_\lambda(1/2)$ .

h. Combining the last two results conclude that for  $2/3 < \lambda < 1$ : Although  $v_\lambda$  is concave, it has discontinuous derivatives at every rational point.

**13. On the notion of guaranteeing.** Consider the game in example 1.5 p. 150:  $K = \{1, 2\}$ ,  $p = (1/2, 1/2)$  and the pay-off matrices

$$G^1 = \begin{pmatrix} 2 & 0 \end{pmatrix} \quad G^2 = \begin{pmatrix} 0 & 1 \\ -1 & 2 \end{pmatrix}$$

The moves of player II are not announced to player I and the signals to player II are  $a$  or  $b$  according to the distributions:

$$Q^1 = \left( \begin{pmatrix} \frac{1}{3}, \frac{2}{3} \\ \frac{2}{3}, \frac{1}{3} \end{pmatrix} \quad \begin{pmatrix} \frac{1}{3}, \frac{2}{3} \\ \frac{2}{3}, \frac{1}{3} \end{pmatrix} \right) \quad Q^2 = \left( \begin{pmatrix} (0, 1) & (0, 1) \\ (\frac{2}{3}, \frac{1}{3}) & (\frac{2}{3}, \frac{1}{3}) \end{pmatrix} \right)$$

a. Prove that:  $\text{Cav } u(p) = \frac{3}{4} - \frac{3}{2} |p - \frac{1}{2}|$  (hence  $\text{Cav } u(1/2) = 3/4$ ).

b. Assume from now on that  $p = 1/2$ . Denote by  $\sigma$  the strategy of player I if  $k = 2$  and by  $\sigma_0$  the strategy consisting of playing  $(1/2, 1/2)$  i.i.d. in all stages.

Prove that given  $\sigma$ , the probability distribution of  $p_n$  is independent of player II's strategy. Denote by  $\mu_\sigma$  the measure induced by  $\sigma$  on  $(H_\infty, \mathcal{H}_\infty^{\text{II}})$ . Use Fubini's Theorem on  $K \times H_\infty$  to obtain:  $\frac{1}{2} \|\mu_\sigma - \mu_{\sigma_0}\| = 2 \mathbb{E} |p_\infty - \frac{1}{2}|$ .

Deduce that if  $\|\mu_\sigma - \mu_{\sigma_0}\| > \delta > 0$ , there exists a  $\tau$  such that  $\limsup \mathbb{E}_{\sigma, \tau}(\bar{g}_n) \leq \frac{3}{4} - \frac{3}{16}\delta$ .

c. Consider now the following strategy  $\tau$  of player II: If the first signal is  $a$ , play  $N_1$  times Right, then  $N_2$  times Left and so on, with  $N_n = 2^{2^n}$ ;  $n = 1, 2, \dots$ . If the first signal is  $b$  start with  $N_1$  times Left etc.

d. Prove that if  $\varphi_{B,m}$  denotes the relative frequency of the times in which  $B$  (bottom row) is played in the  $m^{\text{th}}$  bloc when  $k = 2$ , then for all large enough  $m$ :

$$\|\mu_\sigma - \mu_{\sigma_0}\| < \varepsilon \Rightarrow |\mathbb{E}(\varphi_{B,m}) - 1/2| \leq \frac{9}{2}\varepsilon$$

e. Denote by  $f_m^1$  and  $f_m^2$  the relative frequency of Left in games  $G^1$  and  $G^2$  respectively up to stage  $N_m$  (the end of the  $m^{\text{th}}$  bloc.) Prove that for any play of the game that results from  $\sigma$  and  $\tau$  the expected pay-off up to stage  $N_m$  is:

$$\frac{3}{4} + \mathbb{E}(f_m^1 - f_m^2) \pm 9\varepsilon.$$

f. Conditioning on the first move of player I in  $G^2$  prove that  $\mathbb{E}(f_m^1)$  is near  $1/3$  if  $m$  is even and near  $2/3$  if  $m$  is odd. Also show that

$$\begin{aligned} \mathbb{E}(f_{2n}^2 \mid T) &= 0, & \mathbb{E}(f_{2n}^2 \mid B) &\simeq 2/3 = \Pr(b_1 = a \mid B) \\ \mathbb{E}(f_{2n+1}^2 \mid T) &\simeq 1, & \mathbb{E}(f_{2n+1}^2 \mid B) &\simeq 1/3 \end{aligned}$$

g. Conclude that in all cases (whether the first move in  $G^2$  is  $T$  or  $B$ )

$$\mathbb{E}_{\sigma, \tau}(\liminf \bar{g}_n) = 5/12 \pm 9\varepsilon < 3/4$$

h. Prove a similar result for the following game with full monitoring and pay-off matrices:

$$G^1 = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \quad G^2 = \begin{pmatrix} -1 & 1 \\ 0 & 0 \end{pmatrix}$$

at  $p = 1/2$ .

**14. The conjugate recursive formula.** (De Meyer, 1996b,c) Consider a game  $\Gamma_n$  with full monitoring so that lemma 4.2 p. 236 holds. Denote by  $v_n^*$  (defined on  $\mathbb{R}^K$ ) the conjugate function of  $v_n$  (extended to  $-\infty$  outside  $\Pi = \Delta(K)$ ):

$$v_n^*(\alpha) = \min_{p \in \Pi} \{\langle \alpha, p \rangle - v_n(p)\}.$$

a. Prove that  $v_n^*(\alpha)$  is the value of a game  $\Gamma_n^*(\alpha)$  played as follows: at stage 0, I (in this case the minimiser) chooses  $k$  in  $K$ , II being uninformed. Then the play is as in  $\Gamma_n$  (after the chance move) and the pay-off is  $\alpha^k - \bar{g}_n^k$ .

b. Show that  $v_n^*$  satisfies:

$$v_{n+1}^*(\alpha) = \max_{y \in \Delta(T)} \min_{s \in S} \frac{n}{n+1} v_n^*\left(\frac{n+1}{n}\alpha - \frac{1}{n} \sum_t g_{st} y_t\right).$$

HINT. Deduce from lemma 4.2 p. 236 that:  $v_{n+1}^*(\alpha) =$

$$\frac{1}{n+1} \min_{\theta \in \Theta} \max_{y \in \Delta(T)} \left\{ (n+1) \langle \alpha, \theta_K \rangle - \sum_{s,t,k} g_{st}^k y_t \theta(k, s) - n \sum_s \theta(s) v_n(\theta(\cdot | s)) \right\}$$

where  $\Theta = \Delta(K \times S)$ ,  $\theta_K$  denotes the marginal of  $\theta$  on  $K$ , and  $\theta(\cdot | s)$  is the conditional on  $K$  given  $s$ . Apply then a minmax theorem.

c. Show that I and II have optimal strategies in  $\Gamma^*$  that do not depend on the previous moves of II. Call  $\tilde{\Sigma}_n^*$  and  $\tilde{\mathcal{T}}_n^*$  the corresponding sets, in particular  $\tilde{\Sigma}_n^* = \Delta(K \times S^n)$ .

Deduce that:

$$v_n^*(\alpha) = \max_{\tau \in \tilde{\mathcal{T}}_n^*} \min_{\sigma \in \tilde{\Sigma}_n^*} \left( \sum_{a \in I^n} \sigma(a) \min_k (\alpha^k - E_{a,\tau} \bar{g}_n^k) \right).$$

**15. Optimal strategies in finite games.** Assume standard signalling so that lemma 4.2 p. 236 holds.

a. Show that player I has an optimal strategy that depends only on his own past moves — more precisely on the stage  $m$  and the state variable  $p_m$  computed by II (cf. ex. VEx.6 p. 255).

b. Given  $\tau$  strategy of II in  $\Gamma_n$ , let  $\beta^k(\tau) = \max_{\sigma} E_{\sigma,\tau}^k \bar{g}_n^k$ .

i. Recall that  $v_n(p) = \min_{\tau} \langle \beta(\tau), p \rangle$  and that  $\forall \beta \in B_n = \{ \beta \in \mathbb{R}^K \mid \langle \beta, q \rangle \geq v_n(q) \text{ on } \Pi \}$ ,  $\exists \tau$  such that  $\beta \geq \beta(\tau)$  (cf. cor. 1.5 p. 184).

ii. An optimal strategy of II in  $\Gamma_n(p)$  is thus defined by a sequence of vectors  $\beta$  and mixed moves  $y$  s.t.  $\beta_n \in B_n, \langle \beta_n, p \rangle \leq v_n(p)$ , and inductively, at stage  $n-m+1$ , given  $\beta_m \in B_m$ ,  $y_{n-m+1} = y$  and  $(\beta_{m-1,s})_{s \in S}$  satisfy:  $\beta_{m-1,s} \in B_{m-1}, \forall s$  and  $m\beta_m^k = (m-1)\beta_{m-1,s}^k + \sum_t G_{st}^k y_t$ ,  $\forall k, \forall s$ .

Hence at stage  $n-m+1$ , player II uses  $y_{n-m+1}$  and if I's move is  $s$  the next reference vector is  $\beta_{n-m,s}$ . (See also ex. VIEx.6 p. 317 and VIEx.7 p. 318).

**16. An alternative proof of Theorem 4.1.** (Heuer, 1991b) We again consider the following two-state game (example 4.1 p. 236 above):

$$G^1 = \begin{pmatrix} 3 & -1 \\ -3 & 1 \end{pmatrix} \quad \text{and} \quad G^2 = \begin{pmatrix} 2 & -2 \\ -2 & 2 \end{pmatrix},$$

with  $p = P(k=1)$ .

Define  $b(k, n) = \binom{n}{k} 2^{-n}$ ,  $B(k, n) = \sum_{m \leq k} b(m, n)$ , for  $0 \leq k \leq n$  and  $B(-1, n) = 0$ .

a. Let  $p_{k,n} = B(k-1, n)$ ,  $k = 0, \dots, n+1$ , and prove that  $v_n$  is linear on each interval  $[p_{k,n}, p_{k+1,n}]$  with values  $v_n(p_{k,n}) = \frac{1}{2}b(k-1, n-1)$ .

So:  $v_n(p) = \min_{k=0, \dots, n+1} \langle \beta_{k,n}, p \rangle$ ,  $\forall p$ , with  $\beta_{k,n}^2 = \frac{1}{2}b(k, n-1) - (1 - \frac{2k}{n})B(k, n)$  and  $\beta_{k,n}^1 = \beta_{k,n}^2 + (1 - \frac{2k}{n})$ .

HINT. For player II, use ex. VEx.15 to obtain  $\beta_{k,n+1}$  by playing Left with probability  $\frac{1}{2} - \frac{1}{4}B(k-1, n)$  and then approaching  $\beta_{k,n}$  (resp.  $\beta_{k-1,n}$ ) if Top (resp. Bottom).

For player I it is enough to show that he can obtain  $v_n$  at each  $p_{k,n}$ . Note that

$$v_{n+1}(p_{k,n+1}) = \frac{1}{2} (v_n(p_{k-1,n}) + v_n(p_{k,n}))$$

and use at  $p_{k,n}$  the splitting defined by:

$$x^1(T) = \frac{B(k-1, n-1)}{2B(k-1, n)} \quad \text{and} \quad x^2(T) = \frac{1 - B(k-1, n-1)}{2(1 - B(k-1, n))}$$

b. Deduce theorem 4.1 p. 236.

HINT. Let  $\zeta(m, p) = \min\{k \mid B(k, n) \geq p\}$ . Show that  $(2\zeta(m, p) - n)/\sqrt{n}$  converges to  $x_p$  and use Stirling's formula.

**17. Exhausting information.** Consider the model of sect. 3, and fix a strategy  $\sigma$  of player I. This defines the sequence of posterior probabilities  $p_n$  of player II, everywhere on the set  $B_\sigma$  of histories in  $\bigcup_n B^n$  that are reachable under  $\sigma$ . Let  $D_n = \sum_k \sum_{m \geq n} [p_{m+1}(k) - p_m(k)]^2$ .

For any behavioural strategy  $\tau$  of player II (thus defined everywhere on  $\bigcup_n B^n$ ), let (cf. proof of lemma 2.1 p. 186):

$$X_n^\tau = \mathbb{E}_\tau \left[ \sum_k (p_\infty^2(k) - p_n^2(k)) \mid \mathcal{H}_n^{\text{II}} \right] = \mathbb{E}_\tau [D_n \mid \mathcal{H}_n^{\text{II}}] = \mathbb{E}_\tau \left[ \sum_k (p_\infty(k) - p_n(k))^2 \mid \mathcal{H}_n^{\text{II}} \right].$$

Let  $V_n = \sup_\tau X_n^\tau$ . Then:

- (1)  $X^\tau$  and  $V$  are well defined on  $B_\sigma$ .
- (2)  $V$  is a supermartingale (with values in  $[0, 1]$ ) under  $P_{\sigma,\tau}$  for all  $\tau$ , and conditionally to every  $h \in B_\sigma$ .
- (3) There exists a strategy  $\tau_\sigma$  such that,  $\forall h \in B_\sigma$ ,  $V_n \rightarrow 0$ ,  $P_{\sigma,\tau_\sigma}(\cdot \mid h)$  a.s.
- (4)  $\forall h \in B_\sigma, \forall \tau, \sum_k p(k \mid h) \|P_\tau(\cdot \mid k, h) - P_\tau(\cdot \mid h)\| \leq \sqrt{(\#K)V(h)}$ , where  $\|\cdot\|$  is the norm of measures on  $B_\infty$ .

HINT. For 2, prove first that  $X^\tau$  is a  $P_{\sigma,\tau}$  supermartingale, next construct an approximately optimal — or just sufficient —  $\tau$  for  $h \in B^n$ , given approximately optimal ones in  $B^{n+1}$ . Observe that this argument shows also the  $X_n^\tau$  are a lattice. For the bound, use the first formula for  $X^\tau$ .

For 3, let  $n_0 = 0$ , and given  $n_1, \dots, n_k$  ( $n_i < n_{i+1}$ ) and strategies  $\tau_k$  for the histories of length  $n$  ( $n_{k-1} \leq n < n_k$ ), consider, for  $h \in B_\sigma$  of length  $n_k$ , a strategy  $\tau^h$  for the future such that  $X_{n_k}^{\tau^h}(h) > V_{n_k}(h) - k^{-1}$ , next  $n_{k+1} = \min\{n > n_k \mid \forall h \in B_\sigma \cap B^{n_k}, \mathbb{E}_{\tau^h}[D_{n_k} - D_n \mid h] > V_{n_k}(h) - k^{-1}\}$ , and let  $\tau_{k+1}$  consist of using  $\tau^h$  after  $h$  until stage  $n_{k+1}$ . The  $\tau_k$ 's taken together form a strategy  $\tau_\sigma$ . And for  $h \in B_\sigma \cap B^{n_k}$ , and any strategy  $\tau$  that coincides with  $\tau_\sigma$  between  $n_k$  and  $n_{k+1}$ ,  $V_{n_k}(h) \geq \mathbb{E}_\tau(D_{n_k} \mid h) = \mathbb{E}_{\tau_\sigma}[D_{n_k} - D_{n_{k+1}} \mid h] + \mathbb{E}_{\tau_\sigma}[X_{n_{k+1}}^\tau \mid h] > V_{n_k}(h) - k^{-1} + \mathbb{E}_{\tau_\sigma}[X_{n_{k+1}}^\tau \mid h] \leq k^{-1}$ . Since  $\tau$  is arbitrary, and the  $X_n^\tau$  are directed, this yields  $\mathbb{E}_{\tau_\sigma}[V_{n_{k+1}} \mid h] \leq k^{-1}$ , — hence  $\mathbb{E}_{\tau_\sigma}[V_n \mid h] \leq \ell^{-1}, \forall \ell > k$ . So conditionally to  $h$ ,  $V_n$  is a positive supermartingale with expectation converging to 0. Conclude.

For 4, apply first (conditionally) Fubini's theorem (cf. 2.3 p. 274), next Hölder's inequality to get a bound  $\sqrt{(\#K)X^\tau(h)}$ .

COMMENT 5.5. This is in fact an exercise on (non-stationary) dynamic programming, or statistics.

COMMENT 5.6. One could have used the strict inequality in the definition of  $\tau_k$  to make this strategy in addition completely mixed, hence  $\tau_\sigma$  also — dispensing in this way with the complication of the conditional statement in 3 (and similarly in 2).

But it is in this form that the statement will be used (cf. e.g. sect. 3 p. 408, also sect. 3.a — in the sense that it implies that if one were to change  $\tau_\sigma$  to an arbitrary other strategy for the first  $n$  stages ( $n$  arbitrary), it would still have the same property.

## 6. Appendix

In this Appendix we shall prove lemmas 4.7 p. 241 and 4.8 p. 241 concerning two properties of the normal density function  $\phi(p)$  defined in theorem 4.1 p. 236. We start by examining the derivatives of the functions  $\phi(p)$  and  $x_p$ .

**PROPOSITION 6.1.** *The functions  $\phi(p)$  and  $x_p$  satisfy:*

- (1)  $\phi'(p) = -x_p$ ,
- (2)  $x'_p = 1/\phi(p)$ ,
- (3)  $\phi''(p) = -1/\phi(p) = -x'_p$ ,
- (4)  $\phi^{(3)}(p) = -x_p/\phi^2(p)$ ,
- (5)  $\phi^{(4)}(p) = -(1 + 2x_p^2)/\phi^3(p)$ ,
- (6)  $\phi^{(5)}(p) = -x_p(7 + 6x_p^2)/\phi^4(p)$ ,
- (7)  $\phi^{(6)}(p) = -(4x_p^2 + 7)(1 + 6x_p^2)/\phi^5(p)$ ,
- (8)  $\phi^{(2n)}(p) \leq 0 ; n = 1, 2, \dots$

**PROOF.** Parts 1 to 7 are results of straightforward differentiation. Part 8 will follow if we prove that

$$\phi^{(2n)}(p) = \frac{-1}{\phi^{2n-1}(p)} \sum_{j=0}^{n-1} a_j x_p^{2j}$$

where  $a_j \geq 0$  for  $j = 1, \dots, n-1$ .

We prove this by induction. By 3 it is true for  $n = 1$ . Assume it is true for  $n$  then

$$\begin{aligned} \phi^{(2n+1)}(p) &= \frac{-1}{\phi^{2n}(p)} \left\{ \sum_{j=0}^{n-1} [(2n-1)a_j + 2(j+1)a_{j+1}] x_p^{2j+1} + (2n-1)a_{n-1}x_p^{2n-1} \right\} \\ &= \frac{-1}{\phi^{2n+1}(p)} \sum_{j=0}^{n-1} \beta_j x_p^{2j+1}, \end{aligned}$$

where  $\beta_j \geq 0$ ,  $j = 0, \dots, n-1$ . Consequently

$$\begin{aligned} \phi^{(2n+2)}(p) &= \frac{-1}{\phi^{2n+1}(p)} \left[ 2n \sum_{j=0}^{n-1} \beta_j x_p^{2j+2} + \sum_{j=0}^{n-1} (2j+1)\beta_j x_p^{2j} \right] \\ &= \frac{-1}{\phi^{2n+1}(p)} \sum_{j=0}^{n+1} \gamma_j x_p^{2j}, \end{aligned}$$

where  $\gamma_j \geq 0$ ,  $j = 0, 1, \dots, n+1$ , concluding the proof of the proposition. ■

**PROPOSITION 6.2.** *Define the sequence  $\{p_n\}_1^\infty$  by*

$$\exp\left(-\frac{1}{2}x_{p_n}^2\right) = \frac{1}{n} \quad \text{and} \quad p_n \leq \frac{1}{2}$$

*then there exists  $n_0$  such that for any  $n \geq n_0$*

$$(1) \quad p_n \leq p \leq p'_n \implies \phi(p)/\sqrt{n} \leq \min(p, p').$$

**PROOF.** First by definition  $p = (1/\sqrt{2\pi}) \int_{-\infty}^{x_p} e^{-\frac{1}{2}x^2} dx$ , from which we have

$$p_n \leq p \leq p'_n \iff x_p^2 \leq x_{p_n}^2 \iff \exp\left(-\frac{1}{2}x_p^2\right) \geq \exp\left(-\frac{1}{2}x_{p_n}^2\right),$$

hence  $p_n \leq p \leq p'_n \iff \exp(-\frac{1}{2}x_p^2) \geq 1/n$  and the statement (1) may now be written as

$$\exp(-\frac{1}{2}x_p^2) \geq \frac{1}{n} \implies \frac{1}{\sqrt{n}} \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2}x_p^2} \leq \min\left(\frac{1}{\sqrt{2\pi}} \int_{-\infty}^{x_p} e^{-\frac{1}{2}x^2} dx, \frac{1}{\sqrt{2\pi}} \int_{x_p}^{\infty} e^{-\frac{1}{2}x^2} dx\right).$$

The statement at the right hand side is

$$\frac{1}{\sqrt{n}} \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2}x_p^2} \leq \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{-|x_p|} e^{-\frac{1}{2}x^2} dx.$$

We may therefore consider just, say,  $x_p \leq 0$  and prove (replacing  $x_p$  by  $y$ ) that

$$(2) \quad \frac{1}{\sqrt{n}} \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2}y^2} \leq \frac{1}{\sqrt{2\pi}} \int_{-\infty}^y e^{-\frac{1}{2}x^2} dx$$

holds whenever  $e^{-\frac{1}{2}y^2} \geq 1/n$  and  $y \leq 0$ .

Now (2) is true (for all  $n \geq 1$ ) whenever  $-1 \leq y \leq 0$ . This is because it is true for  $y = -1$  (by direct computation) and the left hand side is concave on  $-1 \leq y \leq 0$  and has smaller slope than the right hand side which is convex on  $-1 \leq y \leq 0$  (cf. Figure 1).

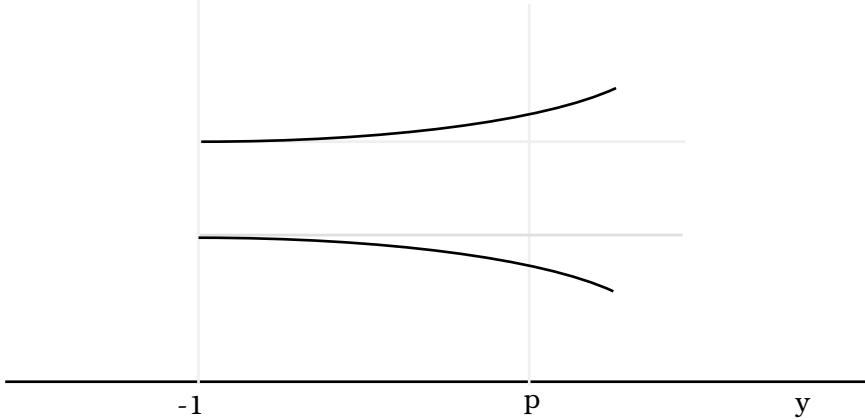


FIGURE 1. An inequality

For  $y < -1$ ,  $de^{-\frac{1}{2}y^2}/dy$  is positive and increasing. Hence at any point  $y < -1$ , the part of the tangent left to  $y$  lies below the line  $e^{-\frac{1}{2}x^2}$ . It intersects the  $x$ -axis at  $y + 1/y$  (cf. Figure 2). Therefore the integral on the right hand side of (2) p. 266 can be bounded

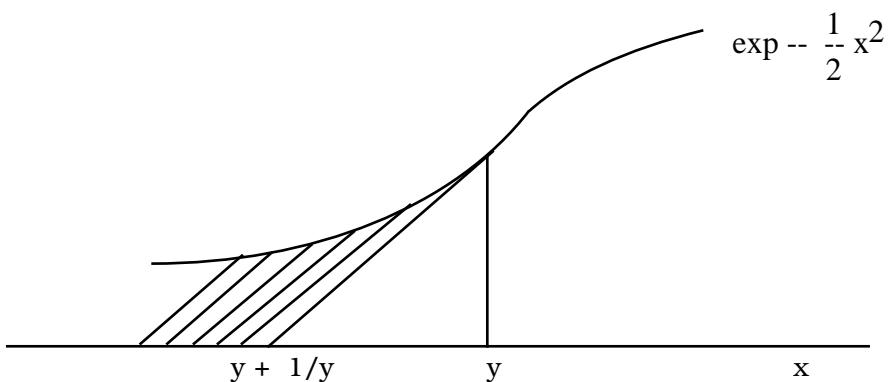


FIGURE 2. The tangents to the normal curve

by

$$\int_{-\infty}^y e^{-\frac{1}{2}x^2} dx \geq -\frac{1}{2y}e^{-\frac{1}{2}y^2},$$

and it suffices to prove that

$$e^{-\frac{1}{2}y^2} \geq \frac{1}{n} \implies -\frac{1}{2y} \geq \frac{1}{\sqrt{n}}.$$

In fact

$$e^{-\frac{1}{2}y^2} \geq \frac{1}{n} \implies |y| = -y \leq \sqrt{2 \ln n} \implies -\frac{1}{2y} \geq \frac{1}{2\sqrt{2 \ln n}}$$

and since  $(\ln n)/n \rightarrow 0$  let  $n_0$  be such that  $n \geq n_0$  implies  $-1/(2y) \geq 1/\sqrt{n}$  and we have thus proved (2) p. 266 for  $n \geq n_0$  completing the proof of prop. 6.2 p. 265. ■

PROOF OF LEMMA 4.7 P. 241. Using prop. 6.1 p. 265 we expand the first term in the left hand side of equation (15) in sect. 4 (lemma 4.7 p. 241) as follows: For some  $0 \leq \delta \leq x$ ,

$$(3) \quad \frac{1}{2}[\phi(p+x) + \phi(p-x)] = \phi(p) + \frac{x^2}{2}\phi''(p) + \frac{x^4}{4!} \frac{1}{2} [\phi^{(4)}(p+\delta) + \phi^{(4)}(p-\delta)]$$

$$(4) \quad = \phi(p) - \frac{x^2}{2\phi(p)} - \frac{x^4}{4!} \frac{1}{2} \left( \frac{1+2x_{p+\delta}^2}{\phi^3(p+\delta)} + \frac{1+2x_{p-\delta}^2}{\phi^3(p-\delta)} \right)$$

Clearly it is enough to prove equation (15) p. 241 for  $n \geq n_0$  for any fixed  $n_0$  and then to modify the constant  $c$  so as to make equation (15) hold for all  $n$ .

Define  $p_n$  by  $\exp(-\frac{1}{2}x_{p_n}^2) = 1/n$  and  $p_n \leq 1/2$ ; then by prop. 6.2 p. 265  $x = \phi(p)/\sqrt{n}$  is in the domain of maximisation in (15) p. 241 for  $n \geq n_0$  hence, denoting the left hand side of equation (15) by  $A$ , we get, using equation (4) for  $x = \phi(p)/\sqrt{n}$ ,

$$A - \phi(p) \geq \frac{\phi(p)}{\sqrt{n+1}} \left[ \sqrt{n} + \frac{1}{2\sqrt{n}} - \sqrt{n+1} \right] \\ - \frac{1}{2} \left( \frac{1+2x_{p+\delta}^2}{\phi^3(p+\delta)} + \frac{1+2x_{p-\delta}^2}{\phi^3(p-\delta)} \right) \frac{\phi^4(p)}{4!n^2\sqrt{n+1}} \sqrt{n}$$

which implies (since the sum of the first two terms in  $\{\dots\}$  is positive and  $\sqrt{n+1} \geq 1$ ):

$$(5) \quad A - \phi(p) \geq -\frac{1}{2} \left( \frac{1+2x_{p+\delta}^2}{\phi^3(p+\delta)} + \frac{1+2x_{p-\delta}^2}{\phi^3(p-\delta)} \right) \frac{\phi^4(p)}{4!n^2}$$

where  $0 \leq \delta \leq \phi(p)/\sqrt{n}$ .

Notice that  $\phi^{(4)}(p) = -(1+2x_p^2)/\phi^3(p)$  is symmetric about  $p = 1/2$  since  $\phi(p)$  is symmetric about  $p = 1/2$  and since  $x_{p'} = -x_p$ . It follows from (5) that for  $p \leq 1/2$ ,

$$A - \phi(p) \geq -\frac{1+2x_{p-\delta}^2}{\phi^3(p-\delta)} \frac{\phi^4(p)}{4!n^2}$$

which is

$$(6) \quad A - \phi(p) \geq -(1+2x_{p-\delta}^2) \frac{1}{\sqrt{2\pi}} \exp\left(\frac{3}{2}x_{p-\delta}^2 - 2x_p^2\right) \frac{1}{4!n^2}$$

Now  $1+2x^2 \leq 8 \exp(x^2/4)$  for all  $x$  hence,

$$(7) \quad A - \phi(p) \geq \frac{-8}{4!n^2\sqrt{2\pi}} \exp\left(\frac{7}{4}x_{p-\delta}^2 - 2x_p^2\right)$$

We now establish the existence of a constant  $\tilde{K}$  such that  $x_p - x_{p-\delta} \leq \tilde{K}/\sqrt{n}$  holds for  $p_n \leq p \leq p'_n$ ,  $p \leq 1/2$  and  $n$  sufficiently large. Since  $0 \leq \delta \leq \phi(p)/\sqrt{n}$  and since  $x_p - x_{p-\delta}$  is monotonically increasing in  $\delta$  we have to show that  $\Delta \leq \tilde{K}/\sqrt{n}$  where  $\Delta = x_p - x_{p-\phi(p)/\sqrt{n}}$ . Letting  $y = x_p \leq 0$  we claim in other words that

$$\frac{\phi(p)}{\sqrt{n}} = \frac{1}{\sqrt{2\pi}} \int_{y-\Delta}^y e^{-\frac{1}{2}x^2} dx \text{ implies } \Delta \leq \frac{\tilde{K}}{\sqrt{n}}$$

In fact for  $-1 \leq y \leq 0$  we have

$$\frac{1}{\sqrt{2\pi n}} \geq \frac{\phi(p)}{\sqrt{n}} = \frac{1}{\sqrt{2\pi}} \int_{y-\Delta}^y e^{-\frac{1}{2}x^2} dx \geq \frac{\Delta}{\sqrt{2\pi e}},$$

which implies  $\Delta \leq \sqrt{e}/\sqrt{n}$ .

For  $y \leq -1$  the tangent to  $(1/\sqrt{2\pi}) \exp(-\frac{1}{2}x^2)$  at  $x = y$  lies below the function and intersects the  $x$  axis at  $y + 1/|y|$  (cf. Figure 3) forming a triangular area  $(-1/2y)(1/\sqrt{2\pi}) \exp(-\frac{1}{2}y^2) = \phi(p)/(2|y|)$ . Now  $p \geq p_n$  implies  $|y| = |x_p| \leq |x_{p_n}| = \sqrt{2 \ln n} \leq 1/2\sqrt{n}$  for  $n$

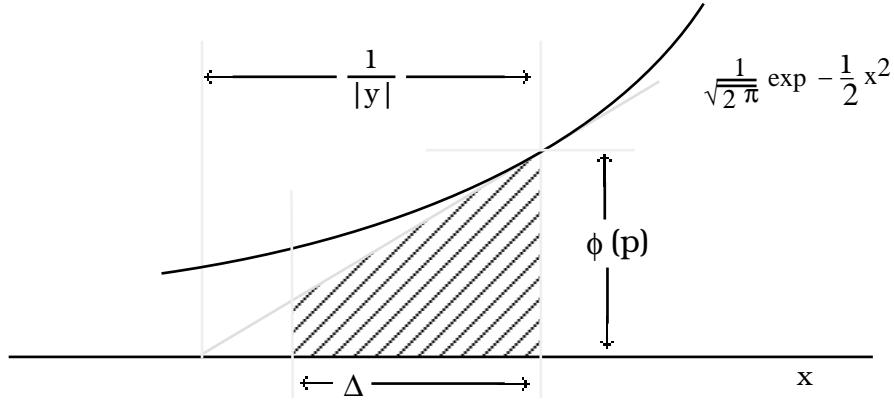


FIGURE 3. The area under the normal curve

sufficiently large, hence the triangular area is  $\geq \phi(p)/\sqrt{n}$  which implies  $\Delta \leq 1/|y|$ . The area of the shaded trapezoid is  $\phi(p)(2 - |y|\Delta)\Delta/2$ , therefore

$$\frac{\phi(p)}{\sqrt{n}} = \frac{1}{\sqrt{2\pi}} \int_{y-\Delta}^y e^{-\frac{1}{2}x^2} dx \geq \phi(p)(2 - |y|\Delta)\frac{\Delta}{2} \geq \phi(p)\frac{\Delta}{2}$$

This completes the proof that  $\Delta \leq \frac{\tilde{K}}{\sqrt{n}}$  for a suitable  $\tilde{K}$  and  $n$  sufficiently large. From this we get

$$x_{p-\delta}^2 = (x_p - \Delta)^2 \leq \left(x_p - \frac{\tilde{K}}{\sqrt{n}}\right)^2 = x_p^2 - \frac{2\tilde{K}x_p}{\sqrt{n}} + \frac{\tilde{K}^2}{n}$$

and

$$(8) \quad \frac{7}{4}x_{p-\delta}^2 - 2x_p^2 = -\frac{1}{4}x_p^2 - \frac{7\tilde{K}x_p}{2\sqrt{n}} + \frac{7\tilde{K}^2}{4n}$$

Since  $x_p < 0$ , the right hand side has a maximum (with respect to  $n$ ) at some  $n_0$ , hence

$$(9) \quad \frac{7}{4}x_{p-\delta}^2 - 2x_p^2 \leq -\frac{1}{4}x_p^2 - \frac{7\tilde{K}x_p}{2\sqrt{n_0}} + \frac{7\tilde{K}^2}{4n_0} \leq K$$

where  $K$  is the maximum of the parabola (in  $x_p$ ) at the right hand side. Combining (7) and (9) we finally obtain the existence of a constant  $C_1 > 0$  such that

$$(10) \quad A \geq \phi(p) - C_1/n^2 \text{ for } n \geq n_0 \text{ and } p_n \leq p \leq p'_n.$$

It remains to establish (10) also for  $p \leq p_n$  or  $p \geq p'_n$ . In this case, by the definition of  $p_n$ :  $\exp(-\frac{1}{2}x^2) \leq 1/n$  and therefore  $\phi(p) \leq 1/(n\sqrt{2\pi})$ . So (choosing  $x = 0$ )

$$\begin{aligned} \frac{1}{\sqrt{n+1}} \max_{0 \leq x \leq p \wedge p'} \left[ \frac{\sqrt{n}}{2} (\phi(p+x) + \phi(p-x)) + x \right] &\geq \frac{\sqrt{n}}{\sqrt{n+1}} \phi(p) \\ &\geq \phi(p) - \frac{\phi(p)}{2n+1} \geq \phi(p) - \frac{1}{n(2n+1)\sqrt{2\pi}} \geq \phi(p) - \frac{C_2}{n^2} \end{aligned}$$

for some constant  $C_2$ . Choose now  $C_3$  such that  $A \geq \phi(p) - C_3/n^2$  for  $1 \leq n \leq n_0$  and finally choose  $c = \max(C_1, C_2, C_3)$ . This completes the proof of lemma 4.7 p. 241. ■

PROOF OF LEMMA 4.8. We have to prove the existence of a constant  $K > 0$  such that for  $0 \leq p \leq 1$

$$(11) \quad \frac{1}{\sqrt{n+1}} \max_{(\xi, \eta) \in S(p)} \left[ \sqrt{n} \left( \frac{\eta}{\xi + \eta} \phi(p+\xi) + \frac{\xi}{\xi + \eta} \phi(p-\eta) \right) + \frac{2\xi\eta}{\xi + \eta} \right] \leq \phi(p) + \frac{K}{n^2}$$

where  $S(p) = \{(\xi, \eta) \mid 0 \leq \xi \leq p' ; 0 \leq \eta \leq p\}$ .

Since  $\phi$  continuous and  $S(p)$  is compact, the maximum in (11) is achieved, say, at  $(\xi_0, \eta_0)$ . Since  $(d\phi/dp)_{1-} = -\infty$ ,  $(d\phi/dp)_{0+} = +\infty$  it follows that  $\xi_0 \neq p'$  and  $\eta_0 \neq p$ . Furthermore we claim that if  $pp' \neq 0$  then  $\xi_0 \neq 0$  and  $\eta_0 \neq 0$ . In fact, denote the function to be maximised in (11) by  $F(\xi, \eta)$ , then  $F(0, \eta) = F(\xi, 0) = \sqrt{n}\phi(p)$  while

$$\max_{(\xi, \eta) \in S(p)} F(\xi, \eta) \geq \max_{0 \leq x \leq p \wedge p'} F(x, x) = \sqrt{n}\phi(p) + \max_{0 \leq x \leq p \wedge p'} ([O(x^2) + x]) > \sqrt{n}\phi(p).$$

We conclude that  $(\xi_0, \eta_0)$  is a local maximum of  $F(\xi, \eta)$  in  $S(p)$ . Equating first partial derivatives to 0 yields

$$(12) \quad \sqrt{n} \frac{\eta_0}{(\xi_0 + \eta_0)^2} [\phi(p - \eta_0) - \phi(p + \xi_0)] + \sqrt{n} \frac{\eta_0}{\xi_0 + \eta_0} \phi'(p + \xi_0) + \frac{2\eta_0^2}{(\xi_0 + \eta_0)^2} = 0$$

and

$$(13) \quad \sqrt{n} \frac{\xi_0}{(\xi_0 + \eta_0)^2} [\phi(p + \xi_0) - \phi(p - \eta_0)] - \sqrt{n} \frac{\xi_0}{\xi_0 + \eta_0} \phi'(p - \eta_0) + \frac{2\xi_0^2}{(\xi_0 + \eta_0)^2} = 0$$

Dividing (12) by  $\eta_0/(\xi_0 + \eta_0)$ , (13) by  $\xi_0/(\xi_0 + \eta_0)$  and adding the results we get

$$\sqrt{n} [\phi'(p + \xi_0) - \phi'(p - \eta_0)] + 2 = 0.$$

Recalling that  $\phi'(p) = -x_p$  we have

$$(14) \quad x_{p+\xi_0} - x_{p-\eta_0} = 2/\sqrt{n}$$

By the mean value theorem

$$x_{p+\xi_0} - x_{p-\eta_0} = [(p + \xi_0) - (p - \eta_0)] x'_{\theta(p+\xi_0)+(1-\theta)(p-\eta_0)}$$

for some  $0 \leq \theta \leq 1$ .

Using (14) and recalling that  $x'_p = 1/\phi(p)$  we get

$$(15) \quad \xi_0 + \eta_0 = (2/\sqrt{n}) \phi(\theta(p + \xi_0) + (1 - \theta)(p - \eta_0))$$

Now

$$\begin{aligned} \frac{\phi(\theta(p + \xi_0) + (1 - \theta)(p - \eta_0))}{\phi(p)} &= \exp\left\{\frac{1}{2}[x_p^2 - x_{p+\theta\xi_0-(1-\theta)\eta_0}^2]\right\} \\ (16) \quad &= \exp\left\{\frac{1}{2}[x_p + x_{p+\theta\xi_0-(1-\theta)\eta_0}][x_p - x_{p+\theta\xi_0-(1-\theta)\eta_0}]\right\} \end{aligned}$$

Since  $x_p$  is monotonically increasing in  $p$  we get from (16)

$$[x_p + x_{p+\theta\xi_0-(1-\theta)\eta_0}][x_p - x_{p+\theta\xi_0-(1-\theta)\eta_0}] \leq (2/\sqrt{n})(2|x_p| + 2/\sqrt{n})$$

and by (15) and (16) therefore

$$(17) \quad \xi_0 + \eta_0 = (2/\sqrt{n})\phi(p) \exp(2|x_p|/\sqrt{n}) \exp(2/n)$$

Denote

$$G(\xi, \eta) = \frac{\eta}{\xi + \eta}\phi(p + \xi) + \frac{\xi}{\xi + \eta}\phi(p - \eta)$$

Expanding  $\phi(p + \xi)$  and  $\phi(p - \eta)$  yields

$$\begin{aligned} (18) \quad G(\xi, \eta) &= \phi(p) + \frac{1}{2}\xi\eta\phi''(p) + \frac{1}{6}\xi\eta(\xi - \eta)\phi'''(p) + \frac{1}{24}\xi\eta(\xi^2 - \xi\eta + \eta^2)\phi^{(4)}(p) \\ &\quad + \frac{1}{120}\left[\frac{\eta\xi^5}{\xi + \eta}\phi^{(5)}(p + \sigma_1\xi) - \frac{\xi\eta^5}{\xi + \eta}\phi^{(5)}(p - \sigma_2\eta)\right] \end{aligned}$$

where  $0 \leq \sigma_1 \leq 1$  and  $0 \leq \sigma_2 \leq 1$ .

Consider the last term in (18) which we denote by  $K(p; \xi, \eta)$ . Since  $\phi^{(5)}$  is decreasing we have by prop. 6.1 p. 265

$$K(p; \xi, \eta) \leq -\frac{1}{120}\xi\eta(\xi^2 + \eta^2)(\xi - \eta)x_p(7 + 6x_p^2)/\phi^4(p).$$

By (17), since  $\max(\xi\eta, \xi^2 + \eta^2) \leq (\xi + \eta)^2$  and  $\xi - \eta \leq \xi + \eta$  we have

$$(19) \quad K(p; \xi_0, \eta_0) \leq -\frac{1}{120}\left[\frac{2}{\sqrt{n}}\phi(p)\exp\left(\frac{2|x_p|}{\sqrt{n}} + \frac{2}{n}\right)\right]^5 \frac{x_p(7 + 6x_p^2)}{\phi^4(p)}$$

$$(20) \quad = \frac{4}{15\sqrt{n^5}}\left[\phi(p)x_p(7 + 6x_p^2)\exp\left(\frac{10|x_p|}{\sqrt{n}} + \frac{10}{n}\right)\right]$$

$$(21) \quad = \frac{4}{15\sqrt{2\pi n^5}}\left[x_p(7 + 6x_p^2)\exp\left(\frac{10|x_p|}{\sqrt{n}} + \frac{10}{n} - \frac{1}{2}x_p^2\right)\right]$$

The last expression is clearly a bounded function of  $x_p$ , hence

$$(22) \quad K(p; \xi_0, \eta_0) \leq \frac{K_1}{n^2}$$

for some constant  $K_1$ .

By (18) and (22), using prop. 6.1 p. 265 we obtain

$$(23) \quad G(\xi_0, \eta_0) \leq \phi(p) - \frac{\xi_0\eta_0}{2\phi(p)} - \frac{\xi_0\eta_0(\xi_0 - \eta_0)x_p}{6\phi^2(p)} - \frac{\xi_0\eta_0(\xi_0^2 - \xi_0\eta_0 + \eta_0^2)(1 + 2x_p^2)}{24\phi^3(p)} + \frac{K_1}{n^2}$$

Therefore

$$(24) \quad \max_{(\xi, \eta) \in S(p)} \left[ G(\xi, \eta) + \frac{1}{\sqrt{n}} \frac{2\xi\eta}{\xi + \eta} \right] \leq \phi(p) + \frac{K_1}{n^2} + \max_{(\xi, \eta) \in S(p)} D(\xi, \eta)$$

where

$$(25) \quad D(\xi, \eta) = \frac{1}{\sqrt{n}} \frac{2\xi\eta}{\xi + \eta} - \frac{\xi\eta}{2\phi(p)} - \frac{\xi\eta(\xi - \eta)x_p}{6\phi^2(p)} - \frac{\xi\eta(\xi^2 - \xi\eta + \eta^2)(1 + 2x_p^2)}{24\phi^3(p)}$$

Observe that  $D(0, \eta) = D(\xi, 0) = 0$  and  $D(\varepsilon, \varepsilon) > 0$  for  $\varepsilon > 0$  sufficiently small. Also  $D(\xi, \eta) \rightarrow -\infty$  as  $\xi \rightarrow \infty$  or  $\eta \rightarrow \infty$ . It follows that  $D$  restricted to the non-negative orthant has a global maximum which also a local maximum. Equating first derivatives of  $D(\xi, \eta)$  to 0 and adding  $\frac{1}{\eta} \frac{\partial D}{\partial \xi} + \frac{1}{\xi} \frac{\partial D}{\partial \eta}$  we get

$$(26) \quad \frac{2}{\sqrt{n}(\xi + \eta)} - \frac{1}{\phi(p)} + \frac{(\xi - \eta)x_p}{2\phi^2(p)} - \frac{(\xi^2 - \xi\eta + \eta^2)(1 + 2x_p^2)}{6\phi^3(p)} = 0$$

Subtracting  $\frac{1}{\eta} \frac{\partial D}{\partial \xi} - \frac{1}{\xi} \frac{\partial D}{\partial \eta}$  we get

$$(27) \quad -\frac{2(\xi - \eta)}{\sqrt{n}(\xi + \eta)^2} - \frac{(\xi + \eta)x_p}{6\phi^2(p)} - \frac{(\xi^2 - \eta^2)(1 + 2x_p^2)}{12\phi^3(p)} = 0$$

Dividing (27) by  $(\xi^2 - \eta^2)$  and eliminating  $(\eta - \xi)$  yields

$$(28) \quad \eta - \xi = \frac{\frac{x_p}{\phi^2(p)}}{\frac{12}{\sqrt{n}(\xi + \eta)^3} + \frac{\frac{1}{2} + x_p^2}{\phi^3(p)}}$$

Replacing  $(\xi^2 - \xi\eta + \eta^2)$  in (26) by  $\frac{1}{4}(\xi + \eta)^2 + \frac{3}{4}(\xi - \eta)^2$  and  $(\xi - \eta)$  by its value according to (28) results in

$$\frac{2}{\sqrt{n}(\xi + \eta)} - \frac{1}{\phi(p)} + \frac{\frac{x_p^2}{\phi^4(p)}}{\frac{24}{\sqrt{n}(\xi + \eta)^3} + \frac{1 + 2x_p^2}{\phi^3(p)}} - \frac{1 + 2x_p^2}{24\phi^3(p)} \left[ (\xi + \eta)^2 + \frac{3 \frac{x_p}{\phi^2(p)}}{\frac{12}{\sqrt{n}(\xi + \eta)^3} + \frac{\frac{1}{2} + x_p^2}{\phi^3(p)}} \right]^2 = 0$$

The first term in the last equation tends to  $+\infty$  as  $(\xi + \eta) \rightarrow 0$ , the last term is always negative and the third is bounded from above by  $[1/\phi(p)]\{\max_x[x/(1 + 2x^2)]\}$  which is  $1/(2\sqrt{2}\phi(p))$ . So if we denote the left hand side by  $L(\xi, \eta)$  we can bound it by

$$(29) \quad L(\xi, \eta) \leq \frac{2}{\sqrt{n}(\xi + \eta)} - \frac{1}{\phi(p)} + \frac{1}{2\sqrt{2}\phi(p)}$$

The right hand side of (29) is non-negative if and only if  $\xi + \eta \leq \alpha\phi(p)/\sqrt{n}$  where  $\alpha = 2/(1 - 1/(2\sqrt{2})) \simeq 3.1$ . It follows therefore that any solution  $(\xi, \eta)$  of (26) and (28) must satisfy

$$(30) \quad \xi + \eta \leq \frac{\alpha\phi(p)}{\sqrt{n}}$$

By (28) we get that at the maximum  $(\xi, \eta)$

$$(31) \quad |\eta - \xi| \leq \frac{|x_p| \phi(p)}{\frac{12n}{\alpha^3} + \frac{1}{2} + x_p^2} < \frac{|x_p| \phi(p)\alpha^3}{12n}$$

Being interested in obtaining an upper bound for the global maximum of  $D$  we replace its two terms by an upper bound at the maximum. The resulting function will have a maximum which is greater or equal that of  $D$ . Now the last term of  $D$  (in (25)) is not positive and as for the third term, by (30) and (31),

$$\xi\eta(\xi - \eta) \frac{x_p}{6\phi^2(p)} \leq \frac{|x_p| \phi(p)\alpha^5}{72n^2} \leq \frac{\alpha^5}{72n^2} \max_x \frac{1}{\sqrt{2\pi}} |x| \exp(-\frac{1}{2}x^2) \leq \frac{K_2}{n^2}.$$

We conclude that

$$(32) \quad \max_{\xi,\eta} D(\xi, \eta) \leq \frac{K_2}{n^2} + \max_{\xi,\eta} D_1(\xi, \eta)$$

where  $K_2$  is a constant and

$$(33) \quad D_1(\xi, \eta) = \frac{2\xi\eta}{\sqrt{n}(\xi + \eta)} - \frac{\xi\eta}{2\phi(p)}$$

Equating first partial derivatives of  $D_1$  to 0 we get

$$\begin{aligned} \frac{1}{\eta} \frac{\partial D_1}{\partial \xi} &= \frac{2\eta}{\sqrt{n}(\xi + \eta)^2} - \frac{1}{2\phi(p)} = 0 \\ \frac{1}{\xi} \frac{\partial D_1}{\partial \eta} &= \frac{2\xi}{\sqrt{n}(\xi + \eta)^2} - \frac{1}{2\phi(p)} = 0 \end{aligned}$$

which imply  $\xi = \eta = \phi(p)/\sqrt{n}$  and hence

$$(34) \quad \max_{\xi,\eta} D_1(\xi, \eta) \leq \frac{\phi(p)}{2n}$$

By (24), (32) and (34)

$$(35) \quad \max_{(\xi,\eta) \in S(p)} \left[ G(\xi, \eta) + \frac{2\xi\eta}{\sqrt{n}(\xi + \eta)} \right] \leq \phi(p) \left( 1 + \frac{1}{2n} \right) + \frac{K_1 + K_2}{n^2}$$

Combining with (11) p. 269 we have now by (35)

$$\begin{aligned} (36) \quad & \frac{1}{\sqrt{n+1}} \max_{(\xi,\eta) \in S(p)} \left[ \sqrt{n} \frac{\eta}{\xi + \eta} \phi(p + \xi) + \sqrt{n} \frac{\xi}{\xi + \eta} \phi(p - \eta) + \frac{2\xi\eta}{\xi + \eta} \right] \\ &= \frac{\sqrt{n}}{\sqrt{n+1}} \max_{(\xi,\eta) \in S(p)} \left[ G(\xi, \eta) + \frac{2\xi\eta}{\sqrt{n}(\xi + \eta)} \right] \leq \phi(p) \left( 1 + \frac{1}{2n} \right) \frac{\sqrt{n}}{\sqrt{n+1}} + \frac{K_1 + K_2}{n^2} \end{aligned}$$

Now notice that  $\sqrt{1 + \frac{1}{n}} \geq 1 + \frac{1}{2n} - \frac{1}{8n^2}$ , therefore

$$\left( 1 + \frac{1}{2n} \right) \frac{\sqrt{n}}{\sqrt{n+1}} - 1 \leq \frac{\frac{1}{8n^2}}{1 + \frac{1}{2n} - \frac{1}{8n^2}} \leq \frac{K_3}{n^2}$$

where  $K_3$  is a constant. It follows that

$$\phi(p) \left( 1 + \frac{1}{2n} \right) \frac{\sqrt{n}}{\sqrt{n+1}} + \frac{K_1 + K_2}{n^2} \leq \phi(p) + \phi(p) \frac{K_3}{n^2} + \frac{K_1 + K_2}{n^2} \leq \phi(p) + \frac{K}{n^2}$$

where  $K$  is a constant. Combined with (36) this concludes the proof of lemma 4.8.  $\blacksquare$

## CHAPTER VI

# Incomplete Information on Both Sides

### 1. Introduction

The case of incomplete information on both sides is where neither player knows completely the state of nature. We can assume w.l.o.g. that the initial signals are chosen according to some probability  $P$  on  $K \times A \times B$  satisfying: for each  $k$ , there exists one and only one pair  $(a, b)$  with  $P(k, a, b) > 0$ . Just take  $K' = K \times A \times B$  as new state space and extend pay-off and signalling matrices on  $K'$  in the obvious way. It follows that the initial signals of I (resp. II) define a partition  $K^I$  (resp.  $K^{II}$ ) of  $K$ . The case in which one of the two partitions  $K^I$  and  $K^{II}$  of  $K$  is  $\{\{1\}, \{2\}, \dots, \{\#K\}\}$  was treated in the last chapter.

No general results are yet available for the whole class of such games. This chapter will be devoted to the sub-case in which  $Q^k$  is **independent of  $k$** , i.e. the information gained at each stage, does not depend on the state of nature and is determined completely by the players' moves at that stage. Omitting the index  $k$  for the state of nature, the transition probability on signals will therefore be denoted by  $Q$  from  $S \times T$  to  $A \times B$ .

We shall in a first part compute the Minmax and Maxmin of the infinitely repeated game. In the second part, we study the asymptotic value  $\lim v_n$  of the finitely repeated games which will be proved to exist always. A formula for  $\lim v_n$  will be proved with a few remarks on the speed of convergence.

We shall give some procedures to solve the functional equations determining  $\lim v_n$  and illustrate them by examples.

### 2. General preparations

In this section, we prove some lemmas which will be needed later in this chapter.

#### 2.a. Definitions and notations.

2.a.1. *Non-Revealing Strategies.*  $\mathcal{K}^I$  (resp.  $\mathcal{K}^{II}$ ) is the  $\sigma$ -field generated by  $K^I$  (resp.  $K^{II}$ ) on  $K$ .

A one-stage strategy of I is called **non-revealing** if for each column of  $Q$ , the marginal probability distribution on  $B$  induced on the letters of that column is independent of the state of nature  $k \in K$ . Formally  $x = (x^k)_{k \in K}$  in  $\Delta(S)^K$  is non-revealing if it is  $\mathcal{K}^I$  measurable and  $\sum_{s \in S} x^k(s) Q_{s,t}(b)$  is independent of  $k$  for all  $t$  in  $T$ ,  $b$  in  $B$ .

The set of non-revealing one-stage strategies of I is denoted by  $\mathbf{NR}^I$ . Similarly  $\mathbf{NR}^{II}$  is the set of non-revealing one-stage strategies of II, i.e., strategies such that for each row of  $Q$ , the marginal probability distribution on  $A$  on the letters of that row does not depend on  $k \in K$ .  $\mathbf{NR}^I$  and  $\mathbf{NR}^{II}$  are closed convex polyhedra, obviously non-empty: take strategies constant on  $K$ .

$D_{\mathbf{NR}}(p) = D(p)$  is the one-stage game in which I and II are restricted to strategies in  $\mathbf{NR}^I$  and  $\mathbf{NR}^{II}$  respectively. The value of  $D(p)$  is denoted by  $u(p)$ .

REMARK 2.1. Note the definition of non-revealing one-stage strategies given above differs slightly from the usual general definition of that concept which is: A one-stage

strategy of I is non-revealing if the marginal probability distribution on  $B$  induced on the letters of any column in  $Q$  is constant a.e. on  $K$  (i.e. it is the same for every  $k$  s.t.  $p^k > 0$ ). With this definition, the set of non-revealing strategies will depend (even discontinuously) on  $p$  and has to be denoted by  $\text{NR}^I(p)$  and  $\text{NR}^{II}(p)$ . However, it is easily seen that the modified  $D_{\text{NR}(p)}(p)$ , though formally different from  $D(p)$ , will have the same value  $u(p)$  (the projection of  $\text{NR}^I$  and  $\text{NR}^{II}(p)$  on  $K(p)$  are the same). Since all results will be formulated in terms of  $u(p)$  we conclude that the two definitions of non-revealing are equivalent. For obvious reasons we prefer the above definition which makes  $\text{NR}^I$  and  $\text{NR}^{II}$  independent of  $p$ .

**REMARK 2.2.** Note that to define non-revealing strategies in terms of posterior probabilities (as in ch. III p. 107) one has to consider the game where player II has no initial information.

**REMARK 2.3.** Remark that  $u(p)$  is continuous in  $p$  on the simplex  $\Pi = \Delta(K)$  of prior probabilities.

### 2.a.2. Concavification.

**DEFINITION 2.1.** A function on the simplex  $\Pi$  of probabilities is said to be **concave with respect to I** (shortly w.r.t. I) if for every  $p = (p^k)_{k \in K}$  in  $\Pi$  it has concave restriction on the subset  $\Pi^I(p)$  of  $\Pi$ , where

$$\Pi^I(p) = \left\{ (\alpha^k p^k)_{k \in K} \mid \alpha^k \geq 0 \ \forall k \in K, \sum_k \alpha^k p^k = 1 \text{ and } (\alpha^k)_{k \in K} \text{ is } \mathcal{K}^I\text{-measurable} \right\}.$$

A function on  $\Pi$  is said to be **convex with respect to II** if for every  $p = (p^k)_{k \in K}$  in  $\Pi$ , it has a convex restriction to the subset  $\Pi^{II}(p)$  of  $\Pi$ , where

$$\Pi^{II}(p) = \left\{ (\beta^k p^k)_{k \in K} \mid \beta^k \geq 0, \ \forall k \in K; \sum_k \beta^k p^k = 1 \text{ and } (\beta^k)_{k \in K} \text{ is } \mathcal{K}^{II}\text{-measurable} \right\}.$$

**REMARK 2.4.** For any  $p \in \Pi$  both  $\Pi^I(p)$  and  $\Pi^{II}(p)$  are convex and compact subsets of  $\Pi$  containing  $p$ , thus justifying the definition.

**DEFINITION 2.2.** Given any function  $g$  on  $\Pi$ , denote by  $\text{Cav}_I g$  the (point-wise) minimal function  $f \geq g$  which is concave w.r.t. I. Similarly denote by  $\text{Vex}_{II} g$  the (point-wise) maximal function  $f \leq g$  which is convex w.r.t. II.  $\text{Cav}_I g$  is called the **concavification** of  $g$  w.r.t. I and  $\text{Vex}_{II} g$  is called the **convexification** of  $g$  w.r.t. II.

**REMARK 2.5.** Note that if  $K^I = \{\{k\}_{k \in K}\}$  and  $K^{II} = \{K\}$ ,  $\text{Cav}_I g$  is the usual  $\text{Cav } g$  on  $\Pi$  and  $\text{Vex}_{II} g$  is  $g$ .

### 2.b. Preliminary results.

**LEMMA 2.3.** Let  $P$  be a positive measure on the product of two measurable spaces  $(X, \mathcal{X}) \otimes (Y, \mathcal{Y})$  that has a density with respect to the product of its marginals. Then

$$\mathbb{E}(\|P(dy | x) - P(dy)\|) = \mathbb{E}(\|P(dx | y) - P(dx)\|).$$

**PROOF.** Write  $P(dx, dy) = f(x, y)P(dx)P(dy)$ .

$$\mathbb{E}(\|P(dy | x) - P(dy)\|) = \int_X \int_Y |f(x, y)P(dy) - P(dy)| P(dx)$$

hence by Fubini's theorem equals:

$$(1) \quad \int_{X \times Y} |f(x, y)P(dy)P(dx) - P(dy)P(dx)|.$$

■

COMMENT 2.6. The quantity appearing in lemma 2.3 is therefore a natural measure of independence between  $X$  and  $Y$ .

Recall that a random variable  $X$  with values in a Banach space  $B$  is **Bochner integrable** if there exists a sequence  $X_n$  of measurable step functions with values in  $B$  such that  $\mathbb{E}\|X - X_n\| \rightarrow 0$  (Dunford and Schwartz, 1958, III.2).

LEMMA 2.4. For any (Bochner-) integrable random variable  $X$  with values in a Banach space  $B$ , and any  $y \in B$ ,

$$\mathbb{E}(\|X - \mathbb{E}(X)\|) \leq 2\mathbb{E}(\|X - y\|).$$

PROOF.  $\mathbb{E}(\|X - \mathbb{E}(X)\|) \leq \mathbb{E}(\|X - y\| + \|\mathbb{E}(X) - y\|) \leq 2\mathbb{E}(\|X - y\|)$  by the triangle inequality and Jensen's inequality (using convexity of the norm). ■

LEMMA 2.5. Let  $E^I$  be any finite set, and  $p = \sum_{e \in E^I} \lambda_e p_e$  where  $p_e \in \Pi^I(p)$ ,  $\lambda \in \Delta(E^I)$ ; then player I has a  $\mathcal{K}^I$ -measurable transition probability from  $K$  to  $E^I$  such that the resulting compound probability on  $K \times E^I$  satisfies:

- The (marginal) probability of  $e$  is  $\lambda_e$ ,  $\forall e \in E^I$
- The conditional probability on  $K$  given  $e \in E^I$  is  $p_e$ .

REMARK 2.7. This is a  $K^I$  measurable version of prop. 1.2 p. 184.

PROOF. For  $e \in E^I$ , let  $p_e = (\alpha_e^k p^k)_{k \in K}$ , then the required transition probability is defined by:  $P(e | k) = \lambda_e \alpha_e^k$  for  $e \in E^I$ ,  $k \in K$ . ■

LEMMA 2.6. Let  $f(p)$  and  $g(p)$  be functions on  $\Pi$  such that  $g(p) \leq \text{Cav}_I f(p)$  and let  $\#E^I = \#K^I$ ; then for any  $p_0 \in \Pi$  and  $\varepsilon > 0$  there are  $p_e \in \Pi^I(p_0)$ ,  $e \in E^I$  and  $\lambda_e \geq 0$  with  $\sum_e \lambda_e = 1$  such that  $\sum_e \lambda_e p_e = p_0$  and  $\sum_e \lambda_e f(p_e) \geq g(p_0) - \varepsilon$ .

If  $f(p)$  is continuous this is also true for  $\varepsilon = 0$ .

PROOF. The proof is an application of Carathéodory's theorem, since  $\Pi^I(p_0)$  is  $(\#K^I - 1)$  (equals to  $(\#E^I - 1)$ ) dimensional. ■

COMMENT 2.8. The two functional equations  $g(p) \leq \text{Cav}_I \min\{u(p), g(p)\}$  and  $g(p) \geq \text{Vex}_{II} \max\{u(p), g(p)\}$  will play a very important rôle in this chapter.

PROPOSITION 2.7. Let  $P = (P_\theta)$  be a transition probability from a probability space  $(\Theta, \mathcal{C}, \mu)$  to a finite measurable space  $(\Omega, \mathcal{A})$  on which a finite collection  $\mathcal{F}$  of measurable functions with values in  $[0, 1]$  is given, then:

$$\inf_{P' \in \text{NR}} \mathbb{E}(\|P_\theta - P'_\theta\|_1) \leq R \max_{f \in \mathcal{F}} \mathbb{E}(|P_\theta(f) - \mathbb{E}(P_\theta(f))|)$$

for some constant  $R$  depending only on  $\Omega$  and  $\mathcal{F}$ .

REMARK 2.9. Here  $P' \in \text{NR}$  means that  $P'$  is a transition probability satisfying  $P'_\theta(f) = \int f(\omega) P'_\theta(d\omega)$  is constant on  $\Theta$  for any  $f$  in  $\mathcal{F}$  and  $\|\cdot\|_1$  stands for  $L_1$  norm on  $\Omega$ .

REMARK 2.10. Lemma 2.5 of ch. V consists of a weaker version of this proposition and could in fact be derived from it.

In order to prove the proposition, we are going to prove the following stronger result of which part 2 implies prop. 2.7.

PROPOSITION 2.8. (1) For each pair of probabilities  $P_1$  and  $P_2$  on  $(\Omega, \mathcal{A})$ ,

$$\min_{\tilde{P}_2 \sim P_2} \|P_1 - \tilde{P}_2\|_1 \leq R \max_{f \in \mathcal{F}} |P_1(f) - P_2(f)|$$

where  $\tilde{P}_2 \sim P_2$  means:  $\tilde{P}_2(f) = P_2(f)$ ,  $\forall f \in \mathcal{F}$ .

$$(2) \min_{P'_\theta \sim \mathbb{E} P_\theta} \mathbb{E}(\|P_\theta - P'_\theta\|_1) \leq \tilde{R} \max_f \min_z \mathbb{E}(|P_\theta(f) - z|).$$

Further, the constants  $R$  and  $\tilde{R}$  can be chosen to depend only on  $\Omega$  and  $\mathcal{F}$ .

PROOF. We apply ex. I.3Ex.4q p. 30 in the following setup: The closed convex polyhedron is the simplex  $\Delta$  of all probability measures on  $\Omega$ .  $E^n$  is the appropriate Euclidian space containing it.  $\varphi$  is the affine transformation defined (on the simplex of probabilities) by:  $\varphi(P) = (P(f))_{f \in \mathcal{F}}$ . Take the maximum norm in the range space of  $\varphi$ . Part 1 of the proposition follows now from the Lipschitz property of  $\varphi^{-1}$ . (Noticing that the left hand side of 1 is less than  $d(G(P_1), G(P_2))$ .)

To prove the second part of the proposition we use the first part for  $P_1 = P_\theta$  and  $P_2 = \mathbb{E}(P_\theta)$  to establish (using e.g. 7.j p. 427) the existence of  $P'_\theta \sim \mathbb{E}(P_\theta)$  such that:

$$\|P_\theta - P'_\theta\|_1 \leq R \max_f |P_\theta(f) - \mathbb{E}(P_\theta(f))|.$$

Taking expectation with respect to  $\theta$  and using lemma 2.4 we obtain:

$$\begin{aligned} \mathbb{E}(\|P_\theta - P'_\theta\|_1) &\leq R(\#F) \max_f \mathbb{E}(|P_\theta(f) - \mathbb{E}(P_\theta(f))|) \\ &\leq \tilde{R} \max_f \min_z \mathbb{E}(|P_\theta(f) - z|), \end{aligned}$$

where  $\tilde{R} = 2R(\#\mathcal{F})$  is again a constant depending on  $\Omega$  and  $\mathcal{F}$  only. This completes the proof of the proposition. ■

REMARK 2.11. Part 2 can in fact be improved (Mertens, 1973) by requiring further that  $\mathbb{E} P'_\theta = \mathbb{E} P_\theta$ .

**2.c. An auxiliary game.** By virtue of Dalkey's theorem (theorem 1.3 p. 53), we can assume from now on, without loss of generality, that no letter of  $B$  appears with positive probability in two different rows of  $Q$  (resp. no letter of  $A$  in two different columns). This situation can be achieved for instance, by replacing a letter  $a$  in the  $s^{\text{th}}$  row of  $Q$  by  $a_s$ . This modification of  $Q$  does not change  $\text{NR}^I$ ,  $\text{NR}^{II}$ ,  $D(p)$  or  $u(p)$  (by their definitions) and does not change  $v_n(p)$ ,  $\underline{v}(p)$  or  $\bar{v}(p)$  (by Dalkey's theorem). However it will enable us to identify rows of  $Q$  with random variables on  $T$  with values in  $A$ . To analyse the game it will be useful to introduce a “lower-game” (and dually an “upper-game”) having a simpler structure:

Define  $\underline{\Gamma}(p)$  as the repeated game obtained from our original  $\Gamma(p)$  by putting:

- (1)  $\underline{S} = S \cup S'$  (where  $S'$  is a copy of  $S$ ) is the action set of player I.
- (2)  $\underline{G}^k \left( \begin{smallmatrix} G^k \\ -|C| \end{smallmatrix} \right)$ ,  $k \in K$ , where  $|C|$  is a  $S' \times T$  matrix having the constant entry  $C = (\max\{|G_{st}^k| \mid s \in S, t \in T, k \in K\})$ .
- (3)  $\underline{Q}$  is a  $\underline{S} \times T$  matrix in which the entries are probability distributions on  $\underline{A} \times \underline{B} = (A \times \{a_0\}) \times (B \cup \tilde{B})$  where  $a_0 \notin A$  and  $\tilde{B}$  is a (large enough) set of signals disjoint from  $B$ , such that the following is satisfied:
  - $\forall s \in S, \forall t \in T, \forall b \in B, \underline{Q}_{s,t}(a, b) = \mathbb{1}_{a=a_0} Q_{s,t}(b)$ .
  - $\forall s \in S', \forall t \in T, \underline{Q}_{s,t}(a, b) = \mathbb{1}_{b=b(s,t,a)} Q_{s,t}(a)$ , where  $b(s, t, a) \in \tilde{B}$  and  $s' \neq s$  or  $t' \neq t$  or  $a' \neq a$  implies  $b(s, t, a) \neq b(s', t', a')$ .

- (4) I is restricted to play each of his additional pure strategies (i.e. those in  $S'$  numbered from  $(\#S + 1)$  to  $2(\#S)$ ) with probability  $\delta/(\#S)$  for some  $\delta > 0$  to be specified later.

In words,  $\underline{\Gamma}(p)$  differs from  $\Gamma(p)$  by the fact that I has to pay an amount  $C$  for hearing the signals in  $A$  induced by  $Q$ . Furthermore he is restricted to use this option of buying information exactly with probability  $\delta$  while with probability  $(1 - \delta)$  he gets no information whatsoever. Whenever he does receive a non-trivial information, his signal is completely known to II. Clearly this modification introduces asymmetry between the players (i.e. gives advantages to II). We do so in order to prove that I can guarantee what we claim to be the maxmin of the game, despite the disadvantageous modifications. Interchanging the rôles of the players in an obvious way would provide the dual statements that establishes minmax.

We will still use letters  $a$ ,  $a \in \underline{A}$ , (resp.  $b$ ,  $b \in \underline{B}$ ) for the signals to player I (resp. II) appearing in the support of  $\underline{Q}$ .

**2.d. The probabilistic structure.** From here on, let  $E^I$  be an index set of the same cardinality as  $K^I$ . An element of  $E^I$  will be denoted by  $e$ . As we shall see soon we will consider a special representation of the behavioural strategies of I in which, in each stage  $n$ , he performs a lottery to choose an element in  $E^I$ . The probability distribution on  $E^I$  may of course depend on all the information available to him. Then he will use a strategy in  $\text{NR}^I$  that may depend only on the outcome of the lottery (i.e., the element of  $E^I$  that was chosen) and the history of the game *not including his type* (the element of  $K^I$ ). This representation of the behavioural strategies of I plays a very crucial rôle in this chapter. In what follows we write formally the probabilistic structure of the game under such a strategy (and a strategy of II).

Let  $\Omega = K \times [E^I \times \underline{S} \times T \times \underline{A} \times \underline{B}]^{\mathbb{N}}$ . The  $n^{\text{th}}$  factor spaces  $E^I, \underline{S},$  etc. of  $\Omega$  will be denoted by  $E_n^I, \underline{S}_n,$  etc. respectively.

Unless otherwise specified, any set  $Z$  will be assumed to be endowed with its discrete  $\sigma$ -field  $\mathcal{P}(Z)$ . Let  $\mathcal{K} = \mathcal{P}(K)$ . The  $n^{\text{th}}$  factor space  $E^I \times \underline{S} \times T \times \underline{A} \times \underline{B}$  of  $\Omega$  represents therefore the five outcomes of stage  $n$ , namely the result of the lottery performed by I before the stage ( $E^I$ ), the pair of actions chosen by the two players ( $\underline{S} \times T$ ) and the pair of signals received at the end of the stage ( $\underline{A} \times \underline{B}$ ).

In what follows we introduce a certain number of  $\sigma$ -fields on  $\Omega$ . Any  $\sigma$ -field on a factor of  $\Omega$  will be identified with the corresponding  $\sigma$ -field on  $\Omega$ . For two  $\sigma$ -fields  $\mathcal{A}$  and  $\mathcal{B}$ ,  $\mathcal{A} \vee \mathcal{B}$  denotes the  $\sigma$ -field generated by  $\mathcal{A}$  and  $\mathcal{B}$ . Similarly, for a notation like  $\bigvee_{i=1}^n \mathcal{A}_i$ . If  $\mathcal{A}$  and  $\mathcal{B}$  are defined on two different spaces  $A$  and  $B$  then  $\mathcal{A} \otimes \mathcal{B}$  denotes the product  $\sigma$ -field on  $A \times B$ .

As a general guide and motivation for the definitions to come we adopt the following notational conventions:

- $\sigma$ -fields generated (among other things) by the signals from  $\underline{A}$  and  $\underline{B}$  will be denoted by the letter  $\mathcal{H}$ .
- A superscript I or II (or both) indicates that among the generators of the  $\sigma$ -field under consideration, are the announced signals of  $\underline{A}$  or  $\underline{B}$  (or both) respectively. Including  $\underline{A}$  will here also imply including  $\mathcal{P}(E^I)$ .
- A subscript I or II indicates that among the generators of the  $\sigma$ -field under consideration is the initial information  $\mathcal{K}^I$  (of I) or  $\mathcal{K}^{II}$  (of II) respectively.

- An index  $n$  indicates as usual that the  $\sigma$ -field under consideration corresponds to the situation before the  $n^{\text{th}}$  move in the game.
- Adding a  $\sim$  on top of a  $\sigma$ -field of index  $n$  indicates that we add  $\mathcal{P}(E_{n+1}^I)$  to the generators of the given  $\sigma$ -field. In other words, adding  $\sim$  corresponds to adding the outcome of the lottery of the next stage as an additional information.

Let us now define formally the  $\sigma$ -fields on  $\Omega$  that we shall need.

Denote by  $\underline{\mathcal{H}}_n^I$  and  $\underline{\mathcal{H}}_n^{II}$  the  $\sigma$ -fields generated by  $\underline{A}$  and  $\underline{B}$  respectively on the  $n^{\text{th}}$  factor space  $\underline{A}_n \times \underline{B}_n$ . Remark that by definition  $\underline{\mathcal{H}}_n^I \subseteq \underline{\mathcal{H}}_n^{II}$ . From these we construct the following  $\sigma$ -fields:

$\mathcal{H}_n^I := \bigvee_{i=1}^{n-1} (\underline{\mathcal{H}}_i^I \otimes \mathcal{P}(E_i^I))$ : All information collected by I before stage  $n$ , excluding the information about his own type.

$\mathcal{H}_n^{II} := \bigvee_{i=1}^{n-1} \underline{\mathcal{H}}_i^{II}$ : All information collected by II before stage  $n$ , excluding the information about his own type.

$\mathcal{H}_{I,n}^I := \mathcal{H}_n^I \vee \mathcal{K}^I$ : All information collected by I before stage  $n$ , including the information about his own type.

$\mathcal{H}_{II,n}^{II} := \mathcal{H}_n^{II} \vee \mathcal{K}^{II}$ : All information collected by II before stage  $n$ , including the information about his own type.

$\tilde{\mathcal{H}}_n^I := \mathcal{H}_n^I \vee \mathcal{P}(E_n^I)$ : All information available to I before move  $n$ , excluding the information about his own type, but including the outcome of the lottery at that stage.

$\tilde{\mathcal{H}}_{I,n}^I := \mathcal{H}_{I,n}^I \vee \mathcal{P}(E_n^I) = \tilde{\mathcal{H}}_n^I \vee \mathcal{K}^I$ : All information available to I before move  $n$ , including the information about his type and the outcome of the lottery.

$\tilde{\mathcal{H}}_{II,n}^{II} := \mathcal{H}_n^{II} \vee \mathcal{K}^{II}$ : All the information in  $\tilde{\mathcal{H}}_n^I$  plus the information about the type of II.

$\mathcal{H}_n^{I,II} := \mathcal{H}_n^I \vee \mathcal{H}_n^{II}$ : All information received by both players not including the information about their types.

$\mathcal{H}_{II,n}^{I,II} := \mathcal{H}_{II,n}^{II} \vee \mathcal{H}_n^I$ : All information received by both players plus the information about the type of II.

Finally, let  $\mathcal{G}$  be the  $\sigma$ -field generated by the moves and the outcomes of the lotteries before stage  $n$ , i.e.:

$$\mathcal{G}_n = [\bigvee_{i=1}^{n-1} \mathcal{P}(\underline{S}_i \times T_i \times E_i^I \times \underline{A}_i \times \underline{B}_i)] \vee \mathcal{K}; \quad \tilde{\mathcal{G}}_n = \mathcal{G}_n \vee \mathcal{P}(E_n^I).$$

By virtue of Dalkey's theorem we may assume that II uses a behavioural strategy, i.e. at every stage  $n$  he uses a transition probability  $\tau^n$  from  $(\Omega, \mathcal{H}_{II,n}^{II})$  to the  $n^{\text{th}}$  factor space  $T_n$ . As for I, the only strategies we will consider for him will be of the following type: At each stage  $n$  he uses first a transition probability from  $(\Omega, \mathcal{H}_{I,n}^I)$  to  $E_n^I$  and then he uses a  $\tilde{\mathcal{H}}_n^I$ -measurable function from  $\Omega$  to  $\text{NR}^I$  which selects a point in  $\underline{S}_n$ .

Given such a pair of strategies  $\sigma$  and  $\tau$  for the two players and the probability distribution  $p = (p^k)$  on  $\mathcal{K}$  of the initial choice of nature, the probability  $P_{\sigma,\tau}$  on  $(\Omega, \mathcal{G}_\infty)$  is completely defined by the following requirements:

- (1) The conditional distribution on  $E_n^I$  given  $\mathcal{G}_n$  is  $\mathcal{H}_{I,n}^I$ -measurable and is given by I's strategy (thus:  $E_n^I$  and  $\mathcal{G}_n$  are conditionally independent given  $\mathcal{H}_{I,n}^I$ ).
- (2)  $\underline{S}_n$  and  $T_n$  are conditionally independent given  $\tilde{\mathcal{H}}_n^I$ .
- (3) The conditional distribution on  $T_n$  given  $\tilde{\mathcal{G}}_n$  is  $\mathcal{H}_{II,n}^{II}$ -measurable and is given by II's strategy (thus  $T_n$  and  $\tilde{\mathcal{H}}_n^I$  are conditionally independent given  $\mathcal{H}_{II,n}^{II}$ ).

- (4) The conditional distribution on  $\underline{S}_n$  given  $\tilde{\mathcal{G}}_n$  is  $\tilde{\mathcal{H}}_{I,n}^I$ -measurable, and is given by I's strategy (thus:  $\underline{S}_n$  and  $\tilde{\mathcal{G}}_n$  are conditionally independent given  $\tilde{\mathcal{H}}_{I,n}^I$ ).
- (5) For any column of  $Q$ , the signal of player II, considered as a random variable on  $\underline{S}_n$  is conditionally independent of  $\tilde{\mathcal{G}}_n$  given  $\tilde{\mathcal{H}}_n^I$  (i.e., the conditional probability of  $b$  given  $\tilde{\mathcal{G}}_n$  is  $\tilde{\mathcal{H}}_n^I$ -measurable).
- (6) The distribution of the pair of signals, given the whole past, including stage 0 and including the last pair of moves  $(s, t)$ , is given by the transition probability  $\underline{Q}_{s,t}$  (as a function of those moves only).

When it is understood which are the underlying strategies  $\sigma$  and  $\tau$ , we shall omit the subscripts in  $P_{\sigma,\tau}$  and write  $(\Omega, \mathcal{G}_\infty, P)$  for the probability space generated by the strategies  $\sigma$  and  $\tau$  via the above six conditions.

**LEMMA 2.9.** *For any pair of strategies  $\sigma$  and  $\tau$  of the above specified types, there exists one and only one probability on  $(\Omega, \mathcal{G}_\infty)$  satisfying 1 to 6. And conversely, to every such probability there corresponds a pair of strategies of the specified type. This is no longer true if any one of the conditions is omitted.*

**PROOF.** See ex. VIEx.1 p. 314. ■

Given an increasing sequence of  $\sigma$ -fields  $(\mathcal{A}_n)_{n \in \overline{\mathbb{N}}}$ , (with  $\mathcal{A}_\infty = \vee_n \mathcal{A}_n$ ) on  $(\Omega, \mathcal{A})$  and a stopping time  $\theta$  w.r.t.  $(\mathcal{A}_n)_{n \in \overline{\mathbb{N}}}$  and with values in  $\overline{\mathbb{N}}$ , we define the  $\sigma$ -field  $\mathcal{A}_\theta$  by:

$$\mathcal{A}_\theta = \{ A \mid A \in \mathcal{A}, A \cap \{\theta \leq n\} \in \mathcal{A}_n \ \forall n \in \overline{\mathbb{N}} \}.$$

**REMARK 2.12.** Note that  $\{\theta = n\} \in \mathcal{A}_n \cap \mathcal{A}_\theta$ .

Define also for any  $(\tilde{\mathcal{H}}_{II,n}^I)$ -stopping time  $\theta$ ,  $\tilde{\mathcal{H}}_{II,\theta-}^I$  as the  $\sigma$ -field generated by  $\mathcal{K}^II$  and the sets  $A \cap \{n < \theta\} \ \forall A \in \tilde{\mathcal{H}}_{II,n}^I \ \forall n \in \overline{\mathbb{N}}$ . It should be thought of as the  $\sigma$ -fields of  $(\tilde{\mathcal{H}}_{II}^I)$ -events strictly before  $\theta$ .

**REMARK 2.13.** Note that  $\theta$  is  $(\tilde{\mathcal{H}}_{II,\theta-}^I)$ -measurable, that  $\tilde{\mathcal{H}}_{II,n-}^I = \tilde{\mathcal{H}}_{II,n-1}^I$  for  $n \geq 1$  and  $\tilde{\mathcal{H}}_{II,1-}^I = \mathcal{K}^II$ , and that the restrictions of  $\tilde{\mathcal{H}}_{II,\theta-}^I$  and  $\tilde{\mathcal{H}}_{II,n-}^I$  to  $\{\theta = n\}$  coincide.

We shall prove now an essential property of the probability space  $(\Omega, \mathcal{G}_\infty, P)$  generated by any pair of strategies  $\sigma$  and  $\tau$  of the above mentioned type.

From here on we will write  $p_n(k)$  for  $P(k \mid \mathcal{H}_n^I)$  and  $\tilde{p}_n(k)$  for  $P(k \mid \tilde{\mathcal{H}}_n^I)$ . The initial probability is thus  $p_1$ .

**PROPOSITION 2.10.** *For any probability satisfying 1 – 6 and for any  $(\tilde{\mathcal{H}}_{II,n}^I)$ -stopping time  $\theta$ ,  $\mathcal{H}_{II,\theta}^{I,II}$  and  $\mathcal{K}$  are conditionally independent given  $\tilde{\mathcal{H}}_{II,\theta-}^I$ .*

**COROLLARY 2.11.** *For any  $n$ , any  $k \in K$  and any  $\omega \in \Omega$ :*

$$\begin{aligned} p_{n+1}(k \mid \mathcal{K}^II)(\omega) &= P(k \mid a_1, \dots, a_n, e_1, \dots, e_n, \kappa^II) \\ &= P(k \mid b_1, \dots, b_n, e_1, \dots, e_n, \kappa^II) \\ &= P(k \mid b_1, \dots, b_{n-1}, e_1, \dots, e_n, \kappa^II) \\ &= P(k \mid a_1, \dots, a_{n-1}, e_1, \dots, e_n, \kappa^II) \\ &= \tilde{p}_n(k \mid \mathcal{K}^II)(\omega) \end{aligned}$$

where  $\kappa^II \in K^II$  and  $e_i \in E_i^I = 1, \dots, n$ .

PROOF OF THE COROLLARY.

- The first and last equalities are definitions.

- The equality of the third and the next to last term is the statement of prop. 2.10 for  $\theta \equiv n + 1$ , remembering that, in  $\underline{A}$ , “given  $b_i$ ” implies also “given  $a_i$ ”,  $i = 1, 2, \dots, n$ , (i.e. the letters  $a$  may be considered as forming a partition of the letters  $b$ ).
- The conditioning  $\sigma$ -fields in the second and in the fourth term are intermediate between those in the third and the next to last term. ■

PROOF OF PROP. 2.10. Let us first prove the proposition for a stopping time  $\theta$  which is constant:  $\theta \equiv n < +\infty$ . We do this by induction on  $n$ . We make again use of the fact that the letters  $a$  in  $\underline{A}$  form a partition of the letters  $b$  in  $\underline{B}$ . By induction hypothesis we have (cf. cor. 2.11):

$$P(k | b_1, \dots, b_n, e_1, \dots, e_n, \kappa^{\text{II}}) = P(k | a_1, \dots, a_n, e_1, \dots, e_n, \kappa^{\text{II}}),$$

and this holds obviously also for  $n = 1$ . On the other hand by condition 1 above:

$$P(e_{n+1} | b_1, \dots, b_n, e_1, \dots, e_n, k) = P(e_{n+1} | a_1, \dots, a_n, e_1, \dots, e_n, k).$$

These two relations imply:

$$P(e_{n+1}, k | b_1, \dots, b_n, e_1, \dots, e_n, \kappa^{\text{II}}) = P(e_{n+1}, k | a_1, \dots, a_n, e_1, \dots, e_n, \kappa^{\text{II}})$$

and thus  $P(k | b_1, \dots, b_n, e_1, \dots, e_{n+1}, \kappa^{\text{II}}) = P(k | a_1, \dots, a_n, e_1, \dots, e_{n+1}, \kappa^{\text{II}})$ , which is the conditional independence of  $\tilde{\mathcal{H}}_{\text{II},n+1}^{\text{I},\text{II}}$  and  $\mathcal{K}$  given  $\tilde{\mathcal{H}}_{\text{II},n+1}^{\text{I}}$ , or in other words,

$$(\star) \quad P(b_1, \dots, b_n | a_1, \dots, a_n, e_1, \dots, e_{n+1}, k) = P(b_1, \dots, b_n | a_1, \dots, a_n, e_1, \dots, e_{n+1}, \kappa^{\text{II}})$$

For any  $b$  in  $\underline{B}$  (resp  $a$  in  $A$ ), denote by  $t(b)$  (resp.  $s(a)$ ) the column (resp. the row) of  $\underline{Q}$  where the signal  $b$  (resp.  $a$ ) has a positive probability. Then

$$\begin{aligned} P(b_{n+1} | \tilde{\mathcal{G}}_{n+1}) &= \mathbb{E}\left(Q_{s_{n+1}, t_{n+1}}(b_{n+1}) | \tilde{\mathcal{G}}_{n+1}\right) \\ &= P(t_{n+1} = t(b_{n+1}) | \tilde{\mathcal{G}}_{n+1}) \mathbb{E}\left(Q_{s_{n+1}, t(b_{n+1})}(b_{n+1}) | \tilde{\mathcal{G}}_{n+1}\right) \\ &= P(t_{n+1} = t(b_{n+1}) | \mathcal{H}_{\text{II},n+1}^{\text{II}}) \mathbb{E}\left(Q_{s_{n+1}, t(b_{n+1})}(b_{n+1}) | \tilde{\mathcal{H}}_{n+1}^{\text{I}}\right) \quad (\text{by 2 and 6}) \\ &= P(b_{n+1} | \tilde{\mathcal{H}}_{\text{II},n+2}^{\text{I},\text{II}}) \quad (\text{by 3, 4 and 5}) \end{aligned}$$

Together with  $(\star)$  this implies

$$P(b_1, \dots, b_{n+1} | a_1, \dots, a_n, e_1, \dots, e_{n+1}, k) = P(b_1, \dots, b_{n+1} | a_1, \dots, a_n, e_1, \dots, e_{n+1}, \kappa^{\text{II}}),$$

which is the conditional independence of  $\mathcal{H}_{\text{II},n+2}^{\text{I},\text{II}}$  and  $\mathcal{K}$  given  $\tilde{\mathcal{H}}_{\text{II},n+1}^{\text{I}}$ .

This completes the proof of the proposition for a constant stopping time  $\theta \equiv n + 1$ . Let now  $\theta$  be any  $(\mathcal{H}_{\text{II},n}^{\text{I}})$ -stopping time. We want to show that  $P(k | \mathcal{H}_{\text{II},\theta}^{\text{I},\text{II}})$  is  $(\tilde{\mathcal{H}}_{\text{II},\theta}^{\text{I}})$ -measurable.

Let  $Z_n = P(k | \mathcal{H}_{\text{II},n}^{\text{I},\text{II}})$ ,  $X_n = P(k | \tilde{\mathcal{H}}_{\text{II},n}^{\text{I}})$ : we have just shown that  $Z_n = X_{n-1}$ . Since  $\theta$  is a  $(\mathcal{H}_{\text{II},n}^{\text{I},\text{II}})$ -stopping time as well, we have  $P(k | \mathcal{H}_{\text{II},\theta}^{\text{I},\text{II}}) = Z_\theta$ , hence equals  $X_{\theta-1}$ .

Thus we only have to show that, if  $X_n$  is an  $\mathcal{A}_n$ -adapted process (for  $\mathcal{A}_n = \tilde{\mathcal{H}}_{\text{II},n}^{\text{I}}$ ), converging a.e. to  $X_\infty$  (martingale convergence theorem), and if  $\theta$  is an  $\mathcal{A}_n$ -stopping time, then  $X_{\theta-1}$  is  $(\mathcal{A}_{\theta-1})$ -measurable: indeed this will imply that  $P(k | \mathcal{H}_{\text{II},\theta}^{\text{I},\text{II}})$  is  $(\tilde{\mathcal{H}}_{\text{II},\theta-1}^{\text{I}})$ -measurable, and since one checks immediately (on generators) that  $\tilde{\mathcal{H}}_{\text{II},\theta-1}^{\text{I}} \subseteq \mathcal{H}_{\text{II},\theta}^{\text{I},\text{II}}$ , it will indeed follow that  $P(k | \mathcal{H}_{\text{II},\theta}^{\text{I},\text{II}}) = P(k | \tilde{\mathcal{H}}_{\text{II},\theta-1}^{\text{I}})$ , hence the result.

Consider thus our adapted process  $X_n$  and let  $Y_n = X_{n-1}\mathbb{1}_{\theta=n}$ . Since  $X_{\theta-1} = \sum_{n \in \bar{\mathbb{N}}} Y_n$ , it suffices to show that  $Y_n$  is  $(\mathcal{A}_{\theta-})$ -measurable. From the characterisation in terms of generators, we know that  $X_{n-1}\mathbb{1}_{\theta \geq n}$  is  $(\mathcal{A}_{\theta-})$ -measurable, hence letting  $n \rightarrow \infty$ ,  $Y_\infty = X_\infty\mathbb{1}_{\theta=\infty}$  is so, and also — for  $X \equiv 1$  — we obtain the  $\mathcal{A}_{\theta-}$ -measurability of  $\theta$ , hence of  $\mathbb{1}_{\theta=n}$ , and therefore of  $Y_n = (X_{n-1}\mathbb{1}_{\theta \geq n})\mathbb{1}_{\theta=n}$ . This completes the proof of prop. 2.10. ■

Let now  $\theta$  be a  $\mathcal{H}_n^I$ -stopping time,  $\omega$  stand for a typical point in  $\tilde{\mathcal{H}}_\theta^I$  (i.e. a mapping from  $\Omega$  to  $\tilde{\mathcal{H}}_\theta^I$  associating to a point in  $\Omega$  the atom of  $\tilde{\mathcal{H}}_\theta^I$  containing it) and for each  $k \in K$ ,  $t \in T$ , let  $\bar{\tau}^k(\omega)(t) = P(t_{\theta+1} = t \mid \omega, k)$ .

PROPOSITION 2.12.  $\bar{\tau}^k(\omega)(t)$  is  $\mathcal{K}^{\text{II}}$ -measurable in  $k$  and  $\mathcal{H}_\theta^I$ -measurable in  $\omega$ .

PROOF.  $P(t_{\theta+1} = t \mid \tilde{\mathcal{H}}_\theta^I \vee \mathcal{K}) = \mathbb{E}[P(t_{\theta+1} = t \mid \mathcal{G}_\theta) \mid \tilde{\mathcal{H}}_\theta^I \vee \mathcal{K}]$ .

Now by condition 3 — extended to stopping times —  $P(t_{\theta+1} = t \mid \mathcal{G}_\theta)$  is  $(\mathcal{H}_{\text{II},\theta}^{\text{II}})$ -measurable. Therefore, by condition 1 — extended to stopping times —, we get that

$$\bar{\tau}^k(\omega)(t) = \mathbb{E}[P(t_{\theta+1} = t \mid \mathcal{H}_{\text{II},\theta}^{\text{II}}) \mid \tilde{\mathcal{H}}_\theta^I \vee \mathcal{K}]$$

is  $(\tilde{\mathcal{H}}_\theta^I \vee \mathcal{K})$ -measurable:

$$\bar{\tau}^k(\omega)(t) = \mathbb{E}[P(t_{\theta+1} = t \mid \mathcal{H}_{\text{II},\theta}^{\text{II}}) \mid \mathcal{H}_\theta^I \vee \mathcal{K}].$$

Note that  $\tilde{\mathcal{H}}_{\text{II},\theta-}^I \subseteq \mathcal{H}_{\text{II},\theta}^I \subseteq \mathcal{H}_{\text{II},\theta}^{I,\text{II}}$  as seen in the proof of prop. 2.10. Hence the corresponding result — in the form “ $P(k \mid \mathcal{H}_{\text{II},\theta}^{I,\text{II}})$  is  $(\tilde{\mathcal{H}}_{\text{II},\theta-}^I)$ -measurable” — remains a fortiori true when replacing  $\tilde{\mathcal{H}}_{\text{II},\theta-}^I$  by  $\mathcal{H}_{\text{II},\theta}^I$ :  $\mathcal{H}_{\text{II},\theta}^{I,\text{II}}$  and  $(\mathcal{K} \vee \mathcal{H}_\theta^I)$  are conditionally independent given  $\mathcal{H}_{\text{II},\theta}^I$ . Since the inner conditional expectation in our last formula for  $\bar{\tau}$  is  $(\mathcal{H}_{\text{II},\theta}^{I,\text{II}})$ -measurable, it follows that  $\bar{\tau}(w)(t)$  is  $(\mathcal{H}_{\text{II},\theta}^I)$ -measurable. In particular, for each  $\omega$ ,  $\bar{\tau}^k$  is  $\mathcal{K}^{\text{II}}$ -measurable w.r.t.  $k$ : hence  $\bar{\tau}$  defines a  $\mathcal{H}_\theta^I$ -measurable map from  $\Omega$  to strategies of player II in the one-shot game. ■

COMMENT 2.14. The interpretation of prop. 2.12 is that, if we imagine that I announces the outcomes of his lotteries to II, and that II uses  $\bar{\tau}$  as “strategy”, everything will be as if both players had the same signalling matrices.

Let  $\omega$  stand for a typical point in  $\tilde{\mathcal{H}}_n^I$ .

LEMMA 2.13. For any pair of strategies  $\sigma$  and  $\tau$  and for any  $n = 1, 2, \dots$ , and  $k \in K$ :

$$\mathbb{E}(|p_{n+1}(k) - \tilde{p}_n(k)| \mid \tilde{\mathcal{H}}_n^I)(\omega) = \tilde{p}_n(k) \frac{\delta}{\#S} \sum_{a \in A} |q_\tau^k(\omega, a) - \sum_{\ell \in K} \tilde{p}_n(\ell) q_\tau^\ell(\omega, a)|$$

where for each  $a$  in  $A$ :  $q_\tau^k(\omega, a) = \sum_{t \in T} \bar{\tau}^k(\omega)(t) \cdot Q_{s(a),t}(a)$ .

PROOF. Recall that by definition  $p_n(k) = P(k \mid \mathcal{H}_n^I)$  and  $\tilde{p}_n(k) = P(k \mid \tilde{\mathcal{H}}_n^I)$ . Thus:

$$p_{n+1}(k)(\omega, a_n) = \frac{\tilde{p}_n(k) P(a_n \mid \omega, k)}{\sum_{\ell \in K} \tilde{p}_n(\ell) P(a_n \mid \omega, \ell)}.$$

But:

$$P(a_n = a \mid \omega, k) = 1 - \delta \quad \text{if } a = a^0.$$

For  $a$  in  $A$  we obtain

$$\begin{aligned} P(a_n = a \mid \omega, k) &= \sum_{s \in S, t \in T} P(s_n = s, t_n = t, a_n = a \mid \omega, k) \\ &= \sum_{S \times T} \underline{Q}_{s,t}(a) \mathbb{E}(P(s_n = s, t_n = t \mid \tilde{G}_n) \mid \omega, k) \\ &= \sum_T \underline{Q}_{s(a),t}(a) \mathbb{E}(P(s_n = s(a) \mid \tilde{G}_n) P(t_n = t \mid \tilde{G}_n) \mid \omega, k) \\ &= (\delta/\#S) \sum_T \underline{Q}_{s(a),t}(a) \bar{\tau}^k(\omega)(t) = (\delta/\#S) q_\tau^k(\omega, a). \end{aligned}$$

Thus

$$P(a_n = a \mid \omega) = \begin{cases} 1 - \delta & \text{if } a = a^0 \\ (\delta/\#S) \sum_\ell \tilde{p}_n(\ell) q_\tau^\ell(\omega, a) & \text{if } a \neq a^0. \end{cases}$$

Taking expectation over all possible values of  $a_n$  concludes the proof of the lemma. ■

LEMMA 2.14. *For any strategy  $\tau$  of  $\Pi$  and any  $n$ :*

$$\inf_{\tilde{\tau} \in \text{NR}^{\text{II}}} \sum_{k \in K} \tilde{p}_n(k) \|\bar{\tau}^k(\omega) - \tilde{\tau}^k(\omega)\|_1 \leq R(\#S/\delta) \mathbb{E} \left( \sum_{k \in K} |p_{n+1}(k) - \tilde{p}_n(k)| \mid \tilde{\mathcal{H}}_n^{\text{I}} \right) (\omega)$$

where  $R$  is a constant that depends only on  $Q$ .

PROOF. Given  $\omega$  in  $\tilde{\mathcal{H}}_n^{\text{I}}$ , each of “strategies”  $\bar{\tau}^k(\omega)$  and  $\tilde{\tau}^k(\omega)$  is a transition probability from  $(K, \mathcal{H}^{\text{II}}, \tilde{p}_n)$  to  $\mathcal{P}(T \times A^S)$ . We can thus rewrite the left hand side of the formula as:

$$\inf_{\tilde{\tau} \in \text{NR}^{\text{II}}} \mathbb{E} \|\bar{\tau}^k(\omega) - \tilde{\tau}^k(\omega)\|_1 \stackrel{\text{def}}{=} (L)$$

and use prop. 2.7 with  $\Theta = K$ ,  $Q = \tilde{p}_n$ ,  $\Omega = T$  and each  $f$  being the probability of a signal  $a$ , i.e.  $f(t) = \underline{Q}_{s(a),t}(a)$ . We obtain

$$\begin{aligned} (L) &\leq R \max_{a \in A} \mathbb{E} |q_\tau^k(\omega, a) - \mathbb{E} q_\tau^k(\omega, a)| \\ &\leq R \sum_{k \in K} \sum_{a \in A} \tilde{p}_n(k) \left| q_\tau^k(\omega, a) - \sum_{\ell \in K} \tilde{p}_n(\ell) q_\tau^\ell(\omega, a) \right| \end{aligned}$$

and the result follows from lemma 2.13. ■

Let  $\rho_n$  denote the conditional expectation of the pay-off of stage  $n$ , given  $\tilde{\mathcal{H}}_n^{\text{I}}$ , in the game  $\Gamma$ .

LEMMA 2.15. *If at some stage  $n$ , player I uses after his lottery an optimal strategy in  $D(q_n)$ ,  $q_n$  being  $\tilde{\mathcal{H}}_n^{\text{I}}$ -measurable, then:*

$$\rho_n \geq u(q_n) - \frac{C(\#S)R}{\delta} \sum_{k \in K} \mathbb{E}(|p_{n+1}(k) - \tilde{p}_n(k)| \mid \tilde{\mathcal{H}}_n^{\text{I}}) - 2\delta C - C \sum_{k \in K} |\tilde{p}_n(k) - q_n^k|$$

(recall that  $C = \max_{s,t,k} |G_{s,t}^k|$  and  $R$  is the constant from Lemma 2.14.)

PROOF. Let  $\underline{\sigma} = \underline{\sigma}^k(\omega)$  be a strategy of I which is optimal in  $D(q_n)$  and let  $\tau$  be the strategy of II, then, since  $\underline{\sigma}G\tau$  is the conditional expected pay-off given  $\tilde{\mathcal{G}}_n$  (i.e. for  $\eta$  in  $\tilde{\mathcal{G}}_n$ , equals to  $\underline{\sigma}^{k(\eta)}(\eta)G^{k(\eta)}\tau^{k(\eta)}(\eta)$ ):

$$\rho_n = \mathbb{E} \left[ \mathbb{E} \left( \underline{\sigma}G\tau \mid \tilde{\mathcal{H}}_n^{\text{I}} \vee \mathcal{K} \right) \mid \tilde{\mathcal{H}}_n^{\text{I}} \right] = \mathbb{E} \left[ \underline{\sigma}G\bar{\tau} \mid \tilde{\mathcal{H}}_n^{\text{I}} \right]$$

or:

$$\rho_n(\omega) = \sum_{k \in K} \tilde{p}_n(k) \underline{\sigma}^k(\omega) G^k \bar{\tau}^k(\omega).$$

Recall (prop. 2.12) that  $\bar{\tau}^k$  is  $\mathcal{K}^{\text{II}}$ -measurable in  $k$ , and that by definition of  $\underline{\Gamma}_n(p)$ ,  $\underline{\sigma}^k$  consists of playing with probability  $(\delta/\#S)$  each of the additional rows, and with probability  $(1 - \delta)$  some strategy  $\sigma^k$  in the upper  $\#S$  rows. Let  $\tilde{\tau}^k(\omega) \in \text{NR}^{\text{II}}$ .

$$\begin{aligned}\rho_n(\omega) &\geq (1 - \delta) \sum_k \tilde{p}_n(k) \sigma^k(\omega) G^k \tilde{\tau}^k(\omega) \\ &\quad - (1 - \delta) C \sum_k \tilde{p}_n(k) \sum_{t \in T} |\bar{\tau}^k(\omega)(t) - \tilde{\tau}^k(\omega)(t)| - \delta C \\ &\geq (1 - \delta) \sum_k q_n^k \sigma^k(\omega) G^k \tilde{\tau}^k(\omega) - C \sum_k |\tilde{p}_n(k) - q_n^k| \\ &\quad - C \sum_k \tilde{p}_n(k) \|\bar{\tau}^k(\omega) - \tilde{\tau}^k(\omega)\|_1 - \delta C,\end{aligned}$$

and therefore since  $\sigma^k(\omega)$  is optimal in  $D(q_n)$ ,

$$\begin{aligned}\rho_n(\omega) &\geq u(q_n) - C \sum_k \tilde{p}_n(k) \|\bar{\tau}^k(\omega) - \tilde{\tau}^k(\omega)\|_1 \\ &\quad - 2\delta C - C \sum_k |\tilde{p}_n(k) - q_n^k|.\end{aligned}$$

Applying lemma 2.14 we obtain the required inequality.  $\blacksquare$

LEMMA 2.16. *For a real valued function  $f$  on  $\Pi$  which is convex w.r.t.  $\Pi$*

$$\mathsf{E}(f(p_{n+1}) \mid \tilde{\mathcal{H}}_n^{\text{I}}) \geq f(\tilde{p}_n) \quad \forall n$$

PROOF. Since by definition  $\mathsf{E}(p_{n+1} \mid \tilde{\mathcal{H}}_n^{\text{I}}) = \tilde{p}_n$ , the proof is just an application of Jensen's inequality to the convex function  $f(p)$  provided we prove that  $p_{n+1} \in \Pi^{\text{II}}(\tilde{p}_n)$ . (E.g. ex. I.3Ex.14biii p. 38 — the required measurability of  $f$  follows since the distribution of  $(\tilde{p}_n, p_{n+1})$  is discrete.) We have thus to show that  $p_{n+1}(k) = g(k)\tilde{p}_n(k)$  with  $g(k)$   $(\mathcal{K}^{\text{II}} \vee \tilde{\mathcal{H}}_n^{\text{I}})$ -measurable.

Now this follows from the explicit expression of  $p_{n+1}(k)$  as given in the proof of lemma 2.13 recalling that  $\bar{\tau}^k(\omega)$  is  $\mathcal{K}^{\text{II}}$ -measurable (prop. 2.12).  $\blacksquare$

### 3. The Infinite Game

**3.a. Minmax and Maxmin.** We are now ready for the first of the two main results of this chapter, namely to prove the existence and to characterise the minmax and the maxmin of the infinite game  $\Gamma_\infty(p)$ .

For any pair of strategies  $\sigma$  and  $\tau$  in  $\Gamma_\infty(p)$  and for any positive integer  $n$  we denote as usual by  $\bar{\gamma}_n(\sigma, \tau)$  the expected average pay-off for the first  $n$  stages, i.e.  $\bar{\gamma}_n(\sigma, \tau) = \mathsf{E}(\frac{1}{n} \sum_{m=1}^n G_{s_m, t_m}^k)$ , where  $\mathsf{E}$  is the expectation with respect to the probability measure induced by  $\sigma, \tau$  and  $p$ .

THEOREM 3.1. *The Minmax of  $\Gamma_\infty$  exists and is given by:*

$$\bar{v}(p) = \mathsf{Vex}_{\text{II}} \mathsf{Cav}_{\text{I}} u(p).$$

Obviously a dual result interchanging the rôles of the players establishes that  $\mathsf{Cav}_{\text{I}} \mathsf{Vex}_{\text{II}} u$  is the Maxmin of  $\Gamma_\infty$ .

PROOF. The proof is split into 3 parts. We first prove some results on strategies in  $\Gamma_\infty$ , then we show that I can defend  $\mathsf{Vex}_{\text{II}} \mathsf{Cav}_{\text{I}} u$  and finally that II can guarantee the same amount (cf. definition 1.2 p. 149).

PART A. Preliminary results.

For any strategy  $\sigma$  of I, for any time  $n$  and for any  $e \in E^I$ , denote by  $\sigma_{n,e}$  the strategy (i.e. the set of transition probabilities) of I that coincides with  $\sigma$  except that at time  $n$ , where  $P(e_n = e) = 1$ . In other words,  $\sigma_{n,e}$  is the same as  $\sigma$  except that at the lottery at stage  $n$ ,  $e$  is chosen deterministically independently of the history.

**LEMMA 3.2.** *For any strategies  $\sigma$  of I and  $\tau$  of II, for any time  $n$  and  $e \in E^I$ , the conditional probability distribution given  $\tilde{\mathcal{H}}_{I,n}^I$  induced by  $P_{\sigma,\tau}$  and  $P_{\sigma_{n,e},\tau}$  on  $\mathcal{G}_\infty$  coincide on  $\{e_n = e\}$ .*

**PROOF.** Let  $k \in \kappa^I \in K^I$ ; we have to show that for any  $m \geq n$  the probability

$$P(k; e_1, s_1, t_1, a_1, b_1, \dots, e_m, s_m, t_m, a_m, b_m, e_{m+1} \mid \kappa^I; e_1, a_1, e_2, a_2, \dots, e_{n-1}, a_{n-1}, e_n)$$

does not depend on whether  $P$  stands for  $P_{\sigma,\tau}$  or  $P_{\sigma_{n,e},\tau}$  (since this means coincidence of the two conditional probabilities on  $\tilde{\mathcal{G}}_m$  for all  $m$ , and hence on  $\mathcal{G}_\infty$ ). Using inductively 1–4 p. 278 this statement can be reduced to the case where  $m = n$ , i.e. to:

$$P(k; e_1, s_1, t_1, a_1, b_1, \dots, e_{n-1}, s_{n-1}, t_{n-1}, a_{n-1}, b_{n-1}, e_n \mid \kappa^I; a_1, e_2, a_2, \dots, e_{n-1}, a_{n-1}, e_n),$$

which equals to

$$P(k; e_1, s_1, t_1, a_1, b_1, \dots, e_{n-1}, s_{n-1}, t_{n-1}, a_{n-1}, b_{n-1} \mid \kappa^I; a_1, e_2, a_2, \dots, e_{n-1}, a_{n-1}, e_n),$$

which by 1 equals

$$P(k; e_1, s_1, t_1, a_1, t_1, \dots, e_{n-1}, s_{n-1}, t_{n-1}, a_{n-1}, b_{n-1} \mid \kappa^I; e_1, a_1, \dots, e_{n-1}, a_{n-1}).$$

The result now follows from the fact that  $P_{\sigma,\tau}$  and  $P_{\sigma_{n,e},\tau}$  coincide on  $\mathcal{G}_n$ . ■

Define now  $\text{NR}_\infty^I$  to be the set of strategies of I such that for every  $n$  and every  $e \in E^I$ ,  $P(e_n = e) \in \{0, 1\}$ .

Given a strategy  $\tau$  of II and  $\eta > 0$ , define also  $\sigma_0 \in \text{NR}_\infty^I$  and  $N$  by:

$$\mathbb{E}_{\sigma_0,\tau} \left[ \sum_K \sum_{n < N} (p_{n+1}(k) - p_n(k))^2 \right] > \sup_{\sigma \in \text{NR}_\infty^I} \mathbb{E}_{\sigma,\tau} \left[ \sum_K \sum_n (p_{n+1}(k) - p_n(k))^2 \right] - \eta$$

Note that for any  $\sigma, \tau$ , and  $p$ ,  $\{p_n\}_{n=1}^\infty$  is a martingale bounded in the simplex  $\Pi$  which implies that the sum of squares  $\sum_K \sum_{n=1}^\infty (p_{n+1}(k) - p_n(k))^2$  has expectation  $\leq 1$ , and hence  $\sigma_0$  is well defined.

In words,  $\text{NR}_\infty^I$  is the set of strategies of I which actually use no lotteries and hence is a sequence of one-stage non-revealing strategies (i.e. in  $\text{NR}^I$ ). Interpreting  $(p_{n+1}(k) - p_n(k))^2$  as a measure of the information revealed at stage  $n$ , we may interpret  $\sigma_0$  as the non-revealing strategy which exhausts the largest (up to  $\eta$ ) amount of information that can be exhausted from  $\tau$  by I without revealing anything itself.

It follows from Hölder's inequality and the definitions of  $\text{NR}_\infty^I$  (which implies that  $\tilde{p}_n = p_n$  a.s.),  $\sigma_0$  and  $N$  that:

**LEMMA 3.3.** *For any strategy  $\sigma$  in  $\text{NR}_\infty^I$  that coincides with  $\sigma_0$  up to stage  $N - 1$  and for any  $n \geq N$ :*

$$\mathbb{E} \left[ \sum_{k \in K} \left| P(k \mid \tilde{\mathcal{H}}_n^I) - P(k \mid \tilde{\mathcal{H}}_N^I) \right| \right] \leq \sqrt{\eta(\#K)}.$$

**LEMMA 3.4.** *Let  $\sigma$  be any strategy of I which coincides with  $\sigma_0$  up to stage  $N - 1$  and such that for all  $\tau$ , for all  $e \in E^I$ , and for all  $n \neq N$ ,  $P(e_n = e) \in \{0, 1\}$ . Then for any  $n \geq N$ :*

$$\mathbb{E} \left( \sum_{k \in K} |\tilde{p}_n(k) - \tilde{p}_N(k)| \right) \leq 2\#K^I \sqrt{\eta(\#K)}.$$

PROOF. As usually let  $P$  be the probability distribution determined by  $\sigma, \tau$  and  $p$ . By lemma 2.3 p. 274 applied conditionally to  $\tilde{\mathcal{H}}_N^I$  (as  $\mathcal{K}$  plays the rôle of  $\mathcal{Y}$  and  $\tilde{\mathcal{H}}_n^I$  plays the rôle of  $\mathcal{X}$ ) we have, denoting  $\|\cdot\|_{\tilde{\mathcal{H}}_n^I}$  simply as  $\|\cdot\|_n$ :

$$\begin{aligned} \mathbb{E} \left( \sum_k \|P(k | \tilde{\mathcal{H}}_n^I) - P(k | \tilde{\mathcal{H}}_N^I)\| \mid \tilde{\mathcal{H}}_N^I \right) &= \sum_k P(k | \tilde{\mathcal{H}}_N^I) \left\| P(\cdot | \tilde{\mathcal{H}}_N^I, k) - P(\cdot | \tilde{\mathcal{H}}_N^I) \right\|_n \\ &= \sum_k P(k | \tilde{\mathcal{H}}_N^I) \left\| P(\cdot | \tilde{\mathcal{H}}_N^I, k) - \sum_{\ell \in K} P(\ell | \tilde{\mathcal{H}}_N^I) P(\cdot | \tilde{\mathcal{H}}_N^I, \ell) \right\|_n \end{aligned}$$

By lemma 2.4 p. 275, applied conditionally to  $\tilde{\mathcal{H}}_N^I$  we get therefore:

$$\begin{aligned} (1) \quad \mathbb{E} \left( \sum_k \|P(k | \tilde{\mathcal{H}}_n^I) - P(k | \tilde{\mathcal{H}}_N^I)\| \mid \tilde{\mathcal{H}}_N^I \right) \\ \leq 2 \sum_k P(k | \tilde{\mathcal{H}}_N^I) \left\| P(\cdot | \tilde{\mathcal{H}}_N^I, k) - \sum_{\ell \in K} P(\ell | \tilde{\mathcal{H}}_N^I) P(\cdot | \tilde{\mathcal{H}}_N^I, \ell) \right\|_n \end{aligned}$$

Let now  $(\omega, e)$  stand for a typical element in  $\tilde{\mathcal{H}}_N^I \times E^I = \tilde{\mathcal{H}}_N^I$  and define

$$X(k, \omega, e) = \left\| P(\cdot | \omega, e, k) - \sum_{\ell} P(\ell | \omega) P(\cdot | \omega, e, \ell) \right\|_n.$$

From the given strategy  $\sigma$  we derive the strategies  $\sigma_e = \sigma_{N,e}, e \in E^I$ . Note that  $\sigma_e \in \text{NR}_\infty^I$  for all  $e \in E^I$ . Using lemma 3.2 we have also that, if  $P'$  denotes  $P_{\sigma_e, \tau}$ :

$$X(k, \omega, e) = \left\| P'(\cdot | \omega, e, k) - \sum_{\ell} P'(\ell | \omega) P'(\cdot | \omega, e, \ell) \right\|_n$$

and by definition of  $\sigma_e$ :

$$(2) \quad X(k, \omega, e) = \left\| P'(\cdot | \omega, k) - \sum_{\ell} P'(\ell | \omega) P'(\cdot | \omega, \ell) \right\|_n.$$

Note that in  $\sigma_e$ , the lottery before move  $N$  is eliminated, therefore conditioning on  $\tilde{\mathcal{H}}_N^I$  or  $\tilde{\mathcal{H}}_N^I$  are equivalent.

Let now:

$$Y(\omega, e) = \sum_k P'(k | \omega) X(k, \omega, e).$$

But  $P' = P_{\sigma_{N,e}, \tau}$  and  $P = P_{\sigma, \tau}$  coincide on  $\mathcal{G}_n$ , therefore

$$Y = \sum_k P(k | \omega) X(k, \omega, e).$$

Rewrite now (1) as:

$$\mathbb{E} \left( \sum_k |\tilde{p}_n(k) - \tilde{p}_N(k)| \mid \omega, e \right) \leq 2 \sum_k P(k | \omega, e) X(k, \omega, e).$$

Taking conditional expectation over  $e$  given  $\omega \in \tilde{\mathcal{H}}_N^I$  we get

$$\begin{aligned} \mathbb{E} \left( \sum_k |\tilde{p}_n(k) - \tilde{p}_N(k)| \mid \omega \right) &\leq 2 \sum_e \sum_k P(k, e | \omega) X(k, \omega, e) \\ &\leq 2 \sum_k P(k | \omega) \sum_e X(k, \omega, e) = 2 \sum_e Y(\omega, e). \end{aligned}$$

But by the definition of  $Y(\omega, e)$  and (2):

$$Y(\omega, e) = \sum_k P'(k | \omega) \left\| P'(\cdot | \omega, k) - \sum_{\ell} P'(\ell | \omega) P'(\cdot | \omega, \ell) \right\|_n.$$

Applying lemma 2.3 p. 274 conditionally to  $\mathcal{H}_N^I$  with  $\mathcal{K}$  in the rôle of  $\mathcal{Y}$  and  $\tilde{\mathcal{H}}_N^I$  in the rôle of  $\mathcal{X}$ , we obtain:

$$\begin{aligned} Y(\omega, e) &= \mathbb{E}\left(\sum_k \left|P'(k \mid \tilde{\mathcal{H}}_n^I) - P'(k \mid \omega)\right| \mid \omega\right) \\ &= \mathbb{E}\left(\sum_k |\tilde{p}'_n(k) - \tilde{p}'_N(k)| \mid \omega\right), \end{aligned}$$

where  $\tilde{p}'_n$  and  $\tilde{p}'_N$  are the probabilities derived from the strategy  $\sigma_e$ , which is a strategy in  $\text{NR}_{\infty}^I$  that coincides with  $\sigma_0$  up to state  $N - 1$ ; therefore by lemma 3.3:

$$\mathbb{E}(Y(\omega, e)) \leq \sqrt{\eta(\#K)} \quad \forall e \in E^I$$

and finally by taking expectations over  $\mathcal{H}_n^I$  in (2) we get:

$$\mathbb{E}\left(\sum_k |\tilde{p}_n(k) - \tilde{p}_N(k)|\right) \leq 2 \sum_e \sqrt{\eta(\#K)} = 2\#K^I \sqrt{\eta(\#K)}$$

which concludes the proof of lemma 3.4. ■

### PART B. Player I can defend $\mathbf{Vex Cav u}$

For any given strategy  $\tau$  of II and  $\eta > 0$ , let  $\sigma_0$  and  $N$  be defined as in lemma 3.3 p. 284 and consider the following strategy  $\sigma$  of I:

- Play  $\sigma_0$  up to stage  $N - 1$ .
- Use the transition probability described in lemma 2.6 p. 275 with  $p = p_N$ ,  $g(p) = \mathbf{Cav}_I u$  and  $\varepsilon = 0$ , to choose  $e_N \in E_N^I$ .
- After stage  $N$  play at every stage independently an optimal strategy in  $D(\tilde{p}_N)$ .

For  $n \geq N$  denote as usual by  $\rho_n$  the conditional expected pay-off at stage  $n$  given  $\tilde{\mathcal{H}}_n^I$ . By lemma 2.15 p. 282 we have:

$$\rho_n \geq u(\tilde{p}_N) - \frac{C(\#S)R}{\delta} \sum_k \mathbb{E}(|p_{n+1}(k) - \tilde{p}_n(k)| \mid \tilde{\mathcal{H}}_n^I) - C \sum_k |\tilde{p}_N(k) - \tilde{p}_n(k)| - 2\delta C$$

Remark that for any pair of strategies  $\{p_1, \tilde{p}_1, p_2, \tilde{p}_2, \dots, p_n, \tilde{p}_n, \dots\}$  is a martingale in II. By construction of  $\sigma$ ,  $\tilde{p}_n = p_n$  for  $n \neq N$ , and for  $n = N$  we have by lemma 2.6 p. 275:

$$(3) \quad \mathbb{E}(u(\tilde{p}_N) \mid \tilde{\mathcal{H}}_N^I) \geq (\mathbf{Cav}_I u)(p_N) \geq (\mathbf{Vex Cav}_I u)(p_N).$$

Now by lemma 2.16 p. 283 applied to  $f(p) = (\mathbf{Vex}_{II} \mathbf{Cav}_I u)(p)$  we have that for all  $n$ ,

$$\mathbb{E}\left((\mathbf{Vex}_{II} \mathbf{Cav}_I u)(p_{n+1}) \mid \tilde{\mathcal{H}}_n^I\right) \geq (\mathbf{Vex}_{II} \mathbf{Cav}_I u)(\tilde{p}_n).$$

It follows that  $(\mathbf{Vex} \mathbf{Cav} u)(p), (\mathbf{Vex} \mathbf{Cav} u)(\tilde{p}_1), \dots, (\mathbf{Vex} \mathbf{Cav} u)(p_m), (\mathbf{Vex} \mathbf{Cav} u)(\tilde{p}_m), \dots, (\mathbf{Vex} \mathbf{Cav} u)(p_N)$  is a submartingale and hence (3) yields:

$$\mathbb{E}(u(\tilde{p}_N)) \geq (\mathbf{Vex}_{II} \mathbf{Cav}_I u)(p).$$

Therefore (using lemma 3.4 p. 284):

$$\mathbb{E}(\rho_n) \geq (\mathbf{Vex}_{II} \mathbf{Cav}_I u)(p) - \frac{C(\#S)R}{\delta} \sum_K \mathbb{E}|p_{n+1}(k) - \tilde{p}_n(k)| - 2C(\#K^I) \sqrt{\eta(\#K)} - 2C\delta$$

Summing on  $n$  from  $N$  to  $N + m$ , dividing by  $N + m$  and recalling that (cf. lemma 2.1 p. 186)  $\sum_{k \in K} \frac{1}{m} \sum_{n=N}^{m+N} \mathbb{E}(|p_{n+1}(k) - \tilde{p}_n(k)|) \leq \sqrt{\frac{\#K-1}{m}}$ , we get:

$$\bar{\gamma}_{N+m}(\sigma, \tau) \geq \frac{-2CN}{N+m} + (\mathbf{Vex}_{II} \mathbf{Cav}_I u)(p) - 2C\delta - 2C(\#K^I) \sqrt{\eta(\#K)} - \frac{C(\#S)R}{\delta} \sqrt{\frac{\#K}{m}}.$$

Finally, for each  $\varepsilon$  we may choose  $\eta$  and  $\delta$  small enough and then  $N_0$  big enough as to have

$$\bar{\gamma}_n(\sigma, \tau) > (\underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} u)(p) - \varepsilon \quad \forall n > N_0.$$

This completes the second part of the proof of the theorem.

### PART C. Player II can guarantee $\underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} u$

The proof of this part is derived from the results on the value of games with lack of information on one side (theorem 3.5 p. 195).

Observe first that if we add as additional columns in  $T$  all extreme points of  $\text{NR}^{\text{II}}$ , and we define the corresponding columns of pay-offs and signals in the obvious way, the game  $\Gamma_\infty(p)$  is actually unchanged. However by doing so, the new set  $\text{NR}^{\text{II}}$  becomes essentially the set of constant strategies (independent of  $k$ ).

Now we make the game less favourable to player II by replacing his signals on these additional columns by a constant letter. In the new game, the distribution on signals is still independent of the state of nature.

Thus if II ignores his private information (i.e.  $\kappa^{\text{II}}$ ) and if for each  $\kappa^{\text{I}} \in K^{\text{I}}$  we let  $q^{\kappa^{\text{I}}} = \sum_{k \in \kappa^{\text{I}}} p^k$  and take as pay-offs  $A^{\kappa^{\text{I}}} = \frac{1}{q^{\kappa^{\text{I}}}} \sum_{k \in \kappa^{\text{I}}} p^k G^k$  (and keep the same distribution on signals) we obtain a game  $\Gamma^*$  with incomplete information on one side, with  $K^{\text{I}}$  as the set of states of nature,  $q$  as initial probability distribution on it and player I informed. In this game consider the set  $\text{NR}(q)$  of non-revealing strategies of player I (as defined in ch. V) and remark that its projection on the support of  $q$  equals the corresponding projection of  $\text{NR}^{\text{I}}$  (as defined sub 2.a.1 p. 273). Indeed this is true even if the additional columns are deleted. Letting  $w'(q)$  (resp.  $w(q)$ ) be the value of the one-shot game where player I plays in  $\text{NR}(q)$ , resp.  $\text{NR}^{\text{I}}$ , we thus have  $w'(q) = w(q)$  for all  $q$ . By theorem 3.5 p. 195, the value of the game  $\Gamma$  is  $\text{Cav } w'$ .

Now by our construction  $w(q) = u(p)$  and  $\text{Cav } w'(q) = \text{Cav}_I u(p)$  so by theorem 3.5, for each  $p$ , II has a strategy  $\tau(p)$  and for each  $\varepsilon > 0$  there is  $N$  s.t.  $\bar{\gamma}_n(\sigma, \tau(p)) < (\text{Cav}_I u)(p) + \varepsilon$ , for all  $n > N$  and for all  $\sigma$  of I.

By lemma 2.6 p. 275 (or rather its dual for II with  $\text{Cav}_I u$  in the place of  $g$  and  $f$ ), II has a transition probability from  $K^{\text{II}}$  to  $E^{\text{II}}$  (of the same cardinality as  $K^{\text{II}}$ ) such that if  $p_e$  is the conditional probability on  $K$  given the outcome  $e \in E^{\text{II}}$  then:

$$\mathsf{E}(\tilde{p}_e) = p \quad \text{and} \quad \mathsf{E}\left((\underset{\text{I}}{\text{Cav}} u)(p_e)\right) = (\underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} u)(p).$$

The desired strategy  $\tau$  of II can now be described as follows: Use the above described transition probability to choose  $e \in E^{\text{II}}$ . If the outcome is  $e$ , play from there on the strategy  $\tau(p_e)$  (to guarantee  $(\text{Cav}_I u)(p_e)$ ). It follows that for each  $\varepsilon > 0$  we have with  $N = \max N(p_e)$ :

$$\bar{\gamma}(\sigma, \tau) < (\underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} u)(p) + \varepsilon, \quad \forall n > N \quad \forall \sigma.$$

This completes the proof of C and hence of theorem 3.1. ■

An immediate consequence of the theorem is:

COROLLARY 3.5.  $\Gamma_\infty(p)$  has a value if and only if

$$\underset{\text{I}}{\text{Cav}} \underset{\text{II}}{\text{Vex}} u = \underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} u$$

and then this is the value.

Exercises VIEx.3 p. 314 and VIEx.4 p. 314 will illustrate that  $\Gamma_\infty$  generally has no value.

### 3.b. Approachability.

3.b.1. *The finite case.* Another consequence of theorem 3.1 p. 283 — and a sharpening, needed in ch. IX — is the following characterisation of approachable vectors. Let, for  $\kappa \in K^I$ ,  $\bar{\gamma}_n(\sigma, \tau, \kappa) = E_{p, \sigma, \tau}(\bar{g}_n | \kappa)$ , where  $\bar{g}_n$  is the average pay-off over the first  $n$  stages.

**DEFINITION 3.6.** Let  $\bar{Z}_p = \{ z \in \mathbb{R}^{K^I} \mid \exists \tau: \forall \sigma, \forall \kappa, \limsup_{n \rightarrow \infty} \bar{\gamma}_n(\sigma, \tau, \kappa) \leq z^\kappa \}$ , and  $\underline{Z}_p = \{ z \in \mathbb{R}^{K^I} \mid \exists \tau: \forall \sigma, \exists \mathcal{L} \text{ (Banach limit)}: \forall \kappa, \mathcal{L}(\bar{\gamma}_n(\sigma, \tau, \kappa)) \leq z^\kappa \}$ .

**REMARK 3.1.** We will shortly show that  $\bar{Z}_p = \underline{Z}_p$ , hence the notation  $Z_p$ , and the name **approachable vectors** for its elements. Also, the existence of  $\mathcal{L}$  (ex. I.2Ex.13 p. 24) just means that the convex hull of the limit points of  $\bar{\gamma}_n(\sigma, \tau, \cdot)$  intersects  $\underline{Z}_p - \mathbb{R}_+^{K^I}$ , while for  $\bar{Z}_p$  one asks inclusion. I.e.,  $\bar{Z}_p$  could be defined in the same way, but with “ $\forall \mathcal{L}$ ”.

**DEFINITION 3.7.**  $W_p = \{ w \in \mathbb{R}^{K^I} \mid \langle \lambda, w \rangle \geq u(\lambda \cdot p) \quad \forall \lambda \in \Delta(K^I) \}$  with the notation  $(\lambda \cdot p)(k) = \sum_{\kappa \in K^I} \lambda_\kappa p(k | \kappa)$ .

**PROPOSITION 3.8.** (1)  $\underline{Z}_p = \bar{Z}_p = Z_p$  is closed and convex; more precisely

$$Z_p = \{ z \in \mathbb{R}^{K^I} \mid \langle \lambda, z \rangle \geq (\text{Vex}_{\text{II}} \text{Cav}_I u)(\lambda \cdot p) \quad \forall \lambda \in \Delta(K^I) \}.$$

(2)  $Z_p$  (and  $W_p$ ) are compactly generated:  $Z_p = (Z_p \cap [-C, C]^{K^I}) + \mathbb{R}_+^{K^I}$ .

(3)  $Z_p = \left\{ \left[ \frac{1}{p(\kappa)} \sum_{e \in E} \pi(e) p_e(\kappa) w_e^\kappa \right]_{\kappa \in K^I} \mid \pi \in \Delta(E); \forall e \in E, (w_e \in W_{p_e} \text{ and } p_e \in \Pi^{\text{II}}(p)); \sum_{e \in E} \pi(e) p_e = p \right\}$  where  $E$  is a set of cardinality  $(\#K^I) + (\#K^{\text{II}})$ .

(4)  $\min_{z \in Z_p} \langle \lambda, z \rangle = (\text{Vex}_{\text{II}} \text{Cav}_I u)(\lambda \cdot p) \quad \forall \lambda \in \Delta(K^I)$ .

**REMARK 3.2.** 3 says in particular that any approachable vector  $z \in Z_p$  can be approached by a “standard strategy” where player II first makes a type-dependent lottery on the set  $E$ , next, given  $e \in E$  and the corresponding posterior  $p_e$ , approaches  $w_e \in W_{p_e}$  with an approachability strategy like in 3.g p. 217 (independent of his type).

**PROOF.** Let the right-hand member sub 3 be denoted by  $Y_p$ . By the arguments of the end of the proof of C, we have  $Y_p \subseteq \bar{Z}_p \subseteq \underline{Z}_p$ . Further, denoting the right-hand member sub 1 by  $X_p$ , we have  $\underline{Z}_p \subseteq X_p$ , by B p. 286 (used at  $p' = \lambda \cdot p$  where there would be strict inequality). So to prove 1 and 3, it suffices to show that  $X_p \subseteq Y_p$ . Finally, the proof of C p. 287 also shows that player II can guarantee the minmax with standard strategies (and with  $\#E \leq \#K^I + \#K^{\text{II}}$ ) — hence 4:  $\min_{z \in Y_p} \langle \lambda, z \rangle = (\text{Vex}_{\text{II}} \text{Cav}_I u)(\lambda \cdot p)$  ( $= \varphi(\lambda)$ , the support function of  $X_p$ ). So if we show that  $Y_p$  is closed and convex,  $X_p \subseteq Y_p$  will follow by (1.21 p.8). For the convexity, observe first that  $Y_p$  would be convex if the cardinality of  $E$  was arbitrary and allowed to vary with the point  $y \in Y_p$  considered: indeed, in that case to obtain a convex combination  $\sum \alpha_i y_i$ , it would suffice for player II to make first a lottery (even type-independent) with weights  $\alpha_i$  to select some  $y_i$ , next to make the correspondent type-dependent lottery to choose  $e \in E_i$ , etc.: but the whole procedure is clearly equivalent to one single type-dependent lottery on  $\bigcup_i E_i$  (disjoint union). But with  $\#E$  arbitrary, the conditions just express that (given  $p_e \in \Pi^{\text{II}}(p)$  and  $w_e \in W_{p_e}$ )

$$\begin{aligned} \sum_e \pi(e) [p_e(\kappa) w_e^\kappa] &= p(\kappa) y^\kappa & \forall \kappa \in K^I, & \text{and} \\ \sum_e \pi(e) p_e(\kappa) &= p(\kappa) & \forall \kappa \in K^{\text{II}} \end{aligned}$$

The second set of equations also implies  $\sum \pi(e) = 1$ , so we have the right-hand member, as a vector in  $\mathbb{R}^{\#K^I + \#K^{\text{II}} - 1}$ , expressed as a convex combination of similar vectors in the

left-hand member: by Carathéodory (ex.I.3Ex.10 p.34, it suffices to have  $\pi(e) > 0$  for  $\#K^I + \#K^{II}$  values of  $e$  (Fenchel would already yield better)). Hence the convexity.

Thus there remains to prove 2, and the closure of  $Y_p$ . But if we show 2 for  $W_p$  it will immediately follow that  $Y_p = \overline{Y}_p + \mathbb{R}_+^{K^I}$ , when  $\overline{Y}_p$  is defined as  $Y_p$  but with, instead of  $W_{p_e}$ ,  $W_{p_e} \cap [-C, C]^{K^I}$ .  $\overline{Y}_p$  is then clearly compact, hence  $Y_p$  closed (and compactly generated), hence 1 and 3 hold, hence  $Z_p (= Y_p)$  is also compactly generated: it remains only to establish 2 for  $W_p$ . This also follows from the proof of theorem 3.1 part C, where we have shown that  $(\text{Cav}_I u)(\lambda \cdot p) = \psi(\lambda)$  was the value  $(\text{Cav } w)(\lambda)$  of a repeated game  $\Gamma$  with incomplete information on one side  $(A^\kappa)_{\kappa \in K^I}$ , where all pay-offs are  $\leq C$  in absolute value: hence — in order to deduce this from the above results, without additional argument — note that  $W_p \subseteq Y_p \subseteq \overline{Z}_p \subseteq X_p$ , for  $\Gamma$ , but for  $\Gamma$ ,  $X_p = W_p$  obviously (the  $\text{Vex}_{II}$  operation is the identity), so it suffices to establish b) for  $\overline{Z}_p$  and  $\Gamma$ , which is a trivial consequence of the definition. ■

**REMARK 3.3.** In the above it was apparently assumed that  $p(\kappa) > 0$ ,  $\forall \kappa \in K^I$ . However, the definitions of  $\overline{Z}_p$ ,  $\underline{Z}_p$  and  $W_p$  depend only on  $[p(k | \kappa)]_{\kappa \in K^I}$ . So the subscripts  $p$  should be interpreted as standing for such a conditional probability on  $K$  given  $\mathcal{K}^I$ . To make the proposition fully correct in those terms, just reinterpret the quantity  $\frac{1}{p(\kappa)} \pi(e) p_e(\kappa)$  sub 3 as  $\sum_{k \in K} p(k | \kappa) q_k(e)$ , with  $q_k \in \Delta(E)$   $\forall k \in K$  and  $q_k(e)$   $\mathcal{K}^{II}$ -measurable  $\forall e \in E$ . [In the proof (convexity of  $Y_p$ ) there is obviously no problem in assuming we have in addition some strictly positive probability on  $K^I$ .]

**COROLLARY 3.9.** If  $z \notin Z_p$ , there exists a compact, convex set  $C$  disjoint from  $z - \mathbb{R}_+^{K^I}$ , such that  $\forall \tau, \exists \sigma, \exists N > 0: \forall n \geq N, \overline{\gamma}_n(\sigma, \tau) \in C$ .

**PROOF.** Choose  $\lambda \in \Delta(K^I)$  with  $\langle \lambda, z \rangle + 2\varepsilon < (\text{Vex}_{II} \text{Cav}_I u)(\lambda \cdot p)$ , take as convex set  $\{x \in [-C, C]^{K^I} \mid \langle \lambda, x \rangle \geq \langle \lambda, z \rangle + \varepsilon\}$ , and use the proof of B. ■

**COROLLARY 3.10.**  $\text{Vex}_{II} \text{Cav}_I u$  is concave w.r.t. I.

**PROOF.** By prop. 3.8.4. ■

**3.b.2. More intrinsically — NR-strategies.** We look now for a different characterisation of  $W_p$  — by a set, with corresponding strategies, which are independent of  $p$ . This leads, at the same time, to extend the previous results to the case of a continuum of types of player I.

We consider thus the case where the set of states of nature is  $K \times I \times J$ , with a probability distribution  $\pi$  on it — and where the sets  $K$  and  $J$  are finite, while  $I$  can be an arbitrary measurable space. The pay-off matrix depends on  $K$ , while  $I$  (resp.  $J$ ) is the initial information of player I (resp. II).

Actually, we can therefore as in the previous chapter identify  $I$  with  $\Delta(K \times J)$  — and will use  $p$  for elements of  $\Delta(K \times J)$ , and  $\pi$  will become a probability distribution on  $P = \Delta(K \times J)$ .  $u$  becomes then a function on  $\Delta(P)$ .

**THEOREM 3.11.** (1) In the above framework, propositions 3.43 up to 3.46 p. 228 remain word for word true (with  $\#I = 1, \Delta(K_i) = P$ ), if in propositions 3.45 part 4 p. 231 and 3.46 part 2 p. 232 the strategies of player II are understood as non-revealing strategies in  $\Gamma_\infty$  and if the conclusion of prop. 3.46 part 2 is weakened as in 2 below.

(2) Assume here for notational simplicity that I's signals inform him of his last move.

- (a) For every non-revealing strategy  $\tau_h$  of player II, there exist random variables  $X_{h,n}: A^n \rightarrow \mathbb{R}^{K \times J}$  such that for all  $n > 0$ , all  $p \in \Delta(K \times J) = P$  and all strategies  $\sigma$  of player I

$$\mathsf{E}_{\sigma, \tau_h}^p(\bar{g}_n \mid \mathcal{H}_{I,n}^I) = \langle p, X_{h,n} \rangle.$$

Note thus  $X_{h,n}^{k,j} = \mathsf{E}_{\sigma, \tau_h}^q(\bar{g}_n^k \mid \mathcal{H}_{I,n}^I) = \mathsf{E}_{\tau_h^j}(\bar{g}_n^k \mid A^n)$  is independent of  $p$  and  $\sigma^p$ .

- (b) For every  $h \in \mathbf{H}$ , let  $\mathbf{L}_h = \{x \in \mathbb{R}^{K \times J} \mid \langle p, x \rangle \leq h(p) \forall p \in P\}$ . Then, for some sequence  $\eta_n$  decreasing to zero and that depends only on  $Q$  and on  $\#J$ , we have, for all  $h \in \mathbf{H}$ , all  $n$ , all strategies  $\sigma$  of player I, all  $p \in P$  and  $j \in J$ , that  $\mathsf{E}_{\sigma^p, \tau_h}[\sup_{m \geq n} d(X_{h,m}, \mathbf{L}_h)] \leq C\eta_n$ .

REMARK 3.4.  $\tau_h$  being non-revealing, the distribution of  $X_h$  under  $\sigma^p$  and  $\tau_h^j$  is independent of  $j \in J$ , hence the notation  $\mathsf{E}_{\sigma^p, \tau_h}$  instead of  $\mathsf{E}_{\sigma^p, \tau_h^j}$  or of  $\mathsf{E}_{\sigma, \tau_h}^{p,j}$ .

PROOF. 1: Use the reduction in the beginning of C. It is clear that the optimal strategy of player II obtained in proposition 3.46.2 after this reduction is non-revealing in  $\Gamma_\infty$ . As for 3.45.4, it is easiest to give a direct proof, using the techniques of the present chapter (prop. 2.12 p. 281 yields that  $\bar{\tau}(\omega)$  is in NR if  $\tau$  is in  $\text{NR}_\infty$ , also lemma 2.15 p. 282) — or, in order to use the statements themselves, first to discretise  $P$ , e.g. as in the proof of prop. 3.35 p. 221, making sure in addition that  $\sigma_p$  does not vary by more than  $\delta_n$  for  $p \in \ell \in L_n$  (where  $\sigma$  is a given NR-strategy of I).

2a: Since  $\tau_h$  is non-revealing, we have  $P_{\sigma, \tau_h}^q(k, j \mid \mathcal{H}_{I,n}^I) = q^{k,j}$ , and hence  $\mathsf{E}_{\sigma, \tau_h}^q(\bar{g}_n \mid \mathcal{H}_{I,n}^I) = \sum_{k,j} q^{k,j} \mathsf{E}_{\sigma^q, \tau_h^j}(\bar{g}_n^k \mid A^n)$ . Let thus  $\sigma_0 = \sigma^q$ ,  $\tau_0 = \tau_h^j$ : there remains to show that  $P_{\sigma_0, \tau_0}((\eta_1, \eta_2) \mid \eta_1)$  is independent of  $\sigma_0$ , for all  $\eta_1 \in A^n$ ,  $\eta_2 \in (T \times B)^n$ , and  $n > 0$ . (Because  $\bar{g}_n^k$  is a function of  $(\eta_1, \eta_2)$ ). Assume thus by induction the statement is proved for all smaller values of  $n$ , and let  $\eta_1 = (\xi_1, s, \alpha)$ ,  $\eta_2 = (\xi_2, t, \beta)$ . Then

$$\begin{aligned} P_{\sigma_0, \tau_0}((\eta_1, \eta_2) \mid \eta_1) &= P_{\sigma_0, \tau_0}(\xi_2, s, t, \alpha, \beta \mid \xi_1) / \sum_{\xi_2, t, \beta} P_{\sigma_0, \tau_0}(\xi_2, s, t, \alpha, \beta \mid \xi_1) \\ &= P_{\tau_0}((\xi_1, \xi_2) \mid \xi_1) P_{\sigma_0}(s \mid \xi_1) P_{\tau_0}(t \mid \xi_2) P((\alpha, \beta) \mid (s, t)) / \sum_{\xi_2, t, b} \dots \end{aligned}$$

and, after dividing numerator and denominator by the common factor  $P_{\sigma_0}(s \mid \xi_1)$  one sees that  $\sigma_0$  no longer appears in the expression.

[Recall from the proof of prop. 3.46.2 p. 232 that the  $\tau_h$  are completely mixed, so the  $X_{h,n}^{k,j}$  are indeed well defined over the whole of  $A^n$ . Observe further that, for fixed  $\omega \in \bigcup_n A^n$ , they take the form  $X_h^{k,j}(\omega) = \sum_{s,t} F_{s,t}^\omega(\tau_h^j) G_{s,t}^k$ , where  $F^\omega(\tau)$  is a probability distribution over  $S \times T$  varying continuously with  $\tau$ .]

2b: Fix a sequence  $\varepsilon_n$ , take the corresponding map  $\tau_h$  from prop. 3.46.2 p. 232 in the auxiliary game  $\Gamma^*$  of C, and the corresponding sequence  $\delta_n$ . Observe the matrix  $Q^*$  of  $\Gamma^*$  — and hence the sequence  $\delta_n$  — depends only on  $Q$  and on  $\#J$ . Fix  $h \in \mathbf{H}$ , and a type-independent strategy  $\sigma$  of player I (the strategy  $\sigma^p$  of the statement, for some fixed  $p$ ). Thus,  $\sigma$  is a transition probability from  $\bigcup_n A^n$  to  $S$ . It clearly suffices to consider such  $\sigma$ . Write shortly  $\tau$  for  $\tau_h$  and similarly  $X$  for  $X_h$ . By definition of the auxiliary game, writing  $T_e$  for the set of extreme NR-strategies,  $\tau$  is a transition probability from  $\bigcup_n [B \cup T_e]^n$  to  $(T \cup T_e)$  — assuming for simplicity that in the original game player II's signal informs him of his action.

$\sigma$ ,  $\tau$  and  $p$  determine a probability distribution on the space  $\Omega_0 = [A \times (B \cup T_e)]^\infty$  — in this representation we use explicitly that both players' signals in the original game

inform them of their actions; in fact, every step consists first of a transition probability from the past to  $S \times (T \cup T_e)$ , next using  $Q$  to go from  $S \times T$  to  $A \times B$  and  $Q^*$  to go from  $S \times T_e$  to  $A \times T_e$  (with the identity on the factor  $T_e$ ).

Consider now the induced probability distribution on  $\Omega_1 = [A \times (B \cup T_e) \times T^J]^\infty = \Omega_0 \times (T^J)^\infty$ , where all factors  $T$  are conditionally independent given  $\Omega_0$ , and depend only on the same dated factor  $A \times (B \cup T_e)$ , with  $\Pr(t_j = t \mid a, b) = 1$  if  $t$  is the column where “ $b$ ” appears, and  $\Pr(t_j = t \mid a, t_e) = t_e^j(t)Q_{st}(a)/v_{e,s}(a)$  — here  $t_e^j$  is the behavioural strategy in state  $j \in J$  induced by the extreme point  $t_e$ ,  $s$  is the row containing “ $a$ ”, and  $v_{e,s}(a) = \sum_t t_e^j(t)Q_{st}(a)$  is independent of  $j$  since  $t_e$  is non-revealing. Define on  $\Omega_1$  the sequence of random variables  $f_n^{k,j}(\omega) = G^k(s_n(\omega), t_n^j(\omega))$ , where  $t_n^j(\omega)$  is the projection of  $\omega$  on the factor  $j \in J$  of the  $n^{\text{th}}$  factor  $T^J$  of  $\Omega$ , and  $s_n(\omega)$  is the row containing the projection  $a_n(\omega)$  of  $\omega$  on the  $n^{\text{th}}$  factor  $A$  of  $\Omega$ . And, as usually,  $\bar{f}_n = \frac{1}{n} \sum_{m \leq n} f_m$ .

Observe now that, on this probability space, the distribution on  $A^\infty$  is the distribution induced in the true game by the strategies  $\sigma$  and  $\tau^j$ , whatever be  $j \in J$  — i.e., it is the distribution used in the statement. And also, for every  $j \in J$ , the distribution of the sequence  $(a_n(\omega), t_n^j(\omega))$  is the distribution induced in the true game by  $\sigma$  and  $\tau^j$  — since we have essentially generated  $t_j$  by its correct conditional distribution, so that  $X_n^{kj} = \mathbb{E}(\bar{f}_n^{kj} \mid A^n)$  (by part 1).

Observe finally that the distribution on  $\Omega_1$  and the sequence  $f_n(\omega)$  in  $(\mathbb{R}^{K \times J})^\infty$  are the distributions of the sequence of pairs of signals and of vector pay-offs generated by the strategies  $\sigma$  and  $\tau$  in the auxiliary game  $\tilde{\Gamma}^*$ , defined as  $\Gamma^*$  but having a random vector pay-off in  $\mathbb{R}^{K \times J}$ , where the joint distribution of the vector pay-off and the pair of signals, given the pair of moves, is defined as follows: first use the signalling matrix  $Q^*$  of  $\Gamma^*$  to select the pair of signals in  $A \times (B \cup T_e)$ ; use then our above defined  $\Pr(t_j = t \mid A \times B \cup T_e)$ , independently for each  $j \in J$ , to select  $t_j$ 's at random, and let the vector pay-off in  $\mathbb{R}^{K \times J}$  be  $G_{s(a), t_j}^k$ , where  $s(a)$  is the row containing  $a \in A$ . Remark that  $\Gamma^*$  is obtained from  $\tilde{\Gamma}^*$  by replacing the random pay-offs by their expectations.

Define now  $\bar{\Gamma}$  from  $\tilde{\Gamma}^*$  by replacing the random pay-offs by their conditional expectations given the pair of signals in  $A \times (B \cup T_e)$ . I.e.  $\bar{\Gamma}$  is described by the same signalling matrix as  $\Gamma^*$ , and by a vector pay-off function  $\bar{G}_{a,t}^{k,j}$  defined on  $A \times \bar{T}$  [with  $\bar{T} = T \cup T_e$  — or on  $A \times (B \cup T_e)$ , recalling that  $b \in B$  determines  $t \in T$  — as well as  $a \in A$  determines  $s \in S$ ]. Note it depends only on  $t \in \bar{T}$ , and not on  $b \in B$ , because when  $b \in B$  appears, the  $t \in T$  is determined, and the corresponding vector pay-off is non-random — the constant  $G_{s,t}^k$ . Thus  $u_n^{kj} = \bar{G}_{a_n, t_n}^{k,j} = \mathbb{E}(g_n^{kj} \mid \Omega_0)$ , and therefore we still have  $X_n = \mathbb{E}(\bar{u}_n \mid A^n)$ , where  $\bar{u}_n = \frac{1}{n} \sum_{m \leq n} u_m$ .

Observe all random variables are now defined on  $\Omega_0$ , we no longer need the bigger and more artificial space  $\Omega_1$ , which was just used to establish the above formula.

We can now analyse this game with vector pay-offs  $\bar{\Gamma}$  as a game with incomplete information on one side, with state space  $\bar{K} = K \times J$ , strategy sets  $S$  and  $\bar{T}$ , the signalling matrix  $Q^*$  and, for each state  $k \in \bar{K}$ , the pay-off function  $\bar{G}^k$  defined on  $A \times \bar{T}$ . A point  $p \in \Delta(\bar{K})$  is initially selected at random, told to player I, and then the true state in  $\bar{K}$  is selected according to  $p$ .

Prop. 3.46 part 2 p. 232 yields then the result, with  $\eta_n = \exp(-n\varepsilon_n) + \delta_n$ . (Equivalently, one could apply remark 3.55 p. 235 immediately to the above game with vector pay-offs).

Observe we used above — for convenience — the fact that the strategy  $\tau$  and the sequence  $\delta_n$  of prop. 3.46.2 are the same for  $\Gamma^*$  and for  $\bar{\Gamma}$ , as follows from the proof of that

proposition. Otherwise, we should have started this proof immediately with  $\bar{\Gamma}$  instead of  $\Gamma^*$ . ■

**COMMENT 3.5.** The reason for weakening the conclusion of prop. 3.46.2 is that in the present setup we have no natural probability distribution over the pay-off  $\bar{f}_n^{kj}$ . Even if we assume complete mixed strategies are given for both players (instead of behavioural strategies: thus we have a probability distribution over player I's action and player II's  $J$ -tuple of actions after every history), the randomness of the signals prevents to obtain the joint distribution — this would require to know for every  $J$ -tuple of  $n$ -histories (or at least for every  $J$ -tuple of pairs of moves), the probability distribution over the corresponding  $J$ -tuple of pairs of signals (i.e., the “mixed strategy of nature”).

**COMMENT 3.6.** One could however now use the present result, together with the techniques of last chapter for building blocs, to construct an(other) strategy  $\tau_h$ , consisting of repeated, longer and longer blocs of the present one, such that the statement of prop. 3.46 part 2 would hold (for vectors in  $\mathbb{R}^{K \times J}$ ), in addition uniformly over all “mixed strategies of nature” as above and all mixed strategies of player II that are compatible with  $\tau_h$  (and e.g. preserve the independence between successive blocs). But this does not seem to be the right way of doing it; there should be a simpler statement, not involving such “mixed strategies of nature”, and implying the above. We are probably still missing part of the structure of the problem.

### 3.b.3. Convexification again.

**DEFINITION 3.12.** Given a function  $v$  on  $\Delta(P)$ , define the convexification  $(\text{Vex } v)(\pi)$  as follows: let  $\bar{\pi}$  be the probability measure on  $P \times K \times J$  induced by  $\pi$ ; decompose it into its marginal  $\lambda^\pi$  on  $J$  and a conditional distribution  $\nu_j^\pi(k, dp)$  on  $K \times P$  given  $J$ . For every  $\lambda \in \Delta(J)$ , let  $\pi(\lambda) \in \Delta(P)$  be the distribution under  $\lambda \otimes \nu^\pi$  of the conditional distribution (under  $\lambda \otimes \nu^\pi$ ) of  $(k, j)$  given  $p \in P$ . Then  $(\text{Vex } v)(\pi)$  is the value at  $\lambda^\pi$  of the convexification of  $v[\pi(\lambda)]$  over  $\Delta(J)$ .

**REMARK 3.7.** The definition is unambiguous, because for all  $j$  where  $\lambda^\pi(j) > 0$ ,  $\nu_j$  is uniquely defined — so the function  $v[\pi(\lambda)]$  is uniquely defined on the face of  $\Delta(J)$  spanned by  $\lambda^\pi$ , hence its convexification is uniquely defined at  $\lambda^\pi$ .

**REMARK 3.8.**  $\pi(\lambda)$  is a (weak $^*$ -) continuous function of  $\lambda$ ; in particular, if  $v$  is (lower semi-) continuous, the convexification is achieved — and then by splitting to  $\#J$  points (ex. I.3Ex.10 p. 34).

**REMARK 3.9.** The function  $\text{Cav } u$  being well defined (and continuous) by Theorem 3.11, the above applies in particular to  $\text{Vex Cav } u$ .

**REMARK 3.10.** The definition translates analytically as follows. Assume  $v$  has been extended by homogeneity of degree 1 to all bounded non-negative measures on  $P$ . For  $\alpha \in \mathbb{R}_+^J$  ( $\alpha_j = \lambda_j / \lambda_j^\pi$  — cf. remark 3.7 for the case  $\lambda_j^\pi = 0$ ), define  $\phi_\alpha: P \rightarrow P$  by  $[\phi_\alpha(p)]_{k,j} = \alpha^j p^{k,j} / (\sum_{k',j'} \alpha^{j'} p^{k',j'})$ . Observe  $\phi_\alpha$  maps convex sets to convex sets, and, if  $\alpha \gg 0$ , has  $\phi_{\alpha^{-1}}$  as inverse (with  $(\alpha^{-1})_j = (\alpha_j)^{-1}$ ). Then  $\pi(\alpha)$  is the image by  $\phi_\alpha$  of the measure having  $\sum_{k,j} \alpha^j p^{k,j}$  as density with respect to  $\pi$  [or equivalently, if  $\alpha \gg 0$ , it is the measure having  $1 / \sum_{k,j} (p^{k,j} / \alpha^j)$  as density with respect to the image  $\phi_\alpha(\pi)$  of  $\pi$ ]. And, with  $g_\pi: \mathbb{R}_+^J \rightarrow \mathbb{R}$  defined by  $g_\pi(\alpha) = v[\pi(\alpha)]$ , one has  $(\text{Vex } v)(\pi) = (\text{Vex } g_\pi)(1, 1, \dots, 1)$ . Observe  $\phi_\alpha$  is well defined only on  $\{p \in P \mid \sum_{k,j} \alpha^j p^{k,j} > 0\}$ , but that the density

$\sum_{k,j} \alpha^j p^{kj}$  vanishes on the complement, so the image  $\pi(\alpha)$  is always well defined,  $\forall \alpha \in \mathbb{R}_+^J$ . Also, in the line of remark 3.7,  $\pi(\alpha)$  does not depend on the coordinates  $\alpha_j$  of  $\alpha$  for which  $\lambda_j^\pi = 0$ .

3.b.4. *More intrinsically — Approachability.* Recall the definition of the lower topology (ex. I.3Ex.15 p. 38) on the cone  $C(P)$  (cf. prop. 3.35 p. 221) of continuous convex functions on a compact convex set  $P$  and of the “weak $^*$ ”-topology on the set of regular probability measures on a compact space  $C$  (1.10 p. 6).

**THEOREM 3.13.** Denote by  $\mathbf{T}$  (resp.  $\mathbf{T}_0$ ) the space of transition probabilities  $\mathbf{t}$  from  $J$  to  $\mathbf{H}$  with the uniform topology (resp.  $\mathbf{H}_0$  with the lower topology). Endow  $\mathbf{T}$  and  $\mathbf{T}_0$  with the weak $^*$  topology (cf. remark after prop. 3.43 p. 228). For every  $\mathbf{t} \in \mathbf{T}$ , let  $\forall p \in P$

$$\begin{aligned} z_{\mathbf{t}}(p) &= \mathbb{E}_{p \otimes \mathbf{t}} h[(p \otimes \mathbf{t})(K \times J \mid h)] \\ &= \int h([p_{k,j} \mathbf{t}_j(dh)])_{(k,j) \in K \times J} \end{aligned}$$

assuming in the latter the functions  $h \in \mathbf{H}$  are extended by homogeneity of degree one to  $\mathbb{R}_+^{K \times J}$  (cf. ex. IIIEx.5 p. 145).

- (1)  $\mathbf{T}$  is compact metric, and the space  $\mathbf{T}_0$  is compact,  $T_1$ , with countable basis, and the inclusion map in  $\mathbf{T}$  has a  $G_\delta$ -image, and is a Borel isomorphism with its image.
- (2) The topology of  $\mathbf{T}_0$  is the weakest topology such that,  $\forall j \in J$ ,  $\int F(h) \mathbf{t}_j(dh)$  is lower semi-continuous whenever  $F(h) = \phi(h(p_1), \dots, h(p_n))$  where  $\phi$  is an increasing continuous function on  $\mathbb{R}^n$ .
- (3)
  - $z$  is convex in  $p \otimes \mathbf{t}$ ;
  - $\forall \mathbf{t} \in \mathbf{T}$ ,  $z_{\mathbf{t}}$  has Lipschitz constant  $3C$  and uniform bound  $C$ ;
  - $z: \mathbf{t} \mapsto z_{\mathbf{t}}$  from  $\mathbf{T}$  or  $\mathbf{T}_0$  to  $C(P)$  with the lower topology has compact graph (in the case of  $\mathbf{T}$ , this is equivalent to the continuity of the map).
- (4)  $\forall \mathbf{t} \in \mathbf{T}$ , the strategy  $\tau$  of player II generated by  $\mathbf{t}$  [via prop. 3.46 part 2 p. 232 (theorem 3.11 p. 289)] is such that,  $\forall \sigma, \forall n, \forall p, \mathbb{E}_{\sigma, \tau}^p(\bar{g}_n) \leq z_{\mathbf{t}}(p) + \eta_n$  (e.g. with  $\eta_n = \delta_n + C \exp(-n\varepsilon_n)$ ).
- (5)  $\forall \pi, (\text{Vex } \text{Cav } u)(\pi) = \min_{\mathbf{t}} \int z_{\mathbf{t}}(p) \pi(dp)$ .
- (6) Every convex function  $\varphi(\pi)$  on  $\Delta(P)$  with  $\varphi \geq \text{Vex } \text{Cav } u$  is minorated by some  $\int z_{\mathbf{t}}(p) \pi(dp)$ .
- (7) The minmax of  $\Gamma(\pi)$  exists and equals  $(\text{Vex } \text{Cav } u)(\pi)$ .

**PROOF.** 1. Since  $\mathbf{H}_0$  has a countable base (prop. 3.43 part 2 p. 228), the weak $^*$  topology on  $\mathbf{T}_0$  also has: let  $O_n$  be a basis of  $\mathbf{H}_0$  — which can be made stable under finite unions. Then the functions  $r_0 + \sum_{i=1 \dots k} r_i \mathbb{1}_{O_{n_i}}$  — where the  $r_i$  are rational ( $r_i > 0$  for  $i > 0$ ) form a sequence of lower semi-continuous functions, such that every bounded lower semi-continuous function is the limit of an increasing subsequence. This shows that the topology is also the coarsest topology for which the functions  $\mu(O_n)$  are lower semi-continuous. Observe also that this topology is  $T_1$ : if  $\mu_2$  belongs to the closure of  $\mu_1$ , we would have that  $\mu_2(F) \geq \mu_1(F)$  for every closed set  $F$ ; but every subset  $D$  of  $\mathbf{H}_0$  which is compact in the  $\mathbf{H}$ -topology is closed (since if  $f \in \overline{D}$ , any ultrafilter that refines the trace on  $D$  of the filter of neighbourhoods of  $f$  has some limit  $\tilde{f}$  in  $D$  in the  $\mathbf{H}$ -topology — so  $f \leq \tilde{f}$ , hence  $f = \tilde{f} \in D$  since  $\tilde{f} \in \mathbf{H}_0$  is minimal). Thus  $\mu_2(K) \geq \mu_1(K)$ , hence  $\mu_2(B) \geq \mu_1(B)$  for any Borel set  $B$ , since  $\mathbf{H}_0$  is a  $G_\delta$  in the compact metric space  $\mathbf{H}$  (prop. 3.43 part 4 p. 228), and therefore  $\mu_1 = \mu_2$  since both are probability measures. To

prove compactness, e.g. use that it has a countable base, and extract for any sequence a subsequence that converges — say to  $\mu$  — in the weak\* topology on  $\mathbf{H}$ , and the image of  $\mu$  by the map  $r$  of prop. 3.43 part 4 (and use the Dudley-Skhorod theorem if desired). The fact that the image of the inclusion map is a  $G_\delta$  follows immediately from prop. 3.43 part 4, while the measurability of the map is shown as follows: there is no loss, e.g. adding a constant to all pay-offs, to assume that all functions  $f \in \mathbf{H}$  are positive. By the Stone-Weierstrass theorem, the functions  $f \mapsto f(p_1)f(p_2) \cdots f(p_n)$  span all continuous functions on  $\mathbf{H}$ . Therefore,  $\mathbf{H}$  being compact metric, the Borel structure on  $\Delta(\mathbf{H})$  is generated by the functions  $\mu \mapsto \int [f(p_1)f(p_2) \cdots f(p_n)]\mu(df)$ . So it suffices to show that such a function is lower semi-continuous on  $\Delta(\mathbf{H}_0)$  — hence that  $f \mapsto f(p_1)f(p_2) \cdots f(p_n)$  is lower semi-continuous on  $\mathbf{H}_0$ , hence,  $f$  being positive, that the evaluation maps  $f \mapsto f(p)$  are lower semi-continuous on  $\mathbf{H}_0$ , which follows immediately from the definition of the topology.

2. Clearly  $F(h)$  is a real valued lower semi-continuous function on  $\mathbf{H}_0$ , hence its integral is lower semi-continuous on  $\mathbf{T}_0$ . Conversely, we have to show that then  $\mathbf{t}_j(O)$  is lower semi-continuous for every open set  $O$ . It suffices to prove this for finite union  $\bigcup_i O_i$  of basic open sets, and  $O_i = \bigcap_j U_{ij}$  where the  $U_{ij}$  belong to some subbase of open sets. By the Lipschitz character of  $h \in \mathbf{H}_0$ , the topology of  $\mathbf{H}_0$  has as subbase the sets  $\{h \mid h(p) > \alpha\}$  for  $p \in P$ ,  $\alpha \in \mathbb{R}$ . Let thus  $U_{ij} = \{h \in \mathbf{H}_0 \mid h(p_{ij}) > \alpha_{ij}\}$ , and let  $F(h) = \max_i \min_j \min[1, (h(p_{ij}) - \alpha_{ij})^+/\varepsilon]$ :  $F$  has the required properties, and increases to  $\mathbb{1}_O$  when  $\varepsilon \downarrow 0$  (the function  $\phi$  is only weakly increasing on  $\mathbb{R}^n$ , but could be made strictly increasing while preserving all properties by subtracting e.g.  $\varepsilon[1 + \exp \sum_1^n x_i]^{-1}$ ).

3. Ex. IIIEx.4 p. 142 yields the convexity of  $z$  in  $p \otimes \mathbf{t}$  and its lower semi-continuity. Further, since  $\|h\| \leq C$  for  $h \in \mathbf{H}$  (theorem 3.11 p. 289, and prop. 3.43 p. 228), we have clearly  $\|z_{\mathbf{t}}\| \leq C$ . In particular, for each  $\mathbf{t}$ ,  $z_{\mathbf{t}}$  is a lower semi-continuous real valued convex function on  $\Delta(P)$ , hence continuous. Further, ex. IIIEx.4 p. 142 show that, to prove that  $z_{\mathbf{t}}$  has Lipschitz constant  $3C$ , it suffices to consider the case where  $\mathbf{t}$  has finite support. For those, it follows immediately from the second formula for  $z_{\mathbf{t}}$ , and from an elementary computation showing that if  $h$  on  $\Delta(P)$  has Lipschitz constant  $C$ , its extension by homogeneity to  $\mathbb{R}_+^{K \times J}$  has Lipschitz constant  $3C$ . Finally, the lower semi-continuity obtained from ex. IIIEx.5 p. 145 yields the compactness of the graph of  $z: \mathbf{T} \rightarrow C(P)$  — i.e. continuity of  $z$  since  $\mathbf{T}$  is compact Hausdorff. There remains to show that it still holds with  $\mathbf{T}_0$  and its weaker topology. Consider an ultrafilter on this graph. It converges say to  $(\mathbf{t}_\infty, z_{\mathbf{t}_\infty})$  in  $\mathbf{T} \times C(P)$  by the previous point. Let  $\underline{\mathbf{t}}$  be the image of  $\mathbf{t}_\infty$  under the map  $r$  of prop. 3.43 part 4: clearly  $\underline{\mathbf{t}} \in \mathbf{T}_0$  (cf. also remark after prop. 3.43 p. 228), and the ultrafilter converges to  $\underline{\mathbf{t}}$  in  $\mathbf{T}_0$  by point 2 since  $\int F(h)\underline{\mathbf{t}}_j(dh) \leq \int F(h)\mathbf{t}_{\infty,j}(dh) = \lim_{\mathcal{U}} \int F(h)\mathbf{t}_j(dh)$  using first the monotonicity and then the continuity properties of  $F$ . Since also  $\lim_{\mathcal{U}} z_{\mathbf{t}} = z_{\mathbf{t}_\infty}$  uniformly, it follows that  $(\underline{\mathbf{t}}, z_{\underline{\mathbf{t}}})$  will be a limit point of  $\mathcal{U}$  in  $\mathbf{T}_0 \times C(P)$  if we prove that  $z_{\mathbf{t}} \leq z_{\mathbf{t}_\infty}$ . Now, using ex. I.3Ex.14bi p. 37 with  $\Omega = \mathbf{H}$ ,  $X(\omega) = (p \otimes \mathbf{t}_\infty)(K \times J \mid h)$ ,  $\mathcal{B}$  the  $\sigma$ -field spanned by  $r(h)$ ,  $g(\omega, x) = [r(h)](x)$ , we get  $z_{\underline{\mathbf{t}}}(p) = \mathsf{E}_{p \otimes \mathbf{t}_\infty}([r(h)][(p \otimes \mathbf{t}_\infty)(K \times J \mid r(h))]) \leq \mathsf{E}_{p \otimes \mathbf{t}_\infty}([r(h)][(p \otimes \mathbf{t}_\infty)(K \times J \mid h)]) \leq z_{\mathbf{t}_\infty}(p)$ . This proves point 3.

4. This is clear after theorem 3.11 — even if we assume player I is in addition informed before his first move of the choice of  $h$  by player II.

5 and 7. We first show that  $\int z_{\mathbf{t}}(p)\pi(dp) \geq (\mathsf{Vex} \mathsf{Cav} u)(\pi) \forall \pi, \forall \mathbf{t}$ . By (4), it suffices therefore to show that, against any strategy  $\mathbf{t}$  of player II, player I has a reply yielding him  $(\mathsf{Vex} \mathsf{Cav} u)(\pi) - \varepsilon$  — i.e., to prove the other half of (7) [the first half will then follow from (4) and (5)]. For every finite Borel partition  $\alpha$  of  $P$ , define  $\Gamma_\alpha(\pi)$  to be the same

game, but where player I is initially informed only of the element of  $\alpha$  containing the true  $p$  instead of the value itself of  $p$ . Clearly it suffices to exhibit, for an appropriate  $\alpha$ , a reply in  $\Gamma_\alpha(\pi)$  yielding  $(\text{Vex Cav } u)(\pi) - \varepsilon$ . Since  $\Gamma_\alpha(\pi)$  is a game with finitely many types for both players, it suffices therefore by theorem 3.1 to show that  $F_\alpha(\pi)$  converges to  $(\text{Vex Cav } u)(\pi)$ , where  $F_\alpha(\pi) = (\text{Vex}_\Pi \text{Cav}_I u_\alpha)(\pi)$ ,  $u_\alpha$  being the  $u$ -function of  $\Gamma_\alpha$ . Further  $F_\alpha(\pi)$  (and  $u_\alpha$ , and  $\text{Cav}_I u_\alpha$ ) is clearly increasing, since a finer partition yields a game more favourable to I. Let  $\delta(\alpha)$  be the maximum diameter of the elements of  $\alpha$ . Clearly  $\delta(\alpha)$  tends to zero. For every  $\alpha$ , denote also by  $\bar{\alpha}$  the same partition, together with the specification of some point  $P_\rho$  in every element  $\rho$  of  $\alpha$  — and denote by  $\Gamma_{\bar{\alpha}}(\pi)$  the same game as above, except that, after an element of  $\bar{\alpha}$  has been selected according to  $\pi$ , the pair in  $K \times J$  is selected according to the specified point in that element of  $\bar{\alpha}$ . Clearly the difference between the pay-off functions of  $\Gamma_\alpha$  and of  $\Gamma_{\bar{\alpha}}$  is  $\leq \delta(\alpha)$ , hence  $\|u_\alpha - u_{\bar{\alpha}}\| \leq \delta(\alpha)$  — and thus also  $\|F_\alpha(\pi) - \text{Vex}_\Pi \text{Cav}_I u_{\bar{\alpha}}(\pi)\| \leq \delta(\alpha)$ . But  $\Gamma_{\bar{\alpha}}(\pi)$  can also be interpreted as the game where  $p \in P$  is selected according to  $\pi$ , and told to player I, next the pair in  $K \times J$  is selected according to the specified point in the partition element containing  $p$ . In this version, it has the same strategy sets as  $\Gamma(\pi)$ , and pay-off functions that differ by less than  $\delta(\alpha)$ , hence also  $\|u - u_{\bar{\alpha}}\| \leq \delta(\alpha)$ . Now the  $\text{Cav}_I$  operation, for  $\Gamma_{\bar{\alpha}}$ , is just the concavification over  $\pi$ , since the map from  $\pi$  to the marginal on player I's types is affine, and everything else is independent of  $\pi$ . Thus we obtain that  $\|\text{Cav } u - \text{Cav}_I u_{\bar{\alpha}}\| \leq \delta(\alpha)$ .

Let now  $f(\pi) = \text{Cav}_I u_{\bar{\alpha}}(\pi)$ . It follows from the above that there only remains to show that  $|(\text{Vex Cav } u)(\pi) - (\text{Vex}_\Pi f)(\pi)| \leq 3\delta(\alpha)$ . But since  $\|\text{Cav } u - f\| \leq \delta(\alpha)$ , we have also  $\|\text{Vex Cav } u - \text{Vex } f\| \leq \delta(\alpha)$ , so it suffices to show that  $\|\text{Vex } f - \text{Vex}_\Pi f\| \leq 2\delta(\alpha)$ .

$(\text{Vex}_\Pi f)(\pi)$  is to be computed in the finite game generated by  $\bar{\alpha}$ , this is the game having as “canonical measure”  $\pi_{\bar{\alpha}} = \sum_{\rho \in \alpha} \pi(\rho) \delta_{P_\rho}$ . Hence (remark 3.10 above),  $(\text{Vex}_\Pi f)(\pi) = (\text{Vex } f)(\pi_{\bar{\alpha}})$ . Thus we are concerned with the difference  $|(\text{Vex } f)(\pi) - (\text{Vex } f)(\pi_{\bar{\alpha}})|$ ; up to an additional error of at most  $2\delta(\alpha)$ , we can now set  $f$  equal to the concave, Lipschitz function  $\text{Cav } u$ . Hence, by remark 3.10 above, we want, letting  $g_\alpha(\lambda) = f[\pi_\alpha(\lambda)]$  for  $\lambda \in \mathbb{R}_+^J$ , that  $g_\alpha$  converges to  $g_\infty$  uniformly on compact sets — given the homogeneity of degree 1, this ensures convergence of the convexifications. Thus we need, by the uniform continuity of  $f$ , that  $\pi_\alpha(\lambda) \rightarrow \pi_\infty(\lambda)$  weak\*, uniformly over compact sets of  $\mathbb{R}_+^J$ . To make this uniformity more precise, observe that  $f$  can be uniformly approximated by  $\min_i \int h_i(p) \pi(dp)$  (prop. 3.44.3 p. 230), where the  $h_i$  are convex and Lipschitz. Thus we need that, for every convex Lipschitz function  $h$  on  $P$ ,  $\int h(p)[\pi_\alpha(\lambda)](dp) \rightarrow \int h(p)[\pi_\infty(\lambda)](dp)$  uniformly over compact sets of  $\mathbb{R}_+^J$ . I.e., (cf. remark 3.10 above),  $\int (h[\phi_\lambda(p)]) (\sum_{k,j} \lambda^j p^{kj}) \pi_\alpha(dp) \rightarrow \int (h[\phi_\lambda(p)]) (\sum_{k,j} \lambda^j p^{kj}) \pi_\infty(dp)$ . We can w.l.o.g. also assume  $h$  is extended by homogeneity of degree 1; this becomes then  $\int h[(\lambda^j p^{kj})_{j,k \in J \times K}] \pi_\alpha(dp) \rightarrow \int h[(\lambda^j p^{kj})_{j,k \in J \times K}] \pi_\infty(dp)$ . [Observe that where  $\phi_\lambda(p)$  was not defined, and the density equal to zero, the new integrand  $h[\dots]$  is also zero.] Since now the integrand is clearly jointly continuous in  $\lambda$  and  $p$ , the result follows immediately, e.g. by Ascoli's theorem.

Observe we obtain the further conclusions that  $\text{Vex Cav } u$  is weak\*-continuous in  $\pi$ , and  $(\text{Vex Cav } u)(\pi) = \text{Vex}_{\lambda=(1,1,\dots,1)} [\min_{h \in \mathbf{H}_0} \int h(\lambda \cdot p) \pi(dp)]$ . And that player I can defend  $\text{Vex Cav } u$  against any strategy  $\mathbf{t}$ .

Since (cf. remark 3.8 p. 292), the convexification is achieved by splitting to a finite set  $E$  of points, with  $\#E = \#J$ , we obtain  $(\text{Vex Cav } u)(\pi) = \sum_{e \in E} \alpha_e \int h_e(\lambda_e \cdot p) \pi(dp)$ , with  $\alpha_e \geq 0$ ,  $\sum_e \alpha_e = 1$ ,  $\sum_e \alpha_e \lambda_e = (1, 1, 1, \dots, 1)$ . Let now  $\xi_{je} = \alpha_e \lambda_{ej}$ , and use the homogeneity of  $h_e$ : we have  $(\text{Vex Cav } u)(\pi) = \int \left\{ \sum_{e \in E} h_e \left[ (\xi_{je} p_{jk})_{j,k \in J \times K} \right] \right\} \pi(dp)$ : interpreting

$\xi_{je}$  as  $\mathbf{t}_j(\{h_e\})$ , we see that our integrand is exactly equal to  $z_{\mathbf{t}}(p)$  (cf. second formula in definition).

Thus we obtain also the other direction of (5), with the additional information that one can choose a minimising  $\mathbf{t}$  with support  $\leq \#J$ . In particular,  $(\text{Vex } \text{Cav } u)$  is concave, and Lipschitz with constant  $3C$ .

6. Take a minimal such  $\varphi$  (Zorn). By lemma 3.36 part 3 p. 223,  $\varphi$  is Lipschitz. Consider now the two-person zero-sum game where player II's strategy set is  $\mathbf{T}$ , player I's is  $\Delta(P)$ , and the pay-off function equals  $\int z_{\mathbf{t}}(p)\pi(dp) - \varphi(\pi)$ :  $\mathbf{T}$  and  $\Delta(P)$  are compact and convex, and the pay-off is concave and continuous (Lipschitz) w.r.t.  $\pi$  [by the Lipschitz property of  $z_{\mathbf{t}}$  (point 3) and of  $\varphi$  (above)], and convex and l.s.c. w.r.t.  $\mathbf{t}$  (point 3 again). Further, by point 5, the maxmin is less or equal to zero ( $\varphi \leq \text{Vex } \text{Cav } u$ ). The conclusion follows then from the minmax theorem (e.g. 1.6 p. 4). ■

COMMENT 3.11. The set  $\mathbf{H}$  of “approachable functions” are those functions of I's type that II can approach in a NR way. This restriction is lifted for the “approachable pay-offs”  $z_{\mathbf{t}}$  for  $\mathbf{t} \in \mathbf{T}$ .

COMMENT 3.12. The end of the proof of point 3 shows that in point 6 one can in addition require  $\mathbf{t} \in \mathbf{T}_0$ . Further, denote by  $\mathbf{H}_e$  the set of extreme points of  $\mathbf{H}_0$  (i.e., the set of extreme points of  $\mathbf{H}$  that belong to  $\mathbf{H}_0$ ). Then, by a standard argument,  $\mathbf{H}_e$  is a  $G_\delta$  in  $\mathbf{H}$  (using prop. 3.43 part 4 p. 228, and the fact that the extreme points of  $\mathbf{H}$  form a  $G_\delta$  — since  $\{(x+y)/2 \mid x \in \mathbf{H}, y \in \mathbf{H}, d(x,y) \geq n^{-1}\}$  is closed). So (still using prop. 3.43 part 4)  $\mathbf{H}_e$  is also Borel in  $\mathbf{H}_0$ , and  $\mathbf{H}_0$  and  $\mathbf{H}$  induce the same (standard) Borel structure on  $\mathbf{H}_e$ . It follows then from point 1 that  $\mathbf{T}_e$ , the transition probabilities from  $J$  to  $\mathbf{H}_e$ , is a well defined subset of  $\mathbf{T}$  and of  $\mathbf{T}_0$ , and a  $G_\delta$  in  $\mathbf{T}$ , so a Borel subset of  $\mathbf{T}_0$ , with  $\mathbf{T}$  and  $\mathbf{T}_0$  inducing the same Borel structure on  $\mathbf{T}_e$ . We claim one can in point 6 in addition require  $\mathbf{t} \in \mathbf{T}_e$ .

Indeed, use Choquet's integral representation theorem together with 7.j p. 427 to construct a (universally measurable) transition probability  $Q_h$  from  $\mathbf{H}$  to its set of extreme points, such that every  $h \in \mathbf{H}$  is the barycentre of  $Q_h$ . Observe this implies that, for  $h \in \mathbf{H}_0$ ,  $Q_h$  is carried by  $\mathbf{H}_0$  (prop. 3.43 part 4 again) — and hence by  $\mathbf{H}_e$ . Hence, for our above  $\mathbf{t} \in \mathbf{T}_0$ ,  $\underline{\mathbf{t}}$  defined as  $\underline{\mathbf{t}}_j(B) = \int_{\mathbf{H}_0} Q_h(B) \mathbf{t}_j(dh)$  belongs to  $\mathbf{T}_e$ . There only remains to show that  $z_{\underline{\mathbf{t}}} \leq z_{\mathbf{t}}$ .

Let  $h_1$  be chosen according to  $\mathbf{t}$ , next  $h_2$  according to  $Q_{h_1}$ . Observe  $h_2$  is conditionally independent of  $K \times J$  given  $h_1$ , so  $(p \otimes \mathbf{t})(K \times J \mid h_1) = (p \otimes \mathbf{t})(K \times J \mid h_1, h_2) = X(h_1)$ . Then  $z_{\mathbf{t}}(p) = \mathbb{E}[h_1(X(h_1))] = \mathbb{E}[h_2(X(h_1))] \geq \mathbb{E}[h_2(\mathbb{E}(X(h_1) \mid h_2))]$ , by ex. I.3Ex.14bi p. 37. Since  $X(h_1) = (p \otimes \mathbf{t})(K \times J \mid h_1, h_2)$  we have  $\mathbb{E}(X(h_1) \mid h_2) = (p \otimes \mathbf{t})(K \times J \mid h_2)$ , and thus our right hand member equals  $z_{\underline{\mathbf{t}}}(p)$ .

COMMENT 3.13.  $\mathbf{H}_e$  is better (more directly) characterised as the set of extreme points of the convex set of all affine (resp. convex) functions that majorate  $u$  (resp.  $\text{Cav } u$ ).

Indeed,  $\mathbf{H}_e$  is clearly contained in the set of affine functions that majorate  $u$ ; and this set is contained in the set of convex functions that majorate  $\text{Cav } u$  by the definition in prop. 3.44.1 p. 230. Now any  $h \in \mathbf{H}_e$  is extreme in the latter set (and therefore in the former): otherwise we would have  $h = \frac{1}{2}(h_1 + h_2)$ , with  $h_1 \neq h_2$ , convex — and hence both affine since their sum is so, and hence  $h_1 \geq \underline{h}_1$ ,  $h_2 \geq \underline{h}_2$  with  $\underline{h}_i \in \mathbf{H}_0$  by prop. 3.43 parts 1 and 2 p. 228. Minimality of  $h$  yields then that both inequalities are equalities, contradicting that  $h$  is an extreme point of  $\mathbf{H}$ . Conversely, let  $h$  be an extreme point of our convex set. By prop. 3.43 parts 1 and 2 (prop. 3.44 part 4) it suffices to show that  $h$  is

minimal in the set — and otherwise, by the same property, we have some affine  $h_0$  in the set, with  $h_0 \leq h$ ,  $h_0 \neq h$ . Then  $h_1 = 2h - h_0$  also belongs to the set, hence  $h = \frac{1}{2}(h_0 + h_1)$  would not be extreme.

COROLLARY 3.14.  $\mathbf{t}$  in points 5 and 6 can be assumed to vary over  $\mathbf{T}_e$ , and in point 5 to have in addition a support of cardinality  $\leq \# J$ .

PROOF. We just proved this for point 6. For point 5, this follows from the end of the proof of point 5, observing that  $\min_{h \in \mathbf{H}_0} \int (\lambda \cdot p) \pi(dp) = \min_{h \in \mathbf{H}_e} \int h(\lambda \cdot p) \pi(dp)$  (using e.g. our above  $Q_h$  — or just the Krein-Milman theorem). ■

COROLLARY 3.15. Any co-final subset of the  $z_t$  ( $t \in \mathbf{T}$ ) — like  $\{z_t \mid t \in \mathbf{T}_0\}$  or  $\{z_t \mid t \in \mathbf{T}_e\}$  — is compact in the lower topology.

PROOF. Continuity of  $z: \mathbf{T} \rightarrow C(P)$  (point 3) and compactness of  $\mathbf{T}$  yield compactness of  $\{z_t \mid t \in \mathbf{T}\}$ . For co-final subsets, the result follows then by definition of the lower topology.  $\mathbf{T}_0$  and  $\mathbf{T}_e$  generate co-final subsets by cor. 3.14. ■

COROLLARY 3.16. In the framework of definition 3.7 and prop. 3.8 p. 288 [thus with  $J$  a partition of  $K$ , so that  $p \in P = \Delta(K)$  completely determines  $\bar{p} \in \Delta(K \times J)$ , and with the  $p(\cdot \mid \kappa)$  for  $\kappa \in K^I$  mutually singular measures on  $K\}$   $\{(h[p(\cdot \mid \kappa)])_{\kappa \in K^I} \in \mathbb{R}^{K^I} \mid h \in \mathbf{H}_0\}$  is co-final in  $W_p$  and similarly for the  $z_t$  ( $t \in \mathbf{T}_0$ ) and the sets  $Z_p$ .

PROOF. Given  $w \in W_p$ , let  $\varphi: \Delta(K) \rightarrow \mathbb{R}$  be defined by  $\varphi[p(\cdot \mid \kappa)] = w^\kappa$  for  $\kappa \in K^I$ ,  $\varphi = +\infty$  elsewhere. For  $\lambda \in \Delta(P)$  we have thus  $\int \varphi(q) \lambda(dq)$  is affine; we claim it is  $\geq u(\lambda)$ . Indeed, this is obvious if  $\lambda\{\varphi = \infty\} > 0$ ; else  $\lambda$  corresponds to  $\bar{\lambda} \in \Delta(K^I)$ , and  $\int \varphi(q) \lambda(dq) = \langle \bar{\lambda}, w \rangle \geq u(\bar{\lambda} \cdot p) = u(\lambda)$ . Hence one direction, by prop. 3.43 part 1 and part 2 p. 228.

There only remains to show that, for  $h \in \mathbf{H}_0$ ,  $w = (h[p(\cdot \mid \kappa)])_{\kappa \in K^I} \in W_p$ . This again follows from  $\langle \bar{\lambda}, w \rangle = \int h(q) \lambda(dq) \geq u(\lambda) = u(\bar{\lambda} \cdot p)$ .

The proof of the second statement is completely similar. ■

COMMENT 3.14. Given this “translation” in cor. 3.16, it is now clear that theorem 3.13 to cor. 3.15 fully contain prop. 3.8.

COMMENT 3.15. In the case of statistically perfect monitoring of player I by player II — cf. comments 3.44 and 3.45 after theorem 3.39 p. 224 and the comment after prop. 3.44 p. 230 —, those comments show that  $u$  becomes then a function on  $\Delta(K \times J)$ , mapping a probability distribution on this simplex to its barycentre, that  $(\text{Cav } u)(\mu) = \max_{\nu \leq \mu} \int_{\Delta(K \times J)} u d\nu$ , and that  $\mathbf{H}_0$  consists now of the minimal convex functions on  $\Delta(K \times J)$  that majorate  $u$ .

COMMENT 3.16. In the case of statistically perfect monitoring of player II by player I,  $u$  becomes a function on  $\Delta(\Delta(K))$  — by mapping a measure  $\mu$  on  $\Delta(K \times J)$  to its image measure on  $\Delta(K)$  by the map from  $\Delta(K \times J)$  to  $\Delta(K)$  mapping every probability measure on  $K \times J$  to its marginal on  $K$ . Indeed, non-revealing strategies of player II are then pay-off equivalent to strategies that are independent of his type  $j \in J$ , so this type becomes irrelevant in the non-revealing game.  $\mathbf{H}_0$  consists then of the minimal convex functions on  $\Delta(K)$  such that  $\int h(p) \mu(dp) \geq u(\mu) \forall \mu \in \Delta(\Delta(K))$  — so  $\text{Cav } u$  is also a function on  $\Delta(\Delta(K))$ .

COMMENT 3.17. When there is statistically perfect monitoring on both sides, the above yields that the results of remark 3.15 become then valid with  $\Delta(K)$  instead of  $\Delta(K \times J)$ .

#### 4. The limit of $v_n(p)$

In this section we prove that  $\lim_{n \rightarrow \infty} v_n(p)$  always exists, where  $v_n(p)$  is the value of the  $n$ -stage game  $\Gamma_n(p)$ . We will also give a formula for  $\lim_{n \rightarrow \infty} v_n(p)$ .

Recall  $\underline{\Gamma}(p)$  the  $\delta$ -perturbation of the game in disadvantage of I (cf. 2.c p. 276). Let  $\underline{v}_{n,\delta}(p)$  be the maxmin of the  $\delta$ -perturbed  $\underline{\Gamma}_n(p)$ .

Let  $\underline{w}_\delta(p) = \liminf_{n \rightarrow \infty} \underline{v}_{n,\delta}(p)$ . Let also  $\underline{w} = \liminf_{\delta \rightarrow 0} \underline{w}_\delta$ . Notice that  $\underline{w}_\delta(p)$  and  $\underline{w}(p)$  have the Lipschitz property since  $u(p)$  and  $v_n(p)$  are uniformly Lipschitz. Moreover:

PROPOSITION 4.1.  $\underline{w}(p)$  is concave w.r.t. I.

PROOF. By theorem 1.1 p. 183, the  $\underline{v}_n(p) \forall n$ , are all concave w.r.t. I. The proposition then follows from the fact that the minimum of two concave functions is concave. ■

LEMMA 4.2. Given any strategy  $\tau$  of II in  $\underline{\Gamma}_n(p)$  there is a strategy  $\sigma$  of I such that the probability  $P_{\sigma,\tau}$  on  $(\Omega, \mathcal{G}_n)$  satisfies: if  $\theta = \min\{m \mid u(p_m) \leq \underline{w}_\delta(p_m)\}$  then

- (1) For  $m \leq \theta$ , I uses at stage  $m$  an optimal strategy in  $D(p_m)$ .
- (2) After stage  $\theta$ , I uses  $\sigma_{p_{\theta+1},n-\theta}$ , where  $\sigma_{p,m}$  is an optimal behavioural strategy in  $\underline{\Gamma}_m(p)$  represented in such a way that conditions 1, 2, 3, 4, 5, 6 of p. 278 hold.

PROOF. The proof is a straightforward construction of the above outlined strategy which consists of computing at each stage  $m$ ,  $p_n(k) = \tilde{p}_m(k)$ , playing optimally in  $D(p_m)$  as long as  $u(p_m) > \underline{w}_\delta(p_m)$  and playing in the last  $(n - \theta)$  stages an optimal behavioural strategy in  $\underline{\Gamma}_{n-\theta}(p_{\theta+1})$  where  $\theta = \min\{m \mid u(p_m) \leq \underline{w}_\delta(p_m)\}$  — making for instance all lotteries a one-to-one mapping between  $K^I$  and  $E_m^I$ . ■

LEMMA 4.3. For any strategy  $\tau$  of II and the corresponding strategy  $\sigma$  of I described in lemma 4.2:

$$\mathbb{E}\left(\frac{1}{n-\theta} \sum_{m=\theta+1}^n G_{s_m, t_m}^k \mid \tilde{\mathcal{H}}_\theta^I\right) \geq \underline{v}_{n-\theta}(p_{\theta+1}).$$

PROOF. The claim of this lemma is that the expected average pay-off from time  $\theta$  to  $n$  is at least the value of  $\underline{v}_{n-\theta}(p_{\theta+1})$ . Intuitively this is so since I plays optimally in that game. The formal proof is the following.

Let  $\omega$  stand for a typical point in  $\tilde{\mathcal{H}}_\theta^I$  and for any  $m$  let  $F_m^k$  be the matrix of  $\underline{\Gamma}_m(p)$  if the true state of nature is  $k$  and let  $\sigma_{p,m}^k$  be an optimal strategy of I in this matrix game. Let  $(\tau^k)_{k \in K}$  be a strategy of II in  $\underline{\Gamma}_{n-\theta}(p_\theta)$ ;  $(\tau^k)_{k \in K}$  is therefore  $\mathcal{K}^{\text{II}}$ -measurable in  $k$  and may depend on the information  $(b_1, \dots, b_\theta)$ .

Now

$$\mathbb{E}\left(\frac{1}{n-\theta} \sum_{m=\theta+1}^n G_{s_m, t_m}^k \mid \mathcal{G}_\theta\right) = \sigma_{p_{\theta+1}, n-\theta}^k F_{n-\theta}^k \tau^k$$

which implies

$$\mathbb{E}\left(\frac{1}{n-\theta} \sum_{m=\theta+1}^n G_{s_m, t_m}^k \mid \tilde{\mathcal{H}}_\theta^I \vee \mathcal{K}\right) = \sigma_{p_{\theta(\omega)+1}, n-\theta(\omega)}^k F_{n-\theta(\omega)}^k \bar{\tau}^k(\omega)$$

where  $\bar{\tau}^k(\omega) = \mathbb{E}(\tau^k(b_1, \dots, b_\theta) \mid \tilde{\mathcal{H}}_\theta^I \vee \mathcal{K})$ . So:

$$\mathbb{E}\left(\frac{1}{n-\theta} \sum_{m=\theta+1}^n G_{s_m, t_m}^k \mid \tilde{\mathcal{H}}_\theta^I\right) = \sum_k p_{\theta+1}(k \mid \omega) \sigma_{p_{\theta(\omega)+1}, n-\theta(\omega)}^k F_{n-\theta(\omega)}^k \bar{\tau}^k(\omega).$$

This is the pay-off in  $\underline{\Gamma}_{n-\theta(\omega)}(p_{\theta(\omega)+1})$  resulting from the optimal strategy  $\sigma_{p_{\theta(\omega)+1}, n-\theta(\omega)}$  of I in that game and from  $(\bar{\tau}^k(\omega))_{k \in K}$ . Since  $\theta$  is a  $\mathcal{H}_n^I$ -stopping time, by prop. 2.12,  $\bar{\tau}^k(\omega)$  is  $\mathcal{H}^{\text{II}}$ -measurable in  $K$  hence also a strategy in the game. The above expectation is thus at least the value of that game, i.e.,  $\underline{v}_{n-\theta}(p_{\theta+1})$ .  $\blacksquare$

PROPOSITION 4.4.  $\underline{w} \geq \text{Vex}_{\text{II}} \max(u, \underline{w})$ .

PROOF. Define  $\Delta(p, n) = (\underline{w}_\delta(p) - \underline{v}_{n,\delta}(p))^+$ . Since  $\underline{w}_\delta$  and  $\underline{v}_{n,\delta}$  are uniformly Lipschitz,  $\Delta(p, n)$  converges to 0 uniformly in  $p$ .

Take now a game  $\underline{\Gamma}_n(p)$  and for any optimal strategy  $\tau$  of II, let I play the strategy described in lemma 4.2; then, if  $\rho = \frac{1}{n} \sum_{m=1}^n \rho_m$ , we have, using lemma 2.15 (with  $q_m = \tilde{p}_m = p_m$ ) and lemma 4.3:

$$\begin{aligned} \mathbb{E}(\mathbb{E}(\rho \mid \mathcal{H}_\theta^I)) &\geq \frac{1}{n} \mathbb{E} \left\{ \sum_{m=1}^\theta \left[ u(p_m) - \frac{C(\#S)R}{\delta} \sum_k \mathbb{E}(|p_{m+1}(k) - p_m(k)| \mid \mathcal{H}_m^I) \right] \right. \\ &\quad \left. + \frac{n-\theta}{n} \underline{v}_{n-\theta}(p_{\theta+1}) - 2\delta C \right\} \\ &= \frac{1}{n} \mathbb{E} \left\{ \sum_{m=1}^\theta u(p_m) + (n-\theta) \underline{w}_\delta(p_{\theta+1}) - (n-\theta) \Delta(p_{\theta+1}, n-\theta) - 2\delta C \right. \\ &\quad \left. - \frac{C\#SR}{\delta} \sum_{m=1}^\theta \left[ \sum_k \mathbb{E}(|p_{m+1}(k) - p_m(k)| \mid \mathcal{H}_m^I) \right] \right\} \end{aligned}$$

Since up to stage  $\theta$ , I uses strategies in  $\text{NR}^I$ ,  $p_m \in \Pi^{\text{II}}(p)$  for  $m = 0, 1, \dots, \theta$  and since, by the definition of  $\theta$ ,  $u(p_m) = \max\{u(p_m), \underline{w}_\delta(p_m)\}$ , we have for the expectation (over  $\mathcal{H}_\theta^I$ ) of the first term:

$$\frac{1}{n} \mathbb{E} \left( \sum_{m=1}^\theta u(p_m) + (n-\theta) \underline{w}_\delta(p_\tau) \right) \geq \text{Vex}_{\text{II}} \max(u, \underline{w}_\delta)(p).$$

By lemma 2.1 p. 186:

$$\frac{1}{n} \sum_{m=1}^\theta \sum_k \mathbb{E} |p_{m+1}(k) - p_m(k)| \leq \sqrt{\frac{\#K-1}{n}}.$$

For any  $0 < N \leq n$ ,

$$\begin{aligned} \mathbb{E} \left( \frac{n-\theta}{n} \Delta(p_{\theta+1}, n-\theta) \right) &\leq \sum_{m=1}^{n-N} [\max_p \Delta(p, n-m)] P(\theta = m) \\ &\quad + \frac{N}{n} \sum_{m=n-N+1}^n [\max_p \Delta(p, n-m)] P(\theta = m) \\ &\leq \max_p \sup_{\ell \geq N} \Delta(p, \ell) + 2C \frac{N}{n}, \end{aligned}$$

therefore

$$\underline{v}_{n,\delta}(p) \geq \mathbb{E}(\rho) \geq \text{Vex}_{\text{II}} \max(u, \underline{w}_\delta)(p) - \frac{C\#SR}{\delta} \sqrt{\frac{\#K}{n}} - \max_p \sup_{\ell \geq N} \Delta(p, \ell) - 2C \frac{N}{n} - 2C\delta$$

Choosing  $N = \sqrt{n}$  and letting  $n$  go to infinity we get  $\underline{w}_\delta(p) \geq \text{Vex}_{\text{II}} \max(u, \underline{w}_\delta)(p) - 2C\delta$ , and from this, our result follows using the Lipschitz property of the  $\underline{w}_\delta$ .  $\blacksquare$

LEMMA 4.5. *The set of functions  $g$  satisfying  $g \geq \text{Cav}_I \text{Vex}_{\text{II}} \max(u, g)$  has a smallest element  $g_0$ , satisfying  $g_0 = \text{Cav}_I \text{Vex}_{\text{II}} \max(u, g_0)$ .*

PROOF. The proof follows from the following simply verified observations:

- $g \equiv \max_{p \in \Pi} u(p)$  is a solution, so the set is non-empty.
- The point-wise  $\inf$  of all solutions is still one.
- If, for a solution  $g$ ,  $g(p_0) = (\text{Cav}_I \text{Vex}_{\Pi} \max(u, g))(p_0) + \varepsilon$  for some  $p_0 \in \Pi$  and  $\varepsilon > 0$ , then the function

$$\tilde{g}(p) = \begin{cases} g(p) & \text{if } p \neq p_0 \\ g(p) - \varepsilon & \text{if } p = p_0 \end{cases}$$

is also a solution — and strictly smaller.  $\blacksquare$

LEMMA 4.6.

$$\sum_{m=1}^n \sum_{k \in K} \mathbb{E} |p_m(k) - \tilde{p}_{m-1}(k)| \leq \sqrt{n\delta} \psi(p) \quad \text{with } \psi(p) = \sum_{k \in K} \sqrt{p^k(1-p^k)}.$$

PROOF. Consider the measure space  $\Omega \times \mathbb{N}$ , with measure  $\lambda = P \times \mu$ , where  $\mu$  is the counting measure on  $\mathbb{N}$ . Consider some fixed  $k$ .

Let  $X(\omega, m) = |p_{m+1}(k) - \tilde{p}_m(k)|$ ,  $m = 1, \dots, n$ ,  $X(\omega, m) = 0$  otherwise. Let  $Y(\omega, m) = \mathbb{1}_{s_m(\omega) \in \underline{S} \setminus S}$ ,  $m = 1, \dots, n$ ,  $Y(\omega, m) = 0$  otherwise. Then  $X = XY$ , and the left hand member is equal to  $\int X d\lambda = \int XY d\lambda \leq \|X\|_2 \|Y\|_2$  by the Cauchy-Schwartz inequality. Since  $\|X\|_2^2 =$

$$\begin{aligned} \mathbb{E} \sum_{m=1}^n (p_{m+1}(k) - \tilde{p}_m(k))^2 &\leq \mathbb{E} \left( \sum_{m=1}^n (p_{m+1}(k) - \tilde{p}_m(k))^2 + \sum_{m=1}^n (\tilde{p}_m(k) - p_m(k))^2 \right) \\ &= \mathbb{E} (p_{n+1}(k) - p_1(k))^2 \leq p^k(1-p^k) \end{aligned}$$

and

$$\|Y\|_2^2 = \sum_{m=1}^n P(s_m(\omega) \in \underline{S} \setminus S) = n\delta$$

the result follows.  $\blacksquare$

PROPOSITION 4.7. Let  $f(p)$  be any function on  $\Pi$  satisfying  $f \leq \text{Vex}_{\Pi} \text{Cav}_I \min(u, f)$  and define  $d(p, n) = (f(p) - \underline{v}_n(p))^+$  then

$$d(p, n) \leq C \left[ \frac{(\#S)R\psi(p)}{\sqrt{n\delta}} + 2\delta \right].$$

In particular (letting  $n \rightarrow \infty$  and then  $\delta \rightarrow 0$ ) we have  $f \leq \underline{w}$ .

PROOF. It is clearly sufficient to prove the proposition for the  $\sup$  of all such functions  $f$  and in view of lemma 4.5 we may therefore assume without loss of generality that  $f(p) = \text{Vex}_{\Pi} \text{Cav}_I \min(u, f)$ . In particular we may assume that  $f$  is convex w.r.t.  $\Pi$ .

By lemma 2.6, for each  $p_n$ , I has a transition probability from  $\mathcal{H}_{I,n}^I$  to  $E_{n+1}^I$  such that if  $\tilde{p}_n(k) = P(k \mid \tilde{H}_n^I)$  then  $f(p_n) \leq \mathbb{E}(\min(u(\tilde{p}_n), f(\tilde{p}_n))) + \varepsilon 2^{-n}$ .

For any  $\tau$  of  $\Pi$  in  $\underline{\Gamma}_n(p)$ , consider the strategy of I that after each stage  $n$  ( $n = 1, \dots, m$ ) uses the above mentioned transition probability and then an optimal strategy in  $D(\tilde{p}_n)$ . By Lemma 2.15, the pay-off at stage  $n$  satisfies

$$\begin{aligned} \rho_n &\geq u(\tilde{p}_n) - \frac{C(\#S)R}{\delta} \sum_{k \in K} \mathbb{E}(|p_{n+1}(k) - \tilde{p}_n(k)| \mid \tilde{\mathcal{H}}_n^I) - 2\delta C \\ \mathbb{E}(\rho_n \mid \mathcal{H}_n^I) &\geq f(p_n) - \frac{C(\#S)R}{\delta} \sum_{k \in K} \mathbb{E}(|p_{n+1}(k) - \tilde{p}_n(k)| \mid \mathcal{H}_n^I) - 2\delta C - \varepsilon 2^{-n} \end{aligned}$$

Since  $f$  is convex w.r.t.  $\Pi$  we have by lemma 2.16 that  $\mathbb{E}(f(p_{n+1}) \mid \tilde{\mathcal{H}}_n^I) \geq f(\tilde{p}_n)$  and also  $\mathbb{E}(f(\tilde{p}_n) \mid \mathcal{H}_n^I) \geq f(p_n) - \varepsilon 2^{-n}$ , hence  $\mathbb{E}(f(p_n)) \geq f(p) - \varepsilon$ .

By lemma 4.6, we obtain thus:

$$\underline{v}_n(p) \geq \mathbb{E} \left( \frac{1}{n} \sum_{m=1}^n \rho_m \right) \geq f(p) - C \left[ \frac{\#SR\psi(p)}{\sqrt{n\delta}} + 2\delta \right] - 2\varepsilon.$$

$\varepsilon$  being arbitrary can be set equal to zero.  $\blacksquare$

COROLLARY 4.8.  $v = \lim v_n$  exists and is equal to both  $\bar{w}$  and  $\underline{w}$ . It satisfies

$$\underset{\text{I}}{\text{Cav}} \underset{\text{II}}{\text{Vex}} \max(u, v) \leq v \leq \underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} \min(u, v)$$

PROOF. By the dual of prop. 4.4,  $\bar{w} \leq \underset{\text{I}}{\text{Cav}} \min(u, \bar{w})$ , so, by prop. 4.1,  $\bar{w} \leq \underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} \min(u, \bar{w})$ , thus, by prop. 4.7,  $\bar{w} \leq \underline{w}$ . Since clearly  $\underline{v}_{n,\delta} \leq v_n \leq \bar{v}_{n,\delta}$  for any  $\delta > 0$ , the corollary follows.  $\blacksquare$

COROLLARY 4.9.

$$\begin{aligned} v_n(p) &\geq v(p) - 3C \sqrt[3]{\frac{[(\#S)(R)\psi(p)]^2}{2n}} \\ v_n(p) &\leq v(p) + 3C \sqrt[3]{\frac{[(\#T)(R')\psi(p)]^2}{2n}}. \end{aligned}$$

REMARK 4.1. Remember that the constant  $R$  (prop. 2.7) depends only on the information structure of the game.

PROOF. By cor. 4.8, we can apply prop. 4.7 with  $f = v$ ; since  $v_n \geq \underline{v}_n$ , we get  $v - v_n \leq C \left[ 2\delta + \frac{(\#S)R\psi(p)}{\sqrt{n\delta}} \right]$  for any  $\delta > 0$ .

Letting now  $\delta = [(\#S)(R)\psi(p)]^{2/3}/(2(2n)^{1/3})$  yields the first inequality. The second is dual.  $\blacksquare$

PROPOSITION 4.10. (1) Consider the functional inequalities:

- (a)  $f \geq \underset{\text{I}}{\text{Cav}} \underset{\text{II}}{\text{Vex}} \max(u, f)$
- (b)  $f \leq \underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} \min(u, f)$ .

Then  $v$  is the smallest solution of 1a and the largest solution of 1b. In particular  $v$  is the only solution of the system (1a, 1b).

(2)  $v$  is the only solution of the system (2a, 2b):

- (a)  $g = \underset{\text{II}}{\text{Vex}} \max(u, g)$
- (b)  $g = \underset{\text{I}}{\text{Cav}} \min(u, g)$ .

REMARK 4.2. It follows that if  $\underset{\text{I}}{\text{Cav}} \underset{\text{II}}{\text{Vex}} u = \underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} u$  then this is also  $v = \lim v_n$  (as it should be, knowing that this is the value of  $\Gamma_\infty(p)$ ). In fact, since  $v$  satisfies (2a, 2b) we have:

$$v = \underset{\text{I}}{\text{Cav}} \underset{\text{II}}{\text{Vex}} \max(u, v) \geq \underset{\text{I}}{\text{Cav}} \underset{\text{II}}{\text{Vex}} u = \underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} u \geq \underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} \min(u, v) = v.$$

PROOF. By cor. 4.8,  $v$  is a solution of 1b. By prop. 4.7, any solution  $f$  of 1b satisfies  $f \leq \underline{w} = v$ , i.e.  $v$  is the largest solution of 1b, and dually it is the smallest solution of 1a. This proves 1.

To prove 2, observe that since  $v$  is convex w.r.t. II,  $v \leq \underset{\text{II}}{\text{Vex}} \max(u, v)$ , but by 1  $v$  satisfies  $v \geq \underset{\text{I}}{\text{Cav}} \underset{\text{II}}{\text{Vex}} \max(u, v) \geq \underset{\text{II}}{\text{Vex}} \max(u, v)$ , hence  $v$  is a solution of 2a (and similarly of 2b). To prove that it is the only solution, remark that any solution of (2a, 2b) is both concave w.r.t. I and convex w.r.t. II and therefore it is also a solution to (1a, 1b). Since by 1  $v$  is the only solution of (1a, 1b), the result follows.  $\blacksquare$

### 5. The functional equations: existence and uniqueness

Denote by  $\mathcal{C}(\Pi)$  the space of all continuous functions on the simplex  $\Pi$ , and by  $U$  the subset of  $\mathcal{C}(\Pi)$  consisting of those functions which are “ $u$ -functions”, i.e. values of  $D(p)$ , for some two-person zero-sum game with incomplete information  $\Gamma(p)$  with full monitoring. Denote by  $\varphi$  the mapping from  $U$  to  $\mathcal{C}(\Pi)$  defined by  $\varphi(u) = v = \lim v_n$  (whatever be the game  $\Gamma(p)$ , such that  $u$  is the value of  $D(p)$ , using prop. 4.10). For the rest of this section,  $\mathcal{C}(\Pi)$  is assumed to be endowed with the topology of uniform convergence.

Recall that a vector lattice is an ordered vector space  $V$  such that the maximum and the minimum of two elements of  $V$  exists (in  $V$ ).

**PROPOSITION 5.1.** (1)  $U$  is a vector lattice containing the affine functions.  
(2)  $U$  is dense in  $\mathcal{C}(\Pi)$ .

**PROOF.** We proceed by the following assertions:

- $\alpha$ )  $U$  contains the affine functions — obvious.
- $\beta$ )  $u \in U \Rightarrow -u \in U$ .

If  $u$  arises from the game with matrices  $(G^k)_{k \in K}$  then  $-u$  arises from the game with matrices  $[-(G^k)']_{k \in K}$ , where the prime denotes transposition.

- $\gamma$ )  $u \in U, \lambda \geq 0 \Rightarrow \lambda u \in U$ .

If  $u$  arises from  $(G^k)_{k \in K}$  then  $\lambda u$  arises from  $(\lambda G^k)_{k \in K}$ .

- $\delta$ )  $u_1 \in U$  and  $u_2 \in U$  imply  $u_1 + u_2 \in U$ .

Let  $u_i$  arise from  $(G^{k,i})_{k \in K}$  with pure strategy sets  $S_i$  for I and  $T_i$  for II;  $i = 1, 2$ . Then  $u_1 + u_2$  arises from  $(D^k)_{k \in K}$ , where  $D^k$  with index set  $S_1 \times S_2$  for rows and  $T_1 \times T_2$  for columns, is defined by:

$$D^k_{(s_1, s_2), (t_1, t_2)} = G^{k,1}_{s_1, t_1} + G^{k,2}_{s_2, t_2}, \quad (s_1, s_2) \in S_1 \times S_2; \quad (t_1, t_2) \in T_1 \times T_2.$$

Assertions  $\alpha$  to  $\delta$  prove that  $U$  is a vector space containing the affine functions. The lattice property follows from:

- $\epsilon$ )  $u \in U \Rightarrow u^+ \in U$ .

If  $u$  arises from  $(G^k)_{k \in K}$  then  $u^+$  arises from the same matrices with an additional row of zeroes to each matrix.

This completes the proof of 1. Part 2 follows from 1 by the Stone-Weierstrass theorem, since the affine functions are clearly separating. ■

**PROPOSITION 5.2.** The map  $\varphi: U \rightarrow \mathcal{C}(\Pi)$  has a unique continuous extension  $\varphi: \mathcal{C}(\Pi) \rightarrow \mathcal{C}(\Pi)$ : this extension is monotone and Lipschitz with constant 1, (or non-expansive, i.e.:  $\|\varphi(f) - \varphi(g)\| \leq \|f - g\|$ ).

**PROOF.** The mapping  $\varphi$  is monotone and non-expansive. Indeed, monotonicity follows from prop. 4.10.1a. It also follows from the same proposition that for any constant  $\varepsilon$ ,  $\varphi(u + \varepsilon) = \varphi(u) + \varepsilon$  (since clearly  $\varphi(u) + \varepsilon$  is a solution for 1a and 1b, and hence it is the only solution). This together with monotonicity implies that  $\varphi$  is non expansive. Since by prop. 5.1.2  $U$  is dense in  $\mathcal{C}(\Pi)$ , it follows that there is a unique continuous extension  $\varphi: \mathcal{C}(\Pi) \rightarrow \mathcal{C}(\Pi)$ , which is monotone and non-expansive. ■

**THEOREM 5.3.** Consider the functional inequalities  $u$ ,  $f$  and  $g$  denoting arbitrary functions on the simplex:

- $$\begin{aligned} (\alpha) \quad & f \geq \underset{\text{I}}{\text{Cav}} \underset{\text{II}}{\text{Vex}} \max(u, f) \\ (\beta) \quad & f \leq \underset{\text{II}}{\text{Vex}} \underset{\text{I}}{\text{Cav}} \min(u, f) \\ (\alpha') \quad & g = \underset{\text{II}}{\text{Vex}} \max(u, g) \\ (\beta') \quad & g = \underset{\text{I}}{\text{Cav}} \min(u, g) \end{aligned}$$

There exists a monotone non-expansive mapping  $\varphi: \mathcal{C}(\Pi) \rightarrow \mathcal{C}(\Pi)$ , such that, for any  $u \in \mathcal{C}(\Pi)$ ,  $\varphi(u)$  is the smallest  $f$  satisfying  $(\alpha)$  and the largest  $f$  satisfying  $(\beta)$  — and thus in particular the only solution  $f$  of the system  $(\alpha), (\beta)$ .

$\varphi(u)$  is also the only solution  $g$  of the system  $(\alpha'), (\beta')$ .

**PROOF.** The operators  $\max$ ,  $\min$ ,  $\text{Cav}_I$  and  $\text{Vex}_{II}$  being monotone and non-expansive, propositions 4.10 and 5.2 imply immediately that  $\varphi(u)$  satisfies  $(\alpha), (\beta), (\alpha')$ , and  $(\beta')$ .

If we prove that  $\varphi(u)$  is the smallest solution of  $(\alpha)$  (and the largest of  $(\beta)$ ), the proof will be completed in the same way it was done in prop. 4.10.

In fact let  $f$  be any solution of  $(\alpha)$  and let  $\{u_n\}_{n=1}^\infty$  be an increasing sequence in  $U$  converging uniformly to  $u$  (such a sequence exists by prop. 5.1). By monotonicity of the operators involved, since  $f$  is a solution of  $(\alpha)$  it is a fortiori a solution of  $f \geq \text{Cav}_I \text{Vex}_{II} \max(u_n, f)$  and so by prop. 4.10.1a,  $f \geq \varphi(u_n)$  for  $n = 1, 2, \dots$ . But by continuity  $\varphi(u_n) \rightarrow_{n \rightarrow \infty} \varphi(u)$  and so  $f \geq \varphi(u)$  i.e.  $\varphi(u)$  is indeed the smallest solution of  $(\alpha)$  (and similarly the largest solution of  $(\beta)$ ), which completes the proof of the theorem. ■

**COMMENT 5.1.** Theorem 5.3 is of a purely functional theoretic nature, and involves no game theory at all. So it should be provable independently of our game theoretical context (cf. ex. VIEx.9 p. 319).

**THEOREM 5.4** (An approximation procedure for  $\varphi(u)$ ). Define  $\underline{v}_0 = -\infty$ ,  $\bar{v}_0 = +\infty$ ,  $\underline{v}_{n+1} = \text{Cav}_I \text{Vex}_{II} \max(u, \underline{v}_n)$ ,  $\bar{v}_{n+1} = \text{Vex}_{II} \text{Cav}_I \min(u, \bar{v}_n)$ ;  $n = 0, 1, 2, \dots$ . Then:  $\{\underline{v}_n\}_{n=1}^\infty$  is monotonically increasing,  $\{\bar{v}_n\}_{n=1}^\infty$  is monotonically decreasing and both sequences converge uniformly to  $\varphi(u)$ .

**REMARK 5.2.**  $\underline{v}_1$  (resp.  $\bar{v}_1$ ) is the  $\supinf$  (resp.  $\inf\sup$ ) of  $\Gamma_\infty$  if  $u(p)$  is the value of  $D(p)$ .

**PROOF.** Since  $u$  is continuous on a compact set, it is uniformly continuous and it is easily checked that the operators  $\max$ ,  $\min$ ,  $\text{Cav}_I$  and  $\text{Vex}_{II}$  preserve the modulus of uniform continuity. It follows that both sequences  $\{\underline{v}_n\}_{n=1}^\infty$  and  $\{\bar{v}_n\}_{n=1}^\infty$  are equicontinuous and obviously bounded.

Let us prove inductively that  $\underline{v}_n \leq \underline{v}_{n+1}$ . It is clearly true for  $n = 0$ . If  $\underline{v}_{n-1} \leq \underline{v}_n$ , then  $\underline{v}_n = \text{Cav}_I \text{Vex}_{II} \max(u, \underline{v}_{n-1}) \leq \text{Cav}_I \text{Vex}_{II} \max(u, \underline{v}_n) = \underline{v}_{n+1}$ .

Let us prove inductively that  $\underline{v}_n \leq \varphi(u)$ . It is clearly true for  $n = 0$ . If  $\underline{v}_{n-1} \leq \varphi(u)$ , then:  $\underline{v}_n = \text{Cav}_I \text{Vex}_{II} \max(u, \underline{v}_{n-1}) \leq \text{Cav}_I \text{Vex}_{II} \max(u, \varphi(u)) \leq \varphi(u)$  (by theorem 5.3).

Let  $\underline{v} = \lim_{n \rightarrow \infty} \underline{v}_n$ . This limit is uniform by the equicontinuity of the sequence. Then  $\underline{v} = \text{Cav}_I \text{Vex}_{II} \max(u, \underline{v})$  and  $\underline{v} \leq \varphi(u)$ . But theorem 5.3 implies  $\underline{v} \geq \varphi(u)$ , hence  $\underline{v} = \varphi(u)$ , i.e.  $\{\underline{v}_n\}_{n=1}^\infty$  converges uniformly to  $\varphi(u)$ . The same arguments apply to the sequence  $\bar{v}_n$ . ■

COROLLARY 5.5. If  $u$  is a continuous function on the simplex  $\Pi$ , both  $\text{Cav}_I \text{Vex}_{II} u$  and  $\text{Vex}_{II} \text{Cav}_I u$  are concave w.r.t. I and convex w.r.t. II.

PROOF. Apply cor. 3.10 and prop. 5.1. ■

## 6. On the speed of convergence of $v_n$

Corollaries 4.8 and 4.9 yield now immediately:

THEOREM 6.1. For any game as described in §1, the sequence  $v_n$  of values of the  $n$  stage games converges and we have — the function  $u$  being as defined in the beginning of §2 and the operator  $\varphi$  according to Theorem 5.3:

$$\lim_{n \rightarrow \infty} v_n(p) = \varphi(u)(p),$$

or, more precisely

$$\begin{aligned} v_n - \varphi(u) &\leq 3C \sqrt[3]{\frac{[(\#T)R\psi(p)]^2}{2n}} \text{ and} \\ \varphi(u) - v_n &\leq 3C \sqrt[3]{\frac{[(\#S)R\psi(p)]^2}{2n}}. \end{aligned}$$

REMARK 6.1. Corollary 5.2 p. 248 yields an example satisfying our assumptions and where  $|v - v_n|$  is of the order  $1/\sqrt[3]{n}$ , hence the bound for the speed of convergence given in theorem 6.1 is the best possible.

A special case for which a smaller error term is valid is the case of “full monitoring”. This is the case in which the information revealed to both players at each stage is just the pair  $(s, t)$  of pure moves chosen — called also standard signalling.

THEOREM 6.2. For games with full monitoring,  $|v_n - v|$  is at most of the order  $1/\sqrt{n}$ . More precisely:

$$|v_n(p) - v(p)| \leq \frac{C \sum_k \sqrt{p^k(1-p^k)}}{\sqrt{n}}$$

and there are games in which  $v_n(p) - v(p) = O(\frac{1}{\sqrt{n}})$ .

PROOF. This result is obtained in exactly the same manner as that in theorem 6.1 with the simplification that we do not have to do the modification of the games  $\Gamma_n$  to  $\underline{\Gamma}_n$  and can work directly with  $\Gamma_n(p)$  and  $v_n(p)$ .

The main changes in the proofs are:

- In lemma 2.13, the right hand side reduces to

$$\tilde{p}_n(k) \sum_{t \in T} \left| \bar{\tau}^k(\omega_n)(t) - \sum_{k \in K} \tilde{p}_n(k) \bar{\tau}^k(\omega_n)(t) \right|$$

Lemma 2.14 becomes superfluous.

- Consequently, in lemma 2.15,  $\#SR/\delta$  is replaced by 1, and the term  $2\delta C$  is replaced by zero.
- From this point a proposition similar to prop. 4.4 can be proved with the same strategy described in the proof there, but this time in view of the changes in lemmas 2.13, 2.14, and 2.15, we get:

$$v_n = E \left( \frac{1}{n} \sum_{m=1}^n \rho_m \right) \geq f(p) - \frac{C \sum_k \sqrt{p^k(1-p^k)}}{\sqrt{n}}.$$

As to the second part of the theorem, the game (with incomplete information on one side described in ch. V), belongs to the family of games under consideration and has  $|v_n - v| = O(\frac{1}{\sqrt{n}})$ .  $\blacksquare$

### 7. Examples

The examples which we shall treat here belong to a special subclass of the general class of games we treated in this chapter. This subclass is special in two aspects:

- We assume *full monitoring*.
- The set of states of nature  $K$  can be arranged in a matrix such that the elements of  $K^I$  are the rows and those of  $K^{II}$  are the columns, thus  $K = K^I \times K^{II}$ . Moreover, the probability distribution  $p$  on  $K$  is such that  $p(i, j) = q_I^i q_{II}^j$  for  $i \in K^I$ ,  $j \in K^{II}$ , where  $q_I = (q_I^i)_{i \in K^I}$  and  $q_{II} = (q_{II}^j)_{j \in K^{II}}$  are two probability vectors on  $K^I$  and  $K^{II}$  respectively. We will call such a probability  $p$  a product probability and denote it as  $p = q_I \times q_{II}$ . We denote by  $P$  the set of product probabilities.  $P$  is therefore a subset of the simplex  $\Pi$  of all probability distributions on  $K = K^I \times K^{II}$ .

The significance of the subclass of games satisfying the second assumption is the following: Think of the elements of  $K^I$  (resp.  $K^{II}$ ) as the possible types of I (resp. II). Thus a state of nature  $(i, j) \in K$  consists of a pair of types  $i \in K^I$  and  $j \in K^{II}$  and after the actual choice of types (i.e. the state of nature) each player knows his own type. The meaning of  $p \in P$  is then that the two types are chosen independently;  $i \in K^I$  is chosen according to  $q_I$  and  $j \in K^{II}$  is chosen independently according to  $q_{II}$ . Or equivalently: I's conditional probability on the types of II is independent of his own type (and vice versa for II). Due to this interpretation, the case  $K = K^I \times K^{II}$  and  $p \in P$  is often referred to as the **independent case**. It is easily seen that if  $p = q_I \times q_{II} \in P$ , then

$$\begin{aligned}\Pi^I(p) &= \{ q_I \times q'_{II} \mid q'_{II} \text{ is any probability vector on } K^{II} \} \subseteq P \\ \Pi^{II}(p) &= \{ q'_I \times q_{II} \mid q'_I \text{ is any probability vector on } K^I \} \subseteq P.\end{aligned}$$

It follows that the operators  $\mathbf{Cav}_I$  and  $\mathbf{Vex}_{II}$  can be carried within the subset  $P$  of product probabilities without having to evaluate the function under consideration (such as  $u(p)$ ) outside  $P$ . (Note however that  $P$  is not a convex set). If we write  $P = Q_I \times Q_{II}$  where  $Q_I$  and  $Q_{II}$  are the simplices of probability vectors on  $K^I$  and  $K^{II}$  respectively, then our concepts for this case become:

- $u$ , the value of the game  $\sum_{(i,j) \in K^I \times K^{II}} q_I(i) q_{II}(j) G^{i,j}$ , is a function  $u(q_I, q_{II})$  on  $Q_I \times Q_{II}$ ;
- $\mathbf{Cav}_I f$  is the concavification of  $f$  w.r.t. the variable  $q_I$  keeping  $q_{II}$  constant, and  $\mathbf{Vex}_{II} f$  is the convexification of  $f$  w.r.t. the variable  $q_{II}$ ; hence also  $\mathbf{Cav}_{q_I}$  and  $\mathbf{Vex}_{q_{II}}$ .
- $v(q_I, q_{II}) = \lim_{n \rightarrow \infty} v_n(q_I, q_{II})$ .

The equations determining  $v$  (i.e. equations (2a and 2b) of prop. 4.10) become now:

$$(1) \quad v = \mathbf{Vex}_{q_{II}} \max(u, v)$$

$$(2) \quad v = \mathbf{Cav}_{q_I} \min(u, v).$$

In the first set of examples, there are two types of each player, i.e.  $\#K^I = \#K^{II} = 2$ . We denote  $q_I = (x, x')$ ;  $q_{II} = (y, y')$ , the probability distribution on  $\mathcal{H}^I$  and  $\mathcal{H}^{II}$  respectively (as usual  $x' = 1 - x$ ,  $y' = 1 - y$ ). All functions involved in the solution such as  $u$ ,  $v$ ,  $\mathbf{Cav} u$ ,  $\mathbf{Vex} u$ , etc. will be described as functions of  $(x, y)$  defined on the unit square  $[0, 1] \times [0, 1]$ .

Even in this very special case we do not have in general an explicit solution of equations (1) and (2). However, it turns out that the most useful result for solving these equations

is the following observation (resulting from the fact that in dimension 1,  $\text{Cav } f(p) > f(p)$  implies that  $f$  is linear at  $p$ ):

- (1) At points  $(x, y)$  where  $u(x, y) > v(x, y)$ , the operation  $\text{Vex}_q$  is non-trivial at that point, hence  $v$  is linear in the  $y$  direction.
- (2) At points  $(x, y)$  where  $u(x, y) < v(x, y)$ ,  $v$  is linear in the  $x$  direction.

In view of the continuity of  $u$  and  $v$  it suffices therefore to find the locus of points  $(x, y)$  at which  $v(x, y) = u(x, y)$ .

Before starting with our examples let us first prove some general properties of the limit function  $v$  which are very useful to compute it. Recall that  $v(x, y)$  is Lipschitz both in  $x$  and in  $y$ . Together with the property that  $v$  is convex in  $x$  and concave in  $y$ , this will imply the existence of directional derivatives for  $v$  at any point in all directions. In fact let  $f: D \rightarrow \mathbb{R}$  be a real valued function defined on some closed polyhedral subset  $D$  of  $R_x \times R_y$ , where  $R_x$  and  $R_y$  are two finite dimensional vector spaces. Assume that  $D$  has a convex section at any  $x \in R_x$  and at any  $y \in R_y$  (i.e.  $(\{x\} \times R_y) \cap D$  and  $(R_x \times \{y\}) \cap D$  are convex sets for each  $(x, y) \in R_x \times R_y$ ). Assume further that:

- $f$  is Lipschitz.
- $f$  is concave on  $(R_x, y) \cap D$  for each  $y \in R_y$  and convex on  $(x, R_y) \cap D$  for each  $x \in R_x$ .

Denote by  $F_x(x_0, y_0)$  the tangent cone to  $f(x, y_0)$  at  $x_0$ , i.e.:

$$F_x(x_0, y_0)(h) = \lim_{\tau \rightarrow 0^+} \frac{f(x_0 + \tau h, y_0) - f(x_0, y_0)}{\tau}, \quad \text{for all } h \in R_x$$

for which  $f(x_0 + \tau h, y_0)$  is defined for sufficiently small  $\tau \geq 0$ . ( $F_y(x_0, y_0)$  is defined similarly.)

**PROPOSITION 7.1.** *Let  $f$  be any function satisfying 1 and 2, then for any  $(x_0, y_0) \in D$  and  $(a, b) \in R_x \times R_y$*

$$(3) \quad |[f(x_0 + a, y_0 + b) - f(x_0, y_0)] - [F_x(x_0, y_0)(a) + F_y(x_0, y_0)(b)]| = O(\|a, b\|)$$

whenever all terms are defined and finite.

**PROOF.** Choose  $\|a, b\|$  sufficiently small so that in addition  $f(x_0 + a, y_0)$  and  $f(x_0, y_0 + b)$  are defined. Let  $\mathcal{S}(\xi, \eta) = f(x_0 + \xi a, y_0 + \eta b) - f(x_0, y_0)$ , then  $\mathcal{S}$  is defined on  $([0, 1], 0) \cup (0, [0, 1]) \cup \{(1, 1)\}$ , and therefore (by properties 1 and 2 of  $f$ ) on at least the whole unit square  $[0, 1] \times [0, 1]$ . Also  $\mathcal{S}$  is concave in  $\xi$  (for each  $\eta$ ) and convex in  $\eta$  (for each  $\xi$ ). We shall first prove that:

$$(4) \quad \lim_{\theta \rightarrow 0^+} \frac{\mathcal{S}(\theta, \theta)}{\theta} = \left( \frac{d\mathcal{S}(\theta, 0)}{d\theta} \right)_{\theta=0^+} + \left( \frac{d\mathcal{S}(0, \theta)}{d\theta} \right)_{\theta=0^+} = F_x(x_0, y_0)(a) + F_y(x_0, y_0)(b).$$

Let  $\alpha = \left( \frac{d\mathcal{S}(\theta, 0)}{d\theta} \right)_{\theta=0^+}$ ;  $\beta = \left( \frac{d\mathcal{S}(0, \theta)}{d\theta} \right)_{\theta=0^+}$  and let  $\{\theta_i\}_{i=1}^\infty$ ,  $\theta_i \geq 0$ ,  $\theta_i \rightarrow 0$  be a sequence for which  $\lim_{i \rightarrow \infty} \frac{\mathcal{S}(\theta_i, \theta_i)}{\theta_i} = d \in \overline{\mathbb{R}}$ . We have to show that for any such sequence  $d = \alpha + \beta$ .

By the definition of the tangents  $\alpha$  and  $\beta$  and by the concavity and convexity of  $\mathcal{S}$  we have:

$$\mathcal{S}(0, \theta) \geq \beta \cdot \theta \text{ and } \mathcal{S}(\theta, 0) \leq \alpha \cdot \theta, \quad \forall \theta \geq 0.$$

For any  $\underline{d} < d$  take  $N$  large enough s.t.  $i \geq N \Rightarrow \frac{\mathcal{S}(\theta_i, \theta_i)}{\theta_i} \geq \underline{d}$ . Considering the cut of  $\mathcal{S}$  at  $\xi = \theta_i$ , the straight line  $\ell$  through  $[(\theta_i, \theta_i), \mathcal{S}(\theta_i, \theta_i)]$  and  $[(\theta_i, 0), \theta_i \alpha]$  is below  $\mathcal{S}$  for  $\eta > \theta_i$  (cf. Figure 1), i.e.  $\mathcal{S}(\theta_i, \eta) \geq \frac{\mathcal{S}(\theta_i, \theta_i) - \theta_i \alpha}{\theta_i} \eta + \theta_i \alpha$  for  $\eta \geq \theta_i$ , which for  $i \geq N$  is

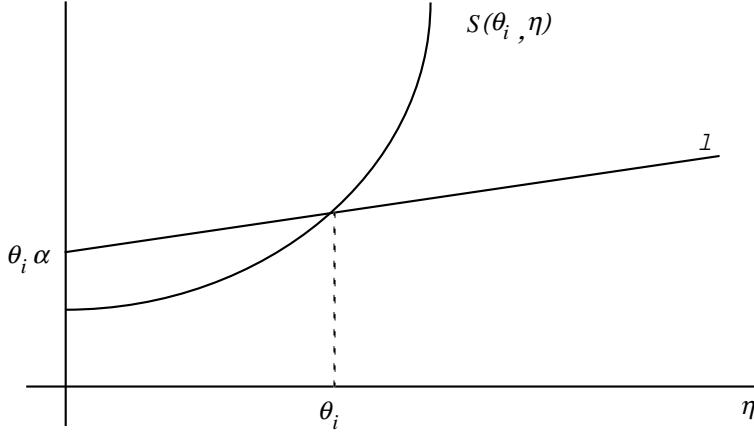


FIGURE 1. An implication of Convexity

$\mathcal{S}(\theta_i, \eta_i) \geq (\underline{d} - \alpha)\eta + \theta_i\alpha$ . Letting  $i \rightarrow \infty$  we get  $\mathcal{S}(0, \eta) \geq (\underline{d} - \alpha)\eta$ . Dividing both sides by  $\eta$  and letting  $\eta \rightarrow 0_+$  we get:  $\beta \geq \underline{d} - \alpha$ , i.e.  $\alpha + \beta \geq \underline{d}$ . Since this is true for any  $\underline{d} < d$  we have  $\alpha + \beta \geq d$ . Dually, using the concavity in  $\xi$  of  $\mathcal{S}(\xi_i, \theta_i)$  we get  $\alpha + \beta \leq d$ , which concludes the proof of Equation (4).

Now recalling the definition of  $\mathcal{S}$  we rewrite (4) as

$$(5) \quad |f(x_0 + \theta a, y_0 + \theta b) - f(x_0, y_0) - [F_x(x_0, y_0)(\theta a) + F_y(x_0, y_0)(\theta b)]| = O(\theta)$$

for each  $(a, b) \in R_x \times R_y$  and  $\theta \geq 0$  for which all terms are defined. Finally (3) follows from (5) in a standard way using the Lipschitz property of  $f$ . ■

For the case  $R_x = R_y = \mathbb{R}$ , i.e. when  $f$  is a function on the real plane we have

COROLLARY 7.2. Any function  $f$  on  $\mathbb{R}^2$  which is concave in  $x$  (for each  $y$ ) and convex in  $y$  (for each  $x$ ) and has the Lipschitz property, has at each point at most four supporting hyperplanes, one on each orthant (with the origin at the point under consideration).

The last corollary applies in particular to the limit value function  $v = \lim_{n \rightarrow \infty} v_n$  in the independent case. However for this function more differentiability properties can be proved by using the additional properties of  $v$ , namely that it is the solution of (1) and (2). For the sake of simplicity we shall state and prove these properties for the case we are presently interested in, namely when  $u$  and  $v$  are functions on the plane. Similar results might be obtained for more general cases by replacing partial derivatives by tangent cones.

PROPOSITION 7.3. At any point  $(x, y)$  where  $u(x, y) = v(x, y)$  and  $\frac{\partial u}{\partial x}$  and  $\frac{\partial u}{\partial y}$  exist:

- (1)  $v$  is differentiable, except for the following cases:
  - (a)  $\frac{\partial v}{\partial x}$  exists and equals  $\frac{\partial u}{\partial x}$ , in which case  $\frac{\partial v}{\partial y}$  may fail to exist.
  - (b)  $\frac{\partial v}{\partial y}$  exists and equals  $\frac{\partial u}{\partial y}$ , in which case  $\frac{\partial v}{\partial x}$  may fail to exist.
- (2) If in addition  $u$  is differentiable at  $(x, y)$ , then

$$(6) \quad \frac{\partial u}{\partial x} \Delta x + \frac{\partial u}{\partial y} \Delta y = \frac{\partial v}{\partial x} \Delta x + \frac{\partial v}{\partial y} \Delta y + o(\|\Delta x, \Delta y\|)$$

for all  $(\Delta x, \Delta y)$  for which  $v(x + \Delta x, y + \Delta y) = u(x + \Delta x, y + \Delta y)$ . When  $\frac{\partial v}{\partial x}$  does not exist, it can be replaced in (6) by any of the directional partial derivatives

$\frac{\partial v}{\partial x^+}$  or  $\frac{\partial v}{\partial x^-}$  ( $\frac{\partial v}{\partial y^+}$  or  $\frac{\partial v}{\partial y^-}$ ). However if either  $\frac{\partial v}{\partial x^+}$  or  $\frac{\partial v}{\partial x^-}$  equals  $\frac{\partial u}{\partial x}$ , this one has to be used in (6) (similarly for the  $y$  direction).

REMARK 7.1. Equation (6) provides a differential equation determining the “line” where  $u = v$  when such a “line” exists:

$$(7) \quad \frac{\partial u}{\partial x} dx + \frac{\partial u}{\partial y} dy = \frac{\partial v}{\partial x} dx + \frac{\partial v}{\partial y} dy.$$

REMARK 7.2. Note that 7.3.1 implies that a partial derivative in the “direction of linearity” can be always substituted for the corresponding partial derivative in (6), i.e. for instance in the  $x$  direction, where  $v$  is convex, we can always take  $\frac{\partial v}{\partial x^+}$  when  $\frac{\partial v}{\partial x^+} \leq \frac{\partial u}{\partial x}$  and  $\frac{\partial v}{\partial x^-}$  when  $\frac{\partial v}{\partial x^-} \geq \frac{\partial u}{\partial x}$ .

PROOF. 1) If either  $\frac{\partial v}{\partial x^+} \neq \frac{\partial u}{\partial x}$  or  $\frac{\partial v}{\partial x^-} \neq \frac{\partial u}{\partial x}$ , then either  $\frac{\partial v}{\partial x^+} < \frac{\partial u}{\partial x}$  or  $\frac{\partial v}{\partial x^-} > \frac{\partial u}{\partial x}$  (since  $v$  is concave in  $x$ ). Consider for instance the first case, i.e.  $\frac{\partial v}{\partial x^+} < \frac{\partial u}{\partial x}$ . Since  $v$  and  $u$  have the Lipschitz property, there exists  $\varepsilon > 0$  such that

$$0 < \xi < \varepsilon \text{ and } \eta \leq \varepsilon \xi \text{ imply } v(x + \xi, y + \eta) < u(x + \xi, y + \eta).$$

Therefore  $v$  is linear in  $y$  in this region (cf. property 1 above), so:

$$v(x + \xi, y + \varepsilon \xi) + v(x + \xi, y - \varepsilon \xi) - 2v(x + \xi, y) = 0$$

and applying prop. 7.1 p. 306 for the left-hand side we obtain:

$$\varepsilon \xi \frac{\partial v}{\partial y^+} - \varepsilon \xi \frac{\partial v}{\partial y^-} + o(\xi) = 0,$$

which implies the existence of  $\frac{\partial v}{\partial y}$ , and hence proving 1 for the first case. The second case is treated in the same way.

2) By prop. 7.1 applied for  $v$  and the differentiability of  $u$  we have:

$$(8) \quad \frac{\partial u}{\partial x} \Delta x + \frac{\partial u}{\partial y} \Delta y = \alpha \Delta x + \beta \Delta y + o(\|\Delta x, \Delta y\|)$$

where  $\alpha = \frac{\partial v}{\partial x^+}$  if  $\Delta x \geq 0$ , similarly for  $\beta$ . When  $v$  is differentiable this is just what is claimed in 2, so assume that  $v$  is not differentiable. Then by 1 we may assume that  $\frac{\partial v}{\partial x^+} = \frac{\partial v}{\partial x^-} = \frac{\partial u}{\partial x}$ , so (8) becomes:

$$\frac{\partial u}{\partial y} \Delta y = \beta \Delta y + o(\|\Delta x, \Delta y\|).$$

Clearly the two possible choices of  $\beta$  are equivalent, provided one takes  $\beta = \frac{\partial u}{\partial y}$  if possible. This completes the proof of prop. 7.3. ■

We are now ready to present some numerical examples.

### EXAMPLE 7.3. (Aumann and Maschler, 1967)

$$K = \{(1, 1), (1, 2), (2, 1), (2, 2)\}$$

$$K^I = \{(1, 1), (1, 2)\}, \{(2, 1), (2, 2)\}$$

$$K^{II} = \{(1, 1), (2, 1)\}, \{(1, 2), (2, 2)\}.$$

When  $K$  is arranged in a matrix with  $K^I$  as the set of rows and  $K^{II}$  as the set of columns, the corresponding pay-off matrices  $G^{ij}$  ( $i = 1, 2; j = 1, 2$ ) are:

$$\begin{array}{c} y' \\ \left( \begin{array}{ccc} -1 & 0 & 2 \\ 1 & 0 & 2 \\ 1 & 0 & 2 \\ -1 & 0 & 2 \\ -1 & 2 & 0 \\ 1 & 2 & 0 \\ -1 & 2 & 0 \\ 1 & 2 & 0 \end{array} \right) \\ x \\ x' \end{array} \quad \begin{array}{c} y \\ \left( \begin{array}{ccc} 1 & 2 & 0 \\ -1 & 2 & 0 \\ 1 & 2 & 0 \\ -1 & 2 & 0 \\ 1 & 0 & 2 \\ -1 & 0 & 2 \\ -1 & 0 & 2 \\ 1 & 0 & 2 \end{array} \right) \end{array}$$

The probability distribution on the rows ( $K^I$ ) is  $q_I = (x, x')$  and the probability distribution on the columns ( $K^{II}$ ) is  $q_{II} = (y, y')$ . The non-revealing game  $D(x, y)$  is:

$$D(x, y) = xyG^{11} + xy'G^{12} + x'yG^{21} + x'y'G^{22}$$

which is:

$$\begin{pmatrix} xy - xy' + x'y - x'y' & 2xy + 2x'y' & 2xy' + 2x'y \\ -xy + xy' - x'y + x'y' & 2xy + 2x'y' & 2xy' + 2x'y \\ xy + xy' - x'y - x'y' & 2xy + 2x'y' & 2xy' + 2x'y \\ -xy - xy' + x'y + x'y' & 2xy + 2x'y' & 2xy' + 2x'y \end{pmatrix}$$

The value  $u(x, y)$  of this game is given in Figure 2.

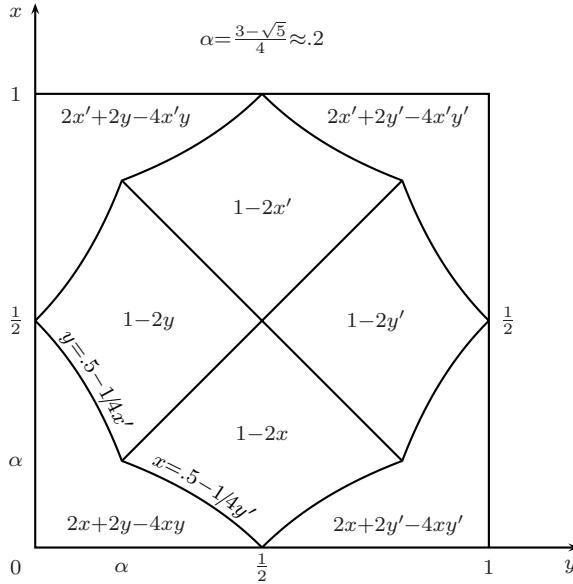


FIGURE 2.  $u(x, y)$  of example 7.3

Note that according to our notation the value of  $u$  at the corners of the square are the values of the corresponding matrices. Note also that  $u$  is symmetric about  $x = \frac{1}{2}$  and about  $y = \frac{1}{2}$ .

Convexifying in the  $y$  direction and then concavifying in the  $x$  direction (all by brute force) yields  $\underline{v} = \text{Cav}_x \text{Vex}_y u$ , in Figure 3, and similarly for  $\bar{v} = \text{Vex}_y \text{Cav}_x u$  in Figure 4.

It is readily seen that for about half the points  $(x, y)$  in the square,  $\text{Cav } \text{Vex } u \neq \text{Vex } \text{Cav } u$  and hence  $v_\infty$  does not exist.

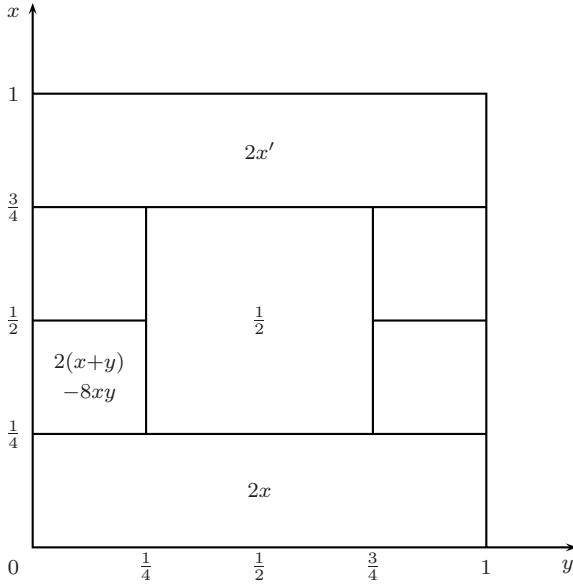


FIGURE 3.  $\underline{v} = \text{Cav}_x \text{Vex}_y u$  for example 7.3

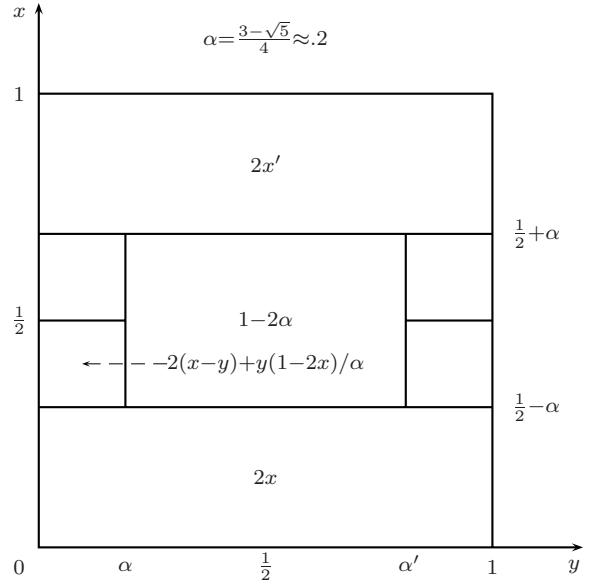


FIGURE 4.  $\bar{v} = \text{Vex}_y \text{Cav}_x u$  for example 7.3

Let us compute now the asymptotic value  $v = \lim_{n \rightarrow \infty} v_n$  by solving the equations:

$$\begin{aligned} v &= \text{Vex}_{\bar{y}} \max(u, v) \\ v &= \text{Cav}_{\bar{x}} \min(u, v). \end{aligned}$$

Recall that  $\text{Cav}_x \text{Vex}_y u \leq v \leq \text{Vex}_y \text{Cav}_x u$ , so that:

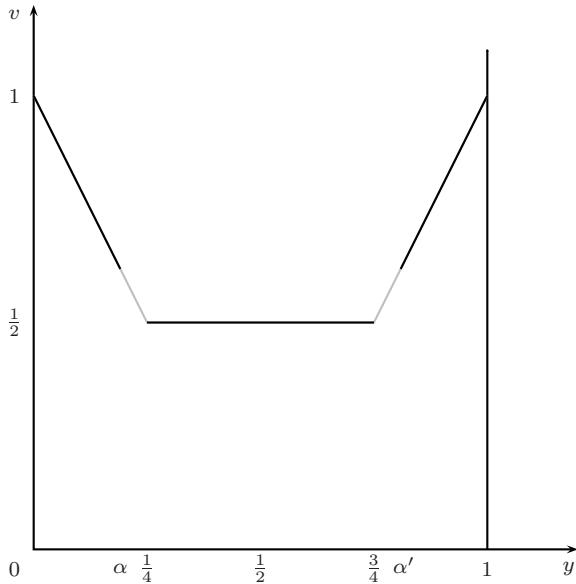
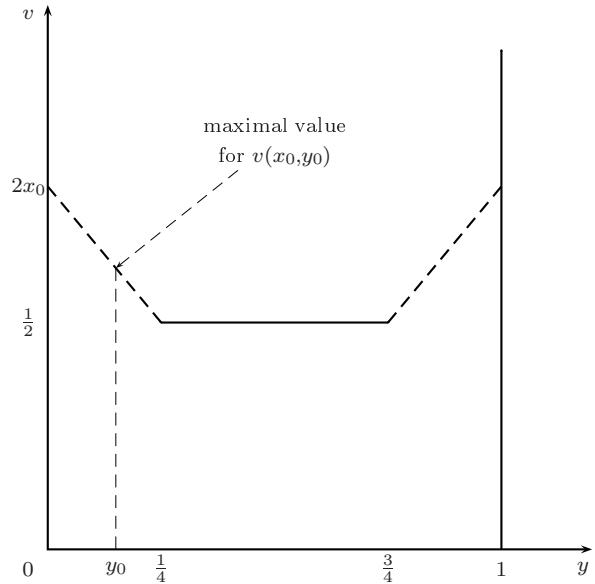
- (1) On  $\{(x, y) \mid y = 0 \text{ or } y = 1 \text{ or } x = \frac{1}{2}, \min(y, y') \leq \alpha\}$ ,  $\text{Cav} \text{Vex} u = \text{Vex} \text{Cav} u = u$ , so  $v = u$  on these segments.
- (2) On  $\{(x, y) \mid x \leq \frac{1}{4} \text{ or } x' \leq \frac{1}{4}\}$ ,  $\text{Cav} \text{Vex} u = \text{Vex} \text{Cav} u$ , hence  $v = \text{Cav} \text{Vex} u = \text{Vex} \text{Cav} u$  there.
- (3) On the set  $\{(x, y) \mid \frac{1}{4} < x < \frac{3}{4}, \frac{1}{4} < y < \frac{3}{4}\}$ ,  $u < \text{Cav} \text{Vex} u$ , so  $u < v$  on this set and hence  $v$  is linear in  $x$  there. Since at the boundaries  $x = \frac{1}{4}$  and  $x = \frac{3}{4}$ ,  $v = \frac{1}{2}$  (cf. 2), we conclude that  $v = \frac{1}{2}$  on this set.
- (4) Consider now the function  $v$  on  $x = \frac{1}{2}$ . From 2 and 3 it follows that (cf. figure 5):

$$v\left(\frac{1}{2}, y\right) = \begin{cases} 1 - 2y & 0 \leq y \leq \alpha \\ \text{not yet determined} & \alpha < y < \frac{1}{4} \\ \frac{1}{2} & \frac{1}{4} \leq y \leq \frac{1}{2} \\ \text{not yet determined} & \frac{1}{2} < y < \alpha' \\ 1 - 2y' & \alpha' \leq y \leq 1 \end{cases}$$

Since  $v\left(\frac{1}{2}, y\right)$  is convex in  $y$  we get:

$$v\left(\frac{1}{2}, y\right) = \begin{cases} 1 - 2y & 0 \leq y \leq \frac{1}{4} \\ \frac{1}{2} & \frac{1}{4} \leq y \leq \frac{1}{2} \\ 1 - 2y' & \frac{1}{2} \leq y \leq 1. \end{cases}$$

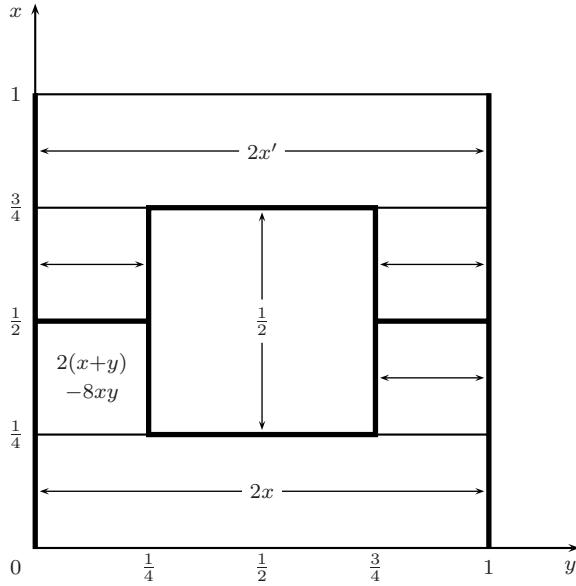
In view of the symmetry of  $v$  about  $x = \frac{1}{2}$  and  $y = \frac{1}{2}$ , it remains now to determine  $v(x, y)$  on the set  $\{(x, y) \mid \frac{1}{4} < x < \frac{1}{2}, 0 < y < \frac{1}{4}\}$ . Take a point  $(x_0, y_0)$  in this set

FIGURE 5.  $v(\frac{1}{2}, y)$  for example 7.3FIGURE 6.  $v(x_0, y)$ , with  $\frac{1}{4} <$ 

and consider first the section  $v(x_0, y)$ . Since  $y_0 = (4y_0)\cdot\frac{1}{4} + (1 - 4y_0)\cdot 0$  and  $v(x_0, y)$  is convex, it follows that (cf. Fig. 6):  $v(x_0, y_0) \leq (4y_0)v(x_0, \frac{1}{4}) + (1 - 4y_0)v(x_0, 0)$ , which is:  $v(x_0, y_0) \leq 2x_0 + 2y_0 - 8x_0y_0$ .

Similarly, using the concavity of  $v(x, y)$  we get  $v(x_0, y_0) \geq 2x_0 + 2y_0 - 8x_0y_0$ , hence in the region under consideration  $v(x, y) = 2x + 2y - 8xy$  and it is obtained as a linear interpolation in the  $y$  direction between  $v(0, y)$  and  $v(\frac{1}{4}, y)$ .

Summing up our construction this function  $v = \lim v_n$  is given by Figure 7. The thick

FIGURE 7.  $v = \lim v_n$  for example 7.3

lines in this figure are the locus of the points  $\{(x, y) \mid v(x, y) = u(x, y)\}$ . The values of  $v$  on the square are obtained by linear interpolation between thick lines in the directions indicated by the arrows.

Note that one can check  $v$ , once it is obtained, by verifying the equations (1) and (2) p. 305, since  $v$  is the unique solution of these equations.

We see that in this first example  $\lim v_n$  coincides with one of its bounds, namely with  $v = \text{Cav Vex } u$ . In our second example we have no longer such a coincidence.

EXAMPLE 7.4. Our second example has the pay-off matrices  $G^{ij}$  ( $i = 1, 2; j = 1, 2$ ):

$$\begin{array}{cc} & \begin{matrix} y' \\ y \end{matrix} \\ \begin{matrix} x \\ x' \end{matrix} & \begin{pmatrix} 1 & -1 & +1 & -1 \\ 0 & 0 & 0 & 0 \end{pmatrix} \quad \begin{pmatrix} 0 & 0 & 0 & 0 \\ -1 & 1 & +1 & -1 \end{pmatrix} \\ & \begin{pmatrix} 0 & 0 & 0 & 0 \\ 1 & -1 & -1 & +1 \end{pmatrix} \quad \begin{pmatrix} -1 & 1 & -1 & +1 \\ 0 & 0 & 0 & 0 \end{pmatrix} \end{array}$$

The non-revealing game  $D(x, y)$  is given by the matrix game:

$$\begin{pmatrix} x - y & y - x & x - y & y - x \\ y' - x & x - y' & x - y' & y' - x \end{pmatrix}$$

Its value  $u$ , the maxmin  $\text{Cav Vex } u$  and the minmax  $\text{Vex Cav } u$  are given in Figures 8, 9

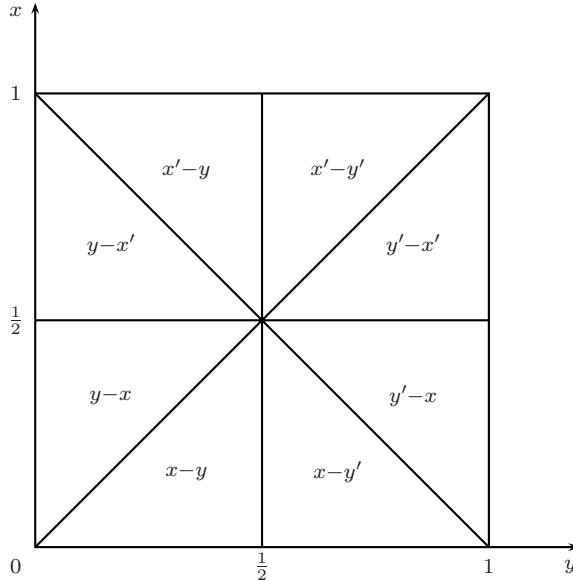
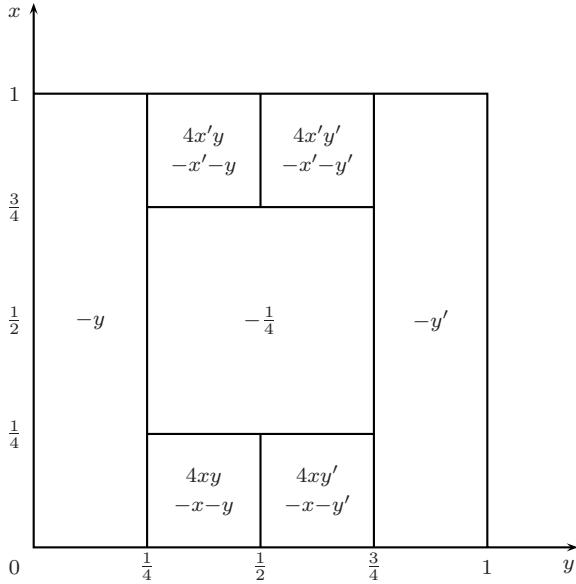
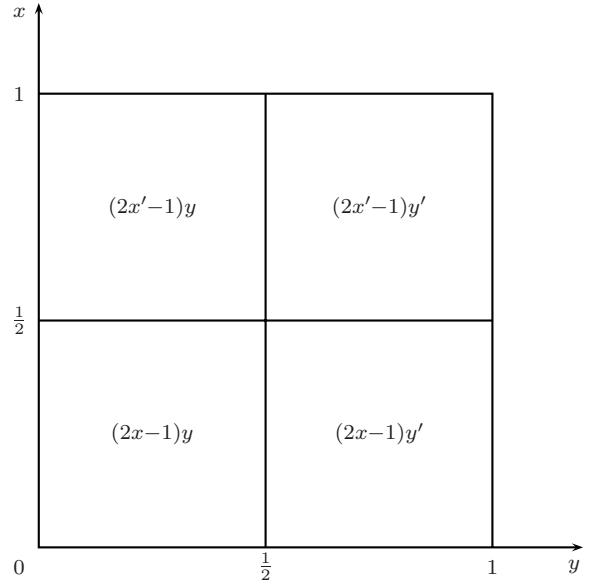


FIGURE 8.  $u(x, y)$  of example 7.4

and 10, and may be verified by the reader as an exercise.

Remark that this game, just as the previous one, has a symmetry about  $x = \frac{1}{2}$  and about  $y = \frac{1}{2}$ . So it suffices to find  $v$  on  $\{(x, y) \mid x \leq \frac{1}{2}, y \leq \frac{1}{2}\}$ . To do this we proceed through the following steps:

- (1) On the segment  $\{(x, \frac{1}{2}) \mid x \leq \frac{1}{4}\}$ ,  $\text{Cav Vex } u = \text{Vex Cav } u = u$ . So  $v = u$  there.
- (2) At  $(0, 0)$ ,  $\text{Cav Vex } u = \text{Vex Cav } u = u$  so  $v(0, 0) = u(0, 0)$ . (This is true for any game and any point-mass on  $K$ .)
- (3) On the triangle  $\{(x, y) \mid 0 < x \leq y < \frac{1}{2}\}$ ,  $u > \text{Vex Cav } u$ , so,  $u > v$  and consequently  $v$  is linear in  $y$  in this region.
- (4) Since  $\text{Cav Vex} = \text{Vex Cav}$  on the boundary of the square,  $v(\cdot, 0) = \text{Cav}_x u(\cdot, 0) = 0$ , so  $v(x, 0) > u(x, 0)$  for  $0 < x < 1$ .

FIGURE 9.  $v = \text{Cav Vex } u$  of example 7.4FIGURE 10.  $\bar{v} = \text{Vex Cav } u$  of ex

From 3 and 4, it follows that a “curve” on which  $v = u$  starts at  $(0,0)$  and lies between  $y = 0$  and the diagonal  $y = x$ .

- (5) From 4,  $v(\frac{1}{2}, \varepsilon) < u(\frac{1}{2}, \varepsilon)$  for  $\varepsilon > 0$  sufficiently small. Hence  $v(\frac{1}{2}, \varepsilon)$  is obtained by linear interpolation in the  $x$  direction, say (using symmetry)  $v(\frac{1}{2}, \varepsilon) = \frac{1}{2}v(x, \varepsilon) + \frac{1}{2}v(x', \varepsilon)$ . But  $v(x, \varepsilon) < 0$  since  $\text{Vex Cav } u < 0$  at  $(x, \varepsilon)$  for  $x \neq \frac{1}{2}$ . It follows that  $v(\frac{1}{2}, \varepsilon) < 0$  and by symmetry  $v(\frac{1}{2}, \varepsilon') < 0$ . From the convexity of  $v(\frac{1}{2}, y)$  it follows that  $v(\frac{1}{2}, \frac{1}{2}) < 0 = u(\frac{1}{2}, \frac{1}{2})$ . This implies that if we denote  $\{(x, \frac{1}{2}) \mid x \leq \xi\}$  the segment on which  $v = u$  and which contains the segment in 1, then  $\xi < \frac{1}{2}$ .

We apply now the differential equation (7) p.308 to determine the curve where  $v = u$  between  $y = 0$  and the line  $y = x$  (cf. Figure 11). By the symmetry of  $v$  about  $x = \frac{1}{2}$ ,  $\frac{\partial v}{\partial x} = 0$ ; also  $\frac{\partial v}{\partial y}$  is the slope of  $v$ , say from  $(x, y)$  to  $(x, \frac{1}{2})$  (where  $v(x, \frac{1}{2}) = x - \frac{1}{2}$ ); so, since  $u(x, y) = y - x$ , (7) yields  $dy - dx = \frac{(x - \frac{1}{2}) - (y - x)}{\frac{1}{2} - y} \cdot dy$  for  $x \leq \xi$ , i.e.:

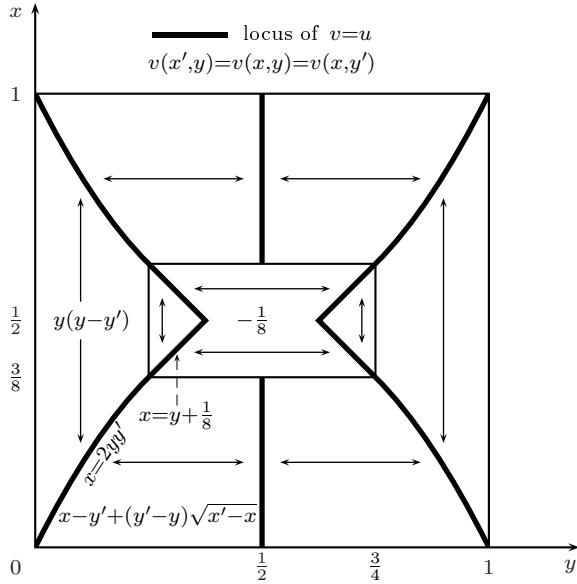
$$\frac{dx}{2 - 4x} = \frac{dy}{1 - 2y}$$

Together with the initial condition  $x(0) = 0$  this yields the curve  $x = 2yy'$ . By linear interpolation between this and the curves  $y = \frac{1}{2}$  and  $x' = 2yy'$  we have:

$$\begin{aligned} v(x, y) &= x - y' + (y' - y)\sqrt{x' - x} \quad \text{for } 0 \leq x \leq \xi, \ x \leq 2yy', \ y \leq \frac{1}{2} \\ &= y(y - y') \quad \text{for } 2yy' \leq x \leq (2yy')', \ 0 \leq y \leq \eta, \ \text{where } \eta \leq \frac{1}{2}, \ \xi = 2\eta\eta' \end{aligned}$$

To determine the point  $\xi, \eta$  we note that above  $\xi$ ,  $u$  is strictly greater than  $v$  on  $y = \frac{1}{2}$ . Hence  $v$  is linear (in  $y$  from the first line to its symmetric w.r.t.  $y = \frac{1}{2}$ ). This implies that  $v(x, \cdot)$  is constant for  $x \geq \xi$  hence  $(\frac{\partial v}{\partial y})_{\xi} = 0$ . This yields:

$$\frac{\partial}{\partial y}(x - y' + (y' - y)\sqrt{x' - x}) = 0,$$

FIGURE 11.  $v = \lim v_n$  for example 7.4

i.e.,  $1 - 2\sqrt{\xi' - \xi} = 0$  or  $\xi = \frac{3}{8}$ ,  $\eta = \frac{1}{4}$ . Beyond the point  $(\xi, \eta)$ ,  $v$  is linear between the curve  $v = u$  and its symmetries about  $x = \frac{1}{2}$  and  $y = \frac{1}{2}$ , in the indicated directions. The equation  $x = x(y)$  of this curve is again obtained by equation (7) which is:

$$dy - dx = 0.$$

Together with the initial condition  $(\xi, \eta)$  this gives  $x = \frac{1}{8} + y$ .

Finally, linear interpolation in the indicated directions gives  $v(x, y) = -\frac{1}{8}$  for  $\frac{3}{8} \leq x \leq \frac{5}{8}$ ,  $\frac{1}{4} \leq y \leq \frac{3}{4}$ . Summing up, the resulting  $v = \lim v_n$  of example 7.4 is given in Figure 11.

### Exercises

1. Prove lemma 2.9 p. 279.

HINT. Define by induction a unique probability on  $(\Omega, \mathcal{G}_n)$  and then let  $n \rightarrow \infty$  using Ionescu-Tulcea's theorem (Neveu, 1970) — cf. prop. 1.6 p. 54.

2. Prove that if  $u(p)$  is either concave w.r.t. I or convex w.r.t. II then  $\Gamma_\infty$  has a value.

3. Consider  $u$  on  $[0, 1]^2$  with  $u(\cdot, y)$  and  $u(x, \cdot)$  piecewise linear satisfying:

$$\begin{array}{lll} u(0, 0) = 0, & u\left(\frac{1}{2}, 0\right) = 1, & u(1, 0) = 0 \\ u(0, \frac{1}{2}) = 1, & u\left(\frac{1}{2}, \frac{1}{2}\right) = 0, & u(1, \frac{1}{2}) = 1 \\ u(0, 1) = 0, & u\left(\frac{1}{2}, 1\right) = 1, & u(1, 1) = 0 \end{array}$$

and prove that  $\text{Vex } \text{Cav } u\left(\frac{1}{2}, \frac{1}{2}\right) = 1$ ,  $\text{Cav } \text{Vex } u\left(\frac{1}{2}, \frac{1}{2}\right) = 0$ .

4. Prove  $\{u \in \mathcal{C}(\Pi) \mid \text{Cav}_I \text{Vex}_{II} u \neq \text{Vex}_{II} \text{Cav}_I u\}$  is open and dense in  $\mathcal{C}(\Pi)$ .

**5. Another Example.** This example again has the same structure as the two examples in sect. 7. It differs from them only by the pay-off matrices  $G^{ij}$  which are now

$$\begin{array}{cc} & y' \\ x & \begin{pmatrix} +1 & -1 \\ 0 & 0 \end{pmatrix} \quad \begin{pmatrix} 0 & 0 \\ -1 & -1 \end{pmatrix} \\ & x' \quad \begin{pmatrix} 0 & 0 \\ -1 & -1 \end{pmatrix} \quad \begin{pmatrix} -1 & +1 \\ 0 & 0 \end{pmatrix} \end{array}$$

HINT. The non-revealing game  $D(x, y)$  is the matrix game:

$$\begin{pmatrix} x-y & y-x \\ -xy-x'y' & -xy-x'y' \end{pmatrix}$$

Verify the functions  $u(x, y)$ ,  $\bar{v}_\infty$  and  $\underline{v}_\infty$  are as in Figures 12, 13 and 14.

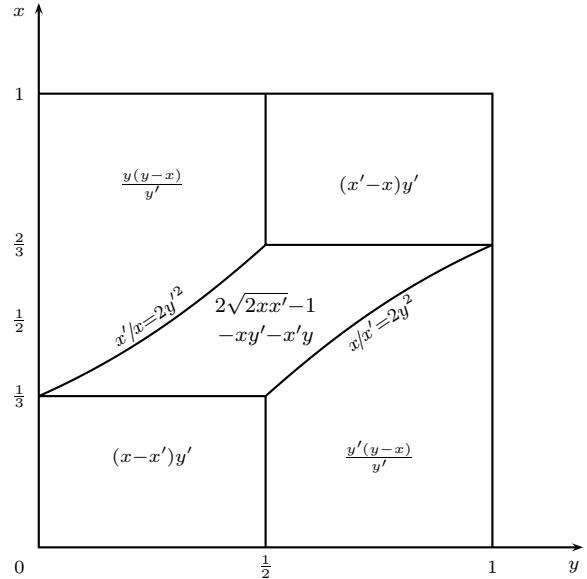
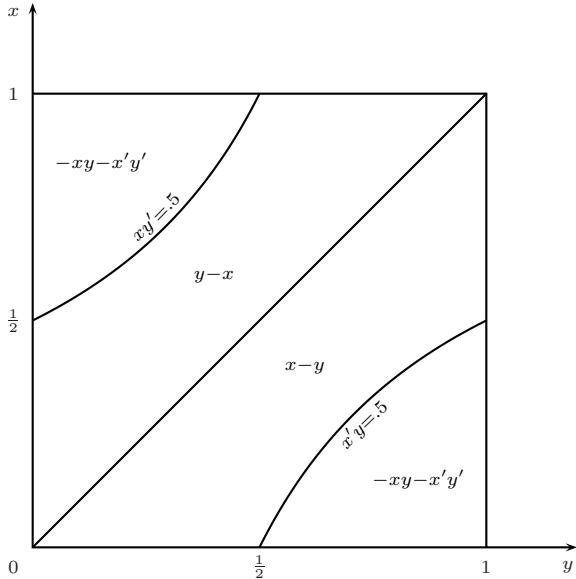
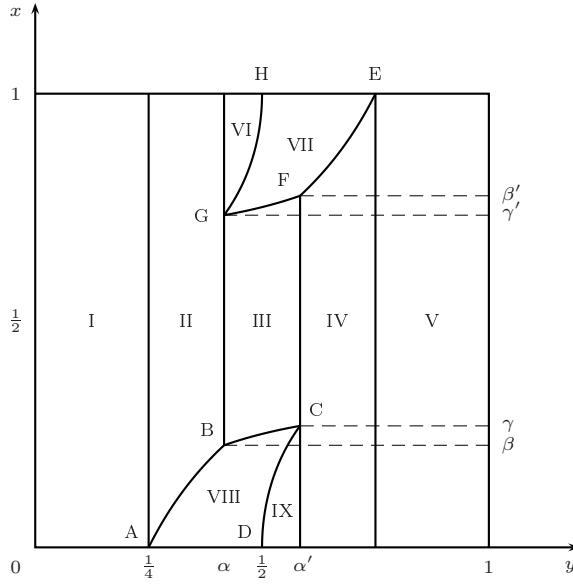


FIGURE 12.  $u(x, y)$  of Example VIEx.5  
FIGURE 13.  $\text{Vex Cav } u$  for example VIEx.5

The equations of the curves and the values of  $\text{Cav Vex } u$  in the various regions of figure 14 are given in figure 15. The values of  $\alpha$ ,  $\beta$  and  $\gamma$  are found by intersecting the corresponding lines and approximately  $\alpha = .416$ ,  $\beta = .225$ ,  $\gamma = .268$ . Note that although  $u$  is symmetric with respect to the main diagonal:  $u(x, y) = u(y, x)$ , the functions  $\text{Cav Vex } u$  and  $\text{Vex Cav } u$  do not have this symmetry because of the difference between the operations  $\text{Cav}_x$  and  $\text{Vex}_y$ . However the game and hence all the functions  $u$ ,  $\text{Cav Vex } u$ ,  $\text{Vex Cav } u$  and  $v$  have the symmetry  $f(x, y) = f(x', y')$ .

To find  $v = \lim v_n$ , proceed by the following steps:

- (1)  $u = \bar{v}_\infty = \underline{v}_\infty$ , and hence  $v = u$  on the segments  $[(0, \frac{1}{2}), (0, 1)]$  and  $[(1, 0), (1, \frac{1}{2})]$ .
- (2)  $u < \text{Cav Vex } u$  and hence  $u < v$  on the lines  $x'y = \frac{1}{2}$  and  $y'x = \frac{1}{2}$ .
- (3)  $u > \text{Vex Cav } u$  and hence  $u > v$  on  $x = y$ ,  $0 < x < 1$ .
- (4)  $u > \text{Vex Cav } u$  and hence  $u > v$  on  $\{(x, \frac{1}{2} - \varepsilon) \mid 0 \leq x \leq \frac{1}{2}\}$  for any small  $\varepsilon > 0$ , and similarly on  $\{(x, \frac{1}{2} + \varepsilon) \mid \frac{1}{2} \leq x \leq 1\}$ .
- (5)  $u = v$  for  $x = y = 0$  and for  $x = y = 1$ .
- (6) For any  $y_0 < \frac{1}{2}$ ,  $u(x, y_0)$  is piecewise linear and from 1 to 5 it follows that it has the following structure:
  - $u(0, y_0) = v(0, y_0)$  and  $u(x, y_0)$  is linearly increasing from  $x = 0$  to  $x = y_0$  and  $u(x, y_0) > v(x, y_0)$  for  $0 < x \leq y_0$ .
  - $u(x, y_0)$  decreases linearly from  $x = y_0$  to  $x = \frac{1}{2y'_0}$  where  $u\left(\frac{1}{2y'_0}, y_0\right) < v\left(\frac{1}{2y'_0}, y_0\right)$ .

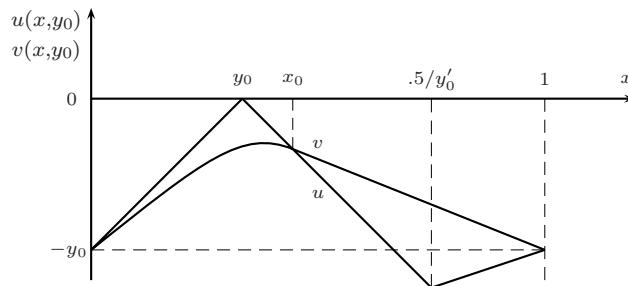
FIGURE 14.  $\underline{v}_\infty = \text{Cav Vex } u$  for example VIEx.5

Region	$\text{Cav Vex } u$	Curve	Equation
I	$-y$	AB	$2x'\sqrt{y} = 1$
II	$(4y - 4\sqrt{y} + 1)x' - y$	EF	$2x\sqrt{y'} = 1$
III	$\frac{1}{4(\sqrt{y} + \sqrt{y'})^2} - \frac{x\sqrt{y'} + x'\sqrt{y}}{\sqrt{y} + \sqrt{y'}}$	BC	$x = \frac{1}{2} - \frac{1}{4}(y + \sqrt{yy'})^{-1}$
IV	$(4y' - 4\sqrt{y'} + 1)x - y'$	FG	$x = \frac{1}{2} + \frac{1}{4}(y' + \sqrt{yy'})^{-1}$
V	$-y'$	GH	$x' = \frac{1}{2}\sqrt{1 - y/y'}$
VI	$[(4y' - 1) - 4\sqrt{y'(y' - y)}]x' - y$		
VII	$(4xx' - 1)y' - x'$	CD	$x = \frac{1}{2}\sqrt{1 - y'/y}$
VIII	$(4xx' - 1)y - x$		
IX	$[(4y - 1) - 4\sqrt{y(y - y')}]x - y'$		

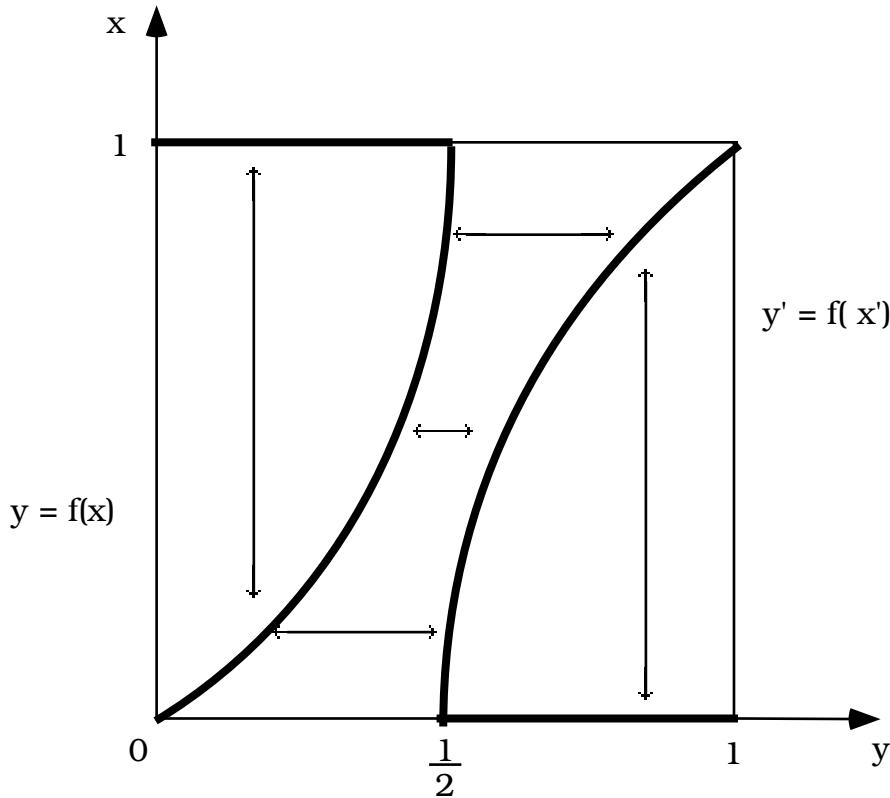
FIGURE 15. The equations of Figure 14

- $u(x, y_0)$  increases linearly from  $x = \frac{1}{2y_0}$  to  $x = 1$ ,  $u(x, y_0) \leq v(x, y_0)$  for  $\frac{1}{2y_0} \leq x \leq 1$  and equality holds only for  $x = 1$ .

Since  $v(x, y_0)$  is concave and continuous on  $0 \leq x \leq 1$ , the relation between  $u(x, y_0)$  and  $v(x, y_0)$  must be of the form given by Figure 16. Conclude that for each  $0 < y_0 < \frac{1}{2}$  there is a unique  $x_0$  for which

FIGURE 16.  $u(x, y_0)$  and  $v(x, y_0)$ ,  $0 < y_0 < \frac{1}{2}$ , for example VIEx.5

$u(x_0, y_0) = v(x_0, y_0)$  and hence there is a unique interior line of  $u = v$  from  $(0, 0)$  to  $(1, \frac{1}{2})$  and of course its image by the transformation  $x \mapsto x'$ ,  $y \mapsto y'$ . The resulting  $\lim v_n$  function is thus given in Figure 17 — without specifying the equation of the lines. For that, cf. (Mertens and Zamir, 1971).

FIGURE 17.  $v = \lim v_n$  for example VIEx.5

### 6. Bilinearity.

a. Consider a  $n$ -stage repeated game  $\Gamma_n(p, q)$  with lack of information on both sides in the independent case with  $K = L \times M$ ,  $p$  probability on  $L$  and  $q$  on  $M$ . By taking the normal form, one gets finite sets of moves, say  $I$  (resp.  $J$ ) for player I (resp. II) and a corresponding pay-off depending on the state say  $c_{ij}^{\ell m}$ . A strategy for player I is thus defined by some vector  $x = (x_i^\ell, \ell \in L)$  where  $x_i^\ell = \Pr(\text{move } i \mid \text{I's type is } \ell)$ , and similarly for player II. Prove that  $V_n(p, q)$  is the value of the following dual linear programming problems:

$$\begin{array}{ll} \max \sum_m q^m u^m & \min \sum_\ell p^\ell u^\ell \\ \sum_{i,\ell} \alpha_i^\ell c_{ij}^{\ell m} \geq u^m & \sum_{j,m} \beta_j^m c_{ij}^{\ell m} \leq u^\ell \\ \sum_i \alpha_i^\ell = p^\ell & \sum_j p_j^m = q^m \\ \alpha_i^\ell \geq 0 & \beta_j^m \geq 0 \end{array} \quad \forall i, \forall \ell \quad \forall j, \forall m$$

and deduce that  $V_n(p, q)$  is concave in  $p$  and convex in  $q$ .

b. Recall from I.3Ex.11eiii p. 36 that a real function  $f$  defined on the product  $C \times D$  of two convex polyhedra is “piecewise bi-linear” if there exists finite partitions of  $C$  (resp.  $D$ ) into convex polyhedra  $C_m$  (resp.  $D_n$ ) such that the restriction of  $f$  to each product  $C_m \times D_n$  is bi-linear. Deduce from the above L.P. formulation in VIEx.6a that  $V_n(p, q)$  is piecewise bi-linear (cf. ex. I.3Ex.11h p. 36).

c. Prove then that in order to compute  $V_n(p, q)$  one can use the following finite algorithm: Compute first  $V_n(0, q)$  and  $V_n(p, 0)$ .

Given  $q_m$  that corresponds to a peak of  $V_n(0, q)$  compute  $V_n(p, q_m)$  and so on (and similarly in the other direction).

When no new peaks are reached extend  $V_n(p, q)$  by bi-linearity.

d. Consider now the dependent case with state space  $K$  and initial probability and partitions  $p, K^I, K^{II}$ . Say that a function  $f$  on  $P$  is I-linear if for all  $p$  its restriction to  $\Pi^I(p)$  is linear and similarly for II-linear. Write  $V_n(p)$  as the value of a linear programming problem as in VIEx.6a to prove that it is I-concave and II-convex. Let  $Q(p) = \{q \in P \mid q^k = \delta\alpha^k\beta^k p^k \text{ where } p_\alpha = (\alpha^k p^k) \in \Pi^I(p) \text{ and } p_\beta = (\beta^k p^k) \in \Pi^{II}(p)\}$  and prove that  $V_n(p)$  is piecewise I-II bi-linear on  $Q(p)$  for each  $p$ .

## 7. Sequential games.

a. *A recursive formula.* Consider  $\Gamma_n(p, q)$  as above (ex. VIEx.6a), but where the players are choosing their moves sequentially, being informed of the previous choice of their opponents. Consider first the reduced game where player I is restricted to use the move  $s$  at stage one (for all states) and let  $V_n^s(p, q)$  be its value. Prove then that  $V_n(p, q) = \text{Cav}_p \max_s V_n^s(p, q)$ .

HINT. Prove that against each first stage strategy  $x_s^\ell$  of player I, player II can decrease the pay-off to  $\sum_s \lambda(s) V_n^s(p^\ell(s), q)$  with  $p^\ell(s) = \Pr_x(\ell \mid s)$  and  $\lambda(s) = \Pr_x(s)$  and use the minmax theorem.

Then deduce by induction that the recursive formula can be rewritten as:

$$nV_n(p, q) = \text{Cav} \max_p \text{Vex} \min_s \left\{ \sum_{\ell, m} p^\ell q^m G_{st}^{\ell m} + (n-1)V_{n-1}(p, q) \right\}.$$

b. *Monotonicity of the values.* Use then the fact that  $\text{Vex}(f + g) \geq \text{Vex } f + \text{Vex } g$  and  $\text{Cav}(f + \text{Cav } f) = 2 \text{Cav } f$  to prove that the sequence  $V_n(p, q)$  is increasing. Deduce that if player II is uninformed, the sequence  $V_n(p)$  is constant (Ponssard and Zamir, 1973).

c. *Speed of convergence.* Let  $f(p, q) = -(\sum_s \min_t \sum_{\ell, m} p^\ell q^m G_{st}^{\ell m} + R_0)$  where  $R_0$  is a constant such that  $v_1 \geq v + f$ . Assume by induction  $nv_n \geq nv + f$  to get

$$(n+1)v_{n+1}(p, q) \geq (n+1) \text{Cav}(\min_p \{u(p, q), v(p, q)\}) + f(p, q)$$

and use VIEx.7b and prop. 4.10 p. 301 to get finally  $|v - v_n| \leq \frac{R}{n}$  for some constant  $R$ .

HINT. Take  $\#L = 1$ ,  $\#M = 2$ ,  $G^1 = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ ,  $G^2 = \begin{pmatrix} 1 & -2 \\ 0 & 1 \end{pmatrix}$  to prove that it is the best bound.

d. *Extend the previous results VIEx.7a, VIEx.7b, VIEx.7c to the dependent case.*

e. *Construction of optimal strategies.* (The length of the game being fixed, we drop the index  $n$ ). For each history  $h$  let  $V^h$  be the value of the “restricted game starting from  $h$ ” and define

$$\begin{aligned} A^h(p) &= \{ \alpha \in \mathbb{R}^L \mid \langle \alpha, q \rangle \leq V^h(p, q) \text{ on } Q \} \\ A^h(p, q) &= \{ \alpha \in A^h(p) \mid \langle \alpha, q \rangle = V^h(p, q) \} \end{aligned}$$

and similarly  $B^h(q)$ ,  $B^h(p, q)$ . Prove using VIEx.7a and ex. VIEx.6b that for any  $\ell \in L(p, q)$  there exists  $(\lambda_s, p(s), \alpha_s)$  for  $s \in S$  such that:

- (1)  $\sum \lambda_s p(s) = p$ ,  $\lambda_s \geq 0$ ,  $\sum \lambda_s = 1$ ,  $V(p, q) = \sum \lambda_s V^s(p(s), q)$
- (2)  $\alpha_s \in A^s(p(s))$
- (3)  $\langle (\sum \lambda_s \alpha_s - \alpha), q' \rangle \geq 0$  for all  $q'$  in  $Q$ .

Deduce that an optimal strategy for player I is to generate  $p(s)$  at stage 1 and to take as new parameter  $\alpha_s$  for his future choice. (Note that this strategy will even be optimal after each history if one chooses after each odd history  $h$ , followed by some move  $t$  of player II, as new state parameter a maximal element in the set  $\alpha' \in A^{ht}(p)$  with  $\alpha' \geq \alpha^h$ .)

**8. Lack of information on  $1\frac{1}{2}$  sides.** Consider the following game  $\Gamma(\lambda, r, s)$  with  $\lambda, r, s$  in  $[0, 1]$ : first,  $t \in \{r, s\}$  is chosen (with  $\Pr(t = r) = \lambda$ ) and this choice is told to player II; then a game with lack of information on one side is played: there are two states of nature, say two pay-off matrices  $A$  and  $B$ , the choice is according to  $t$  and only player I is informed.

a. Write this game as a game with incomplete information on both sides, (dependent case) with 4 states of the world (cf. ch. V) and  $K_I = \{(1, 3), (2, 4)\}$ ,  $K_{II} = \{(1, 2)(3, 4)\}$ ,  $G^1 = G^3 = A$ ,  $G^2 = G^4 = B$  and  $p = (\lambda r, \lambda r', \lambda' s, \lambda' s')$ . (Note that player I knows the true pay-off matrix (state of nature) but not the beliefs of II (state of the world).)

b. Define  $w$  on  $[0, 1]$  by  $w(q) = \text{val}(qA + q'B)$ ,  $\pi(p) = p_1/(p_1 + p_2)$ ,  $\rho(p) = p_3/(p_3 + p_4)$ , and  $\text{Vex}_{[a,b]}$  to be the Vex operator on the interval with end points  $a$  and  $b$ . Prove that:

- $\text{Cav}_I u(p) = \text{Cav } w(p_1 + p_3)$
- $\text{Vex}_{II} u(p) = \text{Vex}_{|\pi(p), \rho(p)|} w(p_1 + p_3)$
- $\text{Vex}_{II} \text{Cav}_I u(p) = \lambda \text{Cav } w(r) + \lambda' \text{Cav } w(s)$
- $\text{Cav}_I \text{Vex}_{II} u(p) = \sup_{t, p^1, p^2} \{ t \text{Vex}_{|\pi(p^1), \rho(p^1)|} w(p_1^1 + p_3^1) + t' \text{Vex}_{|\pi(p^2), \rho(p^2)|} w(p_1^2 + p_3^2) \mid t \in [0, 1], tp^1 + t'p^2 = p, p^i \in \Pi^I(p) \quad i = 1, 2 \}$

c. Example: Let  $A = \begin{pmatrix} 5 & -3 \\ 0 & 0 \end{pmatrix}$ ,  $B = \begin{pmatrix} -3 & 5 \\ 0 & 0 \end{pmatrix}$  and take  $s = 1$ . Verify that  $\bar{v}$ ,  $v$  and  $\underline{v}$  have the shapes given in Figures 18, 19 and 20, hence that there exists a game  $\Gamma(\lambda, r, s)$  with  $\underline{v} < v < \bar{v}$ .

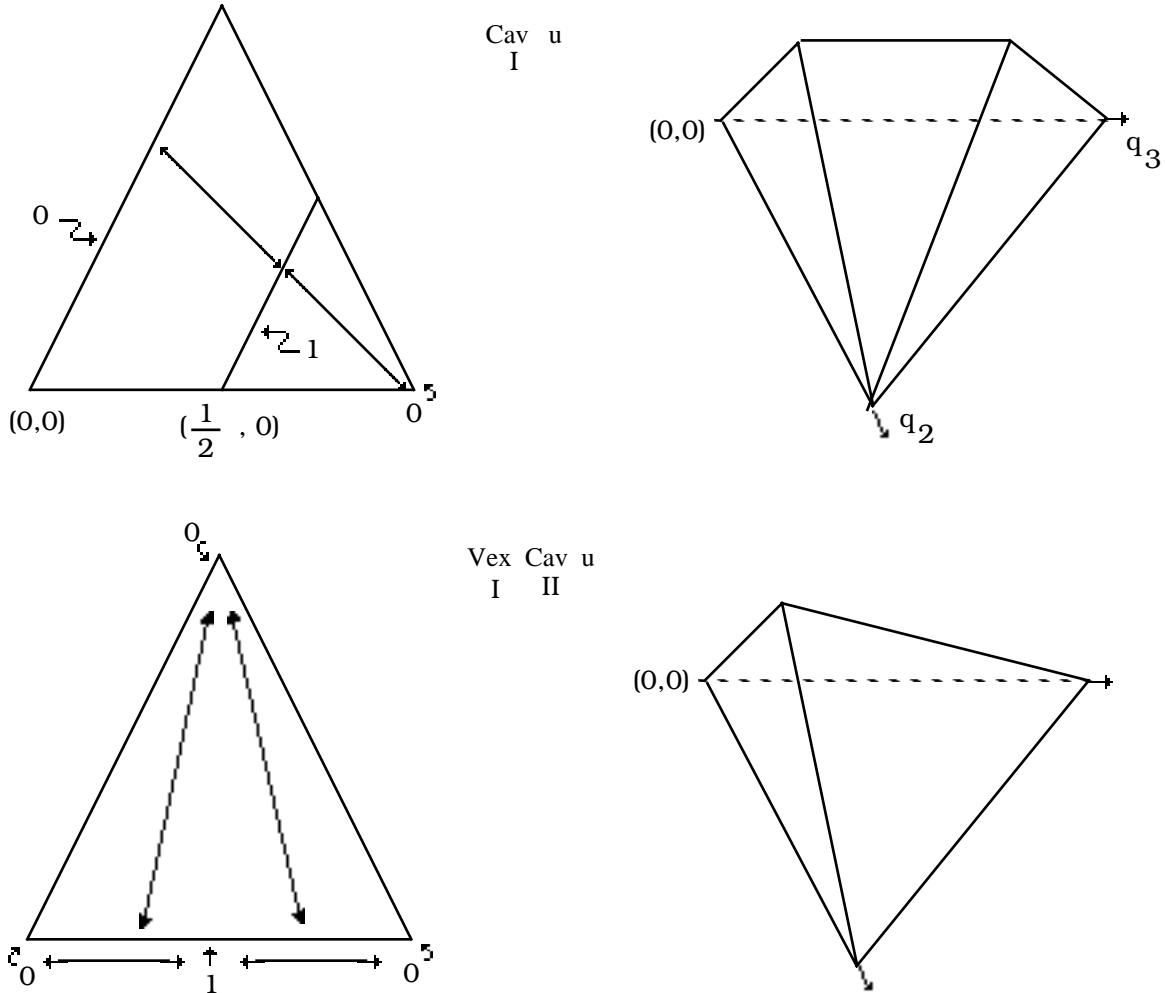
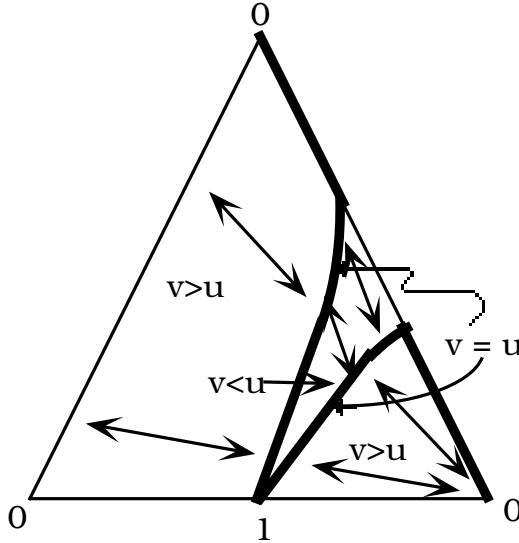


FIGURE 18. The functions  $\text{Cav}_I u$  and  $\text{Vex}_{II} \text{Cav}_I u$

**9. An analytic proof of Theorem 5.3.** Let  $F = \{f \mid f \text{ satisfies } (\alpha)\}$ .

a. Prove that  $F \neq \emptyset$ ,  $\underline{w} = \inf\{f \mid f \in F\}$  belongs to  $F$  and  $\underline{w} = \text{Cav}_I \text{Vex}_{II} \max(u, \underline{w})$ .

FIGURE 19. Partition of  $Q$  according to the relation between  $v$  and  $u$ 

b. Prove that for any real function  $f$  on  $\Pi$ ,  $\text{Cav}_I \text{Vex}_{II} f$  is II-convex.

HINT. Assume  $g$  II-convex. To show  $\text{Cav}_I g$  is II-convex, prove  $\text{Cav}_I g = T^n g$  for  $n$  large enough, with

$$Tg(p) = \sup_{\mu, p^1, p^2} \{ \mu g(p^1) + \mu' g(p^2) \mid p^i \in \pi_I(p), \mu \in [0, 1], \mu p^1 + \mu' p^2 = p \},$$

and that  $T$  preserves II-convexity. For this last point, associate to each triple  $(\mu, p^1, p^2)$  as above and each dual triple  $(\lambda, q^1, q^2)$  (i.e. with  $q^i \in \Pi^{II}(p)$ ,  $\lambda \in [0, 1]$  and  $\lambda q^1 + \lambda' q^2 = p$ ), new variables  $\pi_{ij}$ ,  $\alpha_i$ ,  $\beta_j$ , for  $i = 1, 2$ ,  $j = 1, 2$ , with:

$$\begin{aligned} \pi_{ij} &\in \Pi^{II}(p^i), & \alpha_i \pi_{i1} + \alpha'_i \pi_{i2} &= p^i, & \alpha_i \in [0, 1] & i = 1, 2 \\ \pi_{ij} &\in \Pi^I(p^j), & \beta_j \pi_{1j} + \beta'_j \pi_{2j} &= q^j, & \beta_j \in [0, 1] & j = 1, 2 \end{aligned}$$

and  $\mu \alpha_1 = \lambda \beta_1$ ,  $(1 - \mu) \alpha_2 = \lambda(1 - \beta_2)$ ,  $\mu(1 - \alpha_1) = (1 - \lambda)\beta_2$ ,  $(1 - \mu)(1 - \alpha_1) = (1 - \lambda)(1 - \beta_2)$ .

c. Deduce that  $\underline{w} = \text{Vex}_{II} \max(u, \underline{w})$ . Define  $\{\underline{u}_n\}$  by  $\underline{u}_{n+1} = \text{Cav}_I \text{Vex}_{II} \max(u, \underline{u}_n)$  with  $\underline{u}_0 = -\infty$ , and prove that  $\underline{u}_n$  increases uniformly to  $\underline{w}$ . Introduce similarly  $\overline{w}$  and  $\overline{u}_n$ .

d. Let  $\mathcal{U} = \{u \mid u \text{ can be written as } u(p) = \max_{i \in I} \min_{j \in J} \sum_k a_{ij}^k p^k, I, J \text{ finite sets}\}$ . Prove that, for all  $u$  in  $\mathcal{U}$ ,  $\overline{w} \leq \underline{w}$ .

HINT. Define  $v_0 \equiv 0$  and  $nv_n(p) = \text{Cav}_I \max_i \text{Vex}_{II} \min_j (\sum_k a_{ij}^k p^k + (n-1)v_{n-1}(p))$  and prove:  $v_n \leq \underline{u}_n$ ,  $v_n \geq \overline{u}_n + R/n$  for some constant  $R$ . Compare with ex. VIEx.7c.

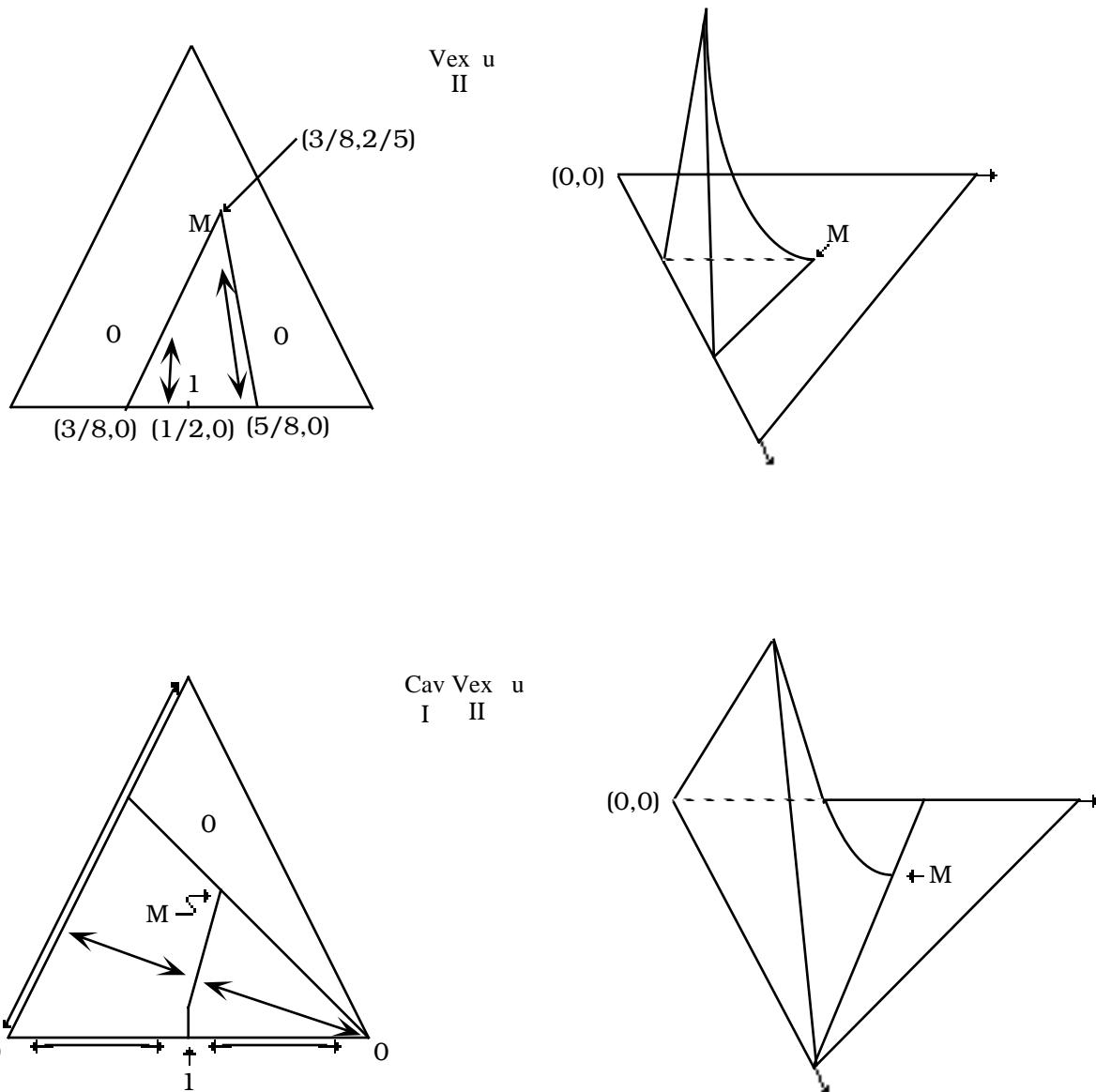
e. Prove that  $\underline{w} \leq \overline{w}$ .

HINT. Show that one obtains the same  $\underline{w}$ ,  $\overline{w}$  when starting with  $u' = \max(u, \underline{u}_1)$  and deduce inductively that  $\overline{w} \geq \underline{u}_n$ , and the result for  $u \in \mathcal{U}$ .

f. Show finally that  $\mathcal{U}$  is dense in  $C(\Pi)$  (compare with prop. 5.1 p. 302) and use prop. 5.2 p. 302.

**10.** Consider a game  $\Gamma$  as defined in ch. V, with corresponding  $v_n$  and  $u$ . Define a game  $\Gamma'_L$  where the set of moves of player I is now  $S \times K$ , the signalling matrices are  $H'^I$ ,  $H'^{II}$  with

$$H'^I((s, k), t) = H^{I,k}(s, t)$$

FIGURE 20. The functions  $\text{Vex}_{\text{II}} u$  and  $\text{Cav}_I \text{Vex}_{\text{II}} u$ 

and similarly for  $H'^{\text{II}}$ , with the same initial information, and as pay-off matrices:

$$G^k((s, k), t) = \begin{cases} -L & \text{if } k' \neq k, \text{ for all } s, t \\ G^k(s, t) & \text{if } k = k' \end{cases}$$

- a. Prove that  $\Gamma'_L$  belongs to the class of ch. VI with associated  $v'_{n,L}$  and  $u'_L$ , hence  $\lim_{n \rightarrow \infty} v'_{n,L} = \text{Cav } u'_L$ .
- b. Show that  $v'_{n,L} \geq v_n$  and that  $\text{Cav } u'_L \searrow \text{Cav } u$  as  $L \rightarrow +\infty$ .
- c. Deduce theorem 3.5 p. 195.

HINT. To avoid circular reasoning, since Theorem 3.5 is apparently used in part C of 3.1 p. 283, proceed as follows: first the convergence of  $v_n$  is established. This yields that  $\lim v_n = \text{Cav } u$  for the games of ch. V, as seen above. This implies immediately  $v_\infty = \text{Cav } u$  for those games, hence theorem 3.5 follows and finally 3.1.

**11. Asymptotically optimal strategies in finite games.** (Heuer, 1991a) Assuming theorem 5.3 p. 303, cf. also ex. VIEx.9 p. 319, we construct strategies that guarantee  $v + O(1/\sqrt{n})$  in  $\Gamma_n$  for the case of standard signalling, hence implying cor. 4.9 p. 301 — on the limit of  $v_n$  and on the speed of convergence.

The basic idea is reminiscent of Blackwell's approachability strategy, (say for II, cf. ex. VEx.2 p. 253), starting with a vector  $\beta$  supporting  $v$ , but then aiming at stage  $r$  to reach  $\beta_r$  (in the remaining  $n - r + 1$ -stage game) so that, given the past pay-offs the average would be  $\beta$ . We will consider the independent case, (cf. examples in sect. 7 p. 305 and ex. VIEx.6–VIEx.7 p. 317). Hence  $K = L \times M$ ,  $p \in \Delta(L)$ ,  $q \in \Delta(M)$  and  $v = \text{Cav}_p \min(u, v) = \text{Vex}_q \max(u, v)$ . Let  $B(q) = \{\beta \in \mathbb{R}^K \mid \langle \beta, p \rangle \geq v(p, q), \forall p \in \Delta(L)\}$ , and  $B(p, q) = \{\beta \in B(q) \mid \langle \beta, p \rangle = v(p, q)\}$  is the supergradient of  $v$  at the point  $(p, q)$ , in the direction of  $p$ .

a. Note that if  $q = \sum_j \lambda_j q_j$ , with  $q_j \in \Delta(M)$  and  $\lambda \in \Delta(J)$ , then  $\sum_j \lambda_j B(q_j) \subseteq B(q)$ .

Prove that if moreover  $v(p, q) = \sum_j \lambda_j v(p, q_j)$ , with  $v(p, q_j) = u(p, q_j)$  and  $v(p, \cdot) < u(p, \cdot)$  on the interior of the convex hull of the  $q_j$ 's, then:  $\sum_j \lambda_j B(q_j) = B(q)$ . (Use the continuity of  $u$  and  $v$  and the comments in sect. 7 p. 305).

b. Recall that  $\bar{\gamma}_n(\sigma, \tau) = \mathbb{E}_{\sigma, \tau}^{p, q} \left( \frac{1}{n} \sum_{r=1}^n g_r \right)$  where  $g_r$  is the pay-off at stage  $r$ . Define  $p_r^\ell = P(\ell \mid \mathcal{H}_r)$ ,  $q_r^m = P(m \mid \mathcal{H}_r)$  and  $\rho_r^\ell = \mathbb{E}_{\sigma, \tau}^{\ell, q} (g_r \mid \mathcal{H}_r)$ , (recall standard signalling). Hence  $\bar{\gamma}_n(\sigma, \tau)$  can also be written as  $\mathbb{E}_{\sigma, \tau} \left( \frac{1}{n} \langle p_n, \sum_r \rho_r \rangle \right)$ .

c. We now define a strategy for II inductively. Given  $(p, q)$  and  $\beta \in B(p, q)$ , let  $\pi_1 = p$ ,  $q_1 = q$ ,  $\xi_1 = \tilde{\beta}_1 = \beta_1 = \tilde{\xi}_1 = \beta$ . At stage 1 player II plays optimally in  $D(p_1, q_1)$ . Then define  $\xi_2$  by  $n\xi_1 = (n-1)\xi_2 + \rho_1$ . Similarly at stage  $r$ , given  $\xi_r$  we consider the following cases:

- (1) if  $\xi_r \in B(q_{r-1})$ ,  $\tau$  is arbitrary at this stage. One puts  $\pi_r = \pi_{r-1}$ ,  $q_r = q_{r-1}$ ,  $\tilde{\xi}_r = \xi_r = \tilde{\beta}_r = \beta_r$ .
- (2) if  $\xi_r \notin B(q_{r-1})$ , let  $\beta_r$  denote its projection on  $B(q_{r-1})$ ; this defines  $\pi_r \in \Delta(L)$ ,  $\pi_r$  proportional to  $\beta_r - \xi_r$ , such that  $\beta_r \in B(\pi_r, q_{r-1})$ .
  - if  $v(\pi_r, q_{r-1}) \geq u(\pi_r, q_{r-1})$ ,  $\tau$  consists of playing optimally in  $D(\pi_r, q_r)$  at that stage, with  $q_r = q_{r-1}$ . Then  $\tilde{\beta}_r = \beta_r$  and  $\tilde{\xi}_r = \xi_r$ .
  - if  $v(\pi_r, q_{r-1}) < u(\pi_r, q_{r-1})$ , use VIEx.11b) to decompose  $\beta_r$  as  $\sum_j \beta_{r,j}$  with  $\beta_{r,j} \in B(\pi_r, q_{r-1,j})$ . In this case  $\tau$  consists of first doing a splitting (prop. 1.2 p. 184) to generate the  $q_{r-1,j}$ 's and then if  $j$  is chosen, in playing optimally in  $D(\pi_r, q_r)$  with, obviously,  $q_r = q_{r-1,j}$ . Then let  $\tilde{\beta}_r = \beta_{r,j}$  and  $\tilde{\xi}_r = \xi_r - \beta_r + \tilde{\beta}_r$  and note that  $\mathbb{E}(\tilde{\xi}_r \mid \xi_r) = \xi_r$ .

Finally let us define  $\xi_{r+1}$  through the following equation:  $(n-r+1)\tilde{\xi}_r = (n-r)\xi_{r+1} + \rho_r$ .

Show that given any non-revealing  $\bar{\sigma}$  one has:  $\mathbb{E}_{\bar{\sigma}, \tau}[\langle (\xi_r - \beta_r), (\tilde{\beta}_r - \rho_r) \rangle] \leq 0$ .

d. Show that

$$\bar{\gamma}_n(\sigma, \tau) - v(p, q) = \frac{1}{n} \mathbb{E}_{\sigma, \tau}(\langle p_n, \rho_n - \xi_n \rangle) \leq \frac{2C}{n} + \frac{1}{n} \mathbb{E}_{\sigma, \tau}(\langle p_n, \beta_n - \xi_n \rangle)$$

and note also that there exists  $\bar{\sigma}$  non-revealing such that:

$$|\mathbb{E}_{\sigma, \tau}(\langle p_n, \beta_n - \xi_n \rangle)| \leq \left( \mathbb{E}_{\bar{\sigma}, \tau}(\|\beta_n - \xi_n\|_2^2) \right)^{1/2}$$

Then prove inductively:

$$\mathbb{E}[\|\xi_r - \beta_r\|^2] \leq (\#L) 4C^2 r / (n+1-r)^2$$

and conclude.

HINT.  $\mathbb{E}[\|\xi_{r+1} - \beta_{r+1}\|^2] \leq \mathbb{E}[\|\xi_{r+1} - \tilde{\beta}_r\|^2]$ ; then use the equality  $\tilde{\xi}_r - \tilde{\beta}_r = \xi_r - \beta_r$  and VIEx.11e).

e. Extend the result to the dependent case.

## 12. A continuum of types on both sides. (After Forges, 1988b)

COMMENT 7.5. In this chapter — notably in sect. 3.b — we could not treat the case of a continuum of types of both sides; only the “approached player” could have a continuum of types. As underscored repeatedly in ch. III, the inability to study the general case at this stage seems at least in part to be due to a lack of study of the fundamental concepts in ch. III — cf. e.g. introduction to sect. 4.b, many remarks after, and most of sect. 4.d, in particular remark 4.11 and the final remarks. In particular, the first natural step towards a generalisation would be to reformulate the known results (in particular Theorem 3.13 p. 293) in the canonical framework of ch. III — and it was shown there (final remarks) that even the basic concept of concavification was not clear.

Yet we need such results in ch. IX — in order to study communication equilibria of games with a *finite* set of types —, for the case where the approaching player has a continuum of types (even though there the approached player could have a finite set of types). Another reason for studying the general case is the basic problem underlying this book: to obtain maxmin, minmax and  $\lim v_n$  for all two-person zero-sum repeated games (ch. IV), or at least in a first time for all two-person zero-sum repeated games with incomplete information (i.e., to generalise this chapter to the case where the signalling function  $Q$  is state dependent). Indeed, as shown in ch. VIII (sect. 2), the general case leads to situations where the number of points in the support of the current posterior consistent probability (ch. IV, sect. 3) grows to infinity, so conceivably one might as well study them from the outset without restricting the support to be finite, and anyway a number of concepts and tools will presumably be required to work directly on  $\mathcal{P}$ : chief candidates for those will be the ones already needed in the particular case of this chapter.

It is thus clear that a satisfactory treatment of the present chapter would require the framework of a continuum of types on both sides.

In a first stage, one could restrict the prior to be absolutely continuous w.r.t. some product measure (i.e., w.r.t. the product of its marginals). This restriction — which is preserved when going from an information scheme to its canonical representation — would keep the complete symmetry between both players in the assumptions of this chapter (hence allowing by duality to cut the number of statements and proofs in two) and still be sufficiently general to encompass all known cases and to lead to the elucidation of the right concepts, while at the same time being technically quite helpful, e.g. in arguments like the one below, or possibly in proving that player I can defend Vex Cav.

We present next a very first result in this direction — not only is it extremely partial, but chiefly it involves the additional restriction of statistically perfect monitoring of player II by player I. But it will suffice for our applications in ch. IX.

a. Consider a game with statistically perfect monitoring of player II by player I. Describe the initial information of the players in the following way:

- the space of types of player II is a measurable space  $(J, \mathcal{J})$ .
- the space of types of player I is a measurable space  $(I, \mathcal{I})$  together with a transition probability  $\theta (= \theta_i(k, dj))$  from  $(I, \mathcal{I})$  to  $(J, \mathcal{J}) \times K$ , and an initial probability measure  $\gamma$  on  $(I, \mathcal{I})$ .

Let as usually — cf. remark 3.16 p. 297 —  $\mathbf{H} = \{ h: \mathbb{R}_+^K \rightarrow \mathbb{R} \mid h \text{ is convex, positively homogeneous of degree 1, } h|_{\Delta(K)} \text{ has Lipschitz constant } C \text{ and is } \leq C, \int_{\Delta(K)} h(\pi) \mu(d\pi) \geq u(\mu) \forall \mu \in \Delta(\Delta(K)) \}$ , and (theorem 3.13)  $\mathbf{T}$  be the set of transition probabilities  $\mathbf{t}$  from  $(J, \mathcal{J})$  to  $\mathbf{H}$ .

Given  $\mathbf{t} \in \mathbf{T}$ , let  $z_{\mathbf{t}}(\theta) = \int_{\mathbf{H}} h \left[ \left( \int_J \mathbf{t}(dh|j) \theta(k, dj) \right)_{k \in K} \right] \quad \forall \theta \in \Delta(K \times J)$ .

- (1)    • player II’s strategy induced by  $\mathbf{t}$  guarantee that type  $i$ ’s pay-off is  $\leq z_{\mathbf{t}}(\theta_i)$   
 •  $z_{\mathbf{t}}(\theta_i)$  is measurable on  $(I, \mathcal{I})$ , and convex in  $\theta$ .
- (2) Assume that the measures  $\theta_i$  ( $i \in I$ ) are dominated — i.e.  $\exists \bar{\theta} \in \Delta(J)$  such that  $\forall i \in I$ ,  $\theta_i$  is absolutely continuous w.r.t.  $\bar{\theta}$ . Then  $\forall \mathcal{L}$  Banach limit,  $\forall \tau_0$  strategy of player II,

$\exists \mathbf{t} \in \mathbf{T}$  s.t.

$$\forall \lambda \in \Delta(I) \quad \sup_{\sigma} \mathcal{L} E_{\sigma, \tau_0}^{\lambda}(\bar{g}_m) \geq \int z_{\mathbf{t}}(\theta_i) \lambda(d_i).$$

HINT. The first point of 1 is as in the text; measurability is standard [observe first that  $i \mapsto \int_J \mathbf{t}(dh|j) \theta_i(k, dj)$  is a transition probability from  $(I, \mathcal{J})$  to  $K \times H$ , hence a measurable map to  $\Delta(K \times H)$  (9.e), and use the lower-semi-continuity of  $\phi$  in IIIEx.4 p. 142]. Similarly convexity follows from that of  $\phi$ .

As for point 2: let  $F_{\sigma}(\lambda) = \mathcal{L} E_{\sigma, \tau_0}^{\lambda}(\bar{g}_m)$ ,  $F(\lambda) = \sup_{\sigma} F_{\sigma}(\lambda)$ ; clearly  $F_{\sigma}$  is affine, so  $F$  is convex. There is no loss to assume  $\mathcal{J}$  separable, by reducing it first to the  $\sigma$ -field generated by  $\tau_0$ . Fix then an increasing sequence of measurable partitions  $\Pi_n$  that generates  $\mathcal{J}$ . For any  $\Pi$ , define  $\Gamma^{\Pi}$  as the same game, but where player II's  $\sigma$ -field  $\mathcal{J}$  is reduced to  $\mathcal{J}^{\Pi}$ , the (finite)  $\sigma$ -field spanned by  $\Pi$ . Viewing  $\tau_0$  as a transition probability from  $(J, \mathcal{J})$  to the pure strategy space in  $\Gamma_{\infty}$ , define also  $\tau_0^{\Pi}$ , the strategy in  $\Gamma^{\Pi}$  corresponding to  $\tau_0$  in  $\Gamma$ , as  $E_{\bar{\theta}}(\tau_0(j) \mid \mathcal{J}^{\Pi})$ . Let finally  $F_{\sigma}^n$  and  $F^n$  be the functions corresponding to  $F_{\sigma}$  and to  $F$  in  $\Gamma^{\Pi_n}$  with  $\tau_0^{\Pi_n}$ .

Observe that, since  $\mathcal{J}^{\Pi_n} \subseteq \mathcal{J}$ ,  $\tau_0^{\Pi_n}$  is also a strategy in  $\Gamma$ ; and clearly, for all  $\lambda, \sigma$  and  $m$ ,  $E_{\sigma, \tau_0^{\Pi_n}}^{\lambda}(\bar{g}_m)$  is the same, say  $F_{\sigma, m}^n(\lambda)$ , whether computed in  $\Gamma$  or in  $\Gamma^{\Pi_n}$ .

- Let  $f_k^n(i, j) = \sum_{B \in \Pi_n} \mathbb{1}_B(j) \theta_i(B \times k) / \bar{\theta}(B \times k)$ , and  $f_k = \liminf_{m \rightarrow \infty} f_k^n$ .  $f^n$  and  $f$  are measurable on  $I \times J$ , and by the martingale convergence theorem  $f^n$  converges to  $f$  a.e. and in  $L_1$  under  $\bar{\theta}(dj)\lambda(di)$ ,  $\forall \lambda$ . And  $f$  is a Radon-Nikodym density of  $\theta_i(k, dj)\lambda(di)$  w.r.t.  $\bar{\theta}(dj)\lambda(di)$ . Now, since  $|\bar{g}_m| \leq C$  uniformly, computation in  $\Gamma$  shows that

$$|F_{\sigma, m}^n(\lambda) - F_{\sigma, m}(\lambda)| \leq C \int \sum_k |f_k^n(i, j) - f_k(i, j)| \bar{\theta}(dj)\lambda(di) \stackrel{\text{def}}{=} CG^n(\lambda),$$

with  $G^n(\lambda)$  converging to zero as seen above. Hence  $|F_{\sigma}^n(\lambda) - F_{\sigma}(\lambda)| \leq CG^n(\lambda)$ , and thus also

$$(9) \quad F^n(\lambda) \rightarrow F(\lambda)$$

- Observe that  $\Gamma^{\Pi}$  is a game with a finite set of types for player II, hence theorem 3.13 is applicable. Denote by  $\pi(\lambda)$  the image measure of  $\lambda$  on  $\Delta(K \times \Pi)$ ; since  $\tau_0^{\Pi}$  is still a strategy in the “semi-canonical” game of theorem 3.13,  $F^{\Pi}(\lambda)$  is a function  $\varphi$  of  $\pi(\lambda)$ . Define  $\varphi$  as  $+\infty$  outside the range of the map  $\pi$ . By point 7, we have  $\varphi \geq \text{Vex Cav}$  on  $\Delta(\Delta(K \times \Pi))$ , and the convexity of  $\varphi$  follows immediately from that of  $F^{\Pi}$  and from the linearity of  $\lambda \mapsto \pi(\lambda)$ . Hence by point 6 there exists  $\forall n \mathbf{t}^n \in \mathbf{T}$  ( $\mathcal{J}^{\Pi_n}$ -measurable) such that

$$(10) \quad \int z_{\mathbf{t}^n}(\theta_i^n) \lambda(di) \leq F^n(\lambda) \quad \forall \lambda$$

where  $\theta_i^n = f_k^n(i, j) \bar{\theta}(dj)$ .

- $\mathbf{T}$  is compact metric since  $\mathbf{H}$  is so and  $\mathcal{J}$  is separable (II.1Ex.17c p. 76). Extract thus if necessary a subsequence such that  $\mathbf{t}^n$  converges, say to  $\mathbf{t}$ . Consider now  $P_i^n = \int_J \mathbf{t}^n(dh|j) f_k^n(i, j) \bar{\theta}(dj) \in \Delta(K \times H)$ . Since  $f^n$  converges in  $L_1$  to  $f$  and  $\mathbf{t}^n$  converges weak\* to  $\mathbf{t}$ , we get that  $P_i^n$  converges weakly to  $P_i$  for all  $i$ . So, by IIIEx.4 p. 142 — using the function  $f(h, p) = h(p)$  —,  $\liminf_{n \rightarrow \infty} \int_H h \left[ \left\{ \int_J P_i^n(k, dh) \right\}_{k \in K} \right] \geq \int_H h \left[ \left\{ \int_J P_i(k, dh) \right\}_{k \in K} \right]$ , i.e.  $\liminf_{n \rightarrow \infty} z_{\mathbf{t}^n}(\theta_i^n) \geq z_{\mathbf{t}}(\theta_i)$ . Hence, by Fatou's lemma, (9) and (10) yield the result.

b. *Particular cases.*

i. Cf. remark 3.17 p. 298.

ii. If player I's information includes the knowledge of the true state of nature (as in ex. VIEx.8), then his posteriors are, whatever  $\mathbf{t}$  is used, concentrated on the true state, so only the values of  $h$  at the vertices of  $\Delta(K)$  matter. Thus any  $h$  can be taken affine (replaced by its concavification). Further randomising over different  $h$ 's serves no purpose, since the posteriors are not affected: one can as well use the average  $h$ . So in this case, a “strategy”  $\mathbf{t}$  of player II is simply a measurable map from  $(J, \mathcal{J})$  to  $\{h \in \mathbb{R}^K \mid \forall p \in \Delta(K), \langle h, p \rangle \geq u(\mu) \forall \mu \in \Delta(\Delta(K)) \text{ with } \bar{\mu} = p\}$ .

[In case there is statistically perfect monitoring on both sides (cf. VIEx.12bi above), this last condition clearly reduces to:  $\langle h, p \rangle \geq u(p) \forall p \in \Delta(K)$ .]

iii. Observe how the case sub VIEx.12bii above a direct generalisation yields of the approachability results in ch. V.



## CHAPTER VII

### Stochastic Games

A stochastic game is a repeated game where the players are at each stage informed of the current state and the previous moves. According to the general model of ch. IV, this means that the signal transmitted to each player  $i$  according to the transition probability  $Q$  includes the new state and the previous moves. It follows that the game is equivalently described by the action sets  $S^i$ ,  $i \in \mathbf{I}$ , the state space  $K$ , a transition probability  $P$  from  $S \times K$  to  $K$  and a pay-off mapping  $g$  from  $S \times K$  to  $\mathbb{R}^{\mathbf{I}}$ .

#### 1. Discounted case

It appears that in this framework our finiteness assumptions are not really used and that we can work with the following more general setup:

- the state space is a measurable space  $(\Omega, \mathcal{A})$ ;
- the action space of player  $i$  is a measurable space  $(S^i, \mathcal{S}^i)$  (with  $S = \prod_i S^i$ );
- $P$  is a transition probability from  $\Omega \times S$  to  $\Omega$ , hence for  $A$  in  $\mathcal{A}$ ,  $P(A \mid \omega, s)$  is the probability that tomorrow's state belongs to  $A$  given today's state  $\omega$  and actions  $s$ .

A strategy for a player is again a transition probability from histories of the form  $(\omega_1, s_1, \dots, s_{n-1}, \omega_n)$  to actions. A **Markov** strategy depends only, at stage  $n$ , on the current state  $\omega_n$ . A **stationary** strategy is one which is a stationary (i.e. time invariant) function of the infinite past — accompanied by a fictitious history before time zero. To force the influence of the far-away past to vanish, one may in addition impose the function to be continuous, say in the product topology.

The main tool when dealing with the discounted case is the following class of one-shot games: Given a vector  $f (= (f^i)_{i \in \mathbf{I}})$  of bounded real-valued measurable functions on  $(\Omega, \mathcal{A})$ , define the single stage game  $\Gamma(f)_\omega$ ,  $\omega \in \Omega$ , with action sets  $S^i$  and (vector) pay-off:

$$\phi(f)_\omega(s) = g(\omega, s) + \int f(\tilde{\omega}) P(d\tilde{\omega} \mid \omega, s).$$

**1.a. Zero-sum case.** Here the basic technique for establishing the existence of a value is based on the “contraction mapping principle”:

LEMMA 1.1. Let  $(E, d)$  be a complete metric space,  $\varepsilon > 0$  and  $f: E \rightarrow E$  such that  $d(f(x), f(y)) \leq (1 - \varepsilon)d(x, y)$  for all  $(x, y)$ . Then  $f$  has a unique fixed point  $\bar{x} \in E$ , and for any  $x \in E$ , the sequence  $f^n(x)$  converges to  $\bar{x}$ .

PROOF.  $d(f^{n+1}(x), f^n(x)) \leq (1 - \varepsilon)d(f^n(x), f^{n-1}(x))$ , hence by induction  $d(f^{n+1}(x), f^n(x)) \leq (1 - \varepsilon)^n d(f(x), x)$ , thus by the triangle inequality

$$d(f^{n+k}(x), f^n(x)) \leq \left( \sum_{i=n}^{\infty} (1 - \varepsilon)^i \right) d(f(x), x) = \frac{(1 - \varepsilon)^n}{\varepsilon} d(f(x), x) :$$

since the right hand member goes to zero with  $n$ , the sequence  $f^n(x)$  is a Cauchy sequence, hence convergent (completeness), say to  $\bar{x}$ . But  $d(f^{n+1}(y), f(\bar{x})) \leq (1 - \varepsilon)d(f^n(y), \bar{x})$ ,

hence going to the limit yields  $d(\bar{y}, f(\bar{x})) \leq (1 - \varepsilon)d(\bar{y}, \bar{x})$ : setting  $y = x$ , hence  $\bar{y} = \bar{x}$  yields  $d(\bar{x}, f(\bar{x})) = 0$ :  $\bar{x}$  is a fixed point; setting then  $y = \bar{y}$  to be any other fixed point yields  $d(\bar{y}, \bar{x}) \leq (1 - \varepsilon)d(\bar{y}, \bar{x})$ , hence  $\bar{y} = \bar{x} : \bar{x}$  is the unique fixed point. ■

The idea about the use of the contraction principle in proving the existence of a value of  $\Gamma_\lambda$  is that any uncertainty about tomorrow's pay-off is reduced by a factor of  $(1 - \lambda)$ ,  $\lambda$  being the discount factor, when evaluated in today's terms. So if one can solve "today's" game for any given pay-offs for the future, one will get a contraction mapping.

The basic requirement is thus that "today's game" has a value for any choice of a "pay-off for the future" in an appropriate complete metric space — and yields a pay-off in the same metric space. So that for any given  $f$  in  $(B, d)$ , a complete metric space of bounded measurable functions on  $(\Omega, \mathcal{A})$ , with  $d$  the uniform distance, our aim is to show that (with the notation of subsection 3.b p. 157):

- (1) For each  $\omega$  in  $\Omega$  the game  $\Gamma(f)_\omega$  has a value, say  $\Psi(f)(\omega)$ .
- (2)  $\Psi(f)$  belongs to  $B$ .
- (3) The games  $\Gamma(f)$  have  $\varepsilon$ -optimal strategies (i.e., strategies that are  $\varepsilon$ -optimal for any  $\omega \in \Omega$ ).

LEMMA 1.2. Assume that the distance on  $B$  is  $d(f_1, f_2) = \sup_{\Omega} |f_1(\omega) - f_2(\omega)|$ , then under 1 and 2, there exists a solution  $V_\lambda \in B$  of  $f = \Psi[(1 - \lambda)f]$ .

PROOF.  $\Psi$  maps  $B$  into  $B$  and is clearly monotone. Since  $\Psi(f + c) = c + \Psi(f)$  for any constant function  $c$ , we have:

$$d(\Psi(f_1), \Psi(f_2)) \leq d(f_1, f_2).$$

Thus  $f \mapsto \Psi[(1 - \lambda)f]$  satisfies lemma 1.1. ■

LEMMA 1.3. Under 1, 2 and 3  $\Gamma_\lambda$  has a value  $\lambda V_\lambda$ . If  $\mu$  is an  $\varepsilon$ -optimal strategy in  $\Gamma((1 - \lambda)V_\lambda)$ , then the corresponding stationary strategy  $\bar{\mu}$  is  $\varepsilon$ -optimal in  $\Gamma_\lambda$ .

PROOF. For any  $f$  in  $B$  and any strategies  $\sigma$  and  $\tau$  in  $\Gamma_\lambda$ ,

$$\mathbb{E}_{\sigma, \tau} \left( \sum_{m \leq n} (1 - \lambda)^{m-1} g_m + (1 - \lambda)^n f(\omega_{n+1}) \right)$$

converges (uniformly) to  $\bar{\gamma}_\lambda(\sigma, \tau)/\lambda$  as  $n$  goes to  $\infty$ .

Conditionally to  $\mathcal{H}_m$  (generated by  $(\omega_1, s_1, \dots, s_{m-1}, \omega_m)$ ) one has, by the definition of  $\bar{\mu}$ :

$$\mathbb{E}_{\bar{\mu}, \tau} (g_m + (1 - \lambda)(V_\lambda(\omega_{m+1}) - \delta) \mid \mathcal{H}_m) \geq V_\lambda(\omega_m) - \varepsilon - (1 - \lambda)\delta.$$

So that taking  $\delta = \varepsilon/\lambda$  one obtains

$$\mathbb{E}_{\bar{\mu}, \tau} \left( \sum_{m \leq n} (1 - \lambda)^{m-1} g_m + (1 - \lambda)^n (V_\lambda(\omega_{n+1}) - \frac{\varepsilon}{\lambda}) \right) \geq V_\lambda - \frac{\varepsilon}{\lambda}$$

hence  $\bar{\gamma}_\lambda(\bar{\mu}, \tau) \geq \lambda V_\lambda - \varepsilon$ . ■

REMARK 1.1. Using similarly  $V_n$  for the non-normalised value  $nv_n$  of  $\Gamma_n$ , one obtains in the same way  $V_{n+1} = \Psi(V_n)$ , with  $V_0 = 0$ .

Two typical illustrations follow — just to illustrate the method; they do not strive for the utmost generality (cf. e.g. ex. VIIEx.18 for the "right" form of the next).

PROPOSITION 1.4. The state space is a standard Borel (cf. App.6 p. 426) space  $(\Omega, \mathcal{A})$ , the action sets are compact metric spaces  $S$  and  $T$ , the pay-off function  $g(\omega, s, t)$  and the transition probability  $P(A \mid \omega, s, t)$  are, for each given  $A \in \mathcal{A}$ , measurable on  $(\Omega \times S \times T)$

and are, for fixed  $\omega$ , separately continuous in  $s$  and in  $t$ . Further  $g$  is bounded. Then the discounted game has a value and  $\mathcal{A}$ -measurable optimal stationary Markov strategies.

**PROOF.** Consider the Banach space  $B$  of bounded measurable functions on  $(\Omega, \mathcal{A})$ . For  $f \in B$  (representing the future (non-normalised) pay-off), today's pay-off is:  $h(\omega, s, t) = \phi((1 - \lambda)f)_\omega(s, t)$ . The assumptions guarantee that this is, like  $g$ , measurable on  $\Omega \times S \times T$  and, for each  $\omega$ , separately continuous in  $s$  and  $t$ . We know from theorem 2.6 p. 17 that, for each  $\omega$ , such a game has a value  $V(\omega)$ . Clearly  $\phi((1 - \lambda)f)$  is bounded, hence  $V$  also. There remains to show measurability of  $V$ .

For  $\mu$  in  $\Delta(S)$ , let  $H(\omega, \mu, t) = \int h(\omega, s, t)\mu(ds)$ : measurability of  $H$  is easy and well known (e.g., just approximate the integrand by a linear combination of indicator functions of sets  $A \times X \times Y$ , with  $A \in \mathcal{A}$ , and  $X \in \mathcal{S}$ ,  $Y \in \mathcal{T}$ ), its linearity and (weak) continuity in  $\mu$  is obvious, and the continuity in  $t$  follows immediately from Lebesgue's bounded convergence theorem. Hence  $F(\omega, \mu) = \min_t H(\omega, \mu, t)$  is measurable (because by continuity it is sufficient to take the infimum over a countable dense set), and upper semi-continuous and concave in  $\mu$ . Similarly for  $\nu$  in  $\Delta(T)$ , we have a lower semi-continuous, convex measurable function  $G(\omega, \nu)$  defined by  $G(\omega, \nu) = \max_s \int h(\omega, s, t)\nu(dt)$ . Hence the graph  $\{(\omega, \mu, \nu) \mid F(\omega, \mu) \geq G(\omega, \nu)\}$  is measurable and has, for each  $\omega$ , compact non-empty values — the corresponding optimal strategy pairs. Such a graph has a measurable selection (7.i p. 427 and 8.b p. 428) yielding thus measurable optimal strategy selections  $(\mu_\omega, \nu_\omega)$ , and also the measurability of  $V(\omega) = F(\omega, \mu_\omega)$  (by composition). ■

The result follows now from lemma 1.3. ■

In the next proposition, we relax the very strong continuity assumption on the transition probabilities as a function of the actions, at the expense of stronger assumptions on the dependence on the state:

**PROPOSITION 1.5.** *Assume the state space  $\Omega$  is metrisable and the action sets  $S$  and  $T$  are compact metric, and that  $g(\omega, s, t)$  and  $\int f(\tilde{\omega})P(\tilde{\omega} \mid \omega, s, t)$  are, for each bounded continuous  $f$ , continuous on  $\Omega \times S$  for fixed  $t$  and on  $\Omega \times T$  for fixed  $s$ . Further  $g$  is bounded. Then the discounted game has a value (which is continuous as a function of the initial state) and Borel-measurable optimal stationary Markov strategies.*

**PROOF.** We show that, under those assumptions, we get a contracting operator on the space  $B$  of bounded continuous functions on  $\Omega$ . We use the notation of the previous proof.  $h$  is separately continuous in  $(\omega, s)$  and  $(\omega, t)$ . Hence  $V(\omega)$  exists. It follows that  $H$  is continuous in  $(\omega, t)$  (Lebesgue's dominated convergence theorem as before). Hence  $F(\omega, \mu)$  is, for fixed  $\mu$ , continuous in  $\omega$  (continuity of  $H(\cdot, \mu, \cdot)$  and compactness of  $T$ ). Hence  $V(\omega) = \sup_\mu F(\omega, \mu)$  is lower semi-continuous. Dually  $V$  is upper semi-continuous, hence continuous, i.e.,  $V \in B$ : we have our contracting operator; the rest of the proof is like in prop. 1.4, using this time that the optimal strategy correspondence is upper semi-continuous. ■

**COMMENT 1.2.** In some sense prop. 1.5 is much better than prop. 1.4: at least, if one were to strengthen the separate continuity property in prop. 1.4 to a joint continuity property, one could immediately construct a separable metrisable topology on  $\Omega$  such that the assumptions of prop. 1.5 would also hold (with joint continuity). So prop. 1.5 “essentially” includes proposition 1.4; but it allows complete flexibility in the transitions — e.g.: next state is a continuous function of current state and actions —, while in proposition 1.4 one is for example constrained to a dominated set of probabilities when a player's action varies.

**1.b. Non-zero-sum case (Finite).** We assume here again that the basic spaces ( $S$  and  $\Omega$ ) are finite. Recall that a subgame perfect equilibrium is an  $\mathbf{I}$ -tuple  $\sigma$  such that after any history  $h$ ,  $h = (\omega_1, s_1, \dots, \omega_n, s_n)$ ,  $\sigma_h$  is an equilibrium in  $\Gamma_\lambda$ , where  $\sigma_h$  is defined on  $h' = (\omega'_1, s'_1, \dots, \omega'_m)$  by  $\sigma_h(h') = \sigma(\omega_1, s_1, \dots, \omega_n, s_n, \omega'_1, s'_1, \dots, \omega'_m)$ .

LEMMA 1.6. Assume that  $x = (x(\cdot | \omega))_{\omega \in \Omega}$  with  $x(\cdot | \omega) \in X (= \prod_i \Delta(S^i))$  form, for each  $\omega$ , a Nash equilibrium of  $\Gamma((1 - \lambda)f)_\omega$  with pay-off  $f(\omega)$ ; then the corresponding stationary strategies  $\bar{x}$  define a subgame perfect equilibrium of  $\Gamma_\lambda$  with pay-off  $\lambda f$ .

PROOF. As in lemma 1.3 p. 328, for any  $\sigma$  and bounded  $h$ ,  $E_\sigma(\sum_{m \leq n} (1 - \lambda)^{m-1} g_m + (1 - \lambda)^n h(\omega_{n+1}))$  converges to  $\bar{\gamma}_\lambda(\sigma)/\lambda$ . By the definition of  $\bar{x}$  one has:

$$E_{\bar{x}}(g_m + (1 - \lambda)f(\omega_{m+1}) \mid \mathcal{H}_m) = f(\omega_m)$$

hence  $\bar{\gamma}_\lambda(\bar{x}) = \lambda f$ . Similarly after each history the future pay-off in  $\Gamma_\lambda$  if  $\bar{x}$  is used is  $\lambda f$ . Hence if  $\sigma'(h)$  say, is a profitable one-stage deviation at  $h$ , against  $\bar{x}$ , then the corresponding first component is a profitable deviation, against  $x$ , in the one-shot game with pay-off  $\lambda\phi((1 - \lambda)f)$ , a contradiction since  $x$  is a Nash equilibrium of  $\Gamma((1 - \lambda)f)$ . The result now follows from the fact that  $\sigma$  is a subgame perfect equilibrium in  $\Gamma_\lambda$  iff there is no profitable one-stage deviation, after any history (ex. IV.4Ex.4 p. 172). ■

PROPOSITION 1.7. *The discounted game  $\Gamma_\lambda$  has a subgame perfect equilibrium in stationary Markov strategies.*

PROOF. Define a correspondence  $\psi$  from  $X^\Omega \times [-C, C]^{\mathbf{I} \times \Omega} = Z$  to itself by:  $\psi(x, f) = \{(y, h) \in Z \mid \text{for each } i \text{ and each } \omega, y^i(\cdot | \omega) \text{ is a best reply against } x \text{ in } \Gamma((1 - \lambda)f)_\omega, \text{ yielding pay-off } h^i(\omega) \text{ to player } i\}$ .  $\psi$  is clearly u.s.c. and convex compact valued hence Kakutani's fixed point theorem (Kakutani, 1941) (e.g., proof of theorem 4.1 p. 39) leads to the existence of fixed points. Now lemma 1.6 applies. ■

COMMENT 1.3. The proof of prop. 1.7 extends clearly to more general setups, e.g. immediately to the case of compact action sets, with pay-offs and transitions depending continuously on the vector of actions.

**1.c. Non-zero-sum case (General).** In fact, one can similarly get the existence of subgame perfect equilibria (i.e. measurable strategies that form, for every initial state  $\omega$ , a subgame perfect equilibrium), even under assumptions hardly stronger than in prop. 1.4:

- The restriction to a standard Borel space  $(\Omega, \mathcal{A})$  is superfluous; an arbitrary measurable space will do.
- One can also allow for compact action sets  $S^i(\omega)$  that vary measurably with  $\omega$ , in the following sense:  $S^i(\omega) \subseteq \overline{S}^i$ ,  $(\overline{S}^i, \mathcal{S}^i)$  is a separable and separating measurable space, each subset  $S^i(\omega)$  is endowed with some compact topology, the  $\sigma$ -field  $\mathcal{S}^i$  is generated by the real valued measurable functions that have a continuous restriction to each set  $S^i(\omega)$ , and  $\{\omega \mid S^i(\omega) \cap O \neq \emptyset\}$  is measurable for each  $O \in \mathcal{S}^i$  whose trace on each set  $S^i(\omega)$  is open.
- The uniformly bounded pay-off functions  $g^i(\omega, s)$  and the transition probability  $P(A \mid \omega, s)$  are measurable (for each  $A \in \mathcal{A}$ ) on the graph of  $S(\omega) = \prod_i S^i(\omega)$ ; and, for each  $\omega \in \Omega$ ,  $g^i(\omega, s)$  and  $P(\cdot \mid \omega, s)$  are continuous functions on  $S(\omega)$  — in the norm topology for  $P$ , i.e.  $s_n \rightarrow s$  implies  $\sup_A |P(A \mid \omega, s_n) - P(A \mid \omega, s)|$  converges to zero.

The basic idea of the proof is somewhat reminiscent of what we did in the zero-sum case, i.e., start with a “large”, compact valued measurable correspondence from state space to pay-off space,  $K_0$  — e.g. the set of all feasible pay-offs. Given a measurable map to compact valued subsets  $K$ , define  $[\Psi(K)]_\omega$  as the set of all Nash equilibrium pay-offs for the uniform closure of all games  $\Gamma((1-\lambda)f)_\omega$ , letting  $f$  vary through all measurable selections from  $K$ . Prove that  $\Psi(K)$  is measurable map to compact subsets. Get in this way inductively a decreasing sequence of measurable maps to compact subsets  $K_n = \Psi^n(K_0)$ , with  $K = \bigcap_n K_n$ :  $K$  is then also measurable; further  $K_{n+1} = \Psi(K_n)$  goes to the limit and yields  $K = \Psi(K)$ . Observe that, at each point  $s \in S(\omega)$ , the set of pay-offs  $\Gamma((1-\lambda)f)_{\omega,s}$  is already closed, when  $f$  varies through all measurable selections from  $K$ . Thus one can choose, first for each  $(\omega, p) \in K$ , a continuous pay-off function  $\gamma_{\omega,p}(s)$  on  $S(\omega)$  which is a uniform limit of functions  $\Gamma[(1-\lambda)f]_{\omega,s}$  ( $f$  measurable selection from  $K$ ), and a Nash equilibrium  $\sigma_{\omega,p}$  of  $\gamma_{\omega,p}$  with pay-off  $p$ ; next for each  $s \in S(\omega)$ , a measurable selection  $f_{\omega,p,s}$  from  $K$  with  $\Gamma[(1-\lambda)f_{\omega,p,s}]_{\omega,s} = \gamma_{\omega,p}(s)$ . Doing all this in a measurable way yields a strategy: if next state is  $\tilde{\omega}$ , just repeat the same thing with  $f_{\omega,p,s}(\tilde{\omega})$  instead of  $p$ .

One can see the close analogy with the previous method — only in the zero-sum case the contracting aspect — i.e., the minmax theorem — insured that the correspondences  $K_n$  would decrease at a rate  $(1-\lambda)$ , hence converge to a single point.

The proof is however technically more heavy; so we refer the reader to (Mertens and Parthasarathy, 1987) for it. There he will also find how the above assumptions can be further relaxed — e.g. the functions  $g^i$  do not need to be uniformly bounded, and the discount factor can be allowed to depend a.o. on the player, on the stage, and on the past sequence of states.

In fact, a much simpler proof (one page) is possible under the following additional assumptions (Mertens and Parthasarathy, 1991):

- the state space  $(\Omega, \mathcal{A})$  is separable;
- the action sets  $S^i(\omega)$  are finite and independent of  $\omega$ ;
- the transition probabilities are dominated by a single measure  $\mu$  on  $(\Omega, \mathcal{A})$ ;
- the pay-off function is bounded, and a fixed discount rate is used.

In the general case, the strategies obtained are neither Markov nor stationary — they only have the very weak stationarity property that strategies are stationary functions of the current state and the currently expected pay-off vector for the future (the “current expectations”). In the particular case above, one obtains somewhat closer to Markov: the behavioural strategies can be chosen such as to be a function only of the current and the previous state. And if in addition the transition probability is nonatomic, one can further obtain (cf. Mertens and Parthasarathy, 1991) stationarity: the function is the same at every period. We give the proof, since it is so simple and contains in germ already several ideas of the general case:

**THEOREM 1.8.** *Under the above assumptions, there exists a subgame perfect stationary equilibrium.*

**PROOF.** Let  $|g^i(\omega, s)| \leq c$ , and  $F_0 = \{f \in [L_\infty(\mu)]^\mathbf{I} \mid \|f_i\|_\infty \leq c \ \forall i \in \mathbf{I}\}$ . For  $f \in F_0$ , let  $N_f(\omega) = \{\text{Nash equilibrium pay-offs of } G_{f,\omega} = \Gamma((1-\lambda)f)_\omega\}$ .  $\emptyset \neq N_f(\omega) \subseteq [-c/\lambda, c/\lambda]^\mathbf{I}$ . Denote by  $\mathcal{N}_f$  the set of all  $\mu$ -measurable selections from the convex hull of  $N_f(\omega)$ . Note that  $\mathcal{N}_f \neq \emptyset$  using a selection theorem (App.7 p. 427). Observe also the correspondence  $f \mapsto \mathcal{N}_f$  from  $F_0$  to itself is convex valued, and weak\*-upper semi-continuous: if  $f_n \xrightarrow{w^*} f$ , then  $G_{f_n,\omega} \rightarrow G_{f,\omega}$  point-wise, so  $\limsup N_{f_n}(\omega) \subseteq N_f(\omega)$ . Thus if  $\varphi_n \in \mathcal{N}_{f_n}$

converges weak<sup>\*</sup> to  $\varphi$ , then  $\varphi$  is the a.e. limit of a sequence of convex combinations of the  $\varphi_n$ , hence  $\varphi \in \mathcal{N}_f$ . It follows then from Fan's fixed point theorem (Fan, 1952) that  $\mathcal{N}$  has a fixed point:  $f_0 \in \mathcal{N}_{f_0}$ . I.e. (Lyapunov)  $\int f_0(\tilde{\omega})P(d\tilde{\omega} | \omega, s)$  is a measurable selection from the graph of  $(\omega, s) \mapsto \int N_{f_0}(\tilde{\omega})P(d\tilde{\omega} | \omega, s)$ . And  $\omega \mapsto N_{f_0}(\omega)$  is a measurable map to compact subsets of  $\mathbb{R}^I$  as the composition of the measurable map  $\omega \mapsto G_{f_0, \omega}$  with the equilibrium correspondence, which is by upper semi-continuity a Borel map from games to compact sets. By the measurable choice theorem in (Mertens, 1987a), it follows thus that there exists a measurable selection  $\psi(\omega, \tilde{\omega}) \in N_{f_0}(\tilde{\omega})$  such that  $G_{f_0(\cdot), \omega} = G_{\psi(\omega, \cdot), \omega}$ . Denote by  $\sigma(p, G)$  a Borel selection of an equilibrium with pay-off  $p$  of the game  $G$ . The stationary equilibrium is now to play  $\sigma[\psi(\omega, \tilde{\omega}), G_{f_0, \tilde{\omega}}]$  at state  $\tilde{\omega}$ , denoting by  $\omega$  the previous state. ■

COMMENT 1.4. Observe that all the trouble w.r.t. the Markovian character of the strategies stems from the nonatomic part of the transitions: under the same assumptions as in the general case above, if one assumes the transition probabilities to be purely atomic, one obtains immediately the existence of subgame perfect equilibria in stationary Markov strategies — e.g. by going to the limit with the result of the remark after prop. 1.7, following an ultrafilter on the increasing net of all finite subsets of  $\Omega$ . [To truncate the game to a finite subset, add e.g. an absorbing state with pay-off zero, which replaces the complement of this finite subset. The argument assumes  $\mathcal{A}$  is the class of all subsets; if  $\mathcal{A}$  was separable, one can always replace it by the class of all subsets on the corresponding quotient space while preserving all assumptions; and it is shown in (Mertens and Partha Sarathy, 1987) how to reduce the general problem to the case where  $\mathcal{A}$  is separable. Measurability of the strategies is anyway almost immaterial here, since the assumptions imply that, for any initial state, only countably many states are reachable.]

COMMENT 1.5. The above remark is well illustrated by the previous proof: one has to convexify the set of Nash equilibrium pay-offs, because a weak<sup>\*</sup>-limit belongs pointwise only to the convex hull of the point-wise limits (it is not because of the fixed point argument, which is not used in the other proofs). So  $f_0(\omega)$  is only a convex combination of equilibrium pay-offs of  $G_{f_0, \omega}$  — and one will play equilibria. So one must select the equilibria, as a function of tomorrow's state  $\tilde{\omega}$ , such as to give today the same game  $G_{f_0, \omega}$ . This uses basically a measurable version of Lyapunov's theorem, because (finiteness of action sets), at  $\omega$  only finitely many measures on  $\tilde{\omega}$  have to be considered. But it is clear that the solution to such a problem depends on the vector measure, i.e. on  $\omega$ : the equilibrium played tomorrow at  $\tilde{\omega}$  will depend on  $\omega$ .

COMMENT 1.6. To illustrate this in still another way, under the same assumptions, there exist stationary Markov sunspot equilibria (extensive form correlated equilibria (cf. 3.b p. 90) with public signals): if one convexifies the set of Nash equilibria, there is no problem, cf. ex. VIIEx.11 p. 348.

1.c.0.i. The assumption of norm-continuity of the transitions as a function of the actions is quite strong. Typically, one needs some form of noise in the model to assure it. The best behaved model where it is not satisfied, and where existence of equilibria is not known, is the following. Take as action sets for each player the one point compactification of the integers, and take the Cantor set as state space. Assume the reward function, and the probability for each Borel subset of the state space, are jointly continuous in state and actions, and use standard discounting. (And assume even further that the transitions are nonatomic to fix ideas — they are then all dominated by a fixed non-atomic probability).

## 2. Asymptotic analysis, finite case: the algebraic aspect

As seen above, we find the stationary Markov equilibria looking for fixed points  $V$  of the operator  $\Psi$  of the previous paragraph, i.e., solutions of  $V = \Psi((1 - \lambda)V)$  (lemma 1.2 p. 328) or  $f$  Nash equilibrium pay-off in  $\Gamma((1 - \lambda)f)$  (lemma 1.6 p. 330).

When state and action sets are finite, this becomes a system of finitely many polynomial equations and inequalities in finitely many variables as defined by the correspondence  $\psi$  (prop. 1.7 p. 330). The system is also polynomial (affine) in  $\lambda$ .

We have thus shown:

**PROPOSITION 2.1.** *The set  $E = \{(\lambda; g_1, \sigma^1; g_2, \sigma^2; \dots) \mid 0 < \lambda \leq 1\}$ , the  $\sigma^i$  form a stationary Markov equilibrium with pay-off vector  $g^i$  for the  $\lambda$ -discounted game} is, for each fixed  $\lambda$ , compact and non-empty and is semi-algebraic (i.e., the set of solutions of a system of polynomial equations and inequalities).*

Moreover one has:

**LEMMA 2.2.** *Any set as in prop. 2.1 contains a subset with the same properties consisting of a singleton, for each fixed  $\lambda$  (i.e., the equilibrium is a semi-algebraic function of  $\lambda$ ).*

**PROOF.** Denote by  $C$  the given set as a subset of some Euclidian space  $\mathbb{R}^k$ . Let  $C_0 = \{(\lambda, x) \in C \mid d(x, 0) \text{ is minimal}\}$ , and  $C_\ell = \{(\lambda, x) \in C_{\ell-1} \mid d(x, e_\ell) \text{ is minimal}\}$  for  $1 \leq \ell \leq k$ , where the  $e_\ell$  are the basis vectors of  $\mathbb{R}^k$ , and  $d$  is the Euclidian distance. If both  $(\lambda, x)$  and  $(\lambda, y)$  belong to  $C_k$ ,  $x$  and  $y$  have the same distance to zero and to all basis vectors, hence  $x = y$ . Since clearly  $C_\ell$  is non-empty and compact for each  $\lambda$  (induction),  $C_k$  is indeed the graph of a function. Also the semi-algebraicity of  $C_\ell$  follows by induction — using 3.4 p. 27 ■

**THEOREM 2.3.** *For any (finite) stochastic game, there exist  $\lambda_0 > 0$ , a positive integer  $M$ , and Puiseux series expansions ( $g$  denotes the normalised pay-off):*

$$g_\omega^i(\lambda) = \sum_{k \geq 0} h_k^{i,\omega} \lambda^{k/M}$$

$$\text{and } \sigma_\omega^i(s)(\lambda) = \sum_{k \geq 0} \alpha_k^{i,\omega,s} \lambda^{k/M}, \quad \forall i \in \mathbf{I}, \forall s \in S_i,$$

such that, for all  $\lambda \in ]0, \lambda_0]$ , the  $\sigma^i(\lambda)$  form a stationary Markov equilibrium, with pay-off vector  $g(\lambda)$ , of the  $\lambda$ -discounted game. (And those functions  $\sigma^i(\lambda), g(\lambda)$  are semi-algebraic on  $[0, \lambda_0]$ .)

**PROOF.** Apply lemma 2.2 to select a semi-algebraic function  $(g(\lambda), \sigma(\lambda))$  from the graph  $E$  of stationary Markov equilibria. Each coordinate of this function is then a real valued semi-algebraic function on  $]0, 1]$  (projection): such functions  $f(\lambda)$  have a Puiseux-series expansion (Ahlfors, 1953) in the neighbourhood of zero — i.e.,  $f(\lambda) = \sum_{k \geq k_0} \varphi_k \lambda^{k/M}$  for some  $k_0 \in \mathbb{Z}$ ,  $M \in \mathbb{N} \setminus \{0\}$ , such that the series is absolutely convergent to  $f(\lambda)$  on some interval  $]0, \lambda_0]$  ( $\lambda_0 > 0$ ). Now each of our coordinates  $f(\lambda)$  is bounded, (remember we use normalised pay-offs), hence one has  $k_0 \geq 0$  — thus we can set  $\varphi_k = 0$  for  $0 \leq k < k_0$ , and use  $k_0 = 0$ .

Replace now all different  $M$ 's by their least common multiple over the different coordinates, and replace the different  $\lambda_0$ 's by their minimum, to obtain the result. ■

**COROLLARY 2.4.** *For such a solution and any  $\lambda_1 < \lambda_0$ ,  $\|(dg/d\lambda, d\sigma/d\lambda)\| \leq A\lambda^{-(M-1)/M}$  for some  $A > 0$  and any  $\lambda \leq \lambda_1$ .*

PROOF. Such an absolutely convergent series can be differentiated term by term in the interior of its radius of convergence. ■

COMMENT 2.1. Every value of  $M$  in the above results can indeed occur: cf. ex. VIIEx.3.

### 3. $\varepsilon$ -optimal strategies in the undiscounted game

**3.a. The Theorem.** We consider here two-person zero-sum games. No finiteness conditions or other are imposed, but we assume the pay-offs to be uniformly bounded, and the values  $v_\lambda(\omega)$  of the discounted games to exist. We will exhibit sufficient conditions for the existence of strategies that guarantee in a strong sense (cf. 1.c p. 149) some function  $v_\infty(\omega)$  (up to  $\varepsilon$ ). The theorem does not require the moves to be announced, only the pay-offs.

THEOREM 3.1. Assume a stochastic game where:

- (1) pay-offs are uniformly bounded;
- (2) the values  $v_\lambda(\omega)$  of the  $\lambda$ -discounted games exist, as well as  $\varepsilon$ -optimal strategies in the sense of 3 p. 328;
- (3)  $\forall \alpha < 1$  there exists a sequence  $\lambda_i$  ( $0 < \lambda_i \leq 1$ ) such that, denoting by  $\|\cdot\|$  the supremum norm over the state space:  $\lambda_{i+1} \geq \alpha \lambda_i$ ,  $\lim_{i \rightarrow \infty} \lambda_i = 0$  and  $\sum_i \|v_{\lambda_i} - v_{\lambda_{i+1}}\| < +\infty$ .

Then the game has a value  $v_\infty$ . More precisely,  $\forall \varepsilon > 0$ ,  $\exists \sigma_\varepsilon$ ,  $\exists N_0$  such that:

$$\forall n \geq N_0, \forall \omega \in \Omega, \forall \tau, E_{\sigma_\varepsilon \tau}^\omega(\bar{g}_n) \geq v_\infty(\omega) - \varepsilon \quad \text{and} \quad E_{\sigma_\varepsilon \tau}^\omega \liminf_{n \rightarrow \infty} (\bar{g}_n) \geq v_\infty(\omega) - \varepsilon$$

and dually for player II.

PROOF. Denote by  $A$  four times the largest absolute value of the pay-offs, and let w.l.o.g.  $0 < \varepsilon \leq A$ ,  $\delta = \varepsilon/(12A)$ .

Take two functions  $L(s)$  and  $\lambda(s)$  of a real variable  $s$ , such that  $0 < \lambda(s) \leq 1$  and  $L(s) > 0$ . Assume  $\exists M > 0$ , such that, for  $s \geq M$ ,  $|\theta| \leq A$  and every  $\omega$

- (1)  $|v(\omega, \lambda(s)) - v_\infty(\omega)| \leq \delta A$
- (2)  $AL(s) \leq \delta s$
- (3)  $|\lambda(s + \theta L(s)) - \lambda(s)| \leq \delta \lambda(s)$
- (4)  $|v(\omega, \lambda(s + \theta L(s))) - v(\omega, \lambda(s))| \leq \delta AL(s) \lambda(s)$
- (5)  $\int_M^\infty \lambda(s) ds \leq \delta A$
- (6)  $\lambda$  is strictly decreasing and  $L$  is integer valued.

Call the above set of conditions  $H(L, \lambda, A, \delta)$ .

**3.b. Proof of the theorem under  $H(L, \lambda, A, \delta)$ .** Define now inductively, using  $g_i$  and  $\omega_i$  for pay-off and state at stage  $i$  ( $i = 1, 2, \dots$ ), starting with  $s_0 \geq M$ :

$$\begin{aligned} \lambda_k &= \lambda(s_k), L_k = L(s_k), B_0 = 1, B_{k+1} = B_k + L_k \\ s_{k+1} &= \max[M, s_k + \sum_{B_k \leq i < B_{k+1}} (g_i - v_\infty(\omega_{B_{k+1}}) + \varepsilon/2)]. \end{aligned}$$

Observe that

$$(7) \quad |s_{k+1} - s_k| \leq AL_k$$

and that, by (5) and (6),  $\lim_{s \rightarrow \infty} s\lambda(s) = 0$ , hence by (2)  $\lim_{s \rightarrow \infty} \lambda(s)L(s) = 0$ , hence

$$(8) \quad \lambda(s)L(s) \leq \delta \quad \text{for } s \geq M \quad \text{by choosing } M \text{ sufficiently large.}$$

Let also  $\ell_k = v(\omega_{B_k}, \lambda_k)$ ,  $\tilde{\ell}_k = v(\omega_{B_{k+1}}, \lambda_k)$ , and note that by (4):

$$(9) \quad |\tilde{\ell}_k - \ell_{k+1}| \leq \delta AL_k \lambda_k.$$

Player I's strategy  $\sigma$  consists in playing for  $B_k \leq i < B_{k+1}$  a  $(\delta AL_k \lambda_k)$ -optimal strategy in the  $\lambda_k$ -discounted game. The following computations are for an arbitrary strategy  $\tau$  of player II and  $\mathbb{E}$  stands for  $\mathbb{E}_{\sigma, \tau}$ .

Hence, denoting by  $\mathcal{H}_k$  the  $\sigma$ -field of all past events at stage  $B_k$ , up to the choice of  $\omega_{B_k}$  included,

$$\ell_k \leq \mathbb{E} \left[ \lambda_k \sum_{0 \leq i < L_k} (1 - \lambda_k)^i g_{B_k+i} + (1 - \lambda_k)^{L_k} \tilde{\ell}_k \mid \mathcal{H}_k \right] + \delta AL_k \lambda_k$$

or, as  $1 - \lambda \sum_{i < L} (1 - \lambda)^i = (1 - \lambda)^L$ ,

$$\mathbb{E} \left[ \tilde{\ell}_k - \ell_k + \lambda_k \sum_{i < L_k} (1 - \lambda_k)^i (g_{B_k+i} - \tilde{\ell}_k) \mid \mathcal{H}_k \right] \geq -\delta AL_k \lambda_k.$$

Using (9),  $1 - \lambda L \leq (1 - \lambda)^i \leq 1$  for  $i < L$ , (8) and (1) we get

$$\mathbb{E} \left[ \ell_{k+1} - \ell_k + \lambda_k \sum_{B_k \leq i < B_{k+1}} (g_i - v_\infty(\omega_{B_{k+1}})) \mid \mathcal{H}_k \right] \geq -4\delta AL_k \lambda_k$$

and hence, since  $s_{k+1} - s_k \geq \sum_{B_k \leq i < B_{k+1}} (g_i - v_\infty(\omega_{B_{k+1}}) + 6A\delta)$ ,

$$\mathbb{E} \left[ \ell_{k+1} - \ell_k + \lambda_k (s_{k+1} - s_k) \mid \mathcal{H}_k \right] \geq 2\delta AL_k \lambda_k$$

Now, using (3) and (7),

$$\lambda_k (s_{k+1} - s_k) - \delta AL_k \lambda_k \leq \int_{s_k}^{s_{k+1}} \lambda(s) ds = t_k - t_{k+1}$$

(letting  $t_k = \int_{s_k}^\infty \lambda(x) dx$  — note that by (5),  $t_k \leq \delta A$ ). Thus  $\mathbb{E}[Y_{k+1} \mid \mathcal{H}_k] \geq Y_k + \delta AL_k \lambda_k$ , for  $Y_k = \ell_k - t_k$  — note that  $|Y_k| \leq A/2$ . Thus  $A \geq \mathbb{E}(Y_k - Y_0) \geq \delta A \mathbb{E}(\sum_{\ell < k} L_\ell \lambda_\ell)$ , hence, by monotone convergence,  $\mathbb{E}(\sum_k L_k \lambda_k) \leq 1/\delta$ , so that

$$(10) \quad \mathbb{E} \left( \sum_k \mathbb{1}_{s_k=M} \right) \leq \frac{1}{\delta \lambda(M)}.$$

Also  $Y$  is a bounded submartingale, hence converges a.e. to  $Y_\infty$ , and by the stopping theorem  $\mathbb{E}(Y_T) \geq Y_0 = \ell_0 - t_0 \geq \ell_0 - \delta A$  for any stopping time  $T$  (including  $\infty$ ).

Let  $k(i)$  denote the stopping time  $\min\{k \mid B_k > i\}$ ,  $k(\infty) = \infty$ ; and let  $\bar{\ell}_i = v_\infty(\omega_{B_{k(i)}})$ . It follows then using that (by (1))  $\bar{\ell}_i \geq \ell_{k(i)} - \delta A \geq Y_{k(i)} - \delta A$ , that

$$(11) \quad \mathbb{E}(\bar{\ell}_i) \geq \bar{\ell}_0 - 3\delta A \quad \text{for all } i = 1, \dots, \infty \quad (\text{letting } \bar{\ell}_\infty = \liminf_{i \rightarrow \infty} \bar{\ell}_i).$$

The definition of  $s_k$  yields

$$s_{k+1} - s_k \leq \sum_{B_k \leq i < B_{k+1}} (g_i - \bar{\ell}_i) + 6\delta AL_k + \mathbb{1}_{s_{k+1}=M} AL_k / 2.$$

By (2) and (7), when  $s_{k+1} = M$ ,  $AL_k \leq \delta M / (1 - \delta)$ ; thus, by summing

$$s_k - s_0 \leq \sum_{i < B_k} (g_i - \bar{\ell}_i) + 6\delta AB_k + \delta M \sum_1^k \mathbb{1}_{s_\ell=M},$$

hence

$$\sum_1^n g_i \geq s_{k(n)} - s_0 + \sum_0^n \bar{\ell}_i - 6\delta An - A(B_{k(n)} - n) - \delta M \sum_0^\infty \mathbb{1}_{s_k=M}.$$

But

$$B_{k(n)} - n \leq L(s_{k(n)-1}) \leq A^{-1}\delta s_{k(n)-1} \leq \delta(A^{-1}s_0 + n)$$

(using a.o. (2)), hence

$$\sum_1^n g_i \geq \sum_1^n \bar{\ell}_i - 7A\delta n - (1 + \delta)s_0 - \delta M \sum_1^\infty \mathbb{1}_{s_k=M}.$$

It follows now from (10) and (11) that, for all  $n$ ,

$$\mathbb{E}(\bar{g}_n) \geq \bar{\ell}_0 - 10A\delta - \frac{K}{n}, \quad \text{with } K = 2s_0 + \frac{M}{\lambda(M)},$$

and hence

$$\mathbb{E}(\bar{g}_n) \geq v_\infty(\omega) - \varepsilon \quad \text{for } n = N_0, N_0 + 1, \dots, +\infty, \quad \text{with } N_0 \geq \frac{6K}{\varepsilon} :$$

the strategy described is  $\varepsilon$ -optimal. ■

**COMMENT 3.1.** Further, those  $\varepsilon$ -optimal strategies are still  $\varepsilon$ -optimal, in the same sense, in any subgame, and they consist in playing in successive blocs  $k$  of length  $L_k$  whichever  $(\delta AL_k\lambda_k)$ -optimal strategy in the modification  $\Gamma(\lambda_k, L_k)$  of the  $\lambda_k$  discounted game where from stage  $L_k$  on, the pay-off  $v_{\lambda_k}(\omega_{L_k})$  is received forever. (One sees immediately that only  $(\delta AL_k\lambda_k)$ -optimality in  $\Gamma(\lambda_k, L_k)$  was used in the proof under  $H(L, \lambda, A, \delta)$ ). The subgame property follows from the fact that in every subgame one uses — with the right proviso for the first bloc — a strategy  $(M, s_n)$  of the same  $\varepsilon$ -optimal type as  $(M, s_0)$  of the full game.

**COMMENT 3.2.** We can already solve the finite case. By cor. 2.4 p. 333,  $\|dv_\lambda/d\lambda\|$  is an integrable function (in the neighbourhood of 0, but point 1 of the proposition below implies it is bounded elsewhere), so the conditions  $H(L, \lambda, A, \delta)$  immediately yield that any functions  $L$  (integer valued) and  $\lambda$  (integrable) that satisfy  $\delta ds/d|\ln \lambda| \geq \|dv_\lambda/d\lambda\|$  (cf. (4)) and  $\delta ds/d|\ln \lambda| \geq AL(s)$  (cf. (3)) will typically work [indeed, for sufficiently smooth functions  $\lambda(s)$  (or  $s(\lambda)$ ), e.g. for power functions, the last condition — (2) — will typically follow from the others]. The cited corollary immediately yields a multitude of such functions. One such particularly simple choice is to take  $\lambda(s) = 1/[s \ln^2 s]$ ;  $L(s) = 1$  (or any choice s.t.  $L(s)/s \rightarrow 0$ ); this yields an extremely simple strategy working for all finite games — the only game-dependent parameter is the link between  $\varepsilon$  and  $M(\varepsilon)$ .

**3.c. End of the Proof.** The theorem will now follow from the next:

**PROPOSITION 3.2.** (1) Under assumptions (1) and (2) of the theorem,  $v_\lambda/\lambda$  is a Lipschitz function of  $1/\lambda$  (with the maximum absolute value of pay-offs  $C$  as Lipschitz constant); in particular  $v_\lambda$  is a Lipschitz function of  $\ln \lambda$ .

In 2, 3 and 4 below,  $v_\lambda$  is an arbitrary function with values in a metric space.

(2) If  $v_\lambda$  is a Lipschitz function of  $\ln(\lambda)$ , then assumption (3) of theorem 3.1 is equivalent to assumption (c\*):

$$(c^*) \quad \exists \lambda_i: \lim \lambda_i = 0, \lim \lambda_{i+1}/\lambda_i = 1 = \lambda_0, \sum_i d(v_{\lambda_{i+1}}, v_{\lambda_i}) < +\infty$$

(3) If  $v_\lambda$  is a Lipschitz function of  $\ln(\lambda)$ , then assumption (c\*) implies the existence of functions  $L$  and  $\lambda$  such that, for any  $A > 0$  and  $\delta > 0$ ,  $H(L, \lambda, A, \delta)$  holds.

- (4) Conversely, if for every  $\delta > 0$  there exists  $L$  and  $\lambda$  such that, for some  $A > 0$ ,  $H(L, \lambda, A, \delta)$  holds, (where conditions (1), (2) and (6) can be deleted), then  $v_\lambda$  satisfies assumption 3 of the theorem.

PROOF. 1)  $v_\lambda/\lambda$  is the value of the game with pay-off function  $\sum_0^\infty (1-\lambda)^i x_i$  ( $|x_i| < C$ ). Thus the pay-off functions corresponding to  $v_\lambda/\lambda$  and  $v_\mu/\mu$  differ by at most  $A|1/\lambda - 1/\mu|$ . Hence the conclusion. It is clear that this Lipschitz condition implies both the boundedness of  $v_\lambda$  and the Lipschitz character of  $v_\lambda$  as a function of  $\ln \lambda$ .

2) It suffices to show that (3)  $\Rightarrow$  (c\*). Add first finitely many values  $\alpha^k$  ( $k = 0, 1, 2, \dots, k_0$ ) in the beginning of the sequence  $\lambda_i$  (with  $k_0 = \max\{j \mid \alpha^j > \lambda_0\}$ ): we have now in addition  $\lambda_0 = 1$ . Let now  $i_0 = 0$ ,  $i_{k+1} = \min\{i \mid \lambda_i < \lambda_{i_k}\}$ ,  $\tilde{\lambda}_k = \lambda_{i_k}$ :  $\tilde{\lambda}_{k+1} \geq \alpha \tilde{\lambda}_k$  is still true, and by the triangle inequality  $\sum \|v_{\tilde{\lambda}_{k+1}} - v_{\tilde{\lambda}_k}\| < \infty$ : we now have in addition that the sequence  $\lambda_i$  is strictly decreasing. Let thus, for  $n \geq 1$ ,  $(x_i^n)_{i=0}^\infty$  satisfy  $x_0^n = 1$ ,  $0 < \ln(x_i^n/x_{i+n}^n) \leq 2^{-n}$ ,  $x_i^n \rightarrow 0$  and  $\sum_i d(v(x_{i+1}^n), v(x_i^n)) < +\infty$ . Choose then  $(a_n)_{n=1}^\infty$  such that  $a_1 = 1$ ,  $\ln(a_N/a_{n+1}) > 1$ , and  $\sum_{n,i} \mathbb{1}_{x_i^n \leq a_n} d(v(x_{i+1}^n), v(x_i^n)) < +\infty$ . Let  $\lambda_i$  be the enumeration in decreasing order of  $\{x_i^n \mid a_{n+1} < x_i^n \leq a_n, n \geq 1, i \geq 0\}$ : obviously  $\lambda_0 = 1$ ,  $\lambda_i \rightarrow 0$ . To verify the other conditions on  $\lambda_i$ , let  $\bar{a}_n = \min\{x_i^{n-1} \mid x_i^{n-1} > a_n\}$ ,  $\underline{a}_n = \max\{x_i^n \mid x_i^n \leq a_n\}$ . As  $\ln(\bar{a}_n/\underline{a}_n) = \ln(\bar{a}_n/a_n) + \ln(a_n/\underline{a}_n) \leq 2^{-n+1} + 2^{-n}$ ,  $\lambda_{i+1}/\lambda_i \rightarrow 1$ , and also  $d(v(\bar{a}_n), v(\underline{a}_n)) \leq K(2^{-n+1} + 2^{-n})$  (Lipschitz property) which is summable; thus

$$\sum_i d(v(\lambda_{i+1}), v(\lambda_i)) \leq \sum_{n=1}^\infty \sum_{x_i^n \leq a_n} d(v(x_{i+1}^n), v(x_i^n)) + \sum_{n=1}^\infty d(v(\bar{a}_n), v(\underline{a}_n)) < +\infty.$$

3) We will do a more general construction but of the same type as the one suggested in remark 3.5 below. To see the ideas of the following proof in a more transparent framework it may be useful to check this case first.

Let  $f(\lambda) = \sum_{\lambda_i \geq \lambda} d(v(\lambda_i), v(\lambda_{i-1}))$ ,  $\Delta[a, b] = f(a) - f(b)$ ,  $\ell_i = \Delta[y_n^{i+1}, y_n^i]$  where  $y_n = n/(n+1)$ ,  $n \geq 1$ . Then  $\sum_{-\infty}^{+\infty} \ell_i \leq \sum_k d(v(\lambda_k), v(\lambda_{k-1})) < +\infty$ . Let also  $\bar{\ell}_i = \sum_{|j| \leq 2} \ell_{i+j}$ ,  $g_n(x) = 2n[\mathbb{1}_{x \leq 1} + \bar{\ell}_i y_n^{-i}]$  for  $y_n^{i+1} < x \leq y_n^i$ , and define  $h_n$  by linear interpolation from the values  $h_n(y_n^i) = n \sum_{j < i} [g_n(y_n^{j+1}) + g_n(y_n^j) + g_n(y_n^{j-1})]$ . Then  $g_n \geq n$  on  $[0, 1]$  and  $h_n$  is continuous, decreasing,  $\geq ng_n$  and integrable [e.g.  $\int_0^\infty h_n(x) dx \leq \sum_i h_n(y_n^i)(y_n^{i-1} - y_n^i) = n(y_n^{-1} - 1) \sum_{j < i} y_n^j [g_n(y_n^{j+1}) + g_n(y_n^j) + g_n(y_n^{j-1})] = n \sum_j y_n^j [g_n(y_n^{j+1}) + g_n(y_n^j) + g_n(y_n^{j-1})] \leq 3ny_n^{-1} \sum_j y_n^j g_n(y_n^j) = 6n(n+1)[1/(1-y_n) + \sum_i \bar{\ell}_i] < +\infty$ ]. Further, for  $y_n^{i+1} < x \leq y_n^i$ ,  $\Delta[xy_n^2, xy_n^{-2}] \leq \Delta[y_n^{i+3}, y_n^{i-2}] \leq \bar{\ell}_i \leq y_n^i g_n(x)/(2n) \leq xg_n(x)/n$ . Moreover, by the linearity on those intervals of  $g_n(x)$  and  $h_n(xy_n^k)$ ,  $h_n(y_n^{i+1}) = h_n(y_n^i) + n[g_n(y_n^{i+1}) + g_n(y_n^i) + g_n(y_n^{i-1})]$  implies that  $h_n(x) + ng_n(x) \leq h_n(xy_n)$  and  $h_n(x) - ng_n(x) \geq h_n(xy_n^{-1})$ . Choose  $(a_n)_{n=1}^\infty$  such that  $\frac{1}{2} \geq a_n > a_{n+1} > 0$ ,  $\inf\{\lambda_i/\lambda_{i-1} \mid \lambda_i \leq 2a_n\} \geq y_n$  and  $\sum_n \int_0^{a_n} h_n(x) dx < +\infty$ . Set  $\bar{h}_n(x) = h_n(x)(2-x/a_n)^+$ . Then  $\bar{h}(x) = \sum_1^\infty \bar{h}_n(x)$  is continuous, decreasing (strictly on  $[0, 2a_1]$ ) and integrable. Hence so is  $\bar{h}^{-1}(s)$  ( $s > 0$ ). Let  $n(x) = \#\{n \mid a_n \geq x\}$ ,  $g_0(x) = 1$ ,  $\bar{g}(x) = g_{n(x)}(x)$ ,  $y_0 = 1/2$ ;  $x_i = xy_{n(x)}^i$ : then  $\bar{g}(x) \geq n(x) \xrightarrow{x \rightarrow 0} \infty$ ,  $\Delta[x_2, x_{-2}] \leq x\bar{g}(x)/n(x)$ , and  $\bar{h}(x) - \bar{h}(x_{-1}) \geq \bar{h}_{n(x)}(x) - \bar{h}_{n(x)}(x_{-1}) \geq h_{n(x)}(x) - h_{n(x)}(xy_{n(x)}^{-1}) \geq n(x)g_{n(x)}(x) = n(x)\bar{g}(x)$  — thus  $\bar{h}(x)/\bar{g}(x) \geq n(x) \xrightarrow{x \rightarrow 0} \infty$ . Similarly  $\bar{h}(x_{+1}) - \bar{h}(x) \geq n(x)\bar{g}(x)$ . Thus  $|\bar{h}(z) - \bar{h}(x)| \leq n(x)\bar{g}(x)$  implies  $x_1 \leq z \leq x_{-1}$ ; so if further  $\lambda_{i+1} \leq z \leq \lambda_i$  then  $x_2 \leq \lambda_{i+1} < \lambda_i \leq x_{-2}$  — because  $\lambda_{i+1} \leq xy_{n(x)}^{-1} \leq a_{n(x)}y_{n(x)}^{-1} \leq 2a_{n(x)}$ , so that  $\lambda_{i+1}/\lambda_i \geq y_{n(x)}$  (assume  $n(x) \geq 1$ ). Therefore if  $\lambda(s)$  is the closest  $\lambda_i$  to  $x = \bar{h}^{-1}(s)$ , and  $L(s) = \bar{g}(x)$ ,

$N(s) = n(x)$ , then  $\lambda(s + \theta L(s)) \in [x_2, x_{-2}]$  for  $|\theta| \leq N(s)$ , hence

$$d[v(\lambda(s + \theta L(s)), v(\lambda(s))] \leq \Delta[x_2, x_{-2}] \leq 2\lambda(s)L(s)/N(s),$$

and

$$\left| \ln \frac{\lambda(s + \theta L(s))}{\lambda(s)} \right| \leq \ln y_{N(s)}^{-4} \leq \frac{4}{N(s)}.$$

Further

$$\frac{L(s)}{s} \leq \frac{\bar{g}(x)}{\bar{h}(x)} \leq \frac{1}{N(s)},$$

and  $s \rightarrow \infty \Rightarrow x \rightarrow 0 \Rightarrow N(s) = n(x) \rightarrow \infty$ . So the above inequalities will give (4), (3) and (2) respectively. Clearly  $\lambda$  is decreasing and integrable, and  $0 < \lambda \leq 1$ , hence (5). Perturb  $\lambda$  slightly such as to make it strictly decreasing (continuity of  $v_\lambda$ ), and replace  $L$  by its integer part (recall  $L(s) \geq N(s) \rightarrow \infty$ ) hence (6). Condition (c\*) implies that the sequence  $v_{\lambda_i}$  is a Cauchy sequence, hence convergent, say to  $v_\infty$ . By the choice of  $\lambda$  this implies (1):  $H(L, \lambda, A, \delta)$  obtains for all  $A > 0, \delta > 0$ .

4) Let  $s_0 = M$ ,  $L_i = AL(s_i)$ ,  $s_{i+1} = s_i + L_i$ ,  $\lambda = \lambda(s_i)$ . Then by (4)  $\sum_i \|v(\lambda_{i+1}) - v(\lambda_i)\| \leq \delta \sum_i \lambda_i L_i = \delta \sum_i \lambda_i (s_{i+1} - s_i) \leq \delta(1 - \delta)^{-1} \sum_i \int_{s_i}^{s_{i+1}} \lambda(s) ds < +\infty$  (by (3) and (5)). Hence  $\lambda_i \rightarrow 0$  ( $L_i \gg 0$ ), and (by (3))  $\lambda_{i+1} \geq (1 - \delta)\lambda_i$ . ■

COMMENT 3.3. Thus one can always use the same functions  $\lambda$  and  $L$  for any given game i.e., independently of  $\varepsilon$ ; taking also  $s_0 = M$  yields then a family of  $\varepsilon$ -optimal strategies where only the parameter  $M$  changes with  $\varepsilon$ . (The arbitrary  $s_0$  was just needed for the subgame property).

COMMENT 3.4. Condition (3) of the theorem is always satisfied when  $v_\lambda$  is of bounded variation or when (for some function  $v_\infty$ )  $\|v_\lambda - v_\infty\|/\lambda$  is integrable. Indeed, this means the integrability of  $\|v_\lambda - v_\infty\|$  as a function of  $\ln \lambda$ ; let then  $\lambda_i$  denote the minimiser of  $\|v_\lambda - v_\infty\|$  in  $[\beta^{i+1}, \beta^i]$  to satisfy condition (3).

COMMENT 3.5. When  $v_\lambda$  is of bounded variation, a much simpler construction of  $L$  and  $\lambda$  is possible than in the proposition: just set  $L = 1$ ,  $\varphi(x) = \int_0^x \|dv_\lambda\|$  (i.e., the total variation of  $v_\lambda$  between 0 and  $x$ ),  $s(\lambda) = \int_\lambda^1 x^{-1} d\sqrt{\varphi(x)} + 1/\sqrt{\lambda}$ , and  $\lambda(s)$  the inverse function. (cf. ex. VIIEx.14 p. 350).

COMMENT 3.6. Under the assumptions of the theorem  $v_\lambda$  converges uniformly to  $v_\infty$ : recall that (c\*) implies the uniform convergence of  $v_{\lambda_i}$  to  $v_\infty$ , and the Lipschitz character of  $v_\lambda$  as function of  $\ln \lambda$  implies that, for  $\lambda_{i+1} \leq \lambda \leq \lambda_i$ ,  $d(v_\lambda, v_{\lambda_i}) \leq K \ln(\lambda_i/\lambda_{i+1}) \rightarrow 0$ . Thus the conclusion, by point 2 of the proposition. In fact the statement itself of the theorem, with  $N_0$  independent of the initial state, implies immediately the uniform convergence of  $v_n$  and of  $v_\lambda$ .

COMMENT 3.7. The function  $\lambda(s)$  constructed in prop. 3.2 p. 336 was strictly decreasing — thus with a continuous inverse  $s(\lambda)$  —, but cannot be made continuous. Indeed the same proof shows that

PROPOSITION 3.3. *The following conditions on  $v_\lambda$  are equivalent:*

- (1) *For some  $A > 0$  and  $\delta$  ( $0 < \delta < 1$ ) there exist  $L$  and  $\lambda$  that satisfy  $H(L, \lambda, A, \delta)$  [without (1), (2) and (6)], where  $\lambda$  is continuous.*
- (2) *There exist  $L$  and  $\lambda$  (continuous) that satisfy  $H(L, \lambda, A, \delta)$  for any  $A > 0$  and  $\delta > 0$ .*
- (3)  $\exists(\lambda_i)_{i=0}^\infty : \inf_i \lambda_{i+1}/\lambda_i > 0 = \lim_i \lambda_i$ ,  $\sum_i \Delta[\lambda_{i+1}, \lambda_i] < +\infty$ .

$$(4) \forall (\lambda_i)_{i=0}^{\infty} : \sup_i \lambda_{i+1}/\lambda_i < 1, \lambda_i > 0 \Rightarrow \sum_i \Delta[\lambda_{i+1}, \lambda_i] < +\infty.$$

Here  $\Delta[\lambda, \mu] = \max_{\lambda \leq \lambda_1 \leq \lambda_2 \leq \mu} d(v_{\lambda_1}, v_{\lambda_2})$ .

PROOF. Show, as in prop. 3.2, that  $1 \Rightarrow 3 \Rightarrow 4 \Rightarrow 2$ . In  $3 \Rightarrow 4$ , just note that the sequence of 4 has only a bounded number of terms between two successive terms in the sequence of 3. In  $4 \Rightarrow 2$ , work directly with the present function  $\Delta$ , do not use the sequence  $\lambda_i$ , and set  $\lambda(s) = \bar{h}^{-1}(s)$ . ■

COMMENT 3.8. Condition 3 or 4 are not implied by the Lipschitz property and (c\*).

### 3.d. Particular cases (finite games, two-person-zero-sum).

3.d.1. When the stochastic game is a normalised form of a game with perfect information, then the games  $\Gamma_\omega((1 - \lambda)v)_\omega$  are also normal forms of games with perfect information, hence have pure strategy solutions, which form a closed, semi-algebraic subset of the set of all solutions: applying theorem 2.3 to those yields that, for some  $\lambda_0 > 0$ , there exists a pure strategy vector such that the corresponding stationary strategies are optimal in the  $\lambda$ -discounted game for all  $\lambda \leq \lambda_0$ . Such strategy vectors (pure or not) are also called “uniformly discount optimal” (Blackwell, 1962).

Observe that the perfect-information case in particular the situation includes where one player is a dummy — i.e., Markov decision processes or dynamic programming (cf. sect. 5 p. 353).

3.d.2. Whenever there exist uniformly discount optimal strategies, the expansion of  $v_\lambda$  is in integer powers of  $\lambda$ :  $v_\lambda$  is in fact a rational function of  $\lambda$ , being the solution of the linear system  $v_\lambda = \lambda g + (1 - \lambda)Pv_\lambda$ , where  $g$  and  $P$  are the single stage expected pay-off and the transition probability generated by the strategy pair.

3.d.3. Whenever there exists a strategy  $\sigma$  in the one-shot game which is  $o(\lambda)$ -optimal in  $\Gamma(\lambda, 1)$ , then for each  $\varepsilon > 0$ , one  $\varepsilon$ -optimal strategy of the theorem will consist in playing this all the time: the corresponding stationary strategy is optimal (in the strong sense of the theorem) in the infinite game. (Recall Remark 3.1 p. 336).

3.d.4. Since the value exists in such a strong sense, it follows in particular that the pay-off  $\underline{v}(\sigma)$  guaranteed by a stationary Markov strategy  $\sigma$  is also completely unambiguous — applying the theorem to the one-person case. Further, the preceding points imply the existence of a pure, stationary Markov best reply, which is best for all  $\lambda \leq \lambda_0$ , and that  $\underline{v}_\lambda(\sigma) \geq \underline{v}(\sigma) - K\lambda$ . One checks similarly that  $\underline{v}_n(\sigma) \geq \underline{v}(\sigma) - K/n$ . (cf. ex. VIIEx.6 p. 347).

It follows in particular that if both players have stationary Markov optimal strategies in  $\Gamma_\infty$ , then  $\|v_\lambda - v_\infty\| \leq K\lambda$ ,  $\|v_n - v_\infty\| \leq K/n$  (and those strategies guarantee such bounds) [and in particular  $\|v_\lambda - \underline{v}_\lambda(\sigma)\| \leq K'\lambda$ , so that the corresponding one-shot strategies are  $O(\lambda)$ -optimal in  $\Gamma(\lambda, 1)$ .]

3.d.5. It follows that 3.d.3 can be improved to: whenever there exists a stationary Markov strategy  $\sigma$  which is  $o(1)$  optimal in  $\Gamma(\lambda)$  (i.e.,  $\|\underline{v}_\lambda(\sigma) - v_\lambda\| \rightarrow 0$  as  $\lambda \rightarrow 0$ ), then  $\sigma$  is optimal in the infinite game. This is an improvement because an easy induction yields that the  $\varepsilon[1 - (1 - \lambda)^L]$ -optimality of  $\sigma$  in  $\Gamma(\lambda, L)$  implies its  $\varepsilon[1 - (1 - \lambda)^{KL}]$ -optimality in  $\Gamma(\lambda, KL)$ , hence its  $\varepsilon$ -optimality in  $\Gamma(\lambda)$ .

3.d.6. Apply the above in the  $N$ -person case for each player, considering all others together as nature: a stationary Markov strategy vector which is,  $\forall \varepsilon > 0$ , an  $\varepsilon$ -equilibrium of  $\Gamma(\lambda)$  for all sufficiently small  $\lambda$ , is also an equilibrium of  $\Gamma_\infty$ .

3.d.7. The perfect information case can always be rewritten (extending the state space and adjusting the discount factor) as a stochastic game where in each state only

one player moves. Actually the same conclusions go through (in the two-person case) assuming only that, in each state, the transition probability depends only on one player's action ("switching control" (Vrieze et al., 1983)): assume e.g. player I controls the transitions at state  $\omega$ ; by theorem 2.3 p. 333, we can assume that the sets  $S_0$  and  $T_0$  of best replies at  $\omega$  for the stationary Markov equilibria  $(\sigma_\lambda, \tau_\lambda)$  are independent of  $\lambda$  for  $\lambda \leq \lambda_0$ .

The equilibrium condition of prop. 1.7 p. 330 takes then the following form at  $\omega$  ( $\sigma$  and  $\tau$  being probabilities on  $S_0$  and  $T_0$ ):

$$\begin{aligned} \sum_s \sigma_{\lambda,\omega}(s)[g_\omega^2(s, t) - g_\omega^2(s, \tilde{t})] &\geq 0 & \forall t \in T_0, \forall \tilde{t} \\ V_{\lambda,\omega}^2 &= \sum_s \sigma_{\lambda,\omega}(s)[g_\omega^2(s, t) + w_{\lambda,\omega}^2(s)] & \text{for some } t \in T_0 \\ w_{\lambda,\omega}(s) &= (1 - \lambda) \sum_{\tilde{\omega}} p(\tilde{\omega} | \omega, s) V_{\lambda,\tilde{\omega}} \\ V_{\lambda,\omega}^1 &= \sum_t \tau_{\lambda,\omega}(t) g_\omega^1(s, t) + w_{\lambda,\omega}^1(s) \quad \forall s \in S_0, & \text{with inequalities for } s \notin S_0. \end{aligned}$$

The first set of inequalities describes a polyhedron of probabilities on  $S_0$ , independent of  $\lambda$ , with extreme points  $\sigma_\omega^1, \dots, \sigma_\omega^k$ . So  $\sigma_{\lambda,\omega} = \sum_{i=1}^k \mu_\omega^\lambda(i) \sigma_\omega^i$ , with  $\mu_\omega^\lambda(i) \geq 0$ . The second inequality takes then the form:  $V_{\lambda,\omega}^2 = \sum_{i=1}^k \mu_\omega^\lambda(i) [G_i(\omega) + (1 - \lambda) \sum_{\tilde{\omega}} q_{i,\omega}(\tilde{\omega}) V_{\lambda,\tilde{\omega}}^2]$ .

Hence, by semi-algebraicity again, there exist, for each  $\omega$ , indices  $i_0$  and  $i_1$  such that  $G_{i_0}(\omega) + (1 - \lambda) \sum_{\tilde{\omega}} q_{i_0,\omega}(\tilde{\omega}) V_{\lambda,\tilde{\omega}}^2 \leq V_{\lambda,\omega}^2 \leq G_{i_1}(\omega) + (1 - \lambda) \sum_{\tilde{\omega}} q_{i_1,\omega}(\tilde{\omega}) V_{\lambda,\tilde{\omega}}^2$  for all sufficiently small  $\lambda$  — say  $\lambda \leq \lambda_0$ . And conversely, for any solution  $v_\lambda^2$  of this pair of inequalities, one obtains the corresponding  $\sigma$ , by a rational computation. Similarly, the last system of inequalities, by eliminating  $\tau$ , yields a system of linear inequalities in the variables  $V_{\lambda,\omega}^1 - (1 - \lambda) \sum_{\tilde{\omega}} p(\tilde{\omega} | \omega, s) V_{\lambda,\tilde{\omega}}^1$ . Putting those inequalities together, for all states, yields a system of linear inequalities in  $V_\lambda^1$  and a similar system in  $V_\lambda^2$ : those systems have a solution for all  $\lambda \leq \lambda_0$ , and any such solution can be extended by a rational computation to a solution in  $\sigma_\lambda, \tau_\lambda$ . Finally, since such a system of linear inequalities has only finitely many possible bases, and each such base is valid in a semi-algebraic subset of  $\lambda$ 's — i.e., a finite union of intervals, one such basis is valid in a whole neighbourhood of zero: inverting it, one obtains solutions  $V_\lambda^1$  and  $V_\lambda^2$  in the field  $K(\lambda)$  of rational fractions in  $\lambda$ , with coefficients in the field  $K$  generated by the data of the game: there exists a stationary Markov equilibrium  $(\sigma, \tau)$  in  $K(\lambda)$ , for all  $\lambda \leq \lambda_1$ ,  $\lambda_1 > 0$ . Further, the above describes a finite algorithm for computing it.

Now, observe there are two cases where we can in addition assume that  $\sigma_\lambda$  is independent of  $\lambda$  — hence equal to its limit  $\sigma_0$  —:

- (1) In the zero-sum case, since every extreme point  $\sigma_\omega^i$  is a best reply of player I, they all yield him the same expected pay-off, and hence — zero-sum assumption — also to player II: hence  $V_{\lambda,\omega}^2 = G_i(\omega) + (1 - \lambda) \sum_{\tilde{\omega}} q_{i,\omega}(\tilde{\omega}) V_{\lambda,\tilde{\omega}}^2$  for all  $i$ , and thus any weights  $\mu_i^\lambda(\omega)$  are satisfactory — in particular constant weights.
- (2) If it is in all states the same player who controls the transitions, then varying  $\mu$  will only vary player II's expected pay-offs — in all states —, and those do not affect any other inequality in the system.

In each of those cases then, it is obvious that, if one were to replace in all states the passive player's  $\tau(\lambda)$  by its limit  $\tau(0)$ , which differs from it (rational fractions...) by at most  $K\lambda$ , the expected pay-offs, under whatever strategies  $\sigma$  of the controlling agents, would vary by at most  $CK\lambda$ , where  $C$  is the maximum absolute value of all pay-offs: indeed, transitions are not affected at all, so the probability distribution on the sequence of states

remains unaffected, and in every state, the current pay-off varies by at most  $CK\lambda$ . Hence by 3.d.6 above,  $(\sigma_0, \tau_0)$  is a stationary Markov equilibrium of  $\Gamma_\infty$ .

So, in those two cases, the rational fraction solutions can be chosen constant in  $\lambda$  (“uniformly discount optimal”) for the controlling player, and are also solutions for  $\lambda = 0$  — in particular, our finite algorithm computes rationally an equilibrium of the undiscounted game.

#### 4. The non-zero-sum two-person undiscounted case

Note first that there may be no uniform equilibrium, hence  $E_\infty$  may be empty. In fact already in the zero-sum case there exist no optimal strategies (cf. ex. VIIEx.4). Thus we define a set  $E_0$  of equilibrium-pay-offs as  $\bigcap_{\varepsilon > 0} E_\varepsilon$  where

$$E_\varepsilon = \{ d \in \mathbb{R}^I \mid \exists \sigma, \tau \text{ and } N \text{ such that: } \forall n \geq N \quad \bar{\gamma}_n^i(\sigma, \tau) \geq d^i - \varepsilon \quad (i = 1, 2), \text{ and} \\ \forall \sigma', \forall \tau' \quad \bar{\gamma}_n^1(\sigma', \tau) \leq d^1 + \varepsilon, \quad \bar{\gamma}_n^2(\sigma, \tau') \leq d^2 + \varepsilon \}.$$

Note that obviously  $E_\infty \subseteq E_0$  and that they coincide for supergames (cf. 4 p. 162).

**4.a. An example.** At first sight sect. 2 and 3 may lead to the idea that a proof of existence of equilibria in the undiscounted game could be obtained along the following lines: by theorem 2.3 p. 333 choose a semi-algebraic, stationary Markov equilibrium  $\sigma(\lambda)$  of the discounted game  $\Gamma_\lambda$ , for  $\lambda \in (0, 1]$ . Prove then that the limit as  $\lambda$  goes to 0 of the corresponding pay-off belongs to  $E_0$  by constructing equilibrium strategies of the kind: at stage  $n$  use  $\sigma(\lambda_n)$  where  $\lambda_n$  is some function of the past history  $h_n$ .

We prove now that this conjecture is false by studying a game where  $E_0$  is non-empty and disjoint from the set of equilibrium pay-offs,  $E_\lambda$ , of any discounted game  $\Gamma_\lambda$ .

**DEFINITION 4.1.** A stochastic game is called with **absorbing states** if all states except one are absorbing: i.e., such that  $P(\omega'; \omega, s) = 0$  if  $\omega' \neq \omega$ .

Obviously we consider then only the game starting from the non-absorbing state,  $\omega_0$  and the state changes at most once during a play (compare with recursive games (sect. 5 p. 175)).

The example is as follows: We consider a two-person game with absorbing states with two strategies for each player and the following pay-off matrix:

$$\begin{pmatrix} (1, 0)^* & (0, 2)^* \\ (0, 1) & (1, 0) \end{pmatrix}$$

where a \* denotes an absorbing pay-off (i.e., the constant pay-off corresponding to an absorbing state) and no \* means that the state is unchanged. Basically as soon as I plays Top the game is over, otherwise the game is repeated.

It is clear that the sets of feasible pay-offs in  $\Gamma_n$ ,  $\Gamma_\lambda$  or  $\Gamma_\infty$  coincide with the convex hull of  $\{(1, 0); (0, 1); (0, 2)\}$ . Moreover player I (resp. II) can guarantee  $1/2$  (resp.  $2/3$ ) (cf. ex. VIIEx.8 p. 348).

Let us write  $V = \{1/2, 2/3\}$  for the threat point and remark that the set of feasible individually rational admissible pay-offs is  $P = \{(\alpha, 2(1 - \alpha)) \mid 1/2 \leq \alpha \leq 2/3\}$  (cf. sect. 4 p. 162). Write  $E_\lambda$  (resp.  $E_n$ ) for the set of equilibrium pay-offs in  $\Gamma_\lambda$  (resp.  $\Gamma_n$ ).

**PROPOSITION 4.2.**  $E_\lambda$  is reduced to  $\{V\}$  for all  $\lambda$  in  $(0, 1]$ .

**REMARK 4.1.** The same result holds for the finitely repeated game, namely  $E_n = \{V\}$  for all  $n$  (cf. ex. VIIEx.8 p. 348).

**PROOF.** Given an equilibrium pair  $(\sigma, \tau)$  of strategies in  $\Gamma_\lambda$ , denote by  $(x, y)$  the corresponding mixed move at stage 1, where  $x$  (resp.  $y$ ) stands for the probability of Top (resp. Left). Remark first that  $x = 1$  is impossible (the only best reply would be  $y = 0$  inducing a non-individually rational pay-off for player I). Similarly we have  $y \neq 1$ .  $\lambda$  being fixed, let us write  $w$  for the maximal equilibrium pay-off of II in  $\Gamma_\lambda$  (compactness) and write again  $(\sigma, \tau)$  for corresponding equilibrium strategies. Define also  $w'$  (resp.  $w''$ ) as the normalised pay-off for player II from stage 2 on, induced by  $(\sigma, \tau)$  after the first stage history Bottom Left (resp. Bottom Right) and note that by the previous remark  $w''$  is also an equilibrium pay-off. Assume first  $x = 0$ .  $y \neq 1$  implies  $w = (1 - \lambda)w''$  which contradicts the choice of  $w$ . On the other hand  $y = 0$  implies that the only best reply is  $x = 0$ . We are thus left with the case  $x, y \in (0, 1)$ . Hence  $w'$  is also an equilibrium pay-off and the equilibrium conditions imply:

$$w = \lambda(1 - x) + (1 - \lambda)(1 - x)w' = 2x + (1 - \lambda)(1 - x)w''$$

Hence:

$$(1 - x)(\lambda + (1 - \lambda)w) \geq w \quad \text{and} \quad (1 - x)(2 - (1 - \lambda)w) \leq 2 - w$$

which implies:

$$(2 - w)(\lambda + (1 - \lambda)w) \geq w(2 - (1 - \lambda)w)$$

and finally:

$$2\lambda \geq 3\lambda w$$

So that  $E_\lambda$  is included in the set  $\{g \in \mathbb{R}^2 \mid g_2 = 2/3\}$  for all  $\lambda \in (0, 1]$ . Consider now an equilibrium pair inducing the maximal pay-off for player I, say  $u$ . Define as above  $u'$  and  $u''$  and write the equilibrium conditions:  $u = y = y(1 - \lambda)u' + (1 - y)(\lambda + (1 - \lambda)u'')$ .

Using the definition of  $u$  this leads to  $u \leq u^2(1 - \lambda) + (1 - u)(\lambda + (1 - \lambda))u$  which gives  $2u\lambda \leq \lambda$ . Hence  $u = 1/2$  and  $E_\lambda = \{V\}$ . ■

We now study  $E_0$ .

**LEMMA 4.3.**  $E_0 \subseteq P$ .

**PROOF.** Since any equilibrium pay-off is feasible and individually rational it is actually sufficient to prove that  $E_0$  is included in the segment  $[(1, 0); (0, 2)]$ . But in fact, if the probability of reaching an absorbing pay-off on the equilibrium path is less than one, player I is essentially playing Bottom after some finite stage. Note that the corresponding feasible pay-offs after this stage are no longer individually rational, hence the contradiction. Formally let  $\theta$  stands for the stopping time corresponding to the first Top and let  $(\sigma, \tau)$  be an equilibrium. Assume  $\rho = \Pr_{\sigma, \tau}(\theta = +\infty) > 0$ . Given  $\eta > 0$ , define  $N$  such that  $\Pr_{\sigma, \tau}(\theta < +\infty \mid \theta \geq N) < \eta$ . The sum of the expected pay-offs after  $N$  conditional on  $\theta \geq N$  is thus at most  $(1 - \eta) + 2\eta$ . So that for  $\eta < 1/18$  one player can get increase his pay-off by  $1/18$  by deviating. The total expected gain is at least  $\rho/18$ , hence for any  $\varepsilon$  equilibrium  $\rho \leq 18\varepsilon$ , and  $2g^I + g^{II} \geq 2 - 18\varepsilon$ . ■

Moreover one can obtain explicitly the set of equilibrium pay-offs:

**PROPOSITION 4.4.**  $E_0 = P$ .

**PROOF.** Take  $g = (\alpha; 2(1 - \alpha))$  in  $P$  with  $1/2 \leq \alpha \leq 2/3$ . We define the equilibrium strategies as follows:

- $\tau$  for player II is to play i.i.d. with  $\Pr(L) = \alpha$ .
- $\sigma$  for player I is to use an  $\varepsilon^2$ -optimal strategy in the zero-sum game  $\tilde{\Gamma}$  with pay-off matrix:

$$\begin{pmatrix} (1 - \alpha_\varepsilon)^* & -\alpha_\varepsilon^* \\ -(1 - \alpha_\varepsilon)/\varepsilon & \alpha_\varepsilon/\varepsilon \end{pmatrix}$$

with  $\alpha_\varepsilon = \alpha - \varepsilon$ .

Hence if  $\tilde{\gamma}$  is the pay-off for player I in  $\tilde{\Gamma}$ , for  $n$  large enough, and any  $\tau'$ :  $\tilde{\gamma}_n(\sigma, \tau') \geq -\varepsilon^2$ , so that  $\rho_{\sigma, \tau'}(\theta < +\infty) \geq 1 - \varepsilon$  and  $\sigma$  is an  $\varepsilon$  best reply to  $\tau$ , since  $\alpha \geq 1/2$ .

On the other hand, we have for  $n$  sufficiently large and any  $\tau'$ :

$$(1) \quad \tilde{\gamma}_n(\sigma, \tau') = \rho((1 - \alpha_\varepsilon)\bar{y} - \alpha_\varepsilon(1 - \bar{y})) + (1 - \rho)(-(1 - \alpha_\varepsilon)\tilde{y} + \alpha_\varepsilon(1 - \tilde{y}))/\varepsilon \geq -\varepsilon^2$$

and

$$(2) \quad \rho((1 - \alpha_\varepsilon)\bar{y} - \alpha_\varepsilon(1 - \bar{y})) \geq -\varepsilon^2$$

where  $\rho = \frac{1}{n} \mathbb{E}[(n - \theta)^+]$ ,  $\rho\bar{y} = \frac{1}{n} \mathbb{E}[(n - \theta)^+ \mathbf{1}_{t_\theta=L}]$ ,  $\tilde{y}(1 - \rho) = \frac{1}{n} \sum_1^n \mathbb{E}\{t_m = L, \theta > m\}$ . So that since  $\gamma_n^{\text{II}}(\sigma, \tau') = \rho(2(1 - \bar{y})) + (1 - \rho)\tilde{y}$  we obtain:

$$\begin{aligned} \gamma_n^{\text{II}}(\sigma, \tau') &\leq 2\rho - (2 - \varepsilon)\rho\bar{y} + \varepsilon^3 + (1 - \rho)\alpha_\varepsilon - \varepsilon\rho\alpha_\varepsilon && \text{by (1)} \\ &\leq 2\rho + (2 - \varepsilon)(\varepsilon^2 - \rho\alpha_\varepsilon) + \varepsilon^3 + (1 - \rho)\alpha_\varepsilon - \varepsilon\rho\alpha_\varepsilon && \text{by (2)} \\ &= \rho(2 - 3\alpha_\varepsilon) + \alpha_\varepsilon + 2\varepsilon^2 \\ &\leq 2(1 - \alpha_\varepsilon) + 2\varepsilon^2 && \text{since } \alpha \leq 2/3 \\ &\leq 2(1 - \alpha) + 2\varepsilon + 2\varepsilon^2 \end{aligned}$$

hence the inclusion  $E_0 \supseteq P$  and the result by the previous lemma 4.3. ■

**COMMENT 4.2.** The two previous propositions show a clear discontinuity in the set of equilibrium pay-offs between  $E_\lambda$  and  $E_0$  as  $\lambda$  goes to 0 (compare with the zero-sum case, e.g. 3.6 p.338). Recall that the above  $\varepsilon$ -equilibria are also  $\varepsilon$ -equilibria in all sufficiently long finite games — hence also in all  $\Gamma_\lambda$  for  $\lambda$  sufficiently small. One is thus lead to wonder whether  $E_0$  is not a more appropriate concept when analysing long games than  $\lim_{\lambda \rightarrow 0} E_\lambda$  — does the latter not rely too much on common knowledge of each other's exact discount factor, and on exact maximisation?

**4.b. Games with absorbing states.** This section shows the existence of equilibria for two-person games with absorbing states (Vrieze and Thuijsman, 1989).

By selecting an equilibrium pay-off for each potential new state and taking expectation w.r.t. the transition probability, we are reduced to the case where the game is described by an  $S \times T$  matrix where the  $(s, t)$  entry consists of a pay-off vector  $G_{s,t}$  in  $\mathbb{R}^2$ , a number  $P_{s,t}$  in  $[0, 1]$  (probability of reaching the set of absorbing states) and a corresponding absorbing pay-off say  $A_{s,t}$  in  $\mathbb{R}^2$ . Given a non-absorbing past, if  $(s, t)$  is the pair of choices at stage  $n$  the pay-off at that stage is  $g_n = G_{s,t}$ . With probability  $P_{s,t}$  the pay-off for all future stages will be  $A_{s,t}$ , and with probability  $(1 - P_{s,t})$  stage  $n + 1$  is like stage  $n$ .

We shall write  $X, Y$  for the set of stationary Markov strategies in  $\Gamma_\infty$  that we identify with the set of mixed actions. By theorem 2.3, for  $\lambda$  small enough there exist equilibrium strategies  $(x_\lambda, y_\lambda)$  in  $\Gamma_\lambda$ , with pay-off  $w_\lambda = \bar{\gamma}_\lambda(x_\lambda, y_\lambda)$  in  $\mathbb{R}^2$ , such that  $(x_\lambda, y_\lambda)$  converges to  $(x_\infty, y_\infty)$  and  $w_\lambda$  to  $w_\infty$  as  $\lambda \rightarrow 0$ . Let  $v_\lambda^I$  (resp.  $v_\infty^I$ ) be the value of the zero-sum game obtained through  $\Gamma_\lambda$  (resp.  $\Gamma_\infty$ ) when player II minimises player I's pay-off and similarly  $v_\lambda^{\text{II}}, v_\infty^{\text{II}}$ . Finally  $v_\infty = (v_\infty^I, v_\infty^{\text{II}})$ .

LEMMA 4.5.  $w_\infty^i \geq v_\infty^i$ ,  $i = \text{I}, \text{II}$ .

PROOF. The equilibrium condition implies  $w_\lambda^i \geq v_\lambda^i$  and theorem 3.1 shows that  $\lim_{\lambda \rightarrow 0} v_\lambda^i = v_\infty^i$ .  $\blacksquare$

LEMMA 4.6.  $w_\infty^{\text{I}} = \lim_{\lambda \rightarrow 0} \bar{\gamma}_\lambda^{\text{I}}(x_\infty, y_\lambda)$ .

PROOF. The equilibrium condition implies  $\bar{\gamma}_\lambda^{\text{I}}(x_\infty, y_\lambda) \leq w_\lambda^{\text{I}}$ . Now for  $\lambda$  small enough the support  $S(x_\infty)$  of  $x_\infty$  in  $S$  is included in  $S(x_\lambda)$  hence the result.  $\blacksquare$

Note that for any pair of stationary Markov strategies  $(x, y)$  the pay-off in  $\Gamma_\infty$  is well defined (as  $\lim_{\lambda \rightarrow 0} \bar{\gamma}_\lambda(x, y)$  or  $\lim_{n \rightarrow \infty} \bar{\gamma}_n(x, y)$ ) and will simply be denoted by  $\gamma(x, y)$ . The vector  $\gamma(x_\infty, y_\infty)$  will play a crucial rôle in the proof: it corresponds to the undiscounted pay-off of the limit of the discounted (optimal) strategies (as opposed to  $w_\infty$  limit of the discounted pay-off of the discounted strategies).

DEFINITION 4.7. Let  $R \subseteq S \times T$  be the set of absorbing entries, i.e., such that  $P_{s,t} > 0$ . Define  $B(x, y) = \sum x_s A_{st} P_{s,t} y_t$ ,  $P(x, y) = \sum x_s P_{s,t} y_t$ . A couple  $(x, y)$  in  $X \times Y$  is **absorbing** if  $(S(x) \times T(y)) \cap R \neq \emptyset$  or equivalently  $P(x, y) > 0$ . In this case let  $A(x, y) = B(x, y)/P(x, y)$ .

Since

$$(3) \quad \bar{\gamma}_\lambda(x, y) = \lambda x G y + (1 - \lambda)[B(x, y) + (1 - P(x, y))\bar{\gamma}_\lambda(x, y)]$$

we obtain that, if  $\lim_{\lambda \rightarrow 0} \tilde{x}_\lambda = x$ ,  $\lim_{\lambda \rightarrow 0} \tilde{y}_\lambda = y$  with  $(x, y)$  absorbing, then:

$$(4) \quad \lim_{\lambda \rightarrow 0} \bar{\gamma}_\lambda(\tilde{x}_\lambda, \tilde{y}_\lambda) = \gamma(x, y) = A(x, y).$$

In particular the equilibrium condition implies that:

$$(5) \quad \text{if } (x_\infty, y) \text{ is absorbing} \quad \gamma^{\text{II}}(x_\infty, y) \leq w_\infty^{\text{II}}$$

and using lemma 4.6 that:

$$(6) \quad \text{if } (x_\infty, y_\lambda) \text{ is non-absorbing for } \lambda \in (0, \bar{\lambda}] \text{ then } \gamma^{\text{I}}(x_\infty, y_\infty) = w_\infty^{\text{I}}.$$

It follows that either:

- Case A:** •  $(x_\infty, y_\infty)$  is absorbing and  $\gamma(x_\infty, y_\infty) = w_\infty$  by (4)
- or  $(x_\infty, y_\lambda)$  and  $(x_\lambda, y_\infty)$  are non-absorbing for  $\lambda \in (0, \bar{\lambda}]$  and again  $\gamma(x_\infty, y_\infty) = w_\infty$  by (6); or:

- Case B:**  $(x_\infty, y_\infty)$  is non-absorbing and  $(x_\infty, y_\lambda)$  or  $(x_\lambda, y_\infty)$  is absorbing for  $\lambda$  sufficiently small.

Assume for example  $(x_\infty, y_\lambda)$  absorbing and define  $y_\lambda''$  as the restriction of  $y_\lambda$  on  $\{t \mid (x_\infty, t)$  is absorbing $\}$  and  $y_\lambda = y_\lambda' + y_\lambda''$ . Obviously  $y_\lambda' \rightarrow y_\infty$ . Let then  $z$  be a limit point in  $Y$  of  $z_\lambda = y_\lambda''/\|y_\lambda''\|$  as  $\lambda \rightarrow 0$ . Since  $\bar{\gamma}_\lambda^{\text{II}}(x_\lambda, z_\lambda) = \bar{\gamma}_\lambda^{\text{II}}(x_\lambda, y_\lambda)$  and  $(x_\infty, z)$  is absorbing we obtain by (4)  $\gamma^{\text{II}}(x_\infty, z) = w_\infty^{\text{II}}$ . Using (3) we get:

$$\bar{\gamma}_\lambda^{\text{I}}(x_\infty, y_\lambda) = \frac{\lambda}{\lambda + (1 - \lambda)P(x_\infty, y_\lambda)} x_\infty G y_\lambda + \left(1 - \frac{\lambda}{\lambda + (1 - \lambda)P(x_\infty, y_\lambda)}\right) \frac{B(x_\infty, y_\lambda)}{P(x_\infty, y_\lambda)}.$$

Note that  $B(x_\infty, y_\lambda) = B(x_\infty, y_\lambda'')$  and  $P(x_\infty, y_\lambda) = P(x_\infty, y_\lambda'')$ . Thus  $A(x_\infty, y_\lambda) = A(x_\infty, y_\lambda'') = A(x_\infty, z_\lambda)$ . Hence if  $\mu$  is a limit point of  $\lambda/(\lambda + (1 - \lambda)P(x_\infty, y_\lambda))$  we obtain, using lemma 4.6 and (4)

$$(7) \quad w_\infty^{\text{I}} = \mu \gamma^{\text{I}}(x_\infty, y_\infty) + (1 - \mu) \gamma^{\text{I}}(x_\infty, z).$$

We can now prove the following:

THEOREM 4.8.  $E_0 \neq \emptyset$ .

PROOF. (1) Assume first that  $\gamma^i(x_\infty, y_\infty) \geq w_\infty^i$ ,  $i = I, II$ . By lemma 4.5  $\gamma^i(x_\infty, y_\infty) \geq v_\infty^i$ ,  $i = I, II$ . The equilibrium strategies consist first of playing  $(x_\infty, y_\infty)$  until some large stage  $N$  and then to punish, i.e., to reduce the pay-off to  $v_\infty^I$ , say, as soon as  $\|\bar{x}_n - x_\infty\| \geq \varepsilon$  for  $n \geq N$  where  $\bar{x}_n$  is the empirical distribution of moves of player I up to stage  $n$ , and symmetrically for II. To prove the equilibrium condition, note that non-absorbing deviations will be observed and punished and by (5) absorbing deviations are not profitable since by stationarity the future expected pay-off, before an absorbing state is still  $\gamma(x_\infty, y_\infty)$ . Hence  $\gamma(x_\infty, y_\infty) \in E_0$ .  
(2) If  $\gamma^I(x_\infty, y_\infty) < w_\infty^I$  then by the above analysis we are in case B where  $(x_\infty, y_\lambda)$  is absorbing for  $\lambda$  small. Then (7) implies that  $\gamma^I(x_\infty, z) > w_\infty^I$  and  $\gamma(x_\infty, z)$  belongs to  $E_0$ : In fact  $\gamma^{II}(x_\infty, z) = w_\infty^{II}$  hence by lemma 4.5  $\gamma(x_\infty, z) \geq v_\infty$  and we now describe the equilibrium strategies:

- Player I uses  $x_\infty$ , until some large stage  $N_1$  and then reduces player II's pay-off to  $v_\infty^{II}$ ;
- Player II plays  $y_\infty$  with probability  $(1 - \delta)$  and  $z$  with probability  $\delta$  as long as  $\bar{x}_n$  is near  $x_\infty$  for  $n \geq N_2$ , and reduces I's pay-off to  $v_\infty^I$  otherwise.

Given  $\varepsilon > 0$ ,  $N_2$ ,  $\delta$  and  $N_1$  are chosen such that:  $P_{x_\infty}(|\bar{x}_n - x_\infty| \geq \varepsilon) \leq \varepsilon$  for  $n \geq N_2$ ,  $(1 - \delta)^{N_2} \leq \varepsilon$ ,  $N_1 \geq N_2$  and  $(1 - \delta)^{N_1} \geq \varepsilon$  so that  $n \geq N_1/\varepsilon$  implies  $\bar{\gamma}_n(\sigma, \tau) \geq \gamma(x_\infty, z) - 4\varepsilon$ . If player II deviates, a non-absorbing pay-off will not be profitable and by (5) the same will be true for an absorbing pay-off.

As for player I, let  $\theta$  be the first stage if any, where  $|\bar{x}_n - x_\infty| \geq \varepsilon$ . For  $n \geq N_1/\varepsilon$ , if  $\theta \geq n$  the pay-off against  $z$  is near  $\gamma^I(x_\infty, z)$  and against  $y_\infty$ , if absorbing less than  $w^I$  by (5). Similarly if  $N_2 < \theta < n$  the absorbing pay-off up to  $\theta$  is near  $\gamma^I(x_\infty, z)$  or less and the pay-off after less than  $w^I$ . Finally if  $\theta \leq N_2$  the pay-off up to  $\theta$  is with probability  $\geq 1 - \varepsilon$  obtained against  $y_\infty$ , hence less than  $w^I$  if absorbing. ■

## Exercises

### 1. Properties of the operator $\Psi$ .

a. Consider a zero-sum stochastic game with finite state space  $K$ . For  $x \in \mathbb{R}^K$ , we want the value  $\Psi(x)$  of  $\Gamma(x)$  to exist: assume thus  $S$  and  $T$  compact, and  $g^k(s, t)$  to be separately u.s.c.-l.s.c., and  $p(k | \ell; s, t)$  to be separately continuous (even if one were interested only in  $x \geq 0$ , so an u.s.c.-l.s.c. assumption would suffice, this assumption would imply separate continuity because probabilities sum to one). Then  $\Psi: \mathbb{R}^K \rightarrow \mathbb{R}^K$  is a well-defined operator.

i.  $\Psi$  is monotone and  $\Psi(x + c) = \Psi(x) + c$  for any constant vector  $c$ . [cf. I.1Ex.1 p. 9].

To obtain additional properties of  $\Psi$ , assume that  $g$  is bounded, so

ii.  $\|\lambda\Psi(x) - \mu\Psi(y)\| \leq M|\lambda - \mu| + \|\lambda x - \mu y\|$

(using systematically the maximum norm on  $\mathbb{R}^K$ )

and that either  $S$  and  $T$  are metrisable (ex. I.1Ex.7b p. 12) or  $g$  is separately continuous (ex. I.2Ex.1 p. 20), so ex. I.1Ex.6 p. 11 and I.1Ex.8 p. 12 are applicable to the functions  $f(\lambda, x) = \lambda\Psi^k(x/\lambda)$  ( $\lambda \geq 0$ ) defined on  $\mathbb{R}_+ \times \mathbb{R}^K$ . Taking thus derivatives at  $(0, x)$  yields

iii.  $\|\Psi(x + y) - A(x) - \Psi_x(y)\| \leq (1 + \|y\|)F[x/(1 + \|y\|)]$

where  $\Psi_x = \Psi_{\lambda x}$  ( $\lambda > 0$ ),  $A$  is positively homogeneous of degree 1,  $F \leq M$  and  $\lim_{\lambda \rightarrow \infty} F(\lambda x) = 0$ .

Also the formula for the derivatives yields that

iv.  $\Psi_x$  has the same properties as  $\Psi = \Psi_0$ , with the same constant  $M$ , and associated  $A_x$ ,  $F_x$ ,  $\Psi_{xy}$  etc. (and  $\Psi_x = \Psi_{x0} = \Psi_{xx}$ ).

[More precisely, for some universal constant  $c$ , VIIEx.1ai–VIIEx.1aiv hold for  $M = c \cdot \sup_{k,s,t} |g^k(s,t)|$ .]

b. A number of other properties follow from the above:

i.  $\|\Psi(y)\| \leq M + \|y\|$  (cf. VIIEx.1a(ii))

ii.  $\|\Psi(x) - A(x)\| \leq 2M$

HINT. Use VIIEx.1bi for  $\Psi_x$  (cf. VIIEx.1a(iv)) and  $y = 0$ , then VIIEx.1a(iii) for  $y = 0$  — actually, a direct proof yields  $\|\Psi(x) - A(x)\| \leq M$ .

iii.  $A(x) = \lim_{\lambda \rightarrow 0} \lambda^{-1} \Psi(\lambda x)$  (from VIIEx.1b(ii) and homogeneity).

iv.  $A$  is monotone,  $A(x+c) = A(x) + c$  (from VIIEx.1b(iii) and VIIEx.1a(i))

v.  $\|A(x+y) - A(x) - A_x(y)\| \leq \|y\| F(x/\|y\|)$ , (from VIIEx.1a(iii), using VIIEx.1b(iii) and VIIEx.1a(iv))

vi.  $\Psi_x(y) = \lim_{\lambda \rightarrow \infty} [\Psi(\lambda x + y) - \lambda A(x)]$  (from VIIEx.1a(iii))

vii.  $A_x(y) = \lim_{\lambda \rightarrow \infty} [A(\lambda x + y) - \lambda A(x)]$  (from VIIEx.1b(v) (and VIIEx.1a(iii)) — besides the formula implied by VIIEx.1b(iii) via VIIEx.1a(iv)).

viii. (“Euler’s formula”)  $A_x(x) = A(x)$ .

In particular, VIIEx.1ai–VIIEx.1a(iv) determine uniquely  $A$  (cf. VIIEx.1b(iii)) and  $\Psi_x$  (cf. VIIEx.1b(v)).

**2.  $\lim v_n$  for games with absorbing states.** (Kohlberg, 1974) Give a direct proof of existence of  $\lim v_n$  for games with absorbing states.

HINT. Let  $\bar{u}(\omega) = g(\omega)$  for  $\omega \neq \omega_0$ ,  $\bar{u}(\omega_0) = u$ . Prove that  $D(u) = \lim_{\lambda \rightarrow \infty} [\Psi^{\omega_0}(\lambda \bar{u}) - (\lambda + 1)u]$  exists for all  $u$  (cf. ex. VIIEx.1 above) and is strictly monotone. Letting  $u_0$  be such that  $(u - u_0)D(u) < 0$  for all  $u \neq u_0$ , prove first that for every  $\varepsilon > 0$ ,  $\limsup v_n > u_0 - \varepsilon$  then that  $\liminf v_n > u_0 - \varepsilon$ ; conclude that  $\lim v_n = u_0$ .

**3.** Consider the following  $(n \times n)$  game with absorbing states (recall that a \* denotes an absorbing pay-off and no \* means that the state is unchanged).

$$\begin{pmatrix} 1^* & \dots & 0^* & \dots & 0^* \\ \vdots & \ddots & & & \vdots \\ 0 & & 1^* & & 0^* \\ \vdots & & & \ddots & \vdots \\ 0 & \dots & 0 & \dots & 1^* \end{pmatrix}$$

and prove that  $v(\lambda) = (1 - \lambda^{1/n})/(1 - \lambda)$ .

**4. Big Match.** (Blackwell and Ferguson, 1968) Consider the following game with absorbing states:  $\begin{pmatrix} 1^* & 0^* \\ 0 & 1 \end{pmatrix}$

a. Prove that  $v_n = v_\lambda \equiv 1/2 \quad \forall n, \forall \lambda$ .

b. Prove that player I has no optimal strategies in  $\Gamma_\infty$ .

c. Prove that player II can guarantee 1/2.

d. Prove that player I cannot guarantee more than 0 with Markov strategies.

e. A strategy  $\varphi$  with **finite memory** (say for I) is defined by an internal (finite) state space, say  $K$  and two functions:  $q$  from  $K \times T$  to  $\Delta(K)$  that determines the new internal state as a function of the old one and of the move of the opponent, and  $f$  from  $K$  to  $\Delta(S)$  that determines the move as a function of the current internal state. Prove that by adding new states one can assume  $f$  to be  $S$  valued. Remark that  $\varphi$  and a stationary strategy  $y$  of player II define a Markov chain on  $K$ , and that we can identify the internal states that induce through  $f$  the action Top. Use then the fact that the ergodic classes depend only on the support of  $y$  to deduce that I cannot guarantee more than 0 with such strategies.

f. Verify that if we modify the original proof of existence of  $v_\infty$  (Blackwell and Ferguson, 1968) (where  $\lambda(s) = 1/s^2$ ,  $N(s) = 1$ ,  $\varepsilon = 0$ ,  $M = 0$ , and  $s_0$  is large) by letting, at stage  $k$ ,  $\bar{\lambda}_k = \inf_{i \leq k} \lambda_i$  and playing  $\sigma_{\bar{\lambda}_k}$  we obtain  $E(g_\infty) = 0$  and  $\liminf E(g_n) < 1/2$ .

g. Verify in the following variant that a lower bound  $M$  is needed for the optimal strategy (I.e., there is an upper bound on  $\lambda_k$ ):

$$\begin{pmatrix} 1^* & 0^* & 0^* \\ 0 & 1 & 1/2^* \end{pmatrix}$$

h. Show that, for a game where the absorbing boxes of the matrix are all those belonging to a fixed subset of rows,  $v_n$  and  $v_\lambda$  are constant.

**5. Need for Uniform Convergence.** Consider the following one-person game:  $\Omega = \mathbb{N} \cup \{\delta\}, S = \mathbb{N}$ .

$$P(n; n+1, s) = 1 \quad \forall n \geq 0, \forall s; P(0; 0, s) = 1 \quad \forall s; P(n; \delta, s) = 1 \text{ iff } n = s.$$

$$\forall s \quad g(\omega, s) = \begin{cases} -1 & \text{if } \omega = 0 \\ 0 & \text{otherwise} \end{cases}$$

Prove that  $\lim v_\lambda(\omega) = \lim v_n(\omega) = v_\infty(\omega) = -1$ , for all  $\omega \neq \delta$ ; and that  $v_\lambda(\delta) = v_n(\delta) = 0$ . Note that, on every history starting from  $\delta$ ,  $\liminf \bar{x}_n$  is  $-1$ , but if  $v_\infty$  exists in our strong sense it should equal  $\lim v_\lambda$ .

Prove that  $\|v_\lambda - \lim v_\lambda\| = 1$ .

Note that if we consider this game as a recursive game (cf. sect. 5 p. 175) it has a well-defined value, everywhere  $-1$ .

**6. Study of  $\lim v_n$ .** (Bewley and Kohlberg, 1976a)

a. Let  $\lambda V_\lambda$  be the value of  $\Gamma_\lambda$  (lemma 1.2).  $V_\lambda = \Psi((1 - \lambda)V_\lambda)$  (lemma 1.2). Assume  $\Omega$  finite. Then there exists  $M > 0$  and vectors  $a_\ell \in \mathbb{R}^\Omega$ ,  $\ell = 0, \dots, M-1$  such that  $V_\lambda = v_\infty/\lambda + \sum_{\ell=0}^{M-1} a_\ell \lambda^{-\ell/M} + O(\lambda^{1/M})$ .

Deduce that  $\Psi(v_\infty(\lambda^{-1} - 1) + \sum_{\ell=0}^{M-1} a_\ell \lambda^{-\ell/M}) = v_\infty/\lambda + \sum_{\ell=0}^{M-1} a_\ell \lambda^{-\ell/M} + O(\lambda^{1/M})$  and that for  $n$  large enough

$$(8) \quad \Psi(v_\infty(n-1) + \sum_{\ell=0}^{M-1} a_\ell n^{\ell/M}) = v_\infty n + \sum_{\ell=0}^{M-1} a_\ell n^{\ell/M} + O(n^{-1/M})$$

Prove then that if  $V_n = nv_n$  is the non-normalised value of  $\Gamma_n$ , one has:

$$\left| V_n - v_\infty n - \sum_{\ell=0}^{M-1} a_\ell n^{\ell/M} \right| \leq \left| V_{n-1} - v_\infty(n-1) - \sum_{\ell=0}^{M-1} a_\ell n^{\ell/M} \right| + O(n^{-1/M})$$

hence that  $\lim v_n = v_\infty$ .

REMARK 4.3. Bewley and Kohlberg prove in fact there is a similar expansion for the value of the  $n$ -stage game:  $v_n = v_\infty + \sum_{k=1}^{M-1} a_k n^{-k/M} + n^{-1} \cdot O(\log n)$ .

The expansion has to stop at that stage: cf. Ex. VIIEx.7 below.

b. Consider now the perfect information case so that  $M = 1$ . Prove that (8) can be strengthened to:

$$\Psi(v_\infty(n-1) + a_0) = v_\infty n + a_0$$

Deduce that  $|V_n - v_\infty n|$  is uniformly bounded.

**7.** (Bewley and Kohlberg, 1976a) Consider the game  $\begin{pmatrix} -1^* & 1 \\ 1 & 0 \end{pmatrix}$ . Let  $x_1 = 1/3, x_{n+1} = x_n + 1/(n+3)$ . Prove that  $V_n \leq x_n$  and  $x_{n+1} - V_{n+1} \leq x_n - V_n + \beta_n$  where  $(\beta_n)$  is a convergent series. Deduce that  $V_n \sim \log n$ , so has no asymptotic expansion in fractional powers of  $n$  — cf. VIIEx.6a above.

### 8.

a. Prove that in the example in sect. 4.a the minmax is  $1/2$  for player I and  $2/3$  for player II (use ex. VIIEx.4).

b. Prove by induction that  $E_n = \{V\}$ .

**9. Feasible pay-offs.** Denote by  $F_\lambda$  resp.  $F_\infty$  the set of feasible pay-offs in  $\Gamma_\lambda$ , resp.  $\Gamma_\infty$ . (Note that they depend on the initial state  $\omega$ ). Prove that  $\text{Co } F_\lambda$  and  $\text{Co } F_\infty$  have finitely many extreme points and that  $\text{Co } F_\lambda$  converges to  $\text{Co } F_\infty$ .

HINT. Consider the one-person dynamic programming problem with pay-off the scalar product  $\langle u, g \rangle$  and let  $u$  vary in  $\mathbb{R}^I$ .

**10. Irreducible Games.** A stochastic game is irreducible if for any vector of (pure) stationary Markov strategies the induced Markov chain on states is irreducible.

a. *Zero Sum Case.* Prove that  $\lim v_n$  exists and is independent of the initial state. Deduce that  $v_\infty$  exists.

b. *Non-Zero-Sum Case.* Prove that  $F_\infty$  is convex and independent of the initial state. Prove that  $F_\lambda$  converges to  $F_\infty$ . Deduce results analogous to theorems 4.1 p. 163 and 4.2 p. 164, and ex. IV.4Ex.3 p. 172 and IV.4Ex.8 p. 173.

### 11. Correlated Equilibrium.

a. (Nowak and Raghavan, 1992) Consider a game as in sect. 1.c p. 330 with moreover  $(\Omega, \mathcal{A})$  separable, and  $P(\cdot | \omega, s)$  dominated by  $\mu$ . Prove the existence of a stationary public extensive form-correlated equilibrium (i.e. generated by a sequence of public i.i.d. random variables).

HINT. Consider the proof of theorem 1.8 p. 331. For proving the measurability of  $\omega \mapsto N_f(\omega)$  (closed subsets of compact metric spaces being endowed with the Hausdorff topology), identify  $G_f(\omega)$  with its graph in  $\overline{S} \times [-c/\lambda, c/\lambda]^I$ , and show

- (1)  $\omega \mapsto G_f(\omega)$  is measurable;
- (2) the map from the space of all such graphs to the corresponding set of Nash equilibria is upper semi-continuous, hence Borel; and
- (3) the continuity of the map from (Graph, Nash Equilibrium) to pay-offs.

Note that  $f \mapsto \mathcal{N}_f$  is still weak upper semi-continuous hence there exists a fixed point  $f_0 \in \mathcal{N}_{f_0}$ . Use then a measurable version of Carathéodory's theorem (Castaing and Valadier, 1977, Th. IV.11) (or Mertens, 1987a) to represent  $f_0$  as  $\sum_{i=1}^{I+1} \lambda_i(\cdot) f_i(\cdot)$  with  $f_i \in N_{f_0}$ . If  $x_n$  is uniform on  $[0, 1]$  play at stage  $n$ , if  $\omega$ , a measurable equilibrium leading to  $f_i(\omega)$  in  $\Gamma((1-\lambda)f_0)\omega$  (using the inverse image of the continuous projection (graph, Nash Equilibrium)  $\rightarrow$  (graph, pay-offs)) if  $x_n \in [\sum_{j=1}^{i-1} \lambda_j(\omega), \sum_{j=1}^i \lambda_j(\omega)]$ .

b. (Duffie et al., 1994)

i. Let  $S$  be a complete separable metric space and  $G$  a correspondence from  $S$  to  $\Delta(S)$ , convex valued and with closed graph.  $J$  measurable in  $S$  is self justified if  $G(s) \cap \Delta(J) \neq \emptyset$  for all  $s \in J$ . For  $J$  self-justified,  $\mu$  is invariant for  $J$  (under  $\pi$ ) if  $\pi$  is a measurable selection from  $G$  and  $\mu$  in  $\Delta(J)$  satisfies

$$\int \pi(s)(B)\mu(ds) = \mu(B)$$

for all  $B$  measurable  $\subseteq J$ . Prove that if  $J$  is compact and self-justified the set of invariant measures is non-empty compact convex: hence there exists an ergodic measure for  $\pi$  (extreme point). Consider the restriction of  $G$  to  $J$  and  $\Delta(J)$  that we still denote by  $G$ . Let  $\overline{G}$  be the graph and  $m_1^{(\nu)}, m_2^{(\nu)}$  the two marginals of  $\nu \in \Delta(\overline{G})$  on  $G$  and  $\Delta(G)$ . If  $\eta \in \Delta(\Delta(J))$  let  $E(\eta)$  denote its mean. Show that  $E \circ m_2 \circ m_1^{-1}$  has a fixed point  $\mu$ : there exists some  $\nu$  such that  $E m_2(\nu) = m_1(\nu) = \mu$ . Thus there exists a transition probability  $P$  from  $J$  to  $\Delta(\Delta(J))$  with  $\mu \otimes P = \nu$ . Show that  $EP(s) \in G(s)$   $\mu$ -a.s.

HINT. Given  $f$  real continuous on  $J$  and  $c$  real, let  $A_2 = \{ \rho \in \Delta(J) \mid \int f d\rho > c \}$  and suppose that  $\mu\{s \mid G(s) \cap A_2 = \emptyset, (EP(s)) \in A_2\} > 0$ . Then  $A_1 = \{s \mid G(s) \cap A_2 = \emptyset, P(s)(A_2) > 0\}$  satisfies  $\mu(A_1) > 0$  and  $0 = \nu(\overline{G} \cap A_1 \times A_2) = \int_{A_1} P(s)(A_2) \mu(ds) > 0$ .

Deduce that there exists a measurable selection  $\Pi$  of  $G$  with  $\mu \circ \Pi = \mu$ . Observe that the set of  $\nu$  satisfying  $Em_2(\nu) = m_1(\nu)$  is compact convex and that every invariant measure  $\mu$  is obtained from some solution  $(\mu \otimes \Pi)$ .

ii. Consider a discounted stochastic game where  $\Omega, S^i$  are metric compact,  $g^i$  is continuous on  $\Omega \times S$  and the transition satisfies:

- norm continuity on  $\Omega \times S$
  - joint mutual absolute continuity, i.e.  $q(\cdot; \omega, s) \ll q(\cdot; \omega', s')$ .

Let  $c = \|g\|/\lambda$  and  $Z = [-c, +c]^{\mathbf{I}}$ . Define  $W = \Omega \times Y$ ,  $Y = \prod_i \Delta(S^i) \times Z$  and a correspondence  $G$  from  $W$  to  $\Delta(W)$  by: for  $\bar{w} = (\bar{\omega}, \bar{x}, \bar{z})$ ,  $G(\bar{w}) = \{\nu \in \Delta(W) \mid \nu = \nu_{\Omega} \otimes \nu^Y$  (marginal on  $\Omega$ , conditional on  $Y$  given  $\omega$ ) such that

- $\nu_\Omega = q(\cdot \mid \bar{\omega}, \bar{x}) \equiv \int q(\cdot \mid \bar{\omega}, s) \bar{x}(ds)$
  - $\bar{x}$  is an equilibrium with pay-off  $\bar{z}$  in the game

$$f^i(x, \overline{\omega}) = g^i(x, \overline{\omega}) + (1 - \lambda) \int z^i d[q(\cdot \mid \overline{\omega}, x) \otimes \nu^Y].$$

Show that  $G$  is convex valued with closed graph. (Use VIIEx.11bi to get the norm continuity of  $q(\cdot | \omega, s)$  and note that for  $\nu_n \in \Delta(A \times B)$ ,  $\rho_n \in \Delta(A)$ ,  $A, B$  compact metric,  $\nu_n \rightarrow \nu$  (weak),  $\nu_{nA} \rightarrow \nu_A$  (norm),  $\rho_n \rightarrow \rho$  (norm),  $\rho \ll \nu_A$  imply  $\rho_n \otimes \nu_n^B \rightarrow \rho \otimes \nu^B$  (weak)).

Deduce the existence of a justified set  $W^* \subseteq W$  for  $G$ . (Consider the sequence of compact  $W = W_0, W_{n+1} = \{w \in W \mid \exists \nu \in G(w), \nu(W_n) = 1\}$ ). (Prove inductively that  $\text{Proj}_\Omega(W_n) = \Omega$  (theorem 4.1 p. 39)).

iii. Deduce the existence of a stationary equilibrium with an ergodic measure in the extended game with state space  $W^*$  and transition from  $W^* \times S$  to  $\Delta(W^*)$  given by

$$q^*(\cdot \mid \omega, x, z; s) \equiv q(\cdot \mid \omega, s) \otimes \Pi(\omega, x, z)$$

(where  $\Pi$  is the selection from  $G$ ). (Let  $\sigma^*(\omega, x, z) = x$ ).

**12.** (Federgruen, 1978) In the following stochastic game with perfect information there are no pure stationary Markov equilibria (the first two coordinates indicate pay-offs, the other transitions):

**13.** (Nowak and Raghavan, 1991) Consider the following game with countable state space  $\mathbb{N}$ . The pay-off is always 0 except in state 1 where it is 1. States 0 and 1 are absorbing. Given state  $\omega = n$  the transition probabilities are deterministic and given by the following  $S \times T$  matrix

$$\begin{pmatrix} 1 & 0 \\ n+1 & n-1 \end{pmatrix}$$

a. Prove that  $v_\infty$  exists and player II has an optimal stationary Markov strategy. Deduce that  $v_\infty(n)$  decreases to  $1/2$ .

HINT. Consider the largest state  $\omega$  where  $v(\omega) \geq \frac{1}{2}$ . Prove that there exists  $\delta > \frac{1}{2}$  such that the probability  $y(\cdot)$  of II playing Left satisfies  $y(\omega') \geq \delta$  for  $\omega' \geq \omega$ .

b. Prove that with stationary Markov strategies player I cannot guarantee more than 0, as the initial state goes to infinity.

HINT. Let  $x(\cdot)$  be the probability of Top. If  $\sum_\omega x(\omega) < \infty$ , II plays Left, and Right otherwise.

COMMENT 4.4. This game can be viewed as a recursive game or as a “positive” game, i.e., with pay-off the sum of the (positive) stage pay-offs. Compare with sect. 5 p. 175 and (Ornstein, 1969) — cf. Ex. VII.5Ex.4a p. 355, cf. also (Nowak, 1985b).

**14.** Prove remark 3.5 p. 338 (use remark 3.2 p. 336).

### 15. Games with no signals on moves.

a. In a stochastic game with finite state and action sets,  $v_n$  and  $v_\lambda$  remain unchanged if the moves are not announced, but  $v_\infty$  may no longer exist.

b. (Coulomb, 1992) Consider a game with absorbing states:  $G$  (resp.  $A$ ) corresponds to the non-absorbing (resp. absorbing) pay-offs and  $P$  is the absorbing transition. Define  $\varphi$  on  $\Delta(S) \times \Delta(T)$  by:

$$\varphi(x, y) = \begin{cases} xGy & \text{if } (x, y) \text{ is non-absorbing } (xPy = 0), \\ (\sum s_y t A_{st} P_{st}) / xPy & \text{otherwise.} \end{cases}$$

Prove that the maxmin  $\underline{v}$  of the game exists and equals  $\max_x \min_y \varphi(x, y)$ .

HINT. To prove that player I cannot get more consider a stage by stage best reply of player II to the conditional strategy of I given non-absorbing past and use the following property:

$$\forall \delta > 0, \exists \varepsilon_0, N: \forall \varepsilon \leq \varepsilon_0, \forall x \exists t: [\varphi(x, t) \leq \underline{v} + \delta \Rightarrow xP_t \leq \varepsilon] \implies [xP_t \leq N\varepsilon \wedge xA_t \leq \underline{v} + 2\delta]$$

**16.** Let  $\Omega = [0, 1]^{\mathbb{N}}$ ,  $X_n$  the  $n^{\text{th}}$  projection and  $\mathcal{F}_n = \sigma(X_1, \dots, X_{n-1})$ . Prove that for all  $t$  in  $[0, 1]$  and all  $\varepsilon > 0$  there exists  $N$  and a randomised stopping time  $\theta$  on  $(\Omega, (\mathcal{F}_n))$  (i.e.  $A_n(\omega) = \Pr(\theta(\omega) \leq n)$  is  $\mathcal{F}_n$ -measurable and  $0 \leq A_n \leq A_{n+1} \leq 1$ ) satisfying

- (1)  $\forall n, \sum_{m \leq n} (A_m - A_{m-1})(X_m - t) \geq -\varepsilon$
- (2)  $\forall n \geq N, \frac{1}{n} \sum_{m=1}^n X_m \geq t + \varepsilon \Rightarrow A_n \geq 1 - \varepsilon$ .

HINT. Obtain  $\theta$  as an  $\varepsilon$ -optimal strategy in a stochastic game.

**17. Stochastic games as normal forms for general repeated games.** Show that sect. 2 p. 153 and 3 p. 156 reduce in effect the asymptotic analysis of a general repeated game to that of a “deterministic stochastic” game, i.e. a stochastic game with perfect information and no moves of nature.

**18. Upper analytic pay-off functions.** (Nowak, 1985b) We use the notations of sect. 1 and of App. A. Assume  $\Omega$ ,  $S$ ,  $T$ ,  $\bar{S} \subseteq \Omega \times S$ ,  $\bar{T} \subseteq \Omega \times T$  and  $C = \{(\omega, s, t) \mid \omega \in \Omega, (\omega, s) \in \bar{S}, (\omega, t) \in \bar{T}\}$  are Borel subsets of polish spaces.  $g$  and  $P(A \mid \cdot)$  are measurable on  $C$  and  $g$  is bounded.

Moreover  $\bar{T}_\omega = \{t \mid (\omega, t) \in \bar{T}\}$  is compact;  $g(\omega, s, \cdot)$  is l.s.c. and  $P(A \mid \omega, s, \cdot)$  is continuous on  $\bar{T}_\omega$ , for all  $\omega$ . Let  $B = \{f: \Omega \rightarrow \mathbb{R} \mid f \text{ bounded and } \forall x \in \mathbb{R} \{f > x\} \in \mathcal{A}\}$  denote the set of bounded upper analytic functions on  $\Omega$ ; we write  $\mathcal{A}$  for analytic sets and  $\mathcal{B}$  for Borel sets.

Then the discounted game has a value (which belongs to  $B$ ); player I has an  $\varepsilon$ -optimal  $\mathcal{A}_{(\sigma,c)}$  measurable strategy and player II an optimal  $\mathcal{A}_{(\sigma,c),s,(\sigma,c)}$  measurable strategy.

- a. Show that, given  $f$  in  $B$ ,  $\Gamma(f)_\omega$  has a value  $\Psi(f)_\omega$  and that  $\Psi(f) \in B$ .

HINT. • Prop. 1.17 p. 7 applied to  $h(\omega, s, t)$  on  $\bar{S}_\omega \times \bar{T}_\omega$  implies the existence of a value  $V(\omega)$ .  
• Let

$$\begin{aligned}\mathcal{M} &= \{(\omega, \mu) \mid \omega \in \Omega, \mu \in \Delta(\bar{S}_\omega)\} \\ \mathcal{N} &= \{(\omega, \nu) \mid \omega \in \Omega, \nu \in \Delta(\bar{T}_\omega)\} \quad \text{and} \\ \mathcal{C} &= \{(\omega, \mu, \nu) \mid \omega \in \Omega, \mu \in \mathcal{M}_\omega, \nu \in \mathcal{N}_\omega\}\end{aligned}$$

Note that these sets are Borel subsets of polish spaces too (9.c p. 429), and that  $\mathcal{N}_\omega$  is compact for all  $\omega$ . Also  $H(\omega, \mu, \nu)$  is upper analytic on  $\mathcal{C}$  (9.f p. 429) and  $H(\omega, \mu, \cdot)$  l.s.c. on  $\mathcal{N}_\omega$ . Recall that  $F(\omega, \mu) = \inf_{t \in \bar{T}_\omega} H(\omega, \mu, t)$  and  $V(\omega) = \sup_{\mu \in \mathcal{M}_\omega} F(\omega, \mu)$ . Hence to obtain  $V \in B$  and the existence of  $\varepsilon$ -optimal  $\mathcal{A}_{(\sigma,c)}$  measurable strategies for I it is enough to show that  $F$  is upper analytic on  $\mathcal{M}$ . Let  $g_n(\omega, s, t) = \inf_{t' \in \bar{T}_\omega} [g(\omega, s, t') + nd(t, t')]$ , where  $d$  is the metric on  $T$ .  $g_n$  is Borel measurable on  $C$  (8.a p. 427) and continuous in  $t$ .  $g_n$  increases to  $g$ , hence  $H_n$  (defined as  $H$  but with  $g_n$ ) increases to  $H$  and by compactness of  $\bar{T}_\omega$ ,  $F_n$  increases to  $F$ . It remains to prove the property for  $F_n$ . Let  $\zeta_k$  be a dense family of Borel selections of  $\bar{T}$  (8.b p. 428). Then  $F_n(\omega, \mu) = \inf_k H_n(\omega, \mu, \zeta_k(\omega))$  and the result follows (7.j p. 427).

- Finally for II's strategy, it suffices to show that  $G(\omega, \nu)$  is  $\mathcal{A}_{(\sigma,c)}^\Omega \otimes \mathcal{B}^{\Delta(T)}$  measurable. In fact then  $\{(\omega, \nu) \mid G(\omega, \nu) = V(\omega)\}$  is also  $\mathcal{A}_{(\sigma,c)}^\Omega \otimes \mathcal{B}^{\Delta(T)}$  measurable. Apply then (7.j p. 427). Using  $g_n$  as above, it is enough to prove the property for  $G_n$ , which is upper analytic. Let then  $\nu_k$  be a dense family of Borel selection of  $\mathcal{N}$ , hence  $f_k(\omega) = G_n(\omega, \nu_k(\omega))$  is upper analytic thus  $\mathcal{A}_{(\sigma,c)}$ -measurable. Finally  $G_n$  being  $n$ -Lipschitz,  $G_n(\omega, \nu) = \sup_k [f_k(\omega) - (n+1)d'(\nu_k, \nu)]$  is  $\mathcal{A}_{(\sigma,c)}^\Omega \otimes \mathcal{B}^{\Delta(T)}$  measurable ( $d'$  being the distance on  $\Delta(T)$  defining the weak topology).

- b. Use then lemmas 1.2 and 1.3 (p. 328) on  $B$ .

**19. An operator solution for stochastic games with limsup pay-off.** (Maitra and Sudderth, 1992) Consider a two-person zero-sum stochastic game with  $\Omega$  countable,  $S$  and  $T$  finite and  $0 \leq g \leq 1$ . The pay-off of the infinitely repeated game is defined on each history by  $g^* = \limsup \bar{g}_n$ .

Then the game has a value.

- a. Reduce the problem to the case where first  $g_n$  depends only on  $\omega_n$  and moreover  $g^* = \limsup g_n$ . It will be convenient to start from  $n = 0$  hence let  $H_n = (\Omega \times S \times T)^n \times \Omega$ .

b. A stop rule  $\theta$  is a stopping time on  $(H_\infty, \mathcal{H}_\infty)$  everywhere finite. To every stop rule  $\theta$  is associated an ordinal number called the index,  $\alpha(\theta)$  and defined inductively by:  $\alpha(0) = 0$  and for  $\theta \neq 0$   $\alpha(\theta) = \sup \{ \alpha(\theta[h_1]) + 1 \mid h_1 \in H_1 \}$  where  $\theta[h_1]$  is  $\theta$  after  $h_1$ , i.e.  $\theta[h_1](h_\infty) = \theta(h_1, h_\infty) - 1$ . Show that for any probability  $P$  on  $(H_\infty, \mathcal{H}_\infty)$  and any function  $u \in \mathcal{U} = \{u: \Omega \rightarrow [0, 1]\}$ , one has:

$$(\star) \quad \int u^* dP = \inf_\theta \sup_{\zeta \geq \theta} \int u(\omega_\zeta) dP$$

c. For every  $u$  and  $n$  we define an auxiliary game  $\Gamma^n(u)$  — leavable for player I after move  $n$  — where in addition to  $\sigma$  I chooses a stop rule  $\theta \geq n$  and the pay-off is  $g_\theta = g(\omega_\theta)$ .

Define an operator  $\Phi$  on  $\mathcal{U}$  by:  $\Phi u(\omega_0) = \text{val}(\mathbf{E}_{\sigma,\tau} u(\omega_1))$  (corresponding to the value of the one-shot game  $\mathcal{G}u$ , i.e. where  $\theta \equiv 1$ ) and a sequence:

$$U_0 = u, \quad U_{n+1} = u \vee \Phi U_n, \quad U = \sup U_n.$$

Deduce that  $U$  is the least function in  $\mathcal{U}$  with  $U \geq u \vee \Phi U$  and that  $U = u \vee \Phi U$ . Prove that the value of  $\Gamma^0(u)$  is  $U$ .

HINT. For player II, prove by induction on  $\alpha(\theta)$  that  $\mathbf{E}_{\sigma,\tau}^\omega u(\omega_\theta) \leq U(\omega)$ , by letting him play at each stage  $n$  optimally in  $\mathcal{G}u$  starting from  $\omega_n$ . Deduce that the value  $\Psi u$  of  $\Gamma^1(u)$  is  $\Phi U$ .

d. The idea is now to approximate in some sense the original game  $\Gamma$  by leavable games  $\Gamma^n$ , with  $n \rightarrow \infty$ , using  $(\star)$ . Define  $Q_0 = \Psi g$  and for each countable ordinal  $Q_\xi = \Psi(g \wedge \inf_{\eta < \xi} Q_\eta)$ , then  $Q = \inf Q_\xi$  and show that  $Q = \Psi(g \wedge Q)$ .

e. Prove that player I can guarantee any  $w \in \mathcal{U}$  satisfying  $w \leq \Psi(g \wedge w)$ .

HINT. Let him play a sequence of  $\varepsilon_n$  strategies in  $\Gamma^1(g \wedge w)$  inducing a strictly increasing sequence  $\theta_n$  of stopping rules with  $\mathbf{E}_{\sigma,\tau}((g \wedge w)(\omega_{\theta_n})) \geq w(\omega_0) - \varepsilon$ .

f. Deduce finally that  $Q$  is the value of  $\Gamma$  by showing by induction that player II can guarantee any  $Q_\xi$ . In particular  $Q$  is the largest solution of  $w = \Psi(g \wedge w)$ .

HINT. Assume that he can guarantee  $\tilde{Q} = \inf_{\eta < \xi} Q_\eta$  with some  $\tilde{\tau}$ . Let  $\tau^*$  be optimal in  $\Gamma^1(g \wedge \tilde{Q})$  and  $m = \inf\{n \mid g(\omega_n) > \tilde{Q}(\omega_n)\}$ . Player II plays  $\tau^*$  up to  $m$  and then  $\tilde{\tau}$ . Show that for any  $\varepsilon > 0$  there exists a stopping rule  $\zeta$  such that  $\theta \geq \zeta$  implies  $\mathbf{E}_{\sigma,\tau} g(\omega_\theta) \leq \mathbf{E}_{\sigma,\tau^*}((g \wedge \tilde{Q})(\omega_{\theta \wedge m})) + \varepsilon$ .

g. The proof of the preceding point could be adapted to show that the minmax of  $\Gamma$ , say  $\bar{V}$  satisfies  $\bar{V} \leq \Psi(g \wedge \bar{V})$ , hence an alternative proof, without using  $Q$ . (Compare with prop. 2.8 p. 83 and ex. II.2Ex.10 p. 88).

h. (Maitra and Sudderth, 1993) The result (assuming the reduction in VIIEx.19a done) extends to a much more general set up:  $\Omega, S, T$  are Borel subsets of polish spaces (the last two can even vary in a measurable way with  $\omega$ ).  $T$  is compact, the transition is continuous in  $t$  and  $g$  is bounded and upper analytic.

HINT. The idea of VIIEx.19g cannot be used since measurability conditions on  $\bar{V}$  are not present. The proof basically follows VIIEx.19d, VIIEx.19e, VIIEx.19f but is much more delicate. First prove that  $U$  is upper analytic (ex. VIIEx.18). Let  $B = \{(\omega, x) \mid g(\omega) > x\}$  and for any real function  $w$  on  $\Omega$  denote by  $E(w)$  its epigraph  $= \{(\omega, x) \mid w(\omega) \leq x\}$ . Given  $C \subseteq \Omega \times [0, 1]$ , define a function  $\chi_C$  on  $\Omega$  by  $\chi_C(\omega) = \sup\{x \in [0, 1] \mid (\omega, x) \in C^c \cap B\}$  so that  $\chi_{E(w)} = u \wedge w$ . One defines now  $\Theta$  from subsets of  $\Omega \times [0, 1]$  to subsets of  $\Omega \times [0, 1]$  by:

$$\Theta(C) = \{(\omega, x) \in \Omega \times [0, 1] \mid [\Psi(\chi_C)](\omega) \leq x\}.$$

Show that  $\Theta$  preserves coanalytic sets and that  $\Theta(E(w)) = E(\Psi(u \wedge v))$ . Moschovakis' theorem (8.c p. 428) or rather its dual then implies, letting  $\Theta^0 = \Theta(\emptyset)$ , that  $\Theta^\infty = \bigcup_{\xi < \omega_1} \Theta^\xi = E(Q) = E(\Psi(u \wedge Q))$  and that  $Q$  is upper analytic. Prove then that  $Q$  is the value and that both players have universally measurable  $\varepsilon$ -optimal strategies.

**20. Solvable states.** (Vieille, 1993) Consider a finite two-person non-zero-sum stochastic game. Define a subset  $R$  of states to be solvable if there exist  $(x, y) \in \Delta(S) \times \Delta(T)$  such that:

- $R$  is ergodic with respect to  $P$  and  $(x, y)$
- $\forall \omega \in R, \forall t \in T, \mathbf{E}_{x,t}^\omega(v_\infty^{\text{II}}) \leq \gamma_R^{\text{II}}(x, y)$ , (and similarly for player I), where  $\gamma_R(x, y)$  is the asymptotic average pay-off on  $R$ .

a. Prove the existence of solvable states, i.e. non-empty solvable subsets.

HINT. Let  $x$  be a limit of optimal strategies of player I for  $v_\lambda^{\text{I}}$  and let  $y$  be a best reply. Show that there exists an ergodic class  $R$  with  $\gamma_R^{\text{I}}(x, y) \geq v_\infty^{\text{I}}$ , where moreover  $R$  is included in the subset of states where  $v_\infty^{\text{I}}$  is maximal.

- b. Show that, if the initial state is solvable, an equilibrium exists.
- c. Show that if  $\#\Omega \leq 3$ , the game has an equilibrium.

## 5. Reminder on dynamic programming

A dynamic programming problem is a one-player stochastic game specified by a state space  $\Omega$ , an action space  $S$ , a set of constraints  $A \subseteq \Omega \times S$  and a transition  $p$  from  $\Omega \times S$  to  $\Omega$ .  $\Omega$  and  $S$  are standard Borel (App.6 p. 426),  $A$  is a Blackwell space (*ibidem*) with  $\text{Proj}_\Omega A = \Omega$  and  $p$  is a Borel transition probability. Let  $H_n = (\Omega \times S)^{n-1} \times \Omega$ ,  $H = \bigcup_{n \geq 1} H_n$ . Strategies are  $\mathcal{U}$ -measurable transitions from  $H$  to  $S$  with  $\sigma(\omega_1, s_1, \dots, s_{n-1}, \omega_n)(A_{\omega_n}) = 1$ , where  $A_{\omega_n} = \{s \mid (\omega_n, s) \in A\}$  and  $\mathcal{U}$  is either the  $\sigma$ -field  $\mathcal{B}_u$  (4.d.1 p. 424) of universally measurable sets ( $\mathcal{B}$  are the Borel sets on  $H$ ) or  $\mathcal{B}_{(s,c)}$  (App.1 p. 421).  $M^*$  (resp.  $M$ , resp.  $SM$ ) is the set of Markovian (resp. pure Markovian, resp. pure stationary Markovian) strategies i.e. satisfying  $\sigma(\omega_1, \dots, \omega_n) = \sigma_n(\omega_n)$  (resp. moreover  $\sigma_n$  pure, resp. moreover  $\sigma_n (= x)$  pure and independent of  $n$ ). It follows that given a  $\mathcal{U}$ -measurable mapping  $f$  (bounded above or below) defined on plays ( $= H_\infty = (\Omega \times S)^\mathbb{N}$ ), for any strategy  $\sigma$  and initial probability  $q \in \Delta(\Omega)$ ,  $\int f d(\sigma \otimes q) = \varphi^q(\sigma)$  is well defined and  $\omega \mapsto \varphi^\omega(\sigma)$  is  $\mathcal{U}$ -measurable. In particular if  $f$  stands for the pay-off  $\tilde{g}$  on plays, we obtain a pay-off function  $\gamma$  on  $\Omega \times \Sigma$ . The value is then  $V^\omega = \sup_\sigma \gamma^\omega(\sigma)$  and given  $\varepsilon \geq 0$ ,  $\sigma$  is  $\varepsilon$ -optimal at  $\omega$  if

$$\gamma^\omega(\sigma) \geq \begin{cases} (V^\omega - \varepsilon) & \text{if } V^\omega < \infty \\ 1/\varepsilon & \text{if } V^\omega = +\infty \end{cases}$$

and  $\sigma$  is  $\varepsilon$ -optimal if it is  $\varepsilon$ -optimal at every  $\omega$ .

We will consider the case where  $\tilde{g} = \lim g_n$  with  $g_n$  defined on  $(\Omega \times S)^n$  and upper analytic (i.e.  $\{g_n > t\}$  (or  $\{g_n \geq t\}$ ) a Blackwell space for all  $t$ ). More precisely let  $g$  be an upper analytic bounded function on  $(\Omega \times S)$  and consider the following cases:<sup>1</sup>

- (D) Discounted Case:  $g_n(\omega_1, \dots, s_n) = \sum_{i=1}^n g(\omega_i, s_i)(1-\lambda)^{i-1}$  (Blackwell, 1965)
- (P) Positive Case:  $g \geq 0$  and  $g_n(\omega_1, \dots, s_n) = \sum_{i=1}^n g(\omega_i, s_i)$  (Blackwell, 1967b)
- (N) Negative Case:  $g \leq 0$  and  $g_n(\omega_1, \dots, s_n) = \sum_{i=1}^n g(\omega_i, s_i)$  (Strauch, 1966)

### Exercises.

#### 1. General properties.

- a. Prove that for any  $\omega$  in  $\Omega$ ,  $\sigma$  in  $\Sigma$ , there exists  $\tau$  in  $M^*$  with  $\gamma^\omega(\sigma) = \gamma^\omega(\tau)$ .

HINT. Denote by  $\rho_n^\omega(\sigma)$  and  $\zeta_n^\omega(\sigma)$  the probabilities induced by  $\sigma$  starting from  $\omega$  on the  $n^{\text{th}}$ -factor  $\Omega$  and  $\Omega \times S$ . Use the von Neumann selection theorem (7.j p. 427) to define  $\tau$  in  $M^*$  such that for all  $B$  Borel in  $\Omega \times S$ :

$$\zeta_n^\omega(\sigma)(B) = \int \tau_n(\tilde{\omega})(B_{\tilde{\omega}}) \rho_n^\omega(\sigma)(d\tilde{\omega}).$$

- b. Given  $\pi$  in  $\Delta(\Omega \times S)$  let  $\hat{\pi}$  be the marginal on  $\Omega$  and define  $C \subseteq \Delta(\Omega) \times \Delta(\Omega \times S)$  by  $C = \{(q, \pi) \mid \hat{\pi} = q, \pi(A) = 1\}$ . Let  $\mathcal{M}$  be the family of sequences  $\{\mu_n\}$ ,  $\mu_n: \Delta(\Omega) \rightarrow \Delta(\Omega \times S)$  with graph in  $C$  and define  $\theta: \Delta(\Omega \times S) \rightarrow \Delta(\Omega)$  by  $\theta(\pi)(K) = \int p(K \mid \omega, s)\pi(d\omega, ds)$  for all Borel  $K$  in  $\Omega$ . Thus given an initial  $q$  on  $\Delta(\Omega)$ ,  $\mu$  in  $\mathcal{M}$  defines (through  $\theta$ ) a distribution on plays. Prove that, given  $\sigma$  in  $M^*$  there exists  $\mu$  in  $\mathcal{M}$  inducing the same  $\zeta_n^\omega$  (for all  $\omega$ ); given  $q$  in  $\Delta(\Omega)$  and  $\mu$  in  $\mathcal{M}$  there exists  $\sigma$  in  $M^*$  inducing the same  $\zeta_n^q$ .

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<sup>1</sup>Other sources include (Blackwell, 1970; Dubins and Savage, 1965; Shreve and Bertsekas, 1979).

c. Deduce that  $V^\omega = \sup_{\mu \in \mathcal{M}} \gamma^\omega(\mu)$ . Denote by  $\Pi$  the set of probabilities on plays that are generated by some  $q$  in  $\Delta(\Omega)$  and  $\mu$  in  $\mathcal{M}$ . Show that  $\Pi$  is Blackwell.

HINT.  $\{\pi \in \Delta(\Omega \times S) \mid \pi(A) = 1\}$  is Blackwell.

Deduce that  $(q, \mu) \mapsto \gamma^q(\mu)$  is upper analytic. Use the von Neumann selection theorem to show that  $\omega \mapsto V^\omega$  is upper analytic (thus  $\mathcal{U}$ -measurable).

d. Given a  $\mathcal{U}$ -transition  $y$  from  $\Omega$  to  $S$ , define  $\Phi_y$  on bounded or positive  $\mathcal{U}$ -measurable mappings  $f$  on  $\Omega$  by:

$$\Phi_y f^\omega = \int (g(\omega, s) + \int f(\tilde{\omega}) p(d\tilde{\omega} \mid \omega, s)) y(\omega)(ds)$$

and

$$\Psi f^\omega = \sup_{s \in A_\omega} \{g(\omega, s) + \int f(\tilde{\omega}) p(d\tilde{\omega} \mid \omega, s)\}$$

Show that:

$$\begin{aligned} (1) & \quad (\text{D}) & \Psi(1 - \lambda)V = V \\ (2) & \quad (\text{N}), (\text{P}) & \Psi V = V \end{aligned}$$

## 2. Existence of $\varepsilon$ -optimal strategies.

a. (D) Prove that (1) has a unique solution.

(P)  $W \geq 0$  and  $\mathcal{U}$ -measurable and  $W \geq \Psi W \Rightarrow W \geq V$

(N)  $W \leq 0$  and  $\mathcal{U}$ -measurable and  $W \leq \Psi W \Rightarrow W \leq V$

HINT. Let  $\varepsilon = \sum \varepsilon_i$  and  $y_i$  a  $\mathcal{U}$ -measurable mapping from  $\Omega$  to  $S$  with:

$$W - \sum_{j \leq i} \varepsilon_j \leq \Phi_{y_i}(\prod_{j < i} \Phi_{y_j})W$$

b. (D) (N) Given  $\sigma$  in  $M^*$ , there exists  $\tau$  in  $M$  with  $\gamma(\sigma) = \gamma(\tau)$ .

HINT. Purify inductively on histories with increasing length.

There exists  $\sigma$  in  $M$   $\varepsilon$ -optimal.

HINT. Use (1) or (2) and point VII.5Ex.2a.

If there exists  $\sigma$  optimal, then there exists  $\tau$  in  $SM$  optimal.

(D) If  $\sigma$  is  $\varepsilon$ -optimal, there exists  $\tau$  in  $SM$ ,  $\varepsilon/\lambda$  optimal (cf. lemma 1.3 p. 328).

c.  $x$  in  $SM$  is optimal implies  $V = \Phi_x V$  and  $\gamma(x) = \Psi\gamma(x)$ .

(D) (N)  $V = \Phi_x V \Rightarrow x$  optimal.

(D) (P)  $\gamma(x) = \Psi\gamma(x) \Rightarrow x$  optimal.

d. Let  $\Psi^\infty(0) = \lim_n \Psi^n(0)$ .

(P) (D)  $\Psi^\infty(0) = V$

(N)  $\Psi^\infty(0) \geq V$  and  $\Psi^\infty(0) = V \iff \Psi(\Psi^\infty(0)) = \Psi^\infty(0)$ .

e. (P) There exists  $\sigma$   $\varepsilon$ -optimal.

HINT. Use VII.5Ex.2d.

## 3. Special cases.

a.  $S$  finite. (D) (N) There exists  $\sigma$  in  $SM$  optimal.

b. (N) No  $\varepsilon$ -optimality in  $SM$  (compare VII.5Ex.2.VII.5Ex.2b): Take  $\Omega = \{0, 1\}, S = \mathbb{N}$ . State 0 is absorbing with (total) pay-off  $-1$ . In state 1 the pay-off is 0 and move  $n$  leads to state 0 with probability  $1/n$ .

$\Psi^\infty(0) \neq V$  (compare VII.5Ex.2d): Take  $\Omega = S = \mathbb{N}$ . The transition is deterministic from state  $n$  to  $n - 1$  with pay-off 0 until state 1 absorbing with (total) pay-off  $-1$ . In state 0 the pay-off is  $-1$  and move  $n$  leads to state  $n$ .

#### 4. (P) ; $\Omega$ countable and $V < +\infty$ .

a.  $\varepsilon$ -optimality. (Ornstein, 1969) Prove that:  $\forall \varepsilon > 0, \exists x \in SM$  with  $\gamma(x) \geq (1 - \varepsilon)V$ .

HINT.

- Prove the result for  $\Omega$  finite, using case (D).

- Given  $K \subseteq \Omega$  denote by  $\gamma_K$  the pay-off in the model where all states in  $\complement K$  are absorbing with zero pay-off. Prove that given  $\omega$  and  $\eta > 0$  there exist  $x$  in  $SM$  and  $K$  finite with  $\gamma_K^\omega(x) > (1 - \eta)V^\omega$ .
- Let  $\delta > 0$  and  $L = \{\omega' \in K \mid V^{\omega'} \geq (1 + \delta)\gamma^{\omega'}(x)\}$ . Prove that  $\gamma_{K \setminus L}^\omega(x) \geq (1 - \eta - 2\eta/\delta)V^\omega$  (decompose the pay-off induced by  $x$  at  $\omega$  before and after the entrance in  $L$ ).
- Given  $\varepsilon > 0$  let  $\delta$  and  $\eta$  satisfy  $1/(1+\delta) \leq 1-\varepsilon/2$  and  $\eta+2\eta/\delta \leq \varepsilon$ . Let  $\mathcal{S}$  be the set of strategies that coincide with  $x$  when the current state is in  $K \setminus L$ . Show that  $\sup_{\sigma \in \mathcal{S}} \gamma^\omega(\sigma) \geq (1 - \varepsilon)V^\omega$ .
- Modify the initial model by imposing the transition induced by  $x$  on  $K \setminus L$ . The new value is at least  $(1 - \varepsilon)$  the old one and any strategy in  $SM$  will give  $(1 - \varepsilon)$  when starting from  $\omega$ .
- Enumerate the elements in  $\Omega$  and repeat the same procedure.

b. Need for  $V$  finite in VII.5Ex.4a. Let  $\Omega = \mathbb{Z} \cup \{\delta\}, S = \mathbb{N}$ . At state 0 move  $n$  leads to state  $n$  then there is a deterministic translation to  $n - 1$  with pay-off 1. On the states in  $-\mathbb{N}$ , the pay-off is 0 and one goes from  $-n$  to 0 with probability  $1/n$  and to the absorbing state  $\delta$  (with pay-off 0) with probability  $1 - 1/n$ . Obviously  $V^\omega = +\infty$  on  $\mathbb{N}$ , hence on  $\mathbb{Z}$  but  $\gamma^n(x) \xrightarrow{n \rightarrow \infty} 0$  for all  $x$  in  $SM$ . ???

c. No additive error term in VII.5Ex.4a. Let  $\Omega = L \cup M, L$  and  $M$  countable and  $S = \{1, 2\}$ . From  $l_n$  move 1 leads to  $l_{n+1}$  (with probability  $1/2$ ) or  $m_0$  (probability  $1/2$ ) with pay-off 0. Move 2 leads to  $m_{2n-1}$ . The pay-off is 1 on  $M$  with deterministic transition from  $m_n$  to  $m_{n-1}$  until  $m_0$  absorbing with pay-off 0. Show that  $V^{l_n} \geq 2^n$  and that for any  $x$  in  $SM$  there exists  $\omega$  in  $L$  with  $\gamma^\omega(x) \leq V^\omega - 1$ .

d. Optimality in VII.5Ex.4a? In the framework of VII.5Ex.4a, prove that if there exists an optimal  $\sigma$ , there exists one in  $SM$ .

HINT. Show that if  $x$  satisfies  $V = \Phi_x V$  and there exists  $t > 0$  with  $\gamma(x) > tV$  then  $x$  is optimal. Use then VII.5Ex.4a.

e. No optimality in VII.5Ex.4a.  $S$  finite and no  $x$  in  $SM$  optimal: Let  $\Omega = \mathbb{N}$ . 0 is absorbing with pay-off 0. In state  $n$  move 1 leads to state  $n + 1$  with pay-off 0 and move 2 leads to state 0 with pay-off  $1 - 1/n$ .

#### 5. Average case.

Let now  $\gamma^\omega(\sigma) = \liminf E_\sigma^\omega(\sum_{i=1}^n g(\omega_i, s_i)/n)$ .

a.  $S$  finite and  $\Omega$  countable.

i. Move 1 leads from  $n$  to  $n + 1$  with pay-off 0. Move 2 leads from  $n$  to  $n$  with pay-off  $1 - 1/n$ . There exists an optimal strategy and no optimal strategy in  $SM$ .

ii. Move 1 leads from  $n$  to  $n + 1$  with pay-off 0. Move 2 leads from  $n$  to  $\delta$  absorbing with pay-off 0 with probability  $\alpha_n$  and to  $-n$  with probability  $1 - \alpha_n$ . The pay-off is 1 on  $-\mathbb{N}$  with deterministic transition from  $-n$  to  $-(n - 1)$ . Let  $\alpha = \prod_n (1 - \alpha_n) > 0$ . Then  $V^0 \geq \alpha/2$  and for all  $x$  in  $SM$ ,  $\gamma^0(x) = 0$ .

b.  $\Omega$  finite and  $S$  countable. 0 is absorbing with pay-off 0. From 1 move  $n$  leads to 0 with probability  $1/n$  and to 1 with pay-off 1 otherwise. There is no  $\varepsilon$ -optimal strategy in  $SM$ .

**6. Adaptive competitive decision.** (Rosenfeld, 1964) Consider the following game  $\Gamma$ : a finite set of  $S \times T$  matrices  $G^k, k \in K$  is given. At stage 0 one element  $k^*$  is selected according to some probability  $p$  known to the players. At each stage  $n$  both players select moves  $s_n, t_n$  and are told the corresponding pay-off  $g_n = G_{s_n t_n}^{k^*}$  and the moves. This is a very special kind of game with incomplete information (cf. ex. VEx.11 p. 260 and sect. 1 p. 273 for related results). Denote by  $v^k$  the value of  $G^k$  and by  $v$  its expectation:  $v = \sum_{k \in K} p^k v^k$ . We consider the infinite game with pay-off function up to stage  $n$   $R_n = n(\bar{g}_n - v^{k^*})$ . Take as state set the set of subsets of  $K$ .

Then there exists a constant  $R$  such that both players can guarantee  $R$  with stationary Markovian ( $SM^*$ ) strategies.

(In particular this implies that  $\bar{\gamma}_n$  converges to  $v$  with a speed of convergence of  $O(1/n)$ ).

a. Show that if  $\sigma$  is in  $SM^*$  for each  $\tau$  there exists  $\tau'$  in  $SM^*$  with  $E_{\sigma, \tau}(R_n) \leq E_{\sigma, \tau'}(R_n)$ .

b. For any non-expansive increasing map  $f: \mathbb{R} \rightarrow \mathbb{R}$  let:

$$S(f) = \begin{cases} -\infty & \text{if } f(x) < x, \forall x \\ +\infty & \text{if } f(x) > x, \forall x \\ x^* \text{ s.t. } f(x^*) = x^*, |x^*| = \min\{|x| \mid f(x) = x\} & \text{otherwise} \end{cases}$$

Prove that  $\lim f^n(0) = S(f)$ . Let  $D_1(f) = \{x \mid x < f(x) \text{ or } x = f(x) \leq 0\}$ . Let  $D_2(f) = \{x \mid x > f(x) \text{ or } x = f(x) \geq 0\}$ . Show that if  $x \in D_1(f)$ ,  $\exists N, n \geq N \Rightarrow f^n(0) \geq x$ .

c. Assume first that only one outcome is unknown. This means that  $G_{st}^k = G_{st}^{k'}$  for all  $k, k'$  and  $(s, t) \neq (1, 1)$ , (these (s,t) are non-revealing). Consider the auxiliary one-shot matrix game  $G(x)$  where  $x$  is real, with pay-off:

$$G_{11}(x) = \sum_k p^k G_{11}^k - v \quad G_{st}(x) = G_{st} - v + x \quad \text{for } (s, t) \neq (1, 1)$$

and denote by  $w(x)$  its value.

i. Prove that if  $x \in D_1(w)$  there exists a  $SM^*$  strategy  $\sigma$  of player I and  $N$  such that  $n \geq N \Rightarrow E_{\sigma, \tau}(R_n) \geq x$ , for all  $\tau$ .

ii. Show that  $\bar{D}_1(w) \cap \bar{D}_2(w) \neq \emptyset$ , (in particular  $D_1 \neq \emptyset$ ).

d. Prove the result by induction on the number of unknown outcomes. Given a revealing pay-off entry  $(s, t)$  let  $\bar{R}(s, t) = \sum_k R(\Gamma \text{ given } G_{st}^k) p^k$  and define a matrix by:

$$G_{st}(x) = \begin{cases} \sum_k p^k G_{st}^k - v + x & \text{if } (s, t) \text{ is non-revealing,} \\ \sum_k p^k G_{st}^k - v + \bar{R}(s, t) & \text{otherwise.} \end{cases}$$

Conclude by following the proof in VII.5Ex.6a.

e. Consider a finite family of games with same strategy spaces, each of which has a value and extend the result with Markovian strategies.

## Part C

# Further Developments



## CHAPTER VIII

### Extensions and Further Results

We study here extensions of the models of ch. VI and VII.

The first 3 sections deal with games with lack of information on both sides where signals and states are correlated: a first class is the symmetric case, a second one corresponds to games with no signals and a third one leads to a first approach of the full monitoring case.

The last section is devoted to a study of specific classes of stochastic games with incomplete information.

#### 1. Incomplete information: the symmetric case

In this section we consider games with incomplete information where the signalling pattern is symmetric so that at each stage both players get the same signals.

Formally we are given a finite collection of  $S \times T$  matrices  $G^k$ ,  $k \in K$ , with initial probability  $p$  in  $\Pi = \Delta(K)$ , and the players have no initial information on the true state  $k$  except  $p$ . We denote by  $A$  the common finite set of signals (the extension to a measurable setup is easy) and by  $A^k$  the corresponding signalling matrices.

Given  $k$  and a couple of moves  $(s, t)$ , a signal  $a$  is announced to both players according to the probability distribution  $A_{s,t}^k$  on  $A$ . Assuming perfect recall means in this framework that for all  $k, k'$  in  $K$ ,  $s \neq s'$  or  $t \neq t'$  implies that  $A_{s,t}^k$  and  $A_{s',t'}^{k'}$  have disjoint support.

Denoting by  $\Gamma(p)$  this game we have the following:

**THEOREM 1.1.**  $\Gamma(p)$  has a value.

**PROOF.** Define first the set of non-revealing moves — or more precisely non-revealing entries — at  $p$  by  $\text{NR}(p) = \{(s, t) \in S \times T \mid p_1 = p\}$  where  $p_1$  is the posterior probability on  $K$  if  $(s, t)$  is played, i.e. for all  $a$  such that  $\Pr_{s,t,p}(a) \equiv A_{st}^p(a) = \sum_k p^k A_{st}^k(a) > 0$ ,  $p_1(a) = p^k A_{st}^k(a)/A_{st}^p(a)$ . Given  $w$ , a bounded real function on  $\Pi$  and  $p$  in  $\Pi$ , we define a stochastic game with absorbing states  $\Gamma^*(w, p)$  (cf. 4.a p. 341) by the matrix:

$$G_{st}(w, p) = \begin{cases} \sum_k p^k G_{s,t}^k & \text{if } (s, t) \in \text{NR}(p) \\ [\mathsf{E}_{s,t,p}(w(p_1))]^* \equiv [\sum_a A_{st}^p(a) w(p_1(a))]^* & \text{otherwise} \end{cases}$$

where as usual a star  $*$  denotes an absorbing entry. By theorem 3.1 p. 334  $\Gamma^*(w, p)$  has a value, say  $Vw(p)$ .

Remark now that if  $\Gamma$  as a value  $v$  on  $\Pi$  then  $v$  is a solution of the recursive equation (cf. 3.2 p. 158):

$$(1) \quad w = Vw$$

The proof that  $\Gamma$  has a value will be done by induction on the number of states in  $K$ . We therefore assume that for all games with strictly less than  $\#K$  states the value exists (and is a solution of the corresponding equation (1)), this being clearly true for  $\#K = 1$ .

Note then that  $v$  is defined and continuous (in fact even Lipschitz) on  $\partial\Pi$ . Let us now prove some properties of the operator  $V$ .

We first remark that  $u \leq w$  on  $\Pi$  implies:

$$(2) \quad \forall p \in \Pi, \quad 0 \leq Vw(p) - Vu(p) \leq \max_{(s,t) \notin \text{NR}(p)} \{G_{st}(w,p) - G_{st}(u,p)\}$$

and that  $V$  is continuous (for the uniform norm). Then we have:

**LEMMA 1.2.** *Let  $u$  be a real continuous function on  $\Pi$  with  $u|_{\partial\Pi} = v$ . Then  $Vu$  is continuous on  $\Pi$  and coincides with  $v$  on  $\partial\Pi$ .*

**PROOF.** It is clear that  $p \in \partial\Pi$  implies  $p_1 \in \partial\Pi$  hence  $G(u,p)$  and  $G(v,p)$  coincide on  $\partial\Pi$  so that  $Vu(p) = v(p)$  on  $\partial\Pi$ . In particular  $Vu(p)$  is continuous there.

Now on  $\Pi \setminus \partial\Pi$ ,  $\text{NR}(p)$  is constant, equal say to  $\text{NR}$ , hence  $Vu$  is again continuous.

Finally if  $p \in \partial\Pi$  with  $\text{NR} \subsetneq \text{NR}(p)$  note first that if  $p^m \rightarrow p$ , with  $p^m \in \Pi \setminus \partial\Pi$  then  $p_1^m(s,t) \rightarrow p_1(s,t) = p$  for  $(s,t)$  in  $\text{NR}(p) \setminus \text{NR}$ . Hence  $G_{st}(u,p^m) \rightarrow E_{s,t,p}(u(p_1)) = u(p) = v(p)$ . Note now that replacing in a game with absorbing pay-off a non-absorbing entry by an absorbing one, equal to the value, does not change the value. So that the value of the new game  $\Gamma'(u,p)$ , where the pay-off is  $v(p)$  for  $(s,t)$  in  $\text{NR}(p) \setminus \text{NR}$  is still  $v(p)$ , hence the continuity. ■

**LEMMA 1.3.** *If player I can guarantee  $u$  in  $\Gamma$ , he can also guarantee  $Vu$ .*

**PROOF.** Let player I use in  $\Gamma(p)$  an  $\varepsilon$ -optimal strategy in  $\Gamma^*(u,p)$  and switch to a strategy that  $\varepsilon$ -guarantees  $u$  in  $\Gamma$ , at the current posterior probability, as soon as one absorbing entry is reached.

Formally, denote by  $g^*$  the pay-off in  $\Gamma^*(u)$ , then there exists  $\sigma^*$  and  $N_0$  such that:

$$E_{\sigma,\tau,p}(\bar{g}_n^*) \geq Vu(p) - \varepsilon, \quad \forall \tau, \forall n \geq N_0$$

But for every  $p$  there exists  $\sigma(p)$  and  $N(p)$  such that:

$$E_{\sigma(p),\tau,p}(\bar{g}_n) \geq u(p) - \varepsilon, \quad \forall \tau, \forall n \geq N(p)$$

We define now  $\theta$  as the entrance time to an absorbing pay-off in  $\Gamma^*$  and  $\sigma$  as:  $\sigma^*$  until  $\theta$  and  $\sigma(p_\theta)$  thereafter. Then:

$$\bar{g}_n = \frac{1}{n} \left\{ \left[ \sum_{m=1}^{\theta \wedge n} g_m^* + (n - \theta \wedge n)u(p_\theta) \right] + \left[ \sum_{m=\theta \wedge n+1}^n g_m - (n - \theta \wedge n)u(p_\theta) \right] \right\}$$

Let  $N_1$  be an upper bound for  $N(p_1)$ . Then for  $n \geq N_0 + (C/\varepsilon)N_1$ , we obtain:

$$(3) \quad E_{\sigma,\tau,p}(\bar{g}_n) \geq Vu(p) - \varepsilon - \varepsilon. \quad \blacksquare$$

We can now prove the theorem.

Player I can obviously guarantee  $u_0(p) = \sup_{q \in \partial\Pi} \{v(q) - C \|p - q\|_1\}$ , which is Lipschitz on  $\Pi$ . Hence by lemma 1.3 p.360 he can guarantee any  $u_n$ ,  $n \geq 0$ , with  $u_{n+1} = \max\{u_n, Vu_n\}$ .  $u_n$  is continuous by lemma 1.2 p.360, hence  $u \equiv \lim \uparrow u_n$  is l.s.c. Player I can guarantee  $u$  that coincides with  $v$  on  $\partial\Pi$  and satisfies  $u \geq Vu$ . Define similarly  $w_0, w_n$  and  $w$ . Obviously  $w \geq u$  on  $\Pi$ . Let us prove the equality, hence the result. Assume not and let:

$$D = \{p \in \Pi \mid w(p) - u(p) = d = \sup_{q \in P} (w(q) - u(q)) > 0\}$$

and note that  $D$  is compact. Consider now some extreme point  $p^*$  of its convex hull. We have then:  $Vw(p^*) - Vu(p^*) \geq d$ . Remark that  $p^* \notin \partial\Pi$  and that  $\text{NR} = S \times T$  implies

$d = 0$  by (2). Hence let  $(s, t) \notin \text{NR}$  and note that the support of the corresponding  $p_1^*$  is not included in  $D$ . Obviously then  $E\{w(p_1) - u(p_1)\} < d$ , hence a contradiction by (2) again. ■

The same proof shows uniqueness:

COROLLARY 1.4. *Let  $u'$  and  $u''$  be two continuous solutions of (1). Then they coincide.*

PROOF. By induction both functions coincide on  $\partial\Pi$ . Let  $u_1 = \min(u', u'')$ ,  $u_2 = \max(u', u'')$ . Then  $u_1 \geq Vu_1$  and  $u_2 \leq Vu_2$ .  $u_1$  and  $u_2$  are continuous and coincide on  $\partial\Pi$ , hence everywhere by the previous proof. ■

## 2. Games with no signals

**2.a. Presentation.** These games were introduced in (Mertens and Zamir, 1976b) under the name “repeated games without a recursive structure”. According to our analysis in ch. IV, the meaning is that the support of the associated consistent probabilities cannot be bounded, hence no analysis based on finite dimensional state variables, like in ch. V and VI or in the previous section, can be achieved.

The description of the game consists again of a finite collection of  $S \times T$  pay-off matrices  $G^k$ ,  $k$  in  $K$ , with some initial probability  $p$  on  $\Pi$  and none of the players is informed upon the initial state. The transition probability on the signals is again defined by a family of matrices  $A^k$ , but where for all  $(s, t)$  in  $S \times T$ ,  $A_{st}^k$  is deterministic, with support in  $\{0, k\}$  (the extension to random signals is simple): either both players receive a “blank” signal (“0”) or the state is completely revealed. We can thus assume in the second case that the pay-off is equal to the value of the revealed game and absorbing from this time on. It is then sufficient to define the strategies on the “blank” histories, hence the name “game with no signals”.

It will be convenient to assume the value of each  $G^k$  to be zero (we subtract the expectation of the value of  $G^k$  from the pay-offs) and to multiply  $G^k$  by  $p^k$ , so that the expected pay-off will be the sum of the conditional expected pay-off, given  $k$  in  $K$ .  $p$  being fixed, we will write  $\Gamma_\infty$  for  $\Gamma_\infty(p)$ .

The analysis of the game will be done as follows:

- We first construct an auxiliary one-shot game  $\bar{G}$  in normal form; basically we consider a class of strategies that mimick some strategies in  $\Gamma_\infty$ , and we define the pay-off in  $\bar{G}$  as the corresponding asymptotic pay-off in  $\Gamma_\infty$ . We prove that  $\bar{G}$  has a value  $v(\bar{G})$ .
- Then we show that the minmax in  $\Gamma_\infty$  exists and equals  $v(\bar{G})$ : in fact the class chosen above is sufficient in the sense that none of the players can do better, for the minmax, than using a strategy in this class.
- Since there are games with no value we consider finally  $v_n$  and  $v_\lambda$  and prove that they converge.

**2.b. An auxiliary game.** We define the one-shot game  $\bar{G}$  by the following strategy sets  $\bar{X}_1$ ,  $\bar{Y}$  and pay-off function  $F$ :

$$\bar{X}_1 = \bigoplus_{S' \subseteq S} \Delta(S') \times \mathbb{N}^{S \setminus S'} \times S', \quad \bar{Y} = \bigoplus_{T' \subseteq T} \Delta(T') \times \mathbb{N}^{T \setminus T'}.$$

Given  $x$  in  $\bar{X}_1$  (resp.  $y$  in  $\bar{Y}$ ) we denote the corresponding subset  $S'$  (resp.  $T'$ ) by  $S^x$  (resp.  $T^y$ ), the first component by  $\alpha^x$  (resp.  $\beta^y$ ), the second by  $c^x$  (resp.  $d^y$ ) and the third by  $s^x$ .

We will also consider  $\alpha^x$  and  $c^x$  as defined on  $S$  with  $\alpha_s^x = 0$  for  $s \notin S^x$  and  $c_{s'}^x = 0$  for  $s' \in S^x$ . Let us denote by  $B^k$  be the subset of non-revealing entries, given  $k$ :

$$B^k = \{ (s, t) \in S \times T \mid A^k(s, t) = 0 \}$$

and by  $B_s^k$  (resp.  $B_t^k$ ) the corresponding sections:

$$B_s^k = \{ t \in T \mid (s, t) \in B^k \}, \quad B_t^k = \{ s \in S \mid (s, t) \in B^k \}.$$

We can now introduce the pay-off  $F(x, y) = \sum_k F^k(x, y)$ , with:

$$F^k(x, y) = \mathbb{1}_{S^x \times T^y \subseteq B^k} \prod_t \alpha^x(B_t^k)^{d_t^y} \prod_s \beta^y(B_s^k)^{c_s^x} (s^x G^k \beta^y)$$

The interpretation of the strategies in  $\bar{G}$  is given by the following strategies in  $\Gamma$ : For  $y$  in  $\bar{Y}$ : play  $\beta^y$  i.i.d. except for  $d(y) = \sum_t d_t^y$  exceptional moves uniformly distributed before some large stage  $N_0$ . On  $d_t^y$  of these stages player II plays the move  $t$ . For  $x$  in  $\bar{X}_1$ ,  $\alpha^x$  and  $c^x$  have similar meanings, but player I uses  $s^x$  after stage  $N_0$ . Note that  $F^k(x, y)$  can be expressed as

$$F^k = \rho^k(x, y) f^k(x, y)$$

where:

$$\rho^k(x, y) = \mathbb{1}_{S^x \times T^y \subseteq B^k} \prod_t \alpha^x(B_t^k)^{d_t^y} \prod_s \beta^y(B_s^k)^{c_s^x}$$

stands for the asymptotic probability that the game is not revealed, given  $k$ , and:

$$f^k(x, y) = s^x G^k \beta^y$$

is the asymptotic pay-off given this event. We also define:

$$X_1 = \{ x \in \bar{X}_1 \mid \alpha_s^x > 0, \forall s \in S^x \}, \quad Y = \{ y \in \bar{Y} \mid \beta_t^y > 0, \forall t \in T^y \}.$$

We begin by proving that  $\bar{G}$  has a value:

**PROPOSITION 2.1.** *The game  $\bar{G}$  has a value  $v(\bar{G})$  and both players have  $\varepsilon$ -optimal strategies with finite support on  $X_1$  and  $Y$ .*

**PROOF.** We first define a topology on  $\bar{X}_1$  and  $\bar{Y}$  for which these sets are compact. Let  $X_1^* = \bigoplus_{S' \subseteq S} \Delta(S') \times \bar{\mathbb{N}}^{S \setminus S'} \times S'$ , where  $\bar{\mathbb{N}} = \mathbb{N} \cup \{\infty\}$ , and define a mapping  $\iota$  from  $X_1^*$  to  $\bar{X}_1$  by:

$$\iota(x^*) = x \quad \text{with} \quad \begin{cases} S^x = S^{x^*} \cup \{ s \mid c_s^{x^*} = \infty \} \\ \alpha_s^x = \alpha_s^{x^*} & \text{for } s \in S^{x^*} \\ = 0 & \text{for } s \in S^x \setminus S^{x^*} \\ c_s^x = c_s^{x^*} & \text{for } s \notin S^x \\ s^x = s^{x^*} \end{cases}$$

Note that  $F(x^*, y) = F(\iota(x^*), y)$  for all  $y$  in  $\bar{Y}$ . Endow now  $\bar{X}_1$  with the strongest topology for which  $\iota$  is continuous.  $X_1^*$  being compact, so will be  $\bar{X}_1$ . A similar construction is done for  $\bar{Y}$ .

$F$  is now uniformly bounded and measurable for any product measure on the product of compacts  $\bar{X}_1 \times \bar{Y}$ . Note that  $F$  is continuous on  $\bar{X}_1$  for each  $y$  in  $Y$  and similarly on  $\bar{Y}$  for each  $x$  in  $X_1$ . Moreover, for each  $x$  in  $\bar{X}_1$  there exists a sequence  $x_m$  in  $X_1$  such that  $F(x_m, y)$  converges to  $F(x, y)$  for all  $y$  in  $\bar{Y}$  and similarly, given  $y$  in  $\bar{Y}$ .

Apply now prop. 2.5 p. 16 to get the result. ■

### 2.c. Minmax and maxmin.

THEOREM 2.2. •  $\bar{v}$  exists and equals  $v(\bar{G})$ .

- Player II has an  $\varepsilon$ -optimal strategy which is a finite mixture of i.i.d. sequences, each one associated with a finite number of exceptional moves, uniformly distributed before some stage  $N_0$ .
- Dual results hold for  $\underline{v}$ .

PROOF. The proof is divided into two parts corresponding to the conditions (i) and (ii) of Definition 1.2 p. 149.

#### PART A. Player I can defend $v(\bar{G})$

Before starting the proof let us present the main ideas: Since we are looking for a best reply, player I can wait long enough to decompose player II's strategy into moves played finitely or infinitely many times. Using then an optimal strategy in  $\bar{G}$ , he can also give a best reply to the behaviour of II "at infinity". Finally this "asymptotic measure" can be approximated by a distribution on finite times and will describe, with the initial strategy in  $\bar{G}$  the reply of I in  $\Gamma_\infty$ .

Given  $\varepsilon > 0$ , assume  $C \geq 1$  and define  $\eta = \varepsilon/9C$ . Let  $\tau$  be a strategy of player II in  $\Gamma_\infty$ . With our conventions in sec. 2.a  $\tau$  can be viewed as a probability on  $T^\infty$ . We will denote in this section this set by  $\Omega$ , with generic element  $\omega = (t_1, \dots, t_n, \dots)$  or  $\omega: \mathbb{N} \rightarrow T$ .

We first introduce some notations:

- Given  $T' \subseteq T$ , let us denote by  $\Omega_{T'}$  the elements  $\omega$  with "support" in  $T'$ , namely:

$$\Omega_{T'} = \{ \omega \in \Omega \mid T(\omega) \equiv \{ t \in T \mid \#\omega^{-1}(t) = \infty \} = T' \}$$

$\{\Omega_{T'}\}_{T' \subseteq T}$  is clearly a partition of  $\Omega$ . Let  $d_t(\omega)$  be the number of exceptional moves:  $d_t(\omega) = \#\omega^{-1}(t)$  for  $t \notin T(\omega)$ , and  $d_t(\omega) = 0$  for  $t \in T(\omega)$ .

- Given  $\chi$ ,  $\eta$ -optimal strategy of player I in  $\bar{G}$  with finite support on  $X_1$  (cf. prop. 2.1 p. 362) we define:

$$\delta = \min\{ \alpha_s^x \mid \chi(x) > 0, s \in S^x \}, \quad \text{thus } \delta > 0, \quad \text{and}$$

$$\bar{c}_s = \max(\{0\} \cup \{ c_s^x \mid \chi(x) > 0 \}), \quad \bar{c} = \sum_s \bar{c}_s$$

- Given  $N_1 \geq N_0$  in  $\mathbb{N}$  and  $T' \subseteq T$  let:

$$\Omega_{T'}^1 = \{ \omega \in \Omega_{T'} \mid \forall t \notin T', \omega^{-1}(t) \cap [N_0, +\infty) = \emptyset \}$$

$$\Omega_{T'}^2 = \{ \omega \in \Omega_{T'} \mid \forall t \in T', \#\{\omega^{-1}(t) \cap [1, N_0]\} > \ln(\eta)/\ln(1 - \delta) \}$$

$$\Omega_{T'}^3 = \{ \omega \in \Omega_{T'} \mid \forall t \in T', \#\{\omega^{-1}(t) \cap [N_0, N_1]\} > \bar{c} \}, \quad \text{and}$$

$$\Omega_{T'}^* = \bigcap_i \Omega_{T'}^i$$

We now choose  $N_0$  and  $N_1$  such that for all  $T' \subseteq T$  with  $\tau(\Omega_{T'}) > 0$ :  $\tau(\Omega_{T'}^i \mid \Omega_{T'}) > 1 - \eta$ ,  $i = 1, 2, 3$ , hence  $\tau(\Omega_{T'}^* \mid \Omega_{T'}) > 1 - 3\eta$ . This means that with high probability there will be no exceptional move after  $N_0$  and that each element in the support of  $\omega$  will appear a large number of times before  $N_0$  and between  $N_0$  and  $N_1$ .

- Given  $\chi$  and any array  $\Theta$  in  $(N_0, +\infty)^{\bar{c}}$  with components  $\{ \theta_{sc} \mid s \in S, 1 \leq c \leq \bar{c}_s \}$  we define the strategy  $\sigma(\Theta)$  of player I in  $\Gamma$  as first choosing  $x \in X_1$  with  $\chi$ , then using  $\sigma[x](\Theta)$ : play  $\alpha^x$  i.i.d. before  $N_0$ , and afterwards play  $s^x$  at each stage, except,  $\forall s \notin S^x, c_s^x$  additional moves  $s$ , played at times  $\theta_{sc}$  ( $1 \leq c \leq \bar{c}_s$ ).

- Let finally  $N_2 = N_1 \vee \max_{\Theta} \theta_{sc}$ .

In the next two lemmas  $\Theta$  is fixed and we will write  $\sigma[x]$  for  $\sigma[x](\Theta)$ .

LEMMA 2.3. Let  $\pi_1^k(\sigma[x], \omega) = \mathbb{1}_{S^x \times T(\omega) \subseteq B^k} \prod_t \alpha^x(B_t^k)^{d_t(\omega)}$  and  $P_1^k(\sigma[x], \omega) = \Pr\{k \text{ is not announced before } N_0 \mid k, \sigma, \omega\}$ . Then:

$$|\pi_1^k(\sigma[x], \omega) - P_1^k(\sigma[x], \omega)| \leq \eta, \text{ for all } \omega \text{ in } \Omega_{T(\omega)}^*.$$

PROOF. Let  $T(\omega) = T'$ .  $\omega$  being in  $\Omega_{T'}^1$  one has:

$$P_1^k(\sigma[x], \omega) = \prod_{t \in T'} \alpha^x(B_t^k)^{\#\{\omega^{-1}(t) \cap [1, N_0]\}} \prod_t \alpha^x(B_t^k)^{d_t(\omega)}.$$

Now if  $S^x \times T' \not\subseteq B^k$ , let  $t$  in  $T'$  with  $S^x \not\subseteq B_t^k$ . Since  $\omega \in \Omega_{T'}^2$  this implies:

$$\alpha^x(B_t^k)^{\#\{\omega^{-1}(t) \cap [1, N_0]\}} < (1 - \delta)^{\ln \eta / \ln(1 - \delta)} = \eta,$$

hence the result.  $\blacksquare$

LEMMA 2.4. Let  $\pi_2^k(\sigma[x], \omega) = \mathbb{1}_{T(\omega) \subseteq B_{s_x}^k} \prod_s \prod_{c=1}^{c_s^x} \mathbb{1}_{\omega(\theta_{sc}) \in B_s^k}$ ,  $P_2^k(\sigma[x], \omega, n) = \Pr\{k \text{ is not announced between } N_0 \text{ and } n \mid k, \sigma, \omega \text{ and } k \text{ is not announced before } N_0\}$ . Assume  $n > N_2$ , then:

$$\pi_2^k(\sigma[x], \omega) = P_2^k(\sigma[x], \omega, n), \text{ for all } \omega \text{ in } \Omega_{T(\omega)}^*.$$

PROOF. Let again  $T' = T(\omega)$  and recall first that  $\omega \in \Omega_{T'}^1$  implies  $\omega(m) \in T'$  for  $m \geq N_0$ . Note also that the moves of player I before and after  $N_0$  are independent hence:

$$\begin{aligned} P_2^k(\sigma[x], \omega, n) &= \Pr\{\omega(n) \in B_{s_m}^k \forall m: N_0 \leq m \leq n \mid k, \sigma, \omega\} \\ &= \prod_s \prod_{c=1}^{c_s^x} \mathbb{1}_{\omega(\theta_{sc}) \in B_s^k} \prod_{N_0 \leq m \leq n, m \notin \Theta} \mathbb{1}_{\omega(m) \in B_{s_x}^k}. \end{aligned}$$

Since  $\omega \in \Omega_{T'}^3$ , the last product is precisely  $\mathbb{1}_{T(\omega) \subseteq B_{s_x}^k}$ .  $\blacksquare$

We obtain thus, writing the pay-off as:  $\gamma_n(\sigma, \tau) = \mathsf{E}_{\chi, \tau}\{\sum_k \Pr(k \text{ not revealed before } n \mid k, \sigma[x], \omega) \mathsf{E}(G_{s_{nt_n}}^k \mid k \text{ is not revealed before } n, \sigma[x], \omega)\}$ , the following majoration:  $|\gamma_n(\sigma, \tau) - \mathsf{E}_{\chi, \tau}\left(\sum_k \pi_1^k(\sigma[x], \omega) \pi_2^k(\sigma[x], \omega, n) \sum_{t \in T(\omega)} G_{s_{xt}}^k \mathbb{1}_{\omega(n)=t}\right)| \leq C\eta + C\tau\{\bigcup_{\omega} (\Omega_{T(\omega)} \setminus \Omega_{T(\omega)}^*)\} \leq 4C\eta$ . Reintroducing explicitly  $\Theta$ , we obtain a minoration of the pay-off:

$$\gamma_n(\sigma, \tau) \geq \mathsf{E}_{\tau}[\varphi\{(\omega(\theta))_{\theta \in \Theta}, \omega(n), \omega\}] - 4C\eta$$

where  $\varphi\{(\omega(\theta))_{\theta \in \Theta}, \omega(n), \omega\} =$

$$\mathsf{E}_{\chi}\left[\sum_k \mathbb{1}_{S^x \times T(\omega) \subseteq B^k} \prod_t \alpha^x(B_t^k)^{d_t(\omega)} \prod_s \prod_{c=1}^{c_s^x} \mathbb{1}_{\omega(\theta_{sc}) \in B_s^k \cap T(\omega)} G_{s_{\omega(n)}}^k \mathbb{1}_{\omega(n) \in T(\omega)}\right]$$

Let  $\tilde{\omega}(n) \in L_{\infty}(\Omega; \tau)^{\#T}$  describe the random move of player II at each stage  $n$  and define  $\Phi: L_{\infty}(\Omega; \tau)^{\#T(\# \Theta + 1)} \rightarrow \mathbb{R}$  by  $\Phi\{(\tilde{\omega}(\theta))_{\theta \in \Theta}, \tilde{\omega}(n)\} = \mathsf{E}_{\tau}[\varphi\{(\omega(\theta))_{\theta \in \Theta}, \omega(n), \omega\}]$ . Then:

$$(1) \quad \gamma_n(\sigma, \tau) \geq \Phi\{(\tilde{\omega}(\theta))_{\theta \in \Theta}, \tilde{\omega}(n)\} - 4C\eta$$

We are now going to study  $\Phi$ . Let  $D$  be the set of limit points of  $\{\tilde{\omega}(n)\}$  in  $L_{\infty}(\Omega; \tau)^{\#T}$  for  $\sigma(L_{\infty}, L_1)$  and denote by  $F$  its closed convex hull. Define finally  $\Phi^*$  from  $F \times F$  to  $\mathbb{R}$  by:  $\Phi^*(f, g) = \Phi\{(f), g\}$  where  $(f)$  is the  $\Theta$ -vector  $f$ . Remark first that:

$$(2) \quad \Phi^*(f, f) = \Phi\{(f), f\} \geq v(\bar{G}) - \eta, \quad \text{for all } f \text{ in } F$$

Indeed note that if  $\tilde{\omega}(\theta) = \tilde{\omega}(n) = f$ , then  $\varphi\{(f(\omega)), f(\omega), \omega\}$  is the pay-off in  $\bar{G}$  induced by  $\chi$  and the pure strategy  $\{T(\omega), f(\omega), d(\omega)\}$  of player II (recall that the support of  $f(\omega)$

is included in  $T(\omega)$ ), thus greater than  $v(\bar{G}) - \eta$ , hence the claim by taking the expectation w.r.t.  $\tau$ .

Now  $F$  is compact convex,  $\Phi^*$  is continuous hence by Theorem 2.6 p. 17  $\Phi^*$  has a value, say  $w$ . Player I has an  $\eta$ -optimal strategy  $\nu$  with finite support on  $F$  and  $\Phi^*$  being affine w.r.t.  $g$ , player II has a pure optimal strategy, say  $g^*$ . Hence, by using (2):

$$(3) \quad \forall g \in F, \quad \eta + \int_F \Phi^*(f, g) \nu(df) \geq w \geq \sup_{f \in F} \Phi^*(f, g^*)$$

$$(4) \quad \geq \Phi^*(g^*, g^*)$$

$$(5) \quad \geq v(\bar{G}) - \eta$$

By convexity,  $F$  is also the closed convex hull of  $D$  in the Mackey topology; hence, since on bounded sets of  $L_\infty$  this topology coincides with the topology of convergence in probability (cf. ex. I.2Ex.16 p. 25), we obtain that every  $f$  in  $F$  is a limit in probability of some sequence of convex combination of elements in the sequence  $\{\tilde{\omega}(n)\}$ . More precisely, for all  $f \in F$ , all  $n$  and all  $j = 1, 2, \dots$ , there exists  $\mu$  with finite support on  $[n, +\infty)$  such that:

$$P\left(\left\|f - \sum_i \mu(i) \tilde{\omega}(i)\right\| > \eta\right) < \eta^j.$$

Define inductively, for  $j = 1, 2, \dots$ , a sequence of  $\eta^j$  approximations  $\mu_j$ , with disjoint supports, starting from  $N_0 = n$ . Let  $J > \bar{c}^2/2\varepsilon$  and  $\bar{\mu} = 1/J \sum_{j=1}^J \mu_j$ . Now we have:

$$(6) \quad P\left(\left\|f - \sum_i \bar{\mu}(i) \tilde{\omega}(i)\right\| > \eta\right) < \eta/(1 - \eta)$$

We then use  $\bar{\mu}$  to select independently the  $\bar{c}$  points of  $\Theta$  in  $[N_0, +\infty)$ . Note that the probability of selecting twice the same  $i$  is smaller than  $\sum_i \binom{\bar{c}}{2} \bar{\mu}^2(i) < \bar{c}^2/2 \max_i \bar{\mu}(i) < \eta$  by the choice of  $J$ . Given  $f$  in  $F$ , let us write  $\sigma^f$  for the strategy using  $\bar{\mu}$  ( $= \bar{\mu}^f$ ) to choose  $\Theta$  and then playing  $\sigma(\Theta)$ . Let  $N^f > N_2$  with  $\bar{\mu}\{N^f, +\infty\} = 0$ . Then for  $n > N^f$  we obtain using (1):

$$\gamma_n(\sigma^f, \tau) \geq E_{\bar{\mu}^f} \Phi\{(\tilde{\omega}(\theta))_{\theta \in \Theta}, \tilde{\omega}(n)\} - 5C\eta.$$

$\Phi$  being linear in each  $\tilde{\omega}(\theta)$  and the  $\theta$  being i.i.d., (6) implies that for  $\eta$  small enough:

$$(7) \quad \gamma_n(\sigma^f, \tau) \geq \Phi\{(f), \tilde{\omega}(n)\} - 7C\eta$$

Let  $N_j$  be a sequence along which  $\gamma_n$  converges to  $\liminf \gamma_n$ , and we still denote by  $N_j$  a subsequence on which  $\tilde{\omega}(n)$  converges  $\sigma(L_\infty, L_1)$  to some  $g \in D \subseteq F$ . We obtain:

$$\liminf \gamma_n(\sigma^f, \tau) = \lim \gamma_{N_j}(\sigma^f, \tau) \geq \lim \Phi\{(f), \tilde{\omega}(N_j)\} - 7C\eta = \Phi^*(f, g) - 7C\eta$$

Let finally  $\sigma^\nu$  denote the strategy of player I that choose first  $f$  according  $\nu$  and then play  $\sigma^f$ . Then:

$$\begin{aligned} \liminf \gamma_n(\sigma^\nu, \tau) &= \liminf \int_F \gamma_n(\sigma^f, \tau) \nu(df) \\ &\geq \int_F \liminf \gamma_n(\sigma^f, \tau) \nu(df) \\ &\geq \int_F \Phi^*(f, g) \nu(df) - 7C\eta \\ &\geq v(\bar{G}) - 9C\eta \quad \text{using (3)} \end{aligned}$$

We have thus proved that, for all  $\varepsilon > 0$ , for all  $\tau$ , there exists  $\sigma$  with:

$$\liminf \gamma_n(\sigma, \tau) \geq v(\bar{G}) - \varepsilon.$$

Obviously the same minoration with  $\bar{\gamma}_n$  follows, hence the proof of A.

### PART B. Player II can guarantee $v(\bar{G})$ .

This part is much easier to explain: a pure strategy of player II in  $\bar{G}$  induces a strategy in  $\Gamma_\infty$  by playing  $\beta$  i.i.d. except at finitely many stages where the exceptional moves are used. We will show that if these moves are uniformly spread on a large number of stages, player I's behaviour can be approximated by a strategy in  $\bar{G}$ , so that the pay-offs in both games will be close. The result will then follow by letting player II make an initial choice according to an optimal strategy in  $\bar{G}$ .

Let  $\varepsilon > 0$  and  $\psi$  be an  $\varepsilon$ -optimal strategy for player II in  $\bar{G}$  with finite support in  $Y$ . Define:  $\bar{d} = \max\{d(y) = \sum_t d_t^y \mid y \in Y, \psi(y) > 0\} + 1$ , and  $\delta = \min\{\beta_t^y \mid y \in Y, t \in T^y \text{ and } \psi(y) > 0\}$ . Choose finally  $N_0 > (|\ln \varepsilon| \cdot \#S \cdot \bar{d} \cdot (\bar{d} + 1)) / \varepsilon^{\bar{d}+1} \delta$  and define a strategy  $\tau = \tau(N_0)$  of player II in  $\Gamma$  as:

- Use first  $\psi$  to select  $y$  in  $Y$  and then play the following  $\tau[y; N_0]$
- Generate random times  $\{\theta_{td} \mid t \notin T^y, 1 \leq d \leq d_t^y\}$ , by choosing independently and uniformly stages in  $[1, N_0]$  (with new choices if some of these stages coincide). Let us write  $\Theta = \{\theta_{td}\}$ .
- Play now  $\beta^y$  i.i.d. except at stages  $\theta_{td}$  where  $t$  is played.

Let  $\omega$  in  $S^\infty$  be a pure strategy of player I in  $\Gamma$  and let us compute the pay-off induced by  $\omega$  and  $\tau$ . For this purpose, we will represent  $\omega$  as a strategy in  $\bar{G}$ ; given  $\omega, N_0$  and some  $r \in \mathbb{N}$  define  $x[\omega; r; N_0]$  as follows: Let  $n_s = \#\{\omega^{-1}(s) \cap [1, N_0]\}$ , then  $S^x = \{s \mid n_s \geq r\}$ ;  $\alpha_s^x = n_s / \sum_{i \in S^x} n_i$ ,  $s \in S^x$ ;  $c_s^x = n_s$ ,  $s \notin S^x$ .

We first approximate the probability  $Q^k(\omega, \tau)$  that the game is not revealed before stage  $N_0$ , given  $k, \omega, \tau$ .

LEMMA 2.5. Let  $N_0 > (|\ln \varepsilon| \cdot \#S \cdot \bar{d} \cdot (\bar{d} + 1)) / \varepsilon^{\bar{d}+1} \delta$  and  $r = ((\bar{d} + 1) |\ln \varepsilon|) / \ln(1 - \delta)$ , then:

$$|Q^k(\omega, \tau[y; N_0]) - \rho^k(x[\omega; r; N_0], y)| \leq 7\varepsilon$$

PROOF.  $\omega$ , hence  $x$ , being fixed, we will write  $S'$  for  $S^x$  and  $\alpha$  (resp.  $c$ ) for  $\alpha^x$  (resp.  $c^x$ ); we will also use  $\zeta(td)$  for  $\omega(\theta_{td})$ ,  $Q$  for  $Q^k(\omega, \tau[y; N_0])$ , and  $\mathbb{E}$  for  $\mathbb{E}_{\tau[y; N_0]}$ . Then we have:

$$Q = \mathbb{E}[\mathbb{1}_{\{\forall t, d, \zeta(td) \in B_t^k\}} \mathbb{1}_{\{\forall n \notin \Theta, 1 \leq n \leq N_0, j_n \in B_{\omega(n)}^k\}}].$$

This can be written as:

$$\begin{aligned} Q &= \mathbb{E}\left[\prod_t \prod_{d=1}^{d_t^y} \mathbb{1}_{t \in B_{\zeta(td)}^k} \prod_{\substack{n \notin \Theta \\ 1 \leq n \leq N_0}} \beta^y(B_{\omega(n)}^k)\right] \\ &= \prod_{n=1}^{N_0} \beta^y(B_{\omega(n)}^k) \mathbb{E}\left[\prod_t \prod_{d=1}^{d_t^y} \mathbb{1}_{t \in B_{\zeta(td)}^k} / \beta^y(B_{\zeta(td)}^k)\right]. \end{aligned}$$

Let us remark that, with  $\mathbb{I} = \mathbb{1}_{\{\forall s, n_s=0 \text{ or } \beta^y(B_s^k) \geq \varepsilon\}}$ ,

$$(8) \quad \left\| Q - \mathbb{I} \prod_s \beta^y(B_s^k)^{n_s} \mathbb{E}\left[\prod_t \prod_{d=1}^{d_t^y} \sum_{s \in B_t^k} \mathbb{1}_{\zeta(td)=s} / \beta^y(B_s^k)\right] \right\| \leq 2\varepsilon$$

In fact, if at some stage  $n$ ,  $\omega(n) = s$  with  $\beta^y(B_s^k) \leq \varepsilon$ , then  $Q \leq 2\varepsilon$  since  $\Pr(n \in \Theta) \leq \bar{d}/N_0 < \varepsilon$ . Now the probability that the random times  $\theta_{td}$  differ from the original i.i.d. choices  $\theta_{td}^*$  is:

$$(9) \quad P\{\exists(t, d) \neq (t', d'), \theta_{td}^* = \theta_{t'd'}^*\} \leq 1 - \prod_{c=1}^{\bar{c}} (1 - c/N_0) \leq 1 - \exp(-\bar{c}^2/N_0) \leq \varepsilon^{\bar{c}+1}$$

Otherwise the  $\theta_{td}$  are i.i.d. and uniformly distributed hence:

$$(10) \quad \mathbb{E} \left[ \prod_t \prod_{d=1}^{d^y} \sum_{s \in B_t^k} \mathbb{1}_{\zeta(td)=s} / \beta^y(B_s^k) \right] = \prod_t \sum_{s \in B_t^k} \left( \frac{n_s/N_0}{\beta^y(B_s^k)} \right)^{d_t^y}$$

Recall that  $\beta^y(B_s^k) < 1$  implies  $\beta^y(B_s^k) \leq 1 - \delta$ , hence  $n_s \geq r$  implies:

$$(11) \quad |\beta^y(B_s^k)^{n_s} - \mathbb{1}_{T^y \subseteq B_s^k}| \leq \varepsilon^{\bar{d}+1}$$

Note finally that by the choice of  $N_0$ ,  $s \notin S^x$  implies  $n_s/N_0 \leq (\varepsilon^{\bar{d}+1})/\#S\bar{d}$ , and that the function  $\prod_t z_t^{d_t^y}$  has Lipschitz constant  $\bar{d}\varepsilon^{-\bar{d}+1}$  on  $0 \leq z_t \leq 1/\varepsilon$ , so that:

$$(12) \quad \left| \prod_t \sum_{s \in B_t^k} \left( \frac{n_s/N_0}{\beta^y(B_s^k)} \right)^{d_t^y} - \prod_t \sum_{s \in S^x \cap B_t^k} \left( \frac{n_s/N_0}{\beta^y(B_s^k)} \right)^{d_t^y} \right| \leq \varepsilon$$

Obviously we have:

$$(13) \quad \mathbb{I} \prod_t \sum_{s \in S^x \cap B_t^k} \left( \frac{n_s/N_0}{\beta^y(B_s^k)} \right)^{d_t^y} \leq \varepsilon^{-\bar{d}}$$

We use now (9–13) to get in (8):

$$\left| Q - \mathbb{I} \prod_{s \notin S^x} \beta^y(B_s^k)^{n_s} \mathbb{1}_{S^x \times T^y \subseteq B^k} \prod_t \sum_{s \in S^x \cap B_t^k} \left( \frac{n_s/N_0}{\beta^y(B_s^k)} \right)^{d_t^y} \right| \leq 5\varepsilon$$

Finally for  $s \notin S^x$ , with  $n_s \neq 0$  and  $\beta^y(B_s^k) < \varepsilon$  the second part above is smaller than  $\varepsilon$ , hence we obtain:

$$(14) \quad \left| Q - \mathbb{1}_{S^x \times T^y \subseteq B^k} \prod_s \beta^y(B_s^k)^{c_s} \prod_t \sum_{s \in S^x \cap B_t^k} (n_s/N_0)^{d_t^y} \right| \leq 6\varepsilon$$

It remains to replace  $n_s/N_0$  by  $\alpha_s = n_s / \sum_{s \in S^x} n_s$ , but  $\sum_{s \in S^x} n_s/N_0 \geq 1 - (\#S/N_0)r > 1 - \varepsilon^{\bar{d}+1}/\bar{d}$  so that  $(\sum_{s \in S^x} n_s/N_0)^{-\bar{d}} < 1 + 2\varepsilon$  for  $\varepsilon$  small enough. Coming back to (14), we finally get:  $|Q - \mathbb{1}_{S^x \times T^y \subseteq B^k} \prod_s \beta^y(B_s^k)^{c_s} \prod_t \alpha(B_t^k)^{d_t^y}| \leq 7\varepsilon$ . ■

Returning to the proof of B, let us now compute the average pay-off in  $\Gamma$  at some stage  $n > N_0/\varepsilon$ . Define for  $s \in S$  and  $j > N_0$ :  $m_s^j = \#\{\omega^{-1}(s) \cap [N_0, j]\}$  and  $n_s^j = \#\{\omega^{-1}(s) \cap [1, j]\}$ . Then we have:

$$n\bar{\gamma}_n(\omega, \tau) \leq CN_0 + \sum_s \sum_{j=N_0}^{n-1} \mathbb{1}_{\omega_{j+1}=s} \mathbb{E}_\tau \sum_k Q^k \left[ \prod_{s' \neq s} (\beta^y(B_{s'}^k))^{m_{s'}^j} \right] \mathbb{1}_{T^y \subseteq B_s^k} s G^k \beta^y + \frac{C}{\delta}$$

where the last term comes from the fact that if  $T^y \not\subseteq B_s^k$ , then the average number of times  $s$  is played before  $k$  is revealed is less than  $1/\delta$ . Using the evaluation of  $Q^k$  in the previous lemma 2.5 we obtain:

$$\begin{aligned} \bar{\gamma}_n(\omega, \tau) &\leq (1/n) \left( \sum_{j=N_0}^{n-1} \sum_s \mathbb{1}_{\omega_{j+1}=s} \mathbb{E}_\tau \sum_k \mathbb{1}_{S^x \times T^y \subseteq B^k} \prod_y (\alpha^x(B_t^k))^{c_t^y} \right. \\ &\quad \left. \left[ \prod_{\substack{s' \notin S^x \\ s' \neq s}} (\beta^y(B_{s'}^k))^{n_{s'}^j} \right] \mathbb{1}_{T^y \subseteq B_s^k} s G^k \beta^y + 10C\varepsilon \right). \end{aligned}$$

It remains to remark that the term  $\mathbb{E}_\tau(\dots)$  equals precisely the pay-off  $F$  in  $\bar{G}$ , corresponding to  $\psi$  and the pure strategy  $x'$  defined by:  $S^{x'} = S^x \cup S$ ;  $\alpha^{x'} = \alpha^x$  on  $S^x$ , and  $\alpha^s = 0$ ;  $c_s^{x'} = n_s^j$  on  $S^{x'}$ ;  $s^{x'} = s$ . By the choice of  $\psi$ ,  $F(x', \psi) \leq v(\bar{G}) + C\varepsilon$ , implying:  $\bar{\gamma}_n(\omega, \tau) \leq \frac{1}{n}((n - N_0)(v(\bar{G}) + C\varepsilon)) + 10C\varepsilon \leq v(\bar{G}) + 12C\varepsilon$ .

This ends the proof of B, hence of theorem 2.2. ■

For examples where  $v(\bar{G}) \neq v(\underline{G})$  ( $\underline{G}$  being obviously defined in a dual way), cf. ex. VIIIEx.2 p. 393. We are thus led to study  $\lim v_n$  and  $\lim v_\lambda$ .

**2.d.  $\lim v_n$  and  $\lim v_\lambda$ .** As in the previous section, the analysis will be done through the comparison with an auxiliary game. Note nevertheless that we cannot use an asymptotic approximation of the pay-offs since the game is basically finite. What we will do is to use a sequence of “approximating games”. (Compare also with sect. 4.a below where a single “limit” game can be constructed).

For each  $L$  in  $\mathbb{N}$  we shall construct a game  $G_L$ . The heuristic interpretation of  $G_L$  is  $\Gamma_n$  played in  $L$  large blocs, on each of which both players are using stationary strategies, except for some singular moves. The strategy sets in  $G_L$  are  $\bar{X}^L$  and  $\bar{Y}^L$ , where as before:

$$\begin{aligned}\bar{X} &= \bigoplus_{S' \subseteq S} \Delta(S') \times \mathbb{N}^{S \setminus S'}, \\ \bar{Y} &= \bigoplus_{T' \subseteq T} \Delta(T') \times \mathbb{N}^{T \setminus T'}.\end{aligned}$$

As in subsection 2.b we will write  $x = (S^x, \alpha^x, c^x)$  and similarly for  $y$ . The probability of getting the signal 0, under  $x$  and  $y$  and given  $k$  is again:

$$\rho^k(x, y) = \mathbb{1}_{S^x \times T^y \subseteq B^k} \prod_t \alpha^x(B_t^k)^{d_t^y} \prod_s \beta^y(B_s^k)^{c_s^x}$$

and the pay-off is  $f^k(x, y) = \alpha^x G^k \beta^y$ . Given  $x = \{x(l)\}$  (resp.  $y = \{y(l)\}$ ) in  $\bar{X}^L$  (resp.  $\bar{Y}^L$ ), we define  $F_L$  by:

$$F_L(x, y) = \sum_k \frac{1}{L} \sum_{l=1}^L \left( \prod_{m=0}^{l-1} \rho^k\{x(m), y(m)\} \right) f^k(x(l)y(l)),$$

with  $\rho^k\{x(0), y(0)\} = 1$ . We introduce also  $X = \{x \in \bar{X} \mid \alpha^x > 0 \text{ on } S^x\}$  and similarly  $Y$ .

Then we have the following result, the proof of which is similar to prop. 2.1:

**PROPOSITION 2.6.**  $G_L$  has a value  $w_L$  and both players have  $\varepsilon$ -optimal strategies with finite support in  $X^L$  and  $Y^L$ .

We can now state:

**THEOREM 2.7.**  $\lim_{n \rightarrow \infty} v_n$  and  $\lim_{L \rightarrow \infty} w_L$  exist and coincide.

**PROOF.**

#### PART A. Sketch of the proof.

Obviously it will be sufficient to show that  $\liminf v_n \geq \limsup w_L$ .

Denote then  $\limsup w_L$  by  $w$ . Given  $\varepsilon > 0$  and  $L_0$  large enough we shall choose  $L \geq L_0$  with  $w_L \geq w - \varepsilon$  and  $\chi_L$  an  $\varepsilon$ -optimal strategy of player I in  $G^L$  with finite support in  $X^L$ . Each  $x$  in the support of  $\chi_L$  will induce a strategy  $\sigma[x]$  of player I in  $\Gamma_n$  for  $n$  large enough. On the other hand each pure strategy  $\tau$  of player II in  $\Gamma_n$  will be represented as a strategy  $y[\tau]$  in  $Y^L$ . We shall then prove that there exists  $N(\varepsilon, L)$  such that  $n \geq N(\varepsilon, L)$  implies:

$$(15) \quad \bar{\gamma}_n(\sigma[x], \tau) \geq F_L(x, y[\tau]) - 8\varepsilon$$

for all  $x \in \text{Supp } \chi_L$  and all  $\tau \in \mathcal{T}$ . Integrating will give:

$$\mathbb{E}_{\chi_L}(\bar{\gamma}_n(\sigma[x], \tau)) \geq F_L(\chi_L, y[\tau]) - 8\varepsilon$$

So that, defining  $\sigma$  in  $\Gamma_n$  as: select  $x$  according to  $\chi_L$  and then play  $\sigma[x]$ , we obtain:

$$v_n \geq w_L - 9\varepsilon \geq w - 10\varepsilon.$$

### PART B. Preliminary results.

Given  $x \in X$  we construct a strategy  $\sigma$  in  $\Gamma_n$  for  $n$  large enough. Let  $c^x = \sum_s c_s^x$ .  $\sigma$  will consist of playing  $\alpha^x$  i.i.d. at each stage except on  $c^x$  stages. These exceptional stages are obtained by using  $c^x$  independent random variables  $\theta_{sc}^*$ ,  $s \in S$ ,  $1 \leq c \leq c_s^x$ , uniformly distributed on  $[1, n]$  (adding new trials if some choices coincide so that the final choices are  $\theta_{sc}$ ). At stage  $\theta_{sc}$ ,  $\sigma$  consists of playing  $s$ . Denote this strategy by  $\sigma[x; n]$ .

Given  $\tau$ , pure strategy of player II in  $\Gamma_n$ , we construct now  $y$  in  $Y$ . Let  $d_t = \#\{j \mid \tau_j = t, 1 \leq j \leq n\}$  and given some  $r$  in  $\mathbb{N}$ , let  $T' = \{t \in T \mid d_t \geq r\}$ . Define  $y$  by  $T^y = T'$ ;  $\beta_t^y = d_t / \sum_{t \in T'} d_t$  for  $t \in T'$ ;  $d_t^y = d_t$  for  $t \notin T'$ . We denote this strategy by  $y[\tau; r; n]$ . Let  $\mathcal{A}_m$  be the event  $\{a_j = 0 \mid 1 \leq j \leq m\}$ , where as usual  $a_j$  is the signal at stage  $j$ .

Then we rewrite lemma 2.5 p. 366 as:

LEMMA 2.8. *Let  $x$  in  $X$  and  $\delta_x = \min\{\alpha_s^x \mid s \in S^x\}$ . Choose  $1/4 > \varepsilon_1 > 0$  and  $r = (c^x + 1) |\ln \varepsilon_1| / \ln(1 - \delta_x)$ . Then  $n \geq |\ln \varepsilon_1| \#T.c^x(c^x + 1) / \varepsilon_1^{c^x+1} \delta_x$  implies:*

$$|\rho^k\{x, y[\tau; r; n]\} - P_{\sigma[x; n], \tau, k}(\mathcal{A}_n)| < 7\varepsilon_1.$$

### PART C. Construction of $\sigma$ in $\Sigma$ and $y$ in $Y^L$ .

Given  $\varepsilon > 0$ , let  $\varepsilon' = \varepsilon/K$  and choose  $L_0$  such that:

$$(16) \quad (1 - \varepsilon')^{\lfloor L_0^{1/2} \rfloor} \leq \varepsilon' \text{ and } \lfloor L_0^{1/2} \rfloor^{-1} \leq \varepsilon'$$

where  $\lfloor \dots \rfloor$  denotes the integral part. Take  $L \geq L_0$  and  $\chi_L$  as in part A. Let  $\delta = \min\{\delta_{x(l)} \mid 1 \leq l \leq L, x = \{x(l)\} \in \text{Supp } \chi_L\}$  and  $\bar{c} = \max\{c^{x(l)} \mid 1 \leq l \leq L, x = \{x(l)\} \in \text{Supp } \chi_L\}$ . Assume  $\varepsilon_1 < \varepsilon'/7L$ ,  $r = (\bar{c} + 1) |\ln \varepsilon_1| / \ln(1 - \delta)$ , and  $N_0 \geq |\ln \varepsilon_1| \#T.\bar{c}(\bar{c} + 1) / \varepsilon_1^{\bar{c}+1} \delta$ . For  $n = NL + n_1$  with  $n_1 \leq L$  and  $N \geq N_0$ , we construct  $\sigma$  in  $\Gamma_n$  by specifying this strategy on each bloc  $l$  of length  $N$ ,  $l = 1, \dots, L$ , (the  $l^{\text{th}}$  bloc consists of the stages  $m \in N(l) \equiv \{(l-1)N + 1, \dots, lN\}$ ) to be  $\sigma(l) \equiv \sigma[x(l); N]$ . Similarly, given  $\tau$  pure strategy of player II in  $\Gamma_n$ , we consider the restriction  $\tau(l)$  of  $\tau$  on each bloc  $l$  and define  $y = \{y(l)\}$  by  $y(l) = y[\tau(l); r; N]$  where (cf. B)  $d_t(l) = \#\{m \in N(l) \mid \tau_m = t\}$ , and so on.

### PART D. Comparison of the pay-offs.

We are going to compare the pay-offs in  $G_L$  for  $x$  and  $y$  and in  $\Gamma_n$  for  $\sigma$  and  $\tau$ . Note first that it is sufficient to approximate on each block and moreover that we can work conditionally on the state being  $k$ , since the mappings  $x \mapsto \sigma$  and  $\tau \mapsto y$  do not depend on  $k$ . Accordingly, we shall drop the index  $k$  until (18). In  $G_L$  the pay-off on bloc  $l$  is:

$$F_l(x, y) = \prod_{m=0}^{l-1} \rho\{x(m), y(m)\} f(x(l)y(l))$$

and in  $\Gamma_n$  the pay-off on the corresponding bloc can be written as:

$$\Phi_l(\sigma, \tau) \equiv \prod_{m=0}^{l-1} Q_m(\sigma, \tau) \varphi_l(\sigma, \tau)$$

where  $Q_m(\sigma, \tau) = P_{\sigma, \tau}\{\mathcal{A}_{mN} \mid \mathcal{A}_{(m-1)N}\}$  and  $\varphi_l(\sigma, \tau)$  is the average expected pay-off on the stages in  $N(l)$  conditionally on  $\mathcal{A}_{(l-1)N}$ . By the choice of  $r$  and  $N_0$  it follows from lemma 2.5 p. 366 that:

$$|\rho\{x(m), y(m)\} - Q_m(\sigma, \tau)| \leq 7\varepsilon_1,$$

hence:

$$(17) \quad \left| \prod_{m=0}^l \rho\{x(m), y(m)\} - \prod_{m=0}^l Q_m(\sigma, \tau) \right| \leq 7L\varepsilon_1 \leq \varepsilon', \quad \text{for all } l.$$

It remains then to compare  $f_l (= f(x(l)y(l)))$  to  $\varphi_l$ . We shall first ignore the blocs where  $Q_m$  is small; in fact after many of such blocs the game will be revealed with a high probability, and both  $F_l$  and  $\Phi_l$  approximately 0. Now on the blocs where  $Q_m$  is large the expected average pay-off given  $\mathcal{A}_{(m-1)N}$  is near the Cesàro mean, hence  $f_m$  near  $\varphi_m$ . Formally let  $M_l = \{m \mid 1 \leq m \leq l, Q_m(\sigma, \tau) \leq 1 - \varepsilon'\}$  and  $m_l = \#M_l$ . Define  $l' = \min\{\{l \mid m_l \geq \lfloor L^{1/2} \rfloor\} \cup \{L+1\}\}$ . It follows from (16) that on  $\{l' \leq L\}$ ,  $\prod_{m=0}^{l'} Q_m \leq \varepsilon'$ , hence using (17), we have for  $l > l'$ :

$$|F_l - \Phi_l| \leq 3\varepsilon'C.$$

The number of blocs in  $M_{l'}$  is at most  $\varepsilon'L$  by (16) hence it remains to consider  $l \in \{1, \dots, l'\} \setminus M_{l'}$ . We have:

$$\varphi_l(\sigma, \tau) = \frac{1}{N} \sum_{m \in N(l)} \lambda_m \mathsf{E}_{\sigma, \tau}(G_{s_m t_m})$$

where  $\lambda_m = \Pr(\mathcal{A}_{m-1} \mid \mathcal{A}_{(l-1)N})$ , hence  $\lambda_m \geq 1 - \varepsilon', \forall m$ , implying:

$$\left| \varphi_l(\sigma, \tau) - \mathsf{E}_{\sigma, \tau} \frac{1}{N} \left( \sum_m G_{s_m t_m} \right) \right| < C\varepsilon'.$$

Using (9) we obtain that  $\{\theta_{sc}\}$  will coincide with  $\{\theta_{sc}^*\}$  with probability greater than  $(1 - \varepsilon')$ . On this event, they define a random subset of  $\bar{c}$  stages in  $N(l)$  such that on its complement:

- player I plays  $\alpha(l)$  i.i.d.
- the expected empirical distribution of  $\tau$  is as on  $N(l)$ , say  $\tau^*(l) = \frac{1}{N} \sum_{m \in N(l)} \tau_m$ .

It follows from the choice of  $r$  and  $N_0$  that:

$$|\tau^*(l) - \beta(l)| \leq \varepsilon'$$

hence:

$$\left| \frac{1}{N} \mathsf{E}_{\sigma, \tau} \left( \sum_m G_{s_m t_m} \right) - \alpha(l) G \beta(l) \right| \leq 3\varepsilon'.$$

So that for all  $l \notin M_{l'}$ , one has:

$$|F_l - \Phi_l| \leq 5\varepsilon'$$

This implies:

$$(18) \quad \left| \sum_{k \in K} \sum_{l \in L} (F_l^k - \Phi_l^k) \right| \leq (5\varepsilon'L + \lfloor L^{1/2} \rfloor)K$$

and finally:

$$|\bar{\gamma}_n(\sigma, \tau) - F_L(x, y)| \leq (5\varepsilon'L + \lfloor L^{1/2} \rfloor)K + 2L/n$$

which implies (15). ■

Basically the same construction will give:

**THEOREM 2.9.**  $\lim v_\lambda$  exists and  $\lim_{\lambda \rightarrow 0} v_\lambda = \lim_{L \rightarrow \infty} w_L$

**PROOF.** Given  $L$  and  $\lambda$  small enough, define  $\{n_l\}, l = 1, \dots, L$ , by requiring  $\sum_{n_l+1}^{n_{l+1}} \lambda(1 - \lambda)^{m-1}$  to be near  $1/L$ . Now  $x$  induces  $\sigma$  with  $\sigma(l) = \sigma[x(l); n_l]$ . Given  $\tau$ , denote by  $\tau(l)$  its restriction to the  $l^{\text{th}}$  bloc of length  $n_l$  and define  $y(l)$  as  $y[\tau(l); r; n_l]$ . The approximations are then very similar to those in the previous section. ■

### 3. A game with state dependent signalling matrices

**3.a. Introduction and notation.** We consider here (not symmetrical) games with lack of information on both sides but where the signalling matrices do depend on the state (compare with ch. VI).

The simplest case is given by the following data:  $K = \{0, 1\}^2 = L \times M$ , we write  $k = (l, m)$ . The probability on  $K$  is the product  $p \otimes q$  of its marginals and we denote by  $P$  and  $Q$  the corresponding simplices. At stage 0, player I is informed about  $l$  and player II about  $m$ . The pay-offs are defined by  $2 \times 2$  pay-offs matrices  $G^{lm}$  and the signalling matrices are the same for both players and given by:

$$\begin{aligned} A^{11} &= \begin{pmatrix} T & L \\ c & d \end{pmatrix} & A^{10} &= \begin{pmatrix} T & R \\ c & d \end{pmatrix} \\ A^{01} &= \begin{pmatrix} B & L \\ c & d \end{pmatrix} & A^{00} &= \begin{pmatrix} B & R \\ c & d \end{pmatrix} \end{aligned}$$

REMARK 3.1. The signals include the moves.

REMARK 3.2. As soon as player I plays Top some “type” is revealed:  $l$  if player II played Left at this stage,  $m$  if he played Right.

REMARK 3.3. The previous game  $\Gamma(p, q)$  has a value,  $v(p, q)$ , as soon as  $p^1 p^0 q^1 q^0 = 0$  by the results of ch. V.

*Notation.*

- We write  $NS$  for the set of non-separating strategies (i.e. strategies that do not depend on the announced type). As in ch. V, and this is one of the main difference with ch. VI,  $NS$  is not included in the set  $NR$  of non-revealing strategies.
- $\theta$  is the stopping time corresponding to the first time I plays Top.
- We denote by  $G(p, q)$  the average of the matrices  $G^k$  and we will also write  $\gamma^{pq}(\sigma, \tau)$  as  $\sum_{lm} p^l q^m \gamma^{lm}(\sigma^l, \tau^m)$  for the pay-off, where  $\sigma^l$  is  $\sigma$  given  $l$  and similarly for  $\tau$ .

**3.b. Minmax.** We prove in this section the existence of the minmax and give an explicit formula for it. As in sect. 2 the construction relies on an auxiliary game.

Let  $G(p)$  be the infinitely repeated stochastic game with lack of information on one side described by:

$$\begin{aligned} p^1 &\quad G^1 = \begin{pmatrix} g_{11}^{1*} & g_{12}^{1*} \\ g_{21}^1 & g_{22}^1 \end{pmatrix} \\ p^0 &\quad G^0 = \begin{pmatrix} g_{11}^{0*} & g_{12}^{0*} \\ g_{21}^0 & g_{22}^0 \end{pmatrix} \end{aligned}$$

where a star (\*) denotes an absorbing pay-off (cf. sect. 4.a p. 341). Player I is informed and we assume full monitoring. Denote by  $w_1(p)$  the value of the one-shot  $G_1(p)$ .

**PROPOSITION 3.1.** *minmax  $G(p)$  exists and equals  $w_1(p)$ .*

**PROOF.** As usual we split the proof into two parts.

**PART A. Player II can guarantee  $w_1(p)$ .**

In fact, let  $y$  be an optimal strategy of player II in  $G(p)$  and define  $\tau$  as: play  $y$  i.i.d. Given  $\sigma$ , strategy of player I and  $\tau$  let:  $z^l(n) = \Pr_{\sigma^l, \tau}(\theta \leq n)$  and  $x^l(n)$  be the vector in

$\Delta(S)$  with first component  $z^l(n)$ . Since the play of player II is independent of  $\sigma$  we easily obtain for the expected pay-off  $\rho_n$  in  $G(p)$  at stage  $n$ :

$$\rho_n(\sigma, \tau) = \sum_l p^l x^l(n) G^l y$$

hence, by the choice of  $y$ ,  $\bar{\rho}_n(\sigma, \tau) \leq w_1(p)$ , for all  $n$  and all  $\sigma$ .

### PART B. Player II cannot get less than $w_1(p)$ .

Given  $\tau$ , strategy of player II in  $G(p)$ , note first that it is enough to define  $\tau$  at stage  $n$  conditionally to  $\{\theta \geq n\}$ , hence  $\tau$  is independent of the moves of player I. We can thus introduce  $y_n = E_\tau(t_n) \in \Delta(T)$ . Given  $x$  optimal strategy of player I in  $G_1(p)$  and  $y$  in  $\Delta(T)$ , define  $\rho'(x, y)$  (resp.  $\rho''(x, y)$ ) to be the absorbing (resp. non-absorbing) component of the pay-off that they induce in  $G_1(p)$ . Formally:

$$\rho'(x, y) = \sum_l p^l x_1^l (G^l y)_1 \quad \rho''(x, y) = \sum_l p^l x_2^l (G^l y)_2$$

Let  $\varepsilon > 0$  and choose  $N$  such that:

$$\rho'(x, y_N) \geq \sup_n \rho'(x, y_n) - \varepsilon$$

Define  $\sigma$  as: Play always Bottom except at stage  $N$  where  $x$  is used. For  $n \geq N$  we get:

$$\begin{aligned} \rho_n(\sigma, \tau) &= \rho'(x, y_N) + \rho''(x, y_n) \\ &\geq \rho'(x, y_n) + \rho''(x, y_n) - \varepsilon \\ &\geq w_1(p) - \varepsilon \end{aligned}$$

hence  $n \geq CN/\varepsilon$  implies  $\bar{\rho}_n(\sigma, \tau) \geq w_1(p) - 2\varepsilon$ . ■

Given  $\alpha = (\alpha^1, \alpha^0)$  and  $\beta = (\beta^1, \beta^0)$  in  $\mathbb{R}^2$ ,  $G(p, q; \alpha, \beta)$  is a game of the previous class with:

$$G^1 = \begin{pmatrix} v(1, q)^* & (q^1 \alpha^1 + q^0 \beta^1)^* \\ G_{21}(1, q) & G_{22}(1, q) \end{pmatrix} \quad G^0 = \begin{pmatrix} ccv(0, q)^* & (q^1 \alpha^0 + q^0 \beta^0)^* \\ G_{21}(0, q) & G_{22}(0, q) \end{pmatrix}$$

Let  $w_1(p, q; \alpha, \beta)$  be the value of  $G_1(p, q; \alpha, \beta)$ , hence by prop. 3.1 the minmax of  $G(p, q; \alpha, \beta)$ .

We introduce two closed convex sets of vector pay-offs in  $\mathbb{R}^2$ :

$$\begin{aligned} \mathcal{A} &= \{ \alpha \mid \alpha^1 r^1 + \alpha^0 r^0 \geq v(r, 1) \text{ for all } r \text{ in } P \} \\ \mathcal{B} &= \{ \beta \mid \beta^1 r^1 + \beta^0 r^0 \geq v(r, 0) \text{ for all } r \text{ in } P \} \end{aligned}$$

REMARK 3.4.  $\mathcal{A}$  (resp.  $\mathcal{B}$ ) corresponds to the affine majoraters of  $v(\cdot, 1)$  (resp.  $v(\cdot, 0)$ ).

THEOREM 3.2.  $\bar{v}(p, q)$  exists on  $P \times Q$  and is given by:

$$\bar{v}(p, q) = \text{Vex} \min_q \min_{\substack{\alpha \in \mathcal{A} \\ \beta \in \mathcal{B}}} w_1(p, q; \alpha, \beta).$$

PROOF. The proof is again divided in two parts corresponding to conditions (i) and (ii) of the definition of the minmax (subsection 1.c p. 149).

### PART C. Player II can guarantee this pay-off.

Since player II knows  $m$ , it is enough by Theorem 1.1 p. 183, to prove that given any  $(\alpha, \beta)$  in  $\mathcal{A} \times \mathcal{B}$ , he can guarantee  $w_1(p, q; \alpha, \beta)$ . Consider now the following class  $\mathcal{T}^*$  of strategies of player II:

- (1) play  $NS$  up to stage  $\theta$ .

- (2) if  $a_\theta = T$  (resp.  $B$ ) play from this stage on optimally in  $\Gamma(1, q)$  (resp.  $\Gamma(0, q)$ ) (cf. sect. 3 p. 191).
- (3) if  $a_\theta = L$  (resp.  $R$ ) play from this stage on a strategy that approaches the vector pay-off  $\alpha$  (resp.  $\beta$ ), (cf. sect. 3 again).

Note that this construction is consistent: for 2, since player II was playing  $NS$  until  $\theta$ , the posterior on  $M$  after this stage is still  $q$ ; as for 3,  $\mathcal{A}$  is precisely the set of vector pay-offs that player II can approach in the game with lack of information on one side defined for  $m = 1$ .

It is now quite clear that, if player II is playing in  $\mathcal{T}^*$ , the original game  $\Gamma(p, q)$  is equivalent to the auxiliary game  $G$  so that by playing an optimal strategy in  $\mathcal{T}^*$  for  $G$ , player II can obtain in  $\Gamma$  minmax  $G$ . In fact, let  $y$  be an optimal strategy of player II in  $G_1(p, q; \alpha, \beta)$  and given  $\varepsilon > 0$ , let  $\tau(1)$  and  $N(1)$  such that:

$$\bar{\gamma}_n^{l1}(\sigma, \tau) \leq \alpha^l + \varepsilon, \quad l = 0, 1, \quad \text{for all } n \geq N(1) \text{ and all } \sigma$$

and define similarly  $\tau(0)$  and  $N(0)$ . Let  $N = \max\{N(1), N(0)\}$  and  $\tau$  be the corresponding strategy in  $\mathcal{T}^*$  where 1 consists of playing  $y$  i.i.d., and 3 is to play  $\tau(1)$  or  $\tau(0)$ .

Given a strategy  $\sigma$  of player I in  $\Gamma$ , let us consider the pay-off  $\bar{\gamma}_n^{pq}(\sigma, \tau)$ . It suffices to majorate  $\bar{\gamma}_n^{1,q}(\sigma^1, \tau)$ , but by definition of  $\tau$  and  $\theta$  one has:

$$P_{\sigma^1, \tau}\{t_n = L \mid n \leq \theta\} = y$$

hence:

$$\mathsf{E}_{q, \sigma^1, \tau}[g_n \mid n \leq \theta] = (G(1, q)y)_2.$$

Moreover, under  $(\sigma^1, \tau)$ ,  $a_\theta$  will take the values  $T, L, R$  with the respective probabilities  $y_1, y_2q^1, y_2q^0$ . Finally we have:

$$\mathsf{E}_{q, \sigma^1, t}[g_n \mid n > \theta \text{ and } a_\theta = T] \leq v(1, q)$$

and if  $n \geq \theta + N$ :

$$\mathsf{E}_{q, \sigma^1, \tau}\left[(n - \theta)^{-1} \sum_{i=\theta}^n g_i \mid a_\theta = L\right] \leq \alpha^1 + \varepsilon$$

From the previous evaluations we deduce that  $n\bar{\gamma}_n^{1,q}(\sigma^1, \tau)$  is bounded by

$$\mathsf{E}_{\sigma^1, \tau}[(\theta \wedge n)(G(1, q)y)_2 + (n - (\theta \wedge n))\{y_1v(1, q) + y_2(q^1\alpha^1 + q^0\beta^1)\}] + n\varepsilon + 2CN.$$

Let now  $F(x, y) \equiv F(x, y; p, q; \alpha, \beta)$  be the pay-off in  $G_1(p, q; \alpha, \beta)$  when  $(x, y)$  is played. Then we have:

$$\bar{\gamma}_n(\sigma, \tau) \leq F(\xi, y) + \varepsilon + 2CN/n$$

where  $\xi^l \in \Delta(S)$  with  $\xi_2^l = (1/n) \mathsf{E}_{\sigma^l, \tau}(\theta \wedge n)$ ,  $l = 1, 0$ . So finally by the choice of  $y$  we get that  $n \geq 2CN/\varepsilon$  implies:

$$\bar{\gamma}_n(\sigma, \tau) \leq w_1(p, q; \alpha, \beta) + 2\varepsilon, \quad \text{for all } \sigma.$$

#### PART D. Player II cannot get less than $\bar{v}$ .

We want here to exhibit good replies of player I. A priori such a strategy, i.e. mainly a distribution of the stopping time  $\theta$  given  $l$ , should depend on the posterior behaviour of player II, which in turn could also be a function of  $\theta$ . We are thus led to use a fixed point — or minmax — argument. Given  $(p, q)$  fixed, let us denote by  $\varphi(\alpha, \beta)$  the set of optimal strategies of player I in  $G_1(p, q; \alpha, \beta)$  defined as above by  $x^l$ ,  $l = 1, 0$ . Denote by  $\psi(x)$  the set of vectors  $(\alpha, \beta)$  in  $\mathcal{A} \times \mathcal{B}$  that minimise  $p^1x_1^1(q^1\alpha^1 + q^0\beta^1) + p^0x_0^0(q^1\alpha^0 + q^0\beta^0)$  or equivalently that minimise the absorbing pay-off given  $x$  in  $G_1$ . Remark that  $(\alpha, \beta) \in \psi(x)$  iff  $\alpha$  is a supporting hyperplane to  $v(., 1)$  at the posterior probability  $\pi$  with  $\pi^1(x) = \Pr\{l = 1 \mid T, x\}$ . Since the correspondences  $\varphi$  and  $\psi$  are u.s.c. and compact-convex valued, it follows that

$\varphi \circ \psi$  has a fixed point, say  $\bar{x}(= \{\bar{x}^l(p, q)\})$ . The construction of  $\sigma$  can now be explained: given  $\tau$ , strategy of player II, player I first plays Bottom until some stage  $N$  after which the martingale  $q_n$  of posterior probabilities on  $M$  is essentially constant. Player I now uses  $\bar{x}(p, q_N)$  to compute the stage where his non-absorbing against  $\tau$  is minimal and plays then once  $\bar{x}$ . Assuming this strategy for player I, a best reply of player II would be to use  $(\alpha, \beta)$  in  $\psi(\bar{x})$ , hence the corresponding pay-off is some  $w_1(p, q_N; \alpha, \beta)$ .

The formal proof follows. Denote by  $b$  the strategy of player I: “play always Bottom”, and given  $\tau$ , define  $N$  such that:

$$(1) \quad \mathbb{E}_{q,b,\tau}[\|q_n - q_N\|] \leq \varepsilon, \quad \forall n \geq N.$$

Use  $\zeta(N)$  for the random variable  $\bar{x}(p, q_N)$ ,  $\pi^N$  for  $\pi(\bar{x}(p, q_N))$ ; and for  $n \geq N$  define:

$$y_1(n) = P_{q,b,\tau}\{t_n = L \mid \mathcal{H}_N\} \quad \text{and} \quad y_1^m(n) = P_{q,b,\tau^m}\{t_n = L \mid \mathcal{H}_N\}$$

So:

$$(2) \quad \mathbb{E}_{q,b,\tau}[\|y(n) - y^m(n)\| \mid \mathcal{H}_N] \leq \mathbb{E}[\|q_{n+1} - q_N\| \mid \mathcal{H}_N] \stackrel{\text{def}}{=} \delta(n+1, N)$$

Denote finally by  $\rho''(p, q; x, y)$  the non-absorbing pay-off in  $G_1(p, q)$ , given  $x$  and  $y$ . For  $\varepsilon > 0$ , define  $N^* \geq N$  such that:

$$(3) \quad \rho''(p, q_N; \zeta(N), y(N^*)) \leq \rho''(p, q_N; \zeta(N), y(n)) + \varepsilon, \quad \forall n \geq N.$$

$\sigma$  is played as follows:

- play  $b$  up to stage  $N^* - 1$ .
- use  $\zeta(N)$  at stage  $N^*$ .
  - if  $s_n = B$ , keep playing  $b$
  - if  $s_n = T$ , use an optimal strategy in the revealed game, namely:
    - \* if  $a_{N^*} = L$ , use a strategy that gives at least  $v(\pi_N, 1)$  at each further stage, and similarly if  $a_{N^*} = R$ . (Recall that since player I was playing  $b$  up to stage  $N^* - 1$  the posterior probability  $p_{N^*}$  is precisely  $\pi_N$ ).
    - \* if  $a_{N^*} = T$ : given  $\varepsilon > 0$ , choose a strategy and some  $\bar{N}$  such that  $n \geq \bar{N}$  implies:

$$\mathbb{E}\left(\sum_{N^*+1}^{N^*+n} g_i \mid \mathcal{H}_{N^*}\right) \geq v(1, q_{N^*}) - \varepsilon, \quad \forall \tau.$$

(Since  $v(p, .)$  is Lipschitz,  $\bar{N}$  can be chosen uniformly w.r.t.  $q_{N^*}$ ).

Before evaluating the pay-off let us compute the probabilities of the different signals. We have:

$$\begin{aligned} P(s_{N^*} = B \mid \mathcal{H}_N) &= p^1 \zeta_2^1(N) + p^0 \zeta_2^0(N), \\ P(a_{N^*} = T \mid \mathcal{H}_N) &= p^1 \zeta_1^1(N) y_1(N^*), \\ P(a_{N^*} = L \mid \mathcal{H}_N) &= (p^1 \zeta_1^1(N) + p^0 \zeta_1^0(N)) q_N^1 y_2^1(N^*). \end{aligned}$$

and analogous formulae for  $B$  and  $R$ .

It follows, using (2), that for  $n \geq \bar{N}$ :

$$\begin{aligned} \mathbb{E}\left(\sum_{N^*+1}^{N^*+n} g_i \mid \mathcal{H}_N\right) &\geq \sum_{i=N^*+1}^{N^*+n} (\rho''(p, q_N; \zeta(N), y(i)) - C\delta(i, N)) \\ &\quad + n [p^1 \zeta_1^1(N) y_1(N^*) v(1, q_{N^*}) + p^0 \zeta_1^0(N) y_1(N^*) v(0, q_{N^*}) - \varepsilon \\ &\quad + (p^1 \zeta_2^1(N) + p^0 \zeta_2^0(N)) y_2(N^*) (q_N^1 v(\pi_N, 1) + q_N^0 v(\pi_N, 0)) - C\delta(N^*, N)] \end{aligned}$$

Using (1), (3) and the choice of  $\zeta$  imply that the right hand side can be minorated by the pay-off in  $G_1$ :

$$\mathbb{E}\left(\sum_{N^*+1}^{N^*+n} g_i \mid \mathcal{H}_N\right) \geq F(\zeta(N), y(N^*); p, q_N; \alpha, \beta) - 2nC\varepsilon - 2n\varepsilon$$

for all  $(\alpha, \beta)$  in  $\mathcal{A} \times \mathcal{B}$ .  $\zeta(N)$  being optimal in  $G_1$  at  $(p, q_N)$ , we obtain, taking expectation and using Jensen's inequality:

$$\mathbb{E}\left(\sum_{N^*+1}^{N^*+n} g_i\right) \geq n \operatorname{Vex} \min_{\alpha, \beta} w_1(p, q; \alpha, \beta) - 2n\varepsilon - 2nC\varepsilon.$$

So that  $n \geq N^* + \bar{N}$  implies  $\bar{\gamma}_n(\sigma, \tau) \geq \operatorname{Vex} \min_{\alpha, \beta} w_1(p, q; \alpha, \beta) - 5\varepsilon(C \vee 1)$ . ■

**3.c. Maxmin.** We prove in this section the existence of the maxmin and give an expression for it.

As in subsection 3.b we introduce the sets of vector pay-offs that player I can obtain, namely:

$$\begin{aligned} \mathcal{C} &= \{ \gamma = (\gamma^1, \gamma^0) \mid \gamma^1 r^1 + \gamma^0 r^0 \leq v(1, r), \text{ for all } r \text{ in } Q \} \\ \mathcal{D} &= \{ \delta = (\delta^1, \delta^0) \mid \delta^1 r^1 + \delta^0 r^0 \leq v(0, r), \text{ for all } r \text{ in } Q \} \end{aligned}$$

Note that,  $v(., q)$  being concave, we obviously have:

$$(4) \quad p^1 \gamma^m + p^0 \delta^m \leq v(p, m), \quad m = 1, 0, \quad \text{for all } p \in P \text{ and all } (\gamma, \delta) \in \mathcal{C} \times \mathcal{D}$$

**3.c.1. Sketch of the proof.** Here player I can do better than using  $NS$  strategies and then concavifying, because player II does not have a reply that could allow him to observe the moves and to play non-revealing until the convergence of the posterior probabilities on  $L$ . Basically a strategy of player I will be described by the distribution of the stopping time  $\theta$  and by the behaviour after  $\theta$ . The second aspect is similar to subsection 3.b, namely play in the game on  $P$ , at the posterior probability induced by his strategy if  $m$  is revealed, or approach some vector pay-off in  $\mathcal{C} \times \mathcal{D}$ , if  $l$  is revealed. This last choice being history-dependent, some minmax argument will also be needed. As for the first part, as long as he is playing Bottom, player I observes the moves of his opponent. Because of (4), we can restrict ourselves to monotone behaviour of the kind: Play Top iff the frequency of Right exceeds some number, say  $z$ , like in the “Big Match” (cf. ex. VIIEx.4 p. 346). It remains to choose this number and this can be done at random in a type-dependent way. A best reply of player II will then be to evaluate these distributions and to play an increasing frequency of Right up to some type-dependent level, say  $u$ , that can also be random.

**3.c.2. Preparations for the proof.** We are thus led to define  $U$  as the set of positive measures  $\mu$  with finite support on  $[0, 1]$  and total mass less than 1 and  $V$  as the set of probability distributions  $\nu$  on  $[0, 1]$ .  $\bar{\mathcal{C}}$  (resp.  $\bar{\mathcal{D}}$ ) is the set of measurable mappings from  $[0, 1]$  to  $\mathcal{C}$  (resp.  $\mathcal{D}$ ). Given  $\mu^1$  and  $\mu^0$  in  $U$  let  $\pi^l$  be a Radon-Nikodym derivative of  $p^l \mu^l$  w.r.t.  $\mu^* \equiv p^1 \mu^1 + p^0 \mu^0$  and  $\rho^l(z) = p^l \mu^l([z, 1]) / \mu^*([z, 1])$ . Note that if  $\mu^1(dz)$  is interpreted as the probability of “playing Top at level  $z$ ” given  $l = 1$ , then  $\pi(z)$  is the posterior probability on  $P$  if  $\theta$  arises at  $z$  and  $\rho(z)$  is the posterior probability, given Bottom up to this level  $z$ . Define also a pay-off function on  $P \times Q \times U^2 \times \bar{\mathcal{C}} \times \bar{\mathcal{D}} \times V^2$  by:

$$\varphi(p, q; \mu^1, \mu^0; \bar{\gamma}, \bar{\delta}; \nu^1, \nu^0) = \sum_m q^m \varphi^m(p; \mu^1, \mu^0; \bar{\gamma}, \bar{\delta}; \nu^m)$$

with  $\varphi^m(p; \mu^1, \mu^0; \bar{\gamma}, \bar{\delta}; \nu^m) =$

$$\begin{aligned} \int_0^1 \nu^m([z, 1]) \{ zv(\pi(z), m) + (1-z)(\pi^1(z)\bar{\gamma}^m(z) + \pi^0(z)\bar{\delta}^m(z)) \} \mu^*(dz) \\ + \int_0^1 \mu^*([z, 1]) J(z, m) \nu^m(dz) \end{aligned}$$

and where  $J(z, m) = \min_{0 \leq u \leq z} \{ ug_{22}(\rho(z), m) + (1-u)g_{21}(\rho(z), m) \}.$

In order to state the result there remains to introduce:

$$\begin{aligned} \underline{\Phi}(p, q; \mu^1, \mu^0) &= \sup_{\bar{\mathcal{C}} \times \bar{\mathcal{D}}} \inf_{V^2} \varphi(p, q; \mu^1, \mu^0; \bar{\gamma}, \bar{\delta}; \nu^1, \nu^0) \\ \bar{\Phi}(p, q; \mu^1, \mu^0) &= \inf_{V^2} \sup_{\bar{\mathcal{C}} \times \bar{\mathcal{D}}} \varphi(p, q; \mu^1, \mu^0; \bar{\gamma}, \bar{\delta}; \nu^1, \nu^0) \end{aligned}$$

Then we have:

PROPOSITION 3.3.  $\underline{\Phi}(p, q; \mu^1, \mu^0) = \bar{\Phi}(p, q; \mu^1, \mu^0)$  on  $P \times Q \times U^2$ .

PROOF. Remark first that  $\varphi$  depends upon  $\bar{\gamma}$  and  $\bar{\delta}$  only through their values at the finitely many points  $\{z_i\}$ ,  $i = 1, \dots, R$  in the union of the supports of  $\mu^0$  and  $\mu^1$ . Hence we can replace  $\bar{\mathcal{C}}$  by the convex compact set  $(\mathcal{C} \cap [-C, C])^R$  and similarly for  $\bar{\mathcal{D}}$ . Note now that  $V$  is convex and  $\varphi$  affine with respect to  $(\bar{\gamma}, \bar{\delta})$ . We can then apply theorem 1.6 p. 4 to get  $\bar{\Phi} = \underline{\Phi}$ . ■

REMARK 3.5. We shall use later the fact that a best reply to  $\nu$  minimises

$$(5) \quad \bar{\gamma}^1(z)q^1\nu^1([z, 1]) + \bar{\gamma}^0(z)q^0\nu^0([z, 1])$$

with  $\bar{\gamma}(z) \in \mathcal{C}$ . It follows that  $\bar{\gamma}(z)$  is a supporting hyperplane for  $v(1, .)$  at the point  $q(z)$  with:  $q^1(z) = q^1\nu^1([z, 1]) / \{q^1\nu^1([z, 1]) + q^0\nu^0([z, 1])\}$ . According to the previous interpretation this corresponds to the posterior probability that player I computes given  $\theta$  and  $z$  (i.e. the conditional on  $M$ , given the event  $\{u \geq z\}$ ).

THEOREM 3.4.  $\underline{v}$  exists on  $P \times Q$  and is given by:

$$\underline{v}(p, q) = \sup_{U^2} \underline{\Phi}(p, q; \mu^1, \mu^0)$$

PROOF.

PART A. Player II can defend  $\underline{v}$ .

Given  $\varepsilon > 0$  and  $\sigma$  strategy of player I, let  $R = 1/\varepsilon$ ,  $z_r = r/R$ ,  $r = 0, \dots, R$  and define:

- $\tau(0) = \text{play always L}$
- $Q^l(0) = P_{\sigma^l, \tau(0)}(\theta < \infty)$ ,  $l = 1, 0$ ,

then  $n(0)$  and  $P^l(0)$  such that

$$P^l(0) = P_{\sigma^l, \tau(0)}(\theta \leq n(0)) \geq Q^l(0) - \varepsilon, \quad l = 1, 0.$$

Introduce now inductively on  $r$ , given  $\tau(r-1)$  and  $n(r-1)$ :

- $\mathcal{T}(r)$ : the set of strategies that coincide with  $\tau(r-1)$  up to stage  $n(r-1)$  and such that  $P_\tau(t_n = R) \leq z_r, \forall n$ .
- $Q^l(r) = \sup_{\tau \in \mathcal{T}(r)} P_{\sigma^l, \tau}(\theta < \infty)$ ,  $l = 1, 0$

then  $\tau(r) \in \mathcal{T}(r)$ ,  $n(r) \geq n(r-1)$  and  $P^l(r)$  such that:

$$P^l(r) = P_{\sigma^l, \tau(r)}(\theta \leq n(r)) \geq Q^l(r) - \varepsilon/2^r, \quad l = 1, 0.$$

Let  $\mu^l$  be the measure in  $U$  with mass  $P^l(r) - P^l(r-1)$  at point  $z_r$ ,  $l = 0, 1$ , and let  $\bar{\nu}$  in  $V^2$  be an  $\varepsilon C$ -optimal strategy for player II for the pay-off  $\bar{\Phi}(p, q; \mu^1, \mu^0)$ . We finally introduce  $\nu$  as an atomic approximation of  $\bar{\nu}$ :  $\nu^m$  has a mass  $\bar{\nu}([z_{r-1}, z_r])$  at point  $z_r$ .

The strategy  $\tau$  of player II is now described as:

choose  $z_r$  according to  $\nu^m$ , given  $m = 0, 1$ .

- play  $\tau(r)$  up to stage  $\theta$  and optimally thereafter in the revealed game.

In order to compute the expected pay-off at some stage  $n \geq n(R)$  we first study the different events induced by  $a_\theta$ . Recall first that by construction the event  $\{\exists r, \tau(r)\}$  is played and  $n \geq \theta > n(r)\}$  has a probability less than  $2\varepsilon$ , so that we will work on its complement. Since  $P_{\sigma^l, \tau(r)}(\theta > n(r)) = \mu^l([z_r, 1])$ , we can compute the posterior probabilities on  $L$  and get:

$$P_{\sigma, \tau(r)}(l = 1 \mid \theta > n(r)) = \rho^1(z_r)$$

Now for  $j \geq r$   $P_{\sigma^l, \tau(j)}\{n(r-1) < \theta \leq n(r)\} = \mu^l([z_{r-1}, z_r])$ , so  $P_{\sigma, \tau(j)}\{l = 1 \mid n(r-1) < \theta \leq n(r)\} = \pi^1(z_r)$ . Similarly for the posteriors on  $M$  induced by  $\nu$ :

$$P\{m = 1 \mid n(r-1) < \theta \leq n(r)\} = P\{m = 1 \mid \tau(j) \text{ is played, with } j \geq r\} = q^1(z_r).$$

Define still  $u_r = P\{t_\theta = R \mid n(r-1) < \theta \leq n(r)\}$  and recall that

$$(6) \quad u_r \leq z_r \quad \text{a.s.}$$

One obtains thus the following description:

- if  $\theta > n$  and player II plays  $\tau(r)$ , his posterior probability on  $L$  is  $\rho(z_r)$  and he can minimise in  $J(z_r, m)$ .
- if  $n(r-1) < \theta \leq n(r)$ ,  $P(t_\theta = R) = u_r$ . The posterior probabilities at stage  $\theta$  on  $L \times M$  are, given T:  $(1, q(z_r))$  and given L:  $(\pi(z_r), 1)$ .

It follows that for  $n$  large enough:

$$\begin{aligned} n \sum_{n(R)}^{n(R)+n} \gamma_n(\sigma, \tau) &\leq \sum_{r=0}^R \left( u_r \left\{ \sum_m q^m \nu^m([z_r, 1]) v(\pi(z_r), m) \right\} \left\{ \sum_l p^l \mu^l([z_{r-1}, z_r]) \right\} \right. \\ &\quad + (1 - u_r) \left\{ \sum_m q^m \nu^m([z_r, 1]) \right\} \left\{ \sum_l p^l \mu^l([z_{r-1}, z_r]) v(l, q(z_r)) \right\} \\ &\quad \left. + \mu^*(z_r, 1) \left\{ \sum_m J(z_r, m) q^m \nu^m([z_{r-1}, z_r]) \right\} \right) + 4C\varepsilon \end{aligned}$$

Using (4), (5) and (6), there exists  $(\bar{\gamma}, \bar{\delta}) \in \bar{\mathcal{C}} \times \bar{\mathcal{D}}$  such that:

$$n \sum_{n(R)}^{n(R)+n} \gamma_n(\sigma, \tau) \leq \varphi(p, q; \mu^1, \mu^0; \bar{\gamma}, \bar{\delta}; \nu^1, \nu^0) + 4C\varepsilon.$$

By the choice of  $\nu$  and  $\bar{\nu}$  it follows that:

$$\begin{aligned} n \sum_{n(R)}^{n(R)+n} \gamma_n(\sigma, \tau) &\leq \varphi(p, q; \mu^1, \mu^0; \bar{\gamma}, \bar{\delta}; \bar{\nu}^1, \bar{\nu}^0) + 4C\varepsilon + C\varepsilon \\ &\leq \bar{\Phi}(p, q; \mu^1, \mu^0) + 6C\varepsilon \end{aligned}$$

Finally there exists  $\bar{N}$  such that  $n \geq \bar{N}$  implies:

$$\bar{\gamma}_n(\sigma, \tau) \leq \bar{\Phi}(p, q; \mu^1, \mu^0) + 7C\varepsilon$$

hence claim A.

**PART B. Player I can guarantee v.**

Let us first choose  $\mu$  that realises  $\sup_{U^2} \underline{\Phi}(p, q; \cdot, \cdot)$  up to  $\varepsilon$  and  $\bar{\gamma}, \bar{\delta}$ ,  $\varepsilon$ -optimal in the corresponding  $\underline{\Phi}(p, q; \mu^1, \mu^0)$ . We shall write  $Z$  for a finite set  $\{z_r \mid r = 0, \dots, R\}$  that contains the support of  $\mu^l, l = 1, 0$  (we assume  $\varepsilon \geq (z_r - z_{r-1}) \geq 0$ ),  $\zeta_r^l$  for  $\mu^l(z_r)$  and  $\gamma_r$ , (resp.  $\delta_r$ ) for  $\bar{\gamma}(z_r)$  (resp.  $\bar{\delta}(z_r)$ ).

Let us consider a family of stochastic games  $\Gamma^*(z)$  described by:

$$\begin{pmatrix} -z^* & (1-z)^* \\ z & -(1-z) \end{pmatrix}$$

and write  $\sigma(z)$  for an  $\varepsilon$ -optimal strategy of player I in it, i.e. such that for all  $n \geq N(z)$  and all  $\tau$ :

$$(7) \quad \bar{\gamma}_n^*(\sigma, \tau) \geq -\varepsilon$$

Let  $N_0 = \max\{N(z_r)\}$ . We first define a family of stopping times by:  $\theta(0) = 1$  and inductively  $\theta(r)$  is  $\theta$  induced by  $\sigma(z_r)$  from stage  $\theta(r-1)$  on. Let also  $\theta(R+1) = \infty$ . The strategy  $\sigma$  of player I is now as follows: given  $l$ , choose  $r \in Z$  with probability  $\zeta_r^l$  or  $R+1$  with probability  $1 - \mu^l(Z)$ . Play then Bottom up to stage  $\theta(r), 0 \leq r \leq R+1$ , and Top at stage  $\theta(r)$ . Obviously after  $\theta$ , player I plays optimally in the revealed game if  $a_\theta = L$  or  $R$  and approaches  $\gamma^r$  (resp.  $\delta_r$ ) if  $T$  (resp.  $B$ ), with an  $(N_1, \varepsilon)$ -strategy. Let  $N = \max\{N_0, N_1\}$ .

We shall now prove that player I can guarantee, for each  $m$ ,  $\inf_z \varphi^m(p; \mu^1, \mu^0; \bar{\gamma}, \bar{\delta}; \omega_z)$  where  $\omega_z$  is the Dirac mass at  $z$ . By the properties of  $\mu$  and  $(\bar{\gamma}, \bar{\delta})$  and the definition of  $\underline{v}$  the result will follow. Given  $n$  and  $\tau$ , strategy of player II we define:

$$\begin{aligned} I(r) &= \{j \mid \theta(r-1) \wedge n < j \leq \theta(r) \wedge n\} & I_r &= \#I(r) \\ u_r &= \mathbb{1}_{t_\theta(r)=R} & \bar{u}_r &= (1/I_r) \sum_{I(r)} \mathbb{1}_{t_j=R} \end{aligned}$$

Hence player I is using  $\sigma(z_r)$  during  $I_r$  stages on the bloc  $I(r)$ , where  $(u_r, \bar{u}_r)$  describe the behaviour of player II. Note also that if  $\{\theta = \theta(r)\}$  with  $a_\theta = L$ , the posterior on  $L$  is  $\pi(z_r)$  (defined through  $\mu$ ) and  $\rho(z_r)$  if  $\{\theta > \theta(r)\}$ . We obtain thus that  $n\bar{\gamma}_n^{p1}(\sigma, \tau^1) \geq$

$$(8) \quad \mathbb{E}_{s,\tau} \left[ \sum_{r=0}^{R+1} I_r \left( \sum_{j=0}^{r-1} \{u_j(p^1\zeta_j^1 + p^0\zeta_j^0)v(\pi(z_j), 1) + (1-u_j)(p^1\zeta_j^1\gamma_j^1 + p^0\zeta_j^0\delta_j^1)\} \mathbb{1}_{I_r \geq N} \right) \right. \\ \left. + [p^1(1 - \sum_1^{r-1} \zeta_j^1) + p^0(1 - \sum_1^{r-1} \zeta_j^0)] [g_{22}(\bar{u}_r\rho(z_{r-1}, 1) + (1-\bar{u}_r)g_{21}(\rho(z_{r-1}, 1))] \right] - C\varepsilon$$

By the choice of  $\sigma(z_r)$  we obtain on  $I(r)$  that  $\mathbb{E}\{I_r \bar{u}_r\} \leq N + \mathbb{E}(I_r)(z_r + \varepsilon)$ , hence the expectation of the last term in (8) is minorated by  $J(z_{r-1}, 1) - ((R+1)N + 2\varepsilon)C$ . Using again (4), we obtain:

$$(9) \quad n\bar{\gamma}_n^{p1}(\sigma, \tau^1) \geq \mathbb{E} \sum_{r=1}^{R+1} I_r U^1(z_r) + D.C - ((R+1)N + 2\varepsilon)C$$

where  $U^1(z)$  stands for  $\varphi^m(p; \mu^1, \mu^0; \bar{\gamma}, \bar{\delta}; \omega_z)$  and  $D = \mathbb{E}\{\sum_r I_r \sum_{j=0}^{r-1} \zeta_j(u_j - z_j)\}$ . Remark that  $u_j - z_j = u_j(1 - z_j) + (1 - u_j)(-z_j)$  is the absorbing pay-off in  $\Gamma^*(z_j)$ , that this pay-off occurs at stage  $\theta(j)$  and last for  $\sum_{j+1}^{R+1} I_r$  stages. By the choice of  $\sigma(z_j)$  we thus have:

$$\mathbb{E} \sum_j I_j(u_j - z_j) \geq -\varepsilon N, \quad \text{so that } D \geq -\varepsilon N(R+1).$$

We get now from (9):  $n\bar{\gamma}_n^{p1}(\sigma, \tau^1) \geq \min_r U^1(z_r) - C(R+1)N - n(R+3)\varepsilon$ .

A similar result for  $\bar{\gamma}_n^{p0}$  implies finally that for all  $\varepsilon > 0$ , there exists  $N^*$  and  $\sigma$  such that, for all  $n \geq N^*$ ,  $\bar{\gamma}_n(\sigma, \tau) \geq \inf_{V^2} \varphi(p, q; \mu^1, \mu^0; \bar{\gamma}, \bar{\delta}; \nu^1, \nu^0) - \varepsilon_0$ , hence the result. ■

#### 4. Stochastic games with incomplete information

We will consider in this section a family of two-person zero-sum stochastic games with incomplete information on one side and full monitoring, described by a set of states  $K$ , a probability  $p$  on  $K$  and for each  $k$  a matrix  $G^k$  with absorbing pay-offs where moreover the set of absorbing entries is independent of  $k$ . One can view the set of states as  $K \times L$ , with incomplete information on one side on  $K$  and complete information on  $L$ ; the first component is chosen at the beginning of the game and kept fixed and the transition on  $L$  is independent of it.

**4.a. A first class.** Here the game is described by matrices  $G^k$  with  $\#S = 2$  where the first line is absorbing.

4.a.1. *Minmax.* As we already saw in sect. 3 when considering games with state dependent signalling matrices, the minmax of the infinitely repeated game equals the value of the one-shot game (prop. 3.1):

PROPOSITION 4.1.  $\bar{v}(p) = v_1(p)$

Before looking at the maxmin let us consider the  $n$ -stage game  $\Gamma_n$ .

4.a.2. *Lim  $v_n$ .* We first remark that the recursive formula (3.2 p. 158) can be written as:

$$(n+1)v_{n+1}(p) = \max_x \min_y \left\{ (n+1) \sum_k p^k x_1^k G_1^k y + \sum_k p^k x_2^k G_2^k y + n\bar{x}_2 v_n(p_2) \right\}$$

where  $G_s^k$  is row  $s$  of the matrix  $G^k$ ,  $y \in Y = \Delta(T)$ ,  $x = \{x^k\}$  with  $x^k$  in  $X = \Delta(S)$ ,  $\bar{x}_2 = \sum_k p^k x_2^k$ , and  $p_2$  is the conditional probability on  $K$ , given  $p$ ,  $x$  and the move  $s = 2$ .

It follows that the value of  $\Gamma_n$  is the same if both players use strategies independent of the histories: From the above formula player I can compute inductively an optimal strategy that depends only on the posterior on  $K$  (i.e. on his previous random choice in  $X^K$  and move in  $S$ ), hence is independent of the moves of II; now against such a strategy, player II cannot do better than conditioning on his own previous moves and finally as soon as player I plays  $s = 1$ , the game is over. It suffices thus to define the strategies at each stage, conditionally to a sequence of  $s = 2$  up to that stage. This remark allows us to approximate  $\Gamma_n$  for  $n$  large by a game in continuous time on  $[0, 1]$  (we do not need a family as in subsection 2.d).

More precisely, a strategy of player II being a sequence  $(y_1, \dots, y_n)$  in  $Y$ , it will be represented by a measurable mapping  $f$  from  $[0, 1]$  to  $Y$ :  $f_t(\omega)$  is the probability of move  $t$  at time  $\omega$ . Similarly, a strategy of player I being a  $K$ -vector of sequences in  $X$  will be described by a family of  $K$  positive Borel measures of mass less than one on  $[0, 1]$ , say  $\rho^k$ , where  $\rho^k([0, \omega])$  denoted by  $\underline{\rho}^k(\omega)$ , is the probability of playing move 1 before time  $\omega$  in game  $k$ . Given  $f$  and  $\rho^k$ , the pay-off in game  $k$  will be absorbing from time  $\omega$  on, with probability  $\rho^k(d\omega)$  and value  $G_1^k f(\omega)$ , and non-absorbing at that time with probability  $1 - \underline{\rho}^k(\omega)$  and value  $G_2^k f(\omega)$ . Denoting by  $F$  and  $Q$  the corresponding sets for  $f$  and each  $\rho^k$ , we now can define a game  $\Gamma(p)$  on  $F \times Q^K$  with pay-off:  $\varphi(f, \{\rho^k\}) = \sum_k p^k \varphi^k(f, \rho^k)$  where

$$\varphi^k(f, \rho^k) = \int_0^1 \left[ (1 - \omega) G_1^k f(\omega) \rho^k(d\omega) + G_2^k f(\omega) (1 - \underline{\rho}^k(\omega)) \ell(d\omega) \right]$$

and  $\ell$  stands for the Lebesgue measure on  $[0, 1]$ .

LEMMA 4.2.  $\Gamma(p)$  has a value  $\nu(p)$ .

PROOF. Denote by  $\underline{\nu}(p)$  and  $\bar{\nu}(p)$  the maxmin and minmax of  $\Gamma(p)$  and similarly  $\underline{\nu}'(p)$  and  $\bar{\nu}'(p)$  when player II's strategy set is reduced to  $F'$ , the set of continuous functions from  $[0, 1]$  to  $Y$ . Since, for  $f$  in  $F'$ ,  $\varphi^k(f, \cdot)$  is continuous when  $Q$  is endowed with the weak topology, for which it is compact, and moreover  $\varphi^k$  is affine in each variable, prop. 1.8 p. 5 implies that  $\underline{\nu}'(p) = \bar{\nu}'(p)$ . Obviously  $\bar{\nu}'(p) \geq \bar{\nu}(p)$ , hence it is enough to prove that  $\underline{\nu}(p) \geq \underline{\nu}'(p)$ . For each  $\rho$  in  $Q$  and  $f$  in  $F$  there exists by Lusin's theorem a sequence in  $F'$  converging to  $f$ , a.e. w.r.t.  $\rho$  and  $\ell$ . Hence the result follows by Lebesgue's dominated convergence theorem. ■

We now prove that  $\Gamma$  is a good representation of  $\lim \Gamma_n$  (and  $\lim \Gamma_\lambda$ ).

**THEOREM 4.3.**  $\lim v_n$  and  $\lim v_\lambda$  exist and both are equal to  $\nu$  on  $\Pi$ .

PROOF. We first show  $\limsup v_n(p) \leq \bar{\nu}'(p)$ . Let  $f$  in  $F'$  be  $\varepsilon$ -optimal for player II in  $\Gamma(p)$  and (uniform continuity) choose  $n$  large enough to have:  $\|f(\omega) - f(\omega')\| \leq \varepsilon$  for  $|\omega - \omega'| \leq 1/n$ . Let player II use in  $\Gamma_n$  the following strategy  $\tau = (y_1, \dots, y_n)$  with  $y_i = f((i-1)/n)$ . By the previous remarks, it is enough to consider a pure strategy  $\sigma$  of player I in  $\Gamma_n$  defined by a sequence of moves in  $S$ ; so let  $i_k + 1$  be the first time it induces the move 1 against  $\tau$  in game  $k$ , and take  $i_k = n$  if only the move 2 is played. The corresponding pay-off is then:

$$n\bar{\gamma}_n^k(\sigma, \tau) = \sum_{i=1}^{i_k} G_2^k f((i-1)/n) + ((n-i_k)/n) G_1^k f(i_k/n)$$

so that  $|\bar{\gamma}_n^k(\sigma, \tau) - \varphi^k(\pi^k, f)| \leq \varepsilon$ , with  $\pi^k$  Dirac mass on  $i_k/n$ . Letting  $\pi = (\pi^k)$  we obtain:

$$\bar{\gamma}_n(\sigma, \tau) \leq \varphi(\pi, f) + \varepsilon \leq \bar{\nu}'(p) + 2\varepsilon,$$

hence the claim.

We now prove  $\liminf v_n(p) \geq \underline{\nu}(p)$ . Take  $\rho$  optimal for player I in  $\Gamma(p)$  (compactness). For each  $k$ , define a sequence  $\sigma^k = (x_1^k, \dots, x_n^k)$  with values in  $X$  such that,  $\theta$  denoting the stopping time at which player I plays 1 for the first time, one has:  $P_{\sigma^k}(\theta \leq i) = \underline{\rho}^k(i/n)$ ,  $i = 1, \dots, n$ , and let player I use  $\sigma = (\sigma^k)$  in  $\Gamma_n(p)$ . For each  $\tau = (y_1, \dots, y_n)$  take a step function  $f$  in  $F$  satisfying  $f(\omega) = y_1$  on  $[0, 1/n]$  and  $f(\omega) = y_i$  on  $((i-1)/n, i/n]$  for  $i = 2, \dots, n$ . We obtain thus:

$$n\bar{\gamma}_n^k(\sigma, \tau) = \sum_{i=1}^n \left[ \{\underline{\rho}^k(i/n) - \underline{\rho}^k((i-1)/n)\} G_1^k y_i(n-i+1) + \{1 - \underline{\rho}^k(i/n)\} G_2^k y_i \right]$$

hence  $|\bar{\gamma}_n^k(\sigma, \tau) - \varphi^k(\rho^k, f)| \leq 2C/n$ , so that for  $n$  large enough:

$$\bar{\gamma}_n(\sigma, \tau) \geq \varphi(\rho, f) - \varepsilon \geq \underline{\nu}(p) - \varepsilon.$$

Finally it is easy to extend these results to  $\Gamma_\lambda$ : just replace the above uniform partition of  $[0, 1)$  by the following:  $\{[\omega_n, \omega_{n+1})\}_{n \in \mathbb{N}}$ , with  $\omega_0 = 0$  and  $\omega_n = \sum_{i=1}^n \lambda(1-\lambda)^{i-1}$ . ■

**4.a.3. Maxmin.** Consider now the maxmin. The main result of this section is the following:

**THEOREM 4.4.**  $\underline{\nu}$  exists and equals  $\nu$  on  $\Pi$ .

**COMMENT 4.1.** This result means that player I can play as well in the infinitely repeated game as in a large finite game or, conversely, that he cannot in large games take advantage of the fact that they are finite. This property also holds for the second class we will consider (cf. subsection 4.b) and in fact is a conjecture for all games where player I's information  $\sigma$ -field is finer than that of his opponent, i.e.  $\mathcal{H}^{\text{II}} \subseteq \mathcal{H}^{\text{I}}$ .

PROOF. The proof is rather long and split into two parts. We first prove:

**PART A. Player II can defend  $\nu$ .**

Notice that the stochastic aspect of the game prevents us to use the same proof as in theorem 3.1, where player II could even guarantee  $\lim v_n$ . A preliminary result amounts to remark that the pay-off in the auxiliary game  $\Gamma(p)$  is the average between 0 and 1 of the expected pay-off at time  $\omega$ . Namely let:

$$\Phi^k(\rho^k, f, \omega) = \int_0^\omega G_1^k f(\omega') \rho^k(d\omega') + (1 - \underline{\rho}^k(\omega)) G_2^k f(\omega),$$

and  $\Phi = \sum_k p^k \Phi^k$ , then one has (recall that  $\ell$  is Lebesgue measure):

$$\text{LEMMA 4.5. } \varphi^k(\rho^k, f) = \int_0^1 \Phi^k(\rho^k, f, \omega) \ell(d\omega).$$

PROOF. Using Fubini's theorem, the initial definition for  $\varphi^k$  and the above formula are both equal to:  $\iint_{0 \leq \omega' \leq \omega \leq 1} [G_1^k f(\omega') \rho^k(d\omega') + (1 - \underline{\rho}^k(\omega)) G_2^k f(\omega)] \ell(d\omega)$ . ■

To construct a uniformly good reply of player II in large games, we shall use the following procedure. Given  $f$  optimal in  $\Gamma$  and  $\sigma$  strategy of I, player II can compute the probability of absorption if he follows  $f$ , hence represent  $\sigma$  as a measure on the path defined by  $f$ . Since, by the choice of  $f$ , the (time) average pay-off is less than  $\nu$ , there exists an initial path on which the pay-off is at most  $\nu$ , and by keeping  $f$  constant thereafter player II can in fact achieve this pay-off. So let us start with an  $\varepsilon$ -optimal strategy  $f$  for player II in  $\Gamma(p)$ . We first remark that  $f$  can be chosen to be a step function, more precisely there exists a finite family of points  $\omega_r$  in  $[0, 1]$  and of values  $f_r$  in  $Y$  with:

- $\omega_1 = 0 \leq \dots \leq \omega_{R+1} = 1$
- $f(\omega) = f_r$  on  $[\omega_r, \omega_{r+1})$ ,  $r = 1, \dots, R$ .

Given  $\sigma$  and  $\varepsilon > 0$ , define inductively strategies  $\tau_r$ , measures  $\bar{\mu}_r^k$  and  $\mu_r^k$  and natural numbers  $N_r$  as follows (recall that  $\theta$  denote the stopping time of the first  $s = 1$ ):

- $\tau_1$  is: play  $f_1$  i.i.d.

For each  $k$  define:  $\bar{\mu}_1^k = \Pr_{\sigma, \tau_1}(\theta < \infty)$ , and let  $\mu_1^k$  and  $N_1$  satisfy:  $\mu_1^k = \Pr_{\sigma, \tau_1}(\theta \leq N_1) \geq \bar{\mu}_1^k - \varepsilon$ .

- Similarly  $\tau_r$  is: play  $\tau_{r-1}$  up to stage  $N_{r-1}$  (included) and then  $f_r$  i.i.d.

Then we let  $\bar{\mu}_r^k = \Pr_{\sigma, \tau_r}(\theta < \infty)$ .  $N_r$  and  $\mu_r^k$  satisfy  $N_r \geq N_{r-1}$  and  $\mu_r^k = \Pr_{\sigma, \tau_r}(\theta \leq N_r) \geq \bar{\mu}_r^k - \varepsilon$ , for all  $k$ . Define positive atomic measures  $\rho^k$  on  $[0, 1]$  by  $\rho^k(\{\omega_r\}) = \mu_r^k - \mu_{r-1}^k$  so that, by definition of  $f$ ,  $\varphi(\rho, f) \leq \nu(p) + \varepsilon$ . This implies that for some  $\omega$  in  $[0, 1)$   $\Phi(\rho, f, \omega) \leq \nu(p) + \varepsilon$ . So let  $r$  be such that  $\omega \in [\omega_r, \omega_{r+1})$  and remark that  $\Phi(\rho, f, \omega) = \Phi(\rho, f, \omega_r)$ .

We now claim that by playing  $\tau_r$ , player II can get an asymptotic pay-off near  $\nu$ . In fact, for  $n \geq N_r$ , the pay-off at stage  $n$  in game  $k$  will be of the form:

$$\gamma_n^k(\sigma, \tau_r) = \sum_{m=1}^r \alpha_m^k G_1^k f_m + (1 - \sum_{m=1}^r \alpha_m^k) G_2^k f_r$$

with  $\alpha_m^k = P_{\sigma, \tau_r}(N_{m-1} < \theta \leq N_m)$  for  $m < r$  (with  $N_0 = 0$ ) and  $\alpha_r^k = P_{\sigma, \tau_r}(N_{r-1} < \theta \leq n)$ . Since  $\mu_r^k - \mu_{r-1}^k \leq \alpha_r^k \leq \mu_r^k - \mu_{r-1}^k + \varepsilon$ , we obtain  $|\gamma_n^k(\sigma, \tau_r) - \Phi^k(\rho, f, \omega_r)| \leq 2C\varepsilon$ , hence, averaging on  $k$ ,  $|\gamma_n(\sigma, \tau_r) - \Phi(\rho, f, \omega_r)| \leq 2C\varepsilon$ , and finally  $\bar{\gamma}_n(\sigma, \tau_r) \leq \nu(p) + 3C\varepsilon + \varepsilon$  for  $n \geq 2N_R/\varepsilon$ .

This proves claim A.

**PART B. Player I can guarantee  $\nu(p)$ .**

The idea of the proof relies basically on two facts: first, there exists a pair of “equalising” strategies  $(\rho, f)$  in  $\Gamma$  such that the pay-off at  $\omega$  is constant; the second point is that player I can adapt his strategy, essentially the stopping time  $\theta$ , to the empirical frequency of moves of II, such that the pay-off in  $\Gamma_n$  corresponds to the one induced by  $\rho$  and  $f$ , if II follows  $f$ , and is less otherwise. We first prove a preliminary result. Essentially it means that given an optimal strategy of player I, there exists a best reply of player II equalising in time, i.e. such that the corresponding pay-off at  $\omega$  is constant between 0 and 1.

**PROPOSITION 4.6.** *Let  $\rho$  be optimal in  $\Gamma(p)$ . There exists  $f \in F$  s.t.  $\Phi(\rho, f, \omega) = \nu(p)$ , for all  $\omega$  in  $[0, 1]$ .*

**PROOF.** Let  $\rho_\varepsilon$  be a non-atomic  $\varepsilon$ -optimal strategy for player I in  $\Gamma(p)$ . We consider an auxiliary game  $\mathcal{G}(\rho_\varepsilon)$  where player I chooses at random a point  $\omega$  in  $[0, 1]$  and player II chooses a function in  $F'$ . The corresponding pay-off is  $\Phi(\rho_\varepsilon, f, \omega)$ . This game has a value  $w_\varepsilon$ . Indeed the strategy set of player I, resp. II, is convex and compact, resp. convex. Moreover the mapping  $f \mapsto \Phi(\rho_\varepsilon, f, \omega)$  is affine and the mapping  $\omega \mapsto \Phi(\rho_\varepsilon, f, \omega)$  is continuous. Obviously one has  $w_\varepsilon \geq \nu - \varepsilon$ , since I can use  $\ell$  to choose  $\omega$  and then the pay-off is precisely  $\varphi(\rho_\varepsilon, f)$ . Let us prove that  $w_\varepsilon \leq \nu$ . In fact, let  $m$  be an optimal (compactness) strategy of I so that  $\int_0^1 \Phi(\rho_\varepsilon, f, \omega) m(d\omega) \geq w_\varepsilon$ , for all  $f$  in  $F'$ . Replacing  $w_\varepsilon$  by  $w_\varepsilon - \delta$ , we can assume that  $\underline{m}(\omega) = m([0, \omega])$  is a strictly increasing continuous function from  $[0, 1]$  to itself with  $\underline{m}(0) = 0$  and  $\underline{m}(1) = 1$ . We can now use  $\underline{m}$  to rescale the time, namely we define  $\tilde{\rho}$  in  $Q$  and  $\tilde{f}$  in  $F$  by  $\tilde{\rho}(\underline{m}(\omega)) = \rho_\varepsilon(\omega)$  and  $\tilde{f}(\underline{m}(\omega)) = f(\omega)$ . Hence we obtain:

$$\int_0^1 \Phi(\rho_\varepsilon, f, \omega) m(d\omega) = \int_0^1 \Phi(\tilde{\rho}, \tilde{f}, \omega) \ell(d\omega)$$

Since  $m$  defines a one-to-one mapping on  $F'$  this gives:  $\varphi(\tilde{\rho}, \tilde{f}) \geq w_\varepsilon - \delta$ , for all  $f$  in  $F'$ , hence  $w_\varepsilon - \delta \leq \nu$ .  $\delta$  being arbitrary, the inequality follows. Let now for each  $\varepsilon = 1/n$ ,  $\rho_\varepsilon = \rho_n$  and let  $f_\varepsilon = f_n$  in  $F'$  with  $\Phi(\rho_n, f_n, \omega) \leq \nu + 1/n$ , for all  $\omega$  in  $[0, 1]$ , and let  $\rho_n$  converge weakly to  $\rho$ :  $\Phi(\rho_n, f_n, \omega)$  converges to  $\Phi(\rho, f_n, \omega)$  for all  $\omega$ . Finally let  $f$  in  $F$  such that  $\Phi(\rho, f_n, \omega)$  converges to  $\Phi(\rho, f, \omega)$  for all  $\omega$ , so that  $\Phi(\rho, f, \omega) \leq \nu$ , hence the equality since  $\rho$  is optimal in  $\Gamma(p)$ . ■

Consider now  $\rho$  and  $f$  as above and let  $\underline{\omega} < 1$  be such that  $\rho^k([\underline{\omega}, 1]) \leq \varepsilon/4$ , for all  $k$ . Note that this implies:

$$(1) \quad \sum_k p^k (1 - \rho^k(\underline{\omega})) G_1^k y \geq \sum_k p^k (1 - \rho^k(\underline{\omega})) G_1^k f(\underline{\omega}) - \varepsilon C/2$$

for all  $y$  in  $Y$ . In fact otherwise, one obtains with  $g = f$  on  $[0, \underline{\omega}]$  and  $= y$  on  $[\underline{\omega}, 1]$  that  $\Phi(\rho, g, \omega) = \Phi(\rho, f, \omega)$  on  $[0, \underline{\omega}]$  and  $\Phi(\rho, g, \omega) < \nu$  on  $[\underline{\omega}, 1]$ , contradicting the optimality of  $\rho$ . Similarly on each atom of  $\sum_k p^k \rho^k$ , say  $\omega$ , one has:

$$(2) \quad \sum_k p^k \rho^k(\{\omega\}) G_1^k f(\omega) = \min_y \sum_k p^k \rho^k(\{\omega\}) G_1^k y$$

otherwise, by modifying  $f$  in some neighbourhood  $O$  of  $\omega$ , one obtains a  $g$  satisfying  $\Phi(\rho, g, \cdot) < \Phi(\rho, f, \cdot)$  on  $O$  and equality a.e. otherwise, contradicting again the optimality of  $\rho$ . Given  $\eta = (1 - \underline{\omega})\varepsilon C/4$ , let us now introduce a partition  $\omega_1, \dots, \omega_{R+1}$  with  $\omega_1 = 0$ ,  $\omega_R = \underline{\omega}$ ,  $\omega_{R+1} = 1$ , and an adapted pair  $\tilde{f}, \tilde{\rho}$  with:

$$(3) \quad \rho^k(\{\omega_r\}) = \tilde{\rho}^k([\omega_r, \omega_{r+1}]) = \tilde{\rho}_r^k \quad \text{and} \quad \tilde{f}(\omega) = \tilde{f}_r \text{ on } [\omega_r, \omega_{r+1}], \text{ with } \tilde{f}_R = f(\underline{\omega})$$

$$(4) \quad \text{if } \sum_k p^k \rho^k(\{\omega\}) > \varepsilon, \text{ then } \omega \in \{\omega_r\}, \tilde{\rho}^k(\{\omega\}) - \rho^k(\{\omega\}) \leq \varepsilon^2 \text{ and } \tilde{f}_r = f(\omega_r)$$

$$(5) \quad \tilde{\rho}^k(\omega) - \underline{\rho}^k(\omega) \leq \varepsilon, \quad \forall \omega$$

$$(6) \quad \tilde{\rho} \text{ is } \eta\text{-optimal in } \Gamma, \quad \text{and } \Phi(\tilde{\rho}, \tilde{f}, \omega) \leq \nu + \eta \text{ on } [0, \underline{\omega}]$$

Obviously  $\Phi(\tilde{\rho}, \tilde{f}, \cdot)$  is piecewise constant on  $[0, 1]$  and equal to:

$$\sum_k p^k \left[ \sum_{j=1}^r \tilde{\rho}_j^k G_1^k \tilde{f}_j + (1 - \sum_{j=1}^r \tilde{\rho}_j^k) G_2^k \tilde{f}_r \right] \quad \text{on } [\omega_r, \omega_{r+1})$$

Now we claim that for any  $y = \{y_r\}_{r < R}$  with  $y_r$  in  $Y$ :

$$(7) \quad \sum_k p^k (1 - \tilde{\rho}^k(\omega_r)) G_2^k y_r \leq \sum_k p^k (1 - \tilde{\rho}^k(\omega_r)) G_2^k \tilde{f}_r + \eta$$

implies:

$$\sum_k p^k \sum_{j=1}^r \tilde{\rho}_j^k G_1^k y_j \geq \sum_k p^k \sum_{j=1}^r \tilde{\rho}_j^k G_1^k \tilde{f}_j - \varepsilon C, \quad \forall r < R.$$

Otherwise defining  $g$  as  $y_j$  on  $[\omega_j, \omega_{j+1})$  if  $\sum_k p^k \tilde{\rho}_j^k G_1^k y_j < \sum_k p^k \tilde{\rho}_j^k G_1^k \tilde{f}_j$  and  $\tilde{f}$  otherwise would give  $\Phi(\tilde{\rho}, g, \omega) < \nu + \eta$  on  $[0, \underline{\omega}]$ , by (6), and  $\Phi(\tilde{\rho}, g, \omega) < \nu + \varepsilon C/2 - \varepsilon C$  on  $[\underline{\omega}, 1)$ , by the choice of  $\rho, \underline{\omega}$  and (3). So that  $\varphi(\tilde{\rho}, g) < \nu + \eta - \varepsilon C(1 - \underline{\omega})/2 < \nu - \eta$  contradicting by (6) the choice of  $\tilde{\rho}$ .

We introduce finally a strategy  $\sigma$  for player I by letting  $\sigma^k$  be: play  $s = 1$  (for the first time) at stage  $\theta_r$ , with probability  $\tilde{\rho}_r^k$ , where the stopping times  $\theta_r$  are inductively defined by the following procedure:

We first consider the probability of absorption. If  $\sum_k p^k \tilde{\rho}_1^k \geq 2\varepsilon$ , this implies that  $\sum_k p^k \rho^k$  is atomic at  $\omega_1$  and we let  $\theta_1 = 1$ . Otherwise, we compute the expected non-absorbing pay-off induced by  $\tilde{f}$ , i.e.  $\psi_1(\tilde{f}) = \sum_k p^k (1 - \tilde{\rho}_1^k) G_2^k \tilde{f}_1 = z_1$ , and we consider an optimal strategy  $\alpha_1$  in an associated stochastic game with absorbing states, pay-off  $\psi_1$  and level  $z_1$  (cf. ex. VIIEx.4 p.346 and VIIEx.19 p.351). More precisely, if  $\theta$  denotes the stopping time of absorption, there exists  $N_1$  such that for all  $n \geq N_1$  and every (pure) strategy of II,

$$(8) \quad \psi_1(\bar{t}_n) \leq z_1 - C\varepsilon \Rightarrow P_{\alpha_1}(\theta \leq n) \geq 1 - \varepsilon, \quad \text{with } \bar{t}_n = (1/n) \sum_{m=1}^n t_m,$$

$$(9) \quad [\mathbb{E}_{\alpha_1}(\psi_1(t_\theta) \mid \theta \leq n) - z_1] P_{\alpha_1}(\theta \leq n) \leq \eta\varepsilon$$

Define similarly, if  $\sum_k p^k \tilde{\rho}_r^k \geq 2\varepsilon$ ,  $\theta_r = \theta_{r-1} + 1$ ; otherwise let  $\psi_r(\tilde{f}) = \sum_k p^k (1 - \sum_{j=1}^r \tilde{\rho}_j^k) G_2^k f(\tilde{\omega}_r) = z_r$ , and let  $\alpha_r$  and  $N_r$  be the corresponding optimal strategies and bound on the number of stages.  $\theta_r$  then follows the law of  $\theta$  induced by  $\alpha_r$  from stage  $\theta_{r-1}$  on.

Let us compute the pay-off given  $\sigma$  and some pure strategy  $\tau$  of II (i.e. a sequence of moves in  $T$ ), at some stage  $n \geq RN$ , with  $N = \max_r N_r$ . We first obtain a (random) sequence of blocs  $B_j, j = 1, 2, \dots, r, \dots$ , where player I uses  $\alpha_1$ , then  $\alpha_2$  and so on. We shall approximate the average pay-off on each of these blocs, except when their length is smaller than  $N$  or when the expectation of occurrence is too small:  $P(\theta_r \leq n - \sum_{j < r} \theta_j) \leq \varepsilon$ . We first notice that on  $B_r, r = 1, \dots, R$ , since the length of the bloc is greater than  $N$  and  $\alpha_r$  is used, (8) implies that with probability greater than  $1 - \varepsilon$ :

$$\psi(t(r)) \geq z_r - C\varepsilon, \quad \text{where } t(r) = (1/\#B_r) \sum_{m \in B_r} t_m$$

Remark now that the expected absorbing pay-off at  $\theta_r$ , hence on  $B_{r+1}$  is of the form:

$$\sum_k p^k \sum_{j \leq r} \tilde{\rho}_j^k G_1^k y_j, \quad \text{with } y_j = \mathbb{E}(t_{\theta_j})$$

If  $\omega_r$  is an atom of  $\sum_k p^k \rho^k$  one obtains  $\sum_k p^k \tilde{\rho}_r^k G_1^k y_r \geq \sum_k p^k \tilde{\rho}_r^k G_1^k \tilde{f}_r - 2\varepsilon^2 C$  by (2) and (4), and there are at most  $1/\varepsilon$  such points. Otherwise  $\alpha_r$  is used and if  $\Pr(\theta_r \leq n - \sum_{j < r} \theta_j) \geq \varepsilon$  one has  $\psi(y_r) \leq z_r + \eta$  by (9), so that the absorbing pay-off is at least

$$\sum_k p^k \sum_{j \leq r} \tilde{\rho}_j^k G_1^k \tilde{f}_j - \varepsilon C \quad \text{by (7)}$$

hence by (5) greater than  $\sum_k p^k \sum_{j \leq r+1} \tilde{\rho}_j^k G_1^k \tilde{f}_j - 3\varepsilon C$ . Finally on  $B_{R+1}$ , using (1) and (3), the new absorbing pay-off is less than  $\varepsilon C/4$  and the non-absorbing one minorated by  $\sum_k p^k (1 - \tilde{\rho}^k(\omega)) G_1^k \tilde{f}(\omega) - \varepsilon C/2$ . Hence the expected pay-off at stage  $n$  is majorated by a convex combination of terms of the form:

$$\Phi(\tilde{\rho}, \tilde{f}, \omega_r) - 2CRN/n - (1/\varepsilon)\varepsilon^2 4C - 8\varepsilon C$$

and the result follows from (6). ■

This completes the proof of theorem 4.4.

**4.b. A second class.** A second family of games that we will study here is given by matrices  $G^k$  with  $S$  lines and  $T$  columns but where the first column is absorbing. It is easily seen (compare exercises of ch. VII) that the values of the (stochastic)  $n$  stage, discounted or infinite game where player I use non-revealing strategies are equal and we will write  $u(p)$  for this common value. As in the previous section we will prove here that  $\underline{v}(p)$ ,  $\lim v_n$  and  $\lim v_\lambda$  exist and are equal. Note that in the current framework also, there is no direct way of proving that player II can defend  $\liminf v_n$ . Nevertheless the proof will be roughly similar to the previous one.

4.b.1. *Maxmin and  $\lim v_n$ .*

**THEOREM 4.7.**  $\underline{v}(p) = \lim v_n = \lim v_\lambda = \mathbf{Cav} u(p)$

**PROOF.** We first remark that player I can concavify as usual (cf. e.g. 1.2 p. 184), so that  $v_n$  and  $v_\lambda$  are greater than  $\mathbf{Cav} u$  and player I can guarantee  $\mathbf{Cav} u$ . (Note that he can even get  $u(p)$  at each stage). We prove now that player II can defend  $\mathbf{Cav} u$ . Basically he will play a best reply to the expected strategy of player I in the non-revealing game at the current posterior after each stage except when the expected variation of this martingale is large. This last event having a small expected frequency the result will follow. Let us first consider  $\Gamma_\infty$  and assume that player I is using  $\sigma$ . Let us write  $\mathcal{T}'$  for the set of player II's strategies with values having a support included in the set  $T \setminus \{1\}$  of non-absorbing columns. Define now  $\tau'$  in  $\mathcal{T}'$  and  $N$  in  $\mathbb{N}$  satisfying:

$$(10) \quad \mathsf{E}_{\sigma, \tau'} \sum_{m=1}^N (p_{m+1} - p_m)^2 \geq \sup_{\tau \in \mathcal{T}'} \mathsf{E}_{\sigma, \tau} \sum_{m=1}^\infty (p_{m+1} - p_m)^2 - \varepsilon$$

and consider  $\tau$  defined by: Play according to  $\tau'$  up to stage  $N$  included and thereafter play at each stage  $m$  a best reply to  $\mathsf{E}(\sigma(h_m) | h_m)$  in  $D(p_m)$ . Note that (10) implies that  $\mathsf{E}_{\sigma, \tau} (\sum_{m=N+1}^\infty (p_{m+1} - p_m)^2) \leq \varepsilon$ . We want to majorate the average pay-off up to some stage  $n \geq N$ . Letting  $\theta'$  be the first time where player II is using move 1 and defining  $\theta = \min(\theta', n)$ , we obtain (as in sect. 4 p. 298):

$$(11) \quad n\bar{\gamma}_n(\sigma, \tau) \leq 2NC + \mathsf{E}\left(\sum_{m=1}^\theta u(p_m) + (n-\theta)u(p_{\theta+1})\right) + CX(p)$$

with  $X(p) = \mathsf{E}(\sum_{N+1}^\theta |p_{m+1} - p_m| + (n-\theta) |p_{\theta+1} - p_\theta|)$ . Since  $\mathsf{E}(\sum_{m=1}^\theta p_m + (n-\theta)p_{\theta+1}) = p$  by the martingale property, the second term is majorated by  $n \mathbf{Cav} u(p)$  (using Jensen's

inequality). As for  $X(p)$ , it can be written as  $\mathbb{E}(\sum_{m=N+1}^{\infty} |p_{m+1} - p_m| Z_m)$ , where

$$Z_m = \begin{cases} 0 & \text{for } m > \theta \\ n - m + 1 & \text{for } m = \theta \text{ and } p_m \text{ constant for } m > \theta \\ 1 & \text{for } m < \theta \end{cases}$$

Hence,

$$\begin{aligned} X(p) &\leq \mathbb{E}\left[\sum(p_{m+1} - p_m)^2 \cdot \sum(Z_m^2)\right]^{1/2} \\ &\leq [\mathbb{E}(\sum(p_{m+1} - p_m)^2)]^{1/2} [\mathbb{E} \sum(Z_m^2)]^{1/2} \\ &\leq \varepsilon^{1/2} n. \end{aligned}$$

We obtain thus  $\bar{\gamma}_n(\sigma, \tau) \leq \text{Cav } u(p) + C(2N/n + \varepsilon^{1/2})$ , and the result follows.

As for  $\Gamma_n$ , recall that for any  $\sigma$  one has:  $\mathbb{E}_{\sigma, \tau} \sum_{m=1}^n (p_{m+1} - p_m)^2 \leq L$ , uniformly in  $\tau$  (cf. 2.1 p. 186), hence the number of stages in  $\Gamma_n$  where  $\mathbb{E}(p_{m+1} - p_m)^2 \geq L/n^{3/4}$  is at most  $n^{3/4}$ . On the other stages, say  $m$  in  $M$ , the probability of the set of histories  $h_m$ , where  $\mathbb{E}((p_{m+1} - p_m)^2 | h_m) \geq L/n^{1/2}$ , is less than  $n^{-1/4}$ . Now define  $\tau$  as being a pure best reply to  $\mathbb{E}(\sigma(h_m) | h_m)$  in  $D(p_m)$  at each stage  $m$  in  $M$  where the variation  $\mathbb{E}((p_{m+1} - p_m)^2 | h_m)$  is smaller than  $L/n^{1/2}$  and any non-absorbing move otherwise. It follows, using the same majoration than in (11) that:

$$n\bar{\gamma}_n(\sigma, \tau) \leq 2C(n^{3/4} + n \cdot n^{-1/4}) + \text{Cav } u(p) + C((nL)^{1/2} + nL^{1/2}/n^{1/4})$$

and this finishes the proof.

Finally for  $\Gamma_\lambda$ , let  $N = \lambda^{-3/4}$ . Then the number of stages where  $\mathbb{E}(p_{m+1} - p_m)^2 \geq L/N$  is less than  $N$ . On the complement of this set of stages the histories on which the conditional quadratic variation is larger than  $L/N^{2/3}$  have a probability less than  $N^{1/3}$ . Now the weight of these stages is at most  $(1 - (1 - \lambda)^{N+1})$ , which is of the order of  $\lambda N$ , hence we obtain a majoration of  $\bar{\gamma}_\lambda(\sigma, \tau) - \text{Cav } u$  by a term of the order of  $\lambda^{1/4}$ . ■

4.b.2. *Minmax.* We assume from now on  $T = \{1, 2\}$ .

As in the previous case and as in the next ones,  $\bar{v}$  will be obtained through an auxiliary game, or more precisely here through a sequence of auxiliary games, as in sect. 3. (For an alternative approach in special cases, leading to an explicit formula, cf. ex. VIIIEx.8, VIIIEx.10 and VIIIEx.11 p. 395). For each  $L$  in  $\mathbb{N}$  define  $G_L$  by the following strategy sets  $Z_L$  for I and  $\mathcal{F}_L$  for II and pay-off  $\Psi_L$  ( $G_L$  should be  $G_L(p)$  but we will keep  $p$  fixed during the whole section, hence we drop it):  $Z_L = (\Delta_L(S))^{K \times L}$ , where  $\Delta_L(S)$  is the triangulation of the simplex  $\Delta(S)$  with mesh  $1/L$  and  $\mathcal{F}_L$  is the set of mappings  $f$  from sequences of length less or equal to  $L$  in  $\Delta_L(S)$  to half space  $s$  in  $\mathbb{R}^S$ . The pay-off corresponding to a pair  $(z, f)$  in  $Z_L \times \mathcal{F}_L$  is  $\Psi_L(z, f) = \sum_k p^k \Psi_L^k(z^k, f)$ , with:

$$\Psi_L^k(z^k, f) = (1/L) \sum_{m=1}^{\ell-1} z_m^k G_1^k + (1 - (\ell-1)/L) z_\ell^k G_2^k$$

where  $\ell = \min(\{m \mid z_m^k \notin f(z_1^k, z_2^k, \dots, z_{m-1}^k)\} \cup \{L+1\})$ .  $G_L$  is thus a matrix game (only finitely many different  $f$  give different pay-offs) with value  $w_L$ . The interpretation is that player I is playing i.i.d. on blocs (his strategy can obviously be assumed independent of the moves of player II since it is enough to define it at each stage  $n$ , on the event  $\theta \geq n$ , where  $\theta$  is the first time II uses  $t = 1$ , recall that  $T = \{1, 2\}$ ), and player II uses on each

bloc an optimal strategy in the stochastic game with absorbing states defined by the half space at that stage.

THEOREM 4.8.  $\lim w_L$  and  $\bar{v}$  exist and are equal.

PROOF.

PART A. I can defend  $\limsup w_L$ .

We will show that given  $z$  for player I in  $G_L$  and  $\tau$  for player II in  $\Gamma_\infty$ , there exists  $f$  and  $\sigma$  such that for  $n$  large enough  $\bar{\gamma}_n(\sigma, \tau)$  is near  $\psi_L(z, f)$ . Given  $\tau$ , strategy of player II, we first can assume that one has for all  $n$  and all strategies of I,  $P(\theta = n) \leq 1/L^2$ ; in fact, it is easy to see that if player II can guarantee  $d$ , she can also guarantee it with such a strategy. We now define a probability induced by  $\tau$  on a set of  $f$  in  $\mathcal{F}_L$  (it is sufficient to describe their intersection with  $\Delta_L(S)$ ). Given  $z$  in  $\Delta_L(S)$  let  $\zeta(z) = P_{\sigma(z), \tau}(\theta < \infty)$ , where  $\sigma(z)$  is play  $z$  i.i.d. Define  $\tau(\emptyset) = \{z \in \Delta_L(S) \mid \zeta(z) \geq 1/L\}$ , and for  $z$  in  $\tau(\emptyset)$ , let  $N(z)$  such that  $P_{\sigma(z), \tau}(\theta \leq N(z)) \in [1/L, 1/L + 1/L^2]$ . If  $z \notin \tau(\emptyset)$ , let  $N(z) = \infty$ .

We first define  $f(\emptyset)$ : On the extreme points of  $\Delta_L(S)$ ,  $\mathbb{1}_{f(\emptyset)} = \mathbb{1}_{\tau(\emptyset)}$ . On the one-dimensional faces of  $\Delta_L(S)$ ,  $z$  belongs to  $f(\emptyset)$  iff there exists an extreme point  $z'$  and some  $z''$  both in  $\tau(\emptyset)$  with  $z$  on  $[z', z'']$ .  $f(\emptyset)$  is now defined on the whole simplex by the half space that coincides with it on the previous one dimensional faces.

We now define  $f(\{z_1\})$ . If  $z_1 \notin f(\emptyset)$ ,  $f(\{z_1\})$  is arbitrary. If  $z_1 \in f \cap \tau(\emptyset)$  we first introduce  $\zeta(z_1, z) = P_{\sigma(z_1, z), \tau}(\theta < \infty)$  where  $\sigma(z_1, z)$  is the strategy of I defined by  $z_1$  i.i.d. up to stage  $N(z_1)$  and then  $z$  i.i.d. As above let  $\tau(z_1) = \{z \mid \zeta(z_1, z) \geq 1/L\}$  and for  $z$  in this set let  $N(z_1, z)$  be such that  $\Pr_{\sigma(z_1, z), \tau}(N(z_1) < \theta \leq N(z_1, z)) \in [1/L, 1/L + 1/L^2]$ . Now  $f(z_1)$  is defined from the set  $\tau(z_1)$  exactly as  $f(\emptyset)$  from  $\tau(\emptyset)$ . Finally if  $z_1 \in f(\emptyset) \setminus \tau(\emptyset)$ , there exists  $(\#S)$  points  $z_1^i$  in  $f \cap \tau(\emptyset)$  such that  $z_1$  is a barycentre,  $z_1 = \sum_i \lambda^i z_1^i$ .  $\tau$  defines then a mixture at  $z_1$ : play  $f(z_1^i)$  with probability  $\lambda^i$ . For the general construction, given a sequence  $(z_1, \dots, z_\ell)$ , consider first the (random) sequence that II has generated: say  $(z_1^{i_1}, \dots, z_j^{i_j}, \dots)$  as long as  $z_{j+1} \in f(z_1^{i_1}, \dots, z_j^{i_j})$ , (otherwise  $f$  is from then on arbitrary), the  $z_j^{i_j}$  being defined from  $z_j$  as above. On each path of length  $j$  we apply the same construction, introducing first  $\zeta$ , then  $\tau$  and  $N$  and finally  $f$ . Note that we have described this way a behavioural strategy  $F_\tau$  on  $\mathcal{F}_L$ . Given any array  $z = (z_1, \dots, z_L)$  in  $\Delta_L(S)$ , we introduce a (non-revealing) strategy  $\sigma_z$  for I in  $\Gamma_\infty$  such that for  $n$  large enough the average pay-off against  $\tau$  if  $k$ ,  $\bar{\gamma}_n^k(\sigma_z, \tau)$  will be near  $\int \psi_L^k(z, f) F_\tau(df) = \psi_L^k(z, F_\tau)$ . If  $z_1 \notin f(\emptyset)$ ,  $z_1$  can also be written as a barycentre of points  $y_1^i \notin \tau(\emptyset) \cup f(\emptyset)$ , say  $z_1 = \sum_i \mu^i y_1^i$ .  $\sigma_z$  is then: play with probability  $\mu^i$ ,  $y_1^i$  i.i.d. for ever. The corresponding expected pay-off is then obviously (at each stage)  $z_1 G_2^k = \psi^k(z, F_\tau)$  up to  $2C/L$  (corresponding to the probability that  $\theta$  will be finite). If  $z_1 \in f(\emptyset)$ , consider the same points  $z_1^i$  introduced above and let  $\sigma_z$  satisfy: with probability  $\lambda^i$  play  $z_1^i$  i.i.d. up to stage  $N(z_1^i)$ .

Now to define  $\sigma_z$  at some further stage we first consider the (random) sequence generated up to now say  $z_1^{i_1}, \dots, z_j^{i_j}$ :

- either  $z_{j+1} \notin f(z_1^{i_1}, \dots, z_j^{i_j})$  then one introduces points  $y_{j+1}^i \notin f \cup \tau(z_1^{i_1}, \dots, z_j^{i_j})$  and  $\sigma_z$  is, from  $N(z_1^{i_1}, \dots, z_j^{i_j})$  on, the corresponding splitting of i.i.d. sequences;
- or  $z_{j+1} \in f(z_1^{i_1}, \dots, z_j^{i_j})$  and one defines the points  $z_{j+1}^i \in \tau \cap f(z_1^{i_1}, \dots, z_j^{i_j})$  as above and  $\sigma_z$  is, with probability  $\lambda^i$  (obviously function of  $(z_1^{i_1}, \dots, z_j^{i_j})$ ), play  $z_{j+1}^i$  i.i.d. from stage  $N(z_1^{i_1}, \dots, z_j^{i_j}) + 1$  to  $N(z_1^{i_1}, \dots, z_{j+1}^i)$ .

Thus for  $n \geq \underline{N} = L \times \max\{N(z_1, \dots, z_L) \mid z \in \Delta_L(S)^L\}$ , one obtains that the average pay-off up to stage  $n$  is an expectation of terms of the form (with  $\ell$  being the first time where  $z_j$  is not in  $f$ ):

$$(1/L) \sum_{j < \ell} z_j^{i_j} G_1^k + (1 - (\ell - 1)/N) y_\ell^i G_2^k, \quad \text{up to } 2C(L \times 1/L^2 + 1/L + 1/L),$$

where the first term corresponds to the error on the absorbing pay-offs, the second is for the non-absorbing pay-off and the third takes care of the stages up to  $N(z_1^{i_1}, \dots, z_{\ell-1}^{i_{\ell-1}})$ . Now the expectation of  $y_\ell^i$  is precisely  $z_\ell$ , and moreover, by construction the probabilities of the sequences  $(z_1^{i_1}, \dots, z_j^{i_j})$  are the same under  $\sigma_z$  and  $F_\tau$ . It follows that for  $n \geq \overline{N}$ :  $\bar{\gamma}_n(\sigma_z, \tau) \geq \psi^k(z, F^\tau) - 6C/L$ . Finally, given  $\varepsilon > 0$ , let  $L_0 > 6/\varepsilon$  such that  $w_{L_0} \geq \limsup w_L - \varepsilon$  and let  $\chi$  be an optimal strategy for I in  $G_{L_0}$ . A strategy  $\sigma$  in  $\Gamma_\infty$  is then defined by: choose  $z = (z^k)$  according to  $\chi$ ; given  $k$  and  $z$ , use the above strategy  $\sigma_{z^k}$ . We thus obtain, for  $n \geq \overline{N}$ :

$$\bar{\gamma}_n(\sigma, \tau) = \mathbb{E}_\chi \sum_k p^k \bar{\gamma}_n(\sigma_{z^k}, \tau) \geq \mathbb{E}_\chi \sum_k p^k \psi_{L_0}^k(z^k, F_\tau) - 6C/L_0$$

hence  $\bar{\gamma}_n(\sigma, \tau) \geq \psi_{L_0}(\chi, F_\tau) - \varepsilon \geq \limsup w_L - 2\varepsilon$ .

#### PART B. II can guarantee $\liminf w_L$ .

We will first represent a strategy  $f$  in  $G_L$  as a strategy  $\tau(f)$  in  $\Gamma_\infty$  and then show that for  $n$  large enough the pay-off it induces against some  $\sigma$  (in  $\Gamma_\infty$ ) is near a pay-off corresponding to  $f$  and some  $z(\sigma)$  in  $G_L$ . The choice of an  $L$  realising the liminf up to some  $\varepsilon$ , and then of a strategy  $\tau$  associated to an optimal mixture (with finite support) in  $G_L$  then implies the result.

We shall proceed as in part B of the previous sect. 4.a.3, and given  $f$ , construct a strategy that corresponds to a sequence of optimal strategies in some auxiliary games with absorbing states. The computations being roughly similar we will mainly describe the procedure. To each half space  $f(z_1, \dots, z_j)$  is associated a strategy  $\tau(z_1, \dots, z_j)$  such that:  $P(\theta \leq n)$  is near 1 as soon as  $n \geq M(z_1, \dots, z_j)$  and the empirical frequency of moves of I up to  $n$  is at a distance at most  $1/L$  from  $f(z_1, \dots, z_j)$ ; moreover  $\mathbb{E}(s_\theta \mid \theta \leq m)$  is with probability near one within  $1/L^2$  of  $f(z_1, \dots, z_j)$ , as soon as the probability of the event  $\{\theta \leq m\}$  is not too small. Given a (pure) strategy of I (i.e. in this case a sequence of moves), let us introduce a sequence of stopping times.  $\theta_1$  follows the law of  $\theta$  under  $\tau(\emptyset)$ . Further  $\omega_1$  is such that  $P(\theta_1 \leq \omega_1)$  is near 1 (and  $\infty$  if no such number exists). Finally if  $\omega_1$  is finite, let  $x_1 = \mathbb{E}(s_\theta \mid \theta_1 \leq \omega_1)$  and choose  $z_1$  as a closest point to  $x_1$  in  $\Delta_L(S)$ .  $\theta_2$  follows the law of  $\theta$  under  $\tau(z_1)$  from stage  $\omega_1 + 1$  on. We define similarly  $\omega_2$ ,  $x_2$  and  $z_2$ , and  $\theta_j$  inductively up to  $j = L$ .  $\tau$  is obtained by choosing  $j$  at random, uniformly in  $\{1, \dots, L\}$  and playing  $t = 1$  for the first time at stage  $\theta_j$ . Hence for  $n$  large enough (like  $L^2 \times \max_z M(z_1, \dots, z_L) = L^2 \overline{M}$ ) the average expected pay-off up to stage  $n$  given  $(s_1, \dots, s_n)$  and  $\tau$  will be, for any  $k$ , near some average of  $\Psi^k(y, f)$ , with  $y$  in  $\Delta_L(S)$ . In fact considering the sequence  $\omega_j$ , one obtains when  $\omega_{j+1} - \omega_j$  is large ( $\geq \overline{M}$ ) that the average frequency of moves of player I during these stages is with probability near one, within  $1/L$  of a point  $y_{j+1}$  not in  $f$ . Moreover with a probability near 1 the sequence  $z_1, \dots, z_j$  is compatible with  $f$ , i.e. one has  $z_i \in f(z_1, \dots, z_{i-1})$ . Since the probability of each event  $\{\theta = \theta_i\}$  is  $1/L$  we obtain that the average pay-off on the stages in  $(\omega_j, \omega_{j+1}]$  is roughly (adding another error term of the order of  $1/L$ ), given  $k$ :  $(1/L) \sum_{i \leq j} z_i G_1^k + (1 - (j+1)/L) y_{j+1} G_2^k$ , hence in the range of  $\Psi^k(\cdot, f)$ .

It remains to see that the total weight of the “small blocs” is at most  $L \times \overline{M}$  so that by taking expectation over  $k$ ,  $\overline{\gamma}_n(\sigma, \tau)$  will be near the range of  $\psi(\cdot, f)$  on  $Z_L$ . ■

**4.c. Minmax: two more examples.** Recall we conjecture that  $\lim v_n$ ,  $\lim v_\lambda$  and  $\max \min \underline{v}$  exist and coincide in games where one player is more informed than his opponent.

The purpose of this section is thus to give examples where on the other hand the existence of the minmax is established, again through an auxiliary game, but moreover with an explicit description.

4.c.1. *Example A.* Let  $G^k, k \in K$  be a finite set of  $2 \times 2$  pay-off matrices of the form  $G^k = \begin{pmatrix} a^{k*} & b^k \\ c^k & d^k \end{pmatrix}$ , where as usual the star \* denotes an absorbing pay-off.

4.c.1.i. An auxiliary game.  $G(p)$  is the (one-shot) game in normal form defined by  $A^K$  (resp.  $B$ ) as strategy set of player I (resp. II) and pay-off  $f$  where:  $A = \overline{\mathbb{N}} \cup \{\partial\}$  ( $\overline{\mathbb{N}}$  is the compactification  $\mathbb{N} \cup \{\infty\}$  of the set of positive integers  $\mathbb{N}$  and  $\partial$  is some isolated point with  $\partial > \infty$ ),  $B = \{0, 1\}^{\mathbb{N}}$  and the pay-off  $f$  is the average of the state pay-off  $f^k$ ,  $f(\alpha, \beta) = \sum_k p^k f^k(\alpha^k, \beta)$  with finally:

$$\begin{aligned} f^k(n, \beta) &= a^k \left(1 - \prod_0^{n-1} \beta'_m\right) + \left(\prod_0^{n-1} \beta'_m\right) (\beta_n c^k + \beta'_n d^k) \quad \text{for } n \in \mathbb{N} \\ f^k(\infty, \beta) &= a^k \left(1 - \prod_0^\infty \beta'_m\right) + \left(\prod_0^\infty \beta'_m\right) d^k \\ f^k(\partial, \beta) &= a^k \left(1 - \prod_0^\infty \beta'_m\right) + \left(\prod_0^\infty \beta'_m\right) b^k \end{aligned}$$

where  $\beta = (\beta_0, \dots, \beta_m, \dots) \in B$ ,  $\beta'_m$  denotes  $1 - \beta_m$ , and  $\prod_0^{-1} = 1$ . Defining  $\theta(\beta) = \min\{m \mid \beta_m = 1 \text{ or } m = \infty\}$ , it is clear that  $\theta(\beta)$  determines the pay-off, hence  $B$  can also be written as  $\overline{\mathbb{N}}$ . Then one has, with  $\xi \in A$ :

$$\begin{aligned} f^k(\xi, m) &= \mathbb{1}_{\xi \leq m-1} d^k + \mathbb{1}_{\xi=m} c^k + \mathbb{1}_{\xi > m} a^k \quad \text{for } m \in \mathbb{N}, \quad \text{and} \\ f^k(\xi, \infty) &= \mathbb{1}_{\xi \leq \infty} d^k + \mathbb{1}_{\xi=\partial} b^k. \end{aligned}$$

We write  $\overline{G}(p)$  for the mixed extension of  $G(p)$  where player I’s strategies are probabilities on  $A^K$  (or as well  $K$  vector of probabilities on  $A$ , since  $f$  is decomposed on  $A^K$ ), say  $\chi$  in  $\Delta(A)^K$ , and player II’s strategies are probabilities with finite support on  $B$ , say  $\Psi$  in  $\Delta(B)$ .

PROPOSITION 4.9.  $\overline{G}(p)$  has a value  $w(p)$ .

PROOF. For each  $\beta$ ,  $f^k(\cdot, \beta)$  is continuous on  $A$ . In fact either  $\beta$  corresponds to some  $m$  in  $\mathbb{N}$  and  $f^k$  is constant on  $\alpha^k > m$ , or to  $\infty$  and  $f^k$  is constant on  $\overline{\mathbb{N}}$ . Moreover  $A^K$  is compact, hence by prop. 1.17 p. 7, the game  $\overline{G}$  has a value (and player I has an optimal strategy). ■

THEOREM 4.10.  $\min \max \Gamma(p)$  exists and  $\overline{v}(p) = w(p)$ .

The proof will follow from the two next lemmas.

Note first that for all  $\Psi \in \Delta(B)$  there exists  $\delta \in [0, 1]^{\mathbb{N}}$  s.t.  $\int f(\alpha, \beta) \Psi(d\beta) = f(\alpha, \delta)$  for all  $\alpha \in A^K$ . In fact by the above remarks  $\Psi$  can be described as the distribution of the stopping time  $\theta$ .  $\delta_n$  is then just the conditional probability on the  $n^{\text{th}}$  factor, given  $\{\theta \geq n\}$ .

LEMMA 4.11. Player II can guarantee  $w$

PROOF. Given  $\varepsilon > 0$ , let  $\Psi$  be an  $\varepsilon/4$  optimal strategy of player II for  $w(p)$  and as above represent  $\Psi$  by some  $\delta$  in  $[0, 1]^{\mathbb{N}}$ . Define now  $\tau$ , strategy of player II in  $\Gamma(p)$  as: play  $\delta_0$  i.i.d. until the first Top of player I, then play  $\delta_1$  i.i.d. until the second Top, then  $\delta_2$ , and so on. Let also  $\rho = \prod_{m=0}^{\infty} \delta'_m$  and define  $N$  such that: if  $\rho = 0$ ,  $\prod_{m=0}^{N-1} \delta'_m < \varepsilon/4$ , and if  $\rho > 0$ ,  $\prod_{m=N}^{\infty} \delta'_m > 1 - \varepsilon/4$ .

Let us majorate  $\bar{\gamma}_n(\sigma, \tau)$  for  $n \geq N$ . It is enough to consider  $\bar{\gamma}_n^k(\sigma^k, \tau)$  for each  $k$ . Since  $\tau$  is independent of the previous moves of player II, we can assume that  $\sigma^k$  has the same property (one can replace at each stage  $\sigma^k$  by its expectation w.r.t.  $\tau$  on  $J$  without changing the pay-off). Moreover we can consider best replies and assume  $\sigma^k$  pure. It follows then that  $\sigma^k$  is completely described by the dates  $M_1, \dots, M_m, \dots$  of the successive Top. We obtain thus, for  $n \in ]M_m, M_{m+1}[$ , with  $M_0 = -1$ :

$$\begin{aligned}\gamma_n^k(\sigma, \tau) &= E_{\sigma^k, \tau}(g_n^k) \\ &= a^k \left(1 - \prod_{\ell=0}^{m-1} \delta'_\ell\right) + \left(\prod_{\ell=0}^{m-1} \delta'_\ell\right) (\delta_m c^k + \delta'_m d^k) \\ &= f^k(m, \delta).\end{aligned}$$

Now for  $n \geq M_N$ , the expected stage pay-off satisfies:

- if  $\rho = 0$ ,  $|\gamma_n^k(\sigma, \tau) - a^k| \leq \varepsilon/2$ , thus  $|\gamma_n^k(\sigma, \tau) - f^k(\infty, \delta)| \leq \varepsilon/2$ .
- if  $\rho > 0$ ,
  - either player I plays Bottom and one has  $|\gamma_n^k(\sigma, \tau) - (a^k(1 - \rho) + \rho d^k)| \leq \varepsilon/2$  by the choice of  $N$ , thus:

$$|\gamma_n^k(\sigma, \tau) - f^k(\infty, \delta)| \leq \varepsilon/2,$$

– or player I plays Top and we obtain similarly:

$$|\gamma_n^k(\sigma, \tau) - f^k(\partial, \delta)| \leq \varepsilon/2.$$

It follows that every expected stage pay-off, except at most  $K \times N$  of them, corresponding to  $\{M_m\}, m = 1, \dots, N$ , for each  $\sigma^k$ , is within  $\varepsilon/2$  of a feasible pay-off against  $\Psi$  in  $\bar{G}(p)$ . Hence  $n \geq 8KN/\varepsilon$  implies:  $\bar{\gamma}_n(\sigma, \tau) \leq w(p) + \varepsilon$ . ■

LEMMA 4.12. *Player I can force  $w$ .*

PROOF. Consider first  $\chi$ , optimal strategy of player I in  $\bar{G}(p)$ . Given  $\varepsilon > 0$ ,  $\chi$ , and a strategy  $\tau$  of player II in  $\Gamma(p)$ , we shall define a strategy  $\sigma$  of player I in  $\Gamma(p)$  by the following procedure (similar to sect. 2): We introduce a family, indexed by  $A$ , of “non-revealing” strategies in  $\Gamma(p)$ , i.e. transition probabilities from  $H$  to  $S$ , say  $\mu_\alpha$ .  $\sigma^k$  will then be: select some  $\alpha$  according to the probability  $\chi^k$  on  $X$  and play  $\mu_\alpha$ . Let  $\eta$  be the stopping time of reaching the absorbing entry:  $\eta = \min\{\{m \mid i_m = T, j_m = L\} \cup \{\infty\}\}$  and define  $N$  such that  $\chi^k(N) \leq \varepsilon/3$  for all  $k$  in  $K$ .

For each  $m \leq N$ , define inductively strategies  $\mu_m$  and times  $L_m$  as follows:

- $\mu_0$  is always Bottom. Given  $\mu_0$  and  $\tau$ , let  $t_n^0 = P_{\mu_0, \tau}(j_n = L)$  and  $L_0$  be the first time  $\ell$  where:

$$\sum_k p^k \chi^k(0) (t_\ell^0 c^k + t_\ell^0 d^k) \leq \inf_n [\sum_k p^k \chi^k(0) (t_n^0 c^k + t_n^0 d^k)] + \varepsilon/3.$$

- $\mu_1$  is Bottom up to stage  $L_0$  (excluded), Top at that stage  $L_0$ , and always Bottom after.

- Similarly given  $\mu_m$ , let  $t_n^m = P_{\mu_m, \tau}(j_n = \text{Left} \mid \eta > L_{m-1})$  and let  $L_m$  be the first  $\ell > L_{m-1}$  where:

$$(12) \quad \sum_k p^k \chi^k(m) (t_\ell^m c^k + t_\ell^{m'} d^k) \leq \inf_n [\sum_k p^k \chi^k(m) (t_n^m c^k + t_n^{m'} d^k)] + \varepsilon/3$$

$\mu_{m+1}$  is then  $\mu_m$  up to stage  $L_m$  (excluded), Top at that stage and Bottom thereafter.

For  $m > N$ , we introduce a new stopping time  $L'$  and a non-revealing strategy  $\mu'$  satisfying:

$$(13) \quad \pi = P_{\mu', \tau}(\eta \leq L') \geq \sup_{\mu \in M} P_{\mu, \tau}(\eta < \infty) - \varepsilon/9$$

where  $M$  is the set of strategies that coincide with  $\mu_N$  up to  $L_{N-1}$  (included). If  $N < m \leq \infty$ , let  $\mu_m = \mu_\infty$ : play  $\mu'$  up to stage  $L'$  (included) and then Bottom for ever. Finally we define  $\mu_\partial$  as: play  $\mu'$  up to stage  $L'$  (included) and then always Top.

Let also  $\delta$  in  $[0, 1]^{\mathbb{N}}$  satisfy:

$$\begin{aligned} \delta_m &= t_{L_m}^m, \text{ for } m < N, \\ \delta_N &= u, \text{ where } \pi = 1 - (\prod_{m=0}^{N-1} \delta'_m) u' \text{ (note that } 1 \geq \pi \geq 1 - \prod_{m=0}^{N-1} \delta'_m) \\ \delta_m &= 0, \text{ for } m > N, \end{aligned}$$

and we shall prove that for  $n \geq L'$ :

$$(14) \quad \gamma_n^p(\sigma, \tau) \geq \int f(\alpha, \delta) \chi(d\alpha) - 2\varepsilon/3$$

In fact we can decompose the above pay-off on the events  $\{\mu_\alpha \text{ is played}\}$ , with  $\alpha$  in  $A$ , so that:

$$\gamma_n^p(\sigma, \tau) = \sum_k p^k \gamma_n^k(\sigma^k, \tau) = \sum_k \sum_\alpha p^k P_{\sigma^k}(\mu_\alpha) \gamma_n^k(\mu_\alpha, \tau) = \sum_x \varphi_n(\alpha, \tau).$$

For  $m \leq N$  we obtain, using (12):

$$\begin{aligned} \sum_k p^k \chi^k(m) [(1 - \prod_0^{m-1} \delta'_\ell) a^k + \prod_0^{m-1} \delta'_\ell (t_n^m c^k + t_n^{m'} d^k)] \\ \geq \sum_k p^k \chi^k(m) [(1 - \prod_0^{m-1} \delta'_\ell) a^k + \prod_0^{m-1} \delta'_\ell (\delta_m c^k + \delta'_m d^k)] - \varepsilon/3 \\ \geq \sum_k p^k \chi^k(m) f^k(m, \delta) - \varepsilon/3. \end{aligned}$$

For  $N < m \leq \infty$ , we get, using (13):

$$\begin{aligned} \sum_k p^k \chi^k(m) [P_{\mu_m, \tau}(\eta \leq n) a^k + (1 - P_{\mu_m, \tau}(\eta \leq n)) (t_n^m c^k + t_n^{m'} d^k)] \\ \geq \sum_k p^k \chi^k(m) (\pi a^k + (1 - \pi) d^k) - \varepsilon/3 \end{aligned}$$

since the choice of  $\mu'$  and  $L'$  implies  $(1 - P_{\mu_m, \tau}(\eta \leq n)) t_n^m \leq \varepsilon/9$ . Similarly, when  $\mu_\partial$  is used the pay-off is at least:  $\sum_k p^k \chi^k(\partial) (\pi a^k + (1 - \pi) b^k) - \varepsilon/3$ .

It follows that for all  $m \in X$ ,  $m \neq N$ ,  $\varphi_n(m, \tau) \geq \sum_k p^k \chi^k(m) f^k(m, \delta) - \varepsilon/3$ . Since, by the choice of  $N, |\varphi_n(N, \tau)|$  as well as  $|\sum_k p^k \chi^k(N) f^k(N, \delta)|$  are bounded by  $\varepsilon/3$ , we obtain (14) by summing. Hence  $n \geq 6L'/\varepsilon$  implies  $\overline{\gamma}_n^p(\sigma, \tau) \geq w(p) - \varepsilon$ .

This completes the proof of theorem 4.10. ■

For a geometric approach with an explicit description, cf. ex. VIIIEx.9–VIIIEx.11.

4.c.2. *Example B.* We consider now the case where  $G^k = \begin{pmatrix} a^{k*} & b^k \\ c^k & d^{k*} \end{pmatrix}$ .

We will prove that  $\bar{v}$  is the value  $w(p)$  of the one-shot game with incomplete information with pay-off matrices

$$B^k = \begin{array}{c|cccccc} & \widetilde{LL} & \widetilde{LLR} & \widetilde{LR} & \widetilde{RL} & \widetilde{RR} & \widetilde{R} \\ \hline \widetilde{T\bar{T}} & a^k & a^k & a^k & a^k & a^k & b^k \\ \widetilde{T\bar{TB}} & a^k & a^k & a^k & a^k & a^k & d^k \\ \widetilde{T\bar{B}} & a^k & a^k & a^k & c^k & d^k & d^k \\ \widetilde{B\bar{T}} & a^k & a^k & b^k & d^k & d^k & d^k \\ \widetilde{B\bar{BT}} & a^k & d^k & d^k & d^k & d^k & d^k \\ \widetilde{B\bar{B}} & c^k & d^k & d^k & d^k & d^k & d^k \end{array} \quad ???$$

THEOREM 4.13.  $w(p) = \bar{v}(p)$ .

PROOF.

#### PART A. Player II can guarantee $w$ .

Let  $(\alpha_1, \alpha_2, \alpha_3; \beta_1, \beta_2, \beta_3)$  be an optimal strategy of player II in the auxiliary game. Player II uses it to choose a column and play according to the following dictionary:  $\widetilde{LL}$  is Left then always Left,  $\widetilde{LR}$  is Left then always Right and  $\widetilde{LLR}$  is Left then  $(x, 1-x)$  i.i.d. for any fixed  $x$  in  $(0, 1)$ , and similarly for the columns starting with  $R$ . Assume that player I plays Top at stage one, then his expected pay-off after stage  $n$ , where  $n$  is such that the stopping time  $\theta$  of reaching a \* is smaller than  $n$  with probability near 1 when II plays  $(x, 1-x)$ , is essentially of the form (with  $\alpha = \sum_i \alpha_i$ ):

$$\begin{aligned} \alpha a^k + \beta_1 a^k + \beta_2 a^k + \beta_3 b^k &\text{ if only Top is played (corresponding to } \widetilde{T\bar{T}}), \\ \alpha a^k + \beta_1 c^k + \beta_2 d^k + \beta_3 d^k &\text{ if only Bottom is played (corresponding to } \widetilde{T\bar{B}}), \\ \text{and } \alpha a^k + \beta_1 a^k + \beta_2 (ya^k + (1-y)d^k) + \beta_3 d^k &\text{ else.} \end{aligned}$$

Note now that if  $a^k \geq d^k$  the above pay-off is maximal for  $y = 1$  and corresponds to  $\widetilde{T\bar{B}}$ . Finally if  $a^k < d^k$  player I can obtain, by playing Bottom then  $(\varepsilon, 1-\varepsilon)$  i.i.d., a pay-off near  $\alpha_1 a^k + \alpha_2 d^k + \alpha_3 d^k + \beta_3 d^k$ , hence better and corresponding to  $\widetilde{B\bar{BT}}$ . Hence by the choice of  $\tau$ , for any  $\eta > 0$ , for  $n$  large enough and for all  $\sigma$ :  $\bar{\gamma}_n(\sigma, \tau) \leq w(p) + \eta$ .

#### PART B. Player I can defend $w$ .

Given  $\tau$  player I can compute the probability that II will play Left if he plays Top always, say  $x$ , hence by playing Top a large number of times then still Top or Bottom he can get if  $k$ , either  $xa^k + (1-x)b^k$  or  $xa^k + (1-x)d^k$ . By playing Top then always Bottom he will obtain for  $n$  large enough some pay-off of the form:  $ya^k + zc^k + (1-z-y)d^k$ , with obviously  $y + z \leq x$ . Define then  $\alpha' = y$ ,  $\beta'_1 = z$ ,  $\beta'_3 = (1-x)$ ; note that if player I starts by Bottom the pay-off is  $d^k$  with probability  $(1-y)$  at stage 1, so that the same analysis starting with Bottom allows us to define  $\alpha'_i$ .

This proves that  $\tau$  gives the same pay-off than a strategy in the auxiliary game, hence if  $\pi = (q_i^k; r_i^k)$  is an optimal strategy of player I in the auxiliary game if  $k$ , we define  $\sigma$  if  $k$  as: with probability  $q = \sum_i q_i$ , play Top at stage 1, with (total) probability  $q_3$  (that corresponds to  $\widetilde{T\bar{B}}$ ) play from then on Bottom; otherwise play Top until stage  $n$  where  $\Pr(\theta \leq n)$  is within  $\varepsilon$  of its supremum, and from then on keep playing Top (with probability  $q_1$  corresponding to  $\widetilde{T\bar{T}}$ ) or play Bottom for ever (with probability  $q_2$  corresponding

to  $\widetilde{TTB}$ ), and similarly with  $r$ . The pay-off corresponding to  $\sigma, \tau$  in  $\Gamma_\infty$  is thus near the pay-off induced by  $\pi$  and some  $(\alpha'_i, \beta'_i)$  in the auxiliary game. This completes the proof of the theorem.  $\blacksquare$

### Exercises

**1. A stochastic game with signals.** (Ferguson et al., 1970) We are given two states of nature with the following pay-off matrices:

$$G^1 = \begin{pmatrix} 1 & 0 \\ 1 & 0 \end{pmatrix} \quad G^2 = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$$

The transition from 1 to 2 is a constant  $(1 - \pi) \in (0, 1)$ , independent of the moves. From state 2, one goes to state 1 iff player I plays Bottom. Player I knows everything, player II is told only the times of the transition from 2 to 1. We will consider  $\Gamma_\lambda$ , the discounted game starting from  $k = 1$ , and write  $v_\lambda$  for its value. Let us take as state variable the number  $m^*$  of stages since the last transition from 2 to 1.

Consider the following class of strategies for player I:

$$\begin{aligned} X = \{x = (x_m) \mid x_m = \Pr(\text{play Top} \mid m^* = m)\} \quad \text{and similarly:} \\ Y = \{y = (y_m) \mid y_m = \Pr(\text{play Right} \mid m^* = m)\} \quad \text{for player II.} \end{aligned}$$

Given  $x$  and  $y$ , let  $U_m$  (resp.  $W_m$ ) be the pay-off of the  $\lambda$ -discounted game starting at  $m^* = m$  and  $k = 1$  (resp.  $k = 2$ ).

a. Prove that:

$$\begin{aligned} U_m &= (1 - y_m)\lambda + (1 - \lambda)(\pi U_{m+1} + (1 - \pi)W_{m+1}) \\ W_m &= (x_m y_m)\lambda + (1 - \lambda)(x_m W_{m+1} + (1 - x_m)U_0) \end{aligned}$$

and that this system has a unique bounded solution with:

$$U_0 = \lambda \frac{\left( \sum_{j=0}^{\infty} (1 - y_j)[\pi(1 - \lambda)]^j + (1 - \pi) \sum_{j=1}^{\infty} \{ \pi^{j-1} \sum_{m=j}^{\infty} [x_m y_m (1 - \lambda)^m \prod_{l=j}^{m-1} x_l] \} \right)}{\left( 1 - (1 - \pi)(1 - \lambda) \sum_{j=1}^{\infty} \{ \pi^{j-1} \sum_{m=j}^{\infty} [(1 - x_m)(1 - \lambda)^m \prod_{l=j}^{m-1} x_l] \} \right)}$$

b. Prove that if the game has a value on  $X \times Y$ , this value is  $v_\lambda$  (and the corresponding strategies are optimal in  $\Gamma_\lambda$ ).

c. Let  $r$  in  $\mathbb{N}$  satisfying  $\pi^{r-1} > 1/2, \pi^r \leq 1/2$  and define  $\bar{x}$  in  $X$  by:

$$\bar{x}_m = \begin{cases} 1 & m < r \\ \pi^r / (1 - \pi^r) & m = r \\ \pi / (2 - \pi) & m > r \end{cases}$$

Show that  $U_0(\bar{x}, y)$  depends only of  $(y_l), 0 \leq l < r$ , that a best reply gives  $y_l = 1, 0 \leq l < r$ , and that the corresponding pay-off is:

$$V = \frac{[1 - (1 - \lambda)^r][1 - \pi(1 - \lambda)] - \lambda\{1 - 2[\pi(1 - \lambda)]^r\}}{[1 - (1 - \lambda)^{r+1}][1 - \pi(1 - \lambda)] + 2(1 - \lambda)^{r+1}\pi^r\lambda}.$$

Define  $\bar{y}$  in  $Y$  by:

$$\bar{y}_m = \begin{cases} 1 & m < r \\ (1 - \lambda)V & m \geq r \end{cases}$$

Prove then that  $W_m(\bar{x}, \bar{y}) = (1 - \lambda)V$ , for  $m \geq r$ .

To show that  $\bar{x}$  is a best reply to  $\bar{y}$ , assume first that  $x$  coincides with  $\bar{x}$  from some stage on and then conclude. Deduce then that  $v_\lambda = V$ .

d. Let  $v_0 = \lim_{\lambda \rightarrow 0} v_\lambda = \frac{r(1-\pi)-(1-2\pi^r)}{(r+1)(1-\pi)+2\pi^r}$  and define  $y^*$  by:

$$y^* = \begin{cases} 1 & m < r \\ v_0 & m \geq r \end{cases}$$

Prove that  $\forall \varepsilon > 0$ ,  $\exists \lambda^*$  such that,  $\forall \lambda \leq \lambda^*$ ,  $\forall \sigma \in \Sigma$ ,  $\forall \tau \in \mathcal{T}$ :

$$\bar{\gamma}_\lambda(\bar{x}, \tau) \geq v_0 - \varepsilon \text{ and } \bar{\gamma}_\lambda(\sigma, y^*) \leq v_0 + \varepsilon.$$

**2. Examples with no value.** (Mertens and Zamir, 1976b; Waternaux, 1983b,a) We consider games as in sect. 2, with  $\#K = 2$ . A  $\sharp$  means that the entry is revealing hence the pay-off thereafter is 0. Note that the results are more precise than theorem 2.2 p. 363 since the auxiliary games  $\overline{G}$  or  $\underline{G}$  are finite; in particular minmax and maxmin are algebraic.

a. Let  $G^1 = \begin{pmatrix} x_{11}^\sharp & x_{12} \\ x_{21} & x_{22} \end{pmatrix}$  and  $G^2 = \begin{pmatrix} y_{11} & y_{12} \\ y_{21} & y_{22}^\sharp \end{pmatrix}$ .

Show that the maxmin is the value of the following game:

$$\overline{G} = \widetilde{G} = \begin{array}{c|ccc} & \widetilde{L} & \widetilde{R} & \widetilde{\beta} \\ \hline \widetilde{T} & y_{11} & x_{12} + y_{12} & \beta y_{11} + \beta' y_{12} \\ \widetilde{B} & x_{21} + y_{21} & x_{22} & \beta x_{21} + \beta' x_{22} \\ \widetilde{1-\varepsilon} & y_{11} & x_{12} & 0 \\ \widetilde{\varepsilon} & y_{21} & x_{22} & 0 \\ B_1 & y_{11} & x_{12} & \beta(\beta y_{11} + \beta' y_{12}) \\ T_1 & y_{21} & x_{22} & \beta'(\beta x_{21} + \beta' x_{22}) \end{array}$$

and obviously a dual result holds for the minmax.  $\widetilde{T}$  corresponds to the strategy always Top, ?? similarly for  $\widetilde{B}, \widetilde{L}, \widetilde{R}$ ;  $\widetilde{\beta}$  is playing Left with probability  $\beta$  i.i.d. and  $\widetilde{1-\varepsilon}$  corresponds to the strategy with support (Top, Bottom) and frequency of Top 1;  $T_n$  stands for  $n$  exceptional moves Top, similarly for Bottom. Finally  $\beta'$  is  $1 - \beta$ .

Show that II has an optimal strategy using a single value of  $\beta$ .

b. In the following example

$$G^1 = \begin{pmatrix} -1^\sharp & 2 \\ 2 & -4 \end{pmatrix} \quad G^2 = \begin{pmatrix} -4 & 2 \\ 2 & -1^\sharp \end{pmatrix}$$

show that  $\bar{v} = -1/2$  with optimal strategies  $(1/4, 1/4, 1/4, 1/4, 0, 0)$  for I and  $(1/4, 1/4, 1/2(\widetilde{1/2}))$  for II in  $\overline{G}$ ;  $\underline{v} = -2/3$  and  $(1/6, 1/6, 2/3(\widetilde{1/2}))$  for I and  $(1/6, 1/6, 0, 0, 1/3, 1/3)$  for II are optimal in  $\underline{G}$ .

c. Let now  $G^1 = \begin{pmatrix} x_{11}^\sharp & x_{12} \\ x_{21} & x_{22} \end{pmatrix}$  and  $G^2 = \begin{pmatrix} y_{11} & y_{12}^\sharp \\ y_{21} & y_{22} \end{pmatrix}$

Show that the minmax is the value of

$$\overline{G} = \widetilde{G} = \begin{array}{c|ccc} & \widetilde{L} & \widetilde{R} & \widetilde{\beta} \\ \hline \widetilde{T} & y_{11} & x_{12} & 0 \\ \widetilde{B} & x_{21} + y_{21} & x_{22} + y_{22} & \beta(x_{21} + y_{21}) + \beta'(x_{22} + y_{22}) \\ \widetilde{\varepsilon} & y_{21} & x_{22} & 0 \\ T_1 & y_{21} & x_{22} & \beta'(\beta x_{21} + \beta' x_{22}) + \beta(\beta y_{21} + \beta' y_{22}) \\ T_2 & y_{21} & x_{22} & \beta'^2(\beta x_{21} + \beta' x_{22}) + \beta^2(\beta y_{21} + \beta' y_{22}) \end{array}$$

and the maxmin the value of

$$\underline{G} = \widetilde{G} = \begin{pmatrix} \widetilde{L} & \widetilde{R} & \widetilde{1-\varepsilon} & \widetilde{\varepsilon} \\ \widetilde{B} & \begin{pmatrix} x_{21} + y_{21} & x_{22} + y_{22} & x_{21} + y_{21} & x_{22} + y_{22} \\ \alpha y_{11} + \alpha' y_{21} & \alpha x_{12} + \alpha' x_{22} & 0 & 0 \end{pmatrix} \end{pmatrix}$$

???

???

d. For  $G^1 = \begin{pmatrix} 7^\sharp & -3 \\ -7 & 3 \end{pmatrix}$  and  $G^2 = \begin{pmatrix} -31 & 11^\sharp \\ 31 & -11 \end{pmatrix}$ , show that  $\bar{v} = 1/4$  with  $(0, 0, 0, 2/3, 1/3)$  optimal for I and  $(0, 0, \widetilde{1/4})$  optimal for II; and  $\underline{v} = 0$ , with  $(0, \widetilde{1/2})$  and  $(0, 0, 0, 1)$  optimal for I and II.

From now on we are in the framework of sect. 4

**3.** Consider the games introduced in subsection 4.a and prove that:

$$w(p) = \inf_{F'} \sum_k p^k \max_{\omega} \Upsilon^k(f, \omega)$$

with:

$$\Upsilon^k(f, \omega) = (1 - \omega)G_1^k f(\omega) + \int_0^\omega G_2^k f(\omega') \ell(d\omega').$$

Assume  $\#K = \#S = \#T = 2$ . Show that  $w(p) = \inf_c [pc + (1 - p)J(c)]$  with

$$J(c) = \max_{\omega} \{ \Upsilon^2(f, \omega) \mid \Upsilon^1(f, \omega') \leq c \text{ for all } \omega' \}$$

Deduce the existence of optimal strategies for II and an explicit formula for them.

**4.** Let  $G^1 = \begin{pmatrix} 1^* & 0^* \\ 0 & 0 \end{pmatrix}$  and  $G^2 = \begin{pmatrix} 0^* & 0^* \\ 0 & 1 \end{pmatrix}$ . Show that:

$$u(p) = \text{Cav } u(p) = p(1 - p)$$

$$\nu(p) = (1 - p)[1 - \exp(-p/(1 - p))]$$

$$f(x) = \begin{cases} L/(1 - x) & \text{on } [0, 1 - L] \\ 1 & \text{on } [1 - L, 1] \end{cases} \quad \text{with } L = \exp(-p/(1 - p)),$$

$$\underline{\rho}^1(x) = \begin{cases} -((1 - p)/p) \ln(1 - x) & \text{on } [0, 1 - L] \\ 1 & \text{on } [1 - L, 1] \end{cases}$$

$$\underline{\rho}^2(x) \equiv 0.$$

Note that  $\underline{v}$  and  $\lim v_n$  are transcendental functions.

In the two following examples  $\rho$  has an atomic part.

**5.** Let  $G^1 = \begin{pmatrix} 4^* & -2^* \\ 0 & 0 \end{pmatrix}$  and  $G^2 = \begin{pmatrix} 0^* & 1^* \\ -1 & 2 \end{pmatrix}$ . Show that:

- for  $0 \leq p \leq 1/7$ :  $u(p) = \nu(p) = v_1(p) = 4p$
- for  $1/7 \leq p \leq 1$ :

$$\nu(p) = (1 - p)(1 - (1/3) \exp[(1 - 7p)/3(1 - p)])$$

$$f(x) = \begin{cases} 1/3 + L/(1 - x) & \text{on } [0, 1 - 6L] \\ 1/2 & \text{on } [1 - 6L, 1] \end{cases}, \quad \text{with } L = \frac{1}{9} \exp \frac{1 - 7p}{3(1 - p)},$$

$$\underline{\rho}^1(x) = \begin{cases} -((1 - p)/2p) \ln(1 - x) & \text{on } [0, 1 - 9L] \\ 1 & \text{on } [1 - 9L, 1] \end{cases},$$

$$\underline{\rho}^2(x) = \begin{cases} 0 & \text{on } [0, 1 - 9L] \\ 1 & \text{on } [1 - 9L, 1] \end{cases}$$

**6.** Let  $G^1 = \begin{pmatrix} 3^* & -1^* \\ 2 & 0 \end{pmatrix}$  and  $G^2 = \begin{pmatrix} -1^* & 1^* \\ 1 & -1 \end{pmatrix}$ . Prove that:

- for  $0 \leq p \leq 1/3$ :

$$\nu(p) = p,$$

$$f(x) = 1/2 \text{ on } [0, 1]$$

$$\underline{\rho}^1(x) = \underline{\rho}^2(x) = 1 - (1 - x)^{1/(1-3p)} \text{ on } [0, 1]$$

- for  $1/3 \leq p \leq 2/3$ :

$$\begin{aligned}\nu(p) &= (1/9)(4 - 3p) \\ f(x) &= \begin{cases} 1/2 - 2/9(1 - x^2) & \text{on } [0, 1/3] \\ 0 & \text{on } [1/3, 1] \end{cases} \\ \underline{\rho}^1(x) &= (2/3 - p)/p \text{ on } [0, 1] \\ \underline{\rho}^2(x) &= \begin{cases} [2(2/3 - p) + (3p - 1)x]/(1 - p) & \text{on } [0, 1/3] \\ 1 & \text{on } [1/3, 1] \end{cases}\end{aligned}$$

- for  $2/3 \leq p \leq 1$ :

$$\begin{aligned}\nu(p) &= p(1 - p) \\ f(x) &= \begin{cases} [1 - p^2/(1 - x)^2]/2 & \text{on } [0, 1 - p] \\ 0 & \text{on } [1 - p, 1], \end{cases} \\ \underline{\rho}^1(x) &= 0 \text{ on } [0, 1] \\ \underline{\rho}^2(x) &= \begin{cases} x/(1 - p) & \text{on } [0, 1 - p] \\ 1 & \text{on } [1 - p, 1] \end{cases}\end{aligned}$$

7. Consider the game of ex. VIIEx.4. Prove that (with  $x' = 1 - x$ ):

$$(n + 1)v_{n+1}(p) = \max_{s,t} \min_x \{(n + 1)psx + p't'x' + n(ps' + p't')v_n(ps'/(ps' + p't'))\}$$

and deduce the following heuristic differential equation for  $\lim v_n(p)$ :

$$y(p)(2 - p) = 1 - p - (1 - p)^2y'(p).$$

Adding the initial conditions leads to:  $\nu(p) = (1 - p)\{1 - \exp(-p/(1 - p))\}$ . Prove by induction, using the recursive formula, that  $v_n \geq \nu - L/n$  and  $v_n \leq \nu + L/n^{1/2}$ , for some  $L$  large enough.

8. Let  $G^1 = \begin{pmatrix} 1^* & 0 \\ 0^* & 0 \end{pmatrix}$  and  $G^2 = \begin{pmatrix} 0^* & 0 \\ 0^* & 1 \end{pmatrix}$ . Prove that:

$$\begin{aligned}\bar{v}(p) &= \inf_{\rho \in Q} \sup_{0 \leq t \leq 1} \left[ p \int_0^1 (1 - s)\rho(ds) + (1 - p)t(1 - \underline{\rho}(t)) \right] \\ &= p[1 - \exp(1 - (1 - p)/p)]\end{aligned}$$

Note that in this example the minmax is a transcendental function.

9. **A Geometric Approach to the Minmax.** Player I can *defend* a set  $D$  in  $\mathbb{R}^K$ , if for every  $\varepsilon > 0$  and every  $\tau$  there exists a strategy  $\sigma$  and a number  $N$ , such that for all  $n \geq N$  there exists  $d$  in  $D$  with  $\bar{\gamma}_n^k(\sigma, \tau) \geq d^k - \varepsilon$ , for all  $k$ . Prove that if player I can defend the half spaces  $H(p, g(p)) = \{t \in \mathbb{R}^K \mid \langle p, t \rangle \geq g(p)\}$  for all  $p \in \Pi$ , then he can also defend  $H(p, \text{Cav } g(p))$ .

Player II can *guarantee* a point  $M$  in  $\mathbb{R}^K$ , if for every  $\varepsilon > 0$  there exists a strategy  $\tau$  and a number  $N$ , such that for every  $\sigma$ ,  $n \geq N$  implies:  $\bar{\gamma}_n^k(\sigma, \tau) \leq M^k + \varepsilon$ , for all  $k$ , (i.e. he can approach  $M - \mathbb{R}_+^K$ ).

Show that the set of points that player II can guarantee is convex.

Denote by  $D_I$  the intersection of the half spaces of the form  $H(p, \alpha)$  that I can defend and by  $D_{II}$  the set of points that II can guarantee.

Prove that:  $\bar{v}(p)$  exists and equals  $\min_{d \in D} \langle p, d \rangle \iff D = D_I = D_{II}$  ( $D$  is then called the *minmax set*).

10. Let  $G^1 = \begin{pmatrix} 1^* & 0 \\ 0^* & 1 \end{pmatrix}$  and  $G^2 = \begin{pmatrix} 0^* & 1 \\ 1^* & 0 \end{pmatrix}$ . Show the minmax set is  $\text{Co}\{M^1, M^2, M\} + \mathbb{R}_+^2$  with:

$$M^1 = (1/2, 1) \quad (\text{player II plays optimal for } k=1)$$

$$M^2 = (1, 1/2) \quad (\text{player II plays optimal for } k=2)$$

$M = (2/3, 2/3)$  (player II plays once  $(1/3, 2/3)$ , then guarantees

$M^1$  (resp.  $M^2$ ) if  $s_1 = \text{Top}$  (resp. Bottom).

11. Let  $G^1 = \begin{pmatrix} 8^* & -2 \\ -4^* & 1 \end{pmatrix}$  and  $G^2 = \begin{pmatrix} -3^* & 2 \\ 6^* & -4 \end{pmatrix}$ . Prove that the minmax set is  $D' \cup D'' + \mathbb{R}_+^2$  with:

$$D' = \{(x, f(x)) \mid x \in [0, 1], f(x) = 3 - (3/4)x - (1/4)x^{1/5}\}, \quad \text{and}$$

$$D'' = \{(g(y), y) \mid y \in [0, 2], g(y) = 4 - (4/3)y - (1/3)(y/2)^{2/5}\}$$

12. Consider the games introduced in sect. 3.a p. 371 and assume  $\#K = 2$ .

Prove that the minmax set is the intersection of the positive half spaces (i.e. of the form  $H(p, \alpha)$ , cf. ex. VIIIEx.9) that contains one of the following five sets: The point  $V$ , the segments  $[C, D]$ ,  $[A, B \vee D]$ ,  $[Q, P]$  or  $[Q', P']$  where  $A$  (resp.  $B, C, D$ ) is the point  $(a^1, a^2)$ ,  $V = (v^1, v^2)$  where  $v^k$  is the value (minmax would suffice) of  $\Gamma_\infty(k)$ ;  $Q = (a^1, c^2)$ ,  $P = (v^1, d^2)$  and similarly  $Q' = (a^2, c^1)$ ,  $P' = (d^1, v^2)$ . (Given two points  $M$  and  $N$ ,  $M \vee N$  is  $(\max\{m^1, n^1\}, \max\{m^2, n^2\})$ ).

Deduce that  $\bar{v}$  is the value of a one-shot matrix game.

13. Consider a game as in sect. 4.c.1 and 4.c.2 where the matrix  $G^k$  is of the form:

$$G^k = \begin{pmatrix} a^{k*} & b^{k*} \\ c^{k*} & d^k \end{pmatrix}$$

Prove that  $\bar{v}$  is the value of the one-shot game with incomplete information and infinite pay-off matrices  $B^k$  with:

$$B^k = \begin{pmatrix} a^k & b^k & b^k & \dots & b^k & b^k \\ c^k & a^k & b^k & \dots & b^k & b^k \\ c^k & c^k & a^k & \dots & b^k & b^k \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ c^k & c^k & c^k & \dots & a^k & b^k \\ c^k & c^k & c^k & \dots & c^k & d^k \end{pmatrix}.$$

Show that for any fixed  $\#K$ , one can replace  $B^k$  by finite matrices.

## CHAPTER IX

### Non-zero-Sum Games with Incomplete Information

#### 1. Equilibria in $\Gamma_\infty$

In this section, we will study equilibria in games with incomplete information. The analysis will be limited to 2 players and some major questions are still open.

As seen in ch. VI, in the zero-sum case, when there is lack of information on both sides, the value may not exist hence there is no equilibrium, even in a weak sense, i.e.  $E_0$  is empty. We are thus led to consider games with lack of information on one side. We will moreover assume full monitoring hence the game  $\Gamma_\infty$  is described by the action sets  $S$  and  $T$ , the state space  $K$  with initial probability  $p$  and, for each  $k$ ,  $S \times T$  vector pay-off matrices  $G^k$  with elements in  $\mathbb{R}^2$ .

We shall write  $u_I(p)$  (resp.  $u_{II}(p)$ ) for the value of the game  $G^I(p) = \sum p^k G^{k,I}$  (resp.  $G^{II}(p) = \sum p^k G^{k,II}$ ) where player I (resp. II) maximises.  $\gamma_n(\sigma, \tau)$  is the expected (vector) pay-off at stage  $n$  and  $\bar{\gamma}_n(\sigma, \tau)$  the corresponding average up to stage  $n$ . Remark that the equilibrium condition for I (the informed player) can equivalently be written by using  $\bar{\gamma}_n^{k,I}$ , conditional pay-off given the true state, i.e.

$$\bar{\gamma}_n^{k,I} = E_{\sigma, \tau}^k(\bar{g}_n^{I,k}) = E_{\sigma, \tau}(\bar{g}_n^{I,k} \mid k) = E_{\sigma^k, \tau}(\bar{g}_n^{I,k})$$

This leads to a vector pay-off for I and we shall also use this formulation.

We will first prove a partial result<sup>1</sup> concerning the existence of uniform equilibria and then give a complete characterisation of  $E_\infty$  in  $\Gamma_\infty$ .

**1.a. Existence.** Following (Aumann et al., 1968), we define a **joint plan** to be a triple  $(R, z, \theta)$  where:

- $R$  (set of signals) is a subset of  $S^m$  for some fixed  $m$ ,
- $z$  (signalling strategy) is a  $K$ -tuple where, for each  $k$ ,  $z^k \in \Delta(R)$  and we can always assume

$$z(r) = \sum_k p^k z^k(r) > 0$$

- $\theta$  (contract) is a  $R$ -tuple where, for each  $r$ ,  $\theta(r) \in \Delta(S \times T)$  (correlated move).

To each joint plan is associated a family of probabilities  $p(r)$  on  $\Pi = \Delta(K)$  where  $p^k(r) = \frac{p^k z^k(r)}{z(r)}$  (conditional probability on  $k$  given  $r$  induced by  $z$ ) and pay-offs:

$$\begin{aligned} a^k(r) &= \sum_{s,t} G_{st}^{k,I} \theta_{st}(r) & b^k(r) &= \sum_{s,t} G_{st}^{k,II} \theta_{st}(t) & \forall k \in K \\ \alpha(r) &= \langle p(r), a(r) \rangle & \beta(r) &= \langle p(r), b(r) \rangle \\ a^k &= \sum_r z^k(r) a^k(r) & \beta &= \sum_r z(r) \beta(r) \end{aligned}$$

Each joint plan will induce a pair of strategies as shown by the following

---

<sup>1</sup>The existence theorem 1.3 p. 399 is valid for any  $K$ , cf. (Simon et al., 1995).

PROPOSITION 1.1. A sufficient condition for a joint plan to generate a uniform equilibrium in  $\Gamma_\infty(p)$  with pay-off  $(a = (a^k); \beta)$  is

- (1)  $\beta(r) \geq \text{Vex } u_{\text{II}}(p(r)) \quad \forall r \in R,$
- (2)  $\langle q, \bar{a} \rangle \geq u_{\text{I}}(q) \quad \forall q \in \Pi, \quad \text{where } \bar{a} \text{ is such that:}$
- (3)  $a^k(r) \leq \bar{a}^k \quad \forall r, \forall k, \text{ and } \sum_k p^k(r) a^k(r) = \sum_k p^k(r) \bar{a}^k \quad \forall r.$

PROOF. Player I uses  $z$  to select a signal  $r$  and plays during  $m$  stages according to  $r$ . After these stages, both players are required to follow an history  $h(r)$  where the empirical distribution of the couple of moves  $(s, t)$  converges to  $\theta_{s,t}(r)$ , (compare with sect. 4 p. 162).

It is thus clear that the corresponding asymptotic pay-off, given  $r$ , will be  $a(r), \beta(r)$  ( $k$ -vector pay-off for player I, real pay-off for player II).

Consider now player II after stage  $m$ , given  $r$ . Since, from this stage on, player I is playing non-revealing, player II posterior on  $K$  will remain precisely  $p(r)$ . It follows that, if player II deviates, by not following  $h(r)$ , player I can use a punishing strategy  $\sigma$  satisfying  $\tilde{g}_n(\sigma, \tau) \leq \text{Vex } u_{\text{II}}(p(r))$ , where  $\tilde{g}_n$  is the conditional expected pay-off given  $r$  at each stage  $n$  following the deviation. To prove this, consider the 0-sum game with pay-off  $G^{\text{II}}$  starting at  $p(r)$  where player I is informed and minimises and use V.3. Condition 1 thus implies that there is no profitable deviation for player II.

Consider now player I. He has two kinds of possible deviations:

- First to send a wrong signal, namely if he is of type  $k$  to use some  $r' \neq r, r' \in R$  with  $p^k(r') = 0 < p^k(r)$ . Note that this deviation is not observable by player II but, by condition 3, player I cannot gain by it, since his (future) pay-off will then be  $a^k(r') \leq \bar{a}^k = a^k(r)$ . Similarly I cannot gain by using another signalling strategy  $z'$ .
- The second possibility is for player I to make a detectable deviation at some stage. We then require player II to approach, from this stage on, the vector  $\bar{a}$ , and this is possible due to 3.c p. 195 and condition 2. It is now easy to see that the above strategies define a uniform equilibrium. ■

REMARK 1.1. For an extension of such construction, cf. the next sect. 1.b.

There remains thus to exhibit a joint plan satisfying 1, 2 and 3, i.e. an equilibrium joint plan (EJP for short). We shall use the following notations:

$$\begin{aligned} Y(p) &= \{ y \in Y = \Delta(T) \mid y \text{ is optimal for player II in } G^{\text{II}}(p) \} \\ f_y(q) &= \max_x xG^{\text{I}}(q)y \\ C(y) &= \{ (q, d) \in \Pi \times \mathbb{R} \mid d \geq f_y(q) \} \\ D &= \{ (q, d) \in \Pi \times \mathbb{R} \mid d \leq \text{Cav } u_{\text{I}}(q) \} \\ D_1 &= \{ (q, d) \in \Pi \times \mathbb{R} \mid d \leq u_{\text{I}}(q) \} \end{aligned}$$

(Note that these sets are closed and the first two are convex). By the definition of  $u_{\text{I}}$ ,  $D_1$  and  $C(y)$  have, for all  $y$ , an intersection with empty interior. On the other hand, if  $D$  and  $C(y)$  have the same property, they can be weakly separated and the minmax theorem (applied to  $S$  and  $T$ ) implies the existence of  $x$  satisfying, for all  $q$ ,  $xG^{\text{I}}(q)y \geq u_{\text{I}}(q)$ . Thus:

LEMMA 1.2. If  $C(y) \cap D$  has an empty interior for some  $y$  in  $Y(p)$ , then there exists an EJP at  $p$ .

PROOF. In fact, take  $R = \emptyset$  (the joint plan is non-revealing),  $\theta = x \otimes y$ , where  $x$  is as above, and note that 2 follows and 3 is void. Now 1 comes from  $y \in Y(p)$  hence  $\beta \geq u_{\text{II}}(p) \geq \text{Vex } u_{\text{II}}(p)$ . ■

Denote by SEJP these special EJP and let  $\Pi_1 = \{ q \in \Pi \mid \text{there exists a SEJP at } q \}$ . It is then clear that  $\Pi_1$  is closed (and equals  $\Pi$  if  $u_I$  is concave in particular if  $\#S = 1$ ).

Now we restrict the analysis to the case  $\#K = 2$  and  $\#S > 1$ .

**THEOREM 1.3.** *Assume  $\#K = 2$ . Then for each  $p$ , there exists an EJP at  $p$ .*

**PROOF.** Assume  $p_0 \in \Pi^\circ \setminus \Pi_1$  and  $y_0 \in Y(p_0)$ . Since  $u_I$  is algebraic and  $C(y_0) \cap D_1$  has an empty interior, the projection of  $C(y_0) \cap D$  on  $\mathbb{R}$  is included in some open interval, say  $(q_1, q_2)$ , on which  $\text{Cav } u_I$  is linear. So that there exists  $c$  in  $\mathbb{R}^2$  with  $\text{Cav } u_I(q) = \langle c, q \rangle$  on  $[q_1, q_2]$  and  $\text{Cav } u_I(q_i) = u_I(q_i)$  for  $i = 1, 2$ . Define  $Q(y) = \{ p \in \Pi \mid f_y(p) - \langle c, p \rangle = \min_q f_y(q) - \langle c, q \rangle \equiv \zeta(y) \}$ . Note that  $Q$  is a u.s.c. convex valued correspondence, that  $\zeta$  is continuous and, as above, for each  $y$ , there exists some  $x$  with, for all  $q$ ,  $xG^I(q)y - \langle c, q \rangle \geq \zeta(y)$ , hence the equality. Finally,  $\zeta(y) \geq 0$  implies that the interior of  $C(y) \cap D$  is empty.

By definition of  $y_0$ , we have  $Q(y_0) \subseteq (q_1, q_2)$  and  $\zeta(y_0) < 0$ .

Let now  $p_1$  and  $p_2$  be in  $\Pi_1 \cup \partial\Pi$  with  $p_0 \in (p_1, p_2)$  and  $(p_1, p_2) \cap \Pi_1 = \emptyset$ . Since  $Y(p)$  is also an u.s.c. convex valued correspondence,  $Q(Y(p_1, p_2))$  is connected, hence, by the choice of the  $p_i$ 's, one has

$$(1) \quad \zeta(y) < 0 \quad \text{and} \quad Q(y) \subseteq (q_1, q_2) \quad \text{for all } y \in Y(p_1, p_2)$$

By compactness, one gets  $y'_1 \in Y(p_1)$  and  $x'_1$  with

$$(2) \quad x'_1 G^I(q) y'_1 - \langle c, q \rangle = \zeta(y'_1) \leq 0 \quad \text{and} \quad Q(y'_1) \cap [q_1, q_2] \neq \emptyset$$

Assume first  $p_1 \in \Pi_1$ . Then, either  $C(y_1) \cap D$  has an empty interior (write  $y_1 = y'_1$ ,  $x_1 = x'_1$ ), or there exists  $y_1$ , a closest point to  $y'_1$  in  $Y(p_1)$  having this property. But then (1) holds for  $y \in [y'_1, y_1]$  and one gets (2) for  $y_1$  and some  $x_1$  implying  $\zeta(y_1) = 0$ . Now, if  $p_1 = (1, 0)$ , let  $y'_1 = y_1$ . Note first that  $f_{y_1} \geq c_1$ , hence we can define a supporting hyperplane  $c^1 = (c_1, c'_2)$  to  $f_{y_1}$  with  $c'_2 \leq c_2$  and  $f_{y_1}(q) = \langle c^1, q \rangle$  for some  $q \in (q_1, q_2)$ . Again this implies the existence of  $x_1$  with  $x_1 G^I(q) y_1 = \langle c^1, q \rangle$ . A dual analysis holds for  $p_2$ . Finally, let  $R = 2$  and define  $z$  to be the splitting strategy (prop. 1.2 p. 184) of player I generating  $p_1$  and  $p_2$ . Define now  $\theta(r)$  as  $x_r \otimes y_r$ , then we have an EJP at  $p$ . In fact, 1 follows from  $y_r \in Y(p_r)$ , 2 holds with  $\bar{a} = c$  and 3 with  $a(r) = c^r$ . ■

**1.b. Characterisation (Hart, 1985).** In the above proof, the splitting property was used to convexify in  $p$  the set of non-revealing EJP leading to the same vector pay-off for player I. Note also that, for a fixed  $p$ , the set of equilibrium pay-offs is convex (the players can use their moves to construct a jointly controlled lottery (Aumann et al., 1968), cf. below). The content of the next result is that a repetition of such operations in fact characterises the set of equilibria. Namely each equilibrium pair generates such a sequence and any such sequence leads to an equilibrium.

We first define the set of non-revealing feasible pay-offs:

$$F = \text{Co}\{ G_{s,t}^k \mid s \in S, t \in T \} \subseteq \mathbb{R}^K \times \mathbb{R}^K$$

and  $W =$

$$\begin{aligned} \{ (a, \beta, p) \in \mathbb{R}^K \times \mathbb{R} \times \Pi \mid & \text{i) } \beta \geq \text{Vex } u_{\Pi}(p) \\ & \text{ii) } \langle q, a \rangle \geq u_I(q) \forall q \in \Pi \\ & \text{iii) } \exists (c, d) \in F \text{ such that } a \geq c \text{ and } \langle p, a \rangle = \langle p, c \rangle, \langle p, d \rangle = \beta \} \end{aligned}$$

(Compare with prop. 1.1 p. 398 and note that these conditions correspond to a “non-revealing” joint plan.)

Before introducing the next set, we first define a “ $W$ -process” starting at  $w$  in  $\mathbb{R}^k \times \mathbb{R} \times \Pi$  to be a bounded martingale  $w_n = (a_n, \beta_n, p_n)$  from an auxiliary space  $(\Omega, \mathcal{F}, Q)$  with an **atomic** filtration  $\mathcal{F}_n$  to  $\mathbb{R}^k \times \mathbb{R} \times \Pi$  with  $w_1 = w$ , satisfying, for each  $n$ , either  $a_{n+1} = a_n$ , or  $p_{n+1} = p_n$  a.e. (i.e.  $(a_n, p_n)$  is a “**bi-martingale**”), and converging to some point  $w_\infty$  in  $W$ . Define now  $W^* = \{w \in \mathbb{R}^K \times \mathbb{R} \times \Pi \mid \text{there exists a } W\text{-process starting at } w\}$ . The main result of this section is:

**THEOREM 1.4.**  $(a, \beta) \in \mathbb{R}^k \times \mathbb{R}$  belongs to  $E_\infty(p)$  iff  $(a, \beta, p) \in W^*$ .

**PROOF.** The proof is divided into two parts:

First, given any  $\mathcal{L}$ -equilibrium with pay-off  $(a, \beta)$ , we show that it induces a  $W$ -process starting at  $(a, \beta, p)$ .

Second, given any  $W$ -process starting at  $(a, \beta, p)$ , we construct a uniform equilibrium in  $\Gamma_\infty(p)$  with pay-off  $(a, \beta)$ .

#### PART A. From equilibrium to $W$ -process.

A heuristic outline of the proof is as follows.  $p_n$  will be the posterior probability on  $K$  induced by  $\sigma$ , hence is a martingale converging to some  $p_\infty$ . Thus player I will be asymptotically playing non-revealing, hence the asymptotic pay-offs corresponding to  $\sigma, \tau$ , say  $c_\infty, d_\infty$ , will be in  $F$ . In the definition of  $W$ , conditions 1 and 2 of prop. 1.1 correspond to the individually rational requirements (cf. 4.a p. 162). Remark that the equilibrium condition on the equilibrium path implies that the sequence of conditional future expected pay-offs define a martingale converging to  $(c_\infty, \langle p_\infty, d_\infty \rangle)$ . Condition 3 indicates that player I cannot cheat (i.e. send a wrong signal) and finally the bi-martingale property will be obtained by conditioning on half histories, i.e. just after a move of player I. Formally, let  $H_{n+1/2} = H_n \times S = (S \times T)^{n-1} \times S$ ,  $\mathbb{N}' = \mathbb{N} + 1/2$ ,  $\mathbb{M} = \mathbb{N} \cup \mathbb{N}'$ . We define thus  $H_m$  and the corresponding  $\sigma$ -algebra  $\mathcal{H}_m$  for  $m$  in  $\mathbb{M}$  and all the martingales will be with respect to this filtration. Fix now a Banach limit  $\mathcal{L}$ , an initial  $p$  in  $\Pi$  and  $\mathcal{L}$ -equilibrium strategies  $(\sigma, \tau)$  inducing the pay-off  $(a, \beta)$  so that  $a^k = \mathcal{L}(\bar{\gamma}_n^{I,k}(\sigma, \tau))$ ,  $\beta = \mathcal{L}(\bar{\gamma}_n^{II}(\sigma, \tau))$ ,  $\langle p, a \rangle = \mathcal{L}(\bar{\gamma}_n^I(\sigma, \tau))$ . Define also  $p_m^k = P_{\sigma, \tau, p}(k \mid \mathcal{H}_m)$ ,  $\alpha_m = \mathcal{L}(\mathbb{E}_{\sigma, \tau, p}(\bar{g}_n^{I,k} \mid \mathcal{H}_m))$ ,  $\beta_m = \mathcal{L}(\mathbb{E}_{\sigma, \tau, p}(\bar{g}_n^{II,k} \mid \mathcal{H}_m))$  and note the following:

- (3)  $p_m$  is a martingale in  $\Pi$  converging to some  $p_\infty$ , and  $p_{n+1/2} = p_{n+1}$  for  $n \in \mathbb{N}$
- (4)  $\alpha_m$  is a bounded martingale converging to some  $\alpha_\infty$  with  $\alpha_0 = \langle p, a \rangle$

The Banach limits commute with conditional expectations, the  $\sigma$ -fields  $\mathcal{H}_m$  being finite:

- (5)  $\beta_m$  is a bounded martingale converging to some  $\beta_\infty$  with  $\beta_0 = \beta$

We now define conditional vector pay-offs with respect to the marginal distribution  $\mathbb{E}'$  induced by  $\sigma, \tau$  and  $p$  on  $H_\infty$  (and not  $H_\infty \times K$ ) conditional to  $\mathcal{H}_m$ . Since  $\tau$  is independent of  $k$ , this amounts to take the expectation with respect to  $\mathbb{E}_p(\sigma \mid \mathcal{H}_m)$ , namely the average, non-revealing strategy of player I, given  $\mathcal{H}_m$ , with value  $\sum_k p_m^k(h_m) \sigma^k(h_m)$  on  $h_m$ . Hence, let  $c_m^k = \mathcal{L}(\mathbb{E}'(\bar{g}_n^{I,k} \mid \mathcal{H}_m))$ ,  $d_m^k = \mathcal{L}(\mathbb{E}'(\bar{g}_n^{II,k} \mid \mathcal{H}_m))$ . Then we have:

- (6)  $c_m$  and  $d_m$  are bounded martingales converging to  $(c_\infty, d_\infty)$  with  $(c_\infty, d_\infty) \in F$  a.s.

The last assertion follows from the fact that, for all  $m$ ,  $(c_m, d_m)$  belongs to the compact convex set  $F$ . Moreover, the convergence of  $p_n$  implies:

$$(7) \quad \alpha_\infty = \langle p_\infty, c_\infty \rangle \quad \text{and} \quad \beta_\infty = \langle p_\infty, d_\infty \rangle$$

PROOF. Remark that, for  $n, m$  in  $\mathbb{N}$ ,  $n \geq m$ ,  $\mathsf{E}(g_n^{I,k} | \mathcal{H}_m) = \mathsf{E}'(\sum_k p_{n+1}^k g_n^{I,k} | \mathcal{H}_m) = \sum_k p_m^k \mathsf{E}'(g_n^{I,k} | \mathcal{H}_m) + \sum_k \mathsf{E}'((p_{n+1}^k - p_m^k) g_n^{I,k} | \mathcal{H}_m)$ . Letting  $\pi_m = \sum_k \sup_{n \geq m} |p_{n+1}^k - p_m^k|$ , we obtain  $|\alpha_m - \langle p_m, c_m \rangle| \leq C \mathsf{E}'(\pi_m | \mathcal{H}_m)$ . Since  $p_n^k$  converges to  $p_\infty^k$ ,  $\theta_m^k = \sup_{n \geq m} |p_{n+1}^k - p_\infty^k|$  is a non-increasing positive sequence converging to zero. Hence  $\mathsf{E}'(\theta_m^k | \mathcal{H}_m)$  is a supermartingale converging to 0. Finally,  $\pi_m \leq 2 \sum_k \theta_m^k$  implies the result. The other equality is proved similarly. ■

We have thus defined some asymptotic pay-offs; we use now the equilibrium properties:

$$(8) \quad \beta_\infty \geq \mathbf{Vex} u_{\text{II}}(p_\infty) \quad \text{a.s.}$$

In fact, given any history  $h_m$ , the posterior probability on  $K$  is  $p_m$  and player II can obtain, from this stage on, an asymptotic pay-off of  $\mathbf{Vex} u_{\text{II}}(p_m)$  (cf. theorem 2.12 p. 190). It follows that, if for some  $h_m$  in  $H_m$  with  $P(h_m) > 0$ , one had  $\beta_m(h_m) < \mathbf{Vex} u_{\text{II}}(p_m)$ , then player II could increase his total pay-off  $\beta_m = \sum P(h_m) \beta_m(h_m)$  by switching to the previous strategy after  $h_m$ . The result follows letting  $m \rightarrow \infty$ , using the continuity of  $u_{\text{II}}$ .

To get a similar property for player I, we have to consider again the vector pay-offs. We start by a sequence that majorates  $c_m$ . Let first  $e_m^k = \sup_{\sigma'} \mathcal{L}(\mathsf{E}_{\sigma',\tau}(\bar{g}_n^{k,\text{I}} | \mathcal{H}_m, k)) = \sup_{\sigma'} \mathcal{L}(\mathsf{E}_{\sigma',\tau}(\bar{g}_n^{k,\text{I}} | \mathcal{H}_m))$ .

Note that the expectation is now given  $k$  and that we can assume  $\sigma'$  non-revealing.

$$(9) \quad e_0^k = a^k \quad e_m \geq c_m$$

$$(10) \quad e_{m+1/2} = \mathsf{E}(e_{n+1} | \mathcal{H}_{n+1/2}) \quad \forall n \in \mathbb{N} \quad e_n^k(h_t) = \max_s e_{n+1/2}^k(h_t, s) \quad \forall k, \forall h_t$$

PROOF. (9) follows from the definition since  $\sigma$  is an equilibrium strategy and the pay-off  $c_m$  is attainable by player I. The first equality in (10) comes from the fact that player I is a dummy between stages  $n + 1/2$  and  $n + 1$ . The second comes from the fact that, for each  $k$ , the player can first choose a move and then play optimally. ■

The previous result proves that  $e_m$  is a supermartingale. To get a martingale, we first introduce a sequence  $\lambda_n$  of random variables in  $[0, 1]$ ,  $n \in \mathbb{N}'$  with  $C - e_n^k = \lambda_{n+1/2}(C - e_{n+1/2}^k) \quad \forall n \in \mathbb{N}$ , and we define finally  $f_m$  by  $C - f_m^k = (C - e_m^k) \prod_{n \leq m} \lambda_{n+1/2} \quad \forall m \in \mathbb{M}$ .  $f_m$  is a bounded martingale converging to some  $f_\infty$  with

$$(11) \quad f_0 = a \quad f_n = f_{n+1/2} \quad \forall n$$

$$(12) \quad f_\infty \geq c_\infty \quad \text{and} \quad \langle p_\infty, f_\infty \rangle = \langle p_\infty, c_\infty \rangle$$

PROOF. (11) follows from the definitions since  $\lambda_n$  is  $\mathcal{H}_n$ -measurable,  $n \in \mathbb{N}'$ . Now  $f_m \geq e_m \geq c_m \quad \forall m \in \mathbb{M}$ .

Finally, by (3) and (11),  $(p_n, f_n)$  is a bi-martingale, so  $\langle p_n, f_n \rangle$  a martingale:

$$\mathsf{E}(\langle p_\infty, f_\infty \rangle) = \langle p, f_0 \rangle = \langle p, a \rangle = \alpha = \mathsf{E}(\alpha_\infty) = \mathsf{E}(\langle p_\infty, c_\infty \rangle)$$

by (4) and (7), hence (12). ■

Remains to check the individual rationality condition for the pay-offs  $f_\infty$  of player I:

$$(13) \quad \langle q, f_\infty \rangle \geq u_{\text{I}}(q) \quad \forall q \in \Pi$$

PROOF. At each stage, given any history, player I can, by playing optimally non-revealing in  $G^I(q)$ , get a stage pay-off greater than  $u_I(q)$ . Thus, for all  $m$ ,

$$\begin{aligned} u_I(q) &\leq \mathcal{L} \sup_{\sigma'} \sum q^k \mathsf{E}_{\sigma'}(\bar{g}_n^{I,k} \mid \mathcal{H}_m) \\ &\leq \langle q, e_m \rangle \end{aligned}$$

Hence the result since  $f_m \leq e_m$ . ■

To conclude,  $(f_m, \beta_m, p_m)$  is the required  $W$ -process using  $(c_\infty, d_\infty)$  and (3), (5), (6), (7), (8), (12) and (13).

#### PART B. Second part: From $W$ -process to equilibrium.

Observe first that  $W^*$  is not enlarged if one only required the  $W$ -process, say  $W_n = (a_n, \beta_n, p)$  on a probability space  $(\Omega, \mathcal{F}_n, Q)$ , to satisfy “ $a_{n+1}$  or  $p_{n+1}$  is constant on each atom of  $F_n$ ” rather than on all of them (adding some intermediate fields reduces to the initial case). Now, the filtration being atomic, it is usefully represented by an oriented tree, where the nodes at distance  $n$  from the origin are the atoms of  $\mathcal{F}_n$ . Such a node, say  $w_n$ , leads to all  $w_{n+1}$  in  $\mathcal{F}_{n+1}$  with  $w_{n+1} \subseteq w_n$  and the corresponding arc has the probability  $Q(w_{n+1} \mid w_n)$ . Further, we can assume that each  $w_n$  has two successors and, moreover,  $Q(w_{n+1} \mid w_n) = 1/2$  for each. In fact, the first assertion is proved by adding intermediate fields (i.e. nodes) with the appropriate probabilities. For the second point, one can replace the arc between  $w_n$  and its successor by an infinite tree corresponding to the first occurrence of 1 in a sequence of i.i.d. centred random variables on  $\{-1, 1\}$ . Write  $Q(w_{n+1} \mid w_n) = \sum_1^\infty \lambda_m / 2^m$  and let  $w_{n+1}$  be associated to the first 1 at stage  $m$  iff  $\lambda_m = 1$ . It is easy to see how to extend the  $W$ -process while keeping all its properties.

We henceforth assume that the  $W$ -process possesses the above properties. Basically, the proof will require both players to “follow the above tree”, namely,

- to realise the transitions from one node to its successor. This is done by signalling strategies of player I if  $p_{n+1} \neq p_n$ , and by jointly controlled lotteries if  $a_{n+1} \neq a_n$ ,
- to play a specified sequence of moves between  $w_n$  and  $w_{n+1}$  in order to realise the required pay-off, roughly  $a_n$  and  $\beta_n$ ,
- to use the parameters at  $w_n$  to punish if a deviation occurs between  $w_n$  and  $w_{n+1}$ .

#### STEP 1. Preliminary results.

We first represent points in  $F$  as pay-offs associated to correlated strategies. For all  $(c, d)$  in  $F$ , there exists  $\theta \in \Delta(S \times T)$  with

$$c^k = \sum_{s,t} G_{st}^{I,k} \theta(s, t) \quad d^k = \sum_{s,t} G_{st}^{II,k} \theta(s, t)$$

and write  $c = G^I \cdot \theta$ ,  $d = G^{II} \cdot \theta$ . We thus obtain, using a measurable selection theorem (cf. 7.j p. 427), that there exists a random variable  $\theta_\infty$  with values in  $\Delta(S \times T)$ , s.t.  $Q$ -a.s.  $a_\infty \geq G^I \cdot \theta_\infty$  and  $\langle p_\infty, a_\infty \rangle = \langle p_\infty, G^I \cdot \theta_\infty \rangle$ ,  $\beta_\infty = \langle p_\infty, G^{II} \cdot \theta_\infty \rangle$ .

Approximate now  $\mathsf{E}(\theta_\infty \mid \mathcal{F}_n)$  by  $\mathcal{F}_n$ -measurable random variables  $\theta_n$  s.t.  $n\theta_n$  is integer-valued and  $\|\theta_n - \mathsf{E}(\theta_\infty \mid \mathcal{F}_n)\| \leq 1/n$ .

We shall also use the following properties. Since  $\langle q, a_\infty \rangle \geq u_I(q)$ , one has, by taking expectation

$$(14) \quad \langle a_n, q \rangle \geq u_I(q) \quad \text{for all } n$$

Similarly, from  $\beta_\infty \geq \mathsf{Vex} u_{II}(p_\infty)$ , one has, using Jensen's inequality

$$\beta_n = \mathsf{E}(\beta_\infty \mid \mathcal{F}_n) \geq \mathsf{E}(\mathsf{Vex} u_{II}(p_\infty) \mid \mathcal{F}_n) \geq \mathsf{Vex} u_{II}(\mathsf{E}(p_\infty \mid \mathcal{F}_n)) \geq \mathsf{Vex} u_{II}(p_n)$$

STEP 2. *Construction of the strategies.*

As in the previous proofs, the strategies will be defined through a cooperative procedure — master plan — and punishments in case of detectable deviation. Stages  $m = n!$  for  $n \in \mathbb{N}$  will be communication stages and related to the transition from one node to its successor. The remaining stages will be devoted to the pay-offs, while the reference node will be kept fixed. We define now inductively histories consistent with the master plan and a mapping  $\zeta$  from histories to nodes. Let us write  $\{s', s''\}, \{t', t''\}$  for two pairs of moves of each player that will be used to communicate. Assume now  $m = n! - 1$ ,  $h_{m-1}$  defined and consistent with the master plan and  $\zeta(h_{m-1}) = w_{n-1}$ . Recall that the successor of  $w_{n-1}$  is a random variable  $w_n$  with equally likely values  $w'_n$  and  $w''_n$ .

To define the behaviour at stage  $m$ , we consider the two cases:

- if  $p_n \neq p_{n-1}$ , player I uses a signalling strategy with support on  $\{s', s''\}$  to generate  $p_n$  (cf. prop. 1.2 p. 184) namely  $\sigma^k(h_{m-1})(s') = p_n^k/2p_{n-1}^k$ , the move of player II is arbitrary and  $\zeta(h_{m-1}, s') = w'_n$ .
- if  $a_n \neq a_{n-1}$ , both players randomise equally on the above moves, namely  $\sigma^k(h_{m-1})(s') = \sigma^k(h_{m-1})(s'') = 1/2$ ,  $\theta(h_{m-1})(t') = \theta(h_{m-1})(t'') = 1/2$ , and  $\zeta(h_{m-1}, s', t') = \zeta(h_{m-1}, s'', t'') = w'_n$ , and  $w''_n$  otherwise.

(Note that  $P(w'_n | w_n) = 1/2$  as soon as one of the players is using the above procedure: No cheating is possible. This is a **jointly controlled lottery**.)

We now look at the pay-off stages (from  $n! + 1$  to  $(n + 1)! - 1$ ). Given  $w_n$ , consider  $\theta_n = \theta_n(w_n)$  introduced in b). The players are then requested to play  $n$ -cycles of pure moves realising  $\theta_n$  during this bloc  $n$ . The node associated to such an history is still  $w_n$ .

Let us finally consider deviations and punishments. If  $m + 1$  is the first stage where a detectable deviation occurs, define  $\zeta(h_{m+1})$  as  $\zeta(h_m)$  where  $h_m$  precedes  $h_{m+1}$  and assume that  $\zeta(h_m) = w_n = (p_n, a_n, \beta_n)$ ;

- if player I is the deviator, player II uses, from stage  $m + 1$  on, a strategy that approaches  $a_n$  (cf. 2.c p. 190 and use (14) above), namely such that, for  $\ell \geq N_0$ ,  $E_{\sigma', \tau}(\rho_\ell | \mathcal{H}_{m+1}, k) \leq a_n^k$ , for all  $\sigma'$ , where  $\rho_\ell$  is the average pay-off of player I between stages  $m + 1$  and  $\ell$ .
- in the case of player II, let player I use a strategy inducing in  $G^{\text{II}}(p_n)$  a stage pay-off less or equal to  $\text{Vex } u_{\text{II}}(p_n)$ . (Note that player I is minimising and use 2.b p. 188).

This ends the definition of  $\sigma$  and  $\tau$ .

STEP 3. *Probabilities and pay-offs.*

We now prove that the probabilities and pay-offs induced by  $(\sigma, \tau)$  on  $\mathcal{H}_\infty \times \wp(K)$  correspond to the arcs and nodes of the initial probability tree. For every  $w_n = (a_n, \beta_n, p_n)$  and  $n! \leq m < (n + 1)! - 1$ ,

$$(15) \quad P_{\sigma, \tau, p}(\zeta(h_m) = w_n) = Q(w_n)$$

$$(16) \quad P_{\sigma, \tau, p}(k | h_m) = p_n^k(\zeta(h_m))$$

PROOF. Both properties are proved by induction. (15) follows from the fact that, at each communication stage, say  $r = \ell!$ ,  $P(h_r, \zeta(h_r) = w'_\ell | \zeta(h_{r-1}) = w_{\ell-1}) = \frac{1}{2}$ . (16) comes from the fact that player I is playing non-revealing during pay-off stages and uses the right signalling strategy at the communication stages. ■

Consider now the pay-offs. If  $\phi_m$  denotes the empirical distribution of moves up to stage  $m$ , and  $\lambda_m = n!/m$ , one has:

$$(17) \quad \forall s, t, m (n! \leq m < (n+1)! - 1), |\phi_m(s, t) - [\lambda_m \theta_{n-1}(s, t) + (1 - \lambda_m) \theta_n(s, t)]| \leq 4/n$$

This follows easily from the fact that most of the stages are in blocs  $n-1$  or  $n$  where both players play  $\theta_{n-1}$  or  $\theta_n$  by cycles.

From the previous properties, we can now deduce

$$(18) \quad \lim \bar{\gamma}_m^{I,k}(\sigma, \tau) = a^k$$

$$(19) \quad \lim \bar{\gamma}_m^{II}(\sigma, \tau) = \beta$$

PROOF. Consider first (19). From (17), we obtain, using (15) and conditioning on each history  $\bar{\gamma}_m^{II}(\sigma, \tau) = \sum_{h_m} P_{\sigma, \tau, p}(h_m) \mathbb{E}(\bar{g}_m^{II,k} | h_m) = \sum_{w_n} P_Q(w_n) \langle p_n, G^{II} \cdot \phi_m \rangle$ , so that  $|\bar{\gamma}_m^{II}(\sigma, \tau) - \mathbb{E}_Q(\langle p_n, G^{II} \cdot (\lambda_m \theta_{n-1} + (1 - \lambda_m) \theta_n) \rangle)| \leq R/n$  with  $R = 4C \cdot \#S \cdot \#T$ .

Hence as  $m \rightarrow \infty$ , we obtain  $\lim \bar{\gamma}_m^{II}(\sigma, \tau) = \mathbb{E}_Q(\langle p_\infty, G^{II} \cdot \theta_\infty \rangle) = \mathbb{E}(\beta_\infty) = \beta$ .

Similarly, one first gets

$$(20) \quad \lim \bar{\gamma}_m^I(\sigma, \tau) = \mathbb{E}_Q(p_\infty, a_\infty) = \langle p, a \rangle$$

(Recall that  $\langle p_n, a_n \rangle$  is a martingale).

To get the equality component-wise, we first obtain, from (17), that

$$(21) \quad \bar{g}_m^{I,k} \leq (\lambda_m a_{n-1}^k + (1 - \lambda_m) a_n^k) + R'/n \quad \text{with } R' = 5C(\#S)(\#T)$$

since  $\bar{g}_m^I = G^I \cdot \phi_m$  and  $a_n = \mathbb{E}(a_\infty | \mathcal{F}_n) \geq \mathbb{E}(G^I \cdot \theta_\infty | \mathcal{F}_n) \geq G^I \cdot \theta_n - 1/n$ . We know that the probability induced by  $\sigma, \tau, p$  on the nodes of the tree (through the mapping  $\zeta$ ) coincides with the initial  $Q$ . Now, when considering only  $a_n^k$ , we have the same property with  $\sigma^k$  and  $\tau$ , for all  $k$ . In fact,  $\sigma$  is non-revealing except at the stages where  $p_n$  changes but then  $a_n$  is constant. We obtain thus  $\bar{\gamma}_m^{I,k} = \mathbb{E}^k(\bar{g}_m^{I,k}) \leq \lambda_m \mathbb{E}^k a_{n-1}^k + (1 - \lambda_m) \mathbb{E}^k a_n^k + R'/n \leq a^k + R'/n$ , hence (18) by using (20).  $\blacksquare$

#### STEP 4. Equilibrium conditions.

Up to now,  $\sigma$  and  $\tau$  are adapted to the tree and define the right probabilities and pay-offs. There remains to check the equilibrium conditions.

Consider first player II after some history  $h_m$  with  $\zeta(h_m) = w_r$ . His posterior probability on  $K$  is precisely  $p_r$  by (16). If  $m+1$  is a signalling stage with  $p_r = p_{r+1}$ , player II can make an non-detectable deviation but without affecting the distribution of outcomes (jointly controlled lottery) and its future pay-off will be the same. In other cases, a deviation will be detected and player II punished.

The pay-off thereafter will be  $\text{Vex } u_I(p_r) \leq \beta_r$ , while, if player II keeps playing the equilibrium strategy, he could expect  $\mathbb{E}(\langle p_n, G^{II} \cdot \theta_n \rangle | \mathcal{H}_r)$  on each of the following blocs  $n$ . For  $n$  large enough, this is near  $\mathbb{E}(\langle p_\infty, G^{II} \cdot \theta_\infty \rangle | \mathcal{H}_r) = \mathbb{E}(\beta_\infty | \mathcal{H}_r) = \beta_r$ . Hence, the above argument shows that  $(\sigma, \tau)$  is an  $\mathcal{L}$ -equilibrium. To get the uniform condition, recall that the punishment strategy of player I is uniform. Since  $\langle p_n, G^{II} \cdot \theta_n \rangle$  converges a.e. to  $\langle p_\infty, G^{II} \cdot \theta_\infty \rangle = \langle p_\infty, d_\infty \rangle = \beta_\infty = \lim \beta_n$ , there exists, for every  $\varepsilon > 0$ , a subset  $\Omega_1$  of  $\Omega$  with probability greater than  $1 - \varepsilon$  on which the convergence is uniform.

Define hence  $N$  so that, on  $\Omega_1$ ,  $n$  and  $n' \geq N$  implies  $|\langle p_n, G^{II} \cdot \theta_n \rangle - \beta_{n'}| \leq \varepsilon$ .

Given any  $\tau'$ , consider the game up to stage  $m \geq (N+1)!$  and denote by  $X+1$  the time of the first deviation (with  $X = m$  if no deviation occurs before  $m$ ). Until stage  $X$ , the pay-offs under  $\tau$  and  $\tau'$  coincide so we have only to consider the pay-offs after. For simplicity, let us write, if  $X+1 \leq m$ ,  $f_\tau(X)$  for  $(1/m) \mathbb{E}_\tau(\sum_{X+1}^m g_n^{II,k} | h_X)$  and similarly

$f_{\tau'}(X)$ . Assume  $n! \leq m < (n+1)!$ , from stage  $X+1$  to  $m$ , the pay-off under  $(\sigma, \tau)$  is mainly according to  $\theta_{n-1}$  or  $\theta_n$  and we have

$$|f_{\tau}(X) - (1/m)(m-X+1)\mathbb{E}_{\sigma,\tau}(\beta_n | h_X)| \leq CP(\Omega_1 | h_X) + R/n$$

Moreover, the conditional expectation, given  $h_X$ , depends only on  $\zeta(h(X)) = w_\ell$ , hence  $\beta_n$  being a martingale  $\mathbb{E}(\beta_n | h_X) = \beta_\ell$ .

On the other hand, under  $(\sigma, \tau')$ , the expected pay-off after stage  $X+1$  is at most  $\beta_\ell$  so that we get finally  $\mathbb{E}_X(f_{\tau'}(X) - f_{\tau}(X)) \leq \varepsilon + R/n + C/m$ , hence, for  $m$  sufficiently large,  $\bar{\gamma}_m^{\text{II}}(\sigma', \tau) \leq \bar{\gamma}_m^{\text{II}}(\sigma, \tau) + 2\varepsilon C$ .

Consider now player I again after some history consistent with the master plan and with  $\zeta(h_m) = w_n$ . We claim that, by deviating, player I cannot get more than  $a_n$ . This is clearly the case if the deviation is observable since player II can approach  $a_n$ . Otherwise, it occurs at some communication stage where either player I cannot cheat or he has to use a signalling strategy, hence  $a_{n+1} = a_n$ . But, as seen above, the future expected pay-off per bloc is at most  $a_n$ . Since  $a_n$  is a martingale with expectation  $a$ , the result will follow. In fact, consider the average pay-off up to stage  $m$ , with  $n! \leq m < (n+1)!$  under some  $\sigma'$  and  $\tau$ . Let  $X+1$  be the time of a first observable deviation as before. Write  $\zeta(h_X) = w_r$  and let us compute first the pay-off until  $X$ : then we have, using (21),  $\mathbb{E}^k((1/m)\sum_1^X g_{\ell}^{I,k}) \leq \mathbb{E}^k((X/m)a_r^k) + \mathbb{E}^k((X/m)((r-1)!/X)(a_{r-1}^k - a_r^k)) + \mathbb{E}^k((X/m)(1/r))5R'$ . Note that the random variable in the last term is always less than  $1/n$ . As, for the second term, if  $r < n$ , the integrand is smaller than  $2C/n$  and, finally, by conditioning on  $h = h_{(n-1)!}$ , one gets  $\mathbb{E}^k(\mathbb{1}_{\{r=n\}}(a_{r-1}^k - a_r^k) | h) = \mathbb{1}_{\{r=n\}} \mathbb{E}^k(a_{r-1}^k - a_r^k | h) = 0$  by the martingale property for  $a_n^k$ , since the laws induced by  $(\sigma, \tau)$  and  $(\sigma', \tau)$  coincide until  $X$ . So that:

$$(22) \quad \mathbb{E}^k((1/m)\sum_1^X g_{\ell}^{I,k}) \leq \mathbb{E}^k((X/m)a_r^k) + (5R' + 2C)/n$$

Finally, for the pay-off after  $X+1$ , we obtain

$$(23) \quad \mathbb{E}_{\sigma',\tau}^k((1/m)\sum_{X+2}^m g_{\ell}^{I,k} | \mathcal{H}_X) \leq (m - (X+2)/m)a_r^k + (2C/\sqrt{m})$$

since player II approaches  $a_r^k$ . (22) and (23) together imply

$$\bar{\gamma}_m^{\text{I},k}(\sigma', \tau) \leq \mathbb{E}_{\sigma',\tau} a_r^k + (C/m) + (5R + 2C)/n + 2C/(\sqrt{m} - 1)$$

and, since  $\mathbb{E}_{\sigma',\tau}(a_r^k) = a^k = \lim_{m \rightarrow \infty} \bar{\gamma}_m^{\text{I},k}(\sigma, \tau)$ , we finally obtain:  $\forall \varepsilon > 0, \exists M : \forall m \geq M$

$$\bar{\gamma}_m^{\text{I},k}(\sigma', \tau) \leq \bar{\gamma}_m^{\text{I},k}(\sigma, \tau) + \varepsilon$$

This finishes the proof of the second part and of theorem 1.4. ■

## 2. Bi-convexity and bi-martingales

This section is devoted to the study of the new tools that were introduced in the previous section, namely bi-martingales and the corresponding processes.  $X$  and  $Y$  are two compact convex subsets of some Euclidian space and  $(Z_n)$  is a bi-martingale with values in  $Z = X \times Y$ , namely there exists a probability space  $(\Omega, \mathcal{F}, P)$  and a filtration with finite fields  $\mathcal{F}_n$ ,  $\mathcal{F}_n \subseteq \mathcal{F}$  such that  $(Z_n)$  is a  $(\mathcal{F}_n)$ -martingale and, for each  $n$ , either  $X_n = X_{n+1}$  or  $Y_n = Y_{n+1}$ . For each  $(\mathcal{F}_n)$  stopping time  $\theta$ ,  $Z^\theta$  is the bi-martingale stopped at  $\theta$  (i.e.  $Z_n^\theta(\omega) = Z_{n \wedge \theta(\omega)}(\omega)$ ) and converges to  $Z_\infty^\theta$ . Given a subset  $A$  of  $Z$ , we now define the following sets:  $A^* (\text{resp. } A_f, \text{ resp. } A_b) = \{z \mid \text{there exists a bi-martingale } Z^n \text{ starting at } z \text{ and a (resp. finite, resp. bounded) stopping time } \theta \text{ such that } Z_\infty^\theta \in A \text{ a.s.}\}$ . Observe that this definition is unambiguous: even if  $Z_\infty^\theta \in A$  a.s. was only interpreted in the sense

that the outer probability of this event is one, by removing the exceptional set one would obtain another probability space with a filtration of finite fields and a bi-martingale on those, with now  $Z_\infty^\theta \in A$  everywhere.

We will give a geometrical characterisation of these sets and prove that they may differ. In particular, this will imply that the number of communicating stages needed to reach an equilibrium pay-off may be unbounded. (See also ex. IXEx.2 p. 419).

We start with some definitions.

A set  $A$  in  $Z$  is **bi-convex** if, for all  $x$ , resp.  $y$ , the sections  $A_x = \{y \mid (x, y) \in A\}$ , resp.  $A_y$  are convex. Similarly, a function on a bi-convex set  $A$  is **bi-convex** if any of its restrictions to  $A_x$  or  $A_y$  is convex. Given  $A$  in  $Z$ ,  $\text{bico } A$  is the smallest bi-convex set containing  $A$  or, equivalently,  $\text{bico } A = \cup A_n$  where the  $A_n$  are constructed inductively from  $A_{n-1}$ , by convexifying along  $X$  or  $Y$  starting from  $A_1 = A$ . (Note that contrary to the convex hull operator  $\text{bico}$  may require an unbounded number of stages of convexification, cf. example 2.2 in (Aumann and Hart, 1986).)

Since  $A_n$  corresponds precisely to the starting point of bi-martingales  $\{Z_m^\theta\}$  with  $\theta \leq n$ , we obtain

**PROPOSITION 2.1.**  $A_b = \text{bico } A$ .

We thus obtain a first distinction between  $A_b$  and  $A_f$  by the example in figure 1:

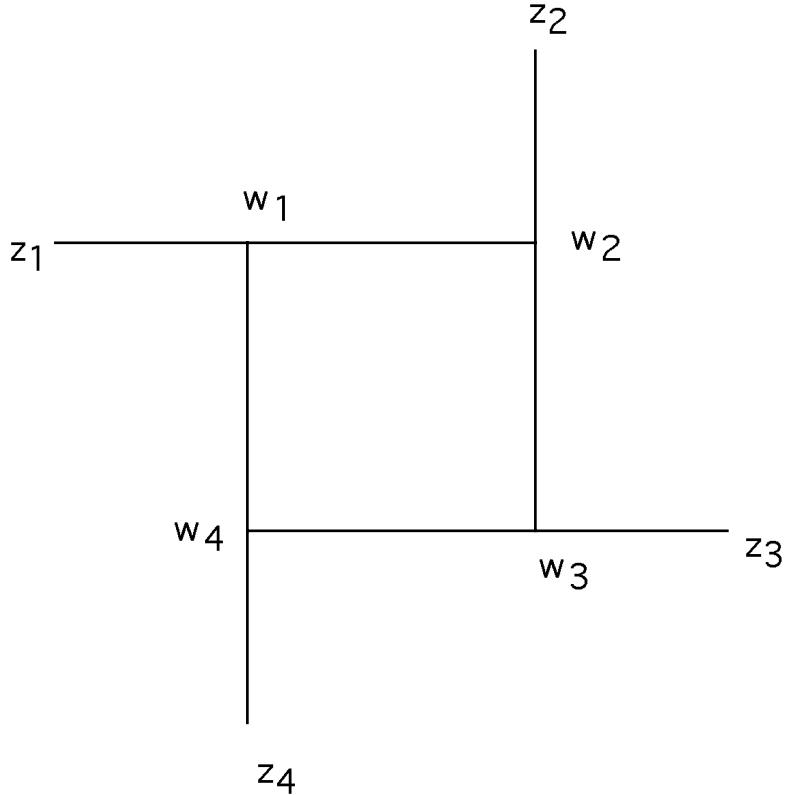


FIGURE 1. An unbounded conversation-protocol

$$A = \{z_i \mid i = 1, \dots, 4\} \quad \text{hence} \quad \text{bico } A = A$$

To prove that  $w_1$  belongs to  $A_f$ , consider the (cyclic) process splitting from  $w_i$  to  $z_i$  and  $w_j$  with  $j \equiv i+1 \pmod 4$ . This induces a bi-martingale and the stopping time corresponding to the entrance in  $A$  is clearly finite a.s.

We consider now separation properties. Given a bi-convex set  $B$ , a set  $A \subseteq B$  and  $z \in B$ , say that  $z$  is (strictly bi-) separated from  $A$  if there exists a bounded bi-convex function  $f$  on  $B$  such that

$$f(z) > \sup\{ f(a) \mid a \in A \} = f^*(A)$$

$\text{ns } B$  is the set of points of  $B$  that cannot be separated from  $A$  (*not separated*). It is easy to see that  $\text{ns } B$  is a bi-convex set containing  $\text{bico } A$ . Note nevertheless that, contrary to the convex case (when the corresponding set  $\text{ns } B$  equals  $\overline{\text{Co } A}$ ), one may have:

- $\overline{\text{bico } A} \subsetneq \text{ns } B$  (cf. example 1 above);
- the separation by bi-affine functions is not sufficient (Aumann and Hart, 1986, example 3.4);
- the resulting set depends on  $B$  (Aumann and Hart, 1986, example 3.5).

We are thus led to apply repeatedly this separating process and we define inductively  $B_1 = Z$  and, for every ordinal  $\alpha$ ,  $B_\alpha = \bigcap_{\beta < \alpha} \text{ns } B_\beta$ . Hence  $B_\alpha$  converges to some set  $C$  denoted by  $\text{bis } A$  (*bi-separated*) which satisfies  $C = \text{ns } C$  and is the largest superset of  $A$  having this property. We are now ready to give the second characterisation.

PROPOSITION 2.2.  $A_f = \text{bis } A$ .

PROOF. Let us show that  $A_f$  shares the characteristic properties of  $\text{bis } A$ . First,  $A_f = \text{ns } A_f$ . In fact, starting from  $z$  in  $A_f$ , consider the associated bi-martingale  $Z_n^\theta$ . Given any bounded bi-convex function  $f$  on  $A_f$ ,  $f(Z_n^\theta)$  is then a bounded sub-martingale, hence  $\theta$  being a.e. finite,

$$f(z) \leq \mathbb{E}(f(Z_\infty^\theta))$$

Since  $Z_\infty^\theta$  belongs to  $A$  a.s., we obtain

$$f(z) \leq f^*(A)$$

so that  $z$  belongs to  $\text{ns } A_f$ .

Assume now that  $A \subseteq B = \text{ns } B$ . Define  $\phi$  on  $B$  by

$$\phi(z) = \inf P(Z_n \notin A, \forall n)$$

where the infimum is taken over all  $B$ -valued bi-martingales starting from  $z$ . (Observe that, since each  $Z_n$  takes only finitely many values, the event considered is measurable, even for non-measurable  $A$ .) Clearly,  $\phi$  is bi-convex and equals 0 on  $A$ , hence  $\phi(z) = 0$ ,  $\forall z \in B = \text{ns } B$ . It follows that

$$1 = 1 - \phi(z) = \sup P(Z_\infty^\theta \in A)$$

where the supremum is taken over all bi-martingales starting from  $z$  and all a.e. finite stopping times  $\theta$ . To prove that the supremum is reached (hence that  $z$  belongs to  $A_f$ ), note that, given any positive  $\rho$  and for each  $z$ , there exists  $m = m(z)$  and an adapted bi-martingale with  $P(Z_m \in A) > \rho$ . If  $A$  is not reached at time  $m$ , start a new process adapted to  $Z_m$  and so on (Recall  $Z_m$  takes only finitely many values). This way,  $A$  will be reached in a finite time with probability one. ■

To obtain  $A^*$ , we have to consider separation by bi-convex functions **continuous** at each point of  $A$ . Thus, as above, we define the set  $\text{nsc } B$  of points of  $B$  that cannot be strictly separated from  $A$  by such functions and inductively  $\text{bisc } A$  as the largest set satisfying  $B = \text{nsc } B$  ( $c$  for continuous).

PROPOSITION 2.3. Assume  $A$  is closed. Then  $A^* = \text{bisc } A$ .

PROOF. First  $A^* = \text{nsc } A^*$ . Given  $z$  in  $A^*$ , an associated bi-martingale  $Z_n$  and a bounded bi-convex function  $f$  on  $A^*$  continuous at every point of  $A$ , one has that  $f(Z_n)$  is a submartingale and  $Z_n \rightarrow Z_\infty \in A$ . Hence,  $f(Z_n) \rightarrow f(Z_\infty)$  so that

$$f(z) \leq Ef(Z_n) \leq Ef(Z_\infty) \leq f^*(A)$$

hence  $z \in \text{nsc } A^*$ . Assume now  $B = \text{nsc } B$ . Let  $d(\cdot, A)$  be the distance to the (closed) set  $A$  and define  $\psi$  on  $B$  by

$$\psi(z) = \inf E[d(Z_\infty, A)]$$

where the infimum is taken over all bi-martingales starting from  $z$ . Note that  $\psi$  is bi-convex. Now  $\psi(\cdot) \leq d(\cdot, A)$  implies that  $\psi$  vanishes and is continuous on  $A$ , hence  $\psi(z) = 0$  for any  $z$  in  $B = \text{nsc } B$ . To prove that the infimum is actually reached (hence  $Z_\infty \in A$  a.e. and  $z \in A^*$ ), we define, for any  $z$  in  $B$  and any  $\rho > 0$ , some  $m = m(z, \rho)$  and an adapted bi-martingale satisfying  $E(d(Z_m, A)) \leq \rho$ . If  $A$  is not reached at stage  $m$  ( $= m_1$ ), start from  $Z_m$  with  $\rho/2$  and so on. This defines inductively an entrance time  $\theta$  in  $A$  and a sequence of stages  $m_n$  with  $E(d(Z_{m_n}^\theta, A)) \leq \rho/n$ , hence  $E(d(Z_\infty^\theta, A)) = 0$ . ■

Finally, to prove that  $A_f$  and  $A^*$  may differ, consider the following example:

EXAMPLE 2.1. (Aumann and Hart, 1986)  $X = Y = [0, 1]^2$ . Define  $T = [0, \varepsilon]$  with  $\varepsilon > 0$  small enough, and points in  $Z = X \times Y$  by:

$$\begin{aligned} b_t &= (1, 3t - 2t^2; 2t, 4t^2) & c_t &= (t, t^2; 1, 3t - 2t^2) \\ d_t &= (2t, 4t^2; 2t, 4t^2) & e_t &= (t, t^2; 2t, 4t^2) \end{aligned}$$

Let  $B = \{b_t\}_{t \in T}$  and similarly  $C, D$  and  $E$ . Define  $A$  as  $B \cup C \cup \{0\}$  with  $0 = ((0, 0); (0, 0))$ . It is easy to see that  $D \cup E$  is in  $A^*$  since one has

$$\begin{aligned} d_t &= (t/(1-t))b_t + ((1-2t)/(1-t))e_t & (y \text{ constant}) \\ e_t &= (t/(1-t))c_t + ((1-2t)/(1-t))d_{t/2} & (x \text{ constant}) \end{aligned}$$

and  $d_t, e_t \xrightarrow[t \rightarrow \infty]{} 0$ . On the other hand  $(D \cup E) \cap A_f = \emptyset$ . In fact, consider  $z$  in  $D \cup E$  and  $(Z_n^\theta)$  an adapted bi-martingale with  $\theta$  a.e. finite. Let  $F_+ = \{z \in Z \mid x_2 > 0 \text{ and } y_2 > 0\}$ ,  $F_0 = \{z \in Z \mid x_2 = 0 \text{ and } y_2 = 0\}$ , and  $F = F_+ \cup F_0$ .  $F$  is clearly convex and contains  $A$ , hence also  $A^*$  and  $A_f$ . In particular,  $Z_n^\theta \in F$  for all  $n$ , but  $z \in F_+$  implies moreover  $Z_n^\theta \in F_+$  since  $X_n$  and  $Y_n$  cannot change simultaneously. This implies  $z \in (A \cap F_+)_f$ . On the other hand on  $A \cap F_+$  one has  $x_1 + y_1 \geq 1$ , hence the same property on  $\text{Co}(A \cap F_+)$ , hence for  $z$ , but no point of  $D \cup E$  satisfies it.

One can also show that  $(A^*)^*$  may differ from  $A^*$  (Aumann and Hart, 1986, example 5.3) and that  $A$  closed does not imply  $A^*$  closed (Aumann and Hart, 1986, example 5.6).

### 3. Correlated equilibrium and communication equilibrium

An  $r$ -device ( $r = 0, 1, \dots, \infty$ ) is a communication device (3.c p.92) where players make inputs only until stage  $r$ . Thus, for  $r = 0$ , we obtain the autonomous devices (3.b p.90) and for  $r = \infty$  we obtain the communication devices. The corresponding standard devices (3.c p.92) have corresponding sets of inputs  $I_n$  (by the players) and of messages

$M_n$  (to the players) relative to stage  $n$  of the game ( $n = 1, 2, \dots$ ), for  $r > 0$ :

$$\begin{aligned} I_0^I &= K & \#I_0^{\text{II}} &= 1 \\ I_n^I &= T & I_n^{\text{II}} &= S & M_n^I &= S & M_n^{\text{II}} &= T & \text{for } 1 \leq n < r \\ \#I_n^I &= 1 & \#I_n^{\text{II}} &= 1 & M_n^I &= S^{(S \times T)^{n-r}} & M_n^{\text{II}} &= T^{(S \times T)^{n-r}} & \text{for } n \geq r \end{aligned}$$

using (the proofs of) theorem 3.7 p. 91 and Corollary 3.16 p. 93. (One could also use Dalkey's theorem (1.3 p. 53) to eliminate redundancies in  $M_n^I$ ,  $M_n^{\text{II}}$  for  $n > r$  — thus  $S^{T^{n-r}}$  and  $T^{S^{n-r}}$  — but the present formulation leads, for  $r \geq 1$ , to a set of pure strategies for the device,  $(S \times T)^H$  with  $H = K \times \bigcup_{n \geq 0} (S \times T)^n$ , which is independent of  $r$  — and the canonical device or equilibrium is a probability distribution over this set). And for  $r = 0$  we could take  $M_n^I = S^{K \times T^{n-1}}$ ,  $M_n^{\text{II}} = T^{S^{n-1}}$ , so we get, as for correlated equilibria, the product of the two pure strategy spaces as strategy space for the device.

As in sect. 3 p. 88, an  $r$ -communication equilibrium is an  $r$ -device together with an “equilibrium” of the corresponding extended game. It is called canonical if it uses a canonical  $r$ -device and if the equilibrium strategies are to report truthfully and to follow the recommendation. As in ch. IV, definition 1.4 and in sect. 1 here, we do not in fact define the equilibria, because of the ambiguity of the pay-off function, but just the corresponding set of pay-offs  $E_\infty \subseteq \mathbb{R}^K \times \mathbb{R}$ , consisting of a vector pay-off to player I and a scalar to player II. Recall  $E_\infty$  “exists” (though conceivably empty) iff any equilibrium pay-off corresponding to any Banach limit is also a uniform equilibrium pay-off. The corresponding set for  $r$ -communication equilibria is denoted  $D_r$ . So  $D_\infty$  is the set of communication equilibria, and  $D_0$  the set of extensive form correlated equilibria.

We first need the following analogue of theorem 3.7 p. 91 and cor. 3.16 p. 93:

**LEMMA 3.1.** *For any Banach limit  $\mathcal{L}$  and any corresponding equilibrium of some  $r$ -device, there exists a corresponding canonical  $\mathcal{L}$ - $r$ -communication equilibrium.*

**PROOF.** Is as in theorem 3.7 p. 91 and Corollary 3.16 p. 93. Observe that the players' personal devices, which handle all communications with the central device and compute the players' strategy — thus receiving from the player as input until stage  $r$  his reported signal in the game and give him as output a recommended action, and which will form the canonical device when taken as whole with the central device, can be taken to receive in fact no inputs from the player from stage  $r$  on, by recommending to him an action conditional on the signals he received in the game from stage  $r$  on. ■

**3.a. Communication equilibrium.** Let  $Z = \{a \in \mathbb{R}^K \mid \langle q, a \rangle \geq u_I(q), \forall q \in \Delta(K)\}$ . Consider the following class of communication devices: Player I first reports  $k \in K$  to the device, next  $(c, d, z) \in F \times Z$  is selected according to some probability distribution  $P^k$  and  $(c, d)$  is transmitted to both players. At every later stage the players' inputs consist of either doing nothing, in which case nothing happens, or of hitting an alarm button, in which case “ $z$ ” is transmitted to both players. Some Borel map  $\varphi$  from  $F$  to  $(S \times T)^\infty$  is fixed yielding for every  $f \in F$  a sequence of moves with this average pay-off. The players' strategies are the following: player I reports the true state of nature  $k$ ; both follow the sequence of moves  $\varphi(c, d)$ , transmitting nothing to the device, until they notice a deviation by their opponent, in that case they hit the alarm button. As soon as “ $z$ ” is announced, player I holds player II's pay-off down to  $(\text{Vex } u_{\text{II}})(p(\cdot \mid c, d, z))$  (theorem 3.5 p. 195), and player II holds player I down to  $z$  (cor. 3.33 p. 218 or lemma 3.36.2 p. 223 for the measurability).

Observe that, with such a device, the game remains a game with incomplete information on one side: Both players know at every stage the full past history, except for the state of nature and player I's first report, which remain his private information. Let  $\overline{D}_\infty = \{(p, (a, \beta), (P^k)_{k \in K}) \mid \text{the above strategies are in equilibrium, with pay-off } (a, \beta) \in \mathbb{R}^K \times \mathbb{R}\}$  (with the understanding that  $P^k$  and  $a^k$  are not to be defined when  $p^k = 0$ ) and  $\overline{D}_\infty^1 = \text{Proj}_{\Delta(K) \times \mathbb{R}^K \times \mathbb{R}}(\overline{D}_\infty)$ ,  $\overline{D}_\infty^2 = \text{Proj}_{[\Delta(F \times Z)]^K}(\overline{D}_\infty)$ . Then we have clearly:

**PROPOSITION 3.2.**  $(p, a, \beta, (P^k)_{k \in K})$  belongs to  $\overline{D}_\infty$  iff

- (I)  $E^k(c^k) = a^k \quad \forall k \in K \quad \text{s.t.} \quad p^k > 0$
- (II)  $E^k \max[c^l, E^k(z^l \mid c, d)] \leq a^l \quad \forall (k, l) \in K \times K \quad \text{s.t.} \quad p^k > 0$
- (III)  $E\langle p(\cdot \mid c, d), d \rangle = \beta$
- (IV)  $\langle p(\cdot \mid c, d), d \rangle \geq E[(\text{Vex } u_{\Pi})(p(\cdot \mid c, d, z)) \mid c, d] \quad \text{a.s.}$

More precisely, for every solution, the corresponding strategies form a uniform equilibrium with pay-off  $(a, \beta)$ .

**REMARK 3.1.** In principle, one needs II only for those  $(k, l)$  such that  $p^k > 0, p^l > 0$ , but one can always set  $z^l = a^l = \max_{s,t} G_{s,t}^{l,l}$  under all  $p^k$  when  $p^l = 0$ , and obtain then equivalently the above system. Observe that now  $a^l$  is defined, and constrained, even for  $p^l = 0$ .

**PROPOSITION 3.3.** Let  $\tilde{D}_\infty$  be the set of all  $(p, a, \beta, \tilde{P}) \in \Delta(K) \times \mathbb{R}^K \times \mathbb{R} \times \Delta[\Delta(K) \times F \times Z]$  such that, denoting by  $\pi, (c, d)$  and  $z$  the random variables under  $\tilde{P}$  which are the 1<sup>th</sup>, 2<sup>th</sup>, and 3<sup>th</sup> projection, and by  $E$  the expectation under  $\tilde{P}$ :

- (I)  $E(\pi^k c^k) = p^k a^k \quad \forall k \in K$
- (II)  $E \max[E(\pi^k c^l \mid c, d), E(\pi^k z^l \mid c, d)] \leq p^k a^l \quad \forall (k, l) \in K \times K$
- (III)  $E\langle \pi, d \rangle = \beta$
- (IV)  $E(\langle \pi, d \rangle \mid c, d) \geq E[(\text{Vex } u_{\Pi})(\pi) \mid c, d] \quad \text{a.s.}$
- (V)  $E(\pi^k) = p^k \quad \forall k \in K$

Also, for  $\tilde{P} \in \Delta(\Delta(K) \times F \times Z)$  define  $\varphi(\tilde{P}) \in [\Delta(F \times Z)]^K$  by  $[\varphi(\tilde{P})]_k$  being the conditional distribution on  $F \times Z$  given  $k$  under  $\tilde{P}$ , the distribution on  $K \times F \times Z$  induced by  $\tilde{P}$ . And for  $p \in \Delta(K)$ ,  $(P^k)_{k \in K} \in \Delta(F \times Z)$ , let  $\Psi(p, (P^k)_{k \in K})$  be the distribution on  $\Delta(K) \times F \times Z$  of  $[p(\cdot \mid f, z), f, z]$  under the distribution on  $K \times F \times Z$  induced by  $p$  and  $(P^k)_{k \in K}$ . Then the map  $(p, a, \beta, \tilde{P}) \mapsto (p, a, \beta, \varphi(\tilde{P}))$  maps  $\tilde{D}_\infty$  onto  $\overline{D}_\infty$  and  $(p, a, \beta, (P^k)_{k \in K}) \mapsto (p, a, \beta, \Psi(p, (P^k)_{k \in K}))$  maps  $\overline{D}_\infty$  to  $\tilde{D}_\infty$ , and yields the identity on  $\overline{D}_\infty$  when composed with the first.

**PROOF.** Is straightforward. Observe that inequalities I and II, being multiplied by  $p^k$ , become now valid for all  $k$ . Use Jensen's inequality to deduce IV in prop. 3.2 from IV here. ■

**COROLLARY 3.4.** One could equivalently replace IV by

$$(IV') \quad E[\langle \pi, d \rangle \mid c, d] \geq E[u_{\Pi}(\pi) \mid c, d]$$

Or even require in addition that, with probability 1,  $u_{\Pi}(\pi) = \text{Vex } u_{\Pi}(\pi)$ .

**PROOF.** Use a measurable version of Carathéodory's theorem (Castaing and Valadier, 1977, Th. IV.11) or (Mertens, 1987a) to replace  $\pi$  by its image by a transition probability  $Q$  from  $\Delta(K)$  to itself such that,  $\forall \pi, Q(\{q \mid u_{\Pi}(q) = \text{Vex } u_{\Pi}(q)\} \mid \pi) = 1$ . ■

REMARK 3.2. Denote by  $Z_e$  the set of extreme points of  $Z$ . By remarks 3.12 and 3.13 p. 296 — which could already have been made after prop. 3.44 p. 230 —  $\text{Co}(Z_e)$  is compact and  $Z = \text{Co}(Z_e) + \mathbb{R}_+^K$ , and there exists a universally measurable (in fact, one could even have Borel measurability here, due to the finite dimensionality of  $Z$ , using a Borel version of Carathéodory's theorem) transition probability  $Q$  from  $Z$  to  $Z_e$  such that  $\int yQ(z, dy) \leq z$ ,  $\forall z \in Z$ . Such a  $Q$  can be used to modify any  $\tilde{P}$  in  $\tilde{D}_\infty$ , without changing the distribution of  $(\pi, c, d)$  — and hence  $(p, a, \beta)$  —, to another one carried by  $Z_e$  (in II, the expectation given any  $(c, d)$  is decreased, and in IV one uses Jensen's inequality — the new  $z$  is less informative). So nothing essential is changed in  $\tilde{D}_\infty$  by requiring  $\tilde{P}$  to be carried by a compact subset  $\tilde{Z}$  of  $Z$  containing  $Z_e$  (like  $\text{Co}(Z_e)$ , or  $\overline{Z}_e$ , or  $Z \cap [-C, C]^K$ ). Similarly then — cf. remark after prop. 3.2 —  $(p, a, \beta, \tilde{P}) \in \tilde{D}_\infty \Leftrightarrow (p, a', \beta, \tilde{P}) \in \tilde{D}_\infty$  with  $(a')^k = \min(a^k, \overline{C}^k)$  where  $\overline{C}^k = \max_{s,t} G_{s,t}^{I,k}$ . Observe that the inequalities I and II always imply that  $a^k \geq \underline{C}^k = \min_{s,t} G_{s,t}^{I,k}$ . So, for any closed, convex subset  $A$  of  $\mathbb{R}^K$  that contains  $\Pi_{k \in K} [\underline{C}^k, \overline{C}^k]$  and has a maximal point, nothing essential would be changed if we were to restrict  $a$  to  $A$  in the definition of  $\overline{D}_\infty$ . Only the coordinates of  $a$  corresponding to zero-probability states are possibly restricted.

PROPOSITION 3.5. (1)  $\tilde{D} = \{(p, a, \beta, \tilde{P}) \in \tilde{D}_\infty \mid a \in A, \tilde{P}(\tilde{Z}) = 1\}$  is compact.  
(2)  $\tilde{D}_\infty$  is convex in the other variables both for  $P$  fixed and for  $a$  fixed.  
(3)  $\tilde{D}_\infty^1 = \text{Proj}_{\Delta(K) \times \mathbb{R}^K \times \mathbb{R}}(\tilde{D}_\infty)$  is unchanged if  $\tilde{P}$  is restricted to have finite support.

More precisely we can assume  $\#\text{Supp}(\tilde{P}) \leq [(\#K)^2 + (\#K) + 1]^2$  and  $\tilde{P}(Z_e) = 1$ .

PROOF. 1) The set is a subset of the compact space  $\Delta(K) \times (A \cap \Pi_{k \in K} [\underline{C}^k, \infty)) \times [-C, C] \times \Delta[\Delta(K) \times F \times \tilde{Z}]$ . It is closed because, in conditions I, III and V, the left-hand expectation is a continuous linear function of  $\tilde{P}$ , and conditions II and IV are, by a monotone class argument, equivalent to  $E[\pi^k c^l f(c, d) + \pi^k z^l (1 - f(c, d))] \leq p^k a^l$  and  $E\{[(\text{Vex } u_\Pi)(\pi) - \langle \pi, d \rangle] f(c, d)\} \leq 0$  for every continuous  $f: F \rightarrow [0, 1]$ , and now the left hand members are continuous linear functionals of  $\tilde{P}$ .

2) follows in the same way.

3) Choose first  $\tilde{P}$  such that  $\tilde{P}(Z_e) = 1$  (cf. remark after Corollary 3.4). Let  $q(d\pi, dz \mid c, d)$  be regular conditional probability distribution under  $\tilde{P}$  on  $\Delta(K) \times Z_e$  given  $F$  (ex. II.1Ex.16c p. 76). Use  $q$  to transform the left hand member of conditions I, II, III and V into expectations of measurable functions of  $(c, d)$ . There are  $(\#K) \times (\#K + 1)$  such independent equations by deleting one from group V, and deleting the inequalities of group II for  $k = l$ , since, given those of group I, those amount to  $E(\pi^k (c^k - z^k) \mid c, d) \geq 0$  a.e. Delete from  $F$  the negligible set where one of those conditional expectations, computed with  $q$ , is negative and also those points where condition IV, computed with  $q$ , does not hold. Let  $F_0$  be the remaining set.

We can change the marginal of  $\tilde{P}$  on  $F_0$ , while keeping the conditional  $q$ , preserving the value of those  $(\#K) \times (\#K + 1)$  left hand members, and such that the new marginal has at most  $(\#K)^2 + (\#K) + 1$  points (of  $F_0$ ) in its support (ex. I.3Ex.10 p. 34). For each of those points, consider now the conditional distribution  $q$  on  $\Delta(K) \times Z_e$  given this point. It can be changed to any other one provided we preserve the expectations of  $\pi^k$  (conditions I, III and V), and of  $(\text{Vex } u_\Pi)(\pi)$ . Then condition IV is also preserved, and we have now already  $\#K$  different expectations, — and for condition II we just have to preserve in addition the expectations of  $\pi^k z^l$ , yielding  $(\#K)^2$  more expectations: Again we can do this with  $(\#K)^2 + (\#K) + 1$  points. This finishes the proof. ■

COMMENT 3.3. 2) shows something more: Let  $D = \{(p, y, \beta, \tilde{P}) \in \Delta(K) \times \mathbb{R}^{K \times K} \times \mathbb{R} \times \Delta(\Delta(K) \times F \times \tilde{Z}) \mid \text{conditions I to V are satisfied, using } y^{kk} \text{ (resp. } y^{kl}) \text{ in the right hand member of condition I (resp. II), and } (\sum_k y^{kl})_{l \in K} \in A\}$ . Then  $D$  is compact convex, and  $\tilde{D} = \{(p, a, \beta, \tilde{P}) \mid (p, (p^k a^l)_{k, l \in K}, \beta, \tilde{P}) \in D\}$ , i.e.  $\tilde{D}$  is the “linear” one-to-one image by  $(p \mapsto p, \sum_k y^{k,l} \mapsto a^l, \beta \mapsto \beta, \tilde{P} \mapsto \tilde{P})$  of  $D \cap \{(p, y, \beta, \tilde{P}) \mid y^{k,l} = p^k \sum_n y^{n,l}\}$ . It follows in particular that, if  $\mu \in \Delta(\overline{D}_\infty^1)$ , and if  $\beta = \int \hat{\beta} \mu(d\hat{\beta})$ ,  $p = \int \hat{p} \mu(d\hat{p})$ , and  $p^k a^l \geq \int \hat{p}^k \hat{a}^l \mu(d\hat{p}, d\hat{a})$  with equality if  $k = l$ , then  $(p, a, \beta) \in \overline{D}_\infty^1$ , (letting  $\tilde{P} = \int \tilde{P}(\hat{p}, \hat{a}, \hat{\beta}) \mu(d\hat{p}, d\hat{a}, d\hat{\beta})$  where  $\tilde{P}(\hat{p}, \hat{a}, \hat{\beta})$  is a measurable selection). Further call  $(p, a, \beta)$  “extreme” if such a  $\mu$  is unique. Then for every  $(p, a, \beta)$  there is such a  $\mu$  carried by at most  $(\#K)^2 + (\#K) + 1$  extreme points (ex. I.3Ex.10 p. 34).

THEOREM 3.6.  $D_\infty$  exists and equals the projection of  $\overline{D}_\infty^1(p) = \{(a, \beta) \mid (p, a, \beta) \in \overline{D}_\infty^1\}$  on  $\mathbb{R}^{K_0} \times \mathbb{R}$ , with  $K_0 = \{k \mid p^k > 0\}$ .

PROOF. The remark after prop. 3.2 shows that it suffices to prove the theorem when  $p^k > 0 \forall k \in K$ . And prop. 3.2 yields that  $\overline{D}_\infty^1(p)$  consists of uniform equilibria. So there remains to show, using lemma 3.1, that any canonical  $\mathcal{L}$ -communication equilibrium pay-off  $(\alpha, \beta)$  belong to  $\overline{D}_\infty^1(p)$ .

Let  $H_n = \prod_{m=1}^n (M_m^I \times M_m^{\text{II}} \times (S \times T)_m \times I_m^I \times I_m^{\text{II}})$  for  $n = 0, \dots, \infty$ , with  $H = H_\infty$ . Our basic probability space is  $K \times H$  with the  $\sigma$ -fields  $\mathcal{K} \times \mathcal{H}_n$ , where  $\mathcal{K}$  and  $\mathcal{H}_n$  are the (finite)  $\sigma$ -fields generated by  $K$  and  $H_n$ .  $K$  is interpreted as  $I_0^I$ , the initial report of player I; the actual pay-off function (state of nature) is indexed by  $l \in L$ , a copy of  $K$  — as in the inequalities in prop. 3.2 and prop. 3.3. We fix  $p$  as initial probability on  $K$ , assuming implicitly that player I always reports the true state of nature. So his strategies have to be specified only after that; we will only use the set  $\Sigma$  of his strategies where he ignores  $K$  and  $L$  and always reports truthfully player II’s move. The strategy of the device,  $p$ , and any pair of strategies  $\sigma$  and  $\tau$  of the players induce a probability  $P_{\sigma, \tau}$  on  $K \times H$ , with conditional  $P_{\sigma, \tau}^k$  on  $H$ . For the equilibrium (i.e., truthful and obedient) strategies  $\sigma_0$  and  $\tau_0$ ,  $P$  and  $P^k$  will stand for  $P_{\sigma_0, \tau_0}$  and  $P_{\sigma_0, \tau_0}^k$ , and a.e. for  $P$ -a.s..  $\mathsf{E}_{\sigma, \tau}$ ,  $\mathsf{E}_{\sigma, \tau}^k$ ,  $\mathsf{E}$  and  $\mathsf{E}^k$  will denote corresponding (conditional) expectation operators. Observe that the posterior probabilities  $P_{\sigma, \tau}(k \mid \mathcal{H}_n)$  do not really depend on the pair  $(\sigma, \tau)$ , hence we can denote them by  $p_n^k(h)$ , which is a well-defined point in  $\Delta(K)$  for every  $h \in H_n$  that is reachable (i.e., under some pair  $\sigma, \tau$  and for some  $k$ ) given the strategy of the device.  $p_n$  is a  $\Delta(K)$ -valued martingale w.r.t.  $\mathcal{H}_n$ , for every  $P_{\sigma, \tau}$ , and converges, say to  $p_\infty$ ,  $P_{\sigma, \tau}$ -a.e. for every  $(\sigma, \tau)$ .

### PART A. Expected pay-offs (conditions I and III)

Observe  $(\bar{g}_m^{I,l}, \bar{g}_m^{\text{II},l})_{l \in L}$  is an  $F$ -valued random variable on  $H_m$  — thus independent of  $K$ . By ex. I.2Ex.13f p. 24, viewing  $(\bar{g}_m^I, \bar{g}_m^{\text{II}})$  as elements of  $L_\infty^{(2L)}, \sigma(L_\infty, L_1)$ , we have an  $F$ -valued  $\mathcal{H}_\infty$ -measurable random variable  $(c_\infty, d_\infty) = \mathcal{L}(\bar{g}_m^I, \bar{g}_m^{\text{II}})$  — and since  $\mathcal{K} \times \mathcal{H}_n$  is finite for  $n < \infty$ , one has thus for all  $n = 0, 1, \dots, \infty$

$$(1) \quad \mathsf{E}[(c_\infty, d_\infty) \mid \mathcal{K} \times \mathcal{H}_n] = \mathcal{L} \mathsf{E}[(\bar{g}_m^I, \bar{g}_m^{\text{II}}) \mid \mathcal{K} \times \mathcal{H}_n]$$

Further, by definition

$$(2) \quad a^l = \mathcal{L} \mathsf{E}(\bar{g}_m^{I,l} \mid k = l) = \mathsf{E}(c_\infty^l \mid k = l) = \mathsf{E}^l(c_\infty^l)$$

and  $\beta = \sum_l p^l \mathbb{E}(\bar{g}_m^{\Pi,l} \mid k = l) = \sum_l p^l \mathbb{E}(d_\infty^l \mid k = l)$ . And  $d_\infty^l$  is  $\mathcal{H}_\infty$ -measurable, so  $\mathbb{E}(p_\infty^l d_\infty^l) = \mathbb{E}[\mathbb{1}_{\{k=l\}} d_\infty^l] = p^l \mathbb{E}[d_\infty^l \mid k = l]$ , hence

$$(3) \quad \beta = \mathbb{E}\langle p_\infty, d_\infty \rangle$$

#### PART B. Equilibrium condition for player I (condition II)

Let  $x_n^l = \sup_{\sigma \in \Sigma} \mathcal{L} \mathbb{E}_{\sigma, \tau_0} [\bar{g}_m^{I,l} \mid \mathcal{K} \times \mathcal{H}_n]$ :  $x^l$  is a supermartingale w.r.t  $(\mathcal{K} \times \mathcal{H}_n)_{n \in N}$ , under  $P_{\sigma, \tau_0}, \forall \sigma \in \Sigma$  — by a standard “dynamic programming” argument. So

$$(4) \quad \mathbb{E}(x_\infty^l \mid \mathcal{K}) \leq x_0^l \leq a^l$$

because player I has no profitable deviation. And the definition of  $x^l$  implies with (1) that

$$(5) \quad x_n^l \geq \mathcal{L}[\mathbb{E}(\bar{g}_m^{I,l} \mid \mathcal{K} \times \mathcal{H}_n)] = \mathbb{E}[c_\infty^l \mid \mathcal{K} \times \mathcal{H}_n], \quad \text{so } x_\infty^l \geq c_\infty^l \text{ a.e.}$$

Let also  $z_n^l = \sup_{\sigma \in \Sigma} \mathcal{L} \mathbb{E}_{\sigma, \tau_0} (\bar{g}_m^{I,l} \mid \mathcal{H}_n)$ :  $z^l$  is a bounded supermartingale w.r.t.  $\mathcal{H}_n$  and  $P_{\sigma, \tau_0}, \forall \sigma \in \Sigma$ . Further  $z = (z^l)_{l \in L}$  is  $Z$ -valued. Indeed, after every history  $h \in H_n$ ,

- (1) After each stage, I is fully informed of the past, save for  $k \in K$ , since  $\Pi$  uses  $\tau_0$ ;
- (2) The pair formed by the device and player II can be considered as a single opponent, whose strategy  $\tau$  first selects  $k \in K$  with probability  $p_n^k(h)$ , next plays  $\tau^k$ ;
- (3) Viewed this way,  $\Pi$  no longer needs I’s signals (which are truthful anyway), and we can further worsen I’s situation by restricting him to strategies  $\sigma \in \Sigma$  that do not listen to the device’s messages after  $h$  — he remains fully informed of  $h$ .
- (4) Therefore  $z_n(h)$  is now at least as large as the best vector pay-off player I can expect against  $\tau$  in a zero-sum game  $(G^l)_{l \in L}$  with lack of information on one side, and standard signalling. By cor. 3.33.1 p. 218,  $z_n(h) \in Z$ .

Further, I, even when actually forgetting  $k$  as here, can try to exhaust the information about  $k$ . Indeed, for this problem, the pay-off function  $G^l$  is irrelevant, so we can assume that I is ignorant of both  $k$  and  $l$ . And our fictitious opponent can view  $k \in K$  as his initial private information, after which he plays  $\tau^k$ . So we have again a game with incomplete information on one side, this time the opponent being the informed player.

Here every stage consists of first the opponent sending a message to I, next both players choosing simultaneously an action [(3) above]. When putting such a stage in normal form, one obtains a case of non-standard signalling, but where I’s signals are equivalent to knowing  $\mathcal{H}_n$  after every stage  $n$  [(1) above]. Since  $\bar{g}_m^{I,l}$  is  $\mathcal{H}_\infty$ -measurable, we get that, for all  $\sigma$

$$p_n^k |\mathbb{E}_{\sigma, \tau_0}(\bar{g}_m^{I,l} \mid \mathcal{K} \times \mathcal{H}_n) - \mathbb{E}_{\sigma, \tau_0}(\bar{g}_m^{I,l} \mid \mathcal{H}_n)| \leq C p_n^k \|P_{\sigma, \tau_0}(\cdot \mid \mathcal{K} \times \mathcal{H}_n) - P_{\sigma, \tau_0}(\cdot \mid \mathcal{H}_n)\|$$

Hence, by the definition of  $x$  and  $z$ , using  $x^{l,k}(h)$  for  $x^l(k, h)$ :

$$p_n^k(h) |x_n^{l,k}(h) - z_n^l(h)| \leq C' \sqrt{v_n(h)}$$

for  $h \in \mathcal{H}_n$ , using ex. VEx.17 p. 264 in the right hand member. Therefore, under the strategy  $\sigma' = \sigma_\tau$  of that exercise, we get  $p_\infty^k x_\infty^{l,k} = p_\infty^k z_\infty^l, P_{\sigma', \tau_0}$ a.s.. Recall finally from the exercise that the strategy  $\sigma'$  can be chosen to coincide with an arbitrary strategy  $\sigma$  until stage  $n$ . So we get, by the supermartingale property of  $x^l$ , that  $\forall l, \forall k, \forall \sigma, \forall n, P_{\sigma, \tau_0}$ -a.s.:

$$\begin{aligned} p_n^k x_n^{l,k} &\geq p_n^k \mathbb{E}_{\sigma', \tau_0}(x_\infty^{l,k} \mid \{k\} \times \mathcal{H}_n) = \mathbb{E}_{\sigma', \tau_0}[p_\infty^k x_\infty^{l,k} \mid \mathcal{H}_n] = \mathbb{E}_{\sigma', \tau_0}[p_\infty^k z_\infty^l \mid \mathcal{H}_n] \\ &= p_n^k \mathbb{E}_{\sigma', \tau_0}[z_\infty^l \mid \mathcal{K} \times \mathcal{H}_n] \end{aligned}$$

hence  $x_n^l \geq E_{\sigma', \tau_0}[z_\infty^l \mid \mathcal{K} \times \mathcal{H}_n] \forall l, \forall n, P_{\sigma, \tau_0}\text{-a.s. } \forall \sigma$ . Thus, with  $Q_n$  the transition probability from  $K \times H$  to  $Z$  defined by  $Q_n(B) = P_{\sigma', \tau_0}(z_\infty \in B \mid \mathcal{K} \times \mathcal{H}_n)$  for every  $B$  Borel in  $Z$ , we obtain  $x_n^l \geq \int_Z z^l Q_n(dz) \text{ a.s.}$ . Select a limit point  $Q_\infty$  in the sense of ex. II.1Ex.17b p. 76, we have then  $x_\infty^l \geq \int_Z z^l Q_\infty(dz) \text{ a.s.}$ , since  $z^l$  is continuous on  $Z$  and inequalities are preserved under  $\sigma(L_\infty, L_1)$  convergence. Denote by  $P_\infty(P_n)$  the probability on  $K \times H \times Z$  generated by  $P$  and  $Q_\infty(Q_n)$ ; the inequality becomes  $x_\infty^l \geq E_\infty(z^l \mid \mathcal{K} \times \mathcal{H}_\infty) \text{ a.s.}$

Together with (4) and (5) this yields

$$(6) \quad E_\infty(\max\{c_\infty^l, E_\infty(z^l \mid \mathcal{K} \times \mathcal{H}_\infty)\} \mid \mathcal{K}) \leq a^l$$

### PART C. Equilibrium condition for Player II (condition IV)

Since the variable  $z$  constructed sub B appears as an information to player II in his equilibrium condition, we have to find a way to give player II access to it. Player I accessed it by using the strategy  $\sigma'$ . So we will here construct a strategy of II that mimics any strategy  $\sigma' \in \Sigma$  of player I, in the sense of giving, against  $\sigma_0$ , the same history to the device as  $(\sigma', \tau_0)$ , and such that player II can reconstruct by some map  $f$ , from the full history under  $(\sigma_0, \tau')$  — which he knows because of  $\sigma_0$  —, the full history under  $(\sigma', \tau_0)$  which determined  $z$ . It suffices to do this for pure strategies  $\sigma' \in \Sigma$ , which can by Dalkey's theorem (1.3 p. 53), be taken as giving player I's next action as a function of his past messages and of player II's past actions. Player II knows both under  $\sigma_0$ , since player I follows the recommendation, so he can compute this next action  $s_n$  recommended by  $\sigma'$ . His strategy  $\tau' = \tau(\sigma')$  consists them of reporting always to the device that player I used this computed action, and of always following the device's recommendation. It is clear that  $(\sigma_0, \tau')$  will then yield the same history to the device — and hence the same actions by the device — as  $(\sigma', \tau_0)$ , and hence our map  $f$  maps a history of the form  $(m_n^I, m_n^{\text{II}}, m_n^I, m_n^{\text{II}}, m_n^{\text{II}}, s_n)_{n=1}^\infty$  — which are the only ones arising with positive probability under any  $(\sigma_0, \tau')$  — to  $(m_n^I, m_n^{\text{II}}, s_n, m_n^{\text{II}}, m_n^{\text{II}}, s_n)_{n=1}^\infty$ . Since this map  $f$  is independent of the specific  $\sigma'$  used, we can now in the same way transform a mixed strategy  $\sigma'$  into a mixed strategy  $\tau'$ , by transforming all underlying pure strategies, and still use  $f$ , thus obtaining  $P_{\sigma', \tau_0}^k(E) = P_{\sigma_0, \tau'}^k(f^{-1}(E)), \forall k \in K, \forall E \in \mathcal{H}_\infty$ . Since  $\sigma'$  can be assumed to coincide with  $\sigma_0$  during the first  $m$  stages, and since then  $\tau'$  also coincides with  $\tau_0$  during those stages, the same result still holds conditionally to  $h \in \mathcal{H}_m$  with  $P(h) > 0$ . In particular, with  $y_\infty = z_\infty \circ f$ :

$$(7) \quad Q_m(B) = P_{\sigma', \tau_0}(z_\infty \in B \mid \mathcal{K} \times \mathcal{H}_m) = P_{\sigma_0, \tau'}(y_\infty \in B \mid \mathcal{K} \times \mathcal{H}_m) \text{ a.s.}$$

Now  $\beta_n = \sup_\tau \mathcal{L} E_{\sigma_0, \tau}[\bar{g}_m^{\text{II}, k} \mid \mathcal{H}_n]$  is a  $P_{\sigma_0, \tau}$ -supermartingale w.r.t.  $\mathcal{H}_n$ , for every  $\tau$ , because under  $\sigma_0$  player II is fully informed of  $\mathcal{H}_n$ . But also player II can play independently of the device in the future; the game starting after  $h \in H_n$  is then an Indexincomplete information on 1 side!zero-sum repeated@zero-sum repeated games withininfinitely repeated game with incomplete information on one side and standard signalling, viewing player I together with the device as a single informed opponent — which thus no longer needs player II's input. The initial distribution on  $K$  is then  $p_n(h)$ , so that  $\beta_n \geq (\text{Vex } u^{\text{II}})(p_n)$ ,  $P_{\sigma_0, \tau}$ -a.s.  $\forall \tau$ , by theorem 3.5 p. 195. Hence  $\beta_\infty \geq (\text{Vex } u^{\text{II}})(p_\infty)$   $P_{\sigma_0, \tau}$ -a.s., so that,  $y_\infty$  being  $\mathcal{H}_\infty$ -measurable, and using Jensen's inequality (ex. I.3Ex.14bi p. 37):

$$\begin{aligned} \beta_n &\geq E_{\sigma_0, \tau'}[\beta_\infty \mid \mathcal{H}_n] \geq E_{\sigma_0, \tau'}[(\text{Vex } u^{\text{II}})[(P_{\sigma_0, \tau'}(k \mid \mathcal{K}_\infty, y_\infty))_{k \in K}] \mid \mathcal{H}_n] \\ &\geq E_{\sigma_0, \tau'}[(\text{Vex } u^{\text{II}})[(P_{\sigma_0, \tau'}(k \mid \mathcal{K}_n, y_\infty))_{k \in K}] \mid \mathcal{H}_n] \end{aligned}$$

For  $m \geq n$ , assume now that  $\sigma'$  coincides with  $\sigma_0$  until  $m$ , hence  $\tau'$  with  $\tau_0$ . By (7), the joint distribution of  $k$  and  $y_\infty$  given  $\mathcal{H}_n$  under  $P_{\sigma_0, \tau'}$  is the same as that of  $k$  and  $z$  given  $\mathcal{H}_n$  under  $P_m (= P.Q_m)$ . So we get  $\beta_n \geq \mathbb{E}_m[(\text{Vex } u^{\text{II}})[(P_m(k | \mathcal{H}_n, z))_{k \in K}] | \mathcal{H}_n]$  using  $\mathbb{E}_m$  for the expectation under  $P_m$ .

Let now  $m \rightarrow \infty$ . Since the conditional distribution under  $P_m$  on  $K \times Z$  given  $h \in H_n$  has that under  $P_\infty$  as a limit point (weak\*), ex. IIIEx.4 p. 142 yields that

$$\beta_n \geq \mathbb{E}_\infty[(\text{Vex } u^{\text{II}})[(P_\infty(k | \mathcal{H}_n, z))_{k \in K}] | \mathcal{H}_n]$$

So, by the supermartingale property of  $\beta$ , we get for  $i \leq n$

$$\beta_i \geq \mathbb{E}_\infty[(\text{Vex } u^{\text{II}})[(P_\infty(k | \mathcal{H}_n, z))_{k \in K}] | \mathcal{H}_i]$$

Hence, when  $n \rightarrow \infty$ , using the martingale property of conditional probabilities, and the bounded convergence theorem:

$$\beta_i \geq \mathbb{E}_\infty[(\text{Vex } u^{\text{II}})[(P_\infty(k | \mathcal{H}_\infty, z))_{k \in K}] | \mathcal{H}_i]$$

and finally

$$\beta_\infty \geq \mathbb{E}_\infty[(\text{Vex } u^{\text{II}})[(P_\infty(k | \mathcal{H}_\infty, z))_{k \in K}] | \mathcal{H}_\infty] \quad P_\infty\text{-a.s.}$$

By definition,

$$\beta_n \geq \mathcal{L} \mathbb{E}[\bar{g}_m^{\text{II}, k} | \mathcal{H}_n] = \sum_k p_n^k \mathbb{E}(d_\infty^k | k, \mathcal{H}_n)$$

so  $\beta_\infty \geq \langle p_\infty, d_\infty \rangle$ . But, by (3),  $\beta = \mathbb{E}\langle p_\infty, d_\infty \rangle \leq \mathbb{E}(\beta_\infty) \leq \beta_0$ , and one must have equality since player II is in equilibrium. So  $\beta_\infty = \langle p_\infty, d_\infty \rangle$  a.s., and

$$(8) \quad \langle p_\infty, d_\infty \rangle \geq \mathbb{E}_\infty[(\text{Vex } u^{\text{II}})[(P_\infty(k | \mathcal{H}_\infty, z))_{k \in K}] | \mathcal{H}_\infty] \quad P_\infty\text{-a.s.}$$

#### PART D. End of the proof.

Let  $\mathcal{B}_\infty$  be the  $\sigma$ -field generated by  $(c_\infty, d_\infty)$ . By Jensen's inequality  $\mathcal{H}_\infty$  can be replaced by  $\mathcal{B}_\infty$  in (6), and in (8) also — when replacing  $p_\infty$  by  $[P_\infty(k | \mathcal{B}_\infty)]_{k \in K}$  in the left hand member — by taking first conditional expectations w.r.t. the intermediate  $\sigma$ -field spanned by  $\mathcal{B}_\infty$  and  $z$ . Now those modified inequalities, together with (2) and (3), can be expressed in terms of the conditional distributions  $P^k$  of the  $F \times Z$  valued random variable  $(c_\infty, d_\infty, z)$  given  $k$ , yielding the conditions of prop. 3.2. ■

**3.b. “Noisy Channels”; characterisation of  $D_r$  ( $0 < r < \infty$ ).** Let  $\Phi$  denote the space of continuous convex functions on  $\Delta(K)$  which are  $\geq u^{\text{II}}$ . A “noisy channel” is a 1-device where player I first reports  $k \in K$  to the device, next  $(c, d, z, \varphi) \in F \times Z \times \Phi$  is selected according to  $P^k$  and  $(c, d, z)$  is transmitted to player II while  $(c, d, \varphi)$  is to player I. The corresponding strategies of the players are to play the sequence of moves associated with  $(c, d)$  by the same Borel map as sub 3.a until the other deviates from it. From that stage on, player I holds player II down to  $\varphi$  (e.g., in the sense of comments 3.44 and 3.45 p. 226) and player II holds player I down to  $z$  (cor. 3.33 p. 218 or prop. 3.46.2 p. 232). Further player I reports truthfully the state of nature. The device is such that those strategies are in equilibrium.

Let  $N = \{(p, a, \beta, (P^k)_{k \in K}) \mid \text{the above strategies are in equilibrium, with pay-off } (a, \beta) \in \mathbb{R}^K \times \mathbb{R}\}$  (with the understanding that  $P^k$  and  $a^k$  are not to be defined when  $p^k = 0$ ). And  $N^1 = \text{Proj}_{\Delta(K) \times \mathbb{R}^K \times \mathbb{R}}(N)$ . Then we have clearly, as in sect. 3.a — with the same remark after the proposition:

PROPOSITION 3.7.  $(p, a, \beta, (P^k)_{k \in K}) \in \Delta(K) \times \mathbb{R}^K \times \mathbb{R} \times [\Delta(F \times Z \times \Phi)]^K$  is in  $N$  iff:

- (I)  $\mathbb{E}^k(c^k) = a^k \quad \forall k \in K \text{ s.t. } p^k > 0$
- (II)  $\mathbb{E}^k \max[c^l, \mathbb{E}^k(z^l | c, d, \varphi)] \leq a^l \quad \forall (k, l) \in K \times K \text{ s.t. } p^k > 0$
- (III)  $\mathbb{E}\langle p(\cdot | c, d, z), d \rangle = \beta$
- (IV)  $\langle p(\cdot | c, d, z), d \rangle \geq \mathbb{E}[\varphi(p(\cdot | c, d, z, \varphi)) | c, d, z] \quad \text{a.s.}$

More precisely, for every solution, the corresponding strategies form a uniform equilibrium with pay-off  $(a, \beta)$ .

PROPOSITION 3.8. Let  $\tilde{N}$  denote the set of all  $(p, a, \beta, P) \in \Delta(K) \times \mathbb{R}^K \times \mathbb{R} \times \Delta[\Delta(K) \times F \times Z \times \Phi]$  s.t., denoting by  $\pi, (c, d), z$  and  $\varphi$  the random variables under  $P$  which are the successive projections, and by  $\mathbb{E}$  the expectation under  $P$ :

- (I)  $\mathbb{E}(\pi^k c^k) = p^k a^k \quad \forall k \in K$
- (II)  $\mathbb{E} \max[\mathbb{E}(\pi^k c^l | c, d, \varphi), \mathbb{E}(\pi^k x^l | c, d, \varphi)] \leq p^k a^l \quad \forall (k, l) \in K \times K$
- (III)  $\mathbb{E}\langle \pi, d \rangle = \beta$
- (IV)  $\mathbb{E}(\langle \pi, d \rangle | c, d, z) \geq \mathbb{E}(\varphi(\pi) | c, d, z) \quad \text{a.s.}$
- (V)  $\mathbb{E} \pi^k = p^k \quad \forall k \in K.$

Then  $\tilde{N}$  equals  $N$  in the same sense as  $\tilde{D}_\infty$  was related to  $\overline{D}_\infty$  in prop. 3.3.

REMARK 3.4. The same remark applies as after prop. 3.3. In addition, for the same reason  $\Phi$  can be replaced by  $\Phi_e$ , and all functions in  $\Phi_e$  have uniform norm and Lipschitz constant  $\leq C$ . So  $\Phi$  can also be replaced by any compact subset  $\tilde{\Phi}$  of  $\Phi$  that contains  $\Phi_e$ , like the closure, or the closed convex hull, or like all functions in  $\Phi$  having uniform norm and Lipschitz constant  $\leq C$ .

- PROPOSITION 3.9.
- (1)  $\overline{N} = \{(p, a, \beta, P) \in \tilde{N} \mid a \in A, P(\tilde{Z} \times \tilde{\Phi}) = 1\}$  is compact.
  - (2)  $\tilde{N}$  (or  $\overline{N}$ ) is convex in the other variable both for  $p$  fixed and for  $a$  fixed.
  - (3) The same remark applies as after prop. 3.5 p. 411.

COMMENT 3.5. One cannot show as in prop. 3.5 p. 411 that one can always select  $P$  with finite support. Indeed the analogous procedure would be to first replace  $\pi$  by its conditional expectation given  $(c, d, z, \varphi)$ , next try to replace the conditional distribution of  $(z, \varphi)$  given  $(c, d)$  by one with finite support and finally to apply Carathéodory's theorem to  $(c, d)$ . But the second step involves both the conditional distribution of  $z$  given  $\varphi$  and that of  $\varphi$  given  $z$ . An example is given in (Forges, 1988b, footnote 3 p. 202) of the difficulties to which such conditions can lead.

THEOREM 3.10.  $D_r$  ( $0 < r < \infty$ ) exists and equals the projection of  $N^1(p) = \{(a, \beta) \mid (p, a, \beta) \in N^1\}$  on  $\mathbb{R}^{K_0} \times \mathbb{R}$ , with  $K_0 = \{k \mid p^k > 0\}$ .

PROOF. Again it suffices to prove the theorem in case  $K_0 = K$ . By prop. 3.7 and lemma 3.1, it suffices to show that any canonical  $\mathcal{L}$ - $r$ -communication equilibrium pay-off  $(a, \beta)$  belongs to  $N^1(p)$ . We will consider only  $n > r$ , i.e. all strategies coincide with  $(\sigma_0, \tau_0)$  until  $r$ , except possibly for player I being untruthful. The notation (including  $\Sigma$ ) is the same as in the previous section, except that on  $H$  we will, in addition to the  $\sigma$ -fields  $\mathcal{H}_n$ , consider  $\mathcal{H}_n^I$  generated by  $\prod_{m \leq n} (M_m^I \times (S \times T)_m \times I_m^I)$  and  $\mathcal{H}_n^{II}$  by  $\prod_{m \leq n} (M_m^{II} \times (S \times T)_m \times I_m^{II})$ , and  $\mathcal{J}_n = \mathcal{H}_n^I \cap \mathcal{H}_n^{II}$ . Finally  $p_n^k$  stands for  $P_{\sigma, \tau}(\{k\} \mid \mathcal{H}_n^{II})$ .

#### PART A. expected pay-offs.

$(c_\infty, d_\infty) \in F$  is defined as previously, and  $\mathcal{J}_\infty$ -measurable. It satisfies

$$(9) \quad \mathbb{E}[(c_\infty, d_\infty) \mid \mathcal{K} \times \mathcal{H}_n] = \mathcal{L} \mathbb{E}[(\bar{g}_m^I, \bar{g}_m^{II}) \mid \mathcal{K} \times \mathcal{H}_n]$$

$$(10) \quad a^l = \mathbb{E}^l(c_\infty^l)$$

$$(11) \quad \text{and: } \beta = \mathbb{E}\langle p_\infty, d_\infty \rangle, \quad d_\infty \text{ being } \mathcal{H}^{\text{II}}\text{-measurable}$$

**PART B. Equilibrium condition for player I.**

Let now  $x_n^l = \sup_\sigma \mathcal{L} \mathbb{E}_{\sigma, \tau_0}(\bar{g}_m^{l,l} \mid \mathcal{K} \times \mathcal{H}_n^{\text{I}})$ . As before we get

$$(12) \quad \mathbb{E}(x_\infty^l \mid \mathcal{K}) \leq a^l$$

$$(13) \quad \text{and: } x_\infty^l \geq c_\infty^l \text{ a.e., } c_\infty^l \text{ being } \mathcal{H}_\infty^{\text{I}}\text{-measurable}$$

Consider now the game starting at stage  $n$  as a zero-sum game with incomplete information, with  $T^{\text{I}} = L \times K \times H_n^{\text{I}}$  as set of types of player I and  $H_n^{\text{II}} \times \prod_{m>n} M_m^{\text{II}} = T^{\text{II}}$  for player II. Further worsen player I's situation by assuming that the device wont send any messages to him after stage  $n$ . Observe that the distribution  $R$  on  $T^{\text{I}} \times T^{\text{II}}$  is independent of the players' actions in the new game, and that the distribution on  $T^{\text{II}}$  given  $T^{\text{I}}$  is independent of  $L$ , while the pay-off matrix  $G^{\text{I},l}$  depends only on  $L$ . Player II's strategy  $\tau_0$  is a strategy in this game. So  $x^l$  has decreased say to  $\bar{x}^l$ ; and  $\bar{x}_n^l = \sup_\sigma \mathcal{L} \mathbb{E}_{\sigma, \tau_0}(\bar{g}_m^{l,l} \mid \mathcal{K} \times \mathcal{L} \times \mathcal{H}_n^{\text{I}})$ . By VIEx.12bii p.324 there exists a  $T^{\text{II}}$ -measurable map  $z_n$  to  $Z$ , such that  $\bar{x}_n^l \geq \mathbb{E}[z_n^l \mid \mathcal{K} \times \mathcal{L} \times \mathcal{H}_n^{\text{I}}] = \mathbb{E}(z_n^l \mid \mathcal{K} \times \mathcal{H}_n^{\text{I}})$  (by the conditional independence property of  $R$ ), and hence  $x_n^l \geq \mathbb{E}(z_n^l \mid \mathcal{K} \times \mathcal{H}_n^{\text{I}})$ ,  $\forall n$ , for some  $\mathcal{H}_\infty^{\text{II}}$ -measurable maps  $z_n$  to  $Z$ .

The supermartingale property of  $x$  yields then  $x_n^l \geq \mathbb{E}(z_i^l \mid \mathcal{K} \times \mathcal{H}_n^{\text{I}}) \quad \forall i \geq n$ , and hence for a  $\sigma(L_\infty, L_1)$  limit point  $z$  we obtain that  $x_\infty^l \geq \mathbb{E}(z^l \mid \mathcal{K} \times \mathcal{H}_\infty^{\text{I}}) \quad \forall n$ , hence  $z$  is an  $\mathcal{H}_\infty^{\text{I}}$ -measurable random variable with values in  $Z$  such that  $x_\infty^l \geq \mathbb{E}(z^l \mid \mathcal{K} \times \mathcal{H}_\infty^{\text{I}})$ . Combining this with (12) and (13), we obtain

$$(14) \quad a^l \geq \mathbb{E}[\max[c_\infty^l, \mathbb{E}(z^l \mid \mathcal{K} \times \mathcal{H}_\infty^{\text{I}})] \mid \mathcal{K}]$$

**PART C. Equilibrium condition for player II.**

$\beta_n = \sup_\tau \mathcal{L} \mathbb{E}_{\sigma_0, \tau}[\bar{g}_m^{\text{II},k} \mid \mathcal{H}_n^{\text{II}}]$  is a supermartingale,  $\forall \tau$ . Let us minorate  $\beta_n$  by considering the game with incomplete information starting at date  $n$  where this time player II receives no more messages after stage  $n$  and player I receives all his future messages immediately — thus  $T^{\text{II}} = H_n^{\text{II}}, T_n^{\text{I}} = K \times H_n^{\text{I}} \times \prod_{m>n} M_m^{\text{I}}$ . Applying now VIEx.12bi p.324 to this new game, we obtain a  $T^{\text{I}}$ -(hence  $\mathcal{K} \times \mathcal{H}_\infty^{\text{I}}$ -)measurable transition probability  $Q_n$  to  $\Phi$ , such that, for the induced probability  $P_n$  on  $K \times H \times \Phi$ , we have

$$\beta_n \geq \mathbb{E}_n[\varphi[(P_n(k \mid \mathcal{H}_n^{\text{II}}, \varphi))_{k \in K}] \mid \mathcal{H}_n^{\text{II}}]$$

So, by the supermartingale property, we get for  $i \leq n$

$$\beta_i \geq \mathbb{E}_n[\varphi[(P_n(k \mid \mathcal{H}_n^{\text{II}}, \varphi))_{k \in K}] \mid \mathcal{H}_i^{\text{II}}]$$

and since, by Jensen's inequality,

$$\mathbb{E}_n[\varphi[(P_n(k \mid \mathcal{H}_n^{\text{II}}, \varphi))_{k \in K}] \mid \mathcal{H}_i^{\text{II}}, \varphi] \geq \varphi[(P_n(k \mid \mathcal{H}_i^{\text{II}}, \varphi))_{k \in K}]$$

we get

$$\beta_i \geq \mathbb{E}_n[\varphi[(P_n(k \mid \mathcal{H}_i^{\text{II}}, \varphi))_{k \in K}] \mid \mathcal{H}_i^{\text{II}}]$$

We can now let  $n \rightarrow \infty$  using ex.IIIEx.4 p.142; taking then conditional expectations given  $j \leq i$  we obtain

$$\beta_j \geq \mathbb{E}_\infty[\varphi[(P_\infty(k \mid \mathcal{H}_i^{\text{II}}, \varphi))_{k \in K}] \mid \mathcal{H}_j^{\text{II}}]$$

Let now  $i \rightarrow \infty$ , using the martingale property of posteriors and the dominated convergence theorem:

$$\beta_j \geq \mathbb{E}_\infty[\varphi[(P_\infty(k | \mathcal{H}_\infty^{\text{II}}, \varphi))_{k \in K}] | \mathcal{H}_j^{\text{II}}]$$

so finally we can let  $j \rightarrow \infty$

$$\beta_\infty \geq \mathbb{E}_\infty[\varphi[(P_\infty(k | \mathcal{H}_\infty^{\text{II}}, \varphi))_{k \in K}] | \mathcal{H}_\infty^{\text{II}}]$$

As in the previous theorem, one obtains now  $\beta_\infty = \langle p_\infty, d_\infty \rangle$  a.e., so

$$(15) \quad \langle p_\infty, d_\infty \rangle \geq \mathbb{E}_\infty[\varphi[(P_\infty(k | \mathcal{H}_\infty^{\text{II}}, \varphi))_{k \in K}] | \mathcal{H}_\infty^{\text{II}}]$$

#### PART D. End of the proof.

Since  $Q_\infty$  is  $\mathcal{K} \times \mathcal{H}_\infty^{\text{I}}$ -measurable, we have that under  $P_\infty$ ,  $\mathcal{H}_\infty^{\text{II}}$  and  $\varphi$  are conditionally independent given  $\mathcal{K} \times \mathcal{H}_\infty^{\text{I}}$ . Therefore, since  $z$  is  $\mathcal{H}_\infty^{\text{II}}$ -measurable,  $\mathbb{E}_\infty[z | \mathcal{K} \times \mathcal{H}_\infty^{\text{I}}, \varphi] = \mathbb{E}[z | \mathcal{K} \times \mathcal{H}_\infty^{\text{I}}]$ . So (14) yields

$$a^l \geq \mathbb{E}_\infty[\max[c_\infty^l, \mathbb{E}_\infty(z^l | \mathcal{K} \times \mathcal{H}_\infty^{\text{I}}, \varphi)] | \mathcal{K}]$$

Jensen's inequality allows to decrease here  $\mathcal{H}_\infty^{\text{I}}$  to the  $\sigma$ -field generated by  $(c_\infty, d_\infty)$ ; similarly in (15) we can first reduce the  $\mathcal{H}_\infty^{\text{II}}$  appearing in  $p_\infty$  and that in the conditional expectation to the  $\sigma$ -field spanned by  $(c_\infty, d_\infty, z)$  (recall  $z$  is  $\mathcal{H}_\infty^{\text{II}}$ -measurable), next use Jensen's inequality as before to do the same replacement for  $\mathcal{H}_\infty^{\text{II}}$  in the conditional probability. This, plus conditions (10) and (12), shows that the conditional distributions  $P^k$  of  $(c_\infty, d_\infty, z, \varphi)$  given  $k$ , under  $P_\infty$ , satisfy the conditions of prop. 3.7.

■

### Exercises

**1. Incomplete information on the opponent's pay-off.** (Shalev, 1988); (Koren, 1992); (Israeli, 1989) Consider a game with incomplete information on one side as in sect. 1 where moreover  $G^{\text{II}}$  is independent of  $k$ .

a. Prove that  $E_\infty(p)$  is the set  $L(p)$  of completely revealing E.J.P. pay-offs, i.e. satisfying  $R = K$  and  $p^k(r) = \mathbb{1}_{\{r\}}(k)$  in prop. 1.1 p. 398. (Observe it depends only on the support of  $p$ ).

HINT. To prove that  $E_\infty \subseteq L$ :

- (1) either use theorem 1.4 p. 400, prop. 2.3 p. 407 and show that  $L$  can be strictly separated from any point in its complement by a bi-convex continuous function; or
- (2) define  $\theta_{st}(k) = \mathcal{L}(\mathbb{E}_{\sigma^k, \tau}(\bar{l}_n(s, t)))$  where  $\bar{l}_n(s, t)$  is the empirical frequency of  $(s, t)$  up to stage  $n$ , and introduce  $a, b, \alpha, \beta$  as in 1.a p. 397.

Note that  $\beta_m = \mathcal{L}(\mathbb{E}_{\sigma, \tau, p}(\bar{g}_n^{\text{II}} | \mathcal{H}_m)) = \mathcal{L}(\mathbb{E}'_{\sigma, \tau}(\bar{g}_n^{\text{II}} | \mathcal{H}_m)) \geq u_{\text{II}}$ .

Show that  $\mathbb{E}_{\sigma^k, \tau}[\mathcal{L}(\mathbb{E}'_{\sigma, \tau}(\bar{g}_n^{\text{II}} | \mathcal{H}_m)) - \mathcal{L}(\mathbb{E}_{\sigma^k, \tau}(\bar{g}_n^{\text{II}}))] \xrightarrow[m \rightarrow \infty]{} 0$ , so  $\beta(k) = G^{\text{II}}[\theta(k)] = \sum_{s,t} G_{st}^{\text{II}} \theta_{st}(k) = \mathcal{L}(\mathbb{E}_{\sigma^k, \tau}(\bar{g}_n^{\text{II}}))$ . Use then prop. 1.1 p. 398.

b. Prove that  $E_\infty(p) \neq \emptyset$  for all  $p$ .

HINT. Use a splitting leading to a S.E.J.P.

c. Prove that  $D_\infty = E_\infty$ .

HINT. Extend the construction of 2 above.

d. Given two matrices  $A$  and  $B$ , consider all the games defined as sub IXEx.1a above with  $G^{1,\text{I}} = A$  and  $G^{2,\text{I}} = B$ . Show that there exists  $\zeta$  (function of the pay-off matrices) such that the projection of  $E_\infty$  on the  $(a^1, \beta)$  components is the set of feasible pay-offs in  $(A, B)$  satisfying  $a^1 \geq \zeta, \beta \geq u_{\text{II}}$ .

Prove that the maximum of such  $\zeta$  is

$$\max_{x \in \Delta(S)} \min_{y \in B(x)} xAy, \quad \text{with } B(x) = \{ y \in \Delta(T) \mid xBy \geq u_{\text{II}} \},$$

and is obtained for  $K = 2$  and  $G^{2,\text{I}} = -B$ .

HINT. Use a fully revealing equilibrium with  $x^k$  maximising  $\min_{y \in B(x)} xG^{k,\text{I}}y$ . For the other inequality, use that  $a$  has to be approachable by  $\text{II}$ .

e. *Incomplete information on both sides.* Consider a non-zero-sum game in the independent case ( $K = L \times M$ ), with standard signalling satisfying moreover:

$$G^{(l,m),\text{I}} = G^{l,\text{I}}, \quad G^{(l,m),\text{II}} = G^{m,\text{II}} \text{ and } \#S \geq \#L, \#T \geq \#M.$$

Prove that:

$$\begin{aligned} E_\infty(p, q) &= \left\{ (a, b) \in \mathbb{R}^l \times \mathbb{R}^m \mid \exists \theta(l, m) \in \Delta(S \times T), c(l, m) \in \mathbb{R}^l, d(l, m) \in \mathbb{R}^m, \text{ s.t. :} \right. \\ &\quad \sum_m q^m G^{l,\text{I}}[\theta(l, m)] = a^l, \quad \sum_l p^l G^{m,\text{II}}[\theta(l, m)] = b^m, \\ &\quad a^l \geq \sum_m q^m \max\{G^{l,\text{I}}[\theta(l', m)], c^l(l', m)\}, \quad \forall l', \\ &\quad b^m \geq \sum_l p^l \max\{G^{m,\text{II}}[\theta(l, m')], d^m(l, m')\}, \quad \forall m', \\ &\quad \left. \forall l, m: \quad \langle c(l, m), \pi \rangle \geq u_{\text{I}}(\pi) \quad \forall \pi \in \Delta(L), \quad \langle d(l, m), \rho \rangle \geq u_{\text{II}}(\rho) \quad \forall \rho \in \Delta(M) \right\} \end{aligned}$$

HINT. To get equilibrium strategies use a joint plan completely revealing at stage one; then player  $\text{II}$  approaches  $c(l, m)$  if  $(l, m)$  is announced and player  $\text{I}$  does not follow  $\theta(l, m)$ . For the other inclusion, define  $\theta(l, m)$  as in a) and introduce:

$$c^l(l', m) = \limsup_t E_{\sigma'^t, \tau^m} \sup_{\sigma'} \mathcal{L}(E'_{\sigma', \tau}(\bar{g}_n^{l,\text{I}} \mid \mathcal{H}_t))$$

f. Show in the following example:

$$G^{1,\text{I}} = \begin{pmatrix} 3 & 0 \\ 0 & 1 \end{pmatrix}, \quad G^{2,\text{I}} = \begin{pmatrix} 3 & 3 \\ 1 & 1 \end{pmatrix}, \quad G^{1,\text{II}} = \begin{pmatrix} 1 & 0 \\ 0 & 3 \end{pmatrix}, \quad G^{2,\text{II}} = \begin{pmatrix} 1 & 3 \\ 1 & 3 \end{pmatrix}$$

that  $E_\infty(p)$  is non-empty iff  $p^1 \leq 1/6$  or  $q^1 \leq 1/6$  or  $1/p^1 + 1/q^1 \geq 10$ .

**2. On the number of revelation stages.** (Forges, 1984) In the notation of sect. 1, let  $W_n$  be obtained from  $W$  by  $n$  steps of “bi-convexification” (each step is first in  $p$  then in  $a$ ).

a. Show that in the following game, where the pay-off is independent of  $\text{I}$ 's move:

$$\begin{aligned} G^1 &= ((10, -10); (4, -3); (4, 0); (0, 5); (10, 8); (0, 9); (8, 10)) \\ G^2 &= ((8, 10); (0, 9); (10, 8); (6, 5); (4, 0); (4, -3); (10, -10)) \end{aligned}$$

$W_2$  differs from  $W_1$ .

b. In the following example:

$$\begin{aligned} G^1 &= ((6, -6); (2, 1); (8, 6); (4, 9); (0, -10); (0, -2); (0, 4); (0, 8); (0, 10)) \\ G^2 &= ((4, 9); (8, 6); (2, 1); (6, -6); (0, 10); (0, 8); (0, 4); (0, -2); (0, -10)) \end{aligned}$$

$W^\star \neq \text{bico } W$  (cf. example of figure 1).

HINT. Look at the points in  $W$  with  $a^1 + a^2 = 10$ .

c. Finally take:

$$\begin{aligned}G^1 &= ((0, -5); (0, -1); (-1, 2); (3, 4); (1, 5)) \\G^2 &= ((1, 5); (3, 4); (-1, 2); (0, -1); (0, -5))\end{aligned}$$

and show that  $W_2 = \text{bico } W$  but that  $W, W_1, W_2$  and  $W^*$  all differ: one obtains the configuration of figure 1.

## APPENDIX A

### Reminder on Analytic Sets<sup>1</sup>

#### 1. Notation

Denote by  $\mathcal{P}$  a collection of subsets — a “paving” — of a set  $X$ . Let  $\mathbb{T} = \mathbb{N}^{\mathbb{N}}$ ,  $\mathbb{T}_n = \mathbb{N}^n$ ,  $\mathbb{T}_f = \bigcup_n \mathbb{T}_n$ . For  $t \in \mathbb{T}$  or  $t \in \mathbb{T}_k$ ,  $k \geq n$ ,  $t_n$  is the natural projection to  $\mathbb{T}_n$ . A (disjoint) Souslin scheme is a map  $t \mapsto P_t$  from  $\mathbb{T}_f$  to  $\mathcal{P}$  (such that  $(\bigcap_n P_{s_n}) \cap (\bigcap_n P_{t_n}) = \emptyset \forall s \neq t \in \mathbb{T}$ ), and has as kernel  $\bigcup_{t \in \mathbb{T}} \bigcap_n P_{t_n}$ .  $\mathcal{P}_s$  (resp.  $\mathcal{P}_{s_d}$ ) is the paving consisting of the kernels of all (disjoint) Souslin schemes.  $\mathcal{P}_\sigma$ ,  $\mathcal{P}_+$ ,  $\mathcal{P}_\delta$ ,  $\mathcal{P}_c$  denote the pavings consisting respectively of the countable unions, the countable disjoint unions, the countable intersections, and the complements of elements of  $\mathcal{P}$ .  $\mathcal{P}_{\sigma\delta} = (\mathcal{P}_\sigma)_\delta$ , and so on.  $\mathcal{P}_{(\alpha)}$ , where  $\alpha$  is a string of one or more of the above operations on pavings, denotes the stabilisation of  $\mathcal{P}$  under the corresponding operations, i.e., the smallest paving  $\mathcal{P}^*$  containing  $\mathcal{P}$  such that  $\mathcal{P}^* = \mathcal{P}_O^*$  for every operation  $O$  in the string  $\alpha$ . If  $X$  is a topological space,  $\mathcal{F}, \mathcal{G}, \mathcal{K}, \mathcal{L}$  will denote the pavings of closed, open, compact and zero sets (i.e., sets  $f^{-1}(0)$  for  $f$  real valued and continuous).

#### 2. Souslin schemes

**2.a.**  $\mathcal{P}_s = \mathcal{P}_{(s,\sigma,\delta)}$ ,  $\mathcal{P}_{s_d} = \mathcal{P}_{(s_d,+,\delta)}$ .

**2.b.**  $\mathcal{P}_s$  is the set of projections on  $X$  of  $(\mathcal{P} \times \mathcal{J})_{\sigma\delta}$  in  $X \times \mathbb{T}$  or  $X \times [0, 1]$ , where  $\mathcal{J}$  is the paving of closed intervals (i.e., on  $\mathbb{T}$  the subsets with a fixed initial segment in  $\mathbb{T}_f$ .)

**2.c.** If  $\emptyset \in \mathcal{P}$ ,  $\mathcal{P}_s$  is the set of projections on  $X$  of  $(\mathcal{P} \times \mathcal{F})_s$  in  $X \times Y$  if  $Y$  is  $K$ -analytic (cf. below).

**2.d. Second Separation Theorem.** (cf. II.2Ex.5 p.85) Assume  $\mathcal{P}_c \subseteq \mathcal{P}_s$ . If  $C_n \in \mathcal{P}_{sc}$  ( $n \in \mathbb{N}$ ),  $\exists D_n \in \mathcal{P}_{sc}$  s.t.  $D_n \subseteq C_n$  and  $D_n$  is a partition of  $\bigcup_n C_n$ . In particular:

Assume  $\mathcal{P}_c \subseteq \mathcal{P}_s$ . If  $A_n \in \mathcal{P}_s$ ,  $A_n \cap A_m = \emptyset$ ,  $\exists B_n \in \mathcal{P}_s \cap \mathcal{P}_{sc}$ :  $A_n \subseteq B_n$ ,  $B_n \cap B_m = \emptyset$ .

**2.e.** Assume  $\mathcal{P}_c \subseteq \mathcal{P}_s$ . Then  $\mathcal{P}_{s_d} \subseteq \mathcal{P}_{sc}$ .

**2.f.** Assume  $\mathcal{P}_c \subseteq \mathcal{P}_{(\sigma,\delta)}$ . Then the  $\sigma$ -field  $\mathcal{P}_{(\sigma,\delta)} \subseteq \mathcal{P}_{s_d}$  and every  $X \in \mathcal{P}_{(\sigma,\delta)}$  is the kernel of a disjoint Souslin scheme  $P_t$  with  $\bigcup_{t_n=s_n} \bigcap_k P_{t_k} \in \mathcal{P}_{(\sigma,\delta)} \forall s_n \in \mathbb{T}_f$ . (This is hard — not really an exercise).

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<sup>1</sup>E.g., (Rogers et al., 1980)

### 3. $K$ -analytic and $K$ -Lusin spaces

A Hausdorff topological space  $X$  is called  $K$ -analytic ( $K$ -Lusin) if it is the image of  $\mathbb{T}$  by a compact valued u.s.c. map  $\Gamma$  (with disjoint values). [ $\Gamma$  is u.s.c. (upper semi-continuous) means  $\{t \mid \Gamma(t) \subseteq O\} \in \mathcal{G}^{\mathbb{T}}$ ,  $\forall O \in \mathcal{G}^X$ ]. Denote by  $\mathcal{A}^X$  ( $\mathcal{L}^X$ ) the paving of  $K$ -analytic ( $K$ -Lusin) subspaces of a Hausdorff space  $X$ .  $\mathcal{Li}^X$  will denote the Lindelöf subsets (every open cover contains a countable subcover).

**3.a.** Images of  $K$ -analytic ( $K$ -Lusin) spaces by compact valued u.s.c. maps (with disjoint values) are  $K$ -analytic ( $K$ -Lusin). Their countable products and closed subspaces are also.  $\mathcal{A}^X \subseteq \mathcal{Li}^X$ .

**3.b.**  $\mathcal{L} = \mathcal{L}_{sd} \subseteq \mathcal{F}_{sd}$ ;  $\mathcal{A} = \mathcal{A}_s \subseteq \mathcal{F}_s$ ;

**3.c.** For  $Y$   $K$ -analytic,  $\text{Proj}_X(\mathcal{F}_s^{X \times Y}) = \mathcal{F}_s^X$  and, if  $X$  and  $Y$  are in addition completely regular:  $\text{Proj}_X(\widetilde{\mathcal{Z}}_s^{X \times Y}) \subseteq \mathcal{Z}_s^X$ , with  $\widetilde{\mathcal{Z}} = \mathcal{Z} \cap \mathcal{Li}$ .

**3.d.**  $\mathcal{Z}_c \subseteq \mathcal{Z}_{+\delta+}$  (hence  $\mathcal{Z}_{(+,\delta)}$  is a  $\sigma$ -field — the Baire  $\sigma$ -field).

A regular Lindelöf space is paracompact. For  $X$  compact,  $\mathcal{Z}_{c,\delta}^X = \mathcal{Li}^X \cap \mathcal{G}_\delta^X$ .

**3.e.** For  $X \in \mathcal{A}^X$ , denote by  $G$  the graph in  $\mathbb{T} \times X$  of the corresponding u.s.c. map.  $G$  is a  $\mathcal{Z}_{c\delta}$ -subset of its Stone-Čech compactification  $\widehat{G}$ .

HINT. First prove  $G$  is  $T_3$ ; by (3.a)  $G$  is Lindelöf; use then (3.d).

[In particular,  $K$ -Lusin spaces are the continuous one-to-one images of the  $\mathcal{K}_{+\delta+}$  subsets of compact spaces, and the  $K$ -analytic spaces the continuous images of  $\mathcal{K}_{\delta}$  subsets of compact spaces.]

**3.f. The first separation theorem.** For  $A_n \in \mathcal{A}^X$ ,  $F \in \mathcal{F}_s^X$ , with  $F \cap (\bigcap_{n \in \mathbb{N}} A_n) = \emptyset$ ,  $\exists B_n \in \mathcal{G}_{(\sigma,\delta)}^X$ ,  $B_n \supseteq A_n$ ,  $F \cap (\bigcap_{n \in \mathbb{N}} B_n) = \emptyset$ . (So one can increase  $F$  to an  $\mathcal{F}_{(\sigma,\delta)}$ -set). One could replace  $\mathcal{G}$  by any paving  $\mathcal{P}$  containing a basis of neighbourhoods of every point — in particular, by  $\mathcal{F}$  if the space is regular, and by  $\mathcal{Z}$  if completely regular. Note the following consequences:

**3.g.**  $\mathcal{L}^X \subseteq \mathcal{F}_{(\sigma,c)}^X$  (the Borel sets)

$$\mathcal{G}^X \subseteq \mathcal{F}_s^X \Rightarrow \mathcal{L}^X \subseteq \mathcal{G}_{(\sigma,\delta)}^X$$

$$X \text{ regular} \Rightarrow \mathcal{L}^X \subseteq \mathcal{F}_{(\sigma,\delta)}^X$$

$$X \text{ completely regular} \Rightarrow \mathcal{L}^X \subseteq \mathcal{F}_{(+\delta)}^X$$

$$X \text{ completely regular and } \mathcal{G}^X \subseteq \mathcal{F}_s^X \Rightarrow \mathcal{L}^X \subseteq \mathcal{Z}_{(+\delta)}^X$$

$X$  completely regular,  $L \in \mathcal{L}^X \Rightarrow [L \in \mathcal{Z}_{(+\delta)}^X \Leftrightarrow \exists f: X \rightarrow \mathbb{R}^N \text{ continuous, } f(L) \cap f(\complement L) = \emptyset]$ .

**3.h.** For  $X \in \mathcal{A}^X$ , one has:

$$\mathcal{F}_s^X \cap \mathcal{F}_{sc}^X = \mathcal{F}_{(\sigma,\delta)}^X \cap \mathcal{G}_{(\sigma,\delta)}^X \text{ (by 3.f) [is a } \sigma\text{-field between Baire and Borel } \sigma\text{-fields.]}$$

In particular,  $\mathcal{G}^X \subseteq \mathcal{F}_s^X \Rightarrow \mathcal{G}_{(\sigma,\delta)}^X = \mathcal{F}_{(\sigma,\delta)}^X$  (= Borel sets,  $= \mathcal{F}_{sd}^X$  by 2.e and 2.f).

$$X \text{ completely regular} \Rightarrow \mathcal{F}_{(+\delta)}^X = \mathcal{F}_{sd}^X (\supseteq \mathcal{F}_s^X \cap \mathcal{F}_{sc}^X = \mathcal{Z}_{(+\delta)}^X).$$

$$\mathcal{Z}_{(+\delta)}^X = \mathcal{Z}_s^X \cap \mathcal{Z}_{sc}^X \text{ (hence } = \mathcal{Z}_{sd}^X \text{ by 2.e) (the Baire sets).}$$

#### 4. Capacities

$E_n$  and  $F$  are Hausdorff spaces.

A *multiplicity* on  $\prod_n E_n$  is a map  $I$  from  $\prod_n \wp(E_n)$  to  $\overline{\mathbb{R}}_+$  which is

- monotone:  $X_n \subseteq Y_n \Rightarrow I(X_1, X_2, \dots) \leq I(Y_1, Y_2, \dots)$
- separately left continuous:  $X_n^k \uparrow X_n \Rightarrow I(X_1, X_2, \dots, X_n^k, X_{n+1}, \dots) \uparrow I(X_1, X_2, \dots, X_n, X_{n+1}, \dots)$
- right continuous: for  $K_n$  compact and  $\varepsilon > 0$  there exist open sets  $U_n \supseteq K_n$  with  $U_n = E_n$  for all but finitely many indices such that  $I(U_1, U_2, \dots) \leq I(K_1, K_2, \dots) + \varepsilon$ .

A *capacity operation* on  $\prod_n E_n$  with values in  $F$  is a map  $I$  from  $\prod_n \wp(E_n)$  to  $\wp(F)$  which is

- monotone
- separately left continuous (i.e.  $\forall f \in F$ ,  $\mathbb{1}_{I(\dots)}(f)$  is a multiplicity), and
- right continuous: for  $K_n$  compact,  $I(K_1, K_2, \dots)$  is compact, and for each of its neighbourhoods  $V$  there exist open sets  $U_n \supseteq K_n$  with  $U_n = E_n$  for all but finitely many indices, such that  $I(U_1, U_2, \dots) \subseteq V$ .

**4.a.** If  $I$  is a capacity operation (multiplicity) on  $\prod_n E_n$ , and the  $J_n$ 's are capacity operations from  $\prod_k E_{n,k}$  to  $E_n$ , then the composition is a capacity operation (multiplicity).

#### 4.b.

4.b.1. If  $\Gamma$  is an u.s.c. map from  $E$  to  $\mathcal{K}^F$ , then  $I_\Gamma(X) = \bigcup_{x \in X} \Gamma(x)$  is a capacity operation from  $E$  to  $F$ . (In particular, any  $K$ -analytic set  $A$  is the form  $I_{\Gamma_A}(\mathbb{T})$ ).

4.b.2.  $\bigcap_n X_n$  and  $\prod_n X_n$  are capacity operations  $J$  and  $P$ . (In particular,  $\mathbb{T} = P(\mathbb{N}, \mathbb{N}, \mathbb{N}, \dots)$ , so  $A = I_A(\mathbb{N}, \mathbb{N}, \mathbb{N}, \dots)$  for  $A \in \mathcal{A}$  and  $I_A = I_{\Gamma_A} \circ P$ .)

4.b.3. Given a capacity operation  $I$  with  $E_n = \mathbb{N}$ , define  $\psi_I: \mathbb{T} \rightarrow \mathcal{K}^F$  by  $\psi_I(n_1, n_2, \dots) = I(\bar{n}_1, \bar{n}_2, \dots)$ , with  $\bar{n} = \{1, 2, \dots, n\}$ :  $I(\mathbb{N}, \mathbb{N}, \dots) = \sup I(\bar{n}_1, \bar{n}_2, \dots)$  — reducing first by evaluation to a multiplicity, then using first left continuity, finally right continuity. Also  $\psi_I$  is u.s.c., so  $I(\mathbb{N}, \mathbb{N}, \mathbb{N}, \dots) = \psi_I(\mathbb{T})$  is  $K$ -analytic.

4.b.4. It follows thus from (4.a here above) that capacity operations map  $K$ -analytic arguments to  $K$ -analytic values, and conversely every  $K$ -analytic set can be obtained in this way, using just  $(\mathbb{N}, \mathbb{N}, \dots)$  as argument.

4.b.5. With  $E_0 = \{1, 2, 3, \dots\}$ ,  $J(X_0, X_1, X_2, \dots) = \bigcup_{n \in X_0} X_n$  is a capacity operation [so  $\bigcup_{n \geq 1} X_i = J(\mathbb{N}, X_1, X_2, \dots)$ ].

4.b.6. With  $E_0 = \mathbb{T}$ ,  $E_t = X$  for  $t \in \mathbb{T}_f$ ,  $I(X_0, \dots, X_t, \dots) = \bigcup_{x \in X_0} \bigcap_n X_{x_n}$  is a capacity operation, yielding the kernel of the Souslin scheme as  $I(\mathbb{T}, \dots, X_t, \dots)$ .

4.b.7. If  $\mu$  is a probability measure on the Borel sets, such that  $\mu(K) = \inf\{\mu(O) \mid K \subseteq O, O \text{ open}\}$  for any compact set  $K$ , the outer measure  $\mu^*$  is a capacity ( $\mu^*(A) = \inf\{\mu(B) \mid A \subseteq B, B \text{ Borel}\}$ , for every subset  $A$ ).

4.b.8. For  $X$  completely regular, denote by  $M_{\sigma,+}^X$  the space of non-negative countably additive bounded measures on the Baire  $\sigma$ -field, which are regular in the sense of inner approximation by zero-sets. For  $\mu \in M_{\sigma,+}^X$ , denote by  $\mu^*$  the corresponding outer measure. Endow  $M_{\sigma,+}^X$  with the weak\*-topology determined by the duality with all bounded continuous functions.

Let  $I(A, B) = \sup_{\mu \in A} \mu^*(B)$  for  $B \subseteq X$ ,  $A \subseteq M_{\sigma,+}^X$ .  $I$  is a bi-capacity.

HINT.  $\mu^*(K) = \inf\{\int f d\mu \mid f > \mathbb{1}_K, f \text{ continuous and bounded}\}$  by regularity. Use Dini's theorem on the compact subset of  $M_{\sigma,+}^X$  to select some  $f_0$ .

**4.c.** For  $I$  a capacity operation or a multicapacity, one has:

$$I(A_1, A_2, \dots) = \sup_{K_n \subseteq A_n} I(K_1, K_2, \dots) = \inf_{B_n \supseteq A_n} I(B_1, B_2, \dots)$$

for  $A_n$   $K$ -analytic,  $K_n$  compact,  $B_n \in \mathcal{G}_{(\sigma,\delta)}$  (or  $B_n \in \mathcal{P}_{(\sigma,\delta)}$  where the paving  $\mathcal{P}$  contains a basis of neighbourhoods of each point) (for a multicapacity, the  $\inf$  is obviously achieved).

HINT. It suffices to consider the case of a multicapacity. The first formula follows then from (4.a) and (4.b.4), which reduce the problem to the case  $A_n = \mathbb{N}$ , which was solved in (4.b.3). The second follows now by showing that  $J(X_1, X_2, \dots) = \inf_{B_n \supseteq X_n} I(B_1, B_2, \dots)$  is a multicapacity, equal to  $I$  on compact sets, hence by the first formula also equal on  $K$ -analytic sets.

[In particular, this yields the first separation theorem (3.f p. 422), at least for  $F \in \mathcal{F}_\sigma$ . It suffices indeed to consider  $F \in \mathcal{F}$ , then  $J_F(Y) = \mathbb{1}_{F \cap Y \neq \emptyset}$  is a capacity, and  $I(Y_1, Y_2, \dots) = \bigcap_n Y_n$  a capacity operation (4.b.2), so the above can be applied to the multicapacity (4.a)  $J_F \circ I$ .]

#### 4.d.

4.d.1. Given a measurable space  $(\Omega, \mathcal{B})$ , define the  $\sigma$ -field  $\mathcal{B}_u$  of universally measurable sets as consisting of those sets which are  $\mu$ -measurable for any probability measure — and then for any measure — on  $(\Omega, \mathcal{B})$ . Then  $\mathcal{B}_u = \mathcal{B}_{us}$ .

HINT. Since  $\mathcal{B}_u = \mathcal{B}_{uu}$ , it suffices to show that  $\mathcal{B}_s \subseteq \mathcal{B}_u$ . Reduce first to the case  $\mathcal{B}$  separable, then to  $(\Omega, \mathcal{B}) = ([0, 1], \mathcal{F}_{(\sigma,\delta)})$ , then use (4.c).

4.d.2. A probability as in (4.b.7) (or (4.b.8) on a  $K$ -analytic space  $X$ ) is regular.

HINT. Reduce by (4.c) to the case of a compact space; use then Riesz's theorem.

4.d.3. If  $A \in \mathcal{A}^X$ ,  $X$  completely regular, and  $L \in \mathcal{A}^{M_{\sigma,+}^X}$  are such that  $\mu(A) = 0 \forall \mu \in L$ , there exists a Baire set  $B$  containing  $A$ , such that  $\mu(B) = 0 \forall \mu \in L$  — using 4.b.8 and 4.c

## 5. Polish, analytic and Lusin spaces

DEFINITION 5.1. A polish space is a regular Hausdorff space which is a continuous, open image of  $\mathbb{T}$ . An analytic (Lusin) space is a  $K$ -analytic ( $K$ -Lusin space where  $\mathcal{G} \subseteq \mathcal{F}_s$  and where the Borel  $\sigma$ -field is separable).

**5.a.**  $X$  is polish iff it is separable and can be endowed with a complete metric. (cf. ex. II.2Ex.8d p. 87) The separable metric spaces are the subspaces of the compact metric spaces. If  $\overline{X} = P$ ,  $P$  Hausdorff, then  $X \in \mathcal{G}_\delta^P$  if  $X$  can be endowed with a complete metric. Conversely, a metrisable  $\mathcal{G}_\delta$  in a compact space or in a complete metric space can be endowed with a complete metric. A polish space is homeomorphic to a closed subspace of  $\mathbb{R}^\mathbb{N}$ ; conversely, a metrisable space which has a proper map to  $\mathbb{R}^\mathbb{N}$  is polish. [A map  $f$  is proper if it is continuous, maps closed sets to closed sets and has compact point inverses.]

**5.b.** The analytic spaces are the continuous, Hausdorff images of  $\mathbb{T}$ . The uncountable Lusin spaces are the continuous, one-to-one images of  $\mathbb{T} \cup \mathbb{N}$ .

HINT. To prove separability of the Borel  $\sigma$ -field, use that  $(X \times X) \setminus \Delta$  ( $\Delta$  is the diagonal) is the continuous image of a separable metric space, hence Lindelöf, so there exists a weaker Hausdorff topology with a countable base — and use 3.h p. 422. In the other direction, show that closed subspaces of  $\mathbb{T}$  are continuous images of  $\mathbb{T}$  — and, if uncountable, continuous one-to-one images of  $\mathbb{T} \cup \mathbb{N}$ . Also, one does not need the Borel  $\sigma$ -field to be separable, it suffices that it contains a separable and separating sub- $\sigma$ -field; this is then generated by a sequence  $O_n$  of open sets and their complements, which separate points, and belong to  $\mathcal{F}_s$  by assumption. Apply then the following lemma:

**Lemma** For a  $K$ -analytic (( $K$ -Lusin) space to be analytic (Lusin), it suffices already that there exists a sequence of pairs  $(C_n^1, C_n^2)$  of disjoint sets in  $\mathcal{F}_{sc}$ , such that  $\forall x_1, x_2 \in X, \exists n: x_1 \in C_n^1, x_2 \in C_n^2$ .

Indeed, let  $K_n^i$  be an u.s.c. map from  $\mathbb{T}$  to  $\mathcal{K}^X$  corresponding to  $(C_n^i)^c$ ;  $L_n: \{1, 2\} \times \mathbb{T} \rightarrow \mathcal{K}^X$ ,  $L_n(i, t) = K_n^i(t)$ ;  $L: S = [\{1, 2\} \times \mathbb{T}]^{\mathbb{N}} \rightarrow \mathcal{K}^{X^{\mathbb{N}}}$ ,  $L(i_1, t_1; i_2, t_2; \dots) = \prod_n L_n(i_n, t_n)$ .  $L$  is u.s.c., and  $L(S) = X^{\mathbb{N}}$ . Denote by  $\Delta$  the diagonal in  $X^{\mathbb{N}}$ , and let  $\psi(s) = L(s) \cap \Delta$ ,  $F = \{s \mid \psi(s) \neq \emptyset\}$ :  $F$  is a closed subset of  $\mathbb{N}^{\mathbb{N}}$ , and  $\psi$  is single-valued on  $S$ , thus a continuous map onto  $\Delta$  (which is homeomorphic to  $X$ ). This proves the  $K$ -analytic case. If  $X$  was  $K$ -Lusin, the above proves it is analytic, hence — cf. supra — there exists a weaker Hausdorff topology with a countable base  $O_k$ . So we can take for  $(C_n^1, C_n^2)$  all pairs of disjoint sets  $O_k$ . Their complements are then  $K$ -Lusin, so we can choose the maps  $K_n^i$  to have disjoint values. The map  $\psi$  is then one-to-one. This proves the lemma, and thereby our claim.

**5.c.** For a  $K$ -analytic (respectively  $K$ -Lusin) space to be analytic (Lusin), it suffices already that there exists a sequence  $U_n$  of open sets, such that for every pair of distinct points  $(x, y)$ ,  $x \in U_n$  and  $y \notin U_n$  for some  $n$ .

HINT. Since the projection from the graph  $G$  of the correspondence is continuous, and one-to-one in the  $K$ -Lusin case, it will suffice to show that  $G$  is analytic (Lusin) —cf. 3.e p. 422. The u.s.c. character of the map means that the projection  $\pi$  from  $G$  to  $\mathbb{T}$  is proper. Viewing  $\mathbb{T}$  homeomorphically as the irrationals in  $[0, 1]$ , this means that the extension  $\hat{\pi}: \widehat{G} \rightarrow [0, 1]$  is such that  $\hat{\pi}[\widehat{G} \setminus G]$  is the set of rationals. Let thus  $V_{2n+1}$  enumerate all open subsets of  $\widehat{G}$  of the form  $\{g \mid \hat{\pi}(g) > q\}$  and of the form  $\{g \mid \hat{\pi}(g) < q\}$  for  $q$  rational. Let also  $V_{2n}$  be the largest open set of  $\widehat{G}$  such that  $V_{2n} \cap G$  equals the inverse image of  $U_n$  in  $G$ . The open sets  $V_n$  have the property that for any pair of distinct points  $(x, y)$  of  $\widehat{G}$ , such that at least one of them lies in  $G$ , one has  $x \in V_n, y \notin V_n$  for some  $n$ . Stabilise the sequence  $V_n$  under finite intersections, and consider then all finite open coverings  $(V_n)_{i=1}^k$  of  $\widehat{G}$ : there are countably many of them. For each such cover, there exists a corresponding continuous partition of unity, i.e., continuous functions  $(f_i)_{i=1}^k$  with  $f_i \geq 0$ ,  $\sum_i f_i = 1$ ,  $\{x \mid f_i(x) > 0\} \subseteq V_n$ . We claim that the resulting countable family of continuous functions separates points of  $G$ : for  $x_1, x_2 \in G, x_1 \neq x_2$ , consider first open sets  $O_1, O_2$ , among the  $V_n$  with  $x_i \in O_i, x_j \notin O_i$ . For each point  $z \in \widehat{G} \setminus (O_1 \cup O_2)$ , let  $O_z^i$  be an open set among the  $V_n$  with  $z \in O_z^i, x_i \notin O_z^i$ , and let  $O_z = O_z^1 \cap O_z^2$ . Extract by compactness a finite subcovering from  $O_1, O_2$  and the  $O_z$ : since  $x_i$  belongs only to  $O_i$ ,  $O_1$  and  $O_2$  will belong to the subcovering, and the corresponding continuous functions  $f_i$  will satisfy  $f_i(x_i) = 1$  — and  $f_i(x_j) = 0$ . Thus we have a one-to-one continuous map  $\varphi$  from  $G$  to  $[0, 1]^{\mathbb{N}}$ . Let  $W_n$  be a basis of open sets in the latter: the sets  $\varphi^{-1}(W_n)$  are open, for each pair of distinct points  $(x_1, x_2)$  in  $G$  there exist two disjoint such sets  $U_1$  and  $U_2$  with  $x_i \in U_i$ . Since also  $G$  is  $K$ -analytic (( $K$ -Lusin)), the assumptions of the lemma sub(5.b) above are satisfied;  $G$  is analytic (Lusin).

**5.d.** For any analytic space  $X$  there exists a weaker analytic topology with countable base. If the space is regular, there exists both a weaker and a stronger metrisable analytic topology. The latter also exists if the space has a countable base, and can in any case be chosen as being an analytic subset of the Cantor space  $\{0, 1\}^{\mathbb{N}}$ . Those topologies can further be chosen such as to leave any given sequence of open sets open.

HINT. For the first, use that  $(X \times X) \setminus \Delta$  is Lindelöf (by (5.b)) (and that regularity implies complete regularity (3.d p. 422); hence, every open set being Lindelöf belongs to  $\mathcal{Z}_c$ ). To get the stronger metrisable topology, use the sequence  $(\mathbb{1}_{B_n})$  as a measurable one-to-one map to  $\{0, 1\}^{\mathbb{N}}$ , with continuous inverse — taking for  $B_n$  either the countable base of the space or the closures of the images of a countable base of  $\mathbb{T}$ . Use (5.e) below to deduce that the image is analytic.

**5.e.** Let  $X$  be  $K$ -analytic, with its bianalytic  $\sigma$ -field  $\mathcal{B} = \mathcal{F}_s \cap \mathcal{F}_{sc} = \mathcal{F}_{(\sigma,\delta)} \cap \mathcal{G}_{(\sigma,\delta)}$  (3.h p. 422). Let  $(E, \mathcal{E})$  be an analytic space with its Borel sets, and  $f: (X, \mathcal{B}) \rightarrow (E, \mathcal{E})$  a measurable map. Then  $f(X)$  is analytic. If  $f$  is one-to-one, it is a Borel isomorphism with  $f(X)$  (so  $X$  is analytic by (5.c) and (5.d)). If furthermore  $X$  is Lusin, then  $f(X)$  is a Borel set, and even Lusin if  $E$  is Lusin.

HINT. Because  $\mathcal{E}$  is separable and separating, the graph  $F$  of  $f$  belongs to  $(\mathcal{B} \times \mathcal{E})_{\sigma\delta}$ , hence is  $K$ -analytic. So  $f(X)$  is  $K$ -analytic, hence analytic by (5.c). This, applied to elements of  $\mathcal{B}$ , together with the first separation theorem, yields also the second statement. If  $E$  and  $X$  are Lusin, then  $F$  is Lusin by the above argument as a Borel subset of the Lusin space  $E \times X$ , so that the one-to-one projection yields that  $f(X)$  is Lusin — and in particular Borel (3.b, 3.g, 3.h). If just  $X$  is Lusin, to show that  $f(X)$  is still Borel, use (5.d) above to change the topology of  $E$ , without changing its Borel sets, first to an analytic topology with countable base, then to a metrisable analytic topology, which is a subspace of a (Lusin) compact metric space.

Remark that the same argument shows that for  $X$  Lusin,  $f$  one-to-one,  $f(X)$  will be Lusin if  $E$  is a Hausdorff space with countable base, or if  $E$  is a regular Hausdorff space which is the continuous image of a separable metric space — more generally if there exists a stronger topology on  $E$  with the same Borel sets, under which it is a subspace of Lusin space. But for  $E$  analytic? The problem reduces to: given a continuous map from  $\mathbb{T}$  onto the Cantor space, is the quotient topology Lusin?

**5.f.** 3.b, 3.g and 3.h imply that, in a Lusin space, the Lusin subsets are the Borel sets. Thus the Lusin subsets of a Hausdorff space are stable under countable unions.

**5.g.** Every bounded Borel measure on an analytic space is regular.

HINT. Use (5.d) and (4.d.2): observe that, in a space where there exists a weaker Hausdorff topology with countable base,  $\mathcal{K} \subseteq \mathcal{G}_\delta$ .

## 6. Blackwell spaces and standard Borel spaces

A Blackwell (standard Borel) space is a measurable space  $(E, \mathcal{E})$  where  $\mathcal{E}$  is the Borel  $\sigma$ -field of an analytic (Lusin) topology on  $E$ . It follows from (5.b) and (5.e) that all uncountable standard Borel spaces are isomorphic, and from (5.d) and (5.e) that every Blackwell space is isomorphic to an analytic subset of the Cantor space.

Given a measurable space  $(E, \mathcal{E})$ , define the equivalence relation  $R$  on  $E$  where two points of  $E$  are equivalent if they are not separated by  $\mathcal{E}$ . Call  $\mathcal{E}$  a Blackwell  $\sigma$ -field if the quotient space of  $(E, \mathcal{E})$  by this equivalence relation is a Blackwell space — the equivalence classes of  $E$  are also called the atoms of  $\mathcal{E}$ .

5.e implies then that, for  $X$   $K$ -analytic with bianalytic  $\sigma$ -field  $\mathcal{B}$ , all separable sub  $\sigma$ -fields of  $\mathcal{B}$  are Blackwell. And such a  $\sigma$ -field  $\mathcal{C}$  contains all elements  $B$  of  $\mathcal{B}$  which are a union of atoms of  $\mathcal{C}$ . [Consider the map from  $(X, B \vee \mathcal{C})$  to  $(X, \mathcal{C})$ ]. In other words, if  $f_1$  and  $f_2$  are two real-valued random variables on  $(X, \mathcal{B})$  — functions such that the sets  $\{x \mid f(x) \geq \alpha\}$  and  $\{x \mid f(x) \leq \alpha\}$  are analytic (i.e.,  $\in \mathcal{F}_s^X$ ) —, and if  $f_2(x)$  is a function of  $f_1(x)$ , then it is a Borel function:  $f_2 = h \circ f_1$  with  $h: \mathbb{R} \rightarrow \mathbb{R}$  Borel measurable. Similarly, 5.e implies that, if  $f$  is a measurable map from a Blackwell space  $(B, \mathcal{B})$  to a separable and separating measurable space  $(E, \mathcal{E})$ , then  $f(B)$  is a Blackwell space,  $f$  is an isomorphism with  $f(B)$  if  $f$  is one-to-one, and if in addition  $B$  is standard Borel then  $f(B) \in \mathcal{E}$ .

## 7. Spaces of subsets

Given a topological space  $X$ , the Hausdorff topology on  $\mathcal{K}^X$  has as basis of open subsets  $\{ K \in \mathcal{K}^X \mid K \subseteq \bigcup_{i \in I} O_i, K \cap O_i \neq \emptyset \forall i \in I \}$  for all finite families  $(O_i)_{i \in I}$  in  $\mathcal{G}^X$ .

Also, the Effros  $\sigma$ -field  $\mathcal{E}^X$  on  $\mathcal{F}^X$  is spanned by the sets  $\{ F \in \mathcal{F}^X \mid F \subseteq F_0 \} \forall F_0 \in \mathcal{F}^X$ .

**7.a.** If  $Y$  is a (closed) (open) subspace of  $X$ , so is  $\mathcal{K}^Y$  in  $\mathcal{K}^X$ . If  $X$  is compact or metrisable, so is  $\mathcal{K}^X$ . Hence if  $X$  is polish, or locally compact, so is  $\mathcal{K}^X$ .

**7.b.** If  $X$  is compact metric,  $\mathcal{E}^X$  is the Borel  $\sigma$ -field of  $\mathcal{K}^X$ .

**7.c.**  $\mathcal{E}^{\mathbb{T}}$  is the Borel  $\sigma$ -field of the topology with as subbase of clopen sets  $\{ F \in \mathcal{F}^{\mathbb{T}} \mid F \text{ is excluded by } h \}$  for  $h \in \mathbb{T}_f$ . Define  $d(F_1, F_2)$  as  $k^{-1}$ , where  $k$  is the smallest integer for which there exists  $h \in \mathbb{T}_f$  with sum of its terms  $\leq k$  such that either  $F_1$  is excluded by  $h$  and  $F_2$  not or vice versa. Show that this distance induces the above topology, and that  $(\mathcal{F}^{\mathbb{T}}, d)$  is complete: the topology is polish, and  $\mathcal{E}^{\mathbb{T}}$  is standard Borel.

**7.d.** If  $f: X_1 \rightarrow X_2$  is continuous, let, for  $F \in \mathcal{F}^{X_1}$ ,  $\phi(F) = \overline{f(F)}$ . Then  $\phi: \mathcal{F}^{X_1} \rightarrow \mathcal{F}^{X_2}$  is measurable. If  $f$  is an inclusion,  $\phi$  is an isomorphism with its image.

**7.e.** If  $f: X_1 \rightarrow X_2$  is open, let, for  $F \in \mathcal{F}^{X_2}$ ,  $\psi(F) = \overline{f^{-1}(F)}$ . Then  $\psi: \mathcal{F}^{X_2} \rightarrow \mathcal{F}^{X_1}$  is measurable. If  $f$  is continuous (so  $\psi(F) = f^{-1}(F)$ ) and onto,  $\psi$  is an isomorphism with its image. If furthermore  $X_1$  has a countable base, then  $\psi(\mathcal{F}^{X_2}) \in \mathcal{E}^{X_1}$ .

**7.f.** Conclude from (7.c) and (7.e) that  $\mathcal{E}^P$  is standard Borel for any polish space  $P$ .

**7.g.** Conclude from (7.d) and (7.f) — and App.6 above — that, if  $S$  is analytic, all separable sub- $\sigma$ -fields of  $\mathcal{E}^S$  are Blackwell.

[Note that this property of a measurable space  $(\Omega, \mathcal{C})$  is sufficient to imply all nice properties mentioned sub F about Blackwell spaces — including that  $\mathcal{C}_s \cap \mathcal{C}_{sc} = \mathcal{C}$  —; it is also this property which was obtained for the  $\sigma$ -field  $\mathcal{B}$  on a  $K$ -analytic space; characterisation: every real-valued random variable on  $(\Omega, \mathcal{C})$  has an analytic range.]

**7.h.** Show that the map  $c: \mathcal{F}^{\mathbb{T}} \rightarrow \mathbb{T}$  which selects in each closed set the lexicographically smallest element is Borel measurable (in fact, u.s.c. in the lexicographic order, which spans the Borel  $\sigma$ -field).

**7.i. Kuratowski-Ryll-Nardzewski selection theorem.** Conclude from (7.e) and (7.h) that, for any polish space  $P$ , there exists a Borel function  $c: \mathcal{F}^P \rightarrow P$  such that  $c(F) \in F \forall F \in \mathcal{F}^P$  s.t.  $F \neq \emptyset$ .

**7.j. von Neumann selection theorem.** Given a measurable space  $(\Omega, \mathcal{C})$  and a Blackwell space  $(B, \mathcal{B})$ , and given a subset  $A \in (\mathcal{C} \times \mathcal{B})_s$ , there exists a map  $\tilde{c}$  from  $\Omega$  to  $B$ , with  $(\omega, \tilde{c}(\omega)) \in A$  whenever possible, and which is  $((\mathcal{C}_s)_{(\sigma,c)} - \mathcal{B})$ -measurable.

HINT. Reduce to  $(\Omega, \mathcal{C})$  separable, then a subset of  $[0, 1]$ , then  $[0, 1]$ .  $B$  can also be viewed as an analytic subset of  $[0, 1]$ , so  $A$  becomes an analytic subset of  $[0, 1] \times [0, 1]$ . Let  $\varphi$  be the continuous map from  $\mathbb{T}$  onto  $A$ , and use  $(\pi_B \circ \varphi)(c[(\pi_\Omega \circ \varphi)^{-1}(\omega)])$  — check that  $\omega \mapsto (\pi_\Omega \circ \varphi)^{-1}(\omega)$  is  $((\mathcal{A}_s)_{(\sigma,c)} - \mathcal{E}^{\mathbb{T}})$ -measurable.

## 8. Some harder results

**8.a.** Assume  $Y$  is a regular Lusin space,  $S \subseteq X \times Y$  is bianalytic for the paving  $\mathcal{P} \times \mathcal{F}^Y$ , where  $\mathcal{P}$  is a paving on  $X$  and  $S_x \in \mathcal{K}_\sigma^Y \forall x \in X$ . Then  $\text{Proj}_X(S) \in \mathcal{P}_s \cap \mathcal{P}_{sc}$ .

**8.b.**

8.b.1. In particular, if also  $S_x \in \mathcal{F}^Y \forall x \in X$ , the map  $x \mapsto S_x$  is Effros measurable (cf. App.7) w.r.t.  $\mathcal{P}_s \cap \mathcal{P}_{sc}$ . 7.i yields then the existence of a  $\mathcal{P}_s \cap \mathcal{P}_{sc}$ -measurable selection, when  $Y$  is polish, or when  $S_x \in \mathcal{K}^Y$ .

8.b.2. Under the assumptions of 8.a, if in addition either  $Y$  is metrisable or  $S_x \in \mathcal{K}^Y \forall x \in X$ , there even exists a sequence of  $\mathcal{P}_s \cap \mathcal{P}_{sc}$ -measurable selections giving at every  $x$  a dense sequence in  $S_x$  — so that, if  $\mathcal{B}$  is a separable  $\sigma$ -field w.r.t. which these selections are measurable, one can construct a sequence of selections  $f_n$  such that the closure of  $\{f_n \mid n \in \mathbb{N}\}$  under point-wise convergence of sequences equals the set of all  $\mathcal{B}$ -measurable selections.

HINT. Use, for some weaker metrisable Lusin topology (sect. App.5), a theorem of Louveau implying that  $S = \bigcup_n S^n$ , where  $S^n$  is bianalytic and  $S_x^n \in \mathcal{K}^Y \forall x$ .

**8.c.** Let  $X$  be a polish space and  $\Theta: \wp(X) \rightarrow \wp(X)$ .

$\Theta$  preserves analytic sets if for any polish space  $W$  and any analytic  $A \subseteq X \times W$ ,

$$\Theta^*(A) = \{(y, w) \in X \times W \mid y \in \Theta(A_w)\}$$

is again analytic in  $X \times W$  (where  $A_w = \{x \in X \mid (x, w) \in A\}$ ).

Assume  $\Theta$  is an **analytic derivation**, i.e.:

- $\Theta$  is increasing
- $\Theta(A) \subseteq A$
- $\Theta$  preserves analytic sets.

Let  $\Theta^0(A) = A$ ,  $\Theta^{\alpha+1}(A) = \Theta(\Theta^\alpha(A))$  and  $\Theta^\beta(A) = \bigcap_{\alpha < \beta} \Theta^\alpha(A)$  if  $\beta$  is a limit ordinal. Let also  $\Theta^\infty(A) = \bigcap_\alpha \Theta^\alpha(A)$ , the largest fixed point of  $\Theta$  included in  $A$ .

MOSCHOVAKIS' THEOREM. (Moschovakis, 1980, 7c.8, p. 414)

- $\Theta^\infty$  is an analytic derivation. In particular  $\Theta^\infty(A)$  is analytic if  $A$  is so.
- If  $A$  is analytic  $\Theta^\infty(A) = \Theta^{x_1}(A)$
- If  $A$  is analytic,  $B$  is coanalytic and  $\Theta^\infty(A) \subseteq B$ , then  $\Theta^\alpha(A) \subseteq B$  for some countable  $\alpha$ .

## 9. Complements on Measure Theory

**9.a.** Any measurable function from a subset of a measurable space to a standard Borel space has a measurable extension to the whole space.

HINT. Take the standard Borel space to be  $[0, 1]$ : consider then first indicator functions, next step functions, finally arbitrary measurable functions (use e.g.  $\liminf$ ).

**9.b.** Assume  $f: X \rightarrow Y$  continuous,  $Y$  Hausdorff, and endow the spaces  $\Delta(X)$  and  $\Delta(Y)$  of regular probability measures on  $X$  and  $Y$  with the weak\* topology (1.10 p. 6). Let  $\bar{f}: \Delta(X) \rightarrow \Delta(Y)$  be the induced continuous map (1.16 p. 7).

9.b.1. If  $f$  is one-to-one, so is  $\bar{f}$ .

9.b.2. If  $f$  is an inclusion [i.e., a homeomorphism with  $f(X)$ ] (and  $f(X) \in \mathcal{F}^Y$ , or  $\mathcal{Z}^Y$ , or  $\mathcal{Z}_{cd}^Y$ ), so is  $\bar{f}$  — and  $\bar{f}(\Delta(X)) = \{\mu \in \Delta(Y) \mid \mu(f(X)) = 1\}$ .

9.b.3. If  $X$  is  $K$ -analytic, and  $f$  onto, so is  $\bar{f}$ .

**9.b.4.** If  $X$  is  $K$ -Lusin, and  $f$  is one-to-one and onto, then  $f$  is a Borel isomorphism.

HINT. For (9.b.1): two different measures differ already on some compact subset.

For (9.b.2), the formula for  $\bar{f}(\Delta(X))$  is clear. Given that  $\bar{f}$  is continuous and one-to-one (9.b.1), it will be an inclusion if we show that the image of a sub-basic open set  $\bar{f}\{\mu \in \Delta(X) \mid \mu(U) > \alpha\}$  with  $U$  open in  $X$  — is open in  $\bar{f}(\Delta(X))$ : this is because  $U = f^{-1}(V)$  for  $V$  open in  $Y$ , so that our set equals  $\{\nu \in \bar{f}(\Delta(X)) \mid \nu(V) > \alpha\}$ . Finally, if  $X$  is closed in  $Y$ ,  $\Delta(X)$  is so in  $\Delta(Y)$  by definition of the topology, if  $f: Y \rightarrow [0, 1]$ ,  $X = f^{-1}(0)$ , then  $\Delta(X) = \{\mu \in \Delta(Y) \mid \mu(f) = 0\}$ , and if  $X = \bigcap_n O_n$ ,  $O_n = \{y \mid f_n(y) > 0\}$  with  $f_n$  continuous with values in  $[0, 1]$ , then  $\Delta(X) = \bigcap_n \Delta(O_n) = \bigcap_n \bigcap_k \{\mu \in \Delta(Y) \mid \lim_{i \rightarrow \infty} \mu(f_{n,i}) > 1 - k^{-1}\}$  with  $f_{n,i} = \min(1, i \cdot f_n)$ , so  $\Delta(X) = \bigcap_n \bigcap_k \{\mu \mid \varphi_{n,k}(\mu) > 0\}$  where  $\varphi_{n,k}(\mu) = \sum_i 2^{-i}[\mu(f_{n,i}) - 1 + k^{-1}]^+$  is continuous.

For (9.b.3), fix  $\mu \in \Delta(Y)$ , and let, for  $A \subseteq X$ ,  $I(A) = \mu^*(f(A))$ : by 4.a, 4.b.1 and 4.b.7,  $I$  is a capacity on  $X$ . Hence by 4.c there exist compact subsets  $K$  of  $X$  such that  $f(K)$  has arbitrarily large measure. So  $\mu = \sum_1^\infty \alpha_n \mu_n$ , with  $\alpha_n \geq 0$  and  $\mu_n \in \Delta(Y)$  carried by some  $f(K_n)$ ,  $K_n$  compact in  $X$ . It suffices thus that  $\mu_n = \bar{f}(\nu_n)$ , for  $\nu_n \in \Delta(K_n)$ : this reduces the problem to the case where  $X$  — and hence  $Y$  — is compact. Then  $\bar{f}(\Delta(Y))$  is compact, as a continuous image of a compact set (1.12 p. 6), and contains all probability measures with finite support, since  $f$  is onto. Since those are dense [by the separation theorem (1.e p. 8)], the result follows.

For (9.b.4), let  $F \in \mathcal{F}^X$ ,  $F$  and hence  $f(F)$  are  $K$ -Lusin (3.a), so  $f(F)$  is Borel (3.g).

**9.c.** If  $X$  is  $K$ -analytic (resp.  $K$ -Lusin, resp. a  $\mathcal{Z}_{c\delta}$  subset of a compact space, resp. analytic, resp. Lusin, resp. polish), so is  $\Delta(X)$ .

HINT. For  $\mathcal{Z}_{c\delta}$  subsets, use 9.b.2 above and 1.12 p. 6. In the  $K$ -analytic ( $K$ -Lusin) case,  $X$  is a continuous (one-to-one) image of a  $\mathcal{Z}_{c\delta}$  subset  $G$  of a compact space (3.e p. 422), so  $\Delta(X)$  has the same property, by 9.b.3, (and 9.b.1) above, and the previous case — hence (3.a)  $\Delta(X)$  is  $K$ -analytic (( $K$ -Lusin). The argument for the latter three cases is the same, given that the compact space can then be taken metrisable (5.a, 5.b).

**9.d.**  $f: X \rightarrow Y$ , with  $Y$  a  $T_2$  space, is *universally measurable* iff it is  $\mu$ -measurable (cf. 1.11 p. 6)  $\forall \mu \in \Delta(X)$ .

$A \subseteq X$  is *universally measurable* iff  $\mathbb{1}_A$  is so (Recall (4.d.1)  $\mathcal{B}_u$  is the  $\sigma$ -field of universally measurable sets). Then:

9.d.1. Universally measurable maps are stable under composition and under countable products (i.e.,  $\prod_n f_n: \prod_n X_n \rightarrow \prod_n Y_n$  is universally measurable if each  $f_n$  is).

9.d.2. If  $f: X \rightarrow Y$  is universally measurable, then  $[(\Delta(f))(\mu)](B) = \mu(f^{-1}(B))$  defines  $\Delta(f): \Delta(X) \rightarrow \Delta(Y)$ .

9.d.3. If  $Y$  is a separable metric space,  $f: X \rightarrow Y$  is universally measurable iff  $f^{-1}(B) \in \mathcal{B}_u$  for every Borel set  $B$ .

**9.e.** Endow the space  $M_X$  of probability measures on a measurable space  $(X, \mathcal{X})$  with the  $\sigma$ -field spanned by the functions  $\mu \mapsto \mu(Y)$  for  $Y \in \mathcal{X}$ .

A transition probability to  $(X, \mathcal{X})$  is then just a measurable map to  $M_X$ .

If  $X$  is analytic, the Borel  $\sigma$ -field on  $\Delta(X)$  coincides with the above defined one for  $\mathcal{X}$  the Borel  $\sigma$ -field on  $X$ . Thus if  $X$  is standard Borel, or Blackwell, so is  $M_X$ .

HINT.  $\mathcal{X}$  being separable if  $X$  is analytic, our  $\sigma$ -field on  $M_X$  will be separable and separating. It is coarser than the Borel  $\sigma$ -field, because  $\mu \mapsto \mu(Y)$  is Borel measurable for the weak\*-topology. Therefore both coincide by (5.e) and the analyticity of  $\Delta(X)$  (9.c above). Use again (9.c) for the final conclusion.

**9.f.** Given a transition probability  $P$  from  $(E, \mathcal{E})$  to  $(F, \mathcal{F})$ , and  $X \in \mathcal{F}_s$ ,  $\{e \in E \mid P_e(X) > \alpha\}$  and  $\{e \in E \mid P_e(X) \geq \alpha\}$  belongs to  $\mathcal{E}_s$ .

REMARK 9.1. Since  $P$  can also be viewed as a transition probability from  $E$  to  $E \times F$ , the same result holds if  $X \in (\mathcal{E} \otimes \mathcal{F})_s$ .

HINT.  $X$  already belongs to  $(\mathcal{F}_0)_s$  for a separable sub  $\sigma$ -field  $\mathcal{F}_0$  of  $\mathcal{F}$ ; for some separable sub  $\sigma$ -field  $\mathcal{E}_0$  of  $\mathcal{E}$ ,  $P$  will still be a transition probability from  $\mathcal{E}_0$  to  $\mathcal{F}_0$ , and it suffices to show the sets are in  $(\mathcal{E}_0)_s$ . There is no loss then in passing to the quotient, so  $E$  and  $F$  can be viewed as subsets of  $[0, 1]$  with the Borel sets.  $P$  can then be viewed as a transition probability from  $E$  to  $[0, 1]$ , (under which  $F$  has outer probability one), so can be extended (by (9.a) and (9.e) above) as a transition probability from  $[0, 1]$  to  $[0, 1]$ .  $X$  is the trace on  $F$  of an analytic subset  $\overline{X}$  of  $[0, 1]$ . Further,  $P_e(X) = P_e(\overline{X})$  for  $e \in E$  — indeed for a Borel set  $B$  of  $[0, 1]$  with  $X \subseteq B$  and  $P_e(X) = P_e(B)$  one has  $P_e(\overline{X} \setminus B) = 0$ , because any compact subset of  $\overline{X} \setminus B$  is negligible:  $P_e(X) \geq P_e(\overline{X})$ . Similarly (with  $B \cap F \subseteq X$ ) one obtains  $P_e(X) \geq P_e(\overline{X})$ . Therefore,  $\{e \in E \mid P_e(X) > \alpha\} = E \cap \{e \in [0, 1] \mid P_e(\overline{X}) > \alpha\}$ : it suffices to prove the result for  $E$  and  $F$  compact metric — or (again 9.e), that if  $X$  is analytic in a compact metric space  $F$ ,  $M_\alpha = \{\mu \in M_F \mid \mu(X) > \alpha\}$  is analytic. Let  $X$  be the projection of a Borel set in  $F \times [0, 1]$ :  $M_\alpha$  is the projection of the Borel set  $\{\mu \in M_{F \times [0, 1]} \mid \mu(B) > \alpha\}$  (7.j), hence analytic.

## 10. \*-Radon Spaces

**10.a.** A  $\tau$ -Radon space is a  $T_2$  space where every probability measure  $\mu$  on the Borel sets satisfying  $\mu(\bigcup_\alpha O_\alpha) = \sup_\alpha \mu(O_\alpha)$  for every increasing net of open sets is regular.

A quasi-Radon space is a  $T_2$  space where, for  $P \in \Delta(\Delta(X))$ , the barycentre  $\bar{P} = \beta(P)$  defined by  $\bar{P}(B) = \int \mu(B)P(d\mu)$  for every Borel set  $B$  (observe  $\mu(B)$  is a Borel function of  $\mu$ ) — is regular.

10.a.1. For the regularity of  $\mu$  (or of  $\bar{P}$ ), it suffices already that  $\sup_{K \in \mathcal{K}} \mu(K) = 1$ .

10.a.2. A  $\tau$ -Radon space is quasi-Radon.

HINT. Observe  $\mu(O_\alpha)$  is l.s.c.

A  $K$ -analytic space is  $\tau$ -Radon.

HINT. Observe  $\mu$  also satisfies  $\mu(K) = \inf\{\mu(O) \mid K \subseteq O \in \mathcal{G}_X\} \forall K \in \mathcal{K}^X$  ( $T_2$ -assumption), hence use 4.d.2 p. 424.

10.a.3. If  $A$  is universally measurable in  $X$ , and  $X$  is quasi- or  $\tau$ -Radon, so is  $A$ .  $\tau$ -Radon subspaces are universally measurable.

10.a.4. For  $f: X \rightarrow Y$  universally measurable,  $\Delta(f)$  — cf. 9.d.2 — is universally measurable if  $X$  is quasi-Radon.

HINT. For  $P \in \Delta(\Delta(X))$ , find  $K_n \in \mathcal{K}^X$  increasing with  $\bar{P}(K_n) \rightarrow 1$  and  $f|_{K_n}$  continuous: then  $\mu(K_n) \nearrow 1$   $P$ -a.e. By Egorov's theorem [or 9.d.3, with  $Y = C(\mathbb{N} \cup \{\infty\})$ ],  $\exists C \in \mathcal{K}^{\Delta(X)}$  s.t.  $P(C) \geq 1 - \varepsilon$  and  $\mu(K_n)|_C$  [is continuous and] converges uniformly to 1. To show the continuity of  $[\Delta(f)]|_C$ , choose  $F$  closed in  $Y$ : we have to show that  $\mu(f^{-1}(F))$  is u.s.c on  $C$ . Since  $\mu(K_n) \rightarrow 1$  uniformly, it suffices to show that  $\mu(K_n \cap f^{-1}(F)) = \mu([(f|_{K_n})^{-1}(F)])$  is u.s.c., which follows from the continuity of  $f|_{K_n}$ .

10.a.5. quasi- or  $\tau$ -Radon spaces are closed under countable products.

10.a.6. For  $X$  quasi-Radon,  $\beta: P \mapsto \bar{P}$  is a continuous map from  $\Delta(\Delta(X))$  to  $\Delta(X)$ .

10.a.7. For  $X$  quasi- or  $\tau$ -Radon, so is  $\Delta(X)$ .

HINT. Assume first  $X$   $\tau$ -Radon, let  $\mu$  be an appropriate (i.e., as in the definition — also called a “ $\tau$ -smooth” measure in the literature) measure on  $\Delta(X)$ . Then  $\bar{\mu}$  is appropriate on  $X$ , hence  $\bar{\mu} \in \Delta(X)$ . Choose  $C_i \in \mathcal{K}^X$  disjoint with  $\bar{\mu}(K_n) \rightarrow 1$ , where  $K_n = \bigcup_{i \leq n} C_i$ . By Egorov's theorem, we have thus that,  $\forall \varepsilon > 0, \exists \delta_n \searrow 0: \mu(C) > 1 - \varepsilon$  with  $C = \{\nu \in \Delta(X) \mid \nu(K_n) \geq 1 - \delta_n, \forall n\}$ .  $C$  is compact in  $\Delta(X)$  — e.g. as the continuous image of the corresponding set on the (locally compact) disjoint union of the  $C_i$ 's. Hence, by (10.a.1),  $\mu \in \Delta(\Delta(X))$ :  $\Delta(X)$  is  $\tau$ -Radon.

Assume now  $X$  quasi-Radon, and fix  $P \in \Delta(\Delta(\Delta(X)))$ . Then  $\bar{P}$  is an appropriate measure on  $\Delta(X)$ , hence  $\bar{\bar{P}}$  on  $X$ . Also, with the continuous map  $\beta$  of 10.a.6 above from  $\Delta(\Delta(X))$  to  $\Delta(X)$ , we have  $[\beta \circ (\Delta(\beta))](P) = \bar{\bar{P}}$  — compute for each Borel set  $B$ . So  $\bar{\bar{P}} \in \Delta(X)$ . Finish now as in the  $\tau$ -Radon case, with  $\bar{\bar{P}}$  instead of  $\mu$ .

**10.b.** A countably Radon space is a  $\tau$ -Radon space with countable base.

10.b.1. In a countably Radon space, every probability measure on the Borel sets is regular (use the countable base to show it is appropriate). (Such spaces are called Radon in (Bourbaki, 1969)).

10.b.2. Countably Radon spaces are closed under countable products, and by taking universally measurable subspaces.

10.b.3. If  $X$  is countably Radon, so is  $\Delta(X)$  — and the Borel  $\sigma$ -field on  $\Delta(X)$  is the smallest making the functions  $\mu \mapsto \mu(B)$  —  $B$  Borel in  $X$  — measurable.

HINT. A sub-basis for the topology on  $\Delta(X)$  are the sets  $\{\mu \mid \mu(O) > \alpha\}$  for  $O$  open in  $X$ . One can then clearly further restrict  $\alpha$  to be rational and  $O$  to be a finite union of basic open sets in  $X$ :  $\Delta(X)$  has a countable sub-basis — hence is countably Radon by 10.a.7. Another consequence is that every open set, and hence every Borel set, belongs to the  $\sigma$ -field spanned by the sub-basic open sets — hence to the  $\sigma$ -field making all functions  $\mu \mapsto \mu(O)$  measurable.

10.b.4. If  $X$  is countably Radon, there exists a stronger topology with the same Borel sets under which it becomes a universally measurable subset of the Cantor space.

HINT. Declare all basic open sets to be clopen. To show universal measurability of the image, choose  $\mu$  giving outer measure 1 to the image:  $\mu$  induces a probability measure on the space itself, which is regular (10.b.1), hence carried by a  $\mathcal{K}_\sigma$ -subset — and the compacts are metrisable by the countable basis assumption, so the  $\mathcal{K}_\sigma$  is Lusin (5.f p. 426), hence (5.e p. 426) it is Borel in the Cantor space: the image has also  $\mu$ -inner probability one.

10.b.5. A map from a topological space to a countably Radon space is universally measurable iff the inverse image of every Borel set (and then also of every universally measurable set) is universally measurable.

HINT. Use 10.b.4 and 9.d.3.

10.b.6. If a map from a countably Radon space to a Hausdorff space is universally measurable, then it is still so for any other countably Radon topology with the same universally measurable sets.

HINT. Apply (10.b.5) to the identity from the space to itself, and 9.d.1. Alternatively, observe such a map is universally measurable iff  $f^{-1}(B) \in \mathcal{B}_u$  for every Borel set  $B$  and,  $\forall \mu \in \Delta(X)$ , there exists a union of a sequence of disjoint compact metric subsets which has measure 1 under  $f(\mu)$  (an analytic subset would already be sufficient) — use (5.d) the fact that a compact analytic space is metrisable.

REMARK 10.1. By (10.b.1), two countably Radon topologies that have the same Borel  $\sigma$ -field also have the same  $\sigma$ -field  $\mathcal{B}_u$ . (10.b.5) and (10.b.6) say that, if they have the same  $\mathcal{B}_u$ , they have the same universally measurable maps to and from all topological spaces, by (10.b.1) they have also the same set  $\Delta(X)$ , with the same universally measurable sets on  $\Delta(X)$  (the sets which are  $\mu$ -measurable for every probability measure  $\mu$  on  $\Delta(X)$  endowed with the  $\sigma$ -field spanned by the maps  $P \mapsto P(B)$ , for  $B$  universally measurable in  $X$ ). Also, the  $\sigma$ -field  $\mathcal{B}_u$  on a countable product depends only on those on the factors (recall the countable base ...).

REMARK 10.2. To make the picture complete, one would still like an example of such a space, which is not countable, and for which there exists no universally measurable isomorphism with  $[0, 1]$  (or a theorem to the opposite effect).

REMARK 10.3. The concept is not quite satisfactory in the sense that analytic spaces have the same properties (5.g for 10.b.1, use 5.d and 10.b.1 to prove 10.b.5 and hence 10.b.6, 10.b.2 for closed, or analytic, subspaces is obvious, and use 9.e). In addition, they lead to a more restrictive set of measurable spaces (sect. App.6), so one would have liked a concept here that would include all analytic spaces — just like  $K$ -analytic spaces are  $\tau$ -Radon (10.a.2).

10.b.2, 10.b.3, 10.b.5 and 10.b.6 remain true for “countably quasi-Radon” spaces, and in 10.b.4 one would just have to drop the “universally measurable”.

## APPENDIX B

### Historical Notes

#### 1. Chapter I

**1.a. Section 1.** The original proof of Sion's theorem (1958) (theorem 1.6) uses the KKM lemma (ex. I.4Ex.18) in  $\mathbb{R}^n$  (which is equivalent to the fixed point theorem). However he wrote “the difficulty lies in the fact that we cannot use a fixed point theorem (due to lack of continuity) nor the separation of disjoint convex sets by a hyperplane (due to lack of convexity)”.

Also in Sion's paper is the proof that his theorem implies Fan's theorem (1953) (cf. the remark after prop. 2.10).

The proof of Sion's theorem using lemma 1.7, as well as the lemma itself and its proof, appear in (Berge, 1966, p. 220, resp. p. 172). In fact he wrote: “Signalons aussi que c'est la démonstration de Sion qui nous a suggéré l'énoncé du théorème de l'intersection”. On the other hand the original proof of the lemma (Berge, 1959) uses also KKM.

A direct proof of prop. 1.17 was obtained by Kneser (1952) (all these results in the Hausdorff case).

A survey and specific results for games on the square can be found in (Yanovskaya, 1974).

**1.b. Section 2.** The proof of ex. I.2Ex.17 is due to Karamata (cf. e.g. Titchmarsh, 1939, p. 227).

**1.c. Section 3.** The original finite minmax theorem (ex. I.3Ex.1) can be proved by “elementary tools” namely the theorem of the alternative (von Neumann and Morgenstern, 1944, pp. 138, 154–155) (also Ville, 1938): the iterated elimination of variables implies the existence of optimal strategies (and a fortiori the value) in the ordered field of coefficients. The first analysis in this framework and theorem 3.8 are due to Weyl (1950), cf. ex. I.3Ex.13. This elementary aspect was used later by Bewley and Kohlberg (1976b) in analysing stochastic games.

Another elementary proof by induction on the size of the matrix is due to Loomis (1946).

**1.d. Section 4.** Theorem 4.1 was proved in the finite case using Kakutani's (1941) fixed point theorem (Nash, 1950) and Brouwer's fixed point theorem in (Nash, 1951). He constructed explicitly an “improving mapping” from  $\sigma$  to  $\tau$  as follows:  $\tau_i(s_i) = \{\sigma_i(s_i) + (F^i(s_i, \sigma_{-i}) - F^i(\sigma))^+\}/\{1 + \sum_{t_i \in S_i} (F^i(t_i, \sigma_{-i}) - F^i(\sigma))^+\}$ .

Glicksberg (1952) extends Kakutani's theorem and the first proof of Nash and hence obtains theorem 4.1. A similar extension of Kakutani's theorem is due to Fan (1952).

Ex. I.4Ex.17 bears a clear relation with the fixed point theorems of Eilenberg and Montgomery (1946) and Begle (1950).

Debreu (1952) uses this result to prove an equilibrium theorem for a game with constraints, namely where the set of feasible outcomes is a subset of the product of the strategy spaces.

Other related topics include purification of strategies and equilibria (Aumann et al., 1983),  $\varepsilon$ -equilibria (cf. e.g. Tijs, 1981), and the study by Blume and Zame (1994) of algebraic-geometrical aspects of the manifold (cf. ex. I.4Ex.4) of equilibria.

## 2. Chapter II

**2.a. Section 1.** The initial definition of Kuhn (1953) extends the approach of von Neumann and Morgenstern (1944, pp. 67–79). In the former, a sequence of dates is associated to the nodes and is public knowledge: hence it is our model of a multistage game, except that only a single player moves at every node. This means, the following is impossible:

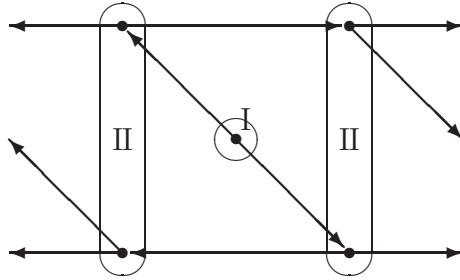


FIGURE 1. Perfect recall, and not multistage.

Isbell's construction (1957) extends Kuhn's definition (which considers only linear games) and essentially corresponds to the notion of tree described in ex. II.1Ex.8 (in the finite case).

An extension of theorem 1.4 to the infinite case is in (Aumann, 1964). Distributional strategies as defined in ex. II.1Ex.16c appear in (Milgrom and Weber, 1985).

**2.b. Section 2.** For the use of games with perfect information in descriptive set theory Jech (cf. 1978, ch. 7), Dellacherie et al. (1992, ch. XXIV).

On topological games, cf. the survey of Telgársky (1987).

**2.c. Section 3.** Correlated equilibria are due to Aumann (1974). The canonical representation appears explicitly in (Aumann, 1987) but was known and used before; similar ideas can be found in the framework of coordination mechanisms (cf. Myerson, 1982). Communication equilibria and extensive form correlated equilibria were first introduced in the framework of non zero-sum repeated games with incomplete information (Forges, 1982a), (Forges, 1985), (Forges, 1988a); then formally defined and studied for themselves in (Forges, 1986a). They also appear in (Myerson, 1986).

**2.d. Section 4.** Blackwell's theorem plays a crucial rôle in games with incomplete information, cf. ch. V and VI.

The fact that any set is either weakly-approachable or weakly-excludable (cf. ex. II.4Ex.2) has been proved by Vieille (1992).

The results of ex. II.4Ex.8 have been extended to game forms in (Abdou and Mertens, 1989).

### 3. Chapter III

Harsanyi (1967) (1968a) (1968b) made precise the difference between games with incomplete information (lack of information about the description of the game), and games with imperfect information (lack of information about the play of the game). In the first framework he remarked that this leads to an infinite hierarchy of beliefs, called “sequential expectation model”. Then he proposed to represent the situation through a consistent probability on the product of the set of types by the set of states, each type being identified with a probability distribution on the product other types  $\times$  states. He also observes that non-consistent situations may occur. In the first case the above reduction amounts to represent a game with incomplete information as a game with imperfect information, adding a move of nature. Then several interpretations are possible: the usual one is that each player’s type is chosen at random. Another way to look at it is to consider each type as a player, the actual players in the game being chosen at random — one in each “group” — according to the initial probability of the types (Selten game).

Böge and Eisele (1979) consider games with unknown utility functions and construct the hierarchy of beliefs under the name “system of complete reflexion” (and refer to a previous construction of Böge called “oracle system”), cf. Ex. IIIEx.3b.

The construction of the universal belief space [i.e., point 3 of theorem 1.1] is due to Mertens and Zamir (1985), and the content of this paper is the basis of this chapter. They assume  $K$  compact and proceed as in Ex. IIIEx.2. Further constructions include (Brandenburger and Dekel, 1993), in the polish case, and (Heifetz, 1993), which is similar to the one exposed here. The present treatment is the first to offer a characterisation [points 1 and 2 of theorem 1.1].

Theorem 2.5 was given informally in (Mertens, 1986b). The rest of this chapter is new.

Lemma 3.2 extends theorem 2 in (Blackwell and Dubins, 1962).

Further relations between common knowledge and belief subspaces can be found in (Vassilakis and Zamir, 1993).

### 4. Chapter IV

The general model of repeated games appears in (Mertens, 1986b) and sect. 2 and 3 are based on this paper.

A specific version of theorem 3.2 has been obtained by Armbruster (private communication, 1983).

**4.a. Section 4.** The “Folk Theorem” (4.1) was already known in the sixties by Aumann and Shapley among others. Theorem 4.2 is in (Sorin, 1986b) (but was probably known before), as well as prop. 4.3.

Most of the results of sect. 4.b can be found in Lehrer except prop. 4.5. The content of 4.b.2 is in (Lehrer, 1992a), 4.b.3 is in (Lehrer, 1992c), 4.b.4 in (Lehrer, 1990).

The notion of subgame (cf. subsection 1.a) has been emphasised by Selten who introduced the notion of subgame perfection: cf. (Selten, 1975) and ex. IV.4Ex.3–IV.4Ex.8.

There has been a tremendous amount of literature on refinement of equilibria and their applications to repeated games, cf. e.g. the surveys by Kohlberg (1990) and by Van Damme (1992). Somehow the structure of specific — namely public perfect — equilibrium pay-offs in discounted games is easier to characterise (in the spirit of ex. IV.4Ex.5), cf. (Fudenberg and Levine, 1994).

???

Ex. IV.5Ex.4 is a typical example of “recursive structure” and can be found in (Blackwell and Girshick, 1954, pp. 69–73).

## 5. Chapter V

Most of the basic results of this chapter (until sect. 3.c) are introduced and proved by Aumann and Maschler (1966) for the full monitoring case and (1968) for the general case. The construction and proof of sect. 3.d are due to Kohlberg (1975b).

Theorem 3.1 is new.

Sect. 3.e–3.h are new.

Section 4 is based on the papers of Mertens and Zamir (1976a) for theorem 4.1, (1977b) for theorem 4.3 and (1995) for sect. 4.c.

The recursive formula (lemma 4.2) appears in (Zamir, 1971–1972).

Section 5 uses the results of Zamir (1973a) for 5.a and 5.b and (1971–1972) for 5.c.

Ex. VEx.5 is due to Mertens and VEx.7 to Sorin (1979). Ex. VEx.10 comes in part from Zamir (1971–1972) and (1973b). Ex. VEx.17 is an extension of Stearns’ measure of information (1967).

## 6. Chapter VI

The basic model and the first results are due to Aumann and Maschler (1967) and Stearns (1967). They consider the independent case with standard signalling and proved theorem 3.1 (existence and computation of the minmax) in this framework. The construction of the strategy exhausting the information (lemma 3.3) is due to Stearns. The extension to the dependent case with general (state independent) information structure is due to Mertens and Zamir (1980) and is exposed in sect. 3.a.

Section 3.b is new.

Previously Mertens and Zamir (1971) proved the existence of  $\lim v_n$  with standard signalling; then Mertens (1971–1972) in the general case: this is the content of sect. 4. Corollary 4.9 gives a better bound ( $n^{1/3}$ ) on the speed of convergence than the original one ( $n^{1/4}$ ).

Section 5 is in Mertens and Zamir (1971) and (1977a).

Part of the results of sect. 7 can be found in (Mertens and Zamir, 1981). Example 7.4 is in (Mertens and Zamir, 1971) as example 8.5.

Ex. VIEx.6 is due to Ponssard and Sorin (1980a), (1980b). Ex. VIEx.7 comes from Sorin (1979) (cf. also Ponssard and Zamir, 1973) and Ponssard and Sorin (1982). Ex. VIEx.8 comes from (Sorin and Zamir, 1985).

The material of ex. VIEx.9 can be found in (Mertens and Zamir, 1977a), and (Sorin, 1984b) for the dependent case.

## 7. Chapter VII

**7.a. Section 1.** Stochastic games were introduced by Shapley (1953) who considered the finite two-person zero-sum case with a strictly positive probability of stopping the game in each state. He proves the existence of a value and of optimal stationary strategies; the result, which basically corresponds to the discounted case was then extended by Takahashi (1962), Parthasarathy (1973), Couwenbergh (1980), Nowak (1985b) among others.

The non-zero-sum discounted case was studied by Fink (1964) and Takahashi (1964), Federgruen (1978), Sobel (1971). For a survey cf. (Parthasarathy and Stern, 1977) and

(Parthasarathy, 1984).

The existence of stationary optimal strategies is still an open problem, cf. (Nowak and Raghavan, 1992) and (Parthasarathy and Sinha, 1989). For  $\varepsilon$ -equilibria cf. (Rieder, ???, 1979), (Whitt, 1980), (Nowak, 1985a); they basically use approximation by games that do have stationary equilibria.

The content of 1.c is due to Mertens and Parthasarathy (1987) and (1991).

**7.b. Section 2.** The first papers devoted to the asymptotic analysis are due to Bewley and Kohlberg (1976b). They worked in the field of Puiseux series to get theorem 2.3, in the spirit of theorem 3.8.

For similar results concerning the limit of finite games, cf. ex. VIIEx.2 and VIIEx.6 and (Bewley and Kohlberg, 1976a).

**7.c. Section 3.** The undiscounted case was introduced by Everett (1957) who proves the existence of a value in the irreducible case (ex. VIIEx.10), (cf. also (Hoffman and Karp, 1966) and in the perfect information case (Liggett and Lippman, 1969)), and suggest an example that was later solved by Blackwell and Ferguson (1968). The ideas introduced for studying this game (“Big Match”, ex. VIIEx.4) were deeply used in further work: existence of the value for games with absorbing states (Kohlberg, 1974), and finally in the general case. This is the content of sect. 3 and is due to Mertens and Neyman (1981).

Some of the results of sect. 3.d can be found in (Bewley and Kohlberg, 1978) and (Vrieze et al., 1983).

The content of sect. 4.a is in (Sorin, 1986a), and of sect. 4.b in (Vrieze and Thuijsman, 1989).

Part of ex. VIIEx.9 was in some unpublished notes of Neyman.

More recent results on dynamic programming problems can be found in (Lehrer and Monderer, 1994), (Lehrer and Sorin, 1992), (Monderer and Sorin, 1993) and (Lehrer, 1993).

## 8. Chapter VIII

Section 1 is due to Forges (1982b) after a first proof in the deterministic case in (Kohlberg and Zamir, 1974).

Section 2 corresponds to a class introduced by Mertens and Zamir (1976b) and solved for the minmax and maxmin in a specific  $2 \times 2$  case. This result was generalised by Waternaux (1983b) to all  $2 \times 2$  games, then to the general case (1983a). This is the content of 2.c. (The 1983b paper contains also a more precise description of optimal strategies).

Part 2.d is due to Sorin (1989).

Section 3 follows (Sorin, 1985b) and the content of sect. 4 can be found in (Sorin, 1984a), (Sorin, 1985a) and (Sorin and Zamir, 1991).

## 9. Chapter IX

The first approach to non-zero-sum games with lack of information on one side is due to Aumann et al. (1968). They introduce the notions of joint plan and of jointly controlled lottery.

They gave a sufficient condition for equilibria as in prop. 1.1 with condition 3 replaced by the stronger condition  $a^k(r) = \bar{a}^k$ . They also exhibit examples of equilibrium pay-offs that require several stages of signalling.

The content of 1.a is in (Sorin, 1983).

Section 1.b follows (Hart, 1985).

The content of sect. 2 is taken from (Aumann and Hart, 1986).

The results of sect. 3 are due to Forges (1988b). More precise results concerning a specific class (information transmission game) can be found in (Forges, 1985): the set of communication equilibrium pay-offs equals the set of (normal form) correlated equilibrium pay-offs. For a parallel study in the Selten representation, cf. (Forges, 1986b).

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